

# PiCore: A Rely-guarantee Framework for Event-based Systems

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	theory <i>Heap</i>	
	imports <i>Main</i>	
	begin	

## 0.1 References

```

definition ref = (UNIV::nat set)

typedef ref = ref by (simp add: ref-def)

code-datatype Abs-ref

lemma finite-nat-ex-max:
  assumes fin: finite (N::nat set)
  shows  $\exists m. \forall n \in N. n < m$ 
using fin
proof (induct)
  case empty
  show ?case by auto
next
  case (insert k N)
  have  $\exists m. \forall n \in N. n < m$  by fact

```

```

then obtain  $m$  where  $m\text{-max}$ :  $\forall n \in N. n < m..$ 
show  $\exists m. \forall n \in \text{insert } k \ N. n < m$ 
proof (rule  $\text{exI}$  [where  $x = \text{Suc } (\text{max } k \ m)$ ])
qed (insert  $m\text{-max}$ , auto simp add:  $\text{max-def}$ )
qed

lemma  $\text{infinite-nat}$ :  $\neg \text{finite } (\text{UNIV}::\text{nat set})$ 
proof
  assume  $\text{fin}$ :  $\text{finite } (\text{UNIV}::\text{nat set})$ 
  then obtain  $m::\text{nat}$  where  $\forall n \in \text{UNIV}. n < m$ 
    by (rule  $\text{finite-nat-ex-max}$  [ $\text{elim-format}$ ]) auto
  moreover have  $m \in \text{UNIV}..$ 
  ultimately show  $\text{False}$  by  $\text{blast}$ 
qed

lemma  $\text{infinite-ref}$  [ $\text{simp}, \text{intro}$ ]:  $\neg \text{finite } (\text{UNIV}::\text{ref set})$ 
proof
  assume  $\text{finite } (\text{UNIV}::\text{ref set})$ 
  hence  $\text{finite } (\text{range } \text{Rep-ref})$ 
    by  $\text{simp}$ 
  moreover
  have  $\text{range } \text{Rep-ref} = \text{ref}$ 
  proof
    show  $\text{range } \text{Rep-ref} \subseteq \text{ref}$ 
      by (simp add:  $\text{ref-def}$ )
    next
    show  $\text{ref} \subseteq \text{range } \text{Rep-ref}$ 
    proof
      fix  $x$ 
      assume  $x$ :  $x \in \text{ref}$ 
      show  $x \in \text{range } \text{Rep-ref}$ 
        by (rule  $\text{Rep-ref-induct}$ ) (auto simp add:  $\text{ref-def}$ )
      qed
    qed
  ultimately have  $\text{finite } \text{ref}$ 
    by  $\text{simp}$ 
  thus  $\text{False}$ 
    by (simp add:  $\text{ref-def infinite-nat}$ )
qed

consts  $\text{Null} :: \text{ref}$ 

definition  $\text{new} :: \text{ref set} \Rightarrow \text{ref}$  where
   $\text{new } A = (\text{SOME } a. a \notin \{\text{Null}\} \cup A)$ 

```

Constant  $\text{Null}$  can be defined later on. Conceptually  $\text{Null}$  and  $\text{new}$  are *fixes* of a locale with  $\text{finite } A \implies \text{new } A \notin A \cup \{\text{Null}\}$ . But since definitions relative to a locale do not yet work in Isabelle2005 we use this workaround to avoid lots of parameters in definitions.

```

lemma new-notin [simp,intro]:
  finite A  $\implies$  new (A)  $\notin$  A
  apply (unfold new-def)
  apply (rule someI2-ex)
  apply (fastforce intro: ex-new-if-finite)
  apply simp
  done

lemma new-not-Null [simp,intro]:
  finite A  $\implies$  new (A)  $\neq$  Null
  apply (unfold new-def)
  apply (rule someI2-ex)
  apply (fastforce intro: ex-new-if-finite)
  apply simp
  done

end
theory aux-lemma
imports Main
begin

lemma mod-div-self: (a::nat) mod b = 0  $\implies$  (a div b) * b = a
by auto

lemma mod-div-mult: (a::nat) mod b = 0  $\implies$  a div b  $\leq$  (c - 1)  $\implies$  a  $\leq$  c * b
- b
  apply(subgoal-tac a  $\leq$  (c - 1) * b)
  apply (simp add: left-diff-distrib')
  by fastforce

lemma mod0-div-self: (a::nat) mod b = 0  $\implies$  b * (a div b) = a by auto

lemma m-mod-div: n mod x = 0  $\implies$  (m::nat) * n div x = m * (n div x)
by auto

lemma pow-mod-0: x  $\geq$  y  $\implies$  (m::nat)  $^$  x mod m  $^$  y = 0
by (simp add: le-imp-power-dvd)

lemma ge-pow-mod-0: (x::nat) > y  $\implies$  4 * n * (4::nat)  $^$  x mod 4  $^$  y = 0
by (metis less-imp-le-nat mod-mod-trivial mod-mult-right-eq mult-0-right pow-mod-0)

lemma div2-eq-minus: x  $\neq$  0  $\wedge$  m  $\geq$  n  $\implies$  (x::nat)  $^$  m div x  $^$  n = x  $^$  (m - n)
by (metis add-diff-cancel-left' div-mult-self1-is-m grOI le-Suc-ex power-add power-not-zero)

lemma pow-lt-mod0: (n::nat) > 0  $\wedge$  (x::nat) > y  $\implies$  (n  $^$  x div n  $^$  y) mod n =
0
by (simp add: div2-eq-minus)

```

**lemma** *mod-div-gt*:

$(m::nat) < n \implies n \bmod x = 0 \implies m \operatorname{div} x < n \operatorname{div} x$   
**by** (*simp add: less-mult-imp-div-less mod-div-self*)

**lemma** *div2-eq-divmul*:  $(a::nat) \operatorname{div} b \operatorname{div} c = a \operatorname{div} (b * c)$   
**by** (*simp add: Divides.div-mult2-eq*)

**lemma** *addr-in-div*:

$(addr::nat) \in \{j2 * M ..< (Suc j2) * M\} \implies addr \operatorname{div} M = j2$   
**by** (*simp add: div-nat-eqI mult.commute*)

**lemma** *divn-mult-n*:  $x > 0 \implies (n::nat) = m \operatorname{div} x * x \implies (\text{if } m \bmod x = 0 \text{ then } m = n \text{ else } n < m \wedge m < n + x \wedge n \bmod x = 0)$

**apply** *auto*

**apply** (*metis div-mult-mod-eq less-add-same-cancel1*)

**by** (*metis add-le-cancel-left div-mult-mod-eq mod-less-divisor not-less*)

**lemma** *mod-minus-0*:

$(m::nat) \leq n \wedge 0 < m \implies a * (4::nat) ^ n \bmod 4 ^ (n - m) = 0$

**by** (*metis diff-le-self mod-mult-right-eq mod-mult-self2-is-0 mult-0 mult-0-right pow-mod-0*)

**lemma** *mod-minus-div-4*:

$(m::nat) \leq n \wedge 0 < m \implies a * (4::nat) ^ n \operatorname{div} 4 ^ (n - m) \bmod 4 = 0$

**by** (*metis add.left-neutral add-lessD1 diff-less m-mod-div mod-0 mod-mult-right-eq*

*mult-0-right nat-less-le pow-lt-mod0 pow-mod-0 zero-less-numeral*)

**lemma** *modn0-xy-n*:  $(n::nat) > 0 \implies x \bmod n = 0 \implies y \bmod n = 0 \implies x < y \implies x + n \leq y$

**by** (*metis Nat.le-diff-conv2 add.commute add.left-neutral add-diff-cancel-left' le-less less-imp-add-positive mod-add-left-eq mod-less not-less*)

**lemma** *divn-multn-addn-le*:  $(n::nat) > 0 \implies y \bmod n = 0 \implies x < y \implies x \operatorname{div} n * n + n \leq y$

**using** *divn-mult-n[of n x div n \* n x] modn0-xy-n*

**apply**(*case-tac x mod n = 0*)

**apply**(*rule subst[where s=x and t=x div n \* n]*) **apply** *metis*

**by** *auto*

**lemma** *div-in-suc*:  $y > 0 \implies n = (x::nat) \operatorname{div} y \implies x \in \{n * y ..< Suc n * y\}$

**by** (*simp add: dividend-less-div-times*)

**lemma** *int1-eq*:  $P \cap \{V\} \neq \{\} \implies P \cap \{V\} = \{V\}$  **by** *auto*

**lemma** *int1-belong*:  $P \cap \{V\} = \{V\} \implies V \in P$  **by** *auto*

**lemma** *two-int-one*:  $P \cap \{V\} \cap \{Va\} \neq \{\} \implies V = Va \wedge \{V\} = P \cap \{V\} \cap$

```

{ Va } by auto

end
theory List-aux
imports aux-lemma
begin

primrec list-updates :: 'a list  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  'a  $\Rightarrow$  'a list where
  list-updates [] i1 i2 v = [] |
  list-updates (x#xs) i1 i2 v =
    (case i1 of 0  $\Rightarrow$  (if i2 > 0 then v # list-updates xs 0 (i2 - 1) v else (v#xs))
  |
    Suc j  $\Rightarrow$  (if i2 > j then x # list-updates xs j (i2 - 1) v else (x#xs)) )

value list-updates [1::nat,2,3,4,5] 9 0 6

lemma length-list-update2 [simp]: length (list-updates l i1 i2 v) = length l
  apply(induct l arbitrary: i1 i2 v)
  apply simp
  apply(case-tac i1)
  apply(case-tac i2) apply simp+
done

lemma list-updates-eq [simp]:  $\llbracket i1 \leq i; i \leq i2; i2 < \text{length } l \rrbracket \Longrightarrow (\text{list-updates } l \ i1 \ i2 \ v)!i = v$ 
  apply(induct l arbitrary: i i1 i2 v)
  apply simp
  apply(case-tac i1) apply auto
  apply(case-tac i2) apply simp
  by (metis (no-types, lifting) One-nat-def Suc-less-SucD diff-Suc-1
    le-SucE le-zero-eq not-less-eq-eq nth-Cons' zero-induct)

lemma list-updates-neq [simp]:  $i < i1 \vee i > i2 \Longrightarrow (\text{list-updates } l \ i1 \ i2 \ v)!i = !!i$ 
  apply(induct l arbitrary: i i1 i2 v)
  apply simp
  apply(case-tac i1) apply simp
  apply(case-tac i2) apply simp apply(case-tac i) apply simp+
done

lemma list-updates-beyond[simp]:  $i1 \geq \text{length } l \Longrightarrow (\text{list-updates } l \ i1 \ i2 \ v) = l$ 
  apply(induct l arbitrary: i1 i2 v)
  apply simp apply(case-tac i1) by auto

lemma list-updates-beyond2[simp]:  $i2 < i1 \Longrightarrow (\text{list-updates } l \ i1 \ i2 \ v) = l$ 
  apply(induct l arbitrary: i1 i2 v)
  apply simp apply(case-tac i1) by auto

lemma list-updates-nonempty[simp]:  $(\text{list-updates } l \ i1 \ i2 \ v) = [] \longleftrightarrow l = []$ 

```

```

by (metis length-greater-0-conv length-list-update2)

lemma list-updates-same-conv:
  i1 < length l ∧ i2 < length l ⇒ ((list-updates l i1 i2 v) = l) = (∀ i. i ≥ i1 ∧ i
≤ i2 ⇒ l ! i = v)
  apply (induct l arbitrary: i1 i2 v)
  apply simp
  apply (case-tac i1 ≤ i2) apply (rule iffI)
  apply (metis list-updates-eq)
  apply (smt length-list-update2 list-updates-eq list-updates-neq not-le-imp-less
nth-equalityI)
  by (metis (mono-tags, lifting) list-updates-beyond2 list-updates-eq not-le-imp-less)

lemma list-updates-append1:
  i2 < length l ⇒ list-updates (l @ t) i1 i2 v = list-updates l i1 i2 v @ t
  apply (induct l arbitrary: i1 i2 v)
  apply simp
  apply (case-tac i1 ≤ i2)
  apply (case-tac i1) apply simp
  apply (case-tac i2) apply simp apply auto[1]
  by (metis list-updates-beyond2 not-less)

primrec list-updates-fstn :: 'a list ⇒ nat ⇒ 'a ⇒ 'a list where
  list-updates-fstn [] n v = [] |
  list-updates-fstn (x#xs) n v =
    (case n of 0 ⇒ x#xs | Suc m ⇒ v#list-updates-fstn xs m v)

primrec list-updates-n :: 'a list ⇒ nat ⇒ nat ⇒ 'a ⇒ 'a list where
  list-updates-n [] i n v = [] |
  list-updates-n (x#xs) i n v =
    (case i of 0 ⇒ list-updates-fstn (x#xs) n v | Suc j ⇒ x#list-updates-n xs j n
v)

value list-updates-n [1::nat,2,3,4,5] 0 9 6

lemma length-list-update-fstn [simp]: length (list-updates-fstn l n v) = length l
  apply (induct l arbitrary: n v)
  apply simp
  apply (case-tac n) apply simp+
done

lemma length-list-update-n [simp]: length (list-updates-n l i n v) = length l
  apply (induct l arbitrary: i n v)
  apply simp
  apply (case-tac i)
  apply (case-tac n) apply simp+
done

```

```

lemma list-updates-fstn-eq [simp]:  $\llbracket i < \text{length } l; i < n \rrbracket \implies (\text{list-updates-fstn } l \ n \ v)!i = v$ 
  apply (induct l arbitrary: i n v) apply simp
  apply (case-tac i)
  apply (case-tac n) apply simp+
  apply (case-tac n) apply simp+
done

lemma list-updates-n-eq [simp]:  $\llbracket i \leq j; j < \text{length } l; j < i + n \rrbracket \implies (\text{list-updates-n } l \ i \ n \ v)!j = v$ 
  apply (induct l arbitrary: i j n v) apply simp
  apply (case-tac i) apply auto
  apply (case-tac n) apply auto
  using less-Suc-eq-0-disj by auto

lemma list-updates-fst0 [simp]:  $\text{list-updates-fstn } l \ 0 \ v = l$ 
  apply (induct l arbitrary: v) by simp+

lemma list-updates-0 [simp]:  $\text{list-updates-n } l \ i \ 0 \ v = l$ 
  apply (induct l arbitrary: i v) apply simp apply (case-tac i) apply simp+
done

lemma list-updates-fstn-neq [simp]:  $j \geq n \implies (\text{list-updates-fstn } l \ n \ v)!j = !j$ 
  apply (induct l arbitrary: j n v) apply simp
  apply (case-tac n) apply simp+
done

lemma list-updates-n-neq [simp]:  $j < i \vee j \geq i + n \implies (\text{list-updates-n } l \ i \ n \ v)!j = !j$ 
  apply (induct l arbitrary: i j n v) apply simp
  apply (case-tac i) apply (case-tac n) apply simp+
  apply (case-tac n) apply simp apply (case-tac j) apply simp apply auto
done

lemma list-updates-n-beyond [simp]:  $i \geq \text{length } l \implies (\text{list-updates-n } l \ i \ n \ v) = l$ 
  apply (induct l arbitrary: i n v)
  apply simp apply (case-tac i) by auto

lemma lst-udptn-set-eq:  $n > 0 \implies \text{list-updates-n } (\text{lst}[jj := \text{TAG}]) (jj \text{ div } n * n) \ n \ \text{TAG1} =$ 
   $\text{list-updates-n } \text{lst } (jj \text{ div } n * n) \ n \ \text{TAG1}$ 
apply (rule nth-equalityI) apply simp
apply clarify
apply (case-tac i = jj)
apply (subgoal-tac i  $\geq jj \text{ div } n * n$ ) prefer 2 apply (metis divn-mult-n less-or-eq-imp-le)
apply (subgoal-tac i  $< jj \text{ div } n * n + n$ ) prefer 2
apply (metis (no-types) add.commute dividend-less-div-times)
apply simp

```



**by** (metis length-list-update length-list-update-n list-updates-n-eq list-updates-n-neq  
not-less nth-list-update-neq)

**thm** list-updates-n.simps

**lemma** list-updates-n-simps2: list-updates-n (a#lst) (Suc ii) m v = a # list-updates-n  
lst ii m v  
**by** fastforce

**lemma** list-updates-n-simps2': ii > 0  $\implies$  list-updates-n (a#lst) ii m v = a #  
list-updates-n lst (ii - 1) m v  
**using** list-updates-n-simps2[of a lst ii - 1 m v] **by** force

**lemma** lst-updt1-eq-upd: list-updates-n lst ii 1 v = lst[ii := v]  
**apply**(induct lst arbitrary: ii) **apply** simp  
**apply**(case-tac ii = 0) **apply** simp  
**using** list-updates-n-simps2'  
**by** (metis One-nat-def Suc-pred list-update-code(3) neq0-conv)

**lemma** list-neq-udpt-neq:  
 $\forall i < \text{length } l. l ! i \neq P \implies$   
 $l' = \text{list-updates-n } l \ s \ n \ Q \implies$   
 $P \neq Q \implies$   
 $\forall i < \text{length } l'. l' ! i \neq P$   
**apply**(induct l' arbitrary:l) **apply** simp  
**by** (metis le-neq-implies-less length-list-update-n list-updates-n-eq list-updates-n-neq  
nat-le-linear)

**lemma** lst-updts-eq-updts-updt:

$1 \leq ii \implies$   
list-updates-n lst st (ii - 1) TAG [st + ii - 1 := TAG] =  
list-updates-n lst st ii TAG  
**apply**(rule nth-equalityI)  
**apply** simp  
**apply** clarsimp **apply**(rename-tac ia)  
**apply**(case-tac ia < st) **using** list-updates-n-neq **apply** simp  
**apply**(case-tac ia  $\geq$  st + ii) **using** list-updates-n-neq **apply** simp  
**apply**(case-tac ia < st + ii - 1) **using** list-updates-n-eq **apply** simp  
**apply**(subgoal-tac ia = st + ii - 1) **prefer** 2  
**apply** force  
**apply**(subgoal-tac length lst = length (list-updates-n lst st ii TAG))  
**prefer** 2 **apply** simp  
**apply**(subgoal-tac length lst = length (list-updates-n lst st (ii - 1) TAG))  
**prefer** 2 **using** length-list-update-n **apply** metis  
**apply**(case-tac ia  $\geq$  length lst) **apply** linarith  
**apply**(subgoal-tac list-updates-n lst st (ii - 1) TAG [st + ii - 1 := TAG] !  
ia = TAG) **prefer** 2

```

    apply (metis nth-list-update-eq)
  apply(subgoal-tac list-updates-n lst st ii TAG ! ia = TAG) prefer 2
  apply (meson list-updates-n-eq not-less)
using One-nat-def by presburger

primrec removes :: 'a list  $\Rightarrow$  'a list  $\Rightarrow$  'a list
where removes [] l = l |
  removes (x#xs) l = removes xs (remove1 x l)

lemma removes-distinct [simp]: distinct l  $\implies$  distinct (removes rs l)
  apply(induct rs arbitrary:l) by auto

lemma removes-length [simp]:  $\llbracket \text{set } rs \subseteq \text{set } l; \text{distinct } l; \text{distinct } rs \rrbracket$ 
   $\implies \text{length } rs + \text{length } (\text{removes } rs \ l) = \text{length } l$ 
  apply(induct rs arbitrary:l)
  apply simp apply auto
  by (metis (no-types, lifting) One-nat-def Suc-pred distinct-remove1
    in-set-remove1 length-pos-if-in-set length-remove1 subset-eq)

lemma removes-empty [simp]: removes rs [] = []
  apply(induct rs) by simp+

lemma removes-subst1 [simp]: set (removes rs l)  $\subseteq$  set l
  apply(induct rs arbitrary: l) apply simp apply simp
  apply(subgoal-tac set (remove1 a l)  $\subseteq$  set l) apply auto[1]
  by (simp add: set-remove1-subset)

lemma removes-subst2 [simp]: distinct l  $\implies$  set (removes (a#rs) l)  $\subseteq$  set (removes
rs l)
  apply simp
  apply(induct rs arbitrary: l a)
  apply auto by (metis (full-types) distinct-remove1 remove1-commute set-mp)

lemma removes-nin [simp]:  $\llbracket x \in \text{set } rs; \text{distinct } l \rrbracket \implies x \notin \text{set } (\text{removes } rs \ l)$ 
  apply(induct rs arbitrary:l x)
  apply simp
  apply simp apply auto
  by (metis DiffE contra-subsetD removes-subst1 set-remove1-eq singletonI)

lemma rmvs-empty:  $a \in \text{set } es \implies \text{removes } es \ [a] = []$ 
apply(induct es) apply simp apply auto
done

lemma rmvs-unchg:  $a \notin \text{set } es \implies \text{removes } es \ [a] = [a]$ 
apply(induct es) apply simp apply auto
done

```

```

lemma rmvs-onemore-same:
  distinct lst  $\implies e \notin \text{set } lst \implies \text{removes } (es@[e]) \text{ } lst = \text{removes } es \text{ } lst$ 
apply (induct es arbitrary:lst)
apply (simp add: remove1-idem)
apply auto
done

```

```

lemma rmvs-rev: removes (es@[e]) lst = remove1 e (removes es lst)
apply (induct es arbitrary:lst) apply simp apply auto
done

```

```

definition inserts xs l  $\equiv l @ xs$ 

```

```

lemma inserts-set-un: set (inserts xs l) = set xs  $\cup$  set l
  by (simp add: inserts-def sup-commute)

```

```

lemma inserts-emp1: set (inserts xs []) = set xs
  using inserts-set-un[of xs []] by auto

```

```

lemma inserts-emp2: set (inserts [] l) = set l
  using inserts-set-un[of [] l] by auto

```

```

lemma list-updt-samelen: length l = length (l[jj := a]) by simp

```

```

lemma list-nhd-in-tl-set: el  $\in$  set l  $\implies el \neq \text{hd } l \implies el \in \text{set } (tl \text{ } l)$ 
  by (metis empty-iff empty-set list.exhaust-sel set-ConsD)

```

```

lemma dist-hd-nin-tl: distinct l  $\implies a \in \text{set } (tl \text{ } l) \implies a \neq \text{hd } l$ 
  by (metis distinct.simps(2) equals0D list.collapse set-empty tl-Nil)

```

```

end

```

```

theory mem-spec
imports Main Heap PiCore-SIMP.picore-SIMP-lemma List-aux
begin

```

## 1 data types and state

```

typedecl Thread

```

```

typedef mempool-ref = ref by (simp add: ref-def)

```

we define memory address as nat

```

type-synonym mem-ref = nat

```

```

abbreviation NULL  $\equiv 0 :: \text{nat}$ 

```

we have a thread scheduler, thread has 3 types. BLOCKED means a thread is waiting for memory and is in wait queue

**datatype** *Thread-State-Type* = *READY* | *RUNNING* | *BLOCKED*

a memory block: a ref to a memor pool, a level index and a block index in this level, a start address “data”. max number of levels is n\_level of a memory pool. So @level should be j n\_levels. The number of blocks at level 0 is n\_max. the max number of blocks at level n is  $n\_max * 4^n$ . the block index should less then this number.

**record** *Mem-block* = *pool* :: *mempool-ref*  
                   *level* :: *nat*  
                   *block* :: *nat*  
                   *data* :: *mem-ref*

BlockState defines the bit info in bitmap. We uses different types, while not 0 or 1 in this design. Then the blockstate could be implemented as 0 or 1, with additional information.

basic states of memory block are ALLOCATED, FREE, DIVIDED and NOEXIST. The levels of bitmap is actually a quad-tree of BlockState. ALLOCATED: the block is allocated to a thread FREE: the block is free DIVIDED: the block is divided, which means is was splited to 4 subblocks NOEXIST: the block is not exist

ALLOCATED and FREE blocks are the leaf blocks of the quad-tree. DIVIDED blocks are inner nodes of the quad-tree. Otherwise is NOEXIST.

we also introduce FREEING and ALLOCATING state to avoid a case that a FREEING block may be allocated by other threads and a ALLOCATING block may be freed by other threads. In OS implementation, the allocating/freeing block is an inner block of alloc/free services, and other threads will not manipulate them. they are used to indicate state of the block which are going to be merged during freeing a block, and the block which is going to be split during allocating a block.

we may remove FREEING/ALLOCATING state later by revising alloc and free syscalls to avoid allocate or free blocks in freeing\_node and allocating\_node.

**datatype** *BlockState* = *ALLOCATED* | *FREE* | *DIVIDED* | *NOEXIST* | *FREEING* | *ALLOCATING*

data stucture at each level, a bitmap and a free block list

**record** *Mem-pool-lvl* =  
                   *bits* :: *BlockState list*  
                   *free-list* :: *mem-ref list*

a memory pool is actually a forest of @n\_max numbers of blocks with size of @max.sz. A block may be split to 4 sub-blocks and so on, at most for

@n\_levels times. Thus, each block may be split as a quad-tree. a memory pool maintains a big memory block, where @buf is the start address of the memory block. The size of a memory pool is @n\_max \* @max\_sz. @max\_sz has a constraint. a small block at last level (level index is @n\_levels - 1) should be aligned by 4 bits, i.e. the size of block at last level should be 4\*n (n > 0). Here, we don't demand  $4^n$ , which is a special case of 4\*n. Thus, @max\_sz should be  $4 * n * 4^n\_levels$ .

@levels maintain the information at each level including a bitmap and a free block list. @wait\_q is a list of threads, which is blocked on this memory pool.

```
record Mem-pool = buf :: mem-ref
                max-sz :: nat
                n-max :: nat
                n-levels :: nat

                levels :: Mem-pool-lvl list
                wait-q :: Thread list
```

The state of memory management consists of thread state, memory pools, and local variables of each thread. In moncore OSs, there is only one currently executing thread @cur, where None means the scheduler has not choose a thread. @tick save a time for the system. @mem\_pools maintains the refs of all memory pools. @mem\_pool\_info shows the detailed information of each memory pool by its ref. we assume that all memory pools are shared by all threads. This is the most relaxed case. The case that some memory pool is only shared by a set of thread is just a special case. Other fields are local vars of each thread used in alloc/free syscalls.

for each thread, we use freeing\_node to maintain the freeing node in free syscall. when free a block, we set it to FREEING, and check if its other 3 partner blocks are also free. If so, we set the 4 blocks to NOEXIST and set their parent block to FREEING, and so on. until that other 3 partner blocks are not all free, then set the FREEING block to FREE. This design avoids the FREEING node is allocated by other threads.

we use allocating\_node to maintain the allocating node in alloc syscall. when alloc a block, we find a free block at the nearest upper level, and set it to ALLOCATING. if size of the block is too big, we split it into 4 child blocks. We set the first child block to ALLOCATING and other 3 blocks to FREE, and so on. until that the size of block is suitable, then set the ALLOCATING block to ALLOCATED. This design avoids the ALLOCATING node is freed by other threads.

```
record State =

  cur :: Thread option
  tick :: nat
```

$thd-state :: Thread \Rightarrow Thread-State-Type$

$mem-pools :: mempool-ref\ set$

$mem-pool-info :: mempool-ref \Rightarrow Mem-pool$

$i :: Thread \Rightarrow nat$   
 $j :: Thread \Rightarrow nat$   
 $ret :: Thread \Rightarrow int$   
 $endt :: Thread \Rightarrow nat$   
 $rf :: Thread \Rightarrow bool$   
 $tmout :: Thread \Rightarrow int$   
 $lsizes :: Thread \Rightarrow nat\ list$   
 $alloc-l :: Thread \Rightarrow int$   
 $free-l :: Thread \Rightarrow int$   
 $from-l :: Thread \Rightarrow int$   
 $blk :: Thread \Rightarrow mem-ref$   
 $nodev :: Thread \Rightarrow mem-ref$   
 $bn :: Thread \Rightarrow nat$   
 $lbn :: Thread \Rightarrow nat$   
 $lsz :: Thread \Rightarrow nat$   
 $block2 :: Thread \Rightarrow mem-ref$   
 $free-block-r :: Thread \Rightarrow bool$   
 $alloc-lsize-r :: Thread \Rightarrow bool$   
 $lvl :: Thread \Rightarrow nat$   
 $bb :: Thread \Rightarrow nat$   
 $block-pt :: Thread \Rightarrow mem-ref$   
 $th :: Thread \Rightarrow Thread$   
 $need-resched :: Thread \Rightarrow bool$   
 $mempoolalloc-ret :: Thread \Rightarrow Mem-block\ option$

$freeing-node :: Thread \Rightarrow Mem-block\ option$

$allocating-node :: Thread \Rightarrow Mem-block\ option$

## 2 specification of events

### 2.1 data types

Since Zephyr uses fine-grained locks for shared memory pools, interleaving among scheduling, syscalls (alloc, free), and clock tick are allowed. Thus, we use 3 event systems to model scheduling, syscalls from threads, and clock tick. Then the whole system is the parallel composition of the three event systems. Actually, we have 1 scheduler, 1 timer, and n threads.

**datatype**  $Core = \mathcal{S} \mid \mathcal{T}\ Thread \mid Timer$

labels for different events

**datatype**  $EL = \text{ScheduleE} \mid \text{TickE} \mid \text{Mem-pool-allocE} \mid \text{Mem-pool-freeE} \mid \text{Mem-pool-defineE}$

data types for event parameters

**datatype**  $Parameter = \text{Thread} \mid \text{Thread} \mid \text{MPRef mempool-ref} \mid \text{MRef mem-ref} \mid \text{Block Mem-block} \mid \text{Natural nat} \mid \text{Integer int}$

**type-synonym**  $\text{EventLabel} = EL \times (Parameter \text{ list} \times Core)$

**definition**  $\text{get-evt-label} :: EL \Rightarrow Parameter \text{ list} \Rightarrow Core \Rightarrow \text{EventLabel} \text{ (-} \Rightarrow \text{-} [30,30,30] \text{ } 20)$

**where**  $\text{get-evt-label } el \text{ ps } k \equiv (el, (ps, k))$

define the waiting mode for alloc. FOREVER means that if allocating fails, the thread will wait forever until allocating succeed. NOWAIT means that if allocating fails, alloc syscall return error immediately. otherwise n  $\neq$  0, means the thread will wait for a timeout n.

**abbreviation**  $\text{FOREVER} \equiv (-1)::int$

**abbreviation**  $\text{NOWAIT} \equiv 0::int$

return CODE for alloc and free syscalls. free syscall always succeed, so it returns OK. alloc syscall may succeed (OK), timeout (ETIMEOUT), fails(ENOMEM), fails due to request too large size (ESIZEERR).

EAGAIN is an inner flag of alloc syscall. After it finds an available block for request, the block may be allocated immediately by other threads. In such a case, alloc will provide EAGAIN and try to allocate again.

We introduce ESIZEERR for Zephyr to avoid a dead loop. We introduce ETIMEOUT for Zephyr for robustness.

**abbreviation**  $\text{EAGAIN} \equiv (-2)::int$

**abbreviation**  $\text{ENOMEM} \equiv (-3)::int$

**abbreviation**  $\text{ESIZEERR} \equiv (-4)::int$

**abbreviation**  $\text{OK} \equiv 0 :: int$

**abbreviation**  $\text{ETIMEOUT} \equiv (-1) :: int$

due to fine-grained lock used by Zephyr, we use a command for each atomic statement in free/alloc syscalls. the statements of syscalls from a thread  $t$  can only be executed when  $t$  is the currently executing thread by the scheduler. We use the AWAIT statement to represent this semantics.

**definition**  $\text{stm} :: Thread \Rightarrow State \text{ com} \Rightarrow State \text{ com} \text{ (-} \blacktriangleright \text{-} [0,0] \text{ } 21)$

**where**  $\text{stm } t \text{ p} = \text{AWAIT 'cur} = \text{Some } t \text{ THEN } p \text{ END}$

## 2.2 aux definitions for events

**definition**  $\text{ALIGN}_4 :: nat \Rightarrow nat$

**where**  $\text{ALIGN}_4 \text{ } n \equiv ((n + 3) \text{ div } 4) * 4$

```

lemma align40:  $n \bmod 4 = 0 \implies \text{ALIGN}_4\ n = n$ 
  unfolding ALIGN4-def by auto

lemma align41:  $n \bmod 4 = 1 \implies \text{ALIGN}_4\ n = n + 3$ 
  unfolding ALIGN4-def
proof -
  assume  $n \bmod 4 = 1$ 
  then have  $(n + 3) \bmod 4 = 0$ 
    by presburger
  then show  $(n + 3) \div 4 * 4 = n + 3$ 
    by fastforce
qed

lemma align42:  $n \bmod 4 = 2 \implies \text{ALIGN}_4\ n = n + 2$ 
  unfolding ALIGN4-def
proof -
  assume  $n \bmod 4 = 2$ 
  then have  $(n + 2) \bmod 4 = 0$ 
    using mod-add-left-eq by presburger
  then show  $(n + 2) \div 4 * 4 = n + 2$ 
    by fastforce
qed

lemma align43:  $n \bmod 4 = 3 \implies \text{ALIGN}_4\ n = n + 1$ 
  unfolding ALIGN4-def
proof -
  assume  $n \bmod 4 = 3$ 
  then have  $(n + 1) \bmod 4 = 0$ 
    using mod-add-left-eq by presburger
  then show  $(n + 1) \div 4 * 4 = n + 1$ 
    by fastforce
qed

lemma align-mod0:  $\text{ALIGN}_4\ n \bmod 4 = 0$ 
  unfolding ALIGN4-def by simp

lemma align4-gt:  $\text{ALIGN}_4\ n \geq n \wedge \text{ALIGN}_4\ n \leq n + 3$ 
  apply (case-tac  $n \bmod 4 = 0$ )
  using align40 apply simp
  apply (case-tac  $n \bmod 4 = 1$ )
  using align41 apply simp
  apply (case-tac  $n \bmod 4 = 2$ )
  using align42 apply simp
  apply (case-tac  $n \bmod 4 = 3$ )
  using align43 apply simp
  by auto

lemma align2-eq-align:  $\text{ALIGN}_4\ (\text{ALIGN}_4\ n) = \text{ALIGN}_4\ n$ 
  unfolding ALIGN4-def by auto

```



Zephyr uses two events: reschedule for free and swap for alloc for context switch

**definition** *reschedule* :: *State com*

**where** *reschedule*  $\equiv$

*'thd-state* := *'thd-state*(*the 'cur* := *READY*);;  
*'cur* := *Some* (*SOME t. 'thd-state t = READY*);;  
*'thd-state* := *'thd-state*(*the 'cur* := *RUNNING*)

**definition** *swap* :: *State com*

**where** *swap*  $\equiv$

*IF* ( $\exists t. 'thd-state\ t = READY$ ) *THEN*  
*'cur* := *Some* (*SOME t. 'thd-state t = READY*);;  
*'thd-state* := *'thd-state*(*the 'cur* := *RUNNING*)  
*ELSE*  
*'cur* := *None*  
*FI*

**definition** *block-num* :: *Mem-pool*  $\Rightarrow$  *mem-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*

**where** *block-num* *p bl sz*  $\equiv$  (*bl* - (*buf p*)) *div sz*

**definition** *clear-free-bit* :: (*mempool-ref*  $\Rightarrow$  *Mem-pool*)  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  (*mempool-ref*  $\Rightarrow$  *Mem-pool*)

**where** *clear-free-bit* *mp-info p l b*  $\equiv$

*mp-info* (*p* := (*mp-info p*) (*levels* := (*levels* (*mp-info p*))  
 $[l := ((levels\ (mp-info\ p))\ !\ l)\ (bits := (bits\ ((levels\ (mp-info\ p))\ !\ l))$   
 $[b := ALLOCATED]]\ \rangle)$

**definition** *set-bit* :: (*mempool-ref*  $\Rightarrow$  *Mem-pool*)  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *BlockState*  $\Rightarrow$  (*mempool-ref*  $\Rightarrow$  *Mem-pool*)

**where** *set-bit* *mp-info p l b st*  $\equiv$

*mp-info* (*p* := (*mp-info p*) (*levels* := (*levels* (*mp-info p*))  
 $[l := ((levels\ (mp-info\ p))\ !\ l)\ (bits := (bits\ ((levels\ (mp-info\ p))\ !\ l))$   
 $[b := st]]\ \rangle)$

**abbreviation** *set-bit-free* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b FREE*

**abbreviation** *set-bit-alloc* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b ALLOCATED*

**abbreviation** *set-bit-divide* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b DIVIDED*

**abbreviation** *set-bit-noexist* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b NOEXIST*

**abbreviation** *set-bit-freeing* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b FREEING*

**abbreviation** *set-bit-allocating* *mp-info p l b*  $\equiv$  *set-bit* *mp-info p l b ALLOCATING*

**definition** *set-bit-s* :: *State*  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *BlockState*  $\Rightarrow$  *State*

**where** *set-bit-s* *s p l b st*  $\equiv$

*s*(*mem-pool-info* := *set-bit* (*mem-pool-info s*) *p l b st*  $\rangle$ )

**lemma** *set-bit-prev-len*:

*length* (*bits* (*levels* (*mp-info p*) ! *l*)) = *length* (*bits* (*levels* ((*set-bit* *mp-info p l b flg*) *p*) ! *l*))

**apply**(*simp* *add:set-bit-def*)

**using** *list-updt-samelen*  
**by** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.select-convs*(1) *Mem-pool-lvl.surjective*  
*Mem-pool-lvl.update-convs*(1) *list-update-beyond* *not-less* *nth-list-update-eq*)

**lemma** *set-bit-prev-len2*:  
 $l \neq t \implies \text{length } (\text{bits } (\text{levels } (\text{mp-info } p) ! l)) = \text{length } (\text{bits } (\text{levels } ((\text{set-bit mp-info } p \ t \ b \ \text{flg}) \ p) ! l))$   
**by** (*simp* *add:set-bit-def*)

**abbreviation** *get-bit* :: (*mempool-ref*  $\Rightarrow$  *Mem-pool*)  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *BlockState*  
**where** *get-bit* *mp-info* *p* *l* *b*  $\equiv$  (*bits* ((*levels* (*mp-info* *p*)) ! *l*)) ! *b*

**abbreviation** *get-bit-s* :: *State*  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *BlockState*  
**where** *get-bit-s* *s* *p* *l* *b*  $\equiv$  *get-bit* (*mem-pool-info* *s*) *p* *l* *b*

**lemma** *set-bit-get-bit-eq*:  
 $l < \text{length } (\text{levels } (\text{mp-info } p)) \implies$   
 $b < \text{length } (\text{bits } (\text{levels } (\text{mp-info } p) ! l)) \implies$   
 $\text{mp-info2} = \text{set-bit mp-info } p \ l \ b \ st \implies$   
 $\text{get-bit mp-info2 } p \ l \ b = st$   
**by** (*simp* *add:set-bit-def*)

**lemma** *set-bit-get-bit-eq2*:  
 $l < \text{length } (\text{levels } ((\text{mem-pool-info } Va) \ p)) \implies$   
 $b < \text{length } (\text{bits } (\text{levels } ((\text{mem-pool-info } Va) \ p) ! l)) \implies$   
 $\text{get-bit-s } (Va \parallel \text{mem-pool-info} := \text{set-bit } (\text{mem-pool-info } Va) \ p \ l \ b \ st)) \ p \ l \ b = st$   
**using** *set-bit-get-bit-eq*  
 $[of \ l \ (\text{mem-pool-info } Va) \ p \ b \ \text{set-bit } (\text{mem-pool-info } Va) \ p \ l \ b \ st]$   
**by** *simp*

**lemma** *set-bit-get-bit-neq*:  
 $p \neq p1 \vee l \neq l1 \vee b \neq b1 \implies$   
 $\text{mp-info2} = \text{set-bit mp-info } p \ l \ b \ st \implies$   
 $\text{get-bit mp-info2 } p1 \ l1 \ b1 = \text{get-bit mp-info } p1 \ l1 \ b1$   
**apply** (*simp* *add:set-bit-def*) **apply** *auto*  
**by** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.select-convs*(1) *Mem-pool-lvl.surjective*  
*Mem-pool-lvl.update-convs*(1) *list-update-beyond* *not-less* *nth-list-update-eq* *nth-list-update-neq*)

**lemma** *set-bit-get-bit-neq2*:  
 $p \neq p1 \vee l \neq l1 \vee b \neq b1 \implies$   
 $\text{get-bit-s } (Va \parallel \text{mem-pool-info} := \text{set-bit } (\text{mem-pool-info } Va) \ p \ l \ b \ st)) \ p1 \ l1 \ b1$   
 $= \text{get-bit-s } Va \ p1 \ l1 \ b1$   
**using** *set-bit-get-bit-neq*  
 $[of \ p \ p1 \ l \ l1 \ b \ b1 \ \text{set-bit } (\text{mem-pool-info } Va) \ p \ l \ b \ st \ \text{mem-pool-info } Va]$

**by** *simp*

**definition** *buf-size* :: *Mem-pool*  $\Rightarrow$  *nat*  
**where** *buf-size* *m*  $\equiv$  *n-max* *m* \* *max-sz* *m*

**definition** *block-fits* :: *Mem-pool*  $\Rightarrow$  *mem-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *bool*  
**where** *block-fits* *p* *b* *bsz*  $\equiv$  *b* + *bsz* < *buf-size* *p* + *buf* *p* + 1

**definition** *block-ptr* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *mem-ref*  
**where** *block-ptr* *p* *lsize* *b*  $\equiv$  *buf* *p* + *lsize* \* *b*

**definition** *partner-bits* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *bool*  
**where** *partner-bits* *p* *l* *b*  $\equiv$  *let* *bits* = *bits* (*levels* *p* ! *l*);  
 $a = (b \text{ div } 4) * 4$  *in*  
 $bits!a = FREE \wedge bits!(a+1) = FREE \wedge bits!(a+2) =$   
 $FREE \wedge bits!(a+3) = FREE$

**lemma** *partbits-div4*:  $a \text{ div } 4 = b \text{ div } 4 \implies \text{partner-bits } p \text{ } l \text{ } a = \text{partner-bits } p \text{ } l \text{ } b$   
**by** (*simp add:partner-bits-def*)

**abbreviation** *noexist-bits* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *nat*  $\Rightarrow$  *bool*  
**where** *noexist-bits* *mp* *ii* *jj*  $\equiv$  (*bits* (*levels* *mp* ! *ii*)) ! *jj* = *NOEXIST*  
 $\wedge$  (*bits* (*levels* *mp* ! *ii*)) ! (*jj* + 1) = *NOEXIST*  
 $\wedge$  (*bits* (*levels* *mp* ! *ii*)) ! (*jj* + 2) = *NOEXIST*  
 $\wedge$  (*bits* (*levels* *mp* ! *ii*)) ! (*jj* + 3) = *NOEXIST*

**definition** *level-empty* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *bool*  
**where** *level-empty* *p* *n*  $\equiv$  *free-list* (*levels* *p* ! *n*) = []

**definition** *head-free-list* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *mem-ref*  
**where** *head-free-list* *p* *l*  $\equiv$  *hd* (*free-list* ((*levels* *p*) ! *l*))

**definition** *rmhead-free-list* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *Mem-pool*  
**where** *rmhead-free-list* *p* *l*  $\equiv$

$p \langle \text{levels} := (\text{levels } p) \rangle$   
 $[l := ((\text{levels } p) ! l) \langle \text{free-list} := \text{tl } (\text{free-list } ((\text{levels } p) ! l)) \rangle] \rangle$

**definition** *remove-free-list* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *mem-ref*  $\Rightarrow$  *Mem-pool*  
**where** *remove-free-list* *p* *l* *b*  $\equiv$

$p \langle \text{levels} := (\text{levels } p) \rangle$   
 $[l := ((\text{levels } p) ! l) \langle \text{free-list} := \text{remove1 } b (\text{free-list } ((\text{levels } p) ! l)) \rangle] \rangle$

**definition** *append-free-list* :: *Mem-pool*  $\Rightarrow$  *nat*  $\Rightarrow$  *mem-ref*  $\Rightarrow$  *Mem-pool*  
**where** *append-free-list* *p* *l* *b*  $\equiv$

$p \langle \text{levels} := (\text{levels } p) \rangle$   
 $[l := ((\text{levels } p) ! l) \langle \text{free-list} := (\text{free-list } ((\text{levels } p) ! l)) @ [b] \rangle] \rangle$

**definition** *in-free-list* :: *mem-ref*  $\Rightarrow$  *mem-ref* *list*  $\Rightarrow$  *bool*

where  $\text{in-free-list } v \text{ fl} \equiv (\exists i < \text{length fl}. \text{fl}[i] = v)$

### 2.3 specification of events

**lemma** *timeout-lm*:  $(\text{timeout} = \text{FOREVER} \vee \text{timeout} = \text{NOWAIT} \vee \text{timeout} > 0) = (\text{timeout} \geq -1)$   
**by** *auto*

**definition** *Mem-pool-alloc* ::  $\text{Thread} \Rightarrow \text{mempool-ref} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow (\text{EventLabel}, \text{Core}, \text{State}, \text{State com option}) \text{ event}$

**where** *Mem-pool-alloc*  $t \ p \ sz \ \text{timeout} =$

*EVENT* *Mem-pool-allocE* [*MPRef*  $p$ , *Natural*  $sz$ , *Integer*  $\text{timeout}$ ]  $\Rightarrow (\mathcal{T} \ t)$   
**WHEN**  
 $p \in \text{'mem-pools}$   
 $(* \wedge \text{'cur} = \text{Some } t *) (* \text{ } t \text{ is the current thread } *) (** \text{ this condition is not stable on rely condition } **)$   
 $\wedge \text{timeout} \geq -1 (* \text{ equiv to } (\text{timeout} = \text{FOREVER} \vee \text{timeout} = \text{NOWAIT} \vee \text{timeout} > 0) *)$   
 $(* \wedge p \in \text{'pools-of-thread } t *) (* \text{ the mem pool } p \text{ is shared in the thread } t *)$   
**THEN**  
 $(t \blacktriangleright \text{'tmout} := \text{'tmout}(t := \text{timeout}));;$   
 $(t \blacktriangleright \text{'end} := \text{'end}(t := 0));;$   
 $(t \blacktriangleright \text{IF } \text{timeout} > 0 \text{ THEN}$   
 $\quad \text{'end} := \text{'end}(t := \text{'tick} + \text{nat } \text{timeout})$   
 $\quad \text{FI});;$   
 $(t \blacktriangleright \text{'mempoolalloc-ret} := \text{'mempoolalloc-ret } (t := \text{None}));;$   
 $(t \blacktriangleright \text{'ret} := \text{'ret}(t := \text{ESIZEERR}));;$   
 $(t \blacktriangleright \text{'rf} := \text{'rf}(t := \text{False}));;$   
**WHILE**  $\neg (\text{'rf } t)$  **DO**  
 $(* ===== \text{start: ret} = \text{pool-alloc}(p, \text{block}, \text{size}); =====$   
 $*)$   
 $(*(t \blacktriangleright \text{'lsizes} := \text{'lsizes}(t := []));*)$   
 $(t \blacktriangleright \text{'blk} := \text{'blk}(t := \text{NULL}));;$   
 $(t \blacktriangleright \text{'alloc-lsize-r} := \text{'alloc-lsize-r}(t := \text{False}));;$   
 $(t \blacktriangleright \text{'alloc-l} := \text{'alloc-l}(t := -1));;$   
 $(t \blacktriangleright \text{'free-l} := \text{'free-l}(t := -1));;$   
 $(t \blacktriangleright \text{'lsizes} := \text{'lsizes}(t := [\text{ALIGN}_4 (\text{max-sz } (\text{'mem-pool-info } p))]);;$   
 $(t \blacktriangleright \text{'i} := \text{'i}(t := 0));;$   
**WHILE**  $\text{'i } t < n\text{-levels } (\text{'mem-pool-info } p) \wedge \neg \text{'alloc-lsize-r } t$  **DO**  
**IF**  $\text{'i } t > 0$  **THEN**  
 $(t \blacktriangleright \text{'lsizes} := \text{'lsizes}(t := \text{'lsizes } t @ [\text{ALIGN}_4 (\text{'lsizes } t ! (\text{'i } t - 1) \text{ div } 4)]));;$   
**FI**;;  
**IF**  $\text{'lsizes } t ! \text{'i } t < sz$  **THEN**  
 $(t \blacktriangleright \text{'alloc-lsize-r} := \text{'alloc-lsize-r}(t := \text{True}))$   
**ELSE**  
 $(t \blacktriangleright \text{'alloc-l} := \text{'alloc-l}(t := \text{int } (\text{'i } t)));;$

```

    IF  $\neg$  level-empty ('mem-pool-info p) ('i t) THEN
      (t  $\blacktriangleright$  'free-l := 'free-l(t := int ('i t)))
    FI;;
    (t  $\blacktriangleright$  'i := 'i(t := 'i t + 1))
  FI
OD;;

IF 'alloc-l t < 0 THEN
  (t  $\blacktriangleright$  'ret := 'ret(t := ESIZEERR))
ELSE
  IF 'free-l t < 0 THEN
    (* block->data = NULL; *)
    (t  $\blacktriangleright$  'ret := 'ret(t := ENOMEM))
  ELSE
    (* ===== start: blk = alloc-block(p, free-l, lsizes[free-l]); *)
    (t  $\blacktriangleright$  ATOMIC
      (* ===== start: block = sys-dlist-get(&p->levels[l].free-list); *)
      IF level-empty ('mem-pool-info p) (nat ('free-l t)) THEN
        'blk := 'blk(t := NULL)
      ELSE
        'blk := 'blk(t := head-free-list ('mem-pool-info p) (nat ('free-l t)));

        (* sys-dlist-remove(node); *)
        'mem-pool-info := 'mem-pool-info (p := rmhead-free-list ('mem-pool-info
p) (nat ('free-l t)))

      FI;;
      (* ===== end: block = sys-dlist-get(&p->levels[l].free-list); *)

      IF 'blk t  $\neq$  NULL THEN
        (* clear-free-bit(p, l, block-num(p, block, lsz)); *)
        'mem-pool-info := set-bit-allocating 'mem-pool-info p (nat ('free-l t))
          (block-num ('mem-pool-info p) ('blk t) ((('lsizes t)!(nat
('free-l t)))));;
        (* set the allocating node info of the thread *)
        'allocating-node := 'allocating-node (t := Some (pool = p, level = nat
('free-l t),
          block = (block-num ('mem-pool-info p) ('blk t) ((('lsizes t)!(nat
('free-l t)))), data = 'blk t ))
      FI
    END);;
    (* ===== end: blk = alloc-block(p, free-l, lsizes[free-l]); *)

    IF 'blk t = NULL THEN
      (t  $\blacktriangleright$  'ret := 'ret (t := EAGAIN))
    ELSE
      FOR (t  $\blacktriangleright$  'from-l := 'from-l(t := 'free-l t));
        (* level-empty ('mem-pool-info p) (nat ('alloc-l t))  $\wedge$  *) 'from-l t <

```

```

'alloc-l t;
  (***** we remove the FOR termination condition "level-empty"
to remove a concurrency BUG here *****)
  (t ► 'from-l := 'from-l(t := 'from-l t + 1)) DO

    (* ===== start: blk = break-block(p, blk, from-l, lsizes); *)
    (t ► ATOMIC
      'bn := 'bn (t := block-num ('mem-pool-info p) ('blk t) (('lsizes
t)!(nat ('from-l t)))));

    'mem-pool-info := set-bit-divide 'mem-pool-info p (nat ('from-l t))
('bn t));

    'mem-pool-info := set-bit-allocating 'mem-pool-info p (nat ('from-l t
+ 1)) (4 * 'bn t));

    (* set the allocating node info of the thread *)
    'allocating-node := 'allocating-node (t := Some (pool = p, level =
nat ('from-l t + 1),
      block = 4 * 'bn t, data = 'blk t ));

    FOR 'i := 'i (t := 1);
      'i t < 4;
      'i := 'i (t := 'i t + 1) DO
        'lbn := 'lbn (t := 4 * 'bn t + 'i t);
        'lsz := 'lsz (t := ('lsizes t) ! (nat ('from-l t + 1)));
        'block2 := 'block2(t := 'lsz t * 'i t + 'blk t);

        (* set-free-bit(p, l + 1, lbn); *)
        'mem-pool-info := set-bit-free 'mem-pool-info p (nat ('from-l t +
1)) ('lbn t));

        IF block-fits ('mem-pool-info p) ('block2 t) ('lsz t) THEN

          (* sys-dlist-append(&p->levels[l + 1].free-list, block2); *)
          'mem-pool-info := 'mem-pool-info (p :=
            append-free-list ('mem-pool-info p) (nat ('from-l t + 1))
('block2 t) )

          FI
        ROF

      END)
    (* ===== end: blk = break-block(p, blk, from-l, lsizes); *)

  ROF;;

  (* finally set the node from allocating to allocated and remove the allocating
node info of the thread *)
  (t ► 'mem-pool-info := set-bit-alloc 'mem-pool-info p (nat ('alloc-l t))

```

```

      (block-num ('mem-pool-info p) ('blk t) (('lsizes t)!(nat
('alloc-l t))));
      'allocating-node := 'allocating-node (t := None)
    );;;

    (t ► 'mempoolalloc-ret := 'mempoolalloc-ret (t :=
      Some (pool = p, level = nat ('alloc-l t),
        block = block-num ('mem-pool-info p) ('blk t) (('lsizes t)!(nat
('alloc-l t))),
        data = 'blk t ));

    (t ► 'ret := 'ret (t := OK))
  FI
FI
FI;;
(* ===== end: ret = pool-alloc(p, block, size); =====
*)

(* IF 'ret t = 0 ∨ timeout = NOWAIT ∨ 'ret t = EAGAIN ∨ 'ret t ≠
ENOMEM THEN *)
(***** we change the IF condition to remove a functional BUG here *****)
IF 'ret t = OK ∨ timeout = NOWAIT ∨ 'ret t = ESIZEERR THEN
  (t ► 'rf := 'rf(t := True));
  IF 'ret t = EAGAIN THEN (*EAGAIN should not export to users*)
    (t ► 'ret := 'ret(t := ENOMEM))
  FI
ELSE
  IF 'ret t = EAGAIN THEN SKIP
  ELSE
    (t ► ATOMIC

      (* -pend-current-thread(&p->wait-q, timeout); *)
      'thd-state := 'thd-state(the 'cur := BLOCKED);
      (* 'cur := None; *)
      'mem-pool-info := 'mem-pool-info(p := 'mem-pool-info p (wait-q :=
wait-q ('mem-pool-info p) @ [the 'cur] ));

      (* -Swap(key); *)
      swap

    END);;

    IF 'tmout t ≠ FOREVER THEN
      (t ► 'tmout := 'tmout (t := int ('endt t) - int 'tick));
      IF 'tmout t < 0 THEN
        (t ► 'rf := 'rf(t := True));
        (t ► 'ret := 'ret (t := ETIMEOUT))
      FI
    FI
  FI

```

FI  
FI  
OD

END

**definition** *Mem-pool-free* :: Thread  $\Rightarrow$  Mem-block  $\Rightarrow$  (EventLabel, Core, State, State com option) event

**where** *Mem-pool-free* t b =

EVENT *Mem-pool-freeE* [Block b]  $\Rightarrow$  ( $\mathcal{T}$  t)

WHEN

*pool* b  $\in$  'mem-pools

$\wedge$  level b < length (levels ('mem-pool-info (pool b)))

$\wedge$  block b < length (bits (levels ('mem-pool-info (pool b))!(level b)))

$\wedge$  data b = block-ptr ('mem-pool-info (pool b)) ((ALIGN4 (max-sz ('mem-pool-info (pool b)))) div (4 ^ (level b))) (block b)

(\* $\wedge$  (bits ((levels ('mem-pool-info (pool b)))(level b)))(block b) = ALLOCATED

$\wedge$  'cur = Some t\*) (\* t is the current thread \*)

(\*  $\wedge$  pool b  $\in$  'pools-of-thread t \*) (\* the mem pool is shared in the thread t \*)

THEN

(\* here we set the bit to FREEING, so that other thread cannot mem-pool-free the same block

it also requires that it can only free ALLOCATED block \*)

(t  $\blacktriangleright$  AWAIT (bits ((levels ('mem-pool-info (pool b)))(level b)))(block b) = ALLOCATED THEN

'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) (level b) (block b));;

'freeing-node := 'freeing-node (t := Some b) (\* set the freeing node of current thread \*)

END);;

(t  $\blacktriangleright$  'need-resched := 'need-resched(t := False));;

(\* (t  $\blacktriangleright$  'lsizes := 'lsizes(t := []));; \*)

(t  $\blacktriangleright$  'lsizes := 'lsizes(t := [ALIGN4 (max-sz ('mem-pool-info (pool b)))]));;

FOR (t  $\blacktriangleright$  'i := 'i(t := 1));

'i t  $\leq$  level b;

(t  $\blacktriangleright$  'i := 'i(t := 'i t + 1)) DO

(t  $\blacktriangleright$  'lsizes := 'lsizes(t := 'lsizes t @ [ALIGN4 ('lsizes t ! ('i t - 1) div 4)]))

ROF;;

(\* === start: free-block(get-pool(block->id.pool), block->id.level, lsizes, block->id.block); \*)

(t  $\blacktriangleright$  'free-block-r := 'free-block-r (t := True));;

(t  $\blacktriangleright$  'bn := 'bn (t := block b));;

(t  $\blacktriangleright$  'lvl := 'lvl (t := level b));;

WHILE 'free-block-r t DO

(t  $\blacktriangleright$  'lsz := 'lsz (t := 'lsizes t ! ('lvl t)));;



```

(t ► 'blk := 'blk (t := block-ptr ('mem-pool-info (pool b)) ('lsz t) ('bn t)));

(t ► ATOMIC

  'mem-pool-info := set-bit-free 'mem-pool-info (pool b) ('lvl t) ('bn t);
  'freeing-node := 'freeing-node (t := None); (* remove the freeing node info
of the thread *)

  IF 'lvl t > 0 ∧ partner-bits ('mem-pool-info (pool b)) ('lvl t) ('bn t) THEN
    FOR 'i := 'i(t := 0);
      'i t < 4;
      'i := 'i(t := 'i t + 1) DO
        'bb := 'bb (t := ('bn t div 4) * 4 + 'i t);
        (*(t ► 'mem-pool-info := clear-free-bit 'mem-pool-info (pool b) ('lvl t)
('bb t));*)
        'mem-pool-info := set-bit-noexist 'mem-pool-info (pool b) ('lvl t) ('bb
t));
        'block-pt := 'block-pt (t := block-ptr ('mem-pool-info (pool b)) ('lsz t)
('bb t));
        IF 'bn t ≠ 'bb t ∧ block-fits ('mem-pool-info (pool b))
          ('block-pt t)
          ('lsz t) THEN

          (* sys-dlist-remove(block-ptr(p, lsz, b)); *)
          'mem-pool-info := 'mem-pool-info ((pool b) :=
            remove-free-list ('mem-pool-info (pool b)) ('lvl t) ('block-pt t))

          FI
          ROF;;

      (
        (* 'j := 'j (t := 'lvl t); (* use lbn and j to store the previous lvl and bn,
or can not give the post condition *)
        'lbn := 'lbn (t := 'bn t); (* since the lbn and j are not used in M-pool-free
*)
        'lvl := 'lvl (t := 'j t - 1);
        'bn := 'bn (t := 'lbn t div 4));*)
        'lvl := 'lvl (t := 'lvl t - 1);
        'bn := 'bn (t := 'bn t div 4);
        (* we add this statement. set the parent node from divided to freeing *)
        'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) ('lvl t) ('bn t);
        (* 'freeing-node := 'freeing-node (t := Some (pool = (pool b), level = ('lvl
t),
          block = ('bn t), data = block-ptr ('mem-pool-info (pool b)) ('lsz
t) ('bn t) ))*)
        'freeing-node := 'freeing-node (t := Some (pool = (pool b), level = ('lvl
t),
          block = ('bn t),
          data = block-ptr ('mem-pool-info (pool b))
            (((ALIGN4 (max-sz ('mem-pool-info (pool b)))) div (4 ^

```

```

('lvl t))))
      ('bn t) [])
    )

ELSE
  IF block-fits ('mem-pool-info (pool b)) ('blk t) ('lsz t) THEN

    (* sys-dlist-append(&p->levels[level].free-list, block); *)
    'mem-pool-info := 'mem-pool-info ((pool b) :=
      append-free-list ('mem-pool-info (pool b)) ('lvl t) ('blk t) )
    FI;;

    'free-block-r := 'free-block-r (t := False)
  FI

END)

OD;;
(* == end: free-block(get-pool(block->id.pool), block->id.level, lsize, block->id.block);
*)

(t ► ATOMIC

  WHILE wait-q ('mem-pool-info (pool b)) ≠ [] DO
    'th := 'th (t := hd (wait-q ('mem-pool-info (pool b))));
    (* -unpend-thread(th); *)
    'mem-pool-info := 'mem-pool-info (pool b := 'mem-pool-info (pool b)
      (wait-q := tl (wait-q ('mem-pool-info (pool b))));
    (* -ready-thread(th); *)
    'thd-state := 'thd-state ('th t := READY);
    'need-resched := 'need-resched(t := True)
  OD;;

  IF 'need-resched t THEN
    reschedule
  FI
END)
END

```

**definition** *Schedule* :: Thread  $\Rightarrow$  (EventLabel, Core, State, State com option) event

**where** *Schedule* t  $\equiv$   
 EVENT *ScheduleE* [Thread t]  $\Rightarrow$   $\mathcal{S}$   
 THEN  
 AWAIT 'thd-state t = READY THEN (\* only schedule the READY threads  
 \*)  
 IF ('cur ≠ None) THEN  
 'thd-state := 'thd-state(the ('cur) := READY);  
 'cur := None

```

    FI;;
    'cur := Some t;;
    'thd-state := 'thd-state(t := RUNNING)
  END
END

```

**definition** *Tick* :: (*EventLabel*, *Core*, *State*, *State com option*) *event*  
**where** *Tick*  $\equiv$   
*EVENT TickE []  $\Rightarrow$  Timer*  
*THEN*  
*'tick := 'tick + 1*  
*END*

**end**

**theory** *invariant*  
**imports** *mem-spec HOL-Eisbach.Eisbach-Tools*  
**begin**

this theory defines the invariant and its lemmas.

### 3 invariants

#### 3.1 defs of invariants

we consider multi-threaded execution on mono-core. A thread is the currently executing thread iff it is in RUNNING state.

**definition** *inv-cur* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-cur s*  $\equiv \forall t. \text{cur } s = \text{Some } t \longleftrightarrow \text{thd-state } s \text{ } t = \text{RUNNING}$

**abbreviation** *dist-list* :: 'a *list*  $\Rightarrow$  *bool*  
**where** *dist-list l*  $\equiv \forall i \ j. i < \text{length } l \wedge j < \text{length } l \wedge i \neq j \longrightarrow !i \neq !j$

the relation of thread state and wait queue. here we dont consider other modules of zephyr, so blocked thread is in wait que of mem pools.

**definition** *inv-thd-waitq* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-thd-waitq s*  $\equiv$   
 ( $\forall p \in \text{mem-pools } s. \forall t \in \text{set } (\text{wait-q } (\text{mem-pool-info } s \text{ } p)). \text{thd-state } s \text{ } t = \text{BLOCKED}$ )

(*\* thread in waitq is BLOCKED \**)  
 $\wedge (\forall t. \text{thd-state } s \text{ } t = \text{BLOCKED} \longrightarrow (\exists p \in \text{mem-pools } s. t \in \text{set } (\text{wait-q } (\text{mem-pool-info } s \text{ } p))))$

(*\* BLOCKED thread is in a waitq \**)  
 $\wedge (\forall p \in \text{mem-pools } s. \text{dist-list } (\text{wait-q } (\text{mem-pool-info } s \text{ } p)))$   
 (*\* threads in a waitq are different with each other, which means a thread could not waiting for the same pool two times \**)

$$\begin{aligned} & \wedge (\forall p \ q. \ p \in \text{mem-pools } s \wedge q \in \text{mem-pools } s \wedge p \neq q \longrightarrow (\nexists t. t \in \text{set } (\text{wait-q} \\ & (\text{mem-pool-info } s \ p))) \\ & \wedge t \in \text{set } (\text{wait-q } (\text{mem-pool-info } s \ q)))) \end{aligned}$$

invariant of configuration of memory pools. its actually a well-formed property for memory configuration. (1) the max size (the size of top-level (level 0) block) is  $4^{n\_levels}$  times of block size of the lowest level.  $4 * n$  means that the block size of the lowest level is aligned with 4. (2) the block number at level 0 ( $n\_max$ )  $\leq 0$ , and the max number of levels is  $n\_levels \leq 0$  (3)  $n\_level$  is equal to the length of levels list. (4) the length of bitmap list at each level is equal to the block number at the same level. Thus, bitmap saves a complete quad-tree with height of  $n\_levels$ . A real memory pool is a top subtree of the complete tree. bits of subnodes of a leaf node (ALLOCATED, FREE, ALLOCATING, FREEING) is NOEXIST.

**abbreviation**  $\text{inv-mempool-info-mp} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{bool}$

**where**  $\text{inv-mempool-info-mp } s \ p \equiv$

$$\begin{aligned} & \text{let } mp = \text{mem-pool-info } s \ p \text{ in} \\ & \quad \text{buf } mp \neq \text{NULL} \wedge (\exists n > 0. \text{max-sz } mp = (4 * n) * (4 ^ n\_levels \ mp)) \\ & \quad \wedge n\_max \ mp > 0 \wedge n\_levels \ mp > 0 \\ & \quad \wedge n\_levels \ mp = \text{length } (\text{levels } mp) \\ & \quad \wedge (\forall i < \text{length } (\text{levels } mp). \text{length } (\text{bits } (\text{levels } mp \ ! \ i)) = (n\_max \ mp) * 4 ^ i) \end{aligned}$$

**definition**  $\text{inv-mempool-info} :: \text{State} \Rightarrow \text{bool}$

**where**  $\text{inv-mempool-info } s \equiv \forall p \in \text{mem-pools } s. \text{inv-mempool-info-mp } s \ p$

**lemma**  $\text{inv-max-sz-gt0}: \text{inv-mempool-info } s \implies \forall p \in \text{mem-pools } s. \text{let } mp = \text{mem-pool-info } s \ p \text{ in } \text{max-sz } mp > 0$

**unfolding**  $\text{inv-mempool-info-def}$  **using**  $\text{neg0-conv}$  **by**  $\text{fastforce}$

invariant between bitmap and free block list at each level. (1) bit of a block is FREE, iff its start address is in free list. the start address is  $\text{buf } mp + j * (\text{max\_sz } mp \text{ div } (4^i))$ , the start address of the mempool + block size at this level \* block index (2) start address of blocks in free list is valid, i.e. it is the start address of some block (index  $n$ ), where  $n$  is in the range of block index at the level (3) start address of blocks in free list are different with each other.

**abbreviation**  $\text{inv-bitmap-freelist-mp} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{bool}$

**where**  $\text{inv-bitmap-freelist-mp } s \ p \equiv$

$$\begin{aligned} & \text{let } mp = \text{mem-pool-info } s \ p \text{ in} \\ & \quad \forall i < \text{length } (\text{levels } mp). \\ & \quad \quad \text{let } bts = \text{bits } (\text{levels } mp \ ! \ i); \\ & \quad \quad \text{fl} = \text{free-list } (\text{levels } mp \ ! \ i) \text{ in} \\ & \quad \quad (\forall j < \text{length } bts. bts \ ! \ j = \text{FREE} \longleftrightarrow \text{buf } mp + j * (\text{max-sz } mp \text{ div } \\ & (4 ^ i)) \in \text{set } fl) \\ & \quad \quad (* \text{ the block corresponding to a free bit iff it is in freelist } *) \\ & \quad \quad \wedge (\forall j < \text{length } fl. (\exists n. n < n\_max \ mp * (4 ^ i) \wedge fl \ ! \ j = \text{buf } mp + n \end{aligned}$$

$*$  ( $\text{max-sz } mp \text{ div } (4 \wedge i)$ ))  
 $(*$  pointers in freelist are head address of blocks  $*$ )  
 $\wedge \text{distinct fl } (*(\forall k j. k < \text{length fl} \wedge j < \text{length fl} \longrightarrow \text{fl!}k = \text{fl!}j \longrightarrow$   
 $k = j) *)$   
 $(*$  pointers in freelist are different with each other  $*$ )

**definition**  $\text{inv-bitmap-freelist} :: \text{State} \Rightarrow \text{bool}$

**where**  $\text{inv-bitmap-freelist } s \equiv$

$\forall p \in \text{mem-pools } s. \text{inv-bitmap-freelist-mp } s \ p$

this invariant represents that a memory pools is forest of valid quad-trees of blocks. parent node of a leaf node (ALLOCATED, FREE, ALLOCATING, FREEING) is an inner node (DIVIDED). parent node of an inner node (DIVIDED) is also a DIVIDED node. child nodes of a NOEXIST node is also NOEXIST nodes. parent node of a NOEXIST node should not be DIVIDE nodes (may be NOEXIST, ALLOCATED, FREE, ALLOCATING, FREEING)

**abbreviation**  $\text{inv-bitmap-mp} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{bool}$

**where**  $\text{inv-bitmap-mp } s \ p \equiv$

$\text{let } mp = \text{mem-pool-info } s \ p \text{ in}$   
 $\forall i < \text{length } (\text{levels } mp).$   
 $\text{let } bts = \text{bits } (\text{levels } mp \ ! \ i) \text{ in}$   
 $(\forall j < \text{length } bts.$   
 $(bts \ ! \ j = \text{FREE} \vee bts \ ! \ j = \text{FREEING} \vee bts \ ! \ j = \text{ALLOCATED}$   
 $\vee bts \ ! \ j = \text{ALLOCATING} \longrightarrow$   
 $(i > 0 \longrightarrow (\text{bits } (\text{levels } mp \ ! \ (i - 1))) \ ! \ (j \text{ div } 4) = \text{DIVIDED})$   
 $\wedge (i < \text{length } (\text{levels } mp) - 1 \longrightarrow \text{noexist-bits } mp \ (i+1) \ (j*4)$   
 $))$   
 $\wedge (bts \ ! \ j = \text{DIVIDED} \longrightarrow i > 0 \longrightarrow (\text{bits } (\text{levels } mp \ ! \ (i - 1))) \ !$   
 $(j \text{ div } 4) = \text{DIVIDED})$   
 $\wedge (bts \ ! \ j = \text{NOEXIST} \longrightarrow i < \text{length } (\text{levels } mp) - 1$   
 $\longrightarrow \text{noexist-bits } mp \ (i+1) \ (j*4))$   
 $\wedge (bts \ ! \ j = \text{NOEXIST} \wedge i > 0 \longrightarrow (\text{bits } (\text{levels } mp \ ! \ (i - 1))) \ ! \ (j$   
 $\text{div } 4) \neq \text{DIVIDED})$

**definition**  $\text{inv-bitmap} :: \text{State} \Rightarrow \text{bool}$

**where**  $\text{inv-bitmap } s \equiv$

$\forall p \in \text{mem-pools } s. \text{inv-bitmap-mp } s \ p$

due to the rule of merge as possible, there should not exist a node with 4 FREE child blocks. In free syscall, 4 free child blocks should be merged to a bigger block.

**abbreviation**  $\text{inv-bitmap-not4free-mp} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{bool}$

**where**  $\text{inv-bitmap-not4free-mp } s \ p \equiv$

$\text{let } mp = \text{mem-pool-info } s \ p \text{ in}$   
 $\forall i < \text{length } (\text{levels } mp).$   
 $\text{let } bts = \text{bits } (\text{levels } mp \ ! \ i) \text{ in}$   
 $(\forall j < \text{length } bts. i > 0 \longrightarrow \neg \text{partner-bits } mp \ i \ j)$

**definition** *inv-bitmap-not4free* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-bitmap-not4free* *s*  $\equiv$   
 $\forall p \in \text{mem-pools } s. \text{inv-bitmap-not4free-mp } s \ p$

blocks at level 0 should not be NOEXIST. If so, the memory pool does not exist. We only allow real memory pools.

**definition** *inv-bitmap0* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-bitmap0* *s*  $\equiv$   
 $\forall p \in \text{mem-pools } s. \text{let bits0} = \text{bits } (\text{levels } (\text{mem-pool-info } s \ p) \ ! \ 0) \text{ in } \forall i < \text{length } \text{bits0}. \text{bits0} \ ! \ i \neq \text{NOEXIST}$

blocks at last level (n\_level - 1) should not be split again, thus should not be DIVIDED

**definition** *inv-bitmapn* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-bitmapn* *s*  $\equiv$   
 $\forall p \in \text{mem-pools } s. \text{let bitsn} = \text{bits } ((\text{levels } (\text{mem-pool-info } s \ p) \ ! \ (\text{length } (\text{levels } (\text{mem-pool-info } s \ p)) - 1)))$   
 $\text{in } \forall i < \text{length } \text{bitsn}. \text{bitsn} \ ! \ i \neq \text{DIVIDED}$

**definition** *mem-block-addr-valid* :: *State*  $\Rightarrow$  *Mem-block*  $\Rightarrow$  *bool*  
**where** *mem-block-addr-valid* *s* *b*  $\equiv$   
 $\text{data } b = \text{buf } (\text{mem-pool-info } s \ (\text{pool } b)) + (\text{block } b) * ((\text{max-sz } (\text{mem-pool-info } s \ (\text{pool } b))) \text{ div } (4 \wedge (\text{level } b)))$

invariants between FREEING/ALLOCATING blocks and freeing/allocating\_node variables.

**definition** *inv-aux-vars* :: *State*  $\Rightarrow$  *bool*  
**where** *inv-aux-vars* *s*  $\equiv$   
 $(\forall t \ n. \text{freeing-node } s \ t = \text{Some } n \longrightarrow \text{get-bit } (\text{mem-pool-info } s) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{FREEING})$   
 $(* \text{freeing node is state of FREEING} *)$   
 $\wedge (\forall n. \text{get-bit } (\text{mem-pool-info } s) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{FREEING} \wedge \text{mem-block-addr-valid } s \ n$   
 $\longrightarrow (\exists t. \text{freeing-node } s \ t = \text{Some } n))$   
 $(* \text{node of state of FREEING is freeing} *)$   
 $\wedge (\forall t \ n. \text{allocating-node } s \ t = \text{Some } n \longrightarrow \text{get-bit } (\text{mem-pool-info } s) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{ALLOCATING})$   
 $(* \text{freeing node is state of FREEING} *)$   
 $\wedge (\forall n. \text{get-bit } (\text{mem-pool-info } s) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{ALLOCATING} \wedge \text{mem-block-addr-valid } s \ n$   
 $\longrightarrow (\exists t. \text{allocating-node } s \ t = \text{Some } n))$   
 $(* \text{node of state of FREEING is freeing} *)$   
 $\wedge (\forall t1 \ t2 \ n1 \ n2. t1 \neq t2 \wedge \text{freeing-node } s \ t1 = \text{Some } n1 \wedge \text{freeing-node } s \ t2 = \text{Some } n2$   
 $\longrightarrow \neg(\text{pool } n1 = \text{pool } n2 \wedge \text{level } n1 = \text{level } n2 \wedge \text{block } n1 = \text{block } n2))$   
 $(* \text{here we only consider the pool, level, and block, not the first addr of the block} *)$

(\* freeing nodes are different each other \*)  
 $\wedge (\forall t1\ t2\ n1\ n2. t1 \neq t2 \wedge \text{allocating-node } s\ t1 = \text{Some } n1 \wedge \text{allocating-node } s\ t2 = \text{Some } n2$   
 $\longrightarrow \neg(\text{pool } n1 = \text{pool } n2 \wedge \text{level } n1 = \text{level } n2 \wedge \text{block } n1 = \text{block } n2))$   
 (\* allocating node are different each other \*)  
 $\wedge (\forall t1\ t2\ n1\ n2. \text{allocating-node } s\ t1 = \text{Some } n1 \wedge \text{freeing-node } s\ t2 = \text{Some } n2$   
 $\longrightarrow \neg(\text{pool } n1 = \text{pool } n2 \wedge \text{level } n1 = \text{level } n2 \wedge \text{block } n1 = \text{block } n2))$

**definition** *inv* :: *State*  $\Rightarrow$  *bool*

**where** *inv* *s*  $\equiv$  *inv-cur* *s*  $\wedge$  *inv-thd-waitq* *s*  $\wedge$  *inv-mempool-info* *s*  
 $\wedge$  *inv-bitmap-freelist* *s*  $\wedge$  *inv-bitmap* *s*  $\wedge$  *inv-aux-vars* *s*  
 $\wedge$  *inv-bitmap0* *s*  $\wedge$  *inv-bitmapn* *s*  $\wedge$  *inv-bitmap-not4free* *s*

**method** *simp-inv* = (*simp* *add:inv-def* *inv-bitmap-def* *inv-bitmap-freelist-def*  
*inv-mempool-info-def* *inv-thd-waitq-def* *inv-cur-def* *inv-aux-vars-def*  
*inv-bitmap0-def* *inv-bitmapn-def*  
*inv-bitmap-not4free-def* *mem-block-addr-valid-def*)

**method** *unfold-inv* = (*unfold* *inv-def* *inv-bitmap-def* *inv-bitmap-freelist-def*  
*inv-mempool-info-def* *inv-thd-waitq-def* *inv-cur-def* *inv-aux-vars-def*  
*inv-bitmap0-def* *inv-bitmapn-def*  
*inv-bitmap-not4free-def* *mem-block-addr-valid-def*)[1]

**lemma** *inv-imp-fl-lt0*:

*inv* *Va*  $\implies$   
 $\forall p \in \text{mem-pools } Va.$   
 $\text{let } mp = \text{mem-pool-info } Va\ p\ \text{in}$   
 $\forall i < \text{length } (\text{levels } mp).$   
 $\forall j < \text{length } (\text{free-list } (\text{levels } mp\ i)). \text{free-list } (\text{levels } mp\ i)\ !\ j > 0$

**apply**(*simp* *add:inv-def* *inv-mempool-info-def* *inv-bitmap-freelist-def*)

**apply**(*simp* *add:Let-def*) **apply** *clarsimp*

**by** *fastforce*

### 3.2 initial state $s_0$

we dont consider *mem\_pool\_init*, only define *s0* to show the state after memory pool initialization.

**axiomatization** *s0::State* **where**

*s0a1*: *cur* *s0* = *None* **and**

*s0a2*: *tick* *s0* = 0 **and**

*s0a3*: *thd-state* *s0* = ( $\lambda t. \text{READY}$ ) **and**

*s0a5*: *mem-pools* *s0*  $\neq \{\}$  **and**

*s0a7*:  $\forall p \in \text{mem-pools } s0. \text{wait-q } (\text{mem-pool-info } s0\ p) = []$  **and**

*s0a6*:  $\forall p \in \text{mem-pools } s0. \text{let } mp = \text{mem-pool-info } s0\ p\ \text{in}$

$$\text{buf } mp > 0 \wedge (\exists n > 0. \text{max-sz } mp = (4 * n) * (4 ^ n\text{-levels } mp))$$

$$\wedge n\text{-max } mp > 0 \wedge n\text{-levels } mp > 1$$

$$\wedge n\text{-levels } mp = \text{length } (\text{levels } mp) \text{ and}$$

$$s0a8: \forall p \in \text{mem-pools } s0. (* \text{ defines level 1 to } n *)$$

$$\text{let } mp = \text{mem-pool-info } s0 \text{ } p \text{ in}$$

$$\forall i. i > 0 \wedge i < \text{length } (\text{levels } mp) \longrightarrow$$

$$\text{length } (\text{bits } (\text{levels } mp ! i)) = n\text{-max } mp * 4 ^ i$$

$$\wedge (\forall j < \text{length } (\text{bits } (\text{levels } mp ! i)). \text{bits } (\text{levels } mp ! i) ! j =$$

$$NOEXIST)$$

$$\wedge \text{free-list } (\text{levels } mp ! i) = [] \text{ and}$$

$$s0a9: \forall p \in \text{mem-pools } s0. (* \text{ defines the level0 } *)$$

$$\text{let } mp = \text{mem-pool-info } s0 \text{ } p;$$

$$lv0 = (\text{levels } mp)!0 \text{ in}$$

$$\text{length } (\text{bits } lv0) = n\text{-max } mp$$

$$\wedge \text{length } (\text{free-list } lv0) = n\text{-max } mp$$

$$\wedge (\forall i < \text{length } (\text{bits } lv0). (\text{bits } lv0)!i = FREE)$$

$$\wedge (\forall i < \text{length } (\text{free-list } lv0). (\text{free-list } lv0) ! i = (\text{buf } mp) + i * \text{max-sz}$$

$$mp)$$

$$\wedge \text{distinct } (\text{free-list } lv0) \text{ and}$$

$$s0a4: \text{freeing-node } s0 = \text{Map.empty} \text{ and}$$

$$s0a10: \text{allocating-node } s0 = \text{Map.empty} \text{ and}$$

$$s0a11: \nexists n. \text{get-bit-s } s0 \text{ (pool } n) \text{ (level } n) \text{ (block } n) = FREEING \text{ and}$$

$$s0a12: \nexists n. \text{get-bit-s } s0 \text{ (pool } n) \text{ (level } n) \text{ (block } n) = ALLOCATING$$

**lemma**  $s0\text{-max-sz-gt0}$ :  $\forall p \in \text{mem-pools } s0. \text{let } mp = \text{mem-pool-info } s0 \text{ } p \text{ in } \text{max-sz } mp > 0$

**using**  $s0a6$  **zero-less-power** **by** **fastforce**

**lemma**  $s0\text{-inv-cur}$ :  $\text{inv-cur } s0$

**by**  $(\text{simp add: inv-cur-def } s0a1 \text{ } s0a3)$

**lemma**  $s0\text{-inv-thdwaitq}$ :  $\text{inv-thd-waitq } s0$

**by**  $(\text{simp add: inv-thd-waitq-def } s0a7 \text{ } s0a3)$

**lemma**  $s0\text{-inv-mempool-info}$ :  $\text{inv-mempool-info } s0$

**apply**  $(\text{simp add: inv-mempool-info-def Let-def})$  **apply**  $\text{clarsimp}$

**apply**  $(\text{rule conjI})$  **apply**  $(\text{metis neq0-conv } s0a6)$

**apply**  $(\text{rule conjI})$  **apply**  $(\text{meson } s0a6)$

**apply**  $(\text{rule conjI})$  **apply**  $(\text{meson } s0a6)$

**apply**  $(\text{rule conjI})$  **using**  $\text{neq0-conv } s0a6$  **apply** **fastforce**

**apply**  $(\text{rule conjI})$  **apply**  $(\text{meson } s0a6)$

**by**  $(\text{metis One-nat-def mult-numeral-1-right neq0-conv numeral-1-eq-Suc-0 power.simps}(1) \text{ } s0a8 \text{ } s0a9)$

**lemma**  $s0\text{-inv-bitmap-freelist}$ :  $\text{inv-bitmap-freelist } s0$

**apply**  $(\text{simp add: inv-bitmap-freelist-def})$

**apply**  $(\text{simp add: Let-def})$  **apply**  $\text{clarsimp}$



```

apply(case-tac i = 0)
apply(rule conjI) apply clarsimp apply (metis nth-mem s0a9)
apply(rule conjI) apply clarsimp apply (metis s0a9)
apply (meson s0a9)

apply(rule conjI) apply clarsimp
apply(subgoal-tac n-levels (mem-pool-info s0 p) = length (levels (mem-pool-info
s0 p))))
  prefer 2 apply (meson s0a6)
  apply(subgoal-tac get-bit-s s0 p i j ≠ FREE)
  prefer 2 apply (metis BlockState.distinct(13) s0a8)
  apply(subgoal-tac set (free-list (levels (mem-pool-info s0 p) ! i)) = {})
  prefer 2 apply (metis all-not-in-conv in-set-conv-nth length-greater-0-conv
neq0-conv not-less-zero s0a8)
  apply simp

apply(rule conjI) apply clarsimp
  apply (metis length-greater-0-conv neq0-conv not-less-zero s0a8)
  apply (metis distinct-conv-nth length-0-conv neq0-conv not-less-zero s0a8)
done

lemma s0-inv-bitmap: inv-bitmap s0
apply(simp add: inv-bitmap-def)
apply(simp add: Let-def) apply clarsimp
  apply(case-tac i = 0)
  apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def partner-bits-def)

  apply(rule conjI) apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def)

  apply(rule conjI) apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def)

  apply(rule conjI) apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def)

  apply(rule conjI) apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def)

  apply(rule conjI) apply clarsimp using s0a6 s0a8 s0a9 apply(simp add:Let-def)

  apply(case-tac i = 1) apply clarsimp using s0a6 s0a8 s0a9 apply(simp
add:Let-def)
    apply(subgoal-tac i > 1) prefer 2 apply simp
    apply(subgoal-tac get-bit-s s0 p (i - Suc NULL) (j div 4) = NOEXIST)
    prefer 2 using s0a8 apply(simp add: Let-def)
    apply(subgoal-tac j div 4 < length (bits (levels (mem-pool-info s0 p) ! (i
- 1)))))
      prefer 2 using s0a6 apply(simp add:Let-def)
      apply(subgoal-tac n-max (mem-pool-info s0 p) > 0)

```

```

      prefer 2 using s0a6 apply(simp add:Let-def)
      apply(simp add: power-eq-if)
      apply auto[1]
    apply simp
  done

lemma s0-inv-bitmap-not4free: inv-bitmap-not4free s0
  apply(simp add: inv-bitmap-not4free-def)
  apply(simp add: Let-def) apply clarsimp
  using s0a6 s0a8 s0a9 apply(simp add:Let-def partner-bits-def)
done

lemma s0-inv-aux-vars: inv-aux-vars s0
  apply(simp add: inv-aux-vars-def Let-def)
  apply(rule conjI) apply (simp add: s0a4)
  apply(rule conjI) apply clarify using s0a11 apply auto[1]
  apply(rule conjI) apply (simp add: s0a10)
  apply(rule conjI) apply clarify using s0a12 apply auto[1]
  apply(rule conjI) apply (simp add: s0a4)
  apply(rule conjI) apply (simp add: s0a10)
  apply (simp add: s0a4 s0a10)
done

lemma s0-inv-bitmap-freelist0: inv-bitmap0 s0
  apply(simp add: inv-bitmap0-def Let-def)
  using s0a9 apply(simp add:Let-def)
done

lemma s0-inv-bitmap-freelistn: inv-bitmapn s0
  apply(simp add: inv-bitmapn-def Let-def)
  using s0a8 apply(simp add:Let-def) apply clarify
  apply(subgoal-tac get-bit-s s0 p (length (levels (mem-pool-info s0 p)) - Suc 0) i
= NOEXIST)
  prefer 2 apply(subgoal-tac length (levels (mem-pool-info s0 p)) > 0)
  prefer 2 using s0a6 apply(simp add:Let-def) apply auto[1]
  using s0a6 apply(simp add:Let-def) apply auto[1]
  apply simp
done

lemma s0-inv: inv s0
  apply(unfold inv-def)
  apply(rule conjI) using s0-inv-cur apply fast
  apply(rule conjI) using s0-inv-thdwaitq apply fast
  apply(rule conjI) using s0-inv-mempool-info apply fast
  apply(rule conjI) using s0-inv-bitmap-freelist apply fast
  apply(rule conjI) using s0-inv-bitmap apply fast
  apply(rule conjI) using s0-inv-aux-vars apply fast
  apply(rule conjI) using s0-inv-bitmap-freelist0 apply fast

```

**apply**(*rule conjI*) **using** *s0-inv-bitmap-freelistn* **apply** *fast*  
**using** *s0-inv-bitmap-not4free* **apply** *fast*  
**done**

### 3.3 lemmas of invariants

**lemma** *inv-bitmap-presv-setbit-0*:

$\neg (x = l \wedge y = b) \implies$   

$$Vb = Va(\text{mem-pool-info} := (\text{mem-pool-info } Va)$$

$$(p := \text{mem-pool-info } Va \text{ } p$$

$$(\text{levels} := \text{levels } (\text{mem-pool-info } Va \text{ } p)$$

$$[l := (\text{levels } (\text{mem-pool-info } Va \text{ } p) \text{ } ! l)(\text{bits} := \text{bits } (\text{levels}$$

$$(\text{mem-pool-info } Va \text{ } p) \text{ } ! l)[b := st])))) \implies$$

$$\text{get-bit-s } Va \text{ } p \text{ } x \text{ } y = \text{get-bit-s } Vb \text{ } p \text{ } x \text{ } y$$
**apply** *simp* **by** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.simps(1)* *Mem-pool-lvl.simps(4)*)

*Mem-pool-lvl.surjective list-update-beyond not-less nth-list-update-eq*  
*nth-list-update-neq*)

**lemma** *inv-bitmap-presv-setbit*:

*inv-bitmap* *Va*  $\implies$   

$$\text{get-bit-s } Va \text{ } p \text{ } l \text{ } b = \text{FREE} \vee \text{get-bit-s } Va \text{ } p \text{ } l \text{ } b = \text{FREEING} \vee \text{get-bit-s } Va \text{ } p \text{ } l \text{ } b$$

$$= \text{ALLOCATED}$$

$$\vee \text{get-bit-s } Va \text{ } p \text{ } l \text{ } b = \text{ALLOCATING} \implies$$

$$st = \text{FREE} \vee st = \text{FREEING} \vee st = \text{ALLOCATED} \vee st = \text{ALLOCATING}$$

$$\implies$$

$$Vb = \text{set-bit-s } Va \text{ } p \text{ } l \text{ } b \text{ } st \implies$$

$$\text{inv-bitmap } Vb$$
**apply**(*simp* *add:inv-bitmap-def*) **apply**(*simp* *add:set-bit-s-def* *set-bit-def*)

**apply**(*simp* *add:Let-def*) **apply** *clarify* **apply**(*rename-tac ii jj*)  
**apply**(*subgoal-tac* *p ∈ mem-pools Va*) **prefer** 2 **apply**(*simp* *add:set-bit-s-def* *set-bit-def*)  
**apply**(*subgoal-tac* *jj < length (bits (levels (mem-pool-info Va p) ! ii))*)  
**prefer** 2 **apply**(*simp* *add:set-bit-s-def* *set-bit-def*)  
**apply** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.simps(1)* *Mem-pool-lvl.simps(4)*)  
*Mem-pool-lvl.surjective list-updt-samelen nth-list-update-eq nth-list-update-neq*)

**apply**(*rule conjI*) **apply** *clarify* **apply**(*rule conjI*) **apply** *clarify*

**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div 4) =*  
*DIVIDED*)  
**prefer** 2 **apply** (*smt* *Mem-pool-lvl.simps(1)* *Mem-pool-lvl.simps(4)* *Mem-pool-lvl.surjective*

*One-nat-def nth-list-update-eq nth-list-update-neq*)  
**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div 4)*  
 $= (\text{bits } (\text{levels } (\text{mem-pool-info } Vb \text{ } p) \text{ } ! (ii - 1))) \text{ } ! (jj \text{ div } 4))$ )  
**prefer** 2 **apply**(*case-tac* *ii - 1 = l ∧ jj div 4 = b*) **apply** *simp* **using** *inv-bitmap-presv-setbit-0*

```

apply simp
apply simp

apply clarify apply(rule conjI)

apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4) = NOEXIST)
prefer 2 apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

      nth-list-update-eq nth-list-update-neq)
apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4)
      = (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj * 4))
prefer 2 apply(case-tac Suc ii = l ∧ jj * 4 = b) apply simp using inv-bitmap-presv-setbit-0
apply simp
apply simp

apply(rule conjI)
apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 1) =
NOEXIST)
prefer 2 apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

      Nat.add-0-right One-nat-def add-Suc-right nth-list-update-eq nth-list-update-neq)

apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 1)
      = (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj * 4 + 1))
prefer 2 apply(case-tac Suc ii = l ∧ jj * 4 + 1 = b) apply simp using
inv-bitmap-presv-setbit-0 apply metis
apply simp

apply(rule conjI)
apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 2) =
NOEXIST)
prefer 2
apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

      add-2-eq-Suc' nth-list-update-eq nth-list-update-neq)
apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 2)
      = (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj * 4 + 2))
prefer 2 apply(case-tac Suc ii = l ∧ jj * 4 + 2 = b) apply simp using
inv-bitmap-presv-setbit-0 apply metis
apply simp

apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 3) =
NOEXIST)
prefer 2
apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

      add-2-eq-Suc' nth-list-update-eq nth-list-update-neq)
apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj * 4 + 3)

```

$$= (\text{bits } (\text{levels } (\text{mem-pool-info } Vb \ p) \ ! \ \text{Suc } ii)) \ ! \ (jj * 4 + 3))$$
**prefer 2 apply**(*case-tac* *Suc ii = l ∧ jj \* 4 + 3 = b*) **apply simp using**  
*inv-bitmap-presv-setbit-0* **apply metis**  
**apply simp**

**apply**(*rule conjI*) **apply clarify** **apply**(*rule conjI*) **apply clarify**

**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! (ii - 1)) ! (jj div 4) =*  
*DIVIDED*))  
**prefer 2 apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

$$\text{One-nat-def } nth\text{-list-update-eq } nth\text{-list-update-neq}$$
**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! (ii - 1)) ! (jj div 4)*  

$$= (\text{bits } (\text{levels } (\text{mem-pool-info } Vb \ p) \ ! \ (ii - 1))) \ ! \ (jj \ \text{div } 4))$$
**prefer 2 apply**(*case-tac* *ii - 1 = l ∧ jj div 4 = b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply clarify** **apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! Suc ii) ! (jj \* 4) = NOEX-*  
*IST*))  
**prefer 2 apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

$$nth\text{-list-update-eq } nth\text{-list-update-neq}$$
**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! Suc ii) ! (jj \* 4)*  

$$= (\text{bits } (\text{levels } (\text{mem-pool-info } Vb \ p) \ ! \ \text{Suc } ii)) \ ! \ (jj * 4))$$
**prefer 2 apply**(*case-tac* *Suc ii = l ∧ jj \* 4 = b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply**(*rule conjI*)  
**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! Suc ii) ! (jj \* 4 + 1) =*  
*NOEXIST*))  
**prefer 2 apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

*Nat.add-0-right One-nat-def add-Suc-right nth-list-update-eq nth-list-update-neq*

**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! Suc ii) ! (jj \* 4 + 1)*  

$$= (\text{bits } (\text{levels } (\text{mem-pool-info } Vb \ p) \ ! \ \text{Suc } ii)) \ ! \ (jj * 4 + 1))$$
**prefer 2 apply**(*case-tac* *Suc ii = l ∧ jj \* 4 + 1 = b*) **apply simp using**  
*inv-bitmap-presv-setbit-0* **apply metis**  
**apply simp**

**apply**(*rule conjI*)  
**apply**(*subgoal-tac* (*bits (levels (mem-pool-info Va p) ! Suc ii) ! (jj \* 4 + 2) =*  
*NOEXIST*))  
**prefer 2**

**apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

*add-2-eq-Suc' nth-list-update-eq nth-list-update-neg*)  
**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj \* 4 + 2)*  
*= (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj \* 4 + 2))*)  
**prefer 2 apply**(*case-tac Suc ii = l ∧ jj \* 4 + 2 = b*) **apply simp using**  
*inv-bitmap-presv-setbit-0* **apply metis**  
**apply simp**

**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj \* 4 + 3) =*  
*NOEXIST*)  
**prefer 2**  
**apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

*add-2-eq-Suc' nth-list-update-eq nth-list-update-neg*)  
**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj \* 4 + 3)*  
*= (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj \* 4 + 3))*)  
**prefer 2 apply**(*case-tac Suc ii = l ∧ jj \* 4 + 3 = b*) **apply simp using**  
*inv-bitmap-presv-setbit-0* **apply metis**  
**apply simp**

**apply**(*rule conjI*) **apply clarify** **apply**(*rule conjI*) **apply clarify**

**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div 4) =*  
*DIVIDED*)  
**prefer 2 apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

*One-nat-def nth-list-update-eq nth-list-update-neg*)  
**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div 4)*  
*= (bits (levels (mem-pool-info Vb p) ! (ii - 1))) ! (jj div 4))*)  
**prefer 2 apply**(*case-tac ii - 1 = l ∧ jj div 4 = b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply clarify** **apply**(*rule conjI*)

**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj \* 4) = NOEX-*  
*IST*)  
**prefer 2 apply** (*smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*

*nth-list-update-eq nth-list-update-neg*)  
**apply**(*subgoal-tac (bits (levels (mem-pool-info Va p) ! Suc ii)) ! (jj \* 4)*  
*= (bits (levels (mem-pool-info Vb p) ! Suc ii)) ! (jj \* 4))*)  
**prefer 2 apply**(*case-tac Suc ii = l ∧ jj \* 4 = b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 1) = *NOEXIST*)

**prefer** 2 **apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*Nat.add-0-right One-nat-def add-Suc-right nth-list-update-eq nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 1)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 1))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l* ∧ *jj* \* 4 + 1 = *b*) **apply** *simp* **using**  
*inv-bitmap-presv-setbit-0* **apply** *metis*  
**apply** *simp*

**apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 2) = *NOEXIST*)

**prefer** 2

**apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*add-2-eq-Suc' nth-list-update-eq nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 2)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 2))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l* ∧ *jj* \* 4 + 2 = *b*) **apply** *simp* **using**  
*inv-bitmap-presv-setbit-0* **apply** *metis*  
**apply** *simp*

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 3) = *NOEXIST*)

**prefer** 2

**apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*add-2-eq-Suc' nth-list-update-eq nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 3)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 3))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l* ∧ *jj* \* 4 + 3 = *b*) **apply** *simp* **using**  
*inv-bitmap-presv-setbit-0* **apply** *metis*  
**apply** *simp*

**apply**(*rule conjI*) **apply** *clarify* **apply**(*rule conjI*) **apply** *clarify*

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* - 1))) ! (*jj* div 4) = *DIVIDED*)

**prefer** 2 **apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*One-nat-def nth-list-update-eq nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* - 1))) ! (*jj* div 4)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! (*ii* - 1))) ! (*jj* div 4))

**prefer** 2 **apply**(*case-tac* *ii* - 1 = *l* ∧ *jj* div 4 = *b*) **apply** *simp* **using** *inv-bitmap-presv-setbit-0*  
**apply** *simp*

**apply** *simp*

**apply** *clarify* **apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4) = *NOEXIST*)

**prefer** 2 **apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective* *nth-list-update-eq* *nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l*  $\wedge$  *jj* \* 4 = *b*) **apply** *simp* **using** *inv-bitmap-presv-setbit-0*

**apply** *simp*

**apply** *simp*

**apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 1) = *NOEXIST*)

**prefer** 2 **apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*Nat.add-0-right* *One-nat-def* *add-Suc-right* *nth-list-update-eq* *nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 1)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 1))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l*  $\wedge$  *jj* \* 4 + 1 = *b*) **apply** *simp* **using** *inv-bitmap-presv-setbit-0*

**apply** *metis*

**apply** *simp*

**apply**(*rule conjI*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 2) = *NOEXIST*)

**prefer** 2

**apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*add-2-eq-Suc'* *nth-list-update-eq* *nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 2)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 2))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l*  $\wedge$  *jj* \* 4 + 2 = *b*) **apply** *simp* **using** *inv-bitmap-presv-setbit-0*

**apply** *metis*

**apply** *simp*

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 3) = *NOEXIST*)

**prefer** 2

**apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*add-2-eq-Suc'* *nth-list-update-eq* *nth-list-update-neq*)

**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! *Suc ii*)) ! (*jj* \* 4 + 3)  
= (*bits* (*levels* (*mem-pool-info* *Vb p*) ! *Suc ii*)) ! (*jj* \* 4 + 3))

**prefer** 2 **apply**(*case-tac* *Suc ii* = *l*  $\wedge$  *jj* \* 4 + 3 = *b*) **apply** *simp* **using**



*inv-bitmap-presv-setbit-0* **apply metis**  
**apply simp**

**apply**(*rule conjI*)

**apply clarify**  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* - 1))) ! (*jj* div 4) =  
 DIVIDED)  
**prefer 2 apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*One-nat-def nth-list-update-eq nth-list-update-neq*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* - 1))) ! (*jj* div 4)  
 = (*bits* (*levels* (*mem-pool-info* *Vb p*) ! (*ii* - 1))) ! (*jj* div 4))  
**prefer 2 apply**(*case-tac* *ii* - 1 = *l*  $\wedge$  *jj* div 4 = *b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply**(*rule conjI*)

**apply clarify**  
**apply**(*rule conjI*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* + 1))) ! (*jj* \* 4) =  
 NOEXIST)  
**prefer 2 apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*Nat.add-0-right One-nat-def add-Suc-right nth-list-update-eq*  
*nth-list-update-neq*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* + 1))) ! (*jj* \* 4)  
 = (*bits* (*levels* (*mem-pool-info* *Vb p*) ! (*ii* + 1))) ! (*jj* \* 4))  
**prefer 2 apply**(*case-tac* *ii* + 1 = *l*  $\wedge$  *jj* \* 4 = *b*) **apply simp using** *inv-bitmap-presv-setbit-0*  
**apply simp**  
**apply simp**

**apply**(*rule conjI*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* + 1))) ! (*jj* \* 4 + 1) =  
 NOEXIST)  
**prefer 2 apply** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*

*Nat.add-0-right One-nat-def add-Suc-right nth-list-update-eq*  
*nth-list-update-neq*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* + 1))) ! (*jj* \* 4 + 1)  
 = (*bits* (*levels* (*mem-pool-info* *Vb p*) ! (*ii* + 1))) ! (*jj* \* 4 + 1))  
**prefer 2 apply**(*case-tac* *ii* + 1 = *l*  $\wedge$  *jj* \* 4 + 1 = *b*) **apply auto[1] using**  
*inv-bitmap-presv-setbit-0* **apply simp**  
**apply simp**

**apply**(*rule conjI*)  
**apply**(*subgoal-tac* (*bits* (*levels* (*mem-pool-info* *Va p*) ! (*ii* + 1))) ! (*jj* \* 4 + 2) =

```

NOEXIST)
  prefer 2 apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

    Nat.add-0-right One-nat-def add-2-eq-Suc' add-Suc-right nth-list-update-eq
    nth-list-update-neq)
  apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii + 1))) ! (jj * 4 + 2)
    = (bits (levels (mem-pool-info Vb p) ! (ii + 1))) ! (jj * 4 + 2))
  prefer 2 apply(case-tac ii + 1 = l ∧ jj * 4 + 2 = b) apply auto[1] using
  inv-bitmap-presv-setbit-0 apply simp
  apply simp

  apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii + 1))) ! (jj * 4 + 3) =
  NOEXIST)
  prefer 2 apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

    Nat.add-0-right One-nat-def add-2-eq-Suc' add-Suc-right nth-list-update-eq
    nth-list-update-neq)
  apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii + 1))) ! (jj * 4 + 3)
    = (bits (levels (mem-pool-info Vb p) ! (ii + 1))) ! (jj * 4 + 3))
  prefer 2 apply(case-tac ii + 1 = l ∧ jj * 4 + 3 = b) apply auto[1] using
  inv-bitmap-presv-setbit-0 apply simp
  apply simp

  apply clarify

  apply(subgoal-tac bits (levels (mem-pool-info Va p) ! ii) ! jj = NOEXIST)
  prefer 2 apply(case-tac ii = l ∧ jj = b) apply auto[1] using inv-bitmap-presv-setbit-0
  apply simp
  apply(subgoal-tac bits (levels (mem-pool-info Va p) ! (ii - 1)) ! (jj div 4) ≠
  DIVIDED)
  prefer 2 apply (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

    One-nat-def nth-list-update-eq nth-list-update-neq)
  apply(subgoal-tac (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div 4)
    = bits (levels (mem-pool-info Va p)
      [l := (levels (mem-pool-info Va p) ! l) (bits := bits (levels
      (mem-pool-info Va p) ! l) [b := st])] !
      (ii - 1)) ! (jj div 4)) prefer 2
  apply(case-tac ii - 1 = l ∧ jj div 4 = b) apply clarsimp
  apply(subgoal-tac bits (levels (mem-pool-info Va p) ! (ii - Suc NULL)) [jj div
  4 := st] ! (jj div 4) = st)
  prefer 2 apply (metis list-update-beyond not-less nth-list-update-eq)
  apply simp
  using inv-bitmap-presv-setbit-0 apply simp
  apply clarsimp

  done

```

**lemma** *inv-bitmap-freelist-fl-bnum-in*:

*inv-bitmap-freelist*  $Va \implies$   
*inv-mempool-info*  $Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $jj < \text{length } (\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) \implies$   
 $\text{block-num } (\text{mem-pool-info } Va \ p)$   
 $((\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) ! jj)$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge ii) < \text{length } (\text{bits } (\text{levels}$   
 $(\text{mem-pool-info } Va \ p) ! ii))$   
**apply**(*simp add:inv-bitmap-freelist-def inv-mempool-info-def block-num-def Let-def*)  
**apply**(*subgoal-tac*  $\exists n. n < n\text{-max } (\text{mem-pool-info } Va \ p) * (4 \wedge ii) \wedge \text{free-list } (\text{levels}$   
 $(\text{mem-pool-info } Va \ p) ! ii) ! jj$   
 $= \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge$   
 $ii))$ )  
**prefer** 2 **apply** *blast*  
**apply**(*subgoal-tac*  $\text{free-list } (\text{levels } (\text{mem-pool-info } Va \ p) ! ii) ! jj \geq \text{buf } (\text{mem-pool-info}$   
 $Va \ p)$ )  
**prefer** 2 **apply** *linarith*  
**using** *nonzero-mult-div-cancel-right by force*

**lemma** *inv-bitmap-freelist-fl-FREE*:

*inv-bitmap-freelist*  $Va \implies$   
*inv-mempool-info*  $Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $jj < \text{length } (\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) \implies$   
 $\text{get-bit-s } Va \ p \ ii \ (\text{block-num } (\text{mem-pool-info } Va \ p))$   
 $((\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) ! jj)$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge ii) = \text{FREE}$   
**apply**(*simp add:inv-bitmap-freelist-def inv-mempool-info-def block-num-def Let-def*)  
**apply**(*subgoal-tac*  $\exists n. n < n\text{-max } (\text{mem-pool-info } Va \ p) * (4 \wedge ii) \wedge \text{free-list } (\text{levels}$   
 $(\text{mem-pool-info } Va \ p) ! ii) ! jj$   
 $= \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge$   
 $ii))$ )  
**prefer** 2 **apply** *blast*  
**by** (*metis add-diff-cancel-left' div-0 mult-0-right neq0-conv nonzero-mult-div-cancel-right*  
*not-less-zero nth-mem*)

**lemma** *inv-buf-le-fl*:

*inv-bitmap-freelist*  $Va \implies$   
*inv-mempool-info*  $Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $jj < \text{length } (\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) \implies$   
 $\text{buf } (\text{mem-pool-info } Va \ p) \leq (\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) ! ii)) ! jj$   
**apply**(*simp add:inv-bitmap-freelist-def inv-mempool-info-def Let-def*)

**apply**(*subgoal-tac*  $\exists n. \text{free-list } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ ii) \ ! \ jj$   
 $= \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^ \wedge$   
 $ii))$   
**prefer** 2 **apply** *blast*  
**by** *linarith*

**lemma** *inv-fl-mod-sz0*:  
*inv-bitmap-freelist*  $Va \implies$   
*inv-mempool-info*  $Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $jj < \text{length } (\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) \ ! \ ii)) \implies$   
 $((\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) \ ! \ ii)) \ ! \ jj - \text{buf } (\text{mem-pool-info } Va \ p))$   
 $\text{mod}$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^ \wedge \ ii) = 0$   
**apply**(*simp add:inv-bitmap-freelist-def inv-mempool-info-def Let-def*)  
**apply**(*subgoal-tac*  $\exists n. \text{free-list } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ ii) \ ! \ jj$   
 $= \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^ \wedge$   
 $ii))$   
**prefer** 2 **apply** *blast*  
**by** *force*

**lemma** *sameinfo-inv-bitmap-mp*:  
 $\text{mem-pool-info } Va \ p = \text{mem-pool-info } Vb \ p \implies \text{inv-bitmap-mp } Va \ p = \text{inv-bitmap-mp}$   
 $Vb \ p$   
**apply**(*simp add: Let-def*)  
**done**

**lemma** *sameinfo-inv-bitmap-freelist-mp*:  
 $\text{mem-pool-info } Va \ p = \text{mem-pool-info } Vb \ p \implies \text{inv-bitmap-freelist-mp } Va \ p =$   
 $\text{inv-bitmap-freelist-mp } Vb \ p$   
**apply**(*simp add: Let-def*)  
**done**

**lemma** *inv-bitmap-presv-mpls-mpi*:  
 $\text{mem-pools } Va = \text{mem-pools } Vb \implies$   
 $\text{mem-pool-info } Va = \text{mem-pool-info } Vb \implies$   
 $\text{inv-bitmap } Va \implies$   
 $\text{inv-bitmap } Vb$   
**by**(*simp add:inv-bitmap-def Let-def*)

**lemma** *inv-bitmap-presv-mpls-mpi2*:  
 $\text{mem-pools } Va = \text{mem-pools } Vb \implies$   
 $(\forall p. \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) = \text{length } (\text{levels } (\text{mem-pool-info } Vb$   
 $p))) \implies$   
 $(\forall p \ ii. \ ii < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p))$   
 $\longrightarrow \text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ ii) = \text{bits } (\text{levels } (\text{mem-pool-info } Vb$

$p) ! ii)) \implies$   
 $\text{inv-bitmap } Va \implies$   
 $\text{inv-bitmap } Vb$   
**by** (simp add: inv-bitmap-def Let-def)

**lemma** inv-bitmap-freeing2free:

$\text{inv-bitmap-mp } V p \implies$   
 $\exists lv \ bl. \text{ bits } (\text{levels } (\text{mem-pool-info } V p) ! lv) ! bl = \text{FREEING}$   
 $\wedge \text{ bits } (\text{levels } (\text{mem-pool-info } V2 p) ! lv) = \text{bits } (\text{levels } (\text{mem-pool-info } V$   
 $p) ! lv) [bl := \text{FREE}]$   
 $\wedge (\forall lv'. lv \neq lv' \longrightarrow \text{ bits } (\text{levels } (\text{mem-pool-info } V2 p) ! lv') = \text{ bits } (\text{levels}$   
 $(\text{mem-pool-info } V p) ! lv')) \implies$   
 $\text{length } (\text{levels } (\text{mem-pool-info } V p)) = \text{length } (\text{levels } (\text{mem-pool-info } V2 p))$   
 $\implies \text{inv-bitmap-mp } V2 p$   
**apply**(simp add:Let-def) **apply** clarify  
**apply**(subgoal-tac length (bits (levels (mem-pool-info V p) ! i)) = length (bits  
(levels (mem-pool-info V2 p) ! i)))  
**prefer** 2 **apply**(case-tac i = lv) **apply** auto[1] **apply** auto[1]  
**apply**(rule conjI) **apply** clarify  
**apply**(rule conjI) **apply** clarify  
**apply**(case-tac i = lv  $\wedge$  j = bl) **apply** clarsimp  
**apply**(subgoal-tac get-bit-s V2 p i j = get-bit-s V p i j)  
**prefer** 2 **apply**(case-tac i = lv) **apply** clarsimp **apply** presburger  
**apply** (metis BlockState.distinct(21) nth-list-update-neq)  
  
**apply** clarify  
**apply**(case-tac i = lv  $\wedge$  j = bl) **apply** clarsimp  
**apply**(subgoal-tac get-bit-s V2 p i j = get-bit-s V p i j)  
**prefer** 2 **apply**(case-tac i = lv) **apply** clarsimp **apply** presburger  
**apply** (smt BlockState.distinct(25) nth-list-update-neq)  
  
**apply**(rule conjI) **apply** clarify  
**apply**(case-tac i = lv  $\wedge$  j = bl) **apply** clarsimp  
**apply**(subgoal-tac get-bit-s V2 p i j = get-bit-s V p i j)  
**prefer** 2 **apply**(case-tac i = lv) **apply** clarsimp **apply** presburger  
**apply** (smt BlockState.distinct(21) BlockState.distinct(25) nth-list-update-neq)  
  
**apply**(rule conjI) **apply** clarify  
**apply**(case-tac i = lv  $\wedge$  j = bl) **apply** clarsimp  
**apply**(subgoal-tac get-bit-s V2 p i j = get-bit-s V p i j)  
**prefer** 2 **apply**(case-tac i = lv) **apply** clarsimp **apply** presburger  
**apply** (smt BlockState.distinct(21) BlockState.distinct(25) nth-list-update-neq)

```

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(21) BlockState.distinct(25) nth-list-update-neq)

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(21) BlockState.distinct(25) nth-list-update-neq)

apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(11) list-update-beyond not-less nth-list-update-eq
nth-list-update-neq)

done

lemma inv-bitmap-allocating2allocate:
  inv-bitmap-mp  $V\ p \implies$ 
     $\exists\ lv\ bl. \text{bits}(\text{levels}(\text{mem-pool-info}\ V\ p) ! lv) ! bl = \text{ALLOCATING}$ 
       $\wedge \text{bits}(\text{levels}(\text{mem-pool-info}\ V2\ p) ! lv) = \text{bits}(\text{levels}(\text{mem-pool-info}\ V$ 
 $p) ! lv) [bl := \text{ALLOCATED}]$ 
       $\wedge (\forall\ lv'. lv \neq lv' \longrightarrow \text{bits}(\text{levels}(\text{mem-pool-info}\ V2\ p) ! lv') = \text{bits}(\text{levels}$ 
 $(\text{mem-pool-info}\ V\ p) ! lv')) \implies$ 
       $\text{length}(\text{levels}(\text{mem-pool-info}\ V\ p)) = \text{length}(\text{levels}(\text{mem-pool-info}\ V2\ p))$ 
       $\implies \text{inv-bitmap-mp}\ V2\ p$ 
  apply(simp add:Let-def) apply clarify
  apply(subgoal-tac  $\text{length}(\text{bits}(\text{levels}(\text{mem-pool-info}\ V\ p) ! i)) = \text{length}(\text{bits}$ 
 $(\text{levels}(\text{mem-pool-info}\ V2\ p) ! i)))$ 
    prefer 2 apply(case-tac  $i = lv$ ) apply auto[1] apply auto[1]
  apply(rule conjI) apply clarify
    apply(rule conjI) apply clarify
      apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
        apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
          prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
        apply (metis BlockState.distinct(23) nth-list-update-neq)

    apply clarify
      apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
        apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
          prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
        apply (smt BlockState.distinct(27) nth-list-update-neq)

  apply(rule conjI) apply clarify

```

```

apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
  apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
    prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
  apply (smt BlockState.distinct(23) BlockState.distinct(27) nth-list-update-neq)

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(23) BlockState.distinct(27) nth-list-update-neq)

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(23) BlockState.distinct(27) nth-list-update-neq)

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(23) BlockState.distinct(27) nth-list-update-neq)

apply(rule conjI) apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(23) BlockState.distinct(27) nth-list-update-neq)

apply clarify
  apply(case-tac  $i = lv \wedge j = bl$ ) apply clarsimp
    apply(subgoal-tac get-bit-s  $V2\ p\ i\ j = \text{get-bit-s}\ V\ p\ i\ j$ )
      prefer 2 apply(case-tac  $i = lv$ ) apply clarsimp apply presburger
    apply (smt BlockState.distinct(3) list-update-beyond not-less nth-list-update-eq
nth-list-update-neq)
done

```

```

lemma inv-bitmap-freelist-presv-setbit-notfree-h:
 $\neg (x = lv \wedge y = bkn) \implies$ 
   $Vb = \text{set-bit-s}\ V\ p\ lv\ bkn\ st \implies$ 
   $\text{get-bit-s}\ V\ p\ x\ y = \text{get-bit-s}\ Vb\ p\ x\ y$ 
apply(simp add:set-bit-s-def set-bit-def)
by (metis (no-types, lifting) Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)
  Mem-pool-lvl.surjective list-update-beyond not-less nth-list-update-eq
nth-list-update-neq)

```

```

lemma inv-bitmap-freelist-presv-setbit-notfree:

```

$p \in \text{mem-pools } V \implies$   
 $\text{inv-mempool-info } V \wedge \text{inv-aux-vars } V \wedge \text{inv-bitmap-freelist } V \implies$   
 $st \neq \text{FREE} \implies$   
 $\text{get-bit-s } V p \text{ lv bkn} \neq \text{FREE} \implies$   
 $\text{inv-bitmap-freelist } (\text{set-bit-s } V p \text{ lv bkn } st)$   
**apply**(simp add:inv-bitmap-freelist-def) **apply**(simp add:set-bit-s-def set-bit-def)  
**apply**(simp add:Let-def)  
**apply** clarsimp **apply**(rename-tac ii)  
**apply**(rule conjI)  
**apply** clarsimp **apply**(rename-tac jj)  
**apply**(subgoal-tac length (bits (levels (mem-pool-info V p)  
 $[lv := (\text{levels } (\text{mem-pool-info } V p) ! lv)(\text{bits} := \text{bits } (\text{levels}$   
 $(\text{mem-pool-info } V p) ! lv)[bkn := st]]) ! ii)$   
 $= \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V p) ! ii)))$  **prefer** 2  
**apply**(case-tac ii = lv) **apply** fastforce **apply** fastforce  
**apply**(case-tac ii = lv  $\wedge$  jj = bkn)  
**using** inv-bitmap-freelist-presv-setbit-notfree-h **apply** force  
**apply**(subgoal-tac (bits (levels (mem-pool-info V p)[lv := (levels (mem-pool-info  
V p) ! lv)  
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V p) ! lv)[bkn := st]]) !$   
ii)) ! jj =  
 $(\text{bits } (\text{levels } (\text{mem-pool-info } V p) ! ii)) ! jj)$  **prefer** 2  
**apply**(case-tac ii  $\neq$  lv) **apply** fastforce  
**apply**(case-tac jj  $\neq$  bkn) **apply** fastforce  
**apply** fastforce  
**apply**(subgoal-tac free-list (levels (mem-pool-info V p)[lv := (levels (mem-pool-info  
V p) ! lv)  
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V p) ! lv)[bkn := st]]) ! ii)$   
 $= \text{free-list } (\text{levels } (\text{mem-pool-info } V p) ! ii))$  **prefer** 2  
**apply**(case-tac ii  $\neq$  lv) **apply** fastforce **apply** fastforce  
**apply** auto[1]

**apply**(rule conjI)  
**apply** clarsimp **apply**(rename-tac jj)  
**apply**(subgoal-tac length (bits (levels (mem-pool-info V p)  
 $[lv := (\text{levels } (\text{mem-pool-info } V p) ! lv)(\text{bits} := \text{bits } (\text{levels}$   
 $(\text{mem-pool-info } V p) ! lv)[bkn := st]]) ! ii)$   
 $= \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V p) ! ii)))$  **prefer** 2  
**apply**(case-tac ii = lv) **apply** fastforce **apply** fastforce  
**apply**(subgoal-tac free-list (levels (mem-pool-info V p)[lv := (levels (mem-pool-info  
V p) ! lv)  
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V p) ! lv)[bkn := st]]) ! ii)$   
 $= \text{free-list } (\text{levels } (\text{mem-pool-info } V p) ! ii))$  **prefer** 2  
**apply**(case-tac ii  $\neq$  lv) **apply** fastforce **apply** fastforce  
**apply** auto[1]

**apply**(subgoal-tac free-list (levels (mem-pool-info V p)[lv := (levels (mem-pool-info  
V p) ! lv)



```

      (|bits := bits (levels (mem-pool-info V p) ! lv)[bkn := st]|) ! ii)
    = free-list (levels (mem-pool-info V p) ! ii) prefer 2
  apply(case-tac ii ≠ lv) apply fastforce apply fastforce
  apply auto[1]
done

end

```

```

theory memory-cover
imports invariant
begin

```

## 4 partition of memory addresses of a pool

```

declare [[smt-timeout = 300]]

```

this theory shows that all memory blocks are a COVER of address space of a memory pool. A COVER means blocks are disjoint and continuous. It means that for any memory address of a memory pool, the address is in the address range of only one block.

Due to algorithm, address range of each block is implicitly derived. address range of a block at level  $ii$  with block index  $jj$  at this level is  $jj * (\text{max\_sz mp div } (4^i i)) \dots jj * (\text{max\_sz mp div } (4^i i))$ .

