

linked list

Singly linked list :-

It is a linear collection of nodes where each node contains data value and a pointer to hold the address of coming after node. The first node is tracked by keeping its address in a pointer called start or home.

To traverse the complete list of to go to a particular node we must have to access the first node. That is the list can be traversed only from beginning to end this is why it is called as one way or singly linked list. The following operation are done on such a list :-

(1) Creation of the list :- As each node contains at least two parts to create the linked list at first we need to define a structure that describes the type of data part and pointer.

After that the home 'pointer' is declared i.e assignment with null to it creates the empty list:

Algorithm :-

Step 1:- Define structure node with members data as any primitive data type and a pointer link as node type.

Step 2:- Declare a pointer variables start of node type.

Step 3:- Let start = null

Step 4:- End

Ex - `typedef struct nd`

```
{ int data;
```

```
struct nd *link;
```

```
} node;
```

(new datatype based on structure)

```
node *start = Null;
```

(12) Traversal of singly linked list —

Since, singly linked list is an one way list, to traverse it we must have to locate the first node and from that we can go ahead through its link part until the list is finished. This traversal is done by assigning a pointer with the start or home then, accessing the allocation pointed by this pointer and assigning the pointers with the link part of the node. The whole process procedure is repeated until the pointer becomes null.

Algorithm

1. Start Traversal to show()
2. Declare pointer Ptr as node type
3. Let, Ptr = start
4. Repeat step 5 & 6 while Ptr != Null
5. Print data of Ptr
6. Let Ptr = link of Ptr
7. End of step 4 while
8. End of function.

Algorithm for inserting 1st node

Step 1 :- start, gnsfirst (val)

Step 2 :- Declare pointer ptr As nodetype

Step 3 :- Allocate memory for node to ptr

Step 4 :- Let data of ptr = val

Step 5 :- let link of ptr = start

Step 6 :- Let start = ptr

Step 7 :- End of function.

* Inserting at the end of list :-

To insert a node the end of an existing list we need to allocate memory location for a node and to assign the next member or linked with the address of deallocated memory but to traverse to the last node starting from the first node.

This operation is done by using the following algorithm.

Step 1:- Start, $s = \text{last}$ (nil) and $\text{ptr} = \text{new}$

Step 2:- Declare pointers ptr and s as node type.

Step 3:- Allocate memory to pointer ptr .

Step 4:- Let data of $\text{ptr} = v$

Step 5:- Let link of $\text{ptr} = \text{Null}$

Step 6:- If $\text{start} = \text{Null}$ then

GA Let $\text{start} = \text{ptr}$

GB Return

Step 7:- Let $s = \text{start}$

Step 8:- while link of $s \neq \text{Null}$

Step 9:- Let $s = \text{link of } s$

Step 10:- End of while

Step 11:- Let $s = \text{ptr}$

Step 12:- Return

* Inserting a node at the given position:

After allocating memory for the new node and assigning its data-part we need to traverse to the node coming before the given position by changing the link part of the pointer starting from start. After that we assign the link part of the new node with connect it to the node currently present at the given position and then we assign the link part of the traversed node with the address of newly allocated node. It will follow the given algorithm.

Step 1 :- Start insmid (n, p)

Step 2 :- Declare pointers ptr and s of node type

Declare counter c = 0

Step 3 :- If start = Null then

3. A Print "List empty"

3. B Return

Step 4 :- If $P = 1$, Then

4. A call `instList(v)`

4. B Return

Step 5 :- Let $S := \text{start}$

Step 6 :- while $S \neq \text{Null}$ And $C < P - 1$

Step 7 :- Let $s = \text{link of } S$

Step 8 :- Let $c = c + 1$

Step 9 :- End of step 6 Loop

Step 10 :- If $S = \text{Null}$ And $C < P - 2$ Then,

10. A Print "Position not found"

10. B Return

Step 11 :- end of if

Step 12 :- Allocate memory to `ptr`

Step 13 :- Let `Data of ptr = v`

Step 14 :- Let `link of ptr = link of s`

Step 15 :- Let `link of s = ptr`

Step 16 :- Return

* Deletion from the linked list :-

Since linked list is a dynamic structure deletion of node is performed physically. We can delete a node from anywhere within a link list. There are different algorithm from the beginning from the end and from the middle. But in any algorithm the middle deletion is performed by holding a node address in pointer, adjusting the links and then releasing the pointers.

* Deletion of first node - To delete the first node first of all we need to check that is there any node or node not, if it is then we assign pointer with the value of start and adjust the start with link the pointer before releasing it.
The algorithm is given below

Step 1:- Start Delfirst()

Step 2:- Declare pointer ptr of node type

Step 3:- If start = Null then

3.A Print "List empty"

3.B Return

Step 4:- Let $\text{ptr} = \text{start}$

Step 5:- Let $\text{start} = \text{link of } \text{ptr}$

Step 6:- Release ptr

Step 7:- Return

* Deletion of last node :- To delete the last node we need to traverse to the second last node because it holds the address of last node as its link. Then we keep the link value in a pointer and assign the link with null to indicate the end. After then we release the pointer.

Algorithm for deletion of last node

Step 1:- Start $\text{DelLast}()$

Step 2:- Declare pointer ptr and s of node type.

Step 3:- If $\text{start} = \text{Null}$ then

3.A Print "List empty"

3.B Return

Step 4:- Let $s = \text{start}$

Step 5:- While link of s & $\text{link} \neq \text{Null}$

Step 6:- Let $s = \text{link of } s$

Step 7 :- end of while loop

Step 8 :- let $\text{ptr} = \text{link of } s$

Step 9 :- let link of $s = \text{Null}$

Step 10 :- Release ptr

Step 11 :- Return

Deletion of node at given position

Step 1 :- Start DelPos(Pos)

Step 2 :- Declare pointer $\text{ptr}, \text{s}, \text{a}$ pointer

and $\text{c} = 1$, and v

Step 3 :- If $\text{start} = \text{Null}$ Then

3.A Print "List is empty"

3.B Return with 0 (zero)

Step 4 :- end of if

Step 5 :- If $\text{Pos} = 1$ Then

5.A call Delfirst()

5.B Return

Step 6 :- End of if

Step 7 :- Let $\text{s} = \text{start}$

Step 8 :- Do step 9 while $\text{c} < \text{Pos}-1$ And
links of $\text{s} \neq \text{Null}$

Step 9 :- Let $\text{s} = \text{link of } \text{s}$

Step 10 :- end of while

Step 11 :- If $\text{c} < \text{Pos}$ Then

11.A Print "given position not found"

11.B Return with 0

Step 12 :- end of if

Step 13 :- Let $\text{ptr} = \text{link of } \text{s}$

Step 14 :- Let dlink of $S = \text{link}$ of ptr

Step 15 :- Let $N = \text{Data}$ of ptr

Step 16 :- Declare ptr

Step 17 :- Return with \checkmark

Doubly linked list :

A doubly linked list is also a collection of nodes scattered within the memory and connected together logically by holding the address of consequence nodes. The difference between SLL and DLL is that it can be traversed in two ways. If from the beginning this two way traversal provides a lot of advantages over the SLL it reduces processing time particularly when we have to search a value and have to append a node.

For the two way traversal, the nodes of DLL holds before and coming after. The nodes of DLL have at least three parts - the data part, the previous link pointer and the next link pointer. The previous