



# Chapter 3 Stacks (1)

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College of Computer Science, CQU

# Outline

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- ❑ **Stack ADT**
- ❑ **Array-Based Stack**
- ❑ **Linked Stack**
- ❑ **Comparison of Array-Based and Linked Stacks**
- ❑ **Applications**

# Stacks

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- ❑ The stack is a list-like structure in which elements may be inserted or removed from only one end, called the top of the stack.
- ❑ All access is restricted to the most recently inserted elements
- ❑ Basic operations are push, pop

# Stacks

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LIFO: Last In, First Out.

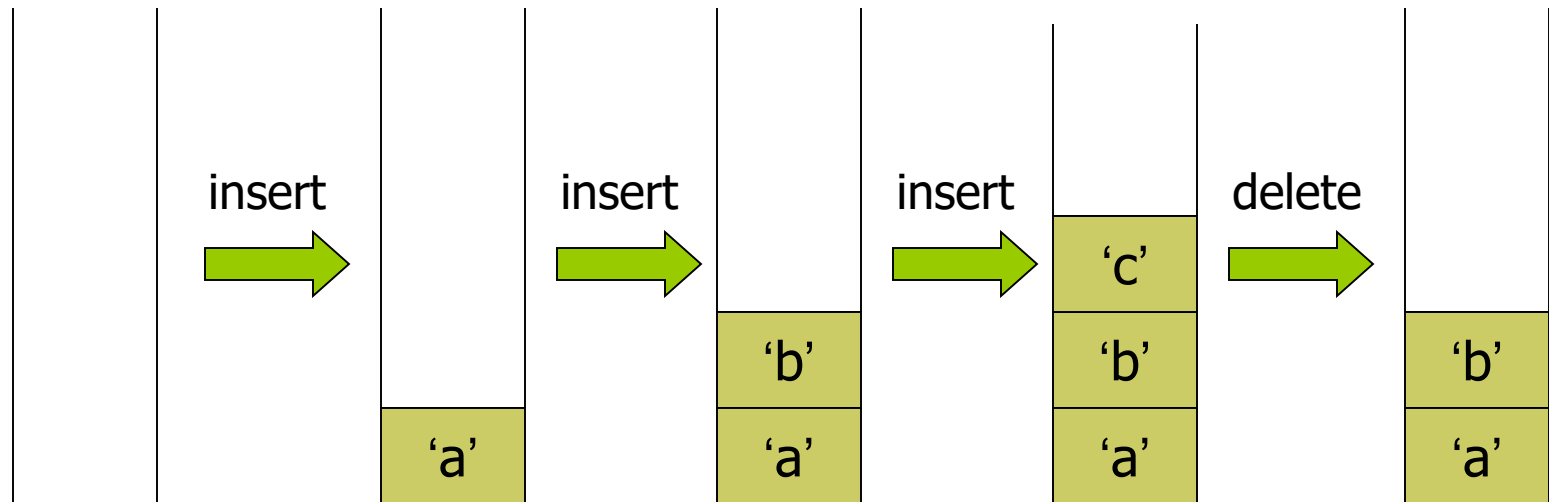
Restricted form of list: Insert and remove only at front of list.

Notation:

- ❑ Insert: PUSH
- ❑ Remove: POP
- ❑ The accessible element is called TOP.

# Example: Stack of Char

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# Stack ADT

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```
// Stack abstract class  
template <typename E> class Stack {  
private:  
    void operator =(const Stack&) {} // Protect assignment  
    Stack(const Stack&) {} // Protect copy constructor  
public:  
    Stack() {} // Default constructor  
    virtual ~Stack() {} // Base destructor  
// Reinitialize the stack. The user is responsible for  
// reclaiming the storage used by the stack elements.  
    virtual void clear() = 0;
```



# Stack ADT

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```
// Push an element onto the top of the stack.  
// it: The element being pushed onto the stack.  
virtual void push(const E& it) = 0;  
// Remove the element at the top of the stack.  
// Return: The element at the top of the stack.  
virtual E pop() = 0;  
// Return: A copy of the top element.  
virtual const E& topValue() const = 0;  
// Return: The number of elements in the stack.  
virtual int length() const = 0;  
}
```

# Array-Based Stack

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- ❑ The array-based stack implementation is essentially a simplified version of the array-based list.
- ❑ The only important design decision to be made is which end of the array should represent the top of the stack.
  - One choice is to make the top be at **position 0** in the array. This implementation is inefficient. (Why? )
  - The other choice is have the top element be at **position n-1** when there are n elements in the stack.
- ❑ For the implementation of Figure 4.18, **top** is defined to be the array index of the first free position in the stack. Thus, an empty stack has top set to 0, the first available free position in the array.