**abbreviation** *addr-in-block mp addr ii jj*  $\equiv$   
 $ii < \text{length } (\text{levels mp}) \wedge jj < \text{length } (\text{bits } (\text{levels mp} ! ii))$   
 $\wedge (\text{bits } (\text{levels mp} ! ii) ! jj = \text{FREE} \vee \text{bits } (\text{levels mp} ! ii) ! jj = \text{FREEING} \vee$   
 $\text{bits } (\text{levels mp} ! ii) ! jj = \text{ALLOCATED} \vee \text{bits } (\text{levels mp} ! ii) ! jj =$   
 $\text{ALLOCATING})$   
 $\wedge \text{addr} \in \{jj * (\text{max\_sz mp div } (4^i i)) \dots jj * (\text{max\_sz mp div } (4^i i))\}$

**abbreviation** *mem-cover-mp* ::  $\text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{bool}$

**where** *mem-cover-mp s p*  $\equiv$

*let mp = mem-pool-info s p in*  $(\forall \text{addr} < n\text{-max mp} * \text{max\_sz mp}. (\exists !(i,j). \text{addr-in-block mp addr } i j))$

**definition** *mem-cover* ::  $\text{State} \Rightarrow \text{bool}$

**where** *mem-cover s*  $\equiv \forall p \in \text{mem-pools } s. \text{mem-cover-mp } s p$

**lemma** *split-div-lemma*:

**assumes**  $0 < n$

**shows**  $n * q \leq m \wedge m < n * \text{Suc } q \longleftrightarrow q = ((m :: \text{nat}) \text{ div } n) \text{ (is ?lhs} \longleftrightarrow \text{?rhs)}$

**proof**

**assume** *?rhs*

**with** *minus-mod-eq-mult-div [symmetric]* **have**  $nq: n * q = m - (m \bmod n)$  **by**  
*simp*

**then have**  $A: n * q \leq m$  **by** *simp*

```

have  $n - (m \bmod n) > 0$  using mod-less-divisor assms by auto
then have  $m < m + (n - (m \bmod n))$  by simp
then have  $m < n + (m - (m \bmod n))$  by simp
with nq have  $m < n + n * q$  by simp
then have  $B: m < n * \text{Suc } q$  by simp
from  $A \ B$  show ?lhs ..
next
  assume  $P: ?lhs$ 
  then show ?rhs
    using div-nat-eqI by blast
qed

lemma align-up-ge-low:
 $sz1 > 0 \implies sz2 > sz1 \implies sz2 \bmod sz1 = 0 \implies (addr::nat) \text{div } sz2 * sz2 + sz2$ 
 $\geq addr \text{div } sz1 * sz1 + sz1$ 
  apply(subgoal-tac  $\exists n>0. sz2 = n * sz1$ ) prefer 2 apply auto[1]
  apply(rule subst[where  $s=addr - addr \bmod sz1$  and  $t=addr \text{div } sz1 * sz1$ ]])
    using minus-mod-eq-div-mult apply auto[1]
  apply(rule subst[where  $s=addr - addr \bmod sz2$  and  $t=addr \text{div } sz2 * sz2$ ]])
    using minus-mod-eq-div-mult apply auto[1]
  apply(subgoal-tac  $sz1 - addr \bmod sz1 \leq sz2 - addr \bmod sz2$ )
    apply(subgoal-tac  $addr \bmod sz1 < sz1$ ) prefer 2 apply simp
    apply(subgoal-tac  $addr \bmod sz2 < sz2$ ) prefer 2 apply simp
    apply simp

  apply clarsimp
  apply(case-tac  $addr \bmod (sz1 * n) \geq sz1 * (n - 1)$ )
    apply(subgoal-tac  $addr \bmod (sz1 * n) = sz1 * (n - 1) + addr \bmod sz1$ )
      prefer 2 using Suc-lessD Suc-pred' mod-less-divisor
        mult-div-mod-eq nat-0-less-mult-iff mod-mult-self4 split-div-lemma
      apply (metis mod-mult2-eq)
    apply clarsimp
  apply (metis (no-types, lifting) Nat.diff-diff-right One-nat-def Suc-lessD add-diff-cancel-left'

    le-add1 less-numeral-extra(3) less-or-eq-imp-le mult commute mult-eq-if)
  by (metis (no-types, lifting) Nat.le-diff-conv2 Suc-lessD add-mono-thms-linordered-semiring(1)

    diff-le-self less-numeral-extra(3) mod-le-divisor mult commute mult-eq-if
    mult-pos-pos nat-le-linear)

lemma addr-exist-block-h1-1:
 $li < ii \implies ii < nl \implies (4::nat) \wedge^{nl \text{div } 4} ii < 4 \wedge^{nl \text{div } 4} li$ 
  apply(rule subst[where  $s=4 \wedge^{nl - ii}$  and  $t=4 \wedge^{nl \text{div } 4} ii$ ]])
    apply (simp add: div2-eq-minus)
  apply(rule subst[where  $s=4 \wedge^{nl - li}$  and  $t=4 \wedge^{nl \text{div } 4} li$ ]])
    apply (simp add: div2-eq-minus)
  apply auto
done

```

**lemma** *mod-time*:  $(x::nat) \bmod m = 0 \implies n * x \bmod (n * m) = 0$

**by** *simp*

**lemma** *addr-exist-block-h1*:

$li < ii \implies$

$\exists n > 0. msz = (4 * n) * (4 ^ nl) \implies$

$ii < nl \implies$

$Suc (addr \div (msz \div 4 ^ ii)) * (msz \div 4 ^ ii)$

$\leq Suc (addr \div (msz \div 4 ^ ii) \div 4 ^ (ii - li)) * (msz \div 4 ^ li)$

**apply**(rule *subst*[**where**  $s = (addr \div (msz \div 4 ^ ii)) * (msz \div 4 ^ ii) + (msz \div 4 ^ ii)$ ])

**and**  $t = Suc (addr \div (msz \div 4 ^ ii)) * (msz \div 4 ^ ii)$  **apply**

*auto*[1]

**apply**(rule *subst*[**where**  $s = (addr \div (msz \div 4 ^ ii) \div 4 ^ (ii - li)) * (msz \div 4 ^ li) + (msz \div 4 ^ li)$ ])

**and**  $t = Suc (addr \div (msz \div 4 ^ ii) \div 4 ^ (ii - li)) * (msz \div 4 ^ li)$  **apply**

*auto*[1]

**apply**(rule *subst*[**where**  $s = addr \div (msz \div 4 ^ li)$  **and**  $t = addr \div (msz \div 4 ^ ii) \div 4 ^ (ii - li)$ ])

**apply**(rule *subst*[**where**  $s = addr \div (msz \div 4 ^ ii * 4 ^ (ii - li))$  **and**  $t = addr \div (msz \div 4 ^ ii) \div 4 ^ (ii - li)$ ])

**using** *div2-eq-divmul*[of  $addr \ msz \div 4 ^ ii \ 4 ^ (ii - li)$ ] **apply** *simp*

**apply**(rule *subst*[**where**  $s = msz \div 4 ^ li$  **and**  $t = msz \div 4 ^ ii * 4 ^ (ii - li)$ ])

**apply**(*subgoal-tac*  $msz \bmod 4 ^ ii = 0$ ) **prefer** 2

**using** *ge-pow-mod-0* **apply** *auto*[1]

**apply** (*smt add-diff-inverse-nat less-imp-le-nat mod-div-self mult.left-commute nonzero-mult-div-cancel-left not-less power-add power-not-zero*)

*rel-simps*(76))

**apply** *fast*

**apply**(rule *align-up-ge-low*[of  $msz \div 4 ^ ii \ msz \div 4 ^ li \ addr$ ])

**apply** (*metis ge-pow-mod-0 mod-div-self nat-0-less-mult-iff zero-less-numeral zero-less-power*)

**apply** *clarsimp* **apply**(*subgoal-tac*  $4 ^ nl \div 4 ^ ii < 4 ^ nl \div 4 ^ li$ )

**prefer** 2 **using** *addr-exist-block-h1-1*[of  $li \ ii \ nl$ ] **apply** *simp*

**using** *m-mod-div pow-mod-0* **apply** *auto*[1]

**apply** *clarsimp* **using** *mod-time*[of  $4 ^ nl \div 4 ^ li \ 4 ^ nl \div 4 ^ ii$ ]

**by** (*smt less-imp-add-positive mod-div-self mod-mult-self1-is-0 mult.left-commute*

*nonzero-mult-div-cancel-left power-add power-not-zero zero-neq-numeral*)

**lemma** *divornoe-imp-div-noe-neigh*:

$\forall li \leq ii. get\_bit\_s \ s \ p \ li \ (jj \div 4 ^ (ii - li)) = DIVIDED \vee get\_bit\_s \ s \ p \ li \ (jj \div 4 ^ (ii - li)) = NOEXIST \implies$

$get\_bit\_s \ s \ p \ NULL \ (jj \div 4 ^ ii) = DIVIDED \implies$

$get\_bit\_s \ s \ p \ ii \ jj = NOEXIST \implies$

$ii > 0 \implies$

$\exists n. n > 0 \wedge n \leq ii \wedge get\_bit\_s \ s \ p \ (n-1) \ (jj \div 4 ^ (ii - (n-1))) = DIVIDED \wedge$

```

      get-bit-s s p n (jj div 4 ^ (ii - n)) = NOEXIST
apply(induction ii arbitrary: jj)
apply simp

apply(case-tac get-bit-s s p ii (jj div 4) = DIVIDED)
  apply auto[1]

apply(subgoal-tac get-bit-s s p ii (jj div 4) = NOEXIST)
  prefer 2
  apply (metis One-nat-def Suc-diff-Suc diff-self-eq-0 lessI less-imp-le-nat power-one-right)

apply(case-tac ii = 0) apply auto[1]

apply(subgoal-tac  $\forall li \leq ii. \text{get-bit-s } s \text{ p } li ((jj \text{ div } 4) \text{ div } 4 ^ (ii - li)) = DIVIDED$ 
   $\vee \text{get-bit-s } s \text{ p } li ((jj \text{ div } 4) \text{ div } 4 ^ (ii - li)) = NOEXIST$ )
  prefer 2 apply clarsimp
  apply (metis Suc-diff-le div-mult2-eq le-SucI power-Suc)

apply(subgoal-tac  $\exists n > NULL. n \leq ii \wedge$ 
   $\text{get-bit-s } s \text{ p } (n - 1) ((jj \text{ div } 4) \text{ div } 4 ^ (ii - (n - 1))) =$ 
   $DIVIDED \wedge$ 
   $\text{get-bit-s } s \text{ p } n ((jj \text{ div } 4) \text{ div } 4 ^ (ii - n)) = NOEXIST$ )
  prefer 2 apply (simp add: Divides.div-mult2-eq)
proof -
  fix iia :: nat and jja :: nat
  assume  $\exists n > NULL. n \leq iia \wedge \text{get-bit-s } s \text{ p } (n - 1) (jja \text{ div } 4 \text{ div } 4 ^ (iia - (n - 1))) = DIVIDED$ 
   $\wedge \text{get-bit-s } s \text{ p } n (jja \text{ div } 4 \text{ div } 4 ^ (iia - n)) = NOEXIST$ 
  then obtain nn :: nat where
    f1:  $NULL < nn \wedge nn \leq iia \wedge \text{get-bit-s } s \text{ p } nn (jja \text{ div } 4 \text{ div } 4 ^ (iia - nn)) = NOEXIST$ 
   $NOEXIST$ 
   $\wedge \text{get-bit-s } s \text{ p } (nn - 1) (jja \text{ div } 4 \text{ div } 4 ^ (iia - (nn - 1))) = DIVIDED$ 
  by meson
  then have f2:  $\text{get-bit-s } s \text{ p } nn (jja \text{ div } 4 ^ \text{Suc } (iia - nn)) = NOEXIST$ 
  by (metis (no-types) div-mult2-eq semiring-normalization-rules(27))
  have f3:  $\text{get-bit-s } s \text{ p } (nn - 1) (jja \text{ div } 4 ^ \text{Suc } (\text{Suc } (iia - nn))) = DIVIDED$ 
  using f1 by (metis (no-types) Suc-diff-eq-diff-pred Suc-diff-le div-mult2-eq
    semiring-normalization-rules(27))
  have  $nn \leq iia \wedge NULL < nn$ 
  using f1 by meson
  then show  $\exists n > NULL. n \leq \text{Suc } iia \wedge \text{get-bit-s } s \text{ p } (n - 1) (jja \text{ div } 4 ^ (\text{Suc } iia - (n - 1))) = DIVIDED$ 
   $\wedge \text{get-bit-s } s \text{ p } n (jja \text{ div } 4 ^ (\text{Suc } iia - n)) = NOEXIST$ 
  using f3 f2 Suc-diff-le le-Suc-eq by auto
qed

```

```

lemma addr-exist-block:
assumes
  p2: inv-bitmap0 s and
  p3: inv-bitmap s and
  p6: inv-mempool-info s and
  p4:  $p \in \text{mem-pools } s$  and
  p7: inv-bitmapn s and
  p5:  $\text{addr} < n\text{-max } (\text{mem-pool-info } s \ p) * \text{max-sz } (\text{mem-pool-info } s \ p)$ 
shows  $\exists i \ j. \text{addr-in-block } (\text{mem-pool-info } s \ p) \ \text{addr } i \ j$ 
proof -
  obtain ii where ii:  $ii = \text{length } (\text{levels } (\text{mem-pool-info } s \ p)) - 1$  by auto
  obtain jj where jj:  $jj = \text{addr div } (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (4 ^ ii))$  by
auto
  have bits-len-nmax:  $\forall i < \text{length } (\text{levels } (\text{mem-pool-info } s \ p)). \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } s \ p) ! i)) = (n\text{-max } (\text{mem-pool-info } s \ p)) * 4 ^ i$ 
  using p6 p4 by (simp add:inv-mempool-info-def Let-def)
  have maxsz:  $\exists n > 0. \text{max-sz } (\text{mem-pool-info } s \ p) = (4 * n) * (4 ^ n\text{-levels } (\text{mem-pool-info } s \ p))$ 
  using p4 p6 apply (simp add:inv-mempool-info-def Let-def) by auto
  have nl-eq-len:  $n\text{-levels } (\text{mem-pool-info } s \ p) = \text{length } (\text{levels } (\text{mem-pool-info } s \ p))$ 
  using p4 p6 by (simp add:inv-mempool-info-def Let-def)
  from ii have ii-len:  $ii < \text{length } (\text{levels } (\text{mem-pool-info } s \ p))$ 
  by (metis diff-less inv-mempool-info-def length-greater-0-conv p4 p6 rel-simps(68))
  from ii p6 have blk-ii:  $\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 ^ ii > 0$ 
  by (smt Euclidean-Division.div-eq-0-iff divisors-zero grOI ii-len less-imp-le-nat m-mod-div maxsz mod-if nl-eq-len pow-mod-0 power-not-zero zero-neq-numeral)
  hence addr-ran:  $\text{addr} \in \{jj * (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (4 ^ ii)) .. < \text{Suc } jj * (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (4 ^ ii))\}$ 
  using jj div-in-suc [of  $\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 ^ ii \ jj \ \text{addr}$ ] by blast
  have jj-lt-maxdiv4ii:  $jj < n\text{-max } (\text{mem-pool-info } s \ p) * 4 ^ ii$ 
  apply (rule subst [where  $s = \text{addr div } (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 ^ ii)$ 
and  $t = jj$ ]) using jj apply fast
  apply (rule subst [where  $s = n\text{-max } (\text{mem-pool-info } s \ p) * \text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 ^ ii)$ 
and  $t = n\text{-max } (\text{mem-pool-info } s \ p) * 4 ^ ii$ ]) using ii-len maxsz
  apply (metis (no-types, lifting) blk-ii ge-pow-mod-0 inv-mempool-info-def m-mod-div
    mod-div-self mod-mult-self1-is-0 neq0-conv nonzero-mult-div-cancel-left p4
p6)
  apply (rule mod-div-gt [of  $\text{addr } n\text{-max } (\text{mem-pool-info } s \ p) * \text{max-sz } (\text{mem-pool-info } s \ p)$ 
     $\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 ^ ii$ ]) using p5 apply fast
  using maxsz nl-eq-len
  apply (metis ge-pow-mod-0 ii-len mod-div-self mod-mult-right-eq mod-mult-self1-is-0
    mult-0-right)

```

```

done
have lvlii-eq-last: levels (mem-pool-info s p) ! ii = last (levels (mem-pool-info s
p))
  apply(subgoal-tac length (levels (mem-pool-info s p)) > 0)
  prefer 2 using p4 p6 ii jj-lt-maxdiv4ii p4 p6 ii-len apply(simp add:inv-mempool-info-def
Let-def)
  using ii apply clarsimp
  by (simp add: last-conv-nth)
  have jj-lt-len-lstbits: jj < length (bits (last (levels (mem-pool-info s p))))
  using ii jj-lt-maxdiv4ii p4 p6 ii-len apply(simp add:inv-mempool-info-def
Let-def)
  apply(subgoal-tac length (bits (levels (mem-pool-info s p) ! ii)) = n-max
(mem-pool-info s p) * 4 ^ ii)
  prefer 2 apply auto[1]
  apply(subgoal-tac levels (mem-pool-info s p) ! ii = last (levels (mem-pool-info
s p)))
  prefer 2 apply(subgoal-tac length (levels (mem-pool-info s p)) > 0)
  prefer 2 using p4 p6 apply(simp add:inv-mempool-info-def Let-def) apply
clarsimp
  apply (simp add: last-conv-nth)
  by fastforce

have  $\exists li \leq ii. \text{addr-in-block } (mem\text{-pool-info } s \text{ } p) \text{ } addr \text{ } li \text{ } (jj \text{ div } 4 \wedge (ii - li))$ 
proof -
{
  assume asm:  $\neg (\exists li \leq ii. \text{addr-in-block } (mem\text{-pool-info } s \text{ } p) \text{ } addr \text{ } li \text{ } (jj \text{ div } 4 \wedge$ 
 $(ii - li)))$ 

  from asm have  $\forall li \leq ii. \neg \text{addr-in-block } (mem\text{-pool-info } s \text{ } p) \text{ } addr \text{ } li \text{ } (jj \text{ div } 4$ 
 $\wedge (ii - li))$  by fast
  moreover
  from ii have ii-len:  $ii < \text{length } (levels (mem\text{-pool-info } s \text{ } p))$ 
  by (metis diff-less inv-mempool-info-def length-greater-0-conv p4 p6 rel-simps(68))
  moreover
  have  $\forall li \leq ii. \text{addr} \in \{jj \text{ div } 4 \wedge (ii - li) * (\text{max-sz } (mem\text{-pool-info } s \text{ } p) \text{ div}$ 
 $4 \wedge li).. $$Suc (jj \text{ div } 4 \wedge (ii - li)) * (\text{max-sz } (mem\text{-pool-info } s \text{ } p) \text{ div}$ 
 $4 \wedge li)\}$ 
  apply(subgoal-tac  $\exists n > 0. \text{max-sz } (mem\text{-pool-info } s \text{ } p) = (4 * n) * (4 \wedge$ 
 $n\text{-levels } (mem\text{-pool-info } s \text{ } p))$ )
  prefer 2 using p4 p6 apply(simp add:inv-mempool-info-def Let-def)
  apply auto[1]
  apply(subgoal-tac  $n\text{-levels } (mem\text{-pool-info } s \text{ } p) = \text{length } (levels (mem\text{-pool-info}
s \text{ } p)) \wedge$ 
 $\text{length } (levels (mem\text{-pool-info } s \text{ } p)) > 0$ )
  prefer 2 using p4 p6 apply(simp add:inv-mempool-info-def Let-def)
  apply auto[1]
  apply clarify apply auto
  apply(subgoal-tac  $jj * (\text{max-sz } (mem\text{-pool-info } s \text{ } p) \text{ div } 4 \wedge ii)$$$ 
```

$\geq jj \text{ div } 4 \wedge (ii - li) * (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 \wedge li))$   
**prefer** 2 **apply**(*case-tac*  $li = ii$ ) **apply** *auto*[1]  
**using** *Divides.div-mult2-eq Groups.mult-ac*(2) *blk-ii add-diff-inverse-nat*  
*calculation*(2)  
*div-mult-self-is-m divisors-zero ge-pow-mod-0 mod-div-self neq0-conv*  
*not-less power-add*  
*semiring-normalization-rules*(17) *split-div-lemma zero-less-numeral*  
*zero-less-power*  
**apply** (*smt div-mult-self1-is-m nat-mult-le-cancel-disj*)  
**using** *addr-ran* **apply** *auto*[1]  
  
**apply**(*subgoal-tac*  $Suc \ jj * (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (4 \wedge ii))$   
 $\leq Suc \ (jj \text{ div } 4 \wedge (ii - li)) * (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } 4 \wedge li))$ )  
**prefer** 2 **apply**(*case-tac*  $li = ii$ ) **apply** *simp*  
**apply**(*rule subst*[**where**  $s = \text{addr div } (\text{max-sz } (\text{mem-pool-info } s \ p) \text{ div } (4 \wedge ii))$  and  $t = jj$ ]) **using** *jj* **apply** *fast*  
**using** *addr-exist-block-h1*[*of - ii max-sz (mem-pool-info s p) n-levels*  
*(mem-pool-info s p) addr*]  
*ii-len nl-eq-len maxsz* **apply** *fastforce*  
**using** *addr-ran ii-len* **apply** *auto*[1]  
**done**  
**moreover**  
**have** *li-len*:  $\forall li \leq ii. jj \text{ div } 4 \wedge (ii - li) < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } s \ p) ! li))$   
**apply** *clarsimp*  
**apply**(*subgoal-tac*  $\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } s \ p) ! li)) = (n\text{-max } (\text{mem-pool-info } s \ p)) * 4 \wedge li)$ )  
**prefer** 2 **using** *p4 p6 ii-len* **apply**(*simp add:inv-mempool-info-def Let-def*)  
**using** *jj maxsz nl-eq-len jj-lt-maxdiv4ii Divides.div-mult2-eq add-diff-cancel-left'*  
*blk-ii div-eq-0-iff gr-implies-not0*  
*le-Suc-ex less-not-refl2 mult.commute mult.left-commute mult-0 mult-is-0*  
*p5 power-add*  
**by** (*smt not-less*)  
**ultimately have**  $\forall li \leq ii. \neg (\text{get-bit-s } s \ p \ li \ (jj \text{ div } 4 \wedge (ii - li)) = \text{FREE} \vee$   
 $\text{get-bit-s } s \ p \ li \ (jj \text{ div } 4 \wedge (ii - li)) = \text{FREEING} \vee$   
 $\text{get-bit-s } s \ p \ li \ (jj \text{ div } 4 \wedge (ii - li)) = \text{ALLOCATED} \vee \text{get-bit-s } s \ p \ li$   
 $(jj \text{ div } 4 \wedge (ii - li)) = \text{ALLOCATING})$   
**by** *auto*  
**hence** *all-dv-ne*:  $\forall li \leq ii. \text{get-bit-s } s \ p \ li \ (jj \text{ div } 4 \wedge (ii - li)) = \text{DIVIDED} \vee$   
 $\text{get-bit-s } s \ p \ li \ (jj \text{ div } 4 \wedge (ii - li)) = \text{NOEXIST}$   
**using** *BlockState.exhaust* **by** *blast*  
**moreover**  
**have** *bit-lvl0*:  $\text{get-bit-s } s \ p \ 0 \ (jj \text{ div } 4 \wedge ii) = \text{DIVIDED}$  **using** *all-dv-ne p2*  
*p4* **apply**(*simp add:inv-bitmap0-def Let-def*)  
**using** *li-len* **by** *fastforce*  
**moreover**  
**have** *bit-lvln*:  $\text{get-bit-s } s \ p \ ii \ jj = \text{NOEXIST}$   
**using** *all-dv-ne p4 p7* **apply**(*simp add:inv-bitmapn-def inv-bitmap-not4free-def*)

```

Let-def)
  using jj-lt-len-lstbits ii lvlii-eq-last
  by (metis One-nat-def diff-self-eq-0 div-by-Suc-0 eq-imp-le power-0)
  ultimately have  $\exists n. n > 0 \wedge n \leq ii \wedge \text{get-bit-s } s \ p \ (n-1) \ (jj \text{ div } 4 \wedge (ii - (n-1))) = DIVIDED \wedge$ 
     $\text{get-bit-s } s \ p \ n \ (jj \text{ div } 4 \wedge (ii - n)) = NOEXIST$ 
    using divornoe-imp-div-noe-neigh[of ii s p jj] by fastforce

  then obtain n where  $n > 0 \wedge n \leq ii \wedge \text{get-bit-s } s \ p \ (n-1) \ (jj \text{ div } 4 \wedge (ii - (n-1))) = DIVIDED \wedge$ 
     $\text{get-bit-s } s \ p \ n \ (jj \text{ div } 4 \wedge (ii - n)) = NOEXIST$  by auto

  moreover
    with p3 have  $\text{get-bit-s } s \ p \ (n - \text{Suc } NULL) \ (jj \text{ div } 4 \wedge (ii - (n - \text{Suc } NULL))) \neq DIVIDED$ 
    apply (simp add: inv-bitmap-def Let-def)
    using Divides.div-mult2-eq One-nat-def Suc-diff-eq-diff-pred Suc-pred
    diff-Suc-eq-diff-pred diff-commute ii less-Suc-eq-le li-len p4 power-minus-mult
    zero-less-diff
    by (smt le-imp-less-Suc zero-le-numeral)
    ultimately have False by simp

} thus ?thesis by auto
qed

thus ?thesis by auto
qed

```

```

lemma div-imp-up-alldiv:
 $\forall i1 \ j1 \ j2. \text{inv-bitmap } s \wedge \text{inv-bitmap0 } s \wedge$ 
   $\text{inv-mem-pool-info } s \wedge$ 
   $p \in \text{mem-pools } s \wedge$ 
   $i1 < \text{length } (\text{levels } (\text{mem-pool-info } s \ p)) \wedge$ 
   $j1 < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } s \ p) \ ! \ i1)) \wedge$ 
   $i2 < \text{length } (\text{levels } (\text{mem-pool-info } s \ p)) \wedge$ 
   $j2 < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } s \ p) \ ! \ i2)) \wedge$ 
   $\text{get-bit-s } s \ p \ i2 \ j2 = DIVIDED \wedge$ 
   $i1 < i2 \wedge$ 
   $j1 = j2 \text{ div } 4 \wedge (i2 - i1) \longrightarrow$ 
   $\text{get-bit-s } s \ p \ i1 \ j1 = DIVIDED$ 
apply (induct i2)
apply simp

apply clarsimp
apply (case-tac i1 = i2)
  apply clarsimp apply (simp add: inv-bitmap-def Let-def)
  apply fastforce

apply (subgoal-tac i1 < i2) prefer 2 apply simp

```



```

apply(subgoal-tac get-bit-s s p i2 (j2 div 4) = DIVIDED) prefer 2
apply(simp add:inv-bitmap-def Let-def) apply fastforce
apply(subgoal-tac get-bit-s s p i1 ((j2 div 4) div 4 ^ (i2 - i1)) = DIVIDED)
prefer 2
apply(subgoal-tac (j2 div 4) div 4 ^ (i2 - i1) < length (bits (levels (mem-pool-info
s p) ! i1)))
prefer 2 apply (simp add: Divides.div-mult2-eq Suc-diff-le)
apply(subgoal-tac j2 div 4 < length (bits (levels (mem-pool-info s p) ! i2)))
prefer 2 apply(simp add:inv-mempool-info-def Let-def)
apply fastforce
apply(subgoal-tac j2 div 4 div 4 ^ (i2 - i1) = j2 div 4 ^ (Suc i2 - i1))
prefer 2
apply (metis Suc-diff-le div-mult2-eq less-or-eq-imp-le power-Suc)
apply fastforce
done

```

**lemma** block-imp-up-alldiv:

```

inv-bitmap s ==> inv-bitmap0 s ==>
  inv-mempool-info s ==>
    p ∈ mem-pools s ==>
      i1 < length (levels (mem-pool-info s p)) ==>
        j1 < length (bits (levels (mem-pool-info s p) ! i1)) ==>
          i2 < length (levels (mem-pool-info s p)) ==>
            j2 < length (bits (levels (mem-pool-info s p) ! i2)) ==>
              (get-bit-s s p i2 j2 = FREE ∨
                get-bit-s s p i2 j2 = FREEING ∨ get-bit-s s p i2 j2 = ALLOCATED ∨ get-bit-s
s p i2 j2 = ALLOCATING) ==>
                i1 < i2 ==>
                  j1 = j2 div 4 ^ (i2 - i1) ==>
                    get-bit-s s p i1 j1 = DIVIDED
apply(subgoal-tac get-bit-s s p (i2 - 1) (j2 div 4) = DIVIDED)
prefer 2 apply(simp add:inv-bitmap-def Let-def)
apply (metis neq0-conv not-less-zero)
apply(case-tac i1 = i2 - 1)
apply simp

apply clarsimp
apply(rule div-imp-up-alldiv[rule-format, of s p i1 j2 div 4 ^ (i2 - i1) i2 - 1 j2
div 4])
apply clarsimp
apply(rule conjI) apply simp
apply(rule conjI) apply(simp add:inv-mempool-info-def Let-def)
  using One-nat-def div-eq-0-iff gr-implies-not0 nat-0-less-mult-iff
apply (metis (no-types, lifting) less-mult-imp-div-less nat-neq-iff power-minus-mult
semiring-normalization-rules(17))
  using Divides.div-mult2-eq Suc-diff-Suc Suc-pred linorder-neqE-nat not-less-eq
not-less-zero power-Suc
  by (metis not-less)

```

**lemma** *addr-in-same-block*:

*inv-bitmap0 s*  $\implies$  *inv-bitmap s*  $\implies$  *inv-mempool-info s*  $\implies$

*p*  $\in$  *mem-pools s*  $\implies$  *addr* < *n-max* (*mem-pool-info s p*) \* *max-sz* (*mem-pool-info s p*)  $\implies$

*addr-in-block* (*mem-pool-info s p*) *addr i1 j1*  $\implies$

*addr-in-block* (*mem-pool-info s p*) *addr i2 j2*  $\implies$

*i1* = *i2*  $\wedge$  *j1* = *j2*

**apply**(*case-tac i1 = i2*)

**apply**(*rule conjI*) **apply** *fast*

**apply** *clarsimp*

**apply**(*case-tac j1 < j2*)

**apply** (*smt Groups.mult-ac*(2) *mult-Suc-right nat-0-less-mult-iff neq0-conv not-le split-div-lemma*)

**apply**(*case-tac j1 > j2*)

**apply** (*smt Groups.mult-ac*(2) *mult-Suc-right nat-0-less-mult-iff neq0-conv not-le split-div-lemma*)

**apply** *simp*

**apply**(*subgoal-tac*  $\exists n > 0. \text{max-sz } (\text{mem-pool-info } s \text{ } p) = (4 * n) * (4 \wedge n\text{-levels } (\text{mem-pool-info } s \text{ } p))$ )

**prefer** 2 **apply**(*simp add:inv-mempool-info-def Let-def*) **apply** *metis*

**apply**(*subgoal-tac length* (*levels* (*mem-pool-info s p*)) = *n-levels* (*mem-pool-info s p*))

**prefer** 2 **apply**(*simp add:inv-mempool-info-def Let-def*)

**apply**(*case-tac i1 < i2*)

**apply**(*subgoal-tac addr div* (*max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i1$ ) = *j1*)

**prefer** 2 **using** *addr-in-div*[*of addr j1 max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i1$ ]

**apply** *simp*

**apply**(*subgoal-tac addr div* (*max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$ ) = *j2*)

**prefer** 2 **using** *addr-in-div*[*of addr j2 max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$ ]

**apply** *simp*

**apply**(*subgoal-tac j1 = j2 div* ( $4 \wedge (i2 - i1)$ )) **prefer** 2

**apply**(*rule subst*[**where** *s=addr div* (*max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$ ) *div*  $4 \wedge (i2 - i1)$  **and** *t=j2 div*  $4 \wedge (i2 - i1)$ ])

**apply** *fast*

**apply**(*rule subst*[**where** *s=addr div* ((*max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$ ) \*  $4 \wedge (i2 - i1)$ )

**and** *t=addr div* (*max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$ ) *div*  $4 \wedge (i2 - i1)$ ])

**using** *div2-eq-divmul*[*of addr max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$   $4 \wedge (i2 - i1)$ ] **apply** *simp*

**apply**(*rule subst*[**where** *s=max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i1$  **and**

*t=max-sz* (*mem-pool-info s p*) *div*  $4 \wedge i2$  \*  $4 \wedge (i2 - i1)$ ])

**apply**(*subgoal-tac max-sz* (*mem-pool-info s p*) *mod* ( $4 \wedge i1$ ) = 0)

**prefer** 2 **apply** (*metis ge-pow-mod-0*)

**apply**(*subgoal-tac max-sz* (*mem-pool-info s p*) *mod* ( $4 \wedge i2$ ) = 0)

```

    prefer 2 apply (metis ge-pow-mod-0)
    apply(smt add-diff-inverse-nat div2-eq-minus less-imp-le-nat m-mod-div minus-div-mult-eq-mod

        minus-mult-div-eq-mod mod-div-self mod-mult-self2-is-0 not-less power-add
        zero-neq-numeral)
    apply fast

    apply(subgoal-tac get-bit-s s p i1 j1 = DIVIDED)
    prefer 2 using block-imp-up-alldiv[of s p i1 j1 i2 j2] apply fast
    apply auto[1]

    apply(case-tac i1 > i2)
    apply(subgoal-tac addr div (max-sz (mem-pool-info s p) div 4 ^ i1) = j1)
    prefer 2 using addr-in-div[of addr j1 max-sz (mem-pool-info s p) div 4 ^ i1]
    apply simp
    apply(subgoal-tac addr div (max-sz (mem-pool-info s p) div 4 ^ i2) = j2)
    prefer 2 using addr-in-div[of addr j2 max-sz (mem-pool-info s p) div 4 ^ i2]
    apply simp

    apply(subgoal-tac j2 = j1 div (4 ^ (i1 - i2))) prefer 2
    apply(rule subst[where s=addr div (max-sz (mem-pool-info s p) div 4 ^ i1) div
        4 ^ (i1 - i2) and t=j1 div 4 ^ (i1 - i2)])
    apply fast
    apply(rule subst[where s=addr div ((max-sz (mem-pool-info s p) div 4 ^ i1) *
        4 ^ (i1 - i2))
        and t=addr div (max-sz (mem-pool-info s p) div 4 ^ i1) div 4 ^
        (i1 - i2)])
    using div2-eq-divmul[of addr max-sz (mem-pool-info s p) div 4 ^ i1 4 ^ (i1 -
        i2)] apply simp
    apply(rule subst[where s=max-sz (mem-pool-info s p) div 4 ^ i2 and
        t=max-sz (mem-pool-info s p) div 4 ^ i1 * 4 ^ (i1 - i2)])
    apply(subgoal-tac max-sz (mem-pool-info s p) mod (4 ^ i1) = 0)
    prefer 2 apply (metis ge-pow-mod-0)
    apply(subgoal-tac max-sz (mem-pool-info s p) mod (4 ^ i2) = 0)
    prefer 2 apply (metis ge-pow-mod-0)
    apply(smt add-diff-inverse-nat div2-eq-minus less-imp-le-nat m-mod-div minus-div-mult-eq-mod

        minus-mult-div-eq-mod mod-div-self mod-mult-self2-is-0 not-less power-add
        zero-neq-numeral)
    apply fast

    apply(subgoal-tac get-bit-s s p i2 j2 = DIVIDED)
    prefer 2 using block-imp-up-alldiv[of s p i2 j2 i1 j1] apply fast
    apply auto[1]

    apply auto
    done

    lemma inv-impl-mem-cover':

```

```

inv-mempool-info s  $\impl$ 
  inv-bitmap0 s  $\impl$  inv-bitmap s  $\impl$  inv-bitmapn s  $\impl$  mem-cover s
apply(simp add: mem-cover-def Let-def)
apply clarify
apply(rule ex-ex1I)
  apply clarsimp using addr-exist-block[of s] apply fastforce
  apply clarsimp using addr-in-same-block[of s] apply force
done

```

```

lemma inv-impl-mem-cover: inv s  $\impl$  mem-cover s
  apply(simp add:inv-def)
  using inv-impl-mem-cover' apply fast
done

```

```

abbreviation divide-noexist-cont' :: State  $\Rightarrow$  mempool-ref  $\Rightarrow$  bool
where divide-noexist-cont' s p  $\equiv$ 
  let mp = mem-pool-info s p in
   $\forall i < \text{length } (\text{levels } mp).$ 
  let bts = bits (levels mp ! i) in
  ( $\forall j < \text{length } bts. (bts ! j = \text{DIVIDED} \longrightarrow i > 0 \longrightarrow (\text{bits } (\text{levels } mp$ 
  ! (i - 1))) ! (j div 4) = DIVIDED)
   $\wedge (bts ! j = \text{NOEXIST} \longrightarrow i < \text{length } (\text{levels } mp) - 1 \longrightarrow \text{noexist-bits}$ 
  mp (i+1) (j*4)) )

```

```

definition divide-noexist-cont :: State  $\Rightarrow$  bool
where divide-noexist-cont s  $\equiv$ 
   $\forall p \in \text{mem-pools } s. \text{divide-noexist-cont}' s p$ 

```

**end**

```

theory rg-cond
imports mem-spec invariant
begin

```

## 5 Rely-guarantee condition of events

### 5.1 defs of rely-guarantee conditions

```

definition lvars-nochange :: Thread  $\Rightarrow$  State  $\Rightarrow$  State  $\Rightarrow$  bool
where lvars-nochange t r s  $\equiv$ 
  i r t = i s t  $\wedge$  j r t = j s t  $\wedge$  ret r t = ret s t
   $\wedge$  endt r t = endt s t  $\wedge$  rf r t = rf s t  $\wedge$  tmout r t = tmout s t
   $\wedge$  lsizes r t = lsizes s t  $\wedge$  alloc-l r t = alloc-l s t  $\wedge$  free-l r t = free-l s t
   $\wedge$  from-l r t = from-l s t  $\wedge$  blk r t = blk s t  $\wedge$  nodev r t = nodev s t
   $\wedge$  bn r t = bn s t  $\wedge$  lbn r t = lbn s t  $\wedge$  lsz r t = lsz s t  $\wedge$  block2 r t = block2 s t
   $\wedge$  free-block-r r t = free-block-r s t  $\wedge$  alloc-lsize-r r t = alloc-lsize-r s t  $\wedge$  lvl r t
  = lvl s t  $\wedge$  bb r t = bb s t
   $\wedge$  block-pt r t = block-pt s t  $\wedge$  th r t = th s t  $\wedge$  need-resched r t = need-resched

```

$s \ t$   
 $\wedge \text{mempoolalloc-ret } r \ t = \text{mempoolalloc-ret } s \ t$   
 $\wedge \text{freeing-node } r \ t = \text{freeing-node } s \ t \wedge \text{allocating-node } r \ t = \text{allocating-node } s \ t$

**definition**  $\text{lvars-nochange-rel} :: \text{Thread} \Rightarrow (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{lvars-nochange-rel } t \equiv \{(s,r). \text{lvars-nochange } t \ s \ r\}$

**definition**  $\text{lvars-nochange-4all} :: (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{lvars-nochange-4all} \equiv \{(s,r). \forall t. \text{lvars-nochange } t \ s \ r\}$

**definition**  $\text{lvars-nochange1} :: \text{Thread} \Rightarrow \text{State} \Rightarrow \text{State} \Rightarrow \text{bool}$   
**where**  $\text{lvars-nochange1 } t \ r \ s \equiv \text{freeing-node } r \ t = \text{freeing-node } s \ t \wedge \text{allocating-node } r \ t = \text{allocating-node } s \ t$

**definition**  $\text{lvars-nochange1-rel} :: \text{Thread} \Rightarrow (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{lvars-nochange1-rel } t \equiv \{(s,r). \text{lvars-nochange1 } t \ s \ r\}$

**definition**  $\text{lvars-nochange1-4all} :: (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{lvars-nochange1-4all} \equiv \{(s,r). \forall t. \text{lvars-nochange1 } t \ s \ r\}$

**lemma**  $\text{lvars-nochange-trans}$ :  
 $\text{lvars-nochange } t \ x \ y \implies \text{lvars-nochange } t \ y \ z \implies \text{lvars-nochange } t \ x \ z$   
**apply**( $\text{simp add:lvars-nochange-def}$ )  
**done**

**lemma**  $\text{lvars-nochange-sym}$ :  
 $\text{lvars-nochange } t \ x \ y \implies \text{lvars-nochange } t \ y \ x$   
**apply**( $\text{simp add:lvars-nochange-def}$ )  
**done**

**lemma**  $\text{lvars-nochange-refl}$ :  
 $\text{lvars-nochange } t \ x \ x$   
**apply**( $\text{simp add:lvars-nochange-def}$ )  
**done**

**lemma**  $\text{lvars-nc-nc1}$ :  $\text{lvars-nochange } t \ r \ s \implies \text{lvars-nochange1 } t \ r \ s$   
**unfolding**  $\text{lvars-nochange-def lvars-nochange1-def}$  **by**  $\text{simp}$

**lemma**  $\text{lv-noch-all1}$ :  $(s,r) \in \text{lvars-nochange-4all}$   
 $\implies (s,r) \in \text{lvars-nochange-rel } t \wedge (\forall t'. t' \neq t \longrightarrow (s,r) \in \text{lvars-nochange-rel } t')$   
**unfolding**  $\text{lvars-nochange-4all-def lvars-nochange-rel-def}$  **by**  $\text{auto}$

**lemma**  $\text{lv-noch-all2}$ :  $(s,r) \in \text{lvars-nochange-rel } t \wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r)$   
 $\implies (s,r) \in \text{lvars-nochange-4all}$   
**unfolding**  $\text{lvars-nochange-4all-def lvars-nochange-rel-def}$  **by**  $\text{auto}$

**definition**  $\text{gvars-nochange} :: \text{State} \Rightarrow \text{State} \Rightarrow \text{bool}$   
**where**  $\text{gvars-nochange } s \ r \equiv \text{cur } r = \text{cur } s \wedge \text{tick } r = \text{tick } s \wedge \text{thd-state } r =$

*thd-state s*

$$\wedge \text{mem-pools } r = \text{mem-pools } s \wedge \text{mem-pool-info } r =$$

*mem-pool-info s*

**definition** *gvars-nochange-rel* :: (State × State) set

**where** *gvars-nochange-rel* ≡ {(s,r). *gvars-nochange s r*}

**definition** *gvars-conf* :: State ⇒ State ⇒ bool

**where** *gvars-conf s r* ≡

*mem-pools r* = *mem-pools s*

∧ (∀ p. *buf (mem-pool-info s p)* = *buf (mem-pool-info r p)*)

∧ *max-sz (mem-pool-info s p)* = *max-sz (mem-pool-info r p)*

∧ *n-max (mem-pool-info s p)* = *n-max (mem-pool-info r p)*

∧ *n-levels (mem-pool-info s p)* = *n-levels (mem-pool-info r p)*

∧ *length (levels (mem-pool-info s p))* = *length (levels (mem-pool-info r p))*

∧ (∀ i. *length (bits (levels (mem-pool-info s p) ! i))*)

= *length (bits (levels (mem-pool-info r p) ! i))*)

**definition** *gvars-conf-stable* :: (State × State) set

**where** *gvars-conf-stable* ≡ {(s,r). *gvars-conf s r*}

**definition** *inv-sta-rely* :: (State × State) set

**where** *inv-sta-rely* ≡ {(s,r). *inv s* ⇒ *inv r*}

**definition** *inv-sta-guar* :: (State × State) set

**where** *inv-sta-guar* ≡ {(s,r). *inv s* ⇒ *inv r*}

**lemma** *glnchange-inv0*:

(a, b) ∈ *lvars-nochange1-4all* ⇒ *cur a* = *cur b* ⇒

*thd-state a* = *thd-state b* ⇒ *mem-pools a* = *mem-pools b* ⇒

*mem-pool-info a* = *mem-pool-info b* ⇒ *inv a* ⇒ *inv b*

**apply**(*simp add:lvars-nochange1-4all-def lvars-nochange1-def inv-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-thd-waitq-def*) **apply** *auto*[1]

**apply**(*rule conjI*) **apply**(*simp add:inv-mempool-info-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap-freelist-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap-def*)

**apply**(*rule conjI*) **apply**(*simp add: inv-aux-vars-def mem-block-addr-valid-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap0-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmapn-def*)

**apply**(*simp add:inv-bitmap-not4free-def*)

**done**

**lemma** *glnchange-inv1*:

(a, b) ∈ *lvars-nochange-4all* ⇒ *cur a* = *cur b* ⇒

*thd-state a* = *thd-state b* ⇒ *mem-pools a* = *mem-pools b* ⇒

*mem-pool-info a* = *mem-pool-info b* ⇒ *inv a* ⇒ *inv b*

```
apply(simp add:lwars-nochange-4all-def lwars-nochange-def)
using glnochange-inv0
apply(simp add:lwars-nochange1-4all-def lwars-nochange1-def)
by metis
```

```

lemma glnochange-inv:
  inv a  $\implies \forall t'. t' \neq t1 \longrightarrow$  lvvars-nochange t' a b
     $\implies$  gvvars-nochange a b  $\implies$  lvvars-nochange t1 a b  $\implies$  inv b
apply(subgoal-tac (a, b)  $\in$  lvvars-nochange-4all)
  apply(simp add: gvvars-nochange-def)
  using glnochange-inv1 apply auto
using lv-noch-all2[of a b t1] apply auto[1]
by(simp add: lvvars-nochange-rel-def)

```

**definition** *Schedule-rely* :: (State  $\times$  State) set  
**where** *Schedule-rely*  $\equiv \{(s, r). \text{inv } s \longrightarrow \text{inv } r\} \cup Id$

**definition** *Schedule-guar* :: (*State* × *State*) set  
**where** *Schedule-guar* ≡  

$$\begin{aligned} & ((* \{ (\circ cur \neq Some\ t \longrightarrow \\ & \quad (\circ cur \neq None \longrightarrow {}^a thd\text{-}state = (\circ thd\text{-}state\ (the\ (\circ cur) := READY))(t := \\ & RUNNING) \wedge {}^a cur = Some\ t) \\ & \quad \wedge (\circ cur = None \longrightarrow {}^a thd\text{-}state = \circ thd\text{-}state\ (t := RUNNING)) \wedge {}^a cur = \\ & Some\ t) \\ & \quad \wedge (\circ cur = Some\ t \longrightarrow {}^a thd\text{-}state = \circ thd\text{-}state \wedge \circ cur = {}^a cur) \} *) \\ & \{ (s, r).\ inv\ s \longrightarrow inv\ r \} \\ & \cap \{ \circ tick = {}^a tick \wedge \circ mem\text{-}pools = {}^a mem\text{-}pools \wedge \circ mem\text{-}pool\text{-}info = {}^a mem\text{-}pool\text{-}info \} \end{aligned}$$
  

$$\cap (\bigcap t.\ lvars\text{-}nochange\text{-}rel\ t) \cup Id$$

**definition** *Schedule-RGCond* :: *Thread*  $\Rightarrow$  (*State*) *PiCore-Hoare.rgformula*  
**where** *Schedule-RGCond* *t*  $\equiv$   
 $RG[\{s. \text{inv } s\},$   
*Schedule-rely*, *Schedule-guar*,  
 $\{s. \text{inv } s\}]$

**definition**  $Tick\text{-}rely :: (State \times State) \text{ set}$   
**where**  $Tick\text{-}rely \equiv \{\text{tick}^o = \text{tick}^a\} \cup Id$

**definition** *Tick-guar* :: (State × State) set  
**where** *Tick-guar* ≡ (⌈<sup>a</sup>tick = <sup>o</sup>tick + 1 ∧ <sup>o</sup>cur = <sup>a</sup>cur ∧ <sup>o</sup>thd-state = <sup>a</sup>thd-state  
 ∧ <sup>o</sup>mem-pools = <sup>a</sup>mem-pools ∧ <sup>o</sup>mem-pool-info = <sup>a</sup>mem-pool-info⌋  
 ∩ (⋂ t. lvars-nochange-rel t)) ∪ Id

**definition** *Tick-RGCond* :: (State) PiCore-Hoare.rgformula  
**where** *Tick-RGCond*  $\equiv$   
 $RG[\{True\}, Tick\text{-}rely, Tick\text{-}guar, \{True\}]$

**abbreviation**  $alloc\text{-}blk\text{-}valid :: State \Rightarrow mempool\text{-}ref \Rightarrow nat \Rightarrow nat \Rightarrow mem\text{-}ref$

$\Rightarrow \text{bool}$   
**where**  $\text{alloc-blk-valid } s \ p \ lv \ bnum \ blkaddr$   
 $\equiv (blkaddr = buf \ (\text{mem-pool-info } s \ p) + bnum * ((\text{max-sz } (\text{mem-pool-info } s \ p)) \text{ div } (4 \wedge lv)))$   
 $\wedge bnum < n\text{-max } (\text{mem-pool-info } s \ p) * (4 \wedge lv))$

**abbreviation**  $\text{alloc-memblk-data-valid} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{Mem-block} \Rightarrow \text{bool}$   
**where**  $\text{alloc-memblk-data-valid } s \ p \ mb \equiv \text{alloc-blk-valid } s \ p \ (\text{level } mb) \ (\text{block } mb) \ (\text{data } mb)$

**definition**  $\text{alloc-memblk-valid} :: \text{State} \Rightarrow \text{mempool-ref} \Rightarrow \text{nat} \Rightarrow \text{Mem-block} \Rightarrow \text{bool}$   
**where**  $\text{alloc-memblk-valid } s \ p \ sz \ mb \equiv$   
 $p = \text{pool } mb \wedge p \in \text{mem-pools } s$   
 $\wedge sz \leq (\text{max-sz } (\text{mem-pool-info } s \ p)) \text{ div } (4 \wedge (\text{level } mb)) * (\text{block size of level } mb + 1 < sz \leq \text{block size of level } mb)$   
 $\wedge (\text{level } mb < n\text{-levels } (\text{mem-pool-info } s \ p) - 1 \longrightarrow sz > (\text{max-sz } (\text{mem-pool-info } s \ p)) \text{ div } (4 \wedge (\text{level } mb + 1)))$   
 $\wedge \text{alloc-memblk-data-valid } s \ p \ mb$

**abbreviation**  $\text{Mem-pool-alloc-pre} :: \text{Thread} \Rightarrow \text{State set}$   
**where**  $\text{Mem-pool-alloc-pre } t \equiv \{s. \text{inv } s \wedge \text{allocating-node } s \ t = \text{None} \wedge \text{freeing-node } s \ t = \text{None}\}$

**definition**  $\text{Mem-pool-alloc-rely} :: \text{Thread} \Rightarrow (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{Mem-pool-alloc-rely } t \equiv$   
 $((\text{lvars-nochange-rel } t \cap \text{gvars-conf-stable}$   
 $\cap \{(s,r). \text{inv } s \longrightarrow \text{inv } r\}$   
 $\cap \{(s,r). (\text{cur } s = \text{Some } t \longrightarrow \text{mem-pool-info } s = \text{mem-pool-info } r$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r))\}) \cup \text{Id})$

**definition**  $\text{Mem-pool-alloc-guar} :: \text{Thread} \Rightarrow (\text{State} \times \text{State}) \text{ set}$   
**where**  $\text{Mem-pool-alloc-guar } t \equiv$   
 $((\text{gvars-conf-stable} \cap$   
 $\{(s,r). (\text{cur } s \neq \text{Some } t \longrightarrow \text{gvars-nochange } s \ r \wedge \text{lvars-nochange } t \ s \ r)$   
 $\wedge (\text{cur } s = \text{Some } t \longrightarrow \text{inv } s \longrightarrow \text{inv } r)$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r)) \}$   
 $\cap \{\text{tick} = \text{tick}\}) \cup \text{Id})$

**definition**  $\text{Mem-pool-alloc-post} :: \text{Thread} \Rightarrow \text{mempool-ref} \Rightarrow \text{nat} \Rightarrow \text{int} \Rightarrow \text{State set}$   
**where**  $\text{Mem-pool-alloc-post } t \ p \ sz \ \text{timeout} \equiv$   
 $\{s. \text{inv } s \wedge \text{allocating-node } s \ t = \text{None} \wedge \text{freeing-node } s \ t = \text{None}$   
 $\wedge (\text{timeout} = \text{FOREVER} \longrightarrow (\text{ret } s \ t = \text{ESIZEERR} \wedge \text{mempoolalloc-ret } s \ t = \text{None})$



$$\begin{aligned}
& \vee \text{ret } s \ t = \text{OK} \wedge (\exists \text{mbk}. \text{mempoolalloc-ret } s \ t = \text{Some} \\
& \text{mbk} \wedge \text{alloc-memblk-valid } s \ p \ sz \ \text{mbk})) \\
& \wedge (\text{timeout} = \text{NOWAIT} \longrightarrow ((\text{ret } s \ t = \text{ENOMEM} \vee \text{ret } s \ t = \text{ESIZEERR}) \\
& \wedge \text{mempoolalloc-ret } s \ t = \text{None}) \\
& \vee (\text{ret } s \ t = \text{OK} \wedge (\exists \text{mbk}. \text{mempoolalloc-ret } s \ t = \text{Some} \\
& \text{mbk} \wedge \text{alloc-memblk-valid } s \ p \ sz \ \text{mbk})) \\
& \wedge (\text{timeout} > 0 \longrightarrow ((\text{ret } s \ t = \text{ETIMEOUT} \vee \text{ret } s \ t = \text{ESIZEERR}) \wedge \\
& \text{mempoolalloc-ret } s \ t = \text{None}) \\
& \vee (\text{ret } s \ t = \text{OK} \wedge (\exists \text{mbk}. \text{mempoolalloc-ret } s \ t = \text{Some } \text{mbk} \\
& \wedge \text{alloc-memblk-valid } s \ p \ sz \ \text{mbk}))))\}
\end{aligned}$$

**definition** *Mem-pool-alloc-RGCond* :: *Thread*  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *int*  $\Rightarrow$  *(State) PiCore-Hoare.rgformula*

**where** *Mem-pool-alloc-RGCond* *t p sz timeout*  $\equiv$   
 $RG[\text{Mem-pool-alloc-pre } t,$   
 $\text{Mem-pool-alloc-rely } t,$   
 $\text{Mem-pool-alloc-guar } t,$   
 $\text{Mem-pool-alloc-post } t \ p \ sz \ \text{timeout}]$

**abbreviation** *Mem-pool-free-pre* :: *Thread*  $\Rightarrow$  *State set*

**where** *Mem-pool-free-pre* *t*  $\equiv \{s. \text{inv } s \wedge \text{allocating-node } s \ t = \text{None} \wedge \text{freeing-node } s \ t = \text{None}\}$

**definition** *Mem-pool-free-rely* :: *Thread*  $\Rightarrow$  *(State  $\times$  State) set*

**where** *Mem-pool-free-rely* *t*  $\equiv$   
 $((\text{lvars-nochange-rel } t \cap \text{gvars-conf-stable}$   
 $\cap \{(s,r). \text{inv } s \longrightarrow \text{inv } r\}$   
 $\cap \{(s,r). (\text{cur } s = \text{Some } t \longrightarrow \text{mem-pool-info } s = \text{mem-pool-info } r$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r))\}) \cup \text{Id})$

**definition** *Mem-pool-free-guar* :: *Thread*  $\Rightarrow$  *(State  $\times$  State) set*

**where** *Mem-pool-free-guar* *t*  $\equiv$   
 $((\text{gvars-conf-stable} \cap$   
 $\{(s,r). (\text{cur } s \neq \text{Some } t \longrightarrow \text{gvars-nochange } s \ r \wedge \text{lvars-nochange } t \ s \ r)$   
 $\wedge (\text{cur } s = \text{Some } t \longrightarrow \text{inv } s \longrightarrow \text{inv } r)$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r)) \}$   
 $\cap \{\text{tick} = \text{tick}\}) \cup \text{Id})$

**definition** *Mem-pool-free-post* :: *Thread*  $\Rightarrow$  *State set*

**where** *Mem-pool-free-post* *t*  $\equiv \{s. \text{inv } s \wedge \text{allocating-node } s \ t = \text{None} \wedge \text{freeing-node } s \ t = \text{None}\}$

**definition** *Mem-pool-free-RGCond* :: *Thread*  $\Rightarrow$  *Mem-block*  $\Rightarrow$  *(State) PiCore-Hoare.rgformula*

**where** *Mem-pool-free-RGCond* *t b*  $\equiv$   
 $RG[\text{Mem-pool-free-pre } t,$   
 $\text{Mem-pool-free-rely } t,$   
 $\text{Mem-pool-free-guar } t,$   
 $\text{Mem-pool-free-post } t]$

## 5.2 stability, subset relations of rely-guarantee conditions

**lemma** *stable-inv-free-rely*:

$(s,r) \in \text{Mem-pool-free-rely } t \implies \text{inv } s \implies \text{inv } r$   
**apply** (*simp add:Mem-pool-free-rely-def*)  
**apply**(*case-tac cur s = Some t*) **apply** *simp*  
**apply**(*subgoal-tac (s, r) ∈ lvars-nochange-4all*)  
**apply**(*simp add:lvars-nochange-4all-def lvars-nochange-def*)  
**apply**(*simp add:inv-def*) **unfolding** *gvars-conf-stable-def gvars-conf-def*  
**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def*) **apply** *auto[1]* **apply** *metis*  
**apply**(*simp add:lvars-nochange-4all-def lvars-nochange-rel-def*)  
**apply** *auto[1]* **apply**(*simp add:lvars-nochange-def*)  
**apply** *auto*

**done**

**lemma** *stable-inv-free-rely1*: *stable*  $\{\text{'inv}\}$  (*Mem-pool-free-rely t*)  
**using** *stable-inv-free-rely* **unfolding** *stable-def* **by** *auto*

**lemma** *stable-inv-alloc-rely*:

$(s,r) \in \text{Mem-pool-alloc-rely } t \implies \text{inv } s \implies \text{inv } r$   
**apply**(*subgoal-tac Mem-pool-alloc-rely t = Mem-pool-free-rely t*)  
**using** *stable-inv-free-rely* **apply** *simp*  
**by** (*simp add:Mem-pool-alloc-rely-def Mem-pool-free-rely-def*)

**lemma** *stable-inv-alloc-rely1*: *stable*  $\{\text{'inv}\}$  (*Mem-pool-alloc-rely t*)  
**using** *stable-inv-alloc-rely* **unfolding** *stable-def* **by** *auto*

**lemma** *stable-inv-sched-rely*:

$(s,r) \in \text{Schedule-rely} \implies \text{inv } s \implies \text{inv } r$   
**apply** (*simp add:Schedule-rely-def*) **by** *auto*

**lemma** *stable-inv-sched-rely1*: *stable*  $\{\text{'inv}\}$  *Schedule-rely*  
**using** *stable-inv-sched-rely* **unfolding** *stable-def* **by** *auto*

**lemma** *free-guar-stb-inv*: *stable*  $\{\text{'inv}\}$  (*Mem-pool-free-guar t*)

**proof** –

{  
**fix** *x*  
**assume** *a0: inv x*  
 {  
**fix** *y*  
**assume** *b0: (x,y) ∈ Mem-pool-free-guar t*  
**hence**  $(x,y) \in \{(s,r). (\text{cur } s \neq \text{Some } t \longrightarrow \text{gvars-nochange } s \ r \wedge \text{lvars-nochange } t \ s \ r)\}$   
 $\wedge (\text{cur } s = \text{Some } t \longrightarrow \text{inv } s \longrightarrow \text{inv } r)$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ s \ r)\}$   
**unfolding** *Mem-pool-free-guar-def gvars-nochange-def lvars-nochange-def* **by**  
*auto*  
**hence**  $(\text{cur } x \neq \text{Some } t \longrightarrow \text{gvars-nochange } x \ y \wedge \text{lvars-nochange } t \ x \ y)$   
 $\wedge (\text{cur } x = \text{Some } t \longrightarrow \text{inv } x \longrightarrow \text{inv } y)$

```

       $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ x \ y)$  by simp
    hence inv y
      apply(case-tac cur x  $\neq$  Some t)
      apply (simp add: gvars-nochange-def lvars-nochange-def) using a0 apply
clarify
      apply(simp add:inv-def)
      apply(rule conjI) apply(simp add:inv-cur-def)
      apply(rule conjI) apply(simp add:inv-thd-waitq-def) apply metis
      apply(rule conjI) apply(simp add:inv-mempool-info-def)
      apply(rule conjI) using inv-bitmap-freelist-def apply metis
      apply(rule conjI) apply(simp add:inv-bitmap-def)
      apply(rule conjI) apply(simp add:inv-aux-vars-def)
      apply(rule conjI) apply metis
      apply(rule conjI) apply(simp add:mem-block-addr-valid-def) apply metis

      apply(rule conjI) apply metis
      apply(rule conjI) apply(simp add:mem-block-addr-valid-def) apply metis

      apply(rule conjI) apply metis
      apply(rule conjI) apply metis
      apply metis
      apply(rule conjI) apply(simp add:inv-bitmap0-def)
      apply(rule conjI) apply(simp add:inv-bitmapn-def)
      apply(simp add:inv-bitmap-not4free-def)
    using a0 by auto
  }
}
}
then show ?thesis by (simp add:stable-def)
qed

lemma alloc-guar-stb-inv: stable  $\{\text{'inv}\}$  (Mem-pool-alloc-guar t)
  apply(subgoal-tac Mem-pool-alloc-guar t = Mem-pool-free-guar t)
  using free-guar-stb-inv apply simp
  by (simp add:Mem-pool-alloc-guar-def Mem-pool-free-guar-def)

lemma sched-guar-stb-inv:
  (s,r) $\in$ Schedule-guar  $\implies$  inv s  $\implies$  inv r
  apply(simp add:Schedule-guar-def)
  apply(erule disjE) by auto

lemma tick-guar-stb-inv:
  (s,r) $\in$ Tick-guar  $\implies$  inv s  $\implies$  inv r
  apply(simp add:Tick-guar-def) apply(erule disjE)
  using glnochange-inv0 lvars-nc-nc1
  unfolding lvars-nochange1-4all-def lvars-nochange-rel-def apply auto[1] apply
blast
  by auto

lemma mem-pool-alloc-pre-stb: stable (Mem-pool-alloc-pre t) (Mem-pool-alloc-rely

```

$t$ )  
**apply**(rule subst[**where**  $t = \{\text{'inv} \wedge \text{'allocating-node } t = \text{None} \wedge \text{'freeing-node } t = \text{None}\}$ ]  
**and**  $s = \{\text{'inv}\} \cap \{\text{'allocating-node } t = \text{None} \wedge \text{'freeing-node } t = \text{None}\}$ ])  
**apply** auto[1]  
**apply**(rule stable-int2) **apply** (simp add: stable-inv-alloc-rely1)  
**apply**(simp add:stable-def Mem-pool-alloc-rely-def gvars-conf-stable-def lvars-nochange-rel-def  
lvars-nochange-def)  
**done**

**lemma** mp-alloc-post-stb: stable (Mem-pool-alloc-post  $t$   $p$  sz timeout) (Mem-pool-alloc-rely  
 $t$ )  
**apply**(simp add:stable-def) **apply**(rule allI) **apply**(rule impI) **apply**(rule allI)  
**apply**(rule impI)  
**apply**(simp add:Mem-pool-alloc-rely-def Mem-pool-alloc-post-def)  
**apply**(rule conjI)  
**apply**(simp add:gvars-conf-stable-def) **unfolding** gvars-conf-def **apply**metis  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def)  
**apply**(case-tac  $x = y$ )  
**apply** simp **apply** clarify  
**apply**(simp add:alloc-memblk-valid-def gvars-conf-def gvars-conf-stable-def)  
**done**

**lemma** mem-pool-free-pre-stb: stable (Mem-pool-free-pre  $t$ ) (Mem-pool-free-rely  $t$ )  
**apply**(rule subst[**where**  $t = \{\text{'inv} \wedge \text{'allocating-node } t = \text{None} \wedge \text{'freeing-node } t = \text{None}\}$ ]  
**and**  $s = \{\text{'inv}\} \cap \{\text{'allocating-node } t = \text{None} \wedge \text{'freeing-node } t = \text{None}\}$ ])  
**apply** auto[1]  
**apply**(rule stable-int2) **apply** (simp add: stable-inv-free-rely1)  
**apply**(simp add:stable-def Mem-pool-free-rely-def gvars-conf-stable-def lvars-nochange-rel-def  
lvars-nochange-def)  
**done**

**lemma** mem-pool-free-post-stb: stable (Mem-pool-free-post  $t$ ) (Mem-pool-free-rely  
 $t$ )  
**using** mem-pool-free-pre-stb **apply**(simp add:Mem-pool-free-post-def)  
**done**

**lemma** allocuar-in-allocrely:  $t1 \neq t2 \implies \text{Mem-pool-alloc-guar } t1 \subseteq \text{Mem-pool-alloc-rely}$   
 $t2$   
**apply** clarify  
**proof** –  
**fix**  $a$   $b$   
**assume**  $p0$ :  $t1 \neq t2$   
**and**  $p1$ :  $(a, b) \in \text{Mem-pool-alloc-guar } t1$   
**hence**  $p2$ :  $(a, b) \in \text{gvars-conf-stable}$   
 $\wedge (\text{cur } a \neq \text{Some } t1 \implies \text{gvars-nochange } a \text{ } b \wedge \text{lvars-nochange } t1 \text{ } a$   
 $b)$

$\wedge (\text{cur } a = \text{Some } t1 \longrightarrow \text{inv } a \longrightarrow \text{inv } b)$   
 $\wedge (\forall t'. t' \neq t1 \longrightarrow \text{lvars-nochange } t' a b)$   
 $\wedge \text{tick } a = \text{tick } b \vee a = b$   
**unfolding** *Mem-pool-alloc-guar-def* **by** *auto*

**from** *p0 p2* **have**  
 $(a, b) \in \text{lvars-nochange-rel } t2 \wedge (a, b) \in \text{gvars-conf-stable}$   
 $\wedge (\text{inv } a \longrightarrow \text{inv } b)$   
 $\wedge (\text{cur } a = \text{Some } t2 \longrightarrow \text{mem-pool-info } a = \text{mem-pool-info } b$   
 $\wedge (\forall t'. t' \neq t2 \longrightarrow \text{lvars-nochange } t' a b))$   
 $\vee a = b$   
**apply** *clarify*  
**apply**(*rule conjI*) **apply**(*simp add:lvars-nochange-rel-def*)  
**apply**(*rule conjI*) **apply** *simp*  
**apply**(*rule conjI*) **apply** *clarify* **using** *glnochange-inv* **apply** *auto*[1]  
**apply** *clarify*  
**apply**(*rule conjI*) **apply**(*simp add:gvars-nochange-def*)  
**by** *auto*

**thus**  $(a, b) \in \text{Mem-pool-alloc-rely } t2$  **unfolding** *Mem-pool-alloc-rely-def* **by**  
*simp*  
**qed**

**lemma** *schedguar-in-allocrely*:  $\text{Schedule-guar} \subseteq \text{Mem-pool-alloc-rely } t2$   
**apply** *clarify*  
**proof** –  
**fix** *a b*  
**assume** *p0*:  $(a, b) \in \text{Schedule-guar}$   
**hence** *p1*:  $(\text{inv } a \longrightarrow \text{inv } b) \wedge \text{tick } a = \text{tick } b \wedge \text{mem-pools } a = \text{mem-pools } b \wedge$   
 $\text{mem-pool-info } a = \text{mem-pool-info } b$   
 $\wedge (a, b) \in (\bigcap t. \text{lvars-nochange-rel } t) \vee a = b$   
**by**(*simp add:Schedule-guar-def*)

**hence**  $(a, b) \in \text{lvars-nochange-rel } t2 \wedge (a, b) \in \text{gvars-conf-stable}$   
 $\wedge (\text{inv } a \longrightarrow \text{inv } b)$   
 $\wedge (\text{cur } a = \text{Some } t2 \longrightarrow \text{mem-pool-info } a = \text{mem-pool-info } b$   
 $\wedge (\forall t'. t' \neq t2 \longrightarrow \text{lvars-nochange } t' a b))$   
 $\vee a = b$   
**apply** *clarify*  
**apply**(*rule conjI*) **apply**(*simp add:lvars-nochange-rel-def*)  
**apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)  
**apply**(*rule conjI*) **apply** *clarify* **apply** *clarify*  
**by**(*simp add:lvars-nochange-rel-def*)

**thus**  $(a, b) \in \text{Mem-pool-alloc-rely } t2$  **by**(*simp add:Mem-pool-alloc-rely-def*)  
**qed**

**lemma** *schedguar-in-tickrely*:  $\text{Schedule-guar} \subseteq \text{Tick-rely}$   
**apply**(*simp add:Schedule-guar-def Tick-rely-def*)

```

by auto

lemma allocguar-in-tickrely: Mem-pool-alloc-guar  $t \subseteq$  Tick-rely
  apply (simp add: Mem-pool-alloc-guar-def Tick-rely-def)
  by auto

lemma tickguar-in-allocrely: Tick-guar  $\subseteq$  Mem-pool-alloc-rely  $t$ 
  apply clarify
  proof -
    fix  $a\ b$ 
    assume  $p0: (a, b) \in$  Tick-guar
    hence  $p1: \text{tick } b = \text{tick } a + 1 \wedge \text{cur } a = \text{cur } b \wedge \text{thd-state } a = \text{thd-state } b$ 
       $\wedge \text{mem-pools } a = \text{mem-pools } b \wedge \text{mem-pool-info } a = \text{mem-pool-info } b$ 
       $\wedge (a, b) \in (\bigcap t. \text{lvars-nochange-rel } t) \vee a = b$ 
    by (simp add: Tick-guar-def)

    hence  $(a, b) \in \text{lvars-nochange-rel } t \wedge (a, b) \in \text{gvars-conf-stable}$ 
       $\wedge (\text{inv } a \longrightarrow \text{inv } b)$ 
       $\wedge (\text{cur } a = \text{Some } t \longrightarrow \text{mem-pool-info } a = \text{mem-pool-info } b$ 
         $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ a \ b))$ 
       $\vee a = b$ 
    apply clarify
    apply (rule conjI) apply (simp add: lvars-nochange-rel-def)
    apply (rule conjI) apply (simp add: gvars-conf-stable-def gvars-conf-def)
    apply (rule conjI) using glnochange-inv0 lvars-nc-nc1 unfolding lvars-nochange-rel-def
    lvars-nochange1-4all-def
    apply auto[1] apply blast
    by auto

    thus  $(a, b) \in$  Mem-pool-alloc-rely  $t$  by (simp add: Mem-pool-alloc-rely-def)
  qed

lemma allocguar-in-schedrely: Mem-pool-alloc-guar  $t \subseteq$  Schedule-rely
  apply (simp add: Mem-pool-alloc-guar-def Schedule-rely-def)
  apply clarify
  apply (case-tac cur  $a = \text{Some } t$ )
  apply simp
  apply clarify
  using glnochange-inv by auto

lemma tickguar-in-schedrely: Tick-guar  $\subseteq$  Schedule-rely
  apply clarify
  proof -
    fix  $a\ b$ 
    assume  $p0: (a, b) \in$  Tick-guar
    thus  $(a, b) \in$  Schedule-rely
      apply (simp add: Tick-guar-def Schedule-rely-def) apply auto
    using glnochange-inv1 by (simp add: lvars-nochange-4all-def lvars-nochange-rel-def)
  
```

qed

end

**theory** *func-cor-lemma*  
**imports** *rg-cond*  
**begin**

**declare**  $[[\text{smt-timeout} = 300]]$

## 6 some lemmas for functional correctness by rely guarantee proof

**lemma** *inv-mempool-info-maxsz-mod4*:

$\text{inv-mempool-info } s \implies \forall p \in \text{mem-pools } s. \text{max-sz } (\text{mem-pool-info } s \ p) \bmod 4 = 0$

**unfolding** *inv-mempool-info-def*

**by** (*metis mod-mult-left-eq mod-mult-self1-is-0 mod-mult-self2-is-0 mult-0*)

**lemma** *inv-mempool-info-maxsz-align4*:

$\text{inv-mempool-info } s \implies \forall p \in \text{mem-pools } s. \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } s \ p)) = \text{max-sz } (\text{mem-pool-info } s \ p)$

**using** *inv-mempool-info-maxsz-mod4 align40* **by** *simp*

**lemma** *inv-maxsz-align4*:

$\text{inv } s \implies \forall p \in \text{mem-pools } s. \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } s \ p)) = \text{max-sz } (\text{mem-pool-info } s \ p)$

**unfolding** *inv-def* **using** *inv-mempool-info-maxsz-align4* **by** *simp*

**lemma** *lsizes-mod4*:

**assumes** *p0*:  $\text{inv } V$

**and** *p1*:  $\forall ii < \text{length } ls. ls \ ! \ ii = \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } V \ p)) \bmod 4 \wedge ii$

**and** *p2*:  $\text{length } ls \leq \text{length } (\text{levels } (\text{mem-pool-info } V \ p))$

**and** *p3*:  $p \in \text{mem-pools } V$

**shows**  $\forall ii < \text{length } ls. (ls \ ! \ ii) \bmod 4 = 0$

**proof** –

{

**fix** *ii*

**assume** *a0*:  $ii < \text{length } ls$

**from** *p0 p3* **have**  $\exists n > 0. \text{max-sz } (\text{mem-pool-info } V \ p) = (4 * n) * (4 \wedge (\text{length } (\text{levels } (\text{mem-pool-info } V \ p))))$

**apply**(*simp add:inv-def inv-mempool-info-def Let-def*) **by** *auto*

**then obtain** *n* **where**  $n > 0 \wedge \text{max-sz } (\text{mem-pool-info } V \ p) = (4 * n) * (4 \wedge$

(length (levels (mem-pool-info V p)))) **by** auto  
**hence** a1:  $n > 0 \wedge \text{max-sz (mem-pool-info V p)} = n * (4 ^ (\text{length (levels (mem-pool-info V p))} + 1))$  **by** auto  
  
**hence** ALIGN4 (max-sz (mem-pool-info V p)) = max-sz (mem-pool-info V p)  
**using** align40 **by** auto  
**with** a0 p1 **have** a2:  $ls ! ii = \text{max-sz (mem-pool-info V p)} \text{ div } 4 ^ ii$  **by** auto  
**with** a1 **have**  $ls ! ii = n * (4 ^ (\text{length (levels (mem-pool-info V p))} + 1)) \text{ div } 4 ^ ii$  **by** simp  
**moreover**  
**from** a0 p2 **have**  $(4 :: \text{nat}) ^ (\text{length (levels (mem-pool-info V p))} + 1) \bmod 4 ^ ii = 0$   
**using** pow-mod-0[of ii length (levels (mem-pool-info V p)) + 1 4] **by** auto  
**ultimately** **have** a3:  $ls ! ii = n * ((4 ^ (\text{length (levels (mem-pool-info V p))} + 1)) \text{ div } 4 ^ ii)$   
**using** m-mod-div **by** auto  
  
**from** a0 p2 **have**  $4 \neq \text{NULL} \wedge ii \leq \text{length (levels (mem-pool-info V p))} + 1$   
**by** linarith  
**hence**  $((4 :: \text{nat}) ^ (\text{length (levels (mem-pool-info V p))} + 1)) \text{ div } 4 ^ ii$   
 $= 4 ^ (\text{length (levels (mem-pool-info V p))} + 1 - ii)$   
**using** div2-eq-minus[of 4 ii (length (levels (mem-pool-info V p)) + 1)] **by** simp  
**hence**  $n * (((4 :: \text{nat}) ^ (\text{length (levels (mem-pool-info V p))} + 1)) \text{ div } 4 ^ ii)$   
 $= n * (4 ^ (\text{length (levels (mem-pool-info V p))} + 1 - ii))$  **by** auto  
**with** a3 **have**  $ls ! ii = n * (4 ^ (\text{length (levels (mem-pool-info V p))} + 1 - ii))$   
**by** auto  
**with** a0 p2 **have**  $ls ! ii \bmod 4 = 0$  **by** auto  
 }  
**then show** ?thesis **by** auto  
**qed**

**lemma** gvars-conf-stb-inv-mpinf:  $(x,y) \in \text{gvars-conf-stable} \implies \text{inv-mempool-info } y \implies \text{inv-mempool-info } x$   
**apply** (simp add: gvars-conf-stable-def gvars-conf-def inv-mempool-info-def)  
**apply** clarify  
**apply** (rule conjI) **apply** metis **apply** (rule conjI) **apply** metis  
**apply** (rule conjI) **apply** metis **apply** (rule conjI) **apply** metis  
**apply** (rule conjI) **apply** metis **apply** metis  
**done**

**lemma** ref-byblkn-self:  
 $R \geq \text{buf (mem-pool-info Va p)} \implies$   
 $(R - \text{buf (mem-pool-info Va p)}) \bmod \text{sz} = 0 \implies$   
 $\text{buf (mem-pool-info Va p)} + \text{block-num (mem-pool-info Va p)} R \text{ sz} * \text{sz} = R$   
**apply** (simp add: block-num-def)  
**apply** (rule subst[where t=(R - buf (mem-pool-info Va p)) div sz \* sz and s=R - buf (mem-pool-info Va p)])



by auto

**lemma** *partnerbits-udptn-notbit-partbits*:  
 $\forall jj < \text{length } lst. \neg (\text{let } a = (jj \text{ div } 4) * 4 \text{ in}$   
 $\quad lst!a = TAG \wedge lst!(a+1) = TAG \wedge lst!(a+2) = TAG \wedge lst!(a+3)$   
 $= TAG) \implies$   
 $TAG \neq TAG2 \implies lst' = \text{list-updates-n } lst \text{ ii } m \text{ TAG2} \implies$   
 $\forall jj < \text{length } lst'. \neg (\text{let } a = (jj \text{ div } 4) * 4 \text{ in}$   
 $\quad lst!a = TAG \wedge lst!(a+1) = TAG \wedge lst!(a+2) = TAG \wedge$   
 $lst!(a+3) = TAG)$   
**apply**(*unfold Let-def*) **apply**(*rule allI, rule impI*)  
**apply**(*case-tac*  $lst'! (jj \text{ div } 4 * 4) = TAG \wedge lst'! (jj \text{ div } 4 * 4 + 1) = TAG$   
 $\wedge lst'! (jj \text{ div } 4 * 4 + 2) = TAG \wedge lst'! (jj \text{ div } 4 * 4 + 3) = TAG)$   
**apply**(*subgoal-tac*  $\text{length } lst = \text{length } lst'$ ) **prefer** 2 **apply** *simp*  
**apply**(*subgoal-tac*  $\neg (lst! (jj \text{ div } 4 * 4) = TAG \wedge lst! (jj \text{ div } 4 * 4 + 1) =$   
 $TAG$   
 $\wedge lst! (jj \text{ div } 4 * 4 + 2) = TAG \wedge lst! (jj \text{ div } 4 * 4 + 3) =$   
 $TAG)$ )  
**prefer** 2 **apply** *presburger*  
**apply**(*case-tac*  $jj \text{ div } 4 * 4 + 3 < ii$ ) **using** *list-updates-n-neq*  
**apply** (*smt One-nat-def add.right-neutral add-Suc-right add-lessD1 numeral-Bit1*  
*numeral-One one-add-one plus-nat.simps(2)*)  
**apply**(*case-tac*  $jj \text{ div } 4 * 4 \geq ii + m$ ) **using** *list-updates-n-neq* **apply** (*smt*  
*le-add1 le-trans*)  
**using** *list-updates-eq* **apply** (*smt One-nat-def Suc-leI add.right-neutral add-Suc-right*  
*add-lessD1*  
*div-mult-mod-eq le-less-trans list-updates-n-beyond list-updates-n-eq list-updates-n-neq*  
*not-le numeral-Bit1 numeral-One one-add-one*)  
**by** *assumption*

end

**theory** *func-cor-other*  
**imports** *func-cor-lemma*  
**begin**

## 7 Functional correctness of Schedule

**lemma** *Schedule-satRG-h1*:  
 $\Gamma \vdash_I \text{Some } (IF \exists y. 'cur = \text{Some } y \text{ THEN } 'thd\text{-state} := 'thd\text{-state}(\text{the } 'cur :=$   
 $READY)); \text{Basic } (cur\text{-update } Map.empty) \text{ FI};$   
 $\text{Basic } (cur\text{-update } (\lambda-. \text{Some } t));$   
 $'thd\text{-state} := 'thd\text{-state}$   
 $(t := RUNNING)) \text{ sat}_p [\llbracket 'inv \rrbracket \cap \llbracket 'thd\text{-state } t = READY \rrbracket \cap$   
 $\{V\}, \{(s, t). s = t\}, UNIV, \llbracket '(Pair \ V) \in \text{Schedule-guar} \rrbracket$   
 $\cap \llbracket 'inv \rrbracket]$

```

apply(case-tac  $\{\text{'inv}\} \cap \{\text{'thd-state } t = \text{READY}\} \cap \{V\} = \{\}$ )
  using Emptyprecond apply auto[1]
  apply simp
  apply(case-tac  $\exists y. \text{cur } V = \text{Some } y$ )

  apply(rule Seq[where  $\text{mid} = \{V(\text{thd-state} := (\text{thd-state } V)(\text{the } (\text{cur } V) := \text{READY}))(\text{cur} := \text{None})(\text{cur} := \text{Some } t)\}$ ]])
  apply(rule Seq[where  $\text{mid} = \{V(\text{thd-state} := (\text{thd-state } V)(\text{the } (\text{cur } V) := \text{READY}))(\text{cur} := \text{None})\}$ ]])
  apply(rule Cond)
  apply(simp add:stable-def)
  apply(rule Seq[where  $\text{mid} = \{V(\text{thd-state} := (\text{thd-state } V)(\text{the } (\text{cur } V) := \text{READY}))\}$ ]])
  apply(rule Basic)
  apply auto[1]
  apply(simp add:stable-def)+
  apply(rule Basic)
  apply auto[1]
  apply(simp add:stable-def)+
  apply(simp add:Skip-def) apply(rule Basic) apply(simp add:stable-def)+

  apply(rule Basic)
  apply auto[1]
  apply(simp add:stable-def)+

  apply(rule Basic)
  apply(simp add:Schedule-guar-def)
  apply(subgoal-tac  $\text{inv } (V(\text{cur} := \text{Some } t, \text{thd-state} := (\text{thd-state } V)(\text{the } (\text{cur } V) := \text{READY}, t := \text{RUNNING}))) \wedge$ 
     $(\forall x. (V, V(\text{cur} := \text{Some } t, \text{thd-state} := (\text{thd-state } V)(\text{the } (\text{cur } V) := \text{READY}, t := \text{RUNNING}))) \in \text{lvars-nochange-rel } x))$ )
  apply simp
  apply(rule conjI) apply(simp add:inv-def) apply clarify
  apply(rule conjI) apply(simp add:inv-cur-def) apply force
  apply(rule conjI) apply(simp add:inv-thd-waitq-def inv-cur-def)
  apply (metis Thread-State-Type.distinct(3) Thread-State-Type.distinct(6))
  apply(rule conjI) apply(simp add:inv-mempool-info-def)
  apply(rule conjI) apply(simp add:inv-bitmap-freelist-def)
  apply(rule conjI) apply(simp add:inv-bitmap-def)
  apply(rule conjI) apply(simp add:inv-aux-vars-def mem-block-addr-valid-def)
  apply(rule conjI) apply(simp add:inv-bitmap0-def)
  apply(rule conjI) apply(simp add:inv-bitmapn-def)
  apply(simp add:inv-bitmap-not4free-def)
  apply auto[1] using lvars-nochange-rel-def lvars-nochange-def apply simp
  apply(simp add: stable-def)+

  apply(rule Seq[where  $\text{mid} = \{V(\text{cur} := \text{Some } t)\}$ ]])
  apply(rule Seq[where  $\text{mid} = \{V\}$ ]])
  apply(rule Cond)

```

```

    apply(simp add:stable-def)
    apply(rule Seq[where mid = {}])
    apply(rule Basic)
      apply auto[1]
      apply(simp add:stable-def)+
    apply(rule Basic)
      apply auto[1]
      apply(simp add:stable-def)+
    apply(simp add:Skip-def) apply(rule Basic) apply(simp add:stable-def)+
    apply(rule Basic)
      apply auto[1]
      apply(simp add:stable-def)+
    apply(rule Basic)
      apply(simp add:Schedule-guar-def)
      apply(subgoal-tac inv (V⟦cur := Some t, thd-state := (thd-state V)(t :=
RUNNING)⟧) ∧
      (∀ x. (V, V⟦cur := Some t, thd-state := (thd-state V)(t := RUNNING)⟧)
∈ lvars-nochange-rel x))
      apply simp
      apply(rule conjI) apply(simp add:inv-def) apply clarify
      apply(rule conjI) apply(simp add:inv-cur-def)
      apply(rule conjI) apply(simp add:inv-thd-waitq-def) apply auto[1]
      apply(rule conjI) apply(simp add:inv-mempool-info-def)
      apply(rule conjI) apply(simp add:inv-bitmap-freelist-def)
      apply(rule conjI) apply(simp add:inv-bitmap-def)
    apply(rule conjI) apply(simp add:inv-aux-vars-def mem-block-addr-valid-def)
      apply(rule conjI) apply(simp add:inv-bitmap0-def)
      apply(rule conjI) apply(simp add:inv-bitmapn-def)
      apply(simp add:inv-bitmap-not4free-def)

    apply auto[1] using lvars-nochange-rel-def lvars-nochange-def apply simp
    apply(simp add:stable-def)+
done

lemma Schedule-satRG: Γ (Schedule t) ⊢ Schedule-RGCond t
  apply(simp add:Evt-sat-RG-def)
  apply (simp add: Schedule-def Schedule-RGCond-def)
  apply(rule BasicEvt)
    apply(simp add:body-def Pref-def Postf-def guard-def
      Relyf-def Guarf-def getrgformula-def)
  apply(rule Await)
    using stable-inv-sched-rely1 apply simp using stable-inv-sched-rely1 apply
simp
    using Schedule-satRG-h1 apply simp

  apply(simp add:Pref-def Relyf-def getrgformula-def)
  using stable-inv-sched-rely1 apply simp

  by(simp add:Guarf-def getrgformula-def Schedule-guar-def)

```

## 8 Functional correctness of Tick

```

lemma Tick-satRG:  $\Gamma$  Tick  $\vdash$  Tick-RGCond
  apply(simp add: Evt-sat-RG-def)
  apply (simp add: Tick-def Tick-RGCond-def Tick-rely-def Tick-guar-def)
  apply(rule BasicEvt)
  apply(simp add: body-def Pref-def Postf-def guard-def
    Relyf-def Guarf-def getrgformula-def)
  apply(rule Basic)
  apply simp
  using lvars-nochange-rel-def lvars-nochange-def apply simp apply auto[1]
  apply(simp add: stable-def)+
  apply(simp add: stable-def Pref-def getrgformula-def Relyf-def) apply auto[1]
  by (simp add: Guarf-def getrgformula-def)

```

**end**

```

theory func-cor-mempoolfree
imports func-cor-lemma
begin

```

## 9 Functional correctness of *k\_mem\_pool\_free*

### 9.1 intermediate conditions and their stable to rely cond

```

abbreviation mp-free-precond1-ext t b  $\equiv$ 
   $\{ \text{pool } b \in \text{'mem-pools} \wedge \text{level } b < \text{length } (\text{levels } (\text{'mem-pool-info } (\text{pool } b)))$ 
   $\wedge \text{block } b < \text{length } (\text{bits } (\text{levels } (\text{'mem-pool-info } (\text{pool } b))! (\text{level } b)))$ 
   $\wedge \text{data } b = \text{block-ptr } (\text{'mem-pool-info } (\text{pool } b)) ((\text{ALIGN}_4 (\text{max-sz } (\text{'mem-pool-info } (\text{pool } b)))) \text{div } (4 \wedge (\text{level } b))) (\text{block } b) \}$ 

```

```

abbreviation mp-free-precond1 t b  $\equiv$ 
  Mem-pool-free-pre t  $\cap$  mp-free-precond1-ext t b

```

```

lemma mp-free-precond1-ext-stb: stable (mp-free-precond1-ext t b) (Mem-pool-free-rely
t)
  apply(simp add: stable-def) apply clarify
  apply(rule conjI) apply(simp add: Mem-pool-free-rely-def gvars-conf-stable-def
gvars-conf-def) apply metis
  apply(rule conjI) apply(simp add: Mem-pool-free-rely-def gvars-conf-stable-def)
unfolding gvars-conf-def apply metis
  apply(rule conjI)
  apply(simp add: Mem-pool-free-rely-def gvars-conf-stable-def) unfolding gvars-conf-def
apply metis
  apply(simp add: block-ptr-def)
  apply(simp add: Mem-pool-free-rely-def gvars-conf-stable-def gvars-conf-def) ap-
ply metis

```

done

**lemma** *mp-free-precond1-stb* : *stable* (*mp-free-precond1* *t b*) (*Mem-pool-free-rely* *t*)  
 apply(*rule stable-int2*)  
 apply(*simp add:mem-pool-free-pre-stb*)  
 apply(*simp add:mp-free-precond1-ext-stb*)  
 done

**abbreviation** *mp-free-precond1-0 t b*  $\equiv$   
 $\{s. \text{inv } s \wedge \text{allocating-node } s \ t = \text{None}\} \cap \text{mp-free-precond1-ext } t \ b$

**lemma** *mp-free-precond1-0-stb*: *stable* (*mp-free-precond1-0* *t b*) (*Mem-pool-free-rely* *t*)  
 apply(*rule stable-int2*)  
 apply(*rule subst*[**where**  $t = \{\text{'inv} \wedge \text{'allocating-node } t = \text{None}\}$   
 and  $s = \{\text{'inv}\} \cap \{\text{'allocating-node } t = \text{None}\}$ ])  
 apply *force*  
 apply(*rule stable-int2*)  
 apply(*simp add:stable-inv-free-rely1*)  
 apply(*simp add:stable-def Mem-pool-free-rely-def*)  
 apply(*simp add:lvars-nochange-rel-def lvars-nochange-def*)  
 apply(*simp add:mp-free-precond1-ext-stb*)  
 done

**abbreviation** *mp-free-precond2-ext t b*  $\equiv \{\text{'freeing-node } t = \text{Some } b\}$

**abbreviation** *mp-free-precond2 t b*  $\equiv$   
 $\text{mp-free-precond1-0 } t \ b \cap \text{mp-free-precond2-ext } t \ b$

**lemma** *mp-free-precond2-ext-stb*: *stable* (*mp-free-precond2-ext* *t b*) (*Mem-pool-free-rely* *t*)

apply(*simp add:stable-def*) apply(*rule allI*) apply(*rule impI*) apply(*rule allI*)  
 apply(*rule impI*)  
 apply(*simp add:Mem-pool-free-rely-def*)  
 apply(*simp add:lvars-nochange-rel-def lvars-nochange-def*) apply *smt*  
 done

**lemma** *mp-free-precond2-stb*: *stable* (*mp-free-precond2* *t b*) (*Mem-pool-free-rely* *t*)  
 apply(*rule stable-int2*)

apply(*simp add:mp-free-precond1-0-stb*)  
 apply(*simp add:mp-free-precond2-ext-stb*)  
 done

**abbreviation** *mp-free-precond3-ext t b*  $\equiv \{\text{'need-resched } t = \text{False}\}$

**abbreviation**  $mp\text{-}free\text{-}precond3\ t\ b \equiv (mp\text{-}free\text{-}precond2\ t\ b) \cap mp\text{-}free\text{-}precond3\text{-}ext\ t\ b$

**lemma**  $mp\text{-}free\text{-}precond3\text{-}ext\text{-}stb : stable\ (mp\text{-}free\text{-}precond3\text{-}ext\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
 $\text{apply}(simp\ add:stable\text{-}def)\ \text{apply}(rule\ allI)\ \text{apply}(rule\ impI)\ \text{apply}(rule\ allI)$   
 $\text{apply}(rule\ impI)$   
 $\text{apply}(simp\ add:Mem\text{-}pool\text{-}free\text{-}rely\text{-}def\ lvars\text{-}nochange\text{-}rel\text{-}def\ lvars\text{-}nochange\text{-}def)$   
**by** *auto*

**lemma**  $mp\text{-}free\text{-}precond3\text{-}stb : stable\ (mp\text{-}free\text{-}precond3\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
 $\text{apply}(rule\ stable\text{-}int2)$   
**using**  $mp\text{-}free\text{-}precond2\text{-}stb$  **apply** *simp*  
**using**  $mp\text{-}free\text{-}precond3\text{-}ext\text{-}stb$  **apply** *simp*  
**done**

**abbreviation**  $mp\text{-}free\text{-}precond4\text{-}ext\ t\ b \equiv \{\!\{ \text{'lsizes}\ t = [ALIGN4\ (max\text{-}sz\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))))] \!\}$

**abbreviation**  $mp\text{-}free\text{-}precond4\ t\ b \equiv mp\text{-}free\text{-}precond3\ t\ b \cap mp\text{-}free\text{-}precond4\text{-}ext\ t\ b$

**lemma**  $mp\text{-}free\text{-}precond4\text{-}ext\text{-}stb :$   
 $stable\ (mp\text{-}free\text{-}precond4\text{-}ext\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
 $\text{apply}(simp\ add:stable\text{-}def)\ \text{apply}(rule\ allI)\ \text{apply}(rule\ impI)\ \text{apply}(rule\ allI)$   
 $\text{apply}(rule\ impI)$   
 $\text{apply}(simp\ add:Mem\text{-}pool\text{-}free\text{-}rely\text{-}def\ ALIGN4\text{-}def)$   
 $\text{apply}(simp\ add:gvars\text{-}conf\text{-}stable\text{-}def\ gvars\text{-}conf\text{-}def)$   
 $\text{apply}(case\text{-}tac\ x = y)\ \text{apply}\ \text{simp}$   
 $\text{apply}\ \text{clarify}\ \text{apply}(simp\ add: lvars\text{-}nochange\text{-}rel\text{-}def\ lvars\text{-}nochange\text{-}def)$   
**done**

**lemma**  $mp\text{-}free\text{-}precond4\text{-}stb : stable\ (mp\text{-}free\text{-}precond4\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
 $\text{apply}(rule\ stable\text{-}int2)$   
**using**  $mp\text{-}free\text{-}precond3\text{-}stb$  **apply** *simp*  
**using**  $mp\text{-}free\text{-}precond4\text{-}ext\text{-}stb$  **apply** *blast*  
**done**

**abbreviation**  $mp\text{-}free\text{-}precond4\text{-}0\text{-}ext\ t\ b \equiv$   
 $\{\!\{ (\forall\ ii < length\ (\text{'lsizes}\ t). \text{'lsizes}\ t\ !\ ii = (ALIGN4\ (max\text{-}sz\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))))\ \text{div}\ (4\ ^\ ii))$   
 $\wedge\ length\ (\text{'lsizes}\ t) > 0 \!\}$

**abbreviation**  $mp\text{-}free\text{-}precond4\text{-}0\ t\ b \equiv mp\text{-}free\text{-}precond3\ t\ b \cap mp\text{-}free\text{-}precond4\text{-}0\text{-}ext\ t\ b$

**lemma**  $mp\text{-}free\text{-}precond4\text{-}0\text{-}ext\text{-}stb :$   
 $stable\ (mp\text{-}free\text{-}precond4\text{-}0\text{-}ext\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
 $\text{apply}(simp\ add:stable\text{-}def)\ \text{apply}(rule\ allI)\ \text{apply}(rule\ impI)\ \text{apply}(rule\ allI)$

```

apply(rule impI)
  apply(simp add:Mem-pool-free-rely-def ALIGN4-def)
  apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(case-tac x = y) apply simp
  apply clarify apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
done

lemma mp-free-precond4-0-stb : stable (mp-free-precond4-0 t b) (Mem-pool-free-rely
t)
apply(rule stable-int2)
  using mp-free-precond3-stb apply simp
  using mp-free-precond4-0-ext-stb apply blast
done

abbreviation mp-free-precond4-1 t b  $\equiv$ 
  mp-free-precond4-0 t b  $\cap$   $\{\text{length } ('lsizes\ t) = 'i\ t\}$ 

lemma mp-free-precond4-1-stb : stable (mp-free-precond4-1 t b) (Mem-pool-free-rely
t)
apply(rule stable-int2)
  using mp-free-precond4-0-stb apply simp
  apply(simp add:stable-def) apply(rule allI) apply(rule impI) apply(rule allI)+
apply(rule impI)
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
done

abbreviation mp-free-precond4-2 t b  $\equiv$ 
  mp-free-precond4-1 t b  $\cap$   $\{\text{'i } t \leq \text{level } b\}$ 

lemma mp-free-precond4-2-stb : stable (mp-free-precond4-2 t b) (Mem-pool-free-rely
t)
apply(rule stable-int2)
  using mp-free-precond4-1-stb apply simp
  apply(simp add:stable-def) apply(rule allI) apply(rule impI) apply(rule allI)+
apply(rule impI)
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
by smt

abbreviation mp-free-precond4-3 t b  $\equiv$ 
  mp-free-precond4-0 t b  $\cap$   $(\{\text{'i } t \leq \text{level } b\} \cap \{\text{length } ('lsizes\ t) = 'i\ t + 1\})$ 

lemma mp-free-precond4-3-stb : stable (mp-free-precond4-3 t b) (Mem-pool-free-rely
t)
apply(rule stable-int2)
  using mp-free-precond4-0-stb apply simp
  apply(simp add:stable-def) apply(rule allI) apply(rule impI) apply(rule allI)+
apply(rule impI)
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)

```

by *smt*

**abbreviation** *mp-free-precond5-ext* *t b*  $\equiv$   
 $\{\!\{(\forall ii < \text{length } ('lsizes\ t).\ 'lsizes\ t\ !\ ii = (\text{ALIGN}_4\ (\text{max-sz } ('mem\text{-pool-info } (pool\ b)))) \div (4 \wedge ii))$   
 $\wedge \text{length } ('lsizes\ t) > \text{level } b\}\!\}$

**abbreviation** *mp-free-precond5* *t b*  $\equiv$  *mp-free-precond3* *t b*  $\cap$  *mp-free-precond5-ext* *t b*

**term** *mp-free-precond5* *t b*

**lemma** *mp-free-precond5-ext-stb*:  
 $\text{stable } (mp\text{-free-precond5-ext } t\ b) \text{ (Mem-pool-free-rely } t)$   
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-free-rely-def ALIGN<sub>4</sub>-def*)  
**apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)  
**apply**(*case-tac x = y*) **apply** *simp*  
**apply** *clarify* **apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def*)  
**done**

**lemma** *mp-free-precond5-stb* :  $\text{stable } (mp\text{-free-precond5 } t\ b) \text{ (Mem-pool-free-rely } t)$   
**apply**(*rule stable-int2*)  
**using** *mp-free-precond3-stb* **apply** *simp*  
**using** *mp-free-precond5-ext-stb* **apply** *blast*  
**done**

**abbreviation** *mp-free-precond6* *t b*  $\equiv$   
 $mp\text{-free-precond5 } t\ b \cap \{\!\{'free\text{-block-r } t = \text{True}\}\!\}$

**lemma** *mp-free-precond6-stb* :  $\text{stable } (mp\text{-free-precond6 } t\ b) \text{ (Mem-pool-free-rely } t)$

**apply**(*rule stable-int2*)  
**using** *mp-free-precond5-stb* **apply** *simp*  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*) +  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**by** *auto*

**abbreviation** *mp-free-precond7* *t b*  $\equiv$   
 $mp\text{-free-precond6 } t\ b \cap \{\!\{'bn\ t = \text{block } b\}\!\}$

**lemma** *mp-free-precond7-stb* :  $\text{stable } (mp\text{-free-precond7 } t\ b) \text{ (Mem-pool-free-rely } t)$

**apply**(*rule stable-int2*)  
**using** *mp-free-precond6-stb* **apply** *simp*  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*) +



**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**by** *smt*

**abbreviation** *mp-free-precond8 t b*  $\equiv$   

$$\begin{aligned} & mp\text{-}free\text{-}precond1\text{-}0\ t\ b \cap \{\!\{ \text{level } b < \text{length } ('lsizes\ t) \\ & \wedge (\forall ii < \text{length } ('lsizes\ t). 'lsizes\ t\ !\ ii = (\text{ALIGN4 } (\text{max-sz } ('mem\text{-}pool\text{-}info \\ & (\text{pool } b)))) \text{div } (4 \wedge ii)) \\ & \wedge 'bn\ t < \text{length } (\text{bits } (\text{levels } ('mem\text{-}pool\text{-}info\ (\text{pool } b))!('lvl\ t))) \\ & \wedge 'bn\ t = (\text{block } b) \text{div } (4 \wedge (\text{level } b - 'lvl\ t)) \\ & \wedge 'lvl\ t \leq \text{level } b \\ & \wedge ('free\text{-}block\text{-}r\ t \longrightarrow \\ & (\exists \text{blk}. 'freeing\text{-}node\ t = \text{Some } \text{blk} \wedge \text{pool } \text{blk} = \text{pool } b \wedge \text{level } \text{blk} = 'lvl\ t \\ & \wedge \text{block } \text{blk} = 'bn\ t) \\ & \wedge 'alloc\text{-}memblk\text{-}data\text{-}valid\ (\text{pool } b)\ (\text{the } ('freeing\text{-}node\ t))) \\ & \wedge (\neg 'free\text{-}block\text{-}r\ t \longrightarrow 'freeing\text{-}node\ t = \text{None}) \\ & (* \wedge ((\text{if } 'freeing\text{-}node\ t \neq \text{None} \text{ then } 'lvl\ t + 1 \text{ else } 0) > 0 \\ & \longrightarrow 'free\text{-}block\text{-}r\ t) *) (* \text{this cond is implied by upper conds } *) \} \} \end{aligned}$$

**abbreviation** *mp-free-precond8-inv t b*  $\alpha \equiv$   

$$mp\text{-}free\text{-}precond8\ t\ b \cap \{\!\{ \alpha = (\text{if } 'freeing\text{-}node\ t \neq \text{None} \text{ then } 'lvl\ t + 1 \text{ else } 0) \}\!\}$$

**lemma** *inv- $\alpha$ gt0-imp-looppre*:  
 $mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ \alpha \cap \{\!\{ \alpha > 0 \}\!\} \subseteq mp\text{-}free\text{-}precond8\ t\ b \cap \{\!\{ 'free\text{-}block\text{-}r\ t \}\!\}$   
**by** *auto*

**lemma** *looppre-imp-exist- $\alpha$ gt0*:  
 $x \in mp\text{-}free\text{-}precond8\ t\ b \cap \{\!\{ 'free\text{-}block\text{-}r\ t \}\!\} \implies \exists \alpha. x \in mp\text{-}free\text{-}precond8\text{-}inv\ t\ b$   
 $\alpha \cap \{\!\{ \alpha > 0 \}\!\}$   
**by** *clarsimp*

**lemma**  $x \in mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ \alpha \cap \{\!\{ \alpha > 0 \}\!\} \implies x \in mp\text{-}free\text{-}precond8\ t\ b$   
 $\cap \{\!\{ 'free\text{-}block\text{-}r\ t \}\!\}$   
**using** *inv- $\alpha$ gt0-imp-looppre*[*of t b  $\alpha$* ]  
 $\text{subsetI}[\text{of } mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ \alpha \cap \{\!\{ \alpha > 0 \}\!\}$   
 $mp\text{-}free\text{-}precond8\ t\ b \cap \{\!\{ 'free\text{-}block\text{-}r\ t \}\!\}]$   
**by** *blast*

**lemma** *loopbody-sat-invterm-imp-inv-post*:  
 $\Gamma \vdash_I P\ sat_p [mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ \alpha \cap \{\!\{ \alpha > 0 \}\!\}, \text{rely}, \text{guar}, mp\text{-}free\text{-}precond8\text{-}inv$   
 $t\ b\ (\alpha - 1)]$   
 $\implies \Gamma \vdash_I P\ sat_p [mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ \alpha \cap \{\!\{ \alpha > 0 \}\!\}, \text{rely}, \text{guar}, mp\text{-}free\text{-}precond8$   
 $t\ b]$   
**using** *Conseq* [*of mp-free-precond8-inv t b  $\alpha \cap \{\!\{ \alpha > 0 \}\!\}$*  mp-free-precond8-inv t b

$\alpha \cap \{\alpha > 0\}$   
 $\text{rely rely guar guar mp-free-precond8-inv } t \text{ } b \text{ } (\alpha - 1)$   
 $\text{mp-free-precond8 } t \text{ } b \text{ } P] \text{ by blast}$

**lemma** *stm8-inv-imp-prepost*:

$(\forall \alpha. \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8-inv } t \text{ } b \text{ } (\alpha - 1)])$   
 $\implies \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$

**apply**(*rule subst*[**where**  $s = \forall v. v \in \text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\} \longrightarrow$   
 $\Gamma \vdash_I P \text{ sat}_p [\{v\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$  **and**  
 $t = \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$ ])  
**using** *allpre-eq-pre*[*of*  $\text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\}$   
 $P \text{ rely guar mp-free-precond8 } t \text{ } b]$  **apply blast**

**apply**(*rule allI*) **apply**(*rule impI*)  
**apply**(*subgoal-tac*  $\exists \alpha. v \in \text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}$ )  
**prefer 2 using** *looppre-imp-exist- $\alpha$ gt0* **apply blast**

**apply**(*erule exE*)

**using** *sat-pre-imp-allinpre*[*of*  $P - \text{rely guar mp-free-precond8 } t \text{ } b]$   
*loopbody-sat-invterm-imp-inv-post* **apply blast**  
**done**

**lemma** *loopbody-sat-invterm-imp-inv-post2*:

$\exists \beta < \alpha. \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8-inv } t \text{ } b \text{ } \beta]$   
 $\implies \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$   
**using** *Conseq* [*of*  $\text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}$   $\text{mp-free-precond8-inv } t \text{ } b$   
 $\alpha \cap \{\alpha > 0\}$   
 $\text{rely rely guar guar mp-free-precond8-inv } t \text{ } b -$   
 $\text{mp-free-precond8 } t \text{ } b \text{ } P]$  **by blast**

**lemma** *stm8-inv-imp-prepost2*:

$(\forall \alpha. \exists \beta < \alpha. \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ } b \text{ } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8-inv } t \text{ } b \text{ } \beta])$   
 $\implies \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$

**apply**(*rule subst*[**where**  $s = \forall v. v \in \text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\} \longrightarrow$   
 $\Gamma \vdash_I P \text{ sat}_p [\{v\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$  **and**  
 $t = \Gamma \vdash_I P \text{ sat}_p [\text{mp-free-precond8 } t \text{ } b \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \text{ } b]$ ])

```

using allpre-eq-pre[of mp-free-precond8 t b  $\cap \{\neg \text{'free-block-r } t\}$ 
    P rely guar mp-free-precond8 t b] apply blast

apply(rule allI) apply(rule impI)
apply(subgoal-tac  $\exists \alpha. v \in \text{mp-free-precond8-inv } t \text{ b } \alpha \cap \{\alpha > 0\}$ )
  prefer 2 using looppre-imp-exist- $\alpha$ gt0 apply blast

apply(erule exE)

using sat-pre-imp-allinpre[of P - rely guar mp-free-precond8 t b]
  loopbody-sat-invterm-imp-inv-post apply blast
done

lemma stm8-loopinv0: mp-free-precond8-inv t b 0  $\subseteq \{\neg \text{'free-block-r } t\}$ 
by auto

lemma stm8-loopinv- $\alpha$ :  $\alpha > 0 \implies \text{mp-free-precond8-inv } t \text{ b } \alpha \subseteq \{\neg \text{'free-block-r } t\}$ 
by auto

lemma inv- $\alpha$ eq0-eq-looppre:
  mp-free-precond8-inv t b 0 = mp-free-precond8 t b  $\cap \{\neg \text{'free-block-r } t\}$ 
by auto

term mp-free-precond8 t b

lemma alloc-memblk-data-valid-stb-free:
  alloc-memblk-data-valid x (pool b) (the (freeing-node x t))  $\implies$ 
    (x, y)  $\in$  lvars-nochange-rel t  $\implies$ 
    (x, y)  $\in$  gvars-conf-stable  $\implies$ 
    alloc-memblk-data-valid y (pool b) (the (freeing-node y t))
  apply(subgoal-tac blk x t = blk y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
  apply(subgoal-tac buf (mem-pool-info x (pool b)) = buf (mem-pool-info y (pool
b)))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac lsize x t = lsize y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
  apply(subgoal-tac free-l x t = free-l y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
  apply(subgoal-tac max-sz (mem-pool-info x (pool b)) = max-sz (mem-pool-info y
(pool b)))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac freeing-node x t = freeing-node y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
  apply (simp add: gvars-conf-def gvars-conf-stable-def)
done

lemma mp-free-precond8-stb : stable (mp-free-precond8 t b) (Mem-pool-free-rely t)

```

```

apply(rule stable-int2) apply(rule stable-int2)

apply(simp add:stable-def)
apply clarify
apply(rule conjI)
  using stable-inv-free-rely apply blast
  apply(simp add: Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt

apply(simp add:stable-def)
apply clarify
apply(rule conjI)
  apply(simp add:Mem-pool-free-rely-def gvars-conf-stable-def gvars-conf-def) ap-
plymetis
  apply(rule conjI)
  apply(simp add:Mem-pool-free-rely-def gvars-conf-stable-def gvars-conf-def) ap-
plymetis
  apply(rule conjI)
  apply(simp add:Mem-pool-free-rely-def gvars-conf-stable-def gvars-conf-def) ap-
plymetis
  apply(simp add: block-ptr-def ALIGN4-def lvars-nochange-rel-def lvars-nochange-def
    gvars-conf-stable-def gvars-conf-def)
  apply(simp add:Mem-pool-free-rely-def gvars-conf-stable-def gvars-conf-def) ap-
plymetis

apply(simp add: Mem-pool-free-rely-def stable-def)
apply clarify
apply(rule conjI) apply clarify
apply(rule conjI)
  apply(simp add: gvars-conf-stable-def gvars-conf-def lvars-nochange-rel-def lvars-nochange-def)
apply(rule conjI) apply(simp add: ALIGN4-def lvars-nochange-rel-def lvars-nochange-def
  gvars-conf-stable-def gvars-conf-def)
apply(rule conjI) apply(simp add: ALIGN4-def lvars-nochange-rel-def lvars-nochange-def
  gvars-conf-stable-def gvars-conf-def) applymetis
  apply(rule conjI) apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
applymetis
  apply(rule conjI) apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
  apply(rule conjI) apply clarify
  apply(rule conjI) apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
applymetis
  apply(simp add: ALIGN4-def lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
  gvars-conf-def)

apply clarify apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply clarify
done

```

```

lemma mp-free-precond8-inv-stb : stable (mp-free-precond8-inv t b  $\alpha$ ) (Mem-pool-free-rely
t)
  apply(rule stable-int2)
  using mp-free-precond8-stb apply fast
  apply(unfold stable-def) apply clarify

  apply(subgoal-tac lvl x t = lvl y t) prefer 2
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac freeing-node x t = freeing-node y t) prefer 2
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
by simp

```

```

lemma mp-free-precond8-inv-presv-rely:
s ∈ mp-free-precond8-inv t b  $\alpha \implies (s,r) \in \text{Mem-pool-free-rely } t \implies \exists \beta \leq \alpha. r \in \text{mp-free-precond8-inv}
t b  $\beta$ 
apply(rule exI[where x= $\alpha$ ])
apply(rule conjI) apply fast
using mp-free-precond8-inv-stb[of t b  $\alpha$ ] apply(unfold stable-def) apply blast
done$ 
```

```

abbreviation mp-free-precond8-1 t b  $\alpha \equiv$ 
  mp-free-precond8-inv t b  $\alpha \cap \{\alpha > 0\}$ 

```

```

lemma mp-free-precond8-1-imp-free-block-r:
mp-free-precond8-1 t b  $\alpha \subseteq \{\text{'free-block-r } t\}$ 
  using stm8-loopinv- $\alpha$  by blast

```

```

lemma mp-free-precond8-1-stb : stable (mp-free-precond8-1 t b  $\alpha$ ) (Mem-pool-free-rely
t)
  apply(rule stable-int2)
  using mp-free-precond8-inv-stb apply blast
  apply(simp add:stable-def)
done

```

```

abbreviation mp-free-precond8-1' t b  $\equiv$ 
  mp-free-precond8 t b  $\cap \{\text{'free-block-r } t\}$ 

```

```

lemma mp-free-precond8-1'-stb : stable (mp-free-precond8-1' t b) (Mem-pool-free-rely
t)
  apply(rule stable-int2)
  using mp-free-precond8-stb apply blast
  apply(simp add:stable-def) apply clarify
  apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)
by smt

```

**abbreviation**  $mp\text{-}free\text{-}precond8\text{-}2\ t\ b\ \alpha \equiv$   
 $mp\text{-}free\text{-}precond8\text{-}1\ t\ b\ \alpha \cap \{\!| 'lsz\ t = 'lsizes\ t\ !\ ('lvl\ t)\ |\}$

**lemma**  $mp\text{-}free\text{-}precond8\text{-}2\text{-}stb : stable\ (mp\text{-}free\text{-}precond8\text{-}2\ t\ b\ \alpha)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
`apply(rule stable-int2)`  
`using mp-free-precond8-1-stb apply blast`  
`apply(simp add:stable-def) apply clarify`  
`apply(simp add:Mem-pool-free-rely-def lvars-nochange-rel-def lvars-nochange-def)`  
`by smt`

**abbreviation**  $mp\text{-}free\text{-}precond8\text{-}3\ t\ b\ \alpha \equiv$   
 $mp\text{-}free\text{-}precond8\text{-}2\ t\ b\ \alpha \cap \{\!| 'blk\ t = block\text{-}ptr\ ('mem\text{-}pool\text{-}info\ (pool\ b)) ('lsz\ t)\ |\}$   
 $( 'bn\ t)\ |\}$

**lemma**  $mp\text{-}free\text{-}precond8\text{-}3\text{-}stb : stable\ (mp\text{-}free\text{-}precond8\text{-}3\ t\ b\ \alpha)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
`apply(rule stable-int2)`  
`using mp-free-precond8-2-stb apply blast`  
`apply(simp add:stable-def block-ptr-def Mem-pool-free-rely-def) apply clarify`  
`apply(case-tac x = y) apply simp apply clarsimp`  
`apply(subgoal-tac blk x t = blk y t)`  
`apply(subgoal-tac lsz x t = lsz y t)`  
`apply(subgoal-tac bn x t = bn y t)`  
`apply(subgoal-tac buf (mem-pool-info x (pool b)) = buf (mem-pool-info y (pool b)))`  
`apply simp`  
`apply(simp add: gvars-conf-stable-def gvars-conf-def)`  
`apply(simp add: lvars-nochange-rel-def lvars-nochange-def)`  
`apply(simp add: lvars-nochange-rel-def lvars-nochange-def)`  
`apply(simp add: lvars-nochange-rel-def lvars-nochange-def) applymetis`  
`done`

**abbreviation**  $mp\text{-}free\text{-}precond9\ t\ b \equiv mp\text{-}free\text{-}precond1\ t\ b$   
**term**  $mp\text{-}free\text{-}precond1\ t\ b$

**lemma**  $mp\text{-}free\text{-}precond9\text{-}stb : stable\ (mp\text{-}free\text{-}precond9\ t\ b)\ (Mem\text{-}pool\text{-}free\text{-}rely\ t)$   
`using mp-free-precond1-stb apply auto[1]`  
`done`

## 9.2 proof of each statement

**lemma**  $mempool\text{-}free\text{-}stm1\text{-}inv\text{-}mempool\text{-}info :$   
 $inv\text{-}mempool\text{-}info\ Va \wedge inv\text{-}bitmap\text{-}freelist\ Va \implies$   
 $block\ b < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ (pool\ b))\ !\ level\ b)) \implies$   
 $level\ b < length\ (levels\ (mem\text{-}pool\text{-}info\ Va\ (pool\ b))) \implies$   
 $pool\ b \in mem\text{-}pools\ Va \implies$   
 $get\text{-}bit\ (mem\text{-}pool\text{-}info\ Va)\ (pool\ b)\ (level\ b)\ (block\ b) = ALLOCATED \implies$

```

inv-mempool-info
  (Va(|mem-pool-info := (mem-pool-info Va)
    (pool b := mem-pool-info Va (pool b)
      (|levels := levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (|bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING])])]),
    freeing-node := freeing-node Va(t ↦ b)))
apply(simp add:inv-mempool-info-def)
apply(rule conjI) apply metis
apply(rule conjI) apply metis
apply(rule conjI) apply metis
apply(rule conjI) apply(simp add:inv-bitmap-freelist-def) apply (simp add:Let-def)
apply auto[1]
apply(rule conjI) apply(simp add:inv-bitmap-freelist-def) apply (simp add:Let-def)
  apply(rule allI) apply(rule impI)
  apply(subgoal-tac (∀ i < length (levels (mem-pool-info Va (pool b))))).
    length (bits (levels (mem-pool-info Va (pool b)) ! i)) = n-max (mem-pool-info
Va (pool b)) * 4 ^ i)
  apply(case-tac i = level b)
  apply auto[1] apply auto[1]
  apply(simp add:Let-def)
done

```

**lemma** mempool-free-stm1-inv-bitmap-freelist:

```

inv-cur Va ∧ inv-thd-waitq Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist Va ∧
inv-bitmap Va ∧ inv-aux-vars Va ⇒
  block b < length (bits (levels (mem-pool-info Va (pool b)) ! level b)) ⇒
  level b < length (levels (mem-pool-info Va (pool b))) ⇒
  pool b ∈ mem-pools Va ⇒
  get-bit (mem-pool-info Va) (pool b) (level b) (block b) = ALLOCATED ⇒
  inv-bitmap-freelist
  (Va(|mem-pool-info := (mem-pool-info Va)
    (pool b := mem-pool-info Va (pool b)
      (|levels := levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (|bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING])])]),
    freeing-node := freeing-node Va(t ↦ b)))
apply(simp add:inv-bitmap-freelist-def)
apply(rule allI) apply(rule impI)+ apply(simp add:Let-def)
apply(rule conjI) apply(rule allI) apply(rule impI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1] apply fastforce
  apply(case-tac i ≠ level b) apply auto[1]
  apply(case-tac j ≠ block b) apply auto[1]
  apply auto[1]

apply(rule conjI) apply(rule allI) apply(rule impI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]

```

```

apply(case-tac  $i \neq \text{level } b$ ) apply auto[1]
apply(case-tac  $j \neq \text{block } b$ ) apply auto[1]
apply auto[1]
apply(simp add:distinct-def)
apply(case-tac  $i = \text{level } b$ ) apply auto[1]
apply auto[1]
done

lemma mempool-free-stm1-inv-bitmap:
  inv-cur  $Va \wedge \text{inv-thd-waitq } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist } Va \wedge$ 
  inv-bitmap  $Va \wedge \text{inv-aux-vars } Va \implies$ 
     $\text{block } b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b)) \implies$ 
     $\text{level } b < \text{length } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b))) \implies$ 
     $\text{pool } b \in \text{mem-pools } Va \implies$ 
     $\text{get-bit } (\text{mem-pool-info } Va) \text{ (pool } b) \text{ (level } b) \text{ (block } b) = \text{ALLOCATED} \implies$ 
    inv-bitmap
    ( $Va \langle \text{mem-pool-info} := (\text{mem-pool-info } Va)$ 
      ( $\text{pool } b := \text{mem-pool-info } Va \text{ (pool } b)$ 
        ( $\text{levels} := \text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$ 
          ( $\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b)$ 
            ( $\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b) \text{ [block } b$ 
               $:= \text{FREEING}] \rangle \rangle \rangle$ ),
      freeing-node  $:= \text{freeing-node } Va(t \mapsto b) \rangle \rangle$ )
    apply(simp add:inv-bitmap-def)
    apply(rule allI) apply(simp add:Let-def) apply(rule impI) apply(rule allI)
apply(rule impI)
    apply(rule conjI) apply(rule impI)
    apply(rule conjI)
    apply(case-tac  $i = \text{level } b \wedge j = \text{block } b$ ) apply auto[1]
    apply(case-tac  $i - 1 = \text{level } b \wedge j \text{ div } 4 = \text{block } b$ )
    apply (metis (no-types, lifting) BlockState.distinct(3) One-nat-def Suc-pred
      lessI nat-neq-iff nth-list-update-neq)
    apply(rule impI)
    apply(subgoal-tac  $\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$ 
      ( $\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b)$ 
        ( $\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b) \text{ [block } b$ 
           $:= \text{FREEING}] \rangle \rangle \text{ ! } i \text{ ! } j$ 
             $= \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! } i \text{ ! } j)$ 
          apply(subgoal-tac  $\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$ 
            ( $\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b)$ 
              ( $\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! level } b) \text{ [block } b$ 
                 $:= \text{FREEING}] \rangle \rangle \text{ ! } (i - \text{Suc } \text{NULL}) \text{ ! } (j \text{ div } 4)$ 
                   $= \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) \text{ ! } (i - \text{Suc } \text{NULL})) \text{ ! } (j$ 
                     $\text{div } 4)$ 
                apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

                  length-list-update nth-list-update-eq nth-list-update-neq)
                apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

```



```

      One-nat-def nth-list-update-eq nth-list-update-neq)
    apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

    apply(rule impI)
    apply(rule conjI)
      apply(case-tac i = level b ∧ j = block b) apply auto[1]
      apply(case-tac Suc i = level b ∧ j * 4 = block b)
    apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
      apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
          = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
      apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! (j * 4)
          = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4))
    apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
    apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
    apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
    apply(rule conjI)
      apply(case-tac i = level b ∧ j = block b) apply auto[1]
      apply(case-tac Suc i = level b ∧ Suc (j * 4) = block b)
    apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
      apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
          = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
      apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (j * 4)
          = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (j * 4))
    apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
    apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
    apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

```

```

apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (Suc (j * 4)) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)][block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)][block b
:= FREEING]] ! (Suc i)) ! Suc (Suc (j * 4))
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (Suc (j *
4))))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ (j * 4 + 3) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)][block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)][block b
:= FREEING]] ! (Suc i)) ! (j * 4 + 3)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4 + 3))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

apply(rule conjI) apply(rule impI)
apply(rule conjI) apply(rule impI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]

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apply(case-tac  $i - 1 = \text{level } b \wedge j \text{ div } 4 = \text{block } b$ )
apply (metis (no-types, lifting) BlockState.distinct(3) One-nat-def Suc-pred
lessI nat-neq-iff nth-list-update-neq)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
[level b := (levels (mem-pool-info Va (pool b)) ! level b)
(bits := bits (levels (mem-pool-info Va (pool b)) ! level b)]block b
:= FREEING]] ! i ! j
= bits (levels (mem-pool-info Va (pool b)) ! i ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
[level b := (levels (mem-pool-info Va (pool b)) ! level b)
(bits := bits (levels (mem-pool-info Va (pool b)) ! level b)]block b
:= FREEING]] ! (i - Suc NULL) ! (j div 4)
= bits (levels (mem-pool-info Va (pool b)) ! (i - Suc NULL) ! (j
div 4))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

One-nat-def nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule impI)
apply(rule conjI)
apply(case-tac  $i = \text{level } b \wedge j = \text{block } b$ ) apply auto[1]
apply(case-tac  $\text{Suc } i = \text{level } b \wedge j * 4 = \text{block } b$ )
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
[level b := (levels (mem-pool-info Va (pool b)) ! level b)
(bits := bits (levels (mem-pool-info Va (pool b)) ! level b)]block b
:= FREEING]] ! i ! j
= bits (levels (mem-pool-info Va (pool b)) ! i ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
[level b := (levels (mem-pool-info Va (pool b)) ! level b)
(bits := bits (levels (mem-pool-info Va (pool b)) ! level b)]block b
:= FREEING]] ! (Suc i) ! (j * 4)
= bits (levels (mem-pool-info Va (pool b)) ! (Suc i) ! (j * 4))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule conjI)
apply(case-tac  $i = \text{level } b \wedge j = \text{block } b$ ) apply auto[1]
apply(case-tac  $\text{Suc } i = \text{level } b \wedge \text{Suc } (j * 4) = \text{block } b$ )
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)

```

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apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
  [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
  = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
  [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (j * 4)
  = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (j * 4))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
  Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule conjI)
apply(case-tac i = level b ∧ j = block b) apply auto[1]
apply(case-tac Suc i = level b ∧ Suc (Suc (j * 4)) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
  [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
  = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
  [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (Suc (j * 4))
  = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (Suc (j *
4))))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
  Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

apply(case-tac i = level b ∧ j = block b) apply auto[1]
apply(case-tac Suc i = level b ∧ (j * 4 + 3) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
  [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
  = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)

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      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! (j * 4 + 3))
      = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4 + 3))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
  Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

apply(rule conjI) apply(rule impI)
apply(rule conjI) apply(rule impI)
apply(case-tac i = level b ∧ j = block b) apply auto[1]
apply(case-tac i - 1 = level b ∧ j div 4 = block b)
apply (metis (no-types, lifting) BlockState.distinct(3) One-nat-def Suc-pred
lessI nat-neq-iff nth-list-update-neq)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (i - Suc NULL)) ! (j div 4)
      = bits (levels (mem-pool-info Va (pool b)) ! (i - Suc NULL)) ! (j
div 4))
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

One-nat-def nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
  Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule impI)
apply(rule conjI)
apply(case-tac i = level b ∧ j = block b) apply auto[1]
apply(case-tac Suc i = level b ∧ j * 4 = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b

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:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
  apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (j * 4) = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
  apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (Suc (j * 4)) = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b

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:= FREEING]] ! (Suc i) ! Suc (Suc (j * 4))
      = bits (levels (mem-pool-info Va (pool b)) ! (Suc i) ! Suc (Suc (j *
4)))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ (j * 4 + 3) = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
      [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
      [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i) ! (j * 4 + 3)
      = bits (levels (mem-pool-info Va (pool b)) ! (Suc i) ! (j * 4 + 3))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(rule conjI) apply(rule impI)
  apply(rule conjI) apply(rule impI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac i - 1 = level b ∧ j div 4 = block b)
  apply (metis (no-types, lifting) BlockState.distinct(3) One-nat-def Suc-pred
lessI nat-neq-iff nth-list-update-neq)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
      [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
      [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b

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:= FREEING]] ! (i - Suc NULL) ! (j div 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (i - Suc NULL) ! (j
div 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    One-nat-def nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
  apply(rule impI)
  apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ j * 4 = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    length-list-update nth-list-update-eq nth-list-update-neq)
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

    nth-list-update-eq nth-list-update-neq)
  apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
    Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
  apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (j * 4) = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

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length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule conjI)
  apply(case-tac i = level b  $\wedge$  j = block b) apply auto[1]
  apply(case-tac Suc i = level b  $\wedge$  Suc (Suc (j * 4)) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! Suc (Suc (j * 4))
      = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (Suc (j *
4)))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(case-tac i = level b  $\wedge$  j = block b) apply auto[1]
  apply(case-tac Suc i = level b  $\wedge$  (j * 4 + 3) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! i) ! j
      = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
      ( $\lambda$ bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]] ! (Suc i)) ! (j * 4 + 3)
      = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4 + 3))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

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      nth-list-update-eq nth-list-update-neq)
    apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(rule conjI)
  apply(rule impI)+
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac i - 1 = level b ∧ j div 4 = block b)
  apply (metis (no-types, lifting) BlockState.distinct(3) One-nat-def Suc-pred
lessI nat-neq-iff nth-list-update-neq)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! (i - Suc NULL)) ! (j div 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (i - Suc NULL)) ! (j div
4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
    apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      One-nat-def nth-list-update-eq nth-list-update-neq)
    apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)

  apply(rule conjI)
  apply(rule impI)+
  apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ j * 4 = block b)
  apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! (Suc i)) ! (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)

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apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (j * 4) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! (Suc i)) ! Suc (j * 4)
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (j * 4))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)
      Mem-pool-lvl.simps(4) nth-list-update-eq nth-list-update-neq)
apply(rule conjI)
  apply(case-tac i = level b ∧ j = block b) apply auto[1]
  apply(case-tac Suc i = level b ∧ Suc (Suc (j * 4)) = block b)
apply (metis BlockState.distinct(5) less-Suc-eq nth-list-update-neq order-less-irrefl)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! i) ! j
    = bits (levels (mem-pool-info Va (pool b)) ! i) ! j)
  apply(subgoal-tac bits (levels (mem-pool-info Va (pool b))
    [level b := (levels (mem-pool-info Va (pool b)) ! level b)
    (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b :=
FREEING]] ! (Suc i)) ! Suc (Suc (j * 4))
    = bits (levels (mem-pool-info Va (pool b)) ! (Suc i)) ! Suc (Suc (j *
4))))
  apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      length-list-update nth-list-update-eq nth-list-update-neq)
apply (metis Mem-pool-lvl.cases Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4)

      nth-list-update-eq nth-list-update-neq)
apply (metis (no-types, lifting) Mem-pool-lvl.cases Mem-pool-lvl.simps(1)

```

$Mem\text{-}pool\text{-}lvl.simps(4) \text{ } nth\text{-}list\text{-}update\text{-}eq \text{ } nth\text{-}list\text{-}update\text{-}neq$   
**apply**(*case-tac*  $i = level \ b \wedge j = block \ b$ ) **apply** *auto*[1]  
**apply**(*case-tac*  $Suc \ i = level \ b \wedge (j * 4 + 3) = block \ b$ )  
**apply** (*metis* *BlockState.distinct*(5) *less-Suc-eq* *nth-list-update-neq* *order-less-irrefl*)  
**apply**(*subgoal-tac* *bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  
 $[level \ b := (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)$   
 $(\downarrow bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)[block \ b :=$   
 $FREEING]) \ ! \ i) \ ! \ j$   
 $= bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ i) \ ! \ j$ )  
**apply**(*subgoal-tac* *bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  
 $[level \ b := (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)$   
 $(\downarrow bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)[block \ b :=$   
 $FREEING]) \ ! \ (Suc \ i) \ ! \ (j * 4 + 3)$   
 $= bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ (Suc \ i) \ ! \ (j * 4 + 3))$ )  
**apply** (*metis* *Mem-pool-lvl.cases* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4)  
 $length\text{-}list\text{-}update \ nth\text{-}list\text{-}update\text{-}eq \ nth\text{-}list\text{-}update\text{-}neq$ )  
**apply** (*metis* *Mem-pool-lvl.cases* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4)  
 $nth\text{-}list\text{-}update\text{-}eq \ nth\text{-}list\text{-}update\text{-}neq$ )  
**apply** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.cases* *Mem-pool-lvl.simps*(1)  
 $Mem\text{-}pool\text{-}lvl.simps(4) \text{ } nth\text{-}list\text{-}update\text{-}eq \ nth\text{-}list\text{-}update\text{-}neq$ )  
**apply**(*rule impI*)+  
**apply**(*case-tac*  $i = level \ b \wedge j = block \ b$ ) **apply** *auto*[1]  
**apply**(*case-tac*  $i - 1 = level \ b \wedge j \text{ div } 4 = block \ b$ ) **apply** *auto*[1]  
**apply**(*subgoal-tac* *bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  
 $[level \ b := (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)$   
 $(\downarrow bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)[block \ b :=$   
 $FREEING]) \ ! \ i) \ ! \ j$   
 $= bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ i) \ ! \ j$ )  
**prefer 2 apply** (*metis* (*no-types*, *lifting*) *Mem-pool-lvl.cases* *Mem-pool-lvl.simps*(1)  
 $Mem\text{-}pool\text{-}lvl.simps(4) \text{ } nth\text{-}list\text{-}update\text{-}eq \ nth\text{-}list\text{-}update\text{-}neq$ )  
**apply**(*subgoal-tac* *bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  
 $[level \ b := (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)$   
 $(\downarrow bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)[block \ b :=$   
 $FREEING]) \ ! \ (i - 1) \ ! \ (j \text{ div } 4)$   
 $= bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ (i - 1) \ ! \ (j \text{ div } 4))$ )  
**prefer 2 apply** (*metis* *Mem-pool-lvl.cases* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4)  
 $nth\text{-}list\text{-}update\text{-}eq \ nth\text{-}list\text{-}update\text{-}neq$ )  
**apply**(*subgoal-tac* *bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  $! \ (i - Suc \ NULL)$ )  
 $! \ (j \text{ div } 4) \neq DIVIDED$ )  
**prefer 2 apply**(*subgoal-tac* *length* (*bits* (*levels* (*mem-pool-info* *Va* (*pool* *b*))  
 $[level \ b := (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level \ b)$   
 $(\downarrow bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ Va \ (pool \ b)) \ ! \ level$   
 $b)[block \ b := FREEING]) \ !$

$i)) = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! i)))$   
**prefer 2 apply**(*case-tac*  $i = \text{level } b$ )  
**apply** *auto*[1] **apply** *auto*[1]  
**apply** *simp*  
  
**apply** *simp*  
**done**

**lemma** *mempool-free-stm1-inv-auxvars*:  
 $\text{inv-cur } Va \wedge \text{inv-thd-waitq } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist } Va \wedge$   
 $\text{inv-bitmap } Va \wedge \text{inv-aux-vars } Va \implies$   
 $\text{block } b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b)) \implies$   
 $\text{level } b < \text{length } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b))) \implies$   
 $\text{pool } b \in \text{mem-pools } Va \implies$   
 $\text{data } b = \text{block-ptr } (\text{mem-pool-info } Va \text{ (pool } b)) (\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \text{ (pool } b))) \text{ div } 4 \wedge \text{level } b) (\text{block } b) \implies$   
 $\text{get-bit } (\text{mem-pool-info } Va) (\text{pool } b) (\text{level } b) (\text{block } b) = \text{ALLOCATED} \implies$   
 $\text{allocating-node } Va \text{ } t = \text{None} \implies$   
 $\text{freeing-node } Va \text{ } t = \text{None} \implies$   
 $\text{inv-aux-vars}$   
 $(Va \langle \text{mem-pool-info} := (\text{mem-pool-info } Va)$   
 $\quad (\text{pool } b := \text{mem-pool-info } Va \text{ (pool } b)$   
 $\quad \langle \text{levels} := \text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$   
 $\quad \quad [\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b)$   
 $\quad \quad \quad \langle \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b) [\text{block } b$   
 $\quad := \text{FREEING}] \rangle \rangle \rangle),$   
 $\quad \text{freeing-node} := \text{freeing-node } Va (t \mapsto b) \rangle \rangle)$   
**apply**(*unfold inv-aux-vars-def*)  
**apply**(*rule conjI*)  
**apply** *clarify*  
**apply**(*case-tac*  $ta = t$ ) **apply** *auto*[1]  
**apply**(*subgoal-tac*  $\neg(\text{pool } n = \text{pool } b \wedge \text{level } n = \text{level } b \wedge \text{block } n = \text{block } b)$ )  
**apply**(*subgoal-tac freeing-node*  
 $(Va \langle \text{mem-pool-info} := (\text{mem-pool-info } Va)$   
 $\quad (\text{pool } b := \text{mem-pool-info } Va \text{ (pool } b)$   
 $\quad \langle \text{levels} := \text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$   
 $\quad \quad [\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b)$   
 $\quad \quad \quad \langle \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b) [\text{block } b := \text{FREEING}] \rangle \rangle \rangle),$   
 $\quad \text{freeing-node} := \text{freeing-node } Va (t \mapsto b) \rangle \rangle) \text{ } ta = \text{freeing-node } Va \text{ } ta)$   
**apply**(*subgoal-tac get-bit*  $(\text{mem-pool-info } Va) (\text{pool } n) (\text{level } n) (\text{block } n) = \text{FREEING})$   
**apply**(*subgoal-tac get-bit*  $(\text{mem-pool-info } Va) (\text{pool } n) (\text{level } n) (\text{block } n) =$   
 $\text{get-bit}$   
 $(\text{mem-pool-info}$   
 $(Va \langle \text{mem-pool-info} := (\text{mem-pool-info } Va)$   
 $\quad (\text{pool } b := \text{mem-pool-info } Va \text{ (pool } b)$   
 $\quad \langle \text{levels} := \text{levels } (\text{mem-pool-info } Va \text{ (pool } b))$   
 $\quad \quad [\text{level } b := (\text{levels } (\text{mem-pool-info } Va \text{ (pool } b)) ! \text{level } b)$

```

      (bits := bits (levels (mem-pool-info Va (pool b)) ! level
b)[block b := FREEING]))),
      freeing-node := freeing-node Va(t ↦ b)))
      (pool n) (level n) (block n))
apply auto[1]
  apply(case-tac ¬ pool n = pool b) apply simp
  apply(case-tac ¬ level n = level b) apply simp
  apply(case-tac ¬ block n = block b) apply simp apply simp
apply auto[1] apply auto[1]

apply(subgoal-tac freeing-node Va ta = Some n) prefer 2 apply auto[1]
apply fastforce

apply(rule conjI)
apply clarify
apply(case-tac ¬(pool n = pool b ∧ level n = level b ∧ block n = block b))
  apply(subgoal-tac get-bit (mem-pool-info Va) (pool n) (level n) (block n) =
FREEING)
    prefer 2 apply auto[1]
    using set-bit-def set-bit-get-bit-neq apply auto[1]
    using set-bit-def set-bit-get-bit-neq apply auto[1]
  apply(subgoal-tac mem-block-addr-valid Va n)
    prefer 2 using mem-block-addr-valid-def apply auto[1]
  apply(subgoal-tac ∃ t'. t' ≠ t ∧ freeing-node Va t' = Some n)
    prefer 2 apply (metis option.discI)
  apply auto[1]

  apply(subgoal-tac data b = data n)
  prefer 2 apply(simp add:block-ptr-def mem-block-addr-valid-def inv-mempool-info-maxsz-align4)
    apply auto[1]

apply(rule conjI)
apply clarify
apply(case-tac ta = t) apply auto[1]
apply(subgoal-tac allocating-node Va ta = Some n)
  prefer 2 apply auto[1]
apply(subgoal-tac get-bit-s Va (pool n) (level n) (block n) = ALLOCATING)
  prefer 2 apply auto[1]
apply(case-tac ¬(pool n = pool b ∧ level n = level b ∧ block n = block b))
  apply(case-tac ¬ pool n = pool b) apply simp
  apply(case-tac ¬ level n = level b) apply force
  apply(case-tac ¬ block n = block b) apply force apply simp
apply fastforce

apply(rule conjI)
apply clarify
apply(case-tac ¬(pool n = pool b ∧ level n = level b ∧ block n = block b))
  apply(subgoal-tac get-bit (mem-pool-info Va) (pool n) (level n) (block n) =
ALLOCATING)

```

```

    prefer 2 apply auto[1]
      using set-bit-def set-bit-get-bit-neq apply auto[1]
      using set-bit-def set-bit-get-bit-neq apply auto[1]
  apply(subgoal-tac mem-block-addr-valid Va n)
    prefer 2 using mem-block-addr-valid-def apply auto[1]
  apply(subgoal-tac  $\exists t'. t' \neq t \wedge \text{allocating-node } Va \ t' = \text{Some } n$ )
    prefer 2 apply (metis option.discI)
  apply auto[1]

  apply(subgoal-tac data b = data n)
  prefer 2 apply(simp add:block-ptr-def mem-block-addr-valid-def inv-mempool-info-maxsz-align4)
    apply auto[1]

  apply(rule conjI)
  apply clarify
  apply(case-tac  $t1 \neq t \wedge t2 \neq t$ )
  apply auto[1]
  apply(case-tac  $t1 = t$ )
  apply clarify
  apply(subgoal-tac freeing-node Va  $t2 = \text{Some } n2$ )
    prefer 2 apply auto[1]
  apply(subgoal-tac  $b = n1$ )
    prefer 2 apply auto[1]
  apply simp

  apply(case-tac  $t2 = t$ )
  apply clarify
  apply(subgoal-tac freeing-node Va  $t1 = \text{Some } n1$ )
    prefer 2 apply auto[1]
  apply(subgoal-tac  $b = n2$ )
    prefer 2 apply auto[1]
  apply fastforce

  apply simp

  apply(rule conjI)
  apply clarify
  apply(case-tac  $t1 \neq t \wedge t2 \neq t$ )
  apply auto[1]
  apply(case-tac  $t1 = t$ )
  apply clarify
  apply(subgoal-tac freeing-node Va  $t2 = \text{Some } n2$ )
    prefer 2 apply auto[1]
  apply(subgoal-tac  $b = n1$ )
    prefer 2 apply auto[1]
  apply simp

  apply(case-tac  $t2 = t$ )
  apply clarify

```

```

apply(subgoal-tac freeing-node Va t1 = Some n1)
  prefer 2 apply auto[1]
apply(subgoal-tac b = n2)
  prefer 2 apply auto[1]
apply fastforce

```

```

apply simp

```

```

apply clarify
apply(case-tac t1 ≠ t ∧ t2 ≠ t)
apply auto[1]
apply(case-tac t1 = t)
apply clarify
apply(subgoal-tac freeing-node Va t2 = Some n2)
  prefer 2 apply auto[1]
apply(subgoal-tac b = n1)
  prefer 2 apply auto[1]
apply simp

```

```

apply(case-tac t2 = t)
apply clarify
apply(subgoal-tac allocating-node Va t1 = Some n1)
  prefer 2 apply auto[1]
apply(subgoal-tac b = n2)
  prefer 2 apply auto[1]
apply fastforce

```

```

apply simp

```

**done**

**lemma** *mempool-free-stm1-inv-lvl0*:

```

inv-cur Va ∧ inv-thd-waitq Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist Va
∧ inv-bitmap Va ∧ inv-aux-vars Va ∧ inv-bitmap0 Va ⇒
block b < length (bits (levels (mem-pool-info Va (pool b)) ! level b)) ⇒
level b < length (levels (mem-pool-info Va (pool b))) ⇒
pool b ∈ mem-pools Va ⇒
data b = block-ptr (mem-pool-info Va (pool b)) (ALIGN4 (max-sz (mem-pool-info
Va (pool b))) div 4 ^ level b) (block b) ⇒
get-bit (mem-pool-info Va) (pool b) (level b) (block b) = ALLOCATED ⇒
allocating-node Va t = None ⇒
freeing-node Va t = None ⇒
inv-bitmap0
  (Va(mem-pool-info := (mem-pool-info Va)
    (pool b := mem-pool-info Va (pool b)
      (levels := levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]]]])),

```



