# MEMORY FORENSICS AND MALWARE ON M1 MACS

# A Thesis

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# Abstract

Malware continues to shift to infection and persistence mechanisms that can only be seen in memory. These techniques evade traditional forensic analysis, requiring memory forensics to examine these malicious programs. Memory forensics allows for the recovery of historical data created by running malware, including information that it tries to hide. Memory analysis capabilities have lagged behind on Apple's new M1 architecture, while the number of malicious programs only grows. To make matters worse, Apple has developed Rosetta 2, the translation layer for running x86\_64 binaries on an M1 Mac. This decision has resulted in all malware compiled for Intel Macs being executable on M1 machines. In this paper, malware will be executed through the Rosetta 2 translation environment, in an effort to document functionality of malware run through it. Afterwards, memory analysis will be performed on select samples to confirm functionality. Finally, the research efforts to bring memory forensics to the M1 will be discussed, along with new artifacts from Rosetta 2 that will be able to be analyzed as a result of these research efforts.

# Chapter 1. Introduction

#### 1.1 Memory Forensics

Memory forensics is the examination of a host through the acquisition and analysis of volatile memory. Acquisition of memory samples is possible through a number of tools, many of which are freely available [10] [43] [63]. These tools often work similarly, by loading the kernel driver and determining the ranges of the RAM's physical address space, then translating physical memory to virtual memory through the use of page tables and copying the result to storage. It is also possible to acquire samples through virtual machines, either by suspending the state of the machine or by taking a snapshot in a virtual machine.

After acquisition, parsing is performed by frameworks that will present the acquired memory sample from the host in a meaningful state for analysis. In order to parse memory dumps, these frameworks implement strategies like walking linked lists and file carving, which is looking for objects that look like files, processes, threads, and more. These strategies result in useful structures and objects being exposed from the OS and applications, often even if they were deallocated or terminated prior to acquisition. These structures are used by analysts, who then examine the state of the host through various plugins and tools looking for evidence from an incident.

Memory forensics is an increasingly important part of an incident response team's efforts to understand an incident forensically. It is a necessary addition to traditional storage forensics, as it allows for the recovery and study of information only available in RAM, such as running and hidden processes, open files, encryption keys, private browsing data, network connections, clipboard data, and kernel and application structures. The ability to get information about deallocated objects and terminated processes is extremely useful. It makes it possible to build timelines that would otherwise be incomplete and find indicators of compromise that otherwise would not be seen.

Memory forensics is also important for malware analysis. Through memory forensics,

an analyst can not only see all of the previously mentioned indicators as it relates to the malware, but also defeat many attempts the malware can use to hide itself from traditional analysis and live forensics tools. An example is Direct Kernel Object Manipulation (DKOM), where a rootkit modifies process lists in kernel memory. This modification will allow a process to hide not only from live analysis and OS subsystems, but also from some memory forensics tools. This technique can be seen in action in Figure 1.1, where a malicious application has removed itself from the process list via manipulation of the \_EPROCESS kernel structure, which is used by the *pslist* plugin and column in that figure. However, it is still detected as active (TRUE in the figure) through the Volatility psscan plugin. That plugin scans the memory dump for \_EPROCESS objects using an associated tag, rather than relying on the process list; this method is slower than walking a list, but more thorough. Plugins that use alternative methods to find processes are also shown, many of which also find the hidden process. Other malicious techniques like DKOM are increasingly commonplace, including memory-only loading, dropping, and execution. Memory forensics is therefore a crucial component, and sometimes the only solution, for a digital forensics investigation.

#### 1.2 Rosetta 2

# 1.2.1 Background and Translation Pipelines

Rosetta 2 bridges previous-generation x86 Intel apps with M1 Macs. It is translation software that automatically converts an executable with only x86\_64 instructions, then runs the translated binary [1]. As a result, it is usually the case that the first launch will be slower while translation occurs, but subsequent launches will have only slight reduction in performance compared to running the same binary on an Intel Mac. This method, called ahead-of-time (AOT) translation by Apple, is better than acting simply as a full emulator for x86 when a program is run, as that would noticeably hinder speed and the user experience every time. It is accomplished by generating a Mach-O file with the file

Offset(P)	Name	PID	pslist	psscan	thrdproc	pspcid	csrss	session	deskthrd ExitTime
0x0a720180	prtk_worker_ser	360	True	True	False	True	True	True	False
0x097e0c68	vmtoolsd.exe	3932	True	True	False	True	True	True	True
0x0a4e0da0	InstallShield L	1292	True	True	False	True	True	True	True
0x0aad9888	services.exe	1152	True	True	False	True	True	True	True
0x09852490	wcescomm.exe	4052	True	True	False	True	True	True	True
0x09ec2da0	wmiprvse.exe	2984	True	True	False	True	True	True	True
0x09f15810	wdfmgr.exe	2316	True	True	False	True	True	True	True
0x09f32960	sshd.exe	2080	True	True	False	True	True	True	True
0x09b49020	CodeMeterCC.exe	1372	True	True	False	True	True	True	True
0x097a6c28	cmd.exe	2164	True	True	False	True	True	True	True
0x09f2e518	oracle.exe	2096	True	True	False	True	True	True	True
0x098185a8	rundl132.exe	3840	True	True	False	True	True	True	True
0x0a926da0	vmacthlp.exe	1332	True	True	False	True	True	True	True
0x0a86bb20	svchost.exe	1228	True	True	False	True	True	True	True
0x0a785da0	winlogon.exe	1056	True	True	False	True	True	True	True
0x0938f370	GoogleUpdate.ex	1104	True	True	False	True	True	True	True
0x09f30020	ctfmon.exe	3952	True	True	False	True	True	True	True
0x0aad9608	svchost.exe	1628	True	True	False	True	True	True	True
0x09f49da0	cygrunsrv.exe	1596	True	True	False	True	True	True	True
0x0a607228	CodeMeter.exe	436	True	True	False	True	True	True	True
0x09f32020	extjob.exe	2072	True	True	False	True	True	True	True
0x0a509b30	vmtoolsd.exe	2476	True	True	False	True	True	True	True
0x0a841a88	svchost.exe	1744	True	True	False	True	True	True	True
0x0a640da0	svchost.exe	1432	True	True	False	True	True	True	True
0x09f14020	svchost.exe	2240	True	True	False	True	True	True	True
0x0aa947b8	svchost.exe	396	True	True	False	True	True	True	True
0x0u76eda0	svchost.exe	564	True	True	False	True	True	True	True
0x09796020	MagicDisc.exe	2260	False	True	False	True	True	True	True
	wuguelt exe	2776	True	True	False	True	True	True	True
0x097f7020	AdobeARM.exe	3804	True	True	False	True	True	True	True
0x0a5fdda0	lsass.exe	1164	True	True	False	True	True	True	True
0x09e4d020	BTTray.exe	2880	True	True	False	True	True	True	True

Figure 1.1. psxview finding a malicious process attempting to hide

extension .aot. It is based on the original x86\_64 executable, and only generated once. Whenever the x86\_64 binary needs to be run, this translated AOT file is mapped into the memory of the executing process. There are security mechanisms in place to ensure the translation artifact has not been tampered with, such as code directory hashes and protection through Security Integrity Protection (SIP).

The other translation pipeline for Rosetta 2 is called just-in-time (JIT). JIT is necessary because some applications generate code at runtime. In the case of an x86\_64 application that does this, Rosetta 2 will encounter code that cannot be translated ahead of time. The JIT translation pipeline takes care of these instances entirely within the process after being transferred control by the kernel.

