

## **Practical Concurrent and Parallel Programming VI**

## **Performance and Scalability**

Jørgen Staunstrup

# Agenda



- Executors and Future (continued from Week 5)
- Scalability, speed-up and loss (of scalablity) classification
   Example: QuickSort
- Lock striping
  - A case study with Hash maps

# From week05 splitting tasks



Body of: <a href="mailto:countPrimesTask">countPrimesTask</a>

```
@Override
  public void run() {
   if ((high-low) < threshold) {
      for (int i=low; i<=high; i++) if (isPrime(i)) lc.increment();
   } else {
      int mid= low+(high-low)/2;
      pool.submit(new countPrimesTask(lc, low, mid, pool, threshold) );
      pool.submit(new countPrimesTask(lc, mid+1, high, pool, threshold) );
   }
}</pre>
```

- Shortcomings:

  1. How to stop?
- 2. Will create too many "small" tasks
- 3. Returning result (# primes)

# Splitting tasks

```
-4
```

```
public void run() {
  if ((high-low) < threshold) { ...
  } else {
    int mid= low+(high-low)/2;
    Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid, pool, threshold) );
    Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high, pool, threshold) );
    ...
}
```

## Combining tasks

```
-5
```

```
public void run() {
  if ((high-low) < threshold) { ...
  } else {
    int mid= low+(high-low)/2;
    Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid, pool, threshold) );
    Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high, pool, threshold) );
    try { f1.get();f2.get(); }
    catch (InterruptedException | ExecutionException e) { }
}
```

Does the order of f1.get and f2.get matter?

## Combining tasks

```
-7
```

```
public void run() {
  if ((high-low) < threshold) {</pre>
    } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, low, mid,
                                 pool, threshold) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, mid+1, high,
                                 pool, threshold) );
      try { f1.get();f2.get(); }
      catch (InterruptedException | ExecutionException e) {     }
       Shortcomings:
                                          How do we get the result?
```

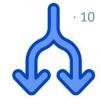
- 1. How to stop?
- 2. Will create too many "small" tasks
- 3. Returning result (# primes)

## Java Executors - Tasks



- Tasks are a central concept for executors
- When designing a program using executors, first think about the tasks to be executed (e.g. <u>countPrimesTask</u>)
  - Like for threads, tasks can be conveniently defined in their own class

## countPrimesTask



```
public class countPrimesTask implements Runnable {
 private final int low;
 private static boolean isPrime(int n) {
 public countPrimesTask(PrimeCounter lc, int low, int high,
   ExecutorService pool, int threshold) {
    this.lc = lc;
  @Override
 public void run() {
```

## Java Executors - Tasks



- Tasks are a central concept for executors
- When designing a program using executors, first think about the tasks to be executed (e.g. <u>countPrimesTask</u>)
  - Like for threads, tasks can be conveniently defined in their own class
- Ideally, tasks should be independent



```
public class countPrimesTask implements Runnable {
 private final int low; private final int high; priva
 private final ExecutorService pool;
                                                         Are countPrimesTasks
 private final PrimeCounter lc;
                                                              independent?
 private static boolean isPrime(int n) {...)
 public countPrimesTask(PrimeCounter lc...) {this.lc= lc;... }
  @Override
 public void run() {
    if ((high-low) < threshold) {</pre>
      for (int i=low; i<=high; i++) if (isPrime(i)) lc.increment();</pre>
   } else {
      int mid= low+(high-low)/2;
      Future<?> f1= pool.submit( new countPrimesTask(lc, ... ) );
      Future<?> f2= pool.submit( new countPrimesTask(lc, ... ) );
      try { f1.get();f2.get();
      catch (InterruptedException | ExecutionException e) {    e.printStackTrace();
                                                  Code in countPrimesTask.java
```

## PrimeCounter



```
class PrimeCounter {
 private int count= 0;
  public synchronized void increment() {
    count= count + 1;
  public synchronized int get() {
    return count;
  public synchronized void setZero() {
    count= 0;
```

