

⊕ Class Performance  $\Rightarrow$  51.

Project  $\Rightarrow$  51.

HW / Assignment  $\Rightarrow$  101.

Programming  $\Rightarrow$  51.

Quiz  $\Rightarrow$  2 out of 2  $\Rightarrow$  201. [ 35 minute conceptual question ]

Exam-1 (Mid)  $\Rightarrow$  251.

Exam-2 (Final)  $\Rightarrow$  301. [ Comprehensive syllabus ]

⊕ Ebook  $\Rightarrow$  Sebersta

⊕ Number Theory

- Research Area

⊕ Q-1  $\Rightarrow$  L8-L9

Mid  $\Rightarrow$  L16

Q-2  $\Rightarrow$  L22-23

⊕ Why the first programming language is formed as sequential?

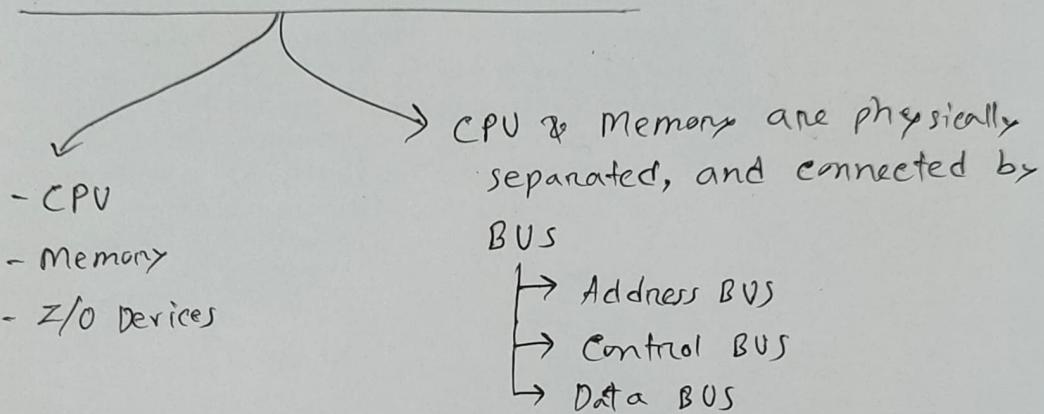
$\Rightarrow$  Computer was task specific and we need to modify hardware to build new algorithm. So, it depends on hardware. And there is no other way.

⊕ Programming language need to be user friendly. Because lot of time spend on debugging rather than coding.

## ④ Programming Languages:

- Sequential C, P
  - Functional Lisp
  - Logical Prolog
  - OOP Java, C++, R
  - Non-OOP C, B
  - Structural Fortran, P
- get some comparative idea

## ⑤ Von-Neumann Architecture



## ⑥ CPU: Central Processing Unit

⇒ Control Unit: Control order of the instruction to be executed. Also, controls the retrieval of desired operands.

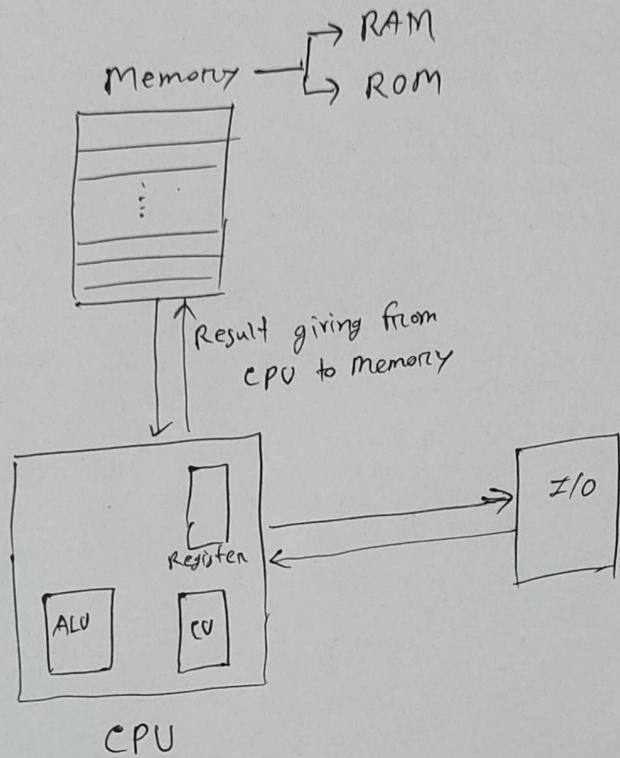
## ⇒ ALU: Arithmetic Logic Unit

- Perform all the logical and mathematical operations.