The typical flow for an x86\_64 binary is that the AOT pipeline will first create an AOT artifact derived from the binary. That artifact is then located when running the app and mapped onto the executing process's memory. JIT is then used for any x86\_64 Mach

objects encountered throughout the execution. Constants are referenced in the original binary, not copied into the translated one. Thus, the original binaries are still required; the AOT artifact is not a standalone that can be moved and executed by itself [39].

#### 1.2.2 Constraints

Rosetta 2 cannot translate kernel extensions or virtual machine apps that virtualize x86\_64 platforms, but it can translate almost anything else compiled for x86\_64 Macs. It does not come pre-installed on M1 Macs, so the first time Rosetta 2 is required a prompt to download it appears. A final constraint is that Rosetta translation covers the entire process, meaning all code or dependencies loaded by the process must also be x86\_64 if the process was started as such [1].

#### 1.2.3 Handling of Mach-O Universal Binaries

The Mach object file format is used as the native format for binaries in macOS, including executables, libraries, and dynamically-loaded code [18]. Mach-O binaries have multiple formats, one of which is universal. Universal binaries support multiple architectures within one Mach-O file, such as PowerPC and x86\_64, or x86\_64 and arm64. For universal executables with both x86\_64 and arm64 code, Rosetta 2 is by default not used at all. The code slice that most closely resembles the architecture of the host will be run, so the arm64 binary will just execute. The user can decide to execute the x86\_64 variant through Rosetta 2, however, if it is expected to load x86\_64 code modules rather than arm64 versions of them. An exception to this is that script-only apps are run under Rosetta 2 translation by default [3]. It is also possible to specify an app be run under x86\_64 by default using the LSArchitecturePriority key [11]. Another consideration is that running an unsigned universal binary from an arm64 shell will result in the process being killed, due to arm64 executables needing to be signed to run [35]. This constraint does not apply to x86\_64 binaries, which are allowed to run completely unsigned [21].

#### 1.3 Research Importance

It may initially seem like an obvious decision for malicious programmers to want to build universal binaries, so that their malware immediately natively runs as arm64 on M1 Macs. However, there are significant factors discussed in the previous section, such as needing to build all dependencies universally, needing to sign validly to run arm64 at all, and not being able to debug arm64 without an M1 Mac themselves. As a result of these constraints, malware authors may actually prefer building only for x86\_64 and letting Rosetta handle the rest. This possibility makes it even more important that the Rosetta 2 translation software and the artifacts it generates are understood.

Last year, data breach incidents had an average cost of over \$4 million to the company [6]. Ransomware attacks exceeded \$600 million in 2021, over double the amount in 2020 [56]. As many businesses and livelihoods have shifted online as a result Covid-19, it is more important than ever that digital forensics analysts have the tools they need to fully investigate and remediate incidents. They currently do not on Apple M1 Macs.

In addition, malware is already being seen compiled with code for M1. Some of these strains are highly infectious, with one sample affecting 30,000 macOS machines [60]. Also, modern malware with memory-only capabilities are known to run on M1 through Rosetta 2. EvilQuest is an example of malware that is tested in this research that attempts in-memory execution of a payload [51]. This example serves to show that there are aspects of malicious attacks that cannot be currently detected on M1 Macs.

#### 1.4 Outline

Chapter 2 will introduce prior works that are related to my research. Chapter 3 will go into detail regarding how the testing environment was set up, how malware samples were chosen and obtained, the methodology for performing analysis on those samples in Chapter 4, and how memory samples were acquired for the analysis performed in Chapter 5. Chapter 4 will present and discuss the results of executing and analyzing malware samples

in the M1 virtual machine environments. Chapter 5 will discuss the malware selected for further analysis, examine those samples through memory forensics, and finally discuss the results. Chapter 6 will present further research that was performed to enable and enhance memory forensics on Apple M1 Macs. Chapter 7 will provide concluding thoughts on the research presented, and also discusses future work that could be done to build upon this work and M1 memory forensics.

# Chapter 2. Related Work

#### 2.1 M1 Forensic Analysis

The research regarding M1 Mac analysis is currently quite sparse. Koh M. Nakagawa has examined M1 Macs, specifically looking at reversing Rosetta 2 [39]. His motivation was to investigate it as an attack surface and potentially exploit it. In the course of his research, he patched Ghidra to correctly disassemble AOT files, reverse engineered the Rosetta 2 runtime binary, detailed how to debug the emulation process at the arm64 instruction level, and more.

The Corellium team have looked at M1 devices at the hardware level. They have presented their reverse engineering research at BlackHat 2021, where the got Linux running on M1 [57].

Joshua Duke has also performed relevant research [37], where the full Volatility suite of plugins was run and the output compared between Intel Macs and M1 Macs. Technical explanations for differences were given, and the future necessary fixes for plugins that broke were documented.

#### 2.2 Memory Analysis on Intel Macs

Modern macOS Userland Runtime Analysis [46] and Memory Analysis of macOS Page Queues [28] are two examples of recent research that perform memory analysis on macOS to better understand internals and find forensically significant artifacts. The Volatility Foundation [58] and Volexity [24] have been a part of advancing memory forensics capabilities for years, including on Macs. The Art of Memory Forensics [40] is the flagship book on memory forensics and performing it. It has over 100 pages dedicated solely to Mac memory forensics, and covers foundational skills as well as extremely advanced ones on that topic.

#### 2.3 Malware Analysis on Intel Macs

Objective by the Sea is a Mac-focused cybersecurity conference where high quality talks, often about malware, are given yearly [16]. Objective-See is a blog maintained by Patrick Wardle, where technical analyses of macOS malware are regularly posted [17]. Patrick Wardle also has a work-in-progress book titled The Art of Mac Malware that covers all aspects of macOS malware and the tools needed to analyze it [64].

Trendmicro has compiled a detailed technical report on the *XCSSET* macOS malware, including reverse-engineering and complete analysis of the binaries and infection mechanisms associated with it [26]. Many other threat intelligence blogs perform similar analyses regularly [12] [20] [14] [25] [22]

# 2.4 Malware on M1

As the number of samples grows, there are increasingly many reports on specific samples of malware with native arm64 code.

Silver Sparrow writeup, which exposed and details the most infectious strand of M1 malware yet [61]. Google's Threat Analysis Group detailed a watering hole attack targeting pro-democracy Hong Kong visitors [32]. Patrick Wardle also has given a talk at Objective by the Sea v4.0, where he reverse engineered the anti-analysis arm logic of M1 malware found in the wild [53].

# Chapter 3. Experimental Setup

This section will discuss the steps taken to create an environment for safe and thorough testing of malware. It will also discuss how the malware samples were chosen and obtained, as well as how samples were analyzed to determine functionality through Rosetta 2 translation. Finally, it will detail the process to acquire a memory image for examination, as this process on M1 Macs is more difficult than with other systems.