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- Kick-off class for the program
- It initializes the Executor service

```
class PrimeCountExecutor {
   private ExecutorService pool;
    ...
   public PrimeCountExecutor () {
      pool= new ForkJoinPool();
      Future<?> done= pool.submit(new countPrimesTask( ... ));

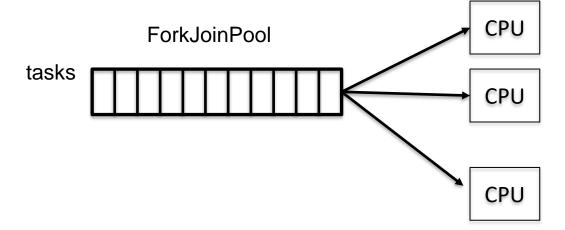
      try { done.get(); }
   }
}
```

https://docs.oracle.com/javase/tutorial/essential/concurrency/forkjoin.html

# Thread pools

```
•17
```

```
class PrimeCountExecutor {
   private ExecutorService pool;
    ...
   public PrimeCountExecutor () {
      pool= new ForkJoinPool(1);
      ...
}
```



# Thread pools

```
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```

```
class PrimeCountExecutor {
   private ExecutorService pool;
   public PrimeCountExecutor () {
     pool= newFixedThreadPool(n);
tasks
                                                 CPU
                                                 CPU
```

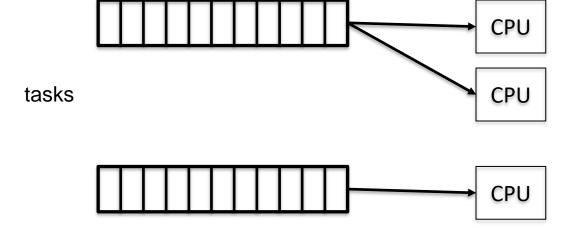
https://docs.oracle.com/javase/tutorial/essential/concurrency/pools.html

# Thread pools

```
. 19
```

```
class PrimeCountExecutor {
   private ExecutorService pool;
    ...
   public PrimeCountExecutor () {
      pool= newWorkStealingPool(x);
      ...
}
```

Quite a few more types of ExecutorService



## ForkJoinPool



https://docs.oracle.com/javase/8/docs/api/java/util/concurrent/ForkJoinPool.html

# Agenda



- Executors and Future
- Scalability, speed-up and loss (of scalablity) classification
   Example: QuickSort
- Lock striping
  - A case study with Hash maps

# Quicksort



1	2	43	78	19	54	33	21	64	52	17	53	
1	2	43	78	19	54	33	21	64	52	17	53	
1	2	43 <b>†</b>	78	19	54	33	21	64	52	<b>17</b>	53	
1	2	17 <b>†</b>	78	19	54	33	21	64	52	<b>43</b>	53	
1	2	17	78 <b>†</b>	19	54	33	21 <b>†</b>	64	52	43	53	
1	2	17	21	1 0	33	5 <i>1</i>	78	61	52	13	53	

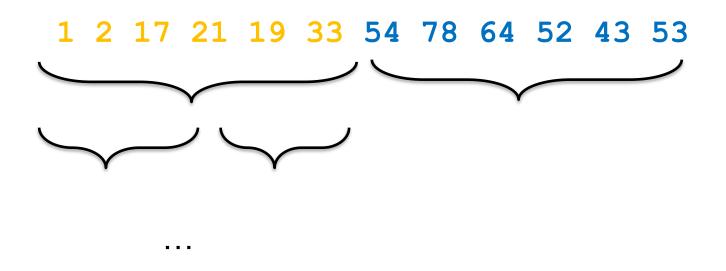
17 21 19 33 54 78 64 52 43 53

sorted independently

Two parts can be

# Distributing work to threads





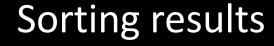
No further splitting when the sorting problem is smaller than a threshold (similarly to what we did for prime counting)

These tasks may differ in size!!

# Quicksort executor (pseudo code)