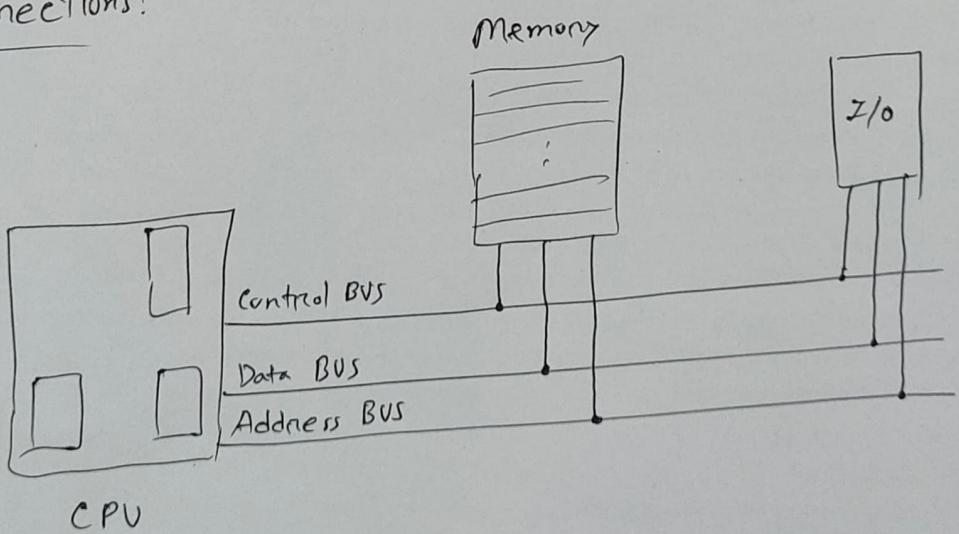
⇒ Register: Temporary storage to store and transfer data upon fetching from main memory.

## ⊕ Von Neumann Architecture:

- CPU and memory are physically separated.



## ⊕ BUS Connections:



## (\*) Von neumann bottleneck

- Fetching , sometimes CPU stay on idle.

CXL connection in data center

(\*) At first programmers need to do code using instruction or short code. Then programming languages arrived.

Using 1 & 0  
- Low level

## (\*) Machine Code:

- Sequence of 1's and 0's

Very complex

xxxx  
4 bits for OP. Code

xxxx xxxx xxxx  
12 bits for operand

⇒ To simplify the process, we got compiler in place

↳ 1951 (Grace Hopper)

## (\*) Memory :

- RAM ⇒ Volatile
- ROM ⇒ Non-Volatile

Female Scientist  
from US Navy

## Compilation:

⇒ Three ways with which a programming language can be implemented.