# 3.1 Environment Setup

The larger batch of malware samples tested for Chapter 4 were executed on virtual machines (VMs) running on Parallels Desktop Pro 17.1.1-51537. The operating system of the M1 Mac VMs was macOS Monterey 12.2, build 21D49, and the operating system of the host M1 Mac was macOS Monterey 12.2.1, build 21D62. The host that the memory was acquired from was running macOS Monterey 12.1.0, build 21C52, in order for the memory sample to be compatible with an already created Volatility profile. Rosetta 2 had to be installed on each VM prior to execution of malware. Command-line Developer Tools were installed as well, as some malware utilized functionality from tools in this suite. On the VMs, Wireshark and various built-in Terminal tools such as ps aux were used to monitor high-level functionality of the malware. Each VM was within what Parallels calls a Host-Only Network [50], where the VM is in a separate subnet that can only connect to a gateway and other VMs in the subnet. Each VM was also set up such that it was the DNS Server under its networking configuration. This setup allowed the observation of DNS queries made during dynamic analysis of the malware. Because there is currently no snapshot functionality on M1 images for Parallels, multiple samples were run on one VM before preparing and using another VM for more samples. For all VMs, System Integrity Protection (SIP) remained enabled. As a result, a few samples identified by macOS as malware were unable to run on the VMs. SIP would need to be disabled because the ability

to disable those protection measures is restricted under it. Samples that were applications could still be run by overriding malware protection, but executables, disk images, and package installers could not. It was not possible to disable SIP on a VM as Recovery Mode was not accessible. The reason Recovery Mode was inaccessible is discussed in Section 3.4.

Three malware samples were also run on the host machine while it was isolated from the internet, so that a memory image could be collected for analysis. Section 3.4 also discusses why this was necessary. The host machine was wiped and rebuilt after the memory dump was acquired and moved off disk to a flash drive.

#### 3.2 Malware Selection

Malware was chosen based on a number of factors, such as impact level, viability of attack vector, and indicators of compromise. Newer malware was also generally given preference to similar malware that was older. For example, *CDDS / MacMa* had both a 2019 and 2021 installer; the 2021 version was the version tested. A majority of samples are thus modern, higher impact malware like backdoors or spyware with persistence mechanisms. In total, roughly 150 malware samples were considered for testing. Those chosen were obtained from the Objective-See blog [13], where many malicious macOS samples dating as far back as the early 2000's are freely available.

# 3.3 Forensic Analysis of Malware

A high-level analysis was performed for each of the selected 16 malware samples to determine which of them functioned through Rosetta 2 translation, and which did not. Samples for this portion were run in the VMs, and clear indicators of compromise, such as Command and Control (C&C) pings or running Launch Agents, were used to determine functionality. For those samples that did not have expected behavior, more analysis was done to determine the cause and whether Rosetta 2 was responsible.

# 3.4 Acquisition of Memory Image

For the malware samples that were chosen for further analysis, a memory sample needed to be obtained. Version 17.1.1-51537 of Parallels Desktop Pro does not support snapshots or suspension of an M1 VM [49], so the virtualization methods of obtaining a memory sample were not possible. In that version of Parallels, boot order configuration options are also not supported. It would be necessary to be able to change the boot order in order to get to Recovery Mode on an M1 VM. If SIP were disabled, it would still be possible via setting Recovery Boot Mode in the terminal. However, in order to disable SIP a user needs to boot into Recovery Mode. Fig 3.1 shows this catch-22. Once in Recovery Mode, it would be possible to allow third-party kernel extensions. This allowance is necessary because memory acquisition tools such as Volexity's Surge Collect Pro software require the use of their signed kernel extension to function. With both the virtualization and Recovery Mode routes not possible for acquiring the memory of a VM, the decision was made to run the malware samples on the host M1 with networking disabled. The host was prepared with Surge and the password-protected zip of malware samples prior to disabling networking. With the host prepared, the samples were run and the memory dump was captured. The dump was moved to a USB and the host was restored to a known good state.

```
Last login: Wed Mar 9 15:42:09 on ttys000
//Users/chrles/Downloads/SysJoker/types-config; exit;
chrles@chrless-Mac ~ % //Users/chrles/Downloads/SysJoker/types-config; exit;
addToStatup

Saving session...
...copying shared history...
...saving history...truncating history files...
...completed.

[Process completed]
   [Restored Mar 9, 2022 at 7:07:20 PM]
Last login: Wed Mar 9 19:07:17 on console
Restored session: Wed Mar 9 16:11:10 CST 2022
chrles@chrless-Mac ~ % sudo nvram "recovery-boot-mode=unused"
Password:
Error setting variable - 'recovery-boot-mode': (iokit/common) not permitted.
chrles@chrless-Mac ~ %
```

Figure 3.1. Setting Recovery Boot Mode requires SIP to be disabled

# Chapter 4. Malware Results

#### 4.1 macOS Security Mechanisms

The malware tested in this research encountered multiple built-in macOS security mechanisms. The below results report on which samples were stopped by which mechanisms. Thus, it is useful to introduce them here.

#### 4.1.1 File Quarantine

File quarantine is a security feature where files downloaded from the internet are marked as quarantined. When quarantined files are first run, a user needs to give permission via an opened dialog box before the file executes. Files downloaded via *curl* and some other methods, however, are not marked as quarantined [30], so this method of protecting users is not without flaws.

# 4.1.2 Gatekeeper

Gatekeeper is a security feature that is intended to prevent apps developed by an unknown or invalid developer from running without additional user consent. This method has been more effective, but ultimately still is susceptible to users being tricked into running the program anyway. Also, if they decide to, malware authors are often easily able to bypass this mechanism by signing their apps with illegitimately obtained but valid developer IDs [52].

#### 4.1.3 Notarization

Notarization is a security feature that automatically scans quarantined apps for issues related to code-signing or malicious content [15]. Once an app, package installer, or disk image passes the checks, that information is published online for Gatekeeper to check when users download and run it. If a user's system encounters a previously unnotarized app, package installer, or disk image, it will inform them and then let them bypass and run it if

they choose [34].

#### 4.1.4 Malware Protection

As a final security measure, Apple also develops an antivirus called XProtect. XProtect is built-in to macOS and uses YARA signatures to detect malware and prevent it from running [19]. When a sample was flagged as malware, it either requires overriding malware protection for that app, or disabling of SIP in order to run executables. EvilQuest is an executable that did not have the "Override Malware Protection" option, while the MacDownloader, XCSSET, and ZuRu applications did. Therefore, it seems the ability to override malware protections exists for applications but not for executables. This observation was not able to be confirmed through documentation, however.

#### 4.2 Malware That Was Functional

# 4.2.1 CDDS / MacMa

#### **Background**

CDDS / MacMa was discovered in November 2021, targeted visitors to a few websites in Hong Kong, and has the ability for executing terminal commands, file transfer, audio recording, and more [32].

#### Execution

After running the 2021 version of the installer, a payload is dropped that triggers warnings such as attempting to control the computer using accessibility features and accessing the microphone. *User Agent* is the executable that is dropped and persists via a Launch Agent. Launch Agents are perhaps the most common method for malware to persist on macOS. Launch Agents can be used to take actions whenever directories change, launch jobs at set intervals, and even run scripts every time a user logs in [5]. Launch Daemons are largely equivalent except they run on behalf of the root user or another specified user [9].

