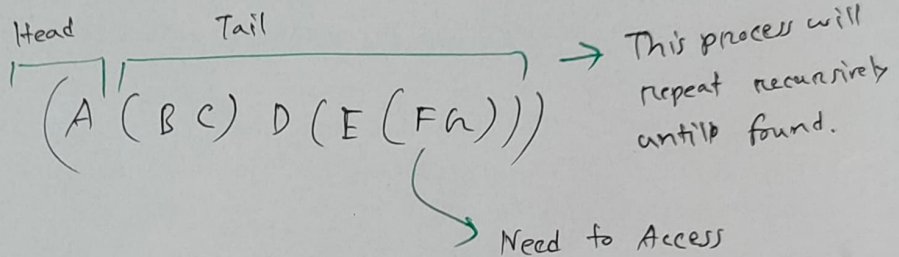


⊗ Linked list will be given

- show the box pointer diagram
- write a code in c++
- how to access the list.



⊗ More on LISP! → 1970 (MIT)

- Purely functional programming language.

- Recursive calling of function
- Lambda expression (~~Anonymous~~ Anonymous)
  - Lambda Calculus
- Not the traditional variable concept
- Introduce Symbolic Programming.

⊗ Composite function

$$\underbrace{\left. \begin{matrix} f(x) \\ g(x) \end{matrix} \right\} \left. \begin{matrix} f(g(x)) \\ g(x) \end{matrix} \right\}}_{\text{Higher order filter design.}} \quad \text{fog, gof}$$

## ⊗ Lambda expression:

$$\left. \begin{aligned} \text{cube}(n) &= x * x * x \\ \text{cube}(2) &= ? \\ &= 8 \end{aligned} \right\}$$

$$(\lambda(x) x * x * x)(2) = 8$$

- ⊗ will be given.
- ⊗ what will be the output?

⊗ Hardware deprecated this language.

## ⊗ Scoping:

- visibility / accessibility of a variable.
- may be in the containing block
- may be from global definition
- sometimes dynamically resolved.
- Benefit: we can reuse a variable name.

⊗ Two type of scoping:

### ① Static:

- resolve from containing block then outer containing block, if not found, look for in the global.

### ② Dynamic:

- Local definition in the point of invocation.  
containing block.





## Scoping:

```
int a=10, b=20;

int main()
{
    int a=5;
    {
        int c;
        c = b/a;
        printf("1.8", c);
    }
}
```

value of b and a missing in the containing block.

value of a found in the parent block but b is still missing.

value of b found on the global definition.



```
sample()
{
```

```
    y = 5;
```

```
    x = 20;
```

```
    x = 20 + y;
```

```
    Test();
```

```
}
```

```
Test()
{
```

```
    printf(x);
```

```
}
```

Point of invocation. from which point function was called. In that point function will search for value of x. Means what was the value of x in that position.



sample 1()

```
{
    printf(x);
}
```

sample 2()

```
{
    double x;
    x = 5.5;
    sample 1();
}
```

x = 0.25;

main()

```
{
    sample 1();
    sample 2();
}
```

⇒ What will be the output for static & dynamic? ~~\*\*\*~~ important for exam

Output:

Static:

0.25

0.25

Dynamic:

0.25

5.5



Home-Task:

void main()

```
{
    m(a);
}
```

void m(x)

```
{
    a = 2;
    a = x - a;
    n(a);
    printf(a);
}
```

a = 3;

void n(x)

```
{
    x = x * a;
    printf(x);
}
```

⊗ What will be the output?

⇒

Static

6

2

Dynamic

4

2



~~ALGOL~~ 58,60!

- concept of data types introduced.
- concept of compound statement:  
$$\begin{array}{l} \{ \text{! begin} \\ \quad \vdots \\ \quad \} \text{! end} \end{array}$$
- Identifiers could be of any length.
- Any dimensional array
- Nested statement.

58 → 60: upgrade

- upgraded
- Motivation: Fortran for IBM
  - Limited portability
  - still universal program was missing.
  - for ~~exp~~ specific computation.

## Syntax Analyzer:

- CFG : Context Free Grammar

↳ BNF  $\rightarrow$  EBNF (modern language use)

For exam:

Functional Programming Basic is important.

L-11/02.10.2024/

Fortran Practice	upto
Next Class Quiz -	Dynamic
	Array

L-12/ 07.10.2024/

## Quiz-1

### \* Syntan Analyzer:

- set of rules that define the 'structure' of the accurate source code.

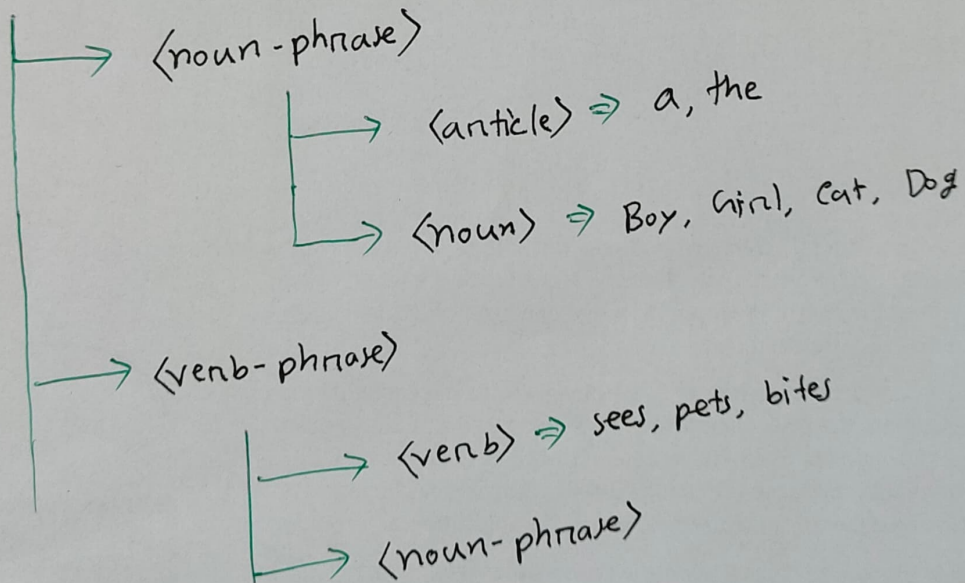
- we use CFG → Content Free Grammar

→ Another one is BNF (Backus-Naur Form)

\* Why should ~~be~~ this be context free?

⇒

<sentence>



\* <Sentence> ⇒ <article> <noun> <verb> <article> <noun>

⇒ a boy sees a dog.



⊗ Making sentences using or following the grammar rules and real meaning is an instance of content sensibility.

In contrast, whatever we can do/make is content flexibility.

⊗ Context free grammar produces all the possible combination and check any thing that is not allowed and doesn't exist in the source code.

⊗ CFG (Context Free Grammar)

⊗ Lexemes:

- a string of characters that are the lowest syntactical unit.
- numerical literals.
- identifiers  $\Rightarrow$  could be a single token
- Reserve words/ keywords (if, else, or, do, while, switch...)

⊗ Tokens:

- A group of/collection of lexemes given a name,
- However, sometimes a single lexeme can also form token.  $\Rightarrow (+)$

⊗ Identifiers:

- name of variable
- name of a function
- a name that identify something.

L-13 / 09.10.2024 /

Online

Fortran Practice