```

      freeing-node := freeing-node Va(t ↦ b))
apply(simp add: inv-bitmap0-def Let-def)
apply clarsimp
apply(case-tac level b = 0)
  apply(case-tac block b = i) apply auto[1] apply simp
  apply simp
done

```

**lemma** mempool-free-stm1-inv-lvl:

```

  inv-cur Va ∧ inv-thd-waitq Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist Va
  ∧ inv-bitmap Va ∧ inv-aux-vars Va ∧ inv-bitmapn Va ⇒
    block b < length (bits (levels (mem-pool-info Va (pool b)) ! level b)) ⇒
    level b < length (levels (mem-pool-info Va (pool b))) ⇒
    pool b ∈ mem-pools Va ⇒
    data b = block_ptr (mem-pool-info Va (pool b)) (ALIGN4 (max-sz (mem-pool-info
  Va (pool b))) div 4 ^ level b) (block b) ⇒
    get-bit (mem-pool-info Va) (pool b) (level b) (block b) = ALLOCATED ⇒
    allocating-node Va t = None ⇒
    freeing-node Va t = None ⇒
    inv-bitmapn
      (Va (mem-pool-info := (mem-pool-info Va)
        (pool b := mem-pool-info Va (pool b)
          (levels := levels (mem-pool-info Va (pool b))
            [level b := (levels (mem-pool-info Va (pool b)) ! level b)
              (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]]))),
    freeing-node := freeing-node Va(t ↦ b))
apply(simp add: inv-bitmapn-def Let-def)
apply clarsimp
apply(case-tac level b = length (levels (mem-pool-info Va (pool b))) - 1)
  apply(case-tac block b = i) apply auto[1] apply simp
  apply simp
done

```

**lemma** mempool-free-stm1-inv-lvl-not4free:

```

  inv-cur Va ∧ inv-thd-waitq Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist Va
  ∧ inv-bitmap Va ∧ inv-aux-vars Va ∧ inv-bitmap-not4free Va ⇒
    block b < length (bits (levels (mem-pool-info Va (pool b)) ! level b)) ⇒
    level b < length (levels (mem-pool-info Va (pool b))) ⇒
    pool b ∈ mem-pools Va ⇒
    data b = block_ptr (mem-pool-info Va (pool b)) (ALIGN4 (max-sz (mem-pool-info
  Va (pool b))) div 4 ^ level b) (block b) ⇒
    get-bit (mem-pool-info Va) (pool b) (level b) (block b) = ALLOCATED ⇒
    allocating-node Va t = None ⇒
    freeing-node Va t = None ⇒
    inv-bitmap-not4free
      (Va (mem-pool-info := (mem-pool-info Va)
        (pool b := mem-pool-info Va (pool b)
          (levels := levels (mem-pool-info Va (pool b))

```

```

[level b := (levels (mem-pool-info Va (pool b)) ! level b)
  (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]])),
  freeing-node := freeing-node Va(t ↦ b))
apply(simp add: inv-bitmap-not4free-def Let-def partner-bits-def)
apply clarsimp
apply(case-tac level b = i) prefer 2 apply auto[1]
  apply(case-tac block b = j div 4 * 4) apply auto[1]
  apply(case-tac block b = j div 4 * 4 + 1) apply auto[1]
  apply(case-tac block b = j div 4 * 4 + 2) apply auto[1]
  apply(case-tac block b = j div 4 * 4 + 3) apply auto[1]

apply simp
done

```

**lemma** mempool-free-sm1-inv:

```

inv Va ⇒
  block b < length (bits (levels (mem-pool-info Va (pool b)) ! level b)) ⇒
  level b < length (levels (mem-pool-info Va (pool b))) ⇒
  pool b ∈ mem-pools Va ⇒
  data b = block_ptr (mem-pool-info Va (pool b)) (ALIGN4 (max-sz (mem-pool-info
Va (pool b))) div 4 ^ level b) (block b) ⇒
  get-bit (mem-pool-info Va) (pool b) (level b) (block b) = ALLOCATED ⇒
  allocating-node Va t = None ⇒
  freeing-node Va t = None ⇒
  inv (Va (mem-pool-info := (mem-pool-info Va)
    (pool b := mem-pool-info Va (pool b)
      (levels := levels (mem-pool-info Va (pool b))
        [level b := (levels (mem-pool-info Va (pool b)) ! level b)
          (bits := bits (levels (mem-pool-info Va (pool b)) ! level b)[block b
:= FREEING]]])),
    freeing-node := freeing-node Va(t ↦ b)))
apply(simp add:inv-def)
apply(rule conjI) apply(simp add:inv-cur-def Mem-pool-free-guar-def)
apply(rule conjI) apply(simp add:inv-thd-waitq-def)
  apply(rule conjI) apply clarify apply metis
  apply clarify apply metis
apply(rule conjI) using mempool-free-sm1-inv-mempool-info apply auto[1]
apply(rule conjI) using mempool-free-sm1-inv-bitmap-freelist apply auto[1]
apply(rule conjI) using mempool-free-sm1-inv-bitmap apply auto[1]
apply(rule conjI) using mempool-free-sm1-inv-auxvars apply auto[1]
apply(rule conjI) using mempool-free-sm1-inv-lvl0 apply auto[1]
apply(rule conjI) using mempool-free-sm1-inv-lvl1 apply auto[1]
  using mempool-free-sm1-inv-lvl-not4free apply auto[1]
done

```

**lemma** mempool-free-sm1-h1:

Mem-pool-free-pre t ∩

```

     $\{\{ \text{pool } b \in 'mem-pools \wedge$ 
     $\text{level } b < \text{length } (\text{levels } ('mem-pool-info (\text{pool } b))) \wedge$ 
     $\text{block } b < \text{length } (\text{bits } (\text{levels } ('mem-pool-info (\text{pool } b)) ! \text{level } b)) \wedge$ 
     $\text{data } b =$ 
     $\text{block-ptr } ('mem-pool-info (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz } ('mem-pool-info$ 
     $(\text{pool } b))) \text{ div } 4 \wedge \text{level } b) (\text{block } b) \}$   $\cap$ 
     $\{\{ 'cur = \text{Some } t \} \cap$ 
     $\{ Va \} \cap$ 
     $\{\{ 'get-bit-s (\text{pool } b) (\text{level } b) (\text{block } b) = \text{ALLOCATED} \} \cap$ 
     $\{ Va \} \neq$ 
     $\{ \} \Rightarrow$ 
     $\Gamma \vdash_I \text{Some } ('mem-pool-info := \text{set-bit-freeing } 'mem-pool-info (\text{pool } b) (\text{level } b)$ 
     $(\text{block } b));;$ 
     $'freeing-node := 'freeing-node(t \mapsto$ 
     $b)) \text{ sat}_p [\text{Mem-pool-free-pre } t \cap$ 
     $\{\{ \text{pool } b \in 'mem-pools \wedge$ 
     $\text{level } b < \text{length } (\text{levels } ('mem-pool-info (\text{pool } b))) \wedge$ 
     $\text{block } b < \text{length } (\text{bits } (\text{levels } ('mem-pool-info (\text{pool } b)) ! \text{level}$ 
     $b)) \wedge$ 
     $\text{data } b =$ 
     $\text{block-ptr } ('mem-pool-info (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz}$ 
     $('mem-pool-info (\text{pool } b))) \text{ div } 4 \wedge \text{level } b)$ 
     $(\text{block } b) \}$   $\cap$ 
     $\{\{ 'cur = \text{Some } t \} \cap$ 
     $\{ Va \} \cap$ 
     $\{\{ 'get-bit-s (\text{pool } b) (\text{level } b) (\text{block } b) = \text{ALLOCATED} \} \cap$ 
     $\{ Va \}, \{(x, y).$ 
     $x = y\}, \text{UNIV}, \{\{ (Pair Va) \in \text{Mem-pool-free-guar } t \} \cap$ 
     $(\{\{ \text{invariant.inv} \wedge 'allocating-node } t = \text{None} \} \cap$ 
     $\{\{ \text{pool } b \in 'mem-pools \wedge$ 
     $\text{level } b < \text{length } (\text{levels } ('mem-pool-info (\text{pool}$ 
     $b))) \wedge$ 
     $\text{block } b < \text{length } (\text{bits } (\text{levels } ('mem-pool-info$ 
     $(\text{pool } b)) ! \text{level } b)) \wedge$ 
     $\text{data } b =$ 
     $\text{block-ptr } ('mem-pool-info (\text{pool } b))$ 
     $(\text{ALIGN4 } (\text{max-sz } ('mem-pool-info (\text{pool}$ 
     $b))) \text{ div } 4 \wedge \text{level } b) (\text{block } b) \}$   $\cap$ 
     $\text{mp-free-precond2-ext } t \ b)]$ 

    apply clarsimp
    apply(rule Seq[where  $\text{mid} = \{ Va \mid \text{mem-pool-info} := \text{set-bit-freeing } (\text{mem-pool-info}$ 
     $Va) (\text{pool } b) (\text{level } b) (\text{block } b) \}$ ]])
    apply(rule Basic)
    apply auto[1] apply(simp add:stable-def)+
    apply(rule Basic)
    apply(simp add:set-bit-def)
    apply(rule conjI)
    apply(simp add:Mem-pool-free-guar-def)
    apply(rule disjI1)

```

```

apply(rule conjI)
  apply(simp add:gvars-conf-stable-def gvars-conf-def) apply auto[1]
  apply(case-tac i = level b) apply auto[1] apply auto[1]

apply(rule conjI)
  using mempool-free-smt1-inv apply auto[1]
apply(simp add:lvars-nochange-def)

apply(rule conjI)
  using mempool-free-smt1-inv apply auto[1]
apply(simp add:block-ptr-def)

apply(simp add:stable-def)+
done

lemma mempool-free-stm1:
   $\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{AWAIT bits } (\text{levels } (' \text{mem-pool-info } (\text{pool } b)) ! \text{level } b) ! \text{block } b$ 
   $= \text{ALLOCATED THEN}$ 
   $' \text{mem-pool-info} := \text{set-bit-freing } ' \text{mem-pool-info } (\text{pool } b) (\text{level } b) (\text{block } b);;$ 
   $' \text{freing-node} := ' \text{freing-node } (t := \text{Some } b)$ 
   $\text{END}) \text{ sat}_p$ 
   $[ \text{mp-free-precond1 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond2}$ 
   $t \ b]$ 
  apply(simp add:stm-def)
  apply(rule Await)

using mp-free-precond1-stb apply auto[1]
using mp-free-precond2-stb apply auto[1]

apply(rule allI)
apply(rule Await)
  apply(simp add:stable-def) apply(auto simp add:stable-def)
  apply(case-tac  $V \neq Va$ ) apply auto[1] using Emptyprecond apply blast
  apply clarsimp
  apply(case-tac  $\text{mp-free-precond1 } t \ b \cap \{ ' \text{cur} = \text{Some } t \} \cap \{ Va \} \cap$ 
     $\{ \text{get-bit } ' \text{mem-pool-info } (\text{pool } b) (\text{level } b) (\text{block } b) = \text{ALLOCATED} \}$ 
   $\cap$ 
     $\{ Va \} = \{ \}$ )
  apply simp using Emptyprecond apply auto[1]
  using mempool-free-stm1-h1 apply force
done

lemma mempool-free-stm2:
   $\Gamma \vdash_I \text{Some } (t \blacktriangleright ' \text{need-resched} := ' \text{need-resched}(t := \text{False})) \text{ sat}_p$ 
   $[ \text{mp-free-precond2 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond3}$ 
   $t \ b]$ 
  apply(simp add:stm-def)

```

```

apply(rule Await)
using mp-free-precond2-stb apply simp
using mp-free-precond3-stb apply simp

apply clarify
apply(rule Basic)
apply(case-tac mp-free-precond2 t b  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} = \{\}$ )
  apply auto[1]
  apply clarsimp
  apply(rule conjI)
  apply(simp add:Guarf-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
  apply(subgoal-tac (V, V( $\llbracket need-resched := (need-resched\ V)(t := False) \rrbracket$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

  apply clarify apply(simp add: lvars-nochange-def)

  apply(subgoal-tac (V, V( $\llbracket need-resched := (need-resched\ V)(t := False) \rrbracket$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

  apply(simp add:stable-def)+

done

lemma mempool-free-stm3:
 $\Gamma \vdash_I Some\ (t \blacktriangleright 'lsizes := 'lsizes(t := [ALIGN_4\ (max-sz\ ('mem-pool-info\ (pool\ b))))))\ sat_p$ 
[mp-free-precond3 t b, Mem-pool-free-rely t, Mem-pool-free-guar t, mp-free-precond4
t b]
apply(simp add:stm-def)
apply(rule Await)
using mp-free-precond3-stb apply simp
using mp-free-precond4-stb apply simp

apply clarify
apply(rule Basic)
apply(case-tac mp-free-precond3 t b  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} = \{\}$ )
  apply auto[1]
  apply clarsimp
  apply(rule conjI)
  apply(simp add:Guarf-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
  apply(subgoal-tac (V, V( $\llbracket lsizes := (lsizes\ V)(t := [ALIGN_4\ (max-sz\ (mem-pool-info\ V\ (pool\ b)))) \rrbracket$ )) $\in$ lvars-nochange1-4all)

```

```

    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

    apply clarify apply(simp add: lvars-nochange-def)
    apply(subgoal-tac (V, V(|lsizes := (lsizes V)(t := [ALIGN4 (max-sz (mem-pool-info
V (pool b)))))))))
        ∈lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:stable-def)+

done

lemma mempool-free-stm41-h1-1: (n::nat) > 0 ∧ (x::nat) mod y = 0 ⇒ n * x
mod y = 0
by auto

lemma mempool-free-stm41-h1:
  assumes p1: i V t ≤ level b
    and p2: length (lsizes V t) = i V t
    and p3: inv V
    and p4: ∀ ii < i V t. lsizes V t ! ii = ALIGN4 (max-sz (mem-pool-info V (pool
b))) div 4 ^ ii
    and p5: lsizes V t ≠ []
    and p6: pool b ∈ mem-pools V
    and p7: level b < length (levels (mem-pool-info V (pool b)))
    and p8: block b < length (bits (levels (mem-pool-info V (pool b)) ! level b))
    and p9: ii = i V t
  shows (lsizes V t @ [ALIGN4 (lsizes V t ! (i V t - Suc NULL) div 4)]) ! ii
    = ALIGN4 (max-sz (mem-pool-info V (pool b))) div 4 ^ ii
proof -
  from p2 p9 have a0: (lsizes V t @ [ALIGN4 (lsizes V t ! (i V t - 1) div 4)]) !
ii
    = ALIGN4 (lsizes V t ! (i V t - 1) div 4)
  by (metis nth-append-length)

  from p2 p4 p5 have lsizes V t ! (i V t - 1) div 4 = ALIGN4 (max-sz
(mem-pool-info V (pool b))) div 4 ^ (i V t - 1) div 4
  by (metis One-nat-def diff-less-mono2 diff-zero length-greater-0-conv zero-less-Suc)
  hence a1: lsizes V t ! (i V t - 1) div 4 = ALIGN4 (max-sz (mem-pool-info V
(pool b))) div 4 ^ (i V t)
  by (metis One-nat-def Suc-pred div-mult2-eq length-greater-0-conv p2 p5
plus-1-eq-Suc power-add power-commutes power-one-right)

  from p6 p3 have ∃ n > 0. max-sz (mem-pool-info V (pool b))
    = (4 * n) * (4 ^ (length (levels (mem-pool-info V (pool b)))))
  apply(simp add:inv-def inv-mempool-info-def Let-def) by auto
  then obtain n where n > 0 ∧ max-sz (mem-pool-info V (pool b))
    = (4 * n) * (4 ^ (length (levels (mem-pool-info V (pool b)))))

```

**by** *auto*  
**hence**  $a2: n > 0 \wedge \text{max-sz} (\text{mem-pool-info } V (\text{pool } b))$   
 $= n * (4 \wedge (\text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) + 1))$  **by** *auto*  
**hence**  $\text{max-sz} (\text{mem-pool-info } V (\text{pool } b)) \bmod 4 = 0$  **by** *simp*  
**hence**  $a3: \text{ALIGN}_4 (\text{max-sz} (\text{mem-pool-info } V (\text{pool } b))) = \text{max-sz} (\text{mem-pool-info } V (\text{pool } b))$   
**using** *align40* **by** *auto*  
**with**  $a1$  **have**  $a4: \text{lsizes } V \ t \ ! \ (i \ V \ t - 1) \ \text{div } 4 = \text{max-sz} (\text{mem-pool-info } V (\text{pool } b)) \ \text{div } 4 \wedge (i \ V \ t)$  **by** *simp*  
  
**from**  $p1 \ p2 \ p7 \ a2$  **have**  $(\text{max-sz} (\text{mem-pool-info } V (\text{pool } b)) \ \text{div } 4 \wedge i \ V \ t) \bmod 4 = 0$   
**apply**  $(\text{subgoal-tac } 4 \wedge (\text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) + 1) \ \text{div } 4 \wedge i \ V \ t \bmod 4 = \text{NULL})$   
**prefer** 2 **using** *pow-lt-mod0*  $[of \ 4 \ i \ V \ t \ \text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) + 1]$  **apply** *auto*  $[1]$   
**apply** *simp* **using** *mempool-free-stm41-h1-1*  
 $[of \ n \ 4 * 4 \wedge \text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) \ \text{div } 4 \wedge i \ V \ t \ 4]$   
**using** *m-mod-div* *mempool-free-stm41-h1-1* *pow-mod-0* **by** *force*  
**with**  $a0 \ a1 \ a3 \ a4 \ p9$  **show** *?thesis* **using** *align40* **by** *simp*  
**qed**

**lemma** *mempool-free-stm41*:

$\Gamma \vdash_I \text{Some } ('lsizes := 'lsizes$   
 $(t := 'lsizes \ t \ @ \ [\text{ALIGN}_4 ('lsizes \ t \ ! \ ('i \ t - 1) \ \text{div } 4)]))$   
 $\text{sat}_p [\text{mp-free-precond}_4\text{-}2 \ t \ b \cap \{\}'cur = \text{Some } t\} \cap \{V\}, \{(s, t). s = t\},$   
 $\text{UNIV},$   
 $\{\}'(Pair \ V) \in \text{Mem-pool-free-guar } t\} \cap (\text{mp-free-precond}_2 \ t \ b \cap \{\}'\neg$   
 $'need-resched \ t\} \cap$   
 $\{\}(\forall \ ii < \text{length } ('lsizes \ t).$   
 $'lsizes \ t \ ! \ ii = \text{ALIGN}_4 (\text{max-sz} ('mem-pool-info (\text{pool } b))) \ \text{div } 4 \wedge$   
 $ii) \wedge 'lsizes \ t \neq []\} \cap$   
 $(\{\}'i \ t \leq \text{level } b\} \cap \{\}\text{length } ('lsizes \ t) = \text{Suc } ('i \ t)\}\})$   
**apply** *(rule Basic)*  
**apply**  $(\text{case-tac } \text{mp-free-precond}_4\text{-}2 \ t \ b \cap \{\}'cur = \text{Some } t\} \cap \{V\} = \{\})$   
**apply** *simp* **apply** *clarify* **apply** *auto*  $[1]$   
**apply**  $(\text{simp add: } \text{Guar}_f\text{-def } \text{gvars-conf-stable-def } \text{gvars-conf-def } \text{Mem-pool-free-guar-def})$   
  
**apply** *(rule disjI1)*  
**apply** *(rule conjI)*  
**apply**  $(\text{subgoal-tac } (V, V(\text{lsizes} := (\text{lsizes } V)(t := \text{lsizes } V \ t \ @ \ [\text{ALIGN}_4 (\text{lsizes } V \ t \ ! \ (i \ V \ t - \text{Suc } \text{NULL}) \ \text{div } 4)])) \in \text{lvars-nochange1-4all})$   
**using** *glnochange-inv0* **apply** *auto*  $[1]$  **apply**  $(\text{simp add: } \text{lvars-nochange1-4all-def } \text{lvars-nochange1-def})$   
**apply**  $(\text{simp add: } \text{lvars-nochange-def})$   
**apply**  $(\text{subgoal-tac } (V, V(\text{lsizes} := (\text{lsizes } V)(t := \text{lsizes } V \ t \ @ \ [\text{ALIGN}_4 (\text{lsizes } V \ t \ ! \ (i \ V \ t - \text{Suc } \text{NULL}) \ \text{div } 4)])) \in \text{lvars-nochange1-4all})$   
**using** *glnochange-inv0* **apply** *auto*  $[1]$  **apply**  $(\text{simp add: } \text{lvars-nochange1-4all-def } \text{lvars-nochange1-def})$

```

    apply(case-tac ii < i V t) apply (simp add: nth-append)
    apply(case-tac ii = i V t)
    using mempool-free-stm4-1-h1 apply metis
    apply simp
  by (simp add: stable-def)+

lemma mempool-free-stm4:
   $\Gamma \vdash_I \text{Some } (FOR (t \blacktriangleright 'i := 'i(t := 1));$ 
     $'i \ t \leq \text{level } b;$ 
     $(t \blacktriangleright 'i := 'i(t := 'i \ t + 1)) \text{ DO}$ 
     $(t \blacktriangleright 'lsizes := 'lsizes(t := 'lsizes \ t \ @ \ [ALIGN_4 \ ('lsizes \ t \ ! \ ('i \ t - 1) \ \text{div}$ 
     $4)))))$ 
     $ROF) \ \text{sat}_p \ [mp\text{-free-precond}_4 \ t \ b, \ \text{Mem-pool-free-rely } t, \ \text{Mem-pool-free-guar } t,$ 
     $mp\text{-free-precond}_5 \ t \ b]$ 
    apply(rule Seq[where mid=mp-free-precond4-1 t b])

  apply(simp add: stm-def)
  apply(rule Await)
  using mp-free-precond4-stb apply simp
  using mp-free-precond4-1-stb apply simp
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-free-precond4 t b  $\cap \ \{\cur = \text{Some } t\} \cap \ \{V\} = \{\}$ )
    apply auto[1]
  apply(simp add: Guar-f-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
  apply auto[1]
  apply(subgoal-tac (V, V( $i := (i \ V)(t := \text{Suc } \text{NULL})$ )) $\in \text{lvars-nochange1-4all}$ )
    using glnochange-inv0 apply auto[1] apply(simp add: lvars-nochange1-4all-def
    lvars-nochange1-def)
  apply(simp add: lvars-nochange-def)
  apply(subgoal-tac (V, V( $i := (i \ V)(t := \text{Suc } \text{NULL})$ )) $\in \text{lvars-nochange1-4all}$ )
    using glnochange-inv0 apply auto[1] apply(simp add: lvars-nochange1-4all-def
    lvars-nochange1-def)
  apply(simp add: stable-def)+

  apply(rule While)
  using mp-free-precond4-1-stb apply simp
  apply auto[1]
  using mp-free-precond5-stb apply simp

  apply(rule Seq[where mid=mp-free-precond4-3 t b])

  apply(simp add: stm-def)
  apply(rule Await)

```



```

using mp-free-precond4-2-stb apply simp
using mp-free-precond4-3-stb apply simp
apply(rule allI)
using mempool-free-stm41 apply simp

apply(simp add:stm-def)
apply(rule Await)
using mp-free-precond4-3-stb apply simp
using mp-free-precond4-1-stb apply simp
apply(rule allI)
apply(rule Basic)
  apply(case-tac mp-free-precond4-3 t b  $\cap \llbracket \text{'cur} = \text{Some } t \rrbracket \cap \{V\} = \{\}$ )
  apply auto[1] apply clarify apply(simp add:gvars-conf-stable-def gvars-conf-def
Mem-pool-free-guar-def) apply auto[1]
  apply(subgoal-tac (V, V( $\llbracket i := (i \ V) (t := \text{Suc } (i \ V \ t) \rrbracket) \in \text{lvars-nochange1-4all}$ ))
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:lvars-nochange-def)
    apply(subgoal-tac (V, V( $\llbracket i := (i \ V) (t := \text{Suc } (i \ V \ t) \rrbracket) \in \text{lvars-nochange1-4all}$ ))
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
      apply(simp add:stable-def) + apply(simp add:Guarf-def Mem-pool-free-guar-def
Id-def)

```

done

**lemma** mempool-free-stm5:

```

 $\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{'free-block-r} := \text{'free-block-r } (t := \text{True}))$ 
 $\text{sat}_p [\text{mp-free-precond5 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond6}$ 
 $t \ b]$ 

```

```

apply(simp add:stm-def)
apply(rule Await)
using mp-free-precond5-stb apply simp
using mp-free-precond6-stb apply simp

apply clarify
apply(rule Basic)
apply(case-tac mp-free-precond5 t b  $\cap \llbracket \text{'cur} = \text{Some } t \rrbracket \cap \{V\} = \{\}$ )
  apply auto[1]
  apply clarsimp
  apply(rule conjI)
  apply(simp add:Guarf-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
    apply(subgoal-tac (V, V( $\llbracket \text{free-block-r} := (\text{free-block-r } V) (t := \text{True}) \rrbracket$ ))
       $\in \text{lvars-nochange1-4all}$ )
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def

```

*lvars-nochange1-def*)

```

    apply clarify apply(simp add: lvars-nochange-def)
    apply(subgoal-tac (V, V( $\lfloor$ free-block-r := (free-block-r V)(t := True) $\rfloor$ ))  $\in$  lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add: lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add: stable-def)+
done

```

**lemma** *mempool-free-stm6*:

$\Gamma \vdash_I \text{Some } (t \blacktriangleright 'bn := 'bn \ (t := \text{block } b))$   
 $\text{sat}_p [\text{mp-free-precond6 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond7 } t \ b]$

```

    apply(simp add: stm-def)
    apply(rule Await)
    using mp-free-precond6-stb apply simp
    using mp-free-precond7-stb apply simp

    apply clarify
    apply(rule Basic)
    apply(case-tac mp-free-precond6 t b  $\cap \lfloor 'cur = \text{Some } t \rfloor \cap \{V\} = \{\}$ )
    apply auto[1]
    apply clarsimp
    apply(rule conjI)
    apply(simp add: Guarf-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
    apply(rule disjI1)
    apply(rule conjI)
    apply(subgoal-tac (V, V( $\lfloor$ bn := (bn V)(t := block b) $\rfloor$ ))  $\in$  lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add: lvars-nochange1-4all-def
lvars-nochange1-def)

```

```

    apply clarify apply(simp add: lvars-nochange-def)
    apply(subgoal-tac (V, V( $\lfloor$ bn := (bn V)(t := block b) $\rfloor$ ))  $\in$  lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add: lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add: stable-def)+
done

```

**lemma** *mempool-free-stm7*:

$\Gamma \vdash_I \text{Some } (t \blacktriangleright 'lvl := 'lvl \ (t := \text{level } b))$   
 $\text{sat}_p [\text{mp-free-precond7 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond8 } t \ b]$

```

    apply(unfold stm-def)
    apply(rule Await)
    using mp-free-precond7-stb apply simp
    using mp-free-precond8-stb[of t b] apply fast

```

```

    apply clarify
    apply(rule Basic)

```

```

apply(case-tac mp-free-precond7  $t \ b \cap \llbracket 'cur = Some \ t \rrbracket \cap \{V\} = \{\}$ )
apply auto[1]
apply clarsimp
apply(rule conjI)
apply(simp add:Guarf-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def)
apply(rule disjI1)
apply(rule conjI)
apply(subgoal-tac ( $V, V \llbracket lvl := (lvl \ V)(t := level \ b) \rrbracket \in lvars-nochange1-4all$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

apply clarify apply(simp add: lvars-nochange-def)
apply(rule conjI)
apply(subgoal-tac ( $V, V \llbracket lvl := (lvl \ V)(t := level \ b) \rrbracket \in lvars-nochange1-4all$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

apply auto[1]
apply (simp add: block-ptr-def inv-maxsz-align4)
apply (metis inv-mempool-info-def inv-def)
apply simp
apply(simp add:stable-def)
using stable-id2 apply metis
done

```

### 9.3 statement 8

**abbreviation** *free-stm8-precond1*  $Va \ t \ b \equiv Va \llbracket mem-pool-info := set-bit-free \ (mem-pool-info \ Va) \ (pool \ b) \ (lvl \ Va \ t) \ (bn \ Va \ t) \rrbracket$

**abbreviation** *free-stm8-precond2*  $Va \ t \ b \equiv (free-stm8-precond1 \ Va \ t \ b) \llbracket freeing-node := (freeing-node \ Va)(t := None) \rrbracket$

**abbreviation** *free-stm8-loopinv1*  $Va \ t \ b \equiv$   
 $\{V. \text{let } minf0 = (mem-pool-info \ Va)(pool \ b);$   
 $lvl0 = (levels \ minf0) ! (lvl \ Va \ t);$   
 $minf1 = (mem-pool-info \ V)(pool \ b);$   
 $lvl1 = (levels \ minf1) ! (lvl \ Va \ t) \text{ in}$   
 $cur \ V = cur \ Va \wedge tick \ V = tick \ Va \wedge thd-state \ V = thd-state \ Va \wedge$   
 $(V, Va) \in gvars-conf-stable$   
 $\wedge (\forall p. p \neq pool \ b \longrightarrow mem-pool-info \ V \ p = mem-pool-info \ Va \ p)$   
 $\wedge (\forall j. j \neq lvl \ Va \ t \longrightarrow (levels \ minf0) ! j = (levels \ minf1) ! j)$   
 $\wedge (bits \ lvl1 = list-updates-n \ (bits \ lvl0) \ ((bn \ Va \ t \ div \ 4) * 4) \ (i \ V \ t) \ NOEXIST)$   
 $\wedge (free-list \ lvl1 = removes \ (map \ (\lambda ii. block-ptr \ minf0 \ (lsz \ Va \ t) \ ((bn \ Va \ t \ div \ 4) * 4 + ii)) \ [0..<(i \ V \ t)]) \ (free-list \ lvl0))$   
 $\wedge (wait-q \ minf0 = wait-q \ minf1)$   
 $\wedge (\forall t'. t' \neq t \longrightarrow lvars-nochange \ t' \ V \ Va)$   
 $\wedge freeing-node \ Va \ t = freeing-node \ V \ t \wedge allocating-node \ Va \ t = allocating-node \ V \ t \wedge free-block-r \ Va \ t = free-block-r \ V \ t$   
 $\wedge bn \ Va \ t = bn \ V \ t \wedge lvl \ Va \ t = lvl \ V \ t \wedge lsz \ Va \ t = lsz \ V \ t \wedge lsizes \ Va \ t = lsizes \ V \ t$

$$\wedge i \ V \ t \leq 4 \}$$

**lemma** *V-free-stm8-loopinv1*:  $i \ V \ t = 0 \implies V \in \text{free-stm8-loopinv1} \ V \ t \ b$   
**by** (*simp add:Let-def gvars-conf-stable-def gvars-conf-def lvvars-nochange-def*)

**abbreviation** *free-stm8-precond3*  $V \ a \ t \ b \equiv \text{free-stm8-loopinv1} \ (\text{free-stm8-precond2} \ V \ a \ t \ b) \ t \ b$

**abbreviation** *free-stm8-precond4*  $V \ a \ t \ b \equiv \text{free-stm8-precond3} \ V \ a \ t \ b \cap \{i \ t = 4\}$

**abbreviation** *free-stm8-precond30*  $V \ a \ t \ b \equiv \text{free-stm8-precond3} \ V \ a \ t \ b \cap \{i \ t < 4\}$

**abbreviation** *free-stm8-precond31*  $V \ t \ b \equiv V(\text{bits} := (\text{bits} \ V) \ (t := (\text{bn} \ V \ t \ \text{div} \ 4) * 4 + i \ V \ t))$

**abbreviation** *free-stm8-precond32*  $V \ t \ b \equiv$

*let* *minf* = *mem-pool-info* *V* (*pool* *b*) *in*

$V(\text{mem-pool-info} := (\text{mem-pool-info} \ V) \ (\text{pool} \ b := \text{minf} \ (\text{levels} := (\text{levels} \ \text{minf})))$

$[\text{lvl} \ V \ t := ((\text{levels} \ \text{minf}) ! (\text{lvl} \ V \ t)) \ (\text{bits} := (\text{bits} \ ((\text{levels} \ \text{minf}) ! (\text{lvl} \ V \ t)))]$

$[\text{bb} \ V \ t := \text{NOEXIST}] \ \rangle \ \rangle \ \rangle$

**abbreviation** *free-stm8-precond33*  $V \ t \ b \equiv$

$V(\text{block-pt} := (\text{block-pt} \ V) \ (t := \text{block-ptr} \ (\text{mem-pool-info} \ V \ (\text{pool} \ b)) \ (\text{lsz} \ V \ t) \ (\text{bb} \ V \ t))$

**abbreviation** *free-stm8-precond34*  $V \ t \ b \equiv$

*let* *minf* = *mem-pool-info* *V* (*pool* *b*) *in*

$V(\text{mem-pool-info} := (\text{mem-pool-info} \ V) \ (\text{pool} \ b := \text{minf} \ (\text{levels} := (\text{levels} \ \text{minf})))$

$[\text{lvl} \ V \ t := ((\text{levels} \ \text{minf}) ! (\text{lvl} \ V \ t)) \ (\text{free-list} := \text{remove1} \ (\text{block-pt} \ V \ t) \ (\text{free-list} \ ((\text{levels} \ \text{minf}) ! (\text{lvl} \ V \ t)))) \ \rangle \ \rangle \ \rangle$

**lemma** *mempool-free-stm8-atombody-h1*:

$\{\text{free-stm8-precond1} \ V \ t \ b\} \subseteq \{(\text{freeing-node-update} \ (\lambda-. \ \text{freeing-node}(t := \text{None}))) \in \{\text{free-stm8-precond2} \ V \ t \ b\}\}$

**by** *fastforce*

**lemma** *block-fits0-h1*:  $\text{maxsz} \ \text{mod} \ mm = 0 \implies aa < nmax * mm \implies$

$\text{maxsz} \ \text{div} \ mm * aa + \text{maxsz} \ \text{div} \ mm < \text{Suc} \ (nmax * \text{maxsz})$

**apply** (*subgoal-tac*  $\text{maxsz} \ \text{div} \ mm * aa \leq \text{maxsz} \ \text{div} \ mm * (nmax * mm - 1)$ )

**prefer** 2 **apply** *auto*[1]

**by** (*smt* *Groups.add-ac*(2) *Groups.mult-ac*(2) *Groups.mult-ac*(3) *One-nat-def* *Suc-leI* *distrib-left*

*le-imp-less-Suc* *mod-div-self* *mult.right-neutral* *mult-0-right* *mult-Suc-right* *mult-less-cancel2* *mult-zero-right* *not-le plus-nat.simps*(2))

**lemma** *block-fits0-h2*:  $(\text{lvt}::\text{nat}) > 0 \implies \text{lvt} \leq \text{lvlb} \implies \text{ivt} < (4::\text{nat}) \implies \text{blockb} < nmax * 4 \wedge \text{lvlb} \implies$

```

    blockb div 4 ^ (lvlb - lvl) div 4 * 4 + ivt < nmax * 4 ^ lvl
  apply(subgoal-tac nmax > 0) prefer 2 using mult-not-zero apply fastforce
  apply(subgoal-tac blockb div 4 ^ (lvlb - lvl) < (nmax * 4 ^ lvl)) prefer 2
  apply(subgoal-tac blockb < nmax * 4 ^ (lvl + (lvlb - lvl)) ∧ nmax * 4 ^ lvl
≠ 0) prefer 2 apply simp
  apply(subgoal-tac ∧ n na nb. ¬ n < na * nb ∨ n div na < nb ∨ nb = NULL)
  prefer 2
  apply (simp add: less-mult-imp-div-less mult.commute)
  apply (metis mult.commute mult.left-commute power-add)
  apply(subgoal-tac blockb div 4 ^ (lvlb - lvl) div 4 * 4 + 4 ≤ nmax * 4 ^ lvl)
  prefer 2 apply(subgoal-tac ∧ x. x < nmax * 4 ^ lvl → x div 4 * 4 + 4 ≤
nmax * 4 ^ lvl)
  prefer 2 apply(case-tac x mod 4 = 0) apply auto[1] apply(rule modn0-xy-n)
  apply auto[1] apply auto[1] apply auto[1] apply auto[1]
  apply auto[1] apply(rule divn-multn-addn-le) apply auto[1] apply
auto[1] apply auto[1]
  apply auto
done

```

**lemma** *block-fits0*:

```

  V ∈ mp-free-precond8-3 t b α ∩ {cur = Some t} ⇒
  vt ∈ free-stm8-precond3 V t b ∩ {i t < 4} ⇒
  {free-stm8-precond2 V t b} ∩ {0 < 'lvl t ∧ partner-bits ('mem-pool-info (pool
b)) ('lvl t) ('bn t)} ≠ {} ⇒
  free-stm8-precond33 (free-stm8-precond32 (free-stm8-precond31 vt t b) t b) t b
  ∈ {block-fits ('mem-pool-info (pool b)) ('block-pt t) ('lsz t)}
  apply(unfold block-fits-def block-ptr-def buf-size-def) apply clarsimp
  apply(rule subst[where s=lsz vt and t=lsz (let minf = mem-pool-info vt (pool
b)
    in vt(bb := (bb vt)(t := bn vt t div 4 * 4 + i vt t), mem-pool-info :=
(mem-pool-info vt)
      (pool b := minf(levels := levels minf[lvl vt t := (levels minf ! lvl vt
t)
        (bits := bits (levels minf ! lvl vt t)[bn vt t div 4 * 4 + i vt t :=
NOEXIST]))))])))
  apply(simp add:Let-def)
  apply(rule subst[where s=bn vt t div 4 * 4 + i vt t and t=bb (let minf =
mem-pool-info vt (pool b)
    in vt(bb := (bb vt)(t := bn vt t div 4 * 4 + i vt t), mem-pool-info :=
(mem-pool-info vt)
      (pool b := minf(levels := levels minf[lvl vt t := (levels minf ! lvl vt
t)
        (bits := bits (levels minf ! lvl vt t)[bn vt t div 4 * 4 + i vt t :=
NOEXIST]))))] t)]
  apply(simp add:Let-def)
  apply(rule subst[where s=n-max (mem-pool-info vt (pool b)) and t=n-max
(mem-pool-info
    (let minf = mem-pool-info vt (pool b)

```

```

      in vt( $\text{bb} := (\text{bb vt})(t := \text{bn vt } t \text{ div } 4 * 4 + i \text{ vt } t)$ ),
      mem-pool-info := (mem-pool-info vt)
      (pool b := minf ( $\text{levels} := \text{levels minf}[lvl \text{ vt } t := (\text{levels}$ 
minf ! lvl vt t)
      ( $\text{bits} := \text{bits} (\text{levels minf ! lvl vt t})[\text{bn vt } t \text{ div } 4 * 4 + i \text{ vt } t$ 
:= NOEXIST])])])
      (pool b)])
      apply(simp add:Let-def)
      apply(rule subst[where s=max-sz (mem-pool-info vt (pool b)) and t=max-sz
(mem-pool-info
      (let minf = mem-pool-info vt (pool b)
      in vt( $\text{bb} := (\text{bb vt})(t := \text{bn vt } t \text{ div } 4 * 4 + i \text{ vt } t)$ ),
      mem-pool-info := (mem-pool-info vt)
      (pool b := minf ( $\text{levels} := \text{levels minf}[lvl \text{ vt } t := (\text{levels}$ 
minf ! lvl vt t)
      ( $\text{bits} := \text{bits} (\text{levels minf ! lvl vt t})[\text{bn vt } t \text{ div } 4 * 4 + i \text{ vt } t$ 
:= NOEXIST])])])])
      (pool b)])
      apply(simp add:Let-def)

      apply(rule subst[where s=n-max (mem-pool-info V (pool b)) and t=n-max
(mem-pool-info vt (pool b))])
      apply(simp add:Let-def set-bit-def gvars-conf-stable-def gvars-conf-def)

      apply(rule subst[where s=max-sz (mem-pool-info V (pool b)) and t=max-sz
(mem-pool-info vt (pool b))])
      apply(simp add:Let-def set-bit-def gvars-conf-stable-def gvars-conf-def)
      apply(rule subst[where s=ALIGN4 (max-sz (mem-pool-info V (pool b))) div 4
^ lvl V t and t=lsz vt t])
      apply(simp add:Let-def) apply metis
      apply(rule subst[where s=block b div 4 ^ (level b - lvl V t) and t=bn vt t])
      apply(simp add:Let-def) apply metis
      apply(rule subst[where s=max-sz (mem-pool-info V (pool b)) and t=ALIGN4
(max-sz (mem-pool-info V (pool b)))]])
      apply(simp add:inv-def) using inv-mempool-info-maxsz-align4[rule-format, of
V pool b] apply metis

      apply(subgoal-tac length (bits ((levels (mem-pool-info V (pool b))) ! level b)) =
(n-max (mem-pool-info V (pool b))) * 4 ^ (level b))
      prefer 2 apply(simp add: inv-def inv-mempool-info-def Let-def)

      apply(subgoal-tac max-sz (mem-pool-info V (pool b)) mod 4 ^ lvl V t = 0)
      prefer 2 apply(subgoal-tac  $\exists n. \text{max-sz} (\text{mem-pool-info V (pool b)}) = (4 * n)$ 
* ( $4 ^ n\text{-levels} (\text{mem-pool-info V (pool b)})$ ))
      prefer 2 apply(simp add:inv-def) using inv-mempool-info-def[rule-format,
of V] apply meson
      apply(subgoal-tac length (levels (mem-pool-info V (pool b))) = n-levels
(mem-pool-info V (pool b)))

```



$$\begin{aligned} & \quad \quad \quad (\text{bits} := \text{bits} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) \\ & ! \text{lvl } vt \ t) [\text{bn } vt \ t := \text{FREE}])]) \\ & \quad \quad \quad \text{else mem-pool-info } V \ p) ! \\ & \quad \quad \quad i))) \implies \\ & \quad \quad \quad \forall j. j \neq \text{lvl } vt \ t \longrightarrow \text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! j = \text{levels} (\text{mem-pool-info } \\ & vt \ (\text{pool } b)) ! j \implies \\ & \quad \quad \quad \text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) = \text{length} (\text{levels} (\text{mem-pool-info } vt \ (\text{pool } \\ & b))) \implies \\ & \quad \quad \quad \text{length} (\text{bits} (\text{levels} (\text{mem-pool-info } vt \ (\text{pool } b)) \\ & \quad \quad \quad [\text{lvl } vt \ t := (\text{levels} (\text{mem-pool-info } vt \ (\text{pool } b)) ! \text{lvl } vt \ t) \\ & \quad \quad \quad (\text{bits} := \text{list-updates-n} (\text{bits} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } \\ & vt \ t) [\text{bn } vt \ t := \text{FREE}]) (\text{bn } vt \ t \text{ div } 4 * 4) (i \ vt \ t) \\ & \quad \quad \quad \text{NOEXIST} \\ & \quad \quad \quad [\text{bn } vt \ t := \text{NOEXIST}])]) ! \\ & \quad \quad \quad ia)) = \\ & \quad \quad \quad \text{length} (\text{bits} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) \\ & \quad \quad \quad [\text{lvl } vt \ t := (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } vt \ t) \\ & \quad \quad \quad (\text{bits} := \text{bits} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } vt \ t) [\text{bn } vt \ t := \\ & \text{FREE}])]) ! \\ & \quad \quad \quad ia)) \end{aligned}$$

**apply**(*case-tac* *ia* < *length* (*levels* (*mem-pool-info* *V* (*pool* *b*))))  
**apply**(*case-tac* *ia* = *lvl* *vt* *t*) **apply** *auto*[1]  
**apply** (*metis* (*no-types*, *lifting*) *nth-list-update-neq*)  
**by** (*smt* *list-eq-iff-nth-eq* *list-update-beyond* *not-less* *nth-list-update-neq*)

**lemma** *mempool-free-stm8-set4partbits-while-one-isuc-h1-2*:

$$\begin{aligned} & \neg \text{free-block-r } vt \ t \longrightarrow \text{freeing-node } V \ t = \text{None} \implies \\ & \quad \text{free-block-r } V \ t = \text{free-block-r } vt \ t \implies \\ & \quad \alpha = (\text{if } \exists y. \text{freeing-node } V \ t = \text{Some } y \text{ then } \text{lvl } V \ t + 1 \text{ else } \text{NULL}) \implies \\ & \quad \text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } vt \ t) = \text{bn } vt \ t \implies \\ & \quad V \in (\text{if } \text{NULL} < (\text{if } \exists y. \text{freeing-node } V \ t = \text{Some } y \text{ then } \text{lvl } V \ t + 1 \text{ else } \text{NULL}) \\ & \text{then } \text{UNIV} \text{ else } \{\}) \implies \\ & \quad \text{free-block-r } V \ t \\ & \text{by } \text{force} \end{aligned}$$

**lemma** *mempool-free-stm8-set4partbits-while-one-isuc-h2*:

$$\begin{aligned} & \text{inv } V \implies \\ & \quad \text{pool } b \in \text{mem-pools } V \implies \\ & \quad \text{lvl } vt \ t < \text{length} (\text{levels} (\text{mem-pool-info } V (\text{pool } b))) \implies \\ & \quad \text{bn } vt \ t < \text{length} (\text{bits} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } vt \ t)) \implies \\ & \quad \text{get-bit-s } V (\text{pool } b) (\text{lvl } vt \ t) (\text{bn } vt \ t) \neq \text{FREE} \implies \\ & \quad \text{buf} (\text{mem-pool-info } V (\text{pool } b)) + \text{max-sz} (\text{mem-pool-info } V (\text{pool } b)) \text{ div } 4 \wedge \\ & \quad \text{lvl } vt \ t * \text{bn } vt \ t \\ & \quad \notin \text{set} (\text{free-list} (\text{levels} (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } vt \ t)) \\ & \text{apply} (\text{simp } \text{add:inv-def } \text{inv-bitmap-freelist-def } \text{Let-def}) \\ & \text{apply} (\text{rule } \text{subst} [\text{where } t = \text{max-sz} (\text{mem-pool-info } V (\text{pool } b)) \text{ div } 4 \wedge \text{lvl } vt \ t * \text{bn} \end{aligned}$$



$vt\ t$  and  
 $s = bn\ vt\ t * (max\text{-}sz\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ div\ 4\ \wedge\ lvl\ vt\ t))$   
**apply simp apply simp**  
**done**

**lemma** *mempool-free-stm8-set4partbits-while-one-isuc-h1-3*:  
 $\forall p. (\forall i. length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ vt\ p)\ !\ i)) =$   
 $length\ (bits\ (levels\ (if\ p = pool\ b$   
 $then\ mem\text{-}pool\text{-}info\ V\ (pool\ b)$   
 $(\!|levels := levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$   
 $[lvl\ vt\ t := (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl$   
 $vt\ t)$   
 $(\!|bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$   
 $! lvl\ vt\ t)[bn\ vt\ t := FREE])])$   
 $else\ mem\text{-}pool\text{-}info\ V\ p)\ !$   
 $i))) \implies$   
 $\forall j. j \neq lvl\ vt\ t \longrightarrow levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ j = levels\ (mem\text{-}pool\text{-}info$   
 $vt\ (pool\ b))\ !\ j \implies$   
 $length\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) = length\ (levels\ (mem\text{-}pool\text{-}info\ vt\ (pool$   
 $b))) \implies$   
 $length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ vt\ (pool\ b))$   
 $[lvl\ vt\ t := (levels\ (mem\text{-}pool\text{-}info\ vt\ (pool\ b))\ !\ lvl\ vt\ t)$   
 $(\!|bits := list\text{-}updates\text{-}n\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl$   
 $vt\ t)[bn\ vt\ t := FREE])\ (bn\ vt\ t\ div\ 4 * 4)\ (i\ vt\ t)$   
 $NOEXIST$   
 $[bn\ vt\ t\ div\ 4 * 4 + i\ vt\ t := NOEXIST],$   
 $free\text{-}list :=$   
 $remove1\ (block\text{-}ptr$   
 $(mem\text{-}pool\text{-}info\ vt\ (pool\ b)$   
 $(\!|levels := levels\ (mem\text{-}pool\text{-}info\ vt\ (pool\ b))$   
 $[lvl\ vt\ t := (levels\ (mem\text{-}pool\text{-}info\ vt\ (pool\ b))\ !\ lvl\ vt$   
 $t)$   
 $(\!|bits := list\text{-}updates\text{-}n\ (bits\ (levels\ (mem\text{-}pool\text{-}info$   
 $V\ (pool\ b))\ !\ lvl\ vt\ t)[bn\ vt\ t := FREE])$   
 $(bn\ vt\ t\ div\ 4 * 4)\ (i\ vt\ t)\ NOEXIST$   
 $[bn\ vt\ t\ div\ 4 * 4 + i\ vt\ t := NOEXIST])])$   
 $(lsz\ vt\ t)\ (bn\ vt\ t\ div\ 4 * 4 + i\ vt\ t))$   
 $(removes\ (map\ (\lambda ii. block\text{-}ptr$   
 $(mem\text{-}pool\text{-}info\ V\ (pool\ b)$   
 $(\!|levels := levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$   
 $[lvl\ vt\ t := (levels\ (mem\text{-}pool\text{-}info\ V\ (pool$   
 $b))\ !\ lvl\ vt\ t)$   
 $(\!|bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V$   
 $(pool\ b))\ !\ lvl\ vt\ t)[bn\ vt\ t := FREE])])$   
 $(lsz\ vt\ t)\ (bn\ vt\ t\ div\ 4 * 4 + ii))$   
 $[NULL..<i\ vt\ t])$   
 $(free\text{-}list\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ vt\ t)))])\ !$   
 $ia)) =$

length (bits (levels (mem-pool-info V (pool b)))  
 [lvl vt t := (levels (mem-pool-info V (pool b)) ! lvl vt t)  
 (bits := bits (levels (mem-pool-info V (pool b)) ! lvl vt t)[bn vt t :=  
 FREE]]) !  
 ia))

**apply**(case-tac ia < length (levels (mem-pool-info V (pool b))))  
**apply**(case-tac ia = lvl vt t) **apply** auto[1]  
**apply** (metis (no-types, lifting) nth-list-update-neq)  
**by** (smt list-eq-iff-nth-eq list-update-beyond not-less nth-list-update-neq)

**lemma** mempool-free-stm8-set4partbits-while-one-isuc-h1:  
 $V \in mp\text{-free-precond8-3 } t \ b \ \alpha \cap \llbracket 'cur = Some \ t \rrbracket \implies$   
 $vt \in free\text{-stm8-precond3 } V \ t \ b \cap \llbracket 'i \ t < 4 \rrbracket \implies$   
 $\{free\text{-stm8-precond2 } V \ t \ b\} \cap \llbracket NULL < 'lvl \ t \wedge partner\text{-bits } ('mem\text{-pool-info}$   
 $(pool \ b)) ('lvl \ t) ('bn \ t) \rrbracket \neq \{\} \implies$   
 $\{let \ vv = free\text{-stm8-precond33 } (free\text{-stm8-precond32 } (free\text{-stm8-precond31 } vt \ t$   
 $b) \ t \ b) \ t \ b \ in$   
 $if \ bn \ vv \ t = bb \ vv \ t \ then \ vv \ else \ free\text{-stm8-precond34 } vv \ t \ b\}$   
 $\subseteq \{s. s[i := (i \ s) \ (t := Suc \ (i \ s \ t))]\} \in free\text{-stm8-precond3 } V \ t \ b\}$   
**apply**(simp add:Let-def set-bit-def)

**apply**(rule conjI)  
**apply** clarsimp  
**apply**(rule conjI)  
**apply**(simp add: gvars-conf-stable-def gvars-conf-def)  
**apply** clarsimp  
**apply**(subgoal-tac length (levels (mem-pool-info V (pool b)))=length (levels  
(mem-pool-info vt (pool b))))  
**prefer** 2 **apply** simp  
**using** mempool-free-stm8-set4partbits-while-one-isuc-h1-1 **apply** blast  
**apply**(rule conjI)  
**apply**(subgoal-tac length (levels (mem-pool-info V (pool b)))=length (levels  
(mem-pool-info vt (pool b))))  
**prefer** 2 **apply**(simp add: gvars-conf-stable-def gvars-conf-def)

**apply**(rule subst[**where** s=list-updates-n (bits (levels (mem-pool-info V (pool  
b)) ! lvl vt t)[bn vt t := FREE]) (bn vt t div 4 \* 4) (i vt t) NOEXIST  
[bn vt t := NOEXIST] **and** t=bits (levels (mem-pool-info vt (pool b))  
[lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)  
(bits := list-updates-n (bits (levels (mem-pool-info V (pool b)) ! lvl vt  
t)[bn vt t := FREE]) (bn vt t div 4 \* 4) (i vt t) NOEXIST  
[bn vt t := NOEXIST])]) !  
lvl vt t)]) **apply** auto[1]  
**using** lst-updts-eq-updts-updt[of Suc (i vt t) bits (levels (mem-pool-info V (pool  
b)) ! lvl vt t)[bn vt t := FREE]  
bn vt t div 4 \* 4 NOEXIST] **apply** auto[1]  
**apply**(rule conjI)  
**apply**(simp add:block-ptr-def)

```

apply(rule subst[where  $s = \text{free-list } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)) \text{ ! } lvl \text{ } vt \text{ } t)$  and
   $t = \text{free-list } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b))$ 
   $[lvl \text{ } vt \text{ } t := (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)) \text{ ! } lvl \text{ } vt \text{ } t)$ 
   $(\text{bits} := \text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))$ 
   $\text{! } lvl \text{ } vt \text{ } t)[bn \text{ } vt \text{ } t := \text{FREE}]) (bn \text{ } vt \text{ } t \text{ div } 4 * 4) (i \text{ } vt \text{ } t) \text{ NOEXIST}$ 
   $[bn \text{ } vt \text{ } t := \text{NOEXIST}]) \text{! } lvl \text{ } vt \text{ } t)])$ 
  apply(case-tac  $lvl \text{ } vt \text{ } t < \text{length } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)))$ ) apply
  auto[1] apply auto[1]
  apply(subgoal-tac removes (map ( $\lambda ii. \text{buf } (\text{mem-pool-info } V \text{ (pool } b)) + \text{lsz } vt$ 
   $t * (bn \text{ } vt \text{ } t \text{ div } 4 * 4 + ii)$ ) [ $\text{NULL}..<i \text{ } vt \text{ } t$ ] @
     $[\text{buf } (\text{mem-pool-info } V \text{ (pool } b)) + \text{lsz } vt \text{ } t * bn \text{ } vt \text{ } t]$ )
     $(\text{free-list } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } vt \text{ } t)) =$ 
    removes (map ( $\lambda ii. \text{buf } (\text{mem-pool-info } V \text{ (pool } b)) + \text{lsz } vt \text{ } t *$ 
     $(bn \text{ } vt \text{ } t \text{ div } 4 * 4 + ii)$ ) [ $\text{NULL}..<i \text{ } vt \text{ } t$ ])
     $(\text{free-list } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } vt \text{ } t)))$ ) apply
  metis
  apply(rule rmvs-onemore-same)
  apply(simp add:inv-def inv-bitmap-freelist-def Let-def)
  apply(subgoal-tac get-bit (mem-pool-info V) (pool b) (lvl vt t) (bn vt t) =
  FREEING) prefer 2
  apply(subgoal-tac free-block-r V t) prefer 2
  using mempool-free-stm8-set4partbits-while-one-isuc-h1-2 apply blast
  apply(subgoal-tac  $\exists blk. \text{freeing-node } V \text{ } t = \text{Some } blk \wedge \text{pool } blk = \text{pool } b \wedge$ 
   $\text{level } blk = lvl \text{ } vt \text{ } t \wedge \text{block } blk = bn \text{ } vt \text{ } t$ )
  prefer 2 apply fast
  apply(simp add:inv-def inv-aux-vars-def) apply metis
  apply(rule subst[where  $s = \text{max-sz } (\text{mem-pool-info } V \text{ (pool } b)) \text{ div } 4 \wedge lvl \text{ } vt \text{ } t$ 
and  $t = \text{lsz } vt \text{ } t$ ])
  using inv-maxsz-align4[rule-format, of V pool b] apply force
  apply(subgoal-tac get-bit (mem-pool-info V) (pool b) (lvl vt t) (bn vt t)  $\neq$ 
  FREE) prefer 2 apply auto[1]
  apply(subgoal-tac  $lvl \text{ } vt \text{ } t < \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)))$ ) prefer
  2 apply force
  using mempool-free-stm8-set4partbits-while-one-isuc-h2 apply blast

apply clarsimp apply(simp add:block-ptr-def lvars-nochange-def)

apply clarsimp
apply(rule conjI)
apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply clarsimp
apply(subgoal-tac  $\text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) = \text{length } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)))$ )
prefer 2 apply simp
using mempool-free-stm8-set4partbits-while-one-isuc-h1-3 apply blast
apply(rule conjI)
apply(subgoal-tac  $\text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) = \text{length } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)))$ )

```

```

prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(rule subst[where s=list-updates-n (bits (levels (mem-pool-info V (pool
b))) ! lvl vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t) NOEXIST
[bn vt t div 4 * 4 + i vt t := NOEXIST] and t=bits (levels
(mem-pool-info vt (pool b))
[lvl vt t :=
(levels (mem-pool-info vt (pool b))
[lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)
(|bits := list-updates-n (bits (levels (mem-pool-info V (pool b)) ! lvl
vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t) NOEXIST
[bn vt t div 4 * 4 + i vt t := NOEXIST])]) !
lvl vt t)
(|free-list :=
remove1 (block-ptr
(mem-pool-info vt (pool b))
(|levels := levels (mem-pool-info vt (pool b))
[lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)
(|bits := list-updates-n (bits (levels (mem-pool-info V
(pool b)) ! lvl vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t)
NOEXIST
[bn vt t div 4 * 4 + i vt t := NOEXIST])])])
(lsz vt t) (bn vt t div 4 * 4 + i vt t))
free-list
(levels (mem-pool-info vt (pool b))
[lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)
(|bits := list-updates-n (bits (levels (mem-pool-info V (pool b)) !
lvl vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t) NOEXIST
[bn vt t div 4 * 4 + i vt t := NOEXIST])]) !
lvl vt t))]) !
lvl vt t)]) apply auto[1]
using lst-updts-eq-updts-updt[of Suc (i vt t) bits (levels (mem-pool-info V (pool
b))) ! lvl vt t)[bn vt t := FREE]
bn vt t div 4 * 4 NOEXIST] apply auto[1]

apply(rule conjI)
apply(simp add:block-ptr-def)
apply(subgoal-tac lvl vt t < length (levels (mem-pool-info vt (pool b)))) prefer
2
apply(subgoal-tac length (levels (mem-pool-info V (pool b)))=length (levels
(mem-pool-info vt (pool b))))
prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply force
apply(rule subst[where s=free-list (levels (mem-pool-info vt (pool b)) ! lvl vt
t) and
t=free-list (levels (mem-pool-info vt (pool b))
[lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)
(|bits := list-updates-n (bits (levels (mem-pool-info V (pool b))
! lvl vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t) NOEXIST
[bn vt t div 4 * 4 + i vt t := NOEXIST])]) ! lvl vt t)])

```

```

apply auto[1]
apply(rule subst[where s=remove1 (buf (mem-pool-info vt (pool b)) + lsz vt
t * (bn vt t div 4 * 4 + i vt t))
```

$$(\text{free-list } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)) ! \text{lvl } vt \text{ } t)) \text{ and}$$

$$t = \text{free-list } (\text{levels } (\text{mem-pool-info } vt \text{ (pool } b)))$$

```

    [lvl vt t :=
      (levels (mem-pool-info vt (pool b))
        [lvl vt t := (levels (mem-pool-info vt (pool b)) ! lvl vt t)
          (bits := list-updates-n (bits (levels (mem-pool-info V (pool
b)) ! lvl vt t)[bn vt t := FREE]) (bn vt t div 4 * 4) (i vt t) NOEXIST
            [bn vt t div 4 * 4 + i vt t := NOEXIST])]) !
        lvl vt t)
      (free-list :=
        remove1 (buf (mem-pool-info vt (pool b)) + lsz vt t * (bn vt
t div 4 * 4 + i vt t))
          (free-list (levels (mem-pool-info vt (pool b)) ! lvl vt t))]) !
        lvl vt t))
    apply auto[1]
    apply(subgoal-tac buf (mem-pool-info vt (pool b)) = buf (mem-pool-info V
(pool b))) prefer 2
    apply(simp add: gvars-conf-stable-def gvars-conf-def)
    using rmvs-rev[of (map ( $\lambda ii.$  buf (mem-pool-info V (pool b)) + lsz vt t * (bn
vt t div 4 * 4 + ii)) [NULL..<i vt t])
      buf (mem-pool-info V (pool b)) + lsz vt t * (bn vt t div 4 * 4 +
i vt t)
      free-list (levels (mem-pool-info V (pool b)) ! lvl vt t) ] apply simp

apply clarsimp apply(simp add:block-ptr-def lvars-nochange-def)
done

```

**lemma** *mempool-free-stm8-set4partbits-while-one-isuc*:

$$V = Va \implies$$

$$V \in \text{mp-free-precond8-3 } t \text{ } b \text{ } \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \implies$$

$$vt \in \text{free-stm8-precond3 } Va \text{ } t \text{ } b \cap \llbracket 'i \text{ } t < 4 \rrbracket \implies$$

$$\{\text{free-stm8-precond2 } V \text{ } t \text{ } b\} \cap \llbracket \text{NULL} < 'lvl \text{ } t \wedge \text{partner-bits } ('mem\text{-pool-info}$$

$$(\text{pool } b)) ('lvl \text{ } t) ('bn \text{ } t) \rrbracket \neq \{\} \implies$$

$$\Gamma \vdash_I \text{Some } ('i := 'i(t := \text{Suc } ('i \text{ } t)))$$

$$\text{sat}_p [\{\text{let } vv = \text{free-stm8-precond33 } (\text{free-stm8-precond32 } (\text{free-stm8-precond31}$$

$$vt \text{ } t \text{ } b) \text{ } t \text{ } b \text{ in}$$

$$\text{if } bn \text{ } vv \text{ } t = bb \text{ } vv \text{ } t \text{ then } vv \text{ else free-stm8-precond34 } vv \text{ } t \text{ } b\}, \{(s, t). s =$$

$$t\}, \text{UNIV}, \text{free-stm8-precond3 } Va \text{ } t \text{ } b]$$

```

apply(rule Basic)
defer 1
apply fast using stable-id2 apply fast using stable-id2 apply fast

using mempool-free-stm8-set4partbits-while-one-isuc-h1[of Va t b  $\alpha$  vt] apply
fast

```

done

**lemma** *mempool-free-stm8-set4partbits-while-one*:

$V = Va \implies$   
 $V \in mp\text{-}free\text{-}precond8\text{-}3\ t\ b\ \alpha \cap \{\cur = Some\ t\} \implies$   
 $vt \in free\text{-}stm8\text{-}precond3\ Va\ t\ b \cap \{i\ t < 4\} \implies$   
 $\{free\text{-}stm8\text{-}precond2\ V\ t\ b\} \cap \{NULL < \text{'lvl}\ t \wedge partner\text{-}bits\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))\ (\text{'lvl}\ t)\ (\text{'bn}\ t)\} \neq \{\} \implies$   
 $\Gamma \vdash_I Some\ (\text{'bb} := \text{'bb}(t := \text{'bn}\ t\ div\ 4 * 4 + \text{'i}\ t));;$   
 $\text{'mem}\text{-}pool\text{-}info := set\text{-}bit\text{-}noexist\ \text{'mem}\text{-}pool\text{-}info\ (pool\ b)\ (\text{'lvl}\ t)\ (\text{'bb}\ t);;$   
 $\text{'block}\text{-}pt := \text{'block}\text{-}ptr\ (t := block\text{-}ptr\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))\ (\text{'lsz}\ t)\ (\text{'bb}\ t));;$   
 $IF\ \text{'bn}\ t \neq \text{'bb}\ t \wedge block\text{-}fits\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))\ (\text{'block}\text{-}pt\ t)\ (\text{'lsz}\ t)$   
 $THEN$   
 $\text{'mem}\text{-}pool\text{-}info := \text{'mem}\text{-}pool\text{-}info(pool\ b := remove\text{-}free\text{-}list\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))\ (\text{'lvl}\ t)\ (\text{'block}\text{-}pt\ t))$   
 $FI;;$   
 $i := i(t := Suc\ (i\ t))$   
 $sat_p\ [\{vt\}, \{(s, t). s = t\}, UNIV, free\text{-}stm8\text{-}precond3\ Va\ t\ b]$   
**apply**(rule Seq[**where**  $mid = \{let\ vv = free\text{-}stm8\text{-}precond33\ (free\text{-}stm8\text{-}precond32\ (free\text{-}stm8\text{-}precond31\ vt\ t\ b)\ t\ b\ in$   
 $if\ bn\ vv\ t = bb\ vv\ t\ then\ vv\ else\ free\text{-}stm8\text{-}precond34\ vv\ t$   
 $b\}\}$ ])  
**apply**(rule Seq[**where**  $mid = \{free\text{-}stm8\text{-}precond33\ (free\text{-}stm8\text{-}precond32\ (free\text{-}stm8\text{-}precond31\ vt\ t\ b)\ t\ b)\}$ ])  
**apply**(rule Seq[**where**  $mid = \{free\text{-}stm8\text{-}precond32\ (free\text{-}stm8\text{-}precond31\ vt\ t\ b)\}$ ])  
**apply**(rule Seq[**where**  $mid = \{free\text{-}stm8\text{-}precond31\ vt\ t\ b\}$ ])  
  
**apply**(rule Basic)  
**apply fast apply fast using stable-id2 apply fast using stable-id2 apply fast**  
  
**apply**(rule Basic)  
**apply**(simp add:Let-def set-bit-def)  
**apply fast using stable-id2 apply fast using stable-id2 apply fast**  
  
**apply**(rule Basic)  
**apply fast apply fast using stable-id2 apply fast using stable-id2 apply fast**  
  
**apply**(rule Cond)  
**using stable-id2 apply fast**  
  
**apply**(rule Basic)  
**apply**(simp add:Let-def remove-free-list-def) **apply auto[1]**

**apply fast using stable-id2 apply fast using stable-id2 apply fast**

**apply**(*case-tac* *bn* (*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*) *t*  
 $\neq$  *bb* (*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*) *t*)  
**apply**(*rule subst*[**where** *s*={*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*}  $\cap$   
 $\neg \llbracket \text{block-fits } ('mem\text{-pool-info } (pool\ b)) ('block\text{-pt } t) ('lsz\ t) \rrbracket$   
**and** *t*={*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*}  $\cap$   
 $\neg \llbracket 'bn\ t \neq 'bb\ t \wedge \text{block-fits } ('mem\text{-pool-info } (pool\ b)) ('block\text{-pt } t) ('lsz\ t) \rrbracket$ )] **apply fast**  
**apply**(*rule subst*[**where** *s*={} **and** *t*={*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*}  $\cap$   
 $\neg \llbracket \text{block-fits } ('mem\text{-pool-info } (pool\ b)) ('block\text{-pt } t) ('lsz\ t) \rrbracket$ )]  
**using** *block-fits1*[*of V t b  $\alpha$  vt*] **apply fast**  
**using** *Emptyprecond* **apply fast**

**apply**(*rule subst*[**where** *s*={*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*}  
**and** *t*={*free-stm8-precond33* (*free-stm8-precond32* (*free-stm8-precond31* *vt t b*) *t b*) *t b*}  $\cap$   
 $\neg \llbracket 'bn\ t \neq 'bb\ t \wedge \text{block-fits } ('mem\text{-pool-info } (pool\ b)) ('block\text{-pt } t) ('lsz\ t) \rrbracket$ )]  
**apply fast**

**apply**(*unfold Skip-def*)  
**apply**(*rule Basic*)  
**using** *mempool-free-stm8-set4partbits-while-one-h1* **apply fast**  
**apply fast**  
**using** *stable-id2* **apply fast**  
**using** *stable-id2* **apply fast**

**apply fast**

**using** *mempool-free-stm8-set4partbits-while-one-isuc*[*of V Va t b  $\alpha$  vt*] **apply fast**  
**done**

**lemma** *mempool-free-stm8-set4partbits-while*:

$V = Va \implies$   
 $V \in mp\text{-free-precond8-3 } t\ b\ \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \implies$   
 $\{free\text{-stm8-precond2 } V\ t\ b\} \cap \llbracket NULL < 'lwl\ t \wedge partner\text{-bits } ('mem\text{-pool-info } (pool\ b)) ('lwl\ t) ('bn\ t) \rrbracket \neq \{\}$   
 $\Gamma \vdash_I \text{Some}('bb := 'bb(t := 'bn\ t\ div\ 4 * 4 + 'i\ t));$   
 $'mem\text{-pool-info} := set\text{-bit-noexist } 'mem\text{-pool-info } (pool\ b) ('lwl\ t) ('bb\ t);;$

$\text{'block-pt} := \text{'block-pt } (t := \text{block-ptr } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lsz } t) (\text{'bb } t));;$   
 $\text{IF } \text{'bn } t \neq \text{'bb } t \wedge \text{block-fits } (\text{'mem-pool-info } (\text{pool } b)) (\text{'block-pt } t) (\text{'lsz } t)$   
 $\text{THEN}$   
 $\text{'mem-pool-info} := \text{'mem-pool-info}(\text{pool } b := \text{remove-free-list } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lvl } t) (\text{'block-pt } t))$   
 $\text{FI};;$   
 $\text{'i} := \text{'i}(t := \text{Suc } (\text{'i } t))$   
 $\text{sat}_p [\text{free-stm8-precond3 } \forall a \ t \ b \cap \llbracket \text{'i } t < 4 \rrbracket, \{(s, t). s = t\}, \text{UNIV}, \text{free-stm8-precond3 } \forall a \ t \ b]$   
**using**  $\text{mempool-free-stm8-set4partbits-while-one}$  [of  $V \ \forall a \ t \ b \ \alpha$ ]  
 $\text{Allprecond}[\text{where } U = \text{free-stm8-precond3 } \forall a \ t \ b \cap \llbracket \text{'i } t < 4 \rrbracket \text{ and}$   
 $P = \text{Some } (\text{'bb} := \text{'bb}(t := \text{'bn } t \text{ div } 4 * 4 + \text{'i } t));$   
 $\text{'mem-pool-info} := \text{set-bit-noexist } \text{'mem-pool-info } (\text{pool } b) (\text{'lvl } t) (\text{'bb } t);;$   
 $\text{'block-pt} := \text{'block-pt } (t := \text{block-ptr } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lsz } t) (\text{'bb } t));;$   
 $\text{IF } \text{'bn } t \neq \text{'bb } t \wedge \text{block-fits } (\text{'mem-pool-info } (\text{pool } b)) (\text{'block-pt } t) (\text{'lsz } t) \text{ THEN}$   
 $\text{'mem-pool-info} := \text{'mem-pool-info}(\text{pool } b := \text{remove-free-list } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lvl } t) (\text{'block-pt } t))$   
 $\text{FI};;$   
 $\text{'i} := \text{'i}(t := \text{Suc } (\text{'i } t)) \text{ and}$   
 $\text{rely} = \{(x, y). x = y\} \text{ and}$   
 $\text{guar} = \text{UNIV} \text{ and } \text{post} = \text{free-stm8-precond3 } \forall a \ t \ b]$   
**apply**  $\text{meson}$   
**done**

**term**  $\text{free-stm8-precond3 } \forall a \ t \ b$

**abbreviation**  $\text{free-stm8-atombody-rest-cond1 } V \ t \ b \equiv V(\llbracket \text{lvl} := (\text{lvl } V)(t := \text{lvl } V \ t - 1) \rrbracket)$

**abbreviation**  $\text{free-stm8-atombody-rest-cond2 } V \ t \ b \equiv V(\llbracket \text{bn} := (\text{bn } V)(t := \text{bn } V \ t \text{ div } 4) \rrbracket)$

**abbreviation**  $\text{free-stm8-atombody-rest-cond3 } V \ t \ b \equiv$

$\text{let } \text{minf} = \text{mem-pool-info } V \ (\text{pool } b) \text{ in}$

$V(\llbracket \text{mem-pool-info} := (\text{mem-pool-info } V) (\text{pool } b := \text{minf } \llbracket \text{levels} := (\text{levels } \text{minf})$

$\llbracket \text{lvl } V \ t := ((\text{levels } \text{minf}) ! (\text{lvl } V \ t)) \rrbracket \llbracket \text{bits} := (\text{bits } ((\text{levels } \text{minf}) ! (\text{lvl } V \ t)))$   
 $\llbracket \text{bn } V \ t := \text{FREEING} \rrbracket \rrbracket)$

**lemma**  $\text{mempool-free-stm8-atombody-rest-one-finalstm-inv-cur}$ :

$V \in \text{mp-free-precond8-3 } t \ b \ \alpha \cap \llbracket \text{'cur} = \text{Some } t \rrbracket \implies$

$(\{ \text{free-stm8-precond2 } V \ t \ b \} \cap \llbracket \text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lvl } t) (\text{'bn } t) \rrbracket \neq \{\} \implies *)$

$V2 \in \text{free-stm8-precond3 } V \ t \ b \cap \llbracket \text{'i } t = 4 \rrbracket \implies$

$x = \text{free-stm8-atombody-rest-cond3 } (V2(\llbracket \text{lvl} := (\text{lvl } V2)(t := \text{lvl } V2 \ t - 1), \text{bn}$



```

:= (bn V2)(t := bn V2 t div 4)) t b ==>
  y = x(freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
t, block = bn x t,
      data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))) ==>
  inv-cur y
apply(rule subst[where s=inv-cur x and t=inv-cur y])
apply(simp add:block-ptr-def inv-cur-def)

apply(simp add:Let-def inv-def inv-cur-def)
apply(subgoal-tac thd-state V t = RUNNING) prefer 2 apply fast
apply clarsimp
done

```

```

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv-thd-waitq:
  V ∈ mp-free-precond8-3 t b α ∩ {cur = Some t} ==>
    (* {free-stm8-precond2 V t b} ∩ {NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)} ≠ {} ==> *)
    V2 ∈ free-stm8-precond3 V t b ∩ {i t = 4} ==>
      x = free-stm8-atombody-rest-cond3 (V2(lvl := (lvl V2)(t := lvl V2 t - 1), bn
:= (bn V2)(t := bn V2 t div 4)) t b ==>
        y = x(freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
t, block = bn x t,
            data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))) ==>
          inv-thd-waitq y
apply(rule subst[where s=inv-thd-waitq x and t=inv-thd-waitq y])
apply(simp add:block-ptr-def inv-thd-waitq-def)

apply(simp add:Let-def inv-def inv-thd-waitq-def)
apply(simp add:set-bit-def)
apply(subgoal-tac mem-pools V = mem-pools V2)
  prefer 2 apply(simp add:gvars-conf-stable-def gvars-conf-def)

```

```

apply(rule conjI)
  apply clarify apply metis

```

```

apply(rule conjI)
  apply clarify apply metis

```

```

apply(rule conjI)
  apply clarify apply metis
apply metis
done

```

```

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h1:
  ∀ p. buf (mem-pool-info V2 p) =
    buf (if p = pool b

```

$$\begin{aligned}
& \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t)([bits := bits \\
& ((levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t)[block b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p) \wedge \\
& max-sz (\textit{mem-pool-info } V2 p) = \\
& max-sz (\textit{if } p = \textit{pool } b \\
& \quad \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t) \\
& \quad ([bits := bits (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t)[block \\
& b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p) \wedge \\
& n-max (\textit{mem-pool-info } V2 p) = \\
& n-max (\textit{if } p = \textit{pool } b \\
& \quad \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t) \\
& \quad ([bits := bits (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t)[block \\
& b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p) \wedge \\
& n-levels (\textit{mem-pool-info } V2 p) = \\
& n-levels (\textit{if } p = \textit{pool } b \\
& \quad \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t) \\
& \quad ([bits := bits (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V \\
& t)[block b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p) \wedge \\
& length (levels (\textit{mem-pool-info } V2 p)) = \\
& length (levels (\textit{if } p = \textit{pool } b \\
& \quad \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t) \\
& \quad ([bits := bits (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V \\
& t)[block b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p))) \wedge \\
& (\forall i . length (bits (levels (\textit{mem-pool-info } V2 p) ! i)) = \\
& \quad length (bits (levels (\textit{if } p = \textit{pool } b \\
& \quad \textit{then mem-pool-info } V (\textit{pool } b) \\
& \quad ([levels := levels (\textit{mem-pool-info } V (\textit{pool } b)) \\
& \quad [lvl V t := (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V t) \\
& \quad ([bits := bits (levels (\textit{mem-pool-info } V (\textit{pool } b)) ! lvl V \\
& t)[block b \operatorname{div} 4 ^ {(\textit{level } b - lvl V t)} := FREE]]))]]) \\
& \quad \textit{else mem-pool-info } V p)! \\
& \quad i)))) \implies \\
& ia < length (levels (\textit{mem-pool-info } V (\textit{pool } b))) \implies
\end{aligned}$$

$length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ ia)) = length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ ia))$   
**apply** *auto*  
**apply**(*case-tac* *lvl V t = ia*) **apply** *auto*[1] **apply** *auto*[1]  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h2*:  
 $ia < length\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) \implies$   
 $length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b)))$   
 $[lvl\ V2\ t - Suc\ NULL := (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ NULL))$   
 $NULL])$   
 $([bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ NULL))[bn$   
 $V2\ t\ div\ 4 := FREEING]])\ !$   
 $ia)) = length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ ia))$   
**apply**(*case-tac* *lvl V2 t - Suc NULL = ia*)  
**apply**(*case-tac* *ia < length\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b)))*) **apply** *auto*  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h3*:  
 $mem\text{-}pools\ V = mem\text{-}pools\ V2 \implies$   
 $p \in mem\text{-}pools\ V2 \implies$   
 $\forall p \in mem\text{-}pools\ V2.$   
 $NULL < buf\ (mem\text{-}pool\text{-}info\ V\ p) \wedge$   
 $(\exists n > NULL. max\text{-}sz\ (mem\text{-}pool\text{-}info\ V\ p) = 4 * n * 4 \wedge n\text{-}levels\ (mem\text{-}pool\text{-}info$   
 $V\ p)) \wedge$   
 $NULL < n\text{-}max\ (mem\text{-}pool\text{-}info\ V\ p) \wedge$   
 $NULL < n\text{-}levels\ (mem\text{-}pool\text{-}info\ V\ p) \wedge$   
 $n\text{-}levels\ (mem\text{-}pool\text{-}info\ V\ p) = length\ (levels\ (mem\text{-}pool\text{-}info\ V\ p)) \wedge$   
 $(\forall i < length\ (levels\ (mem\text{-}pool\text{-}info\ V\ p)). length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V$   
 $p)\ !\ i)) = n\text{-}max\ (mem\text{-}pool\text{-}info\ V\ p) * 4 \wedge i) \implies$   
 $mem\text{-}pools\ V = mem\text{-}pools\ V2 \implies$   
 $pool\ b \in mem\text{-}pools\ V2 \implies levels\ (mem\text{-}pool\text{-}info\ V\ p) \neq []$   
**apply** *auto*  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h1-1'*:  
 $\forall j. j \neq lvl\ V\ t \longrightarrow$   
 $levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$   
 $[lvl\ V\ t := (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)$   
 $([bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)[block\ b\ div\ 4 \wedge$   
 $(level\ b - lvl\ V\ t) := FREE]])\ !$   
 $j =$   
 $levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ j \implies$   
 $bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b))\ !\ lvl\ V\ t) =$   
 $list\text{-}updates\text{-}n$   
 $(bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$   
 $[lvl\ V\ t := (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)$   
 $([bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)[block\ b\ div\ 4 \wedge$

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(level b - lvl V t) := FREE]] !
  lvl V t))
  (block b div 4 ^ (level b - lvl V t) div 4 * 4) 4 NOEXIST ==>
length (bits (levels (mem-pool-info V (pool b)) ! ia)) =
length (bits (levels (mem-pool-info V2 (pool b))
  [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2
t - Suc NULL))
  (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL))][bn V2 t div 4 := FREEING]] !
  ia))
apply(rule subst[where s=length (bits (levels (mem-pool-info V2 (pool b))!ia))
and t=length (bits (levels (mem-pool-info V2 (pool b))
  [lvl V2 t - Suc 0 := (levels (mem-pool-info V2 (pool b)) ! (lvl V2
t - Suc 0))
  (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc 0))][bn V2 t div 4 := FREEING]] !
  ia)))]
apply(case-tac ia = lvl V2 t - Suc 0)
apply(case-tac ia < length (levels (mem-pool-info V2 (pool b))))
apply auto[1] apply auto[1] apply auto[1]

apply(case-tac ia = lvl V t)
apply(subgoal-tac length (list-updates-n
  (bits (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4
^ (level b - lvl V t) := FREE]] !
    ia))
    (block b div 4 ^ (level b - lvl V t) div 4 * 4) 4 NOEXIST) = length (bits
(levels (mem-pool-info V (pool b))
  [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
  (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4
^ (level b - lvl V t) := FREE]] !
  ia))))
prefer 2 using length-list-update-n apply fast
apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
  [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
  (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b
div 4 ^ (level b - lvl V t) := FREE]] !
  ia)) = length (bits (levels (mem-pool-info V (pool b)) ! ia)))
prefer 2 apply(case-tac ia = lvl V t )
apply(case-tac ia < length (levels (mem-pool-info V (pool b))))
apply auto[1] apply auto[1] apply auto[1]
apply auto[1]

apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
  [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
  (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level

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$b - lvl\ V\ t) := FREE]]] !$   
 $ia)) = length\ (bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! ia))$   
**prefer 2 apply**( $case\ tac\ ia = lvl\ V\ t$ ) **apply**( $case\ tac\ ia < length\ (levels$   
 $(mem\ pool\ info\ V\ (pool\ b))))$   
**apply**  $auto[1]$  **apply**  $auto[1]$  **apply**  $auto[1]$   
**apply**( $subgoal\ tac\ levels\ (mem\ pool\ info\ V\ (pool\ b))$   
 $[lvl\ V\ t := (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! lvl\ V\ t)$   
 $(bits := bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! lvl\ V\ t)[block\ b\ div\ 4 \wedge$   
 $(level\ b - lvl\ V\ t) := FREE]]] !$   
 $ia = levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! ia)$   
**prefer 2 apply fast**  
**apply auto**  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h1-1:*

$\forall j. j \neq lvl\ V\ t \longrightarrow$   
 $levels\ (set\ bit\ free\ (mem\ pool\ info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level$   
 $b - lvl\ V\ t))\ (pool\ b)) ! j$   
 $= levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! j \implies$   
 $bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! lvl\ V\ t) =$   
 $list\ updates\ n\ (bits\ (levels\ (set\ bit\ free\ (mem\ pool\ info\ V)\ (pool\ b)\ (lvl\ V\ t)$   
 $(block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))\ (pool\ b)) ! lvl\ V\ t))$   
 $(block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t)\ div\ 4 * 4)\ 4\ NOEXIST \implies$   
 $length\ (bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! ia)) =$   
 $length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))$   
 $[lvl\ V2\ t - Suc\ NULL := (levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! (lvl\ V2$   
 $t - Suc\ NULL))$   
 $(bits := bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! (lvl\ V2\ t - Suc$   
 $NULL)))[bn\ V2\ t\ div\ 4 := FREEING]]] !$   
 $ia))$   
**apply**( $rule\ subst[where\ s=length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))!ia))$   
**and**  $t=length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))$   
 $[lvl\ V2\ t - Suc\ 0 := (levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! (lvl\ V2$   
 $t - Suc\ 0))$   
 $(bits := bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b)) ! (lvl\ V2\ t -$   
 $Suc\ 0)))[bn\ V2\ t\ div\ 4 := FREEING]]] !$   
 $ia))])$   
**apply**( $case\ tac\ ia = lvl\ V2\ t - Suc\ 0$ )  
**apply**( $case\ tac\ ia < length\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))))$   
**apply**  $auto[1]$  **apply**  $auto[1]$  **apply**  $auto[1]$

**apply**( $simp\ add:set\ bit\ def$ )  
**apply**( $case\ tac\ ia = lvl\ V\ t$ )  
**apply**( $subgoal\ tac\ length\ (list\ updates\ n$   
 $(bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b))$   
 $[lvl\ V\ t := (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! lvl\ V\ t)$   
 $(bits := bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b)) ! lvl\ V\ t)[block\ b\ div\ 4$   
 $\wedge (level\ b - lvl\ V\ t) := FREE]]] !$   
 $ia))$

```

    (block b div 4 ^ (level b - lvl V t) div 4 * 4) 4 NOEXIST) = length (bits
(levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4
^ (level b - lvl V t) := FREE]] !
    ia)))
  prefer 2 using length-list-update-n apply fast
  apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b
div 4 ^ (level b - lvl V t) := FREE]] !
    ia)) = length (bits (levels (mem-pool-info V (pool b)) ! ia)))
  prefer 2 apply(case-tac ia = lvl V t)
  apply(case-tac ia < length (levels (mem-pool-info V (pool b))))
  apply auto[1] apply auto[1] apply auto[1]
  apply auto[1]

  apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level
b - lvl V t) := FREE]] !
    ia)) = length (bits (levels (mem-pool-info V (pool b)) ! ia)))
  prefer 2 apply(case-tac ia = lvl V t) apply(case-tac ia < length (levels
(mem-pool-info V (pool b))))
  apply auto[1] apply auto[1] apply auto[1]
  apply(subgoal-tac levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^
(level b - lvl V t) := FREE]] !
    ia = levels (mem-pool-info V2 (pool b)) ! ia)
  prefer 2 apply auto[1]
  apply auto
done

```

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info*:

```

V ∈ mp-free-precond8-3 t b α ∩ {cur = Some t} ⇒
(* {free-stm8-precond2 V t b} ∩ {NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)} ≠ {} ⇒ *)
V2 ∈ free-stm8-precond3 V t b ∩ {i t = 4} ⇒
x = free-stm8-atombody-rest-cond3 (V2(lvl := (lvl V2)(t := lvl V2 t - 1), bn
:= (bn V2)(t := bn V2 t div 4)) t b ⇒
y = x(freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
t, block = bn x t,
    data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))) ⇒
inv-mempool-info y
  apply(rule subst[where s=inv-mempool-info x and t=inv-mempool-info y])
  apply(simp add:block-ptr-def inv-mempool-info-def)

```

```

apply(simp add:Let-def inv-def inv-mempool-info-def)
apply(simp add:set-bit-def)
apply(simp add:gvars-conf-stable-def gvars-conf-def)

apply(subgoal-tac mem-pools V = mem-pools V2)
  prefer 2 apply simp

apply clarify
apply(rule conjI) apply clarify
  apply(rule conjI) apply metis
  apply(rule conjI)
    apply(subgoal-tac length (levels (mem-pool-info V (pool b))) > 0) prefer 2
apply metis apply fast
  apply clarify
    apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b)) ! ia))
      = length (bits (levels (mem-pool-info V2 (pool b)) ! ia)))
    prefer 2 using mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h1
apply blast
    apply(subgoal-tac length (bits (levels (mem-pool-info V2 (pool b))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) !
(lvl V2 t - Suc NULL))
      (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
      [bn V2 t div 4 := FREEING]!)) !
      ia)) = length (bits (levels (mem-pool-info V2 (pool b)) ! ia)))
    prefer 2 using mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h2
apply blast
  apply metis

apply clarify
apply(rule conjI)
  apply metis
  using mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h3 apply
blast
done

```

```

lemma free-stm8-atombody-rest-one-finalstm-VV2-len:
   $\forall p. \text{length (levels (mem-pool-info V2 p))} =$ 
     $\text{length (levels (if } p = \text{pool } b$ 
       $\text{then mem-pool-info V (pool } b)$ 
         $(\text{levels := levels (mem-pool-info V (pool } b))$ 
           $[\text{lvl V } t := (\text{levels (mem-pool-info V (pool } b)) ! \text{lvl V } t)$ 
             $(\text{bits := bits (levels (mem-pool-info V (pool } b)) !$ 
               $\text{lvl V } t)[\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl V } t) := \text{FREE}]!)]$ 
             $\text{else mem-pool-info V } p)) \implies$ 
     $\text{length (levels (mem-pool-info V } p)) = \text{length (levels (mem-pool-info V2 } p))}$ 
by auto

```

**lemma** *free-stm8-atombody-rest-one-finalstm-bits-len:*

$lvl\ V\ t = lvl\ V2\ t \implies$

$p = pool\ b \implies$

$length\ (bits\ (levels\ (if\ p = pool\ b$

$\quad then\ mem\text{-}pool\text{-}info\ V\ (pool\ b)$

$\quad \quad (levels := levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))$

$\quad \quad \quad [lvl\ V\ t := (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)$

$\quad \quad \quad (bits := bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V\ t)$

$\quad \quad \quad [block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V\ t) := FREE]]])$

$\quad else\ mem\text{-}pool\text{-}info\ V\ p)\ !$

$\quad (lvl\ V2\ t)\ ) = length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2$

$t))$

**apply**(*case-tac*  $lvl\ V2\ t < length\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)))$ )

**apply** *auto*

**done**

**lemma** *free-stm8-atombody-rest-one-finalstm-ltlen:*

$lvl\ V2\ t > 0 \implies$

$lvl\ V2\ t = lvl\ V\ t \implies$

$length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$

$\quad = (n\text{-}max\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) * 4\ ^\wedge\ lvl\ V2\ t \implies$

$block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V2\ t) < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t)) \implies$

$block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V2\ t)\ div\ 4 * 4 + 4$

$\leq length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$

**apply**(*rule* *divn-multn-addn-le*[*of*  $4\ length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$ ]

$block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V2\ t)])$

**apply** *simp* **apply** *simp* **apply** *simp*

**done**

**lemma** *free-stm8-atombody-rest-one-finalstm-jj:*

$lvl\ V2\ t > 0 \implies$

$lvl\ V2\ t = lvl\ V\ t \implies$

$length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$

$\quad = (n\text{-}max\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) * 4\ ^\wedge\ lvl\ V2\ t \implies$

$block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V2\ t) < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t)) \implies$

$jj \in \{block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V\ t)\ div\ 4 * 4 ..<$

$block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V\ t)\ div\ 4 * 4 + 4\} \implies$

$jj < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$

**apply** *clarsimp*

**apply**(*subgoal-tac*  $block\ b\ div\ 4\ ^\wedge\ (level\ b - lvl\ V2\ t)\ div\ 4 * 4 + 4$

$\leq length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ !\ lvl\ V2\ t))$ )

**prefer** 2

**apply**(*rule* *free-stm8-atombody-rest-one-finalstm-ltlen*)



**apply** *simp+*  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap'-h1*:  

$$\begin{aligned} & \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t) = \\ & \quad \text{list-updates-n} \\ & \quad (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \\ & \quad \quad [lvl \text{ } V \text{ } t := (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t) \\ & \quad \quad \quad (\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t)[\text{block } b \text{ div } 4 \wedge \\ & \quad \quad \quad (\text{level } b - lvl \text{ } V \text{ } t) := FREE]]]) \text{ !} \\ & \quad \quad lvl \text{ } V \text{ } t)) \\ & \quad (\text{block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V \text{ } t) \text{ div } 4 * 4) (i \text{ } V2 \text{ } t) \text{ NOEXIST} \implies \\ & i \text{ } V2 \text{ } t = 4 \implies \\ & \text{level } b < \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) \implies \\ & lvl \text{ } V \text{ } t = lvl \text{ } V2 \text{ } t \implies \\ & lvl \text{ } V2 \text{ } t \leq \text{level } b \implies \\ & ia = lvl \text{ } V \text{ } t \implies \\ & ia > 0 \implies \\ & p = \text{pool } b \implies \\ & \text{block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V2 \text{ } t) < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } \\ & b)) \text{ ! } lvl \text{ } V2 \text{ } t)) \implies \\ & \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V2 \text{ } t)) \\ & \quad = (n\text{-max } (\text{mem-pool-info } V \text{ (pool } b))) * 4 \wedge lvl \text{ } V2 \text{ } t \implies \\ & \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) = \text{length } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } \\ & b))) \implies \\ & jj \in \{\text{block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V \text{ } t) \text{ div } 4 * 4 .. < \\ & \quad \text{block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V \text{ } t) \text{ div } 4 * 4 + 4\} \implies \\ & \text{get-bit-s } (V2(\text{mem-pool-info } := (\text{mem-pool-info } V2)) \\ & \quad (\text{pool } b := \text{mem-pool-info } V2 \text{ (pool } b)) \\ & \quad \quad (\text{levels } := \text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \\ & \quad \quad \quad [lvl \text{ } V2 \text{ } t - \text{Suc } NULL := (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } (lvl \text{ } V2 \\ & t - \text{Suc } NULL)) \\ & \quad \quad \quad (\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } (lvl \text{ } V2 \text{ } t - \text{Suc } \\ & NULL)) \\ & \quad \quad \quad [bn \text{ } V2 \text{ } t \text{ div } 4 := FREEING]])) p \text{ } ia \text{ } jj = \text{NOEXIST} \\ & \text{apply}(\text{rule subst}[\text{where } s = \text{get-bit-s } V2 \text{ } p \text{ } ia \text{ } jj]) \\ & \text{apply}(\text{subgoal-tac list-updates-n} \\ & \quad (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \\ & \quad \quad [lvl \text{ } V \text{ } t := (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t) \\ & \quad \quad \quad (\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t)[\text{block } b \text{ div } 4 \wedge \\ & \quad \quad \quad (\text{level } b - lvl \text{ } V \text{ } t) := FREE]]]) \text{ !} \\ & \quad \quad lvl \text{ } V \text{ } t)) \\ & \quad (\text{block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V \text{ } t) \text{ div } 4 * 4) (i \text{ } V2 \text{ } t) \text{ NOEXIST} ! \\ & jj = \\ & \text{NOEXIST}) \text{ prefer } 2 \\ & \text{apply}(\text{rule list-updates-n-eq}[\text{of block } b \text{ div } 4 \wedge (\text{level } b - lvl \text{ } V \text{ } t) \text{ div } 4 * 4 \text{ } jj \\ & \quad \text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \\ & \quad \quad [lvl \text{ } V \text{ } t := (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t) \\ & \quad \quad \quad (\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } lvl \text{ } V \text{ } t)[\text{block } b \text{ div } \end{aligned}$$

```

4 ^ (level b - lvl V t) := FREE]] !
    lvl V t) i V2 t NOEXIST])
apply fastforce
apply(rule subst[where s=length (bits (levels (mem-pool-info V (pool b)) ! lvl
V t))
    and t=length (bits (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    [bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)
    [block b div 4 ^ (level b - lvl V t) := FREE]] !
    lvl V t))]) apply force
apply(rule free-stm8-atombody-rest-one-finalstm-jj)
apply fast apply fast apply presburger apply argo
apply fast apply force

```

```

apply argo
by fastforce

```

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap'-h2:

get-bit-s

```

(V2 (mem-pool-info := (mem-pool-info V2)
(pool b := mem-pool-info V2 (pool b)
[levels := levels (mem-pool-info V2 (pool b))
[lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl
V2 t - Suc NULL))
[bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
[bn V2 t div 4 := FREEING]])))))
p ia jj =
FREE ∨
get-bit-s
(V2 (mem-pool-info := (mem-pool-info V2)
(pool b := mem-pool-info V2 (pool b)
[levels := levels (mem-pool-info V2 (pool b))
[lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl
V2 t - Suc NULL))
[bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
[bn V2 t div 4 := FREEING]])))))
p ia jj =
FREEING ∨
get-bit-s
(V2 (mem-pool-info := (mem-pool-info V2)
(pool b := mem-pool-info V2 (pool b)
[levels := levels (mem-pool-info V2 (pool b))
[lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl
V2 t - Suc NULL))
[bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))

```

$$[bn \ V2 \ t \ \text{div} \ 4 := FREEING] \rangle \rangle \rangle \rangle$$

$$p \ \text{ia} \ jj =$$

$$ALLOCATED \vee$$

$$\text{get-bit-s}$$

$$(V2 \langle mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info \ V2)$$

$$(pool \ b := mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$\langle levels := levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$[lvl \ V2 \ t - Suc \ NULL := (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$\langle bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$[bn \ V2 \ t \ \text{div} \ 4 := FREEING] \rangle \rangle \rangle \rangle$$

$$p \ \text{ia} \ jj =$$

$$ALLOCATING \implies$$

$$\text{get-bit-s}$$

$$(V2 \langle mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info \ V2)$$

$$(pool \ b := mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$\langle levels := levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$[lvl \ V2 \ t - Suc \ NULL := (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$\langle bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$[bn \ V2 \ t \ \text{div} \ 4 := FREEING] \rangle \rangle \rangle \rangle$$

$$p \ \text{ia} \ jj =$$

$$NOEXIST \implies$$

$$\text{get-bit-s}$$

$$(V2 \langle mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info \ V2)$$

$$(pool \ b := mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$\langle levels := levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$$

$$[lvl \ V2 \ t - Suc \ NULL := (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$\langle bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc \ NULL))$$

$$[bn \ V2 \ t \ \text{div} \ 4 := FREEING] \rangle \rangle \rangle \rangle$$

$$p \ (\text{ia} - 1) \ (jj \ \text{div} \ 4) =$$

$$DIVIDED$$

**by force**

**axiomatization where** *mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap*:

$$V \in mp\text{-}free\text{-}precond8\text{-}3 \ t \ b \ \alpha \cap \langle \langle 'cur = Some \ t \rangle \rangle \implies$$

$$(* \{ free\text{-}stm8\text{-}precond2 \ V \ t \ b \} \cap \langle \langle NULL < 'lvl \ t \wedge partner\text{-}bits \ ('mem\text{-}pool\text{-}info \ (pool \ b)) \ ('lvl \ t) \ ('bn \ t) \rangle \rangle \neq \{ \} \implies *)$$

$$V2 \in free\text{-}stm8\text{-}precond3 \ V \ t \ b \cap \langle \langle 'i \ t = 4 \rangle \rangle \implies$$

$$x = free\text{-}stm8\text{-}atombody\text{-}rest\text{-}cond3 \ (V2 \langle lvl := (lvl \ V2)(t := lvl \ V2 \ t - 1), \ bn := (bn \ V2)(t := bn \ V2 \ t \ \text{div} \ 4) \rangle) \ t \ b \implies$$

$$y = x \langle freeing\text{-}node := (freeing\text{-}node \ x) \ (t := Some \ \langle pool = pool \ b, \ level = lvl \ x \ t, \ block = bn \ x \ t, \ data = block\text{-}ptr \ (mem\text{-}pool\text{-}info \ x \ (pool \ b)) \ (ALIGN4 \ (max\text{-}sz$$

$(\text{mem-pool-info } x \text{ (pool } b))) \text{ div } 4 \wedge \text{lvl } x \text{ t) (bn } x \text{ t)}) \implies$   
 $\text{inv-bitmap } y$

**axiomatization where** *mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap-freelist*:

$V \in \text{mp-free-precond8-3 } t \ b \ \alpha \cap \{\text{'cur} = \text{Some } t\} \implies$   
 $(\{ \text{free-stm8-precond2 } V \ t \ b \} \cap \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits ( 'mem-pool-info$   
 $(\text{pool } b)) \text{ ('lvl } t) \text{ ('bn } t)\} \neq \{\} \implies *)$   
 $V2 \in \text{free-stm8-precond3 } V \ t \ b \cap \{\text{'i } t = 4\} \implies$   
 $x = \text{free-stm8-atombody-rest-cond3 } (V2 \text{ (lvl := (lvl } V2) (t := \text{lvl } V2 \ t - 1), \text{ bn$   
 $:= (\text{bn } V2) (t := \text{bn } V2 \ t \text{ div } 4))) \ t \ b \implies$   
 $y = x \text{ (freeing-node := (freeing-node } x) \text{ (t := Some (pool = pool } b, \text{ level = lvl } x$   
 $t, \text{ block = bn } x \ t,$   
 $\text{data = block-ptr (mem-pool-info } x \text{ (pool } b)) \text{ (ALIGN4 (max-sz$   
 $(\text{mem-pool-info } x \text{ (pool } b))) \text{ div } 4 \wedge \text{lvl } x \text{ t) (bn } x \text{ t)}) \implies$   
 $\text{inv-bitmap-freelist } y$

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-len-bits1*:

$\forall j. j \neq \text{lvl } V \ t \longrightarrow \text{levels (mem-pool-info } V \text{ (pool } b)) ! j = \text{levels (mem-pool-info$   
 $V2 \text{ (pool } b)) ! j \implies$   
 $\text{bits (levels (mem-pool-info } V2 \text{ (pool } b)) ! \text{lvl } V \ t) =$   
 $\text{list-updates-n}$   
 $(\text{bits (levels (mem-pool-info } V \text{ (pool } b))$   
 $\text{[lvl } V \ t := (\text{levels (mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V \ t)$   
 $\text{[bits := bits (levels (mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V \ t)] \text{block } b \text{ div } 4 \wedge$   
 $(\text{level } b - \text{lvl } V \ t) := \text{FREE}]) !$   
 $\text{lvl } V \ t))$   
 $(\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) \text{ div } 4 * 4) (i \ V2 \ t) \text{ NOEXIST} \implies$   
 $(i \ V2 \ t) = 4 \implies$   
 $\text{length (bits (levels (mem-pool-info } V \text{ (pool } b)) ! (\text{lvl } V2 \ t - \text{Suc NULL}))}$   
 $= \text{length (bits (levels (mem-pool-info } V2 \text{ (pool } b)) ! (\text{lvl } V2 \ t - \text{Suc NULL}))}$   
**apply**(rule subst[**where**  $s = \text{length (bits (levels (mem-pool-info } V2 \text{ (pool } b))$   
 $\text{[lvl } V2 \ t - \text{Suc NULL} := (\text{levels (mem-pool-info } V2 \text{ (pool } b)) ! (\text{lvl}$   
 $V2 \ t - \text{Suc NULL})}$   
 $\text{[bits := bits (levels (mem-pool-info } V2 \text{ (pool } b)) ! (\text{lvl } V2 \ t - \text{Suc}$   
 $\text{NULL})] \text{bn } V2 \ t \text{ div } 4 := \text{FREEING}]) !$   
 $(\text{lvl } V2 \ t - \text{Suc NULL}))$  **and**  $t = \text{length (bits (levels (mem-pool-info$   
 $V \text{ (pool } b)) ! (\text{lvl } V2 \ t - \text{Suc NULL}))}$ )]  
**using** *mempool-free-stm8-atombody-rest-one-finalstm-h1-1'*[of  $V \ t \ b \ V2 \text{ lvl } V2 \ t -$   
 $\text{Suc NULL}$ ] **apply** auto[1]  
**apply**(case-tac  $\text{lvl } V2 \ t - \text{Suc NULL} < \text{length (levels (mem-pool-info } V2 \text{ (pool}$   
 $b)))$ )  
**apply** auto  
**done**

**lemma** *lm11*:

$\text{lvl } V2 \ t \leq \text{level } b \wedge \text{level } b > 0 \wedge \text{level } b < \text{length (levels (mem-pool-info } V \text{ (pool$

$b))) \wedge$   
 $0 < lvl\ V2\ t \implies$   
 $block\ b < n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b \implies$   
 $block\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t) = bn\ V2\ t \implies$   
 $bn\ V2\ t\ div\ 4 < n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge (lvl\ V2\ t - Suc\ NULL)$   
**apply**(rule subst[**where**  $s=block\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t)\ div\ 4$  **and**  $t=bn\ V2\ t\ div\ 4$ ])  
**apply simp**  
  
**apply**(rule subst[**where**  $s=n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ div\ 4$   
 $\wedge (level\ b - lvl\ V2\ t)\ div\ 4$   
**and**  $t=n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge (lvl\ V2\ t - Suc\ NULL)$ ])  
**apply** (smt Groups.mult-ac(2) Groups.mult-ac(3) One-nat-def add-diff-cancel-left'  
div-mult-self1-is-m  
le-Suc-ex power-add power-minus-mult zero-less-numeral zero-less-power)  
  
**apply**(subgoal-tac  $n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ mod\ 4 \wedge (level\ b - lvl\ V2\ t) = 0$ )  
**prefer 2 using** mod-minus-0[of lvl V2 t level b  $n-max\ (mem-pool-info\ V\ (pool\ b))$ ]  
**apply fast**  
  
**apply**(subgoal-tac  $block\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t) < n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t)$ )  
**prefer 2 using** mod-div-gt[of block b  $n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ 4 \wedge (level\ b - lvl\ V2\ t)$ ]  
**apply fast**  
**apply**(subgoal-tac  $n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t)\ mod\ 4 = 0$ )  
**prefer 2 using** mod-minus-div-4[of lvl V2 t level b  $n-max\ (mem-pool-info\ V\ (pool\ b))$ ]  
**apply fast**  
  
**using** mod-div-gt[of block b  $div\ 4 \wedge (level\ b - lvl\ V2\ t)$   
 $n-max\ (mem-pool-info\ V\ (pool\ b)) * 4 \wedge level\ b\ div\ 4 \wedge (level\ b - lvl\ V2\ t)\ 4$ ]  
**apply fast**  
**done**

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-aux-vars-h2:  
 $pool\ b \in mem-pools\ V \implies$   
 $inv-mempool-info\ V \implies$   
 $level\ b < length\ (levels\ (mem-pool-info\ V\ (pool\ b))) \implies$   
 $block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t) = bn\ V2\ t \implies$   
 $block\ b < length\ (bits\ (levels\ (mem-pool-info\ V\ (pool\ b))\ !\ level\ b)) \implies$   
 $lvl\ V\ t = lvl\ V2\ t \implies$   
 $lvl\ V2\ t \leq level\ b \implies$   
 $0 < lvl\ V2\ t \implies$   
 $length\ (levels\ (mem-pool-info\ V\ (pool\ b))) = length\ (levels\ (mem-pool-info\ V2\ (pool\ b))) \implies$

```

length (bits (levels (mem-pool-info V (pool b)) ! (lvl V2 t - Suc NULL)))
= length (bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc NULL)))
 $\implies$ 
bits (levels (mem-pool-info V2 (pool b))
  [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) !
    (lvl V2 t - Suc NULL))
  (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
    Suc NULL)) [bn V2 t div 4 := FREEING] !
    (lvl V2 t - Suc NULL)) !
  (bn V2 t div 4) =
  FREEING
apply (subgoal-tac lvl V2 t - Suc 0 < length (levels (mem-pool-info V2 (pool b))))
prefer 2 apply auto[1]
apply (subgoal-tac bn V2 t div 4 < length (bits (levels (mem-pool-info V2 (pool b))
! (lvl V2 t - Suc NULL))))
prefer 2

  apply (subgoal-tac level b > 0) prefer 2 apply auto[1]
  apply (subgoal-tac n-max (mem-pool-info V (pool b)) * 4 ^ (lvl V2 t - Suc
    NULL)
    = length (bits (levels (mem-pool-info V (pool b)) ! (lvl V2 t -
    Suc NULL))))
  prefer 2 apply (simp add: inv-mempool-info-def Let-def)
  apply (metis inv-mempool-info-def lm11)
apply auto
done

```

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-len-lvl*:

```

(V2, V (mem-pool-info := (mem-pool-info V)
  (pool b := mem-pool-info V (pool b))
  (levels := levels (mem-pool-info V (pool b))
    [lvl V2 t := (levels (mem-pool-info V (pool b)) ! lvl V2 t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V2 t) [bn V2
    t := FREE] ! ! !)),
  freeing-node := (freeing-node V)(t := None) !))
 $\in$  gvars-conf-stable  $\implies$ 
length (levels (mem-pool-info V2 (pool b))) = length (levels (mem-pool-info V
(pool b)))
apply (simp add: gvars-conf-stable-def gvars-conf-def)
done

```

**axiomatization where** *mempool-free-stm8-atombody-rest-one-finalstm-inv-aux-vars*:

```

V  $\in$  mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = Some t \rrbracket \implies$ 
 $\{ free-stm8-precond2 V t b \} \cap \llbracket NULL < 'lvl t \wedge partner-bits ('mem-pool-info$ 
(pool b)) ('lvl t) ('bn t) \rrbracket  $\neq \{ \} \implies$ 
V2  $\in$  free-stm8-precond3 V t b  $\cap \llbracket 'i t = 4 \rrbracket \implies$ 
x = free-stm8-atombody-rest-cond3 (V2 (lvl := (lvl V2)(t := lvl V2 t - 1), bn

```

$$:= (bn\ V2)(t := bn\ V2\ t\ div\ 4))\ t\ b \implies$$

$$y = x(\text{freeing-node} := (\text{freeing-node}\ x)\ (t := \text{Some}\ (\text{pool} = \text{pool}\ b,\ \text{level} = \text{lvl}\ x\ t,\ \text{block} = bn\ x\ t),$$

$$\text{data} = \text{block-ptr}\ (\text{mem-pool-info}\ x\ (\text{pool}\ b))\ (\text{ALIGN}_4\ (\text{max-sz}\ (\text{mem-pool-info}\ x\ (\text{pool}\ b)))\ \text{div}\ 4\ \wedge\ \text{lvl}\ x\ t)\ (bn\ x\ t))) \implies$$

$$\text{inv-aux-vars}\ y$$

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case1-h1:*

$$NULL < \text{length}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))) \implies$$

$$\text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ NULL)[ia := \text{FREEING}] =$$

$$\text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))$$

$$[NULL := (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ NULL)$$

$$(\text{bits} := \text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ NULL)[ia := \text{FREEING}])])$$

$$!$$

$$NULL)$$

**by** *auto*

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case1:*

$$\text{pool}\ b \in \text{mem-pools}\ V2 \implies$$

$$\text{inv}\ V \implies$$

$$NULL < \text{lvl}\ V2\ t \implies$$

$$\text{pool}\ b \in \text{mem-pools}\ V \implies$$

$$*(V2,\ V(\text{mem-pool-info} := \text{set-bit-free}\ (\text{mem-pool-info}\ V)\ (\text{pool}\ b)\ (\text{lvl}\ V2\ t)\ (bn\ V2\ t)),$$

$$\text{freeing-node} := (\text{freeing-node}\ V)(t := \text{None}))$$

$$\in \text{gvars-conf-stable} \implies$$

$$\forall p. p \neq \text{pool}\ b \longrightarrow \text{mem-pool-info}\ V2\ p = \text{set-bit-free}\ (\text{mem-pool-info}\ V)\ (\text{pool}\ b)\ (\text{lvl}\ V2\ t)\ (bn\ V2\ t)\ p \implies *$$

$$\forall j. j \neq \text{lvl}\ V2\ t \longrightarrow$$

$$\text{levels}\ (\text{set-bit-free}\ (\text{mem-pool-info}\ V)\ (\text{pool}\ b)\ (\text{lvl}\ V2\ t)\ (bn\ V2\ t)\ (\text{pool}\ b))\ !$$

$$j =$$

$$\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ j \implies$$

$$\text{lvl}\ V2\ t \leq \text{level}\ b \implies$$

$$\text{lvl}\ V\ t = \text{lvl}\ V2\ t \implies$$

$$\forall i < \text{length}\ (\text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V\ (\text{pool}\ b))\ !\ NULL)).\ \text{get-bit-s}\ V\ (\text{pool}\ b)\ NULL\ i \neq \text{NOEXIST} \implies$$

$$ia < \text{length}\ (\text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))$$

$$[\text{lvl}\ V2\ t - \text{Suc}\ NULL := (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !$$

$$(\text{lvl}\ V2\ t - \text{Suc}\ NULL))$$

$$(\text{bits} := \text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !$$

$$(\text{lvl}\ V2\ t - \text{Suc}\ NULL))[bn\ V2\ t\ div\ 4 := \text{FREEING}])])\ !$$

$$NULL)) \implies$$

$$\text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))$$

$$[\text{lvl}\ V2\ t - \text{Suc}\ NULL := (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ (\text{lvl}\ V2\ t - \text{Suc}\ NULL))$$

$$(\text{bits} := \text{bits}\ (\text{levels}\ (\text{mem-pool-info}\ V2\ (\text{pool}\ b))\ !\ (\text{lvl}\ V2\ t - \text{Suc}\ NULL)))[bn\ V2\ t\ div\ 4 := \text{FREEING}])])\ !$$

$$NULL)\ !$$

```

    ia =
    NOEXIST  $\implies$ 
    False
  apply(simp add:set-bit-def)
  apply(subgoal-tac levels (mem-pool-info V (pool b)) ! (lvl V2 t - 1)
    = levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - 1)) prefer 2
    apply(subgoal-tac levels (mem-pool-info V (pool b))
      [lvl V2 t := (levels (mem-pool-info V (pool b)) ! lvl V2 t)
        (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V2 t)[bn V2 t :=
FREE])] ! (lvl V2 t - 1)
      = levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - 1)) prefer
2 apply auto[1]
  apply auto[1]
  apply auto[1]

  apply(case-tac lvl V2 t - Suc 0 = 0)
    apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b)) ! 0))
      = length (bits (levels (mem-pool-info V2 (pool b)) ! 0))) prefer 2
    apply auto[1]
    apply(subgoal-tac length (levels (mem-pool-info V2 (pool b))) > 0) prefer 2 ap-
ply auto[1]
    apply(subgoal-tac ia < length (bits (levels (mem-pool-info V (pool b)) ! 0)))
prefer 2
    apply(rule subst[where s=length (bits (levels (mem-pool-info V2 (pool b)) !
NULL)) and
      t=length (bits (levels (mem-pool-info V (pool b)) !
NULL))])
      apply auto[1]
      apply(rule subst[where t=length (bits (levels (mem-pool-info V2 (pool b)) !
NULL)) and
        s=length (bits (levels (mem-pool-info V2 (pool b))
[lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) !
(lvl V2 t - Suc NULL))
(bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))[bn V2 t div 4 := FREEING])] !
        NULL)) ])
      apply auto[1]
      apply auto[1]

  apply(case-tac ia = bn V2 t div 4)
    apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b))
[lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t
- Suc NULL))
(bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL))[bn V2 t div 4 := FREEING])] !
      NULL) ! ia = FREEING) prefer 2
    apply(rule subst[where s=0 and t=lvl V2 t - Suc 0]) apply metis
    apply(rule subst[where s=ia and t=bn V2 t div 4]) apply metis
    apply(rule subst[where s=bits (levels (mem-pool-info V2 (pool b)) ! NULL)[ia

```



```

:= FREEING] and
    t=bits (levels (mem-pool-info V2 (pool b))
    [NULL := (levels (mem-pool-info V2 (pool b)) ! NULL) (bits := bits (levels
(mem-pool-info V2 (pool b)) ! NULL)
    [ia := FREEING]] ! NULL)) apply auto[1] apply auto[1]
    apply (metis BlockState.distinct(25))

    apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t
- Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL)) [bn V2 t div 4 := FREEING]] !
    NULL) ! ia ≠ NOEXIST) prefer 2
    apply(rule subst[where s=0 and t=lvl V2 t - Suc 0]) apply metis
    apply(rule subst[where s=bits (levels (mem-pool-info V2 (pool b)) ! NULL) [bn
V2 t div 4 := FREEING] and
    t=bits (levels (mem-pool-info V2 (pool b))
    [NULL := (levels (mem-pool-info V2 (pool b)) ! NULL) (bits := bits (levels
(mem-pool-info V2 (pool b)) ! NULL)
    [bn V2 t div 4 := FREEING]] ! NULL)) apply auto[1] apply
auto[1]
    apply fast

    apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t
- Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL)) [bn V2 t div 4 := FREEING]] !
    NULL) ! ia ≠ NOEXIST) prefer 2
    apply(rule subst[where s=levels (mem-pool-info V2 (pool b)) ! NULL and
t=levels (mem-pool-info V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t
- Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL)) [bn V2 t div 4 := FREEING]] !
    NULL]) apply simp
    apply(subgoal-tac length (bits (levels (mem-pool-info V2 (pool b)) ! NULL)) =
length (bits (levels (mem-pool-info V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) !
(lvl V2 t - Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL)) [bn V2 t div 4 := FREEING]] !
    NULL)) ) prefer 2 apply simp
    apply presburger
    apply fast
    done

```

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case2*:  
 $p \in \text{mem-pools } V2 \implies$   
 $\text{inv } V \implies$   
 $\text{NULL} < \text{lvl } V2 \ t \implies$   
 $\text{pool } b \in \text{mem-pools } V \implies$   
 $(V2, V \llbracket \text{mem-pool-info} := \text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V2 \ t) (\text{bn } V2 \ t),$   
 $\text{freeing-node} := (\text{freeing-node } V)(t := \text{None}) \rrbracket)$   
 $\in \text{gvars-conf-stable} \implies$   
 $\forall p. p \neq \text{pool } b \longrightarrow \text{mem-pool-info } V2 \ p = \text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V2 \ t) (\text{bn } V2 \ t) \ p \implies$   
 $\text{level } b < \text{length } (\text{lsizes } V2 \ t) \implies$   
 $p \neq \text{pool } b \implies$   
 $\text{ia} < \text{length } (\text{bits } (\text{levels } (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V2 \ t) (\text{bn } V2 \ t) \ p) ! \text{NULL})) \implies$   
 $\text{get-bit } (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V2 \ t) (\text{bn } V2 \ t)) \ p \ \text{NULL} \ \text{ia} = \text{NOEXIST} \implies \text{False}$   
**apply**(*simp add:set-bit-def*)  
**apply**(*subgoal-tac*  $\forall i < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ p) ! 0)). (\text{bits } (\text{levels } (\text{mem-pool-info } V \ p) ! 0)) ! i \neq \text{NOEXIST}$ )  
**prefer 2 apply**(*subgoal-tac mem-pools V2 = mem-pools V*) **prefer 2 apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)  
**apply**(*simp add:inv-def inv-bitmap0-def Let-def*)  
**apply auto**  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0*:  
 $V \in \text{mp-free-precond8-3 } t \ b \ \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \implies$   
 $\{ \text{free-stm8-precond2 } V \ t \ b \} \cap \llbracket \text{NULL} < 'lvl \ t \wedge \text{partner-bits } ('mem\text{-pool-info } (\text{pool } b)) ('lvl \ t) ('bn \ t) \rrbracket \neq \{ \} \implies$   
 $V2 \in \text{free-stm8-precond3 } V \ t \ b \cap \llbracket 'i \ t = 4 \rrbracket \implies$   
 $x = \text{free-stm8-atombody-rest-cond3 } (V2 \llbracket \text{lvl} := (\text{lvl } V2)(t := \text{lvl } V2 \ t - 1), \text{bn} := (\text{bn } V2)(t := \text{bn } V2 \ t \text{ div } 4) \rrbracket) \ t \ b \implies$   
 $y = x \llbracket \text{freeing-node} := (\text{freeing-node } x) (t := \text{Some } \llbracket \text{pool} = \text{pool } b, \text{level} = \text{lvl } x \ t, \text{block} = \text{bn } x \ t,$   
 $\text{data} = \text{block-ptr } (\text{mem-pool-info } x (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } x (\text{pool } b))) \text{ div } 4 \wedge \text{lvl } x \ t) (\text{bn } x \ t) \rrbracket \rrbracket \implies$   
 $\text{inv-bitmap0 } y$   
**apply**(*simp add:inv-bitmap0-def Let-def*)  
**apply clarify**

**apply**(*rule conjI*)  
**apply clarsimp**  
**apply**(*subgoal-tac*  $\forall i < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! 0)). (\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! 0)) ! i \neq \text{NOEXIST}$ )  
**prefer 2 apply**(*simp add:inv-def inv-bitmap0-def Let-def*)  
**apply**(*subgoal-tac levels*  $(\text{mem-pool-info } V (\text{pool } b)) ! (\text{lvl } V2 \ t - 1)$   
 $= \text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 \ t - 1))$  **prefer 2**  
**apply**(*subgoal-tac levels*  $(\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V2 \ t) (\text{bn } V2 \ t))$

```

V2 t) (pool b)) ! (lvl V2 t - 1)
      = levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - 1)) prefer
2 apply auto[1]
  apply(simp add:set-bit-def)

using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case1 apply blast

apply clarsimp
using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case2 apply blast
done

term mp-free-precond8-3 t b  $\alpha$ 
term free-stm8-precond2 V t b
term free-stm8-precond3 V t b

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0-case1:
pool b  $\in$  mem-pools V2  $\implies$ 
  inv V  $\implies$ 
  NULL < lvl V2 t  $\implies$ 
  pool b  $\in$  mem-pools V  $\implies$ 
  level b < length (levels (mem-pool-info V (pool b)))  $\implies$ 
  (V2, V( $\lfloor$ mem-pool-info := set-bit-free (mem-pool-info V) (pool b) (lvl V2 t) (bn
  V2 t),
    freeing-node := (freeing-node V)(t := None) $\rfloor$ )
   $\in$  gvars-conf-stable  $\implies$ 
   $\forall j. j \neq \text{lvl V2 t} \longrightarrow$ 
    levels (set-bit-free (mem-pool-info V) (pool b) (lvl V2 t) (bn V2 t) (pool b)) !
  j =
    levels (mem-pool-info V2 (pool b)) ! j  $\implies$ 
    bits (levels (mem-pool-info V2 (pool b)) ! lvl V2 t) =
      list-updates-n (bits (levels (set-bit-free (mem-pool-info V) (pool b) (lvl V2
  t) (bn V2 t) (pool b)) ! lvl V2 t))
        (bn V2 t div 4 * 4) 4 NOEXIST  $\implies$ 
    lvl V2 t  $\leq$  level b  $\implies$ 
    ia < length (bits (levels (mem-pool-info V2 (pool b))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) !
  (lvl V2 t - Suc NULL))
        ( $\lfloor$ bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
  Suc NULL) $\rfloor$ )
      [bn V2 t div 4 := FREEING] $\rfloor$ ) !
      (length (levels (mem-pool-info V2 (pool b)) - Suc NULL)))  $\implies$ 
    bits (levels (mem-pool-info V2 (pool b))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
  Suc NULL))
        ( $\lfloor$ bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
  NULL) $\rfloor$ )[bn V2 t div 4 := FREEING] $\rfloor$ ) !
      (length (levels (mem-pool-info V2 (pool b)) - Suc NULL)) ! ia = DIVIDED
   $\implies$ 

```

```

False
apply(simp add:set-bit-def)
apply(subgoal-tac length (levels (mem-pool-info V2 (pool b)))
      = length (levels (mem-pool-info V (pool b)))) prefer 2
using mempool-free-stm8-atombody-rest-one-finalstm-len-lvls apply blast

apply(subgoal-tac let bitsn = bits ((levels (mem-pool-info V (pool b)) ! (length
(levels (mem-pool-info V (pool b))) - 1)))
      in  $\forall i < \text{length } \text{bitsn}. \text{bitsn } ! i \neq \text{DIVIDED}$ ) prefer 2
apply(simp add:inv-def inv-bitmapn-def)

apply(case-tac lvl V2 t = length (levels (mem-pool-info V2 (pool b))) - Suc 0)
apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b)))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
      ( $\text{bits} := \text{bits} (\text{levels} (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 t - \text{Suc } \text{NULL}))$ )[bn V2 t div 4 := FREEING]]) !
      (length (levels (mem-pool-info V2 (pool b))) - Suc NULL) ! ia  $\neq \text{DIVIDED}$ )

apply auto[1]
apply(rule subst[where s=bits (levels (mem-pool-info V2 (pool b)) ! (length
(levels (mem-pool-info V2 (pool b))) - Suc NULL))
      and t=bits (levels (mem-pool-info V2 (pool b)))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool
b)) ! (lvl V2 t - Suc NULL))
      ( $\text{bits} := \text{bits} (\text{levels} (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 t - \text{Suc } \text{NULL}))$ )[bn V2 t div 4 := FREEING]]) !
      (length (levels (mem-pool-info V2 (pool b))) - Suc NULL)])

apply auto[1]
apply(unfold Let-def)[1]
apply(subgoal-tac  $\forall i < \text{length } (\text{bits} (\text{levels} (\text{mem-pool-info } V2 (\text{pool } b)) ! \text{lvl } V2 t)). \text{get-bit-s } V2 (\text{pool } b) (\text{lvl } V2 t) i \neq \text{DIVIDED}$ )
apply auto[1]

apply(rule list-neq-udpt-neq[of bits (levels (mem-pool-info V (pool b)) ! lvl V2 t) DIVIDED
      bits (levels (mem-pool-info V2 (pool b)) ! lvl V2 t) (bn V2 t div 4 * 4) 4 NOEXIST])
apply auto[1]
using lst-udptn-set-eq[of 4 bits (levels (mem-pool-info V (pool b)) ! lvl V2 t)
bn V2 t FREE NOEXIST]
apply auto[1]
apply blast

apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b)))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
      ( $\text{bits} := \text{bits} (\text{levels} (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 t - \text{Suc } \text{NULL}))$ )]

```

$NULL)) [bn \ V2 \ t \ \text{div } 4 := FREEING]] !$   
 $(length \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))) - Suc \ NULL)) ! ia \neq DIVIDED)$

**apply fast**  
**apply**(rule subst[**where**  $s = levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) !$   
 $(length \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))) - Suc \ NULL)$  **and**  $t = levels$   
 $(mem\text{-}pool\text{-}info \ V2 \ (pool \ b))$   
 $[lvl \ V2 \ t - Suc \ NULL := (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t$   
 $- Suc \ NULL))$   
 $(bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (lvl \ V2 \ t - Suc$   
 $NULL)) [bn \ V2 \ t \ \text{div } 4 := FREEING]] !$   
 $(length \ (levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b))) - Suc \ NULL)])$   
**apply auto[1]**  
**apply**(unfold Let-def)[1]  
**apply**(subgoal-tac  $levels \ (mem\text{-}pool\text{-}info \ V \ (pool \ b)) ! (length \ (levels \ (mem\text{-}pool\text{-}info$   
 $V \ (pool \ b))) - 1) =$   
 $levels \ (mem\text{-}pool\text{-}info \ V2 \ (pool \ b)) ! (length \ (levels \ (mem\text{-}pool\text{-}info$   
 $V \ (pool \ b))) - 1)$ )  
**prefer 2 apply** (metis One-nat-def)  
**by** (metis One-nat-def Suc-diff-1 inv-mempool-info-def invariant.inv-def not-less  
nth-list-update-neq)

**lemma mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl-n-case2:**  
 $p \in mem\text{-}pools \ V2 \implies$   
 $inv \ V \implies$   
 $NULL < lvl \ V2 \ t \implies$   
 $level \ b < length \ (levels \ (mem\text{-}pool\text{-}info \ V \ (pool \ b))) \implies$   
 $(V2, V (mem\text{-}pool\text{-}info := set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool \ b) \ (lvl \ V2 \ t) \ (bn$   
 $V2 \ t),$   
 $freeing\text{-}node := (freeing\text{-}node \ V)(t := None)))$   
 $\in gvars\text{-}conf\text{-}stable \implies$   
 $\forall p. p \neq pool \ b \longrightarrow mem\text{-}pool\text{-}info \ V2 \ p = set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool$   
 $b) \ (lvl \ V2 \ t) \ (bn \ V2 \ t) \ p \implies$   
 $lvl \ V2 \ t \leq level \ b \implies$   
 $p \neq pool \ b \implies$   
 $ia < length \ (bits \ (levels \ (set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool \ b) \ (lvl \ V2 \ t) \ (bn$   
 $V2 \ t) \ p) !$   
 $(length \ (levels \ (set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool \ b) \ (lvl \ V2 \ t)$   
 $(bn \ V2 \ t) \ p)) - Suc \ NULL))) \implies$   
 $get\text{-}bit \ (set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool \ b) \ (lvl \ V2 \ t) \ (bn \ V2 \ t)) \ p$   
 $(length \ (levels \ (set\text{-}bit\text{-}free \ (mem\text{-}pool\text{-}info \ V) \ (pool \ b) \ (lvl \ V2 \ t) \ (bn \ V2 \ t) \ p))$   
 $- Suc \ NULL) \ ia =$   
 $DIVIDED \implies$   
 $False$   
**apply**(simp add:set-bit-def)  
**apply**(subgoal-tac  $\forall i < length \ (bits \ ((levels \ (mem\text{-}pool\text{-}info \ V \ p) ! (length \ (levels$   
 $(mem\text{-}pool\text{-}info \ V \ p)) - 1))))$ .  
 $bits \ ((levels \ (mem\text{-}pool\text{-}info \ V \ p) ! (length \ (levels \ (mem\text{-}pool\text{-}info$   
 $V \ p)) - 1))) ! i \neq DIVIDED)$

**prefer 2 apply**(subgoal-tac mem-pools V2 = mem-pools V) **prefer 2 apply**(simp add:gvars-conf-stable-def gvars-conf-def)  
**apply**(simp add:inv-def inv-bitmapn-def Let-def)  
**apply auto**  
**done**

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl:

$V \in mp\text{-free-precond8-3 } t \ b \ \alpha \cap \{\cur = \text{Some } t\} \implies$   
 $\{\text{free-stm8-precond2 } V \ t \ b\} \cap \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info } (\text{pool } b)) (\text{'lvl } t) (\text{'bn } t))\} \neq \{\} \implies$   
 $V2 \in \text{free-stm8-precond3 } V \ t \ b \cap \{\text{'i } t = 4\} \implies$   
 $x = \text{free-stm8-atombody-rest-cond3 } (V2 (\text{lvl} := (\text{lvl } V2)(t := \text{lvl } V2 \ t - 1), \text{bn} := (\text{bn } V2)(t := \text{bn } V2 \ t \ \text{div } 4))) \ t \ b \implies$   
 $y = x (\text{freeing-node} := (\text{freeing-node } x) (t := \text{Some } (\text{pool} = \text{pool } b, \text{level} = \text{lvl } x \ t, \text{block} = \text{bn } x \ t),$   
 $\text{data} = \text{block-ptr } (\text{mem-pool-info } x (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } x (\text{pool } b))) \ \text{div } 4 \wedge \text{lvl } x \ t) (\text{bn } x \ t))) \implies$   
 $\text{inv-bitmapn } y$   
**apply**(simp add:inv-bitmapn-def Let-def)  
**apply clarify**

**apply**(rule conjI)  
**apply clarsimp**  
**using** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl-case1 **apply blast**  
  
**apply clarsimp**  
**using** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl-case2 **apply blast**

**done**

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl-not4free-case1-h1:

$\text{lvl } V2 \ t - \text{Suc NULL} = \text{ia} \implies$   
 $\text{lvl } V2 \ t - \text{Suc NULL} < \text{length } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b))) \implies$   
 $\text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc NULL}))[\text{bn } V2 \ t \ \text{div } 4 := \text{FREEING}] =$   
 $\text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b))$   
 $\quad (\text{levels} := \text{levels } (\text{mem-pool-info } V2 (\text{pool } b))$   
 $\quad [\text{lvl } V2 \ t - \text{Suc NULL} := (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) !$   
 $(\text{lvl } V2 \ t - \text{Suc NULL}))$   
 $\quad (\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! (\text{lvl } V2 \ t -$   
 $\text{Suc NULL}))[\text{bn } V2 \ t \ \text{div } 4 := \text{FREEING}])) !$   
 $\text{ia})$

**by** simp

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl-not4free-case1-h2:

$\text{ia} < \text{length } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b))) \implies$   
 $\text{jj} < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b))))$

$$[lvl\ V2\ t - Suc\ NULL := (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !$$
  

$$(lvl\ V2\ t - Suc\ NULL))$$
  

$$(\{bits := bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t -$$
  

$$Suc\ NULL)) [bn\ V2\ t\ div\ 4 := FREEING]\}) !$$
  

$$ia)) \implies$$
  

$$NULL < ia \implies$$
  

$$length\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))) = length\ (levels\ (mem\ pool\ info\ V$$
  

$$(pool\ b))) \implies$$
  

$$length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ ia)) = length\ (bits\ (levels$$
  

$$(mem\ pool\ info\ V\ (pool\ b))\ !\ ia)) \implies$$
  

$$\forall jj < length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ ia)).$$
  

$$\neg (let\ bits = bits\ (levels\ (mem\ pool\ info\ V\ (pool\ b))\ !\ ia); a = jj\ div\ 4 * 4$$
  

$$in\ bits\ !\ a = FREE \wedge bits\ !\ (a + 1) = FREE \wedge bits\ !\ (a + 2) = FREE$$
  

$$\wedge bits\ !\ (a + 3) = FREE) \implies$$
  

$$lvl\ V2\ t - Suc\ NULL = ia \implies$$
  

$$levels\ (mem\ pool\ info\ V\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ NULL) = levels\ (mem\ pool\ info$$
  

$$V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ NULL) \implies$$
  

$$\neg (let\ bits = bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ NULL)) [bn$$
  

$$V2\ t\ div\ 4 := FREEING]; a = jj\ div\ 4 * 4$$
  

$$in\ bits\ !\ a = FREE \wedge bits\ !\ (a + 1) = FREE \wedge bits\ !\ (a + 2) = FREE \wedge$$
  

$$bits\ !\ (a + 3) = FREE)$$
  
**apply**(*unfold Let-def*)  
**apply**(*rule subst[where s=list-updates-n (bits (levels (mem-pool-info V2 (pool b))*  

$$!\ (lvl\ V2\ t - Suc\ NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ and$$
  

$$t=bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)) [bn\ V2\ t\ div\ 4 := FREEING]]$$
  
**using** *lst-updt1-eq-upd* **apply** *fast*  
**apply**(*subgoal-tac length (list-updates-n (bits (levels (mem-pool-info V2 (pool b))*  

$$!\ (lvl\ V2\ t - Suc\ NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING$$
  

$$= length\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ ia)))$$
  
**prefer** 2 **using** *length-list-update-n* **apply** *fast*  
**apply**(*subgoal-tac*  $\forall jj < length\ (list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool$   

$$b))\ !\ (lvl\ V2\ t - Suc\ NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING).$$
  

$$\neg (let\ a = jj\ div\ 4 * 4$$
  

$$in\ list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ !\ a = FREE \wedge$$
  

$$list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ !\ (a + 1) = FREE \wedge$$
  

$$list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ !\ (a + 2) = FREE \wedge$$
  

$$list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ !\ (a + 3) = FREE))$$
  
**prefer** 2  
**apply**(*rule partnerbits-udptn-notbit-partbits[of bits (levels (mem-pool-info V2*  

$$(pool\ b))\ !\ ia)\ FREE\ FREEING$$
  

$$list-updates-n\ (bits\ (levels\ (mem\ pool\ info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc$$
  

$$NULL)))\ (bn\ V2\ t\ div\ 4)\ 1\ FREEING\ (bn\ V2\ t\ div\ 4)\ 1]$$
  
**apply**(*unfold Let-def*)[1] **apply** *metis*  
**apply** *blast*

**apply** *fast*  
**apply**(*unfold Let-def*)  
**apply**(*subgoal-tac length (bits (levels (mem-pool-info V2 (pool b))*  
 $[lvl\ V2\ t - Suc\ 0 := (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ 0))$   
 $(\llbracket bits := bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t - Suc\ 0)) \rrbracket [bn\ V2\ t\ div\ 4 := FREEING])\ !$   
 $ia)) = length\ (bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ ia))$ )  
**prefer** 2  
**using** *mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h2* **ap-**  
**ply** *fast*  
**by** *metis*

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1-h3*:  
 $lvl\ V2\ t - Suc\ NULL \neq ia \implies$   
 $bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ ia) =$   
 $bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))$   
 $(\llbracket levels := levels\ (mem-pool-info\ V2\ (pool\ b))$   
 $[lvl\ V2\ t - Suc\ NULL := (levels\ (mem-pool-info\ V2\ (pool\ b))\ !$   
 $(lvl\ V2\ t - Suc\ NULL)$   
 $(\llbracket bits := bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ (lvl\ V2\ t -$   
 $Suc\ NULL) \rrbracket [bn\ V2\ t\ div\ 4 := FREEING])\ !$   
 $ia)$   
**by** *auto*

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1*:  
 $pool\ b \in mem-pools\ V2 \implies$   
 $inv\ V \implies$   
 $NULL < lvl\ V2\ t \implies$   
 $partner-bits\ (set-bit-free\ (mem-pool-info\ V)\ (pool\ b)\ (lvl\ V2\ t)\ (bn\ V2\ t)\ (pool\ b))$   
 $(lvl\ V2\ t)\ (bn\ V2\ t) \implies$   
 $pool\ b \in mem-pools\ V \implies$   
 $level\ b < length\ (levels\ (mem-pool-info\ V\ (pool\ b))) \implies$   
 $(V2,\ V(\llbracket mem-pool-info := set-bit-free\ (mem-pool-info\ V)\ (pool\ b)\ (lvl\ V2\ t)\ (bn$   
 $V2\ t),$   
 $freeing-node := (freeing-node\ V)(t := None)\rrbracket)$   
 $\in gvars-conf-stable \implies$   
 $block\ b < length\ (bits\ (levels\ (mem-pool-info\ V\ (pool\ b))\ !\ level\ b)) \implies$   
 $\forall j. j \neq lvl\ V2\ t \longrightarrow$   
 $levels\ (set-bit-free\ (mem-pool-info\ V)\ (pool\ b)\ (lvl\ V2\ t)\ (bn\ V2\ t)\ (pool\ b))\ !$   
 $j =$   
 $levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ j \implies$   
 $level\ b < length\ (lsizes\ V2\ t) \implies$   
 $bits\ (levels\ (mem-pool-info\ V2\ (pool\ b))\ !\ lvl\ V2\ t) =$   
 $list-updates-n\ (bits\ (levels\ (set-bit-free\ (mem-pool-info\ V)\ (pool\ b)\ (lvl\ V2\ t)\ (bn$   
 $V2\ t)\ (pool\ b))\ !\ lvl\ V2\ t)$   
 $(bn\ V2\ t\ div\ 4 * 4)\ NOEXIST \implies$   
 $lvl\ V2\ t \leq level\ b \implies$