As an example of a Launch Agent, CDDS's can be seen in Figure 6.1. User Agent

running accordingly is shown in Figure 4.2

```
CDDS — vim ~/Library/LaunchAgents/com.UserAgent.va.plist — 93×26
<?xml version="1.0" encoding="UTF-8"?>
<!DOCTYPE plist PUBLIC "-//Apple//DTD PLIST 1.0//EN" "http://www.apple.com/DTDs/PropertyList-</p>
1.0.dtd">
<dict>
        <key>Label</key>
        <string>com.UserAgent.va</string>
        <key>LimitLoadToSessionType</key>
        <string>Aqua</string>
        <key>ProgramArguments</key>
        <array>
                <string>/Users/chrles/Library/Preferences/UserAgent/lib/UserAgent</string>
                <string>-runMode</string>
                <string>ifneeded</string>
        </array>
        <key>RunAtLoad</key>
        <true/>
        <key>StartInterval</key>
        <integer>600</integer>
        <key>ThrottleInterval</key>
        <integer>2</integer>
        <key>WorkingDirectory</key>
        <string>/Users/chrles/Library/Preferences/UserAgent/lib</string>
</dict>
</plist>
 ~/Library/LaunchAgents/com.UserAgent.va.plist" 24L, 707B
```

Figure 4.1. CDDS ensuring persistence through a Launch Agent

# 4.2.2 DazzleSpy

#### **Background**

DazzleSpy was discovered in January 2022, was distributed via a compromised Hong Kong pro-democracy website, and supports fully interactive RDP sessions, the ability to dump the Keychain on vulnerable systems, and self-deletion [44]. This sample comes from a similar campaign to CDDS, with the main difference being the payloads (CDDS and DazzleSpy) that are delivered.

#### Execution

After running *DazzleSpy*, connection attempts to its probable c2 were observed. In addition, a Launch Agent was loaded that will launch the executable each time the user logs in.

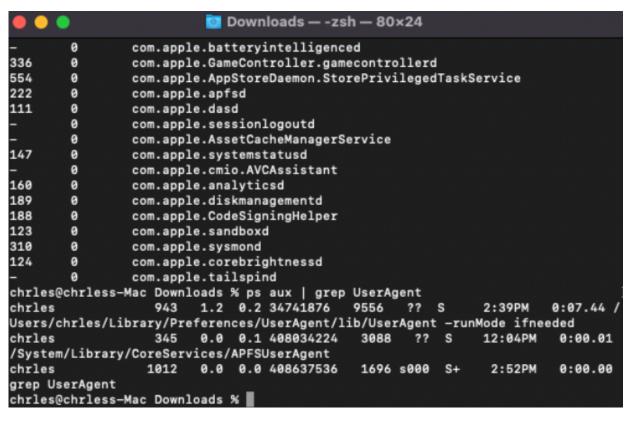


Figure 4.2. The *User Agent* binary running via a Launch Agent with arguments specified

#### 4.2.3 EvilQuest

#### Background 3 4 1

EvilQuest was discovered in June 2020, was distributed predominately through pirated programs, and has the ability to ransom a host, infect files to run its payload, search for and kill processes, and harvest data [51].

#### Execution

The malware bundled with the pirated program was identified as such by macOS, so it was prevented from running. As discussed in prior sections, it is not possible to fully disable malware protections for executables with the current version of Parallels for M1 Mac VMs. As this sample was of particular interest due to its distribution vector, capabilities, and relative recency, it was selected for further analysis and run on the host machine before a memory image was collected. Just from initial observation, persistence was set, test binaries were infected, and a fake process named *Avast* was killed. Networking was unable

to be tested due to running on a host with it completely disabled but based on other functionality it is likely to have attempted to connect to C&C.

Figure 4.3 shows both a running Launch Agent and file infected from *EvilQuest*. The file infection is marked by the ending hex 0xdeadface.

```
user@cyber2-2 darwin % ps aux
                   1529
                          0.1
                               0.0 408666384
                                                7792
                                                        ??
                                                            Ss
                                                                  6:38PM
                                                                            0:00.02
root
sudo /Library/AppQuest/com.apple.questd --silent
                   1530
                          0.1
                               0.0 34260360
                                               2668
                                                       ??
                                                           S
                                                                 6:38PM
                                                                           0:00.01
Library/AppQuest/com.apple.questd --silent
                          0.0
                               0.0 408628368
                                                1616 s000
                                                            S+
                                                                  6:38PM
                                                                            0:00.00
                   1533
grep silent
user@cyber2-2 darwin % hexdump ~/Downloads/testBinary |
0021950 65 61 64 65 72 00 5f 68 65 6c 6c 6f
                                              00 5f
0021960 69 6e 00 5f 6e 75 6d 00 5f 6e 75 6d 32 00 5f 6e
        75 6d 62 65
                     72 5f
                           00 5f
                                 5f
                                    5f
                                        73
                                           74
                                              72 63
0021980
        5f
           63
              68
                 6b
                     00 5f
                           66
                              72
                                 65
                                    65
                                        00
                                           5f
                                              67
                                                 65
0021990
        68
           61
              72
                 00
                     5f
                        6d 61 6c
                                 6c
                                    6f
                                           00
                                        63
                                              5f
                                                 70
00219a0 6e 74 66
                 00
                    64 79 6c 64
                                 5f
                                    73 74
                                           75
                                              62
                                                 5f
00219b0 6e 64 65 72 00 5f 5f 64 79
                                    6c 64 5f
                                              70 72 69 76
       61 74 65 00 00 00 00 00 03 70 57 01 00 ce fa ad
00219c0
00219d0
00219d1
user@cyber2-2 darwin % 📗
```

Figure 4.3. EvilQuest file infection and persistence

#### 4.2.4 FinSpy

# Background

The macOS version of *FinSpy* was discovered by Amnesty International in Fall 2019. The original Windows version was distributed through a backdoored version of Adobe update [27]. The macOS version is distributed as a trojaned application with a Turkish name [48] and is used as spyware.

#### Execution

After the disk image was mounted and the installer app was run, a malicious installer was dropped and executed at the same time a network request to download a legitimate version of Adobe AIR was made. The malicious installer then detected it was running inside virtualization software and stopped running. As the focus of this portion of the research was confirming high-level functionality through Rosetta 2 translation, using a debugger to

get around this check was outside of scope.

# 4.2.5 Hydromac

#### **Background**

Hydromac was discovered in the middle of 2021 and was distributed via the Tarmac and Bundlore malwares [41].

#### Execution

After giving execution permissions, it runs and attempts to connect out to C&C. It is thought to be a downloader/stager beyond that, as it has the ability to check for installed anti-virus programs and can download and execute programs.

# 4.2.6 Komplex

# Background

Komplex was discovered in late 2016, is associated with APT28 Sofacy, likely targets individuals working in aerospace, is distributed as an attachment in an email, and has functionality including remote execution and information collection about the host it runs on [31].

# Execution

Upon running Komplex, a PDF in Russian appearing to relate to a space program opened. It also dropped and ran an executable that pinged Google checking for connectivity, as well as creating a Launch Agent that launches that script at startup. However, an additional script that was supposed to be responsible for loading the Launch Agent did not get created or run, so persistence was ultimately not set. This was confirmed by rebooting the system and noting that the payload did not start running. The VM this malware ran on was isolated from the internet, so the connectivity check failed. Parallels Host-Only mode did not seem to allow a VM to see other VMs on the same subnet, though the documentation implies it should. As a result, using a program like INetSIM [8] on another VM to return 200's to the malware did not seem possible, and was out of scope since

functionality was already observed.

#### 4.2.7 MacDownloader

#### **Background**

MacDownloader was discovered in early 2017, is associated with Iranian APTs Charming Kitten and Flying Kitten, was distributed as both an Adobe Flash installer and Bit-Defender adware remover, and has the ability to harvest and upload sensitive information [7].

#### Execution

After running it, it completes a fake installation while asking the user for their credentials. A file is created in the temp directory that is populated with sensitive information including username and password, Keychain files, running processes, and more. However, traffic was not seen to its known domain.

In Figure 4.4, the harvested info from a *MacDownloader* execution can be seen.

