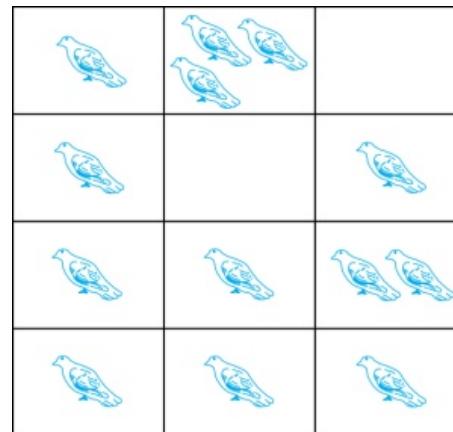


# Session 63: The Pigeonhole Principle

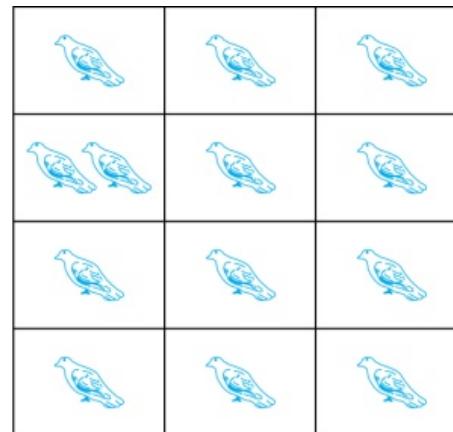
- The Pigeonhole Principle
- The Generalized Pigeonhole Principle

# The Pigeonhole Principle

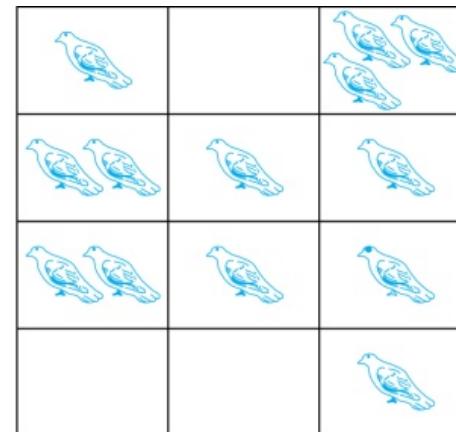
If a flock of 13 pigeons lives in a set of 12 pigeonholes, one of the pigeonholes must have more than 1 pigeon.



(a)



(b)



(c)

# The Pigeonhole Principle

**Pigeonhole Principle:** If  $k$  is a positive integer and  $k + 1$  objects are placed into  $k$  boxes, then at least one box contains two or more objects.

**Proof:** We use a proof by contraposition.

- Suppose none of the  $k$  boxes has more than one object.
- Then the total number of objects would be at most  $k$ .
- This contradicts the assumption that we have  $k + 1$  objects.  $\square$

# Using the Pigeonhole Principle

**Corollary:** A function  $f$  from a set with  $k+1$  elements to a set with  $k$  elements is not one-to-one.

$$f : P \rightarrow H, |P| = k+1, |H| = k$$

put  $x \in P$  into box  $y \in H$  if  $f(x) = y$

since we have  $k+1$  "pigeons" and only  $k$  "holes"  
at least one hole contain more than 1 element.

Therefore we have  $x_1, x_2, x_1 \neq x_2$  with  $f(x_1) = f(x_2)$

# The Generalized Pigeonhole Principle

**The Generalized Pigeonhole Principle:** If  $N$  objects are placed into  $k$  boxes, then there is at least one box containing at least  $\lceil N/k \rceil$  objects.