```

    ia < length (levels (mem-pool-info V2 (pool b))) ==>
    jj < length (bits (levels (mem-pool-info V2 (pool b)))
      [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl
V2 t - Suc NULL))
      (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
      [bn V2 t div 4 := FREEING]]]) !
    ia)) ==>
    NULL < ia ==>
    partner-bits
    (mem-pool-info V2 (pool b)
    (levels := levels (mem-pool-info V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc NULL))
    [bn V2 t div 4 := FREEING]]]))
    ia jj ==>
    False
  apply(simp add:set-bit-def)
  apply(subgoal-tac length (levels (mem-pool-info V2 (pool b)))
    = length (levels (mem-pool-info V (pool b)))) prefer 2
    using mempool-free-stm8-atombody-rest-one-finalstm-len-lvls apply blast

  apply(subgoal-tac ¬ partner-bits (mem-pool-info V2 (pool b) (levels := levels (mem-pool-info
V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL))
    [bn V2 t div 4 := FREEING]]]))
    ia jj) apply fast

  apply(subgoal-tac length (bits (levels (mem-pool-info V2 (pool b)) ! ia)) = length
(bits (levels (mem-pool-info V (pool b)) ! ia)))
    prefer 2 apply(case-tac lvl V2 t = ia) apply(simp add:set-bit-def) apply
presburger
  apply(subgoal-tac ∀ jj < length (bits (levels (mem-pool-info V2 (pool b)) ! ia)). ¬
partner-bits (mem-pool-info V (pool b)) ia jj)
    prefer 2 apply(simp add:inv-def inv-bitmap-not4free-def Let-def)

  apply(case-tac lvl V2 t = ia)
    apply(rule subst[where s=partner-bits (mem-pool-info V2 (pool b)) ia jj and
t=partner-bits (mem-pool-info V2 (pool b) (levels := levels (mem-pool-info
V2 (pool b))
    [lvl V2 t - Suc NULL := (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc NULL))
    (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc
NULL))
    [bn V2 t div 4 := FREEING]]))]
    ia jj]) apply(simp add:partner-bits-def Let-def)

```

```

apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b)) ! lvl V2 t) =
  list-updates-n (bits (levels (mem-pool-info V (pool b)) ! lvl V2 t)) (bn V2 t div
4 * 4) 4 NOEXIST) prefer 2
  using lst-udptn-set-eq[of 4 bits (levels (mem-pool-info V (pool b)) ! lvl V2 t)
bn V2 t FREE NOEXIST] apply simp

apply(unfold partner-bits-def)[1]
apply(subgoal-tac  $\neg$  (let a = jj div 4 * 4
  in list-updates-n (bits (levels (mem-pool-info V (pool b)) ! ia)) (bn V2 t div 4
* 4) 4 NOEXIST ! a = FREE  $\wedge$ 
  list-updates-n (bits (levels (mem-pool-info V (pool b)) ! ia)) (bn V2 t div 4
* 4) 4 NOEXIST ! (a + 1) = FREE  $\wedge$ 
  list-updates-n (bits (levels (mem-pool-info V (pool b)) ! ia)) (bn V2 t div 4
* 4) 4 NOEXIST ! (a + 2) = FREE  $\wedge$ 
  list-updates-n (bits (levels (mem-pool-info V (pool b)) ! ia)) (bn V2 t div 4
* 4) 4 NOEXIST ! (a + 3) = FREE))
prefer 2
apply(rule partnerbits-udptn-notbit-partbits[rule-format, of bits (levels (mem-pool-info
V (pool b)) ! ia) FREE NOEXIST
  list-updates-n (bits (levels (mem-pool-info V (pool b)) ! ia)) (bn V2 t div
4 * 4) 4 NOEXIST
  bn V2 t div 4 * 4 4 jj])
apply(unfold Let-def)[1] apply presburger
apply blast apply fast apply force
apply(unfold Let-def)[1] apply presburger

```

```

apply(case-tac lvl V2 t - Suc 0 = ia)
apply(unfold partner-bits-def)
apply(rule subst[where s=bits (levels (mem-pool-info V2 (pool b)) ! (lvl V2 t -
Suc 0))[bn V2 t div 4 := FREEING]
  and t=bits (levels (mem-pool-info V2 (pool b))
    (levels := levels (mem-pool-info V2 (pool b))
      [lvl V2 t - Suc 0 := (levels (mem-pool-info V2 (pool
b)) ! (lvl V2 t - Suc 0))
      (bits := bits (levels (mem-pool-info V2 (pool b)) !
(lvl V2 t - Suc 0))[bn V2 t div 4 := FREEING])])]) !
  ia)])
apply(subgoal-tac lvl V2 t - Suc NULL < length (levels (mem-pool-info V2
(pool b)))) prefer 2 apply blast
using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1-h1
apply blast

```

```

apply(subgoal-tac levels (mem-pool-info V (pool b)) ! (lvl V2 t - Suc 0) = levels
(mem-pool-info V2 (pool b)) ! (lvl V2 t - Suc 0))
prefer 2 apply presburger
using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1-h2 ap-
ply blast

```

**apply**(rule subst[**where**  $s = \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } ia)$  **and**  
 $t = \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b))$   
 $\quad (\text{levels } := \text{levels } (\text{mem-pool-info } V2 \text{ (pool } b))$   
 $\quad [\text{lvl } V2 \text{ } t - \text{Suc } NULL := (\text{levels } (\text{mem-pool-info } V2$   
 $\text{(pool } b)) \text{ ! } (\text{lvl } V2 \text{ } t - \text{Suc } NULL))$   
 $\quad (\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } (\text{lvl } V2 \text{ } t - \text{Suc } NULL))$   
 $\text{[bn } V2 \text{ } t \text{ div } 4 := \text{FREEING}]]) \text{ ! } ia])$   
**using** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1-h3 **ap-**  
**ply** blast

**apply**(rule subst[**where**  $s = \text{levels } (\text{mem-pool-info } V \text{ (pool } b)) \text{ ! } ia$  **and**  $t = \text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } ia]$   
**apply** metis  
**apply**(unfold Let-def)

**apply**(subgoal-tac length (bits (levels (mem-pool-info V2 (pool b))  
 $[\text{lvl } V2 \text{ } t - \text{Suc } NULL := (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } (\text{lvl } V2 \text{ } t - \text{Suc } NULL))$   
 $(\text{bits } := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \text{ (pool } b)) \text{ ! } (\text{lvl } V2 \text{ } t - \text{Suc } NULL))$   
 $\text{[bn } V2 \text{ } t \text{ div } 4 := \text{FREEING}]) \text{ ! } ia)) =$   
length (bits (levels (mem-pool-info V2 (pool b)) ! ia))) **prefer** 2  
**using** mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info-h2 **apply**  
blast  
**by** metis

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case2:  
 $p \in \text{mem-pools } V2 \implies$   
 $\text{inv } V \implies$   
 $NULL < \text{lvl } V2 \text{ } t \implies$   
 $\text{partner-bits } (\text{set-bit-free } (\text{mem-pool-info } V) \text{ (pool } b) \text{ (lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t) \text{ (pool } b))$   
 $(\text{lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t) \implies$   
 $\text{pool } b \in \text{mem-pools } V \implies$   
 $\text{level } b < \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) \implies$   
 $(V2, V[\text{mem-pool-info} := \text{set-bit-free } (\text{mem-pool-info } V) \text{ (pool } b) \text{ (lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t),$   
 $\text{freeing-node} := (\text{freeing-node } V)(t := \text{None}]) \implies$   
 $\in \text{gvars-conf-stable} \implies$   
 $\text{lvl } V2 \text{ } t \leq \text{level } b \implies$   
 $\forall p. p \neq \text{pool } b \implies \text{mem-pool-info } V2 \text{ } p = \text{set-bit-free } (\text{mem-pool-info } V) \text{ (pool } b) \text{ (lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t) \text{ } p \implies$   
 $p \neq \text{pool } b \implies$   
 $ia < \text{length } (\text{levels } (\text{set-bit-free } (\text{mem-pool-info } V) \text{ (pool } b) \text{ (lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t) \text{ } p)) \implies$   
 $jj < \text{length } (\text{bits } (\text{levels } (\text{set-bit-free } (\text{mem-pool-info } V) \text{ (pool } b) \text{ (lvl } V2 \text{ } t) \text{ (bn } V2 \text{ } t) \text{ } p)) \text{ ! } ia))$

```

t) p) ! ia))  $\implies$ 
  NULL < ia  $\implies$  partner-bits (set-bit-free (mem-pool-info V) (pool b) (lvl V2 t)
  (bn V2 t) p) ia jj  $\implies$  False
apply(simp add:set-bit-def)
apply(subgoal-tac length (levels (mem-pool-info V2 (pool b)))
      = length (levels (mem-pool-info V (pool b)))) prefer 2
using mempool-free-stm8-atombody-rest-one-finalstm-len-lvls apply blast

apply(subgoal-tac  $\neg$  partner-bits (mem-pool-info V p) ia jj) apply fast

apply(subgoal-tac p  $\in$  mem-pools V) prefer 2 apply(simp add:gvars-conf-stable-def
gvars-conf-def)
apply(subgoal-tac  $\forall$  jj < length (bits (levels (mem-pool-info V p) ! ia)).  $\neg$  partner-bits
(mem-pool-info V p) ia jj)
prefer 2 apply(simp add:inv-def inv-bitmap-not4free-def Let-def)

by blast

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free:
  V  $\in$  mp-free-precond8-3 t b  $\alpha \cap \{\text{'cur} = \text{Some } t\} \implies$ 
  {free-stm8-precond2 V t b}  $\cap \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info}$ 
  (pool b)) (\text{'lvl } t) (\text{'bn } t)\} \neq \{\} \implies
  V2  $\in$  free-stm8-precond3 V t b  $\cap \{\text{'i } t = 4\} \implies$ 
  x = free-stm8-atombody-rest-cond3 (V2(lvl := (lvl V2)(t := lvl V2 t - 1), bn
  := (bn V2)(t := bn V2 t div 4)) t b  $\implies$ 
  y = x(freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
  t, block = bn x t,
  data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz (mem-pool-info
  x (pool b)) div 4 ^ lvl x t) (bn x t)))  $\implies$ 
  inv-bitmap-not4free y
apply(simp add:inv-bitmap-not4free-def Let-def)
apply clarify
apply(rule conjI)
apply clarsimp
using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case1 ap-
ply blast

apply clarsimp
using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvls-not4free-case2 ap-
ply blast
done

```

```

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv':
  V  $\in$  mp-free-precond8-3 t b  $\alpha \cap \{\text{'cur} = \text{Some } t\} \implies$ 
  {free-stm8-precond2 V t b}  $\cap \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info}$ 
  (pool b)) (\text{'lvl } t) (\text{'bn } t)\} \neq \{\} \implies
  V2  $\in$  free-stm8-precond3 V t b  $\cap \{\text{'i } t = 4\} \implies$ 
  x = free-stm8-atombody-rest-cond3 (V2(lvl := (lvl V2)(t := lvl V2 t - 1), bn

```

```

:= (bn V2)(t := bn V2 t div 4)) t b ==>
  y = x(⟦freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
t, block = bn x t,
      data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))⟧) ==>
  inv y
apply(rule subst[where s=inv-cur y ∧ inv-thd-waitq y ∧ inv-mempool-info y
  ∧ inv-bitmap-freelist y ∧ inv-bitmap y ∧ inv-aux-vars y
  ∧ inv-bitmap0 y ∧ inv-bitmapn y ∧ inv-bitmap-not4free y and t=inv
y])
using inv-def[of y] apply fast

apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-cur[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-thd-waitq[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-mempool-info[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap-freelist[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-bitmap[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-aux-vars[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl0[of
V t b α V2 x y] apply fast
apply(rule conjI) using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl n[of
V t b α V2 x y] apply fast
  using mempool-free-stm8-atombody-rest-one-finalstm-inv-lvl s-not4free[of
V t b α V2 x y] apply fast

done

lemma mempool-free-stm8-atombody-rest-one-finalstm-inv:
  V ∈ mp-free-precond8-3 t b α ∩ ⟦'cur = Some t⟧ ==>
  {free-stm8-precond2 V t b} ∩ ⟦NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)⟧ ≠ {} ==>
  V2 ∈ free-stm8-precond3 V t b ∩ ⟦'i t = 4⟧ ==>
  x = free-stm8-atombody-rest-cond3 (V2(⟦lvl := (lvl V2)(t := lvl V2 t - 1), bn
:= (bn V2)(t := bn V2 t div 4)⟧) t b ==>
  x(⟦freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x t,
block = bn x t,
      data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))⟧)
  ∈ ⟦'inv⟧
using mempool-free-stm8-atombody-rest-one-finalstm-inv'[of V t b α V2 x
  x(⟦freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x
t, block = bn x t,

```

$data = block\_ptr (mem\_pool\_info\ x\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ x\ (pool\ b)))\ \text{div}\ 4\ \wedge\ lvl\ x\ t)\ (bn\ x\ t))\ \}\}$  **apply fast**  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h2:*

$V \in mp\_free\_precond8-3\ t\ b\ \alpha \cap \{\cur = Some\ t\} \implies$   
 $\{\text{free-stm8-precond2}\ V\ t\ b\} \cap \{\text{NULL} < \text{'}lvl\ t \wedge partner\_bits\ (\text{'}mem\_pool\_info\ (pool\ b))\ (\text{'}lvl\ t)\ (\text{'}bn\ t))\} \neq \{\} \implies$   
 $V2 \in free\_stm8\_precond3\ V\ t\ b \cap \{\text{'}i\ t = 4\} \implies$   
 $x = free\_stm8\_atombody\_rest\_cond3\ (V2\ (lvl := (lvl\ V2)\ (t := lvl\ V2\ t - 1),\ bn := (bn\ V2)\ (t := bn\ V2\ t\ \text{div}\ 4)))\ t\ b \implies$   
 $x\ (freeing\_node := (freeing\_node\ x)\ (t := Some\ (pool = pool\ b,\ level = lvl\ x\ t,\ block = bn\ x\ t),$   
 $data = block\_ptr (mem\_pool\_info\ x\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ x\ (pool\ b)))\ \text{div}\ 4\ \wedge\ lvl\ x\ t)\ (bn\ x\ t))\ \}\}$   
 $\in \{\text{'}allocating\_node\ t = None\}$   
**by** (*simp add:Let-def block-ptr-def*)

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h1-2:*

$V \in mp\_free\_precond8-3\ t\ b\ \alpha \cap \{\cur = Some\ t\} \implies$   
 $\{\text{free-stm8-precond2}\ V\ t\ b\} \cap \{\text{NULL} < \text{'}lvl\ t \wedge partner\_bits\ (\text{'}mem\_pool\_info\ (pool\ b))\ (\text{'}lvl\ t)\ (\text{'}bn\ t))\} \neq \{\} \implies$   
 $V2 \in free\_stm8\_precond3\ V\ t\ b \cap \{\text{'}i\ t = 4\} \implies$   
 $x = free\_stm8\_atombody\_rest\_cond3\ (V2\ (lvl := (lvl\ V2)\ (t := lvl\ V2\ t - 1),\ bn := (bn\ V2)\ (t := bn\ V2\ t\ \text{div}\ 4)))\ t\ b \implies$   
 $y = x\ (freeing\_node := freeing\_node\ x\ (t \mapsto$   
 $(pool = pool\ b,\ level = lvl\ x\ t,\ block = bn\ x\ t,$   
 $data = block\_ptr (mem\_pool\_info\ x\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ x\ (pool\ b)))\ \text{div}\ 4\ \wedge\ lvl\ x\ t)\ (bn\ x\ t))\ \}\} \implies$   
 $y \in \{\text{'}(Pair\ V)$   
 $\in \{(s, r). (\text{cur}\ s \neq Some\ t \longrightarrow gvars\_nochange\ s\ r \wedge lvars\_nochange\ t\ s\ r) \wedge$   
 $(\text{cur}\ s = Some\ t \longrightarrow invariant.in\ s \longrightarrow invariant.in\ r) \wedge (\forall t'. t' \neq t \longrightarrow lvars\_nochange\ t'\ s\ r))\}$   
**apply**(*subgoal-tac* (*cur*  $V \neq Some\ t \longrightarrow gvars\_nochange\ V\ y \wedge lvars\_nochange\ t\ V\ y$ )  $\wedge$   
 $(\text{cur}\ V = Some\ t \longrightarrow invariant.in\ V \longrightarrow invariant.in\ y) \wedge (\forall t'. t' \neq t \longrightarrow lvars\_nochange\ t'\ V\ y))$ )  
**prefer 2**  
**apply**(*rule conjI*)  
**apply**(*subgoal-tac* *cur*  $V = Some\ t$ ) **prefer 2** **apply fast** **apply fast**  
**apply**(*rule conjI*)  
**apply**(*rule impI*) + **using** *mempool-free-stm8-atombody-rest-one-finalstm-inv'* [*of*  
 $V\ t\ b\ \alpha\ V2\ x\ y]$  **apply fast**  
**apply**(*rule allI*) **apply**(*rule impI*) **apply**(*simp add:lvars-nochange-def Let-def*)  
**apply fast**  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h1-h1* :

$\forall j. j \neq \text{lvl } V \ t \longrightarrow \text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ j = \text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ j \implies$

$\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ \text{lvl } V \ t) =$

$\text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)[\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}])$

$(\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) \ \text{div } 4 * 4) \ 4 \ \text{NOEXIST} \implies$

$\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ ia)) =$

$\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b))$

$[\text{lvl } V2 \ t - \text{Suc } \text{NULL} := (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$

$(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))[\text{bn } V2 \ t \ \text{div } 4 := \text{FREEING}]) \ !$

$ia))$

**apply**(*rule subst*[**where**  $s = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ ia))$ ])

**and**  $t = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b))$

$[\text{lvl } V2 \ t - \text{Suc } 0 := (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ (\text{lvl } V2 \ t - \text{Suc } 0))$

$(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) \ ! \ (\text{lvl } V2 \ t - \text{Suc } 0))[\text{bn } V2 \ t \ \text{div } 4 := \text{FREEING}]) \ !$

$ia))])$

**apply**(*case-tac*  $ia = \text{lvl } V2 \ t - \text{Suc } 0$ )

**apply**(*case-tac*  $ia < \text{length } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)))$ )

**apply** *auto*[1] **apply** *auto*[1] **apply** *auto*[1]

**apply**(*case-tac*  $ia = \text{lvl } V \ t$ )

**apply**(*subgoal-tac*  $\text{length } (\text{list-updates-n}$

$(\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b))$

$[\text{lvl } V \ t := (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)$

$(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)[\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}]) \ !$

$ia))$

$(\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) \ \text{div } 4 * 4) \ 4 \ \text{NOEXIST}) = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b))$

$[\text{lvl } V \ t := (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)$

$(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)[\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}]) \ !$

$ia))$

$)$

**prefer** 2 **using** *length-list-update-n* **apply** *fast*

**apply**(*subgoal-tac*  $\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b))$

$[\text{lvl } V \ t := (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)$

$(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ \text{lvl } V \ t)[\text{block } b \ \text{div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}]) \ !$

$ia)) = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) \ ! \ ia))$

$)$

**prefer** 2 **apply**(*case-tac*  $ia = \text{lvl } V \ t$ )

**apply**(*case-tac*  $ia < \text{length } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)))$ )

**apply** *auto*[1] **apply** *auto*[1] **apply** *auto*[1]

**apply** *auto*[1]

```

apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
  [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level
b - lvl V t) := FREE]])) !
  ia)) = length (bits (levels (mem-pool-info V (pool b)) ! ia)))
prefer 2 apply(case-tac ia = lvl V t) apply(case-tac ia < length (levels
(mem-pool-info V (pool b))))
apply auto[1] apply auto[1] apply auto[1]
apply auto[1]
done

```

**lemma** mempool-free-stm8-atombody-rest-one-finalstm-h1:

```

V ∈ mp-free-precond8-3 t b α ∩ {cur = Some t} ⇒
{free-stm8-precond2 V t b} ∩ {NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)} ≠ {} ⇒
V2 ∈ free-stm8-precond3 V t b ∩ {i t = 4} ⇒
x = free-stm8-atombody-rest-cond3 (V2 (lvl := (lvl V2)(t := lvl V2 t - 1), bn
:= (bn V2)(t := bn V2 t div 4))) t b ⇒
x (freeing-node := (freeing-node x) (t := Some (pool = pool b, level = lvl x t,
block = bn x t,
  data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t) (bn x t))))
  ∈ {'(Pair V) ∈ Mem-pool-free-guar t}
apply(unfold Mem-pool-free-guar-def)
apply(rule pairv-rId)
apply(rule pairv-IntI) apply(rule pairv-IntI)

```

```

apply(unfold gvars-conf-stable-def gvars-conf-def)[1]
apply clarify apply(simp add:Let-def set-bit-def)
apply clarify using mempool-free-stm8-atombody-rest-one-finalstm-h1-h1[of V t
b V2] apply blast

```

```

using mempool-free-stm8-atombody-rest-one-finalstm-h1-2[of V t b α V2 x
x (freeing-node := freeing-node x(t ↦
  (pool = pool b, level = lvl x t, block = bn x t,
    data = block-ptr (mem-pool-info x (pool b)) (ALIGN4 (max-sz
(mem-pool-info x (pool b))) div 4 ^ lvl x t)
    (bn x t)))))] apply fast

```

```

apply(simp add:Let-def)
done

```



**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-I1*:

$x \in \llbracket \text{'invariant.inv} \rrbracket \implies$

$x \in \llbracket \text{'allocating-node } t = \text{None} \rrbracket \implies$

$x \in \llbracket \text{'invariant.inv} \wedge \text{'allocating-node } t = \text{None} \rrbracket$

**by** *auto*

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h3-h1*:

*inv*  $V \wedge$

*pool*  $b \in \text{mem-pools } V2 \wedge$

*level*  $b < \text{length } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) \wedge \text{lvl } V2 \ t \leq \text{level } b \wedge \text{NULL}$   
 $< \text{lvl } V2 \ t \implies$

*mem-pools*  $V = \text{mem-pools } V2 \wedge$

*lvl*  $V \ t = \text{lvl } V2 \ t \implies$

*n-max*  $(\text{mem-pool-info } V (\text{pool } b)) * 4 \wedge (\text{lvl } V2 \ t - \text{Suc } \text{NULL}) =$

*length*  $(\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$

**apply**(*simp add:inv-def inv-mempool-info-def Let-def*) **apply** *auto*[1]

**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h3*:

*invariant.inv*  $V \wedge$

*pool*  $b \in \text{mem-pools } V2 \wedge$

*level*  $b < \text{length } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) \wedge$

*block*  $b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) ! \text{level } b) \wedge$

*level*  $b < \text{length } (\text{lsizes } V2 \ t) \wedge$

*bn*  $V \ t < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) ! \text{lvl } V2 \ t) \wedge$

*lvl*  $V2 \ t \leq \text{level } b \wedge$

$\text{NULL} < \text{lvl } V2 \ t \implies$

*mem-pools*  $V = \text{mem-pools } V2 \wedge$

$(\forall p.$

$(\forall i. \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ p)) ! i)) =$

$\text{length } (\text{bits } (\text{levels } (\text{if } p = \text{pool } b$

$\text{then mem-pool-info } V (\text{pool } b)$

$(\text{levels} := \text{levels } (\text{mem-pool-info } V (\text{pool } b))$

$[\text{lvl } V \ t := (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) !$

$\text{lvl } V \ t)$

$([\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))$

$b)) ! \text{lvl } V \ t][\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}]]])$

$\text{else mem-pool-info } V \ p) !$

$i))))) \wedge$

$(\forall p. p \neq \text{pool } b \longrightarrow \text{mem-pool-info } V2 \ p = \text{mem-pool-info } V \ p) \wedge$

$(\forall j. j \neq \text{lvl } V \ t \longrightarrow \text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! j = \text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! j) \wedge$

$\text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! \text{lvl } V \ t) =$

$\text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! \text{lvl } V \ t)[\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}])$

$(\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) \text{ div } 4 * 4) (i \ V2 \ t) \text{ NOEXIST} \wedge$

$\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) = \text{bn } V2 \ t \wedge$

$\text{lvl } V \ t = \text{lvl } V2 \ t \wedge \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } V (\text{pool } b))) \text{ div } 4 \wedge \text{lvl } V$

$t = \text{lsz } V2 \ t \wedge \text{lsizes } V \ t = \text{lsizes } V2 \ t \wedge i \ V2 \ t \leq 4 \wedge i \ V2 \ t = 4 \implies$   
 $x = V2(\text{lvl} := (\text{lvl } V2)(t := \text{lvl } V2 \ t - \text{Suc } \text{NULL}), \text{bn} := (\text{bn } V2)(t := \text{bn } V2 \ t \text{ div } 4),$   
 $\text{mem-pool-info} := (\text{mem-pool-info } V2)$   
 $(\text{pool } b := \text{mem-pool-info } V2 \ (\text{pool } b))$   
 $(\text{levels} := \text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)))$   
 $[\text{lvl } V2 \ t - \text{Suc } \text{NULL} := (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))[\text{bn } V2 \ t \text{ div } 4 := \text{FREEING}]])) \implies$   
 $\text{bn } V2 \ t \text{ div } 4$   
 $< \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)))$   
 $[\text{lvl } V2 \ t - \text{Suc } \text{NULL} := (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))[\text{bn } V2 \ t \text{ div } 4 := \text{FREEING}]] !$   
 $(\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$   
**apply**(rule subst[**where**  $s = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$  **and**  $t =$   
 $\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)))$   
 $[\text{lvl } V2 \ t - \text{Suc } \text{NULL} := (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } \text{NULL}))[\text{bn } V2 \ t \text{ div } 4 := \text{FREEING}]] !$   
 $(\text{lvl } V2 \ t - \text{Suc } \text{NULL}))$  ])  
**apply**(subgoal-tac  $\forall j. j \neq \text{lvl } V \ t \longrightarrow \text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) ! j =$   
 $\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! j$ )  
**prefer 2 apply fast**  
**apply**(subgoal-tac  $\text{bits } (\text{levels } (\text{mem-pool-info } V2 \ (\text{pool } b)) ! \text{lvl } V \ t) =$   
 $\text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) ! \text{lvl } V \ t)[\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) := \text{FREE}]$   
 $(\text{block } b \text{ div } 4 \wedge (\text{level } b - \text{lvl } V \ t) \text{ div } 4 * 4) (i \ V2 \ t) \text{ NOEXIST}$ )  
**prefer 2 apply fast**  
**using** mempool-free-stm8-atombody-rest-one-finalstm-h1-h1[*of*  $V \ t \ b \ V2 \ (\text{lvl } V2 \ t - \text{Suc } \text{NULL})$ ] **apply** auto[1]

**apply**(rule subst[**where**  $s = (\text{n-max } (\text{mem-pool-info } V \ (\text{pool } b))) * 4 \wedge (\text{lvl } V2 \ t - \text{Suc } 0)$   
**and**  $t = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) ! (\text{lvl } V2 \ t - \text{Suc } 0)))$ ]  
**using** mempool-free-stm8-atombody-rest-one-finalstm-h3-h1 **apply fast**

**apply**(subgoal-tac  $\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \ (\text{pool } b)) ! \text{level } b))$   
 $= (\text{n-max } (\text{mem-pool-info } V \ (\text{pool } b))) * 4 \wedge \text{level } b$ )  
**prefer 2 apply**(simp add:inv-def inv-mempool-info-def Let-def)

**apply**(rule lm11[*of*  $V2 \ t \ b \ V$ ]) **apply simp apply simp applymetis**  
**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm-h4*:

$(\neg \text{free-block-r } V2 \ t \longrightarrow \text{freeing-node } V \ t = \text{None}) \wedge$   
 $\alpha = (\text{if } \exists y. \text{freeing-node } V \ t = \text{Some } y \text{ then } \text{lvl } V \ t + 1 \text{ else } \text{NULL}) \wedge$   
 $V \in (\text{if } \text{NULL} < \alpha \text{ then } \text{UNIV} \text{ else } \{\}) \implies \text{free-block-r } V2 \ t$

**apply** *auto*

**done**

**lemma** *mempool-free-stm8-atombody-rest-one-finalstm*:

$V \in \text{mp-free-precond8-3 } t \ b \ \alpha \cap \{\cur = \text{Some } t\} \implies$   
 $\{\text{free-stm8-precond2 } V \ t \ b\} \cap \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info}$   
 $(\text{pool } b)) (\text{'lvl } t) (\text{'bn } t)\} \neq \{\} \implies$   
 $V2 \in \text{free-stm8-precond3 } V \ t \ b \cap \{i \ t = 4\} \implies$   
 $\{\text{free-stm8-atombody-rest-cond3 } (V2(\text{lvl} := (\text{lvl } V2)(t := \text{lvl } V2 \ t - 1), \text{bn} :=$   
 $(\text{bn } V2)(t := \text{bn } V2 \ t \text{ div } 4)) \ t \ b\}$   
 $\subseteq \{\text{'(freeing-node-update}$   
 $(\lambda-. \text{'freeing-node}(t \mapsto$   
 $(\text{pool} = \text{pool } b, \text{level} = \text{'lvl } t, \text{block} = \text{'bn } t,$   
 $\text{data} = \text{block-ptr } (\text{'mem-pool-info } (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz}$   
 $(\text{'mem-pool-info } (\text{pool } b))) \text{ div } 4 \wedge \text{'lvl } t) (\text{'bn } t)))\}$   
 $\in \{\text{'(Pair } V) \in \text{Mem-pool-free-guar } t\} \cap \text{mp-free-precond8-inv } t \ b \ (\alpha - 1)\}$

**apply** *(rule subsetI)*

**apply** *(subgoal-tac x = free-stm8-atombody-rest-cond3 (V2(lvl := (lvl V2)(t := lvl V2 t - 1), bn := (bn V2)(t := bn V2 t div 4))) t b)*

**prefer** 2 **apply** *fast*

**apply** *(subgoal-tac x(lfreeing-node := (freeing-node x) (t := Some (lpool = pool b, level = lvl x t, block = bn x t,  
 $\text{data} = \text{block-ptr } (\text{mem-pool-info } x \ (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz}$   
 $(\text{mem-pool-info } x \ (\text{pool } b))) \text{ div } 4 \wedge \text{lvl } x \ t) (\text{bn } x \ t)))$*

$(\alpha - 1))$

**apply** *blast*

**apply** *(rule IntI)*

**using** *mempool-free-stm8-atombody-rest-one-finalstm-h1* *[of V t b  $\alpha$  V2]* **apply** *meson*

**apply** *(rule IntI)*

**apply** *(rule IntI)*

**apply** *(rule IntI)*

**apply** *(rule mempool-free-stm8-atombody-rest-one-finalstm-I1)*

**using** *mempool-free-stm8-atombody-rest-one-finalstm-inv* *[of V t b  $\alpha$  V2]* **apply** *meson*

**using** *mempool-free-stm8-atombody-rest-one-finalstm-h2* *[of V t b  $\alpha$  V2]* **apply** *meson*

```

apply(simp add:Let-def gvars-conf-stable-def gvars-conf-def block-ptr-def set-bit-def)
apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b)) ! level b)) =
      length (bits (levels (mem-pool-info V2 (pool b))
        [lvl V2 t - Suc NULL := (levels (mem-pool-info V2
(pool b)) ! (lvl V2 t - Suc NULL))
        (bits := bits (levels (mem-pool-info V2 (pool b)) ! (lvl
V2 t - Suc NULL))[bn V2 t div 4 := FREEING]]] !
        level b)) )
  prefer 2 apply(subgoal-tac  $\forall j. j \neq \text{lvl } V \ t \longrightarrow \text{levels (mem-pool-info } V \ (\text{pool } b)) \ ! \ j$ 
    = levels (mem-pool-info V2 (pool b)) ! j) prefer 2 apply fast
  apply(subgoal-tac bits (levels (mem-pool-info V2 (pool b)) ! lvl V t) =
    list-updates-n (bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div
4 ^ (level b - lvl V t) := FREE])
    (block b div 4 ^ (level b - lvl V t) div 4 * 4) (i V2 t) NOEXIST)
  prefer 2 apply fast
  using mempool-free-stm8-atombody-rest-one-finalstm-h1-h1[of V t b V2 level b]
apply argo
apply auto[1]

```

```

apply(simp add:Let-def gvars-conf-stable-def gvars-conf-def block-ptr-def set-bit-def)
apply(rule conjI)
  using mempool-free-stm8-atombody-rest-one-finalstm-h3 apply blast

```

```

apply(rule conjI)
  apply(rule subst[where s=lvl V t and t=lvl V2 t]) apply fast
  apply (metis Nat.add-diff-assoc div-mult2-eq plus-1-eq-Suc power-add power-commutes
power-one-right)

```

```

apply(rule conjI)
  apply (metis Suc-pred le-imp-less-Suc nat-le-linear not-less)

```

```

apply(rule conjI)
  apply(rule subst[where s=max-sz (mem-pool-info V (pool b)) and t=ALIGN4
(max-sz (mem-pool-info V (pool b)))]
  using inv-maxsz-align4 apply auto[1]
  apply clarify apply(rule conjI)
  apply fast

```

```

  apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b)) ! level b)) =
(n-max (mem-pool-info V (pool b))) * 4 ^ level b)
  prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def)
  apply(subgoal-tac level b > 0 ) prefer 2 apply auto[1]
  apply(subgoal-tac block b < n-max (mem-pool-info V (pool b)) * 4 ^ level b)
prefer 2 apply argo
  apply(subgoal-tac block b div 4 ^ (level b - lvl V2 t) = bn V2 t) prefer 2 apply
metis
  using lm11[of V2 t b V] apply meson

```

```

using mempool-free-stm8-atombody-rest-one-finalstm-h4[of  $V2\ t\ V$ ] apply fast

apply(simp add:Let-def)
apply clarsimp
apply auto
done

term free-stm8-precond2  $V\ t\ b$ 

lemma mempool-free-stm8-atombody-rest-one:
   $V \in mp\text{-}free\text{-}precond8\text{-}3\ t\ b\ \alpha \cap \{\!\{cur = Some\ t\}\!\} \implies$ 
   $\{\!\{free\text{-}stm8\text{-}precond2\ V\ t\ b\}\!\} \cap \{\!\{NULL < 'lvl\ t \wedge partner\text{-}bits\ ('mem\text{-}pool\text{-}info$ 
  (pool b)\)\ ('lvl' t) ('bn' t)\}\!\} \neq \{\!\{\}\!\} \implies
   $V2 \in free\text{-}stm8\text{-}precond3\ V\ t\ b \cap \{\!\{i\ t = 4\}\!\} \implies$ 
   $\Gamma \vdash_I Some\ ('lvl := 'lvl(t := 'lvl\ t - 1));;$ 
   $'bn := 'bn(t := 'bn\ t\ div\ 4);;$ 
   $'mem\text{-}pool\text{-}info := set\text{-}bit\text{-}freeing\ 'mem\text{-}pool\text{-}info\ (pool\ b)\ ('lvl\ t)\ ('bn\ t);;$ 
   $'freeing\text{-}node := 'freeing\text{-}node(t \mapsto (pool = pool\ b,\ level = 'lvl\ t,\ block = 'bn$ 
  t,
   $data = block\text{-}ptr\ ('mem\text{-}pool\text{-}info\ (pool\ b))\ (ALIGN4\ (max\text{-}sz\ ('mem\text{-}pool\text{-}info$ 
  (pool b)))  $div\ 4 \wedge 'lvl\ t)\ ('bn\ t)))$ 
   $sat_p\ [\{V2\}, \{(s, t). s = t\}, UNIV,$ 
   $\{\!\{(Pair\ V) \in Mem\text{-}pool\text{-}free\text{-}guar\ t\}\!\} \cap mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ (\alpha - 1)]$ 
apply(rule Seq[where mid={free-stm8-atombody-rest-cond3 (free-stm8-atombody-rest-cond2
  (free-stm8-atombody-rest-cond1  $V2\ t\ b$ ) t b) t b})}]])
apply(rule Seq[where mid={free-stm8-atombody-rest-cond2 (free-stm8-atombody-rest-cond1
   $V2\ t\ b$ ) t b})}]])
apply(rule Seq[where mid={free-stm8-atombody-rest-cond1  $V2\ t\ b$ })}]])

apply(rule Basic)
apply fast apply fast using stable-id2 apply fast using stable-id2 apply fast

apply(rule Basic)
apply fast apply fast using stable-id2 apply fast using stable-id2 apply fast

apply(rule Basic)
apply(simp add:set-bit-def Let-def) apply fast using stable-id2 apply fast us-
ing stable-id2 apply fast

apply(rule Basic)
apply(rule subst[where s=bn  $V2$  and t=bn ( $V2(\text{!}lvl := (lvl\ V2)(t := lvl\ V2\ t$ 
   $- 1))$ )]])
apply auto[1]
using mempool-free-stm8-atombody-rest-one-finalstm[of  $V\ t\ b\ \alpha\ V2$ ] apply meson

```

**apply fast using stable-id2 apply fast using stable-id2 apply fast**  
**done**

**lemma** *mempool-free-stm8-atombody-rest*:

$V \in mp\text{-}free\text{-}precond8\text{-}3 \ t \ b \ \alpha \cap \{\cur = Some \ t\} \implies$   
 $\{free\text{-}stm8\text{-}precond2 \ V \ t \ b\} \cap \{NULL < 'lvl \ t \wedge partner\text{-}bits \ ('mem\text{-}pool\text{-}info$   
 $(pool \ b)) \ ('lvl \ t) \ ('bn \ t)\} \neq \{\} \implies$   
 $\Gamma \vdash_I Some \ ('lvl := 'lvl(t := 'lvl \ t - 1));;$   
 $'bn := 'bn(t := 'bn \ t \text{ div } 4);;$   
 $'mem\text{-}pool\text{-}info := set\text{-}bit\text{-}freeing \ 'mem\text{-}pool\text{-}info \ (pool \ b) \ ('lvl \ t) \ ('bn \ t);;$   
 $'freeing\text{-}node := 'freeing\text{-}node(t \mapsto (pool = pool \ b, level = 'lvl \ t, block = 'bn$   
 $t,$   
 $data = block\text{-}ptr \ ('mem\text{-}pool\text{-}info \ (pool \ b)) \ (ALIGN4 \ (max\text{-}sz \ ('mem\text{-}pool\text{-}info$   
 $(pool \ b))) \text{ div } 4 \wedge 'lvl \ t) \ ('bn \ t)))$   
 $sat_p \ [free\text{-}stm8\text{-}precond3 \ V \ t \ b \cap \{i \ t = 4\}, \{(s, t). s = t\}, UNIV,$   
 $\{(Pair \ V) \in Mem\text{-}pool\text{-}free\text{-}guar \ t\} \cap mp\text{-}free\text{-}precond8\text{-}inv \ t \ b \ (\alpha - 1)]$   
**using** *mempool-free-stm8-atombody-rest-one*[*of*  $V \ t \ b \ \alpha$ ]  
 $Allprecond[\textbf{where } U = free\text{-}stm8\text{-}precond3 \ V \ t \ b \cap \{i \ t = 4\} \textbf{ and}$   
 $P = Some \ ('lvl := 'lvl(t := 'lvl \ t - 1));;$   
 $'bn := 'bn(t := 'bn \ t \text{ div } 4);;$   
 $'mem\text{-}pool\text{-}info := set\text{-}bit\text{-}freeing \ 'mem\text{-}pool\text{-}info \ (pool \ b) \ ('lvl \ t)$   
 $('bn \ t);;$   
 $'freeing\text{-}node := 'freeing\text{-}node(t \mapsto (pool = pool \ b, level = 'lvl \ t,$   
 $block = 'bn \ t,$   
 $data = block\text{-}ptr \ ('mem\text{-}pool\text{-}info \ (pool \ b)) \ (ALIGN4 \ (max\text{-}sz$   
 $('mem\text{-}pool\text{-}info \ (pool \ b))) \text{ div } 4 \wedge 'lvl \ t) \ ('bn \ t))) \textbf{ and}$   
 $rely = \{(x, y). x = y\} \textbf{ and}$   
 $guar = UNIV \textbf{ and } post = \{(Pair \ V) \in Mem\text{-}pool\text{-}free\text{-}guar \ t\} \cap$   
 $mp\text{-}free\text{-}precond8\text{-}inv \ t \ b \ (\alpha - 1)]$   
**apply** *meson*  
**done**

**abbreviation** *free-stm8-bd2-cond1*  $V \ t \ b \equiv V(j := (j \ V)(t := lvl \ V \ t))$

**abbreviation** *free-stm8-bd2-cond2*  $V \ t \ b \equiv V(lbn := (lbn \ V)(t := bn \ V \ t))$

**abbreviation** *free-stm8-bd2-cond3*  $V \ t \ b \equiv V(lvl := (lvl \ V)(t := j \ V \ t - 1))$

**abbreviation** *free-stm8-bd2-cond4*  $V \ t \ b \equiv V(bn := (bn \ V)(t := lbn \ V \ t \text{ div } 4))$

**abbreviation** *free-stm8-bd2-cond5*  $V \ t \ b \equiv$

$let \ minf = mem\text{-}pool\text{-}info \ V \ (pool \ b) \ in$

$V(mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info \ V) \ (pool \ b := minf \ (levels := (levels \ minf)$

$[lvl \ V \ t := ((levels \ minf) ! (lvl \ V \ t)) \ (bits := (bits \ ((levels \ minf) ! (lvl \ V \ t)))$

$[bn \ V \ t := FREEING])])$

**lemma** *mempool-free-stm8-atombody-else-blockfit*:

$V \in mp\text{-}free\text{-}precond8\text{-}3 \ t \ b \ \alpha \cap \{\cur = Some \ t\} \implies$

$free\text{-}stm8\text{-}precond2 \ V \ t \ b \in \{block\text{-}fits \ ('mem\text{-}pool\text{-}info \ (pool \ b)) \ ('blk \ t) \ ('lsz \ t)\}$

**apply** (*simp add: block-fits-def block-ptr-def buf-size-def set-bit-def*)

**apply** (*rule subst[where s = max-sz (mem-pool-info V (pool b)) and t = ALIGN4*

```

(max-sz (mem-pool-info V (pool b))))
  apply(simp add: inv-def) using inv-mempool-info-maxsz-align4[rule-format, of
V pool b] apply metis

  apply(subgoal-tac length (bits ((levels (mem-pool-info V (pool b))) ! level b)) =
(n-max (mem-pool-info V (pool b))) * 4 ^ (level b))
  prefer 2 apply(simp add: inv-def inv-mempool-info-def Let-def)

  apply(subgoal-tac max-sz (mem-pool-info V (pool b)) mod 4 ^ lvl V t = 0)
  prefer 2 apply(subgoal-tac ∃ n. max-sz (mem-pool-info V (pool b)) = (4 * n)
* (4 ^ n-levels (mem-pool-info V (pool b))))
  prefer 2 apply(simp add: inv-def) using inv-mempool-info-def[rule-format,
of V] apply meson
  apply(subgoal-tac length (levels (mem-pool-info V (pool b))) = n-levels
(mem-pool-info V (pool b)))
  prefer 2 apply(simp add: inv-def inv-mempool-info-def) apply metis
  apply(simp add: inv-def inv-mempool-info-def)
  using ge-pow-mod-0[of lvl V t n-levels (mem-pool-info V (pool b))]
  apply (metis add-diff-inverse-nat add-lessD1 ge-pow-mod-0 le-antisym nat-less-le)

  apply(subgoal-tac block b div 4 ^ (level b - lvl V t) < n-max (mem-pool-info V
(pool b)) * 4 ^ lvl V t)
  prefer 2 apply (metis (no-types, lifting) add-lessD1 inv-mempool-info-def
invariant.inv-def le-Suc-ex)

  apply(rule block-fits0-h1[of max-sz (mem-pool-info V (pool b)) 4 ^ lvl V t
block b div 4 ^ (level b - lvl V t) n-max (mem-pool-info V (pool b))])
  apply blast apply blast
done

lemma mempool-free-stm8-atombody-else-inv-mempool-info:
inv-mempool-info V ⇒
  inv-mempool-info
    (V (freeing-node := (freeing-node V)(t := None),
      mem-pool-info := (set-bit-free (mem-pool-info V) (pool b) (lvl V t) (block
b div 4 ^ (level b - lvl V t)))
      (pool b := append-free-list (set-bit-free (mem-pool-info V) (pool b) (lvl V
t) (block b div 4 ^ (level b - lvl V t)) (pool b)) (lvl V t)
      (block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz
(mem-pool-info V (pool b))) div 4 ^ lvl V t) (block b div 4 ^ (level b - lvl V t)))),
      free-block-r := (free-block-r V)(t := False)))
  apply(simp add: inv-mempool-info-def append-free-list-def set-bit-def) apply clarify
  apply(rule conjI) apply meson
  apply(rule conjI) apply meson
  apply(rule conjI) apply meson
  apply(rule conjI) apply meson
  apply(rule conjI) apply meson
  apply clarify

```

**apply**(*subgoal-tac* ( $\forall i < \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)))$ ).  
 $\text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! i)) = n\text{-max}$   
 $(\text{mem-pool-info } V \text{ (pool } b)) * 4 ^ i$ ))  
**prefer 2 apply**(*simp add:Let-def*)  
**apply**(*case-tac*  $i = \text{lvl } V t$ )  
**by auto**

**lemma** *mempool-free-stm8-atombody-else-inv-bitmap-freelist:*  
 $\text{inv-mempool-info } V \wedge \text{inv-bitmap-freelist } V \wedge \text{inv-aux-vars } V \implies$   
 $\text{level } b < \text{length } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))) \implies$   
 $\text{block } b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! \text{level } b)) \implies$   
 $\text{block } b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t) < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V t)) \implies$   
 $\text{lvl } V t \leq \text{level } b \implies$   
 $\text{freeing-node } V t = \text{Some } \text{blka} \implies$   
 $\text{pool } \text{blka} = \text{pool } b \implies$   
 $\text{level } \text{blka} = \text{lvl } V t \implies$   
 $\text{block } \text{blka} = \text{block } b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t) \implies$   
 $\text{inv-bitmap-freelist}$   
 $(V(\text{freeing-node} := (\text{freeing-node } V)(t := \text{None}),$   
 $\text{mem-pool-info} := (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V t) (\text{block } b$   
 $\text{div } 4 ^ (\text{level } b - \text{lvl } V t)))$   
 $(\text{pool } b := \text{append-free-list } (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (\text{lvl } V$   
 $t) (\text{block } b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t)) (\text{pool } b)) (\text{lvl } V t)$   
 $(\text{block-ptr } (\text{mem-pool-info } V \text{ (pool } b)) (\text{ALIGN4 } (\text{max-sz}$   
 $(\text{mem-pool-info } V \text{ (pool } b)) \text{ div } 4 ^ \text{lvl } V t) (\text{block } b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t))))),$   
 $\text{free-block-r} := (\text{free-block-r } V)(t := \text{False}))$ )  
**apply**(*simp add:inv-bitmap-freelist-def append-free-list-def set-bit-def block-ptr-def*)  
**apply clarify**  
**apply**(*simp add:Let-def*)  
**apply**(*rule subst[where s=max-sz (mem-pool-info V (pool b)) and t=ALIGN4*  
 $(\text{max-sz } (\text{mem-pool-info } V \text{ (pool } b)))]$ )  
**apply** (*metis inv-mempool-info-maxsz-align4*)

**apply**(*rule conjI*) **apply clarify** **apply**(*rename-tac ii jj*)  
**apply**(*case-tac*  $ii \neq \text{lvl } V t$ ) **apply force**  
**apply**(*case-tac*  $jj = \text{block } b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t)$ )  
**apply clarsimp**

**apply**(*subgoal-tac*  $\text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! ii) ! jj = \text{bits}$   
 $(\text{levels } (\text{mem-pool-info } V \text{ (pool } b))$   
 $[\text{lvl } V t := (\text{levels } (\text{mem-pool-info } V \text{ (pool } b))$   
 $[\text{lvl } V t := (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V t)$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V t)[\text{block}$   
 $b \text{ div } 4 ^ (\text{level } b - \text{lvl } V t) := \text{FREE}]] ! \text{lvl } V t)$   
 $(\text{free-list} := \text{free-list } (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) [\text{lvl } V t$   
 $:= (\text{levels } (\text{mem-pool-info } V \text{ (pool } b)) ! \text{lvl } V t)$



$$t)[block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t) := FREE]] ! lvl\ V\ t) @$$

$$[buf\ (mem\ pool\ info\ V\ (pool\ b)) + ALIGN4\ (max\ sz\ (mem\ pool\ info\ V\ (pool\ b))) \div 4 \wedge lvl\ V\ t * (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))]] !$$

$$ii) ! jj)$$
**prefer 2 apply fastforce**
**apply**(subgoal-tac length (bits (levels (mem-pool-info V (pool b))!ii)) = length (bits (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t) (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] ! lvl V t) (free-list := free-list (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t) (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] ! lvl V t) @ [buf (mem-pool-info V (pool b)) + max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t \* (block b div 4 ^ (level b - lvl V t))]] ! ii)))
**prefer 2 apply fastforce**
**apply**(subgoal-tac free-list (levels (mem-pool-info V (pool b)) ! ii) @ [buf (mem-pool-info V (pool b)) + max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t \* (block b div 4 ^ (level b - lvl V t))] = free-list (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t) (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] ! lvl V t) (free-list := free-list (levels (mem-pool-info V (pool b)) [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t) (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] ! lvl V t) @ [buf (mem-pool-info V (pool b)) + max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t \* (block b div 4 ^ (level b - lvl V t))]] ! ii))
**prefer 2 apply clarsimp**
**apply**(case-tac bits (levels (mem-pool-info V (pool b)) ! ii) ! jj = FREE)
**apply**(subgoal-tac (buf (mem-pool-info V (pool b)) + jj \* (max-sz (mem-pool-info V (pool b)) div 4 ^ ii) \in set (free-list (levels (mem-pool-info V (pool b)) ! ii))))
**prefer 2 apply simp**
**apply clarsimp**
**apply**(subgoal-tac (buf (mem-pool-info V (pool b)) + jj \* (max-sz (mem-pool-info V (pool b)) div 4 ^ ii) \notin set (free-list (levels (mem-pool-info V (pool b)) ! ii))))
**prefer 2 apply simp**
**apply**(subgoal-tac buf (mem-pool-info V (pool b)) + max-sz (mem-pool-info

```

V (pool b)) div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))
    ≠ buf (mem-pool-info V (pool b)) + jj * (max-sz (mem-pool-info
V (pool b)) div 4 ^ ii))
    prefer 2 apply(subgoal-tac max-sz (mem-pool-info V (pool b)) div 4 ^ lvl
V t > 0)

    prefer 2 apply(simp add:inv-mempool-info-def Let-def)
    apply(subgoal-tac ∃ n>NULL. max-sz (mem-pool-info V (pool
b)) = 4 * n * 4 ^ n-levels (mem-pool-info V (pool b)))
    prefer 2 apply auto[1]
    apply(subgoal-tac lvl V t < n-levels (mem-pool-info V (pool b)))
    prefer 2 apply auto[1]
    apply(metis divisors-zero ge-pow-mod-0 grOI mod0-div-self
mult-0-right power-not-zero zero-neq-numeral)
    apply auto[1]

    apply(subgoal-tac buf (mem-pool-info V (pool b)) + jj * (max-sz (mem-pool-info
V (pool b)) div 4 ^ ii)
    ∉ set (free-list (levels (mem-pool-info V (pool b)) ! ii) @
    [buf (mem-pool-info V (pool b)) +
    max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t
    * (block b div 4 ^ (level b - lvl V t))]))
    prefer 2 apply auto[1]

    apply auto[1]

apply(rule conjI)
apply clarify apply(rename-tac ii jj)
    apply(case-tac ii ≠ lvl V t) apply force
    apply(subgoal-tac free-list (levels (mem-pool-info V (pool b)) ! ii) @
    [buf (mem-pool-info V (pool b)) +
    max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t *
    (block b div 4 ^ (level b - lvl V t))] = free-list (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) !
    lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] !
    lvl V t)
    (free-list := free-list (levels (mem-pool-info V (pool
b)) [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool
b)) ! lvl V t)[block b div 4 ^ (level b - lvl V t) := FREE]] ! lvl V t) @
    [buf (mem-pool-info V (pool b)) +
    max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t
    * (block b div 4 ^ (level b - lvl V t))]) ! ii))
    prefer 2 apply clarsimp
    apply(case-tac jj < length (free-list (levels (mem-pool-info V (pool b)) ! ii)))
    apply(simp add: nth-append)
    apply(case-tac jj = length (free-list (levels (mem-pool-info V (pool b)) ! ii)))

```

**apply**(*subgoal-tac* (*free-list* (*levels* (*mem-pool-info* *V* (*pool b*)) ! *ii*) @  
 [buf (*mem-pool-info* *V* (*pool b*)) +  
 max-sz (*mem-pool-info* *V* (*pool b*)) div 4 ^ lvl *V t* \* (*block*  
*b* div 4 ^ (*level b* - lvl *V t*))] ! *jj*  
 = buf (*mem-pool-info* *V* (*pool b*)) +  
 max-sz (*mem-pool-info* *V* (*pool b*)) div 4 ^ lvl *V t* \* (*block*  
*b* div 4 ^ (*level b* - lvl *V t*)))  
**prefer** 2 **apply** *clarsimp*  
**apply**(*subgoal-tac* *block b* div 4 ^ (*level b* - lvl *V t*) < *n-max* (*mem-pool-info*  
*V* (*pool b*)) \* 4 ^ *ii*)  
**prefer** 2 **apply** (*metis inv-mempool-info-def*)  
**apply** *auto*[1]  
**apply** *auto*[1]

**apply**(*subgoal-tac* *distinct* (*free-list* (*levels* (*mem-pool-info* *V* (*pool b*)) ! *i*)))  
**prefer** 2 **apply** *auto*[1] **apply**(*rename-tac* *ii*)  
**apply**(*case-tac* *ii* ≠ lvl *V t*) **apply** *force*  
**apply**(*subgoal-tac* *free-list* (*levels* (*mem-pool-info* *V* (*pool b*)) ! *ii*) @  
 [buf (*mem-pool-info* *V* (*pool b*)) +  
 max-sz (*mem-pool-info* *V* (*pool b*)) div 4 ^ lvl *V t* \*  
 (*block b* div 4 ^ (*level b* - lvl *V t*))] = *free-list* (*levels* (*mem-pool-info* *V* (*pool b*))  
 [lvl *V t* := (*levels* (*mem-pool-info* *V* (*pool b*))  
 [lvl *V t* := (*levels* (*mem-pool-info* *V* (*pool b*)) ! lvl *V t*)  
 (bits := bits (*levels* (*mem-pool-info* *V* (*pool b*)) ! lvl *V*  
*t*)[*block b* div 4 ^ (*level b* - lvl *V t*) := FREE]] !  
 lvl *V t*)  
 (free-list := *free-list* (*levels* (*mem-pool-info* *V* (*pool b*))  
 [lvl *V t* := (*levels* (*mem-pool-info* *V* (*pool b*)) ! lvl *V t*)  
 (bits := bits (*levels* (*mem-pool-info* *V* (*pool b*)) ! lvl  
*V t*)[*block b* div 4 ^ (*level b* - lvl *V t*) := FREE]]] ! lvl *V t*) @  
 [buf (*mem-pool-info* *V* (*pool b*)) +  
 max-sz (*mem-pool-info* *V* (*pool b*)) div 4 ^ lvl *V t* \*  
 (*block b* div 4 ^ (*level b* - lvl *V t*))] ! *ii*))  
**prefer** 2 **apply** *clarsimp*

**apply**(*subgoal-tac* buf (*mem-pool-info* *V* (*pool b*)) +  
 max-sz (*mem-pool-info* *V* (*pool b*)) div 4 ^ lvl *V t* \* (*block b* div  
 4 ^ (*level b* - lvl *V t*))  
 ∉ *set* (*free-list* (*levels* (*mem-pool-info* *V* (*pool b*)) ! *ii*)))  
**prefer** 2 **apply**(*subgoal-tac* *get-bit* (*mem-pool-info* *V*) (*pool b*) (lvl *V t*) (*block*  
*b* div 4 ^ (*level b* - lvl *V t*)) = *FREEING*)  
**prefer** 2 **apply**(*simp add:inv-aux-vars-def*) **apply** *metis*  
**apply** (*metis BlockState.distinct*(15) *semiring-normalization-rules*(7))  
**apply** *auto*  
**done**

**lemma**  
*pool b* ∈ *mem-pools* *V* ⇒

$lvl\ V\ t \leq level\ b \implies$   
 $level\ b < length\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) \implies$   
 $block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t) < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) \ !\ lvl\ V\ t) \implies$   
 $V2 = V(\text{freeing-node} := (\text{freeing-node}\ V)(t := None),$   
 $mem\text{-}pool\text{-}info := (\text{set-bit-free}\ (mem\text{-}pool\text{-}info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t)))$   
 $(pool\ b := \text{append-free-list}\ (\text{set-bit-free}\ (mem\text{-}pool\text{-}info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))\ (pool\ b))\ (lvl\ V\ t)$   
 $(block\text{-ptr}\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ (ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)))\ div\ 4 \wedge lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))))),$   
 $free\text{-}block\text{-}r := (\text{free-block-r}\ V)(t := False)) \implies$   
 $\exists\ lv\ bl. bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b)) \ !\ lv) = bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)) \ !\ lv)\ [bl := FREE]$   
 $\wedge (\forall\ lv'. lv \neq lv' \longrightarrow bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b)) \ !\ lv') = bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)) \ !\ lv'))$   
**apply**(simp add:append-free-list-def set-bit-def block-ptr-def)  
**apply**(subgoal-tac bits (levels (mem-pool-info V2 (pool b)) ! lvl V t) = bits (levels (mem-pool-info V (pool b)) ! lvl V t) [block b div 4 ^ (level b - lvl V t) := FREE])  
**prefer 2 apply auto[1]**  
**apply**(subgoal-tac  $\forall\ lv'. lvl\ V\ t \neq lv' \longrightarrow bits\ (levels\ (mem\text{-}pool\text{-}info\ V2\ (pool\ b)) \ !\ lv') = bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)) \ !\ lv')$ )  
**prefer 2 apply clarify apply auto[1]**  
**apply**(rule exI[where x=lvl V t],auto)  
**done**

**lemma** mempool-free-stm8-atombody-else-inv-bitmap:

$inv\text{-}bitmap\ V \wedge inv\text{-}aux\text{-}vars\ V \implies$   
 $pool\ b \in mem\text{-}pools\ V \implies$   
 $level\ b < length\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))) \implies$   
 $block\ b < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)) \ !\ level\ b)) \implies$   
 $block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t) < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)) \ !\ lvl\ V\ t)) \implies$   
 $lvl\ V\ t \leq level\ b \implies$   
 $freeing\text{-}node\ V\ t = Some\ blka \implies$   
 $pool\ blka = pool\ b \implies$   
 $level\ blka = lvl\ V\ t \implies$   
 $block\ blka = block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t) \implies$   
 $V2 = V(\text{freeing-node} := (\text{freeing-node}\ V)(t := None),$   
 $mem\text{-}pool\text{-}info := (\text{set-bit-free}\ (mem\text{-}pool\text{-}info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t)))$   
 $(pool\ b := \text{append-free-list}\ (\text{set-bit-free}\ (mem\text{-}pool\text{-}info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))\ (pool\ b))\ (lvl\ V\ t)$   
 $(block\text{-ptr}\ (mem\text{-}pool\text{-}info\ V\ (pool\ b))\ (ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ V\ (pool\ b)))\ div\ 4 \wedge lvl\ V\ t)\ (block\ b\ div\ 4 \wedge (level\ b - lvl\ V\ t))))),$   
 $free\text{-}block\text{-}r := (\text{free-block-r}\ V)(t := False)) \implies$   
 $inv\text{-}bitmap\ V2$   
**apply**(unfold inv-bitmap-def) **apply** clarify

**apply**(*case-tac*  $p = \text{pool } b$ )  
**apply**(*subgoal-tac*  $\exists lv \ bl. \text{ bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv) ! bl = \text{FREEING}$   
 $\wedge \text{ bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! lv) = \text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv) [bl := \text{FREE}]$   
 $\wedge (\forall lv'. lv \neq lv' \longrightarrow \text{ bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! lv') = \text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv') )$ )  
**prefer 2 apply**(*simp add:append-free-list-def set-bit-def block-ptr-def*)  
**apply**(*subgoal-tac*  $\text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv \ V \ t) ! (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t)) = \text{FREEING}$ )  
**prefer 2 apply**(*simp add:inv-aux-vars-def*) **apply** *metis*  
**apply**(*subgoal-tac*  $\text{bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! lv \ V \ t) = \text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv \ V \ t) [\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t) := \text{FREE}]$ )  
**prefer 2 apply** *auto*[1]  
**apply**(*subgoal-tac*  $\forall lv'. lv \ V \ t \neq lv' \longrightarrow \text{ bits } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)) ! lv') = \text{bits } (\text{levels } (\text{mem-pool-info } V (\text{pool } b)) ! lv')$ )  
**prefer 2 apply** *clarify* **apply** *auto*[1]  
**apply**(*rule exI*[**where**  $x = lv \ V \ t$ ]) **apply** *auto*[1]  
  
**apply**(*subgoal-tac*  $\text{length } (\text{levels } (\text{mem-pool-info } V (\text{pool } b))) = \text{length } (\text{levels } (\text{mem-pool-info } V2 (\text{pool } b)))$ )  
**prefer 2 apply**(*simp add:append-free-list-def set-bit-def block-ptr-def*)  
  
**apply**(*subgoal-tac*  $\text{inv-bitmap-mp } V (\text{pool } b)$ ) **prefer 2 apply**(*simp add:inv-bitmap-def*)  
  
**apply**(*rule subst*[**where**  $s = V2$  **and**  $t = V(\text{freeing-node} := (\text{freeing-node } V)(t := \text{None}),$   
 $\text{mem-pool-info} := (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t)))$   
 $(\text{pool } b := \text{append-free-list } (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t)) (\text{pool } b)) (lv \ V \ t)$   
 $(\text{block-ptr } (\text{mem-pool-info } V (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } V (\text{pool } b))) \text{ div } 4 \wedge lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t))))$ ,  
 $\text{free-block-r} := (\text{free-block-r } V)(t := \text{False})$ ]) **apply** *fast*  
  
**using** *inv-bitmap-freeing2free*[*of*  $V \text{ pool } b \ V2$ ] **apply** *fast*  
  
**apply**(*subgoal-tac*  $\text{mem-pool-info } V \ p = \text{mem-pool-info}$   
 $(V(\text{freeing-node} := (\text{freeing-node } V)(t := \text{None}),$   
 $\text{mem-pool-info} := (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t)))$   
 $(\text{pool } b := \text{append-free-list } (\text{set-bit-free } (\text{mem-pool-info } V) (\text{pool } b) (lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t)) (\text{pool } b)) (lv \ V \ t)$   
 $(\text{block-ptr } (\text{mem-pool-info } V (\text{pool } b)) (\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } V (\text{pool } b))) \text{ div } 4 \wedge lv \ V \ t) (\text{block } b \text{ div } 4 \wedge (\text{level } b - lv \ V \ t))))$ ,  
 $\text{free-block-r} := (\text{free-block-r } V)(t := \text{False})$ )  
**apply** *fast*

```

      free-block-r := (free-block-r V)(t := False))) p)
    prefer 2 apply(simp add:append-free-list-def set-bit-def block-ptr-def)
    apply(subgoal-tac mem-pools V = mem-pools (V(|freeing-node := (freeing-node
V)(t := None),
      mem-pool-info := (set-bit-free (mem-pool-info V) (pool b)
(lvl V t) (block b div 4 ^ (level b - lvl V t)))
      (pool b := append-free-list (set-bit-free (mem-pool-info V)
(pool b) (lvl V t) (block b div 4 ^ (level b - lvl V t)) (pool b)) (lvl V t)
      (block-ptr (mem-pool-info V (pool b)) (ALIGN4
(max-sz (mem-pool-info V (pool b))) div 4 ^ lvl V t) (block b div 4 ^ (level b - lvl
V t))))),
      free-block-r := (free-block-r V)(t := False)))
    prefer 2 apply(simp add:append-free-list-def set-bit-def block-ptr-def)
    by (smt BlockState.distinct(13))

```

**lemma** mempool-free-stm8-atombody-else-inv-aux-vars:

```

inv-mempool-info V ∧ inv-aux-vars V ⇒
  allocating-node V t = None ⇒
  pool b ∈ mem-pools V ⇒
  level b < length (levels (mem-pool-info V (pool b))) ⇒
  block b < length (bits (levels (mem-pool-info V (pool b)) ! level b)) ⇒
  block b div 4 ^ (level b - lvl V t) < length (bits (levels (mem-pool-info V (pool
b)) ! lvl V t)) ⇒
  bn V t = block b div 4 ^ (level b - lvl V t) ⇒
  lvl V t ≤ level b ⇒
  freeing-node V t = Some blka ⇒
  pool blka = pool b ⇒
  level blka = lvl V t ⇒
  block blka = block b div 4 ^ (level b - lvl V t) ⇒
  inv-aux-vars
  (V(|freeing-node := (freeing-node V)(t := None),
    mem-pool-info := (set-bit-free (mem-pool-info V) (pool b) (lvl V t) (block b
div 4 ^ (level b - lvl V t)))
    (pool b := append-free-list (set-bit-free (mem-pool-info V) (pool b) (lvl V
t) (block b div 4 ^ (level b - lvl V t)) (pool b)) (lvl V t)
    (block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz
(mem-pool-info V (pool b))) div 4 ^ lvl V t) (block b div 4 ^ (level b - lvl V t))))),
    free-block-r := (free-block-r V)(t := False)))
  apply(simp add:inv-aux-vars-def append-free-list-def set-bit-def block-ptr-def) ap-
  ply clarify
  apply(rule subst[where s=max-sz (mem-pool-info V (pool b)) and t=ALIGN4
(max-sz (mem-pool-info V (pool b)))])
  apply (metis inv-mempool-info-maxsz-align4)

```

```

  apply(rule conjI) apply clarify
  apply(subgoal-tac ¬(pool n = pool blka ∧ level n = level blka ∧ block n = block
blka))

```

```

    prefer 2 apply blast
    apply(case-tac pool n ≠ pool blk a) apply auto[1]
    apply(case-tac level n ≠ level blk a) apply (metis (no-types, lifting) nth-list-update-neq)

    apply(case-tac block n ≠ block blk a)
    apply(subgoal-tac bits (levels (mem-pool-info V (pool b))
      [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
        (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b
div 4 ^ (level b - lvl V t) := FREE],
      free-list :=
        free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
        [buf (mem-pool-info V (pool b)) +
          max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t * (block b
div 4 ^ (level b - lvl V t))] !
        level n) !
      block n = bits (levels (mem-pool-info V (pool b)) ! level n) ! block n)
    prefer 2 apply(subgoal-tac bits (levels (mem-pool-info V (pool b))
      [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
        (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b
div 4 ^ (level b - lvl V t) := FREE],
      free-list :=
        free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
        [buf (mem-pool-info V (pool b)) +
          max-sz (mem-pool-info V (pool b)) div 4 ^ lvl V t * (block b
div 4 ^ (level b - lvl V t))] !
      level n = bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b
div 4 ^ (level b - lvl V t) := FREE])
    prefer 2 apply auto[1] apply auto[1]
    apply metis
    apply fast

    apply(rule conjI) apply clarify
    apply(rule conjI) apply clarify
    apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! level n) ! block n
= FREEING
      ∧ (lvl V t ≠ level n ∨ block b div 4 ^ (level b - lvl V t) ≠ block
n))
    prefer 2 apply(case-tac lvl V t = level n) apply(case-tac block b div 4 ^
(level b - lvl V t) = block n)
    apply clarsimp apply clarsimp apply clarsimp
    apply(subgoal-tac mem-block-addr-valid V n)
    prefer 2 apply(simp add:mem-block-addr-valid-def)
    apply(subgoal-tac blk a ≠ n)
    prefer 2 apply metis
    apply (metis option.inject)
    apply clarify
    apply(subgoal-tac mem-block-addr-valid V n)
    prefer 2 apply(simp add:mem-block-addr-valid-def)

```

```

apply (metis option.inject)

apply(rule conjI) apply clarify
  apply(subgoal-tac get-bit-s V (pool n) (level n) (block n) = ALLOCATING)
prefer 2 apply blast
  apply(case-tac lvl V t = level n) apply(case-tac block b div 4 ^ (level b - lvl V t) = block n)
  apply metis apply clarsimp apply clarsimp

apply clarify
apply(rule conjI) apply clarify
  apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! level n) ! block n = ALLOCATING)
  prefer 2 apply(case-tac lvl V t = level n) apply(case-tac block b div 4 ^ (level b - lvl V t) = block n)
  apply clarsimp apply clarsimp apply clarsimp
  apply(subgoal-tac mem-block-addr-valid V n)
  prefer 2 apply(simp add:mem-block-addr-valid-def)
  apply metis

apply clarify
apply(subgoal-tac mem-block-addr-valid V n)
  prefer 2 apply(simp add:mem-block-addr-valid-def)
apply metis
done

lemma mempool-free-stm8-atombody-else-inv-bitmap0:
inv-mempool-info V ∧ inv-bitmap0 V ⇒
  allocating-node V t = None ⇒
    pool b ∈ mem-pools V ⇒
      level b < length (levels (mem-pool-info V (pool b))) ⇒
        block b < length (bits (levels (mem-pool-info V (pool b)) ! level b)) ⇒
          block b div 4 ^ (level b - lvl V t) < length (bits (levels (mem-pool-info V (pool b)) ! lvl V t)) ⇒
            bn V t = block b div 4 ^ (level b - lvl V t) ⇒
              lvl V t ≤ level b ⇒
                freeing-node V t = Some blka ⇒
                  pool blka = pool b ⇒
                    level blka = lvl V t ⇒
                      block blka = block b div 4 ^ (level b - lvl V t) ⇒
                        inv-bitmap0
                        (V(freeing-node := (freeing-node V)(t := None),
                          mem-pool-info := (set-bit-free (mem-pool-info V) (pool b) (lvl V t) (block b div 4 ^ (level b - lvl V t)))
                          (pool b := append-free-list (set-bit-free (mem-pool-info V) (pool b) (lvl V t) (block b div 4 ^ (level b - lvl V t)) (pool b)) (lvl V t)
                          (block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz (mem-pool-info V (pool b))) div 4 ^ lvl V t) (block b div 4 ^ (level b - lvl V t)))),

```



```

    free-block-r := (free-block-r V)(t := False))
apply(simp add:inv-bitmap0-def inv-mempool-info-def append-free-list-def set-bit-def
block-ptr-def ALIGN4-def Let-def) apply clarsimp
apply(subgoal-tac get-bit-s V (pool b) 0 i ≠ NOEXIST) prefer 2
    apply(subgoal-tac length (bits (levels (mem-pool-info V (pool b))
[lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
(bits := bits (levels (mem-pool-info V (pool b)) ! lvl V
t)[block b div 4 ^ (level b - lvl V t) := FREE],
free-list :=
free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
[buf (mem-pool-info V (pool b)) +
(max-sz (mem-pool-info V (pool b)) + 3) div 4 * 4
div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))] !
0)) = length (bits (levels (mem-pool-info V (pool b)) ! 0)))
prefer 2
    apply(case-tac lvl V t = 0) apply clarsimp
    apply(rule subst[where s=bits (levels (mem-pool-info V (pool b)) ! 0)[block b
div 4 ^ (level b - lvl V t) := FREE]
and t=bits (levels (mem-pool-info V (pool b))
[0 := (levels (mem-pool-info V (pool b)) ! 0)
(bits := bits (levels (mem-pool-info V (pool b)) ! NULL)[block
b div 4 ^ level b := FREE],
free-list :=
free-list (levels (mem-pool-info V (pool b)) ! NULL) @
[buf (mem-pool-info V (pool b)) + (max-sz (mem-pool-info
V (pool b)) + 3) div 4 * 4 * (block b div 4 ^ level b)] !
0)]) apply auto[1] apply auto[1]
    apply auto[1]
    apply auto[1]

apply(case-tac lvl V t = 0)
    apply(case-tac i = block b div 4 ^ (level b - lvl V t))
apply auto[1] apply auto[1] apply auto[1]
done

```

**lemma** mempool-free-stm8-atombody-else-inv-bitmapn:

```

inv-mempool-info V ∧ inv-bitmapn V ⇒
  allocating-node V t = None ⇒
  pool b ∈ mem-pools V ⇒
  level b < length (levels (mem-pool-info V (pool b))) ⇒
  block b < length (bits (levels (mem-pool-info V (pool b)) ! level b)) ⇒
  block b div 4 ^ (level b - lvl V t) < length (bits (levels (mem-pool-info V (pool
b)) ! lvl V t)) ⇒
  bn V t = block b div 4 ^ (level b - lvl V t) ⇒
  lvl V t ≤ level b ⇒
  freeing-node V t = Some blka ⇒
  pool blka = pool b ⇒
  level blka = lvl V t ⇒
  block blka = block b div 4 ^ (level b - lvl V t) ⇒

```

```

inv-bitmapn
  (V (freeing-node := (freeing-node V) (t := None),
    mem-pool-info := (set-bit-free (mem-pool-info V) (pool b) (lvl V t) (block b
div 4 ^ (level b - lvl V t)))
    (pool b := append-free-list (set-bit-free (mem-pool-info V) (pool b) (lvl V
t) (block b div 4 ^ (level b - lvl V t)) (pool b)) (lvl V t)
    (block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz
(mem-pool-info V (pool b))) div 4 ^ lvl V t) (block b div 4 ^ (level b - lvl V t))),
    free-block-r := (free-block-r V) (t := False)))
apply (simp add: inv-bitmapn-def inv-mempool-info-def append-free-list-def set-bit-def
block-ptr-def ALIGN4-def Let-def) apply clarsimp

apply (subgoal-tac get-bit-s V (pool b) (length (levels (mem-pool-info V (pool b)))
- Suc 0) i ≠ DIVIDED) prefer 2
  apply (subgoal-tac length (bits (levels (mem-pool-info V (pool b))
    [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
    (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V
t) [block b div 4 ^ (level b - lvl V t) := FREE],
    free-list :=
      free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
        [buf (mem-pool-info V (pool b)) +
          (max-sz (mem-pool-info V (pool b)) + 3) div 4 * 4
div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))] !
          (length (levels (mem-pool-info V (pool b))) - Suc 0)))
        = length (bits (levels (mem-pool-info V (pool b)) ! (length (levels
(mem-pool-info V (pool b))) - Suc 0)))) prefer 2
    apply (case-tac lvl V t = (length (levels (mem-pool-info V (pool b))) - Suc 0))
apply clarsimp
    apply (rule subst[where s=bits (levels (mem-pool-info V (pool b)) ! (length
(levels (mem-pool-info V (pool b))) - Suc 0))
      [block b div 4 ^ (level b - lvl V t) := FREE]
      and t=bits (levels (mem-pool-info V (pool b))
        [(length (levels (mem-pool-info V (pool b))) - Suc 0) :=
          (levels (mem-pool-info V (pool b)) ! (length (levels (mem-pool-info
V (pool b))) - Suc 0))
        (bits := bits (levels (mem-pool-info V (pool b)) ! NULL) [block
b div 4 ^ level b := FREE],
        free-list :=
          free-list (levels (mem-pool-info V (pool b)) ! NULL) @
            [buf (mem-pool-info V (pool b)) + (max-sz (mem-pool-info
V (pool b)) + 3) div 4 * 4 * (block b div 4 ^ level b)] !
            (length (levels (mem-pool-info V (pool b))) - Suc 0))]) apply
auto[1] apply auto[1]
    apply auto[1]

apply (case-tac lvl V t = (length (levels (mem-pool-info V (pool b))) - Suc 0))
  apply (case-tac i = block b div 4 ^ (level b - lvl V t))
apply auto[1] apply auto[1] apply auto[1]
done

```

**lemma** *mempool-free-stm8-atombody-else-inv-bitmap-not4free*:

*lvl V t = NULL*  $\vee$

$\neg$  *partner-bits* (*set-bit-free* (*mem-pool-info V*) (*pool b*) (*lvl V t*) (*block b div 4*  $\wedge$  (*level b* - *lvl V t*)) (*pool b*)) (*lvl V t*)

(*block b div 4*  $\wedge$  (*level b* - *lvl V t*))  $\implies$

*inv-mempool-info V*  $\wedge$  *inv-bitmap-not4free V*  $\implies$

*allocating-node V t = None*  $\implies$

*pool b*  $\in$  *mem-pools V*  $\implies$

*level b* < *length* (*levels* (*mem-pool-info V* (*pool b*)))  $\implies$

*block b* < *length* (*bits* (*levels* (*mem-pool-info V* (*pool b*)) ! *level b*))  $\implies$

*block b div 4*  $\wedge$  (*level b* - *lvl V t*) < *length* (*bits* (*levels* (*mem-pool-info V* (*pool b*)) ! *lvl V t*))  $\implies$

*bn V t = block b div 4*  $\wedge$  (*level b* - *lvl V t*)  $\implies$

*lvl V t*  $\leq$  *level b*  $\implies$

*freeing-node V t = Some blka*  $\implies$

*pool blka = pool b*  $\implies$

*level blka = lvl V t*  $\implies$

*block blka = block b div 4*  $\wedge$  (*level b* - *lvl V t*)  $\implies$

*inv-bitmap-not4free*

(*V* (*freeing-node* := (*freeing-node V*)(*t* := *None*),

*mem-pool-info* := (*set-bit-free* (*mem-pool-info V*) (*pool b*) (*lvl V t*) (*block b div 4*  $\wedge$  (*level b* - *lvl V t*)))

(*pool b* := *append-free-list* (*set-bit-free* (*mem-pool-info V*) (*pool b*) (*lvl V t*) (*block b div 4*  $\wedge$  (*level b* - *lvl V t*)) (*pool b*)) (*lvl V t*)

(*block-ptr* (*mem-pool-info V* (*pool b*)) (*ALIGN4* (*max-sz* (*mem-pool-info V* (*pool b*)) *div 4*  $\wedge$  *lvl V t*) (*block b div 4*  $\wedge$  (*level b* - *lvl V t*))))),

*free-block-r* := (*free-block-r V*)(*t* := *False*)))

**apply**(*simp add:inv-bitmap-not4free-def inv-mempool-info-def append-free-list-def set-bit-def block-ptr-def ALIGN4-def Let-def*)

**apply** *clarsimp*

**apply**(*case-tac lvl V t = 0*)

**apply** *clarsimp*

**apply**(*simp add:partner-bits-def Let-def*) **apply** *auto*[1]

**apply** *clarsimp*

**apply**(*case-tac i = lvl V t*)

**apply**(*simp add:partner-bits-def Let-def*)

**apply** *clarsimp*

**apply**(*case-tac j div 4 = block b div 4*  $\wedge$  (*level b* - *lvl V t*) *div 4*)

**apply** *auto*[1]

**apply** *auto*[1]

**apply**(*simp add:partner-bits-def Let-def*)

**apply** *clarsimp*

**done**

**lemma** *mempool-free-stm8-atombody-else-inv*:

$lvl\ V\ t = NULL \vee$

$\neg partner\_bits\ (set\_bit\_free\ (mem\_pool\_info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4\ \wedge$

$(level\ b - lvl\ V\ t))\ (pool\ b))\ (lvl\ V\ t)$

$(block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t)) \implies$

$inv\ V \implies$

$allocating\_node\ V\ t = None \implies$

$pool\ b \in mem\_pools\ V \implies$

$level\ b < length\ (levels\ (mem\_pool\_info\ V\ (pool\ b))) \implies$

$block\ b < length\ (bits\ (levels\ (mem\_pool\_info\ V\ (pool\ b))\ !\ level\ b)) \implies$

$data\ b = block\_ptr\ (mem\_pool\_info\ V\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ V\ (pool\ b)))\ div\ 4\ \wedge\ level\ b)\ (block\ b) \implies$

$level\ b < length\ (lsizes\ V\ t) \implies$

$\forall ii < length\ (lsizes\ V\ t).\ lsizes\ V\ t\ !\ ii = ALIGN4\ (max\_sz\ (mem\_pool\_info\ V\ (pool\ b)))\ div\ 4\ \wedge\ ii \implies$

$block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t) < length\ (bits\ (levels\ (mem\_pool\_info\ V\ (pool\ b))\ !\ lvl\ V\ t)) \implies$

$bn\ V\ t = block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t) \implies$

$lvl\ V\ t \leq level\ b \implies$

$free\_block\_r\ V\ t \implies$

$lsz\ V\ t = ALIGN4\ (max\_sz\ (mem\_pool\_info\ V\ (pool\ b)))\ div\ 4\ \wedge\ lvl\ V\ t \implies$

$blk\ V\ t = block\_ptr\ (mem\_pool\_info\ V\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ V\ (pool\ b)))\ div\ 4\ \wedge\ lvl\ V\ t)\ (block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t)) \implies$

$cur\ V = Some\ t \implies$

$data\ blka = buf\ (mem\_pool\_info\ V\ (pool\ b)) + block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t)$

$* (max\_sz\ (mem\_pool\_info\ V\ (pool\ b))\ div\ 4\ \wedge\ lvl\ V\ t) \implies$

$block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t) < n\_max\ (mem\_pool\_info\ V\ (pool\ b)) * 4\ \wedge\ lvl\ V\ t \implies$

$freeing\_node\ V\ t = Some\ y \implies$

$pool\ y = pool\ b \implies$

$level\ y = lvl\ V\ t \implies$

$block\ y = block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t) \implies$

$inv$

$(V(|freeing\_node := (freeing\_node\ V)(t := None),$

$mem\_pool\_info := (set\_bit\_free\ (mem\_pool\_info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t)))$

$(pool\ b := append\_free\_list\ (set\_bit\_free\ (mem\_pool\_info\ V)\ (pool\ b)\ (lvl\ V\ t)\ (block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t))\ (pool\ b))\ (lvl\ V\ t)$

$(block\_ptr\ (mem\_pool\_info\ V\ (pool\ b))\ (ALIGN4\ (max\_sz\ (mem\_pool\_info\ V\ (pool\ b)))\ div\ 4\ \wedge\ lvl\ V\ t)\ (block\ b\ div\ 4\ \wedge (level\ b - lvl\ V\ t))))),$

$free\_block\_r := (free\_block\_r\ V)(t := False)|)$

**apply**(*simp add:inv-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def Mem-pool-free-guar-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-thd-waitq-def append-free-list-def set-bit-def*)

**apply** *smt*

**apply**(*rule conjI*) **using** *mempool-free-stm8-atombody-else-inv-mempool-info* **ap-**

**ply** *blast*

**apply**(*rule conjI*) **using** *mempool-free-stm8-atombody-else-inv-bitmap-freelist* **ap-**

```

ply blast
  apply(rule conjI) using mempool-free-stm8-atombody-else-inv-bitmap apply blast
  apply(rule conjI) using mempool-free-stm8-atombody-else-inv-aux-vars apply
blast
  apply(rule conjI) using mempool-free-stm8-atombody-else-inv-bitmap0 apply
blast
  apply(rule conjI) using mempool-free-stm8-atombody-else-inv-bitmapn apply
blast
    using mempool-free-stm8-atombody-else-inv-bitmap-not4free apply
blast
done

```

**lemma** *mp-free-stm8-intI*:

```

{V} ⊆ {'}(free-block-r-update (λ-. 'free-block-r(t := False))) ∈ A} ⇒
  {V} ⊆ {'}(free-block-r-update (λ-. 'free-block-r(t := False))) ∈ B} ⇒
  {V} ⊆ {'}(free-block-r-update (λ-. 'free-block-r(t := False))) ∈ A ∩ B}
by auto

```

**lemma** *mempool-free-stm8-atombody-else-h1*:

```

V ∈ {'}free-block-r t} ∩ mp-free-precond8-3 t b α ∩ {'}cur = Some t} ⇒
  {free-stm8-precond2 V t b} ∩ - {'}NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)} ≠ {} ⇒
  {let V2 = free-stm8-precond2 V t b in V2 {mem-pool-info := (mem-pool-info V2)
(pool b := append-free-list (mem-pool-info V2 (pool b)) (lvl V2 t) (blk
V2 t))}}
  ⊆ {'}(free-block-r-update (λ-. 'free-block-r(t := False)))
  ∈ {'}(Pair V) ∈ Mem-pool-free-guar t} ∩ mp-free-precond8-inv t b 0}
apply(rule mp-free-stm8-intI)

```

```

apply(simp add:Mem-pool-free-guar-def)
apply(rule disjI1) apply(rule conjI)
  apply(simp add:gvars-conf-stable-def gvars-conf-def append-free-list-def set-bit-def
block-ptr-def) apply clarify
  apply(rename-tac ii blk)
  apply(case-tac lvl V t = ii)
  apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div
4 ^ (level b - lvl V t) := FREE] =
    bits (levels (mem-pool-info V (pool b))
[lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
[bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block
b div 4 ^ (level b - lvl V t) := FREE],
free-list := free-list (levels (mem-pool-info V (pool b)) ! lvl
V t) @
[buf (mem-pool-info V (pool b)) + ALIGN4 (max-sz
(mem-pool-info V (pool b)))]
div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))]) !
ii))
  prefer 2 apply auto[1]
  apply (metis length-list-update)

```

```

apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! ii) =
      bits (levels (mem-pool-info V (pool b))
        [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
          (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block
            b div 4 ^ (level b - lvl V t) := FREE],
            free-list := free-list (levels (mem-pool-info V (pool b)) ! lvl
              V t) @
              [buf (mem-pool-info V (pool b)) + ALIGN4 (max-sz
                (mem-pool-info V (pool b)))
                div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))]] !
              ii))
      prefer 2 apply auto[1]
      apply auto[1]
apply(rule conjI) apply clarsimp
      using mempool-free-stm8-atombody-else-inv[of V t b ] apply metis

apply clarify apply(simp add:lvars-nochange-def)

apply(rule mp-free-stm8-intI)
apply clarsimp
apply(rule conjI)
      using mempool-free-stm8-atombody-else-inv[of V t b ] apply metis
apply(rule conjI)
      apply(simp add:append-free-list-def set-bit-def block-ptr-def)
apply(rule conjI)
      apply(simp add:append-free-list-def set-bit-def block-ptr-def)
      apply(case-tac lvl V t = level b)
      apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! lvl V t)[block b div
        4 ^ (level b - lvl V t) := FREE] =
        bits (levels (mem-pool-info V (pool b))
          [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
            (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V
              t)[block b div 4 ^ (level b - lvl V t) := FREE],
              free-list :=
                free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
                [buf (mem-pool-info V (pool b)) +
                  ALIGN4 (max-sz (mem-pool-info V (pool b))) div
                    4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))]] !
                level b))
        prefer 2 apply auto[1]
        apply (metis length-list-update)
        apply(subgoal-tac bits (levels (mem-pool-info V (pool b)) ! level b) =
          bits (levels (mem-pool-info V (pool b))
            [lvl V t := (levels (mem-pool-info V (pool b)) ! lvl V t)
              (bits := bits (levels (mem-pool-info V (pool b)) ! lvl V
                t)[block b div 4 ^ (level b - lvl V t) := FREE],
                free-list :=
                  free-list (levels (mem-pool-info V (pool b)) ! lvl V t) @
                  [buf (mem-pool-info V (pool b)) +
                    ALIGN4 (max-sz (mem-pool-info V (pool b)))
                    div 4 ^ lvl V t * (block b div 4 ^ (level b - lvl V t))]] !
                  level b))
          prefer 2 apply auto[1]
          apply (metis length-list-update)
          apply(rule conjI)
          using mempool-free-stm8-atombody-else-inv[of V t b ] apply metis

```

$ALIGN_4 \ (max\text{-}sz \ (mem\text{-}pool\text{-}info \ V \ (pool \ b))) \ div$   
 $4 \wedge lvl \ V \ t \ * \ (block \ b \ div \ 4 \wedge (level \ b - lvl \ V \ t)) \ !$   
 $level \ b))$

**prefer 2 apply auto[1]**  
**apply metis**  
**apply(rule conjI)**  
**apply(simp add:append-free-list-def set-bit-def block-ptr-def)**  
**apply(rule conjI)**  
**apply(simp add:append-free-list-def set-bit-def block-ptr-def)**  
**apply(simp add:append-free-list-def set-bit-def block-ptr-def)**

**apply clarsimp**  
**done**

**lemma mempool-free-stm8-atombody-else':**

$V \in \{\text{'free-block-r } t\} \cap mp\text{-free-precond8-3 } t \ b \ \alpha \cap \{\text{'cur} = \text{Some } t\} \implies$   
 $\Gamma \vdash_I \text{Some } (IF \text{ block-fits } (\text{'mem-pool-info } (pool \ b)) (\text{'blk } t) (\text{'lsz } t)) \text{ THEN}$   
 $\text{'mem-pool-info} := \text{'mem-pool-info}(pool \ b := \text{append-free-list } (\text{'mem-pool-info}$   
 $(pool \ b)) (\text{'lvl } t) (\text{'blk } t))$   
 $FI;;$   
 $\text{'free-block-r} := \text{'free-block-r } (t := \text{False})$   
 $sat_p \ [\{\text{free-stm8-precond2 } V \ t \ b\} \cap - \ \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info}$   
 $(pool \ b)) (\text{'lvl } t) (\text{'bn } t)\}],$   
 $\{(s, t). s = t\}, \text{UNIV}, \{\text{'(Pair } V) \in \text{Mem-pool-free-guar } t\}$   
 $\cap mp\text{-free-precond8-inv } t \ b \ 0]$   
**apply(case-tac  $\{\text{free-stm8-precond2 } V \ t \ b\} \cap - \ \{\text{NULL} < \text{'lvl } t \wedge \text{partner-bits}$   
 $(\text{'mem-pool-info } (pool \ b)) (\text{'lvl } t) (\text{'bn } t)\} = \{\}$ )**  
**using Emptyprecond[of Some (IF block-fits ( $\text{'mem-pool-info } (pool \ b)) (\text{'blk } t)$   
 $(\text{'lsz } t))$  THEN**  
 $\text{'mem-pool-info} := \text{'mem-pool-info}(pool \ b := \text{append-free-list } (\text{'mem-pool-info}$   
 $(pool \ b)) (\text{'lvl } t) (\text{'blk } t))$   
 $FI;;$   
 $\text{'free-block-r} := \text{'free-block-r } (t := \text{False}) \ \{(s, t). s = t\}$   
 $\text{UNIV } \{\text{'(Pair } V) \in \text{Mem-pool-free-guar } t\} \cap mp\text{-free-precond8-inv } t \ b \ 0]$  **apply**  
**metis**

**apply(rule Seq[where mid={let V2 = free-stm8-precond2 V t b in V2(mem-pool-info**  
 $:= (\text{mem-pool-info } V2)$   
 $(pool \ b := \text{append-free-list } (\text{mem-pool-info } V2 \ (pool \ b)) (\text{lvl } V2 \ t) (\text{blk } V2$   
 $t))\}])$

**apply(rule Cond)**  
**using stable-id2 apply fast**

**apply(rule subst[where s={ } and t={free-stm8-precond2 V t b}  $\cap - \ \{\text{NULL}$   
 $< \text{'lvl } t \wedge \text{partner-bits } (\text{'mem-pool-info } (pool \ b)) (\text{'lvl } t) (\text{'bn } t)\} \cap$   
 $- \ \{\text{block-fits } (\text{'mem-pool-info } (pool \ b)) (\text{'blk } t)$   
 $(\text{'lsz } t)\}])$**   
**using mempool-free-stm8-atombody-else-blockfit[of V t b  $\alpha$ ] apply fast**

**apply**(*rule Basic*) **apply**(*simp add:Let-def*)  
**apply** *auto*[1] **apply** *fast* **using** *stable-id2* **apply** *fast* **using** *stable-id2* **apply**  
*fast*

**apply**(*rule subst*[**where**  $s=\{\}$  **and**  $t=\{\text{free-stm8-precond2 } V \ t \ b\} \cap - \llbracket \text{NULL} < 'lwl \ t \wedge \text{partner-bits } ('mem\text{-pool-info } (pool \ b)) ('lwl \ t) ('bn \ t) \rrbracket \cap$   
 $- \llbracket \text{block-fits } ('mem\text{-pool-info } (pool \ b)) ('blk \ t) ('lsz \ t) \rrbracket$ ]  
**using** *mempool-free-stm8-atombody-else-blockfit*[*of*  $V \ t \ b \ \alpha$ ] **apply** *fast*  
**using** *Emptyprecond* **apply** *blast*

**apply** *fast*

**apply**(*rule Basic*)  
**using** *mempool-free-stm8-atombody-else-h1*[*of*  $V \ t \ b \ \alpha$ ] **apply** *fast*

**apply** *fast*  
**using** *stable-id2* **apply** *fast* **using** *stable-id2* **apply** *fast*  
**done**

**lemma** *mempool-free-stm8-atombody-else*:

$V \in mp\text{-free-precond8-3 } t \ b \ \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \implies$   
 $\Gamma \vdash_I \text{Some } (IF \text{ block-fits } ('mem\text{-pool-info } (pool \ b)) ('blk \ t) ('lsz \ t) \text{ THEN}$   
 $'mem\text{-pool-info} := 'mem\text{-pool-info}(pool \ b := \text{append-free-list } ('mem\text{-pool-info}$   
 $(pool \ b)) ('lwl \ t) ('blk \ t))$   
 $FI;;$   
 $'free\text{-block-r} := 'free\text{-block-r } (t := \text{False})$   
 $\text{sat}_p [\{\text{free-stm8-precond2 } V \ t \ b\} \cap - \llbracket \text{NULL} < 'lwl \ t \wedge \text{partner-bits } ('mem\text{-pool-info}$   
 $(pool \ b)) ('lwl \ t) ('bn \ t) \rrbracket,$   
 $\{(s, t). s = t\}, \text{UNIV}, \llbracket (Pair \ V) \in \text{Mem-pool-free-guar } t \rrbracket$   
 $\cap mp\text{-free-precond8-inv } t \ b \ 0]$   
**apply**(*subgoal-tac*  $V \in \llbracket 'free\text{-block-r } t \rrbracket \cap mp\text{-free-precond8-3 } t \ b \ \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket$ )  
**prefer** 2 **apply**(*subgoal-tac*  $mp\text{-free-precond8-1 } t \ b \ \alpha = \llbracket 'free\text{-block-r } t \rrbracket \cap$   
 $mp\text{-free-precond8-1 } t \ b \ \alpha$ )  
**prefer** 2 **using** *mp-free-precond8-1-imp-free-block-r*[*of*  $t \ b \ \alpha$ ] *Int-absorb1*[*of*  
 $mp\text{-free-precond8-1 } t \ b \ \alpha \llbracket 'free\text{-block-r } t \rrbracket$ ] **apply** *metis*  
**apply** *auto*[1]  
**using** *mempool-free-stm8-atombody-else'*[*of*  $V \ t \ b \ \alpha$ ] **apply** *metis*  
**done**

**lemma** *mempool-free-stm8-atombody-rest-extpost*:

$V \in mp\text{-free-precond8-3 } t \ b \ \alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \implies$   
 $\{\text{free-stm8-precond2 } V \ t \ b\} \cap \llbracket \text{NULL} < 'lwl \ t \wedge \text{partner-bits } ('mem\text{-pool-info}$   
 $(pool \ b)) ('lwl \ t) ('bn \ t) \rrbracket \neq \{\}$   
 $\Gamma \vdash_I \text{Some } ('lwl := 'lwl(t := 'lwl \ t - 1));;$



```

    'bn := 'bn(t := 'bn t div 4);;
    'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) ('lvl t) ('bn t);;
    'freeing-node := 'freeing-node(t ↦ (pool = pool b, level = 'lvl t, block = 'bn
t,
    data = block-ptr ('mem-pool-info (pool b)) (ALIGN4 (max-sz ('mem-pool-info
(pool b))) div 4 ^ 'lvl t) ('bn t)))
    satp [free-stm8-precond3 V t b ∩ {i t = 4}, {(s, t). s = t}, UNIV,
    {(Pair V) ∈ Mem-pool-free-guar t} ∩ (mp-free-precond8-inv t b (α - 1) ∪
mp-free-precond8-inv t b 0)]
apply(rule Conseq[of free-stm8-precond3 V t b ∩ {i t = 4} free-stm8-precond3 V
t b ∩ {i t = 4}
    {(s, t). s = t} {(s, t). s = t} UNIV UNIV
    {(Pair V) ∈ Mem-pool-free-guar t} ∩ mp-free-precond8-inv t b (α -
1)
    {(Pair V) ∈ Mem-pool-free-guar t} ∩ (mp-free-precond8-inv t b (α
- 1) ∪ mp-free-precond8-inv t b 0)
    Some ('lvl := 'lvl(t := 'lvl t - 1));;
    'bn := 'bn(t := 'bn t div 4);;
    'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) ('lvl t) ('bn
t));;
    'freeing-node := 'freeing-node(t ↦ (pool = pool b, level = 'lvl t,
block = 'bn t,
    data = block-ptr ('mem-pool-info (pool b)) (ALIGN4 (max-sz
('mem-pool-info (pool b))) div 4 ^ 'lvl t) ('bn t))))
apply fast apply fast apply fast apply auto[1]
using mempool-free-stm8-atombody-rest apply blast
done

```

**lemma** mempool-free-stm8-atombody-else-extpost:

```

V ∈ mp-free-precond8-3 t b α ∩ {cur = Some t} ⇒
Γ ⊢I Some (IF block-fits ('mem-pool-info (pool b)) ('blk t) ('lsz t) THEN
    'mem-pool-info := 'mem-pool-info(pool b := append-free-list ('mem-pool-info
(pool b)) ('lvl t) ('blk t))
    FI;;
    'free-block-r := 'free-block-r (t := False) )
    satp [{free-stm8-precond2 V t b} ∩ - {NULL < 'lvl t ∧ partner-bits ('mem-pool-info
(pool b)) ('lvl t) ('bn t)},
    {(s, t). s = t}, UNIV, {(Pair V) ∈ Mem-pool-free-guar t}
    ∩ (mp-free-precond8-inv t b (α - 1) ∪ mp-free-precond8-inv t b 0)]
apply(rule Conseq[of {free-stm8-precond2 V t b} ∩ - {NULL < 'lvl t ∧ partner-bits
('mem-pool-info (pool b)) ('lvl t) ('bn t)}
    {free-stm8-precond2 V t b} ∩ - {NULL < 'lvl t ∧ partner-bits
('mem-pool-info (pool b)) ('lvl t) ('bn t)}
    {(s, t). s = t} {(s, t). s = t} UNIV UNIV
    {(Pair V) ∈ Mem-pool-free-guar t} ∩ mp-free-precond8-inv t b 0
    {(Pair V) ∈ Mem-pool-free-guar t} ∩ (mp-free-precond8-inv t b (α
- 1) ∪ mp-free-precond8-inv t b 0)
    Some (IF block-fits ('mem-pool-info (pool b)) ('blk t) ('lsz t) THEN
    'mem-pool-info := 'mem-pool-info(pool b := append-free-list

```

```

('mem-pool-info (pool b)) ('lvl t) ('blk t))
  FI;;
  'free-block-r := 'free-block-r (t := False) )])
apply fast apply fast apply fast apply auto[1]
using mempool-free-stm8-atombody-else apply blast
done

lemma mempool-free-stm8-atombody:
   $\Gamma \vdash_I$  Some ('mem-pool-info := set-bit-free 'mem-pool-info (pool b) ('lvl t) ('bn
  t));
  'freeing-node := 'freeing-node(t := None);
  IF NULL < 'lvl t  $\wedge$  partner-bits ('mem-pool-info (pool b)) ('lvl t) ('bn t)
  THEN 'i := 'i(t := 0);
  WHILE 'i t < 4
  DO 'bb := 'bb(t := 'bn t div 4 * 4 + 'i t);
  'mem-pool-info := set-bit-noexist 'mem-pool-info (pool b) ('lvl t) ('bb
  t));
  'block-pt := 'block-pt(t := block-ptr ('mem-pool-info (pool b)) ('lsz t)
  ('bb t));
  IF 'bn t  $\neq$  'bb t  $\wedge$ 
  block-fits ('mem-pool-info (pool b)) ('block-pt t)
  ('lsz t) THEN 'mem-pool-info := 'mem-pool-info
  (pool b := remove-free-list ('mem-pool-info (pool b)) ('lvl
  t) ('block-pt t)) FI;;
  'i := 'i(t := Suc ('i t))
  OD;;
  ('lvl := 'lvl(t := 'lvl t - 1); 'bn := 'bn(t := 'bn t div 4);
  'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) ('lvl t) ('bn t);
  'freeing-node := 'freeing-node(t  $\mapsto$ 
  (pool = pool b, level = 'lvl t, block = 'bn t,
  data = block-ptr ('mem-pool-info (pool b)) (ALIGN4 (max-sz ('mem-pool-info
  (pool b))) div 4 ^ 'lvl t) ('bn t)))
  ELSE IF block-fits ('mem-pool-info (pool b)) ('blk t)
  ('lsz t) THEN 'mem-pool-info := 'mem-pool-info
  (pool b := append-free-list ('mem-pool-info (pool b)) ('lvl t)
  ('blk t)) FI;;
  'free-block-r := 'free-block-r(t := False)
  FI) satp [mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \cap \{V\} \cap UNIV \cap \{Va\}$ ,
  {(s, t). s = t}, UNIV,
   $\llbracket '(Pair Va) \in UNIV \rrbracket \cap (\llbracket '(Pair V) \in Mem\text{-pool-free-guar } t \rrbracket$ 
   $\cap (mp\text{-free-precond8-inv } t b (\alpha - 1) \cup mp\text{-free-precond8-inv } t b 0))]$ 

apply(rule subst[where s=mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \cap \{V\}$ 
 $\cap \{Va\}$ 
and t=mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = \text{Some } t \rrbracket \cap \{V\} \cap UNIV \cap \{Va\}$ ])
apply blast
apply(rule subst[where s= $\llbracket '(Pair V) \in Mem\text{-pool-free-guar } t \rrbracket \cap (mp\text{-free-precond8-inv}$ 
  t b ( $\alpha - 1$ ))
and t= $\llbracket '(Pair Va) \in UNIV \rrbracket \cap (\llbracket '(Pair V) \in Mem\text{-pool-free-guar } t \rrbracket \cap$ 

```

```

(mp-free-precond8-inv t b ( $\alpha - 1$ ))))))
  apply blast

apply(case-tac V  $\neq$  Va)
  apply(rule subst[where s={} and t=mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur =$ 
Some t $\rrbracket \cap \{V\} \cap \{Va\}$ ]])
  apply fast using Emptyprecond[of - {(s, t). s = t} UNIV
   $\llbracket '(Pair\ Va) \in UNIV \rrbracket \cap (\llbracket '(Pair\ V) \in Mem-pool-free-guar\ t \rrbracket \cap (mp-free-precond8-inv$ 
t b ( $\alpha - 1$ )  $\cup mp-free-precond8-inv\ t\ b\ 0$ ))] apply auto[1]

apply(case-tac mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} \cap \{Va\} =$ 
{})
  using Emptyprecond apply metis

apply(rule subst[where s={V} and t=mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur =$ 
Some t $\rrbracket \cap \{V\} \cap \{Va\}$ ]])
  using two-int-one[of mp-free-precond8-3 t b  $\alpha \cap \llbracket 'cur = Some\ t \rrbracket\ V\ Va$ ] apply
fast
apply(subgoal-tac V  $\in mp-free-precond8-3\ t\ b\ \alpha \cap \llbracket 'cur = Some\ t \rrbracket$ )
  prefer 2 apply fast

apply(rule Seq[where mid={free-stm8-precond2 V t b}])
apply(rule Seq[where mid={free-stm8-precond1 V t b}])

apply(rule Basic)
  apply force
  apply fast using stable-id2 apply fast using stable-id2 apply fast

apply(rule Basic)
  using mempool-free-stm8-atombody-h1 apply fast
  using stable-id2 apply fast using stable-id2 apply fast using stable-id2 apply
fast

apply(rule Cond)
  using stable-id2 apply fast
  apply(case-tac {free-stm8-precond2 V t b}  $\cap \llbracket NULL < 'lwl\ t \wedge partner-bits$ 
('mem-pool-info (pool b)) ('lwl t) ('bn t) $\rrbracket = \{\}$ )
  using Emptyprecond apply metis

apply(rule Seq[where mid=free-stm8-precond4 Va t b])
apply(rule Seq[where mid=free-stm8-precond3 Va t b])

apply(rule Basic)
apply simp apply(simp add:gvars-conf-stable-def gvars-conf-def lvars-nochange-def)

```

```

    apply fast using stable-id2 apply fast using stable-id2 apply fast

    apply(rule While)
      using stable-id2 apply fast apply(simp add:Let-def) apply auto[1] using
stable-id2 apply fast
      using mempool-free-stm8-set4partbits-while[of V Va t b  $\alpha$ ] apply fast
      apply fast

    apply(rule subst[where  $s = \llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}free\text{-}guar\ t \rrbracket$ 
       $\cap (mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ (\alpha - 1) \cup mp\text{-}free\text{-}precond8\text{-}inv$ 
 $t\ b\ 0)$ 
      and  $t = \llbracket '(Pair\ Va) \in UNIV \rrbracket \cap (\llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}free\text{-}guar$ 
 $t \rrbracket$ 
       $\cap (mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ (\alpha - 1) \cup mp\text{-}free\text{-}precond8\text{-}inv$ 
 $t\ b\ 0))$ ])
      apply auto[1]
      using mempool-free-stm8-atombody-rest-extpost[of V t b  $\alpha$ ] apply fast
      apply(rule subst[where  $s = \llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}free\text{-}guar\ t \rrbracket$ 
       $\cap (mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ (\alpha - 1) \cup mp\text{-}free\text{-}precond8\text{-}inv$ 
 $t\ b\ 0)$ 
      and  $t = \llbracket '(Pair\ Va) \in UNIV \rrbracket \cap (\llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}free\text{-}guar$ 
 $t \rrbracket$ 
       $\cap (mp\text{-}free\text{-}precond8\text{-}inv\ t\ b\ (\alpha - 1) \cup mp\text{-}free\text{-}precond8\text{-}inv$ 
 $t\ b\ 0))$ ])
      apply auto[1]
      using mempool-free-stm8-atombody-else-extpost[of V t b  $\alpha$ ] apply fast

    apply fast
  done

lemma  $\{(s, t). s = t\} = Id$  by auto

abbreviation st8-while-body  $t\ b \equiv$ 
  ( $t \blacktriangleright 'lsz := 'lsz\ (t := 'lsizes\ t!\ ('lwl\ t))$ );;
  ( $t \blacktriangleright 'blk := 'blk\ (t := block\text{-}ptr\ ('mem\text{-}pool\text{-}info\ (pool\ b))\ ('lsz\ t)\ ('bn\ t))$ );;

  ( $t \blacktriangleright ATOMIC$ 

     $'mem\text{-}pool\text{-}info := set\text{-}bit\text{-}free\ 'mem\text{-}pool\text{-}info\ (pool\ b)\ ('lwl\ t)\ ('bn\ t)$ ;;
     $'freeing\text{-}node := 'freeing\text{-}node\ (t := None)$ ;; (* remove the freeing node info
of the thread *)

    IF  $'lwl\ t > 0 \wedge partner\text{-}bits\ ('mem\text{-}pool\text{-}info\ (pool\ b))\ ('lwl\ t)\ ('bn\ t)$  THEN
      FOR  $'i := 'i(t := 0)$ ;
         $'i\ t < 4$ ;
         $'i := 'i(t := 'i\ t + 1)$  DO
           $'bb := 'bb\ (t := ('bn\ t\ div\ 4) * 4 + 'i\ t)$ ;;
          ( $* (t \blacktriangleright 'mem\text{-}pool\text{-}info := clear\text{-}free\text{-}bit\ 'mem\text{-}pool\text{-}info\ (pool\ b)\ ('lwl\ t)\ ('bb$ 
 $t))$ );;*)

```

```

    'mem-pool-info := set-bit-noexist 'mem-pool-info (pool b) ('lvl t) ('bb t);;
    'block-pt := 'block-pt (t := block-ptr ('mem-pool-info (pool b)) ('lsz t) ('bb
t));;
    IF 'bn t ≠ 'bb t ∧ block-fits ('mem-pool-info (pool b))
        ('block-pt t)
        ('lsz t) THEN

        (* sys-dlist-remove(block-ptr(p, lsz, b)); *)
        'mem-pool-info := 'mem-pool-info ((pool b) :=
            remove-free-list ('mem-pool-info (pool b)) ('lvl t) ('block-pt t))
        FI
    ROF;;

    (
    (* 'j := 'j (t := 'lvl t);; (* use lbn and j to store the previous lvl and bn, or
can not give the post condition *)
    'lbn := 'lbn (t := 'bn t);; (* since the lbn and j are not used in M-pool-free
*)
    'lvl := 'lvl (t := 'j t - 1);;
    'bn := 'bn (t := 'lbn t div 4);;*)
    'lvl := 'lvl (t := 'lvl t - 1);;
    'bn := 'bn (t := 'bn t div 4);;
    (* we add this statement. set the parent node from divided to freeing *)
    'mem-pool-info := set-bit-freeing 'mem-pool-info (pool b) ('lvl t) ('bn t);;
    (* 'freeing-node := 'freeing-node (t := Some (pool = (pool b), level = ('lvl
t),
        block = ('bn t), data = block-ptr ('mem-pool-info (pool b)) ('lsz t)
('bn t) ))*)
    'freeing-node := 'freeing-node (t := Some (pool = (pool b), level = ('lvl t),
        block = ('bn t),
        data = block-ptr ('mem-pool-info (pool b))
            (((ALIGN4 (max-sz ('mem-pool-info (pool b)))) div (4 ^ ('lvl
t))))))
        ('bn t) ))
    )

ELSE
    IF block-fits ('mem-pool-info (pool b)) ('blk t) ('lsz t) THEN

        (* sys-dlist-append(&p->levels[level].free-list, block); *)
        'mem-pool-info := 'mem-pool-info ((pool b) :=
            append-free-list ('mem-pool-info (pool b)) ('lvl t) ('blk t) )
        FI;;

        'free-block-r := 'free-block-r (t := False)
    FI

END)

```

**lemma** *mp-free-precond8-inv-0-stb* :  
*stable* (*mp-free-precond8-inv* *t b* ( $\alpha - 1$ )  $\cup$  *mp-free-precond8-inv* *t b* 0) (*Mem-pool-free-rely* *t*)  
**apply**(*rule stable-un2*)  
**using** *mp-free-precond8-inv-stb*[*of t b*  $\alpha - 1$ ] **apply** *fast*  
**using** *mp-free-precond8-inv-stb*[*of t b* 0] **apply** *fast*  
**done**

**lemma** *mempool-free-stm8-body-terminate*:  
 $\Gamma \vdash_I \text{Some } (st8\text{-while-body } t \ b)$   
 $sat_p [mp\text{-free-precond8-inv } t \ b \ \alpha \cap \{\alpha > 0\}, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar}$   
 $t,$   
 $mp\text{-free-precond8-inv } t \ b \ (\alpha - 1) \cup mp\text{-free-precond8-inv } t \ b \ 0]$   
**apply**(*rule Seq*[**where** *mid*=*mp-free-precond8-3* *t b*  $\alpha$ ])  
**apply**(*rule Seq*[**where** *mid*=*mp-free-precond8-2* *t b*  $\alpha$ ])

**apply**(*unfold stm-def*)[1]  
**apply**(*rule Await*)  
**using** *mp-free-precond8-1-stb*[*of t b*  $\alpha$ ] **apply** *blast*  
**using** *mp-free-precond8-2-stb*[*of t b*  $\alpha$ ] **apply** *blast*

**apply**(*rule allI*)  
**apply**(*rule Basic*)  
**apply**(*case-tac mp-free-precond8-1* *t b*  $\alpha \cap \{\text{cur} = \text{Some } t\} \cap \{V\} = \{\}$ )  
**apply** *auto*[1] **apply** *clarsimp* **apply**(*rule conjI*)  
**apply**(*simp add: Guar<sub>f</sub>-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def*)  
**apply**(*rule disjI1*)  
**apply**(*rule conjI*)  
**apply**(*subgoal-tac* ( $V, V \{lsz := (lsz \ V)(t := lsizes \ V \ t \ ! \ (lvl \ V \ t))\} \in lvars\text{-nochange1-4all}$ )  
**using** *glnochange-inv0* **apply** *auto*[1] **apply**(*simp add:lvars-nochange1-4all-def*  
*lvars-nochange1-def*)  
**apply**(*simp add:lvars-nochange-def*)  
**apply**(*subgoal-tac* ( $V, V \{lsz := (lsz \ V)(t := lsizes \ V \ t \ ! \ (lvl \ V \ t))\} \in lvars\text{-nochange1-4all}$ )  
**using** *glnochange-inv0* **apply** *auto*[1] **apply**(*simp add:lvars-nochange1-4all-def*  
*lvars-nochange1-def*)  
**apply** *fast* **using** *stable-id2* **apply** *fast* **using** *stable-id2* **apply** *fast*

**apply**(*unfold stm-def*)[1]  
**apply**(*rule Await*)  
**using** *mp-free-precond8-2-stb* **apply** *blast*  
**using** *mp-free-precond8-3-stb* **apply** *blast*  
**apply**(*rule allI*)  
**apply**(*rule Basic*)  
**apply**(*case-tac mp-free-precond8-2* *t b*  $\alpha \cap \{\text{cur} = \text{Some } t\} \cap \{V\} = \{\}$ )  
**apply** *auto*[1] **apply** *clarsimp* **apply**(*rule conjI*)  
**apply**(*simp add: Guar<sub>f</sub>-def gvars-conf-stable-def gvars-conf-def Mem-pool-free-guar-def*)

```

apply(rule disjI1)
apply(rule conjI)
apply(subgoal-tac (V, V( $\lfloor$ blk := (blk V)
  (t := block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz (mem-pool-info
V (pool b))) div 4 ^ lvl V t)
    (block b div 4 ^ (level b - lvl V t)))))) $\in$ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply(simp add:lvars-nochange-def)
apply(subgoal-tac (V, V( $\lfloor$ blk := (blk V)
  (t := block-ptr (mem-pool-info V (pool b)) (ALIGN4 (max-sz (mem-pool-info
V (pool b))) div 4 ^ lvl V t)
    (block b div 4 ^ (level b - lvl V t)))))) $\in$ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply fast using stable-id2 apply fast using stable-id2 apply fast

```

```

apply(unfold stm-def)[1]
apply(rule Await)
using mp-free-precond8-3-stb apply blast
using mp-free-precond8-inv-0-stb[of t b  $\alpha$ ] apply fast
apply(rule allI)
apply(rule Await)
using stable-id2 apply blast using stable-id2 apply blast
apply clarify using mempool-free-stm8-atombody[of b t  $\alpha$ ] apply auto[1]

```

done

**lemma** *loopbody-sat-invterm-imp-inv-post'*:  
 $\Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ b } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8-inv}$   
 $t \text{ b } (\alpha - 1) \cup \text{mp-free-precond8-inv } t \text{ b } 0]$   
 $\implies \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ b } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar}, \text{mp-free-precond8}$   
 $t \text{ b}]$   
**using** *Conseq* [of mp-free-precond8-inv t b  $\alpha \cap \{\alpha > 0\}$  mp-free-precond8-inv t b  
 $\alpha \cap \{\alpha > 0\}$   
 $\text{rely rely guar guar mp-free-precond8-inv } t \text{ b } (\alpha - 1) \cup \text{mp-free-precond8-inv}$   
 $t \text{ b } 0$   
 $\text{mp-free-precond8 } t \text{ b } \text{Some } P]$  **by** blast

**lemma** *stm8-inv-imp-prepost2'*:  
 $(\forall \alpha. \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-free-precond8-inv } t \text{ b } \alpha \cap \{\alpha > 0\}, \text{rely}, \text{guar},$   
 $\text{mp-free-precond8-inv } t \text{ b } (\alpha - 1) \cup \text{mp-free-precond8-inv } t \text{ b}$   
 $0])$   
 $\implies \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-free-precond8 } t \text{ b } \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8}$   
 $t \text{ b}]$

```

apply(rule subst[where s= $\forall v. v \in \text{mp-free-precond8 } t \text{ b } \cap \{\text{'free-block-r } t\} \longrightarrow$ 

```

$\Gamma \vdash_I \text{Some } P \text{ sat}_p [\{v\}, \text{rely}, \text{guar}, \text{mp-free-precond8 } t \ b] \text{ and}$   
 $t = \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-free-precond8 } t \ b \cap \{\text{'free-block-r } t\}, \text{rely}, \text{guar}, \text{mp-free-precond8}$   
 $t \ b])]$   
**using** *allpre-eq-pre*[*of* *mp-free-precond8* *t b*  $\cap \{\text{'free-block-r } t\}$   
*Some* *P* *rely* *guar* *mp-free-precond8* *t b*] **apply** *blast*

**apply**(*rule allI*) **apply**(*rule impI*)  
**apply**(*subgoal-tac*  $\exists \alpha. v \in \text{mp-free-precond8-inv } t \ b \ \alpha \cap \{\alpha > 0\}$ )  
**prefer** 2 **using** *loopp-re-imp-exist- $\alpha$ gt0* **apply** *blast*

**apply**(*erule exE*)

**using** *sat-pre-imp-allinpre*[*of* *Some* *P* - *rely* *guar* *mp-free-precond8* *t b*]  
*loopbody-sat-invterm-imp-inv-post'* **apply** *blast*  
**done**

**lemma** *mempool-free-stm8-body*:

$\Gamma \vdash_I \text{Some } (st8\text{-while-body } t \ b)$   
 $\text{sat}_p [\text{mp-free-precond8 } t \ b \cap \{\text{'free-block-r } t\}, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar}$   
 $t, \text{mp-free-precond8 } t \ b]$

**using** *stm8-inv-imp-prepost2'*[*of* (*st8-while-body* *t b*) *t b* *Mem-pool-free-rely* *t* *Mem-pool-free-guar*  
*t*]

*mempool-free-stm8-body-terminate*[*of* *t b*] **apply** *fast*

**done**

**lemma** *mempool-free-stm8*:

$\Gamma \vdash_I \text{Some } (WHILE \ \text{'free-block-r } t \ DO$   
 $st8\text{-while-body } t \ b$   
 $OD)$   
 $\text{sat}_p [\text{mp-free-precond8 } t \ b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{mp-free-precond9}$   
 $t \ b]$

**apply**(*rule While*)

**using** *mp-free-precond8-stb*[*of* *t b*] **apply** *blast*

**apply** *simp-inv* **apply** *auto*[1]

**using** *mp-free-precond9-stb*[*of* *t b*] **apply** *auto*[1]

**using** *mempool-free-stm8-body*[*of* *t b*] **apply** *fast*

**apply**(*simp add: Mem-pool-free-guar-def*)

**done**

## 9.4 statement 9

**abbreviation** *stm9-precond-while*  $Va \ t \ b$

$\equiv \{V. \text{inv } V \wedge \text{cur } V = \text{cur } Va \wedge \text{tick } V = \text{tick } Va \wedge (V, Va) \in \text{gvars-conf-stable}$

$\wedge \text{freeing-node } V \ t = \text{freeing-node } Va \ t \wedge \text{allocating-node } V \ t = \text{allocating-node}$



$Va\ t$   
 $\wedge (\forall p. \text{levels} (\text{mem-pool-info } V\ p) = \text{levels} (\text{mem-pool-info } Va\ p))$   
 $\wedge (\forall p. p \neq \text{pool } b \longrightarrow \text{mem-pool-info } V\ p = \text{mem-pool-info } Va\ p)$   
 $\wedge (\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ V\ Va)$

**lemma** *va-precond-while*:  $\text{inv } Va \implies Va \in \text{stm9-precond-while } Va\ t\ b$   
**by** (*simp add:gvars-conf-stable-def gvars-conf-def lvars-nochange-def*)

**lemma** *mempool-free-stm9-resch-inv-help*:

$\text{cur } V = \text{Some } t \implies \text{thd-state } V\ t = \text{RUNNING} \implies$   
 $(\text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta = \text{READY}) = t \implies$   
 $V = V \langle \text{cur} := \text{Some } (\text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta = \text{READY}),$   
 $\text{thd-state} := (\text{thd-state } V)(t := \text{READY}, \text{SOME } ta. ta \neq t \longrightarrow \text{thd-state}$   
 $V\ ta = \text{READY} := \text{RUNNING}) \rangle$

**apply** *auto*

**proof** –

**assume** *a1*:  $\text{thd-state } V\ t = \text{RUNNING}$

**assume** *a2*:  $\text{cur } V = \text{Some } t$

**have**  $(\text{thd-state } V)(t := \text{RUNNING}) = \text{thd-state } V$

**using** *a1* **by** (*metis fun-upd-triv*)

**then show**  $V = V \langle \text{cur} := \text{Some } t, \text{thd-state} := (\text{thd-state } V)(t := \text{RUNNING}) \rangle$

**using** *a2* **by** *simp*

**qed**

**lemma** *mempool-free-stm9-resch-inv*:

$\text{cur } V = \text{Some } t \implies \text{inv } V \implies \text{inv } (V \langle \text{cur} := \text{Some } (\text{SOME } ta. ta \neq t \longrightarrow$   
 $\text{thd-state } V\ ta = \text{READY}),$

$\text{thd-state} := (\text{thd-state } V)(t := \text{READY}, \text{SOME } ta. ta \neq t \longrightarrow \text{thd-state}$   
 $V\ ta = \text{READY} := \text{RUNNING}) \rangle$

**apply**(*subgoal-tac*  $\text{thd-state } V\ t = \text{RUNNING}$ )

**apply**(*case-tac*  $(\text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta = \text{READY}) = t$ )

$\text{apply}(\text{subgoal-tac } V = V \langle \text{cur} := \text{Some } (\text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta$   
 $= \text{READY}),$

$\text{thd-state} := (\text{thd-state } V)(t := \text{READY}, \text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta$   
 $= \text{READY} := \text{RUNNING}) \rangle$

**apply** *simp using mempool-free-stm9-resch-inv-help[of V t]* **apply** *auto[1]*

$\text{apply}(\text{subgoal-tac } \text{thd-state } V\ (\text{SOME } ta. ta \neq t \longrightarrow \text{thd-state } V\ ta = \text{READY})$   
 $= \text{READY})$

**apply**(*simp add:inv-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def*) **apply** *auto[1]*

**apply**(*rule conjI*) **apply**(*simp add:inv-thd-waitq-def*) **apply** *auto[1]*

**apply**(*rule conjI*) **apply**(*simp add:inv-mempool-info-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap-freelist-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap-def*)

**apply**(*rule conjI*) **apply**(*simp add: inv-aux-vars-def mem-block-addr-valid-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmap0-def*)

**apply**(*rule conjI*) **apply**(*simp add:inv-bitmapn-def*)

**apply**(*simp add:inv-bitmap-not4free-def*)

**apply** (*metis* (*mono-tags*, *lifting*) *someI-ex*)

**apply**(*simp add:inv-def inv-cur-def*) **apply** *auto*[1]

**done**

**lemma** *mempool-free-stm9-ifpart-one*:

$Va \in mp\text{-free-precond9 } t \ b \cap \{\text{'cur} = \text{Some } t\} \implies$   
 $V \in stm9\text{-precond-while } Va \ t \ b \cap \{\text{wait-}q \ (\text{'mem-pool-info } (pool \ b)) = []\} \implies$   
 $\Gamma \vdash_I \text{Some } (IF \ \text{'need-resched } t \ THEN \text{reschedule } FI \ )$   
 $sat_p \ [\{V\}, \{(x, y). x = y\}, UNIV, \{\text{'(Pair } Va) \in Mem\text{-pool-free-guar } t\} \cap$   
 $Mem\text{-pool-free-post } t]$

**apply**(*rule Cond*)

**apply**(*simp add:stable-def*)

**apply**(*simp add:reschedule-def*)

**apply**(*rule Seq*[**where**  $mid = \{V \ (\text{thd-state} := (\text{thd-state } V)(\text{the } (cur \ V) := READY))\}$ ])

$\{\text{cur} := \text{Some } (SOME \ t. (\text{thd-state } (V \ (\text{thd-state} := (\text{thd-state } V)(\text{the } (cur \ V) := READY)))) \ t = READY)\}$

**apply**(*rule Seq*[**where**  $mid = \{V \ (\text{thd-state} := (\text{thd-state } V)(\text{the } (cur \ V) := READY))\}$ ])

**apply**(*rule Basic*)

**apply** *auto*[1] **apply**(*simp add:stable-def*)+

**apply**(*rule Basic*)

**apply** *auto*[1] **apply**(*simp add:stable-def*)+

**apply**(*rule Basic*)

**apply** *auto*[1] **apply**(*simp add:Mem-pool-free-guar-def*) **apply**(*rule disjI1*)

**apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)

**apply**(*rule conjI*) **using** *mempool-free-stm9-resch-inv* **apply** *auto*[1]

**apply**(*simp add:lvars-nochange-def*) **apply**(*simp add:Mem-pool-free-post-def*)

**using** *mempool-free-stm9-resch-inv* **apply** *auto*[1] **apply** *auto*[1] **apply**(*simp add:stable-def*)+

**apply**(*simp add:Skip-def*)

**apply**(*rule Basic*) **apply** *auto*[1] **apply**(*simp add:Mem-pool-free-guar-def*)

**apply**(*rule disjI1*) **apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)

**apply**(*simp add:lvars-nochange-def*)

**apply**(*simp add:Mem-pool-free-post-def*)

**apply**(*simp add:stable-def*)+

**done**

**lemma** *mempool-free-stm9-ifpart*:

$Va \in mp\text{-free-precond9 } t \ b \cap \{\text{'cur} = \text{Some } t\} \implies$   
 $\Gamma \vdash_I \text{Some } (IF \ \text{'need-resched } t \ THEN \text{reschedule } FI \ )$   
 $sat_p \ [stm9\text{-precond-while } Va \ t \ b \cap \{\text{wait-}q \ (\text{'mem-pool-info } (pool \ b)) = []\},$   
 $\{(x, y). x = y\}, UNIV, \{\text{'(Pair } Va) \in Mem\text{-pool-free-guar } t\} \cap$   
 $Mem\text{-pool-free-post } t]$

**using** *mempool-free-stm9-ifpart-one*[of  $Va\ t\ b$ ]  
*Allprecond*[**where**  $U = \text{stm9-precond-while}\ Va\ t\ b \cap \llbracket \text{wait-}q\ (\text{'mem-pool-info}\ (pool\ b)) = [] \rrbracket$  **and**  
 $P = \text{Some}\ (IF\ \text{'need-resched}\ t\ THEN\ \text{reschedule}\ FI)$  **and**  $\text{rely} = \{(x, y). x = y\}$  **and**  
 $\text{guar} = UNIV$  **and**  $\text{post} = \llbracket \text{'(Pair}\ Va) \in \text{Mem-pool-free-guar}\ t \rrbracket$   
 $\cap \text{Mem-pool-free-post}\ t]$   
**by** *blast*

**lemma** *mempool-free-stm9-loopbody-one*:

$Va \in \text{mp-free-precond9}\ t\ b \cap \llbracket \text{'cur} = \text{Some}\ t \rrbracket \implies$   
 $Vb \in \text{stm9-precond-while}\ Va\ t\ b \cap \llbracket \text{wait-}q\ (\text{'mem-pool-info}\ (pool\ b)) \neq [] \rrbracket \implies$   
 $\Gamma \vdash_I \text{Some}\ (\text{'th} := \text{'th}(t := \text{hd}\ (\text{wait-}q\ (\text{'mem-pool-info}\ (pool\ b)))));;$   
 $\text{'mem-pool-info} := \text{'mem-pool-info}$   
 $(pool\ b := \text{'mem-pool-info}\ (pool\ b)) \llbracket \text{wait-}q := \text{tl}\ (\text{wait-}q\ (\text{'mem-pool-info}\ (pool\ b))) \rrbracket);;$   
 $\text{'thd-state} := \text{'thd-state}(\text{'th}\ t := \text{READY});;$   
 $\text{'need-resched} := \text{'need-resched}(t := \text{True})$   
 $\text{sat}_p\ [\{Vb\}, \{(x, y). x = y\}, UNIV, \text{stm9-precond-while}\ Va\ t\ b]$   
**apply**(*rule Seq*[**where**  $\text{mid} = \{Vb \llbracket \text{th} := (\text{th}\ Vb)\ (t := \text{hd}\ (\text{wait-}q\ ((\text{mem-pool-info}\ Vb)\ (pool\ b)))) \rrbracket\}$ ]  
 $\llbracket \text{mem-pool-info} := (\text{mem-pool-info}\ Vb)$   
 $(pool\ b := (\text{mem-pool-info}\ Vb)\ (pool\ b)) \llbracket \text{wait-}q := \text{tl}\ (\text{wait-}q$   
 $((\text{mem-pool-info}\ Vb)\ (pool\ b))) \rrbracket \rrbracket$   
 $\llbracket \text{thd-state} := (\text{thd-state}\ Vb)(\text{hd}\ (\text{wait-}q\ ((\text{mem-pool-info}\ Vb)$   
 $(pool\ b))) := \text{READY}) \rrbracket \rrbracket$

**apply**(*rule Seq*[**where**  $\text{mid} = \{Vb \llbracket \text{th} := (\text{th}\ Vb)\ (t := \text{hd}\ (\text{wait-}q\ ((\text{mem-pool-info}\ Vb)\ (pool\ b)))) \rrbracket\}$ ]  
 $\llbracket \text{mem-pool-info} := (\text{mem-pool-info}\ Vb)$   
 $(pool\ b := (\text{mem-pool-info}\ Vb)\ (pool\ b)) \llbracket \text{wait-}q := \text{tl}\ (\text{wait-}q$   
 $((\text{mem-pool-info}\ Vb)\ (pool\ b))) \rrbracket \rrbracket$

**apply**(*rule Seq*[**where**  $\text{mid} = \{Vb \llbracket \text{th} := (\text{th}\ Vb)\ (t := \text{hd}\ (\text{wait-}q\ ((\text{mem-pool-info}\ Vb)\ (pool\ b)))) \rrbracket \rrbracket$

**apply**(*rule Basic*) **apply** *auto*[1] **apply** *simp* **apply**(*simp add:stable-def*)+  
**apply**(*rule Basic*) **apply** *auto*[1] **apply** *simp* **apply**(*simp add:stable-def*)+  
**apply**(*rule Basic*) **apply** *auto*[1] **apply** *simp* **apply**(*simp add:stable-def*)+

**apply**(*rule Basic*) **apply** *clarify* **apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def*  
*gvars-conf-def*)

**apply**(*simp add:inv-def*)  
**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def inv-thd-waitq-def*)  
**apply**(*rule conjI*) **apply**(*simp add:inv-thd-waitq-def*) **apply** *clarify*  
**apply**(*rule conjI*) **apply** *clarify* **apply** (*rule conjI*) **apply** *clarify* **apply**(*rule*  
*conjI*) **apply** *clarify*  
**apply** (*smt List.nth-tl Nitpick.size-list-simp*(2) *Suc-mono gr-implies-not0*  
*hd-conv-nth in-set-conv-nth length-pos-if-in-set lessI list.set-sel*(1))