Figure 4.4. *MacDownloader*'s harvested info

#### 4.2.8 Realtime-Spy

**Background** 

This *Realtime-Spy* sample is from early 2018, was distributed via the *realtime-spy-mac[.]com* website, and has the ability to upload screen recordings, keystrokes, and other user activity it has collected [47].

#### Execution

While the sample runs, it can be seen reaching out to its namesake domain and attempting to screen record.

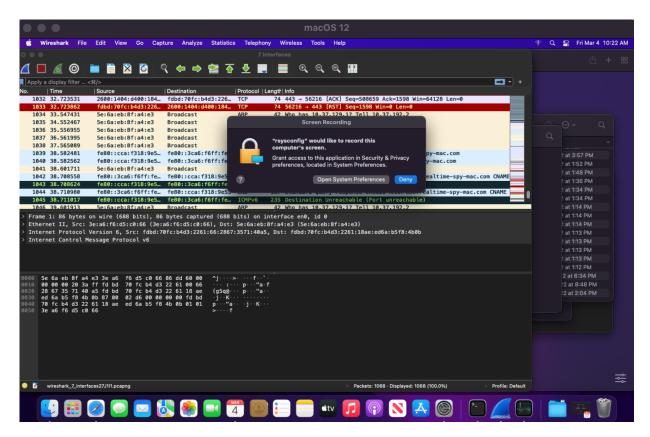


Figure 4.5. Realtime-Spy attempting to screen record

#### **4.2.9** Ventir

# Background

Ventir was discovered back in 2014 and has the ability to execute commands, exfiltrate files, and keylog on a fully infected system [45].

#### Execution

It drops 5 executables and removes itself, creates and loads a launch agent for persis-

tence, and runs one of its exe's. However, the keylogging functionality fails as the binary responsible is a universal binary compiled for i386 and PowerPC!

#### 4.2.10 XCSSET

#### Background

XCSSET was discovered in early or mid 2020, spreads through infected XCode projects, and attempts to steal credentials and data through trojanized browsers [38].

#### Execution

Running either of the associated scripts fails, as they are fingerprinted as malware and prevented from executing. Running the infected XCode application after overriding malware protection, however, results in a connection attempt to C&C infrastructure. None of the other infection markers, such as processes being killed or fake apps being created, are observed due to the failed connection attempt.

#### 4.2.11 ZuRu

# Background

ZuRu was discovered in September 2021. It spread through the search engine Baidu, where it paid to show at the top of the search results for "iTerm2" via a sponsored link. As a result, it appeared above even the legitimate iTerm2, likely tricking many users. Its functionality includes downloading and executing additional files that exfiltrate information and leverage remote access [54].

#### Execution

An initial run was terminated at launch because it was not able to load in a custom malicious dynamically linked library. After investigation, it was determined that the library wasn't loaded in due to macOS expecting it to be arm64 rather than the x86\_64 library it found. The library was expected to be arm64 because this application was a universal binary, so by default it ran as arm64 on an M1 Mac. After selecting "Open using Rosetta", it did not crash. It attempted to connect out to C&C but stopped running because it did

not receive a response.

Using a Terminal, it is also possible to run the app without the "Open using Rosetta" checkbox selected. After navigating the application's directory to the actual program, it can be executed through the *arch* command, which allows for specification of the architecture with which to execute a universal binary. Both methods are equivalent and successfully run the application. This successful execution means that Rosetta 2 handles both translation of x86\_64 code, as well as any x86\_64 dependencies. However, when an app is started as arm64, even if a translatable x86\_64 dependency is found, the app will fail to run. The implication then is that if a malware author is going to build malware as a universal binary, any dependencies within the app will need to be universal as well. This implication is confirmed through Apple documentation. Otherwise, their app will fail on M1 Macs, since arm64 is the default architecture that will run in a universal binary on M1 Macs [2]. As a result, macOS will expect dependencies that a binary loads in to be arm64 as well. The other way to prevent a crash would be for the author to force execution as x86\_64 on an infected host, but at that point there would be no reason to have also built as arm64 in the first place.

In Figure 4.6, ZuRu's crash report when run as arm64 can be seen. It includes errors regarding incompatible architecture due to finding an x86\_64 dependency it expected to be arm64.

#### 4.3 Malware That Was Not Functional

Malware was determined to be nonfunctional if it needed significant modification or to be handheld through bypassing multiple security mechanisms for different stages of its payloads. As examples, *Applejeus* was caught by Gatekeeper, then had a second stage get immediately caught by malware protection. Another example, *SysJoker*, needed to have its file extension modified on top of being provided executable permissions to run correctly. This modification changed the entire attack vector of [29], as its file extension is thought

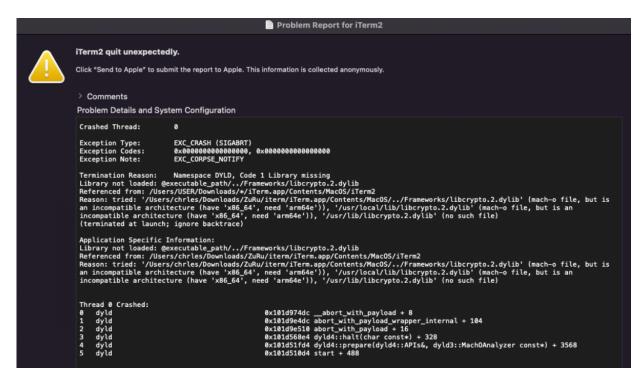


Figure 4.6. ZuRu crashing due to an incompatible dependency

to play an important part of the initial infection vector, either as a malicious media player attachment or as a TypeScript file inside of an infected package.

# 4.3.1 Applejeus / MacLoader

# Background

Applejeus was discovered around the middle of 2018, is associated with a North Korean APT, was distributed through phishing emails and trojanized programs [33], and has functionality including execution of memory-only payloads.

#### Execution

The initial package installer ran, but did not function as intended. It gave root permissions to the malicious executable unioncryptoupdater and moved it to a separate location, but that executable was not launched. This failure seems to be due to unioncryptoupdater being identified as malware (MACOS.8d038b3) and prevented from running. The same occurs with the CrashReporter malware that another Applejeus package installer tries to drop and run as a Launch Daemon.

#### **4.3.2** Crisis

# Background

Crisis was discovered in 2016, appeared to be sold as a part of government surveillance tools, use Adobe-related phishing as an infection method, and spies on users via audio recording, stealing browser data, and more [42].

#### Execution

It failed to run due to being a 32-bit (i386) executable, which is incompatible on an M1 even with Rosetta 2 installed.

In Figure 4.7, Rosetta being unable to translate the *Crisis* i386 executable can be seen.

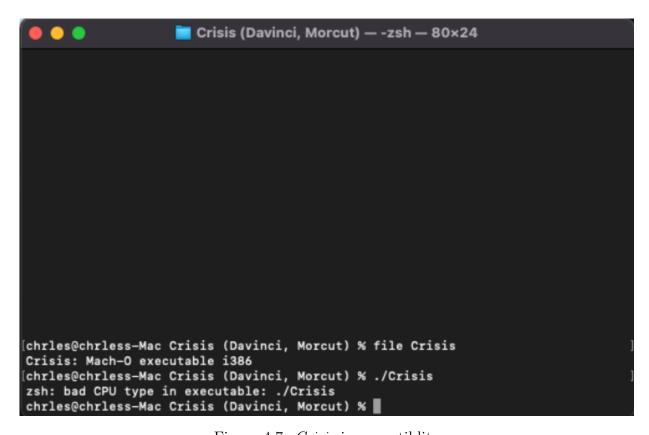


Figure 4.7. Crisis incompatiblity

# 4.3.3 Shlayer

# **Background**

Shlayer was originally discovered in 2018, with new variants found through 2021. It

is distributed through malicious sites intended to be discovered through search engine optimization, that then redirect to fake Adobe Flash installation pages. A recent variant from last year had the ability to run only with a user double-click, as it utilized 0-days to bypass most Apple security measures described in section 4.1. After infection, it only attempts to download adware [36].

