李佳

2013-06-21

密

严肃考纪

公平竞争、诚实守信、

重庆大学《算法分析与设计》课程试卷

、 ○R巻

2012~2013 学年 第二学期

开课学院: <u>计算机学院</u> 课程号: <u>18016435</u> 考试日期: <u>2013-06-25</u>

考试方式: ○开卷 ⊙闭卷 ○其他

考试时间: 120 分钟

题 号	A	В	C	D	E	F	G	Н	I	J	总 分
得 分											

注: 答案须填写在试题页后的答案纸上, 用英文或中文解答。

A. (20 points)

- (1) With <u>substitution</u>(替换法), prove that the asymptotic complexity of $T(n) = T(n/2) + \Theta(n)$ comes to $\Theta(n \log n)$.
- (2) For the polynomial $f(n) = 3n^3 + 2^n + 3^n + 6n^4 + 79$, what is the tight bound (Θ)? And explain your answer <u>briefly</u>(简要地).
- B. (25 points) Given an array $A = \langle 4, 8, 1, 3, 7, 9, 20, 21, 13 \rangle$, try to employ the heapsort to sort all integers in A. Note that the heapsort consists of two steps: <u>max-heap building</u> (最大堆构成) and <u>sorting</u>(排序).
 - (1) In a binary heap, what is the index of the *k*-th element's PARENT? What is the index of the *k*-th element's LEFT child?
 - (2) Draw the <u>initial state</u>(初始状态) of the heap.
 - (3) Draw the final state(最终状态) of the heap after max-heap building.
 - (4) Count the exact number of <u>swapping</u>(交换) between two nodes (integers) in the heap during the max-heap building.
 - (5) What is the <u>asymptotic(</u>渐进的) computational complexity of the heapsort? Give a <u>brief explanation(</u>简要的说明).
- C. (10 points) In activity selection problem, suppose that <u>instead of selecting</u>(不选择) the first activity to finish, we always <u>select(选择)</u> the <u>last activity to start(</u>最后开始的活动) that

is compatible with all previously selected activities. Describe how this approach is a greedy algorithm, and prove that it <u>yields(</u>输出) an <u>optimal solution(</u>最优解).

D. (20 points) For input sequence $X_{[1..m]}$ and $Y_{[1..n]}$, we define c[i,j] = length of Longest Common Sequence of $X_{[1..i]}$ and $Y_{[1..j]}$; By using methodology of Dynamic Programming, we got c[i,j]'s recursive formula(递推表达式):

$$c[i,j] = \begin{cases} 0 & i = 0 \text{ or } j = 0 \\ i,j > 0 \text{ and } x_i = y_j \\ i,j > 0 \text{ and } x_i \neq y_j \end{cases}$$

- (1) Fill out(填写) the above formula.
- (2) Give a <u>memoized version(</u>备忘录方法) of DP algorithm to compute c[m,n] in O(mn) time. The algorithm should be given in <u>pseudo code(</u>伪代码) or any popular programming language(c, java, c# etc).

E. (25 points) Work out the <u>maximum flow(最大流)</u> of the next flow network using Edmond-Karp algorithm (通过 BFS 找出增广路径). That is, at <u>each iteration(每次循环)</u>, **draw** the <u>residual network(</u>残值网络) induced by the current flow network, **find** a <u>shortest path</u> (最短路径) from source s to sink t by a breadth-first search (BFS) in the residual network, then **augment**(增加) the flows of each edge on the path.