```

apply clarify apply (meson list.set-sel(2)) apply clarify apply (metis list.set-sel(1))

apply(rule conjI) apply clarify apply (metis hd-Cons-tl set-ConsD)
apply(rule conjI) apply clarify apply (metis (no-types, lifting) List.nth-tl
Nitpick.size-list-simp(2))
One-nat-def Suc-mono length-tl nat.inject
apply clarify apply(rule conjI) apply clarify apply (metis list.set-sel(2))
apply clarify apply(rule conjI) apply clarify apply (metis list.set-sel(2))
apply clarify apply metis
apply(rule conjI) apply(simp add: inv-mempool-info-def) apply auto[1]
apply(rule conjI) apply(simp add: inv-bitmap-freelist-def)
apply(rule conjI) apply(simp add: inv-bitmap-def)
apply(rule conjI) apply(simp add: inv-aux-vars-def mem-block-addr-valid-def)
apply(rule conjI) apply metis apply metis
apply(rule conjI) apply(simp add: inv-bitmap0-def)
apply(rule conjI) apply(simp add: inv-bitmapn-def)
apply(simp add: inv-bitmap-not4free-def partner-bits-def)

apply(rule conjI) apply auto[1]
apply(rule conjI) apply auto[1]
apply(rule conjI) apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(rule conjI) apply auto[1]
apply(rule conjI) apply force
apply(rule conjI)
apply clarify apply(simp add: lvars-nochange-def)
apply(simp add: lvars-nochange-def)
by (simp add: stable-def)+

```

**lemma** *mempool-free-stm9-loopbody*:

```

Va ∈ mp-free-precond9 t b ∩  $\{\text{'cur} = \text{Some } t\} \implies$ 
   $\Gamma \vdash_I \text{Some } (\text{'th} := \text{'th}(t := \text{hd } (\text{wait-q } (\text{'mem-pool-info } (\text{pool } b)))));;$ 
    'mem-pool-info := 'mem-pool-info
    (pool b := 'mem-pool-info (pool b)(wait-q := tl (wait-q ('mem-pool-info
(pool b))))));;
    'thd-state := 'thd-state('th t := READY);;
    'need-resched := 'need-resched(t := True) )
  satp [stm9-precond-while Va t b ∩  $\{\text{wait-q } (\text{'mem-pool-info } (\text{pool } b)) \neq []\}$ ,
     $\{(x, y). x = y\}$ , UNIV, stm9-precond-while Va t b]
using mempool-free-stm9-loopbody-one
  Allprecond[where U=stm9-precond-while Va t b ∩  $\{\text{wait-q } (\text{'mem-pool-info$ 
(pool b))  $\neq []\}$  and
    P= Some ('th := 'th(t := hd (wait-q ('mem-pool-info (pool b)))));;
    'mem-pool-info := 'mem-pool-info(pool b := 'mem-pool-info (pool
b)(wait-q := tl (wait-q ('mem-pool-info (pool b))))));;
    'thd-state := 'thd-state('th t := READY);;
    'need-resched := 'need-resched (t := True))
  and rely= $\{(x, y). x = y\}$  and guar= UNIV and post= stm9-precond-while
Va t b]

```

**by** *blast*

**lemma** *mempool-free-stm9-body-loopinv*:

$Va \in mp\text{-}free\text{-}precond9\ t\ b \cap \{\cur = Some\ t\} \implies$   
 $\Gamma \vdash_I Some\ (WHILE\ wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b)) \neq \emptyset$   
 $DO\ \text{'th} := \text{'th}(t := hd\ (wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))));$   
 $\text{'mem}\text{-}pool\text{-}info := \text{'mem}\text{-}pool\text{-}info$   
 $(pool\ b := \text{'mem}\text{-}pool\text{-}info\ (pool\ b) \parallel wait\text{-}q := tl\ (wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info$   
 $(pool\ b)))));$   
 $\text{'thd}\text{-}state := \text{'thd}\text{-}state(\text{'th}\ t := READY);$   
 $\text{'need}\text{-}resched := \text{'need}\text{-}resched(t := True)$   
 $OD;$   
 $IF\ \text{'need}\text{-}resched\ t\ THEN\ reschedule\ FI\ )$   
 $sat_p\ [stm9\text{-}precond\text{-}while\ Va\ t\ b,\ \{(x, y). x = y\},\ UNIV,\ \{s. (Va, s) \in Mem\text{-}pool\text{-}free\text{-}guar$   
 $t\} \cap Mem\text{-}pool\text{-}free\text{-}post\ t]$   
**apply**(*rule Seq* **where** *mid* = *stm9-precond-while* *Va t b*  $\cap \{\wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info$   
 $(pool\ b)) = \emptyset\}$ )  
**apply**(*rule While*)  
**apply**(*simp add:stable-def*)  
**apply** *auto*[1]  
**apply**(*simp add:stable-def*)  
**using** *mempool-free-stm9-loopbody*[*of Va t b*] **apply** *simp*  
**apply** *simp*

**using** *mempool-free-stm9-ifpart* **by** *blast*

**lemma** *mempool-free-stm9-body*:

$mp\text{-}free\text{-}precond9\ t\ b \cap \{\inv\} \cap \{\cur = Some\ t\} \cap \{Va\} \neq \{\} \implies$   
 $\Gamma \vdash_I Some\ (WHILE\ wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b)) \neq \emptyset$   
 $DO\ \text{'th} := \text{'th}(t := hd\ (wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b))));$   
 $\text{'mem}\text{-}pool\text{-}info := \text{'mem}\text{-}pool\text{-}info(pool\ b := \text{'mem}\text{-}pool\text{-}info\ (pool\ b) \parallel wait\text{-}q$   
 $:= tl\ (wait\text{-}q\ (\text{'mem}\text{-}pool\text{-}info\ (pool\ b)))));$   
 $\text{'thd}\text{-}state := \text{'thd}\text{-}state(\text{'th}\ t := READY);$   
 $\text{'need}\text{-}resched := \text{'need}\text{-}resched(t := True)$   
 $OD;$   
 $IF\ \text{'need}\text{-}resched\ t\ THEN\ reschedule\ FI\ )$   
 $sat_p\ [mp\text{-}free\text{-}precond9\ t\ b \cap \{\cur = Some\ t\} \cap \{Va\},$   
 $\{(s, t). s = t\},\ UNIV,\ \{\text{'(Pair}\ Va) \in Mem\text{-}pool\text{-}free\text{-}guar\ t\} \cap Mem\text{-}pool\text{-}free\text{-}post$   
 $t]$   
**apply**(*subgoal-tac inv Va*) **prefer** 2 **apply** *simp*  
**apply**(*subgoal-tac Va*  $\in mp\text{-}free\text{-}precond9\ t\ b \cap \{\inv\} \cap \{\cur = Some\ t\}$ )  
**prefer** 2 **apply** *simp*  
**using** *mempool-free-stm9-body-loopinv*[*of Va t b*] *va-precond-while*[*of Va t b*]  
 $Conseq[\textbf{where}\ pre = \{Va\}\ \textbf{and}\ pre' = stm9\text{-}precond\text{-}while\ Va\ t\ b\ \textbf{and}\ rely = \{(x,$   
 $y). x = y\}\ \textbf{and}\ rely' = \{(x, y). x = y\}$   
 $\textbf{and}\ guar = UNIV\ \textbf{and}\ guar' = UNIV\ \textbf{and}\ post' = \{\text{'(Pair}\ Va) \in$   
 $Mem\text{-}pool\text{-}free\text{-}guar\ t\} \cap Mem\text{-}pool\text{-}free\text{-}post\ t$   
 $\textbf{and}\ post = \{\text{'(Pair}\ Va) \in Mem\text{-}pool\text{-}free\text{-}guar\ t\} \cap Mem\text{-}pool\text{-}free\text{-}post\ t$

```

and  $P = \text{Some } (\text{WHILE } \text{wait-}q \text{ (}'\text{mem-pool-info (pool } b)\text{)}) \neq []$ 
  DO  $\text{'th} := \text{'th}(t := \text{hd } (\text{wait-}q \text{ (}'\text{mem-pool-info (pool } b)\text{)}))$ ;
     $\text{'mem-pool-info} := \text{'mem-pool-info}$ 
     $(\text{pool } b := \text{'mem-pool-info (pool } b) \parallel \text{wait-}q := \text{tl } (\text{wait-}q \text{ (}'\text{mem-pool-info}$ 
       $(\text{pool } b)\text{)}))$ ;
     $\text{'thd-state} := \text{'thd-state}(\text{'th } t := \text{READY})$ ;
     $\text{'need-resched} := \text{'need-resched}(t := \text{True})$ 
  OD;
  IF  $\text{'need-resched } t$  THEN  $\text{reschedule } FI$  ]
apply force
done

```

**lemma** *mempool-free-stm9*:

```

 $\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{ATOMIC}$ 
  WHILE  $\text{wait-}q \text{ (}'\text{mem-pool-info (pool } b)\text{)}) \neq []$  DO
     $\text{'th} := \text{'th}(t := \text{hd } (\text{wait-}q \text{ (}'\text{mem-pool-info (pool } b)\text{)}))$ ;
     $(* \text{-unpend-thread}(th); *)$ 
     $\text{'mem-pool-info} := \text{'mem-pool-info (pool } b := \text{'mem-pool-info (pool } b)$ 
       $(\text{wait-}q := \text{tl } (\text{wait-}q \text{ (}'\text{mem-pool-info (pool } b)\text{)}))$ ;
     $(* \text{-ready-thread}(th); *)$ 
     $\text{'thd-state} := \text{'thd-state}(\text{'th } t := \text{READY})$ ;
     $\text{'need-resched} := \text{'need-resched}(t := \text{True})$ 
  OD;

  IF  $\text{'need-resched } t$  THEN
     $\text{reschedule}$ 
  FI
END)
 $\text{sat}_p [\text{mp-free-precond9 } t \text{ } b, \text{Mem-pool-free-rely } t, \text{Mem-pool-free-guar } t, \text{Mem-pool-free-post}$ 
 $t]$ 
apply(simp add:stm-def)
apply(rule Await)
using mp-free-precond9-stb apply auto[1]
apply (simp add: mem-pool-free-post-stb)

apply(rule allI)
apply(rule Await)
apply(simp add:stable-def) apply(simp add:stable-def)
apply(rule allI)
apply(case-tac  $V = Va$ ) apply simp
apply(case-tac  $\text{mp-free-precond9 } t \text{ } b \cap \{\text{'inv}\} \cap \{\text{'cur} = \text{Some } t\} \cap \{Va =$ 
 $\{\}$ )
  apply simp apply (simp add: Emptyprecond stable-id2)
  apply clarify using mempool-free-stm9-body apply force
  apply simp apply (simp add: Emptyprecond stable-id2)
done

```

## 9.5 final proof

```

lemma Mempool-free-satRG:  $\Gamma$  (Mem-pool-free  $t$   $b$ )  $\vdash$  Mem-pool-free-RGCond  $t$   $b$ 
  apply (simp add:Evt-sat-RG-def)
  apply (simp add:body-def Pref-def Postf-def guard-def
    Relyf-def Guarf-def getrgformula-def)
  apply (simp add: Mem-pool-free-def Mem-pool-free-RGCond-def)
  apply (rule BasicEvt)
  apply (simp add:body-def Pref-def Postf-def guard-def
    getrgformula-def)
  apply (rule Seq[where mid=mp-free-precond9 t b])
  apply (rule Seq[where mid=mp-free-precond8 t b])
  apply (rule Seq[where mid=mp-free-precond7 t b])
  apply (rule Seq[where mid=mp-free-precond6 t b])
  apply (rule Seq[where mid=mp-free-precond5 t b])
  apply (rule Seq[where mid=mp-free-precond4 t b])
  apply (rule Seq[where mid=mp-free-precond3 t b])
  apply (rule Seq[where mid=mp-free-precond2 t b])

  using mempool-free-stm1 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm2 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm3 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm4 [ $of$   $t$   $b$ ] apply force
  using mempool-free-stm5 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm6 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm7 [ $of$   $t$   $b$ ] apply fast
  using mempool-free-stm8 [ $of$   $t$   $b$ ] apply force
  using mempool-free-stm9 [ $of$   $t$   $b$ ] apply force

  apply (simp add:body-def Pref-def Postf-def guard-def
    Relyf-def Guarf-def getrgformula-def)
  using mem-pool-free-pre-stb apply fast

  apply (simp add:getrgformula-def Mem-pool-free-guar-def)
done

end

theory func-cor-mempoolalloc
imports func-cor-lemma
begin

```

## 10 Functional correctness of $k\_mem\_pool\_alloc$

### 10.1 intermediate conditions and their stable to rely cond

**abbreviation**  $mp\_alloc\_precond1\ t\ p\ tm \equiv$   
 $Mem\_pool\_alloc\_pre\ t \cap \{p \in 'mem\_pools \wedge tm \geq -1\}$

**lemma**  $mp\_alloc\_precond1\_ext\_stb$ :  $stable\ (\{p \in 'mem\_pools \wedge tm \geq -1\})\ (Mem\_pool\_alloc\_rely\ t)$

$\mathbf{apply}(simp\ add:stable-def)\ \mathbf{apply}(rule\ allI)\ \mathbf{apply}(rule\ impI)\ \mathbf{apply}(rule\ allI)$   
 $\mathbf{apply}(rule\ impI)$   
 $\mathbf{apply}(simp\ add:Mem\_pool\_alloc\_rely-def)$   
 $\mathbf{apply}(simp\ add:gvars-conf-stable-def)$   
 $\mathbf{unfolding}\ gvars-conf-def\ \mathbf{apply}\ metis$   
**done**

**lemma**  $mp\_alloc\_precond1\_stb$ :  $stable\ (mp\_alloc\_precond1\ t\ p\ tm)\ (Mem\_pool\_alloc\_rely\ t)$

$\mathbf{apply}(rule\ stable-int2)$   
 $\mathbf{apply}(simp\ add:mem\_pool\_alloc\_pre-stb)$   
 $\mathbf{apply}(simp\ add:mp\_alloc\_precond1\_ext-stb)$   
**done**

**abbreviation**  $mp\_alloc\_precond2\ t\ p\ tm \equiv$   
 $mp\_alloc\_precond1\ t\ p\ tm \cap \{tmout\ t = tm\}$

**lemma**  $mp\_alloc\_precond2\_ext\_stb$ :  $stable\ (\{tmout\ t = tm\})\ (Mem\_pool\_alloc\_rely\ t)$

$\mathbf{apply}(simp\ add:stable-def)\ \mathbf{apply}(rule\ allI)\ \mathbf{apply}(rule\ impI)\ \mathbf{apply}(rule\ allI)$   
 $\mathbf{apply}(rule\ impI)$   
 $\mathbf{apply}(simp\ add:Mem\_pool\_alloc\_rely-def)$   
 $\mathbf{apply}(simp\ add:lvars-nochange-rel-def\ lvars-nochange-def)\ \mathbf{apply}\ smt$   
**done**

**lemma**  $mp\_alloc\_precond2\_stb$ :  $stable\ (mp\_alloc\_precond2\ t\ p\ tm)\ (Mem\_pool\_alloc\_rely\ t)$

$\mathbf{apply}(rule\ stable-int2)$   
 $\mathbf{apply}(simp\ add:mp\_alloc\_precond1-stb)$   
 $\mathbf{apply}(simp\ add:mp\_alloc\_precond2\_ext-stb)$   
**done**

**abbreviation**  $mp\_alloc\_precond3\ t\ p\ tm \equiv$   
 $mp\_alloc\_precond2\ t\ p\ tm \cap \{end\ t = 0\}$

**lemma**  $mp\_alloc\_precond3\_ext\_stb$ :  $stable\ (\{end\ t = 0\})\ (Mem\_pool\_alloc\_rely\ t)$

$\mathbf{apply}(simp\ add:stable-def)\ \mathbf{apply}(rule\ allI)\ \mathbf{apply}(rule\ impI)\ \mathbf{apply}(rule\ allI)$   
 $\mathbf{apply}(rule\ impI)$   
 $\mathbf{apply}(simp\ add:Mem\_pool\_alloc\_rely-def)$   
 $\mathbf{apply}(simp\ add:lvars-nochange-rel-def\ lvars-nochange-def)\ \mathbf{apply}\ smt$   
**done**



**lemma** *mp-alloc-precond3-stb*: *stable (mp-alloc-precond3 t p tm) (Mem-pool-alloc-rely t)*  
**apply**(*rule stable-int2*)  
**apply**(*simp add:mp-alloc-precond2-stb*)  
**apply**(*simp add:mp-alloc-precond3-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond4 t p tm*  $\equiv$   
*mp-alloc-precond2 t p tm*  $\cap \{\text{'end} t \geq 0\}$

**lemma** *mp-alloc-precond4-ext-stb*: *stable ( $\{\text{'end} t \geq 0\}$ ) (Mem-pool-alloc-rely t)*  
**apply**(*simp add:stable-def*)  
**done**

**lemma** *mp-alloc-precond4-stb*: *stable (mp-alloc-precond4 t p tm) (Mem-pool-alloc-rely t)*  
**apply**(*rule stable-int2*)  
**apply**(*simp add:mp-alloc-precond2-stb*)  
**using** *mp-alloc-precond4-ext-stb* **apply** *auto*  
**done**

**abbreviation** *mp-alloc-precond5 t p tm*  $\equiv$   
*mp-alloc-precond4 t p tm*  $\cap \{\text{'mempoolalloc-ret } t = \text{None}\}$

**lemma** *mp-alloc-precond5-ext-stb*: *stable ( $\{\text{'mempoolalloc-ret } t = \text{None}\}$ ) (Mem-pool-alloc-rely t)*  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*  
**done**

**lemma** *mp-alloc-precond5-stb*: *stable (mp-alloc-precond5 t p tm) (Mem-pool-alloc-rely t)*  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond4-stb* **apply** *auto[1]*  
**apply**(*simp add:mp-alloc-precond5-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond6 t p tm*  $\equiv$   
*mp-alloc-precond5 t p tm*  $\cap \{\text{'ret } t = \text{ESIZEERR}\}$

**lemma** *mp-alloc-precond6-ext-stb*: *stable ( $\{\text{'ret } t = \text{ESIZEERR}\}$ ) (Mem-pool-alloc-rely t)*  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*

done

**lemma** *mp-alloc-precond6-stb*: *stable (mp-alloc-precond6 t p tm) (Mem-pool-alloc-rely t)*  
 apply(rule *stable-int2*)  
 using *mp-alloc-precond5-stb* apply auto[1]  
 apply(simp add:*mp-alloc-precond6-ext-stb*)  
 done

**abbreviation** *mp-alloc-precond7-ext t p sz timeout*  $\equiv$   
 $\{s. (rf\ s\ t \longrightarrow (timeout = FOREVER \longrightarrow (ret\ s\ t = ESIZEERR \wedge mempoolalloc-ret\ s\ t = None$   
 $\vee ret\ s\ t = OK \wedge (\exists\ mblk. mempoolalloc-ret\ s\ t = Some$   
 $mblk \wedge alloc-memblk-valid\ s\ p\ sz\ mblk)))$   
 $\wedge (timeout = NOWAIT \longrightarrow ((ret\ s\ t = ENOMEM \vee ret\ s\ t = ESIZEERR)$   
 $\wedge mempoolalloc-ret\ s\ t = None)$   
 $\vee (ret\ s\ t = OK \wedge (\exists\ mblk. mempoolalloc-ret\ s\ t =$   
 $Some\ mblk \wedge alloc-memblk-valid\ s\ p\ sz\ mblk)))$   
 $\wedge (timeout > 0 \longrightarrow ((ret\ s\ t = ETIMEOUT \vee ret\ s\ t = ESIZEERR) \wedge$   
 $mempoolalloc-ret\ s\ t = None)$   
 $\vee (ret\ s\ t = OK \wedge (\exists\ mblk. mempoolalloc-ret\ s\ t = Some$   
 $mblk \wedge alloc-memblk-valid\ s\ p\ sz\ mblk))))$   
 $\wedge (\neg rf\ s\ t \longrightarrow mempoolalloc-ret\ s\ t = None)$   
 $\wedge (timeout = FOREVER \longrightarrow tmout\ s\ t = FOREVER)\}$

**abbreviation** *mp-alloc-precond7 t p sz timeout*  $\equiv$   
*mp-alloc-precond1 t p timeout*  $\cap$  *mp-alloc-precond7-ext t p sz timeout*

**abbreviation** *mp-alloc-precond7-inv t p sz timeout*  $\alpha \equiv$   
*mp-alloc-precond7 t p sz timeout*  
 $\cap \{ \alpha = (if\ 'rf\ t \vee 'mempoolalloc-ret\ t \neq None\ then\ 0\ (*\ if\ timeout = 0$   
 $(NOWAIT),\ rf\ is\ true\ *)$   
 $else\ if\ timeout > 0\ then\ 'end\ t - 'tick$   
 $(*\ in\ rely\ cond,\ tick' \geq tick,\ thus\ convergent\ \beta \leq \alpha,\ not\ < \alpha,\ thus$   
 $not\ absolutely\ convergent\ *)$   
 $else\ 1)$   
 $(*\ cannot\ find\ convergent\ \alpha\ for\ FOREVER,\ so\ just\ set\ 1\ *) \}$

**lemma** *mp-alloc-precond7-ext-stb*: *stable (mp-alloc-precond7-ext t p sz timeout)*  
*(Mem-pool-alloc-rely t)*  
 apply(simp add:*stable-def*) apply(rule *allI*) apply(rule *impI*) apply(rule *allI*)  
 apply(rule *impI*)  
 using *mp-alloc-post-stb*  
 apply(simp add:*Mem-pool-alloc-rely-def*)  
 apply(simp add:*lvars-nochange-rel-def lvars-nochange-def*)  
 apply(case-tac  $x = y$ )  
 apply simp apply clarify  
 apply(simp add:*alloc-memblk-valid-def gvars-conf-def gvars-conf-stable-def*)

**done**

**lemma** *mp-alloc-precond7-stb: stable (mp-alloc-precond7 t p sz timeout) (Mem-pool-alloc-rely t)*  
  **apply**(rule *stable-int2*)  
  **using** *mp-alloc-precond1-stb* **apply** *auto[1]*  
  **apply**(simp *add:mp-alloc-precond7-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-0 t p sz tm*  $\equiv$   
  *mp-alloc-precond7 t p sz tm*  $\cap \{\neg 'rf\ t\}$

**lemma** *mp-alloc-precond1-0-ext-stb: stable  $\{\neg 'rf\ t\}$  (Mem-pool-alloc-rely t)*  
**apply**(simp *add:stable-def*) **apply**(rule *allI*) **apply**(rule *impI*) **apply**(rule *allI*)  
**apply**(rule *impI*)  
  **apply**(simp *add:Mem-pool-alloc-rely-def*)  
  **apply**(simp *add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*  
**done**

**lemma** *mp-alloc-precond1-0-stb: stable (mp-alloc-precond1-0 t p sz tm) (Mem-pool-alloc-rely t)*  
  **apply**(rule *stable-int2*)  
  **using** *mp-alloc-precond7-stb* **apply** *auto[1]*  
  **apply**(simp *add:mp-alloc-precond1-0-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-1 t p sz tm*  $\equiv$   
  *mp-alloc-precond1-0 t p sz tm*  $(*\cap \{\neg 'blk\ t = 0\}*)$

**lemma** *mp-alloc-precond1-1-stb: stable (mp-alloc-precond1-1 t p sz tm) (Mem-pool-alloc-rely t)*  
  **using** *mp-alloc-precond1-0-stb* **by** *simp*

**abbreviation** *mp-alloc-precond1-2 t p sz tm*  $\equiv$   
  *mp-alloc-precond1-1 t p sz tm*  $\cap \{\neg 'alloc-lsize-r\ t = False\}$

**lemma** *mp-alloc-precond1-2-stb: stable (mp-alloc-precond1-2 t p sz tm) (Mem-pool-alloc-rely t)*  
  **apply**(rule *stable-int2*)  
  **using** *mp-alloc-precond1-1-stb* **apply** *auto[1]*  
  **apply**(simp *add:stable-def*) **apply**(rule *allI*) **apply**(rule *impI*) **apply**(rule *allI*)  
**apply**(rule *impI*)  
  **apply**(simp *add:Mem-pool-alloc-rely-def*)  
  **apply**(simp *add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*  
**done**

**abbreviation** *mp-alloc-precond1-3 t p sz tm*  $\equiv$   
  *mp-alloc-precond1-2 t p sz tm*  $\cap \{\neg 'alloc-l\ t = -1\}$

**lemma** *mp-alloc-precond1-3-ext-stb*: *stable*  $\{\text{'alloc-l } t = -1\}$  (*Mem-pool-alloc-rely* *t*)  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*  
**done**

**lemma** *mp-alloc-precond1-3-stb*: *stable* (*mp-alloc-precond1-3 t p sz tm*) (*Mem-pool-alloc-rely* *t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-2-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-3-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-4 t p sz tm*  $\equiv$   
*mp-alloc-precond1-3 t p sz tm*  $\cap \{\text{'free-l } t = -1\}$

**lemma** *mp-alloc-precond1-4-ext-stb*: *stable*  $\{\text{'free-l } t = -1\}$  (*Mem-pool-alloc-rely* *t*)  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def*) **apply** *smt*  
**done**

**lemma** *mp-alloc-precond1-4-stb*: *stable* (*mp-alloc-precond1-4 t p sz tm*) (*Mem-pool-alloc-rely* *t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-3-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-4-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-5 t p sz tm*  $\equiv$   
*mp-alloc-precond1-4 t p sz tm*  $\cap \{\text{'lsizes } t = [\text{ALIGN}_4 \text{ (max-sz ('mem-pool-info } p))]\}$

**lemma** *mp-alloc-precond1-5-ext-stb*: *stable*  $\{\text{'lsizes } t = [\text{ALIGN}_4 \text{ (max-sz ('mem-pool-info } p))]\}$  (*Mem-pool-alloc-rely* *t*)  
**apply**(*simp add:stable-def*) **apply**(*rule allI*) **apply**(*rule impI*) **apply**(*rule allI*)  
**apply**(*rule impI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *simp* **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
*gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-5-stb*: *stable* (*mp-alloc-precond1-5 t p sz tm*) (*Mem-pool-alloc-rely* *t*)

$t$ )  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-4-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-5-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-6-ext t p sz tm*  $\equiv$   
 $\{\!\{ (\forall ii < \text{length } ('lsizes\ t),\ 'lsizes\ t\ !\ ii = (\text{ALIGN4 } (\text{max-sz } ('mem-pool-info\ p))))$   
 $\text{div } (4 \wedge ii))$   
 $\wedge \text{length } ('lsizes\ t) \leq n\text{-levels } ('mem-pool-info\ p)$   
 $\wedge ('i\ t = 0 \longrightarrow 'alloc-l\ t = -1 \wedge 'free-l\ t = -1 \wedge \text{length } ('lsizes\ t) = 1)$   
 $\wedge 'i\ t \leq n\text{-levels } ('mem-pool-info\ p)$   
 $\wedge -1 \leq 'free-l\ t \wedge 'free-l\ t \leq \text{int } ('i\ t) - 1 \wedge 'free-l\ t \leq 'alloc-l\ t$   
 $\wedge 'alloc-l\ t = \text{int } ('i\ t) - 1$   
 $\wedge ('alloc-l\ t \geq 0 \longrightarrow (\forall ii.\ ii \leq \text{nat } ('alloc-l\ t) \longrightarrow 'lsizes\ t\ !\ ii \geq \text{sz}))$   
 $\wedge (\neg 'alloc-lsize-r\ t \longrightarrow ('i\ t = 0 \longrightarrow \text{length } ('lsizes\ t) = 1) \wedge ('i\ t > 0 \longrightarrow$   
 $\text{length } ('lsizes\ t) = 'i\ t))$   
 $\wedge ('alloc-lsize-r\ t \longrightarrow \text{length } ('lsizes\ t) = 'i\ t + 1 \wedge 'i\ t < n\text{-levels}$   
 $('mem-pool-info\ p) \wedge 'lsizes\ t\ !\ ('i\ t) < \text{sz}) \!\}$

**abbreviation** *mp-alloc-precond1-6 t p sz tm*  $\equiv$   
 $\text{mp-alloc-precond1-1 } t\ p\ \text{sz } tm \cap \text{mp-alloc-precond1-6-ext } t\ p\ \text{sz } tm$

**abbreviation** *mp-alloc-lsizeloop- $\alpha$ -cond t p  $\alpha$*   $\equiv$   
 $\{\!\{ \alpha = (\text{if } 'alloc-lsize-r\ t\ (*'lsizes\ t\ !\ ('i\ t) < \text{sz } *)$   
 $\text{then } 0 \text{ else } n\text{-levels } ('mem-pool-info\ p) - 'i\ t) \!\}$

**abbreviation** *mp-alloc-lsizestm-loopinv t p sz tm  $\alpha$*   $\equiv$   
 $\text{mp-alloc-precond1-6 } t\ p\ \text{sz } tm \cap \text{mp-alloc-lsizeloop-}\alpha\text{-cond } t\ p\ \alpha$

**abbreviation** *mp-alloc-lsizestm-loopcond t p*  $\equiv \{\!\{ 'i\ t < n\text{-levels } ('mem-pool-info\ p) \wedge \neg 'alloc-lsize-r\ t \!\}$

**lemma** *lsizestm-loopinv-imp-precond:*  
 $\text{mp-alloc-lsizestm-loopinv } t\ p\ \text{sz } tm\ \alpha \subseteq \text{mp-alloc-precond1-6 } t\ p\ \text{sz } tm$   
**by** *auto*

**lemma** *lsizestm-loopinv- $\alpha$ gt0-imp-loopcond:*  
 $\text{mp-alloc-lsizestm-loopinv } t\ p\ \text{sz } tm\ \alpha \cap \{\!\{ \alpha > 0 \!\} \subseteq \text{mp-alloc-lsizestm-loopcond } t\ p$   
**by** *clarsimp*

**lemma** *lsizestm-loopinv- $\alpha$ eq0-imp-notloopcond:*  
 $\text{mp-alloc-lsizestm-loopinv } t\ p\ \text{sz } tm\ \alpha \cap \{\!\{ \alpha = 0 \!\} \subseteq \neg \text{mp-alloc-lsizestm-loopcond } t\ p$   
**by** *clarsimp*

**lemma** *lsizestm-loopinv- $\alpha$ eq0-imp-notloopcond2:*

$mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ 0 \subseteq -\ mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p$   
**by** *clarsimp*

**lemma** *lsizestm-pre-loopcond-imp-loopinv- $\alpha$ gt0*:  
 $x \in mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ tm \cap mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p \implies$   
 $\exists \alpha. x \in mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ \alpha \cap \{\alpha > 0\}$   
**by** *clarsimp*

**lemma** *lsizestm-pre-notloopcond-imp-loopinv- $\alpha$ eq0*:  
 $x \in mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ tm \cap -\ mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p \implies$   
 $x \in mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ 0$   
**apply** *clarsimp*  
**apply**(*rule conjI*)  
**apply** *clarify* **apply** *simp*  
**apply** *clarify* **apply** *simp*  
**done**

**lemma** *lsizestm-pre-notloopcond-imp-loopinv- $\alpha$ eq0'*:  
 $mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ tm \cap -\ mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p$   
 $\subseteq mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ 0$   
**apply** *clarsimp*  
**apply**(*rule conjI*) **apply** *clarify*  
**apply**(*rule conjI*) **apply** *clarify*  
**apply**(*rule conjI*) **apply** *clarify*  
**apply**(*rule conjI*) **apply** *clarify*  
**apply**(*rule conjI*) **apply** *clarify*  
**apply** *clarify* **apply** *simp*  
**done**

**lemma** *lsizestm-pre-notloopcond-eq-loopinv- $\alpha$ eq0*:  
 $mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ tm \cap -\ mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p$   
 $= mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ 0$   
**apply**(*rule subset-antisym*)  
**using** *lsizestm-pre-notloopcond-imp-loopinv- $\alpha$ eq0'*[*of t p tm sz*] **apply** *blast*  
**apply**(*rule Int-greatest*)  
**using** *lsizestm-loopinv-imp-precond*[*of t p tm sz 0*] **apply** *blast*  
**using** *lsizestm-loopinv- $\alpha$ eq0-imp-notloopcond2*[*of t p tm sz*] **apply** *blast*  
**done**

**lemma** *lsizeloop-inv-cond-eq- $\alpha$ gt0*:  
 $mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ \alpha \cap mp\text{-}alloc\text{-}lsizestm\text{-}loopcond\ t\ p$   
 $= mp\text{-}alloc\text{-}lsizestm\text{-}loopinv\ t\ p\ sz\ tm\ \alpha \cap \{\alpha > 0\}$   
**apply**(*rule subset-antisym*)  
**apply**(*rule Int-greatest*)  
**apply** *fast*  
**apply** *clarify* **apply** *auto*[1]  
**apply**(*rule Int-greatest*)  
**apply** *fast*  
**apply** *clarsimp*

**done**

**lemma** *mp-alloc-precond1-6-ext-stb*: *stable (mp-alloc-precond1-6-ext t p sz tm) (Mem-pool-alloc-rely t)*  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto[1]* **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**apply**(*rule conjI*) **apply** *clarify* **apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**apply** *clarify* **apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-6-stb*: *stable (mp-alloc-precond1-6 t p sz tm) (Mem-pool-alloc-rely t)*  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-1-stb* **apply** *auto[1]*  
**using** *mp-alloc-precond1-6-ext-stb* **apply** *auto*  
**done**

**lemma** *mp-alloc-lsizeloop-α-cond-stb*: *stable (mp-alloc-lsizeloop-α-cond t p α) (Mem-pool-alloc-rely t)*  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*) **apply** *auto*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*) +  
**done**

**lemma** *mp-alloc-lsizestm-loopinv-stb*: *stable (mp-alloc-lsizestm-loopinv t p sz tm α) (Mem-pool-alloc-rely t)*  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-6-stb* **apply** *fast*  
**using** *mp-alloc-lsizeloop-α-cond-stb* **apply** *fast*  
**done**

**lemma** *mp-alloc-lsizestm-loopinv-presv-rely*:  
 $s \in \text{mp-alloc-lsizestm-loopinv } t \text{ } p \text{ } sz \text{ } tm \text{ } \alpha \implies (s, r) \in \text{Mem-pool-alloc-rely } t \implies \exists \beta \leq \alpha.$   
 $r \in \text{mp-alloc-lsizestm-loopinv } t \text{ } p \text{ } sz \text{ } tm \text{ } \beta$   
**apply**(*rule exI[where x=α]*)  
**apply**(*rule conjI*) **apply** *fast*  
**using** *mp-alloc-lsizestm-loopinv-stb[of t p tm sz α]* **apply**(*unfold stable-def*) **apply** *meson*  
**done**

**abbreviation** *mp-alloc-precond1-6-1 t p sz tm α*  $\equiv$   
 $\text{mp-alloc-lsizestm-loopinv } t \text{ } p \text{ } sz \text{ } tm \text{ } \alpha \cap \text{mp-alloc-lsizestm-loopcond } t \text{ } p$

**lemma** *mp-alloc-precond1-6-1-ext-stb*: *stable (mp-alloc-lsizestm-loopcond t p) (Mem-pool-alloc-rely t)*

$t$ )  
**apply**(simp add:stable-def) **apply**(rule allI) **apply**(rule impI) **apply**(rule allI)  
**apply**(rule impI)  
**apply**(simp add:Mem-pool-alloc-rely-def)  
**apply**(case-tac  $x=y$ ) **apply** auto[1] **apply** clarify  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
gvars-conf-def)  
**done**

**lemma** mp-alloc-precond1-6-1-stb: stable (mp-alloc-precond1-6-1  $t$   $p$  sz  $tm$   $\alpha$ ) (Mem-pool-alloc-rely  
 $t$ )  
**apply**(rule stable-int2)  
**using** mp-alloc-lsizestm-loopinv-stb **apply** auto[1]  
**apply**(simp add:mp-alloc-precond1-6-1-ext-stb)  
**done**

**abbreviation** mp-alloc-precond1-6-10  $t$   $p$  sz  $tm$   $\alpha \equiv$   
 $mp\text{-}alloc\text{-}precond1\text{-}6\text{-}1\ t\ p\ sz\ tm\ \alpha \cap \{\ 'i\ t > 0 \}$

**lemma** mp-alloc-precond1-6-10-ext-stb: stable ( $\{\ 'i\ t > 0 \}$ ) (Mem-pool-alloc-rely  $t$ )  
**apply**(simp add:stable-def) **apply**(rule allI) **apply**(rule impI) **apply**(rule allI)  
**apply**(rule impI)  
**apply**(simp add:Mem-pool-alloc-rely-def)  
**apply**(case-tac  $x=y$ ) **apply** auto[1] **apply** clarify  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
gvars-conf-def)  
**done**

**lemma** mp-alloc-precond1-6-10-stb: stable (mp-alloc-precond1-6-10  $t$   $p$  sz  $tm$   $\alpha$ )  
(Mem-pool-alloc-rely  $t$ )  
**apply**(rule stable-int2)  
**using** mp-alloc-precond1-6-1-stb **apply** auto[1]  
**apply**(simp add:mp-alloc-precond1-6-10-ext-stb)  
**done**

**abbreviation** mp-alloc-precond1-6-11  $t$   $p$  sz  $tm$   $\alpha \equiv$   
 $mp\text{-}alloc\text{-}precond1\text{-}6\text{-}1\ t\ p\ sz\ tm\ \alpha \cap -\{\ 'i\ t > 0 \}$

**lemma** mp-alloc-precond1-6-11-ext-stb: stable ( $-\{\ 'i\ t > 0 \}$ ) (Mem-pool-alloc-rely  
 $t$ )  
**apply**(simp add:stable-def) **apply**(rule allI) **apply**(rule impI) **apply**(rule allI)  
**apply**(rule impI)  
**apply**(simp add:Mem-pool-alloc-rely-def)  
**apply**(case-tac  $x=y$ ) **apply** auto[1] **apply** clarify  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
gvars-conf-def)  
**done**

**lemma** mp-alloc-precond1-6-11-stb: stable (mp-alloc-precond1-6-11  $t$   $p$  sz  $tm$   $\alpha$ )



(Mem-pool-alloc-rely t)  
 apply(rule stable-int2)  
 using mp-alloc-precond1-6-1-stb apply auto[1]  
 apply(simp add:mp-alloc-precond1-6-11-ext-stb)  
 done

**abbreviation** mp-alloc-precond1-6-2-ext t p sz tm  $\alpha \equiv$   
 $\{\!(\forall ii < \text{length } ('lsizes\ t). \ 'lsizes\ t\ !\ ii = (\text{ALIGN}_4\ (\text{max-sz } ('mem\text{-pool-info } p)))$   
 $\text{div } (4 \wedge ii))$   
 $\wedge \text{length } ('lsizes\ t) \leq n\text{-levels } ('mem\text{-pool-info } p)$   
 $\wedge ('i\ t = 0 \longrightarrow 'alloc\text{-l } t = -1 \wedge 'free\text{-l } t = -1 \wedge \text{length } ('lsizes\ t) = 1)$   
 $\wedge 'i\ t \leq n\text{-levels } ('mem\text{-pool-info } p)$   
 $\wedge -1 \leq 'free\text{-l } t \wedge 'free\text{-l } t \leq \text{int } ('i\ t) - 1 \wedge 'free\text{-l } t \leq 'alloc\text{-l } t$   
 $\wedge 'alloc\text{-l } t = \text{int } ('i\ t) - 1$   
 $\wedge ('alloc\text{-l } t \geq 0 \longrightarrow (\forall ii. ii \leq \text{nat } ('alloc\text{-l } t) \longrightarrow 'lsizes\ t\ !\ ii \geq \text{sz}))$   
 $\wedge (\neg 'alloc\text{-lsize-r } t \longrightarrow ('i\ t = 0 \longrightarrow \text{length } ('lsizes\ t) = 1) \wedge ('i\ t > 0 \longrightarrow$   
 $\text{length } ('lsizes\ t) = 'i\ t + 1))$   
 $(\ast \text{ here } 'i\ t + 1 \text{ is different from } mp\text{-alloc-precond1-6-ext, } \ast)$   
 $\wedge ('alloc\text{-lsize-r } t \longrightarrow \text{length } ('lsizes\ t) = 'i\ t + 1 \wedge 'i\ t < n\text{-levels}$   
 $('mem\text{-pool-info } p) \wedge 'lsizes\ t\ !\ ('i\ t) < \text{sz})\!\}$   
 $\cap mp\text{-alloc-lsizeloop-}\alpha\text{-cond } t\ p\ \alpha$

**abbreviation** mp-alloc-precond1-6-2 t p sz tm  $\alpha \equiv$   
 $mp\text{-alloc-precond1-2 } t\ p\ \text{sz } tm \cap mp\text{-alloc-precond1-6-2-ext } t\ p\ \text{sz } tm\ \alpha$

**lemma** mp-alloc-precond1-6-2-ext-stb: stable (mp-alloc-precond1-6-2-ext t p sz tm  
 $\alpha$ ) (Mem-pool-alloc-rely t)  
 apply(rule stable-int2)  
 apply(simp add:stable-def) apply clarify  
 apply(simp add:Mem-pool-alloc-rely-def)  
 apply(case-tac x=y) apply auto[1] apply clarify  
 apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
 gvars-conf-def)  
 apply(rule conjI) apply clarify apply(simp add:lvars-nochange-rel-def lvars-nochange-def  
 gvars-conf-stable-def gvars-conf-def)  
 apply clarify apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
 gvars-conf-def)  
 using mp-alloc-lsizeloop- $\alpha$ -cond-stb apply fast  
 done

**lemma** mp-alloc-precond1-6-2-stb: stable (mp-alloc-precond1-6-2 t p sz tm  $\alpha$ ) (Mem-pool-alloc-rely  
 t)  
 apply(rule stable-int2)  
 using mp-alloc-precond1-2-stb apply auto[1]  
 using mp-alloc-precond1-6-2-ext-stb apply auto  
 done

**abbreviation** mp-alloc-precond1-6-20 t p sz tm  $\alpha \equiv$   
 $mp\text{-alloc-precond1-6-2 } t\ p\ \text{sz } tm\ \alpha \cap \{\!'lsizes\ t\ !\ 'i\ t < \text{sz}\!\}$

**lemma** *mp-alloc-precond1-6-20-ext-stb*: *stable* ( $\{\text{'lsizes } t ! 'i\ t < sz\}$ ) (*Mem-pool-alloc-rely* *t*)  
**apply** (*simp add:stable-def*) **apply** *clarify*  
**apply** (*simp add:Mem-pool-alloc-rely-def*)  
**apply** (*case-tac x=y*) **apply** *auto[1]* **apply** *clarify*  
**apply** (*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
*gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-6-20-stb*: *stable* (*mp-alloc-precond1-6-20 t p sz tm*  $\alpha$ )  
*(Mem-pool-alloc-rely t)*  
**apply** (*rule stable-int2*)  
**using** *mp-alloc-precond1-6-2-stb* **apply** *auto[1]*  
**apply** (*simp add:mp-alloc-precond1-6-20-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-6-21 t p sz tm*  $\alpha \equiv$   
*mp-alloc-precond1-6-2 t p sz tm*  $\alpha \cap - \{\text{'lsizes } t ! 'i\ t < sz\}$

**lemma** *mp-alloc-precond1-6-21-ext-stb*: *stable* ( $-\ \{\text{'lsizes } t ! 'i\ t < sz\}$ ) (*Mem-pool-alloc-rely* *t*)  
**apply** (*simp add:stable-def*) **apply** *clarify*  
**apply** (*simp add:Mem-pool-alloc-rely-def*)  
**apply** (*case-tac x=y*) **apply** *auto[1]* **apply** *clarify*  
**apply** (*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
*gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-6-21-stb*: *stable* (*mp-alloc-precond1-6-21 t p sz tm*  $\alpha$ )  
*(Mem-pool-alloc-rely t)*  
**apply** (*rule stable-int2*)  
**using** *mp-alloc-precond1-6-2-stb* **apply** *auto[1]*  
**apply** (*simp add:mp-alloc-precond1-6-21-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-6-21-1-ext t p sz tm*  $\alpha \equiv$   
 $\{\text{'lsizes } t ! ii = (ALIGN4\ (max\text{-}sz\ ('mem\text{-}pool\text{-}info\ p)))\ \text{div}\ (4\ ^\wedge\ ii)\}$   
 $\wedge\ length\ ('lsizes\ t) \leq n\text{-}levels\ ('mem\text{-}pool\text{-}info\ p)$   
 $\wedge\ ('i\ t = 0 \longrightarrow 'alloc\text{-}l\ t = 0 \wedge 'free\text{-}l\ t = -1 \wedge length\ ('lsizes\ t) = 1)$   
 $\wedge\ 'i\ t \leq n\text{-}levels\ ('mem\text{-}pool\text{-}info\ p)$   
 $\wedge\ -1 \leq 'free\text{-}l\ t \wedge 'free\text{-}l\ t \leq int\ ('i\ t) - 1 \wedge 'free\text{-}l\ t \leq 'alloc\text{-}l\ t$   
 $\wedge\ 'alloc\text{-}l\ t = int\ ('i\ t)$   
 $\wedge\ ('alloc\text{-}l\ t \geq 0 \longrightarrow (\forall ii. ii < nat\ ('alloc\text{-}l\ t) \longrightarrow 'lsizes\ t ! ii \geq sz))$   
 $\wedge\ (\neg 'alloc\text{-}l\text{-}r\ t \longrightarrow ('i\ t = 0 \longrightarrow length\ ('lsizes\ t) = 1) \wedge ('i\ t > 0 \longrightarrow$   
 $length\ ('lsizes\ t) = 'i\ t + 1))$   
 $\wedge\ ('alloc\text{-}l\text{-}r\ t \longrightarrow length\ ('lsizes\ t) = 'i\ t + 1 \wedge 'i\ t < n\text{-}levels$   
 $('mem\text{-}pool\text{-}info\ p) \wedge 'lsizes\ t ! ('i\ t) < sz)$

$\wedge \neg 'lsizes\ t ! 'i\ t < sz \} \cap mp\_alloc\_lsizeloop\_alpha\_cond\ t\ p\ \alpha$

**abbreviation**  $mp\_alloc\_precond1-6-21-1\ t\ p\ sz\ tm\ \alpha \equiv$   
 $mp\_alloc\_precond1-2\ t\ p\ sz\ tm \cap mp\_alloc\_precond1-6-21-1-ext\ t\ p\ sz\ tm\ \alpha$

**lemma**  $mp\_alloc\_precond1-6-21-1-ext-stb$ : *stable* ( $mp\_alloc\_precond1-6-21-1-ext\ t\ p\ sz\ tm\ \alpha$ ) (*Mem-pool-alloc-rely*  $t$ )  
**apply**(*rule stable-int2*)  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto*[1] **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**apply** *clarify* **apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**using**  $mp\_alloc\_lsizeloop\_alpha\_cond-stb$  **apply** *fast*  
**done**

**lemma**  $mp\_alloc\_precond1-6-21-1-stb$ : *stable* ( $mp\_alloc\_precond1-6-21-1\ t\ p\ sz\ tm\ \alpha$ ) (*Mem-pool-alloc-rely*  $t$ )  
**apply**(*rule stable-int2*)  
**using**  $mp\_alloc\_precond1-2-stb$  **apply** *auto*[1]  
**using**  $mp\_alloc\_precond1-6-21-1-ext-stb$  **apply** *auto*  
**done**

**abbreviation**  $mp\_alloc\_precond1-6-21-2-ext\ t\ p\ sz\ tm\ \alpha \equiv$   
 $\{(\forall ii < length\ ('lsizes\ t). 'lsizes\ t ! ii = (ALIGN4\ (max-sz\ ('mem-pool-info\ p))) \div (4\ ^\wedge\ ii))$   
 $\wedge length\ ('lsizes\ t) \leq n-levels\ ('mem-pool-info\ p)$   
 $\wedge ('i\ t = 0 \longrightarrow 'alloc-l\ t = 0 \wedge length\ ('lsizes\ t) = 1)$   
 $\wedge 'i\ t \leq n-levels\ ('mem-pool-info\ p)$   
 $\wedge -1 \leq 'free-l\ t \wedge 'free-l\ t \leq int\ ('i\ t) \wedge 'free-l\ t \leq 'alloc-l\ t$   
 $\wedge 'alloc-l\ t = int\ ('i\ t)$   
 $\wedge ('alloc-l\ t \geq 0 \longrightarrow (\forall ii. ii < nat\ ('alloc-l\ t) \longrightarrow 'lsizes\ t ! ii \geq sz))$   
 $\wedge (\neg 'alloc-lsize-r\ t \longrightarrow ('i\ t = 0 \longrightarrow length\ ('lsizes\ t) = 1) \wedge ('i\ t > 0 \longrightarrow$   
 $length\ ('lsizes\ t) = 'i\ t + 1))$   
 $\wedge ('alloc-lsize-r\ t \longrightarrow length\ ('lsizes\ t) = 'i\ t + 1 \wedge 'i\ t < n-levels$   
 $( 'mem-pool-info\ p) \wedge 'lsizes\ t ! ('i\ t) < sz)$   
 $\wedge \neg 'lsizes\ t ! 'i\ t < sz \} \cap mp\_alloc\_lsizeloop\_alpha\_cond\ t\ p\ \alpha$

**abbreviation**  $mp\_alloc\_precond1-6-21-2\ t\ p\ sz\ tm\ \alpha \equiv$   
 $mp\_alloc\_precond1-2\ t\ p\ sz\ tm \cap mp\_alloc\_precond1-6-21-2-ext\ t\ p\ sz\ tm\ \alpha$

**lemma**  $mp\_alloc\_precond1-6-21-2-ext-stb$ : *stable* ( $mp\_alloc\_precond1-6-21-2-ext\ t\ p\ sz\ tm\ \alpha$ ) (*Mem-pool-alloc-rely*  $t$ )  
**apply**(*rule stable-int2*)  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto*[1] **apply** *clarify*

```

    apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
    apply clarify apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
using mp-alloc-lsizeloop-α-cond-stb apply fast
done

```

```

lemma mp-alloc-precond1-6-21-2-stb: stable (mp-alloc-precond1-6-21-2 t p sz tm
α) (Mem-pool-alloc-rely t)
    apply(rule stable-int2)
    using mp-alloc-precond1-2-stb apply auto[1]
    using mp-alloc-precond1-6-21-2-ext-stb apply auto
done

```

```

abbreviation mp-alloc-precond1-7 t p sz tm ≡
    mp-alloc-precond1-6 t p sz tm ∩ {i t ≥ n-levels ('mem-pool-info p) ∨ 'alloc-lsize-r
t }

```

```

lemma mp-alloc-precond1-7-ext-stb: stable ({i t ≥ n-levels ('mem-pool-info p) ∨
'alloc-lsize-r t }) (Mem-pool-alloc-rely t)
    apply(simp add:stable-def) apply clarify
    apply(rule conjI)
    apply clarify
    apply(simp add:Mem-pool-alloc-rely-def)
    apply(case-tac x=y) apply simp apply clarify
    apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)

```

```

    apply clarify
    apply(simp add:Mem-pool-alloc-rely-def)
    apply(case-tac x=y) apply simp apply clarify
    apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
done

```

```

lemma mp-alloc-precond1-7-stb: stable (mp-alloc-precond1-7 t p sz tm) (Mem-pool-alloc-rely
t)
    apply(rule stable-int2)
    using mp-alloc-precond1-6-stb apply auto[1]
    apply(simp add:mp-alloc-precond1-7-ext-stb)
done

```

```

abbreviation mp-alloc-precond1-70-ext t p sz tm ≡
    {(∀ ii < length ('lsizes t). 'lsizes t ! ii = (ALIGN4 (max-sz ('mem-pool-info
p))) div (4 ^ ii))
    ∧ length ('lsizes t) ≤ n-levels ('mem-pool-info p)
    ∧ 'alloc-l t < int (n-levels ('mem-pool-info p))
    ∧ -1 ≤ 'free-l t ∧ 'free-l t ≤ 'alloc-l t
    ∧ ('alloc-l t = -1 ∧ 'free-l t = -1 ∧ length ('lsizes t) = 1

```

$$\begin{aligned} & \vee ('alloc-l\ t \geq 0 \wedge (\forall ii. ii \leq \text{nat } ('alloc-l\ t) \longrightarrow 'lsizes\ t\ !\ ii \geq sz) \\ & \wedge (('alloc-l\ t = \text{int } (\text{length } ('lsizes\ t)) - 1) \wedge \text{length } ('lsizes\ t) = \\ n\text{-levels } ('mem\text{-pool-info } p) \\ & \vee 'alloc-l\ t = \text{int } (\text{length } ('lsizes\ t)) - 2 \wedge 'lsizes\ t\ !\ \text{nat } ('alloc-l \\ & t + 1) < sz))) \} \end{aligned}$$

**abbreviation**  $mp\text{-alloc-precond1-70}\ t\ p\ sz\ tm \equiv$   
 $mp\text{-alloc-precond1-1}\ t\ p\ sz\ tm \cap mp\text{-alloc-precond1-70-ext}\ t\ p\ sz\ tm$

**lemma**  $mp\text{-alloc-precond1-70-ext-stb}$ :  $\text{stable } (mp\text{-alloc-precond1-70-ext}\ t\ p\ sz\ tm)$   
 $(Mem\text{-pool-alloc-rely}\ t)$   
**apply**( $\text{simp add:stable-def}$ ) **apply**  $\text{clarify}$   
**apply**( $\text{simp add:Mem-pool-alloc-rely-def}$ )  
**apply**( $\text{case-tac } x=y$ ) **apply**  $\text{auto}[1]$  **apply**  $\text{clarify}$   
**apply**( $\text{simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def}$   
 $\text{gvars-conf-def}$ )  
**done**

**lemma**  $mp\text{-alloc-precond1-70-stb}$ :  $\text{stable } (mp\text{-alloc-precond1-70}\ t\ p\ sz\ tm)$   $(Mem\text{-pool-alloc-rely}\ t)$   
**apply**( $\text{rule stable-int2}$ )  
**using**  $mp\text{-alloc-precond1-1-stb}$  **apply**  $\text{auto}[1]$   
**using**  $mp\text{-alloc-precond1-70-ext-stb}$  **apply**  $\text{auto}$   
**done**

**lemma**  $precnd17-bl-170$ :  $mp\text{-alloc-precond1-7}\ t\ p\ sz\ tm \subseteq mp\text{-alloc-precond1-70}\ t\ p\ sz\ tm$   
**apply**  $\text{clarify}$  **apply**( $\text{case-tac } i\ x\ t = 0$ )  
**apply**  $\text{clarify}$  **apply**  $\text{auto}[1]$   
**apply**  $\text{clarify}$   
**apply**( $\text{rule IntI}$ ) **apply**  $\text{auto}[1]$  **apply**  $\text{clarify}$   
**apply**( $\text{rule conjI}$ ) **apply**  $\text{simp}$   
**apply**( $\text{rule conjI}$ ) **apply**  $\text{simp}$   
**apply**( $\text{rule conjI}$ ) **apply**  $\text{simp}$   
**apply**  $\text{simp}$   
**apply**( $\text{case-tac alloc-lsize-r } x\ t$ ) **apply**  $\text{auto}$   
**done**

**abbreviation**  $mp\text{-alloc-precond1-70-1}\ t\ p\ sz\ tm \equiv$   
 $mp\text{-alloc-precond1-70}\ t\ p\ sz\ tm \cap \{ 'alloc-l\ t < 0 \}$

**lemma**  $mp\text{-alloc-precond1-70-1-ext-stb}$ :  $\text{stable } (\{ 'alloc-l\ t < 0 \})$   $(Mem\text{-pool-alloc-rely}\ t)$   
**apply**( $\text{simp add:stable-def}$ ) **apply**  $\text{clarify}$   
**apply**( $\text{simp add:Mem-pool-alloc-rely-def}$ )  
**apply**( $\text{case-tac } x=y$ ) **apply**  $\text{auto}[1]$  **apply**  $\text{clarify}$   
**apply**( $\text{simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def}$   
 $\text{gvars-conf-def}$ )  
**done**

**lemma** *mp-alloc-precond1-70-1-stb*: *stable* (*mp-alloc-precond1-70-1 t p sz tm*) (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-70-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-70-1-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-70-2 t p sz tm*  $\equiv$   
*mp-alloc-precond1-70 t p sz tm*  $\cap - \{\text{'alloc-l } t < 0\}$

**lemma** *mp-alloc-precond1-70-2-ext-stb*: *stable* ( $- \{\text{'alloc-l } t < 0\}$ ) (*Mem-pool-alloc-rely t*)  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto*[1] **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-70-2-stb*: *stable* (*mp-alloc-precond1-70-2 t p sz tm*) (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-70-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-70-2-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-70-2-1 t p sz tm*  $\equiv$   
*mp-alloc-precond1-70-2 t p sz tm*  $\cap \{\text{'free-l } t < 0\}$

**lemma** *mp-alloc-precond1-70-2-1-ext-stb*: *stable* ( $\{\text{'free-l } t < 0\}$ ) (*Mem-pool-alloc-rely t*)  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto*[1] **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-70-2-1-stb*: *stable* (*mp-alloc-precond1-70-2-1 t p sz tm*) (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-70-2-stb* **apply** *auto*[1]  
**apply**(*simp add:mp-alloc-precond1-70-2-1-ext-stb*)  
**done**

**abbreviation** *mp-alloc-precond1-70-2-2 t p sz tm*  $\equiv$   
*mp-alloc-precond1-70-2 t p sz tm*  $\cap - \{\text{'free-l } t < 0\}$

**lemma** *mp-alloc-precond1-70-2-2-ext-stb*: *stable* ( $-\{ \text{'free-l } t < 0 \}$ ) (*Mem-pool-alloc-rely* *t*)  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *auto[1]* **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
*gvars-conf-def*)  
**done**

**lemma** *mp-alloc-precond1-70-2-2-stb*: *stable* (*mp-alloc-precond1-70-2-2 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-70-2-stb* **apply** *auto[1]*  
**apply**(*simp add:mp-alloc-precond1-70-2-2-ext-stb*)  
**done**

**lemma** *alloc-memblk-data-valid-stb*:  
 $\text{blk } x \ t = \text{buf } (\text{mem-pool-info } x \ p) +$   
 $\text{block-num } (\text{mem-pool-info } x \ p) (\text{blk } x \ t) (\text{lsizes } x \ t ! \text{nat } (\text{free-l } x \ t)) *$   
 $(\text{max-sz } (\text{mem-pool-info } x \ p) \text{div } 4 ^ \text{nat } (\text{free-l } x \ t)) \implies$   
 $\text{block-num } (\text{mem-pool-info } x \ p) (\text{blk } x \ t) (\text{lsizes } x \ t ! \text{nat } (\text{free-l } x \ t)) < n\text{-max}$   
 $(\text{mem-pool-info } x \ p) * 4 ^ \text{nat } (\text{free-l } x \ t) \implies$   
 $\text{allocating-node } x \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{free-l } x \ t), \text{block} = \text{block-num } (\text{mem-pool-info } x$   
 $p) (\text{blk } x \ t) (\text{lsizes } x \ t ! \text{nat } (\text{free-l } x \ t)),$   
 $\text{data} = \text{blk } x \ t) \implies$   
 $(x, y) \in \text{lvars-nochange-rel } t \implies$   
 $(x, y) \in \text{gvars-conf-stable} \implies$   
 $\text{alloc-memblk-data-valid } y \ p \ (\text{the } (\text{allocating-node } y \ t))$   
**apply**(*subgoal-tac blk x t = blk y t*)  
**prefer** 2 **apply**(*simp add: lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p)*)  
**prefer** 2 **apply**(*simp add: gvars-conf-stable-def gvars-conf-def*)  
**apply**(*subgoal-tac lsizes x t = lsizes y t*)  
**prefer** 2 **apply**(*simp add: lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*subgoal-tac free-l x t = free-l y t*)  
**prefer** 2 **apply**(*simp add: lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p)*)  
**prefer** 2 **apply**(*simp add: gvars-conf-stable-def gvars-conf-def*)  
**apply**(*subgoal-tac allocating-node x t = allocating-node y t*)  
**prefer** 2 **apply**(*simp add: lvars-nochange-rel-def lvars-nochange-def*)  
**apply** (*simp add: gvars-conf-def gvars-conf-stable-def*)  
**done**

**abbreviation** *mp-alloc-precond2-1-ext t p sz tm*  $\equiv$   
 $\{ (\text{'blk } t = \text{NULL} \wedge \text{'allocating-node } t = \text{None})$   
 $\vee (\text{'blk } t > \text{NULL} \wedge \text{'alloc-memblk-data-valid } p \ (\text{the } (\text{'allocating-node } t))$   
 $\wedge \text{'allocating-node } t = \text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{'free-l } t),$   
 $\text{block} = (\text{block-num } (\text{'mem-pool-info } p) (\text{'blk } t))$

$((\text{'lsizes } t)!(\text{nat } (\text{'free-l } t))))$ ,  
 $\text{data} = \text{'blk } t \mid$   
 $\wedge (\exists n. n < n\text{-max } (\text{'mem-pool-info } p) * (4 \wedge (\text{nat } (\text{'free-l } t))))$   
 $\wedge \text{'blk } t = \text{buf } (\text{'mem-pool-info } p) + n * (\text{max-sz } (\text{'mem-pool-info } p))$   
 $\text{div } (4 \wedge (\text{nat } (\text{'free-l } t))))))\}$

**abbreviation**  $\text{mp-alloc-precond2-1 } t \text{ } p \text{ } sz \text{ } tm \equiv$   
 $\{s. \text{inv } s\} \cap \{\text{'freeing-node } t = \text{None}\} \cap \{p \in \text{'mem-pools} \wedge tm \geq -1\} \cap$   
 $\text{mp-alloc-precond7-ext } t \text{ } p \text{ } sz \text{ } tm \cap \{\neg \text{'rf } t\}$   
 $\cap \text{mp-alloc-precond1-70-ext } t \text{ } p \text{ } sz \text{ } tm \cap -\{\text{'alloc-l } t < 0\}$   
 $\cap -\{\text{'free-l } t < 0\} \cap \text{mp-alloc-precond2-1-ext } t \text{ } p \text{ } sz \text{ } tm$

**term**  $\text{mp-alloc-precond2-1 } t \text{ } p \text{ } sz \text{ } tm$

**lemma**  $\text{mp-alloc-freenode-stb}$ :

$\text{stable } \{\text{'freeing-node } t = \text{None}\} (\text{Mem-pool-alloc-rely } t)$   
**apply**( $\text{simp add: stable-def Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**done**

**lemma**  $\text{mp-alloc-precond2-1-ext-stb}$ :  $\text{stable } (\text{mp-alloc-precond2-1-ext } t \text{ } p \text{ } sz \text{ } tm) (\text{Mem-pool-alloc-rely } t)$

**apply**( $\text{simp add: stable-def}$ ) **apply**  $\text{clarify}$   
**apply**( $\text{rule conjI}$ ) **apply**  $\text{clarify}$  **apply**( $\text{simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{rule impI}$ ) **apply**( $\text{rule allI}$ ) **apply**( $\text{rule impI}$ ) **apply**( $\text{rule disjI2}$ )  
**apply**( $\text{subgoal-tac buf } (\text{mem-pool-info } x \text{ } p) = \text{buf } (\text{mem-pool-info } y \text{ } p)$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def}$ )  
**apply**  $\text{metis}$   
**apply**( $\text{subgoal-tac free-l } x \text{ } t = \text{free-l } y \text{ } t$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac max-sz } (\text{mem-pool-info } x \text{ } p) = \text{max-sz } (\text{mem-pool-info } y \text{ } p)$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac blk } x \text{ } t = \text{blk } y \text{ } t$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac allocating-node } x \text{ } t = \text{allocating-node } y \text{ } t$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac lsize } x \text{ } t = \text{lsize } y \text{ } t$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac n-max } (\text{mem-pool-info } x \text{ } p) = \text{n-max } (\text{mem-pool-info } y \text{ } p)$ )  
**prefer 2** **apply**( $\text{simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def}$ )  
**apply**  $\text{smt}$   
**apply**( $\text{subgoal-tac block-num } (\text{mem-pool-info } x \text{ } p) (\text{blk } x \text{ } t) (\text{lsize } x \text{ } t ! \text{nat } (\text{free-l } x \text{ } t))$ )



$= \text{block-num } (\text{mem-pool-info } y \ p) \ (\text{blk } y \ t) \ (\text{lsizes } y \ t \ ! \ \text{nat } (\text{free-l } y \ t)))$   
**prefer 2 apply**(simp add: block-num-def Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)  
**apply smt**  
**done**

**lemma mp-alloc-precond2-1-stb**: stable (mp-alloc-precond2-1 t p sz tm) (Mem-pool-alloc-rely t)  
**apply**(rule stable-int2) **apply**(rule stable-int2) **apply**(rule stable-int2) **apply**(rule stable-int2)  
**apply**(rule stable-int2) **apply**(rule stable-int2) **apply**(rule stable-int2) **apply**(rule stable-int2)  
**apply** (simp add: stable-inv-alloc-rely1)  
**apply** (simp add: mp-alloc-freenode-stb)  
**apply** (simp add: mp-alloc-precond1-ext-stb)  
**apply** (simp add: mp-alloc-precond7-ext-stb)  
**apply** (simp add: mp-alloc-precond1-0-ext-stb)  
**using** mp-alloc-precond1-70-ext-stb **apply** blast  
**apply** (simp add: mp-alloc-precond1-70-2-ext-stb)  
**apply** (simp add: mp-alloc-precond1-70-2-2-ext-stb)  
**using** mp-alloc-precond2-1-ext-stb **by** blast

**abbreviation** mp-alloc-precond2-1-0 t p sz tm  $\equiv$   
mp-alloc-precond2-1 t p sz tm  $\cap \{\text{'blk } t = \text{NULL}\}$

**lemma mp-alloc-precond2-1-0-ext-stb**: stable ( $\{\text{'blk } t = \text{NULL}\}$ ) (Mem-pool-alloc-rely t)  
**apply**(simp add:stable-def) **apply** clarify  
**apply**(simp add:Mem-pool-alloc-rely-def)  
**apply**(case-tac x=y) **apply** auto[1] **apply** clarify  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def gvars-conf-def)  
**done**

**lemma mp-alloc-precond2-1-0-stb**: stable (mp-alloc-precond2-1-0 t p sz tm) (Mem-pool-alloc-rely t)  
**apply**(rule stable-int2)  
**using** mp-alloc-precond2-1-stb **apply** auto[1]  
**apply**(simp add:mp-alloc-precond2-1-0-ext-stb)  
**done**

**abbreviation** mp-alloc-precond2-1-1 t p sz tm  $\equiv$   
mp-alloc-precond2-1 t p sz tm  $\cap \neg\{\text{'blk } t = \text{NULL}\}$

**term** mp-alloc-precond2-1-1 t p sz tm

**lemma mp-alloc-precond2-1-1-ext-stb**: stable ( $\neg\{\text{'blk } t = \text{NULL}\}$ ) (Mem-pool-alloc-rely t)

**apply**(simp add:stable-def) **apply** clarify  
**apply**(simp add:Mem-pool-alloc-rely-def)  
**apply**(case-tac x=y) **apply** auto[1] **apply** clarify  
**apply**(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def  
gvars-conf-def)  
**done**

**lemma** mp-alloc-precond2-1-1-stb: stable (mp-alloc-precond2-1-1 t p sz tm) (Mem-pool-alloc-rely  
t)  
**apply**(rule stable-int2)  
**using** mp-alloc-precond2-1-1-stb **apply** auto[1]  
**apply**(simp add:mp-alloc-precond2-1-1-ext-stb)  
**done**

**abbreviation** mp-alloc-precond2-1-1-loopinv-ext t p sz tm  $\equiv$   
 $\neg \{ \text{'blk } t = \text{NULL} \} \cap \{ \text{'from-l } t \leq \text{'alloc-l } t \wedge \text{'from-l } t \geq \text{'free-l } t \wedge \text{'allocating-node}$   
 $t = \text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{'from-l } t),$   
 $\text{block} = \text{block-num } (\text{'mem-pool-info } p) (\text{'blk } t) ((\text{'lsizes}$   
 $t)!(\text{nat } (\text{'from-l } t))),$   
 $\text{data} = \text{'blk } t \}$   
 $\wedge \text{'alloc-memblk-data-valid } p (\text{the } (\text{'allocating-node } t))$   
 $\wedge (\exists n. n < n\text{-max } (\text{'mem-pool-info } p) * (4 \wedge (\text{nat } (\text{'from-l } t))))$   
 $\wedge \text{'blk } t = \text{buf } (\text{'mem-pool-info } p) + n * (\text{max-sz } (\text{'mem-pool-info } p))$   
 $\text{div } (4 \wedge (\text{nat } (\text{'from-l } t)))) \}$

**abbreviation** mp-alloc-precond2-1-1-loopinv t p sz tm  $\equiv$   
 $\{s. \text{inv } s\} \cap \{ \text{'freeing-node } t = \text{None} \} \cap \{ p \in \text{'mem-pools} \wedge \text{tm} \geq -1 \} \cap$   
 $\text{mp-alloc-precond7-ext } t p \text{ sz } \text{tm} \cap \{ \neg \text{'rf } t \}$   
 $\cap \text{mp-alloc-precond1-70-ext } t p \text{ sz } \text{tm} \cap \neg \{ \text{'alloc-l } t < 0 \}$   
 $\cap \neg \{ \text{'free-l } t < 0 \} \cap \text{mp-alloc-precond2-1-1-loopinv-ext } t p \text{ sz } \text{tm}$

**lemma** alloc-memblk-data-valid-stb2:  
 $\text{blk } x \text{ } t = \text{buf } (\text{mem-pool-info } x \text{ } p) +$   
 $\text{block-num } (\text{mem-pool-info } x \text{ } p) (\text{blk } x \text{ } t) (\text{lsizes } x \text{ } t ! \text{nat } (\text{from-l } x \text{ } t)) *$   
 $(\text{max-sz } (\text{mem-pool-info } x \text{ } p) \text{ div } 4 \wedge \text{nat } (\text{from-l } x \text{ } t)) \implies$   
 $\text{block-num } (\text{mem-pool-info } x \text{ } p) (\text{blk } x \text{ } t) (\text{lsizes } x \text{ } t ! \text{nat } (\text{from-l } x \text{ } t)) < n\text{-max}$   
 $(\text{mem-pool-info } x \text{ } p) * 4 \wedge \text{nat } (\text{from-l } x \text{ } t) \implies$   
 $\text{allocating-node } x \text{ } t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } x \text{ } t), \text{block} = \text{block-num } (\text{mem-pool-info } x$   
 $p) (\text{blk } x \text{ } t) (\text{lsizes } x \text{ } t ! \text{nat } (\text{from-l } x \text{ } t)),$   
 $\text{data} = \text{blk } x \text{ } t) \implies$   
 $(x, y) \in \text{lvars-nochange-rel } t \implies$   
 $(x, y) \in \text{gvars-conf-stable} \implies$   
 $\text{alloc-memblk-data-valid } y \text{ } p (\text{the } (\text{allocating-node } y \text{ } t))$   
**apply**(subgoal-tac blk x t = blk y t)  
**prefer** 2 **apply**(simp add:lvars-nochange-rel-def lvars-nochange-def)  
**apply**(subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p))  
**prefer** 2 **apply**(simp add:gvars-conf-stable-def gvars-conf-def)

```

apply(subgoal-tac lsizes x t = lsizes y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac from-l x t = from-l y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac allocating-node x t = allocating-node y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply (simp add: gvars-conf-def gvars-conf-stable-def)
done

lemma mp-alloc-precond2-1-1-loopinv-ext-stb: stable (mp-alloc-precond2-1-1-loopinv-ext
t p sz tm) (Mem-pool-alloc-rely t)
  apply(rule stable-int2)
  apply (simp add: mp-alloc-precond2-1-1-ext-stb)

  apply(simp add:stable-def) apply clarify
  apply(subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply metis
  apply(subgoal-tac from-l x t = from-l y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply smt
  apply(subgoal-tac blk x t = blk y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac allocating-node x t = allocating-node y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac lsizes x t = lsizes y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac n-max (mem-pool-info x p) = n-max (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply smt
  apply(subgoal-tac block-num (mem-pool-info x p) (blk x t) (lsizes x t ! nat (from-l
x t))
    = block-num (mem-pool-info y p) (blk y t) (lsizes y t ! nat (from-l
y t)))
  prefer 2 apply(simp add: block-num-def Mem-pool-alloc-rely-def lvars-nochange-rel-def
lvars-nochange-def)
  apply(case-tac x=y) apply auto[1]
  apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def lvars-nochange-rel-def
lvars-nochange-def)
  apply smt

```

done

**lemma** *mp-alloc-precond2-1-1-loopinv-stb*: *stable* (*mp-alloc-precond2-1-1-loopinv* *t* *p* *sz* *tm*) (*Mem-pool-alloc-rely* *t*)

**apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule stable-int2*)

**apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule stable-int2*)

**apply** (*simp add: stable-inv-alloc-rely1*)  
**apply** (*simp add: mp-alloc-freenode-stb*)  
**apply** (*simp add: mp-alloc-precond1-ext-stb*)  
**apply** (*simp add: mp-alloc-precond7-ext-stb*)  
**apply** (*simp add: mp-alloc-precond1-0-ext-stb*)  
**using** *mp-alloc-precond1-70-ext-stb* **apply** *blast*  
**apply** (*simp add: mp-alloc-precond1-70-2-ext-stb*)  
**apply** (*simp add: mp-alloc-precond1-70-2-2-ext-stb*)  
**using** *mp-alloc-precond2-1-1-loopinv-ext-stb* **apply** *auto*[1]  
done

**abbreviation** *mp-alloc-precond2-1-2* *t* *p* *sz* *tm*  $\equiv$

$\{s. \text{inv } s\} \cap \{\text{'freeing-node } t = \text{None}\} \cap \{p \in \text{'mem-pools} \wedge tm \geq -1\} \cap$   
 $mp\text{-alloc-precond7-ext } t \text{ } p \text{ } sz \text{ } tm \cap \{\neg \text{'rf } t\}$   
 $\cap mp\text{-alloc-precond1-70-ext } t \text{ } p \text{ } sz \text{ } tm \cap - \{\text{'alloc-l } t < 0\} \cap - \{\text{'free-l } t < 0\}$   
 $\cap - \{\text{'blk } t = \text{NULL}\}$   
 $\cap \{\text{'allocating-node } t = \text{Some } (\text{'pool} = p, \text{level} = \text{nat } (\text{'alloc-l } t),$   
 $\text{block} = \text{block-num } (\text{'mem-pool-info } p) (\text{'blk } t) ((\text{'lsizes}$   
 $t)!(\text{nat } (\text{'alloc-l } t))),$   
 $\text{data} = \text{'blk } t \} \wedge \text{'alloc-memblk-data-valid } p \text{ (the}$   
 $(\text{'allocating-node } t)) \}$

**term** *mp-alloc-precond2-1-1-loopinv* *t* *p* *sz* *tm*

**term** *mp-alloc-precond2-1-2* *t* *p* *sz* *tm*

**lemma** *alloc-memblk-data-valid-stb3*:

*blk* *x* *t* = *buf* (*mem-pool-info* *x* *p*) +  
*block-num* (*mem-pool-info* *x* *p*) (*blk* *x* *t*) (*lsizes* *x* *t* ! *nat* (*alloc-l* *x* *t*)) \*  
(*max-sz* (*mem-pool-info* *x* *p*) *div* 4 ^ *nat* (*alloc-l* *x* *t*))  $\implies$   
*block-num* (*mem-pool-info* *x* *p*) (*blk* *x* *t*) (*lsizes* *x* *t* ! *nat* (*alloc-l* *x* *t*)) < *n-max*  
(*mem-pool-info* *x* *p*) \* 4 ^ *nat* (*alloc-l* *x* *t*)  $\implies$   
*allocating-node* *x* *t* =  
*Some* (*'pool* = *p*, *level* = *nat* (*alloc-l* *x* *t*), *block* = *block-num* (*mem-pool-info* *x*  
*p*) (*blk* *x* *t*) (*lsizes* *x* *t* ! *nat* (*alloc-l* *x* *t*)),  
*data* = *blk* *x* *t*)  $\implies$   
(*x*, *y*)  $\in$  *lvars-nochange-rel* *t*  $\implies$   
(*x*, *y*)  $\in$  *gvars-conf-stable*  $\implies$   
*alloc-memblk-data-valid* *y* *p* (*the* (*allocating-node* *y* *t*))

```

apply(subgoal-tac blk x t = blk y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac lsize x t = lsize y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac alloc-l x t = alloc-l y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac allocating-node x t = allocating-node y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply (simp add: gvars-conf-def gvars-conf-stable-def)
done

```

**lemma** *mp-alloc-precond2-1-2-stb: stable (mp-alloc-precond2-1-2 t p sz tm) (Mem-pool-alloc-rely t)*

```

apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2) apply(rule
stable-int2)
apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2) apply(rule
stable-int2) apply(rule stable-int2)
apply (simp add: stable-inv-alloc-rely1)
apply (simp add: mp-alloc-freenode-stb)
apply (simp add: mp-alloc-precond1-ext-stb)
apply (simp add: mp-alloc-precond7-ext-stb)
apply (simp add: mp-alloc-precond1-0-ext-stb)
using mp-alloc-precond1-70-ext-stb apply blast
apply (simp add: mp-alloc-precond1-70-2-ext-stb)
apply (simp add: mp-alloc-precond1-70-2-2-ext-stb)
apply (simp add: mp-alloc-precond2-1-1-ext-stb)

```

```

apply(simp add: stable-def) apply clarify
apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)

```

```

apply(case-tac x=y) apply auto[1] apply clarify
apply(simp add: block-num-def lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
  apply(subgoal-tac blk x t = blk y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac lsize x t = lsize y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac alloc-l x t = alloc-l y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)
apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p))
  prefer 2 apply(simp add: gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac allocating-node x t = allocating-node y t)
  prefer 2 apply(simp add: lvars-nochange-rel-def lvars-nochange-def)

```

```

apply(subgoal-tac n-max (mem-pool-info x p) = n-max (mem-pool-info y p))
prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)

apply(subgoal-tac block-num (mem-pool-info x p) (blk x t) (lsizes x t ! nat (alloc-l
x t))
      = block-num (mem-pool-info y p) (blk y t) (lsizes y t ! nat (alloc-l
y t)))
prefer 2 apply(simp add: block-num-def Mem-pool-alloc-rely-def lvars-nochange-rel-def
lvars-nochange-def)
by (metis Mem-block.select-convs(2) Mem-block.select-convs(3) Mem-block.select-convs(4)
option.sel)

```

```

abbreviation mp-alloc-precond2-1-3 t p sz tm ≡
  mp-alloc-precond1-70-2-2 t p sz tm ∩ ¬⟦'blk t = NULL⟧
  ∩ ⟦'alloc-blk-valid p (nat ('alloc-l t)) (block-num ('mem-pool-info p) ('blk t)
  (('lsizes t)!(nat ('alloc-l t))))
  ('blk t) ∧ 'allocating-node t = None⟧

```

```

lemma mp-alloc-precond2-1-3-stb: stable (mp-alloc-precond2-1-3 t p sz tm) (Mem-pool-alloc-rely
t)

```

```

apply(rule stable-int2) apply(rule stable-int2)
using mp-alloc-precond1-70-2-2-stb apply auto[1]

apply(simp add:stable-def) apply clarify
apply(simp add:Mem-pool-alloc-rely-def)
apply(case-tac x=y) apply auto[1] apply clarify
apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)

apply(simp add:stable-def) apply clarify
apply(simp add:Mem-pool-alloc-rely-def)
apply(case-tac x=y) apply auto[1] apply clarify
apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def block-num-def)
done

```

```

abbreviation mp-alloc-precond2-1-4 t p sz tm ≡
  mp-alloc-precond1 t p tm ∩ ⟦¬'rf t⟧ ∩ ⟦(tm = FOREVER ⟶ 'tmout t =
  FOREVER)⟧
  ∩ {s. (∃ mblk. mempoolalloc-ret s t = Some mblk ∧ alloc-memblk-valid s p sz
  mblk)}

```

```

lemma mp-alloc-precond2-1-4-stb: stable (mp-alloc-precond2-1-4 t p sz tm) (Mem-pool-alloc-rely
t)

```

```

apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2)

```

```

using mp-alloc-precond1-stb apply auto[1]
apply(simp add:stable-def Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply(simp add:stable-def Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply auto[1]
  apply(simp add:stable-def) apply clarify
    apply(simp add:Mem-pool-alloc-rely-def)
    apply(case-tac x=y) apply simp apply clarify
    apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
      apply(case-tac x=y) apply simp apply clarify
      apply(simp add:alloc-memblk-valid-def lvars-nochange-rel-def lvars-nochange-def
gvars-conf-stable-def gvars-conf-def)
        applymetis
done

```

**abbreviation** *mp-alloc-precond1-8* *t p sz tm*  $\equiv$   
 $mp\text{-}alloc\text{-}precond1\ t\ p\ tm \cap \{\neg rf\ t\} \cap \{(tm = FOREVER \longrightarrow 'tmout\ t = FOREVER)\}$   
 $\cap \{s. (ret\ s\ t = OK \wedge (\exists\ mblk. mempoolalloc\text{-}ret\ s\ t = Some\ mblk \wedge alloc\text{-}memblk\text{-}valid\ s\ p\ sz\ mblk))$   
 $\vee ((ret\ s\ t = ESIZEERR \vee ret\ s\ t = EAGAIN \vee ret\ s\ t = ENOMEM) \wedge mempoolalloc\text{-}ret\ s\ t = None)\}$

**lemma** *mp-alloc-precond1-8-stb*: *stable* (*mp-alloc-precond1-8* *t p sz tm*) (*Mem-pool-alloc-rely* *t*)  
 apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2)  
 using mp-alloc-precond1-stb apply auto[1]  
 apply(simp add:stable-def Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)  
 apply (smt mem-Collect-eq mp-alloc-precond2-ext-stb stable-def)

```

apply(simp add:stable-def) apply clarify
apply(rule conjI)
  apply clarify
  apply(rule conjI)
    apply(simp add:Mem-pool-alloc-rely-def)
    apply(case-tac x=y) apply simp apply clarify
    apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
      apply(simp add:Mem-pool-alloc-rely-def)
      apply(case-tac x=y) apply simp apply clarify
      apply(simp add:alloc-memblk-valid-def lvars-nochange-rel-def lvars-nochange-def
gvars-conf-stable-def gvars-conf-def)
        applymetis
        apply clarify
        apply(simp add:Mem-pool-alloc-rely-def)
        apply(case-tac x=y) apply simp apply clarify
        apply(simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
          done

```

**abbreviation**  $mp\_alloc\_precond1-8-0\ t\ p\ sz\ tm \equiv$   
 $mp\_alloc\_precond1\ t\ p\ tm \cap \{ (tm = FOREVER \longrightarrow 'tmout\ t = FOREVER) \}$   
 $\cap \{ s. (ret\ s\ t = OK \wedge (\exists\ mblk. mempoolalloc-ret\ s\ t = Some\ mblk \wedge alloc-memblk-valid$   
 $s\ p\ sz\ mblk))$   
 $\vee ((ret\ s\ t = ESIZEERR \vee ret\ s\ t = EAGAIN \vee ret\ s\ t = ENOMEM) \wedge$   
 $mempoolalloc-ret\ s\ t = None) \}$

**lemma**  $mp\_alloc\_precond1-8-0-stb$ : *stable* ( $mp\_alloc\_precond1-8-0\ t\ p\ sz\ tm$ ) (*Mem-pool-alloc-rely*  
 $t$ )

**apply**(*rule stable-int2*) **apply**(*rule stable-int2*)  
**using**  $mp\_alloc\_precond1-stb$  **apply** *auto*[1]  
**apply** (*smt mem-Collect-eq mp-alloc-precond2-ext-stb stable-def*)  
  
**apply**(*simp add:stable-def*) **apply** *clarify*  
**apply**(*rule conjI*)  
**apply** *clarify*  
**apply**(*rule conjI*)  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *simp* **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
 $gvars-conf-def$ )  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *simp* **apply** *clarify*  
**apply**(*simp add:alloc-memblk-valid-def lvars-nochange-rel-def lvars-nochange-def*  
 $gvars-conf-stable-def gvars-conf-def$ )  
**apply** *metis*  
**apply** *clarify*  
**apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*case-tac x=y*) **apply** *simp* **apply** *clarify*  
**apply**(*simp add:lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def*  
 $gvars-conf-def$ )  
**done**

**abbreviation**  $mp\_alloc\_precond1-8-1\ t\ p\ sz\ tm \equiv$   
 $mp\_alloc\_precond1-8\ t\ p\ sz\ tm$   
 $\cap \{ 'ret\ t = OK \vee tm = NOWAIT \vee 'ret\ t = ESIZEERR \}$

**lemma**  $mp\_alloc\_precond1-8-1-stb$ : *stable* ( $mp\_alloc\_precond1-8-1\ t\ p\ sz\ tm$ ) (*Mem-pool-alloc-rely*  
 $t$ )

**apply**(*rule stable-int2*)  
**using**  $mp\_alloc\_precond1-8-stb$  **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
 $lvars-nochange-def$ )  
**apply** *auto*  
**done**

**abbreviation**  $mp\_alloc\_precond1-8-1-1\ t\ p\ sz\ tm \equiv$   
 $mp\_alloc\_precond1-8-0\ t\ p\ sz\ tm \cap \{ 'ret\ t = OK \vee tm = NOWAIT \vee 'ret\ t =$



$ESIZEERR\} \cap \{\text{'rf } t = \text{True}\}$

**lemma** *mp-alloc-precond1-8-1-1-stb*: *stable* (*mp-alloc-precond1-8-1-1* *t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*) **apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-0-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**apply** *auto*[1]  
**apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-1-2* *t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-1-1* *t p sz tm*  $\cap \{\text{'ret } t = \text{EAGAIN}\}$

**lemma** *mp-alloc-precond1-8-1-2-stb*: *stable* (*mp-alloc-precond1-8-1-2* *t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-1-1-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-1-3* *t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-1-1* *t p sz tm*  $\cap \neg\{\text{'ret } t = \text{EAGAIN}\}$

**lemma** *mp-alloc-precond1-8-1-3-stb*: *stable* (*mp-alloc-precond1-8-1-3* *t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-1-1-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2* *t p sz tm*  $\equiv$   
*mp-alloc-precond1-8* *t p sz tm*  
 $\cap \neg\{\text{'ret } t = \text{OK} \vee tm = \text{NOWAIT} \vee \text{'ret } t = \text{ESIZEERR}\}$

**lemma** *mp-alloc-precond1-8-2-stb*: *stable* (*mp-alloc-precond1-8-2* *t p sz tm*) (*Mem-pool-alloc-rely*  
*t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2-1* *t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2* *t p sz tm*  $\cap \{\text{'ret } t = \text{EAGAIN}\}$

**lemma** *mp-alloc-precond1-8-2-1-stb*: *stable* (*mp-alloc-precond1-8-2-1 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2-2 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2 t p sz tm*  $\cap$   $\neg \{\! \{ 'ret\ t = EAGAIN \} \}$

**lemma** *mp-alloc-precond1-8-2-2-stb*: *stable* (*mp-alloc-precond1-8-2-2 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2-3 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2-2 t p sz tm*  $\cap \{\! \{ 'tmout\ t \neq FOREVER \} \} \cap \{\! \{ tm > 0 \} \}$

**lemma** *mp-pred1823-eq*: *mp-alloc-precond1-8-2-3 t p sz tm* = *mp-alloc-precond1-8-2-2*  
*t p sz tm*  $\cap \{\! \{ 'tmout\ t \neq FOREVER \} \}$   
**by** *auto*

**lemma** *mp-alloc-precond1-8-2-3-stb*: *stable* (*mp-alloc-precond1-8-2-3 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*) **apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-2-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**apply** *simp*  
**done**

**abbreviation** *mp-alloc-precond1-8-2-20 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2-2 t p sz tm*  $\cap \neg \{\! \{ 'tmout\ t \neq FOREVER \} \}$

**lemma** *mp-alloc-precond1-8-2-20-stb*: *stable* (*mp-alloc-precond1-8-2-20 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-2-stb* **apply** *auto*[1]  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2-4 t p sz tm*  $\equiv$  *mp-alloc-precond1-8-2-2 t p sz*  
*tm*  $\cap \{\! \{ tm > 0 \} \}$

**lemma** *mp-alloc-precond1-8-2-4-stb*: *stable* (*mp-alloc-precond1-8-2-4 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-2-stb* **apply** *auto[1]*  
**apply**(*unfold stable-def*) **apply** *simp*  
**done**

**abbreviation** *mp-alloc-precond1-8-2-40 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2-4 t p sz tm*  $\cap \{\text{'tmout } t < 0\}$

**lemma** *mp-alloc-precond1-8-2-40-stb*: *stable* (*mp-alloc-precond1-8-2-40 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-4-stb* **apply** *blast*  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**term** *mp-alloc-precond1-8-2-40 t p sz tm*

**abbreviation** *mp-alloc-precond1-8-2-41 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-2-4 t p sz tm*  $\cap \neg\{\text{'tmout } t < 0\}$

**lemma** *mp-alloc-precond1-8-2-41-stb*: *stable* (*mp-alloc-precond1-8-2-41 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-2-4-stb* **apply** *blast*  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def*  
*lvars-nochange-def*)  
**done**

**abbreviation** *mp-alloc-precond1-8-2-5 t p sz tm*  $\equiv$   
*mp-alloc-precond1-8-0 t p sz tm*  $\cap \{\text{'tm} > 0\} \cap \neg\{\text{'ret } t = \text{OK} \vee \text{'tm} = \text{NOWAIT}$   
 $\vee \text{'ret } t = \text{ESIZEERR}\} \cap \neg\{\text{'ret } t = \text{EAGAIN}\}$   
 $\cap \{\text{'tmout } t < 0\} \cap \{\text{'rf } t\}$   
**term** *mp-alloc-precond1-8-2-5 t p sz tm*

**lemma** *mp-alloc-precond1-8-2-5-stb*: *stable* (*mp-alloc-precond1-8-2-5 t p sz tm*)  
 (*Mem-pool-alloc-rely t*)  
**apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule stable-int2*) **apply**(*rule*  
*stable-int2*) **apply**(*rule stable-int2*)  
**using** *mp-alloc-precond1-8-0-stb* **apply** *blast*  
**apply**(*unfold stable-def*) **apply**(*simp add:Mem-pool-alloc-rely-def*)  
**apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**apply**(*simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def*)  
**done**

## 10.2 proof of each statement

### 10.3 stm1

**lemma** *mp-alloc-stm1-lm0*:

*cur V = Some t  $\implies$  inv V  $\implies$*   
*V( $\llbracket$ lsizes := (lsizes V)(t := lsizes V t @ [ALIGN<sub>4</sub> (lsizes V t ! (i V t - Suc*  
*NULL) div 4]) $\rrbracket$ )*  
 *$\in \llbracket$ '(Pair V)  $\in$  Mem-pool-alloc-guar t $\rrbracket$*   
**apply** *auto* **apply** (*simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def*  
*lvars-nochange-def*)  
**apply** (*rule disjI1*)  
**apply** (*subgoal-tac (V, V( $\llbracket$ lsizes := (lsizes V)(t := lsizes V t @ [ALIGN<sub>4</sub> (lsizes*  
*V t ! (i V t - Suc NULL) div 4]) $\rrbracket$ ) $\in$ lvars-nochange1-4all)*)  
**using** *glnchange-inv0* **apply** *auto[1]* **apply** (*simp add:lvars-nochange1-4all-def*  
*lvars-nochange1-def*)  
**done**

**lemma** *mp-alloc-stm1-lm1*: *mp-alloc-precond1-6-10 t p sz timeout  $\alpha \cap \llbracket$ 'cur =*  
*Some t $\rrbracket \cap \{V\}$*

$\subseteq \llbracket$ '(*lsizes-update* ( $\lambda$ -. *'lsizes*(t := *'lsizes* t @ [ALIGN<sub>4</sub> (*'lsizes* t ! (*'i* t  
*- Suc NULL) div 4]) $\rrbracket$ ))*

$\in \llbracket$ '(Pair V)  $\in$  Mem-pool-alloc-guar t $\rrbracket \cap$   
*mp-alloc-precond1-6-2 t p sz timeout  $\alpha$*  $\rrbracket$

**apply** *clarify*  
**apply** (*rule IntI*) **using** *mp-alloc-stm1-lm0* **apply** *blast*  
**apply** (*rule IntI*) **prefer** 2 **apply** *clarsimp*  
**apply** (*subgoal-tac lsizes (V( $\llbracket$ lsizes := (lsizes V)(t := lsizes V t @ [ALIGN<sub>4</sub> (lsizes*  
*V t ! (i V t - Suc NULL) div 4]) $\rrbracket$ ) t*  
 $=$  *lsizes V t @ [ALIGN<sub>4</sub> (lsizes V t ! (i V t - Suc NULL) div*  
*4])))  
**prefer** 2 **apply** *auto[1]*  
**apply** (*simp add: subst[where s=lsizes (V( $\llbracket$ lsizes := (lsizes V)(t := lsizes V t*  
 $\cap$  *[ALIGN<sub>4</sub> (lsizes V t ! (i V t - Suc NULL) div 4]) $\rrbracket$ ) t*  
 $\text{and } t=\text{lsizes } V t @ [ALIGN_4 (lsizes V t ! (i V t - Suc NULL)$   
*div 4]))))**

**apply** (*rule conjI*) **apply** *clarify*  
**apply** (*case-tac ii < length (lsizes V t)*) **apply** (*metis nth-append*)  
**apply** (*case-tac ii = length (lsizes V t)*)  
**apply** (*subgoal-tac (lsizes V t @ [ALIGN<sub>4</sub> (ALIGN<sub>4</sub> (max-sz (mem-pool-info*  
*V p)) div 4 ^ (i V t - Suc NULL) div 4]) ! ii*  
 $=$  *ALIGN<sub>4</sub> (ALIGN<sub>4</sub> (max-sz (mem-pool-info V p)) div 4*  
 $\wedge$  *(i V t - Suc NULL) div 4))  
**prefer** 2 **apply** (*meson nth-append-length*)  
**apply** (*subgoal-tac ALIGN<sub>4</sub> (max-sz (mem-pool-info V p)) div 4 ^ (i V t -*  
*Suc NULL) div 4*  
 $=$  *ALIGN<sub>4</sub> (max-sz (mem-pool-info V p)) div 4 ^ ii*)  
**prefer** 2 **apply** (*metis Divides.div-mult2-eq One-nat-def power-minus-mult*  
*zero-le-numeral*)*

```

apply(subgoal-tac (ALIGN4 (max-sz (mem-pool-info V p)) div 4 ^ ii) mod
4 = 0)
  prefer 2 apply(subgoal-tac  $\exists n > 0. \text{max-sz}(\text{mem-pool-info } V p) = (4 * n) * (4 \wedge n\text{-levels}(\text{mem-pool-info } V p))$ )
    prefer 2 apply(simp add:inv-def inv-mempool-info-def) apply metis
    apply (metis (no-types, lifting) inv-maxsz-align4 less-imp-le-nat
      m-mod-div mod-mult-self1-is-0 mult.assoc pow-mod-0)
    apply (metis align40)
    apply linarith
    apply (simp add: le-nat-iff nth-append)

apply(rule IntI) prefer 2 apply clarify apply auto[1]
apply(rule IntI) prefer 2 apply clarify apply auto[1]
apply(rule IntI) prefer 2 apply clarify apply auto[1]
apply(rule IntI) apply simp apply simp-inv apply metis
apply simp
done

lemma mp-alloc-stm1-lm:
   $\Gamma \vdash_I \text{Some} (IF \text{'i t} > 0 \text{ THEN}$ 
    ( $t \blacktriangleright \text{'lsizes} := \text{'lsizes}(t := \text{'lsizes } t @ [\text{ALIGN}_4 (\text{'lsizes } t ! (\text{'i t} - 1) \text{div } 4))$ )))
    FI) satp [mp-alloc-precond1-6-1 t p sz timeout  $\alpha$ , Mem-pool-alloc-rely t,
      Mem-pool-alloc-guar t, mp-alloc-precond1-6-2 t p sz timeout  $\alpha$ ]
apply(rule Cond)
using mp-alloc-precond1-6-1-stb apply simp

apply(unfold stm-def)
apply(rule Await)
using mp-alloc-precond1-6-10-stb[of t p timeout sz  $\alpha$ ] apply fast
using mp-alloc-precond1-6-2-stb apply simp
apply(rule allI)
apply(rule Basic)
apply(case-tac mp-alloc-precond1-6-10 t p sz timeout  $\alpha \cap \{\text{'cur} = \text{Some } t\}$ 
 $\cap \{V\} = \{\}$ )
  apply auto[1]
  using mp-alloc-stm1-lm1[of t p timeout sz  $\alpha$ ] apply auto[1]
  apply simp using stable-id2 apply auto[1]
  using stable-id2 apply auto[1]

apply(unfold Skip-def)
apply(rule Basic) apply clarify
apply simp
apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]
using mp-alloc-precond1-6-11-stb[of t p timeout sz  $\alpha$ ] apply fast
using mp-alloc-precond1-6-2-stb apply fast
apply(simp add:Mem-pool-alloc-guar-def)
done

```

## 10.4 stm2

**lemma** *mp-alloc-stm2-lm2-1*:

$cur\ V = Some\ t \implies inv\ V \implies V(\llbracket alloc\text{-}lsize\text{-}r := (alloc\text{-}lsize\text{-}r\ V)(t := True) \rrbracket)$   
 $\in \llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}alloc\text{-}guar\ t \rrbracket$   
**apply** *auto* **apply**(*simp* *add:Mem-pool-alloc-guar-def* *gvars-conf-stable-def* *gvars-conf-def* *lvars-nochange-def*)  
**apply**(*rule* *disjI1*)  
**apply**(*subgoal-tac* ( $V, V(\llbracket alloc\text{-}lsize\text{-}r := (alloc\text{-}lsize\text{-}r\ V)(t := True) \rrbracket) \in lvars\text{-}nochange1\text{-}4all$ ))  
**using** *glnochange-inv0* **apply** *auto*[1] **apply**(*simp* *add:lvars-nochange1-4all-def* *lvars-nochange1-def*)  
**done**

**lemma** *mp-alloc-stm2-lm2*:

$mp\text{-}alloc\text{-}precond1\text{-}6\text{-}20\ t\ p\ sz\ timeout\ \alpha \cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\}$   
 $\subseteq \llbracket '(alloc\text{-}lsize\text{-}r\text{-}update\ (\lambda\cdot. 'alloc\text{-}lsize\text{-}r(t := True)))$   
 $\in \llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}alloc\text{-}guar\ t \rrbracket \cap mp\text{-}alloc\text{-}lsize\text{-}stm\text{-}loop\text{-}inv\ t\ p$   
 $sz\ timeout\ 0 \rrbracket$   
**apply** *clarify*  
**apply**(*rule* *IntI*)  
**using** *mp-alloc-stm2-lm2-1* **apply** *simp*  
**apply**(*rule* *IntI*) **prefer** 2  
**apply**(*case-tac*  $i\ V\ t = 0$ ) **apply**(*simp* *add:inv-def* *inv-mempool-info-def*) **apply** *simp*  
**apply**(*rule* *IntI*) **prefer** 2 **apply** *auto*[1]  
**apply**(*rule* *IntI*) **prefer** 2 **apply** *simp*  
**apply**(*rule* *IntI*) **prefer** 2 **apply**(*simp* *add:alloc-memblk-valid-def*)  
**apply** *simp* **apply**(*subgoal-tac* ( $V, V(\llbracket alloc\text{-}lsize\text{-}r := (alloc\text{-}lsize\text{-}r\ V)(t := True) \rrbracket) \in lvars\text{-}nochange1\text{-}4all$ ))  
**using** *glnochange-inv0* **apply** *auto*[1] **apply**(*simp* *add:lvars-nochange1-4all-def* *lvars-nochange1-def*)  
**done**

**lemma** *mp-alloc-stm2-lm4-1*:

$cur\ V = Some\ t \implies inv\ V \implies V(\llbracket alloc\text{-}l := (alloc\text{-}l\ V)(t := int\ (i\ V\ t)) \rrbracket) \in$   
 $\llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}alloc\text{-}guar\ t \rrbracket$   
**apply** *auto* **apply**(*simp* *add:Mem-pool-alloc-guar-def* *gvars-conf-stable-def* *gvars-conf-def* *lvars-nochange-def*)  
**apply**(*rule* *disjI1*)  
**apply**(*subgoal-tac* ( $V, V(\llbracket alloc\text{-}l := (alloc\text{-}l\ V)(t := int\ (i\ V\ t)) \rrbracket) \in lvars\text{-}nochange1\text{-}4all$ ))  
**using** *glnochange-inv0* **apply** *auto*[1] **apply**(*simp* *add:lvars-nochange1-4all-def* *lvars-nochange1-def*)  
**done**

**lemma** *mp-alloc-stm2-lm4*:

$mp\text{-}alloc\text{-}precond1\text{-}6\text{-}21\ t\ p\ sz\ timeout\ \alpha \cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\}$   
 $\subseteq \llbracket '(alloc\text{-}l\text{-}update\ (\lambda\cdot. 'alloc\text{-}l(t := int\ ('i\ t))))$   
 $\in \llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}alloc\text{-}guar\ t \rrbracket \cap$   
 $mp\text{-}alloc\text{-}precond1\text{-}6\text{-}21\text{-}1\ t\ p\ sz\ timeout\ \alpha \rrbracket$   
**apply** *clarify*  
**apply**(*rule* *IntI*)

```

    using mp-alloc-stm2-lm4-1 apply simp
  apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply simp apply simp
  apply(rule IntI) prefer 2 apply simp
  apply(rule IntI) prefer 2 apply simp
  apply(rule IntI) prefer 2 apply(simp add:alloc-memblk-valid-def)
  apply simp
  apply(subgoal-tac (V, V(alloc-l := (alloc-l V)(t := int (i V t)))) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

```

done

**lemma** mp-alloc-stm2-lm5-1-1:

```

  cur V = Some t ⇒ inv V ⇒ V(free-l := (free-l V)(t := int (i V t))) ∈
  { (Pair V) ∈ Mem-pool-alloc-guar t }
  apply auto apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def
lvars-nochange-def)
  apply(rule disjI1)
  apply(subgoal-tac (V, V(free-l := (free-l V)(t := int (i V t)))) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

```

done

**lemma** mp-alloc-stm2-lm5-1:

```

  mp-alloc-precond1-6-21-1 t p sz timeout α ∩ { cur = Some t } ∩ { V }
  ⊆ { (free-l-update (λ-. 'free-l(t := int ('i t))))
    ∈ { (Pair V) ∈ Mem-pool-alloc-guar t } ∩
    mp-alloc-precond1-6-21-2 t p sz timeout α }
  apply clarify
  apply(rule IntI)
    using mp-alloc-stm2-lm5-1-1 apply simp
  apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply simp apply simp
  apply(rule IntI) prefer 2 apply simp
  apply(rule IntI) prefer 2 apply simp
  apply(rule IntI) prefer 2 apply(simp add:alloc-memblk-valid-def)
  apply simp
  apply(subgoal-tac (V, V(free-l := (free-l V)(t := int (i V t)))) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

```

done

**lemma** mp-alloc-stm2-lm5:

```

  Γ ⊢I Some (t ► 'free-l := 'free-l (t := int ('i t)))
  satp [mp-alloc-precond1-6-21-1 t p sz timeout α, Mem-pool-alloc-rely t,
    Mem-pool-alloc-guar t, mp-alloc-precond1-6-21-2 t p sz timeout α]
  apply(unfold stm-def)

```

```

apply(rule Await)
  using mp-alloc-precond1-6-21-1-stb apply simp
  using mp-alloc-precond1-6-21-2-stb apply simp
  apply(rule allI)
  apply(rule Basic)
    apply(case-tac mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha \cap \llbracket 'cur = Some$ 
 $t \rrbracket \cap \{V\} = \{\}$ )
      apply auto[1] using mp-alloc-stm2-lm5-1[of t p timeout sz  $\alpha$ ] apply auto[1]
      apply simp using stable-id2 apply auto[1]
      using stable-id2 apply auto[1]
done

```

```

lemma mp-alloc-stm2-lm6:
   $\Gamma \vdash_I Some SKIP sat_p [mp-alloc-precond1-6-21-1 t p sz timeout \alpha, Mem-pool-alloc-rely$ 
 $t,$ 
     $Mem-pool-alloc-guar t, mp-alloc-precond1-6-21-2 t p sz timeout \alpha]$ 
  apply(unfold Skip-def)
  apply(rule Basic)
    apply clarify apply(rule IntI) apply(rule IntI) apply(rule IntI) apply(rule
IntI) apply(rule IntI)
    apply simp+ apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def
gvars-conf-def lvars-nochange-def) apply auto[1]
    using mp-alloc-precond1-6-21-1-stb apply simp
    using mp-alloc-precond1-6-21-2-stb apply simp
done

```

```

lemma mp-alloc-stm2-lm7-1:
   $cur V = Some t \implies inv V \implies V \llbracket i := (i V)(t := (i V t) + 1) \rrbracket \in \llbracket '(Pair V)$ 
 $\in Mem-pool-alloc-guar t \rrbracket$ 
  apply auto apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def
lvars-nochange-def)
  apply(rule disjI1)
  apply(subgoal-tac (V, V  $\llbracket i := (i V)(t := (i V t) + 1) \rrbracket \in lvars-nochange1-4all$ )
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
done

```

```

lemma mp-alloc-stm2-lm7:
   $mp-alloc-precond1-6-21-2 t p sz timeout \alpha \cap \llbracket 'cur = Some t \rrbracket \cap \{V\}$ 
 $\subseteq \llbracket '(i-update (\lambda-. 'i(t := Suc ('i t))))$ 
 $\in \llbracket '(Pair V) \in Mem-pool-alloc-guar t \rrbracket \cap$ 
 $mp-alloc-lsizestm-loopinv t p sz timeout (\alpha - 1) \rrbracket$ 
  apply clarify
  apply(rule IntI)
  using mp-alloc-stm2-lm7-1 apply simp
  apply(rule IntI) prefer 2
  apply(case-tac i V t = 0) apply simp
  apply(simp add:inv-def inv-mempool-info-def)
  apply(rule IntI) apply(rule IntI) apply(rule IntI) apply(rule IntI)

```



```

apply clarsimp apply(subgoal-tac ( $V, V \langle i := (i \ V)(t := (i \ V \ t) + 1) \rangle \rangle \in \textit{lvars-nochange1-4all}$ )
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply clarsimp
  apply clarsimp
  apply clarsimp
  apply clarsimp
  apply(rule conjI) apply auto[1]
  apply(rule conjI) apply clarify apply(case-tac  $ii < i \ V \ t$ ) apply auto[1]
  apply(case-tac  $ii = i \ V \ t$ ) apply simp apply simp
  apply clarsimp
done

```

**lemma** *subset-un-I1*[*intro*]:  $A \subseteq B \implies A \subseteq B \cup C$  **by** *auto*

**lemma** *subset-un-I2*[*intro*]:  $A \subseteq C \implies A \subseteq B \cup C$  **by** *auto*

**lemma** *mp-alloc-stm2-lm8*:

$$\begin{aligned}
P \subseteq \{ \lambda \cdot ( \textit{alloc-lsize-r-update} \ (\lambda \cdot \textit{alloc-lsize-r}(t := \textit{True})) ) \\
\quad \in \{ \lambda \cdot (\textit{Pair} \ V) \in \textit{Mem-pool-alloc-guar} \ t \} \cap B \} \implies \\
P \subseteq \{ \lambda \cdot ( \textit{alloc-lsize-r-update} \ (\lambda \cdot \textit{alloc-lsize-r}(t := \textit{True})) ) \\
\quad \in \{ \lambda \cdot (\textit{Pair} \ V) \in \textit{Mem-pool-alloc-guar} \ t \} \cap (A \cup B) \}
\end{aligned}$$

**apply** *auto*

**done**

**lemma** *mp-alloc-stm2-lm9*:

$$\begin{aligned}
P \subseteq \{ \lambda \cdot ( \textit{i-update} \ (\lambda \cdot \textit{i}(t := \textit{Suc} \ (\textit{i} \ t))) ) \\
\quad \in \{ \lambda \cdot (\textit{Pair} \ V) \in \textit{Mem-pool-alloc-guar} \ t \} \cap A \} \implies \\
P \subseteq \{ \lambda \cdot ( \textit{i-update} \ (\lambda \cdot \textit{i}(t := \textit{Suc} \ (\textit{i} \ t))) ) \\
\quad \in \{ \lambda \cdot (\textit{Pair} \ V) \in \textit{Mem-pool-alloc-guar} \ t \} \cap (A \cup B) \}
\end{aligned}$$

**by** *auto*

**lemma** *mp-alloc-stm2-lm*:

$$\begin{aligned}
& \Gamma \vdash_I \textit{Some} \ (\textit{IF} \ \textit{'lsizes} \ t \ ! \ \textit{'i} \ t < \textit{sz} \ \textit{THEN} \\
& \quad t \blacktriangleright \textit{'alloc-lsize-r} := \textit{'alloc-lsize-r}(t := \textit{True}) \\
& \quad \textit{ELSE} \ (t \blacktriangleright \textit{'alloc-l} := \textit{'alloc-l}(t := \textit{int} \ (\textit{'i} \ t))) ; \\
& \quad \textit{IF} \ \neg \textit{level-empty} \ (\textit{'mem-pool-info} \ p) \ (\textit{'i} \ t) \ \textit{THEN} \\
& \quad \quad t \blacktriangleright \textit{'free-l} := \textit{'free-l}(t := \textit{int} \ (\textit{'i} \ t)) \\
& \quad \textit{FI} ; \\
& \quad (t \blacktriangleright \textit{'i} := \textit{'i}(t := \textit{Suc} \ (\textit{'i} \ t))) \\
& \textit{FI}) \textit{sat}_p \ [\textit{mp-alloc-precond1-6-2} \ t \ p \ \textit{sz} \ \textit{timeout} \ \alpha, \ \textit{Mem-pool-alloc-rely} \ t, \\
& \quad \textit{Mem-pool-alloc-guar} \ t, \ \textit{mp-alloc-lsizestm-loopinv} \ t \ p \ \textit{sz} \ \textit{timeout} \ (\alpha - 1) \\
& \quad \cup \ \textit{mp-alloc-lsizestm-loopinv} \ t \ p \ \textit{sz} \ \textit{timeout} \ 0]
\end{aligned}$$

**apply**(*rule Cond*)

**using** *mp-alloc-precond1-6-2-stb* **apply** *simp*

**apply**(*unfold stm-def*)[1]

**apply**(*rule Await*)

**using** *mp-alloc-precond1-6-20-stb* **apply** *simp*

```

apply(rule stable-un2)
  using mp-alloc-lsizestm-loopinv-stb apply fast
  using mp-alloc-lsizestm-loopinv-stb apply fast
apply(rule allI)
apply(rule Basic)
  apply(case-tac mp-alloc-precond1-6-20 t p sz timeout  $\alpha \cap \{\text{'cur} = \text{Some } t\}$ 
 $\cap \{V\} = \{\}$ )
    apply auto[1]
    apply(rule mp-alloc-stm2-lm8[of - t - mp-alloc-lsizestm-loopinv t p sz timeout
0 mp-alloc-lsizestm-loopinv t p sz timeout ( $\alpha - 1$ )])
      using mp-alloc-stm2-lm2[of t p timeout sz  $\alpha$ ] apply meson
      apply simp using stable-id2 apply auto[1]
      using stable-id2 apply auto[1]

apply(rule Seq[where mid=mp-alloc-precond1-6-21-2 t p sz timeout  $\alpha$ ])
apply(rule Seq[where mid=mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha$ ])

apply(unfold stm-def)[1]
apply(rule Await)
  using mp-alloc-precond1-6-21-stb apply simp
  using mp-alloc-precond1-6-21-1-stb apply simp
  apply(rule allI)
  apply(rule Basic)
    apply(case-tac mp-alloc-precond1-6-21 t p sz timeout  $\alpha \cap \{\text{'cur} = \text{Some } t\}$ 
 $\cap \{V\} = \{\}$ )
      apply auto[1] using mp-alloc-stm2-lm4[of t p timeout sz] apply presburger
      apply simp using stable-id2 apply fast
      using stable-id2 apply fast

apply(rule Cond)
  using mp-alloc-precond1-6-21-1-stb apply simp

  using Conseq[where pre=mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha \cap \{\neg$ 
level-empty ('mem-pool-info p) ('i t)}
  and pre'=mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha$  and rely=Mem-pool-alloc-rely
t
  and rely'=Mem-pool-alloc-rely t and guar=Mem-pool-alloc-guar t and guar'=Mem-pool-alloc-guar
t
  and post=mp-alloc-precond1-6-21-2 t p sz timeout  $\alpha$  and post'=mp-alloc-precond1-6-21-2
t p sz timeout  $\alpha$ 
    and P=Some (t  $\blacktriangleright$  'free-l := 'free-l (t := int ('i t)))
    mp-alloc-stm2-lm5[of t p timeout sz] apply fast
  using Conseq[where pre=mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha \cap \neg$ 
level-empty ('mem-pool-info p) ('i t)}
  and pre'=mp-alloc-precond1-6-21-1 t p sz timeout  $\alpha$  and rely=Mem-pool-alloc-rely
t
  and rely'=Mem-pool-alloc-rely t and guar=Mem-pool-alloc-guar t and guar'=Mem-pool-alloc-guar

```

$t$   
**and**  $post = mp\_alloc\_precond1-6-21-2\ t\ p\ sz\ timeout\ \alpha$  **and**  $post' = mp\_alloc\_precond1-6-21-2$   
 $t\ p\ sz\ timeout\ \alpha$   
**and**  $P = Some\ SKIP$   
 $mp\_alloc\_stm2-lm6[of\ t\ p\ timeout\ sz]$  **apply fast**  
**apply**( $simp\ add:Mem-pool-alloc-guar-def$ )

**apply**( $unfold\ stm-def$ )[1]  
**apply**( $rule\ Await$ )  
**using**  $mp\_alloc\_precond1-6-21-2-stb$  **apply simp**  
**apply**( $rule\ stable-un2$ )  
**using**  $mp\_alloc-lsizestm-loopinv-stb$  **apply fast**  
**using**  $mp\_alloc-lsizestm-loopinv-stb$  **apply fast**  
**apply**( $rule\ allI$ )  
**apply**( $rule\ Basic$ )  
**apply**( $case-tac\ mp\_alloc\_precond1-6-21-2\ t\ p\ sz\ timeout\ \alpha \cap \{\}'cur = Some$   
 $t\} \cap \{V\} = \{\}$ )  
**apply**  $auto[1]$   
**apply**( $rule\ mp\_alloc\_stm2-lm9[of\ -\ t - mp\_alloc-lsizestm-loopinv\ t\ p\ sz\ timeout$   
 $(\alpha - 1)\ mp\_alloc-lsizestm-loopinv\ t\ p\ sz\ timeout\ 0]$ )  
**using**  $mp\_alloc\_stm2-lm7[of\ t\ p\ timeout\ sz\ \alpha]$  **apply meson**  
**apply simp using stable-id2 apply fast**  
**using stable-id2 apply fast**

**apply**( $simp\ add:Mem-pool-alloc-guar-def$ )  
**done**

**term**  $\{\}'(Pair\ Va) \in Mem-pool-alloc-guar\ t\} \cap mp\_alloc\_precond2-1\ t\ p\ sz\ timeout$

## 10.5 lsize while loop

**abbreviation**  $alloc-lsize-loopbody\ t\ p\ sz \equiv$   
 $IF\ \prime i\ t > 0\ THEN$   
 $(t \blacktriangleright \prime lsize := \prime lsize(t := \prime lsize\ t @ [ALIGN4\ (\prime lsize\ t ! (\prime i\ t - 1)\ div\ 4)]))$   
 $FI;;$   
 $IF\ \prime lsize\ t ! \prime i\ t < sz\ THEN$   
 $(t \blacktriangleright \prime alloc-lsize-r := \prime alloc-lsize-r(t := True))$   
 $ELSE$   
 $(t \blacktriangleright \prime alloc-l := \prime alloc-l(t := int\ (\prime i\ t)))$   
 $IF\ \neg level-empty\ (\prime mem-pool-info\ p)\ (\prime i\ t)\ THEN$   
 $(t \blacktriangleright \prime free-l := \prime free-l(t := int\ (\prime i\ t)))$   
 $FI;;$   
 $(t \blacktriangleright \prime i := \prime i(t := \prime i\ t + 1))$   
 $FI$

**lemma**  $lsize-loop-body-terminate$ :

$\Gamma \vdash_I \text{Some } (\text{alloc-lsize-loopbody } t \ p \ sz)$   
 $\text{sat}_p [\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ \alpha \cap \{\alpha > 0\}, \text{Mem-pool-alloc-rely } t,$   
 $\text{Mem-pool-alloc-guar } t,$   
 $\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ (\alpha - 1) \cup \text{mp-alloc-lsizestm-loopinv } t \ p$   
 $\text{sz } tm \ 0]$

**apply**(rule Seq[**where**  $\text{mid} = \text{mp-alloc-precond1-6-2 } t \ p \ sz \ tm \ \alpha$ ])

**apply**(rule subst[**where**  $s = \text{mp-alloc-precond1-6-1 } t \ p \ sz \ tm \ \alpha$  **and**  
 $t = \text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ \alpha \cap \{\alpha > 0\}$ ])  
**using**  $\text{lsizeloop-inv-cond-eq-}\alpha\text{gt0}$ [of  $t \ p \ tm \ sz \ \alpha$ ] **apply** fast  
**using**  $\text{mp-alloc-stm1-lm}$ [of  $t \ p \ tm \ sz \ \alpha$ ] **apply** blast

**using**  $\text{mp-alloc-stm2-lm}$  **apply** simp

**done**

**lemma**  $\text{lsizeloopbody-sat-invterm-imp-inv-post}$ :

$\Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{pre}, \text{rely}, \text{guar},$   
 $\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ (\alpha - 1) \cup \text{mp-alloc-lsizestm-loopinv } t$   
 $p \ sz \ tm \ 0]$   
 $\implies \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{pre}, \text{rely}, \text{guar}, \text{mp-alloc-precond1-6 } t \ p \ sz \ tm]$   
**apply**(rule Conseq [of pre pre  
 $\text{rely } \text{rely } \text{guar } \text{guar } \text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ (\alpha - 1) \cup$   
 $\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ 0$   
 $\text{mp-alloc-precond1-6 } t \ p \ sz \ tm \ \text{Some } P]$ )

**apply** fast+

**done**

**lemma**  $\text{lsizeloopbody-term-imp-prepost}$ :

$(\forall \alpha. \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ \alpha \cap \{\alpha > 0\}, \text{rely},$   
 $\text{guar},$   
 $\text{mp-alloc-lsizestm-loopinv } t \ p \ sz \ tm \ (\alpha - 1) \cup \text{mp-alloc-lsizestm-loopinv } t \ p$   
 $\text{sz } tm \ 0])$   
 $\implies \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-alloc-precond1-6 } t \ p \ sz \ tm \cap \text{mp-alloc-lsizestm-loopcond}$   
 $t \ p,$   
 $\text{rely}, \text{guar}, \text{mp-alloc-precond1-6 } t \ p \ sz \ tm]$

**apply**(rule subst[**where**  $s = \forall v. v \in \text{mp-alloc-precond1-6 } t \ p \ sz \ tm \cap \text{mp-alloc-lsizestm-loopcond}$   
 $t \ p \implies$

$\Gamma \vdash_I \text{Some } P \text{ sat}_p [\{v\}, \text{rely}, \text{guar}, \text{mp-alloc-precond1-6 } t \ p \ sz \ tm]$  **and**  
 $t = \Gamma \vdash_I \text{Some } P \text{ sat}_p [\text{mp-alloc-precond1-6 } t \ p \ sz \ tm \cap \text{mp-alloc-lsizestm-loopcond}$   
 $t \ p,$   
 $\text{rely}, \text{guar}, \text{mp-alloc-precond1-6 } t \ p \ sz \ tm]$ )

**using**  $\text{allpre-eq-pre}$ [of  $\text{mp-alloc-precond1-6 } t \ p \ sz \ tm \cap \text{mp-alloc-lsizestm-loopcond}$   
 $t \ p$

$\text{Some } P \text{ rely } \text{guar } \text{mp-alloc-precond1-6 } t \ p \ sz \ tm]$

**apply** *meson*

**apply**(*rule allI*) **apply**(*rule impI*)  
**apply**(*subgoal-tac*  $\exists \alpha. v \in mp\text{-}alloc\text{-}lsize\text{-}stm\text{-}loopinv\ t\ p\ sz\ tm\ \alpha \cap \{\alpha > 0\}$ )  
**prefer** 2 **using** *lsizestm-pre-loopcond-imp-loopinv- $\alpha$ gt0*[*of* - *t p tm sz*] **apply** *meson*  
*meson*

**apply**(*erule exE*)  
**using** *sat-pre-imp-allinpre*[*of* *Some P - rely guar mp-alloc-precond1-6 t p sz tm*]  
*lsize-loopbody-sat-invtterm-imp-inv-post*[*of* *P - rely guar t p tm sz*] **apply** *meson*  
**done**

**lemma** *lsize-loop-body-satprepost*:  
 $\Gamma \vdash_I \text{Some } (alloc\text{-}lsize\text{-}loopbody\ t\ p\ sz)$   
 $sat_p [mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ timeout \cap mp\text{-}alloc\text{-}lsize\text{-}stm\text{-}loopcond\ t\ p,$   
 $Mem\text{-}pool\text{-}alloc\text{-}rely\ t, Mem\text{-}pool\text{-}alloc\text{-}guar\ t, mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz$   
 $timeout]$   
**using** *lsize-loopbody-term-imp-prepost*[*of* *alloc-lsize-loopbody t p sz t p timeout sz*  
 $Mem\text{-}pool\text{-}alloc\text{-}rely\ t\ Mem\text{-}pool\text{-}alloc\text{-}guar\ t]$   
*lsize-loop-body-terminate*[*of* *t sz p timeout*] **apply** *fast*  
**done**

**lemma** *lsize-loop-stm*:  
 $\Gamma \vdash_I \text{Some } (WHILE\ 'i\ t < n\text{-}levels\ ('mem\text{-}pool\text{-}info\ p) \wedge \neg 'alloc\text{-}lsize\text{-}r\ t\ DO$   
 $alloc\text{-}lsize\text{-}loopbody\ t\ p\ sz$   
 $OD)\ sat_p [mp\text{-}alloc\text{-}precond1\text{-}6\ t\ p\ sz\ timeout, Mem\text{-}pool\text{-}alloc\text{-}rely\ t,$   
 $Mem\text{-}pool\text{-}alloc\text{-}guar\ t, mp\text{-}alloc\text{-}precond1\text{-}7\ t\ p\ sz\ timeout]$   
**apply**(*rule While*)  
**using** *mp-alloc-precond1-6-stb* **apply** *simp*  
**apply**(*rule Int-greatest*) **apply**(*rule Int-greatest*) **apply**(*rule Int-greatest*) **ap-**  
**ply**(*rule Int-greatest*)  
**apply**(*rule Int-greatest*)  
**apply** *force+*  
**using** *mp-alloc-precond1-7-stb* **apply** *simp*

**using** *lsize-loop-body-satprepost*[*of* *t sz p timeout*] **apply** *fast*

**apply**(*simp add:Mem-pool-alloc-guar-def*)  
**done**

## 10.6 stm3

**lemma** *mp-alloc-stm3-lm3-1*:  $ii < n\text{-}levels\ mp \implies$   
 $length\ (levels\ mp) = n\text{-}levels\ mp \implies$   
 $free\text{-}list\ (levels\ mp\ !\ ii) = [] \implies$   
 $rmhead\text{-}free\text{-}list\ mp\ ii = mp$   
**apply**(*simp add:rmhead-free-list-def*)  
**by** (*metis Mem-pool.surjective Mem-pool.update-convs*(5) *Mem-pool-lvl.surjective*)

*Mem-pool-lvl.update-convs(2) list-update-id)*

**lemma** *mp-alloc-stm3-lm3-2:*

*head-free-list mp ii = NULL*  $\implies$   
*ii < n-levels mp*  $\implies$   
*NULL < buf mp*  $\implies$   
 $\forall i < n\text{-levels } mp.$   
 $\forall j < \text{length } (\text{free-list } (\text{levels } mp ! i)).$   
 $\text{buf } mp \leq \text{free-list } (\text{levels } mp ! i) ! j \implies$   
 $\text{length } (\text{levels } mp) = n\text{-levels } mp \implies$   
 $\text{free-list } (\text{levels } mp ! ii) \neq [] \implies$   
*False*  
**apply**(*simp add:head-free-list-def*)  
**apply**(*subgoal-tac hd (free-list (levels mp ! ii))  $\neq$  NULL*)  
**apply** *simp*  
**using** *hd-conv-nth* **by** *force*

**lemma** *mp-alloc-stm3-lm3:*

$Va \in \text{mp-alloc-precond1-70-2-2 } t \text{ } p \text{ } sz \text{ } timeout \cap \{\text{'cur} = \text{Some } t\} \implies$   
*(if level-empty (mem-pool-info Va p) (nat (free-l Va t)) then*  
 $\{V. V = Va(\text{blk} := (\text{blk } Va)(t := \text{NULL})) \wedge \text{level-empty } (\text{mem-pool-info } Va \text{ } p)$   
 $(\text{nat } (\text{free-l } Va \text{ } t))\}$   
*else*  
 $\{V. V = Va(\text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t))),$   
 $\text{mem-pool-info} := (\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t)))\}$   
 $\wedge \neg \text{level-empty } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t))\}) \cap$   
 $\neg \{\text{'blk } t \neq \text{NULL}\}$   
 $\subseteq \{\text{'id} \in \{\text{'(Pair } Va) \in \text{Mem-pool-alloc-guar } t\} \cap \text{mp-alloc-precond2-1 } t \text{ } p \text{ } sz \text{ } timeout\}$   
**apply** *clarsimp apply meson*

**apply**(*unfold Mem-pool-alloc-guar-def*)[1] **apply**(*rule UnI1*) **apply** *simp*  
**apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)  
**apply**(*rule conjI*)  
**apply**(*subgoal-tac (Va, Va(\text{blk} := (\text{blk } Va)(t := \text{NULL}))) \in \text{lvars-nochange1-4all}*)  
**using** *glnochange-inv0 apply auto[1] apply (simp add:lvars-nochange1-4all-def*  
*lvars-nochange1-def)*  
**apply**(*simp add:lvars-nochange-def*)

**apply**(*subgoal-tac (Va, Va(\text{blk} := (\text{blk } Va)(t := \text{NULL}))) \in \text{lvars-nochange1-4all}*)  
**using** *glnochange-inv0 apply auto[1] apply (simp add:lvars-nochange1-4all-def*  
*lvars-nochange1-def)*

**apply** *clarsimp*

**apply**(*subgoal-tac nat (free-l Va t)  $\geq$  0  $\wedge$  nat (free-l Va t) < n-levels (mem-pool-info*