# Array-Based Stack

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```
// Array-based stack implementation
private:
    int size;          // Maximum size of stack
    int top;           // Index for top element
    Elem *listArray;   // Array holding elements
public:
    AStack(int size =defaultSize) // Constructor
        { maxSize = size; top = 0; listArray = new E[size]; }
    AStack() { delete [] listArray; } // Destructor
    void clear() { top = 0; } // Reinitialize
    void push(const E& it) { // Put "it" on stack
        Assert(top != maxSize, "Stack is full");
        listArray[top++] = it;
    }
}
```



# Array-Based Stack

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```
E pop() { // Pop top element
    Assert(top != 0, "Stack is empty");
    return listArray[--top];
}

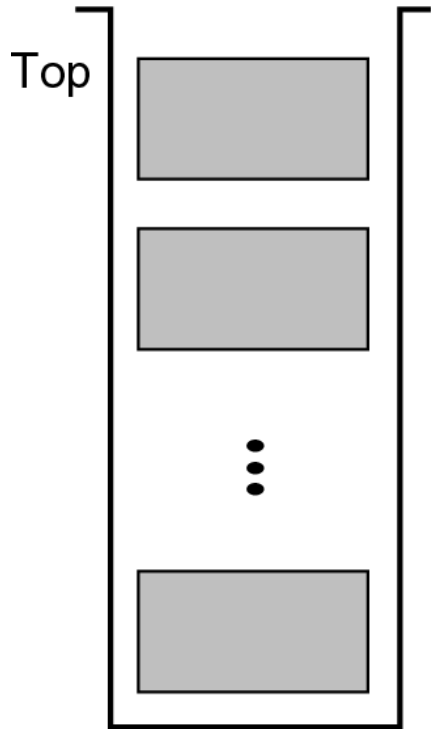
const E& topValue() const { // Return top element
    Assert(top != 0, "Stack is empty");
    return listArray[top-1];
}

int length() const { return top; } // Return length
};
```

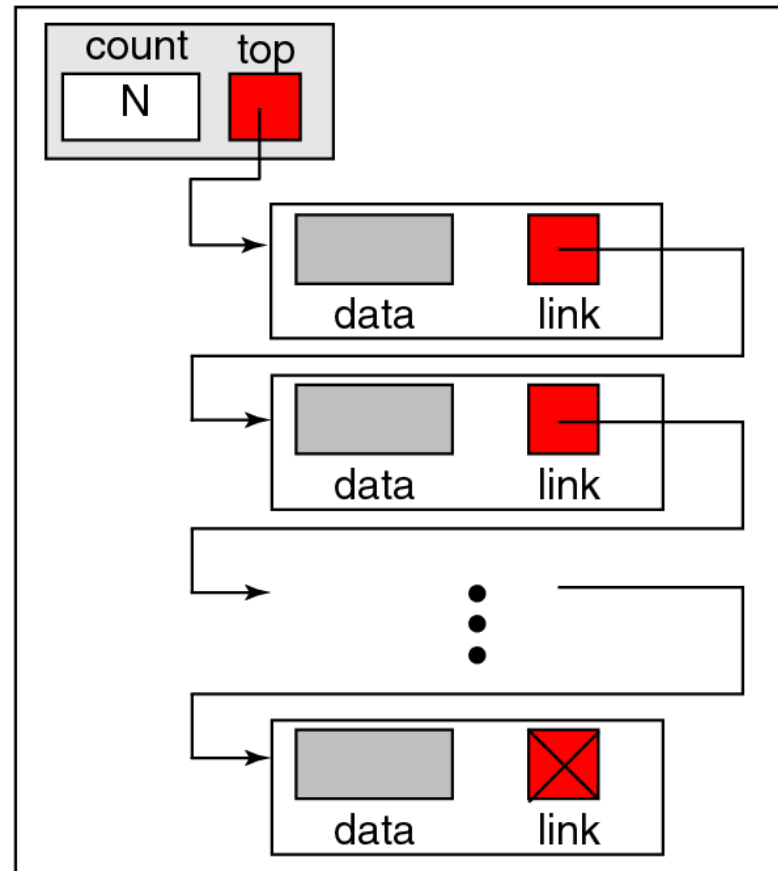
Issues:

- ❑ Which end is the top?
- ❑ What is the cost of the operations?