### 1. Compiler:

- Program translated to machine code.
- C/C++

### 2. Interpreter:

by a software

- Program is interpreted not translated.
- Python / Java Script

### 3. Hybrid Approach:

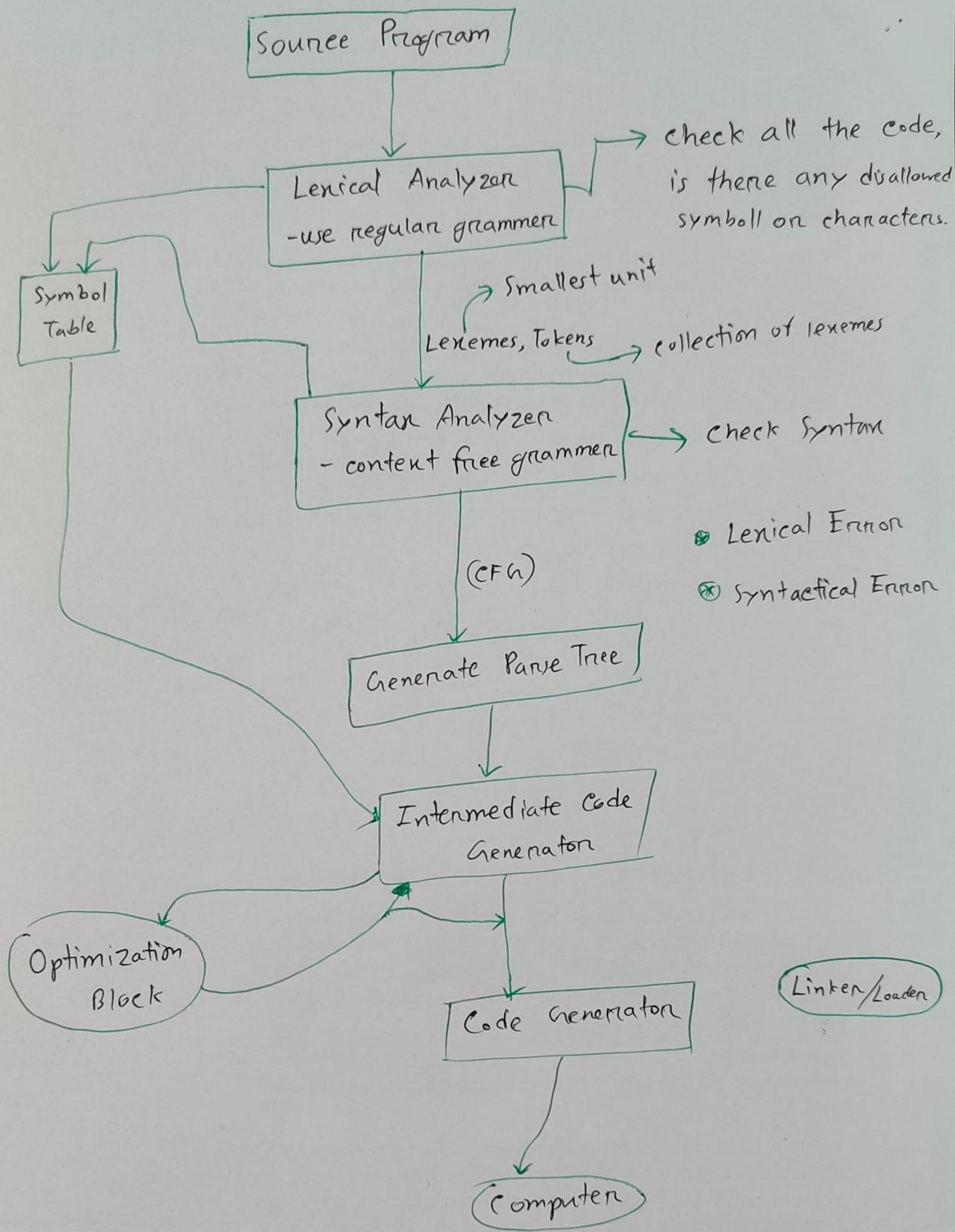
- compromised approach, 1 & 2 combined.
- Java

## Compiler based implementation:

⇒ Translator, that translate high-level instruction directly to the machine code.

⇒ Done in two step:

- generates an intermediate code, which structure is very similar to assembly language code.
- an assembler is in place to do the final translation to machine code.



L-03 / 02.09.2029 /

Restart

No content

⊗ Language construction (Implementation):

LLVM  
Compiler

- compiler based
- interpreter based
- Hybrid

⊗ We need to optimize our code in the intermediate code generator phase, because it's harder to optimize the code in machine language formate.

⊗ Intermediate Code Generation:

⇒ TAC - Three Address Code

- Maximum 3, not more than 3 operands

- use output from syntax analyzer ⇒ AST

⊗ So, AST is the input of Intermediate code.

⇒ Very close to assembly language code:

$t_1 = t_1 + t_2$

operator

operand

destination

$\otimes \quad a = b * c + b * d ;$

$\Rightarrow$

$$\begin{aligned} -t_1 &= b * c ; \\ -t_2 &= b * d ; \\ -t_3 &= -t_1 + -t_2 ; \\ a &= -t_3 ; \end{aligned} \quad \left. \begin{array}{l} \\ \\ \end{array} \right\} \text{b used twice.}$$

$\otimes \quad x = (-b + \sqrt{b^2 - 4 * a * c}) / (2 * a) ;$

$\Rightarrow$  TAC

$$t_1 = b * b ;$$

$$t_2 = 4 * a ;$$

$$t_3 = t_2 * c ;$$

$$t_4 = t_1 - t_3 ;$$

$$t_5 = \sqrt{t_4} ;$$

$$t_6 = 0 - b ;$$

$$t_7 = t_5 + t_6 ;$$

$$t_8 = 2 * a ;$$

$$t_9 = t_7 / t_8 ;$$

$$x = t_9 ;$$

Flow control TAC is important for exam

$\otimes$  Variable assignment in TAC

variable = constant

var1 = var2

var1 = var2 OP var3

var1 = constant OP var2

$\otimes$  Permitted operators:

$+, -, *, /, \%$

$\otimes$  Boolean & others:

$<, ==, ||, \neq$

$$\begin{array}{l|l} Y = -X ; & Y = -1 * n ; \\ Y = 0 - n ; & \end{array}$$

(\*)  $b = (x <= y);$

$\Rightarrow TAC:$

$t_0 = x < y;$

$t_1 = x = y;$

$t_2 = t_0 \parallel t_1;$

(\*) Flow Control:

$\Rightarrow$  two way to control the flow

1. ~~GOTO~~ GOTO "Label"

2. IFZ value GOTO "Label"



int x;

int y;

while ( $x < y$ ) {

$x = x * 2;$   
}

$y = x;$

$\Rightarrow TAC:$

L0:  $t_0 = x < y;$

IFZ  $t_0$  GOTO L1;

$x = x * 2;$

GOTO L0;

L1:

$y = x;$

## TAC - Three Address Code

- Three operand at most
- Flow control
  - GOTO Label
  - IFZ value GOTO Label

for exam  
array  
with  
Practice



```
int x;
int y;
while (x < y) {
    x = x * 2;
}
y = x;
```

### TAC

L0:

$t_0 = x < y ;$

IFZ  $t_0$  GOTO L1;

$y = x * 2 ;$

GOTO L0;

L1:

$y = x ;$

CS143: Compiler



```
void main() {
    int x, y;
    int i = x * x + y * y;
    while (i > 5) {
        i = i - x;
    }
}
```

### TAC

main:

BEGIN Function "Label";

$t_0 = x * x ;$

$t_1 = y * y ;$

$i = t_0 + t_1 ;$

L0:

$t_2 = 5 < i ;$

IFZ  $i$  GOTO L1;

$i = i - x ;$

GOTO L0;

L1:

END FUNCTION;



Practice more with for loop  
& array, nested if-else.

 TAC is very high level assembly language.

TAC :  $C = A + B;$

ADD A, B, C ;

ADD C, A, E ;

Assembly :

[LD A], R1 ;

LD B, R2 ;

ADD R1, R2, R3 ;

[ST R3, C] ;

[LD C], R4 ;

[LD A], R5 ;

ADD R4, R5, R6 ;

ST R6, E ;

 We need to optimize these redundant use.

## Optimization Process :

- Address mode selection
  - helps to improve instruction
- Peephole optimization
  - Remove redundant expression
- Common sub-expression elimination
  - remove redundant computation
- Register allocation
  - reduce number of register allocation

 Graph theory, graph coloring for register allocation.

$\Rightarrow$  available in Kenneth Rosen's book.