```

Va p))
prefer 2 apply linarith
apply(subgoal-tac buf (mem-pool-info Va p) > 0)
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply meson
apply(subgoal-tac  $\forall i < \text{length } (\text{levels } (\text{mem-pool-info } Va p)).$ 
 $\forall j < \text{length } (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! i)).$ 
 $\text{buf } (\text{mem-pool-info } Va p) \leq (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! i))$ 
! j)
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def
Let-def)
apply clarify apply (metis Suc-leI lessI not-le trans-le-add1)
apply(subgoal-tac  $\text{length } (\text{levels } (\text{mem-pool-info } Va p)) = n\text{-levels } (\text{mem-pool-info } Va p)$ )
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply metis
apply(subgoal-tac  $\forall j < \text{length } (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! \text{nat } (\text{free-l } Va t))).$ 
 $\text{buf } (\text{mem-pool-info } Va p) \leq \text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! \text{nat } (\text{free-l } Va t)) ! j$ )
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply(rule conjI)
apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply clarsimp
apply(rule conjI)
apply(simp add:gvars-conf-stable-def gvars-conf-def rmhead-free-list-def) ap-
ply clarsimp
apply(case-tac  $\text{nat } (\text{free-l } Va t) \neq i$ ) apply simp apply simp
apply(rule conjI)
apply(case-tac  $\text{free-list } ((\text{levels } (\text{mem-pool-info } Va p)) ! (\text{nat } (\text{free-l } Va t))))$ 
= []
apply(subgoal-tac  $\text{rmhead-free-list } (\text{mem-pool-info } Va p) (\text{nat } (\text{free-l } Va t))$ 
= mem-pool-info Va p)
apply simp apply(subgoal-tac (Va, Va( $\text{blk} := (\text{blk } Va)(t := \text{NULL})$ )) $\in \text{lvars-nochange1-4all}$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
using mp-alloc-stm3-lm3-1[of  $\text{nat } (\text{free-l } Va t) \text{ mem-pool-info } Va p$ ] apply
meson
using mp-alloc-stm3-lm3-2[of  $\text{mem-pool-info } Va p \text{ nat } (\text{free-l } Va t)$ ] apply
meson
apply(simp add:lvars-nochange-def)
apply(rule conjI)
apply(case-tac  $\text{free-list } ((\text{levels } (\text{mem-pool-info } Va p)) ! (\text{nat } (\text{free-l } Va t))))$  =
[]
apply(subgoal-tac  $\text{rmhead-free-list } (\text{mem-pool-info } Va p) (\text{nat } (\text{free-l } Va t))$  =
mem-pool-info Va p)
apply simp apply(subgoal-tac (Va, Va( $\text{blk} := (\text{blk } Va)(t := \text{NULL})$ )) $\in \text{lvars-nochange1-4all}$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
using mp-alloc-stm3-lm3-1[of  $\text{nat } (\text{free-l } Va t) \text{ mem-pool-info } Va p$ ] apply

```

```

meson
  using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
metis
  apply(rule conjI) apply(simp add:rmhead-free-list-def)
  apply(rule conjI) apply(simp add:rmhead-free-list-def)
  apply(rule conjI) apply(simp add:rmhead-free-list-def) apply(simp add:rmhead-free-list-def)

  apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply simp
  apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
  apply(subgoal-tac (Va, Va(⟦blk := (blk Va)(t := NULL)⟧)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply(simp add:lvars-nochange-def)

  apply(subgoal-tac (Va, Va(⟦blk := (blk Va)(t := NULL)⟧)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

  apply clarsimp
  apply(subgoal-tac nat (free-l Va t) ≥ 0 ∧ nat (free-l Va t) < n-levels (mem-pool-info
Va p))
  prefer 2 apply linarith
  apply(subgoal-tac buf (mem-pool-info Va p) > 0)
  prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply meson
  apply(subgoal-tac ∀ i < length (levels (mem-pool-info Va p)).
    ∀ j < length (free-list (levels (mem-pool-info Va p) ! i)).
      buf (mem-pool-info Va p) ≤ (free-list (levels (mem-pool-info Va p) ! i))
! j)
  prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def
Let-def)
  apply clarify apply (metis Suc-leI lessI not-less trans-le-add1)
  apply(subgoal-tac length (levels (mem-pool-info Va p)) = n-levels (mem-pool-info
Va p))
  prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply metis
  apply(subgoal-tac ∀ j < length (free-list (levels (mem-pool-info Va p) ! nat (free-l
Va t)))).
    buf (mem-pool-info Va p) ≤ free-list (levels (mem-pool-info Va p) ! nat
(free-l Va t)) ! j)
  prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
  apply(rule conjI)
  apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply clarsimp
  apply(rule conjI)
  apply(simp add:gvars-conf-stable-def gvars-conf-def rmhead-free-list-def) ap-

```



```

ply clarsimp
  apply(case-tac nat (free-l Va t) ≠ i) apply simp apply simp
  apply(rule conjI)
  apply(case-tac free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va t))))
= []
  apply(subgoal-tac rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
= mem-pool-info Va p
  apply simp apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  using mp-alloc-stm3-lm3-1[of nat (free-l Va t) mem-pool-info Va p] apply
meson
  using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
meson
  apply(simp add:lvars-nochange-def)
  apply(rule conjI)
  apply(case-tac free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va t)))) =
[]
  apply(subgoal-tac rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t))) =
mem-pool-info Va p
  apply simp apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  using mp-alloc-stm3-lm3-1[of nat (free-l Va t) mem-pool-info Va p] apply
meson
  using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
metis
  apply(rule conjI) apply(simp add:rmhead-free-list-def)
  apply(rule conjI) apply(simp add:rmhead-free-list-def)
  apply(rule conjI) apply(simp add:rmhead-free-list-def) apply(simp add:rmhead-free-list-def)

  apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply simp
  apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
  apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply(simp add:lvars-nochange-def)

  apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

  apply clarsimp
  apply(subgoal-tac nat (free-l Va t) ≥ 0 ∧ nat (free-l Va t) < n-levels (mem-pool-info
Va p))

```

```

prefer 2 apply linarith
apply(subgoal-tac buf (mem-pool-info Va p) > 0)
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply meson
apply(subgoal-tac  $\forall i < \text{length } (\text{levels } (\text{mem-pool-info } Va p)).$ 
 $\forall j < \text{length } (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! i)).$ 
 $\text{buf } (\text{mem-pool-info } Va p) \leq (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! i))$ 
! j)
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def
Let-def)
apply clarify apply (metis Suc-leI lessI not-less trans-le-add1)
apply(subgoal-tac  $\text{length } (\text{levels } (\text{mem-pool-info } Va p)) = n\text{-levels } (\text{mem-pool-info } Va p)$ )
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply metis
apply(subgoal-tac  $\forall j < \text{length } (\text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! \text{nat } (\text{free-l } Va t)))$ .
 $\text{buf } (\text{mem-pool-info } Va p) \leq \text{free-list } (\text{levels } (\text{mem-pool-info } Va p) ! \text{nat } (\text{free-l } Va t)) ! j$ )
prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
apply(rule conjI)
apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply clarsimp
apply(rule conjI)
apply(simp add:gvars-conf-stable-def gvars-conf-def rmhead-free-list-def) ap-
ply clarsimp
apply(case-tac  $\text{nat } (\text{free-l } Va t) \neq i$ ) apply simp apply simp
apply(rule conjI)
apply(case-tac  $\text{free-list } ((\text{levels } (\text{mem-pool-info } Va p)) ! (\text{nat } (\text{free-l } Va t))))$ 
= []
apply(subgoal-tac  $\text{rmhead-free-list } (\text{mem-pool-info } Va p) (\text{nat } (\text{free-l } Va t))$ 
= mem-pool-info Va p)
apply simp apply(subgoal-tac (Va, Va( $\text{blk} := (\text{blk } Va)(t := \text{NULL})$ ))  $\in \text{lvars-nochange1-4all}$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
using mp-alloc-stm3-lm3-1[of  $\text{nat } (\text{free-l } Va t)$  mem-pool-info Va p] apply
meson
using mp-alloc-stm3-lm3-2[of mem-pool-info Va p  $\text{nat } (\text{free-l } Va t)$ ] apply
meson
apply(simp add:lvars-nochange-def)
apply(rule conjI)
apply(case-tac  $\text{free-list } ((\text{levels } (\text{mem-pool-info } Va p)) ! (\text{nat } (\text{free-l } Va t)))) =$ 
[])
apply(subgoal-tac  $\text{rmhead-free-list } (\text{mem-pool-info } Va p) (\text{nat } (\text{free-l } Va t)) =$ 
mem-pool-info Va p)
apply simp apply(subgoal-tac (Va, Va( $\text{blk} := (\text{blk } Va)(t := \text{NULL})$ ))  $\in \text{lvars-nochange1-4all}$ )
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
using mp-alloc-stm3-lm3-1[of  $\text{nat } (\text{free-l } Va t)$  mem-pool-info Va p] apply
meson

```

```

    using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
metis
    apply(rule conjI) apply(simp add:rmhead-free-list-def)
    apply(rule conjI) apply(simp add:rmhead-free-list-def)
    apply(simp add:rmhead-free-list-def)

    apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply simp
    apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
    apply(rule conjI)
    apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:lvars-nochange-def)

    apply(subgoal-tac (Va, Va(|blk := (blk Va)(t := NULL)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

    apply clarsimp
    apply(subgoal-tac nat (free-l Va t) ≥ 0 ∧ nat (free-l Va t) < n-levels (mem-pool-info
Va p))
    prefer 2 apply linarith
    apply(subgoal-tac buf (mem-pool-info Va p) > 0)
    prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
    apply meson
    apply(subgoal-tac ∀ i < length (levels (mem-pool-info Va p)).
    ∀ j < length (free-list (levels (mem-pool-info Va p) ! i)).
    buf (mem-pool-info Va p) ≤ (free-list (levels (mem-pool-info Va p) ! i))
! j)
    prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def
Let-def)
    apply clarify apply (metis Suc-leI lessI not-less trans-le-add1)
    apply(subgoal-tac length (levels (mem-pool-info Va p)) = n-levels (mem-pool-info
Va p))
    prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
    apply metis
    apply(subgoal-tac ∀ j < length (free-list (levels (mem-pool-info Va p) ! nat (free-l
Va t)))).
    buf (mem-pool-info Va p) ≤ free-list (levels (mem-pool-info Va p) ! nat
(free-l Va t)) ! j)
    prefer 2 apply(simp add:inv-def inv-mempool-info-def inv-bitmap-freelist-def)
    apply(rule conjI)
    apply(unfold Mem-pool-alloc-guar-def)[1] apply(rule UnI1) apply clarsimp
    apply(rule conjI)
    apply(simp add:gvars-conf-stable-def gvars-conf-def rmhead-free-list-def) ap-

```

```

ply clarsimp
  apply(case-tac nat (free-l Va t) ≠ i) apply simp apply simp
  apply(rule conjI)
  apply(case-tac free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va t))))
= []
  apply(subgoal-tac rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
= mem-pool-info Va p
  apply simp apply(subgoal-tac (Va, Va(⟦blk := (blk Va)(t := NULL)⟧) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  using mp-alloc-stm3-lm3-1[of nat (free-l Va t) mem-pool-info Va p] apply
meson
  using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
meson
  apply(simp add:lvars-nochange-def)
  apply(rule conjI)
  apply(case-tac free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va t)))) =
[]
  apply(subgoal-tac rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t))) =
mem-pool-info Va p
  apply simp apply(subgoal-tac (Va, Va(⟦blk := (blk Va)(t := NULL)⟧) ∈ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  using mp-alloc-stm3-lm3-1[of nat (free-l Va t) mem-pool-info Va p] apply
meson
  using mp-alloc-stm3-lm3-2[of mem-pool-info Va p nat (free-l Va t)] apply
metis
  apply(rule conjI) apply(simp add:rmhead-free-list-def)
  apply(rule conjI) apply(simp add:rmhead-free-list-def) apply(simp add:rmhead-free-list-def)

done

```

```

lemma mp-alloc-stm3-lm2-1:
  length (bits (levels mp ! ii)) =
    length (bits (levels mp [ii := (levels mp ! ii) (⟦free-list := fl⟧)
! ii)
      (⟦bits := bits (levels mp [ii := (levels mp ! ii) (⟦free-list := fl⟧) ! ii)
      [jj := ALLOCATING]⟧) ! ii))
  apply(case-tac ii < length (levels mp))
  apply simp
  apply auto
done

```

```

lemma mp-alloc-stm3-lm2-2:
  length (bits (levels mp ! ii)) =
    length (bits (levels mp [ii := (levels mp ! ii) (⟦free-list := fl, bits := bits (levels

```

```

mp ! ii) [jj := ALLOCATING]] ! ii))
  apply(case-tac ii < length (levels mp))
  apply simp apply auto
done

```

**lemma** *mp-alloc-stm3-body-meminfo*:

```

pa ≠ p ⇒
  set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info Va
p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t !
nat (free-l Va t)))
    pa = (mem-pool-info Va) pa
  by(simp add: set-bit-def)

```

**lemma** *mp-alloc-stm3-body-minf-buf*:

```

buf (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
  (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t
! nat (free-l Va t)))
  p) = buf (mem-pool-info Va p)
  by (simp add: set-bit-def rmhead-free-list-def)

```

**lemma** *mp-alloc-stm3-body-minf-maxsz*:

```

max-sz (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
  (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t
! nat (free-l Va t)))
  p) = max-sz (mem-pool-info Va p)
  by (simp add: set-bit-def rmhead-free-list-def)

```

**lemma** *mp-alloc-stm3-body-minf-nmax*:

```

n-max (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
  (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t
! nat (free-l Va t)))
  p) = n-max (mem-pool-info Va p)
  by (simp add: set-bit-def rmhead-free-list-def)

```

**lemma** *mp-alloc-stm3-body-minf-nlvl*:

```

n-levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info

```

$$\begin{aligned}
& \text{Va } p) \text{ (nat (free-l Va t))} \text{)) } p \\
& \quad (\text{nat (free-l Va t)}) \\
& \quad (\text{block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))} \\
& \quad \quad (\text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}) (\text{lsizes Va t} \\
& \quad \quad \text{! nat (free-l Va t))}) \\
& \quad \quad p) = \text{n-levels (mem-pool-info Va p)} \\
& \text{by (simp add: set-bit-def rmhead-free-list-def)}
\end{aligned}$$

**lemma** *mp-alloc-stm3-body-len-lvls:*

$$\begin{aligned}
& \text{length (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info} \\
& \text{Va p) (nat (free-l Va t))} \text{)) } p \\
& \quad (\text{nat (free-l Va t)}) \\
& \quad (\text{block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))} \\
& \quad \quad (\text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}) (\text{lsizes Va t ! nat} \\
& \quad \quad (\text{free-l Va t})) \\
& \quad \quad p)) = \text{length (levels (mem-pool-info Va p))} \\
& \text{by (simp add: set-bit-def rmhead-free-list-def)}
\end{aligned}$$

**lemma** *mp-alloc-stm3-body-len-bits:*

$$\begin{aligned}
& \text{length (bits (levels (set-bit-allocating} \\
& \quad ((\text{mem-pool-info Va})(p := \text{rmhead-free-list (mem-pool-info Va p) (nat} \\
& \quad (\text{free-l Va t})) \text{)) } p \\
& \quad (\text{nat (free-l Va t)}) \\
& \quad (\text{block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))} \\
& \quad \quad (\text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}) (\text{lsizes Va t !} \\
& \quad \quad \text{nat (free-l Va t))}) \\
& \quad \quad p) \text{ !} \\
& \quad ii) = \text{length (bits ((levels (mem-pool-info Va p))!ii))} \\
& \text{apply (simp add: set-bit-def rmhead-free-list-def block-num-def head-free-list-def)} \\
& \text{by (smt Mem-pool-lvl.select-convs(1) Mem-pool-lvl.surjective Mem-pool-lvl.update-convs(1)}
\end{aligned}$$

$$\begin{aligned}
& \text{Mem-pool-lvl.update-convs(2) list-update-beyond list-updt-samelen not-less} \\
& \text{nth-list-update-eq nth-list-update-neq)}
\end{aligned}$$

**lemma** *mp-alloc-stm3-body-frlst-otherlvl:*

$$\begin{aligned}
& ii \neq \text{nat (free-l Va t)} \implies \\
& \quad \text{free-list (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info} \\
& \quad \text{Va p) (nat (free-l Va t))} \text{)) } p \\
& \quad \quad (\text{nat (free-l Va t)}) \\
& \quad \quad (\text{block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l} \\
& \quad \quad \text{Va t))}) \\
& \quad \quad \quad (\text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}) \\
& \quad \quad \quad (\text{ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l} \\
& \quad \quad \quad \text{Va t))}) \\
& \quad \quad p) \text{ ! } ii) = \text{free-list (levels (mem-pool-info Va p) ! ii)} \\
& \text{by (simp add: set-bit-def rmhead-free-list-def block-num-def head-free-list-def)}
\end{aligned}$$

**lemma** *mp-alloc-stm3-body-frlst-samelvl:*

$$ii < \text{length (levels (mem-pool-info Va p))} \implies ii = \text{nat (free-l Va t)} \implies$$

$$\begin{aligned} & \text{free-list (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info} \\ & \text{Va p) (nat (free-l Va t)))) p} \\ & \quad (\text{nat (free-l Va t)}) \\ & \quad (\text{block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l} \\ & \text{Va t))}) \\ & \quad (\text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}) \\ & \quad (\text{ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l} \\ & \text{Va t))}) \\ & \quad p) ! ii) = \text{tl (free-list (levels (mem-pool-info Va p) ! ii))} \\ & \text{by (simp add: set-bit-def rmhead-free-list-def block-num-def head-free-list-def)} \end{aligned}$$

**lemma** *mp-alloc-stm3-lm2-inv-1-1*:  $(jj::\text{nat}) \neq (a - b) \text{ div } c \implies$   
 $(a - b) \text{ mod } c = 0 \implies$   
 $c \neq 0 \implies$   
 $b + jj * c \neq a$  **by** *auto*

**lemma** *mp-alloc-stm3-lm2-inv-1-2*:  
 $\exists n > 0. (a::\text{nat}) = 4 * n * 4 ^ b \implies$   
 $ii < b \implies 0 < a \text{ div } 4 ^ ii$   
**by** (*smt div-eq-0-iff divisors-zero less-imp-le-nat m-mod-div mod-if not-gr0 pow-mod-0*  
*power-not-zero zero-neq-numeral*)

**lemma** *mp-alloc-stm3-lm2-inv-1*:  
 $\neg \text{level-empty (mem-pool-info Va p) ii} \implies p \in \text{mem-pools Va} \implies$   
 $\text{inv-mempool-info Va} \implies$   
 $\forall ii < \text{length (lsizes Va t)}. \text{lsizes Va t ! ii} = \text{ALIGN4 (max-sz (mem-pool-info Va}$   
 $\text{p)) div 4 ^ ii} \implies$   
 $\forall p \in \text{mem-pools Va.}$   
 $\forall i < \text{length (levels (mem-pool-info Va p))}. \quad$   
 $(\forall j < \text{length (bits (levels (mem-pool-info Va p) ! i))).$   
 $(\text{get-bit-s Va p i j} = \text{FREE}) =$   
 $(\text{buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va p) div 4 ^}$   
 $i)$   
 $\in \text{set (free-list (levels (mem-pool-info Va p) ! i))}) \wedge$   
 $(\forall j < \text{length (free-list (levels (mem-pool-info Va p) ! i))).$   
 $\exists n < n\text{-max (mem-pool-info Va p) * 4 ^ i.}$   
 $\text{free-list (levels (mem-pool-info Va p) ! i) ! j} =$   
 $\text{buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4}$   
 $^ i)) \wedge$   
 $\text{distinct (free-list (levels (mem-pool-info Va p) ! i))} \implies$   
 $\text{length (lsizes Va t)} \leq \text{length (levels (mem-pool-info Va p))} \implies$   
 $ii < \text{length (lsizes Va t)} \implies$   
 $\text{length (bits (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list}$   
 $(\text{mem-pool-info Va p) ii)) p ii}$   
 $((\text{head-free-list (mem-pool-info Va p) ii} - \text{buf (rmhead-free-list}$   
 $(\text{mem-pool-info Va p) ii)) div}$   
 $\text{lsizes Va t ! ii)}$   
 $p) !$

```

ii)) =
length (bits (levels (mem-pool-info Va p) ! ii)) ==>
jj < length (bits (levels (set-bit-allocating
                           (λa. if a = p then rmhead-free-list (mem-pool-info Va p)
                           ii else mem-pool-info Va a) p ii
                           ((head-free-list (mem-pool-info Va p) ii - buf (rmhead-free-list
(mem-pool-info Va p) ii)) div
                           lsizes Va t ! ii)
                           p) !
                           ii)) ==>
nat (free-l Va t) = ii ==>
jj ≠ (head-free-list (mem-pool-info Va p) ii - buf (rmhead-free-list (mem-pool-info
Va p) ii)) div lsizes Va t ! ii ==>
buf (mem-pool-info Va p) + jj * (max-sz (mem-pool-info Va p) div 4 ^ ii) ≠
head-free-list (mem-pool-info Va p) ii
apply(subgoal-tac buf (rmhead-free-list (mem-pool-info Va p) ii) = buf (mem-pool-info
Va p))
prefer 2 apply(simp add:rmhead-free-list-def)
apply(subgoal-tac lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info Va p)) div
4 ^ ii)
prefer 2 apply metis
apply(subgoal-tac ALIGN4 (max-sz (mem-pool-info Va p)) = max-sz (mem-pool-info
Va p))
prefer 2 using inv-mempool-info-maxsz-align4 apply blast
apply(subgoal-tac (head-free-list (mem-pool-info Va p) ii - buf (rmhead-free-list
(mem-pool-info Va p) ii)) mod lsizes Va t ! ii = 0)
prefer 2 apply(simp add:inv-mempool-info-def head-free-list-def Let-def)
apply(subgoal-tac ∃ n. hd (free-list (levels (mem-pool-info Va p) ! ii)) =
buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ ii))
prefer 2 apply(simp add:level-empty-def)
apply(subgoal-tac ∀ j < length (free-list (levels (mem-pool-info Va p) ! ii)).
(∃ n < n-max (mem-pool-info Va p) * 4 ^ ii. free-list (levels (mem-pool-info
Va p) ! ii) ! j =
buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4
^ ii)))
prefer 2 apply (simp add: linorder-not-less)
apply (metis (full-types, hide-lams) hd-conv-nth length-greater-0-conv)
apply (metis (no-types, hide-lams) add-diff-cancel-left' mod-mult-self2-is-0)

apply(subgoal-tac lsizes Va t ! ii ≠ 0)
prefer 2 apply(simp add:inv-mempool-info-def Let-def)
apply(subgoal-tac ∃ n > 0. max-sz (mem-pool-info Va p) = 4 * n * 4 ^ n-levels
(mem-pool-info Va p))
prefer 2 apply blast
apply(subgoal-tac length (levels (mem-pool-info Va p)) = n-levels (mem-pool-info
Va p))
prefer 2 apply simp
using mp-alloc-stm3-lm2-inv-1-2[of max-sz (mem-pool-info Va p) n-levels (mem-pool-info

```



$Va\ p)\ ii]$   
**apply** (*metis* (*no-types*, *lifting*) *add-lessD1 le-eq-less-or-eq less-imp-add-positive*)  
**using** *mp-alloc-stm3-lm2-inv-1-1* [*of jj head-free-list (mem-pool-info Va p) ii buf*  
*(rmhead-free-list (mem-pool-info Va p) ii) lsize Va t ! ii*]  
**apply** *auto*[1]  
**done**

**lemma** *mp-alloc-stm3-lm2-inv-2*:  
 $(a::nat) + jj * b \neq c \implies \exists n. a + n * b = c \implies$   
 $(c - a) \text{ div } b \neq jj$  **by** *auto*

**lemma** *mp-alloc-stm3-lm2-inv-thd-waitq*:  
*inv-thd-waitq Va*  $\implies$   
*inv-thd-waitq*  
 $(Va(\text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t)))),$   
 $\text{mem-pool-info} :=$   
 $\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info}$   
 $Va\ p) (\text{nat } (\text{free-l } Va\ t))))\ p (\text{nat } (\text{free-l } Va\ t)))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))$   
 $(\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va\ p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } Va\ t))),$   
 $\text{allocating-node} := \text{allocating-node } Va(t \mapsto$   
 $(\text{pool} = p, \text{level} = \text{nat } (\text{free-l } Va\ t),$   
 $\text{block} = \text{block-num}$   
 $(\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list}$   
 $(\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))\ p (\text{nat } (\text{free-l } Va\ t)))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l}$   
 $Va\ t)))) (\text{head-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))$   
 $(\text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va\ p)) \text{ div } 4 \wedge \text{nat } (\text{free-l}$   
 $Va\ t))))$   
 $p)$   
 $(\text{head-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))) (\text{ALIGN4}$   
 $(\text{max-sz } (\text{mem-pool-info } Va\ p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } Va\ t)),$   
 $\text{data} = \text{head-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))))$   
**apply**(*simp add:inv-thd-waitq-def*)  
**apply**(*rule conjI*) **apply**(*simp add: rmhead-free-list-def*  
 $\text{head-free-list-def set-bit-def block-num-def}$ )  
**apply**(*rule conjI*) **apply**(*simp add: rmhead-free-list-def*  
 $\text{head-free-list-def set-bit-def block-num-def}$ ) **apply** *metis*  
**apply**(*rule conjI*) **apply**(*simp add: rmhead-free-list-def*  
 $\text{head-free-list-def set-bit-def block-num-def}$ )  
**apply**(*simp add: rmhead-free-list-def*  
 $\text{head-free-list-def set-bit-def block-num-def}$ ) **apply** *metis*  
**done**

**lemma** *mp-alloc-stm3-lm2-inv-aux-vars-1*:  
 $\neg (\text{pool } n = p \wedge \text{level } n = \text{nat } (\text{free-l } Va\ t) \wedge \text{block } n =$   
 $\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va\ p) (\text{nat } (\text{free-l } Va\ t))))$

$$\begin{aligned}
& (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t))) \implies \\
& \text{get-bit-s } (Va \mid \text{mem-pool-info } := \\
& \quad \text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list} \\
& (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t)))) \\
& \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \\
& t)))))) \\
& (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{get-bit-s } Va \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n)
\end{aligned}$$

**apply**(rule subst[**where**  $t = \text{get-bit-s}$   
 $(Va \mid \text{mem-pool-info } :=$   
 $\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info}$   
 $Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \ (\text{nat } (\text{free-l } Va \ t))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t))))$   
 $(\text{pool } n) \ (\text{level } n) \ (\text{block } n)$  **and**  $s = \text{get-bit-s}$   
 $(Va \mid \text{mem-pool-info } :=$   
 $\text{set-bit-allocating } (\text{mem-pool-info } Va) \ p \ (\text{nat } (\text{free-l } Va \ t))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t))))$   
 $(\text{pool } n) \ (\text{level } n) \ (\text{block } n)]$ )  
**apply**(simp add:rmhead-free-list-def set-bit-def)  
**apply** (smt Mem-pool-lvl.select-convs(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

$\text{Mem-pool-lvl.update-convs}(2) \text{ linorder-not-less list-update-beyond nth-list-update-eq}$   
 $\text{nth-list-update-neq})$   
**apply**(simp add:rmhead-free-list-def set-bit-def)  
**apply** (smt Mem-pool-lvl.select-convs(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective  
 $\text{Mem-pool-lvl.update-convs}(2) \text{ linorder-not-less list-update-beyond nth-list-update-eq}$   
 $\text{nth-list-update-neq})$   
**done**

**lemma** mp-alloc-stm3-lm2-inv-aux-vars-2:

$\text{inv-mempool-info } Va \implies$   
 $\neg \text{level-empty } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)) \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $\text{pool } n = p \wedge$   
 $\text{level } n = \text{nat } (\text{free-l } Va \ t) \wedge$   
 $\text{block } n =$   
 $\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \ (\text{head-free-list}$   
 $(\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t)) \implies$

```

get-bit-s (Va | mem-pool-info :=
  set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
  (nat (free-l Va t))
  (block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info Va
p) ! nat (free-l Va t)) ! NULL)
  (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))))
(pool n) (level n) (block n) =
get-bit
(set-bit-allocating (mem-pool-info Va) p (nat (free-l Va t))
(block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info Va p) !
nat (free-l Va t)) ! NULL)
(max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))))
p (nat (free-l Va t))
(block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info Va p) !
nat (free-l Va t)) ! NULL)
(max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t)))
apply(rule subst[where t=p and s=pool n]) apply simp
apply(rule subst[where t=nat (free-l Va t) and s=level n]) apply simp
apply(rule subst[where t=(block-num (mem-pool-info Va (pool n)) (free-list (levels
(mem-pool-info Va (pool n)) ! level n) ! NULL)
(max-sz (mem-pool-info Va (pool n)) div 4 ^ level n)) and s=block n])
apply(simp add:level-empty-def block-num-def rmhead-free-list-def head-free-list-def)
apply (metis hd-conv-nth inv-mempool-info-maxsz-align4)
apply(rule subst[where t=get-bit-s
(Va | mem-pool-info :=
  set-bit-allocating ((mem-pool-info Va)(pool n := rmhead-free-list (mem-pool-info
Va (pool n)) (level n))) (pool n) (level n)
  (block n))])
  (pool n) (level n) (block n) and
  s=get-bit (set-bit-allocating ((mem-pool-info Va)(pool n := rmhead-free-list
(mem-pool-info Va (pool n)) (level n))) (pool n) (level n)
  (block n)) (pool n) (level n) (block n))] apply auto[1]
apply(simp add:rmhead-free-list-def set-bit-def)
apply(case-tac level n ≥ length (levels (mem-pool-info Va (pool n))))
apply auto
done

```

**lemma** mp-alloc-stm3-lm2-inv-aux-vars:

```

¬ level-empty (mem-pool-info Va p) (nat (free-l Va t)) ⇒
  allocating-node Va t = None ⇒
  freeing-node Va t = None ⇒
  inv-cur Va ∧ inv-thd-waitq Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist Va
∧ inv-bitmap Va ∧ inv-aux-vars Va ⇒
  p ∈ mem-pools Va ⇒
  ETIMEOUT ≤ timeout ⇒
  timeout = ETIMEOUT ⇒ tmout Va t = ETIMEOUT ⇒
  ¬ rf Va t ⇒
  ∀ ii < length (lsizes Va t). lsize Va t ! ii = ALIGN4 (max-sz (mem-pool-info Va

```

$p)) \text{ div } 4 \wedge ii \implies$   
 $\text{length } (lsizes \text{ Va } t) \leq n\text{-levels } (mem\text{-pool-info } \text{Va } p) \implies$   
 $\text{alloc-l } \text{Va } t < \text{int } (n\text{-levels } (mem\text{-pool-info } \text{Va } p)) \implies$   
 $\text{free-l } \text{Va } t \leq \text{alloc-l } \text{Va } t \implies$   
 $\neg \text{free-l } \text{Va } t < OK \implies$   
 $\text{alloc-l } \text{Va } t = \text{int } (\text{length } (lsizes \text{ Va } t)) - 1 \wedge \text{length } (lsizes \text{ Va } t) = n\text{-levels}$   
 $(mem\text{-pool-info } \text{Va } p) \vee$   
 $\text{alloc-l } \text{Va } t = \text{int } (\text{length } (lsizes \text{ Va } t)) - 2 \wedge lsize \text{ Va } t ! \text{ nat } (\text{alloc-l } \text{Va } t +$   
 $1) < sz \implies$   
 $\text{length } (lsizes \text{ Va } t) \leq \text{length } (levels (mem\text{-pool-info } \text{Va } p)) \implies$   
 $\text{nat } (\text{free-l } \text{Va } t) < \text{length } (lsizes \text{ Va } t) \implies$   
 $inv\text{-aux-vars}$   
 $(\text{Va} \backslash blk := (blk \text{ Va})(t := \text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va}$   
 $t))))),$   
 $mem\text{-pool-info} :=$   
 $set\text{-bit-allocating } ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list } (mem\text{-pool-info}$   
 $\text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p (\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } \text{Va } t))),$   
 $allocating\text{-node} := allocating\text{-node } \text{Va}(t \mapsto$   
 $(pool = p, level = \text{nat } (\text{free-l } \text{Va } t),$   
 $block = \text{block-num}$   
 $(set\text{-bit-allocating } ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list}$   
 $(mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p$   
 $(\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l}$   
 $\text{Va } t))))$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l}$   
 $\text{Va } t))))$   
 $p)$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l}$   
 $\text{Va } t))),$   
 $data = \text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))))$   
**apply**(*unfold inv-aux-vars-def*)  
**apply**(*rule conjI*)  
  
**apply** *clarify*  
**apply**(*subgoal-tac freeing-node*  
 $(\text{Va} \backslash blk := (blk \text{ Va})(t := \text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va}$   
 $t))))),$   
 $mem\text{-pool-info} :=$   
 $set\text{-bit-allocating } ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list}$   
 $(mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p$   
 $(\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va}$   
 $t))))$

$(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t)))$ ,  
 $\text{allocating-node} := \text{allocating-node } Va(t \mapsto$   
 $\quad (\text{pool} = p, \text{level} = \text{nat } (\text{free-l } Va \ t),$   
 $\quad \text{block} = \text{block-num}$   
 $\quad (\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list}$   
 $\quad (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$   
 $\quad p \ (\text{nat } (\text{free-l } Va \ t)))$   
 $\quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat}$   
 $\quad (\text{free-l } Va \ t))))$   
 $\quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $\quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat}$   
 $\quad (\text{free-l } Va \ t)))$   
 $\quad p)$   
 $\quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $\quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat}$   
 $\quad (\text{free-l } Va \ t))),$   
 $\text{data} = \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $= \text{freeing-node } Va)$   
**prefer 2 apply simp**  
**apply**(*subgoal-tac get-bit-s*  $Va \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{FREEING}$ )  
**prefer 2 apply auto[1]**  
**apply**(*case-tac*  $(\text{pool } n) = p \wedge (\text{level } n) = \text{nat } (\text{free-l } Va \ t)$   
 $\wedge (\text{block } n) = (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l}$   
 $Va \ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va$   
 $t))))$   
**apply**(*subgoal-tac get-bit-s*  $Va \ p \ (\text{nat } (\text{free-l } Va \ t)) \ (\text{block-num } (\text{mem-pool-info}$   
 $Va \ p)$   
 $((\text{free-list } ((\text{levels } (\text{mem-pool-info } Va \ p)) \ ! \ (\text{nat } (\text{free-l } Va$   
 $t)))) \ ! \ 0)$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge (\text{nat } (\text{free-l } Va \ t))))$   
 $= \text{FREE}$ )  
**prefer 2 apply**(*simp add:level-empty-def*) **using** *inv-bitmap-freelist-fl-FREE*  
**apply auto[1]**  
**apply**(*subgoal-tac block-num*  $(\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l}$   
 $Va \ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t)))$   
 $= (\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{free-list } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ \text{nat}$   
 $(\text{free-l } Va \ t)) \ ! \ \text{NULL})$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge \text{nat } (\text{free-l } Va \ t))))$   
**prefer 2 apply**(*simp add:rmhead-free-list-def head-free-list-def block-num-def*  
*level-empty-def*)  
**apply** (*metis hd-conv-nth inv-mempool-info-maxsz-align4*)  
**apply simp**

**apply**(*subgoal-tac get-bit-s*  
 $(Va \backslash blk := (blk \ Va)(t := head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t))),$   
 $mem-pool-info :=$   
 $set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))) \ p$   
 $(nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va$   
 $t))))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN_4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat \ (free-l \ Va$   
 $t))),$   
 $allocating-node := allocating-node \ Va(t \mapsto$   
 $\backslash pool = p, \ level = nat \ (free-l \ Va \ t),$   
 $block = block-num$   
 $(set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t))))$   
 $p \ (nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat$   
 $(free-l \ Va \ t))))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN_4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat$   
 $(free-l \ Va \ t)))$   
 $p)$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN_4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat$   
 $(free-l \ Va \ t))),$   
 $data = head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(pool \ n) \ (level \ n) \ (block \ n) = get-bit-s \ Va \ (pool \ n) \ (level \ n) \ (block \ n))$   
**apply simp**  
**apply**(*subgoal-tac get-bit-s*  
 $(Va \backslash blk := (blk \ Va)(t := head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t))),$   
 $mem-pool-info :=$   
 $set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))) \ p$   
 $(nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va$   
 $t))))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN_4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat \ (free-l \ Va$   
 $t))),$   
 $allocating-node := allocating-node \ Va(t \mapsto$   
 $\backslash pool = p, \ level = nat \ (free-l \ Va \ t),$   
 $block = block-num$   
 $(set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t))))$   
 $p \ (nat \ (free-l \ Va \ t))$

```

(block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
p)
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)),
data = head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(pool n) (level n) (block n)
= get-bit-s (Va| mem-pool-info :=
set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
(nat (free-l Va t))
(block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))) (pool n) (level n) (block n))
prefer 2 apply force
apply(frule mp-alloc-stm3-lm2-inv-aux-vars-1) apply simp

apply(rule conjI)

apply clarify
apply(subgoal-tac get-bit-s
(Va| blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l
Va t))),
mem-pool-info :=
set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
(nat (free-l Va t))
(block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))),
allocating-node := allocating-node Va(t ↦
|pool = p, level = nat (free-l Va t),
block = block-num
(set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t))))
p (nat (free-l Va t))
(block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t))))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))

```

$$\begin{aligned}
& p) \\
& (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat} \\
& (\text{free-l } Va \ t)), \\
& \text{data} = \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& (\text{pool } n) \ (\text{level } n) \ (\text{block } n) \\
& = \text{get-bit-s } (Va \mid \text{mem-pool-info} := \\
& \quad \text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list} \\
& (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t)))) \\
& \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \\
& t)))) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n)) \\
& \text{prefer } 2 \text{ apply force} \\
& \text{apply}(\text{case-tac } (\text{pool } n) = p \wedge (\text{level } n) = \text{nat } (\text{free-l } Va \ t) \\
& \wedge (\text{block } n) = (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l} \\
& Va \ t)))) \\
& \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \\
& t)))) \\
& \text{apply}(\text{subgoal-tac } \text{get-bit-s } (Va \mid \text{mem-pool-info} := \\
& \quad \text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list} \\
& (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t)))) \\
& \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \\
& t)))) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) = \text{ALLOCATING}) \\
& \text{apply simp} \\
& \text{apply}(\text{subgoal-tac } \text{get-bit } (\text{set-bit-allocating } (\text{mem-pool-info } Va) \ p \ (\text{nat } (\text{free-l} \\
& Va \ t)) \\
& \quad (\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{free-list } (\text{levels} \\
& (\text{mem-pool-info } Va \ p) \ ! \ \text{nat } (\text{free-l } Va \ t)) \ ! \ \text{NULL}) \\
& \quad (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \\
& t)))) \\
& \quad p \ (\text{nat } (\text{free-l } Va \ t)) \\
& \quad (\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{free-list } (\text{levels } (\text{mem-pool-info} \\
& Va \ p) \ ! \ \text{nat } (\text{free-l } Va \ t)) \ ! \ \text{NULL}) \\
& \quad (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t))) = \\
& \text{ALLOCATING}) \\
& \text{prefer } 2 \\
& \text{apply}(\text{rule } \text{set-bit-get-bit-eq}[\text{of } \text{nat } (\text{free-l } Va \ t) \ \text{mem-pool-info } Va \ p \\
& \quad \text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{free-list } (\text{levels } (\text{mem-pool-info} \\
& Va \ p) \ ! \ \text{nat } (\text{free-l } Va \ t)) \ ! \ \text{NULL}) \\
& \quad (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t)) \ \text{set-bit-allocating} \\
& (\text{mem-pool-info } Va) \ p \ (\text{nat } (\text{free-l } Va \ t))
\end{aligned}$$



```

      (block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info Va p) ! nat
(free-l Va t)) ! NULL)
      (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))))))
apply simp apply(simp add:level-empty-def) using inv-bitmap-freelist-ft-bnum-in[of
Va p nat (free-l Va t) 0]
apply (meson le-trans length-greater-0-conv linorder-not-less) apply simp

```

```

apply(subgoal-tac get-bit-s (Va| mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))) (pool n) (level n) (block n)
      = get-bit (set-bit-allocating (mem-pool-info Va) p (nat (free-l Va t))
      (block-num (mem-pool-info Va p) (free-list (levels
(mem-pool-info Va p) ! nat (free-l Va t)) ! NULL)
      (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va
t))))
      p (nat (free-l Va t))
      (block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info
Va p) ! nat (free-l Va t)) ! NULL)
      (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))))

```

```

apply simp
apply(rule subst[where t=block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))
      and s=block-num (mem-pool-info Va p) (free-list (levels
(mem-pool-info Va p) ! nat (free-l Va t)) ! NULL)
      (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t)))]
apply(simp add:level-empty-def block-num-def rmhead-free-list-def head-free-list-def)
apply (metis hd-conv-nth inv-mempool-info-maxsz-align4)
using mp-alloc-stm3-lm2-inv-aux-vars-2[of Va p t] apply blast

```

```

apply(subgoal-tac get-bit (mem-pool-info Va) (pool n) (level n) (block n) =
FREEING)
prefer 2
apply(subgoal-tac get-bit-s
      (Va| mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))

```

```

      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))))))
    (pool n) (level n) (block n) = get-bit (mem-pool-info Va) (pool n) (level n)
(block n))
  prefer 2 using mp-alloc-stm3-lm2-inv-aux-vars-1[of - p Va t] apply blast
  apply simp
  apply(subgoal-tac mem-block-addr-valid Va n)
  prefer 2 apply(simp add:mem-block-addr-valid-def)
  apply (metis mp-alloc-stm3-body-meminfo mp-alloc-stm3-body-minf-buf mp-alloc-stm3-body-minf-maxsz)
  apply(subgoal-tac  $\exists t. \text{freeing-node } Va\ t = \text{Some } n$ ) prefer 2 apply metis
  apply(subgoal-tac  $\forall ta. \text{freeing-node } Va\ ta = \text{freeing-node}$ 
    (Va(blk := (blk Va))(t := head-free-list (mem-pool-info Va p) (nat
(free-l Va t)))),
    mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))),
    allocating-node := allocating-node Va(t  $\mapsto$ 
      (pool = p, level = nat (free-l Va t),
      block = block-num
      (set-bit-allocating
      ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))))) ta)
  prefer 2 apply force
  apply metis

  apply(rule conjI)

  apply clarify

```

**apply**(*subgoal-tac get-bit-s*  
 $(Va \llbracket blk := (blk \ Va)(t := head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$ ),  
 $mem-pool-info :=$   
 $set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))) \ p$   
 $(nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l$   
 $Va \ t)))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat \ (free-l \ Va$   
 $t)))$ ,  
 $allocating-node := allocating-node \ Va(t \mapsto$   
 $\llbracket pool = p, level = nat \ (free-l \ Va \ t),$   
 $block = block-num$   
 $(set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list$   
 $(mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t))))$   
 $p \ (nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat$   
 $(free-l \ Va \ t)))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va$   
 $t)))$   
 $(ALIGN4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat$   
 $(free-l \ Va \ t)))$   
 $p)$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat$   
 $(free-l \ Va \ t))$ ,  
 $data = head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)) \rrbracket \rrbracket$   
 $(pool \ n) \ (level \ n) \ (block \ n)$   
 $= get-bit-s \ (Va \llbracket mem-pool-info :=$   
 $set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list \ (mem-pool-info$   
 $Va \ p) \ (nat \ (free-l \ Va \ t)))) \ p$   
 $(nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat \ (free-l \ Va$   
 $t)))) \rrbracket \rrbracket \ (pool \ n) \ (level \ n) \ (block \ n))$   
**prefer 2 apply force**  
**apply**(*subgoal-tac get-bit-s* ( $Va \llbracket mem-pool-info :=$   
 $set-bit-allocating \ ((mem-pool-info \ Va)(p := rmhead-free-list \ (mem-pool-info$   
 $Va \ p) \ (nat \ (free-l \ Va \ t)))) \ p$   
 $(nat \ (free-l \ Va \ t))$   
 $(block-num \ (rmhead-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(head-free-list \ (mem-pool-info \ Va \ p) \ (nat \ (free-l \ Va \ t)))$   
 $(ALIGN4 \ (max-sz \ (mem-pool-info \ Va \ p)) \ div \ 4 \ ^ \ nat \ (free-l \ Va$   
 $t)))) \rrbracket \rrbracket \ (pool \ n) \ (level \ n) \ (block \ n)$   
 $= get-bit-s \ (Va \llbracket mem-pool-info := set-bit-allocating \ (mem-pool-info$   
 $Va) \ p$

```

      (nat (free-l Va t))
      (block-num (mem-pool-info Va p)
        (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))))) (pool n) (level n) (block n))
prefer 2 apply(simp add:rmhead-free-list-def set-bit-def block-num-def)
apply (smt Mem-pool-lvl.select-conv(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective

linorder-not-less list-update-beyond nth-list-update-eq nth-list-update-neq)

apply(subgoal-tac get-bit-s (Va| mem-pool-info := set-bit-allocating (mem-pool-info
Va) p
      (nat (free-l Va t))
      (block-num (mem-pool-info Va p)
        (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))))) (pool n) (level n) (block n) = ALLOCATING)
apply simp
apply(case-tac t = ta)

apply(subgoal-tac (pool n) = p ∧ (level n) = nat (free-l Va t)
  ∧ (block n) = block-num (mem-pool-info Va p)
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
prefer 2 apply(rule conjI) apply auto[1] apply(rule conjI) apply auto[1]
apply(subgoal-tac (block n) = block-num (set-bit-allocating ((mem-pool-info
Va)(p := rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t))))
  p (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
  p)
  (head-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
prefer 2 apply auto[1]
apply(subgoal-tac block-num (set-bit-allocating ((mem-pool-info Va)(p :=
rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t))))
  p (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^

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nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))
      = block-num (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
    prefer 2 apply (simp add:level-empty-def block-num-def set-bit-def rmhead-free-list-def)
    apply simp
    apply (subgoal-tac nat (free-l Va t) < length (levels (mem-pool-info Va p)))
    prefer 2 apply simp
    apply (subgoal-tac block-num (mem-pool-info Va p) (head-free-list (mem-pool-info
Va p) (nat (free-l Va t))))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))
      < length (bits (levels (mem-pool-info Va p) ! nat (free-l
Va t))))
    prefer 2 apply (rule subst[where t=ALIGN4 (max-sz (mem-pool-info Va
p)) and s=max-sz (mem-pool-info Va p)])
    apply (metis inv-mempool-info-maxsz-align4)
    apply (frule inv-bitmap-freelist-fl-bnum-in[of Va p nat (free-l Va t) 0])
    apply simp apply simp apply simp apply (simp add:level-empty-def)
    apply (simp add:level-empty-def head-free-list-def) apply (metis hd-conv-nth)

    using set-bit-get-bit-eq2[of nat (free-l Va t) Va p block-num (mem-pool-info
Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))]
    ALLOCATING] apply metis

    apply (subgoal-tac allocating-node
      (Va[blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l
Va t))),
      mem-pool-info :=
        set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))),
      allocating-node := allocating-node Va(t ↦
      [pool = p, level = nat (free-l Va t),

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      block = block-num
      (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t))))
      p (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t))))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      ta = allocating-node Va ta) prefer 2 apply force
      apply(subgoal-tac get-bit (mem-pool-info Va) (pool n) (level n) (block n) =
ALLOCATING)
      prefer 2 apply metis
      apply(subgoal-tac block-num (mem-pool-info Va p)
      ((free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va
t)))) ! 0)
      (max-sz (mem-pool-info Va p) div 4 ^ (nat (free-l Va t)))
      = block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
      prefer 2 apply(simp add:block-num-def rmhead-free-list-def head-free-list-def)
      apply (simp add: hd-conv-nth inv-mempool-info-maxsz-align4 level-empty-def)
      apply(case-tac (pool n) = p ∧ (level n) = nat (free-l Va t)
      ∧ (block n) = (block-num (rmhead-free-list (mem-pool-info Va
p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t))))
      apply(subgoal-tac get-bit-s Va p (nat (free-l Va t)) (block-num (mem-pool-info
Va p)
      ((free-list ((levels (mem-pool-info Va p)) ! (nat (free-l Va
t)))) ! 0)
      (max-sz (mem-pool-info Va p) div 4 ^ (nat (free-l Va t))))
      = FREE)
      prefer 2 apply(simp add:level-empty-def) using inv-bitmap-freelist-fl-FREE[of
Va p nat (free-l Va t) 0]
      apply auto[1]

apply simp

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```

apply(subgoal-tac get-bit-s (Va (mem-pool-info :=
  set-bit-allocating (mem-pool-info Va) p (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))))
  (pool n) (level n) (block n) = get-bit-s Va (pool n) (level n) (block n))
prefer 2 apply (metis set-bit-get-bit-neq2)
apply(rule subst[where t=ALIGN4 (max-sz (mem-pool-info Va p)) and
s=max-sz (mem-pool-info Va p)])
apply (metis inv-mempool-info-maxsz-align4)
apply (simp add: hd-conv-nth level-empty-def)
apply (smt nat-less-iff nth-equalityI set-bit-get-bit-eq set-bit-get-bit-neq
set-bit-prev-len zle-int)

apply(rule conjI)

apply clarify
apply(case-tac (pool n) = p ∧ (level n) = nat (free-l Va t)
  ∧ (block n) = (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))))
apply(subgoal-tac allocating-node
  (Va (blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat
(free-l Va t))),
  mem-pool-info :=
    set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))),
  allocating-node := allocating-node Va(t ↦
    (pool = p, level = nat (free-l Va t),
    block = block-num
    (set-bit-allocating
    ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t))))

```

$$\begin{aligned}
& p) \\
& (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^ \ \text{nat} \\
& (\text{free-l } Va \ t)), \\
& \text{data} = \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t))) t = \\
& \text{Some } n) \\
& \text{prefer } 2 \ \text{apply}(\text{rule subst}[\text{where } t = \text{allocating-node} \\
& \quad (Va \ (\text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat} \\
& (\text{free-l } Va \ t))), \\
& \quad \text{mem-pool-info} := \\
& \quad \quad \text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list} \\
& (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l} \\
& Va \ t)))) \\
& \quad \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& \quad \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^ \ \text{nat } (\text{free-l} \\
& Va \ t))), \\
& \text{allocating-node} := \text{allocating-node } Va(t \mapsto \\
& \quad (\text{pool} = p, \text{level} = \text{nat } (\text{free-l } Va \ t), \\
& \quad \text{block} = \text{block-num} \\
& \quad \quad (\text{set-bit-allocating} \\
& \quad \quad \quad ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info} \\
& Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad \quad \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad \quad \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \\
& (\text{nat } (\text{free-l } Va \ t))) \\
& \quad \quad \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t))) \\
& \quad \quad \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^ \ \text{nat} \\
& \text{nat } (\text{free-l } Va \ t))) \\
& p) \\
& (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \\
& (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^ \ \text{nat} \\
& (\text{free-l } Va \ t)), \\
& \text{data} = \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t))) t \ \text{and} \ s = \text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{free-l } Va \ t), \\
& \text{block} = \text{block-num} \\
& \quad (\text{set-bit-allocating} \\
& \quad \quad ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info} \\
& Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \\
& \quad \quad (\text{nat } (\text{free-l } Va \ t)) \\
& \quad \quad (\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \\
& (\text{nat } (\text{free-l } Va \ t))) \\
& \quad \quad (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \\
& t))) \\
& \quad \quad (\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^ \ \text{nat} \\
& \text{nat } (\text{free-l } Va \ t)))
\end{aligned}$$



```

      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))])])
    apply force
    apply(rule subst[where t=block-num
      (set-bit-allocating
        ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
        (nat (free-l Va t))
        (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
        (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))
      and s=block-num (rmhead-free-list (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))])
    apply(simp add: set-bit-def rmhead-free-list-def block-num-def)
    apply(simp add: mem-block-addr-valid-def)
    apply(subgoal-tac buf (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)))
      p) +
      block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)) *
      (max-sz (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))
      p) div
4 ^ nat (free-l Va t)) = head-free-list (mem-pool-info Va p) (nat (free-l Va

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t)))
  apply auto[1]
  apply(rule subst[where t=buf (set-bit-allocating ((mem-pool-info Va)(p :=
rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)))
    p) and s=buf (mem-pool-info Va p)])
  apply(simp add:set-bit-def block-num-def rmhead-free-list-def)
  apply(rule subst[where t=block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
    ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4
^ nat (free-l Va t)))
    and s=block-num (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))))])
  apply(simp add:set-bit-def block-num-def rmhead-free-list-def)
  apply(rule subst[where t=max-sz (set-bit-allocating ((mem-pool-info Va)(p
:= rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (mem-pool-info Va p) (head-free-list (mem-pool-info Va
p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t)))
    p) and s=max-sz (mem-pool-info Va p)])
  apply(simp add:set-bit-def block-num-def rmhead-free-list-def)
  apply(rule subst[where t=ALIGN4 (max-sz (mem-pool-info Va p)) and
s=max-sz (mem-pool-info Va p)])
  apply (metis inv-mempool-info-maxsz-align4)

  apply(rule ref-byblkn-self[of Va p head-free-list (mem-pool-info Va p) (nat
(free-l Va t)) (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))])
  apply(simp add:level-empty-def head-free-list-def)
  using inv-buf-le-fl[of Va p nat (free-l Va t) 0]
  apply (smt hd-conv-nth length-greater-0-conv nat-less-iff zle-int)
  apply(simp add:level-empty-def head-free-list-def)
  using inv-fl-mod-sz0[of Va p nat (free-l Va t) 0]
  apply (smt hd-conv-nth le-eq-less-or-eq le-trans length-greater-0-conv nat-eq-iff
nat-less-iff)
  apply auto[1]

  apply(subgoal-tac get-bit (mem-pool-info Va) (pool n) (level n) (block n) =
ALLOCATING)
  prefer 2

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apply(subgoal-tac get-bit-s
  (Va (mem-pool-info :=
    set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
    (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))))))
  (pool n) (level n) (block n) = get-bit (mem-pool-info Va) (pool n) (level n)
(block n))
  prefer 2 using mp-alloc-stm3-lm2-inv-aux-vars-1[of - p Va t] apply blast
apply force

apply(subgoal-tac  $\exists ta. ta \neq t \wedge$  allocating-node Va ta = Some n)
  prefer 2 apply(subgoal-tac mem-block-addr-valid Va n) apply metis
  apply(simp add:mem-block-addr-valid-def)
apply (metis mp-alloc-stm3-body-meminfo mp-alloc-stm3-body-minf-buf mp-alloc-stm3-body-minf-maxsz)

apply auto[1]

apply(rule conjI)

apply clarify
apply(subgoal-tac  $\forall t. \text{freeing-node}$ 
  (Va (blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l Va
t))),
    mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
        (nat (free-l Va t))
        (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))),
      allocating-node := allocating-node Va(t  $\mapsto$ 
        (pool = p, level = nat (free-l Va t),
          block = block-num
            (set-bit-allocating
              ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
                (nat (free-l Va t))
                (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
                  (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
                    p)
                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l

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Va t)),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
      t = freeing-node Va t)
prefer 2 apply force
apply auto[1]

apply(rule conjI)

apply clarify
apply(case-tac t = t1)
apply(subgoal-tac get-bit-s Va (pool n1) (level n1) (block n1) = FREE)
prefer 2
apply(subgoal-tac pool n1 = p ∧ level n1 = nat (free-l Va t) ∧ block n1 =
block-num
      (set-bit-allocating
      ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
prefer 2 apply auto[1]
apply(subgoal-tac block-num
      (set-bit-allocating
      ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)) = block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
prefer 2 apply(simp add: set-bit-def rmhead-free-list-def block-num-def)
apply(subgoal-tac block-num (rmhead-free-list (mem-pool-info Va p) (nat

```

```

(free-l Va t)))
                                (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
                                ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4
^ nat (free-l Va t))) =
                                block-num (mem-pool-info Va p)
                                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                                ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))))
    prefer 2 apply(simp add:set-bit-def block-num-def rmhead-free-list-def)
    apply(simp add:level-empty-def head-free-list-def)
    using inv-bitmap-freelist-fl-FREE[of Va p nat (free-l Va t) 0]
    apply (smt hd-conv-nth inv-mempool-info-maxsz-align4 le-trans length-greater-0-conv
linorder-not-less)
    apply(subgoal-tac get-bit-s Va (pool n2) (level n2) (block n2) = ALLOCATING)
    prefer 2 apply auto[1]
    apply auto[1]

    apply(case-tac t = t2)
    apply(subgoal-tac get-bit-s Va (pool n2) (level n2) (block n2) = FREE)
    prefer 2
    apply(subgoal-tac pool n2 = p ^ level n2 = nat (free-l Va t) ^ block n2 =
block-num
                                (set-bit-allocating
                                ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
                                (nat (free-l Va t))
                                (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
                                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                                (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))
                                p)
                                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                                (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
    prefer 2 apply auto[1]
    apply(subgoal-tac block-num
                                (set-bit-allocating
                                ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
                                (nat (free-l Va t))
                                (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
                                (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
                                (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
                                p)
                                p)

```

```

      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)) = block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
    prefer 2 apply(simp add: set-bit-def rmhead-free-list-def block-num-def)
    apply(subgoal-tac block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t))))
      (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
      ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4
^ nat (free-l Va t))) =
      block-num (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))))
    prefer 2 apply(simp add: set-bit-def block-num-def rmhead-free-list-def)
    apply(unfold level-empty-def head-free-list-def)[1]
    using inv-bitmap-freelist-fl-FREE[of Va p nat (free-l Va t) 0]
    apply (smt hd-conv-nth inv-mempool-info-maxsz-align4 le-trans length-greater-0-conv
linorder-not-less)
    apply(subgoal-tac get-bit-s Va (pool n1) (level n1) (block n1) = ALLOCATING)
    prefer 2 apply(subgoal-tac allocating-node Va t1 = Some n1) prefer 2 apply
auto[1]
    apply blast
    apply auto[1]

    apply(subgoal-tac allocating-node Va t1 = Some n1)
    prefer 2 apply auto[1]
    apply(subgoal-tac allocating-node Va t2 = Some n2)
    prefer 2 apply auto[1]
    apply auto[1]

    apply clarify
    apply(case-tac t = t1)
    apply(subgoal-tac get-bit-s Va (pool n1) (level n1) (block n1) = FREE)
    prefer 2
    apply(subgoal-tac pool n1 = p ^ level n1 = nat (free-l Va t) ^ block n1 =
block-num
      (set-bit-allocating
      ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t))))

```

```

      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)))

      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
    prefer 2 apply auto[1]
    apply(subgoal-tac block-num
      (set-bit-allocating
        ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p)
(nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)) = block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va
t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
    prefer 2 apply(simp add: set-bit-def rmhead-free-list-def block-num-def)
    apply(subgoal-tac block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
      ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4
^ nat (free-l Va t))) =
      block-num (mem-pool-info Va p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      ((ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))))
    prefer 2 apply(simp add:set-bit-def block-num-def rmhead-free-list-def)
    apply(simp add:level-empty-def head-free-list-def)
    using inv-bitmap-freelist-fl-FREE[of Va p nat (free-l Va t) 0]
    apply (smt hd-conv-nth inv-mempool-info-maxsz-align4 le-trans length-greater-0-conv
linorder-not-less)
    apply(subgoal-tac get-bit-s Va (pool n2) (level n2) (block n2) = FREEING)
    prefer 2 apply auto[1]
    apply auto[1]

    apply(subgoal-tac allocating-node Va t1 = Some n1)
    prefer 2 apply auto[1]

```

```

apply(subgoal-tac allocating-node Va t2 = Some n2)
prefer 2 apply auto[1]
apply auto[1]
done

lemma mp-alloc-stm3-lm2-inv-bitmap0:
inv-mempool-info Va  $\wedge$  inv-bitmap0 Va  $\implies$ 
  p  $\in$  mem-pools Va  $\implies$ 
    inv-bitmap0
      (Va( $\llbracket$ blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l Va t))))),
      mem-pool-info :=
        set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
          Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
          (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
            (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
              (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))),
          allocating-node := allocating-node Va(t  $\mapsto$ 
            ( $\llbracket$ pool = p, level = nat (free-l Va t),
              block = block-num
                (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
                  (mem-pool-info Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
                    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
                      Va t))) (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
                          Va t))))
                      p)
                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
              (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)),
              data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
apply(simp add:set-bit-def)
apply(rule subst[where s=inv-bitmap0
  (Va( $\llbracket$ mem-pool-info := (mem-pool-info Va)
    (p := rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t))
      ( $\llbracket$ levels := levels (rmhead-free-list (mem-pool-info Va p) (nat (free-l
        Va t))))
        [nat (free-l Va t) :=
          (levels (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
            ! nat (free-l Va t))
            ( $\llbracket$ bits := bits (levels (rmhead-free-list (mem-pool-info Va p) (nat
              (free-l Va t))) !
                nat (free-l Va t))
              [block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
                Va t)))
                (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
                  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
                    Va t))] :=
                    ALLOCATING[ $\llbracket$ ]]))
apply(simp add:inv-bitmap0-def)

```



```

apply(subgoal-tac length (levels (mem-pool-info Va p)) > 0) prefer 2
  apply(simp add:inv-def inv-mempool-info-def Let-def) apply fastforce

apply(subgoal-tac  $\forall i < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va\ p) ! 0)).$ 
  (bits (levels (mem-pool-info Va p) ! 0)) ! i  $\neq \text{NOEXIST}$ )
  prefer 2 apply(simp add:inv-def inv-bitmap0-def) apply metis

apply(case-tac nat (free-l Va t) = 0)
  apply(simp add:inv-bitmap0-def Let-def rmhead-free-list-def block-num-def)
  apply clarsimp
  apply(case-tac i = (head-free-list (mem-pool-info Va p) NULL – buf (mem-pool-info
Va p)) div
    ALIGN4 (max-sz (mem-pool-info Va p)))
    apply(subgoal-tac bits (levels (mem-pool-info Va p))
      [NULL := (levels (mem-pool-info Va p)
        [NULL := (levels (mem-pool-info Va p) ! NULL)
          (free-list := tl (free-list (levels (mem-pool-info Va p) !
NULL)))] !
        NULL)
        (bits := bits (levels (mem-pool-info Va p)
          [NULL := (levels (mem-pool-info Va p) ! NULL)
            (free-list := tl (free-list (levels (mem-pool-info Va p) !
NULL)))] !
          NULL)
          [(head-free-list (mem-pool-info Va p) NULL – buf (mem-pool-info
Va p)) div
            ALIGN4 (max-sz (mem-pool-info Va p)) :=
            ALLOCATING)] ! i = ALLOCATING) prefer 2
      apply(rule subst[where s=(bits (levels (mem-pool-info Va p) ! 0))
        [(head-free-list (mem-pool-info Va p) NULL – buf (mem-pool-info
Va p)) div
          ALIGN4 (max-sz (mem-pool-info Va p)) := ALLOCATING]])
        apply fastforce
        apply simp
        apply force

apply(subgoal-tac bits (levels (mem-pool-info Va p) ! NULL)
  [(head-free-list (mem-pool-info Va p) NULL – buf (mem-pool-info Va p))
div
  ALIGN4 (max-sz (mem-pool-info Va p)) := ALLOCATING] ! i  $\neq \text{NOEX-}$ 
IST) prefer 2
  apply force
  apply simp

apply(simp add:inv-bitmap0-def Let-def rmhead-free-list-def block-num-def)
done

```

**lemma** *mp-alloc-stm3-lm2-inv-bitmapn*:  
*inv-mempool-info Va*  $\wedge$  *inv-bitmapn Va*  $\implies$   
 $p \in \text{mem-pools } Va \implies$   
*inv-bitmapn*  
 $(Va \llbracket \text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$ ,  
 $\text{mem-pool-info} :=$   
 $\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \ (\text{nat } (\text{free-l } Va \ t))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t)))$ ,  
 $\text{allocating-node} := \text{allocating-node } Va(t \mapsto$   
 $\llbracket \text{pool} = p, \text{level} = \text{nat } (\text{free-l } Va \ t),$   
 $\text{block} = \text{block-num}$   
 $(\text{set-bit-allocating } ((\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ p \ (\text{nat } (\text{free-l } Va \ t))$   
 $(\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) \ (\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t)))$   
 $p)$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \ (\text{ALIGN}_4$   
 $(\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t)),$   
 $\text{data} = \text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))) \rrbracket \rrbracket \rrbracket$   
**apply**(*simp add:set-bit-def*)  
**apply**(*rule subst[where s=inv-bitmapn*  
 $(Va \llbracket \text{mem-pool-info} := (\text{mem-pool-info } Va)$   
 $(p := \text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))$   
 $\llbracket \text{levels} := \text{levels } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $\text{nat } (\text{free-l } Va \ t) :=$   
 $(\text{levels } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t))))$   
 $! \text{nat } (\text{free-l } Va \ t))$   
 $\llbracket \text{bits} := \text{bits } (\text{levels } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))) !$   
 $\text{nat } (\text{free-l } Va \ t))$   
 $[\text{block-num } (\text{rmhead-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{head-free-list } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)))$   
 $(\text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \ ^\wedge \ \text{nat } (\text{free-l } Va \ t)) :=$   
 $\text{ALLOCATING} \rrbracket \rrbracket \rrbracket \rrbracket \rrbracket \rrbracket$   
**apply**(*simp add:inv-bitmapn-def*)  
**apply**(*subgoal-tac length (levels (mem-pool-info Va p)) > 0*) **prefer 2**  
**apply**(*simp add:inv-def inv-mempool-info-def Let-def*) **apply fastforce**  
**apply**(*subgoal-tac  $\forall i < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) ! (\text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) - \text{Suc } 0)))$ .*  
 $(\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) ! (\text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) - \text{Suc } 0)))$

```

Va p)) - Suc 0))) ! i ≠ DIVIDED)
  prefer 2 apply(simp add:inv-def inv-bitmapn-def) apply metis

apply(case-tac nat (free-l Va t) = length (levels (mem-pool-info Va p)) - Suc 0)
  apply(simp add:inv-bitmapn-def Let-def rmhead-free-list-def block-num-def)
  apply clarsimp
  apply(case-tac i = (head-free-list (mem-pool-info Va p) (length (levels (mem-pool-info
Va p)) - Suc NULL) -
    buf (mem-pool-info Va p)) div
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ (length (levels (mem-pool-info
Va p)) - Suc NULL)))
  apply(subgoal-tac bits (levels (mem-pool-info Va p) ! (length (levels (mem-pool-info
Va p)) - Suc NULL))
    [(head-free-list (mem-pool-info Va p) (length (levels (mem-pool-info Va p))
- Suc NULL) -
      buf (mem-pool-info Va p)) div
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ (length (levels (mem-pool-info
Va p)) - Suc NULL)) :=
        ALLOCATING] ! i ≠ DIVIDED) prefer 2
  apply(rule subst[where s=(bits (levels (mem-pool-info Va p) ! 0))
    [(head-free-list (mem-pool-info Va p) NULL - buf (mem-pool-info
Va p)) div
      ALIGN4 (max-sz (mem-pool-info Va p)) := ALLOCATING]])
  apply fastforce
  apply simp
  apply force

  apply(subgoal-tac bits (levels (mem-pool-info Va p) ! NULL)
    [(head-free-list (mem-pool-info Va p) NULL - buf (mem-pool-info Va p))
div
  ALIGN4 (max-sz (mem-pool-info Va p)) := ALLOCATING] ! i ≠ DI-
VIDED) prefer 2
  apply force
  apply simp

apply(simp add:inv-bitmapn-def Let-def rmhead-free-list-def block-num-def)
done

lemma mp-alloc-stm3-lm2-inv-bitmap-not4free:
inv-mempool-info Va ∧ inv-bitmap-not4free Va ⇒
  p ∈ mem-pools Va ⇒
  inv-bitmap-not4free
    (Va⟦blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l Va
t))),
    mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p (nat (free-l Va t))

```

```

      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))),
    allocating-node := allocating-node Va(t ↦
      (pool = p, level = nat (free-l Va t),
        block = block-num
          (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
          (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))) (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
          (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
            p)
          (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
apply(rule subst[where s=inv-bitmap-not4free (Va(mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p :=
        rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) p (nat
(free-l Va t))
        (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
        (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))))))
apply(simp add: inv-bitmap-not4free-def Let-def partner-bits-def set-bit-def rmhead-free-list-def
block-num-def)

```

```

apply(simp add: inv-bitmap-not4free-def Let-def partner-bits-def set-bit-def rmhead-free-list-def
block-num-def)

```

```

apply clarsimp

```

```

apply(case-tac nat (free-l Va t) = i) prefer 2 apply auto[1]

```

```

apply(subgoal-tac bits (levels (mem-pool-info Va p)
[nat (free-l Va t) :=
  (levels (mem-pool-info Va p)
[nat (free-l Va t) := (levels (mem-pool-info Va p) ! nat (free-l Va
t))
  (free-list := tl (free-list (levels (mem-pool-info Va p) ! nat (free-l
Va t))))]) !
  nat (free-l Va t))
  (bits := bits (levels (mem-pool-info Va p)
[nat (free-l Va t) := (levels (mem-pool-info Va p) ! nat
(free-l Va t))
  (free-list := tl (free-list (levels (mem-pool-info Va
p) ! nat (free-l Va t))))]) !
  nat (free-l Va t))
  [(head-free-list (mem-pool-info Va p) (nat (free-l Va t)) - buf
(mem-pool-info Va p)) div

```

$(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) :=$   
 $ALLOCATING[] !$   
 $i = bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ p)\ !\ i)\ [(head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) - buf\ (mem\text{-}pool\text{-}info\ Va\ p))\ div$   
 $(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) :=$   
 $ALLOCATING[]\ prefer\ 2\ apply\ simp$   
**apply** *simp*  
**apply**  $(case\text{-}tac\ (head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) - buf\ (mem\text{-}pool\text{-}info\ Va\ p))\ div$   
 $(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) = j\ div\ 4\ * 4)$   
**apply** *auto*[1]  
**apply**  $(case\text{-}tac\ (head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) - buf\ (mem\text{-}pool\text{-}info\ Va\ p))\ div$   
 $(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) = Suc\ (j\ div\ 4\ * 4))$   
**apply**  $(subgoal\text{-}tac\ Suc\ (j\ div\ 4\ * 4) < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ p)\ !\ i)))\ prefer\ 2$   
**apply**  $(metis\ list\text{-}update\text{-}beyond\ not\text{-}less)$   
**apply** *auto*[1]  
**apply**  $(case\text{-}tac\ (head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) - buf\ (mem\text{-}pool\text{-}info\ Va\ p))\ div$   
 $(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) = j\ div\ 4\ * 4 + 2)$   
**apply**  $(subgoal\text{-}tac\ j\ div\ 4\ * 4 + 2 < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ p)\ !\ i)))\ prefer\ 2$   
**apply**  $(metis\ list\text{-}update\text{-}beyond\ not\text{-}less)$   
**apply** *auto*[1]  
**apply**  $(case\text{-}tac\ (head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) - buf\ (mem\text{-}pool\text{-}info\ Va\ p))\ div$   
 $(ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t)) = j\ div\ 4\ * 4 + 3)$   
**apply**  $(subgoal\text{-}tac\ j\ div\ 4\ * 4 + 3 < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ p)\ !\ i)))\ prefer\ 2$   
**apply**  $(metis\ list\text{-}update\text{-}beyond\ not\text{-}less)$   
**apply** *auto*[1]  
**apply** *simp*  
**done**

**lemma** *mp-alloc-stm3-lm2-inv-mempool-info:*

*inv-mempool-info*  $Va \wedge$

$p \in mem\text{-}pools\ Va \implies$

$\forall ii < length\ (lsizes\ Va\ t). lsizes\ Va\ t\ !\ ii = ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ ii \implies$

$length\ (lsizes\ Va\ t) \leq n\text{-}levels\ (mem\text{-}pool\text{-}info\ Va\ p) \implies$

$\neg free\text{-}l\ Va\ t < OK \implies$

```

nat (free-l Va t) < length (lsizes Va t) ==>
inv-mempool-info
(Va (blk := (blk Va) (t := head-free-list (mem-pool-info Va p) (nat (free-l Va t))),
  mem-pool-info :=
    set-bit-allocating ((mem-pool-info Va) (p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))),
  allocating-node := allocating-node Va (t ↦
    (pool = p, level = nat (free-l Va t),
    block = block-num
      (set-bit-allocating ((mem-pool-info Va) (p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))) (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))))
    p)
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (ALIGN4
(max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)),
    data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
apply (simp add: inv-mempool-info-def)
apply (simp add: rmhead-free-list-def
  head-free-list-def set-bit-def block-num-def)
apply (rule conjI) apply metis
apply (rule conjI) apply metis
apply (rule conjI) apply metis
apply (rule conjI) apply metis
apply clarsimp apply (simp add: Let-def)
apply (case-tac nat (free-l Va t) = i)
apply (subgoal-tac length (bits (levels (mem-pool-info Va p) ! (nat (free-l Va
t)))))
      = n-max (mem-pool-info Va p) * 4 ^ (nat (free-l Va t)))
prefer 2 apply metis
using mp-alloc-stm3-lm2-2 [where ii = nat (free-l Va t) and mp = mem-pool-info
Va p and
  fl = tl (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))) and
  jj = (hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))) - buf
(mem-pool-info Va p)) div
  lsizes Va t ! nat (free-l Va t)] apply metis
apply simp
done

```

**lemma** mp-alloc-stm3-lm2-inv-bitmap-freelist:  
 $\neg \text{level-empty } (\text{mem-pool-info } Va \ p) \ (\text{nat } (\text{free-l } Va \ t)) \implies$   
 $\text{inv-bitmap-freelist } Va \wedge \text{inv-mempool-info } Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va$

$p)) \text{ div } 4 \wedge ii \implies$   
 $\text{length } (lsizes \text{ Va } t) \leq n\text{-levels } (mem\text{-pool-info } \text{Va } p) \implies$   
 $\text{alloc-l } \text{Va } t < \text{int } (n\text{-levels } (mem\text{-pool-info } \text{Va } p)) \implies$   
 $\text{free-l } \text{Va } t \leq \text{alloc-l } \text{Va } t \implies$   
 $\neg \text{free-l } \text{Va } t < OK \implies$   
 $\text{length } (lsizes \text{ Va } t) \leq \text{length } (levels (mem\text{-pool-info } \text{Va } p)) \implies$   
 $\text{nat } (\text{free-l } \text{Va } t) < \text{length } (lsizes \text{ Va } t) \implies$   
 $\text{inv-bitmap-freelist}$   
 $(\text{Va}(\text{blk} := (\text{blk } \text{Va})(t := \text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))),$   
 $\text{mem-pool-info} :=$   
 $\text{set-bit-allocating } ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p (\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } \text{Va } t))),$   
 $\text{allocating-node} := \text{allocating-node } \text{Va}(t \mapsto$   
 $(\text{pool} = p, \text{level} = \text{nat } (\text{free-l } \text{Va } t),$   
 $\text{block} = \text{block-num}$   
 $(\text{set-bit-allocating } ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p (\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) (\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } \text{Va } t)))$   
 $\text{Va } t)))$   
 $p)$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))$   
 $(ALIGN_4 (\text{max-sz } (mem\text{-pool-info } \text{Va } p)) \text{ div } 4 \wedge \text{nat } (\text{free-l } \text{Va } t)),$   
 $\text{data} = \text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))))$   
**apply**(*simp add:inv-bitmap-freelist-def*)  
**apply** *clarify*  
**apply**(*case-tac pa  $\neq$  p*) **apply**(*simp add:Let-def*)  
**using** *mp-alloc-stm3-body-meminfo* **apply** *smt*  
**apply**(*simp add:Let-def*)  
**apply**(*rule subst[where  $t = \text{length } (levels (set-bit-allocating ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p$*   
 $(\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $\text{Va } t)))$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))) (lsizes$   
 $\text{Va } t ! \text{nat } (\text{free-l } \text{Va } t)))$   
 $p))$  **and**  $s = \text{length } (levels (mem\text{-pool-info } \text{Va } p)))$   
**using** *mp-alloc-stm3-body-len-lvls* **apply** *metis*  
**apply**(*rule subst[where  $t = \text{buf } (set-bit-allocating ((mem\text{-pool-info } \text{Va})(p := \text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t)))) p$*   
 $(\text{nat } (\text{free-l } \text{Va } t))$   
 $(\text{block-num } (\text{rmhead-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))))$   
 $(\text{head-free-list } (mem\text{-pool-info } \text{Va } p) (\text{nat } (\text{free-l } \text{Va } t))) (lsizes \text{ Va } t$   
 $! \text{nat } (\text{free-l } \text{Va } t)))$

```

      p) and s=buf (mem-pool-info Va p)])
    using mp-alloc-stm3-body-minf-buf apply metis
  apply(rule subst[where t=n-max (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t
! nat (free-l Va t)))
      p) and s=n-max (mem-pool-info Va p)])
    using mp-alloc-stm3-body-minf-nmax apply metis
  apply(rule subst[where t=max-sz (set-bit-allocating ((mem-pool-info Va)(p :=
rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t
! nat (free-l Va t)))
      p) and s=max-sz (mem-pool-info Va p)])
    using mp-alloc-stm3-body-minf-maxsz apply metis

  apply clarify apply(rename-tac pa ii)
  apply(subgoal-tac length (bits (levels (set-bit-allocating
      ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (lsizes Va t ! nat (free-l Va t)))
      p ! ii)=length (bits ((levels (mem-pool-info Va p))!ii)))
    prefer 2 using mp-alloc-stm3-body-len-bits apply metis

  apply(rule conjI)
  apply clarify apply(rule iffI) apply(rename-tac pa ii jj)

  apply(case-tac nat (free-l Va t) = ii)

  apply(case-tac jj = (block-num (rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes
Va t ! nat (free-l Va t))))

  apply(subgoal-tac get-bit (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))) (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (lsizes Va t ! nat (free-l Va t))))
      p ii jj = ALLOCATING)
    prefer 2 apply(simp add: set-bit-def rmhead-free-list-def)
    apply (metis BlockState.distinct(17))

```



**apply**(*subgoal-tac* *get-bit* (*mem-pool-info* *Va*) *p ii jj = FREE*)  
**prefer 2 apply**(*simp add: set-bit-def rmhead-free-list-def*)  
**apply**(*subgoal-tac* *buf* (*mem-pool-info* *Va p*) + *jj \* (max-sz (mem-pool-info* *Va*  
*p) div 4 ^ ii*)  
 $\in \text{set (free-list (levels (mem-pool-info Va p) ! ii))}$   
**prefer 2 apply** (*metis mp-alloc-stm3-body-len-lvls*)  
**apply**(*subgoal-tac* *buf* (*mem-pool-info* *Va p*) + *jj \* (max-sz (mem-pool-info* *Va*  
*p) div 4 ^ ii*)  
 $\neq \text{head-free-list (mem-pool-info Va p) (nat (free-l Va t))}$   
**prefer 2 apply**(*simp add:block-num-def*) **using** *mp-alloc-stm3-lm2-inv-1*  
**apply** *simp*  
**apply**(*simp add: set-bit-def rmhead-free-list-def head-free-list-def*)  
**using** *list-nhd-in-tl-set* **apply** *metis*

**apply**(*subgoal-tac* *get-bit* (*mem-pool-info* *Va*) *p ii jj = FREE*)  
**prefer 2 apply**(*simp add: set-bit-def rmhead-free-list-def*)  
**apply**(*subgoal-tac* *buf* (*mem-pool-info* *Va p*) + *jj \* (max-sz (mem-pool-info* *Va*  
*p) div 4 ^ ii*)  
 $\in \text{set (free-list (levels (mem-pool-info Va p) ! ii))}$   
**prefer 2 apply** (*metis mp-alloc-stm3-body-len-lvls*)  
**apply**(*simp add: set-bit-def rmhead-free-list-def head-free-list-def*)

**apply**(*rename-tac* *pa ii jj*)  
**apply**(*subgoal-tac* *length* (*levels* (*set-bit-allocating* ((*mem-pool-info* *Va*)(*p := rmhead-free-list*  
(*mem-pool-info* *Va p*) (*nat (free-l Va t)*)))) *p*  
(*nat (free-l Va t)*)  
(*block-num* (*rmhead-free-list* (*mem-pool-info* *Va p*) (*nat (free-l*  
*Va t)*)))  
(*head-free-list* (*mem-pool-info* *Va p*) (*nat (free-l Va t)*))) (*lsizes*  
*Va t ! nat (free-l Va t)*))  
*p)) = length (levels (mem-pool-info Va p))*)  
**prefer 2 using** *mp-alloc-stm3-body-len-lvls* **apply** *metis*  
**apply**(*case-tac* *nat (free-l Va t) = ii*)

**apply**(*subgoal-tac* *buf* (*mem-pool-info* *Va p*) + *jj \* (max-sz (mem-pool-info* *Va*  
*p) div 4 ^ ii*)  
 $\in \text{set (tl (free-list (levels (mem-pool-info Va p) ! ii))}$   
**prefer 2 using** *mp-alloc-stm3-body-minf-buf mp-alloc-stm3-body-minf-maxsz*  
*mp-alloc-stm3-body-frlst-samelvl* **apply** *metis*  
**apply**(*subgoal-tac* *buf* (*mem-pool-info* *Va p*) + *jj \* (max-sz (mem-pool-info* *Va*  
*p) div 4 ^ ii*)  
 $\in \text{set (free-list (levels (mem-pool-info Va p) ! ii))}$   
**prefer 2 apply**(*metis list.set-sel(2) tl-Nil*)  
**apply**(*subgoal-tac* *get-bit* (*mem-pool-info* *Va*) *p ii jj = FREE*)  
**prefer 2 apply** *metis*  
**apply**(*subgoal-tac* *block-num* (*rmhead-free-list* (*mem-pool-info* *Va p*) (*nat (free-l*

```

Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t)) ≠ jj)
  prefer 2
  apply(subgoal-tac buf (mem-pool-info Va p) + jj * (max-sz (mem-pool-info
Va p) div 4 ^ ii)
      ≠ head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  prefer 2 apply(subgoal-tac distinct (free-list (levels (mem-pool-info Va p)
! ii)))
  prefer 2 apply metis
  apply(simp add:head-free-list-def)
  using dist-hd-nin-tl apply (metis (mono-tags, hide-lams) le-eq-less-or-eq
le-trans linorder-not-less)
  apply(simp add:block-num-def)
  apply(subgoal-tac buf (rmhead-free-list (mem-pool-info Va p) ii) = buf
(mem-pool-info Va p))
  prefer 2 apply(simp add:rmhead-free-list-def)
  apply(subgoal-tac ∃ n. head-free-list (mem-pool-info Va p) ii =
      buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ ii))
  prefer 2 apply(simp add:head-free-list-def level-empty-def)
  apply (smt add-lessD1 hd-conv-nth le-eq-less-or-eq length-greater-0-conv
less-imp-add-positive)
  using mp-alloc-stm3-lm2-inv-2 apply (metis inv-mempool-info-maxsz-align4)
  apply(simp add: set-bit-def rmhead-free-list-def head-free-list-def)

  apply(subgoal-tac buf (mem-pool-info Va p) + jj * (max-sz (mem-pool-info Va
p) div 4 ^ ii)
      ∈ set (free-list (levels (mem-pool-info Va p) ! ii)))
  prefer 2 apply (metis mp-alloc-stm3-body-frlst-otherlvl mp-alloc-stm3-body-minf-buf
mp-alloc-stm3-body-minf-maxsz)
  apply(subgoal-tac get-bit (mem-pool-info Va) p ii jj = FREE)
  prefer 2 apply(simp add: set-bit-def rmhead-free-list-def)
  apply(simp add: set-bit-def rmhead-free-list-def head-free-list-def)

apply(rule conjI)

  apply clarify
  apply(rename-tac pa ii jj)
  apply(subgoal-tac length (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes
Va t ! nat (free-l Va t)))
      p)) = length (levels (mem-pool-info Va p)))

```

```

prefer 2 using mp-alloc-stm3-body-len-lvls apply metis
apply(case-tac nat (free-l Va t) = ii)

apply(subgoal-tac (free-list
  (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))))) p
  (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))))
  p) ! ii)) = (tl (free-list (levels (mem-pool-info Va p) ! ii))))
prefer 2 apply(simp add:level-empty-def set-bit-def rmhead-free-list-def head-free-list-def)

apply(subgoal-tac tl (free-list (levels (mem-pool-info Va p) ! ii)) ! jj = (free-list
(levels (mem-pool-info Va p) ! ii)) ! Suc jj)
prefer 2 apply(rule List.nth-tl)
apply(subgoal-tac length (tl (free-list (levels (mem-pool-info Va p) ! ii))) =
length (free-list
  (levels (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))))) p
  (nat (free-l Va t))
  (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t))))
  (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
  (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t))))
  p) ! ii)))
prefer 2 apply simp
apply metis
apply(subgoal-tac ( $\exists n. n < n\text{-max} \text{ (mem-pool-info Va p) } * (4 \wedge ii) \wedge$  (free-list
(levels (mem-pool-info Va p) ! ii)) ! Suc jj =
  buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4
^ ii)))

using mp-alloc-stm3-body-minf-buf mp-alloc-stm3-body-minf-nmax mp-alloc-stm3-body-minf-maxsz
apply metis
apply(subgoal-tac Suc jj < length (free-list (levels (mem-pool-info Va p) !
ii)))
prefer 2 apply(subgoal-tac jj < length (tl (free-list (levels (mem-pool-info
Va p) ! ii))))
prefer 2 apply metis
apply(simp add:level-empty-def)
apply metis

using mp-alloc-stm3-body-minf-buf mp-alloc-stm3-body-minf-maxsz mp-alloc-stm3-body-minf-nmax

```

*mp-alloc-stm3-body-frlst-otherlvl* **apply** *metis*

**apply**(*case-tac* *nat* (*free-l* *Va* *t*) = *ii*)  
**apply**(*subgoal-tac* (*free-list*  
(*levels* (*set-bit-allocating* ((*mem-pool-info* *Va*)(*p* := *rmhead-free-list*  
(*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))))) *p*  
(*nat* (*free-l* *Va* *t*))  
(*block-num* (*rmhead-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l*  
*Va* *t*)))  
(*head-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))  
(*ALIGN4* (*max-sz* (*mem-pool-info* *Va* *p*)) *div* 4 ^ *nat* (*free-l*  
*Va* *t*)))  
(*p* ! *ii*)) = (*tl* (*free-list* (*levels* (*mem-pool-info* *Va* *p*) ! *ii*))))  
**prefer** 2 **apply**(*simp* *add:level-empty-def* *set-bit-def* *rmhead-free-list-def* *head-free-list-def*)  
**apply**(*subgoal-tac* *distinct* (*free-list* (*levels* (*mem-pool-info* *Va* *p*) ! *ii*)))  
**prefer** 2 **apply** *simp*  
**using** *distinct-tl* **apply** *metis*  
  
**apply**(*subgoal-tac* *distinct* (*free-list* (*levels* (*mem-pool-info* *Va* *p*) ! *ii*)))  
**prefer** 2 **apply** (*metis* *mp-alloc-stm3-body-len-lvls*)  
**using** *mp-alloc-stm3-body-frlst-otherlvl* **apply** *metis*  
**done**

**lemma** *mp-alloc-stm3-lm2-inv-bitmap*:

$\neg$  *level-empty* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*))  $\implies$   
*inv-mempool-info* *Va*  $\wedge$  *inv-bitmap-freelist* *Va*  $\wedge$  *inv-bitmap* *Va*  $\implies$   
*p*  $\in$  *mem-pools* *Va*  $\implies$   
*length* (*lsizes* *Va* *t*)  $\leq$  *n-levels* (*mem-pool-info* *Va* *p*)  $\implies$   
*alloc-l* *Va* *t* < *int* (*n-levels* (*mem-pool-info* *Va* *p*))  $\implies$   
*free-l* *Va* *t*  $\leq$  *alloc-l* *Va* *t*  $\implies$   
 $\neg$  *free-l* *Va* *t* < *OK*  $\implies$   
*length* (*lsizes* *Va* *t*)  $\leq$  *length* (*levels* (*mem-pool-info* *Va* *p*))  $\implies$   
*nat* (*free-l* *Va* *t*) < *length* (*lsizes* *Va* *t*)  $\implies$   
*inv-bitmap*  
(*Va* ! *blk* := (*blk* *Va*)(*t* := *head-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*))),  
*mem-pool-info* :=  
*set-bit-allocating* ((*mem-pool-info* *Va*)(*p* := *rmhead-free-list* (*mem-pool-info*  
*Va* *p*) (*nat* (*free-l* *Va* *t*)))) *p* (*nat* (*free-l* *Va* *t*))  
(*block-num* (*rmhead-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))  
(*head-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))  
(*ALIGN4* (*max-sz* (*mem-pool-info* *Va* *p*)) *div* 4 ^ *nat* (*free-l* *Va* *t*))),  
*allocating-node* := *allocating-node* *Va*(*t*  $\mapsto$   
(*pool* = *p*, *level* = *nat* (*free-l* *Va* *t*),  
*block* = *block-num*  
(*set-bit-allocating* ((*mem-pool-info* *Va*)(*p* := *rmhead-free-list*  
(*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))) *p* (*nat* (*free-l* *Va* *t*))  
(*block-num* (*rmhead-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l*  
*Va* *t*))) (*head-free-list* (*mem-pool-info* *Va* *p*) (*nat* (*free-l* *Va* *t*)))

```

    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l
Va t)))
    p)
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (ALIGN4
(max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)),
    data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
apply(subgoal-tac inv-bitmap (set-bit-s Va p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)))
    ALLOCATING))
prefer 2
apply(subgoal-tac get-bit-s Va p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat
(free-l Va t))) = FREE)
prefer 2 apply(simp add:level-empty-def)
apply(subgoal-tac (block-num (mem-pool-info Va p) (free-list (levels (mem-pool-info
Va p) ! nat (free-l Va t)) ! NULL)
    (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t)))
= (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t))))
using inv-bitmap-freelist-ft-FREE[of Va p nat (free-l Va t) 0] apply simp
apply(simp add:block-num-def rmhead-free-list-def head-free-list-def)
apply (simp add:hd-conv-nth inv-mempool-info-maxsz-align4)
using inv-bitmap-presv-setbit[of Va p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va
t))) ALLOCATING set-bit-s Va p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)))
    ALLOCATING] apply simp

apply(rule inv-bitmap-presv-mpls-mpi2[of (set-bit-s Va p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
(head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (free-l Va t)))
    ALLOCATING) (Va\blk := (blk Va)(t := head-free-list (mem-pool-info Va
p) (nat (free-l Va t))),
    mem-pool-info :=
    set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
    (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
    (head-free-list (mem-pool-info Va p) (nat (free-l Va t))))

```

$(ALIGN_4 (max\text{-}sz (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l\ Va\ t))),$   
 $allocating\text{-}node := allocating\text{-}node\ Va(t \mapsto$   
 $(pool = p, level = nat\ (free\text{-}l\ Va\ t),$   
 $block = block\text{-}num$   
 $(set\text{-}bit\text{-}allocating\ ((mem\text{-}pool\text{-}info\ Va)(p := rmhead\text{-}free\text{-}list$   
 $(mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t))))\ p$   
 $(nat\ (free\text{-}l\ Va\ t))$   
 $(block\text{-}num\ (rmhead\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l$   
 $Va\ t))))$   
 $(head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)))$   
 $(ALIGN_4 (max\text{-}sz (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l$   
 $Va\ t))))$   
 $p)$   
 $(head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)))$   
 $(ALIGN_4 (max\text{-}sz (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ nat\ (free\text{-}l$   
 $Va\ t))),$   
 $data = head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t))))))$   
**apply**(simp add: set-bit-s-def set-bit-def block-num-def rmhead-free-list-def head-free-list-def)  
**apply**(simp add: set-bit-s-def set-bit-def block-num-def rmhead-free-list-def head-free-list-def)  
**apply** clarsimp **apply**(simp add: set-bit-s-def set-bit-def block-num-def rmhead-free-list-def  
head-free-list-def)  
**apply** (smt Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective  
Mem-pool-lvl.update-convs(2)  
linorder-not-less list-update-beyond nth-list-update-eq nth-list-update-neq)  
**by** simp

**lemma** mp-alloc-stm3-lm2-inv:

$(*NULL < head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) \implies *)$   
 $\neg level\text{-}empty\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va\ t)) \implies$   
 $inv\ Va \implies$   
 $allocating\text{-}node\ Va\ t = None \implies$   
 $freeing\text{-}node\ Va\ t = None \implies$   
 $p \in mem\text{-}pools\ Va \implies$   
 $ETIMEOUT \leq timeout \implies$   
 $timeout = ETIMEOUT \longrightarrow tmout\ Va\ t = ETIMEOUT \implies$   
 $\neg rf\ Va\ t \implies$   
 $\forall ii < length\ (lsizes\ Va\ t). lsizes\ Va\ t\ !\ ii = ALIGN_4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va$   
 $p))\ div\ 4\ \wedge\ ii \implies$   
 $length\ (lsizes\ Va\ t) \leq n\text{-}levels\ (mem\text{-}pool\text{-}info\ Va\ p) \implies$   
 $alloc\text{-}l\ Va\ t < int\ (n\text{-}levels\ (mem\text{-}pool\text{-}info\ Va\ p)) \implies$   
 $free\text{-}l\ Va\ t \leq alloc\text{-}l\ Va\ t \implies$   
 $\neg free\text{-}l\ Va\ t < 0 \implies$   
 $alloc\text{-}l\ Va\ t = int\ (length\ (lsizes\ Va\ t)) - 1 \wedge length\ (lsizes\ Va\ t) = n\text{-}levels$   
 $(mem\text{-}pool\text{-}info\ Va\ p) \vee$   
 $alloc\text{-}l\ Va\ t = int\ (length\ (lsizes\ Va\ t)) - 2 \wedge lsizes\ Va\ t\ !\ nat\ (alloc\text{-}l\ Va\ t +$   
 $1) < sz \implies$   
 $inv\ (Va(blk := (blk\ Va)(t := head\text{-}free\text{-}list\ (mem\text{-}pool\text{-}info\ Va\ p)\ (nat\ (free\text{-}l\ Va$   
 $t))))),$

```

      mem-pool-info :=
      set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes Va t !
nat (free-l Va t))),
      allocating-node := allocating-node Va(t ↦
      (pool = p, level = nat (free-l Va t),
      block = block-num
      (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p
      (nat (free-l Va t))
      (block-num (rmhead-free-list (mem-pool-info Va p) (nat (free-l
Va t)))
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t)))
      (lsizes Va t ! nat (free-l Va t)))
      p)
      (head-free-list (mem-pool-info Va p) (nat (free-l Va t))) (lsizes
Va t ! nat (free-l Va t))),
      data = head-free-list (mem-pool-info Va p) (nat (free-l Va t))))))
apply(subgoal-tac nat (free-l Va t) < length (levels (mem-pool-info Va p)))
prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def)
apply(subgoal-tac length (lsizes Va t) ≤ length (levels (mem-pool-info Va p)))
prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def)
apply(subgoal-tac nat (free-l Va t) < length (lsizes Va t))
prefer 2 apply linarith
apply(simp add:inv-def)
apply(rule conjI)
apply(simp add:inv-cur-def)
apply(rule conjI)
using mp-alloc-stm3-lm2-inv-thd-waitq apply fast
apply(rule conjI)
using mp-alloc-stm3-lm2-inv-mempool-info apply fast
apply(rule conjI)
using mp-alloc-stm3-lm2-inv-bitmap-freelist apply fast
apply(rule conjI) using mp-alloc-stm3-lm2-inv-bitmap apply simp
apply(rule conjI) using mp-alloc-stm3-lm2-inv-aux-vars apply simp
apply(rule conjI) using mp-alloc-stm3-lm2-inv-bitmap0 apply simp
apply(rule conjI) using mp-alloc-stm3-lm2-inv-bitmapn apply simp
using mp-alloc-stm3-lm2-inv-bitmap-not4free apply simp
done

```

**lemma** mp-alloc-stm3-lm2-3-1:  
 $(a::nat) \bmod b = 0 \implies c * b * (a \text{ div } b) = c * a$  **by** auto

**lemma** mp-alloc-stm3-lm2-3:

```

¬ level-empty (mem-pool-info Va p) (nat (free-l Va t)) ⇒
  inv Va ⇒
  alloc-l Va t < int (n-levels (mem-pool-info Va p)) ⇒
  free-l Va t ≤ alloc-l Va t ⇒
  p ∈ mem-pools Va ⇒
  ¬ free-l Va t < 0 ⇒
  max-sz (mem-pool-info Va p) = ALIGN4 (max-sz (mem-pool-info Va p)) ⇒
  let fl = hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))); mp =
mem-pool-info Va p
  in ∃ n < n-max mp * 4 ^ nat (free-l Va t). fl = buf mp + n * (max-sz mp div 4
^ nat (free-l Va t)) ⇒
  hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t)))
  < buf (mem-pool-info Va p) + n-max (mem-pool-info Va p) * max-sz (mem-pool-info
Va p)
apply(subgoal-tac nat (free-l Va t) < length (levels (mem-pool-info Va p)))
prefer 2 apply(simp add: inv-def inv-mempool-info-def Let-def)
apply(subgoal-tac hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))) ≥
buf (mem-pool-info Va p))
prefer 2 apply(simp add: inv-def) using inv-buf-le-fl[of Va p nat (free-l Va t)
0]
apply (simp add: hd-conv-nth level-empty-def)
apply (simp add: hd-conv-nth level-empty-def Let-def)
apply clarify

apply(subgoal-tac max-sz (mem-pool-info Va p) mod (4 ^ nat (free-l Va t)) = 0)
prefer 2 apply (metis ge-pow-mod-0 inv-mempool-info-def inv-def)
apply(subgoal-tac n * (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))
< n-max (mem-pool-info Va p) * max-sz (mem-pool-info Va p))
prefer 2 apply(subgoal-tac n-max (mem-pool-info Va p) * 4 ^ nat (free-l Va t)
* (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))
= n-max (mem-pool-info Va p) * max-sz (mem-pool-info Va p))
prefer 2
using mp-alloc-stm3-lm2-3-1[of max-sz (mem-pool-info Va p) 4 ^ nat (free-l
Va t) n-max (mem-pool-info Va p)] apply auto[1]
apply(subgoal-tac n * (max-sz (mem-pool-info Va p) div 4 ^ nat (free-l Va t))
< (n-max (mem-pool-info Va p) * 4 ^ nat (free-l Va t)) * (max-sz
(mem-pool-info Va p) div 4 ^ nat (free-l Va t)))
prefer 2 apply (metis inv-mempool-info-def inv-def mp-alloc-stm3-lm2-inv-1-2
mult-less-mono1)
apply linarith

apply simp
done

```

**lemma** mp-alloc-stm3-lm2-5:

```

¬ level-empty (mem-pool-info Va p) (nat (free-l Va t)) ⇒
  inv Va ⇒
  alloc-l Va t < int (n-levels (mem-pool-info Va p)) ⇒

```



$free-l\ Va\ t \leq alloc-l\ Va\ t \implies$   
 $p \in mem-pools\ Va \implies$   
 $\neg free-l\ Va\ t < 0 \implies$   
 $(hd\ (free-list\ (levels\ (mem-pool-info\ Va\ p)\ !\ nat\ (free-l\ Va\ t))) - buf\ (mem-pool-info\ Va\ p))\ div$   
 $(max-sz\ (mem-pool-info\ Va\ p)\ div\ 4\ \wedge\ nat\ (free-l\ Va\ t))$   
 $< n-max\ (mem-pool-info\ Va\ p)\ * 4\ \wedge\ nat\ (free-l\ Va\ t)$   
**apply**(subgoal-tac nat (free-l Va t) < length (levels (mem-pool-info Va p)))  
**prefer 2 apply**(simp add:inv-def inv-mempool-info-def Let-def)  
**apply**(subgoal-tac hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t)))  $\geq$   
buf (mem-pool-info Va p))  
**prefer 2 apply**(simp add: inv-def) **using** inv-buf-le-fl[of Va p nat (free-l Va t)  
0]  
**apply** (simp add: hd-conv-nth level-empty-def)  
**apply** (simp add: hd-conv-nth level-empty-def)  
**by** (metis block-num-def inv-bitmap-freelist-fl-bnum-in inv-mempool-info-def length-greater-0-conv  
inv-def)

**lemma** mp-alloc-stm3-lm2-4:

$inv\ Va \wedge$   
 $p \in mem-pools\ Va \wedge$   
 $free-list\ (levels\ (mem-pool-info\ Va\ p)\ !\ nat\ (free-l\ Va\ t)) \neq [] \implies$   
 $nat\ (free-l\ Va\ t) < length\ (levels\ (mem-pool-info\ Va\ p)) \implies NULL < hd\ (free-list$   
 $(levels\ (mem-pool-info\ Va\ p)\ !\ nat\ (free-l\ Va\ t)))$   
**using** inv-imp-fl-lt0 **apply**(simp add:Let-def)  
**by** (simp add: hd-conv-nth)

**lemma** mp-alloc-stm3-lm2:

$Va \in mp-alloc-precond1-70-2-2\ t\ p\ sz\ timeout \cap \{\cur = Some\ t\} \implies$   
 $(*head-free-list\ (mem-pool-info\ Va\ p)\ (nat\ (free-l\ Va\ t)) \neq NULL \implies*)$   
 $\neg level-empty\ (mem-pool-info\ Va\ p)\ (nat\ (free-l\ Va\ t)) \implies$   
 $(*\{Va \langle blk := (blk\ Va)(t := NULL) \rangle, Va \langle blk := (blk\ Va)(t := head-free-list$   
 $(mem-pool-info\ Va\ p)\ (nat\ (free-l\ Va\ t)) \rangle\},$   
 $mem-pool-info := (mem-pool-info\ Va)(p := rmhead-free-list\ (mem-pool-info$   
 $Va\ p)\ (nat\ (free-l\ Va\ t)) \rangle \rangle \cap$   
 $\{NULL < \cur\ t\} =$   
 $\{Va \langle blk := (blk\ Va)(t := head-free-list\ (mem-pool-info\ Va\ p)\ (nat\ (free-l\ Va$   
 $t)) \rangle\},$   
 $mem-pool-info := (mem-pool-info\ Va)(p := rmhead-free-list\ (mem-pool-info$   
 $Va\ p)\ (nat\ (free-l\ Va\ t)) \rangle \rangle \implies*)$   
 $\{let\ vb = Va \langle blk := (blk\ Va)(t := head-free-list\ (mem-pool-info\ Va\ p)\ (nat$   
 $(free-l\ Va\ t)) \rangle\},$   
 $mem-pool-info := (mem-pool-info\ Va)(p := rmhead-free-list$   
 $(mem-pool-info\ Va\ p)\ (nat\ (free-l\ Va\ t)) \rangle \rangle$   
 $in\ \cur (=)\ (vb \langle mem-pool-info :=$   
 $set-bit-allocating\ (mem-pool-info\ vb)\ p\ (nat\ (free-l\ vb\ t))$   
 $(block-num\ (mem-pool-info\ vb\ p)\ (blk\ vb\ t)\ (lsizes\ vb\ t\ !\ nat\ (free-l$   
 $vb\ t)) \rangle \rangle) \wedge$

```

    ¬ level-empty (mem-pool-info Va p) (nat (free-l Va t)))}
  ⊆ { (allocating-node-update
    (λ-. 'allocating-node(t ↦
      (pool = p, level = nat ('free-l t), block = block-num ('mem-pool-info
p) ('blk t) ('lsizes t ! nat ('free-l t)),
      data = 'blk t))))
    ∈ { (Pair Va) ∈ Mem-pool-alloc-guar t} ∩ mp-alloc-precond2-1 t p sz
timeout}
apply(subgoal-tac head-free-list (mem-pool-info Va p) (nat (free-l Va t)) ≠ NULL)
prefer 2
apply(subgoal-tac (nat (free-l Va t)) < length (levels (mem-pool-info Va p)))
prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def) apply force
apply(simp add:head-free-list-def level-empty-def) using mp-alloc-stm3-lm2-4
apply simp

apply clarsimp
apply(rule conjI)
apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def)
apply(rule disjI1)
apply(rule conjI)
apply clarify
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def
set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def
set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def
set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def
set-bit-def)
apply clarify apply(simp add: rmhead-free-list-def head-free-list-def set-bit-def
block-num-def)
apply(case-tac nat (free-l Va t) = i)
using mp-alloc-stm3-lm2-1 [of mem-pool-info Va p nat (free-l Va t)
tl (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t)))
(hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))) - buf
(mem-pool-info Va p)) div
lsizes Va t ! nat (free-l Va t)] apply meson

apply simp
apply(rule conjI)
using mp-alloc-stm3-lm2-inv apply simp

apply clarsimp apply(simp add:lvars-nochange-def)

apply(rule conjI)
using mp-alloc-stm3-lm2-inv apply simp

```

```

apply(rule conjI) apply clarsimp apply(simp add: rmhead-free-list-def head-free-list-def
set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def set-bit-def)
apply(rule conjI) apply(simp add: rmhead-free-list-def head-free-list-def set-bit-def)

apply(simp add:block-num-def)
apply(rule subst[where t=buf (set-bit-allocating ((mem-pool-info Va)(p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))) p (nat (free-l Va t))
((head-free-list (mem-pool-info Va p) (nat (free-l Va t)) -
buf (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) div
lsizes Va t ! nat (free-l Va t))
p) and s=buf (mem-pool-info Va p)])
apply(simp add:head-free-list-def rmhead-free-list-def set-bit-def)
apply(rule subst[where t=max-sz (set-bit-allocating ((mem-pool-info Va)(p :=
rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) p
(nat (free-l Va t))
((head-free-list (mem-pool-info Va p) (nat (free-l Va t)) -
buf (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) div
lsizes Va t ! nat (free-l Va t))
p) and s=max-sz (mem-pool-info Va p)])
apply(simp add:head-free-list-def rmhead-free-list-def set-bit-def)
apply(simp add:head-free-list-def)
apply(subgoal-tac lsizes Va t ! nat (free-l Va t) = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ (nat (free-l Va t)))
prefer 2 apply auto[1]
apply(subgoal-tac max-sz (mem-pool-info Va p) = ALIGN4 (max-sz (mem-pool-info
Va p)))
prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def)
apply (metis align40 mod-mult-self1-is-0 semiring-normalization-rules(17))
apply(subgoal-tac let fl = hd (free-list (levels (mem-pool-info Va p) ! nat (free-l
Va t)));
mp = (mem-pool-info Va p) in
(∃ n. n < n-max mp * (4 ^ (nat (free-l Va t)))
∧ fl = buf mp + n * (max-sz mp div (4 ^ (nat (free-l Va
t))))))
prefer 2 apply(simp add:inv-def inv-bitmap-freelist-def level-empty-def Let-def)
apply(subgoal-tac (nat (free-l Va t)) < length (levels (mem-pool-info Va p)))
prefer 2 apply (simp add: inv-mempool-info-def Let-def) apply (smt in-set-conv-nth
list.set-sel(1))
apply(rule conjI)
apply (metis add-diff-cancel-left' div-mult-self-is-m mult-is-0 neq0-conv)

apply(rule subst[where t=n-max (set-bit-allocating (λa. if a = p then rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)) else mem-pool-info Va a)
p (nat (free-l Va t))
((hd (free-list (levels (mem-pool-info Va p) ! nat (free-l Va t))) -
buf (rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))) div

```

```

      lsizes Va t ! nat (free-l Va t))
    p) and s=n-max (mem-pool-info Va p)])
  apply(simp add:rmhead-free-list-def set-bit-def)

  apply(rule conjI)
  using mp-alloc-stm3-lm2-5 apply metis
  using mp-alloc-stm3-lm2-3 apply(simp add:Let-def)
done

lemma head-free-list (mem-pool-info Va p) (nat (free-l Va t)) ≠ NULL ⇒
  { Va(⟦blk := (blk Va)(t := NULL)⟧), Va (⟦blk := (blk Va)(t := head-free-list
(mem-pool-info Va p) (nat (free-l Va t)))⟧),
    mem-pool-info := (mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))⟧} ∩
  { NULL < 'blk t } =
  { Va(⟦blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l Va
t)))⟧),
    mem-pool-info := (mem-pool-info Va)(p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))⟧}
by simp

lemma mp-alloc-stm3-lm1-1:
  inv Va ⇒ p ∈ mem-pools Va ⇒ nat (free-l Va t) < length (levels (mem-pool-info
Va p)) ⇒
  ¬ level-empty (mem-pool-info Va p) (nat (free-l Va t)) ⇒
  { V. V = Va(⟦blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l
Va t)))⟧),
    mem-pool-info := (mem-pool-info Va)
      (p := rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))⟧
    ∧ blk V t ≠ NULL }
  = { V. V = Va(⟦blk := (blk Va)(t := head-free-list (mem-pool-info Va p) (nat
(free-l Va t)))⟧),
    mem-pool-info := (mem-pool-info Va)
      (p := rmhead-free-list (mem-pool-info Va p) (nat (free-l Va t)))⟧ }
  apply(rule equalityI) apply(rule subsetI)
  apply clarsimp
  apply(rule subsetI)
  apply clarsimp
  apply(simp add: head-free-list-def)
  by (simp add: level-empty-def mp-alloc-stm3-lm2-4)

lemma mp-alloc-stm3-lm1:
  mp-alloc-precond1-70-2-2 t p sz timeout ∩ { 'cur = Some t } ∩ { Va } = { Va }
  ⇒ Γ ⊢I Some (IF level-empty ('mem-pool-info p) (nat ('free-l t)) THEN
    'blk := 'blk(t := NULL)
  ELSE
    'blk := 'blk(t := head-free-list ('mem-pool-info p) (nat ('free-l t))));

```

$(* \text{ sys-dlist-remove}(\text{node}); *)$   
 $'\text{mem-pool-info} := '\text{mem-pool-info } (p := \text{rmhead-free-list } (' \text{mem-pool-info } p) (\text{nat } (' \text{free-l } t)))$

$FI;;$   
 $(* ===== \text{end: block} = \text{sys-dlist-get}(\&p \rightarrow \text{levels}[l].\text{free-list}); *)$

$IF '\text{blk } t \neq \text{NULL} \text{ THEN}$   
 $(* \text{clear-free-bit}(p, l, \text{block-num}(p, \text{block}, \text{lsz})); *)$   
 $'\text{mem-pool-info} := \text{set-bit-allocating } '\text{mem-pool-info } p (\text{nat } (' \text{free-l } t))$   
 $(\text{block-num } (' \text{mem-pool-info } p) (' \text{blk } t) ((\text{'lsizes } t)!(\text{nat } (' \text{free-l } t)))));;$   
 $(* \text{set the allocating node info of the thread } *)$   
 $'\text{allocating-node} := '\text{allocating-node } (t := \text{Some } (\text{pool} = p, \text{level} = \text{nat } (' \text{free-l } t),$   
 $\text{block} = (\text{block-num } (' \text{mem-pool-info } p) (' \text{blk } t) ((\text{'lsizes } t)!(\text{nat } (' \text{free-l } t))))), \text{data} = '\text{blk } t))$   
 $FI) \text{ sat}_p [\text{mp-alloc-precond1-70-2-2 } t \text{ } p \text{ sz timeout} \cap \{\text{'cur} = \text{Some } t\} \cap \{Va\},$   
 $\{(s, t). s = t\}, \text{UNIV}, \{\text{'(Pair } Va) \in \text{Mem-pool-alloc-guar } t\} \cap \text{mp-alloc-precond2-1 } t \text{ } p \text{ sz timeout}]$

$\text{apply}(\text{subgoal-tac } Va \in \text{mp-alloc-precond1-70-2-2 } t \text{ } p \text{ sz timeout} \cap \{\text{'cur} = \text{Some } t\})$

$\text{prefer } 2 \text{ apply auto}[1]$   
 $\text{apply}(\text{rule Seq}[\text{where } \text{mid} =$   
 $\text{if level-empty } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t)) \text{ then}$   
 $\{V. V = Va(\text{blk} := (\text{blk } Va)(t := \text{NULL})) \wedge \text{level-empty } (\text{mem-pool-info } Va \text{ } p)$   
 $(\text{nat } (\text{free-l } Va \text{ } t)))\}$   
 $\text{else}$   
 $\{V. V = Va(\text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t))))),$   
 $\text{mem-pool-info} := (\text{mem-pool-info } Va)(p := \text{rmhead-free-list } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t))))\}$   
 $\wedge \neg \text{level-empty } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t)))\}])$

$\text{apply}(\text{rule Cond})$   
 $\text{apply}(\text{simp add:stable-def})$

$\text{apply}(\text{rule Basic})$   
 $\text{apply auto}[1] \text{ apply simp apply}(\text{simp add:stable-def}) \text{ apply}(\text{simp add:stable-def})$

$\text{apply}(\text{rule Seq}[\text{where } \text{mid} = \{V. V = Va(\text{blk} := (\text{blk } Va)(t := \text{head-free-list } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t))))\}$   
 $\wedge \neg \text{level-empty } (\text{mem-pool-info } Va \text{ } p) (\text{nat } (\text{free-l } Va \text{ } t)))\}])$

```

apply(rule Basic)
apply auto[1] apply simp apply(simp add:stable-def) apply(simp add:stable-def)
apply(rule Basic)
  apply clarify

apply auto[1] apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply simp

apply(rule Cond)
  apply(simp add:stable-def)
  apply(case-tac  $\neg$  level-empty (mem-pool-info Va p) (nat (free-l Va t)))
  prefer 2
  apply(subgoal-tac { Va( $\llbracket$  blk := (blk Va)(t := NULL) $\rrbracket$ ), Va
    ( $\llbracket$  blk := (blk Va)(t := head-free-list (mem-pool-info Va
p) (nat (free-l Va t))) $\rrbracket$ ),
    mem-pool-info := (mem-pool-info Va)
    (p := rmhead-free-list (mem-pool-info Va p) (nat
(free-l Va t))) $\rrbracket$  }  $\cap$ 
    ( $\neg$  level-empty (mem-pool-info Va p) (nat (free-l Va t))) $\rrbracket$ 
= { })
  prefer 2 apply auto[1]
  using Emptyprecond[where P=Some ('mem-pool-info := set-bit-allocating
'mem-pool-info p (nat ('free-l t))
    (block-num ('mem-pool-info p) ('blk t) (('lsizes t)!(nat
('free-l t)))));;
    (* set the allocating node info of the thread *)
    'allocating-node := 'allocating-node (t := Some ( $\llbracket$ pool = p, level = nat
('free-l t),
    block = (block-num ('mem-pool-info p) ('blk t) (('lsizes t)!(nat
('free-l t)))) $\rrbracket$ , data = 'blk t  $\rrbracket$ ))
    and rely={ (x,y). x=y } and guar=UNIV
    and post= $\llbracket$  ('Pair Va)  $\in$  Mem-pool-alloc-guar t $\rrbracket$   $\cap$  mp-alloc-precond2-1
t p sz timeout]
  apply meson apply auto[1]

  apply(rule subst[where t= (if level-empty (mem-pool-info Va p) (nat (free-l
Va t))
    then  $\llbracket$  '(=) (Va( $\llbracket$  blk := (blk Va)(t := NULL) $\rrbracket$ )  $\wedge$ 
level-empty (mem-pool-info Va p) (nat (free-l Va t))) $\rrbracket$ 
    else  $\llbracket$  '(=) (Va( $\llbracket$  blk := (blk Va)(t := head-free-list
(mem-pool-info Va p) (nat (free-l Va t))) $\rrbracket$ ,
    mem-pool-info := (mem-pool-info Va)
    (p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t))) $\rrbracket$ )  $\wedge$ 
     $\neg$  level-empty (mem-pool-info Va p) (nat (free-l Va
t))) $\rrbracket$   $\cap$ 
     $\llbracket$  'blk t  $\neq$ 

```

```

      NULL}} and s = { V. V = Va⟦blk := (blk Va)(t :=
head-free-list (mem-pool-info Va p) (nat (free-l Va t))),
      mem-pool-info := (mem-pool-info Va)
      (p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))⟩
      ∧ ¬ level-empty (mem-pool-info Va p) (nat
(free-l Va t))
      ∧ blk V t ≠ NULL}})

  apply auto[1]
  apply(rule subst[where t= { V. V = Va⟦blk := (blk Va)(t := head-free-list
(mem-pool-info Va p) (nat (free-l Va t))),
      mem-pool-info := (mem-pool-info Va)
      (p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))⟩
      ∧ ¬ level-empty (mem-pool-info Va p) (nat
(free-l Va t))
      ∧ blk V t ≠ NULL} and s = { V. V = Va⟦blk
:= (blk Va)(t := head-free-list (mem-pool-info Va p) (nat (free-l Va t))),
      mem-pool-info := (mem-pool-info Va)
      (p := rmhead-free-list (mem-pool-info
Va p) (nat (free-l Va t)))⟩
      ∧ ¬ level-empty (mem-pool-info Va p) (nat
(free-l Va t))}]]

  apply(subgoal-tac (nat (free-l Va t)) < length (levels (mem-pool-info Va p)))
  prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def) apply
force
  apply simp
  using mp-alloc-stm3-lm1-1 apply force

  apply(rule Seq[where mid=
    { V. let vb = Va⟦blk := (blk Va)(t := head-free-list (mem-pool-info Va p)
(nat (free-l Va t))),
      mem-pool-info := (mem-pool-info Va) (p := rmhead-free-list
(mem-pool-info Va p) (nat (free-l Va t)))⟩
      in V = vb⟦ mem-pool-info := set-bit-allocating (mem-pool-info vb) p (nat
(free-l vb t))
      (block-num (mem-pool-info vb p) (blk vb t) ((lsizes vb t)!(nat
(free-l vb t))))⟩
      ∧ ¬ level-empty (mem-pool-info Va p) (nat (free-l Va t))}]]
  apply(rule Basic) apply clarsimp
  apply simp apply(simp add:stable-def) apply(simp add:stable-def)
  apply(rule Basic)
  using mp-alloc-stm3-lm2 apply meson
  apply simp apply(simp add:stable-def)
  using stable-id2[of ⦿'(Pair Va) ∈ Mem-pool-alloc-guar t} ∩ mp-alloc-precond2-1
t p sz timeout]
  apply meson

  apply(unfold Skip-def)

```

```

    apply(rule Basic)
    using mp-alloc-stm3-lm3 apply meson
    apply simp apply(simp add:stable-def)
    using stable-id2[of  $\{(Pair\ Va) \in Mem\text{-pool}\text{-alloc}\text{-guar}\ t\} \cap mp\text{-alloc}\text{-precond2}\text{-1}$ 
    t p sz timeout]
    apply meson
    apply simp

done

lemma mp-alloc-stm3-lm:
 $\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{ATOMIC}$ 
    (* ===== start: block = sys-dlist-get(&p->levels[l].free-list); *)
    IF level-empty ('mem-pool-info p) (nat ('free-l t)) THEN
        'blk := 'blk(t := NULL)
    ELSE
        'blk := 'blk(t := head-free-list ('mem-pool-info p) (nat ('free-l t)));

    (* sys-dlist-remove(node); *)
    'mem-pool-info := 'mem-pool-info (p := rmhead-free-list ('mem-pool-info
p) (nat ('free-l t)))

    FI;;
    (* ===== end: block = sys-dlist-get(&p->levels[l].free-list); *)

    IF 'blk t  $\neq$  NULL THEN
        (* clear-free-bit(p, l, block-num(p, block, lsz)); *)
        'mem-pool-info := set-bit-allocating 'mem-pool-info p (nat ('free-l t))
            (block-num ('mem-pool-info p) ('blk t) ((lsizes t)!(nat
('free-l t))));
        (* set the allocating node info of the thread *)
        'allocating-node := 'allocating-node (t := Some (pool = p, level = nat
('free-l t),
            block = (block-num ('mem-pool-info p) ('blk t) ((lsizes t)!(nat
('free-l t)))), data = 'blk t ))
        FI
    END)
    sat_p [mp-alloc-precond1-70-2-2 t p sz timeout, Mem-pool-alloc-rely t, Mem-pool-alloc-guar
t,
        mp-alloc-precond2-1 t p sz timeout]
    apply(simp add:stm-def)
    apply(rule Await)
    using mp-alloc-precond1-70-2-2-stb apply simp
    using mp-alloc-precond2-1-stb apply simp
    apply(clarify)
    apply(rule Await)
    using stable-id2 apply fast using stable-id2 apply fast

```



```

    apply clarify
    apply(case-tac V = Va) prefer 2 apply simp using Emptyprecond apply
auto[1]
    apply simp
    apply(case-tac mp-alloc-precond1-70-2-2 t p sz timeout
       $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{Va\} = \{\}$ )
    using Emptyprecond apply auto[1]
    apply(subgoal-tac mp-alloc-precond1-70-2-2 t p sz timeout
       $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{Va\} = \{Va\}$ )
    prefer 2 using int1-eq[where P=mp-alloc-precond1-70-2-2 t p sz timeout
 $\cap \llbracket 'cur = Some\ t \rrbracket$ ] apply meson
    using mp-alloc-stm3-lm1[of t p timeout sz] apply auto[1]
done

term mp-alloc-precond1-70-2-2 t p sz timeout
term mp-alloc-precond2-1 t p sz timeout

```

## 10.7 stm4

**abbreviation**  $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}0\ t\ p\ sz\ tm \equiv$   
 $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\ t\ p\ sz\ tm \cap \llbracket 'from\text{-}l\ t < 'alloc\text{-}l\ t \rrbracket$

**lemma**  $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}0\text{-}stb$ : *stable* ( $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}0\ t\ p\ sz\ tm$ ) (*Mem-pool-alloc-rely* t)  
 apply(rule stable-int2)  
 using mp-alloc-precond2-1-1-loopinv-stb apply auto[1]  
 apply(simp add:stable-def) apply clarify  
 apply(simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)

```

    apply(case-tac x=y) apply auto[1] apply clarify
    apply(simp add: block-num-def lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
done

```

**abbreviation**  $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}1'\ t\ p\ sz\ tm \equiv$   
 $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\ t\ p\ sz\ tm \cap \llbracket 'from\text{-}l\ t \leq 'alloc\text{-}l\ t \wedge 'from\text{-}l\ t \geq 'free\text{-}l\ t$   
 $\wedge 'allocating\text{-}node\ t = Some\ (\llbracket pool = p, level = nat\ ('from\text{-}l\ t + 1),$   
 $block = block\text{-}num\ ('mem\text{-}pool\text{-}info\ p)\ ('blk\ t)\ (('lsizes$   
 $t)\!(nat\ ('from\text{-}l\ t + 1))),$   
 $data = 'blk\ t\ \rrbracket \rrbracket$

**abbreviation**  $mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}1\ t\ p\ sz\ tm \equiv$   
 $\{s.\ inv\ s\} \cap \llbracket 'freeing\text{-}node\ t = None \rrbracket \cap \llbracket p \in 'mem\text{-}pools \wedge tm \geq -1 \rrbracket \cap$   
 $mp\text{-}alloc\text{-}precond7\text{-}ext\ t\ p\ sz\ tm \cap \llbracket \neg 'rf\ t \rrbracket$

```

     $\cap$  mp-alloc-precond1-70-ext t p sz tm  $\cap$   $\neg \{ \text{'alloc-l } t < 0 \}$ 
     $\cap$   $\neg \{ \text{'free-l } t < 0 \}$   $\cap$   $\neg \{ \text{'blk } t = \text{NULL} \}$   $\cap$   $\{ \text{'from-l } t < \text{'alloc-l } t \}$ 
     $\cap$   $\{ \text{'from-l } t \leq \text{'alloc-l } t \wedge \text{'from-l } t \geq \text{'free-l } t$ 
       $\wedge$  allocating-node t = Some (pool = p, level = nat ('from-l t + 1),
        block = block-num ('mem-pool-info p) ('blk t) ('lsizes
t)(nat ('from-l t + 1))),
        data = 'blk t  $\}$ 
     $\wedge$  alloc-memblk-data-valid p (the ('allocating-node t))
     $\wedge$  ( $\exists n. n < n\text{-max}$  ('mem-pool-info p) * ( $4 \wedge (\text{nat } ('from-l t + 1))$ ))
     $\wedge$  'blk t = buf ('mem-pool-info p) + n * (max-sz ('mem-pool-info p)
div ( $4 \wedge (\text{nat } ('from-l t + 1))$ ))  $\}$ 

```

```

term mp-alloc-precond2-1-1-loopinv-0 t p sz tm
term mp-alloc-precond2-1-1-loopinv-1 t p sz tm
term mp-alloc-precond2-1-1-loopinv-1' t p sz tm

```

**lemma** *mp-alloc-precond2-1-1-loopinv-1'-stb*: *stable* (*mp-alloc-precond2-1-1-loopinv-1'* *t p sz tm*) (*Mem-pool-alloc-rely t*)

```

  apply(rule stable-int2)
  using mp-alloc-precond2-1-1-stb apply auto[1]
  apply(simp add:stable-def) apply clarify
  apply(simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)

  apply(case-tac x=y) apply auto[1] apply clarify
  apply(simp add: block-num-def lvars-nochange-rel-def lvars-nochange-def gvars-conf-stable-def
gvars-conf-def)
done

```

**lemma** *mp-alloc-precond2-1-1-loopinv-1-stb*: *stable* (*mp-alloc-precond2-1-1-loopinv-1* *t p sz tm*) (*Mem-pool-alloc-rely t*)

```

  apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2)
  apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2)
  apply(rule stable-int2) apply(rule stable-int2) apply(rule stable-int2)
  apply(rule stable-int2)
  apply (simp add: stable-inv-alloc-rely1)
  apply (simp add: mp-alloc-freenode-stb)
  apply (simp add: mp-alloc-precond1-ext-stb)
  apply (simp add: mp-alloc-precond7-ext-stb)
  apply (simp add: mp-alloc-precond1-0-ext-stb)
  using mp-alloc-precond1-70-ext-stb apply blast
  apply (simp add: mp-alloc-precond1-70-2-ext-stb)
  apply (simp add: mp-alloc-precond1-70-2-2-ext-stb)
  using mp-alloc-precond2-1-1-ext-stb apply blast
  apply(simp add:stable-def) apply clarify
  apply(simp add:Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt

```