```
class QuicksortTask implements Runnable {
 Task p; // low and high boundaries
 ExecutorService pool;
  @Override public void run() { qsort(p, pool, ...); }
 public static void qsort(Task p, ExecutorService pool, ...) {
     //split task in two: Low and High
   if (Low.size>= threshold) pool.submit( new QuicksortTask( pLow, pool, ... ))
     else Quicksort(pLow); //sequential sort
   if (High.size>= threshold) pool.submit( new QuicksortTask( pHigh, pool, ... ))
     else Quicksort(pHigh);
     See complete code in: code-lecture/ .../QuickSortTask.java
```



new ForkJoinPool(4)



Executor  $8.5 \, \mathrm{s}$ 2 4.8 s

2.6 s

8 2.2 s 16 2.2 s

1/8: 3.9

11.2 s 2 6.4 s

Threads

4 3.8 s

8 3.2 s 16 3.5 s

1/8: 3.9

Sorting 100\_000 number

# Count prime results



E	kecutor
1	120.6

2 68.0 s

S

- 4 37.7 s
- 8 32.2 s
- 16 32.4 s
- 1/8: 3.7

Threads 1 126.7 s

2 82.4 s

4 47.7 s 8 38.2 s

16 37.2 s

1/8: 3.9

range 2.. 1\_000\_000

# Quicksort



1	2	43	78	19	54	33	21	64	52	17	53	
1	2	43	78	19	54	33	21	64	52	17	53	
1	2	43 †	78	19	54	33	21	64	52	<b>17</b>	53	
1	2	17 <b>†</b>	78	19	54	33	21	64	52	<b>43</b>	53	
1	2	17	78 <b>†</b>	19	54	33	21 <b>†</b>	64	52	43	53	
1	2	17	21	19	33	54	78	64	52	43	53	

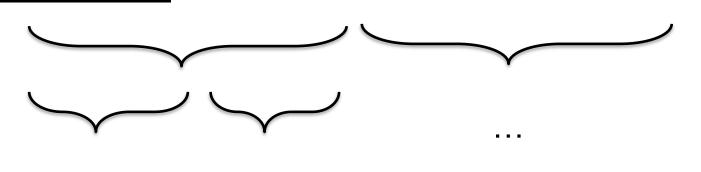
1 2 17 21 19 33 54 78 64 52 43 53

Two parts can be

sorted independently

# Termination



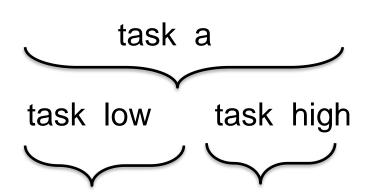






How do we know when all task are done?

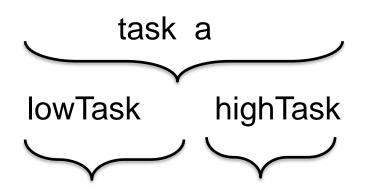
# Termination



When can task a finish?

## Termination





lowTask.get() returns when lowTask
has finished
highTask.get() returns when highTask
has finished



```
@Override
public void run() { // modified Quicksort using Executor tasks
  int a= low; int b= high;
  Future<?> lowTask= null; Future<?> highTask= null;
  if (a < b) { ... // split array in two independent part
    if ((j-a)>= threshold)
       lowTask= pool.submit(new QuickSortTask(arr, a, j, pool, threshold));
    else // all remaining work done without starting more tasks
       SearchAndSort.qsort(arr, a, j);
    if ((b-i)>= threshold)
       highTask= pool.submit(new QuickSortTask(arr, i, b, pool, threshold));
    else // all remaining work done without starting more tasks
       SearchAndSort.qsort(arr, i, b);
  //Waiting for longest running subtask to finish
  try {
    if (lowTask != null ) lowTaskF.get();
    if (highTask != null) highTaskF.get();
   } catch (InterruptedException | ExecutionException e) { e.printStackTrace(); }
```

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#### We use Mark8Setup to measure runtime

