



# Team Contest Reference

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## 1 ds

## 2 graph

## 3 math

### 3.1 geometry lib

```

1 // this library has been copied from https://github.
  com/SuprDewd/T-414-AFLV
2 #include <complex>
3 using namespace std;
4 #define P(p) const point &p
5 #define L(p0, p1) P(p0), P(p1)
6 #define C(p0, r) P(p0), double r
7 #define PP(pp) pair<point,point> &pp
8 typedef complex<double> point;
9 const double pi = acos(-1.0);
10 const double EPS = 1e-9;
11 double dot(P(a), P(b)) {
12     return real(conj(a) * b);
13 }
14 double cross(P(a), P(b)) {
15     return imag(conj(a) * b);

```

```

16 }
17 point rotate(P(p), double radians = pi / 2, P(about) =
  point(0,0)) {
18     return (p - about) * exp(point(0, radians)) +
      about;
19 }
20 point proj(P(u), P(v)) {
21     return dot(u, v) / dot(u, u) * u;
22 }
23 point normalize(P(p), double k = 1.0) {
24     return abs(p) == 0 ? point(0,0) : p / abs(p) * k;
25 }
26 bool parallel(L(a, b), L(p, q)) {
27     return abs(cross(b - a, q - p)) < EPS;
28 }
29 double ccw(P(a), P(b), P(c)) {
30     return cross(b - a, c - b);
31 }
32 bool collinear(P(a), P(b), P(c)) { return abs(ccw(a, b
  , c)) < EPS; }
33 double angle(P(a), P(b), P(c)) {
34     return acos(dot(b - a, c - b) / abs(b - a) / abs(c
  - b));
35 }
36 bool intersect(L(a, b), L(p, q), point &res, bool
  segment = false) {

```

```

37 // NOTE: check for parallel/collinear lines before
    calling this function
38 point r = b - a, s = q - p;
39 double c = cross(r, s), t = cross(p - a, s) / c, u
    = cross(p - a, r) / c;
40 if (segment && (t < 0-EPS || t > 1+EPS || u < 0-
    EPS || u > 1+EPS))
41     return false;
42 res = a + t * r;
43 return true;
44 }
45 point closest_point(L(a, b), P(c), bool segment =
    false) {
46     if (segment) {
47         if (dot(b - a, c - b) > 0) return b;
48         if (dot(a - b, c - a) > 0) return a;
49     }
50     double t = dot(c - a, b - a) / norm(b - a);
51     return a + t * (b - a);
52 }
53
54 typedef vector<point> polygon;
55 #define MAXN 1000
56 point hull[MAXN];
57 bool cmp(const point &a, const point &b) {
58     return abs(real(a) - real(b)) > EPS ?
59         real(a) < real(b) : imag(a) < imag(b); }
60 int convex_hull(vector<point> p) {
61     int n = p.size(), l = 0;
62     sort(p.begin(), p.end(), cmp);
63     for (int i = 0; i < n; i++) {
64         if (i > 0 && p[i] == p[i - 1])
65             continue;
66         while (l >= 2 && ccw(hull[l - 2], hull[l - 1],
67             p[i]) >= 0)
68             l--;
69         hull[l++] = p[i];
70     }
71     int r = l;
72     for (int i = n - 2; i >= 0; i--) {
73         if (p[i] == p[i + 1])
74             continue;
75         while (r - l >= 1 && ccw(hull[r - 2], hull[r -
76             1], p[i]) >= 0)
77             r--;
78         hull[r++] = p[i];
79     }
80     return l == 1 ? 1 : r - 1;
81 }

```

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## 4 misc

## 5 Math

### 5.1 Tree

Diameter: BFS from any node, then BFS from last visited node.  
Max dist is then the diameter. Center: Middle vertex in second step from above.

### 5.2 Divisability Explanation

$D \mid M \Leftrightarrow D \mid \text{digit\_sum}(M, k, \text{alt})$ , refer to table for values of  $D, k, \text{alt}$ .

### 5.3 Combinatorics

- Variations (ordered):  $k$  out of  $n$  objects (permutations for  $k = n$ )
  - without repetition:  
 $M = \{(x_1, \dots, x_k) : 1 \leq x_i \leq n, x_i \neq x_j \text{ if } i \neq j\}$ ,  
 $|M| = \frac{n!}{(n-k)!}$
  - with repetition:  
 $M = \{(x_1, \dots, x_k) : 1 \leq x_i \leq n\}$ ,  $|M| = n^k$
- Combinations (unordered):  $k$  out of  $n$  objects
  - without repetition:  $M = \{(x_1, \dots, x_n) : x_i \in \{0, 1\}, x_1 + \dots + x_n = k\}$ ,  $|M| = \binom{n}{k}$
  - with repetition:  $M = \{(x_1, \dots, x_n) : x_i \in \{0, 1, \dots, k\}, x_1 + \dots + x_n = k\}$ ,  $|M| = \binom{n+k-1}{k}$
- Ordered partition of numbers:  $x_1 + \dots + x_k = n$  (i.e.  $1+3 = 3+1 = 4$  are counted as 2 solutions)
  - #Solutions for  $x_i \in \mathbb{N}_0$ :  $\binom{n+k-1}{k-1}$
  - #Solutions for  $x_i \in \mathbb{N}$ :  $\binom{n-1}{k-1}$
- Unordered partition of numbers:  $x_1 + \dots + x_k = n$  (i.e.  $1+3 = 3+1 = 4$  are counted as 1 solution)
  - #Solutions for  $x_i \in \mathbb{N}$ :  $P_{n,k} = P_{n-k,k} + P_{n-1,k-1}$  where  $P_{n,1} = P_{n,n} = 1$
- Derangements (permutations without fixed points):  $!n = n! \sum_{k=0}^n \frac{(-1)^k}{k!} = \lfloor \frac{n!}{e} + \frac{1}{2} \rfloor$

### 5.4 Polynomial Interpolation

#### 5.4.1 Theory

Problem: for  $\{(x_0, y_0), \dots, (x_n, y_n)\}$  find  $p \in \Pi_n$  with  $p(x_i) = y_i$  for all  $i = 0, \dots, n$ .