**Proof:** We use a proof by contraposition.

- Suppose that none of the boxes contains more than  $\lceil N/k \rceil - 1$  objects.
- Since  $\lceil N/k \rceil < N/k + 1$ , the total number of objects is at most

$$k \left( \left\lceil \frac{N}{k} \right\rceil - 1 \right) < k \left( \left( \frac{N}{k} + 1 \right) - 1 \right) = N,$$

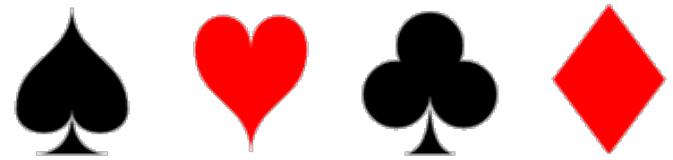
- This is a contradiction because there are a total of  $N$  objects .  $\square$

# Example

Among 100 people there are at least  $\lceil 100/12 \rceil = 9$  who were born in the same month.

(assume for each month we have at most 8 people ,  
since there are 12 months , the total number of people  
cannot be more than  $12 * 8 = 96$  )

# Example



The 4 suits of cards

How many cards must be selected from a standard deck of 52 cards to guarantee that at least three cards of the same suit are chosen?

k boxes : 4 suits

N objects : the chosen cards

$\lceil \frac{N}{k} \rceil \geq 3$ , which is the smallest N, such that  $\lceil \frac{N}{4} \rceil \geq 3^2$ .

$N = 12$  ?  $\lceil \frac{12}{4} \rceil = 3$  but

$N = 11$   $\lceil \frac{11}{4} \rceil = 3$  hence

$N = 9$   $\lceil \frac{9}{4} \rceil = 3$ , 9 is the smallest integer N such that  $\lceil \frac{N}{4} \rceil \geq 3$

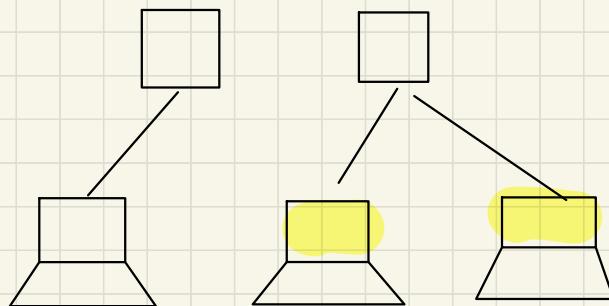
# Summary

- The Pigeonhole Principle
  - Counting functions
- The Generalized Pigeonhole Principle

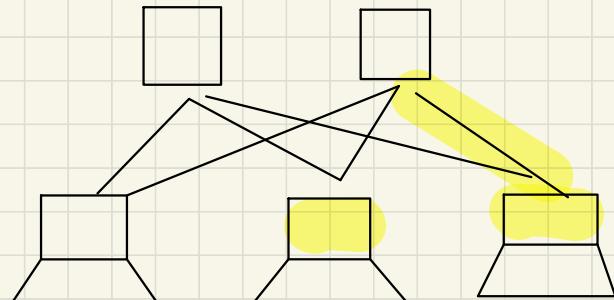
Desktop Cabling Problem     $d$  computers,  $p$  printers,  $d > p$

How many cables  $c$  are needed, connecting computers and printers, such that whatever subset of  $p$  computers you select, each of the selected computer will have separate cable to a printer?

Example :