```
private static void runSize(ExecutorService pool, int pSize, int threshold, int n) {
  final int[] intArray= fillIntArray(pSize);
  Benchmark.Mark8Setup ("Quicksort Executor", String.format("%2d", n),
    new Benchmarkable() {
      public void setup() {
        shuffle(intArray);
      public double applyAsDouble(int i) {
        Future<?> done= pool.submit(new QuickSortTask(intArray, 0, pSize-1, pool, threshold));
        PoolFinish (done);
        //testSorted(intArray); //only needed while testing
        return 0.0;
```

#### Code in PoolSortingBenchmarkable.java



- Starvation loss
  - Minimize the time that the task pool is empty
- Separation loss (best threshold)
  - Find a good threshold to distribute workload evenly
- Saturation loss (locking common data structure)
  - Minimize high thread contention in the problem
- Braking loss
  - Stop all tasks as soon as the problem is solved

Møller-Nielsen, P and Staunstrup, J, Problem-heap. A paradigm for multiprocessor algorithms. *Parallel Computing*, 4:63-74, 1987

## Shut down



The ExecutorService can be shut down.

```
// Executor body
...
...
pool.shutdown();
```

The challenge is often when to shut down

After shutdown the pool cannot be reused, but you may assign it a new value (of type **ExecutorService**)



Both are used to specify the code of a thread.

- Runnable cannot return a result
  - Overrides run()
- <u>Callable</u> returns a result (via a Future)
  - Overrides call()

Could Callables use shared data as well?

As illustrated by the Quicksort and countPrimes examples, Runnables may use shared data (e.g., to deliver a result)

Futures are an example of message passing (coming weeks)

#### Submit vs Execute



Both are used to spawn a task

- pool.execute cannot return a result
   may complicate determining when to finish
- pool.submit returns a result (via a Future)

https://docs.oracle.com/javase/8/docs/api/java/util/concurrent/Executor.html

## dotProduct – Task class

Given two vectors x, y of equal size, their dot product equals

```
public class DotProductTask implements Callable<Integer> {
    final int pos;
    final int[] x, y;
    public DotProductTask(int[] x, int[] y, int pos) {
        this.x
                 = x;
        this.y = y;
        this.pos = pos;
    @Override
    public Integer call() {
        return x[pos] * y[pos];
```

```
\sum_{i} x_{i}y_{i}
We use futures to multiple each vector position
```

## dotProduct – Future execution



Given two vectors x, y of equal size, their dot product equals  $\sum_{i=1}^{n} x_i y_i$ 

```
List<DotProductTask> tasks = new ArrayList<DotProductTask>();
// Randomly initialize arrays ...
// Create the list of tasks (Futures) to execute
for (int i = 0; i < N; i++)
    tasks.add(new DotProductTask(x,y,i));
. . .
// Add all futures to the execution pool at once
List<Future<Integer>> futures = pool.invokeAll(tasks);
for (Future<Integer> f : futures) {
     result += f.get(); // Wait for each future to be executed
                        // and add partial result
                                                                 Code in FuturesDotProduct.java
pool.shutdown(); // We are sure to be done, so we shut down the pool
```

# What limits performance?

# -41

#### **CPU-bound**

- Eg. counting prime numbers
- To speed up, add more CPUs (cores) (exploitation)

#### Input/output-bound

- Eg. I/O or reading from network
- To speed up, use more threads/tasks (inherent)

#### **Synchronization-bound (Saturation loss)**

- Eg. Algorithm using shared data structure
- To speed up, improve shared data structure (Rest of this lecture)

# Agenda



- Executors and Future
- Scalability, speed-up and loss (of scalablity) classification Example: QuickSort
- Lock striping
  - A case study with Hash maps

#### Scalability of Java Collections



A *collection* is simply an object that groups multiple elements into a single unit

Package: java.util

Examples: ArrayList, HashMap, TreeSet, ...

https://docs.oracle.com/javase/tutorial/collections/intro/index.html

Methods: add, remove, size, contains, ...

Many of the classes have thread-safe/concurrent implementations

https://www.baeldung.com/java-synchronized-collections

### Example: synchronizedCollection