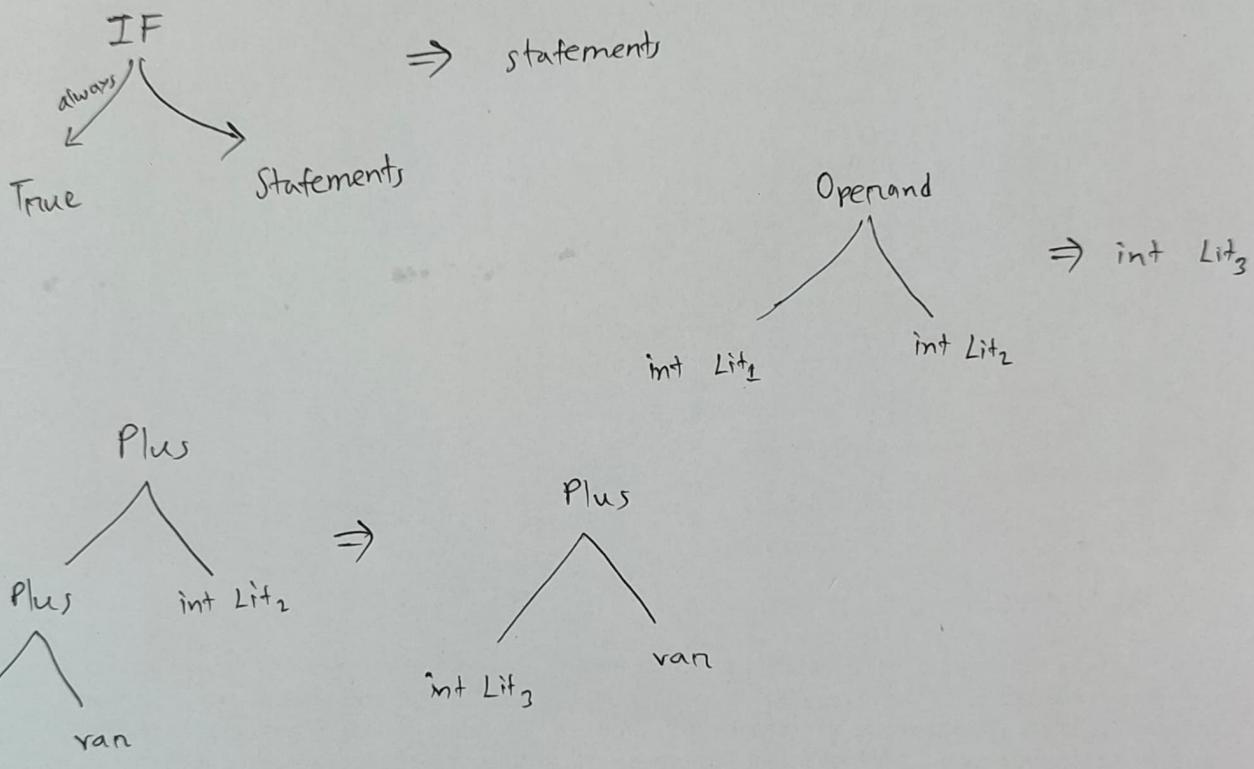


## Peephole optimization:

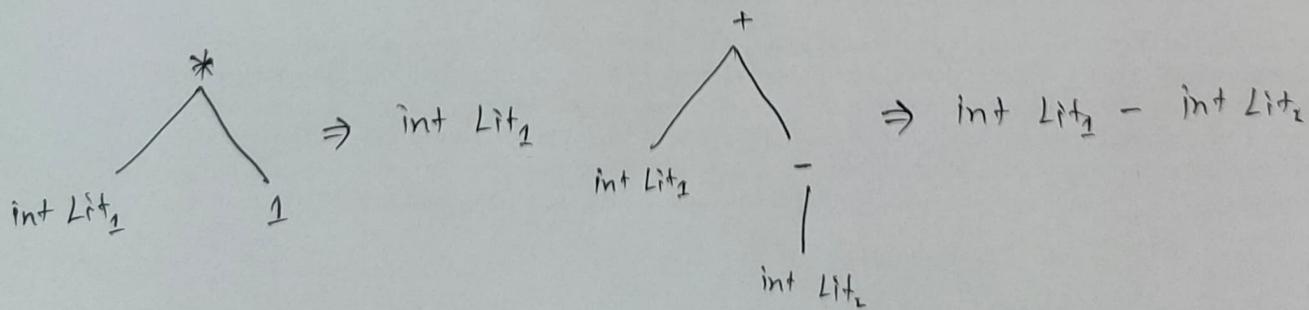
- could be done on
  - the AST level
  - the IR level
  - the generated code level



## AST level:



## Constant folding:



## Constant Propagation:

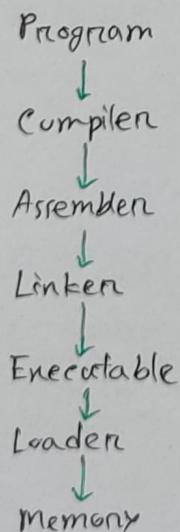
```
int a = 30;  
int b = 9 - (a/5);  
int c;  
c = b * 4;  
if (c > 10){  
    c = c - 10;  
}  
return c * (60/a);
```

```
int a = 30;  
int b = 3;  
int c = 12;  
if (c > 10){  
    c = c - 10;  
}  
return c * 2;
```

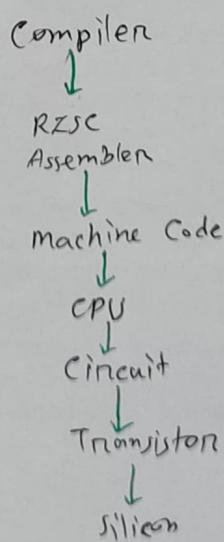
 if we know the calculation, we need to do it during compilation.