#### Execution

The disk image that contains the installer, AdobeFlashPlayer.dmg, was flagged as malicious and unable to be mounted. The installer package within, Player.pkg, was not flagged as malware and ran, but didn't appear functional at all. This lack of functionality is unsurprising because the malware relied on 0-days that were patched on the OS build it was run on. The way that the macOS security mechanisms were bypassed were not highly technical, however. They essentially were the result of leaving out an Info.plist file of an app bundle that was script-based. These exploits serve as a good reminder that, while M1 Macs seem hardened against malware executing without user permission, it is difficult for a system to be truly secure.

#### 4.3.4 Silver Sparrow

#### **Background**

Silver Sparrow was discovered in February 2021, is distributed in installer packages, is a universal binary, and has an unknown final payload. The widest held opinion is that the second stage installs adware, based on similarities between data from the c2, the arg 'upbuchupsf', and an affiliate code often used by adware. It also has roughly 30,000 detections on macOS [61].

#### Execution

Shortly after running, the executable encounters an error evaluating JavaScript. Two expected files, *version.json* and *version.plist*, are created but they are empty. Two expected shell scripts, *agent.sh* and *very.sh*, are not created. It thus appears to fail. The same JavaScript error occurs when running as both arm64 and through Rosetta 2 on x86\_64, so

the failure is not thought to be due to a translation error.

In Figure 4.8, the UDP stream of Silver Sparrow's installer package can be seen.

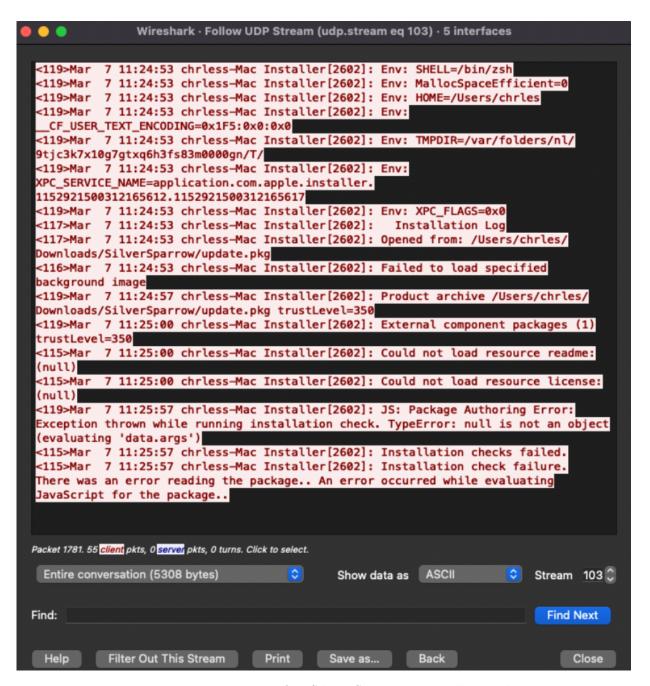


Figure 4.8. UDP stream for Silver Sparrow's installer package

# 4.3.5 SysJoker

*Background* 

SysJoker was discovered in December 2021. Examining CC behavior has led to the thought that it's used only in targeted attacks. However, it is also possible its extension implies it is TypeScript to be used as a part of infected software packages. It is a universal binary, and has the ability to execute commands and files based on remote C&C instructions [?].

#### Execution

When double clicked normally, the file opens as a playable media file and no infection markers are seen. Persistence is not set, the second binary is not dropped and run, and there is no attempt to connect to its command and control infrastructure. However, after converting it from its extension .ts and giving it execution permissions, it is immediately recognized by finder as a universal Mach-O binary. Running that binary then leads to it functioning correctly. The modified binary was also run through Terminal to force launching as x86.64, and the same functionality was observed. Persistence was set via a launch agent that runs the dropped binary textitupdateMacOs at startup. Attempts to connect to a Google drive to generate its c2 were also observed. Because a signed universal binary is being run on an M1 Mac, the .ts version will default to running the arm64 slice if it runs at all. Therefore, despite not knowing the root cause of the inability to function as .ts, it can be said it is not an error or lack of functionality with Rosetta 2 translation.

Table 4.2 covers the functionality that was observed to determine a successful run. The functionality was grouped into 3 columns: Persistence, Networking, and Harvesting. These correspond to the malware maintaining long-term existence on a system, attempting to connect out from a system, or harvesting information from the system, respectively.

Table 4.3 and Table 4.1 are similar, but their columns differ slightly. This difference is due to no functional malware sample encountering Notarization, and no non-functional sample encountering File Quarantine.

Table 4.1. The functional malware and the protections that were overridden

Malware Sample	File Quarantine	Gatekeeper	Malware Protection
CDDS / MacMa	X		
DazzleSpy	X		
EvilQuest		X	X
FinSpy		X	
Hydromac		X	
Komplex	X		
MacDownloader			X
Realtime-Spy		X	
Ventir		X	
XCSSET			X
ZuRu		X	

Table 4.2. The functional malware and observed high-level behavior

Malware Sample	Persistence	Networking	Harvesting
CDDS / MacMa	X		X
DazzleSpy	X		
EvilQuest	X	?	
FinSpy		X	
Hydromac		X	
Komplex		X	
MacDownloader			X
Realtime-Spy		X	X
Ventir	X		
XCSSET		X	
ZuRu		X	

Table 4.3. The non-functional malware and the protections that were overridden

Malware Sample	Gatekeeper	Notarization	Malware Protection
AppleJeus / MacLoader	X		X
Crisis	X		
Shlayer		X	X
Silver Sparrow		X	
SysJoker	X		

#### 4.4 Discussion

#### 4.4.1 Rosetta 2 Translation

From the above results, it is clear that Rosetta 2 is willing to translate x86\_64 malware, even in cases where the sample is unsigned, not notarized, and fingerprinted as malware. It does this translation well, as none of the compatible samples malfunctioned as a result of Rosetta 2 failing to translate properly. However, given that the user consistently had to give permissions, and occasionally override malware protections, it is reasonable that it translated the executables. Rosetta 2 is translating what it is told to translate, so it is functioning correctly. That said, because x86\_64 malware translated through Rosetta 2 does not have to be signed, while arm64 malware would have to be signed in some capacity, and all dependencies have to be ported to arm64 as well, it may be preferable to malware authors to avoid preparing arm64 versions of their programs at all. In most cases with new malware tested, and even with much of the older malware tested, overriding malware protection was not necessary because the binaries were not fingerprinted as malware by XProtect, despite being around for over 5 years in some cases. As a result, users usually would only need to be socially engineered into overriding a warning from Gatekeeper, at which point they would be infected. Running unsigned code is a vector that would not be possible without Rosetta 2, since arm64 code requires a signature to run at all.

#### 4.4.2 Viability of Infection

Most Macs are infected by applications or executables that require user interaction [52]. Malware authors employ a large number of methods to convince users to infect themselves.