```

  apply(simp add:stable-def) apply clarify

```

```

  apply(subgoal-tac buf (mem-pool-info x p) = buf (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply metis
  apply(subgoal-tac from-l x t + 1 = from-l y t + 1)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply metis
  apply(subgoal-tac blk x t = blk y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac allocating-node x t = allocating-node y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac lsizes x t = lsizes y t)
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def lvars-nochange-rel-def lvars-nochange-def)
apply smt
  apply(subgoal-tac n-max (mem-pool-info x p) = n-max (mem-pool-info y p))
  prefer 2 apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def)
apply metis
  apply(subgoal-tac block-num (mem-pool-info x p) (blk x t) (lsizes x t ! nat (from-l
x t + 1)))
    = block-num (mem-pool-info y p) (blk y t) (lsizes y t ! nat (from-l
y t + 1)))
  prefer 2 apply(simp add: block-num-def Mem-pool-alloc-rely-def lvars-nochange-rel-def
lvars-nochange-def)
  apply(case-tac x=y) apply auto[1]
  apply(simp add: Mem-pool-alloc-rely-def gvars-conf-stable-def gvars-conf-def lvars-nochange-rel-def
lvars-nochange-def)
  apply smt
done

```

**abbreviation** *mp-alloc-stm4-pre-precond1*  $Va\ t\ p \equiv$

$Va(\lambda bn := (bn\ Va)\ (t := block\_num\ (mem\_pool\_info\ Va\ p)\ (blk\ Va\ t)\ ((lsizes\ Va\ t)!(nat\ (from-l\ Va\ t))))))$

**abbreviation** *mp-alloc-stm4-pre-precond2*  $Va\ t\ p \equiv$

$Va(\lambda mem\_pool\_info := set\_bit\_divide\ (mem\_pool\_info\ Va)\ p\ (nat\ (from-l\ Va\ t))\ (bn\ Va\ t))$

**abbreviation** *mp-alloc-stm4-pre-precond3*  $Va\ t\ p \equiv$

$Va(\lambda mem\_pool\_info := set\_bit\_allocating\ (mem\_pool\_info\ Va)\ p\ (nat\ (from-l\ Va\ t\ +\ 1))\ (4 * bn\ Va\ t))$

**abbreviation** *mp-alloc-stm4-pre-precond4*  $Va\ t\ p \equiv$

$Va(\lambda allocating\_node := (allocating\_node\ Va)(t := Some\ (\lambda pool = p,\ level = nat\ (from-l\ Va\ t\ +\ 1)),$

$block = 4 * bn\ Va\ t,\ data = blk\ Va\ t\ \lambda))$

**abbreviation** *mp-alloc-stm4-pre-precond5*  $Va\ t\ p \equiv Va(\lambda i := (i\ Va)\ (t := 1))$

**definition** *mp-alloc-stm4-pre-precond-f*  $Va\ t\ p$  (\* we use definition here. abbreviation leads slow parsing in lemmas \*)

$$\begin{aligned} &\equiv mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond5 \\ &\quad (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond4 \\ &\quad (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond3 \\ &\quad (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond2 \\ &\quad (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond1\ Va\ t\ p)\ t\ p)\ t\ p)\ t\ p\ t\ p \end{aligned}$$

**abbreviation** *mp-alloc-stm4-loopinv*  $Va\ t\ mp$

$$\begin{aligned} &\equiv \{ V. cur\ V = cur\ Va \wedge tick\ V = tick\ Va \wedge thd\text{-}state\ V = thd\text{-}state\ Va \wedge \\ &\quad (V, Va) \in gvars\text{-}conf\text{-}stable \\ &\quad \wedge (\forall p. p \neq mp \longrightarrow mem\text{-}pool\text{-}info\ V\ p = mem\text{-}pool\text{-}info\ Va\ p) \\ &\quad \wedge wait\text{-}q\ (mem\text{-}pool\text{-}info\ V\ mp) = wait\text{-}q\ (mem\text{-}pool\text{-}info\ Va\ mp) \\ &\quad \wedge (\forall t'. t' \neq t \longrightarrow lvars\text{-}nochange\ t'\ V\ Va) \\ &\quad \wedge (\forall jj. jj \neq nat\ (from\text{-}l\ Va\ t + 1) \longrightarrow levels\ (mem\text{-}pool\text{-}info\ V\ mp)\ !\ jj = \\ &\quad levels\ (mem\text{-}pool\text{-}info\ Va\ mp)\ !\ jj) \\ &\quad \wedge (bits\ (levels\ (mem\text{-}pool\text{-}info\ V\ mp)\ !\ nat\ (from\text{-}l\ Va\ t + 1)) \\ &\quad = list\text{-}updates\text{-}n\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ Va\ mp)\ !\ nat\ (from\text{-}l\ Va\ t \\ &\quad + 1)))\ (bn\ Va\ t * 4 + 1)\ (i\ V\ t - 1)\ FREE) \\ &\quad \wedge (free\text{-}list\ (levels\ (mem\text{-}pool\text{-}info\ V\ mp)\ !\ nat\ (from\text{-}l\ Va\ t + 1)) \\ &\quad = inserts\ (map\ (\lambda ii. (lsizes\ Va\ t)\ !\ (nat\ (from\text{-}l\ Va\ t + 1)) * ii + blk \\ &\quad V\ t)\ [1..<i\ V\ t]) \\ &\quad (free\text{-}list\ (levels\ (mem\text{-}pool\text{-}info\ Va\ mp)\ !\ nat\ (from\text{-}l\ Va\ t + 1)))) \\ &\quad \wedge j\ V = j\ Va \wedge ret\ V = ret\ Va \wedge endt\ V = endt\ Va \wedge rf\ V = rf\ Va \wedge \\ &\quad tmout\ V = tmout\ Va \\ &\quad \wedge lsizes\ V = lsizes\ Va \wedge alloc\text{-}l\ V = alloc\text{-}l\ Va \wedge free\text{-}l\ V = free\text{-}l\ Va \\ &\quad \wedge from\text{-}l\ V = from\text{-}l\ Va \wedge blk\ V = blk\ Va \wedge nodev\ V = nodev\ Va \\ &\quad \wedge bn\ V = bn\ Va \wedge alloc\text{-}lsize\text{-}r\ V = alloc\text{-}lsize\text{-}r\ Va \wedge lvl\ V = lvl\ Va \wedge bb \\ &\quad V = bb\ Va \\ &\quad \wedge block\text{-}pt\ V = block\text{-}pt\ Va \wedge th\ V = th\ Va \wedge need\text{-}resched\ V = need\text{-}resched \\ &\quad Va \\ &\quad \wedge mempoolalloc\text{-}ret\ V = mempoolalloc\text{-}ret\ Va \wedge freeing\text{-}node\ V = freeing\text{-}node \\ &\quad Va \\ &\quad \wedge allocating\text{-}node\ V = allocating\text{-}node\ Va \\ &\quad \wedge i\ V\ t > 0 \wedge i\ V\ t \leq 4 \} \end{aligned}$$

**lemma** *in-mp-alloc-stm4-loopinv*:  $i\ V\ t = 1 \implies V \in mp\text{-}alloc\text{-}stm4\text{-}loopinv\ V\ t\ mp$

**apply** *simp*  
**apply**(*rule conjI*) **apply**(*simp add:gvars-conf-stable-def gvars-conf-def*)  
**apply**(*rule conjI*) **apply**(*simp add:lvars-nochange-def*)  
**apply**(*simp add:inserts-def*)  
**done**

**abbreviation** *mp-alloc-stm4-while-precond1*  $V\ t\ p \equiv$

$$V(\text{!}bn := (bn\ V)\ (t := 4 * bn\ V\ t + i\ V\ t))$$

**abbreviation** *mp-alloc-stm4-while-precond2*  $V\ t\ p \equiv$

$$V(\text{!}lsz := (lsz\ V)\ (t := lsizes\ V\ t\ !\ nat\ (from\text{-}l\ V\ t + 1)))$$

**abbreviation** *mp-alloc-stm4-while-precond3*  $V\ t\ p \equiv$

$V(\text{block2} := (\text{block2 } V) (t := \text{lsz } V t * i V t + \text{blk } V t))$   
**abbreviation**  $\text{mp-alloc-stm4-while-precond4 } V t p \equiv$   
 $V(\text{mem-pool-info} := \text{set-bit-free } (\text{mem-pool-info } V) p (\text{nat } (\text{from-l } V t + 1)) (\text{lbn } V t))$   
**abbreviation**  $\text{mp-alloc-stm4-while-precond5 } V t p \equiv$   
 $V(\text{mem-pool-info} := (\text{mem-pool-info } V) (p := \text{append-free-list } (\text{mem-pool-info } V p) (\text{nat } (\text{from-l } V t + 1)) (\text{block2 } V t)))$

**lemma**  $\text{mp-alloc-stm4-pre-in}$ :  
 $\{ \text{mp-alloc-stm4-while-precond4}$   
 $(\text{mp-alloc-stm4-while-precond3}$   
 $(\text{mp-alloc-stm4-while-precond2}$   
 $(\text{mp-alloc-stm4-while-precond1 } V t p) t p) t p) t p \} \cap$   
 $\{ \text{block-fits } (' \text{mem-pool-info } p) (' \text{block2 } t) (' \text{lsz } t) \}$   
 $\subseteq \{ \text{mp-alloc-stm4-while-precond5 } t p$   
 $\in \{ \text{mp-alloc-stm4-while-precond5}$   
 $(\text{mp-alloc-stm4-while-precond4}$   
 $(\text{mp-alloc-stm4-while-precond3}$   
 $(\text{mp-alloc-stm4-while-precond2}$   
 $(\text{mp-alloc-stm4-while-precond1 } V t p) t p) t p) t p) t p \}$

**by** *auto*

**lemma**  $\text{mp-alloc-stm4-lsizes}$ :  $\text{lsizes } Va t = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma**  $\text{mp-alloc-stm4-pre-froml}$ :  $\text{from-l } Va t = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma**  $\text{mp-alloc-stm4-pre-buf}$ :  $\text{buf } (\text{mem-pool-info } Va q) = \text{buf } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) q)$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def set-bit-def*)

**lemma**  $\text{mp-alloc-stm4-nmax}$ :  $n\text{-max } (\text{mem-pool-info } Va q) = n\text{-max } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) q)$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def set-bit-def*)

**lemma**  $\text{mp-alloc-stm4-pre-maxsz}$ :  $\text{max-sz } (\text{mem-pool-info } Va q) = \text{max-sz } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) q)$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def set-bit-def*)

**lemma**  $\text{mp-alloc-stm4-blk}$ :  $\text{blk } Va = \text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va t p)$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma**  $\text{mp-alloc-stm4-blockfit-help-1}$ :  
 $p \in \text{mem-pools } Va \implies \text{inv } Va \implies$   
 $(\forall ii < \text{length } (\text{lsizes } Va t). \text{lsizes } Va t ! ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va p)) \text{ div } 4 \wedge ii) \implies$

```

length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) ⇒
ALIGN4 (max-sz (mem-pool-info Va p)) = max-sz (mem-pool-info Va p) ⇒
Suc (nat (from-l Va t)) < length (lsizes Va t) ⇒
max-sz (mem-pool-info Va p) = 4 ^ nat (from-l Va t) * lsizes Va t ! nat (from-l
Va t)
apply(simp add:inv-def inv-mempool-info-def)
apply(subgoal-tac ∃ n>0. max-sz (mem-pool-info Va p) = 4 * n * 4 ^ n-levels
(mem-pool-info Va p))
prefer 2 apply meson
apply(subgoal-tac max-sz (mem-pool-info Va p) mod (4 ^ nat (from-l Va t)) =
0)
prefer 2 apply clarsimp using ge-pow-mod-0[of nat (from-l Va t) n-levels
(mem-pool-info Va p)] apply auto[1]
using mod0-div-self[of max-sz (mem-pool-info Va p) 4 ^ nat (from-l Va t)] apply
simp
done

```

**lemma** mp-alloc-stm4-blockfit-help2:

```

ii ≥ 0 ⇒ (a::nat) mod 4 ^ nat (ii+1) = 0 ⇒ a div 4 ^ nat (ii+1) * 4 = a
div 4 ^ (nat ii)
apply(subgoal-tac nat (ii+1) = nat ii + 1) prefer 2 apply auto[1]
by auto

```

**lemma** mp-alloc-stm4-blockfit-help3:

```

inv Va ⇒
p ∈ mem-pools Va ⇒
nmax > 0 ⇒
maxsz mod frml = 0 ⇒
n < nmax * frml ⇒
blk Va t = buf (mem-pool-info Va p) + n * (maxsz div frml) ⇒
maxsz div frml + blk Va t
≤ nmax * maxsz + buf (mem-pool-info Va p)
apply(case-tac n = 0)
apply (metis (no-types, lifting) Nat.add-0-right add-le-mono1 div-le-dividend div-mult-self1-is-m
le-trans mult-is-0)

```

```

apply(subgoal-tac blk Va t ≤ buf (mem-pool-info Va p) + (nmax * frml - 1) *
(maxsz div frml))
prefer 2 apply auto[1]
by (smt add.left-commute add-diff-cancel-left' le-diff-conv mp-alloc-stm3-lm2-3-1
mult-eq-if nat-add-left-cancel-le not-less-zero)

```

**lemma** mp-alloc-stm4-blockfit-help4:

```

inv Va ∧
p ∈ mem-pools Va ∧
length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) ∧
alloc-l Va t < int (n-levels (mem-pool-info Va p)) ∧

```

```

from-l Va t < alloc-l Va t ∧
¬ free-l Va t < OK ∧
free-l Va t ≤ from-l Va t ⇒
  ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (from-l Va t + 1) * 4 =
  ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat (from-l Va t)
apply(subgoal-tac ALIGN4 (max-sz (mem-pool-info Va p)) mod 4 ^ nat (from-l
Va t + 1) = 0)
prefer 2
apply(rule subst[where t=ALIGN4 (max-sz (mem-pool-info Va p)) and s=max-sz
(mem-pool-info Va p)])
apply (metis inv-maxsz-align4)
apply(subgoal-tac nat (from-l Va t + 1) < n-levels (mem-pool-info Va p))
prefer 2 apply (smt nat-less-iff)
apply(simp add:inv-def inv-mempool-info-def Let-def)
apply(subgoal-tac n-levels (mem-pool-info Va p) = length (levels (mem-pool-info
Va p)))
prefer 2 apply simp
apply (metis ge-pow-mod-0)
apply(rule mp-alloc-stm4-blockfit-help-2)
apply (metis int-nat-eq linorder-not-less nat-int neq0-conv zless-nat-conj)
apply simp
done

```

**lemma** mp-alloc-stm4-blockfit-help:

```

Va ∈ mp-alloc-precond2-1-1-loopinv-0 t p sz timeout ∩ {cur = Some t} ∧ ii <
4
⇒ block-fits (mem-pool-info Va p) (lsizes Va t ! nat (from-l Va t + 1) * ii +
blk Va t)

```

```

      (lsizes Va t ! nat (from-l Va t + 1))
apply(simp add:block-fits-def block-num-def buf-size-def)

apply(case-tac alloc-l Va t = ETIMEOUT ∧ free-l Va t = ETIMEOUT ∧ length
(lsizes Va t) = Suc NULL)
apply simp apply simp

```

```

apply(subgoal-tac lsizes Va t ! nat (from-l Va t + 1) = ALIGN4 (max-sz
(mem-pool-info Va p)) div 4 ^ nat (from-l Va t + 1))
prefer 2 apply(subgoal-tac nat (from-l Va t + 1) < length (lsizes Va t))
prefer 2 apply (smt nat-less-iff)
apply simp
apply(subgoal-tac lsizes Va t ! nat (from-l Va t) = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ nat (from-l Va t))
prefer 2 apply(subgoal-tac nat (from-l Va t) < length (lsizes Va t))
prefer 2 apply (smt nat-less-iff)
apply simp

```

```

apply(rule subst[where t=lsizes Va t ! nat (from-l Va t + 1) * ii + blk Va t +
lsizes Va t ! nat (from-l Va t + 1)
and s=ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ nat

```

$(\text{from-l } Va \ t + 1) * ii + \text{blk } Va \ t + \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1))$   
**apply simp**  
 $\text{apply}(\text{subgoal-tac } (\text{blk } Va \ t - \text{buf } (\text{mem-pool-info } Va \ p)) \text{ div } \text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t))$   
 $< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t))$   
**prefer 2 apply force**  
 $\text{apply}(\text{subgoal-tac } \text{blk } Va \ t \geq \text{buf } (\text{mem-pool-info } Va \ p))$   
**prefer 2 apply force**  
 $\text{apply}(\text{subgoal-tac } (\text{blk } Va \ t - \text{buf } (\text{mem-pool-info } Va \ p)) \text{ mod } (\text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) = 0)$   
**prefer 2 apply** (*metis diff-add-inverse inv-maxsz-align4 mod-mult-self2-is-0*)  
  
 $\text{apply}(\text{subgoal-tac } \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1) * ii + \text{blk } Va \ t +$   
 $\text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1)$   
 $\leq \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1) * 4$   
 $+ \text{blk } Va \ t)$   
**prefer 2 apply simp**  
  
 $\text{apply}(\text{subgoal-tac } \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1) * 4$   
 $= \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t))$   
**prefer 2 using mp-alloc-stm4-blockfit-help4 apply simp**  
  
 $\text{apply}(\text{subgoal-tac } \text{ALIGN}_4 (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) + \text{blk } Va \ t$   
 $\leq n\text{-max } (\text{mem-pool-info } Va \ p) * \text{max-sz } (\text{mem-pool-info } Va \ p) +$   
 $\text{buf } (\text{mem-pool-info } Va \ p))$   
**apply simp**  
**apply clarify**  
 $\text{apply}(\text{subgoal-tac } n\text{-max } (\text{mem-pool-info } Va \ p) > 0)$   
**prefer 2 apply** (*metis gr0I mult-is-0 not-less-zero*)  
 $\text{apply}(\text{subgoal-tac } \text{max-sz } (\text{mem-pool-info } Va \ p) \text{ mod } 4 \wedge \text{nat } (\text{from-l } Va \ t) = 0)$   
**prefer 2 apply** ( $\text{subgoal-tac } \exists n > 0. \text{max-sz } (\text{mem-pool-info } Va \ p) = (4 * n) * (4 \wedge n\text{-levels } (\text{mem-pool-info } Va \ p))$ )  
**prefer 2 apply** (*metis inv-mempool-info-def inv-def*)  
**apply** (*smt ge-pow-mod-0 of-nat-less-imp-less zless-nat-conj zless-nat-eq-int-zless*)  
  
**using mp-alloc-stm4-blockfit-help3**[*of Va p n-max (mem-pool-info Va p) max-sz (mem-pool-info Va p) 4 \wedge nat (from-l Va t) - t*]  
**by** (*metis inv-maxsz-align4*)

**lemma**  $a \geq b \implies c * d \geq e \implies (a::\text{nat}) - b \leq c * d - e \implies a + e - b \leq c * d$



```

by (simp add: Nat.le-diff-conv2)

lemma a + b > c  $\implies$  a + b - c  $\leq$  d  $\implies$  e + f < b  $\implies$  e + a + f - Suc c <
d
by simp

lemma (x::nat) > y  $\implies$   $\exists n. (4::nat) \wedge x = 4 \wedge y * n$ 
  apply(subgoal-tac 4  $\wedge x = 4 \wedge y * 4 \wedge (x - y)$ ) prefer 2 apply auto
  by (metis add-diff-inverse-nat less-imp-le-nat not-less power-add)

lemma int-empt1: ( $\forall v. v \in P \longrightarrow v \notin Q$ )  $\implies P \cap Q = \{\}$  by auto

lemma mp-alloc-stm4-blockfit1-1:
  allocating-node Va t =
    Some ( $\llbracket pool = p, level = nat (from-l Va t), block = block-num (mem-pool-info$ 
    Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t)),
    data = blk Va t  $\rrbracket \wedge$  data (the (allocating-node Va t)) = buf (mem-pool-info
    Va p) +
    block (the (allocating-node Va t)) * (max-sz (mem-pool-info Va p) div 4  $\wedge$  level
    (the (allocating-node Va t)))
     $\wedge$  block (the (allocating-node Va t)) < n-max (mem-pool-info Va p) * 4  $\wedge$  level
    (the (allocating-node Va t))  $\implies$ 
    alloc-blk-valid Va p (nat (from-l Va t)) (block-num (mem-pool-info Va p) (blk Va
    t) (lsizes Va t ! nat (from-l Va t))) (blk Va t)
  apply(simp add:block-num-def) apply auto
done

lemma mp-alloc-stm4-blockfit1:
  Va  $\in$  mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap \llbracket 'cur = Some t \rrbracket \implies$ 
  V  $\in$  mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p  $\cap \llbracket 'i t <$ 
  4  $\rrbracket \implies$ 
   $\forall v. v \in \{mp-alloc-stm4-while-precond4$ 
    (mp-alloc-stm4-while-precond3
    (mp-alloc-stm4-while-precond2
    (mp-alloc-stm4-while-precond1 V t p) t p) t p) t p  $\} \longrightarrow$ 
    v  $\notin$  -  $\llbracket block-fits ('mem-pool-info p) ('block2 t) ('lsz t) \rrbracket$ 
  apply simp
  apply(simp add:block-fits-def buf-size-def set-bit-def)
  apply(subgoal-tac lsizes (mp-alloc-stm4-pre-precond-f Va t p) t = lsizes Va t)
  prefer 2 using mp-alloc-stm4-lsizes apply metis
  apply(subgoal-tac from-l (mp-alloc-stm4-pre-precond-f Va t p) t = from-l Va t)
  prefer 2 using mp-alloc-stm4-pre-froml apply metis
  apply(subgoal-tac buf (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) =
  buf (mem-pool-info Va p))
  prefer 2 using mp-alloc-stm4-pre-buf apply metis
  apply(subgoal-tac n-max (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p)

```

```

= n-max (mem-pool-info Va p))
  prefer 2 using mp-alloc-stm4-nmax apply metis
  apply(subgoal-tac max-sz (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p)
p) = max-sz (mem-pool-info Va p))
  prefer 2 using mp-alloc-stm4-pre-maxsz apply metis
  apply(subgoal-tac blk (mp-alloc-stm4-pre-precond-f Va t p) = blk Va)
  prefer 2 using mp-alloc-stm4-blk apply metis

  apply(subgoal-tac buf (mem-pool-info V p) = buf (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p))
  prefer 2 apply (simp add:gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac n-max (mem-pool-info V p) = n-max (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p))
  prefer 2 apply (simp add:gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac max-sz (mem-pool-info V p) = max-sz (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p))
  prefer 2 apply (simp add:gvars-conf-stable-def gvars-conf-def)

  apply(rule subst[where t=buf (mem-pool-info V p) and s=buf (mem-pool-info
Va p)]) apply simp
  apply(rule subst[where t=n-max (mem-pool-info V p) and s=n-max (mem-pool-info
Va p)]) apply simp
  apply(rule subst[where t=max-sz (mem-pool-info V p) and s=max-sz (mem-pool-info
Va p)]) apply simp
  apply(rule subst[where t=lsizes (mp-alloc-stm4-pre-precond-f Va t p) t and
s=lsizes Va t]) apply simp
  apply(rule subst[where t=from-l (mp-alloc-stm4-pre-precond-f Va t p) t and
s=from-l Va t]) apply simp
  apply(rule subst[where t=blk (mp-alloc-stm4-pre-precond-f Va t p) and s=blk
Va]) apply simp

  using mp-alloc-stm4-blockfit-help [of Va t p timeout sz i V t] apply(unfold
block-fits-def buf-size-def) apply simp
  apply(case-tac alloc-l Va t = ETIMEOUT  $\wedge$  free-l Va t = ETIMEOUT  $\wedge$  length
(lsizes Va t) = Suc NULL)
  apply simp
  apply simp
  apply(subgoal-tac alloc-blk-valid Va p (nat (from-l Va t)) (block-num (mem-pool-info
Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))) (blk Va t))
  prefer 2 using mp-alloc-stm4-blockfit1-1[of Va t p] apply argo apply metis
done

```

**lemma** mp-alloc-stm4-blockfit:

```

Va  $\in$  mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap$   $\{i'cur = Some t\} \implies$ 
V  $\in$  mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p  $\cap$   $\{i' t <$ 
4  $\} \implies$ 
  {mp-alloc-stm4-while-precond4
  (mp-alloc-stm4-while-precond3

```

```

      (mp-alloc-stm4-while-precond2
      (mp-alloc-stm4-while-precond1 V t p) t p) t p} ∩
      – {block-fits ('mem-pool-info p) ('block2 t) ('lsz t)} = {}
using mp-alloc-stm4-blockfit1 [of Va t p timeout sz V]
int-empt1 [of {mp-alloc-stm4-while-precond4
      (mp-alloc-stm4-while-precond3
      (mp-alloc-stm4-while-precond2
      (mp-alloc-stm4-while-precond1 V t p) t p) t p) t p} – {block-fits ('mem-pool-info
p) ('block2 t) ('lsz t)}]]
apply meson

done

term mp-alloc-precond2-1-1-loopinv-0 t p sz timeout ∩ { 'cur = Some t}

term mp-alloc-precond2-1-1-loopinv-0 t p sz timeout ∩ { 'cur = Some t}
term mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p ∩ { 'i t < 4}
term {mp-alloc-stm4-while-precond4
      (mp-alloc-stm4-while-precond3
      (mp-alloc-stm4-while-precond2
      (mp-alloc-stm4-while-precond1 V t p) t p) t p) t p} ∩
      – {block-fits ('mem-pool-info p) ('block2 t) ('lsz t)}

lemma mp-alloc-stm4-inv-mif-buf: buf (mem-pool-info Va pa) = buf (mem-pool-info
(mp-alloc-stm4-pre-precond-f Va t p) pa)
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(simp add: set-bit-def)
done

lemma mp-alloc-stm4-inv-mif-maxsz: max-sz (mem-pool-info Va pa) = max-sz (mem-pool-info
(mp-alloc-stm4-pre-precond-f Va t p) pa)
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(simp add: set-bit-def)
done

lemma mp-alloc-stm4-inv-mif-nmax: n-max (mem-pool-info Va pa) = n-max (mem-pool-info
(mp-alloc-stm4-pre-precond-f Va t p) pa)
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(simp add: set-bit-def)
done

lemma mp-alloc-stm4-inv-mif-nlvl: n-levels (mem-pool-info Va pa) = n-levels (mem-pool-info
(mp-alloc-stm4-pre-precond-f Va t p) pa)
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(simp add: set-bit-def)
done

lemma mp-alloc-stm4-inv-mif-len: length (levels (mem-pool-info Va pa)) = length

```

```

(levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) pa))
  apply(simp add:mp-alloc-stm4-pre-precond-f-def)
  apply(simp add: set-bit-def)
done

```

```

lemma mp-alloc-stm4-inv-bits-len: length (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) pa) ! ii))
  = length (bits (levels (mem-pool-info Va pa) ! ii))

```

```

  apply(simp add:mp-alloc-stm4-pre-precond-f-def)
  apply(case-tac p ≠ pa) apply(simp add: set-bit-def)
  apply simp
  apply(case-tac nat (from-l Va t) = ii) apply simp
    apply (metis set-bit-prev-len set-bit-prev-len2)
  by (metis set-bit-prev-len set-bit-prev-len2)

```

```

lemma inserts-comm:
inserts ilst lst @ [v] = inserts (ilst @ [v]) lst
  by (simp add: inserts-def)

```

```

lemma mp-alloc-stm4-while-isucc'':
nat (from-l Vb t + 1) < length (levels (mem-pool-info Vb p)) ⇒
  V ∈ mp-alloc-stm4-loopinv Vb t p ∩ {i | i < 4} ⇒
  i V t > 0 ⇒
  Γ ⊢I Some (i := i(t := Suc (i t))) satp
  [{mp-alloc-stm4-while-precond5
    (mp-alloc-stm4-while-precond4
      (mp-alloc-stm4-while-precond3
        (mp-alloc-stm4-while-precond2
          (mp-alloc-stm4-while-precond1 V t p) t p) t p) t p)},
    {(s, t). s = t}, UNIV, mp-alloc-stm4-loopinv Vb t p]
apply(rule Basic)

```

```

apply clarsimp
apply(rule conjI)
  apply(simp add:gvars-conf-stable-def gvars-conf-def) apply clarsimp
  apply(rule conjI) apply clarsimp
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply clarsimp apply(simp add:append-free-list-def set-bit-def)
    apply(case-tac ia ≠ nat (from-l Vb t + 1)) apply auto[1]
    apply(case-tac ia < length (levels (mem-pool-info Vb p))) apply fastforce
apply fastforce
  apply clarsimp

```

```

    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply(rule conjI) apply(simp add:append-free-list-def set-bit-def)
    apply clarsimp apply(simp add:append-free-list-def set-bit-def)
  apply(rule conjI)
    apply(simp add:append-free-list-def set-bit-def)
  apply(rule conjI)
    apply(simp add:append-free-list-def set-bit-def)
  apply(rule conjI)
    apply clarsimp
    apply(simp add:append-free-list-def set-bit-def lvars-nochange-def)
  apply(rule conjI)
    apply(simp add:append-free-list-def set-bit-def)
  apply(subgoal-tac length (levels (mem-pool-info Vb p)) = length (levels (mem-pool-info
    V p)))
    prefer 2 apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac nat (from-l Vb t + 1) < length (levels (mem-pool-info Vb p)))
    prefer 2 apply(simp add:inv-mempool-info-def)
  apply(rule conjI)
    apply(simp add:append-free-list-def set-bit-def)
    using lst-updts-eq-updts-updt[of i V t bits (levels (mem-pool-info Vb p) ! nat
      (from-l Vb t + 1))
      Suc (bn Vb t * 4) FREE]
    apply (simp add: semiring-normalization-rules(7))
    apply(simp add:append-free-list-def set-bit-def)
    using inserts-comm apply fast

```

**apply fast using stable-id2 apply fast using stable-id2 apply fast**  
**done**

**lemma** mp-alloc-stm4-while-isucc:

$Va \in \text{mp-alloc-precond2-1-1-loopinv-0 } t \ p \ \text{sz timeout} \cap \{\text{'cur} = \text{Some } t\} \implies$   
 $V \in \text{mp-alloc-stm4-loopinv } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ p \cap \{\text{'i } t <$   
 $4\} \implies$

$\Gamma \vdash_I \text{Some } (\text{'i} := \text{'i}(t := \text{Suc } (\text{'i } t))) \text{ sat}_p$   
 $[\{\text{mp-alloc-stm4-while-precond5}$   
 $(\text{mp-alloc-stm4-while-precond4}$   
 $(\text{mp-alloc-stm4-while-precond3}$   
 $(\text{mp-alloc-stm4-while-precond2}$   
 $(\text{mp-alloc-stm4-while-precond1 } V \ t \ p) \ t \ p) \ t \ p) \ t \ p\},$   
 $\{(s, t). s = t\}, \text{UNIV}, \text{mp-alloc-stm4-loopinv } (\text{mp-alloc-stm4-pre-precond-f } Va$   
 $t \ p) \ t \ p]$

**apply**(rule mp-alloc-stm4-while-isucc')

**apply** clarsimp

**apply**(rule subst[where s=from-l Va t and t=from-l (mp-alloc-stm4-pre-precond-f  
 Va t p) t])

**using** mp-alloc-stm4-pre-froml **apply** blast

```

apply(rule subst[where  $s = \text{length} (\text{levels} (\text{mem-pool-info } Va \ p))$  and
 $t = \text{length} (\text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f } Va \ t$ 
 $p) \ p))$ ])
  using mp-alloc-stm4-inv-mif-len apply blast
apply(subgoal-tac n-levels ( $\text{mem-pool-info } Va \ p = \text{length} (\text{levels} (\text{mem-pool-info}$ 
 $Va \ p))$ ))
  prefer 2 apply(simp add:inv-def inv-mempool-info-def) apply metis
apply linarith

apply assumption

apply clarsimp
done

lemma mp-alloc-stm4-while-help:
   $Va \in \text{mp-alloc-precond2-1-1-loopinv-0 } t \ p \ sz \ \text{timeout} \cap \{\!| \text{'cur} = \text{Some } t \|\} \implies$ 
 $V \in \text{mp-alloc-stm4-loopinv} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ p \cap \{\!| \text{'i } t <$ 
 $4 \|\} \implies$ 
 $\Gamma \vdash_I \text{Some} (\text{'lbn} := \text{'lbn}(t := 4 * \text{'bn } t + \text{'i } t));$ 
 $\text{'lsz} := \text{'lsz}(t := \text{'lsizes } t ! \text{nat} (\text{'from-l } t + 1));$ 
 $\text{'block2} := \text{'block2}(t := \text{'lsz } t * \text{'i } t + \text{'blk } t);$ 
 $\text{'mem-pool-info} := \text{set-bit-free } \text{'mem-pool-info } p (\text{nat} (\text{'from-l } t + 1)) (\text{'lbn}$ 
 $t);$ 
  IF block-fits ( $\text{'mem-pool-info } p$ ) ( $\text{'block2 } t$ ) ( $\text{'lsz } t$ ) THEN
 $\text{'mem-pool-info} := \text{'mem-pool-info}(p := \text{append-free-list} (\text{'mem-pool-info } p)$ 
 $(\text{nat} (\text{'from-l } t + 1)) (\text{'block2 } t))$ 
  FI;
 $\text{'i} := \text{'i}(t := \text{Suc} (\text{'i } t))$ 
  satp [{V}, {(s, t). s = t}, UNIV, mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f
  Va t p) t p]
apply(rule Seq[where mid={mp-alloc-stm4-while-precond5
  (mp-alloc-stm4-while-precond4
  (mp-alloc-stm4-while-precond3
  (mp-alloc-stm4-while-precond2
  (mp-alloc-stm4-while-precond1 V t p) t p) t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-while-precond4
  (mp-alloc-stm4-while-precond3
  (mp-alloc-stm4-while-precond2
  (mp-alloc-stm4-while-precond1 V t p) t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-while-precond3
  (mp-alloc-stm4-while-precond2
  (mp-alloc-stm4-while-precond1 V t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-while-precond2
  (mp-alloc-stm4-while-precond1 V t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-while-precond1 V t p}])

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

```

```

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Cond)
apply(simp add:stable-def)
apply(rule Basic)
using mp-alloc-stm4-pre-in apply blast apply simp apply(simp add:stable-def)
apply(simp add:stable-def)

apply(unfold Skip-def)
apply(rule subst where  $t = \{mp\text{-}alloc\text{-}stm4\text{-}while\text{-}precond4$ 
  (mp-alloc-stm4-while-precond3
  (mp-alloc-stm4-while-precond2
  (mp-alloc-stm4-while-precond1  $V\ t\ p$ )  $t\ p$ )  $t\ p$ )  $t\ p\}$   $\cap$ 
   $-\ \{\!\!\{block\text{-}fits\ (\text{'mem-pool-info}\ p)\ (\text{'block2}\ t)\ (\text{'lsz}\ t)\}\!\!\}$  and  $s = \{\!\!\{\}\!\!\}$ ])
using mp-alloc-stm4-blockfit[of  $Va\ t\ p\ timeout\ sz\ V$ ] apply metis
using Emptyprecond apply metis
apply simp

using mp-alloc-stm4-while-isucc[of  $Va\ t\ p\ timeout\ sz\ V$ ] apply fast
done

lemma mp-alloc-stm4-while-1:  $\{\!\!\{4 \leq \text{'i}\ t\}\!\!\} = -\ \{\!\!\{\text{'i}\ t < 4\}\!\!\}$  by auto

term mp-alloc-precond2-1-1-loopinv-0  $t\ p\ sz\ timeout \cap \{\!\!\{\text{'cur} = Some\ t\}\!\!\}$ 

lemma mp-alloc-stm4-while:
   $Va \in mp\text{-}alloc\text{-}precond2\text{-}1\text{-}1\text{-}loopinv\text{-}0\ t\ p\ sz\ timeout \cap \{\!\!\{\text{'cur} = Some\ t\}\!\!\} \implies$ 
   $\Gamma \vdash_I Some\ (WHILE\ \text{'i}\ t < 4\ DO$ 
     $\text{'lbn} := \text{'lbn}\ (t := 4 * \text{'bn}\ t + \text{'i}\ t);;$ 
     $\text{'lsz} := \text{'lsz}\ (t := (\text{'lsizes}\ t) ! (nat\ (\text{'from-l}\ t + 1)));;$ 
     $\text{'block2} := \text{'block2}\ (t := \text{'lsz}\ t * \text{'i}\ t + \text{'blk}\ t);;$ 

     $(* set\text{-}free\text{-}bit(p, l + 1, lbn); *)$ 
     $\text{'mem-pool-info} := set\text{-}bit\text{-}free\ \text{'mem-pool-info}\ p\ (nat\ (\text{'from-l}\ t + 1))\ (\text{'lbn}$ 
   $t);;$ 

  IF block-fits ('mem-pool-info p) ('block2 t) ('lsz t) THEN

     $(* sys\text{-}dlist\text{-}append(\&p \rightarrow levels[l + 1].free\text{-}list, block2); *)$ 
     $\text{'mem-pool-info} := \text{'mem-pool-info}\ (p :=$ 

```

```

    append-free-list ('mem-pool-info p) (nat ('from-l t + 1)) ('block2
t) )
    FI;;
    'i := 'i(t := Suc ('i t))
    OD) satp [mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p, {(s,
t). s = t}, UNIV,
    mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p ∩ {('i t
≥ 4}]]
  apply(rule While)
  apply(simp add:stable-def)
  apply(rule subst[where t = - {('i t < 4} and s = {4 ≤ 'i t}]) using mp-alloc-stm4-while-1[of
t] apply simp
  apply simp
  apply(simp add:stable-def)
  using mp-alloc-stm4-while-help[of Va t p timeout sz ]
  Allprecond[of mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p
∩ {('i t < 4}]]
  Some ('lbn := 'lbn(t := 4 * 'bn t + 'i t));
  'lsz := 'lsz(t := 'lsizes t ! nat ('from-l t + 1));
  'block2 := 'block2(t := 'lsz t * 'i t + 'blk t);
  'mem-pool-info := set-bit-free 'mem-pool-info p (nat ('from-l t + 1)) ('lbn
t));
  IF block-fits ('mem-pool-info p) ('block2 t)
  ('lsz t) THEN 'mem-pool-info := 'mem-pool-info
  (p := append-free-list ('mem-pool-info p) (nat ('from-l t +
1)) ('block2 t)) FI;;
  'i := 'i(t := Suc ('i t)) {(s, t). s = t} UNIV
  mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p] apply
clarsimp
  apply force
done

```

**lemma** *mp-alloc-stm4-pre-precond-f-in-mp-alloc-stm4-loopinv*:  
 $mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p \in mp\text{-}alloc\text{-}stm4\text{-}loopinv\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t\ p$   
 apply(subgoal-tac i (mp-alloc-stm4-pre-precond-f Va t p) t = 1)  
 using in-mp-alloc-stm4-loopinv apply meson  
 apply(simp add:mp-alloc-stm4-pre-precond-f-def)  
done

**lemma** *mp-alloc-stm4-mempools*:  $(x, mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p) \in gvars\text{-}conf\text{-}stable$   
 $\implies mem\text{-}pools\ x = mem\text{-}pools\ Va$   
 by (simp add:mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def)

**lemma** *mp-alloc-stm4-mempools2*:  $mem\text{-}pools\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p) = mem\text{-}pools\ x \implies mem\text{-}pools\ x = mem\text{-}pools\ Va$   
 by (simp add:mp-alloc-stm4-pre-precond-f-def)



**lemma** *mp-alloc-stm4-inv-cur*:

*cur*  $Va = \text{Some } t \implies \forall ta. (t = ta) = (\text{thd-state } Va \text{ } ta = \text{RUNNING}) \implies$   
 $(\text{cur } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) = \text{Some } ta) = (\text{thd-state } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } ta = \text{RUNNING})$   
**by** (*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-inv-thd-state*:  $(x, \text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \in \text{gvars-conf-stable} \implies$

$\text{thd-state } x = \text{thd-state } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \implies$   
 $\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \text{ } pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } pa \implies$   
 $\text{wait-q } (\text{mem-pool-info } x \text{ } p) = \text{wait-q } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } p) \implies$   
 $\text{inv-thd-waitq } Va \implies \text{inv-thd-waitq } x$   
**apply** (*subgoal-tac*  $\forall q \in \text{mem-pools } x. \text{wait-q } (\text{mem-pool-info } x \text{ } q)$   
 $= \text{wait-q } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } q)$ )  
**prefer 2 apply clarify apply** (*case-tac*  $q = p$ ) **apply simp apply simp**  
**apply** (*subgoal-tac*  $\forall q \in \text{mem-pools } Va. \text{wait-q } (\text{mem-pool-info } Va \text{ } q)$   
 $= \text{wait-q } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } q)$ )  
**prefer 2 apply clarify apply** (*simp add:mp-alloc-stm4-pre-precond-f-def*)  
**apply** (*simp add: set-bit-def*)  
**apply** (*subgoal-tac*  $\text{thd-state } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) = \text{thd-state } Va$ )  
**prefer 2 apply** (*simp add:mp-alloc-stm4-pre-precond-f-def*)  
**apply** (*subgoal-tac*  $\text{mem-pools } x = \text{mem-pools } Va$ )  
**prefer 2 apply** (*simp add:mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def*  
*gvars-conf-def*)  
**apply** (*simp add:inv-thd-waitq-def*)  
**apply clarify**  
**by blast**

**lemma** *inv-mpinfo-inv-mpinfo-stm4*:

*inv-mempool-info*  $Va \implies \text{inv-mempool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p)$   
**apply** (*simp add:inv-mempool-info-def mp-alloc-stm4-pre-precond-f-def*)  
**apply** (*simp add:Let-def*) **apply clarify**  
**apply** (*rule conjI*)  
**apply** (*simp add: set-bit-def*) **apply auto**[1]  
**apply** (*rule conjI*)  
**apply** (*simp add: set-bit-def*)  
**apply** (*rule conjI*) **apply clarify apply auto**[1] **apply clarify apply auto**[1]  
**apply** (*rule conjI*) **apply** (*simp add: set-bit-def*) **apply auto**[1]  
**apply** (*rule conjI*) **apply** (*simp add: set-bit-def*) **apply auto**[1]  
**apply** (*rule conjI*) **apply** (*simp add: set-bit-def*) **apply auto**[1]  
**apply clarify apply** (*rename-tac*  $pa \text{ } ii$ )  
**apply** (*subgoal-tac*  $\text{length } (\text{bits } (\text{levels } ((\text{set-bit-divide } (\text{mem-pool-info } Va) \text{ } p) (\text{nat } (\text{from-l } Va \text{ } t))))$   
 $(\text{block-num } (\text{mem-pool-info } Va \text{ } p) (\text{blk } Va \text{ } t) (\text{lsizes } Va \text{ } t ! \text{ nat } (\text{from-l } Va \text{ } t)))) \text{ } pa) !$   
 $ii)) = \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \text{ } pa) ! ii)))$ )

```

prefer 2 apply (metis (no-types, lifting) fun-upd-apply set-bit-def
  set-bit-prev-len set-bit-prev-len2)
apply(subgoal-tac length (bits (levels (set-bit-allocating
  (set-bit-divide (mem-pool-info Va) p (nat (from-l Va t))
    (block-num (mem-pool-info Va p) (blk Va t) (lsizes
  Va t ! nat (from-l Va t))))))
    p (nat (from-l Va t + 1))
    (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes
  Va t ! nat (from-l Va t)))) pa) !
  ii)) = length (bits (levels ((set-bit-divide (mem-pool-info Va)
  p (nat (from-l Va t))
    (block-num (mem-pool-info Va p) (blk Va t) (lsizes
  Va t ! nat (from-l Va t)))))) pa) !
  ii)))
prefer 2 apply(case-tac ii = nat (from-l Va t + 1))
using set-bit-prev-len set-bit-def apply auto[1]
using set-bit-def apply auto[1]
apply(subgoal-tac n-max (set-bit-allocating
  (set-bit-divide (mem-pool-info Va) p (nat (from-l Va t))
    (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat
  (from-l Va t))))))
  p (nat (from-l Va t + 1)) (4 * block-num (mem-pool-info Va p)
  (blk Va t) (lsizes Va t ! nat (from-l Va t)))) pa) *
  4 ^ ii = n-max (mem-pool-info Va pa) * 4 ^ ii)
prefer 2 apply(simp add: set-bit-def)
apply(subgoal-tac length (levels (set-bit-allocating
  (set-bit-divide (mem-pool-info Va) p (nat (from-l Va t))
    (block-num (mem-pool-info Va p) (blk Va t) (lsizes
  Va t ! nat (from-l Va t))))))
    p (nat (from-l Va t + 1))
    (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes
  Va t ! nat (from-l Va t)))) pa))
  = length (levels (mem-pool-info Va pa)))
prefer 2 apply(simp add: set-bit-def)
apply metis
done

```

```

lemma mp-alloc-stm4-inv-mempool-info:
  (x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable ⇒
    inv-mempool-info Va ⇒ inv-mempool-info x
apply(subgoal-tac inv-mempool-info (mp-alloc-stm4-pre-precond-f Va t p))
prefer 2 using inv-mpinfo-inv-mpinfo-stm4 apply simp
using gvars-conf-stb-inv-mpinf apply simp
done

```

```

lemma mp-alloc-stm4-lvl-len:
  p ∈ mem-pools Va ⇒ (x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable
  ⇒

```

$length\ (levels\ (mem\ pool\ info\ x\ pa)) = length\ (levels\ (mem\ pool\ info\ Va\ pa))$   
**apply**(simp add: mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def)

**apply**(simp add: set-bit-def)  
**done**

**lemma** mp-alloc-stm4-maxsz:  
 $p \in mem\ pools\ Va \implies (x, mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p) \in gvars\ conf\ stable$   
 $\implies$   
 $max\ sz\ (mem\ pool\ info\ x\ pa) = max\ sz\ (mem\ pool\ info\ Va\ pa)$   
**apply**(simp add: mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def)

**apply**(simp add: set-bit-def)  
**done**

**lemma** mp-alloc-stm4-buf:  
 $p \in mem\ pools\ Va \implies (x, mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p) \in gvars\ conf\ stable$   
 $\implies$   
 $buf\ (mem\ pool\ info\ x\ pa) = buf\ (mem\ pool\ info\ Va\ pa)$   
**apply**(simp add: mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def)

**apply**(simp add: set-bit-def)  
**done**

**lemma** mp-alloc-stm4-pres-mpinfo:  
 $pa \neq p \longrightarrow mem\ pool\ info\ Va\ pa = mem\ pool\ info\ (mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p)\ pa$   
**apply**(simp add: mp-alloc-stm4-pre-precond-f-def set-bit-def)  
**done**

**lemma** mp-alloc-stm4-froml:  
 $from\ l\ x = from\ l\ (mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p) \implies$   
 $from\ l\ x = from\ l\ Va$   
**apply**(simp add: mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def)

**done**

**lemma** mp-alloc-stm4-pre-precond-f-lvars-nochange:  
 $t' \neq t \implies lvars\ nochange\ t'\ Va\ (mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p)$   
**apply**(simp add: lvars-nochange-def mp-alloc-stm4-pre-precond-f-def)  
**done**

**lemma** mp-alloc-stm4-pre-precond-f-tick:  
 $tick\ Va = tick\ (mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p)$   
**by**(simp add: mp-alloc-stm4-pre-precond-f-def)

**lemma** mp-alloc-stm4-pre-precond-f-def-frnode:  
 $freeing\ node\ Va = freeing\ node\ (mp\ alloc\ stm4\ pre\ precondition\ f\ Va\ t\ p)$

**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-mpls:*

$p \in \text{mem-pools } Va \implies (x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable}$   
 $\implies p \in \text{mem-pools } x$

**apply**(*simp add:mp-alloc-stm4-pre-precond-f-def gvars-conf-stable-def gvars-conf-def*)  
**done**

**lemma** *mp-alloc-stm4-pre-precond-f-rf:*

$\text{rf } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \implies \text{rf } Va \ t$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-ret:*

$\text{mempoolalloc-ret } Va = \text{mempoolalloc-ret } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-tmout:*

$\text{tmout } Va = \text{tmout } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-lsz:*

$\text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) = \text{lsizes } Va$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-alloc-l:*

$\text{alloc-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) = \text{alloc-l } Va$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-from-l:*

$\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) = \text{from-l } Va$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-free-l:*

$\text{free-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) = \text{free-l } Va$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *mp-alloc-stm4-pre-precond-f-blk:*

$\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) = \text{blk } Va$   
**by**(*simp add:mp-alloc-stm4-pre-precond-f-def*)

**lemma** *same-level-set-bit-l:i1≠i' ⟹*

$\text{levels } ((\text{set-bit } \text{mp-info } p \ i' \ j' \ b) \ p)!i1 = \text{levels } (\text{mp-info } p)!i1$

**unfolding** *set-bit-def*

**by** *auto*

**lemma** *same-bit-set-bit-l:i1≠i' ⟹*

$\text{bits } (\text{levels } ((\text{set-bit } \text{mp-info } p \ i' \ j' \ b) \ p)!i1) = \text{bits } (\text{levels } (\text{mp-info } p)!i1)$

```

using same-level-set-bit-l
by auto

lemma same-bit-set-bit-j:
   $j1 \neq j' \implies$ 
  bits (levels ((set-bit mp-info p i' j' b) p)!i1)!j1 = bits (levels (mp-info p)!i1)!j1
  apply (simp add: set-bit-get-bit-neq set-bit-def)
  by (metis (no-types, lifting) Mem-pool-lvl.select-convs(1) Mem-pool-lvl.surjective
  Mem-pool-lvl.update-convs(1)
  list-update-beyond not-less nth-list-update-eq nth-list-update-neq)

lemma set-bit-set-bit:
   $l1 \neq l2 \vee b1 \neq b2 \implies$ 
  set-bit-s (set-bit-s Va p l1 b1 st1) p l2 b2 st2 =
  set-bit-s ((set-bit-s Va p l2 b2 st2)) p l1 b1 st1
  unfolding set-bit-s-def set-bit-def
  apply auto
  apply (cases b1=b2) apply auto
  apply (simp add: list-update-swap)
  apply (simp add: list-update-swap)
  apply (cases l1=l2) apply auto
  apply (cases l1 < length (levels (mem-pool-info Va p)))
  by (auto simp add: list-update-swap)

lemma get-bit-set-bit-set-bit:
  assumes a0:l1 ≠ l2 ∨ b1 ≠ b2 and
    a1:l1 < length (levels ((mem-pool-info Va) p)) and
    a2:b1 < length (bits (levels ((mem-pool-info Va) p)!l1))
  shows get-bit-s (set-bit-s (set-bit-s Va p l1 b1 st1) p l2 b2 st2) p l1 b1 = st1
proof–
  have a1':l1 < length (levels ((mem-pool-info (set-bit-s Va p l2 b2 st2)) p))
  using a1 unfolding set-bit-s-def set-bit-def by auto
  have a2':b1 < length (bits (levels ((mem-pool-info (set-bit-s Va p l2 b2 st2)) p)
  !l1))
  using a2 unfolding set-bit-s-def set-bit-def apply auto
  by (metis (no-types, lifting) Mem-pool-lvl.select-convs(1) Mem-pool-lvl.surjective
  Mem-pool-lvl.update-convs(1) a1
  length-list-update nth-list-update-eq nth-list-update-neq)
  show ?thesis using set-bit-get-bit-eq2[OF a1' a2'] set-bit-set-bit[OF a0] unfold-
ing set-bit-s-def by auto
qed

lemma mp-alloc-stm4-pre-precond-f-same-bits:assumes
  a0:i1=(nat (from-l Va t)) and
  a1:i2 = (nat (from-l Va t + 1)) and
  a2:Va' = mp-alloc-stm4-pre-precond-f Va t p
  shows  $\forall i j. \neg((i=i1) \vee (i=i2)) \longrightarrow$ 

```

$get-bit (mem-pool-info Va') p i j = get-bit (mem-pool-info Va) p i j$   
**using**  $a0 a1 a2$  *set-bit-get-bit-neq* **unfolding** *mp-alloc-stm4-pre-precond-f-def*  
**by** *auto*

**lemma** *same-bit-mp-alloc-stm4-pre-precond-f*:  
 $i1 = (nat (from-l Va t)) \implies$   
 $j1 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))) \implies$   
 $i2 = (nat (from-l Va t + 1)) \implies$   
 $j2 = (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))) \implies$   
 $Va' = mp-alloc-stm4-pre-precond-f Va t p \implies$   
 $\forall i j. \neg((i=i1 \wedge j=j1) \vee (i=i2 \wedge j=j2)) \longrightarrow$   
 $get-bit (mem-pool-info Va') p i j = get-bit (mem-pool-info Va) p i j$   
**using** *set-bit-get-bit-neq*  
**apply** (*auto simp add:mp-alloc-stm4-pre-precond-f-def*)  
**by** *metis+*

**lemma** *same-bit-mp-alloc-stm4-pre-precond-f1*:  
**assumes**  
 $a1: \neg((l = (nat (from-l Va t)) \wedge b = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t)))) \vee$   
 $(l = (nat (from-l Va t + 1)) \wedge b = (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t)))))$   
**shows**  $(get-bit-s Va p l b = get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p l b)$   
**using**  $a1$  *same-bit-mp-alloc-stm4-pre-precond-f* **by** *metis*

**lemma** *same-bit-mp-alloc-stm4-pre-precond-f11*:  
**assumes**  $a0: (l = (nat (from-l Va t)) \wedge b = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))))$  **and**  
 $a1: l \geq length (levels (mem-pool-info Va p)) \vee$   
 $b \geq length (bits (levels (mem-pool-info Va p) ! l))$   
**shows**  $get-bit-s Va p l b = get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p l b$   
**using**  $a0 a1$  **unfolding** *mp-alloc-stm4-pre-precond-f-def set-bit-def*  
**apply** *auto*  
**by** (*metis (no-types, lifting) Mem-pool-lvl.simps(1) Mem-pool-lvl.simps(4) Mem-pool-lvl.surjective*)

*list-update-beyond not-less nth-list-update-eq nth-list-update-neq*

**lemma** *same-bit-mp-alloc-stm4-pre-precond-f12*:  
**assumes**  $a0: (l = (nat (from-l Va t + 1)) \wedge b = 4 * (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))))$  **and**  
 $a1: l \geq length (levels (mem-pool-info Va p)) \vee$   
 $b \geq length (bits (levels (mem-pool-info Va p) ! l))$   
**shows**  $get-bit-s Va p l b = get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p l b$   
**using**  $a0 a1$  **unfolding** *mp-alloc-stm4-pre-precond-f-def set-bit-def*  
**apply** *auto*

**apply** (*metis* *list-update-beyond nth-list-update-neq*)  
**by** (*smt* *Mem-pool-lvl.simps*(1) *Mem-pool-lvl.simps*(4) *Mem-pool-lvl.surjective*  
*length-list-update list-update-beyond*  
*not-less nth-list-update-eq nth-list-update-neq*)

**lemma** *same-bit-mp-alloc-stm4-pre-precond-f2*:

**assumes**  $a1:l \geq \text{length } (\text{levels } (\text{mem-pool-info } Va \ p)) \vee$   
 $b \geq \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ l))$   
**shows**  $\text{get-bit-s } Va \ p \ l \ b = \text{get-bit-s } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p$   
 $l \ b$   
**apply** (*cases*  $\neg ((l = (\text{nat } (\text{from-l } Va \ t + 1))) \wedge$   
 $b = 4 * (\text{block-num } (\text{mem-pool-info } Va \ p) (\text{blk } Va \ t) (\text{lsizes } Va \ t \ ! \ \text{nat}$   
 $(\text{from-l } Va \ t)))) \vee$   
 $(l = (\text{nat } (\text{from-l } Va \ t)) \wedge b = (\text{block-num } (\text{mem-pool-info } Va \ p) (\text{blk } Va$   
 $t) (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))))) \vee$   
**using** *same-bit-mp-alloc-stm4-pre-precond-f1* **apply** *fast*  
**using**  $a1$  **by** (*auto* *intro*: *same-bit-mp-alloc-stm4-pre-precond-f11* *same-bit-mp-alloc-stm4-pre-precond-f12*)

**lemma** *same-bit-mp-alloc-stm4-pre-precond-divided*:

**assumes**  $a0:(l = (\text{nat } (\text{from-l } Va \ t)) \wedge b = (\text{block-num } (\text{mem-pool-info } Va \ p) (\text{blk}$   
 $Va \ t) (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t))))$  **and**  
 $a1:l < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p))$  **and**  
 $a2:b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ l))$  **and**  
 $a3:(\text{from-l } Va \ t) \geq 0$   
**shows**  $\text{get-bit-s } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p \ l \ b = \text{DIVIDED}$   
**proof**–  
**have**  $l \neq \text{nat } (\text{from-l } Va \ t + 1) \vee b \neq 4 * b$  **using**  $a0 \ a3$  **by** *fastforce*  
**then show** *?thesis* **using**  $a0 \ a1 \ a2 \ a3$  *set-bit-get-bit-eq2*  
**unfolding** *mp-alloc-stm4-pre-precond-f-def*  
**using** *set-bit-get-bit-neq* **by** *auto*  
**qed**

**lemma** *same-bit-mp-alloc-stm4-pre-precond-allocating*:

**assumes**  $a0:(l = (\text{nat } (\text{from-l } Va \ t + 1)) \wedge b = 4 * (\text{block-num } (\text{mem-pool-info } Va$   
 $p) (\text{blk } Va \ t) (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t))))$  **and**  
 $a1:l < \text{length } (\text{levels } (\text{mem-pool-info } Va \ p))$  **and**  
 $a2:b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ l))$  **and**  
 $a3:(\text{from-l } Va \ t) \geq 0$   
**shows**  $\text{get-bit-s } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p \ l \ b = \text{ALLOCATING}$   
**proof**–  
**let**  $?Va = \text{set-bit-s } Va \ p \ (\text{nat } (\text{from-l } Va \ t))$   
 $(\text{block-num } (\text{mem-pool-info } Va \ p) (\text{blk } Va \ t) (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va$   
 $t))) \text{ DIVIDED}$   
**have**  $a1':l < \text{length } (\text{levels } (\text{mem-pool-info } ?Va \ p))$   
**using**  $a1$  **unfolding** *set-bit-s-def* *set-bit-def* **by** *auto*  
**moreover have**  $a2':b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } ?Va \ p) \ ! \ l))$   
**using**  $a2$  **unfolding** *set-bit-s-def* *set-bit-def*  
**by** (*simp* *add*:  $a0 \ a3$  *eq-nat-nat-iff*)  
**ultimately show** *?thesis*

**using**  $a0$  *set-bit-get-bit-eq2*  $a3$   
**unfolding** *mp-alloc-stm4-pre-precond-f-def set-bit-s-def*  
**using** *set-bit-get-bit-eq* **by** *auto*  
**qed**

**lemma** *eq-free-list-set-bit-l*:  
 $free-list (levels ((set-bit mp-info p i' j' b) p)!i1) = free-list (levels (mp-info p)!i1)$   
**unfolding** *set-bit-def*  
**apply** (*cases i' < length (levels (mp-info p)); auto*)  
**by** (*metis (no-types, lifting) Mem-pool-lvl.select-convs(2) Mem-pool-lvl.surjective Mem-pool-lvl.update-convs(1) nth-list-update-eq nth-list-update-neq*)

**lemma** *eq-free-list-mp-alloc-stm4-pre-precond-f*:  
 $free-list (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) ! l) =$   
 $free-list (levels (mem-pool-info Va p) ! l)$   
**unfolding** *mp-alloc-stm4-pre-precond-f-def* **using** *eq-free-list-set-bit-l*  
**by** *auto*

**lemma** *mp-alloc-stm4-pre-precond-f-i*:  $(i (mp-alloc-stm4-pre-precond-f Va t p)) t =$   
 $Suc\ 0 \wedge$   
 $(\forall t'. t' \neq t \longrightarrow (i (mp-alloc-stm4-pre-precond-f Va t p)) t' = (i Va) t')$   
**unfolding** *mp-alloc-stm4-pre-precond-f-def* **by** *force*

**lemma** *mp-alloc-stm4-pre-precond-f-bn*:  
 $(bn (mp-alloc-stm4-pre-precond-f Va t p)) t =$   
 $block-num (mem-pool-info Va p) (blk Va t) ((lsizes Va t)!(nat (from-l Va t)))$   
**unfolding** *mp-alloc-stm4-pre-precond-f-def* **by** *force*

**lemma** *mp-alloc-stm4-pre-precond-f-allocating*:  
 $(allocating-node (mp-alloc-stm4-pre-precond-f Va t p)) t =$   
 $Some (\downarrow pool = p, level = nat (from-l Va t + 1),$   
 $block = 4 * block-num (mem-pool-info Va p) (blk Va t) ((lsizes Va t)!(nat$   
 $(from-l Va t))),$   
 $data = blk Va t \downarrow)$   
**unfolding** *mp-alloc-stm4-pre-precond-f-def*  
**by** *auto*

**lemma** *get-bit-x-l-b*:  
**assumes**  $a0: (l = (nat (from-l (Va::State) t)) \wedge b = (block-num (mem-pool-info$   
 $Va p) (blk Va t) (lsizes Va t ! nat (from-l Va t))))$  **and**  
 $a1: (from-l Va t) \geq 0$  **and**  
 $a2: \forall jj. jj \neq nat (from-l Va t + 1) \longrightarrow$   
 $levels (mem-pool-info x p) ! jj =$   
 $levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) ! jj$  **and**  
 $a4: l < length (levels (mem-pool-info Va p))$  **and**  
 $a5: b < length (bits (levels (mem-pool-info Va p) ! l))$   
**shows**  $get-bit-s x p l b = DIVIDED$



**using** *a0 a2 a1 a4 a5 same-bit-mp-alloc-stm4-pre-precond-divided* **by** *auto*

**lemma** *get-bit-x-l1-b4*:

**assumes** *a0*: $l = (\text{nat } (\text{from-l } (Va::\text{State}) t + 1)) \wedge b = 4 * (\text{block-num } (\text{mem-pool-info } Va p) (\text{blk } Va t) (\text{lsizes } Va t ! \text{nat } (\text{from-l } Va t))))$  **and**  
*a1*: $(\text{from-l } Va t) \geq 0$  **and**  
*a3*: $\text{bits } (\text{levels } (\text{mem-pool-info } x p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t + 1)) =$   
 $\text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t + 1)))$   
 $(\text{Suc } (bn (\text{mp-alloc-stm4-pre-precond-f } Va t p) t * 4)) \ 3 \text{ FREE and}$   
*a4*: $l < \text{length } (\text{levels } (\text{mem-pool-info } Va p))$  **and**  
*a5*: $b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va p) ! l))$   
**shows** *get-bit-s x p l b = ALLOCATING*  
**using** *a0 a1 a3 a4 a5 same-bit-mp-alloc-stm4-pre-precond-allocating*  
*mp-alloc-stm4-pre-precond-f-bn*  
*mp-alloc-stm4-pre-froml*  
**by** (*metis lessI list-updates-n-neq mult.commute*)

**lemma** *get-bit-x-l1-b41*:

**assumes** *a0*: $l = (\text{nat } (\text{from-l } (Va::\text{State}) t + 1)) \wedge$   
 $(b = \text{Suc}(4 * (\text{block-num } (\text{mem-pool-info } Va p) (\text{blk } Va t) (\text{lsizes } Va t !$   
 $\text{nat } (\text{from-l } Va t)))) \vee$   
 $b = \text{Suc}(\text{Suc}(4 * (\text{block-num } (\text{mem-pool-info } Va p) (\text{blk } Va t) (\text{lsizes } Va$   
 $t ! \text{nat } (\text{from-l } Va t)))) \vee$   
 $b = \text{Suc}(\text{Suc}(\text{Suc}(4 * (\text{block-num } (\text{mem-pool-info } Va p) (\text{blk } Va t) (\text{lsizes } Va$   
 $t ! \text{nat } (\text{from-l } Va t))))))$  **and**  
*a1*: $(\text{from-l } Va t) \geq 0$  **and**  
*a2*: $\forall jj. jj \neq \text{nat } (\text{from-l } Va t + 1) \longrightarrow$   
 $\text{levels } (\text{mem-pool-info } x p) ! jj =$   
 $\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) p) ! jj$  **and**  
*a3*: $\text{bits } (\text{levels } (\text{mem-pool-info } x p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t + 1)) =$   
 $\text{list-updates-n } (\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va t p) p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va t p) t + 1)))$   
 $(\text{Suc } (bn (\text{mp-alloc-stm4-pre-precond-f } Va t p) t * 4)) \ 3 \text{ FREE and}$   
*a4*: $l < \text{length } (\text{levels } (\text{mem-pool-info } Va p))$  **and**  
*a5*: $b < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va p) ! l))$   
**shows** *get-bit-s x p l b = FREE*  
**using** *a0 a1 a3 a4 a5*  
**apply** *auto*  
**using** *mp-alloc-stm4-pre-precond-f-bn*  
*mp-alloc-stm4-pre-froml mp-alloc-stm4-inv-bits-len Suc-numeral add-2-eq-Suc*  
*add-Suc-right*  
**by** (*smt add.commute lessI less-add-same-cancel2*)

*list-updates-n-eq mult.commute nat-less-le neq0-conv not-le semiring-norm(5))+*

**lemma** *get-bit-x-l1-b41'*:

**assumes**  $a0:l=(nat\ (from-l\ (Va::State)\ t + 1)) \wedge$   
 $(b=(4*(block-num\ (mem-pool-info\ Va\ p)\ (blk\ Va\ t)\ (lsizes\ Va\ t\ !\ nat$   
 $(from-l\ Va\ t)))) + 1 \vee$   
 $b=(4*(block-num\ (mem-pool-info\ Va\ p)\ (blk\ Va\ t)\ (lsizes\ Va\ t\ !\ nat$   
 $(from-l\ Va\ t))))+2 \vee$   
 $b=(4*(block-num\ (mem-pool-info\ Va\ p)\ (blk\ Va\ t)\ (lsizes\ Va\ t\ !\ nat$   
 $(from-l\ Va\ t))))+3)$  **and**  
 $a1:(from-l\ Va\ t) \geq 0$  **and**  
 $a2:\forall jj. jj \neq nat\ (from-l\ Va\ t + 1) \longrightarrow$   
 $levels\ (mem-pool-info\ x\ p)\ !\ jj =$   
 $levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)\ !\ jj$  **and**  
 $a3:bits\ (levels\ (mem-pool-info\ x\ p)\ !\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f$   
 $Va\ t\ p)\ t + 1)) =$   
 $list-updates-n$   
 $(bits\ (levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)\ !$   
 $nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)))$   
 $(Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4))\ 3\ FREE$  **and**  
 $a4:l < length\ (levels\ (mem-pool-info\ Va\ p))$  **and**  
 $a5:b < length\ (bits\ (levels\ (mem-pool-info\ Va\ p)\ !\ l))$   
**shows** *get-bit-s x p l b = FREE*  
**using**  $a0\ a1\ a3\ a4\ a5$   
**apply** *auto*  
**using** *mp-alloc-stm4-pre-precond-f-bn*  
 $mp-alloc-stm4-pre-froml\ mp-alloc-stm4-inv-bits-len$   
**apply** (*metis add.right-neutral less-not-refl list-updates-n-eq mult.commute nat-add-left-cancel-less*  
*not-less zero-less-numeral*)  
**by** (*smt add.commute add-Suc less-Suc-eq less-add-same-cancel2 list-updates-n-eq*  
*mp-alloc-stm4-inv-bits-len mp-alloc-stm4-pre-froml mp-alloc-stm4-pre-precond-f-bn*  
*mult.commute nat-less-le numeral-3-eq-3*) +

**lemma** *get-bit-x-stm4-pre-eq*:

**assumes**  
 $a0:\forall jj. jj \neq nat\ (from-l\ (Va::State)\ t + 1) \longrightarrow$   
 $levels\ (mem-pool-info\ x\ p)\ !\ jj =$   
 $levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)\ !\ jj$  **and**  
 $a1:bits\ (levels\ (mem-pool-info\ x\ p)\ !\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f$   
 $Va\ t\ p)\ t + 1)) =$   
 $list-updates-n$   
 $(bits\ (levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)\ !$   
 $nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)))$   
 $(Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4))\ 3\ FREE$  **and**  
 $a2:l = nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)$  **and**  
 $a3:b1 = (Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4))$  **and**  
 $a4:b2 = Suc\ (Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4))$  **and**  
 $a5:b3 = Suc\ (Suc\ (Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4)))$   
**shows**  $\forall i\ j. \neg((i=l \wedge j=b1) \vee (i=l \wedge j=b2) \vee (i=l \wedge j=b3)) \longrightarrow$

```

    get-bit-s x p i j = get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p i j
using a0 a1 a2 a3 a4 a5 apply clarsimp
apply (auto simp add: mp-alloc-stm4-pre-precond-f-froml)
by (metis (no-types) add-2-eq-Suc' add-Suc-right eval-nat-numeral(3)
    less-Suc-eq list-updates-n-neq not-less)

lemma same-bit-mp-alloc-stm4-pre-precond-f-x:
assumes a0:  $\forall jj. jj \neq \text{nat (from-l (Va::State) t + 1)} \longrightarrow$ 
    levels (mem-pool-info x p) ! jj =
    levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) ! jj and
    a1: bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f
    Va t p) t + 1)) =
    list-updates-n
    (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))
    (Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4)) 3 FREE and
    a2: i1 = (nat (from-l Va t)) and
    a3: j1 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
    Va t))) and
    a4: i2 = nat (from-l Va t + 1) and
    a5: j2 = (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
    Va t))) and
    a6: j3 = Suc(4 * (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat
    (from-l Va t)))) and
    a7: j4 = Suc(Suc(4 * (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t !
    nat (from-l Va t))))) and
    a8: j5 = Suc(Suc(Suc(4 * (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va
    t ! nat (from-l Va t)))))
shows  $\forall i j. \neg((i=i1 \wedge j=j1) \vee (i=i2 \wedge j=j2) \vee (i=i2 \wedge j=j3) \vee (i=i2 \wedge j=j4) \vee (i=i2 \wedge j=j5)) \longrightarrow$ 
    get-bit-s x p i j = get-bit-s Va p i j
using a0 a1 a2 a4 a5 a6 a7 a8 get-bit-x-stm4-pre-eq
    same-bit-mp-alloc-stm4-pre-precond-f
proof -
  {fix i j
    assume a00:  $\neg((i=i1 \wedge j=j1) \vee (i=i2 \wedge j=j2) \vee (i=i2 \wedge j=j3) \vee (i=i2 \wedge j=j4) \vee (i=i2 \wedge j=j5))$ 
    then have get-bit-s Va p i j =
      get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p i j
    using same-bit-mp-alloc-stm4-pre-precond-f1 a2 a3 a4 a5
    by auto
    also have get-bit-s x p i j =
      get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p i j
    using a00 a0 a1 a2 a4 a5 a6 a7 a8 get-bit-x-stm4-pre-eq[OF a0 a1]
      mp-alloc-stm4-pre-precond-f-bn
    by (metis mult.commute)
    finally have get-bit-s x p i j = get-bit-s Va p i j .
  } thus ?thesis by fastforce

```

qed

**lemma** *same-bit-mp-alloc-x-va*:

**assumes**  
 $a0: \forall jj. jj \neq \text{nat } (\text{from-l } Va \ t \ + \ 1) \longrightarrow$   
 $\text{levels } (\text{mem-pool-info } x \ p) \ ! \ jj =$   
 $\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) \ ! \ jj$  **and**  
 $a1: \text{bits } (\text{levels } (\text{mem-pool-info } x \ p) \ ! \ \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))))$   
 $(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ * \ 4))) \ \exists \text{ FREE and}$   
 $a2: \neg((l = (\text{nat } (\text{from-l } Va \ t))) \wedge b = (\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{blk } Va \ t)$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))) \vee$   
 $(l = (\text{nat } (\text{from-l } Va \ t \ + \ 1))) \wedge b = (4 * \text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{blk } Va \ t)$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))) \vee$   
 $(l = (\text{nat } (\text{from-l } Va \ t \ + \ 1))) \wedge b = \text{Suc}((4 * \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{blk } Va \ t) \ (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))))) \vee$   
 $(l = (\text{nat } (\text{from-l } Va \ t \ + \ 1))) \wedge b = \text{Suc}(\text{Suc}((4 * \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{blk } Va \ t) \ (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))))) \vee$   
 $(l = (\text{nat } (\text{from-l } Va \ t \ + \ 1))) \wedge b = \text{Suc}(\text{Suc}(\text{Suc}((4 * \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{blk } Va \ t) \ (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)))))))))$   
**shows**  $(\text{get-bit-s } x \ p \ l \ b = \text{get-bit-s } Va \ p \ l \ b)$   
**using** *same-bit-mp-alloc-stm4-pre-precond-f-x*[OF *a0 a1*] *a2*  
**by** *auto*

**lemma** *free-list-x*:

**assumes** *a0*: *free-list*  $(\text{levels } (\text{mem-pool-info } x \ p) \ ! \ \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1) \ * \$   
 $ii \ +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))))$   
**shows**  
 $\text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) \ ! \ \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))) = (\text{free-list}$   
 $(\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1))))@$   
 $[\text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ + \ 1) \ * \$

```

1 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t, lsize (mp-alloc-stm4-pre-precond-f
Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
2 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t,
lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
3 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t]
using a0
by (simp add: numeral-3-eq-3 numeral-Bit0 inserts-def)
lemma listx1:
jj = length (free-list (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p)
p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) ==>
free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
(free-list (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))@
[lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 1 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t,
lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 2 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t,
lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 3 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t] ==>
free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1))! jj =
lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 1 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t
by auto

```

```

lemma listx3:
jj = Suc(Suc (length (free-list (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))) ==>
free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
(free-list (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))@
[lsize (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 1 +
blk (mp-alloc-stm4-pre-precond-f Va t p) t,

```

$lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 2 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t,$   
 $lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 3 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t] \implies$   
 $free\_list (levels (mem\_pool\_info x p) ! nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va$   
 $t p) t + 1))! jj =$   
 $lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 3 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t$   
**by** (*simp add: nth-append*)

**lemma** *set-free-x-va*: **assumes** *a0*: $free\_list (levels (mem\_pool\_info x p) ! nat (from\_l$   
 $(mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1)) =$

*inserts*  
 $(map (\lambda ii. lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) *$   
 $ii +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t)$   
 $[Suc NULL..<4])$   
 $(free\_list$   
 $(levels (mem\_pool\_info (mp\_alloc\_stm4\_pre\_precond\_f Va t p) p) !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1)))$   
**shows**  $set (free\_list (levels (mem\_pool\_info x p) ! nat (from\_l Va t + 1))) =$   
 $set (free\_list (levels (mem\_pool\_info Va p) ! nat (from\_l Va t + 1))) \cup$   
 $\{lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 1 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t, lsizes (mp\_alloc\_stm4\_pre\_precond\_f$   
 $Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 2 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t,$   
 $lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t + 1) * 3 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t\}$

**proof**—

**have**  $free\_list (levels (mem\_pool\_info x p) ! nat (from\_l Va t + 1)) =$   
 $(free\_list$   
 $(levels (mem\_pool\_info Va p) !$   
 $nat (from\_l Va t + 1)))@$   
 $[lsizes Va t !$   
 $nat (from\_l Va t + 1) * 1 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t, lsizes (mp\_alloc\_stm4\_pre\_precond\_f$   
 $Va t p) t !$   
 $nat (from\_l Va t + 1) * 2 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t,$   
 $lsizes (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t !$   
 $nat (from\_l Va t + 1) * 3 +$   
 $blk (mp\_alloc\_stm4\_pre\_precond\_f Va t p) t\}$

```

      2 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t,
      lsizes Va t !
      nat (from-l Va t + 1) *
      3 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t]
  using free-list-x[OF a0]
  by (metis eq-free-list-mp-alloc-stm4-pre-precond-f mp-alloc-stm4-pre-froml mp-alloc-stm4-pre-precond-f-lsz)
  then show ?thesis using mp-alloc-stm4-pre-froml mp-alloc-stm4-pre-precond-f-lsz
    by (metis empty-set list.simps(15) set-append)
qed

```

**lemma** *free-list-x-subset-va*:

**assumes**  $a0$ :free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)) =

*inserts*  
 (map ( $\lambda ii.$  lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !  
 nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) \*  
 ii +  
 blk (mp-alloc-stm4-pre-precond-f Va t p) t)  
 [Suc NULL.. $4$ ])

(free-list  
 (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !  
 nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))  
**shows** set (free-list (levels (mem-pool-info Va p) ! nat (from-l Va t + 1)))  $\subseteq$   
 set (free-list (levels (mem-pool-info x p) ! nat (from-l Va t + 1)))

**proof**–

**have** free-list (levels (mem-pool-info x p) ! nat (from-l Va t + 1)) =  
 (free-list  
 (levels (mem-pool-info Va p) !  
 nat (from-l Va t + 1)))@  
 [lsizes Va t !  
 nat (from-l Va t + 1) \*  
 1 +  
 blk (mp-alloc-stm4-pre-precond-f Va t p) t,lsizes (mp-alloc-stm4-pre-precond-f  
 Va t p) t !

nat (from-l Va t + 1) \*  
 2 +  
 blk (mp-alloc-stm4-pre-precond-f Va t p) t,  
 lsizes Va t !  
 nat (from-l Va t + 1) \*  
 3 +  
 blk (mp-alloc-stm4-pre-precond-f Va t p) t]

**using** free-list-x[OF a0]  
**by** (metis eq-free-list-mp-alloc-stm4-pre-precond-f mp-alloc-stm4-pre-froml mp-alloc-stm4-pre-precond-f-lsz)  
**then show** ?thesis *by auto*  
**qed**

**lemma** *free-level-x-va*:

```

assumes
  a0:  $\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$ 
    levels (mem-pool-info x p) ! jj =
    levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) ! jj
shows  $\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$ 
  free-list (levels (mem-pool-info x p) !jj) = free-list (levels (mem-pool-info Va
p) ! jj)
by (simp add: assms eq-free-list-mp-alloc-stm4-pre-precond-f)

lemma mp-alloc-stm4-pre-precond-f-bitmap-not-free:
assumes a0: (get-bit-s Va p l b  $\neq$  FREE) and
  a1: l < length (levels (mem-pool-info Va p)) and
  a2: b < length (bits (levels (mem-pool-info Va p) ! l)) and
  a3: (from-l Va t)  $\geq$  0
shows (get-bit-s (mp-alloc-stm4-pre-precond-f Va t p) p l b  $\neq$  FREE)
using a0 a1 a2 a3 same-bit-mp-alloc-stm4-pre-precond-divided same-bit-mp-alloc-stm4-pre-precond-allocating
  same-bit-mp-alloc-stm4-pre-precond-f1 BlockState.distinct(11) BlockState.distinct(17)

proof -
let ?i1 = (nat (from-l Va t)) and
  ?j1 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
  ?i2 = (nat (from-l Va t + 1)) and
  ?j2 = (4 * block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t)))
have i1orj1: ?i1  $\neq$  ?i2  $\vee$  ?j1  $\neq$  ?j2 using a3 by auto
{assume a00:  $\neg((l = ?i1 \wedge b = ?j1) \vee (l = ?i2 \wedge b = ?j2))$ 
then have ?thesis using same-bit-mp-alloc-stm4-pre-precond-f1 a0
by auto
}
moreover {assume a00: (l = ?i1  $\wedge$  b = ?j1)
have ?thesis
using same-bit-mp-alloc-stm4-pre-precond-divided[OF a00 a1 a2 a3]
by auto
}
moreover {assume a00: (l = ?i2  $\wedge$  b = ?j2)
have ?thesis
using same-bit-mp-alloc-stm4-pre-precond-allocating[OF a00 a1 a2 a3]
by auto
}
ultimately show ?thesis by auto
qed

```

```

lemma mp-alloc-stm4-pre-inv-nmax: n-max (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) pa) * 4 ^ ii =
  n-max (mem-pool-info Va pa) * 4 ^ ii
unfolding mp-alloc-stm4-pre-precond-f-def set-bit-def

```



**by** *auto*

**lemma** *allocating-next-notexists:inv-bitmap*  $Va \implies$   
 $p \in \text{mem-pools } Va \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va p)) \implies$   
 $jj < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va p) ! ii)) \implies$   
 $\text{get-bit-s } Va p ii jj = \text{ALLOCATING} \implies$   
 $ii < \text{length } (\text{levels } (\text{mem-pool-info } Va p)) - 1 \longrightarrow \text{noexist-bits } (\text{mem-pool-info } Va p) (ii + 1) (jj * 4)$   
**unfolding** *inv-bitmap-def inv-mempool-info-def Let-def*  
**by** *auto*

**lemma** *block-n:*

**assumes**  
 $a0: \text{lsizes } Va t ! \text{nat } (\text{from-l } Va t) = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va p))$   
 $\text{div } 4 \wedge \text{nat } (\text{from-l } Va t)$  **and**  
 $a1: p \in \text{mem-pools } Va$  **and**  
 $a2: \text{inv-mempool-info } Va$  **and**  
 $a3: \text{blk } Va t = \text{buf } (\text{mem-pool-info } Va p) + n * (\text{max-sz } (\text{mem-pool-info } Va p))$   
 $\text{div } 4 \wedge \text{nat } (\text{from-l } Va t)$  **and**  
 $a4: \text{alloc-l } Va t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va p))$  **and**  
 $a5: \text{from-l } Va t < \text{alloc-l } Va t$  **and**  
 $a6: OK \leq \text{from-l } Va t$   
**shows**  $(\text{block-num } (\text{mem-pool-info } Va p) (\text{blk } Va t) (\text{lsizes } Va t ! \text{nat } (\text{from-l } Va t))) = n$   
**proof**–  
**have**  $(\exists n > \text{NULL}. \text{max-sz } (\text{mem-pool-info } Va p) = 4 * n * 4 \wedge n\text{-levels } (\text{mem-pool-info } Va p)) \wedge$   
 $\text{NULL} < n\text{-max } (\text{mem-pool-info } Va p) \wedge$   
 $\text{NULL} < n\text{-levels } (\text{mem-pool-info } Va p) \wedge n\text{-levels } (\text{mem-pool-info } Va p) =$   
 $\text{length } (\text{levels } (\text{mem-pool-info } Va p))$   
**using**  $a2 a1$  **unfolding** *inv-mempool-info-def Let-def* **by** *auto*  
**then show** *?thesis using assms mp-alloc-stm3-lm2-inv-1-2 inv-mempool-info-maxsz-align4 [OF a2]* *nat-less-iff*  
**unfolding** *block-num-def Let-def* **apply** *auto*  
**by** *(smt less-numeral-extra(3))*  
**qed**

**definition** *addr::nat  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  nat*

**where**  $\text{addr } m\text{-size } \text{init } l n \equiv \text{init} + n * (m\text{-size } \text{div } 4 \wedge l)$

**definition** *next-addr :: nat  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  nat  $\Rightarrow$  nat*

**where**  $\text{next-addr } m\text{-size } c\text{-addr } l n \equiv (m\text{-size } \text{div } 4 \wedge (l + 1)) * n + c\text{-addr}$

**lemma** *next-level-addr:*

```

assumes
  a0:∃ m. m-size = m*4 ^ p and
  a1:p > l+1
shows next-addr m-size (addr m-size init l n) l ch = addr m-size init (l+1) (n*4
+ ch)
unfolding next-addr-def addr-def
proof(induct ch)
  case 0
  then show ?case
  apply auto
  by (smt One-nat-def a0 a1 dvd-mult-div-cancel dvd-triv-right less-imp-le-nat
mult.commute mult.left-commute nonzero-mult-div-cancel-left power-Suc0-right
power-add power-le-dvd power-not-zero zero-neq-numeral)
next
  case (Suc ch)
  obtain m where m-size = m*(4::nat) ^ p using a0 by auto
  then show ?case using Suc a1 by auto
qed

```

```

lemma next-level-addr-eq1:
assumes
  a0:∃ m. m-size = m*4 ^ p and
  a1:p > l+1
shows next-addr m-size (addr m-size init l n) l 0 = addr m-size init l n
using next-level-addr[OF a0 a1] unfolding next-addr-def addr-def
by linarith

```

```

lemma next-level-addr-eq:
assumes
  a0:∃ m. m-size = m*4 ^ p and
  a1:p > l+1
shows addr m-size init (l + 1) (n * 4) = addr m-size init l n
using next-level-addr[OF a0 a1] next-level-addr-eq1[OF a0 a1]
by auto

```

```

lemma diff-n-m-addr: assumes a0:n≠m and a1:m-size ≥ 4 ^ l
shows addr m-size init l n ≠ addr m-size init l m
using a0 a1 unfolding addr-def
by (auto simp add: Euclidean-Division.div-eq-0-iff)

```

```

lemma lsizes-addr:
assumes a0:p ∈ mem-pools Va and
  a1:∀ ii<length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
  a2:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
  a3:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
Va and
  a4:l+1 < length (lsizes Va t)
shows ∀ j. (lsizes Va t ! (l+1)) * j +

```

```

    (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^
l)) =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) (l + 1)
    ((block-num (mem-pool-info Va p) (buf (mem-pool-info Va p) +
n * (max-sz (mem-pool-info Va p) div 4 ^ l))
    (lsizes Va t ! l))*4 + j)
proof -
  {fix j
  let ?blk = (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ l))
  obtain m where max-sz:max-sz (mem-pool-info Va p) = 4 * m * 4 ^ n-levels
(mem-pool-info Va p)
  using a3 a0 unfolding inv-mempool-info-def Let-def by auto
  have b1:?blk =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) l
    (block-num (mem-pool-info Va p) ?blk (lsizes Va t ! l))
  using a4 a0 a1 a3
  unfolding addr-def block-num-def apply auto
  by (metis add commute add-lessD1 div-mult-self-is-m
inv-mempool-info-maxsz-align4 plus-1-eq-Suc)
  have b2:(lsizes Va t ! (l+1)) * j + ?blk =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p))
(l+1)
    ((block-num (mem-pool-info Va p) ?blk (lsizes Va t ! l))*4 +
j)
  using assms a4 inv-mempool-info-maxsz-align4 max-sz b1 next-level-addr
  unfolding next-addr-def
  by (smt le-eq-less-or-eq le-less-trans)
  }thus ?thesis by auto
qed

```

```

lemma free-list-updates-inv1:
  assumes a0:p ∈ mem-pools Va and
  a1:¬ free-l Va t < OK and
  a2:free-l Va t ≤ from-l Va t and
  a3:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
  a4:from-l Va t < alloc-l Va t and
  a5:alloc-l Va t = int (length (lsizes Va t)) - 1 ∧ length (lsizes Va t) = n-levels
(mem-pool-info Va p) ∨
  alloc-l Va t = int (length (lsizes Va t)) - 2 ∧ lsizes Va t ! nat (alloc-l Va t + 1)
< sz and
  a6:block-num (mem-pool-info Va p)
  (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
  (lsizes Va t ! nat (from-l Va t))
  < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
  a7:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)) and

```

$a8:(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable}$  **and**  
 $a9:\text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a10:\text{freeing-node } x = \text{freeing-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a11:\text{allocating-node } x = \text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a12:\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ pa$  **and**  
 $a13:\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\quad \text{levels } (\text{mem-pool-info } x \ p) ! jj = \text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj$  **and**  
 $a14:\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t ! ii = \text{ALIGN}_4 \ (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge ii$  **and**  
 $a15:i \ x \ t = 4$  **and**  
 $a16:\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a17:\text{length } (\text{lsizes } Va \ t) \leq n\text{-levels } (\text{mem-pool-info } Va \ p)$  **and**  
 $a18:\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\quad \text{list-updates-n}$   
 $\quad (\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\quad \quad \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$   
 $\quad (\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE}$  **and**  
 $a19:$   
 $\text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $\quad (\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\quad \quad \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $\quad \quad ii +$   
 $\quad \quad \text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $\quad [\text{Suc } \text{NULL}..<4])$   
 $\quad (\text{free-list}$   
 $\quad \quad (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\quad \quad \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$  **and**  
 $a20:\text{allocating-node } Va \ t =$   
 $\text{Some } (\downarrow \text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
 $\quad \text{block} = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $\quad (\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t))))$   
 $\quad (\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)),$   
 $\quad \text{data} = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$  **and**  
 $a21:\text{inv-aux-vars } Va \wedge \text{inv-bitmap } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist } Va$  **and**  
 $a22:ii < \text{length } (\text{levels } (\text{mem-pool-info } x \ p))$  **and**  
 $a23:\text{blk } x = \text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a24:\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$   
**shows**  $\forall j < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! ii)).$   
 $\quad (\text{get-bit-s } x \ p \ ii \ j = \text{FREE}) =$   
 $\quad (\text{buf } (\text{mem-pool-info } x \ p) + j * (\text{max-sz } (\text{mem-pool-info } x \ p) \text{ div } 4 \wedge ii)$   
 $\quad \in \text{set } (\text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) ! ii)))$

```

proof–
{
  let ?i1=(nat (from-l Va t)) and
    ?j1= (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (4*block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t)))
  fix j
  let ?mp = mem-pool-info x p
  let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
  assume a00:j<length ?bts
  then have a00':j < length (bits (levels (mem-pool-info Va p) ! ii))
    using mp-alloc-stm4-inv-bits-len
    by (metis a13 a18 length-list-update-n)
  have inv-bitmap1:(∀ j<length (bits (levels (mem-pool-info Va p) ! ii)).
    (get-bit-s Va p ii j = FREE) =
    (buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va p) div 4 ^
ii)
    ∈ set (free-list (levels (mem-pool-info Va p) ! ii))))
    using a21 a0 a22 mp-alloc-stm4-lvl-len[OF a0 a8]
    unfolding Let-def inv-bitmap-freelist-def
    by fastforce+
  have from-l-gt0:0 ≤ from-l Va t using a1 a2 by linarith
  have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
Va p))
    using mp-alloc-stm4-lvl-len[OF a0 a8] by simp
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
    using mp-alloc-stm4-maxsz[OF a0 a8] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
    using mp-alloc-stm4-buf[OF a0 a8] by simp
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a9] by auto
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a8] by auto
  have lsizes-x-va:lsizes x = lsizes Va
    by (simp add: a16 mp-alloc-stm4-pre-precond-f-lsz)
  have len-eq:length (bits (levels (mem-pool-info x p) ! ii)) =
    length (bits (levels (mem-pool-info Va p) ! ii))
    using a22 a8 mp-alloc-stm4-inv-bits-len
    unfolding gvars-conf-stable-def gvars-conf-def
    by fastforce
  then have get-bits-va:(get-bit-s Va p ii j = FREE) =
    (buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va p) div 4
^ ii)
    ∈ set (free-list (levels (mem-pool-info Va p) ! ii)))
    using inv-bitmap1 a00 by auto
  have inv-mempool-info-mp Va p
    using a21 mem-pools a0 unfolding inv-mempool-info-def by auto

```

```

note inv-mempool=this[simplified Let-def]
have a19':ii < length (levels (mem-pool-info Va p))
  using a22 mp-alloc-stm4-inv-mif-len
  by (simp add: len-levels)
{ assume a03:ii≠?i1 ∧ ii≠?i2
  then have eq-get-bit-i-j:get-bit-s x p ii j = get-bit-s Va p ii j
  using same-bit-mp-alloc-x-va[OF a13[simplified a9[simplified mp-alloc-stm4-froml[OF
a9], THEN sym]] a18] by fast
  moreover have free-list (levels (mem-pool-info x p) ! ii) =
    free-list (levels (mem-pool-info Va p) ! ii)
  using free-level-x-va[OF a13] a03 a9 from-l by metis
  ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
ii) ∈ set ?fl)
    using get-bits-va eq-get-bit-i-j
    by (simp add: buf maxsz)
}
moreover { assume a03:ii=?i1
  then have free:free-list (levels (mem-pool-info x p) ! ii) =
    free-list (levels (mem-pool-info Va p) ! ii)
  using free-level-x-va[OF a13] a03 a9 from-l from-l-gt0 by auto
  { assume a04:j≠?j1
    then have eq-get-bit-i-j:get-bit-s x p ii j = get-bit-s Va p ii j
    using same-bit-mp-alloc-x-va[OF a13[simplified a9[simplified from-l, THEN
sym]] a18]
      a03 from-l-gt0
    by (simp add: eq-nat-nat-iff)
    then have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∈
set ?fl)
      using free by (simp add: buf get-bits-va maxsz)
  }
moreover { assume a04:j=?j1
  then have (?bts ! j = DIVIDED)
  using get-bit-x-l-b a03 a13 a00 a22 len-levels
    len-eq a13 from-l from-l-gt0
  by (simp add: a9)
  moreover have buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va
p) div 4 ^ ii) ∉
    set (free-list (levels (mem-pool-info Va p) ! ii))
  using get-bits-va a03 a20 a21 a04 a7 unfolding inv-aux-vars-def
  by (metis BlockState.distinct(17) Mem-block.select-convs(1) Mem-block.select-convs(2)
Mem-block.select-convs(3))
  ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
ii) ∈ set ?fl)
    by (simp add: buf free maxsz)
  }
ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
ii) ∈ set ?fl)
  by auto
}

```

```

moreover { assume  $a03:ii=?i2$ 
  then have  $block-n:(block-num (mem-pool-info Va p)$ 
     $(blk Va t) (lsizes Va t ! nat (from-l Va t))) = n$ 
  proof–
    have  $lsizes Va t ! nat (from-l Va t) =$ 
       $ALIGN4 (max-sz (mem-pool-info Va p)) div 4 \wedge$ 
       $(nat (from-l Va t))$ 
    using  $a14 lsizes-x-va a16 a1 a2 a4 a5 a9 from-l$  by auto
    thus  $?thesis$  using  $block-n a21 a0 a0 a7 a3 a4 from-l-gt0$ 
    by blast
  qed
  obtain  $m$  where  $max-sz:max-sz (mem-pool-info Va p) = 4 * m * 4 \wedge n-levels$ 
     $(mem-pool-info Va p)$ 
    using  $a21 a0$  unfolding inv-mempool-info-def Let-def by auto
    have  $ls:4 \wedge ii dvd 4 * m * 4 \wedge n-levels (mem-pool-info Va p)$  using  $a03 a22$ 
    by  $(metis dvd-triv-right inv-mempool len-levels less-imp-le-nat power-le-dvd)$ 
    have  $b0:buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va p) div 4$ 
       $\wedge ii) =$ 
       $addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) ii j$ 
    unfolding addr-def by auto
    have  $suc-from-l-lt-lsize:(nat (from-l Va t)) + 1 < length (lsizes Va t)$ 
    using  $a4 a5 from-l-gt0$  by linarith
    have  $b2:\forall j. (lsizes Va t ! nat (from-l Va t + 1)) * j + blk Va t =$ 
       $addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p))$ 
     $(nat (from-l Va t + 1))$ 
       $((block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat$ 
     $(from-l Va t))) * 4 + j)$ 
    using  $lsizes-addr[OF a0 a14 a17 a21 suc-from-l-lt-lsize] a7 from-l-gt0 block-n$ 
    by  $(simp add: Suc-nat-eq-nat-zadd1 add.commute)$ 
    then have  $b2:\forall j. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !$ 
       $nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * j +$ 
       $blk (mp-alloc-stm4-pre-precond-f Va t p) t =$ 
       $addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p))$ 
     $(nat (from-l Va t + 1))$ 
       $((block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat$ 
     $(from-l Va t))) * 4 + j)$ 
    by  $(metis mp-alloc-stm4-blk mp-alloc-stm4-pre-precond-f-froml mp-alloc-stm4-pre-precond-f-lsz)$ 
    { assume  $a04:j \neq ?j2 \wedge j \neq Suc ?j2 \wedge j \neq Suc (Suc ?j2) \wedge j \neq Suc (Suc (Suc$ 
       $?j2))$ 
    then have  $eq-get-bit-i-j:get-bit-s x p ii j = get-bit-s Va p ii j$ 
    using  $same-bit-mp-alloc-x-va[OF a13[simplified a9[simplified from-l, THEN$ 
     $sym]] a18]$ 
       $a03 from-l-gt0$ 
    by  $(simp add: eq-nat-nat-iff)$ 
    { assume  $get-bit-s Va p ii j = FREE$ 
    then have  $(buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va$ 
     $p) div 4 \wedge ii)$ 
       $\in set (free-list (levels (mem-pool-info Va p) ! ii)))$ 
    using get-bits-va by blast

```

```

    then have (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∈ set ?fl)
      using free-list-x-subset-va[OF a19] a03 buf maxsz by fastforce
  }
  moreover {
    assume get-bit-s Va p ii j ≠ FREE
    then have not-in-free-Va: (buf (mem-pool-info Va p) + j * (max-sz
(mem-pool-info Va p) div 4 ^ ii)
      ∉ set (free-list (levels (mem-pool-info Va p) ! ii)))
    using get-bits-va by blast
    then have (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∉ set ?fl)
    proof-
      have ∀ k. k < 4 → (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ≠
lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
        nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
k +
          blk (mp-alloc-stm4-pre-precond-f Va t p) t
      using diff-n-m-addr b2 b0 a03 a04 Euclidean-Division.div-eq-0-iff buf
inv-mempool
        ls max-sz maxsz
      by (smt Groups.mult-ac(2) add.right-neutral add-2-eq-Suc' add-Suc-right

          dvd-div-mult-self less-2-cases less-Suc-eq div-greater-zero-iff
mult-is-0 neq0-conv numeral-Bit0 power-not-zero)
    then show ?thesis using buf maxsz not-in-free-Va set-free-x-va[OF a19,
simplified] a03
      apply auto
      by presburger
    qed
    then have (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∉ set ?fl)
      using a03 buf maxsz a04 set-free-x-va[OF a19] by auto
  } ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4
^ ii) ∈ set ?fl)
    using eq-get-bit-i-j by auto
  }
  moreover { assume a04:j=?j2
    then have a03':ii = nat (from-l x t + 1) ∧
      j = 4*block-num (mem-pool-info x p) (blk x t) (lsizes x t ! nat
(from-l x t))
    using a22 a23 a24 from-l buf from-l a03
    unfolding block-num-def
    by (simp add: mp-alloc-stm4-pre-precond-f-blk mp-alloc-stm4-pre-precond-f-lsz)

    have (?bts ! j = ALLOCATING)
    using from-l-gt0 a22 a00 a03 len-levels a04 a18 from-l get-bit-x-l1-b4 len-eq
    by (metis a04 a18 get-bit-x-l1-b4 len-eq)
    then have bts-j-not-free:(?bts ! j ≠ FREE)
    by auto
    moreover have not-in-free-Va:buf (mem-pool-info Va p) + j * (max-sz
(mem-pool-info Va p) div 4 ^ ii) ∉

```



```

      set (free-list (levels (mem-pool-info Va p) ! ii))
proof-
  have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
    using a20 a21 a7 unfolding inv-aux-vars-def
    by (metis (no-types) Mem-block.select-convs(1)
        Mem-block.select-convs(2) Mem-block.select-convs(3))
  have noexist-bits (mem-pool-info Va p) (?i1 + 1) (?j1 * 4)
proof-
  have ?i1 < length (levels (mem-pool-info Va p)) - 1
    using a19' from-l-gt0 a3 a4 inv-mempool by auto
  moreover have ?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    using calculation
    by (simp add: a6 a7 inv-mempool)
  ultimately show ?thesis
    using alloc-i1-j1 a21 a19' a00' a0 a03 a04
    unfolding Let-def inv-bitmap-def
    by (smt One-nat-def Suc-pred inv-mempool less-Suc-eq)
qed
then have (get-bit-s Va p ii j = NOEXIST)
  using a03 a04 from-l-gt0
  by (simp add: mult.commute nat-add-distrib)
then show ?thesis using get-bits-va
  by simp
qed
have (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∉ set ?fl)
proof-
  have ∀ k. k > 0 ∧ k < 4 ⟶ (buf ?mp + j * (max-sz ?mp div 4 ^ ii)) ≠
    lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * k +
    blk (mp-alloc-stm4-pre-precond-f Va t p) t
  using diff-n-m-addr b2 b0 a03 a04 buf inv-mempool
    maxsz
  by (metis a19' add-cancel-right-right div-greater-zero-iff
    mp-alloc-stm3-lm2-inv-1-2 mult.commute neq0-conv)
  then show ?thesis using buf maxsz not-in-free-Va set-free-x-va[OF a19,
simplified] a03
    by auto
qed
then have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∈
set ?fl)
  using bts-j-not-free by auto
}
moreover {
  assume a04:j=Suc ?j2 ∨ j = Suc (Suc ?j2) ∨ j = Suc (Suc (Suc ?j2))
  then have a04':j=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t !
nat (from-l Va t)) * 4 + 1) ∨
    j = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat
(from-l Va t)) * 4 + 2) ∨
    j = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat

```

```

(from-l Va t)) * 4 + 3)
  by auto
  have (?bts ! j = FREE)
    using get-bit-x-l1-b41[OF conjI[OF a03 a04] from-l-gt0 - a18 a19' a00']
    using a13 a9 from-l by auto
  moreover have buf ?mp + j * (max-sz ?mp div 4 ^ ii) ∈ set ?fl
    using a03 a04'[simplified] set-free-x-va[OF a19, simplified b2 buf[THEN
sym] maxsz[THEN sym] ]
    using b0[simplified buf[THEN sym] maxsz[THEN sym] a03] by auto
  ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
ii) ∈ set ?fl) by auto
} ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
^ ii) ∈ set ?fl)
  by auto
} ultimately have (?bts ! j = FREE) = (buf ?mp + j * (max-sz ?mp div 4 ^
ii) ∈ set ?fl) by auto
} then show ?thesis by auto
qed

```

**lemma** *free-list-updates-inv2*:

```

  assumes a0:p ∈ mem-pools Va and
  a1:¬ free-l Va t < OK and
  a2:free-l Va t ≤ from-l Va t and
  a3:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
  a4:from-l Va t < alloc-l Va t and
  a5: alloc-l Va t = int (length (lsizes Va t)) - 1 ∧ length (lsizes Va t) = n-levels
(mem-pool-info Va p) ∨
  alloc-l Va t = int (length (lsizes Va t)) - 2 ∧ lsizes Va t ! nat (alloc-l Va t + 1)
< sz and
  a6:block-num (mem-pool-info Va p)
  (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
  (lsizes Va t ! nat (from-l Va t))
  < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
  a7:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)) and
  a8:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and
  a9:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) and
  a10:∀ jj. jj ≠ nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) →
    levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p) ! jj and
  a11:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
  a12:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p) and
  a13:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
  a14:
  free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
  inserts

```

```

(map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
      nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
      ii +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t)
  [Suc NULL..<4])
(free-list
  (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
a15:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
Va and
a16:ii < length (levels (mem-pool-info x p)) and
a17:blk x = blk (mp-alloc-stm4-pre-precond-f Va t p) and
a18:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p)
shows ∀j<length (free-list (levels (mem-pool-info x p) ! ii)).
  ∃ n<n-max (mem-pool-info x p) * 4 ^ ii.
    free-list (levels (mem-pool-info x p) ! ii) ! j =
    buf (mem-pool-info x p) + n * (max-sz (mem-pool-info x p) div 4 ^ ii)

proof-
{ let ?i1=(nat (from-l Va t)) and
  ?j1= (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
  ?i2 = (nat (from-l Va t + 1)) and
  ?j2 = (4*block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t)))
  let ?mp = mem-pool-info x p
  let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
  fix j
  assume a00:j<length ?fl
  have inv-bitmap2:(∀j<length (free-list (levels (mem-pool-info Va p) ! ii)).
    ∃ n<n-max (mem-pool-info Va p) * 4 ^ ii.
      free-list (levels (mem-pool-info Va p) ! ii) ! j =
      buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4
^ ii))
    using a15 a0 a16 mp-alloc-stm4-lvl-len[OF a0 a8]
    unfolding Let-def inv-bitmap-freelist-def
    by fastforce+
  have from-l-gt0:0 ≤ from-l Va t using a1 a2 by linarith
  have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
Va p))
    using mp-alloc-stm4-lvl-len[OF a0 a8] by simp
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
    using mp-alloc-stm4-maxsz[OF a0 a8] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
    using mp-alloc-stm4-buf[OF a0 a8] by simp
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a9] by auto
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a8] by auto

```

```

have lsize-x-va:lsize x = lsize Va
  by (simp add: a12 mp-alloc-stm4-pre-precond-f-lsz)
have len-eq:length (bits (levels (mem-pool-info x p) ! ii)) =
  length (bits (levels (mem-pool-info Va p) ! ii))
using a16 a8 mp-alloc-stm4-inv-bits-len
unfolding gvars-conf-stable-def gvars-conf-def
by fastforce
have inv-mempool-info-mp Va p
  using a15 mem-pools a0 unfolding inv-mempool-info-def by auto
note inv-mempool=this[simplified Let-def]
have a15':ii < length (levels (mem-pool-info Va p))
  using a16 mp-alloc-stm4-inv-mif-len
  by (simp add: len-levels)
have nmax: n-max (mem-pool-info x p) = n-max (mem-pool-info Va p)
  using a8 unfolding gvars-conf-stable-def gvars-conf-def apply auto
  by (metis mp-alloc-stm4-nmax)
{ assume a03:ii ≠ ?i2
  then have ∃ n < n-max ?mp * 4 ^ ii. ?fl ! j = buf ?mp + n * (max-sz ?mp div
4 ^ ii)
    using a0 a00 a10 buf eq-free-list-mp-alloc-stm4-pre-precond-f
      inv-bitmap2 maxsz nmax
  by (simp add: eq-free-list-mp-alloc-stm4-pre-precond-f mp-alloc-stm4-pre-precond-f-froml
)
}
moreover {
  assume a03:ii = ?i2
  then have suc-from-l-lt-lsize:(nat (from-l Va t)) + 1 < length (lsize Va t)
    using a4 a5 from-l-gt0 by linarith
  then have lsize-i:lsize Va t ! nat (from-l Va t) =
    ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
      (nat (from-l Va t))
    using a11 add-lessD1 suc-from-l-lt-lsize by blast
  then have block-n:(block-num (mem-pool-info Va p)
    (blk Va t) (lsize Va t ! nat (from-l Va t))) = n
    using block-n a0 a3 a4 from-l-gt0 a15 a7 by blast
  have lsize-ii:lsize Va t ! ii =
    ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ ii
    using a11 from-l-gt0 suc-from-l-lt-lsize
    by (simp add: a03)
  {assume a04:j < length (free-list (levels (mem-pool-info Va p) ! ii))
    then have free-list (levels (mem-pool-info Va p) ! ii) ! j = ?fl ! j
    using a14[simplified mp-alloc-stm4-pre-precond-f-froml eq-free-list-mp-alloc-stm4-pre-precond-f]
a03
    unfolding inserts-def
    by (simp add: nth-append)
  moreover have ∃ n < n-max (mem-pool-info Va p) * 4 ^ ii.
    free-list (levels (mem-pool-info Va p) ! ii) ! j =
    buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ ii)
    using a04 inv-bitmap2 by fastforce
  }
}

```

```

ultimately have  $\exists n < n\text{-max} \text{ ?mp} * 4^{\wedge ii}. \text{ ?fl} ! j = \text{buf ?mp} + n * (\text{max-sz}$ 
 $\text{ ?mp} \text{ div } 4^{\wedge ii})$ 
  using buf maxsz nmax by auto
}
moreover { assume a04:j = length (free-list (levels (mem-pool-info Va p) !
ii))
  then have fl-lsizes: ?fl ! j = lsizes Va t ! nat (from-l Va t + 1) * 1 + blk
Va t
    using free-list-x[OF a14] a03 a9 eq-free-list-mp-alloc-stm4-pre-precond-f
nth-append-length
    by (metis (no-types, lifting) a18 from-l lsizes-x-va mp-alloc-stm4-blk)
    let ?nb = (block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4^nat
      (from-l Va t))))
      (lsizes Va t ! nat (from-l Va t)) * 4 + 1)
    have eq-suc-from-l:nat (from-l Va t + 1) = nat (from-l Va t) + 1 using
from-l-gt0 by auto
    from fl-lsizes[simplified a7 this] have ?fl ! j =
      addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) (nat
      (from-l Va t) + 1)
      (block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4^nat
      (from-l Va t))))
      (lsizes Va t ! nat (from-l Va t)) * 4 + 1)
    using spec[OF lsizes-addr[OF a0 a11 a13 a15 suc-from-l-lt-lsize], of 1]
    by auto
  moreover have length (bits (levels (mem-pool-info x p) ! ii)) =
    n-max (mem-pool-info x p) * 4^ii using a15 a0
    unfolding inv-mempool-info-def Let-def
    by (simp add: a15' len-eq nmax)
  moreover have ?nb < n-max (mem-pool-info x p) * 4^ii
    using a6 a03 a0 nmax lsize-ii lsize-i eq-suc-from-l
    inv-mempool-info-maxsz-align4[OF conjunct1[OF conjunct2[OF
conjunct2[OF a15]]], simplified a0]
    unfolding block-num-def
    by auto
  ultimately have
    ?nb < n-max ?mp * 4^ii  $\wedge$  ?fl ! j = buf ?mp + ?nb * (max-sz ?mp div 4
^ ii)
    using block-n a7 buf nmax maxsz a6 a03 eq-suc-from-l unfolding addr-def
    by auto
  then have  $\exists n < n\text{-max} \text{ ?mp} * 4^{\wedge ii}. \text{ ?fl} ! j = \text{buf ?mp} + n * (\text{max-sz ?mp}$ 
 $\text{ div } 4^{\wedge ii})$  by auto
}
moreover { assume a04:j = Suc (length (free-list (levels (mem-pool-info Va
p) ! ii)))
  then have fl-lsizes: ?fl ! j = lsizes Va t ! nat (from-l Va t + 1) * 2 + blk
Va t
    using free-list-x[OF a14] a03 a9 eq-free-list-mp-alloc-stm4-pre-precond-f

```

```

    nth-append-length a18 from-l lsizes-x-va mp-alloc-stm4-blk
  by (metis add.right-neutral add-Suc-right nth-Cons-Suc nth-append-length-plus)
  let ?nb = (block-num (mem-pool-info Va p)
    (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
      (from-l Va t))))
    (lsizes Va t ! nat (from-l Va t)) * 4 + 2)
  have eq-suc-from-l:nat (from-l Va t + 1) = nat (from-l Va t) + 1 using
from-l-gt0 by auto
  from fl-lsizes[simplified a7 this] have ?fl ! j =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) (nat
      (from-l Va t) + 1)
    (block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
        (from-l Va t))))
    (lsizes Va t ! nat (from-l Va t)) * 4 + 2)
  using spec[OF lsizes-addr[OF a0 a11 a13 a15 suc-from-l-lt-lsize], of 2 n]
  by fastforce
  moreover have length (bits (levels (mem-pool-info x p) ! ii)) =
    n-max (mem-pool-info x p) * 4 ^ ii using a15 a0
  unfolding inv-mempool-info-def Let-def
  by (simp add: a15' len-eq nmax)
  moreover have ?nb < n-max (mem-pool-info x p) * 4 ^ ii
  using a6 a03 a0 nmax lsize-ii lsize-i eq-suc-from-l
    inv-mempool-info-maxsz-align4[OF conjunct1[OF conjunct2[OF
conjunct2[OF a15]]], simplified a0]
  unfolding block-num-def
  by auto
  ultimately have
    ?nb < n-max ?mp * 4 ^ ii ∧ ?fl ! j = buf ?mp + ?nb * (max-sz ?mp div 4
    ^ ii)
  using block-n a7 buf nmax maxsz a6 a03 eq-suc-from-l unfolding addr-def
  by auto
  then have ∃ n < n-max ?mp * 4 ^ ii. ?fl ! j = buf ?mp + n * (max-sz ?mp
div 4 ^ ii) by auto
}
  moreover { assume a04:j = Suc (Suc (length (free-list (levels (mem-pool-info
Va p) ! ii))))
    then have fl-lsizes:?fl ! j = lsizes Va t ! nat (from-l Va t + 1) * 3 + blk
Va t
      using free-list-x[OF a14] a03 a9 eq-free-list-mp-alloc-stm4-pre-precond-f
      nth-append-length a18 from-l lsizes-x-va mp-alloc-stm4-blk
    by (metis add.right-neutral add-Suc-right nth-Cons-Suc nth-append-length-plus)
    let ?nb = (block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
        (from-l Va t))))
      (lsizes Va t ! nat (from-l Va t)) * 4 + 3)
    have eq-suc-from-l:nat (from-l Va t + 1) = nat (from-l Va t) + 1 using
from-l-gt0 by auto
    from fl-lsizes[simplified a7 this] have ?fl ! j =

```

```

      addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p)) (nat
(from-l Va t) + 1)
    (block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
      (lsizes Va t ! nat (from-l Va t)) * 4 + 3)
    using spec[OF lsizes-addr[OF a0 a11 a13 a15 suc-from-l-lt-lsize], of 3 n]
    by fastforce
  moreover have length (bits (levels (mem-pool-info x p) ! ii)) =
    n-max (mem-pool-info x p) * 4 ^ ii using a15 a0
  unfolding inv-mempool-info-def Let-def
  by (simp add: a15' len-eq nmax)
  moreover have ?nb < n-max (mem-pool-info x p) * 4 ^ ii
  using a6 a03 a0 nmax lsize-ii lsize-i eq-suc-from-l
    inv-mempool-info-maxsz-align4[OF conjunct1[OF conjunct2[OF
conjunct2[OF a15]]], simplified a0]
  unfolding block-num-def
  by auto
  ultimately have
    ?nb < n-max ?mp * 4 ^ ii ∧ ?fl ! j = buf ?mp + ?nb * (max-sz ?mp div 4
^ ii)
  using block-n a7 buf nmax maxsz a6 a03 eq-suc-from-l unfolding addr-def
  by auto
  then have ∃ n < n-max ?mp * 4 ^ ii. ?fl ! j = buf ?mp + n * (max-sz ?mp
div 4 ^ ii) by auto
} ultimately have ∃ n < n-max ?mp * 4 ^ ii. ?fl ! j = buf ?mp + n * (max-sz
?mp div 4 ^ ii)
  using a00 free-list-x[OF a14,
    simplified eq-free-list-mp-alloc-stm4-pre-precond-f
mp-alloc-stm4-pre-precond-f-froml] a03
  by fastforce
} ultimately have ∃ n < n-max ?mp * 4 ^ ii. ?fl ! j = buf ?mp + n * (max-sz
?mp div 4 ^ ii)
  by auto
} then show ?thesis by auto
qed

```

**lemma** next-block-less-length-bits:

**assumes**

a0:  $n < \text{length } (\text{bits } (\text{levels } pi ! ii))$  **and**

a1:  $(ii+1) < \text{length } (\text{levels } pi)$  **and**

a2:  $(\forall i < \text{length } (\text{levels } pi)).$

$\text{length } (\text{bits } (\text{levels } pi ! i)) = n\text{-max } pi * 4 ^ i$

**shows**  $4 * n + 3 < \text{length } (\text{bits } (\text{levels } pi ! (ii+1)))$

**proof** –

**have**  $n < n\text{-max } pi * 4 ^ ii$  **using** a0 a1 a2 **by** auto

**moreover have**  $\text{length } (\text{bits } (\text{levels } pi ! (ii+1))) = n\text{-max } pi * 4 ^ (ii+1)$  **using**  
a1 a2 **by** auto

**ultimately show** ?thesis **by** auto

qed

**lemma** *distinct-lists*: **assumes**

*a0*:distinct *l1* **and**

*a1*:distinct *l2* **and**

*a2*: $\forall e \in \text{set } l2. e \notin \text{set } l1$

**shows** distinct (*l1* @ *l2*)

**using** *assms*

**by**(*induct l1, auto*)

**lemma** *free-list-updates-inv3*:

**assumes** *a0*: $p \in \text{mem-pools } Va$  **and**

*a1*: $\neg \text{free-l } Va \ t < OK$  **and**

*a2*: $\text{free-l } Va \ t \leq \text{from-l } Va \ t$  **and**

*a3*: $\text{alloc-l } Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p))$  **and**

*a4*: $\text{from-l } Va \ t < \text{alloc-l } Va \ t$  **and**

*a5*: $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 1 \wedge \text{length } (\text{lsizes } Va \ t) = n\text{-levels } (\text{mem-pool-info } Va \ p) \vee$

$\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t ! \text{ nat } (\text{alloc-l } Va \ t + 1) < sz$  **and**

*a6*: $\text{block-num } (\text{mem-pool-info } Va \ p)$

$(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{ nat } (\text{from-l } Va \ t)))$

$(\text{lsizes } Va \ t ! \text{ nat } (\text{from-l } Va \ t))$

$< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 ^ \text{ nat } (\text{from-l } Va \ t)$  **and**

*a7*: $\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{ nat } (\text{from-l } Va \ t))$  **and**

*a8*: $(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable}$  **and**

*a9*: $\text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**

*a10*: $\text{freeing-node } x = \text{freeing-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**

*a11*: $\text{allocating-node } x = \text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**

*a12*: $\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ pa$  **and**

*a13*: $\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$

$\text{levels } (\text{mem-pool-info } x \ p) ! jj = \text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj$  **and**

*a14*: $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t ! ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 ^ ii$  **and**

*a15*: $i \ x \ t = 4$  **and**

*a16*: $\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**

*a17*: $\text{length } (\text{lsizes } Va \ t) \leq n\text{-levels } (\text{mem-pool-info } Va \ p)$  **and**

*a18*: $\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{ nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$

$\text{list-updates-n}$

$(\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$

$\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$

$(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE}$  **and**

*a19*:



```

free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
inserts
  (map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
    ii +
    blk (mp-alloc-stm4-pre-precond-f Va t p) t)
    [Suc NULL..<4])
  (free-list
    (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
      nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
a20:allocating-node Va t =
Some (pool = p, level = nat (from-l Va t),
  block = block-num (mem-pool-info Va p)
    (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
    (lsizes Va t ! nat (from-l Va t)),
  data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t))) and
a21:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
Va and
a22:ii < length (levels (mem-pool-info x p)) and
a23:blk x = blk (mp-alloc-stm4-pre-precond-f Va t p) and
a24:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p)
shows distinct (free-list (levels (mem-pool-info x p) ! ii))
proof-
{
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (4*block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t)))
  let ?mp = mem-pool-info x p
  let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
  have inv-bitmap1:(∀j<length (bits (levels (mem-pool-info Va p) ! ii)).
    (get-bit-s Va p ii j = FREE) =
    (buf (mem-pool-info Va p) + j * (max-sz (mem-pool-info Va p) div 4 ^
ii)
    ∈ set (free-list (levels (mem-pool-info Va p) ! ii)))) and
    inv-bitmap2:(∀j<length (free-list (levels (mem-pool-info Va p) ! ii)).
    ∃ n<n-max (mem-pool-info Va p) * 4 ^ ii.
    free-list (levels (mem-pool-info Va p) ! ii) ! j =
    buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4
^ ii)) and
    inv-bitmap3:distinct (free-list (levels (mem-pool-info Va p) ! ii))
  using a21 a0 a22 mp-alloc-stm4-lvl-len[OF a0 a8]
  unfolding Let-def inv-bitmap-freelist-def
  by fastforce+

```

```

have from-l-gt0:  $0 \leq \text{from-l } Va \ t$  using a1 a2 by linarith
have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
Va p))
  using mp-alloc-stm4-lvl-len[OF a0 a8] by simp
have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
  using mp-alloc-stm4-maxsz[OF a0 a8] by simp
have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
  using mp-alloc-stm4-buf[OF a0 a8] by simp
have from-l:from-l x = from-l Va
  using mp-alloc-stm4-froml[OF a9] by auto
have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
  using from-l-gt0 by auto
have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a8] by auto
have lsizes-x-va:lsizes x = lsizes Va
  by (simp add: a16 mp-alloc-stm4-pre-precond-f-lsz)
have len-eq:length (bits (levels (mem-pool-info x p) ! ii)) =
  length (bits (levels (mem-pool-info Va p) ! ii))
using a22 a8 mp-alloc-stm4-inv-bits-len
unfolding gvars-conf-stable-def gvars-conf-def
by fastforce
have inv-mempool-info-mp Va p
  using a21 mem-pools a0 unfolding inv-mempool-info-def by auto
note inv-mempool=this[simplified Let-def]
have a22':ii < length (levels (mem-pool-info Va p))
  using a22 mp-alloc-stm4-inv-mif-len
  by (simp add: len-levels)
{ assume a03:ii ≠ ?i2
  have free-list (levels (mem-pool-info x p) ! ii) =
    free-list (levels (mem-pool-info Va p) ! ii)
    using free-level-x-va[OF a13] a03 a9 from-l by metis
  then have distinct (free-list (levels (mem-pool-info x p) ! ii))
    using inv-bitmap3 by auto
}
moreover { assume a03:ii = ?i2
  then have block-n:(block-num (mem-pool-info Va p)
    (blk Va t) (lsizes Va t ! nat (from-l Va t))) = n
  proof–
    have lsizes Va t ! nat (from-l Va t) =
      ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
      (nat (from-l Va t))
    using a14 lsizes-x-va a16 a1 a2 a4 a5 a9 from-l by auto
    thus ?thesis using block-n a21 a0 a0 a7 a3 a4 from-l-gt0
    by blast
  qed
then have get-bit-s Va p ( nat (from-l Va t) ) n = ALLOCATING
  using a20 a13 a21 a7 unfolding inv-aux-vars-def
by (metis Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))
moreover have n-len: n < length (bits (levels (mem-pool-info Va p) ! nat

```

```

(from-l Va t )))
  using a03 a22' a6 a7 inv-mempool local.block-n by auto
  ultimately have noexist-bits (mem-pool-info Va p) ii (n * 4)
    using allocating-next-notexists[OF conjunct1[OF conjunct2[OF a21]] a0 - ]
a21 a0 a03 a21 a0 a22'
  from-l-gt0 from-l-suc inv-mempool by auto
  then have get-bit-s Va p ii (n*4 + 1) ≠ FREE ∧
    get-bit-s Va p ii (n*4 + 2) ≠ FREE ∧
    get-bit-s Va p ii (n*4 + 3) ≠ FREE
  by (simp add: mult.commute)
  moreover have n*4 + 3 < length (bits (levels (mem-pool-info Va p) ! nat
(from-l Va t + 1)))
  using a03 a22' n-len inv-mempool from-l-gt0 next-block-less-length-bits from-l-suc
  by simp
  ultimately have not-in-freelist:(buf (mem-pool-info Va p) + (n*4 + 1) *
(max-sz (mem-pool-info Va p) div 4 ^ ii)
    ∉ set (free-list (levels (mem-pool-info Va p) ! ii))) ∧
    (buf (mem-pool-info Va p) + (n*4 + 2) * (max-sz (mem-pool-info Va p)
div 4 ^ ii)
    ∉ set (free-list (levels (mem-pool-info Va p) ! ii))) ∧
    (buf (mem-pool-info Va p) + (n*4 + 3) * (max-sz (mem-pool-info Va p)
div 4 ^ ii)
    ∉ set (free-list (levels (mem-pool-info Va p) ! ii)))
  using inv-bitmap1 a03
  by (metis (no-types, lifting) add-lessD1 numeral-3-eq-3
    one-add-one plus-1-eq-Suc semiring-normalization-rules(21))
  obtain m where max-sz:max-sz (mem-pool-info Va p) = 4 * m * 4 ^ n-levels
(mem-pool-info Va p)
  using a21 a0 unfolding inv-mempool-info-def Let-def by auto
  have ls:4 ^ ii dvd 4 * m * 4 ^ n-levels (mem-pool-info Va p) using a03 a22
  by (metis dvd-triv-right inv-mempool len-levels less-imp-le-nat power-le-dvd)

  have suc-from-l-lt-lsize:(nat (from-l Va t)) + 1 < length (lsizes Va t)
  using a4 a5 from-l-gt0 by linarith
  have b2:∀ j. (lsizes Va t ! nat (from-l Va t + 1)) * j + blk Va t =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p))
(nat (from-l Va t + 1))
    ((block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat
(from-l Va t))) * 4 + j)
  using lsizes-addr[OF a0 a14 a17 a21 suc-from-l-lt-lsize] a7 from-l-gt0 block-n
  by (simp add: Suc-nat-eq-nat-zadd1 add.commute)
  then have b2:∀ j. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * j +
    blk (mp-alloc-stm4-pre-precond-f Va t p) t =
    addr (max-sz (mem-pool-info Va p)) (buf (mem-pool-info Va p))
(nat (from-l Va t + 1))
    ((block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat
(from-l Va t))) * 4 + j)
  by (metis mp-alloc-stm4-blk mp-alloc-stm4-pre-precond-f-froml mp-alloc-stm4-pre-precond-f-lsz)

```

```

    then have distinct (free-list (levels (mem-pool-info x p) ! ii))
  proof-
    have h1:distinct [lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l
      (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 1 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t,
      lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l (mp-alloc-stm4-pre-precond-f
        Va t p) t + 1) * 2 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t,
      lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l (mp-alloc-stm4-pre-precond-f
        Va t p) t + 1) * 3 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t] using b2 a03 a22' inv-mempool
      mp-alloc-stm3-lm2-inv-1-2 unfolding addr-def
      by (smt add-diff-cancel-left' distinct-length-2-or-more
          distinct-singleton mult-cancel-right nat-less-le num.distinct(3)
          num.distinct(5)
          numeral-eq-iff numeral-eq-one-iff semiring-norm(85))
    have h2:∀ e∈set [lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l
      (mp-alloc-stm4-pre-precond-f Va t p) t + 1) * 1 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t,
      lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l (mp-alloc-stm4-pre-precond-f
        Va t p) t + 1) * 2 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t,
      lsizes (mp-alloc-stm4-pre-precond-f Va t p) t ! nat (from-l (mp-alloc-stm4-pre-precond-f
        Va t p) t + 1) * 3 +
      blk (mp-alloc-stm4-pre-precond-f Va t p) t].
      e ∉ set (free-list (levels (mem-pool-info Va p) ! ii))
      using b2 a03 not-in-freelist local.block-n unfolding addr-def apply auto
      by (metis (no-types) not-in-freelist semiring-normalization-rules(12))
    show ?thesis

    using distinct-lists[OF inv-bitmap3 h1 h2] free-list-x[OF a19]
    by (metis a03 eq-free-list-mp-alloc-stm4-pre-precond-f mp-alloc-stm4-pre-froml)

  qed
} ultimately have distinct (free-list (levels (mem-pool-info x p) ! ii))
  by auto

} then show ?thesis by auto
qed

```

**lemma** *mp-alloc-stm4-inv-bitmap-freelist*:  
 assumes  $a0:p \in \text{mem-pools } Va$  and  
 $a1:\neg \text{free-l } Va \ t < OK$  and  
 $a2:\text{free-l } Va \ t \leq \text{from-l } Va \ t$  and  
 $a3:\text{alloc-l } Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p))$  and

```

a4:from-l Va t < alloc-l Va t and
a4':alloc-l Va t = int (length (lsizes Va t)) - 1 ∧ length (lsizes Va t) = n-levels
(mem-pool-info Va p) ∨
alloc-l Va t = int (length (lsizes Va t)) - 2 ∧ lsizes Va t ! nat (alloc-l Va t + 1)
< sz and
a5:block-num (mem-pool-info Va p)
(buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
(lsizes Va t ! nat (from-l Va t))
< n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
a6:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)) and
a7:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and
a8:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) and
a9:freeing-node x = freeing-node (mp-alloc-stm4-pre-precond-f Va t p) and
a10:allocating-node x = allocating-node (mp-alloc-stm4-pre-precond-f Va t p) and
a11:∀ pa. pa ≠ p → mem-pool-info x pa = mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) pa and
a12:∀ jj. jj ≠ nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) →
levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p) ! jj and
a12':∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
a12'':i x t = 4 and
a12''':lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p) and
a12''':length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
a13:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
list-updates-n
(bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))
(Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4)) 3 FREE and
a14:
free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
inserts
(map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
ii +
blk (mp-alloc-stm4-pre-precond-f Va t p) t)
[Suc NULL..<4])
(free-list
(levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
a15: inv-mempool-info Va and
a16:inv-bitmap-freelist Va and
a17:allocating-node Va t =
Some (pool = p, level = nat (from-l Va t),
block = block-num (mem-pool-info Va p)

```

```

      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
      (lsizes Va t ! nat (from-l Va t)),
      data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t))) and
a18:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
Va and
a19:blk x = blk (mp-alloc-stm4-pre-precond-f Va t p)
shows inv-bitmap-freelist x
proof -
{ fix p'
  assume a00:p'∈mem-pools x
  {assume p≠p'
    moreover have mem-pool-info x p' = mem-pool-info Va p'
      using mp-alloc-stm4-pres-mpinfo
      by (metis a11 calculation)
    ultimately have inv-bitmap-freelist-mp x p'
      using a18 a00 mp-alloc-stm4-lvl-len[OF a0 a7] mp-alloc-stm4-maxsz[OF
a0 a7]
      mp-alloc-stm4-buf[OF a0 a7] mp-alloc-stm4-froml[OF a8] mp-alloc-stm4-mempools[OF
a7]
      by (simp add: inv-bitmap-freelist-def Let-def)
  }
  moreover { assume eq-p:p=p'
    let ?mp = mem-pool-info x p'
    have inv-mempool-info-mp Va p'
      using a15 eq-p mp-alloc-stm4-mempools[OF a7] a00 unfolding inv-mempool-info-def
by auto
    note inv-mempool=this[simplified Let-def]
    {fix i
      assume a01:i<length (levels ?mp)
      then have inv-bitmap1:(∀ j<length (bits (levels (mem-pool-info Va p') ! i)).
        (get-bit-s Va p' i j = FREE) =
        (buf (mem-pool-info Va p') + j * (max-sz (mem-pool-info Va
p') div 4 ^ i)
        ∈ set (free-list (levels (mem-pool-info Va p') ! i)))) and
      inv-bitmap2:(∀ j<length (free-list (levels (mem-pool-info Va p') !
i)).
        ∃ n<n-max (mem-pool-info Va p') * 4 ^ i.
        free-list (levels (mem-pool-info Va p') ! i) ! j =
        buf (mem-pool-info Va p') + n * (max-sz (mem-pool-info Va
p') div 4 ^ i)) and
      inv-bitmap3:distinct (free-list (levels (mem-pool-info Va p') ! i))
      using a16 eq-p mp-alloc-stm4-mempools[OF a7] a00 a01 mp-alloc-stm4-lvl-len[OF
a0 a7]
      unfolding Let-def inv-bitmap-freelist-def
      by fastforce+
    let ?bts = bits (levels ?mp ! i) and ?fl = free-list (levels ?mp ! i)
    have f1:(∀ j<length ?bts. (?bts ! j = FREE) = (buf ?mp + j * (max-sz

```