Solution:  $p(x) = \sum_{i=0}^n \gamma_{0,i} \prod_{j=0}^{i-1} (x - x_j)$  where  $\gamma_{j,k} = y_j$  for  $k = 0$

and  $\gamma_{j,k} = \frac{\gamma_{j+1,k-1} - \gamma_{j,k-1}}{x_{j+k} - x_j}$  otherwise.

Efficient evaluation of  $p(x)$ :  $b_n = \gamma_{0,n}$ ,  $b_i = b_{i+1}(x - x_i) + \gamma_{0,i}$  for  $i = n - 1, \dots, 0$  with  $b_0 = p(x)$ .

### 5.5 Fibonacci Sequence

#### 5.5.1 Binet's formula

$$\begin{pmatrix} f_n \\ f_{n+1} \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 1 \end{pmatrix}^n \begin{pmatrix} 0 \\ 1 \end{pmatrix} \Rightarrow f_n = \frac{1}{\sqrt{5}}(\phi^n - \tilde{\phi}^n) \text{ where } \phi = \frac{1+\sqrt{5}}{2} \text{ and } \tilde{\phi} = \frac{1-\sqrt{5}}{2}.$$

#### 5.5.2 Generalization

$$g_n = \frac{1}{\sqrt{5}}(g_0(\phi^{n-1} - \tilde{\phi}^{n-1}) + g_1(\phi^n - \tilde{\phi}^n)) = g_0 f_{n-1} + g_1 f_n \text{ for all } g_0, g_1 \in \mathbb{N}_0$$

### 5.5.3 Pisano Period

Both  $(f_n \bmod k)_{n \in \mathbb{N}_0}$  and  $(g_n \bmod k)_{n \in \mathbb{N}_0}$  are periodic.

### 5.6 Reihen

$$\sum_{i=1}^n i = \frac{n(n+1)}{2}, \sum_{i=1}^n i^2 = \frac{n(n+1)(2n+1)}{6}, \sum_{i=1}^n i^3 = \frac{n^2(n+1)^2}{4}$$

$$\sum_{i=0}^n c^i = \frac{c^{n+1}-1}{c-1}, c \neq 1, \sum_{i=0}^{\infty} c^i = \frac{1}{1-c}, \sum_{i=1}^n c^i = \frac{c}{1-c}, |c| < 1$$

$$\sum_{i=0}^n i c^i = \frac{n c^{n+2} - (n+1) c^{n+1} + c}{(c-1)^2}, c \neq 1, \sum_{i=0}^{\infty} i c^i = \frac{c}{(1-c)^2}, |c| < 1$$

### 5.7 Binomialkoeffizienten

$$\binom{n}{k} = \binom{n-1}{k} + \binom{n-1}{k-1}, \quad \binom{n}{m} \binom{m}{k} = \binom{n}{k} \binom{n-k}{m-k},$$

$$\binom{m+n}{r} = \sum_{k=0}^r \binom{m}{k} \binom{n}{r-k} \text{ and in general, } n_1 + \dots + n_p = \sum_{k_1+\dots+k_p=m} \binom{n_1}{k_1} \dots \binom{n_p}{k_p}$$

### 5.8 Catalanzahlen

$$C_n = \frac{1}{n+1} \binom{2n}{n} = \frac{(2n)!}{(n+1)!n!}$$

$$C_0 = 1, C_{n+1} = \sum_{k=0}^n C_k C_{n-k}, C_{n+1} = \frac{4n+2}{n+2} C_n$$

### 5.9 Geometrie

**Polygonfläche:**  $A = \frac{1}{2}(x_1 y_2 - x_2 y_1 + x_2 y_3 - x_3 y_2 + \dots + x_{n-1} y_n - x_n y_{n-1} + x_n y_1 - x_1 y_n)$

### 5.10 Zahlentheorie

**Chinese Remainder Theorem:** Es existiert eine Zahl  $C$ , sodass:  
 $C \equiv a_1 \pmod{n_1}, \dots, C \equiv a_k \pmod{n_k}, \text{ggT}(n_i, n_j) = 1, i \neq j$

Fall  $k = 2$ :  $m_1 n_1 + m_2 n_2 = 1$  mit EEA finden.

Lösung ist  $x = a_1 m_2 n_2 + a_2 m_1 n_1$ .

Allgemeiner Fall: iterative Anwendung von  $k = 2$

**Eulersche  $\varphi$ -Funktion:**  $\varphi(n) = n \prod_{p|n} (1 - \frac{1}{p}), p \text{ prim}$

$\varphi(p) = p - 1, \varphi(pq) = \varphi(p)\varphi(q), p, q \text{ prim}$

$\varphi(p^k) = p^k - p^{k-1}, p, q \text{ prim}, k \geq 1$

**Eulers Theorem:**  $a^{\varphi(n)} \equiv 1 \pmod{n}$

**Fermats Theorem:**  $a^p \equiv a \pmod{p}, p \text{ prim}$

### 5.11 Faltung

$$(f * g)(n) = \sum_{m=-\infty}^{\infty} f(m)g(n-m) = \sum_{m=-\infty}^{\infty} f(n-m)g(m)$$