L-06 / 11.02.2024 /

## In C,



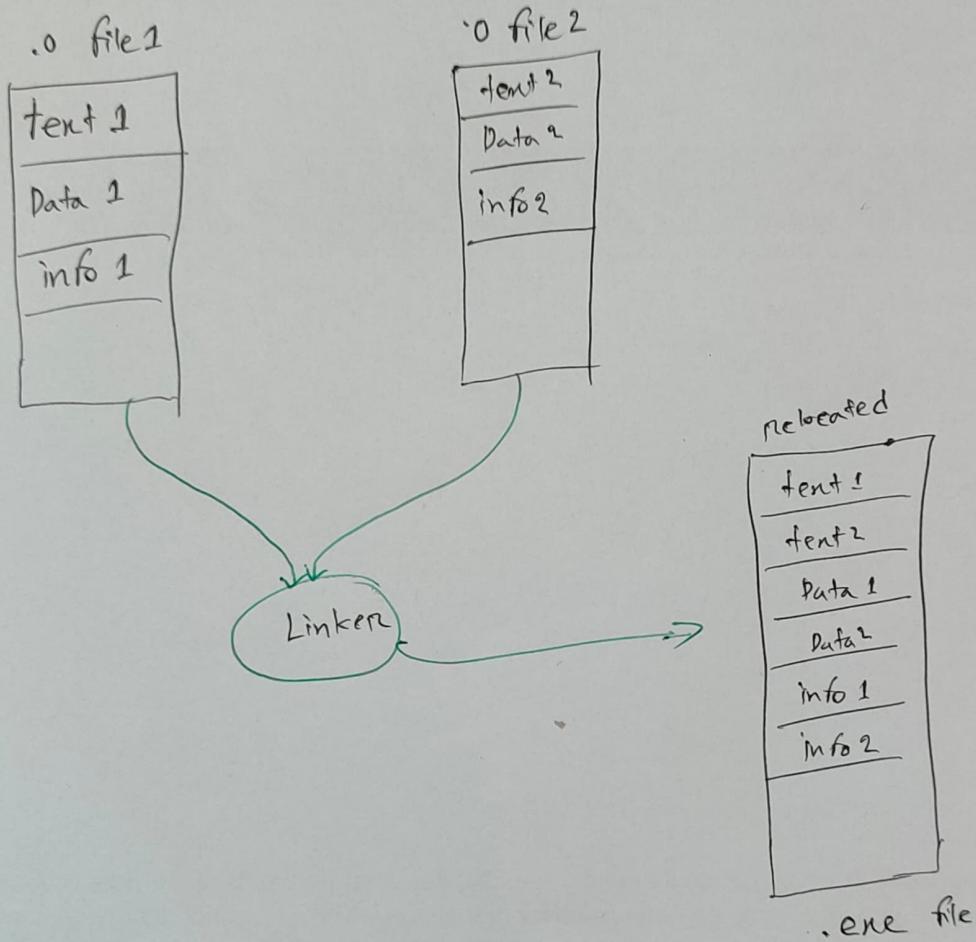
In C,



int x = 10;  
x = 2 \* x / 15;

ADDI x5, x0, 10  
MULI x5, x5, 2  
ADDI x5, x5, 15

001001  
10011011  
100010010  
— —



## \* Object file:

⇒ **Text**: Main code from the instruction

⇒ **Data**: Binary representation of the static data

⇒ **Symbol Table**: Contains list of labels of these files and static data that can be referenced.



## Task Performed by the Linker:

- Take text segment of each .o files and include those in the executable files.
- Same for the data segment.
- Resolve References.

(\*) Linker produces the executable files.

(\*) Linker must solve,

- 'printf' routine (Which was unknown to assembler)
  - because its an external object
  - put 'zero' address
- Linker does the memory re-arrangement
  - address resolution.

## Language Design Criteria:

- availability (Price)
- Politics

"C" is popular due to availability and it is included in Union OS  
Advantage: Very fast  
Disadvantage: Learning is difficult.

2-07 / 18.09.2024

❶ extern float, sin() 

extern printf(), scanf()

 stdio.c



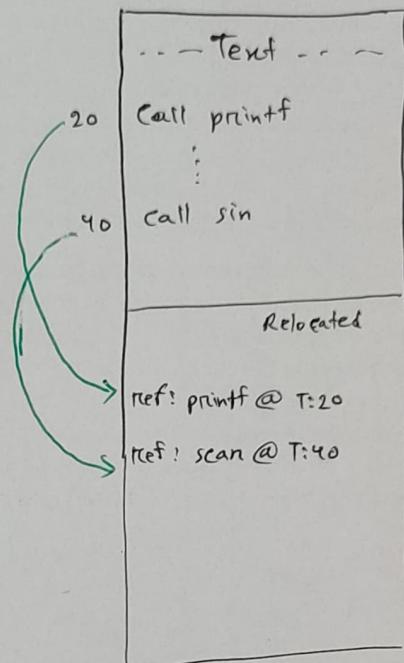
main() {

printf(---);

scanf(---);

result = sin();

}



## ❷ Efficiency:

⇒ System Level

⇒ Programmers Perspective

- 90% time we spend in debugging the code.
- maintainability
- how well the features of the PL are integrated.
- regularity
- not too many surprises

### Pseudocode

```

if (x > 0) {
    num_solutions = 2;
    r1 = sqrt(x);
    r2 = -r1;
}
else if (x == 0) {
    num_solutions = 1;
    r1 = 0.0
}
else
    num_solutions = 0;

```

### Python

```

if x > 0.0:
    num_solutions = 2
    r1 = sqrt(x)
    r2 = -r1
elif x == 0.0:
    num_solutions = 1
    r1 = 0.0
else
    num_solutions = 0

```

### ⊗ Regularity:

- Generality
- Orthogonality ~~xxx~~
- Uniformity

### ⊗ Uniformity:

- Refers to consistency in appearance
- similar things should behave in a similar way

```

int function () {
    --- -
}

```

```

class name {
}

```

⇒ Violation of uniformity.

