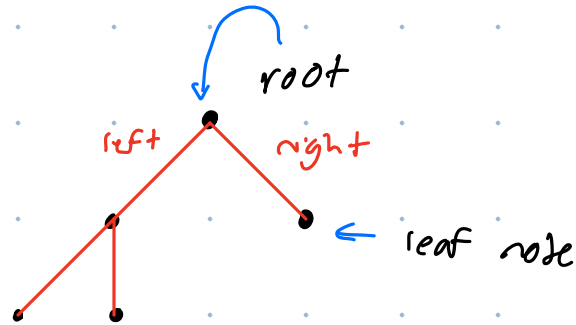


CS 21 LECTURE #9

Assignment 7 hint: for comments - take in whole string, find '#' symbol, clear string, repeat.

Binary Trees!



- A linked structure.
- Follows tree structure.
 - ↳ Binary search tree property.
- Does more operations.

Examples

- Build a P.Q, is N in there? → inefficient...
- Build hash table, output in ascending order... → inefficient...

Use a binary search tree!

Operations for BST:

- Insert
 - Delete
 - Search
- } Basic Dictionary functions...
- min/max

- Predecessor/ Successor
- Traversals

Typical performance of BST function is $O(h)$

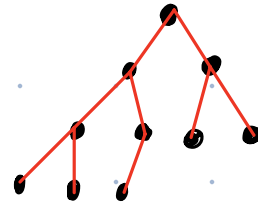
h = height of tree.

hope for $h \approx O(\log(n))$

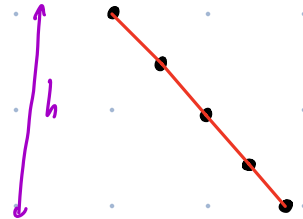
worst case $h \approx O(n)$

Happens w/ sorted data
COMMON! Sadly...

Best case:



Worst Case



How do we stay away from sorted data?

Randomize the data!

But if input one at a time, there's nothing we can do...

Use another data structure (Balancing BST)

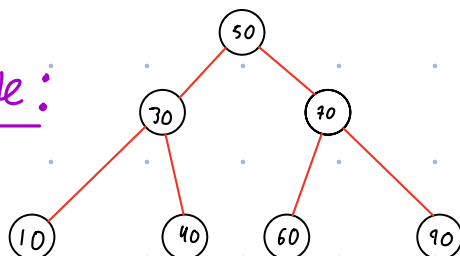
BST Property

$key(x.right) > key(x)$

$key(x.left) < key(x)$

* Assume Unique Keys
(Note for assignment) ↑

For Example:



Traversal: $O(n)$ \rightarrow 'touch' n nodes

- A 'walk' of the nodes that touches each node once in some order.

3 traversals:

- pre order: this, left, right
- in order: left, this, right
- post order: left, right, this

Memorize!

Left always comes before right
& "this" is in the expected
spot of the 3 based on the
name of the traversal

Pseudo code if in order traversal

\rightarrow MUST BE A PUBLIC WRAPPER FUNCTION
BC USER DOESNT HAVE ACCESS TO SENSITIVE
TREE DATA!

\swarrow node reference / pointer (to root I believe)

inorder(x) {

if (x == null) { return }

inorder(x.left);

print x;

// or whatever

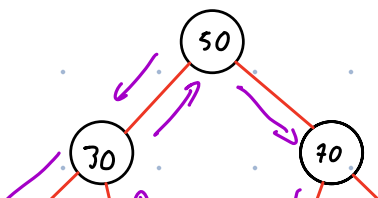
inorder(x.right);

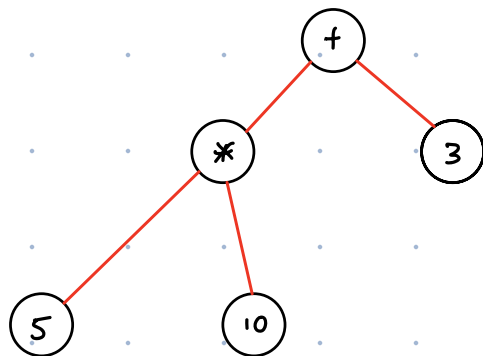
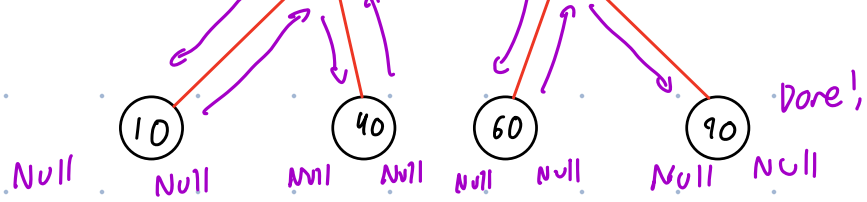
}

