

## BINARY SEARCH TREE

"DYNAMIC SET OPERATIONS"

DICTIONARY: INSERT, SEARCH, DELETE

WHAT IF WE WANT MORE:

- MIN/MAX
- ELEMENT BEFORE OR AFTER
- TRAVERSAL

WE WANT  
 $O(\lg n)$   
 WORST CASE:  
 $O(n)$

BEST:



WORST:



ENDS UP LIKE  
LINKED LIST

TYPICALLY:

$O(h)$  WHERE  $h$  IS HEIGHT

TREES WITH SELF-HEALING

- ① AVL TREE
- ② RED-BLACK TREE

BST → BINARY TREE THAT RESPECTS THE BST PROPERTY

RELATES THE VALUES OF A NODE TO IT'S CHILDREN  
 ASSUMING UNIQUE KEYS:

$x.\text{LEFT}.\text{KEY} < x.\text{KEY}$

$x.\text{RIGHT}.\text{KEY} > x.\text{KEY}$



DSA [04.10.24]

## BST FUNCTIONS:

→ RECURSIVE?

→ MAY HAVE INTERNAL POINTER OR REFERENCE AS PARAM.

→ WE MAY NEED A PRIVATE AND PUBLIC PART FOR EACH FUNCTION.

→ MIGHT WANT WRAPPER FUNCTIONS

PRIVATE:

`NODE* SEARCH(NODE* NODE, INT KEY) // O(h)`

IF (`NODE == NULL` || `NODE->KEY == KEY`) // FOUND & NOT FOUND

RETURN `NODE`;

IF (`KEY < NODE->KEY`)

RETURN `SEARCH(NODE->LEFT, KEY)`;

ELSE

RETURN `SEARCH(NODE->RIGHT, KEY)`;

}

PUBLIC:

`BOOL FOUND(INT KEY)`

RETURN (`SEARCH(ROOT, KEY) != NULL`);

}

PRIVATE:

`NODE MIN(NODE* NODE)` { // O(h)

IF (`NODE == NULL`) // BOUNDS

RETURN `NULL`;

WHILE (`NODE->LEFT != NULL`)

`NODE = NODE->LEFT`;

RETURN `NODE`;

}

`NODE MAX(NODE* NODE)`; // SAME AS MIN BUT MOVE RIGHT.

5.

`SUCCESSOR(NODE* NODE)`

IF (`NODE == NULL`) RETURN `NULL`; // BOUNDS CHECK

IF (`NODE->RIGHT != NULL`) RETURN `MIN(NODE->RIGHT)`; // CASE #1



# DSA

```
NODE SUCCESSOR (NODE * NODE) {  
    IF (NODE == NULL) RETURN NULL;  
    IF (NODE.RIGHT != NULL) } CASE #1.  
        RETURN MIN (NODE.RIGHT);  
    NODE TEMP = NODE.PARENT;  
    WHILE (TEMP != NULL && NODE != TEMP.RIGHT) {  
        NODE = TEMP  
        TEMP = TEMP.PARENT  
    }  
    RETURN TEMP  
}
```

FIND LOWEST ANCESTOR WHOSE LEFT CHILD IS ALSO AN ANCESTOR OF MINE.