## Generality:

- Avoids special cases in the availability or use of language construct.
- Example: nested function not possible in 'C'.
  - ⇒ in 'C'  
"==" is used to compare equality.  
However, two arrays can't be comparing  
'==' in 'C'.

## Orthogonality:

- Language constructs don't behave differently in different contexts.

⇒ C, C++ : values of all data types can be returned from a function except arrays.

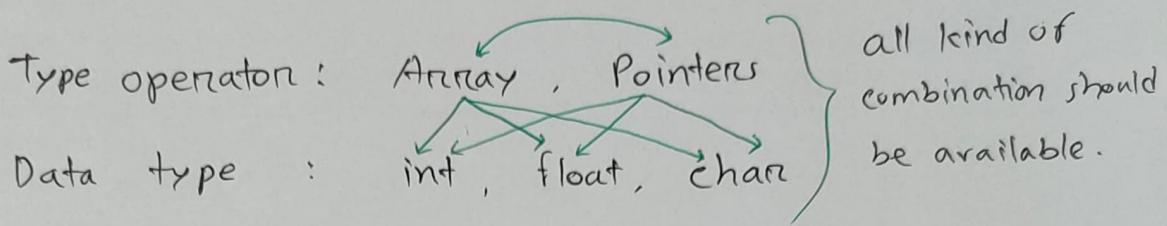
⇒ ALGOL - Algorithmic Language

⇒ first general purpose language

⇒ one of the most orthogonality

⇒ Also too much orthogonality

### Orthogonality:



⇒ Any combination should be valid language construct.  
- too much is not good.

### Array of pointers code:

```
#include <stdio.h>

int main() {
    int var1 = 10;
    int var2 = 20;
    int var3 = 30;
    int *ptr_array[3] = { &var1, &var2, &var3 }
}
```

Define array of pointer and write an iterative code to show it. \*\*\*



## ALGOL : Algorithmic Language - 1960

- conditional assignment

..... = .....

$\Rightarrow$  Changes abruptly at the LHS.

- causing enormous complexity.

- didn't get the popularity it expected.



## Expressiveness!

- easy to write code

- complex process & structure to be implemented concisely.

while (\*b != 0) {

\*a = \*b;

a++;

b++;

}

while (\*a++ = \*b++)

# include <stdio.h>

void main() {

int a;

float b, c, d;

a = 750;

b = a / 350;

c = 750;

d = c / 350;

printf ("%.2f %.2f", b, d);

}

$\Rightarrow$  What would be the output?

$\Rightarrow$  2.00 2.14

$$\begin{aligned} \text{Ans: float } f &= \frac{\text{float}(a)}{\text{float}(b)} \\ &= \frac{(\text{float}) a}{(\text{float}) b} \end{aligned}$$



## Operator Overloading:

a = 10;

b = 5;

a + b = 15;

↳ addition

↳ '+' symbol may be used in string concatenation.

⇒ use of '&' → logical AND

↳ unary operator

↳ has to do with address

=	+	-
==	++	--
!=		



Types

Type checking

} Security & Reliability

L-9 / 25. 09. 2024

Project Idea

1. Bengali Programming Language design for secondary level  
Students.

- Compiler base
- conditional statement
- Iterative (For loop)
- Basic File handling
- Array

2. Vectorized / Parallelization of Language construct.
3. Domain specific Programming Language.  
↳ Graph based PL.

## Reliability!

- if program perform as per the specification in all conditions.
- Type checking
- Exception Handling
- Aliasing
- Readability & Maintainability

## Type checking:

- Done during compilation phase or during run time.
  - ↳ faster
  - ↳ preferred at the early phase of implementation.

```
=> void test-function (int a) {  
    : --  
    : --  
}  
  
        }  
int main(void){  
    : --  
    : --  
    test-function (700);  
}
```

↳ Type Error



## Exception handling:

- intercepting run-time errors, and take corrective measures
- Divided by 'zero' issue
- C++, Java : Exception management ~~not~~ available.



## Aliasing:

- Same memory location being accessed through more than one referencing.
- Allowed but very detrimental in ~~reliability~~ reliability issue.