Pre order Printed: 50, 30, 10, 40, 70, 60, 90

In Order printed: 10, 30, 40, 50, 60, 70, 90

Post order Printed: 10, 40, 30, 60, 90, 70, 50





infix Pre order Printed: + * 5 10 3
 infix In Order printed: 5 * 10 + 3
 postfix Post order Printed: 5 10 * 3 +

DONT MAKE THIS MISTAKE: NOTICE

IN ORDER HAD IN ORDER CALL. IN POST ORDER & PRE ORDER FUNCTIONS, CALL THEMSELVES RECURSIVELY DONT CALL IN ORDER IN THE OTHER FUNCTIONS LOL...

Search → MUST BE A PUBLIC WRAPPER FUNCTION BC USER DOESNT HAVE ACCESS TO SENSITIVE TREE DATA!

recursive version

Search (x, k) {
 ↗ node reference
 ↘ key

if (x == null || x.key == k) return x;

```

    if (k < x.key) {
        return search(x.left, k);
    }
    return search(x.right, k);
}

```

Iterative version

```

search(x, k) {
    while (x != null && x.key != k) {
        if (k < x.key) {
            x = x.left;
        } else {
            x = x.right;
        }
    }
    return x;
}

```

Min

$O(h)$

MUST BE A PUBLIC WRAPPER FUNCTION
BC USER DOESNT HAVE ACCESS TO SENSITIVE
TREE DATA!

CHECK IF X IS NULL

Min (x) {

while (x.left != null) {

x = x.left;

}

return x;

}

Note ref

IN WRAPPER FUNCTION!

Max

O(h)

MUST BE A PUBLIC WRAPPER FUNCTION
BC USER DOESNT HAVE ACCESS TO SENSITIVE
TREE DATA!

CHECK IF X IS NULL
IN WRAPPER FUNCTION!

Max (x)

while (x.right != null) {

x = x.right;

}

return x;

}

MUST BE A PUBLIC WRAPPER FUNCTION
BC USER DOESNT HAVE ACCESS TO SENSITIVE
TREE DATA!

Successor

Note

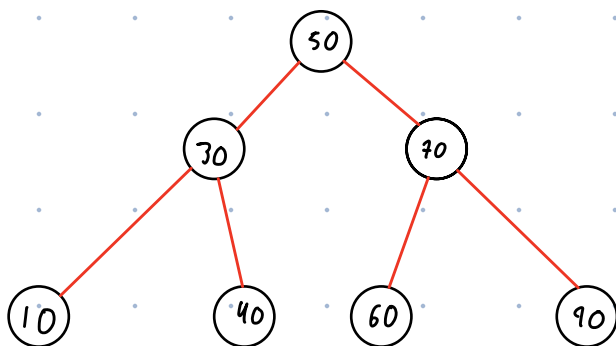
CHECK IF X IS NULL
IN WRAPPER FUNCTION!

Successor (x) {

```

if (x.right != null) {
    return min(x.right);
}
t = x.parent;
while (t != null && x == t.parent) {
    x = t;
    t = t.parent;
}
return t;
}

```



$S(30) \rightarrow 40$

$S(10) \rightarrow 30$

$S(50) \rightarrow 60$

$S(90) \rightarrow \text{null} \dots$

Insert
 $O(h)$

key value

insert(k) {

n = new node(k)

// node constructor left right & parent
autojects to null

p = null; // previous

t = root; // temp

while (t != null) {

p = t;

if (k < t.key) {

t = t.left;

} else {

t = t.right;

}

} // t always null here.

n.parent = p

if (p == null) {

root = n;

return;

}

if (k < p.key) {

p.left = n;

} else {

p.right = n;

}

}