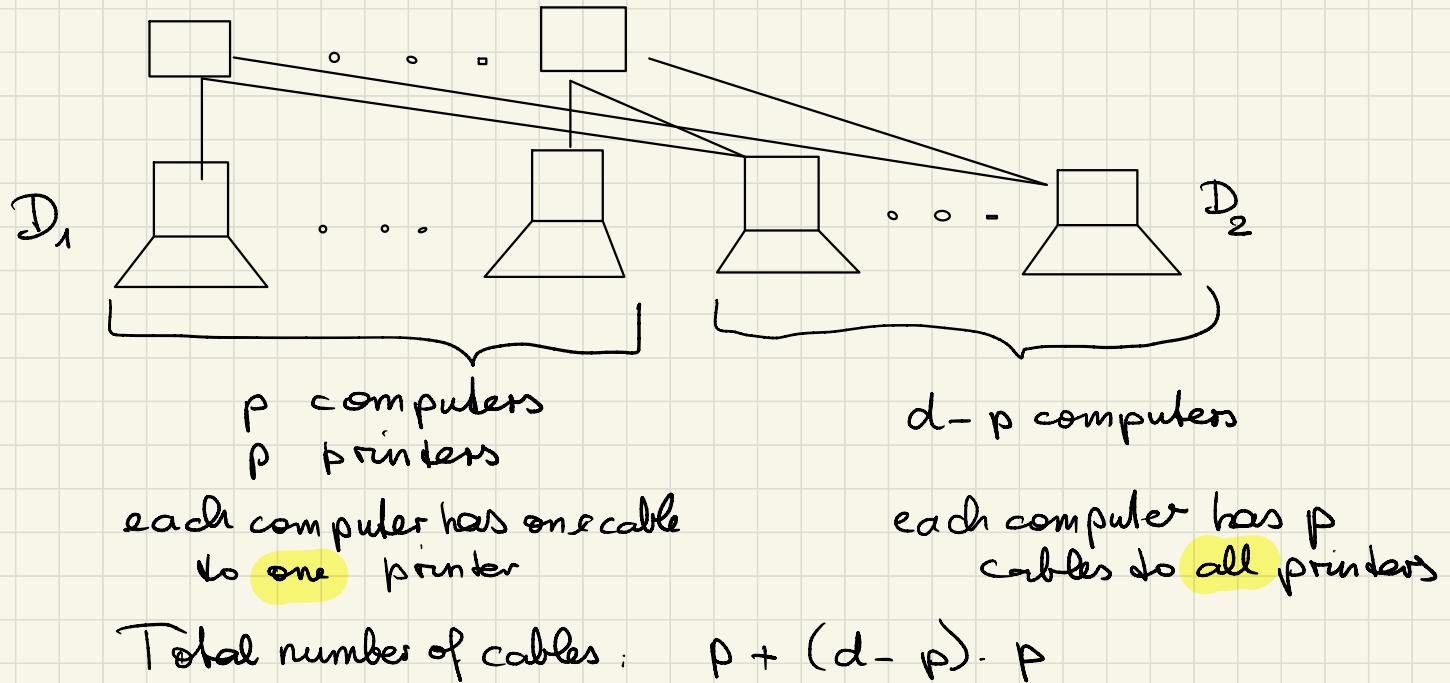


Too few cables, the selected computers have not their own cables



Too many cables, we could remove one cable and still a solution exists in all cases

## A solution that works



Theorem : 1. This cabling solves the cabling problem

2. This cabling uses the minimal number of cables

## Proof Theorem part 1

Let  $S$  be the set of selected computers,  $|S| = p$

$S$  contains  $k$  computers from  $D_1$  ( $|S \cap D_1| = k$ )

These are connected using the available cable

This leaves  $p - k$  unconnected printers

$S$  contains  $p - k$  computers from  $D_2$  ( $|S \cap D_2| = p - k$ )

Since computers in  $D_2$  have cables to all printers,

they can be connected to the remaining  $p - k$  printers.

## Proof of Theorem part 2

assume we have only  $p + (d-p) \cdot p - 1$  cables

then there must be a printer  $pr$  that is connected to

$$\left\lfloor \frac{p + (d-p) \cdot p - 1}{p} \right\rfloor = \left\lfloor (d-p) + \underbrace{\left(1 - \frac{1}{p}\right)}_{<1} \right\rfloor = d-p \text{ computers at most}$$

By pigeonhole principle: otherwise all printers have at least  $(d-p)+1$

cables, and since we have  $p$  printers, we have in total  $p \cdot (d-p) + p$  cables  
which is a contradiction.

Consider all computers not connected to  $pr$ , there are at least  $p$  of those

This gives us  $p$  computers, that are connected to  $p-1$  printers only

Therefore no solution for the calling problem exists.

Therefore a solution with less than  $p + (d-p) \cdot p$  cables cannot exist.