```
import java.util.*;
public class syncCollectionExample {
 public static void main(String[] args) {    new syncCollectionExample(); }
 public String getLast(ArrayList<String> 1) {
    int last= 1.size()-1;
    return l.get(last);
                                                                           Thread-safe?
 public static void delete(ArrayList<String> 1) {
    int last= 1.size()-1;
    1.remove(last);
 public syncCollectionExample() {
    ArrayList<String> a= new ArrayList<String>();
    a.add("A"); ...
    Collection<String> synColl = Collections.synchronizedCollection(a);
```

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# Thread-safety (from week 3)



It is very important to note that for a program p:

#### Making the synchronized ArrayList thread safe



```
import java.util.*;
public class syncCollectionExample {
  public static void main(String[] args) {
                                            new syncCollectionExample(); }
  public String getLast(ArrayList<String> 1) {
    synchronized(1) {
      int last= 1.size()-1;
      return l.get(last);
  public static void delete(ArrayList<String> 1) {
    synchronized(1) {
      int last= 1.size()-1;
      1.remove(last);
  public syncCollectionExample() {
```

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# What if the data structure is huge?



and used by many threads?

for example:

a bank

Facebook updates

•••

Would not work if everything is "synchronized"

What can we do?

Reduce locking !!

#### Example: A huge HashMap

Key value pairs: <k1, v1>, <k2, v2>, ...

```
class HashMap<K,V> {
    ... // data structure
    public V get(K k) { ... }
    public V put(K k, V v) { ... }
    public boolean containsKey(K k) { ... }
    public int size() { return cachedSize; }
    public V remove(K k) { ... }
    ...
}
```

How to make it thread-safe? (without making all the methods synchronized)

У	Value	
ter	20487612	
na	51251218	
na	34458318	
lger	89545010	
a	94959500	

Pet

An

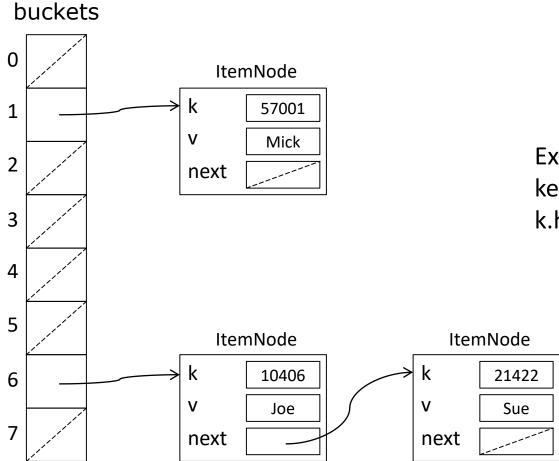
Ler

Ho

Lisa

### HashMap implementation

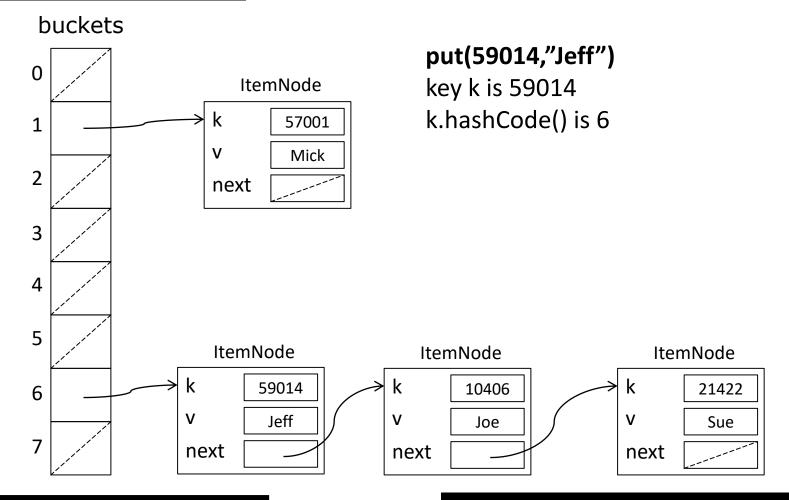




Example get(10406) key k is 10406 k.hashCode() is 6

# HaspMap put





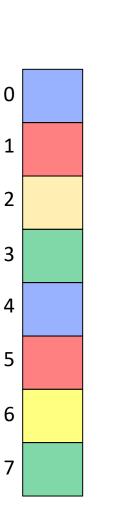
```
• 54
```