# Linked Stack



Conceptual view



Linked list implementation

# Linked Stack

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- ❑ The linked stack implementation is quite simple. The freelist of Section 4.1.2 is an example of a linked stack.
- ❑ Elements are inserted and removed only from the head of the list. A header node is not used because no special-case code is required for lists of zero or one elements.
- ❑ The only data member is **top**, a pointer to the first (top) link node of the stack.



# Linked Stack

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```
// Linked stack implementation
private:
    Link<Elem>* top; // Pointer to first elem
    int size;        // Count number of elems
public:
    LStack(int sz =defaultSize) // Constructor
    { top = NULL; size = 0; }
    LStack() { clear(); } // Destructor
```



# Linked Stack

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```
void clear() { // Reinitialize
    while (top != NULL) { // Delete link nodes
        Link<E>* temp = top;
        top = top->next;
        delete temp;
    }
    size = 0;
}

void push(const E& it) { // Put "it" on stack
    top = new Link<E>(it, top);
    size++;
}
```



# Linked Stack

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```
E pop() { // Remove "it" from stack
    Assert(top != NULL, "Stack is empty");
    E it = top->element;
    Link<E>* ltemp = top->next;
    delete top;
    top = ltemp;
    size--;
    return it;
}
```



# Linked Stack

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```
const E& topValue() const { // Return top value
    Assert(top != 0, "Stack is empty");
    return top->element;
}
int length() const { return size; } // Return length
};
```

**What is the cost of the operations?**

**How do space requirements compare to the array-based stack implementation?**





# Comparison of Array-Based and Linked Stacks

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- ❑ All operations for the array-based and linked stack implementations take constant time
- ❑ Total space required: the analysis is similar to that done for list implementations
- ❑ When multiple stacks are to be implemented, it is possible to take advantage of the one-way growth of the array-based stack. This can be done by using a single array to store two stacks, as illustrated by Figure 4.20



# Stack Applications

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# Application 1

## Data conversion

- A simple algorithm uses the following law:

$$N = (N / d) \times d + N \% d$$

eg: Convert decimal to binary

$$(100)_{10} = (1100100)_2$$

- divide decimal number by 2 continuously
  - record the remainders until the quotient becomes 0
  - print the remainders in backward order.
- Stack is an ideal structure to implement this process.
  - How to implement it?

2	100	
2	50	0
2	25	0
2	12	1
2	6	0
2	3	0
2	1	1
	0	1

# Applications 2

## Balancing Symbols

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- ❑ Make an empty stack. Read characters until end of file.
- ❑ If the character is an opening symbol, push it onto stack.
- ❑ If it is a closing symbol, then if the stack is empty, report an error. Otherwise, pop the stack.
- ❑ If the symbol popped is not the corresponding opening symbol, then report an error.
- ❑ At the end of file, if the stack is not empty, report an error.

# Applications 2

## Balancing Symbols

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- Stack can be used to check a program for balanced symbols (such as {}, (), []).
- Example: {}() is legal, {(}) is not (so simply counting symbols does not work).
- When a closing symbol is seen, it matches the most recently seen unclosed opening symbol. Therefore, a stack will be appropriate.

# Example

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- How do we know the following parentheses are nested correctly ?

$7 - ((X * ((X + Y) / (J - 3)) + Y) / (4 - 2.5))$

1. There are an equal number of right and left parentheses.
2. Every right parenthesis is preceded by a matching left parenthesis.

$((A+B)$     or     $A+B($                       violate 1

$)A+B(-C$     or     $(A+B))-(C+D$                       violate 2

# Application 3

## Line Editor

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- For example ,the following two lines has been accepted from keyboard

```
Whli## ilr# e( s#* s)
```

```
outcha@putchar(*s=#++);
```

- They will be treated as:

```
while (*s) putchar(*s++);
```