These methods include but are not limited to: packaging malware with shareware, including malware with pirated applications, masquerading as fake applications, spreading through malicious attachments, and using search engine optimization or sponsored links to have malicious web results show highly.

The impact and viability of these samples seems to be lowered due to the general lack of signing and requiring of Rosetta 2 to be installed. However, many users will happily run something anyway when presented with a Gatekeeper warning, especially if any of the above methods are used. Indeed, almost all of the tested malware samples utilized one or more of those techniques for initial infection, many to great success.

#### 4.4.3 Conclusions

Four malware samples were tested that were simply applications. XCSSET, MacDownloader, and Realtime-Spy were the three that were functional. All of them were x86\_64, and completely unsigned. The one that did not immediately work, ZuRu, was a universal binary and signed. It is not surprising it was signed, since a universal binary would be expected to default to arm64 on an M1. ZuRu crashed, however, when run as arm64 because the malicious dynamically linked library it loads in, library it. 2.dylib, was only built targeting x86\_64. When an app is run as arm64, arm64 is expected to be the architecture for everything else the app needs. If ZuRu is executed as x86\_64 instead by manually going into the information for the app and checking "Run with Rosetta", it runs without crashing. The same Gatekeeper warning occurs whether run as x86\_64 or arm64. The difference is that as just x86\_64, it would have worked without user intervention required to manually specify the proper architecture. This has an interesting implication: life might be easier for malware staying with x86 64, rather than building as universal or arm64. Malware authors porting to a universal binary have many considerations to port, including apps, plugins, custom frameworks, static and dynamic libraries, build and command-line tools, launch agents and daemons [3]. In addition, it's not possible to debug the arm64 slice of a universal binary on an Intel Mac. Therefore, if authors do not have an M1, they are essentially building on good faith. Finally, to run the arm64 slice, the program would need to be signed. The downside to  $x86\_64$  only, however, is that it assumes users with M1 Macs will have Rosetta 2 installed or be willing to install it when prompted. When comparing downsides for both, it is reasonable to think that some malware authors will keep building solely for  $x86\_64$  as long as Rosetta 2 lets them.

# Chapter 5. Memory Analysis

#### 5.1 Background

Samples were run on a host running macOS Monterey 12.1.0 with networking disabled. Command-line Developer Tools and Rosetta 2 were installed for malware functionality, and Surge 21.12.14 was installed to acquire memory from the host. Analysis was performed with an unreleased version of Volatility 2 with M1 Mac support. Volatility, the new version of it, and efforts to get it functional are discussed in Chapter 6.

## 5.2 Malware chosen

Three samples were chosen for further analysis, all due to observed or purported functionality. CDDS was chosen because it was found recently and its harvesting functionality was of interest. EvilQuest was chosen because it is also relatively modern, and the infection of other files, killing of processes, and attempted memory-only execution seemed to make for good analysis opportunities. It was fingerprinted as malware by XProtect, so SIP was disabled on the host so that malware protection could be overridden and the executable run. Realtime-Spy was chosen because it is an example of commercial spyware and displayed significant functionality during execution that could be investigated further.

#### 5.3 Memory Analysis

The first step was to enumerate the processes in the memory sample and identify any suspicious ones. This step was accomplished through the *mac\_pstree* and *mac\_tasks* Volatility plugins. The *mac\_pstree* plugin not only identifies processes, but also constructs a process tree that shows parent/child relationships between them. The *mac\_tasks* plugin was used to line up processes with metadata like execution time to determine suspicious processes.

With this knowledge, the process tree was used to identify another previously unknown

Figure 5.1. Output from mac\_pstree identifying relaunch process as a child of rtcfq

process named relaunch that was a child of a suspected malicious process named rtcfg. Based on its name, relaunch is suspected to be used for persistence. Another process, com.apple.questd, was also found and was the Launch Agent that EvilQuest loaded in. Using the process identifiers (PIDs) gathered from the previous plugin output, mac\_lsof was run next. The mac\_lsof plugin lists the open file descriptors of a process. Using this plugin with the malicious PIDs, another clear indicator of compromise was discovered. As shown in Figure 5.2, the system.rtcfg folder was discovered. Looking up this directory yields results

regarding the *Realtime-Spy* malware and its keylogging capabilities [55]. Such reports also

confirm the hypothesis that the relaunch process is used for persistence throughout the

current session.

Figure 5.2. Output from mac\_lsof on known malicious PIDs

An indicator of compromise was found for the *CDDS* malware in a slightly different way. The *mac\_list\_files* plugin, which lists all files in a memory dump, was run against the memory sample and had its output redirected to a file. The list of files was then searched for Rosetta 2 artifacts that would be generated by the *CDDS* malware, with the

understanding that if any hits were found then it would be proven the malware had been translated and run on the host. Figure 5.3 shows the output of the search, where multiple AOT files are found for malicious binaries associated with *CDDS*. These binaries could be dumped using an updated version of  $mac\_procdump$  for further analysis. The update to that plugin and one other is discussed in detail in Chapter 6.

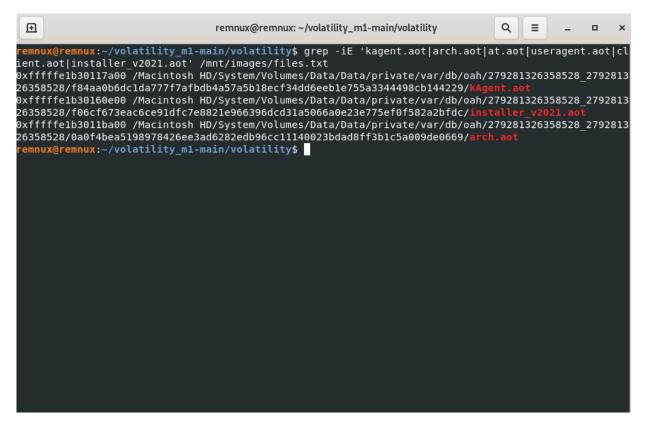


Figure 5.3. Artifacts from Rosetta 2 translation for the CDDS malware sample

As an aside, the process CarbonComponentScannerXPC has no official documentation online, and only some posts on forums asking what it is. Those forums speculate it is related to Rosetta 2, but none confirmed it. Interest was piqued, however, in determining if this process was a part of Rosetta 2, as it relates somewhat to the research in this paper. The process was dumped using  $mac\_procdump$ , then the strings [23] util was run on that dumped executable. As can be seen in Figure 5.4, there were many hits for strings that are related to the Rosetta 2 runtime.

```
∄
                                      remnux@remnux: ~/volatility_m1-main/volatility
              x:~/volatility_m1-main/volatility$ strings -n 8 dumpDir/task.509.0x7ff7ffca1000.dmp
       'x86|jit|aot|translat|rosetta'
   -E
         version
integer overflow in aot metadata load command fragment list
integer overflow in aot metadata load x86 image path
    instruction
interpreter
                  allocation too small
offsets.
              instruction length.has value()
is valid
               thread
        te indirect branch
    _gpr_state_from_arm_state
expected ARM LR to be in <mark>translat</mark>ed code: lr_abi_kind=%hhu arm_pc=0x%llx arm_lr=0x%llx
expected ARM LR to be in <mark>translat</mark>ed code: lr_abi_kind=%hhu arm_pc=0x%llx _saved_lr=0x%llx
expected next ARM PC to be in tr
                                       nslated code: abi kind=%hhu new arm pc=0x%llx
lr abi info.u.tra
                         ted code.instruction extents.kind == InstructionOffsetKind::Syscall
expected saved LR to be in <mark>translat</mark>ed code: lr_abi_kind=%hhu saved_lr=0x%llx arm_pc=0x%llx arm_lr=0x
    rip.has value()
did not find an >
                       pc. abi kind=%hhu arm pc=0x%llx arm lr=0x%llx
emulated forward to an arm pc that isn't in trans
                                                          lated code. arm pc=0x%llu abi kind=%hhu emulation
interval=[0x%llu,0x%llu) instruction interval=[0x%llu, 0x%llu) x86 rip=0x%llu
    fpr state from arm state
arm_gpr_state_from
arm fpr state from
find
              tion in tree
               base address <= x86 address
metadata->
iterate fragments in x86 interval holding lock
         page permission
mprotect failed on jit
                          write fault: %lld
handle jit breakpoint
```