```
static class ItemNode<K,V> {
   private final K k;
   private V v;
   private ItemNode<K,V> next;
   public ItemNode(K k, V v, ItemNode<K,V> next) { ... }
}
```

```
class SynchronizedMap<K,V> {
  private ItemNode<K,V>[] buckets; // guarded by this
  private int cachedSize; // guarded by this
  public synchronized V get(K k) { ... }
  public synchronized boolean containsKey(K k) { ... }
  public synchronized int size() { return cachedSize; }
  public synchronized V put(K k, V v) { ... }
  public synchronized V remove(K k) { ... }
}
```

# Improving scalability – Lock striping

- Guarding the table with a single lock works
- -... but does not scale well (actually **very** badly)
- Idea: Each bucket could have its own lock
- In practice
- -use fewer, to illustrate we use 4, locks
- -guard every 4<sup>th</sup> bucket with the same lock
- -locks[0] guards bucket 0, 4, 8, ... (stripe 0)
- -locks[1] guards bucket 1, 5, 9, ... (stripe 1) et
- -With high probability
- -two operations will work on different stripes
- hence will take different locks
- Less lock contention, better scalability



3

4

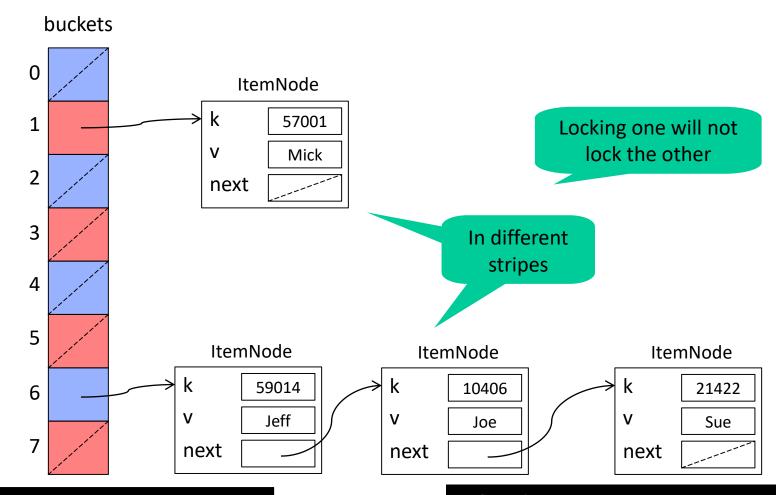
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#### Bucket idea



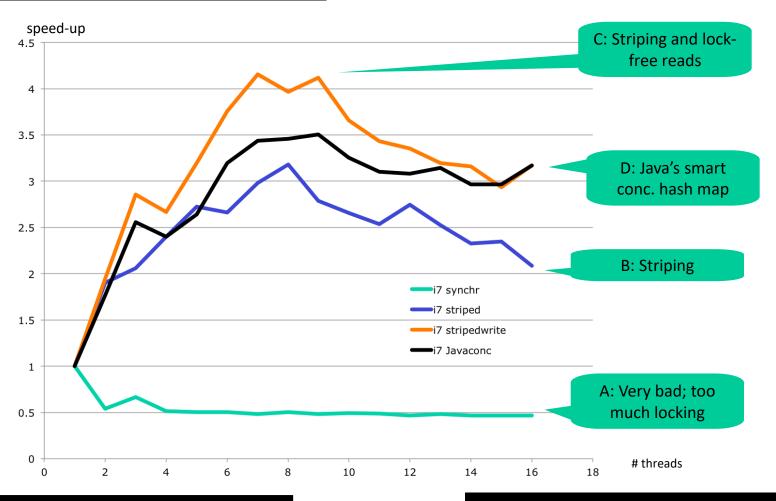


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# Reducing locking





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### Ultimate scalability



A web-shop, Facebook, ...

We must give up thread safety,

but still maintain some sort of consistency

Week 13