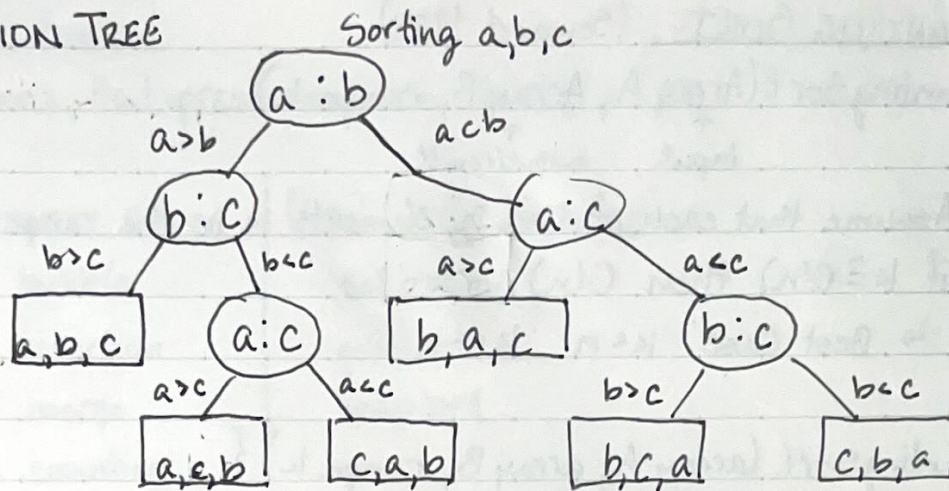


## DECISION TREE



- minimum of  $n!$  leaves
  - height  $h = 2^h$  leaves
  - $h \geq \lg(n!)$
  - $h \geq (n \lg n)$
- ↳ This is why comparison-based sorts are lower-bounded at  $O(n \log n)$

Valgrind ./p1 input  
 ↑ To test for mem  
 leaks



## COUNTING SORT (Seward 1954)

CountingSort(Array A, Array B, range k)

↑                    ↑  
input            output/result

- Assume that each of the n elements is in the range  $0 \dots k$   
if  $k \approx O(n)$  then  $O(n)$  sort  
↳ Best Case:  $k < n$      $k \ll n$

Code:

```
countingSort(array A, array B, range k) {
```

Voting  
Section

```
  newarray C[k+1] = {0};
```

```
  for (i=0; i < A.length(); i++) {
```

```
    C[A[i]]++;
```

```
  }
```

Sum the  
votes to

```
  for (i=1; i <= k; i++) {
```

```
    C[i] += C[i-1];
```

```
  }
```

find indexes

The sort

(based on

accumulated

knowledge)

↑

```
  for (i=A.length()-1; i >= 0; i--) {
```

```
    B[C[A[i]]-1] = A[i];
```

```
    C[A[i]]--;
```

```
  }
```

```
}
```

MTG REF!

Example

A: 

0 <sub>A</sub>	3 <sub>A</sub>	1 <sub>A</sub>	2 <sub>A</sub>	2 <sub>A</sub>	1 <sub>A</sub>	0 <sub>A</sub>	0 <sub>A</sub>	1 <sub>A</sub>	2 <sub>A</sub>	3 <sub>A</sub>	3 <sub>A</sub>
----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------	----------------

0    1    2    3    4    5    6    7    8    9    10    11

k=3    0...3

C: 

0	0	0	0
---	---	---	---

 =>    1    1

B: 

--	--	--	--	--	--	--	--	--	--	--	--

0    1    2    3    4    5    6    7    8    9    10    11



STABLE SORT

- A sort that preserves order in case of ties

<u>STABLE SORTS</u>	<u>(UN)(NON) STABLE</u>
bubble	selection
insertion	quicksort
merge	heapsort
counting	
bucket	

RADIX SORT (Comrie 1929)

- 1 "digit" at a time from lsb to msb

```

radixSort(A, d) {
  for (i=1; i=d; i++) {
    Use an "O(n)" stable sort to sort array i++
  }
}

```

Array  
 d = # of 'digits' (Can be char)  
 k = possible range for each digit

BUCKET SORT ISAB + SINGLETON 1956

```

bucketSort(A, n) {
  for (i=0; i<n; i++) {
    insert A[i] into bucket B[A[i]]
  }
  for (i=0; i<n; i++) {
    sort bucket B[i] with insertion sort
  }
  combine B[0], B[1], ..., B[n-1] back into array
}

```

\* Won't be using bucket or radix, will use a mix of the two



\* Having more

buckets is better RADIX BUCKET SORT ← IMPORTANT $n = A.length()$ Array  
size $d = \# \text{ of "digits"}$  $k = \text{radix of "digits"}$  $\text{radixBucketSort}(A, n, d, k) \{$ 

C++ rec

Py list

Java<sup>array</sup> list

radix part

int r

```

    new Buckets[k]; // k is 10 for range of digits
    for(i=1; i<=d; i++) { // ith digit from lsd to msd
        for(j=0; j<n; j++) {
            r = ith digit from A[j];
            bucket[r].insertBack(A[j]);
            // for j
        }
        for(j=0; j<k; j++) {
            add items from Bucket[j] into A
            clear Bucket[j];
        } // for i
    }

```

Example SSN ints  
 digits  $d=9$  (digits are 0-9)  
 radix  $k=10$

$O(d(n+k))$   
 Best when  
 $d \ll n$   $k \ll n$   
 then  $O(n)$

Do front → RADIX BUCKET SORT EXAMPLE

to back,

correct way

212 012 023 231 233 310 313 021 111 101

 $d=3$   $k=4$ 

3 passes starting on

is for bottom

of bucket to

go first

units digits

tens digits

hundreds digits

BUCKETS