## History



## Von Neumann Architecture:

- shared-program architecture
- conditional control transfer (Branching - conditional statement)
  - sub-routine
  - looping



## First programming language

⇒ Short Code - 1949

- Instruction were written as sequence of 1's and 0's.

Then designed a compiler.

⇒ First compiler - 1951 - by Grace Hopper.

Then came, Assembly Language.

- mnemonic based language

↳ MOV, LD/LDZ, MUL/MULZ

- close to machine code

- assembler does the conversion to machine code.

Then Fortran, and then LISP.

LISP: List Processing Language

- second best programming language.

- 1958 - John McCarthy

- basic & only data type was LIST

- introduced symbolic programming

⇒ n or y | OR(x,y)

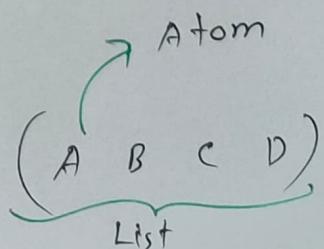
- used in AI application

descendent such as SCHEME

LIST:

- Lists

- Atom ⇒ identifier / Numeric Letters.



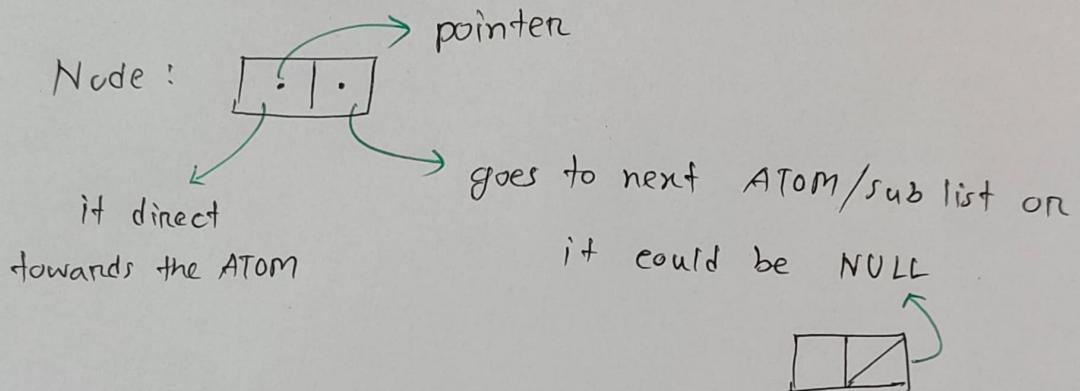
$(A (B C) D (E (F G)))$   $\Rightarrow$  List within a List  
- nested list

$\Rightarrow$  stored sparsely in memory

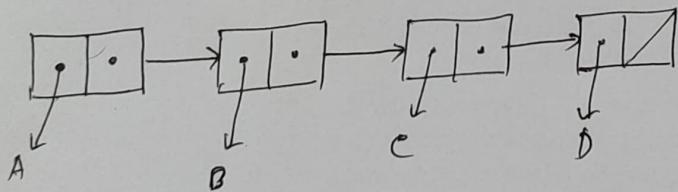
- two pointers for each atom.

How to represent:

- Box pointers diagram

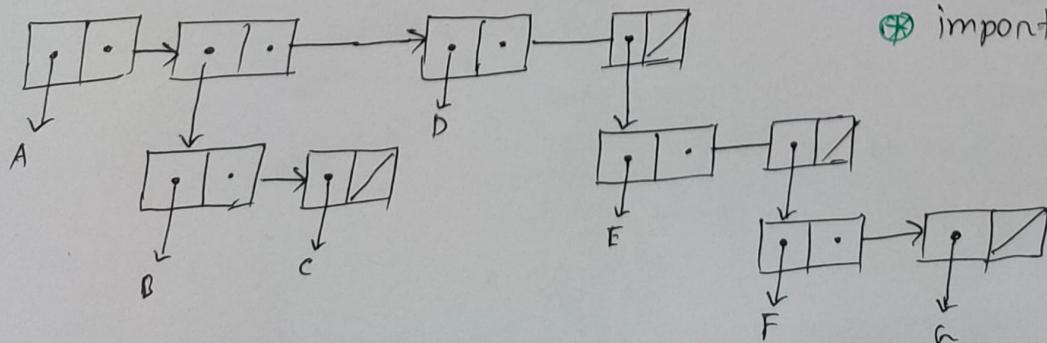


$\Rightarrow$



⊗

$(A (B C) D (E (F G)))$



⊗ important for exam.