Figure 5.4. Output from mac\_procdump for the CarbonComponentScannerXPC process

#### 5.4 Analysis Discussion

As proven by the above analysis, the chosen samples are indeed functioning and running. Indicators of compromise were found for each of the three samples analyzed. Therefore, memory forensics has provided an avenue for confirmation that Rosetta 2 translates malicious x86\_64 programs. Using Volatility, indicators of compromise such as running processes, open file descriptors, Rosetta 2 artifacts, and strings from dumped processes were plainly seen. Memory forensics was extremely useful in finding indication these samples ran, and even information about a process unrelated to malware, allowing for insights that would often not otherwise be possible. These analysis capabilities would not exist without the prior contributions made to get Volatility working on M1 Macs.

# Chapter 6. Other Contributions

## 6.1 Volatility

Volatility is the defacto open-source framework that analysts utilize to perform memory forensics [59]. Through the use of plugins and OS-specific profiles, data from a memory sample is parsed and extracted before being presented to the user. Volatility offers fairly robust support for Windows, macOS, and Linux. However, when Apple released M1 Macs, Volatility had no immediate way to support analysis of them, as there was more work to be done than just using Apple's kernel debug kits to create a new profile [37].

## 6.2 Volatility on M1 Macs

An address space, which is used by Volatility for address translation and mimics the memory management unit of the CPU, for arm64 did not exist prior to M1. One would be necessary for Volatility to support M1 memory forensics, however. An arm64 address space was made by Tsahi Zidenberg and pushed to Github [62], and this was used as the basis for M1 address translation.

Next, Volatility requires knowledge of the location of the kernel page tables and the value of the kernel address space layout randomization. The former is used for address translation and the latter is necessary because the offset provided in the Volatility profile will not be the same as the actual offset in memory. These values could be automatically obtained on Intel Macs, but are currently necessary for an analyst to find themselves via the metadata of a memory acquisition tool such as Surge.

With the above fixes and one more fix for userland virtual address translation, plugins could now be tested. Plugins were tested in bulk and output of each plugin was studied. There were many plugins that required modifications to become functional on M1. Some fixes were as simple as adding arm64 architecture to the list of valid architectures for the plugins to check. There were also multiple plugins that relied on an Intel-specific data

structure to determine information about a process. These plugins now use a new API that is functional on both architectures. A final breaking change was kernel symbols and types that were moved or were missing from M1 Macs. Multiple plugins had equivalent symbols or work-arounds that led to the same functionality. Some, however, currently have not been able to be updated for M1 as a result of missing symbols. Overall, almost all Volatility plugins are now fully functional on M1.

## 6.3 Rosetta 2 Analysis and Plugins

Rosetta 2 technical analysis began by writing a test program that generated artifacts in memory and had a dependency for a test library. The program was then run, and the memory mappings of the running processes were then viewed through the *vmmap* command. Through the memory mappings it could be seen that there were two mappings associated with each process, including the test library. The *mac\_procdump* and *mac\_librarydump* plugins were tested and only grabbed the Intel executables, not the second mapping for the AOT artifacts for each. Both plugins were updated to be aware of and extract the AOT artifacts, adding valuable artifacts that investigators would now be aware of. These artifacts are particularly helpful, as they are proof of Intel binaries being executed on an M1. It is also significant in that these artifacts persist through reboots, even if the original executables are deleted.

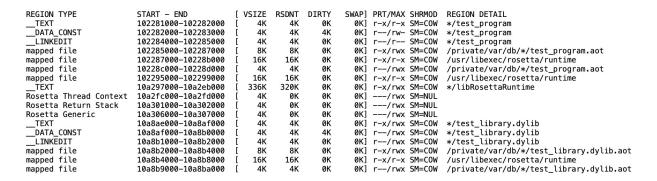


Figure 6.1. vmmap showing dual mappings for the test program and library

A new plugin, mac\_rosetta, was also developed which automatically extracts and lists

all AOT files. Figure 6.2 shows the output of this plugin against the sample program and library. The extraction of AOT files could theoretically be done on a live machine, but the folder that contains them is protected by SIP. On a live machine, therefore, the only way to access that folder is by disabling SIP, which involves rebooting the machine. After disabling SIP, a user would then manually navigate into that folder to grab AOT files for analysis. While reasonable to do on an individual level, that approach quickly becomes infeasible as a part of live incident response with many potential machines involved. These enhanced and new plugins, as well as all plugins updated to be functional for M1 Macs, represent a large increase in the ability to forensically analyze an M1.

Pid	Name	Base	AOT Path
417	CarbonComponentS	0x0000000102fa0000	<pre>Data/private/var/db/oah/[snip]/CarbonComponentScannerXPC.aot</pre>
2277	test_program	0x0000000102285000	<pre>Data/private/var/db/oah/[snip]/test_program.aot</pre>
2277	test_program	0x000000010a8b2000	<pre>Data/private/var/db/oah/[snip]/test_library.dylib.aot</pre>
2277	test_program	0x00000002025ed000	Data/private/var/db/oah/[snip]/dyld.aot

Figure 6.2. mac\_rosetta extracting AOT artifacts

# Chapter 7. Conclusion and Future Work

#### 7.1 Conclusion

Apple made the decision to develop Rosetta 2 to ease the transition from Intel architecture to M1 architecture. A side effect of this decision was that all previous x86\_64 malware would run on their new machines. This research explored the consequences of this reality, by determining the functionality of malicious x86\_64 binaries on M1 Macs and the viability of continuing to run them over arm64 versions. Functionality was explored through two angles: high-level dynamic malware analysis and memory forensics. Both angles had their results thoroughly examined. Finally, the research effort to bring memory analysis to M1 was discussed, as that made the memory forensics in Chapter 5 possible.

#### 7.2 Future Work

One constraint to the analysis performed in Chapter 4 was that connectivity checks were not bypassed. Getting around these checks would be possible either by debugging the process at the instruction level and skipping over the checks, or by using a program like INetSim [8] that will return 200's for any connection attempt. Either method would have allowed further functionality to be observed for multiple of the samples tested.

Another contribution would be to investigate more samples on M1 using memory forensics. As there were only three samples on which memory forensics were performed, there is potentially much more to be learned through continued analysis of a wider variety of samples.

A research effort that is happening concurrently is documenting APIs commonly used by malware on Intel, and determining whether or not the same functionality exists on M1. This effort will take a deeper dive into the exposed macOS functionality, and will lead to a better understanding of how malware will need to function to achieve its goals on M1 devices. As a result, these APIs will be better understood, defended against, and identified in digital forensics investigations.

The research in this paper was also performed using Volatility 2 due to its wider adoption. Porting the efforts detailed in Chapter 6 to Volatility 3 [4] will be necessary as it starts to become more used over Volatility 2.

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