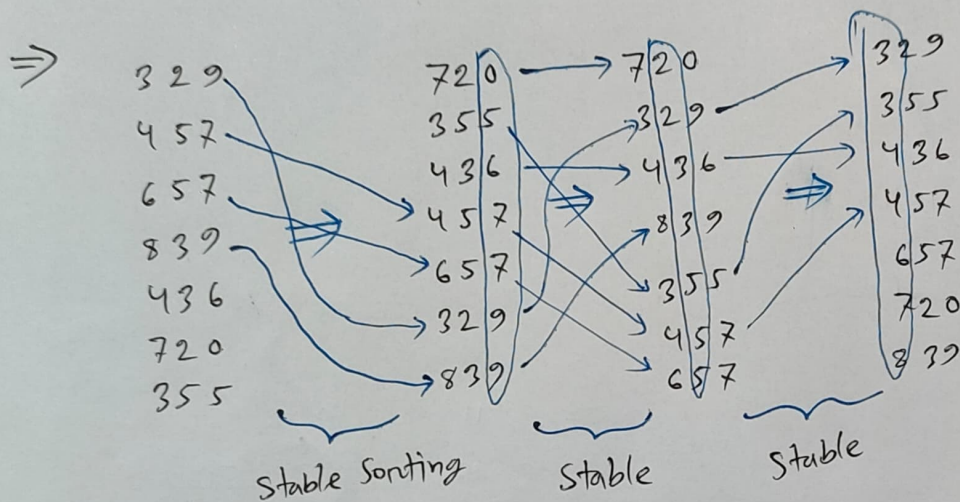


Radix Sort

- used by the card sorting machines.
- solves the problem of cards - counterintuitively - by sorting on the least significant digit first.



⇒ Algorithm:

RADIX-SORT(A, n, d)

number of maximum digit

$\theta(d)$ { for $i = 1$ to d
 use a stable sort to sort array $A[1:n]$ on digit i

Counting Sort ⇒ Best θ in this case, because k is known and value is 9.

$\theta(n+k)$ ← Merge Sort

Insertion Sort

Total time ⇒ $\theta(d(n+k))$
 $= \theta(dn)$

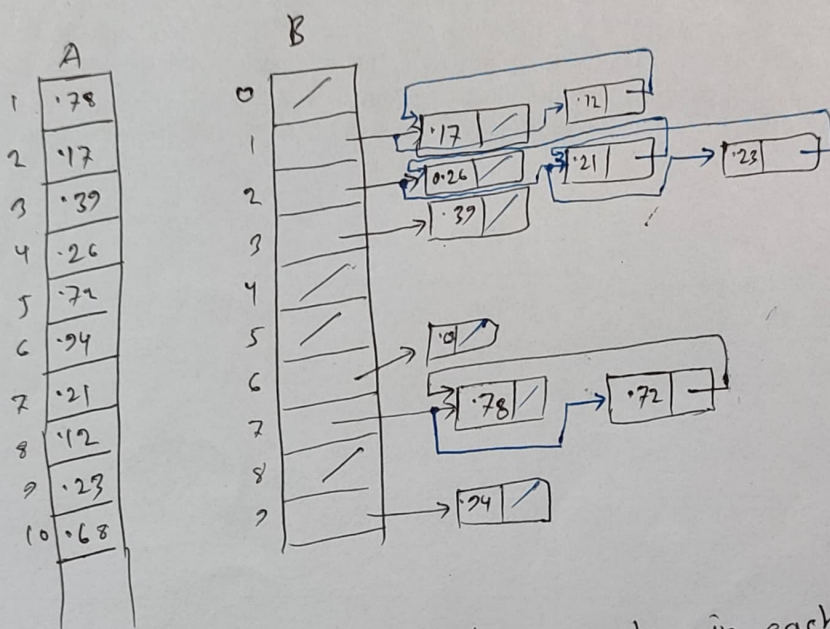


Bucket Sort

- assumes, input is drawn from a uniform distribution.
over the interval $[0, 1)$.

- divid in n equal-sized subintervals or buckets.

⇒ Operations:



Now nodes in each bucket will be
sort.

⇒ Algorithm:

BUCKET-SORT (A, n)

let $B[0:n-1]$ be a new array

$O(n)$ ← for $i = 0$ to $n-1$
make $B[i]$ an empty list

$O(n)$ ← for $i = 1$ to n
insert $A[i]$ into list $B[\lfloor n \cdot A[i] \rfloor]$

$\sum_{i=0}^{n-1} O(n_i)$ ← for $i = 0$ to $n-1$
sort list $B[i]$ with insertion sort $O(n_i)$

concatenate the list $B[0:n-1]$ together in order

$O(n)$ ← return the concatenated lists

$$\begin{aligned}
 \Rightarrow E[T(n)] &= E\left[\theta(n) + \sum_{i=0}^{n-1} O(n_i)\right] \\
 &= \theta(n) + \sum_{i=0}^{n-1} E[O(n_i)] \\
 &= \theta(n) + \sum_{i=0}^{n-1} O[E(n_i)]
 \end{aligned}$$

Here,

n elements for n buckets,

Probability, $p = \frac{1}{n}$ to fall each element onto a particular bucket.

From binomial distribution,

$$\text{mean } E[n_i] = np = 1$$

$$\begin{aligned}
 \text{variance } \text{Var}[n_i] &= np(1-p) \quad ; \quad p = \frac{1}{n} \\
 &= 1 - \frac{1}{n}
 \end{aligned}$$

$$\begin{aligned}
 \therefore E[n_i] &= \text{Var}[n_i] + E^2[n_i] \\
 &= 1 - \frac{1}{n} + 1^2 \\
 &= 2 - \frac{1}{n}
 \end{aligned}$$

$$\begin{aligned}
 \therefore E[T(n)] &= \theta(n) + n \times O\left(2 - \frac{1}{n}\right) \\
 &= \theta(n)
 \end{aligned}$$

A

Quiz - 2
Chapter - 6, 7, 8
31.03.2024

H.W. \Rightarrow 8.2 - 1, 3
8.3 - 1, 2, 5
8.4 - 1, 2