```

?mp div 4 ^ i) ∈ set ?fl))
  using assms free-list-updates-inv1 a00 a01 eq-p by blast
  have f2: (∀ j < length ?fl. ∃ n < n-max ?mp * 4 ^ i. ?fl ! j = buf ?mp + n *
(max-sz ?mp div 4 ^ i))
  using assms free-list-updates-inv2 a00 a01 eq-p by blast
  have f3: distinct ?fl using assms free-list-updates-inv3 a00 a01 eq-p by blast

  note conjI[OF f1 conjI[OF f2 f3]]
} then have inv-bitmap-freelist-mp x p' by auto
}
ultimately have inv-bitmap-freelist-mp x p' by auto
}
thus ?thesis unfolding inv-bitmap-freelist-def by auto
qed

```

**lemma noexists-eq-bits: assumes**  
 $a0: \forall j. j \geq jj \wedge j \leq \text{Suc}(\text{Suc}(\text{Suc } jj)) \longrightarrow$   
 $\text{get-bit-s } x \ p \ ii \ j = \text{get-bit-s } Va \ p \ ii \ j$  **and**  
 $a1: \text{noexist-bits } (\text{mem-pool-info } Va \ p) \ ii \ jj$   
**shows**  $\text{noexist-bits } (\text{mem-pool-info } x \ p) \ ii \ jj$   
**using**  $a0 \ a1$   
**by** *simp*

**lemma mp-alloc-stm4-inv-bitmap1:**  
**assumes**  
 $a0: \text{inv } Va$  **and**  
 $a1: p \in \text{mem-pools } Va$  **and**  
 $a2: \forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 ^ ii$  **and**  
 $a4: \text{alloc-l } Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p))$  **and**  
 $a5: \neg \text{free-l } Va \ t < OK$  **and**  
 $a6: \text{free-l } Va \ t \leq \text{from-l } Va \ t$  **and**  
 $a7: \text{allocating-node } Va \ t =$   
 $\text{Some } (\downarrow \text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
 $\text{block} = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)),$   
 $\text{data} = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{nat } (\text{from-l } Va \ t)))$  **and**  
 $a8: n = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{nat } (\text{from-l } Va \ t) = \text{NULL}$  **and**  
 $a9: \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t))$

$< n\text{-max} \text{ (mem-pool-info } Va \text{ } p) * 4 \wedge \text{nat (from-l } Va \text{ } t) \text{ and}$   
 $a10:\text{from-l } Va \text{ } t < \text{alloc-l } Va \text{ } t \text{ and}$

$a11:n < n\text{-max} \text{ (mem-pool-info } Va \text{ } p) * 4 \wedge \text{nat (from-l } Va \text{ } t) \text{ and}$   
 $a12:\text{blk } Va \text{ } t = \text{buf (mem-pool-info } Va \text{ } p) + n * (\text{max-sz (mem-pool-info } Va \text{ } p)$   
 $\text{div } 4 \wedge \text{nat (from-l } Va \text{ } t)) \text{ and}$   
 $a13:(x, \text{mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \in \text{gvars-conf-stable and}$   
 $a14:\forall jj. jj \neq \text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1) \longrightarrow$   
 $\text{levels (mem-pool-info } x \text{ } p) ! jj = \text{levels (mem-pool-info (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } p) ! jj \text{ and}$   
 $a15:\text{bits (levels (mem-pool-info } x \text{ } p) ! \text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } p) !$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1)))$   
 $(\text{Suc (bn (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t * 4)) \text{ } 3 \text{ FREE and}$   
 $a16:\text{free-list (levels (mem-pool-info } x \text{ } p) ! \text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t !$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1) *$   
 $ii +$   
 $\text{blk (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t)$   
 $[\text{Suc NULL..} < 4])$   
 $(\text{free-list}$   
 $(\text{levels (mem-pool-info (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } p) !$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ } t + 1))) \text{ and}$   
 $a17:\text{lsizes } x = \text{lsizes (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ and}$   
 $a18:\text{from-l } x = \text{from-l (mp-alloc-stm4-pre-precond-f } Va \text{ } t \text{ } p) \text{ and}$   
 $a01:ii < \text{length (levels (mem-pool-info } x \text{ } p)) \text{ and}$   
 $a02:jj < \text{length (bits (levels (mem-pool-info } x \text{ } p) ! ii))$   
**shows**  $(\text{get-bit-s } x \text{ } p \text{ } ii \text{ } jj = \text{FREE} \vee \text{get-bit-s } x \text{ } p \text{ } ii \text{ } jj = \text{FREEING} \vee \text{get-bit-s } x$   
 $p \text{ } ii \text{ } jj = \text{ALLOCATED} \vee \text{get-bit-s } x \text{ } p \text{ } ii \text{ } jj = \text{ALLOCATING} \longrightarrow$   
 $(\text{NULL} < ii \longrightarrow \text{get-bit-s } x \text{ } p \text{ } (ii - 1) (jj \text{ div } 4) = \text{DIVIDED}) \wedge$   
 $(ii < \text{length (levels (mem-pool-info } x \text{ } p)) - 1 \longrightarrow \text{noexist-bits (mem-pool-info}$   
 $x \text{ } p) (ii + 1) (jj * 4)))$   
**proof**–  
**let**  $?mp = \text{mem-pool-info } x \text{ } p$   
**have**  $\text{inv:inv-aux-vars } Va \wedge \text{inv-bitmap } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist}$   
 $Va$   
**using**  $a0 \text{ unfolding inv-def by auto}$   
**have**  $\text{from-l-gt0:} 0 \leq \text{from-l } Va \text{ } t \text{ using } a6 \text{ } a5 \text{ by linarith}$   
**have**  $\text{len-levels:length (levels (mem-pool-info } x \text{ } p)) = \text{length (levels (mem-pool-info}$   
 $Va \text{ } p))$   
**using**  $\text{mp-alloc-stm4-lvl-len[OF } a1 \text{ } a13] \text{ by simp}$   
**have**  $\text{maxsz:max-sz (mem-pool-info } x \text{ } p) = \text{max-sz (mem-pool-info } Va \text{ } p)$   
**using**  $\text{mp-alloc-stm4-maxsz[OF } a1 \text{ } a13] \text{ by simp}$   
**have**  $\text{buf:buf (mem-pool-info } x \text{ } p) = \text{buf (mem-pool-info } Va \text{ } p)$   
**using**  $\text{mp-alloc-stm4-buf[OF } a1 \text{ } a13] \text{ by simp}$



```

have from-l:from-l x = from-l Va
  using mp-alloc-stm4-froml[OF a18] by auto
have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
  using from-l-gt0 by auto
have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a13] by auto
have lsizes-x-va:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a17
  by auto
let ?i1=(nat (from-l Va t)) and
  ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) and
  ?i2=(nat (from-l Va t + 1)) and
  ?j2=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))*4) and
  ?i1'=(nat (from-l Va t)) - 1 and
  ?j1'=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) div 4 and
  ?i2'=(nat (from-l Va t)) + 2
let ?j20'=?j2*4 and ?j21'=(?j2+1)*4 and ?j22'=(?j2+2)*4 and
  ?j23'=(?j2+3)*4 and ?j24'=(?j2+4)*4
let ?mp=mem-pool-info x p
have inv-mempool-info-mp Va p
  using a1 mem-pools inv unfolding inv-mempool-info-def by auto
note inv-mempool=this[simplified Let-def]
have i1-len:?i1 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have i2-len:?i2 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
  by (metis i1-len a9 a12 a1 inv inv-mempool-info-def)
have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
  using i1-len i2-len j1-len inv-mempool from-l-suc
  by simp
let ?bts = bits (levels ?mp ! ii)
let ?btsva = (bits (levels (mem-pool-info Va p) ! ii))
have a01':ii < length (levels (mem-pool-info Va p))
  using a01 len-levels by auto
then have inv-bitmap1:
  ∀ j < length (bits (levels (mem-pool-info Va p) ! ii)).
    (?btsva ! j = FREE ∨ ?btsva ! j = FREEING ∨ ?btsva ! j = ALLOCATED
  ∨ ?btsva ! j = ALLOCATING →
      (ii > 0 → (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (j div
4) = DIVIDED)
      ∧ (ii < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
(mem-pool-info Va p) (ii+1) (j*4) ))
    ∧ (?btsva ! j = DIVIDED → ii > 0 → (bits (levels (mem-pool-info Va

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p) ! (ii - 1))) ! (j div 4) = DIVIDED)
  ∧ (?btsva ! j = NOEXIST → ii < length (levels (mem-pool-info Va p))
- 1
    → noexist-bits (mem-pool-info Va p) (ii+1) (j*4))
  ∧ (?btsva ! j = NOEXIST ∧ ii > 0 → (bits (levels (mem-pool-info Va
p) ! (ii - 1))) ! (j div 4) ≠ DIVIDED)
  using inv mem-pools a1
  unfolding Let-def inv-bitmap-def
  by blast
have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
  using a7 a0 a12 unfolding inv-aux-vars-def invariant inv-def
  by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))

  then have alloc-predi1-j1:?i1 > 0 → get-bit-s Va p (?i1 - 1) (?j1 div 4) =
DIVIDED
    using inv-bitmap1 i1-len j1-len inv a1 unfolding Let-def inv-bitmap-def by
blast
  have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
    using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i1-len j1-len
      alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
  by (smt One-nat-def Suc-pred add commute inv-mempool nat-add-left-cancel-less
plus-1-eq-Suc)
  have nexisti3:?i2 < length (levels (mem-pool-info Va p)) - 1 →
    noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j23'
  proof-
  { assume ?i2 < length (levels (mem-pool-info Va p)) - 1
    then have a00:∀ j < length (bits (levels (mem-pool-info Va p) ! ?i2)).
      get-bit-s Va p ?i2 j = NOEXIST → noexist-bits (mem-pool-info Va
p) ?i2' (j * 4)
      using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i2-len
        from-l-suc by auto
    then have noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
      noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
      noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
      noexist-bits (mem-pool-info Va p) ?i2' ?j23'
    using j2-len nexisti2 Suc-lessD
    by (smt One-nat-def add commute add-2-eq-Suc' add-Suc-right numeral-3-eq-3
plus-1-eq-Suc)
  }
  thus ?thesis by fastforce
qed
let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
have a02':jj < length (bits (levels (mem-pool-info Va p) ! ii))
  using a02 a13 unfolding gvars-conf-def gvars-conf-stable-def

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    by (simp add: mp-alloc-stm4-inv-bits-len)
  have eq-len:length (bits (levels (mem-pool-info x p) ! ii)) =
    length (bits (levels (mem-pool-info Va p) ! ii))
  using mp-alloc-stm4-inv-bits-len a14 a15 length-list-update-n
  by metis
  have inv-va:(?btsva ! jj = FREE  $\vee$  ?btsva ! jj = FREEING  $\vee$  ?btsva ! jj =
    ALLOCATED  $\vee$  ?btsva ! jj = ALLOCATING  $\longrightarrow$ 
    (ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div
    4) = DIVIDED)
     $\wedge$  (ii < length (levels (mem-pool-info Va p)) - 1  $\longrightarrow$  noexist-bits
    (mem-pool-info Va p) (ii+1) (jj*4) ))
     $\wedge$  (?btsva ! jj = DIVIDED  $\longrightarrow$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va
    p) ! (ii - 1))) ! (jj div 4) = DIVIDED)
     $\wedge$  (?btsva ! jj = NOEXIST  $\longrightarrow$  ii < length (levels (mem-pool-info Va p))
    - 1
     $\longrightarrow$  noexist-bits (mem-pool-info Va p) (ii+1) (jj*4))
     $\wedge$  (?btsva ! jj = NOEXIST  $\wedge$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p)
    ! (ii - 1))) ! (jj div 4)  $\neq$  DIVIDED)
  using inv-bitmap1 a02' by auto
  { assume a05: $\neg$ ((ii=?i1  $\wedge$  jj=?j1)  $\vee$ 
    (ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj< ?j2+4)  $\vee$ 
    (ii=?i2'  $\wedge$  jj $\geq$  ?j20'  $\wedge$  jj< ?j24')  $\vee$ 
    (?i1 > 0  $\wedge$  ii = (?i1 - 1)  $\wedge$  jj = ?j1 div 4))
  then have a050': $\neg$ (ii=?i1  $\wedge$  jj=?j1) and
    a051':  $\neg$ (ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj< ?j2 + 4) and
    a052': $\neg$ (ii=?i2'  $\wedge$  jj $\geq$  ?j20'  $\wedge$  jj< ?j24') and
    a053':  $\neg$ (?i1 > 0  $\wedge$  ii = (?i1 - 1)  $\wedge$  jj = ?j1 div 4)
  by force+
  have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
  a18], THEN sym]] a15, of ii jj]
  using a050' a051' by auto
  have eq-get-bit-i1-j1:ii>0  $\longrightarrow$  get-bit-s x p (ii-1) (jj div 4) = get-bit-s Va p
  (ii-1) (jj div 4)
  proof-
  { assume a06:ii>0
  then have  $\neg$ ((ii - 1) = ?i1  $\wedge$  jj div 4 = ?j1)
  using a050' a051' from-l-suc by fastforce
  moreover have  $\forall j. j \geq ?j2 \wedge j \leq ?j2+3 \longrightarrow \neg((ii - 1) = ?i2 \wedge jj \text{ div } 4 =$ 
  j)
  using a051' a052' from-l-gt0 by fastforce
  ultimately have get-bit-s x p (ii-1) (jj div 4) = get-bit-s Va p (ii-1) (jj
  div 4)
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
  a18], THEN sym]] a15,
  of ii - 1 jj div 4] by auto
  } thus ?thesis by auto qed
  have eq-get-bit-i2-j2: $\forall j. j \geq (jj * 4) \wedge j \leq \text{Suc}(\text{Suc}(\text{Suc}(jj * 4))) \longrightarrow$ 
  get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j

```

```

proof-
{ fix j
  assume a00:  $j \geq (jj * 4) \wedge j \leq (jj * 4) + 3$ 
  then have n1:  $\neg((ii + 1) = ?i1 \wedge j = ?j1)$ 
    using a053' from-l-suc by auto
  have n2:  $\forall j. j \geq ?j2 \wedge j \leq ?j2 + 3 \longrightarrow \neg((ii + 1) = ?i2 \wedge jj * 4 = j)$ 
    using a050' from-l-gt0 by fastforce
  have get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18], THEN sym]] a15,
    of ii + 1 j] n1 n2 a00
    apply (cases j=jj*4) by auto
  } thus ?thesis by auto
qed
{ assume a06: get-bit-s x p ii jj = FREE  $\vee$ 
  get-bit-s x p ii jj = FREEING  $\vee$ 
  get-bit-s x p ii jj = ALLOCATED  $\vee$ 
  get-bit-s x p ii jj = ALLOCATING
  { assume a07: NULL < ii
    then have get-bit-s x p (ii - 1) (jj div 4) = DIVIDED
      using a06 a07 eq-get-bit-i1-j1 eq-get-bit-i-j
      using inv-va by auto
    }
  moreover {
    assume a07:  $ii < \text{length}(\text{levels}(\text{mem-pool-info } x \ p)) - 1$ 
    then have ilen:  $ii < \text{length}(\text{levels}(\text{mem-pool-info } Va \ p)) - 1$ 
      by (simp add: len-levels)
    have get-bit-s Va p ii jj = FREE  $\vee$ 
      get-bit-s Va p ii jj = FREEING  $\vee$ 
      get-bit-s Va p ii jj = ALLOCATED  $\vee$ 
      get-bit-s Va p ii jj = ALLOCATING using eq-get-bit-i-j a06 by auto

    then have noexist-bits (mem-pool-info Va p) (ii + 1) (jj * 4)
      using ilen inv-va
      by simp
    then have noexist-bits (mem-pool-info x p) (ii + 1) (jj * 4)
      using eq-get-bit-i2-j2 by (simp add: numeral-3-eq-3)
    }
  ultimately have ?thesis by auto
} then have ?thesis by auto
}
moreover {
  assume a06:  $(ii = ?i1 \wedge jj = ?j1)$ 
  then have get-bit-s x p ii jj = DIVIDED
    using get-bit-x-l-b a14 a18 from-l from-l-gt0 i1-len j1-len by presburger
  then have ?thesis by auto
}
moreover {
  assume a06:  $(ii = ?i2 \wedge jj \geq ?j2 \wedge jj < ?j2 + 4)$ 

```

```

then have a06':jj=?j2  $\vee$  jj=?j2+1  $\vee$  jj=?j2+2  $\vee$  jj=?j2 + 3 by auto
{ assume a07:NULL < ii
  { assume a08:jj=?j2
    then have get-bit:get-bit-s x p ii jj = ALLOCATING
    using a02 a06 a15 eq-len get-bit-x-l1-b4 i2-len from-l-gt0 i1-len j1-len
    by (metis mult.commute)
    then have get-bit-s x p (ii-1) (jj div 4) = DIVIDED
    using a06 a08 get-bit-x-l-b a14 a18 from-l from-l-gt0 i1-len j1-len
    by (simp add: a18 i1-len j1-len from-l-suc)
  }
  moreover {
    assume a07:jj $\neq$ ?j2
    have a07':jj div 4 = ?j1 using a06 a07 by auto
    have get-bit-s x p ii jj = FREE
    using a06 a02 a15 a07 from-l mp-alloc-stm4-inv-bits-len a18 mp-alloc-stm4-pre-precond-f-bn
    by (auto simp add: mp-alloc-stm4-pre-precond-f-bn)
    have get-bit-s x p (ii-1) (jj div 4) = DIVIDED
    using a06 a07' a14 a18 from-l from-l-gt0 i1-len j1-len
    by (simp add: a18 get-bit-x-l-b i1-len j1-len from-l-suc)
  }
  ultimately have get-bit-s x p (ii-1) (jj div 4) = DIVIDED by fastforce
}
moreover { assume a07:ii < length (levels (mem-pool-info x p)) - 1
  then have get-s: $\forall j. j \geq (jj * 4) \wedge j \leq \text{Suc}(\text{Suc}(\text{Suc}(jj * 4))) \longrightarrow$ 
    get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18],
    THEN sym]] a15, of ii + 1 jj*4] a06
    by (metis Suc-1 Suc-eq-plus1 a14 a18 add.right-neutral add-Suc-right
add-left-cancel
    from-l from-l-suc same-bit-mp-alloc-stm4-pre-precond-f1 zero-neg-numeral)
  then have noexist-bits (mem-pool-info x p) (ii + 1) (jj*4)
    using a07[simplified len-levels] a06 inv-va neristi2
    noexists-eq-bits[OF get-s] a06'
    by fastforce
}
ultimately have ?thesis by fastforce
}
moreover {
  assume a06: (ii=?i2'  $\wedge$  jj $\geq$  ?j20'  $\wedge$  jj < ?j24')
  then have a06':jj=?j20'  $\vee$  jj=?j20'+1  $\vee$  jj=?j20'+2  $\vee$  jj=?j20'+3  $\vee$ 
    jj=?j21'  $\vee$  jj=?j21'+1  $\vee$  jj=?j21'+2  $\vee$  jj=?j21'+3  $\vee$ 
    jj=?j22'  $\vee$  jj=?j22'+1  $\vee$  jj=?j22'+2  $\vee$  jj=?j22'+3  $\vee$ 
    jj=?j23'  $\vee$  jj=?j23'+1  $\vee$  jj=?j23'+2  $\vee$  jj=?j23' + 3
    by presburger
  then have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18],
    THEN sym]] a15, of ii jj] using a06

```

```

    by (simp add: from-l-suc)
    have i2-lt-length: ?i2 < length (levels (mem-pool-info Va p)) - 1 using a06
a01
    by (simp add: len-levels)
    { assume a07: get-bit-s x p ii jj = FREE ∨
      get-bit-s x p ii jj = FREEING ∨
      get-bit-s x p ii jj = ALLOCATED ∨
      get-bit-s x p ii jj = ALLOCATING
      have get-bit-s Va p ii jj = NOEXIST
        using a07 a06 inv-va nexisti3[simplified i2-lt-length] a06'
        by auto
      then have get-bit-s x p ii jj = NOEXIST using eq-get-bit-i-j by auto
    } then have ?thesis by auto
  }
  moreover {
    assume a06: (?i1 > 0 ∧ ii = (?i1 - 1) ∧ jj = ?j1 div 4)
    then have eq-get-bit-i-j: get-bit-s x p ii jj = get-bit-s Va p ii jj
    using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18],
                                                                    THEN sym]] a15, of ii jj]
    by linarith
    then have get-bit-divided: get-bit-s x p ii jj = DIVIDED using a06 alloc-predi1-j1
  by simp
    { assume a06: get-bit-s x p ii jj = FREE ∨ get-bit-s x p ii jj = FREEING ∨
      get-bit-s x p ii jj = ALLOCATED ∨ get-bit-s x p ii jj = ALLOCATING
      then have ?thesis using get-bit-divided by auto
    } then have ?thesis by fastforce
  }
  ultimately show ?thesis by fastforce
qed

```

**lemma** mp-alloc-stm4-inv-bitmap2:

```

  assumes
    a0: inv Va and
    a1: p ∈ mem-pools Va and
    a2: ∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
    a3: length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
    a4: alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a5: ¬ free-l Va t < OK and
    a6: free-l Va t ≤ from-l Va t and
    a7: allocating-node Va t =
      Some (pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t)),
        data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t))) and

```

$a8:n = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL and}$   
 $a9:\text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t))$   
 $< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \text{ and}$   
 $a10:\text{from-l } Va \ t < \text{alloc-l } Va \ t \text{ and}$   
 $a11:\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) \text{ and}$   
 $a12:(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable and}$   
 $a13:\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels } (\text{mem-pool-info } x \ p) ! jj = \text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj \text{ and}$   
 $a14:\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$   
 $(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE and}$   
 $a15:\text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $ii +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))) \text{ and}$   
 $a16:\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a17:\text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a01:ii < \text{length } (\text{levels } (\text{mem-pool-info } x \ p)) \text{ and}$   
 $a02:jj < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! ii)) \text{ and}$   
 $a03:\text{get-bit-s } x \ p \ ii \ jj = \text{DIVIDED and}$   
 $a04:0 < ii$   
**shows**  $\text{get-bit-s } x \ p \ (ii - 1) \ (jj \text{ div } 4) = \text{DIVIDED}$   
**proof** –  
**let**  $?mp = \text{mem-pool-info } x \ p$   
**have**  $\text{inv}:\text{inv-aux-vars } Va \wedge \text{inv-bitmap } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist } Va$   
**using**  $a0$  **unfolding**  $\text{inv-def}$  **by**  $\text{auto}$   
**have**  $\text{from-l-gt0}:0 \leq \text{from-l } Va \ t$  **using**  $a6 \ a5$  **by**  $\text{linarith}$   
**have**  $\text{len-levels}:\text{length } (\text{levels } (\text{mem-pool-info } x \ p)) = \text{length } (\text{levels } (\text{mem-pool-info } Va \ p))$

```

    using mp-alloc-stm4-lvl-len[OF a1 a12] by simp
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
    using mp-alloc-stm4-maxsz[OF a1 a12] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
    using mp-alloc-stm4-buf[OF a1 a12] by simp
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a17] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a12] by auto
  have lsize-x-va:lsize x = lsize Va using mp-alloc-stm4-pre-precond-f-lsz a16
    by auto
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))*4) and
    ?i1' = (nat (from-l Va t)) - 1 and
    ?j1' = (block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))) div 4 and
    ?i2' = (nat (from-l Va t)) + 2
  let ?j20' = ?j2 * 4 and ?j21' = (?j2+1) * 4 and ?j22' = (?j2+2)*4 and
    ?j23' = (?j2+3)*4 and ?j24' = (?j2+4)*4
  let ?mp = mem-pool-info x p
  have inv-mempool-info-mp Va p
    using a1 mem-pools inv unfolding inv-mempool-info-def by auto
  note inv-mempool=this[simplified Let-def]
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis i1-len a9 a11 a1 inv inv-mempool-info-def)
  have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp
  let ?bts = bits (levels ?mp ! ii)
  let ?btsva = (bits (levels (mem-pool-info Va p) ! ii))
  have a01':ii < length (levels (mem-pool-info Va p))
    using a01 len-levels by auto
  then have inv-bitmap1:
     $\forall j < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } Va \ p) \ ! \ ii)).$ 
     $(?btsva \ ! \ j = \text{FREE} \vee ?btsva \ ! \ j = \text{FREEING} \vee ?btsva \ ! \ j = \text{ALLOCATED}$ 
 $\vee ?btsva \ ! \ j = \text{ALLOCATING} \longrightarrow$ 

```



```

(ii > 0 → (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (j div
4) = DIVIDED)
  ∧ (ii < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
(mem-pool-info Va p) (ii+1) (j*4) ))
  ∧ (?btsva ! j = DIVIDED → ii > 0 → (bits (levels (mem-pool-info Va
p) ! (ii - 1))) ! (j div 4) = DIVIDED)
  ∧ (?btsva ! j = NOEXIST → ii < length (levels (mem-pool-info Va p))
- 1
    → noexist-bits (mem-pool-info Va p) (ii+1) (j*4))
  ∧ (?btsva ! j = NOEXIST ∧ ii > 0 → (bits (levels (mem-pool-info Va
p) ! (ii - 1))) ! (j div 4) ≠ DIVIDED)
  using inv mem-pools a1
  unfolding Let-def inv-bitmap-def
  by blast
have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
  using a7 a0 a11 unfolding inv-aux-vars-def invariant.inv-def
  by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))

then have alloc-predi1-j1:?i1 > 0 → get-bit-s Va p (?i1 - 1) (?j1 div 4) =
DIVIDED
  using inv-bitmap1 i1-len j1-len inv a1 unfolding Let-def inv-bitmap-def by
blast
have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
  using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i1-len j1-len
  alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
  by (smt One-nat-def Suc-pred add commute inv-mempool nat-add-left-cancel-less
plus-1-eq-Suc)
have nexisti3:?i2 < length (levels (mem-pool-info Va p)) - 1 →
  noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
  noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
  noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
  noexist-bits (mem-pool-info Va p) ?i2' ?j23'
proof-
{ assume ?i2 < length (levels (mem-pool-info Va p)) - 1
  then have a00:∀ j < length (bits (levels (mem-pool-info Va p) ! ?i2)).
    get-bit-s Va p ?i2 j = NOEXIST → noexist-bits (mem-pool-info Va
p) ?i2' (j * 4)
    using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i2-len
    from-l-suc by auto
  then have noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j23'
    using j2-len nexisti2 Suc-lessD
  by (smt One-nat-def add commute add-2-eq-Suc' add-Suc-right numeral-3-eq-3
plus-1-eq-Suc)
}

```

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    thus ?thesis by fastforce
qed
let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
have a02':jj < length (bits (levels (mem-pool-info Va p) ! ii))
  using a02 a12 unfolding gvars-conf-def gvars-conf-stable-def
  by (simp add: mp-alloc-stm4-inv-bits-len)
have eq-len:length (bits (levels (mem-pool-info x p) ! ii)) =
  length (bits (levels (mem-pool-info Va p) ! ii))
  using mp-alloc-stm4-inv-bits-len a13 a14 length-list-update-n
  by metis
have inv-va:(?btsva ! jj = FREE  $\vee$  ?btsva ! jj = FREEING  $\vee$  ?btsva ! jj =
  ALLOCATED  $\vee$  ?btsva ! jj = ALLOCATING  $\longrightarrow$ 
  (ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div
  4) = DIVIDED)
   $\wedge$  (ii < length (levels (mem-pool-info Va p)) - 1  $\longrightarrow$  noexist-bits
  (mem-pool-info Va p) (ii+1) (jj*4))
   $\wedge$  (?btsva ! jj = DIVIDED  $\longrightarrow$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va
  p) ! (ii - 1))) ! (jj div 4) = DIVIDED)
   $\wedge$  (?btsva ! jj = NOEXIST  $\longrightarrow$  ii < length (levels (mem-pool-info Va p))
  - 1
   $\longrightarrow$  noexist-bits (mem-pool-info Va p) (ii+1) (jj*4))
   $\wedge$  (?btsva ! jj = NOEXIST  $\wedge$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p)
  ! (ii - 1))) ! (jj div 4)  $\neq$  DIVIDED)
  using inv-bitmap1 a02' by auto
{ assume a05: $\neg$ ((ii=?i1  $\wedge$  jj=?j1)  $\vee$ 
  (ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj < ?j2+4)  $\vee$ 
  (ii=?i2'  $\wedge$  jj $\geq$  ?j20'  $\wedge$  jj < ?j24')  $\vee$ 
  (?i1 > 0  $\wedge$  ii = (?i1 - 1)  $\wedge$  jj = ?j1 div 4))
  then have a050': $\neg$ ((ii=?i1  $\wedge$  jj=?j1) and
    a051':  $\neg$ ((ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj < ?j2 + 4) and
    a052': $\neg$ ((ii=?i2'  $\wedge$  jj $\geq$  ?j20'  $\wedge$  jj < ?j24') and
    a053':  $\neg$ ((i1 > 0  $\wedge$  ii = (i1 - 1)  $\wedge$  jj = ?j1 div 4))
    by force+
  have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
  a17], THEN sym]] a14, of ii jj]
  using a050' a051' by auto
  have eq-get-bit-i1-j1:ii>0  $\longrightarrow$  get-bit-s x p (ii-1) (jj div 4) = get-bit-s Va p
  (ii-1) (jj div 4)
  proof-
  { assume a06:ii>0
    then have  $\neg$ ((ii - 1) = ?i1  $\wedge$  jj div 4 = ?j1)
      using a050' a051' from-l-suc by fastforce
    moreover have  $\forall j. j \geq ?j2 \wedge j \leq ?j2+3 \longrightarrow \neg((ii - 1) = ?i2 \wedge jj \text{ div } 4 =$ 
    j)
      using a051' a052' from-l-gt0 by fastforce
    ultimately have get-bit-s x p (ii-1) (jj div 4) = get-bit-s Va p (ii-1) (jj
    div 4)
    using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF

```

```

a17], THEN sym]] a14,
  of ii - 1 jj div 4] by auto
} thus ?thesis by auto qed
have eq-get-bit-i2-j2:∀ j. j ≥ (jj * 4) ∧ j ≤ Suc(Suc(Suc(jj * 4))) →
  get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
proof-
{ fix j
  assume a00:j ≥ (jj * 4) ∧ j ≤ (jj * 4)+3
  then have n1:¬((ii + 1) = ?i1 ∧ j = ?j1)
    using a053' from-l-suc by auto
  have n2:∀ j. j ≥ ?j2 ∧ j ≤ ?j2+3 → ¬((ii + 1) = ?i2 ∧ jj * 4 = j)
    using a050' from-l-gt0 by fastforce
  have get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
  using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
a17], THEN sym]] a14,
    of ii + 1 j] n1 n2 a00
  apply (cases j=jj*4) by auto
} thus ?thesis by auto
qed
then have ?thesis
  using a03 a04 eq-get-bit-i1-j1 eq-get-bit-i-j inv-va by auto
}
moreover {
  assume a06:(ii=?i1 ∧ jj=?j1)
  then have ?thesis using a03 a04
    by (metis Suc-eq-plus1 Suc-pred a13 a17
      add commute add-2-eq-Suc' add-cancel-right-right
      alloc-predi1-j1 from-l from-l-suc
      same-bit-mp-alloc-stm4-pre-precond-f zero-neq-numeral)
}
moreover {
  assume a06:(ii=?i2 ∧ jj ≥ ?j2 ∧ jj < ?j2+4)
  then have a06':jj=?j2 ∨ jj=?j2+1 ∨ jj=?j2+2 ∨ jj=?j2+3 by auto
  have l1:(ii - 1) = ?i1 ∧ (jj div 4) = ?j1
    using a2 a03 a04 a01 a02 a02' a06 a13 a14 a06'
      from-l-gt0 from-l-suc
  by (metis add commute add-mult-distrib2
    diff-add-inverse2 div-nat-eqI mult commute
    nat-mult-1-right plus-1-eq-Suc)
  then have ?thesis using get-bit-x-l-b[OF l1] a13
    from-l from-l-gt0 from-l-suc i1-len
    by (simp add: a17 j1-len l1)
}
moreover {
  assume a06:(ii=?i2' ∧ jj ≥ ?j20' ∧ jj < ?j24')
  then have a06':jj=?j20' ∨ jj=?j20'+1 ∨ jj=?j20'+2 ∨ jj=?j20'+3 ∨
    jj=?j21' ∨ jj=?j21'+1 ∨ jj=?j21'+2 ∨ jj=?j21'+3 ∨
    jj=?j22' ∨ jj=?j22'+1 ∨ jj=?j22'+2 ∨ jj=?j22'+3 ∨
    jj=?j23' ∨ jj=?j23'+1 ∨ jj=?j23'+2 ∨ jj=?j23'+3

```

```

    by presburger
  then have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
a17],
                                                                    THEN sym]] a14, of ii jj] using a06

    by (simp add: from-l-suc)
  moreover have i2-lt-length:?i2 < length (levels (mem-pool-info Va p)) - 1
using
  a06[simplified len-levels] a01[simplified len-levels]
  by simp
  then have get-bit-s Va p ii jj = NOEXIST
  using a06 a01 a06 inv-va nexisti3 a06' a06[simplified len-levels] a01[simplified
len-levels]
  by auto
  ultimately have ?thesis
  using a03 by auto
}
moreover {
  assume a06: (?i1 > 0 ∧ ii = (?i1 - 1) ∧ jj = ?j1 div 4)
  then have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
a17],
                                                                    THEN sym]] a14, of ii jj]

    by linarith
  have ?thesis using a03 a04 a13 a17 from-l inv-va same-bit-mp-alloc-stm4-pre-precond-f1
    calculation(1) calculation(2) calculation(3) calculation(4)
  by (smt Suc-pred add-diff-cancel-left' int-nat-eq inv-va of-nat-Suc plus-1-eq-Suc)

}
ultimately show ?thesis by fastforce
qed

```

**lemma** mp-alloc-stm4-inv-bitmap3:

```

  assumes
    a0:inv Va and
    a1:p ∈ mem-pools Va and
    a2:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
    a4:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a5:¬ free-l Va t < OK and
    a6:free-l Va t ≤ from-l Va t and
    a7:allocating-node Va t =
      Some (pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t)),
        data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t))) and

```

$a8:n = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL} \text{ and}$   
 $a9:\text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t))$   
 $< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \text{ and}$   
 $a10:\text{from-l } Va \ t < \text{alloc-l } Va \ t \text{ and}$

$a11:n < n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \text{ and}$   
 $a12:\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) \text{ and}$   
 $a13:(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \text{ and}$   
 $a14:\forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels } (\text{mem-pool-info } x \ p) ! jj = \text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj \text{ and}$   
 $a15:\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$   
 $(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE} \text{ and}$   
 $a16:\text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $ii +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))) \text{ and}$   
 $a17:\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a18:\text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a01:ii < \text{length } (\text{levels } (\text{mem-pool-info } x \ p)) \text{ and}$   
 $a02:jj < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! ii)) \text{ and}$   
 $a03:\text{get-bit-s } x \ p \ ii \ jj = \text{NOEXIST} \text{ and}$   
 $a04:ii < \text{length } (\text{levels } (\text{mem-pool-info } x \ p)) - 1$   
**shows**  $\text{noexist-bits } (\text{mem-pool-info } x \ p) \ (ii + 1) \ (jj * 4)$   
**proof** –  
**let**  $?mp = \text{mem-pool-info } x \ p$   
**have**  $\text{inv}:\text{inv-aux-vars } Va \wedge \text{inv-bitmap } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist } Va$   
**using**  $a0$  **unfolding**  $\text{inv-def}$  **by**  $\text{auto}$   
**have**  $\text{from-l-gt0}:0 \leq \text{from-l } Va \ t$  **using**  $a6 \ a5$  **by**  $\text{linarith}$

```

have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
Va p))
  using mp-alloc-stm4-lvl-len[OF a1 a13] by simp
have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
  using mp-alloc-stm4-maxsz[OF a1 a13] by simp
have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
  using mp-alloc-stm4-buf[OF a1 a13] by simp
have from-l:from-l x = from-l Va
  using mp-alloc-stm4-froml[OF a18] by auto
have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
  using from-l-gt0 by auto
have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a13] by auto
  have lsize-x-va:lsize x = lsize Va using mp-alloc-stm4-pre-precond-f-lsz a17
  by auto
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))) and
    ?i2=(nat (from-l Va t + 1)) and
    ?j2=(block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))*4) and
    ?i1'=(nat (from-l Va t)) - 1 and
    ?j1'=(block-num (mem-pool-info Va p) (blk Va t) (lsize Va t ! nat (from-l Va
t))) div 4 and
    ?i2'=(nat (from-l Va t)) + 2
  let ?j20'=?j2 * 4 and ?j21'=(?j2+1) * 4 and ?j22'=(?j2+2)*4 and
    ?j23'=(?j2+3)*4 and ?j24'=(?j2+4)*4
  let ?mp = mem-pool-info x p
have inv-mempool-info-mp Va p
  using a1 mem-pools inv unfolding inv-mempool-info-def by auto
note inv-mempool=this[simplified Let-def]
have i1-len:?i1 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have i2-len:?i2 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
  by (metis i1-len a9 a12 a1 inv inv-mempool-info-def)
have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
  using i1-len i2-len j1-len inv-mempool from-l-suc
  by simp
let ?bts = bits (levels ?mp ! ii)
let ?btsva = (bits (levels (mem-pool-info Va p) ! ii))
have a01':ii < length (levels (mem-pool-info Va p))
  using a01 len-levels by auto
then have inv-bitmap1:
   $\forall j < \text{length} (\text{bits} (\text{levels} (\text{mem-pool-info } Va \ p) ! ii)).$ 

```

```

    (?btsva ! j = FREE ∨ ?btsva ! j = FREEING ∨ ?btsva ! j = ALLOCATED
    ∨ ?btsva ! j = ALLOCATING →
      (ii > 0 → (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (j div
      4) = DIVIDED)
      ∧ (ii < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
      (mem-pool-info Va p) (ii+1) (j*4) ))
      ∧ (?btsva ! j = DIVIDED → ii > 0 → (bits (levels (mem-pool-info Va
      p) ! (ii - 1))) ! (j div 4) = DIVIDED)
      ∧ (?btsva ! j = NOEXIST → ii < length (levels (mem-pool-info Va p))
      - 1
      → noexist-bits (mem-pool-info Va p) (ii+1) (j*4))
      ∧ (?btsva ! j = NOEXIST ∧ ii > 0 → (bits (levels (mem-pool-info Va
      p) ! (ii - 1))) ! (j div 4) ≠ DIVIDED)
    using inv mem-pools a1
    unfolding Let-def inv-bitmap-def
    by blast
  have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
    using a7 a0 a12 unfolding inv-aux-vars-def invariant.inv-def
    by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))

  then have alloc-predi1-j1:?i1 > 0 → get-bit-s Va p (?i1 - 1) (?j1 div 4) =
  DIVIDED
    using inv-bitmap1 i1-len j1-len inv a1 unfolding Let-def inv-bitmap-def by
    blast
  have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
    using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
    i1-len j1-len
    alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
  by (smt One-nat-def Suc-pred add commute inv-mempool nat-add-left-cancel-less
  plus-1-eq-Suc)
  have nexisti3:?i2 < length (levels (mem-pool-info Va p)) - 1 →
    noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j23'
  proof-
    { assume ?i2 < length (levels (mem-pool-info Va p)) - 1
      then have a00:∀ j < length (bits (levels (mem-pool-info Va p) ! ?i2)).
        get-bit-s Va p ?i2 j = NOEXIST → noexist-bits (mem-pool-info Va
        p) ?i2' (j * 4)
        using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
        i2-len
        from-l-suc by auto
      then have noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
        noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
        noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
        noexist-bits (mem-pool-info Va p) ?i2' ?j23'
      using j2-len nexisti2 Suc-lessD
      by (smt One-nat-def add commute add-2-eq-Suc' add-Suc-right numeral-3-eq-3)
    }

```

```

plus-1-eq-Suc)
}
thus ?thesis by fastforce
qed
let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
have a02':jj < length (bits (levels (mem-pool-info Va p) ! ii))
  using a02 a13 unfolding gvars-conf-def gvars-conf-stable-def
  by (simp add: mp-alloc-stm4-inv-bits-len)
have eq-len:length (bits (levels (mem-pool-info x p) ! ii)) =
  length (bits (levels (mem-pool-info Va p) ! ii))
  using mp-alloc-stm4-inv-bits-len a14 a15 length-list-update-n
  by metis
have inv-va:(?btsva ! jj = FREE  $\vee$  ?btsva ! jj = FREEING  $\vee$  ?btsva ! jj =
  ALLOCATED  $\vee$  ?btsva ! jj = ALLOCATING  $\longrightarrow$ 
    (ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div
4) = DIVIDED)
     $\wedge$  (ii < length (levels (mem-pool-info Va p)) - 1  $\longrightarrow$  noexist-bits
(mem-pool-info Va p) (ii+1) (jj*4) ))
     $\wedge$  (?btsva ! jj = DIVIDED  $\longrightarrow$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va
p) ! (ii - 1))) ! (jj div 4) = DIVIDED)
     $\wedge$  (?btsva ! jj = NOEXIST  $\longrightarrow$  ii < length (levels (mem-pool-info Va p))
- 1
     $\longrightarrow$  noexist-bits (mem-pool-info Va p) (ii+1) (jj*4))
     $\wedge$  (?btsva ! jj = NOEXIST  $\wedge$  ii > 0  $\longrightarrow$  (bits (levels (mem-pool-info Va p)
! (ii - 1))) ! (jj div 4)  $\neq$  DIVIDED)
  using inv-bitmap1 a02' by auto
{ assume a05: $\neg$ ((ii=?i1  $\wedge$  jj=?j1)  $\vee$ 
  (ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj < ?j2+4)  $\vee$ 
  (?i1 > 0  $\wedge$  ii = (?i1 - 1)  $\wedge$  jj = ?j1 div 4))
  then have a050': $\neg$ ((ii=?i1  $\wedge$  jj=?j1) and
    a051':  $\neg$ ((ii=?i2  $\wedge$  jj $\geq$  ?j2  $\wedge$  jj < ?j2 + 4) and
    a053':  $\neg$ ((?i1 > 0  $\wedge$  ii = (?i1 - 1)  $\wedge$  jj = ?j1 div 4))
    by force+
  have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18], THEN sym]] a15, of ii jj]
  using a050' a051' by auto
  have eq-get-bit-i2-j2: $\forall j. j \geq (jj * 4) \wedge j \leq \text{Suc}(\text{Suc}(\text{Suc}(jj * 4))) \longrightarrow$ 
    get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
  proof-
  { fix j
    assume a00:j $\geq$  (jj * 4)  $\wedge$  j $\leq$  (jj * 4)+3
    then have n1: $\neg$ ((ii + 1) = ?i1  $\wedge$  j = ?j1)
      using a053' from-l-suc by auto
    have n2: $\forall j. j \geq ?j2 \wedge j \leq ?j2+3 \longrightarrow \neg((ii + 1) = ?i2 \wedge jj * 4 = j)$ 
      using a050' from-l-gt0 by fastforce
    have get-bit-s x p (ii+1) j = get-bit-s Va p (ii+1) j
    using same-bit-mp-alloc-x-va[OF a14[simplified a18[simplified mp-alloc-stm4-froml[OF
a18], THEN sym]] a15,

```



```

      of  $ii + 1\ j]$   $n1\ n2\ a00$ 
      apply (cases  $j=jj*4$ ) by auto
    } thus ?thesis by auto
  qed
  then have ?thesis
    using  $a04\ len\ levels\ eq\ get\ bit\ i\ j\ a03\ inv\ va$  by (simp add: numeral-3-eq-3)
}
moreover {
  assume  $a06:(ii=?i1 \wedge jj=?j1)$ 
  then have  $get\ bit\ s\ x\ p\ ii\ jj = DIVIDED$ 
    using  $get\ bit\ x\ l\ b\ a14\ a18\ from\ l\ from\ l\ gt0\ i1\ len\ j1\ len$  by presburger
  then have ?thesis using  $a03$  by auto
}
moreover {
  assume  $a06:(ii=?i2 \wedge jj \geq ?j2 \wedge jj < ?j2+4)$ 
  then have  $a06':jj=?j2 \vee jj=?j2+1 \vee jj=?j2+2 \vee jj=?j2+3$  by auto
  then have  $get\ s:\forall j. j \geq (jj * 4) \wedge j \leq Suc(Suc(Suc(jj * 4))) \longrightarrow$ 
     $get\ bit\ s\ x\ p\ (ii+1)\ j = get\ bit\ s\ Va\ p\ (ii+1)\ j$ 
    using  $same\ bit\ mp\ alloc\ x\ va[OF\ a14[simplified\ a18[simplified\ mp\ alloc\ stm4\ from\ l[OF\ a18],$ 
       $THEN\ sym]]\ a15, of\ ii + 1\ jj*4]\ a06$ 
    by (metis  $Suc\ 1\ Suc\ eq\ plus1\ a14\ a18\ add.\ right\ neutral\ add\ Suc\ right\ add\ left\ cancel$ 
       $from\ l\ from\ l\ suc\ same\ bit\ mp\ alloc\ stm4\ pre\ precond\ f1\ zero\ neq\ numeral$ )
  then have ?thesis
    using  $a04[simplified\ len\ levels]\ a06\ inv\ va\ nexisti2$ 
     $noexists\ eq\ bits[OF\ get\ s]\ a06'$ 
    by fastforce
}
moreover {
  assume  $a06:(?i1 > 0 \wedge ii = (?i1 - 1) \wedge jj = ?j1\ div\ 4)$ 
  then have  $eq\ get\ bit\ i\ j: get\ bit\ s\ x\ p\ ii\ jj = get\ bit\ s\ Va\ p\ ii\ jj$ 
    using  $same\ bit\ mp\ alloc\ x\ va[OF\ a14[simplified\ a18[simplified\ mp\ alloc\ stm4\ from\ l[OF\ a18],$ 
       $THEN\ sym]]\ a15, of\ ii\ jj]$ 
    by linarith
  then have  $get\ bit\ divided: get\ bit\ s\ x\ p\ ii\ jj = DIVIDED$  using  $a06\ alloc\ predi1\ j1$ 
  by simp
  then have ?thesis using  $get\ bit\ divided\ a03$  by auto
}
ultimately show ?thesis by fastforce
qed

lemma  $mp\ alloc\ stm4\ inv\ bitmap4$ :
  assumes
     $a0:inv\ Va$  and
     $a1:p \in mem\ pools\ Va$  and
     $a2:\forall ii < length\ (lsizes\ Va\ t). lsize\ Va\ t\ ii = ALIGN4\ (max\ sz\ (mem\ pool\ info\ Va\ p))\ div\ 4 \wedge ii$  and

```

$a3: \text{length } (lsizes \text{ Va } t) \leq n\text{-levels } (mem\text{-pool-info } \text{Va } p) \text{ and}$   
 $a4: \text{alloc-l } \text{Va } t < \text{int } (n\text{-levels } (mem\text{-pool-info } \text{Va } p)) \text{ and}$   
 $a5: \neg \text{free-l } \text{Va } t < OK \text{ and}$   
 $a6: \text{free-l } \text{Va } t \leq \text{from-l } \text{Va } t \text{ and}$   
 $a7: \text{allocating-node } \text{Va } t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } \text{Va } t),$   
 $\text{block} = \text{block-num } (mem\text{-pool-info } \text{Va } p)$   
 $(\text{buf } (mem\text{-pool-info } \text{Va } p) + n * (\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t)))$   
 $(lsizes \text{ Va } t ! \text{nat } (\text{from-l } \text{Va } t)),$   
 $\text{data} = \text{buf } (mem\text{-pool-info } \text{Va } p) + n * (\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t))) \text{ and}$   
 $a8: n = \text{block-num } (mem\text{-pool-info } \text{Va } p)$   
 $(\text{buf } (mem\text{-pool-info } \text{Va } p) + n * (\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t)))$   
 $(lsizes \text{ Va } t ! \text{nat } (\text{from-l } \text{Va } t)) \vee$   
 $\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t) = \text{NULL} \text{ and}$   
 $a9: \text{block-num } (mem\text{-pool-info } \text{Va } p)$   
 $(\text{buf } (mem\text{-pool-info } \text{Va } p) + n * (\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t)))$   
 $(lsizes \text{ Va } t ! \text{nat } (\text{from-l } \text{Va } t))$   
 $< n\text{-max } (mem\text{-pool-info } \text{Va } p) * 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t) \text{ and}$   
 $a10: \text{from-l } \text{Va } t < \text{alloc-l } \text{Va } t \text{ and}$   
 $a11: \text{blk } \text{Va } t = \text{buf } (mem\text{-pool-info } \text{Va } p) + n * (\text{max-sz } (mem\text{-pool-info } \text{Va } p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } \text{Va } t)) \text{ and}$   
 $a12: (x, \text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) \in \text{gvars-conf-stable} \text{ and}$   
 $a13: \forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1) \longrightarrow$   
 $\text{levels } (mem\text{-pool-info } x p) ! jj = \text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) p) ! jj \text{ and}$   
 $a14: \text{bits } (\text{levels } (mem\text{-pool-info } x p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1)))$   
 $(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t * 4)) \text{ } 3 \text{ FREE} \text{ and}$   
 $a15: \text{free-list } (\text{levels } (mem\text{-pool-info } x p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1) *$   
 $ii +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) t + 1))) \text{ and}$   
 $a16: \text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) \text{ and}$   
 $a17: \text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } \text{Va } t p) \text{ and}$   
 $a01: ii < \text{length } (\text{levels } (mem\text{-pool-info } x p)) \text{ and}$

```

a02:jj < length (bits (levels (mem-pool-info x p) ! ii)) and
a03:get-bit-s x p ii jj = NOEXIST and
a04:0 < ii
shows get-bit-s x p (ii - 1) (jj div 4) ≠ DIVIDED
proof-
  let ?mp = mem-pool-info x p
  have inv:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
    Va
  using a0 unfolding inv-def by auto
  have from-l-gt0:0 ≤ from-l Va t using a6 a5 by linarith
  have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
    Va p))
  using mp-alloc-stm4-lvl-len[OF a1 a12] by simp
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
  using mp-alloc-stm4-maxsz[OF a1 a12] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
  using mp-alloc-stm4-buf[OF a1 a12] by simp
  have from-l:from-l x = from-l Va
  using mp-alloc-stm4-froml[OF a17] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
  using from-l-gt0 by auto
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
    a12] by auto
  have lsizes-x:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a16
    by auto
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
    t))) and
    ?i2=(nat (from-l Va t + 1)) and
    ?j2=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
    t))*4) and
    ?i1'=(nat (from-l Va t)) - 1 and
    ?j1'=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
    t))) div 4 and
    ?i2'=(nat (from-l Va t)) + 2
  let ?j20'=?j2 * 4 and ?j21'=(?j2+1) * 4 and ?j22'=(?j2+2)*4 and
    ?j23'=(?j2+3)*4 and ?j24'=(?j2+4)*4
  let ?mp = mem-pool-info x p
  have inv-mempool-info-mp Va p
  using a1 mem-pools inv unfolding inv-mempool-info-def by auto
  note inv-mempool=this[simplified Let-def]
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
  using a10 a1 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
  by (metis i1-len a9 a11 a1 inv inv-mempool-info-def)

```

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have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
  using i1-len i2-len j1-len inv-mempool from-l-suc
  by simp
let ?bts = bits (levels ?mp ! ii)
let ?btsva = (bits (levels (mem-pool-info Va p) ! ii))
have a01':ii < length (levels (mem-pool-info Va p))
  using a01 len-levels by auto
then have inv-bitmap1:
   $\forall j < \text{length (bits (levels (mem-pool-info Va p) ! ii))}.$ 
   $(?btsva ! j = \text{FREE} \vee ?btsva ! j = \text{FREEING} \vee ?btsva ! j = \text{ALLOCATED}$ 
 $\vee ?btsva ! j = \text{ALLOCATING} \longrightarrow$ 
   $(ii > 0 \longrightarrow (\text{bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (j \text{ div } 4) = \text{DIVIDED})$ 
   $\wedge (ii < \text{length (levels (mem-pool-info Va p))} - 1 \longrightarrow \text{noexist-bits}$ 
 $(\text{mem-pool-info Va p}) (ii+1) (j*4))$ 
   $\wedge (?btsva ! j = \text{DIVIDED} \longrightarrow ii > 0 \longrightarrow (\text{bits (levels (mem-pool-info Va}$ 
 $p) ! (ii - 1))) ! (j \text{ div } 4) = \text{DIVIDED})$ 
   $\wedge (?btsva ! j = \text{NOEXIST} \longrightarrow ii < \text{length (levels (mem-pool-info Va p))}$ 
 $- 1$ 
   $\longrightarrow \text{noexist-bits (mem-pool-info Va p) (ii+1) (j*4)})$ 
   $\wedge (?btsva ! j = \text{NOEXIST} \wedge ii > 0 \longrightarrow (\text{bits (levels (mem-pool-info Va}$ 
 $p) ! (ii - 1))) ! (j \text{ div } 4) \neq \text{DIVIDED})$ 
  using inv mem-pools a1
  unfolding Let-def inv-bitmap-def
  by blast
have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
  using a7 a0 a11 unfolding inv-aux-vars-def invariant.inv-def
  by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))

then have alloc-predi1-j1:?i1 > 0  $\longrightarrow$  get-bit-s Va p (?i1 - 1) (?j1 div 4) =
DIVIDED
  using inv-bitmap1 i1-len j1-len inv a1 unfolding Let-def inv-bitmap-def by
blast
have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
  using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i1-len j1-len
  alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
  by (smt One-nat-def Suc-pred add.commute inv-mempool nat-add-left-cancel-less
plus-1-eq-Suc)
have nexisti3:?i2 < length (levels (mem-pool-info Va p)) - 1  $\longrightarrow$ 
  noexist-bits (mem-pool-info Va p) ?i2' ?j20'  $\wedge$ 
  noexist-bits (mem-pool-info Va p) ?i2' ?j21'  $\wedge$ 
  noexist-bits (mem-pool-info Va p) ?i2' ?j22'  $\wedge$ 
  noexist-bits (mem-pool-info Va p) ?i2' ?j23'
proof -
{ assume ?i2 < length (levels (mem-pool-info Va p)) - 1
  then have a00: $\forall j < \text{length (bits (levels (mem-pool-info Va p) ! ?i2))}.$ 
  get-bit-s Va p ?i2 j = NOEXIST  $\longrightarrow$  noexist-bits (mem-pool-info Va

```

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p) ?i2' (j * 4)
  using a1 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i2-len
  from-l-suc by auto
  then have noexist-bits (mem-pool-info Va p) ?i2' ?j20' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j21' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j22' ∧
    noexist-bits (mem-pool-info Va p) ?i2' ?j23'
  using j2-len nexisti2 Suc-lessD
  by (smt One-nat-def add.commute add-2-eq-Suc' add-Suc-right numeral-3-eq-3
plus-1-eq-Suc)
}
thus ?thesis by fastforce
qed
let ?bts = bits (levels ?mp ! ii) and ?fl = free-list (levels ?mp ! ii)
have a02':jj < length (bits (levels (mem-pool-info Va p) ! ii))
  using a02 a12 unfolding gvars-conf-def gvars-conf-stable-def
  by (simp add: mp-alloc-stm4-inv-bits-len)
have eq-len:length (bits (levels (mem-pool-info x p) ! ii)) =
  length (bits (levels (mem-pool-info Va p) ! ii))
  using mp-alloc-stm4-inv-bits-len a13 a14 length-list-update-n
  by metis
have inv-va:(?btsva ! jj = FREE ∨ ?btsva ! jj = FREEING ∨ ?btsva ! jj =
ALLOCATED ∨ ?btsva ! jj = ALLOCATING →
  (ii > 0 → (bits (levels (mem-pool-info Va p) ! (ii - 1))) ! (jj div
4) = DIVIDED)
  ∧ (ii < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
(mem-pool-info Va p) (ii+1) (jj*4))
  ∧ (?btsva ! jj = DIVIDED → ii > 0 → (bits (levels (mem-pool-info Va
p) ! (ii - 1))) ! (jj div 4) = DIVIDED)
  ∧ (?btsva ! jj = NOEXIST → ii < length (levels (mem-pool-info Va p))
- 1
  → noexist-bits (mem-pool-info Va p) (ii+1) (jj*4))
  ∧ (?btsva ! jj = NOEXIST ∧ ii > 0 → (bits (levels (mem-pool-info Va p)
! (ii - 1))) ! (jj div 4) ≠ DIVIDED)
  using inv-bitmap1 a02' by auto
{ assume a05:¬((ii=?i1 ∧ jj=?j1) ∨
  (ii=?i2 ∧ jj≥ ?j2 ∧ jj< ?j2+4) ∨
  (ii=?i2' ∧ jj≥ ?j20' ∧ jj< ?j24'))
  then have a050':¬(ii=?i1 ∧ jj=?j1) and
    a051': ¬(ii=?i2 ∧ jj≥ ?j2 ∧ jj< ?j2 + 4) and
    a052':¬(ii=?i2' ∧ jj≥ ?j20' ∧ jj< ?j24')
  by force+
  have eq-get-bit-i-j:get-bit-s x p ii jj = get-bit-s Va p ii jj
  using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
a17], THEN sym]] a14, of ii jj]
  using a050' a051' by auto
  have eq-get-bit-i1-j1:ii>0 → get-bit-s x p (ii-1) (jj div 4) = get-bit-s Va p
(ii-1) (jj div 4)

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proof-
{ assume a06:ii>0
  then have  $\neg((ii - 1) = ?i1 \wedge jj \text{ div } 4 = ?j1)$ 
    using a050' a051' from-l-suc by fastforce
  moreover have  $\forall j. j \geq ?j2 \wedge j \leq ?j2+3 \longrightarrow \neg((ii - 1) = ?i2 \wedge jj \text{ div } 4 =$ 
j)
    using a051' a052' from-l-gt0 by fastforce
  ultimately have  $\text{get-bit-s } x \ p \ (ii-1) \ (jj \text{ div } 4) = \text{get-bit-s } Va \ p \ (ii-1) \ (jj$ 
div 4)
    using same-bit-mp-alloc-x-va[OF a13[simplified a17[simplified mp-alloc-stm4-froml[OF
a17], THEN sym]] a14,
      of ii -1 jj div 4] by auto
  } thus ?thesis by auto qed
  then have ?thesis
    using a03 a04 eq-get-bit-i1-j1 eq-get-bit-i-j inv-va by auto
}
moreover {
  assume a06:(ii=?i1  $\wedge$  jj=?j1)
  then have  $\text{get-bit-s } x \ p \ ii \ jj = \text{DIVIDED}$ 
    using get-bit-x-l-b a13 a17 from-l from-l-gt0 i1-len j1-len by presburger
  then have ?thesis using a03 by auto
}
moreover {
  assume a06: (ii=?i2  $\wedge$  jj  $\geq$  ?j2  $\wedge$  jj < ?j2+4)
  then have a06':jj=?j2  $\vee$  jj=?j2+1  $\vee$  jj=?j2+2  $\vee$  jj=?j2+3 by auto
  { assume a08:jj=?j2
    then have  $\text{get-bit:}\text{get-bit-s } x \ p \ ii \ jj = \text{ALLOCATING}$ 
      using a02 a06 a14 eq-len get-bit-x-l1-b4 a04 i2-len from-l-gt0 i1-len j1-len

      by (metis mult.commute)
    then have ?thesis using a03 by auto
  }
  moreover {
    assume a07:jj $\neq$ ?j2
    have  $a07':jj \text{ div } 4 = ?j1$  using a06 a07 by auto
    have  $\text{get-bit-s } x \ p \ ii \ jj = \text{FREE}$ 
    using a06 a02 a14 a07 from-l mp-alloc-stm4-inv-bits-len a17 mp-alloc-stm4-pre-precond-f-bn
      by (auto simp add: mp-alloc-stm4-pre-precond-f-bn)
    then have ?thesis using a03 by auto
  }
  ultimately have ?thesis using a06 by fastforce
}
moreover {
  assume a06: (ii=?i2'  $\wedge$  jj  $\geq$  ?j20'  $\wedge$  jj < ?j24')
  then have a06':jj=?j20'  $\vee$  jj=?j20'+1  $\vee$  jj=?j20'+2  $\vee$  jj=?j20'+3  $\vee$ 
    jj=?j21'  $\vee$  jj=?j21'+1  $\vee$  jj=?j21'+2  $\vee$  jj=?j21'+3  $\vee$ 
    jj=?j22'  $\vee$  jj=?j22'+1  $\vee$  jj=?j22'+2  $\vee$  jj=?j22'+3  $\vee$ 
    jj=?j23'  $\vee$  jj=?j23'+1  $\vee$  jj=?j23'+2  $\vee$  jj=?j23'+3
    by presburger
}

```

```

have ij:(ii-1 = ?i2 ∧ (jj div 4) ≥ ?j2 ∧ (jj div 4) ≤ ?j2 + 3)
  using a04 a06 from-l-gt0 by auto
{ assume a08:(jj div 4) = ?j2
  then have get-bit:get-bit-s x p (ii-1) (jj div 4) = ALLOCATING
    using ij a02 a14 eq-len get-bit-x-l1-b4 a04 i2-len from-l-gt0 i1-len j1-len
    by (metis Suc-lessD j2-len mult.commute)
  then have ?thesis using a03 by auto
}
moreover {
  assume a07:(jj div 4) ≠ ?j2
  then have ii-1 = ?i2 ∧ (jj div 4 = Suc ?j2 ∨ jj div 4 = Suc (Suc ?j2) ∨
jj div 4 = Suc (Suc (Suc ?j2)))
    using ij by auto
  then have get-bit-s x p (ii-1) (jj div 4) = FREE
    using ij a01 a02 i2-len j2-len
    get-bit-x-l1-b41[OF - from-l-gt0[simplified from-l a17]
a13[simplified a17[THEN sym] from-l] a14, of ii-1 jj
div 4]
    by (metis Suc-lessD mult.commute)
  then have ?thesis using a03 by auto
}
ultimately have ?thesis using a06 by fastforce
}
ultimately show ?thesis by fastforce
qed

```

**lemma** *mp-alloc-stm4-inv-bitmap*:

```

assumes
a0:inv Va and
a1:freeing-node Va t = None and
a2:p ∈ mem-pools Va and
a3:ETIMEOUT ≤ timeout and
a4:timeout = ETIMEOUT ⟶ tmout Va t = ETIMEOUT and
a5:¬ rf Va t and
a6:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
a7:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
a8:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
a9:¬ free-l Va t < OK and
a10:NULL < buf (mem-pool-info Va p) ∨ NULL < n ∧ NULL < max-sz (mem-pool-info
Va p) div 4 ^ nat (from-l Va t) and
a11:free-l Va t ≤ from-l Va t and
a12:allocating-node Va t =
Some (pool = p, level = nat (from-l Va t),
block = block-num (mem-pool-info Va p)
(buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
(lsizes Va t ! nat (from-l Va t)),
data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div

```

$4 \wedge \text{nat} (\text{from-l } Va \ t)) \parallel \text{and}$   
 $a13:n = \text{block-num} (\text{mem-pool-info } Va \ p)$   
 $(\text{buf} (\text{mem-pool-info } Va \ p) + n * (\text{max-sz} (\text{mem-pool-info } Va \ p) \text{div } 4 \wedge \text{nat}$   
 $(\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat} (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz} (\text{mem-pool-info } Va \ p) \text{div } 4 \wedge \text{nat} (\text{from-l } Va \ t) = \text{NULL} \text{ and}$   
 $a14:\text{block-num} (\text{mem-pool-info } Va \ p)$   
 $(\text{buf} (\text{mem-pool-info } Va \ p) + n * (\text{max-sz} (\text{mem-pool-info } Va \ p) \text{div } 4 \wedge \text{nat}$   
 $(\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat} (\text{from-l } Va \ t))$   
 $< n\text{-max} (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat} (\text{from-l } Va \ t) \text{ and}$   
 $a15:\text{from-l } Va \ t < \text{alloc-l } Va \ t \text{ and}$   
 $a16:\text{cur } Va = \text{Some } t \text{ and}$   
 $a17:n < n\text{-max} (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat} (\text{from-l } Va \ t) \text{ and}$   
 $a18:\text{blk } Va \ t = \text{buf} (\text{mem-pool-info } Va \ p) + n * (\text{max-sz} (\text{mem-pool-info } Va \ p)$   
 $\text{div } 4 \wedge \text{nat} (\text{from-l } Va \ t)) \text{ and}$   
 $a19:\text{mempoolalloc-ret } Va \ t = \text{None} \text{ and}$   
 $a20:\forall ii \leq \text{nat} (\text{alloc-l } Va \ t). \text{sz} \leq \text{lsizes } Va \ t ! ii \text{ and}$   
 $a21:\text{alloc-l } Va \ t = \text{int} (\text{length} (\text{lsizes } Va \ t)) - 1 \wedge \text{length} (\text{lsizes } Va \ t) = n\text{-levels}$   
 $(\text{mem-pool-info } Va \ p) \vee$   
 $\text{alloc-l } Va \ t = \text{int} (\text{length} (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t ! \text{nat} (\text{alloc-l } Va \ t + 1)$   
 $< \text{sz} \text{ and}$   
 $a22:i \ x \ t = 4 \text{ and}$   
 $a23:\text{cur } x = \text{cur} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a24:\text{tick } x = \text{tick} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a25:\text{thd-state } x = \text{thd-state} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a26:(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \text{ and}$   
 $a27:\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ pa \text{ and}$   
 $a28:\text{wait-q} (\text{mem-pool-info } x \ p) = \text{wait-q} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ p) \text{ and}$   
 $a29:\forall t'. t' \neq t \longrightarrow \text{lvars-nochange } t' \ x (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a30:\forall jj. jj \neq \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels} (\text{mem-pool-info } x \ p) ! jj = \text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ p) ! jj \text{ and}$   
 $a31:\text{bits} (\text{levels} (\text{mem-pool-info } x \ p) ! \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va$   
 $t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits} (\text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))))$   
 $(\text{Suc} (\text{bn} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE} \text{ and}$   
 $a32:\text{free-list} (\text{levels} (\text{mem-pool-info } x \ p) ! \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map} (\lambda ii. \text{lsizes} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $ii +$   
 $\text{blk} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$



```

(free-list
  (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
    nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
a33:j x = j (mp-alloc-stm4-pre-precond-f Va t p) and
a34:ret x = ret (mp-alloc-stm4-pre-precond-f Va t p) and
a35:endt x = endt (mp-alloc-stm4-pre-precond-f Va t p) and
a36:rf x = rf (mp-alloc-stm4-pre-precond-f Va t p) and
a37:tmout x = tmout (mp-alloc-stm4-pre-precond-f Va t p) and
a38:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p) and
a39:alloc-l x = alloc-l (mp-alloc-stm4-pre-precond-f Va t p) and
a40:free-l x = free-l (mp-alloc-stm4-pre-precond-f Va t p) and
a41:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) and
a42:blk x = blk (mp-alloc-stm4-pre-precond-f Va t p) and
a43:nODEV x = nodev (mp-alloc-stm4-pre-precond-f Va t p) and
a44:bn x = bn (mp-alloc-stm4-pre-precond-f Va t p) and
a45:alloc-lsize-r x = alloc-lsize-r (mp-alloc-stm4-pre-precond-f Va t p) and
a46:lvl x = lvl (mp-alloc-stm4-pre-precond-f Va t p) and
a47:bb x = bb (mp-alloc-stm4-pre-precond-f Va t p) and
a48:block-pt x = block-pt (mp-alloc-stm4-pre-precond-f Va t p) and
a49:th x = th (mp-alloc-stm4-pre-precond-f Va t p) and
a50:need-resched x = need-resched (mp-alloc-stm4-pre-precond-f Va t p) and
a51:mempoolalloc-ret x = mempoolalloc-ret (mp-alloc-stm4-pre-precond-f Va t p)
and
a52:freeing-node x = freeing-node (mp-alloc-stm4-pre-precond-f Va t p) and
a53:allocating-node x = allocating-node (mp-alloc-stm4-pre-precond-f Va t p)
shows inv-bitmap x
proof-
  let ?mp = mem-pool-info x p
  have inv:inv-aux-vars Va  $\wedge$  inv-bitmap Va  $\wedge$  inv-mempool-info Va  $\wedge$  inv-bitmap-freelist
  Va
    using a0 unfolding inv-def by auto
  have from-l-gt0:0  $\leq$  from-l Va t using a11 a9 by linarith
  have len-levels:length (levels (mem-pool-info x p)) = length (levels (mem-pool-info
  Va p))
    using mp-alloc-stm4-lvl-len[OF a2 a26] by simp
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
    using mp-alloc-stm4-maxsz[OF a2 a26] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
    using mp-alloc-stm4-buf[OF a2 a26] by simp
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a41] by auto
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
  a26] by auto
  have lsizes-x-va:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a38
  by auto
  have from-l-gt0:OK  $\leq$  from-l Va t using a11 a9 by linarith
  { fix p'
    assume a00:p'  $\in$  mem-pools x
    let ?i1=(nat (from-l Va t)) and

```

```

      ?j1 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
      ?i2 = (nat (from-l Va t + 1)) and
      ?j2 = (4*block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
      ?i1' = (nat (from-l Va t)) - 1 and
      ?j1' = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) div 4 and
      ?i2' = (nat (from-l Va t)) + 2 and
      ?j2' = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) * 16

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```

have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
using a12 a0 a18 unfolding inv-aux-vars-def invariant.inv-def
by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2)
Mem-block.select-convs(3))
{assume p≠p'
moreover have mem-pool-info x p' = mem-pool-info Va p'
using mp-alloc-stm4-pres-mpinfo
by (metis a27 calculation)
ultimately have inv-bitmap-mp x p'
using a00 inv len-levels maxsz buf from-l mem-pools
by (simp add: inv-bitmap-def Let-def)
}
moreover { assume eq-p:p=p'
let ?mp = mem-pool-info x p
have inv-mempool-info-mp Va p
using eq-p mem-pools a00 inv unfolding inv-mempool-info-def by auto
note inv-mempool=this[simplified Let-def]
{fix i
assume a01:i<length (levels ?mp)
let ?bts = bits (levels ?mp ! i)
let ?btsva = (bits (levels (mem-pool-info Va p) ! i))
have a01':i < length (levels (mem-pool-info Va p))
using a01 len-levels by auto
then have inv-bitmap1:
  ∀ j < length (bits (levels (mem-pool-info Va p) ! i)).
    (?btsva ! j = FREE ∨ ?btsva ! j = FREEING ∨ ?btsva ! j =
ALLOCATED ∨ ?btsva ! j = ALLOCATING →
      (i > 0 → (bits (levels (mem-pool-info Va p) ! (i - 1))) ! (j
div 4) = DIVIDED)
      ∧ (i < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
(mem-pool-info Va p) (i+1) (j*4))
      ∧ (?btsva ! j = DIVIDED → i > 0 → (bits (levels (mem-pool-info
Va p) ! (i - 1))) ! (j div 4) = DIVIDED)
      ∧ (?btsva ! j = NOEXIST → i < length (levels (mem-pool-info Va
p)) - 1
      → noexist-bits (mem-pool-info Va p) (i+1) (j*4))
      ∧ (?btsva ! j = NOEXIST ∧ i > 0 → (bits (levels (mem-pool-info

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Va p) ! (i - 1))) ! (j div 4) ≠ DIVIDED)
  using inv eq-p mem-pools a00
  unfolding Let-def inv-bitmap-def
  by blast
  let ?bts = bits (levels ?mp ! i) and ?fl = free-list (levels ?mp ! i)
  have f1:∀j < length ?bts.
    (?bts ! j = FREE ∨ ?bts ! j = FREEING ∨ ?bts ! j = ALLOCATED
  ∨ ?bts ! j = ALLOCATING →
    (i > 0 → (bits (levels (mem-pool-info x p) ! (i - 1))) ! (j div
4) = DIVIDED)
    ∧ (i < length (levels (mem-pool-info x p)) - 1 → noexist-bits
(mem-pool-info x p) (i+1) (j*4) ))
    ∧ (?bts ! j = DIVIDED → i > 0 → (bits (levels (mem-pool-info x
p) ! (i - 1))) ! (j div 4) = DIVIDED)
    ∧ (?bts ! j = NOEXIST → i < length (levels (mem-pool-info x p))
- 1
    → noexist-bits (mem-pool-info x p) (i+1) (j*4))
    ∧ (?bts ! j = NOEXIST ∧ i > 0 → (bits (levels (mem-pool-info x
p) ! (i - 1))) ! (j div 4) ≠ DIVIDED)
  proof-
  { fix j
    assume a02:j < length ?bts
    then have a02':j < length (bits (levels (mem-pool-info Va p) ! i))
      using a26 unfolding gvars-conf-def gvars-conf-stable-def
      by (simp add: mp-alloc-stm4-inv-bits-len)
    have eq-len:length (bits (levels (mem-pool-info x p) ! i)) =
      length (bits (levels (mem-pool-info Va p) ! i))
      using mp-alloc-stm4-inv-bits-len a30 a31 length-list-update-n
      by metis
    have inv-va:(?btsva ! j = FREE ∨ ?btsva ! j = FREEING ∨ ?btsva ! j
= ALLOCATED ∨ ?btsva ! j = ALLOCATING →
      (i > 0 → (bits (levels (mem-pool-info Va p) ! (i - 1))) ! (j
div 4) = DIVIDED)
      ∧ (i < length (levels (mem-pool-info Va p)) - 1 → noexist-bits
(mem-pool-info Va p) (i+1) (j*4) ))
      ∧ (?btsva ! j = DIVIDED → i > 0 → (bits (levels (mem-pool-info
Va p) ! (i - 1))) ! (j div 4) = DIVIDED)
      ∧ (?btsva ! j = NOEXIST → i < length (levels (mem-pool-info Va
p)) - 1
      → noexist-bits (mem-pool-info Va p) (i+1) (j*4))
      ∧ (?btsva ! j = NOEXIST ∧ i > 0 → (bits (levels (mem-pool-info
Va p) ! (i - 1))) ! (j div 4) ≠ DIVIDED)
    using inv-bitmap1 a02' eq-p by auto
    let ?goal1 = (?bts ! j = FREE ∨ ?bts ! j = FREEING ∨ ?bts ! j =
ALLOCATED ∨ ?bts ! j = ALLOCATING →
      (i > 0 → (bits (levels (mem-pool-info x p) ! (i - 1))) ! (j div
4) = DIVIDED)
      ∧ (i < length (levels (mem-pool-info x p)) - 1 → noexist-bits
(mem-pool-info x p) (i+1) (j*4) ))

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    let ?goal2 = (?bts ! j = DIVIDED → i > 0 → (bits (levels (mem-pool-info
x p) ! (i - 1))) ! (j div 4) = DIVIDED)
    let ?goal3 = (?bts ! j = NOEXIST → i < length (levels (mem-pool-info
x p)) - 1
    → noexist-bits (mem-pool-info x p) (i+1) (j*4))
    let ?goal4 = (?bts ! j = NOEXIST ∧ i > 0 → (bits (levels (mem-pool-info
x p) ! (i - 1))) ! (j div 4) ≠ DIVIDED)
    have ?goal1 using eq-p
    mp-alloc-stm4-inv-bitmap1[OF a0 a2 a6 a8 a9 a11 a12 a13 a14 a15 a17
a18 a26 a30 a31 a32 a38 a41 a01 a02]
    by auto
    moreover have ?goal2
    using mp-alloc-stm4-inv-bitmap2[OF a0 a2 a6 a7 a8 a9 a11 a12 a13 a14
a15 a18 a26 a30 a31 a32 a38 a41 a01 a02]
    by auto
    moreover have ?goal3
    using mp-alloc-stm4-inv-bitmap3[OF a0 a2 a6 a8 a9 a11 a12 a13 a14
a15 a17 a18 a26 a30 a31 a32 a38 a41 a01 a02]
    by auto
    moreover have ?goal4 using mp-alloc-stm4-inv-bitmap4[OF a0 a2 a6 a7
a8 a9 a11 a12 a13 a14 a15 a18 a26 a30 a31 a32 a38 a41 a01 a02]
    by auto
    ultimately have ?goal1 ∧ ?goal2 ∧ ?goal3 ∧ ?goal4
    by blast
  } thus ?thesis by auto
qed
} then have inv-bitmap-mp x p' using eq-p by auto
} ultimately have inv-bitmap-mp x p' by fastforce
} then show ?thesis unfolding inv-bitmap-def by auto
qed

```

**lemma** *mp-alloc-stm4-inv-aux-vars1*:

```

  assumes
    a0:inv Va and
    a1:freeing-node Va t = None and
    a2:p ∈ mem-pools Va and
    a3:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
    a4:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
    a5:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a6:¬ free-l Va t < OK and
    a7:free-l Va t ≤ from-l Va t and
    a8:allocating-node Va t =
      Some (pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t))),

```

```

      data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)) and
a9:block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
      (lsizes Va t ! nat (from-l Va t))
      < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
a10:from-l Va t < alloc-l Va t and
a11:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p)
div 4 ^ nat (from-l Va t)) and
a12:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and
a13:∀ pa. pa ≠ p → mem-pool-info x pa = mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) pa and
a14:∀ jj. jj ≠ nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) →
      levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p) ! jj and
a15:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
t p) t + 1)) =
list-updates-n
      (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))
      (Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4)) 3 FREE and
a16:free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f
Va t p) t + 1)) =
inserts
      (map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
ii +
blk (mp-alloc-stm4-pre-precond-f Va t p) t)
[Suc NULL..<4])
      (free-list
      (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
a17:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p) and
a18:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) and
a19:freeing-node x = freeing-node (mp-alloc-stm4-pre-precond-f Va t p) and
a20:allocating-node x = allocating-node (mp-alloc-stm4-pre-precond-f Va t p) and
a21:freeing-node x t' = Some m
shows get-bit-s x (pool m) (level m) (block m) = FREEING
proof-
have inv:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
Va
using a0 unfolding inv-def by auto
have from-l-gt0:0 ≤ from-l Va t using a7 a6 by linarith
have inv-aux-v: (∀ t n. freeing-node Va t = Some n →
get-bit (mem-pool-info Va) (pool n) (level n) (block n) = FREEING)
using a0 unfolding inv-def inv-aux-vars-def
by blast
let ?i1=(nat (from-l Va t)) and

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```

    ?j1 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))*4)
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a12] by auto
  have inv-mempool-info-mp Va p
    using a2 mem-pools inv unfolding inv-mempool-info-def Let-def by auto
  note inv-mempool=this[simplified Let-def]
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a18] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a10 a2 a5 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a10 a2 a5 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis a0 a2 a9 a11 i1-len inv-mempool-info-def invariant.inv-def)
  have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp
  have lsizes-x-va:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a17
    by auto
  { assume t' = t
    then have ?thesis
      using a1 a19 a21
      by (metis mp-alloc-stm4-pre-precond-f-def-frnode option.distinct(1))
    }
  moreover {assume a00:t'≠t
    then have freeing-node (mp-alloc-stm4-pre-precond-f Va t p) t' = freeing-node
Va t'
      unfolding mp-alloc-stm4-pre-precond-f-def by auto
      then have eq-alloc:freeing-node Va t' = freeing-node x t'
        using a19 by auto
      then have t2-same-allocating-node-Va:freeing-node Va t' = Some m
        using a0 a21 a19
      unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def inv-aux-vars-def

      by auto
      then have diff-t:¬(pool m = p ∧ level m = ?i1 ∧ block m = ?j1)
        using a00 a21 a8 inv unfolding inv-aux-vars-def
      by (metis Mem-block.simps(1) Mem-block.simps(2) Mem-block.simps(3) a11)
      {
        assume pool m ≠ p

```

```

    then have ?thesis using a0 a13 a21 eq-alloc mp-alloc-stm4-pres-mpinfo
      unfolding inv-aux-vars-def invariant.inv-def
      by metis
  } note not-pool-p-allocating = this
  moreover {
    assume a01:pool m = p
    have bit-m-va-alloc:get-bit (mem-pool-info Va) (pool m) (level m) (block m)
= FREEING
      using a21 eq-alloc inv-aux-va by presburger
    have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
      using mp-alloc-stm4-maxsz[OF a2 a12] by simp
    have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
      using mp-alloc-stm4-buf[OF a2 a12] by simp
    have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
      using a8 a0 a11 unfolding inv-aux-vars-def invariant.inv-def
      by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2)
Mem-block.select-convs(3))
    have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
      using a2 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i1-len j1-len
      alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
    by (smt One-nat-def Suc-pred add.commute inv-mempool nat-add-left-cancel-less
plus-1-eq-Suc)
    { assume a02:(level m = ?i1 ∧ block m = ?j1)
      then have ?thesis using diff-t a01 by auto
    }
  }
  moreover {
    assume a02:¬(level m = ?i1 ∧ block m = ?j1)
    { assume a03:¬(level m = ?i2 ∧ (block m) ≥ ?j2 ∧ (block m) < ?j2 + 4)
      then have eq-get-bit-i-j:get-bit-s x p (level m) (block m) =
        get-bit-s Va p (level m) (block m)
      using same-bit-mp-alloc-x-va[OF a14[simplified
a18[simplified mp-alloc-stm4-froml[OF a18], THEN sym]] a15, of level
m block m]
      a01 a02 by auto
      then have ?thesis using a01 a20 inv-aux-va not-pool-p-allocating
a21 eq-alloc inv-aux-va by force
    }
  }
  moreover {
    assume a03:(level m = ?i2 ∧ (block m) ≥ ?j2 ∧ (block m) < ?j2 + 4)
    then have block m = ?j2 ∨ block m = ?j2 + 1 ∨ block m = ?j2 + 2 ∨
block m = ?j2 + 3
      by auto
    then have ?thesis using bit-m-va-alloc nexisti2 a01 a03 by auto
  } ultimately have ?thesis by fastforce
} ultimately have ?thesis by fastforce
} ultimately show ?thesis by auto
qed

```

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lemma mp-alloc-stm4-inv-aux-vars2:
  assumes
    a0:inv Va and
    a1:freeing-node Va t = None and
    a2:p ∈ mem-pools Va and
    a3:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a4: $\neg$  free-l Va t < OK and
    a5:free-l Va t ≤ from-l Va t and
    a6:block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
        (from-l Va t)))
      (lsizes Va t ! nat (from-l Va t))
       $< n\text{-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t)}$  and
    a7:from-l Va t < alloc-l Va t and
    a8:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
       $4 ^ nat (from-l Va t))$  and
    a9:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and
    a10: $\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \text{ } pa = \text{mem-pool-info (mp-alloc-stm4-pre-precond-f
      Va t p) pa and
    a11: $\forall jj. jj \neq nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) \longrightarrow$ 
      levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
      Va t p) p) ! jj and
    a12:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
      t p) t + 1)) =
      list-updates-n
      (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
        nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))
      (Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4))) 3 FREE and
    a13:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) and
    a14:freeing-node x = freeing-node (mp-alloc-stm4-pre-precond-f Va t p) and
    a54:get-bit-s x (pool m) (level m) (block m) = FREEING ∧ mem-block-addr-valid
      x m
  shows  $(\exists t. \text{freeing-node } x \text{ } t = \text{Some } m)$ 
proof–
  have inv:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist
    Va
    using a0 unfolding inv-def by auto
    have from-l-gt0:0 ≤ from-l Va t using a5 a4 by linarith
    have block-valid-va:mem-block-addr-valid Va m
      using a2 a9 a54 mp-alloc-stm4-buf mp-alloc-stm4-maxsz
      unfolding mem-block-addr-valid-def by auto
    have inv-aux-va:( $\forall n. \text{get-bit (mem-pool-info Va) (pool n) (level n) (block n) =$ 
      FREEING ∧ mem-block-addr-valid Va n
       $\longrightarrow (\exists t. \text{freeing-node Va } t = \text{Some } n))$ 
      using a0 unfolding inv-def inv-aux-vars-def
      by blast
    {assume (pool m) ≠ p$ 
```



```

    then have get-bit-s Va (pool m) (level m) (block m) = get-bit-s x (pool m)
(level m) (block m)
    using a10
    by (metis mp-alloc-stm4-pres-mpinfo)
    then have ?thesis using a54 inv-aux-va block-valid-va a14 mp-alloc-stm4-pre-precond-f-def-frnode
    by metis
  }
  moreover{
    assume a01:pool m = p
    let ?i1=(nat (from-l Va t)) and
    ?j1= (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))*4)
    have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a9] by auto
    have inv-mempool-info-mp Va p
    using a2 mem-pools inv unfolding inv-mempool-info-def Let-def by auto
    note inv-mempool=this[simplified Let-def]
    have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a13] by auto
    have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
    have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a7 a2 a3 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
    have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a7 a2 a3 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
    have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis a0 a2 a6 a8 i1-len inv-mempool-info-def invariant.inv-def)
    have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp
    { assume a02:¬(((level m)=?i1 ∧ (block m)=?j1) ∨
((level m)=?i2 ∧ (block m)≥ ?j2 ∧ (block m)< ?j2+4))
    then have a020':¬((level m)=?i1 ∧ (block m)=?j1) and
    a021': ¬((level m)=?i2 ∧ (block m)≥ ?j2 ∧ (block m)< ?j2 + 4)
    by force+
    then have eq-get-bit-i-j:get-bit-s x p (level m) (block m) = get-bit-s Va p
(level m) (block m)
    using same-bit-mp-alloc-x-va[OF a11[simplified
a13[simplified mp-alloc-stm4-froml[OF a13], THEN sym]] a12, of level
m block m]
    using a020' a021' by auto
    then have ?thesis using a01 a54 inv-aux-va
    block-valid-va a14 mp-alloc-stm4-pre-precond-f-def-frnode

```

```

    by metis
  }
  moreover{
    assume a02:((level m)=?i1 ∧ (block m)=?j1)
    then have get-bit-s x p ?i1 ?j1 = DIVIDED
    by (simp add: a11 from-l-gt0 from-l-suc i1-len j1-len mp-alloc-stm4-pre-precond-f-froml

          same-bit-mp-alloc-stm4-pre-precond-divided)
    then have ?thesis using a54 a02 a01 by auto
  }
  moreover{
    assume a02:(level m)=?i2 ∧ (block m)≥ ?j2 ∧ (block m)< ?j2+4
    then have get-bit-s x p ?i2 ?j2 = ALLOCATING
      using i2-len j2-len a12 get-bit-x-l1-b4[OF - from-l-gt0 a12, of ?i2 ?j2]
    by (metis (no-types, lifting) add-2-eq-Suc' add-Suc-right add-lessD1 mult.commute)
    moreover {
      assume a07:(block m)≠?j2
      then have (level m) = ?i2 ∧ ((block m) = Suc ?j2 ∨
        (block m) = Suc (Suc ?j2) ∨ (block m) = Suc (Suc (Suc ?j2)))
        using a02 by auto
      then have get-bit-s x p (level m) (block m) = FREE
        using a02 i2-len j2-len
        get-bit-x-l1-b41[OF - from-l-gt0[simplified from-l a13]
          a11[simplified a13[THEN sym] from-l] a12, of level m
            block m]
        by (metis Suc-lessD mult.commute)
    }
    ultimately have ?thesis using a54 a02 a01 by fastforce
  } ultimately have ?thesis by auto
} ultimately show ?thesis by auto
qed

```

**lemma** *mp-alloc-stm4-inv-aux-vars3*:

```

  assumes
    a0:inv Va and
    a1:freeing-node Va t = None and
    a2:p ∈ mem-pools Va and
    a3:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a4:¬ free-l Va t < OK and
    a5:free-l Va t ≤ from-l Va t and
    a6:allocating-node Va t =
      Some (|pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
            4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t)),
        data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
            4 ^ nat (from-l Va t)))|) and
    a7:block-num (mem-pool-info Va p)

```

(buf (mem-pool-info Va p) + n \* (max-sz (mem-pool-info Va p) div 4 ^ nat  
 (from-l Va t)))  
 (lsizes Va t ! nat (from-l Va t))  
 < n-max (mem-pool-info Va p) \* 4 ^ nat (from-l Va t) **and**  
 a8:from-l Va t < alloc-l Va t **and**  
 a9:blk Va t = buf (mem-pool-info Va p) + n \* (max-sz (mem-pool-info Va p) div  
 4 ^ nat (from-l Va t)) **and**  
 a10:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable **and**  
 a11:∀ pa. pa ≠ p → mem-pool-info x pa = mem-pool-info (mp-alloc-stm4-pre-precond-f  
 Va t p) pa **and**  
 a12:∀ jj. jj ≠ nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) →  
 levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f  
 Va t p) p) ! jj **and**  
 a13:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va  
 t p) t + 1)) =  
 list-updates-n  
 (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !  
 nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))  
 (Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t \* 4)) 3 FREE **and**  
 a14:free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f  
 Va t p) t + 1)) =  
 inserts  
 (map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !  
 nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) \*  
 ii +  
 blk (mp-alloc-stm4-pre-precond-f Va t p) t)  
 [Suc NULL..<4])  
 (free-list  
 (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !  
 nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) **and**  
 a15:lsizes x = lsizes (mp-alloc-stm4-pre-precond-f Va t p) **and**  
 a16:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p) **and**  
 a17:allocating-node x = allocating-node (mp-alloc-stm4-pre-precond-f Va t p) **and**  
 a18:allocating-node x t' = Some m  
**shows** get-bit-s x (pool m) (level m) (block m) = ALLOCATING  
**proof**–  
**have** inv:inv-aux-vars Va ∧ inv-bitmap Va ∧ inv-mempool-info Va ∧ inv-bitmap-freelist  
 Va  
**using** a0 **unfolding** inv-def **by** auto  
**have** from-l-gt0:0 ≤ from-l Va t **using** a5 a4 **by** linarith  
**have** inv-aux-va:(∀ t n. allocating-node Va t = Some n →  
 get-bit (mem-pool-info Va) (pool n) (level n) (block n) = ALLOCATING)  
**using** a0 **unfolding** inv-def inv-aux-vars-def  
**by** blast  
**let** ?i1=(nat (from-l Va t)) **and**  
 ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va  
 t))) **and**  
 ?i2 = (nat (from-l Va t + 1)) **and**  
 ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va

```

t))*4)
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a10] by auto
  have inv-mempool-info-mp Va p
    using a2 mem-pools inv unfolding inv-mempool-info-def Let-def by auto
  note inv-mempool=this[simplified Let-def]
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a16] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a8 a2 a3 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a8 a2 a3 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis a0 a2 a7 a9 i1-len inv-mempool-info-def invariant.inv-def)
  have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp
  have lsizes-x-va:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a15
    by auto
  {assume a00:t' ≠ t
  then have allocating-node (mp-alloc-stm4-pre-precond-f Va t p) t' = allocating-node
Va t'
    unfolding mp-alloc-stm4-pre-precond-f-def by auto
  then have eq-alloc:allocating-node Va t' = allocating-node x t'
    using a17 by auto
  then have diff-t:¬(pool m = p ∧ level m = ?i1 ∧ block m = ?j1)
    using a00 a18 a6 inv unfolding inv-aux-vars-def
    by (metis Mem-block.simps(1) Mem-block.simps(2) Mem-block.simps(3) a9)
  {
    assume pool m ≠ p
    then have ?thesis
      by (metis a11 a18 eq-alloc inv-aux-va mp-alloc-stm4-pres-mpinfo)
  } note not-pool-p-allocating = this
  moreover {
    assume a01:pool m = p
    have bit-m-va-alloc:get-bit (mem-pool-info Va) (pool m) (level m) (block m)
= ALLOCATING
      using a18 eq-alloc inv-aux-va by presburger
    have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
      using mp-alloc-stm4-maxsz[OF a2 a10] by simp
    have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
      using mp-alloc-stm4-buf[OF a2 a10] by simp
    have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
      using a6 a0 a9 unfolding inv-aux-vars-def invariant.inv-def

```

```

    by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2)
Mem-block.select-convs(3))
    have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
    using a2 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
i1-len j1-len
    alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
    by (smt One-nat-def Suc-pred add commute inv-mempool nat-add-left-cancel-less
plus-1-eq-Suc)
    { assume a02:(level m = ?i1 ∧ block m = ?j1)
      then have ?thesis using diff-t a01 by auto
    }
    moreover {
      assume a02:¬(level m = ?i1 ∧ block m = ?j1)
      { assume a03:¬(level m = ?i2 ∧ (block m) ≥ ?j2 ∧ (block m) < ?j2 + 4)
        then have eq-get-bit-i-j:get-bit-s x p (level m) (block m) =
          get-bit-s Va p (level m) (block m)
          using same-bit-mp-alloc-x-va[OF a12[simplified
a16[simplified mp-alloc-stm4-froml[OF a16], THEN sym]] a13, of level
m block m]
          a01 a02 by auto
        then have ?thesis using a01 a17 inv-aux-va not-pool-p-allocating
a18 eq-alloc inv-aux-va by force
      }
      moreover {
        assume a03:(level m = ?i2 ∧ (block m) ≥ ?j2 ∧ (block m) < ?j2 + 4)
        then have block m = ?j2 ∨ block m = ?j2 + 1 ∨ block m = ?j2 + 2 ∨
block m = ?j2 + 3
          by auto
        then have ?thesis using bit-m-va-alloc nexisti2 a01 a03 by auto
      }
      ultimately have ?thesis by fastforce
    } ultimately have ?thesis by fastforce
  } ultimately have ?thesis by auto
}
moreover {
  assume t'=t
  then have (pool m) = p ∧ (level m) = ?i2 ∧ (block m) = ?j2
  by (metis Mem-block.simps(1) Mem-block.simps(2) Mem-block.simps(3) a17
a18
    mp-alloc-stm4-pre-precond-f-allocating mult.commute option.sel)
  then have ?thesis using get-bit-x-l1-b4[OF - from-l-gt0 a13 i2-len ] j2-len
  by (metis Suc-lessD mult.commute)
}
ultimately show ?thesis by auto
qed

lemma mp-alloc-stm4-inv-aux-vars4:
  assumes
    a0:inv Va and

```

$a1$ :freeing-node  $Va\ t = None$  **and**  
 $a2$ : $p \in mem\text{-}pools\ Va$  **and**  
 $a3$ : $\forall ii < length\ (lsizes\ Va\ t). lsizes\ Va\ t\ !\ ii = ALIGN4\ (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p))\ div\ 4\ \wedge\ ii$  **and**  
 $a4$ :alloc-l  $Va\ t < int\ (n\text{-}levels\ (mem\text{-}pool\text{-}info\ Va\ p))$  **and**  
 $a5$ : $\neg free\text{-}l\ Va\ t < OK$  **and**  
 $a6$ :free-l  $Va\ t \leq from\text{-}l\ Va\ t$  **and**  
 $a7$ :allocating-node  $Va\ t =$   
 $Some\ (\downarrow pool = p,\ level = nat\ (from\text{-}l\ Va\ t),$   
 $\quad block = block\text{-}num\ (mem\text{-}pool\text{-}info\ Va\ p)$   
 $\quad (buf\ (mem\text{-}pool\text{-}info\ Va\ p) + n * (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p)\ div\ 4\ \wedge\ nat\ (from\text{-}l\ Va\ t)))$   
 $\quad (lsizes\ Va\ t\ !\ nat\ (from\text{-}l\ Va\ t)),$   
 $\quad data = buf\ (mem\text{-}pool\text{-}info\ Va\ p) + n * (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p)\ div\ 4\ \wedge\ nat\ (from\text{-}l\ Va\ t)))$  **and**  
 $a8$ :block-num  $(mem\text{-}pool\text{-}info\ Va\ p)$   
 $(buf\ (mem\text{-}pool\text{-}info\ Va\ p) + n * (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p)\ div\ 4\ \wedge\ nat\ (from\text{-}l\ Va\ t)))$   
 $(lsizes\ Va\ t\ !\ nat\ (from\text{-}l\ Va\ t))$   
 $< n\text{-}max\ (mem\text{-}pool\text{-}info\ Va\ p) * 4\ \wedge\ nat\ (from\text{-}l\ Va\ t)$  **and**  
 $a9$ :from-l  $Va\ t < alloc\text{-}l\ Va\ t$  **and**  
 $a10$ :blk  $Va\ t = buf\ (mem\text{-}pool\text{-}info\ Va\ p) + n * (max\text{-}sz\ (mem\text{-}pool\text{-}info\ Va\ p)\ div\ 4\ \wedge\ nat\ (from\text{-}l\ Va\ t))$  **and**  
 $a11$ :alloc-l  $Va\ t = int\ (length\ (lsizes\ Va\ t)) - 1 \wedge length\ (lsizes\ Va\ t) = n\text{-}levels\ (mem\text{-}pool\text{-}info\ Va\ p) \vee$   
 $alloc\text{-}l\ Va\ t = int\ (length\ (lsizes\ Va\ t)) - 2 \wedge lsizes\ Va\ t\ !\ nat\ (alloc\text{-}l\ Va\ t + 1)$   
 $< sz$  **and**  
 $a12$ :( $x,\ mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p$ )  $\in gvars\text{-}conf\text{-}stable$  **and**  
 $a13$ : $\forall pa. pa \neq p \longrightarrow mem\text{-}pool\text{-}info\ x\ pa = mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ pa$  **and**  
 $a14$ : $\forall jj. jj \neq nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1) \longrightarrow$   
 $levels\ (mem\text{-}pool\text{-}info\ x\ p)\ !\ jj = levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)\ !\ jj$  **and**  
 $a15$ :bits  $(levels\ (mem\text{-}pool\text{-}info\ x\ p)\ !\ nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)) =$   
 $list\text{-}updates\text{-}n$   
 $(bits\ (levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)\ !$   
 $\quad nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)))$   
 $(Suc\ (bn\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t * 4))\ 3\ FREE$  **and**  
 $a16$ :free-list  $(levels\ (mem\text{-}pool\text{-}info\ x\ p)\ !\ nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)) =$   
 $inserts$   
 $(map\ (\lambda ii. lsizes\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t\ !$   
 $\quad nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1) *$   
 $\quad ii +$   
 $\quad blk\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t)$   
 $[Suc\ NULL..<4])$   
 $(free\text{-}list$   
 $\quad (levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)\ !$

$\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$  and  
 $a17:\text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  and  
 $a18:\text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  and  
 $a19:\text{blk } x = \text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  and  
 $a20:\text{allocating-node } x = \text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  and  
 $a21:\text{get-bit-s } x \ (\text{pool } m) \ (\text{level } m) \ (\text{block } m) = \text{ALLOCATING} \wedge \text{mem-block-addr-valid}$   
 $x \ m$   
**shows**  $(\exists t. \text{allocating-node } x \ t = \text{Some } m)$   
**proof**–  
**have**  $\text{inv}:\text{inv-aux-vars } Va \wedge \text{inv-bitmap } Va \wedge \text{inv-mempool-info } Va \wedge \text{inv-bitmap-freelist}$   
 $Va$   
**using**  $a0$  **unfolding**  $\text{inv-def}$  **by**  $\text{auto}$   
**have**  $\text{from-l-gt0}:0 \leq \text{from-l } Va \ t$  **using**  $a6 \ a5$  **by**  $\text{linarith}$   
**have**  $\text{block-valid-va}:\text{mem-block-addr-valid } Va \ m$   
**using**  $a2 \ a12 \ a21 \ \text{mp-alloc-stm4-buf} \ \text{mp-alloc-stm4-maxsz}$   
**unfolding**  $\text{mem-block-addr-valid-def}$  **by**  $\text{auto}$   
**have**  $\text{data-m}:\text{data } m =$   
 $\text{buf } (\text{mem-pool-info } x \ (\text{pool } m)) + (\text{block } m) * ((\text{max-sz } (\text{mem-pool-info}$   
 $x \ (\text{pool } m))) \text{div } (4 \wedge (\text{level } m)))$   
**using**  $a21$  **unfolding**  $\text{mem-block-addr-valid-def}$  **by**  $\text{auto}$   
**have**  $\text{inv-aux-va}:(\forall n. \text{get-bit } (\text{mem-pool-info } Va) \ (\text{pool } n) \ (\text{level } n) \ (\text{block } n) =$   
 $\text{ALLOCATING} \wedge \text{mem-block-addr-valid } Va \ n$   
 $\longrightarrow (\exists t. \text{allocating-node } Va \ t = \text{Some } n))$   
**using**  $a0$  **unfolding**  $\text{inv-def} \ \text{inv-aux-vars-def}$   
**by**  $\text{blast}$   
**{ assume**  $a00:(\text{pool } m) \neq p$   
**then obtain**  $t'$  **where**  $\text{allocating-node } Va \ t' = \text{Some } m$  **using**  $\text{inv-aux-va}$   
**by**  $(\text{metis } a13 \ a21 \ \text{block-valid-va} \ \text{mp-alloc-stm4-pres-mpinfo})$   
**moreover have**  $t' \neq t$  **using**  $a2$  **unfolding**  $\text{inv-def} \ \text{inv-aux-vars-def}$   
**using**  $a00 \ a7$  **calculation** **by**  $\text{auto}$   
**then have**  $\text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t' = \text{allocating-node}$   
 $Va \ t'$   
**unfolding**  $\text{mp-alloc-stm4-pre-precond-f-def}$  **by**  $\text{auto}$   
**then have**  $\text{eq-alloc}:\text{allocating-node } Va \ t' = \text{allocating-node } x \ t'$   
**using**  $a20$  **by**  $\text{auto}$   
**ultimately have**  $?thesis$  **by**  $\text{auto}$   
**}**  
**moreover{**  
**assume**  $a01:\text{pool } m = p$   
**let**  $?i1=(\text{nat } (\text{from-l } Va \ t))$  **and**  
 $?j1=(\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{blk } Va \ t) \ (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va$   
 $t)))$  **and**  
 $?i2 = (\text{nat } (\text{from-l } Va \ t + 1))$  **and**  
 $?j2 = (\text{block-num } (\text{mem-pool-info } Va \ p) \ (\text{blk } Va \ t) \ (\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va$   
 $t))*4)$   
**have**  $\text{mem-pools}:\text{mem-pools } x = \text{mem-pools } Va$  **using**  $\text{mp-alloc-stm4-mempools}[OF$   
 $a12]$  **by**  $\text{auto}$   
**have**  $\text{inv-mempool-info-mp } Va \ p$   
**using**  $a2 \ \text{mem-pools} \ \text{inv}$  **unfolding**  $\text{inv-mempool-info-def} \ \text{Let-def}$  **by**  $\text{auto}$

```

note inv-mempool=this[simplified Let-def]
have from-l:from-l x = from-l Va
  using mp-alloc-stm4-froml[OF a18] by auto
have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
  using from-l-gt0 by auto
have i1-len:?i1 < length (levels (mem-pool-info Va p))
  using a9 a2 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have i2-len:?i2 < length (levels (mem-pool-info Va p))
  using a9 a2 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
  by auto
have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
  by (metis a0 a2 a8 a10 i1-len inv-mempool-info-def invariant.inv-def)
have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
  using i1-len i2-len j1-len inv-mempool from-l-suc
  by simp
have lsizes-x-va:lsizes x = lsizes Va
  by (simp add: a17 mp-alloc-stm4-pre-precond-f-lsz)
have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
  using mp-alloc-stm4-buf[OF a2 a12] by simp
have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
  using mp-alloc-stm4-maxsz[OF a2 a12] by simp
{ assume a02:¬(((level m)=?i1 ∧ (block m)=?j1) ∨
  ((level m)=?i2 ∧ (block m)≥ ?j2 ∧ (block m)< ?j2+4))
  then have a020':¬((level m)=?i1 ∧ (block m)=?j1) and
    a021': ¬((level m)=?i2 ∧ (block m)≥ ?j2 ∧ (block m)< ?j2 + 4)
  by force+
  then have eq-get-bit-i-j:get-bit-s x p (level m) (block m) = get-bit-s Va p
    (level m) (block m)
    using same-bit-mp-alloc-x-va[OF a14[simplified
      a18[simplified mp-alloc-stm4-froml[OF a18], THEN sym]] a15, of level
m block m]
    using a020' a021' by auto
  then have get-bit-s Va (pool m) (level m) (block m) = ALLOCATING ∧
    mem-block-addr-valid Va m
    using a01 a21 block-valid-va by auto
  then obtain t' where allocating-node Va t' = Some m using inv-aux-va by
    auto
  moreover have t'≠t using a02 a7 a10 calculation by auto
  then have allocating-node (mp-alloc-stm4-pre-precond-f Va t p) t' = allocating-node
    Va t'
    unfolding mp-alloc-stm4-pre-precond-f-def by auto
  then have eq-alloc:allocating-node Va t' = allocating-node x t'
    using a20 by auto
  ultimately have ?thesis by auto
}
moreover{
  assume a02:((level m)=?i1 ∧ (block m)=?j1)

```



```

then have get-bit-s x p ?i1 ?j1 = DIVIDED
by (simp add: a14 from-l-gt0 from-l-suc i1-len j1-len mp-alloc-stm4-pre-precond-f-froml

      same-bit-mp-alloc-stm4-pre-precond-divided)
then have ?thesis using a21 a02 a01 by auto
}
moreover{
  assume a02:(level m)=?i2  $\wedge$  (block m) $\geq$  ?j2  $\wedge$  (block m)< ?j2+4
  then have block-n:(block-num (mem-pool-info Va p)
    (blk Va t) (lsizes Va t ! nat (from-l Va t))) = n
  proof–
    have lsizes Va t ! nat (from-l Va t) =
      ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^
      (nat (from-l Va t))
    using a3 lsizes-x-va a5 a6 a9 a11 a5 from-l by auto
    thus ?thesis using block-n inv a2 a10 a4 a9 from-l-gt0
    by blast
  qed
  then have get-bit-s x p ?i2 ?j2 = ALLOCATING
    using i2-len j2-len a15 get-bit-x-l1-b4[OF - from-l-gt0 a15, of ?i2 ?j2]
  by (metis (no-types, lifting) add-2-eq-Suc' add-Suc-right add-lessD1 mult.commute)

  { assume a03:block m = ?j2
    then have m = (pool = p, level = ?i2, block = ?j2,
      data = buf (mem-pool-info x p) +
      ?j2 * (max-sz (mem-pool-info x p) div 4 ^ ?i2)
    ) using data-m a03 a02 a01 by auto
    moreover have blk x t = buf (mem-pool-info x p) +
      ?j1 * ((max-sz (mem-pool-info x p) div 4 ^ ?i1))
    using a10[simplified buf[THEN sym] maxsz[THEN sym]] block-n a19
    mp-alloc-stm4-blk
    by metis
    then have allocating-node x t = Some (pool = p, level = ?i2, block = ?j2,
      data = buf (mem-pool-info x p) +
      ?j1 * (max-sz (mem-pool-info x p) div 4 ^
    ?i1))
    using a20 a19 mp-alloc-stm4-blk mp-alloc-stm4-pre-precond-f-allocating
    by (metis mult.commute)
    ultimately have ?thesis using buf maxsz next-level-addr-eq unfolding
    addr-def
    by (metis from-l-suc i2-len inv-mempool)
  }
  moreover {
    assume a07:(block m) $\neq$ ?j2
    then have (level m) = ?i2  $\wedge$  ((block m) = Suc ?j2  $\vee$ 
      (block m) = Suc (Suc ?j2)  $\vee$  (block m) = Suc (Suc (Suc ?j2)))
    using a02 by auto
    then have get-bit-s x p (level m) (block m) = FREE
    using a02 i2-len j2-len

```

```

      get-bit-x-l1-b41 [OF - from-l-gt0 [simplified from-l a18]
        a14 [simplified a18 [THEN sym] from-l] a15, of level m
    block m]
      by (metis Suc-lessD mult.commute)
      then have ?thesis using a21 a01 by auto
    }
  ultimately have ?thesis by auto
}
ultimately have ?thesis by auto
} ultimately show ?thesis by auto
qed

```

**lemma** *mp-alloc-stm4-inv-aux-vars5*:

```

  assumes
    a0:inv Va and
    a1:freeing-node x = freeing-node (mp-alloc-stm4-pre-precond-f Va t p) and
    a2: t1 ≠ t2 ∧ freeing-node x t1 = Some n1 ∧ freeing-node x t2 = Some n2
  shows ¬ (pool n1 = pool n2 ∧ level n1 = level n2 ∧ block n1 = block n2)
  proof-
    have t1 ≠ t2 ∧ freeing-node Va t1 = Some n1 ∧ freeing-node Va t2 = Some n2
      using a1 a2
    by (metis mp-alloc-stm4-pre-precond-f-def-frnode)
    then have ¬ (pool n1 = pool n2 ∧ level n1 = level n2 ∧ block n1 = block n2)
      using a0 unfolding inv-def inv-aux-vars-def by auto
    then show ?thesis
      by blast
  qed

```

**lemma** *mp-alloc-stm4-inv-aux-vars6*:

```

  assumes
    a0:inv Va and
    a1:freeing-node Va t = None and
    a2:p ∈ mem-pools Va and
    a3:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
    a4:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a5:¬ free-l Va t < OK and
    a6:free-l Va t ≤ from-l Va t and
    a7:allocating-node Va t =
      Some (|pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
            4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t)),
        data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
            4 ^ nat (from-l Va t)))|) and
    a8:block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
        (from-l Va t)))
      (lsizes Va t ! nat (from-l Va t))

```

$< n\text{-max} \text{ (mem-pool-info Va p) } * 4 \wedge \text{nat (from-l Va t) and}$   
 $a9\text{:from-l Va t} < \text{alloc-l Va t and}$   
 $a10\text{:blk Va t} = \text{buf (mem-pool-info Va p) } + n * (\text{max-sz (mem-pool-info Va p)}$   
 $\text{div } 4 \wedge \text{nat (from-l Va t)}) \text{ and}$   
 $a11\text{:}(x, \text{mp-alloc-stm4-pre-precond-f Va t p}) \in \text{gvars-conf-stable and}$   
 $a12\text{:}\forall pa. pa \neq p \longrightarrow \text{mem-pool-info x pa} = \text{mem-pool-info (mp-alloc-stm4-pre-precond-f}$   
 $\text{Va t p) pa and}$   
 $a13\text{:}\forall jj. jj \neq \text{nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)} \longrightarrow$   
 $\text{levels (mem-pool-info x p) ! jj} = \text{levels (mem-pool-info (mp-alloc-stm4-pre-precond-f}$   
 $\text{Va t p) p) ! jj and}$   
 $a14\text{:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va}$   
 $\text{t p) t + 1))} =$   
 $\text{list-updates-n}$   
 $(\text{bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !}$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))})$   
 $(\text{Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4)) } 3 \text{ FREE and}$   
 $a15\text{:free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f}$   
 $\text{Va t p) t + 1))} =$   
 $\text{inserts}$   
 $(\text{map (}\lambda ii. \text{lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !}$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *}$   
 $ii +$   
 $\text{blk (mp-alloc-stm4-pre-precond-f Va t p) t)}$   
 $[\text{Suc NULL..} < 4])$   
 $(\text{free-list}$   
 $(\text{levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !}$   
 $\text{nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))}) \text{ and}$   
 $a16\text{:lsizes x} = \text{lsizes (mp-alloc-stm4-pre-precond-f Va t p) and}$   
 $a17\text{:from-l x} = \text{from-l (mp-alloc-stm4-pre-precond-f Va t p) and}$   
 $a18\text{:allocating-node x} = \text{allocating-node (mp-alloc-stm4-pre-precond-f Va t p) and}$   
 $a19\text{:t1} \neq t2 \wedge \text{allocating-node x t1} = \text{Some n1} \wedge \text{allocating-node x t2} = \text{Some n2}$   
**shows**  $\neg (\text{pool n1} = \text{pool n2} \wedge \text{level n1} = \text{level n2} \wedge \text{block n1} = \text{block n2})$   
**proof**–  
**have**  $\text{inv:inv-aux-vars Va} \wedge \text{inv-bitmap Va} \wedge \text{inv-mempool-info Va} \wedge \text{inv-bitmap-freelist}$   
 $\text{Va}$   
**using**  $a0$  **unfolding**  $\text{inv-def}$  **by**  $\text{auto}$   
**have**  $\text{from-l-gt0:0} \leq \text{from-l Va t}$  **using**  $a6$   $a5$  **by**  $\text{linarith}$   
**let**  $?i1 = (\text{nat (from-l Va t)})$  **and**  
 $?j1 = (\text{block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l}$   
 $\text{Va t))})$  **and**  
 $?i2 = (\text{nat (from-l Va t + 1)})$  **and**  
 $?j2 = (\text{block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l}$   
 $\text{Va t))} * 4)$   
**have**  $\text{mem-pools:mem-pools x} = \text{mem-pools Va}$  **using**  $\text{mp-alloc-stm4-mempools[OF}$   
 $a11]$  **by**  $\text{auto}$   
**have**  $\text{inv-mempool-info-mp Va p}$   
**using**  $a2$   $\text{mem-pools inv}$  **unfolding**  $\text{inv-mempool-info-def}$   $\text{Let-def}$  **by**  $\text{auto}$   
**note**  $\text{inv-mempool=this[simplified Let-def]}$   
**have**  $\text{from-l:from-l x} = \text{from-l Va}$

```

    using mp-alloc-stm4-froml[OF a17] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a9 a2 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a9 a2 a4 from-l-gt0 inv unfolding inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis a0 a2 a8 a10 i1-len inv-mempool-info-def invariant.inv-def)
  have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
    ?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp
  have lsizes-x-va:lsizes x = lsizes Va using mp-alloc-stm4-pre-precond-f-lsz a16
    by auto
  have maxsz:max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p)
    using mp-alloc-stm4-maxsz[OF a2 a11] by simp
  have buf:buf (mem-pool-info x p) = buf (mem-pool-info Va p)
    using mp-alloc-stm4-buf[OF a2 a11] by simp
  have alloc-i1-j1:get-bit-s Va p ?i1 ?j1 = ALLOCATING
    using a7 a0 a10 unfolding inv-aux-vars-def invariant.inv-def
    by (metis (no-types) Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3))

  have nexisti2:noexist-bits (mem-pool-info Va p) ?i2 ?j2
    using a2 conjunct1[OF conjunct2[OF inv], simplified Let-def inv-bitmap-def]
  i1-len j1-len
    alloc-i1-j1 from-l-suc i2-len i1-len j1-len a1
    by (smt One-nat-def Suc-pred add commute inv-mempool nat-add-left-cancel-less
    plus-1-eq-Suc)
  { assume t≠t1 and t≠t2
    then have ?thesis
      using a0 a19 a18 inv-aux-vars-def
      unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def by force
    }
  moreover {
    assume a00:t=t1
    then have t2≠t using a19 by auto
    then have t2-same-allocating-node-Va:allocating-node Va t2 = Some n2
      using a0 a19 a18
      unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def inv-aux-vars-def
  by force
    then have get-bit-n2:get-bit-s Va (pool n2) (level n2) (block n2) = ALLOCAT-
    ING
      using a0 a19 a18 inv-aux-vars-def
      unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def by force
    have ¬ (pool n1 = pool n2 ∧ level n1 = level n2 ∧ block n1 = block n2) =
      (pool n1 = pool n2 ⟶ ¬(level n1 = level n2 ∧ block n1 = block n2))

```

```

    by auto
  moreover {
    assume a02:pool n1 = pool n2
    have n1 = (pool = p, level = ?i2, block = ?j2,
              data = blk Va t
            ) using a19
    by (simp add: a00 a18 mp-alloc-stm4-pre-precond-f-allocating mult.commute)

    then have ¬(level n1 = level n2 ∧ block n1 = block n2)
      using get-bit-n2 a02 nexisti2 by auto
  }
  then have ?thesis by auto
}
moreover {
  assume a00:t=t2
  then have t1≠t using a19 by auto
  then have t2-same-allocating-node-Va:allocating-node Va t1 = Some n1
    using a0 a19 a18 inv-aux-vars-def
  unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def by force
  then have get-bit-n2:get-bit-s Va (pool n1) (level n1) (block n1) = ALLOCAT-
ING
    using a0 a19 a18 inv-aux-vars-def
  unfolding mp-alloc-stm4-pre-precond-f-def invariant.inv-def by force
  have ¬(pool n1 = pool n2 ∧ level n1 = level n2 ∧ block n1 = block n2) =
    (pool n1 = pool n2 → ¬(level n1 = level n2 ∧ block n1 = block n2))
    by auto
  moreover {
    assume a02:pool n1 = pool n2
    have n2 = (pool = p, level = ?i2, block = ?j2,
              data = blk Va t
            ) using a19
    by (simp add: a00 a18 mp-alloc-stm4-pre-precond-f-allocating mult.commute)

    then have ¬(level n1 = level n2 ∧ block n1 = block n2)
      using get-bit-n2 a02 nexisti2 by auto
  }
  then have ?thesis by auto
}
ultimately show ?thesis by auto
qed

```

**lemma** *mp-alloc-stm4-inv-aux-vars7*:  
**assumes**  
*a0:inv Va and*  
*a1:freeing-node Va t = None and*  
*a2:p ∈ mem-pools Va and*  
*a3:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info Va p)) div 4 ^ ii and*  
*a4:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and*

$a5: \text{alloc-l } Va \ t < \text{int } (n\text{-levels } (mem\text{-pool-info } Va \ p)) \text{ and}$   
 $a6: \neg \text{free-l } Va \ t < OK \text{ and}$   
 $a7: \text{free-l } Va \ t \leq \text{from-l } Va \ t \text{ and}$   
 $a8: \text{allocating-node } Va \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
 $\text{block} = \text{block-num } (mem\text{-pool-info } Va \ p)$   
 $(\text{buf } (mem\text{-pool-info } Va \ p) + n * (\text{max-sz } (mem\text{-pool-info } Va \ p) \text{ div}$   
 $4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)),$   
 $\text{data} = \text{buf } (mem\text{-pool-info } Va \ p) + n * (\text{max-sz } (mem\text{-pool-info } Va \ p) \text{ div}$   
 $4 \wedge \text{nat } (\text{from-l } Va \ t))) \text{ and}$   
 $a9: n = \text{block-num } (mem\text{-pool-info } Va \ p)$   
 $(\text{buf } (mem\text{-pool-info } Va \ p) + n * (\text{max-sz } (mem\text{-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat}$   
 $(\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (mem\text{-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL} \text{ and}$   
 $a10: \text{block-num } (mem\text{-pool-info } Va \ p)$   
 $(\text{buf } (mem\text{-pool-info } Va \ p) + n * (\text{max-sz } (mem\text{-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat}$   
 $(\text{from-l } Va \ t)))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t))$   
 $< n\text{-max } (mem\text{-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \text{ and}$   
 $a11: \text{from-l } Va \ t < \text{alloc-l } Va \ t \text{ and}$   
 $a12: \text{blk } Va \ t = \text{buf } (mem\text{-pool-info } Va \ p) + n * (\text{max-sz } (mem\text{-pool-info } Va \ p)$   
 $\text{div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) \text{ and}$   
 $a23: (x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \text{ and}$   
 $a14: \forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ pa \text{ and}$   
 $a15: \forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels } (mem\text{-pool-info } x \ p) ! jj = \text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ p) ! jj \text{ and}$   
 $a16: \text{bits } (\text{levels } (mem\text{-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va$   
 $t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$   
 $(\text{Suc } (bn (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE} \text{ and}$   
 $a17: \text{free-list } (\text{levels } (mem\text{-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f}$   
 $Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $ii +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (mem\text{-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))) \text{ and}$   
 $a18: \text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$   
 $a19: \text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \text{ and}$

*a20:freeing-node*  $x = \text{freeing-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
*a21:allocating-node*  $x = \text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
*a22:allocating-node*  $x \ t1 = \text{Some } n1 \wedge \text{freeing-node } x \ t2 = \text{Some } n2$   
**shows**  $\neg (\text{pool } n1 = \text{pool } n2 \wedge \text{level } n1 = \text{level } n2 \wedge \text{block } n1 = \text{block } n2)$   
**proof** –  
 {**assume**  $\text{pool } n1 = \text{pool } n2$   
   **moreover have**  $\text{get-bit-s } x \ (\text{pool } n1) \ (\text{level } n1) \ (\text{block } n1) = \text{ALLOCATING}$   
     **using**  $\text{mp-alloc-stm4-inv-aux-vars3}$  **assms by** *blast*  
   **moreover have**  $\text{get-bit-s } x \ (\text{pool } n2) \ (\text{level } n2) \ (\text{block } n2) = \text{FREEING}$   
     **using**  $\text{mp-alloc-stm4-inv-aux-vars1}$  **assms by** *blast*  
   **ultimately have** *?thesis* **by** *auto*  
 } **thus** *?thesis* **by** *auto*  
**qed**

**lemma** *mp-alloc-stm4-inv-aux-vars*:

**assumes**  
*a0:inv*  $Va$  **and**  
*a1:freeing-node*  $Va \ t = \text{None}$  **and**  
*a2:p*  $\in \text{mem-pools } Va$  **and**  
*a3:* $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge ii$  **and**  
*a4:length*  $(\text{lsizes } Va \ t) \leq n\text{-levels } (\text{mem-pool-info } Va \ p)$  **and**  
*a5:alloc-l*  $Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p))$  **and**  
*a6:* $\neg \text{free-l } Va \ t < \text{OK}$  **and**  
*a7:free-l*  $Va \ t \leq \text{from-l } Va \ t$  **and**  
*a8:allocating-node*  $Va \ t =$   
    $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
      $\text{block} = \text{block-num } (\text{mem-pool-info } Va \ p)$   
        $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t))))$   
      $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)),$   
      $\text{data} = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$  **and**  
*a9:n*  $= \text{block-num } (\text{mem-pool-info } Va \ p)$   
    $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
    $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)) \vee$   
    $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL}$  **and**  
*a10:block-num*  $(\text{mem-pool-info } Va \ p)$   
    $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))$   
    $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t))$   
    $< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t)$  **and**  
*a11:from-l*  $Va \ t < \text{alloc-l } Va \ t$  **and**  
*a12:blk*  $Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t))$  **and**  
*a13:* $(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable}$  **and**  
*a14:* $\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ pa$  **and**

$a15: \forall jj. jj \neq \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels} (\text{mem-pool-info } x \ p) ! jj = \text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj$  **and**  
 $a16: \text{bits} (\text{levels} (\text{mem-pool-info } x \ p) ! \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits} (\text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$   
 $(\text{Suc} (\text{bn} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE}$  **and**  
 $a17: \text{free-list} (\text{levels} (\text{mem-pool-info } x \ p) ! \text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map} (\lambda ii. \text{lsizes} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) * ii +$   
 $\text{blk} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels} (\text{mem-pool-info} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat} (\text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)))$  **and**  
 $a18: \text{lsizes } x = \text{lsizes} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a19: \text{from-l } x = \text{from-l} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a20: \text{freeing-node } x = \text{freeing-node} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a21: \text{allocating-node } x = \text{allocating-node} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a22: \text{alloc-l } Va \ t = \text{int} (\text{length} (\text{lsizes } Va \ t)) - 1 \wedge \text{length} (\text{lsizes } Va \ t) = n\text{-levels}$   
 $(\text{mem-pool-info } Va \ p) \vee$   
 $\text{alloc-l } Va \ t = \text{int} (\text{length} (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t ! \text{nat} (\text{alloc-l } Va \ t + 1)$   
 $< sz$  **and**  
 $a23: \text{blk } x = \text{blk} (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$   
**shows**  $\text{inv-aux-vars } x$  **unfolding**  $\text{inv-aux-vars-def}$   
**using**  
 $\text{mp-alloc-stm4-inv-aux-vars1} [OF \text{ assms}(1-9, 11-22)]$   
 $\text{mp-alloc-stm4-inv-aux-vars2} [OF \text{ assms}(1-3, 6-8, 11-17, 20-21)]$   
 $\text{mp-alloc-stm4-inv-aux-vars3} [OF \text{ assms}(1-3, 6-7, 8-9, 11-20, 22)]$   
 $\text{mp-alloc-stm4-inv-aux-vars4} [OF \text{ assms}(1-4, 6-9, 11-13, 23, 14-20, 24, 22)]$   
 $\text{mp-alloc-stm4-inv-aux-vars5} [OF \text{ assms}(1, 21)]$   
 $\text{mp-alloc-stm4-inv-aux-vars6} [OF \text{ assms}(1-3, 5-9, 11-20, 22)]$   $\text{mp-alloc-stm4-inv-aux-vars7} [OF$   
 $\text{assms}(1-22)]$   
**by** *auto*

**lemma**  $\text{mp-alloc-stm4-inv-bitmap0}$ :

**assumes**  
 $a0: \text{inv } Va$  **and**  
 $a1: \text{freeing-node } Va \ t = \text{None}$  **and**  
 $a2: p \in \text{mem-pools } Va$  **and**  
 $a3: \forall ii < \text{length} (\text{lsizes } Va \ t). \text{lsizes } Va \ t ! ii = \text{ALIGN4} (\text{max-sz} (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge ii$  **and**  
 $a4: \text{length} (\text{lsizes } Va \ t) \leq n\text{-levels} (\text{mem-pool-info } Va \ p)$  **and**  
 $a5: \text{alloc-l } Va \ t < \text{int} (n\text{-levels} (\text{mem-pool-info } Va \ p))$  **and**



$a6: \neg \text{free-l } Va \ t < OK$  **and**  
 $a7: \text{free-l } Va \ t \leq \text{from-l } Va \ t$  **and**  
 $a8: \text{allocating-node } Va \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
 $\text{block} = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)),$   
 $\text{data} = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t)))$  **and**  
 $a9: n = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL}$  **and**  
 $a10: \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t))))$   
 $(\text{lsizes } Va \ t ! \text{nat } (\text{from-l } Va \ t))$   
 $< n\text{-max } (\text{mem-pool-info } Va \ p) * 4 ^ \wedge \text{nat } (\text{from-l } Va \ t)$  **and**  
 $a11: \text{from-l } Va \ t < \text{alloc-l } Va \ t$  **and**  
 $a12: \text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 ^ \wedge \text{nat } (\text{from-l } Va \ t))$  **and**  
 $a13: (x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable}$  **and**  
 $a14: \forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ pa$  **and**  
 $a15: \forall jj. jj \neq \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) \longrightarrow$   
 $\text{levels } (\text{mem-pool-info } x \ p) ! jj = \text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) ! jj$  **and**  
 $a16: \text{bits } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{list-updates-n}$   
 $(\text{bits } (\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))))$   
 $(\text{Suc } (\text{bn } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t * 4)) \ 3 \text{ FREE}$  **and**  
 $a17: \text{free-list } (\text{levels } (\text{mem-pool-info } x \ p) ! \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)) =$   
 $\text{inserts}$   
 $(\text{map } (\lambda ii. \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1) *$   
 $ii +$   
 $\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
 $[\text{Suc } \text{NULL}..<4])$   
 $(\text{free-list}$   
 $(\text{levels } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p) !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1))))$  **and**  
 $a18: \text{lsizes } x = \text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a19: \text{from-l } x = \text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**  
 $a20: \text{freeing-node } x = \text{freeing-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p)$  **and**

```

a21:allocating-node x = allocating-node (mp-alloc-stm4-pre-precond-f Va t p) and
a22:alloc-l Va t = int (length (lsizes Va t)) - 1  $\wedge$  length (lsizes Va t) = n-levels
(mem-pool-info Va p)  $\vee$ 
alloc-l Va t = int (length (lsizes Va t)) - 2  $\wedge$  lsizes Va t ! nat (alloc-l Va t + 1)
< sz and
a23:blk x = blk (mp-alloc-stm4-pre-precond-f Va t p)
shows inv-bitmap0 x
proof(simp add: inv-bitmap0-def Let-def )
{ fix p' j
  assume a00:p'  $\in$  mem-pools x
  assume a01:j < length (bits (levels (mem-pool-info x p') ! NULL))
  { assume p'  $\neq$  p
    then have get-bit-s x p' NULL j  $\neq$  NOEXIST
    by (metis a0 a00 a01 a13 a14 inv-bitmap0-def
      invariant.inv-def mp-alloc-stm4-mempools mp-alloc-stm4-pres-mpinfo)
  }
  moreover { assume a02:p' = p
    let ?i1=(nat (from-l Va t)) and
    ?j1= (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))*4)
    have from-l-gt0:(nat (from-l Va t + 1)) > 0
    using a6 a7 by linarith
    have zero-lt-len-levels:0 < length (levels (mem-pool-info x p))
    by (metis a0 a13 a2 inv-mempool-info-def invariant.inv-def mp-alloc-stm4-lvl-len)
    then have len-eq:length (bits (levels (mem-pool-info x p) ! 0)) =
      length (bits (levels (mem-pool-info Va p) ! 0))
    using a13 mp-alloc-stm4-inv-bits-len
    unfolding gvars-conf-stable-def gvars-conf-def
    by fastforce
    have from-l-gt0:0  $\leq$  from-l Va t using a7 a6 by linarith
    { assume a04:j = ?j1
      then have get-bit-s x p' NULL j  $\neq$  NOEXIST using a00 a01 a02
        get-bit-x-l-b a13
        mp-alloc-stm4-lvl-len[OF a2 a13] len-eq mp-alloc-stm4-froml[OF a19]
        from-l-gt0 a19 a0 a15 a2 same-bit-mp-alloc-stm4-pre-precond-divided
      unfolding inv-bitmap0-def inv-def apply auto
      using mp-alloc-stm4-pre-precond-f-same-bits zero-lt-len-levels
      by (smt BlockState.distinct(19) nat-0-iff )
    }
    moreover {
      assume a04:j  $\neq$  ?j1
      then have eq-get-bit-i-j:get-bit-s x p 0 j = get-bit-s Va p 0 j
      using same-bit-mp-alloc-x-va[OF a15[simplified a19[simplified mp-alloc-stm4-froml[OF
a19], THEN sym]] a16, of 0 j]
      using from-l-gt0 by auto
      then have get-bit-s x p' NULL j  $\neq$  NOEXIST

```

```

    using a0 unfolding inv-def inv-bitmap0-def a00 a01
    by (metis a01 a02 a2 len-eq)
  } ultimately have get-bit-s x p' NULL j ≠ NOEXIST by auto
  } ultimately have get-bit-s x p' NULL j ≠ NOEXIST by auto
} then show  $\forall p \in \text{mem-pools } x.$ 
 $\forall i < \text{length } (\text{bits } (\text{levels } (\text{mem-pool-info } x p) ! \text{NULL})).$ 
 $\text{get-bit-s } x p \text{ NULL } i \neq \text{NOEXIST}$  by auto
qed

```

**lemma** *mp-alloc-stm4-inv-bitmapn*:

```

  assumes
    a0:inv Va and
    a1:p ∈ mem-pools Va and
    a2:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a3:¬ free-l Va t < OK and
    a4:free-l Va t ≤ from-l Va t and
    a5:block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
        (from-l Va t)))
      (lsizes Va t ! nat (from-l Va t))
      < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
    a6:from-l Va t < alloc-l Va t and
    a7:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
      4 ^ nat (from-l Va t)) and
    a8:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and
    a9:∀ pa. pa ≠ p ⟶ mem-pool-info x pa = mem-pool-info (mp-alloc-stm4-pre-precond-f
      Va t p) pa and
    a10:∀ jj. jj ≠ nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) ⟶
      levels (mem-pool-info x p) ! jj = levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
        Va t p) p) ! jj and
    a11:bits (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f Va
      t p) t + 1)) =
      list-updates-n
        (bits (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
          nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1)))
        (Suc (bn (mp-alloc-stm4-pre-precond-f Va t p) t * 4)) 3 FREE and
    a12:free-list (levels (mem-pool-info x p) ! nat (from-l (mp-alloc-stm4-pre-precond-f
      Va t p) t + 1)) =
      inserts
        (map (λii. lsizes (mp-alloc-stm4-pre-precond-f Va t p) t !
          nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1) *
          ii +
          blk (mp-alloc-stm4-pre-precond-f Va t p) t)
          [Suc NULL..<4])
        (free-list
          (levels (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p) p) !
            nat (from-l (mp-alloc-stm4-pre-precond-f Va t p) t + 1))) and
    a13:from-l x = from-l (mp-alloc-stm4-pre-precond-f Va t p)
  shows inv-bitmapn x

```

```

proof(simp add: inv-bitmapn-def Let-def )
{ fix p' j
  let ?k = (length (levels (mem-pool-info x p')) - Suc 0)
assume a00:p' ∈ mem-pools x
assume a01:j < length (bits (levels (mem-pool-info x p') ! ?k))
{ assume p' ≠ p
  then have get-bit-s x p' ?k j ≠ DIVIDED
    using a00 a01 a0 a8 a9 mp-alloc-stm4-mempools mp-alloc-stm4-pres-mpinfo
    unfolding inv-bitmapn-def inv-def
    by (metis One-nat-def)
  }
moreover { assume a02:p' = p
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l Va
t))*4)
  have from-l-gt0:(nat (from-l Va t + 1)) > 0
    using a3 a4 by linarith
  have zero-lt-len-levels:0 < length (levels (mem-pool-info x p))
  by (metis a0 a8 a1 inv-mempool-info-def invariant.inv-def mp-alloc-stm4-lvl-len)
  then have len-eq:length (bits (levels (mem-pool-info x p) ! 0)) =
    length (bits (levels (mem-pool-info Va p) ! 0))
    using a8 mp-alloc-stm4-inv-bits-len
    unfolding gvars-conf-stable-def gvars-conf-def
    by fastforce
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a8] by auto
  have inv-mempool-info-mp Va p
    using a1 mem-pools a0 unfolding inv-def inv-mempool-info-def Let-def by
auto
  note inv-mempool=this[simplified Let-def]
  have from-l:from-l x = from-l Va
    using mp-alloc-stm4-froml[OF a13] by auto
  have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
    using from-l-gt0 by auto
  have i1-len:?i1 < length (levels (mem-pool-info Va p))
    using a6 a1 a2 from-l-gt0 a0 unfolding inv-def inv-mempool-info-def Let-def
    by auto
  have i2-len:?i2 < length (levels (mem-pool-info Va p))
    using a0 a6 a1 a2 from-l-gt0 unfolding inv-def inv-mempool-info-def Let-def
    by auto
  have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
    by (metis a0 a1 a5 a7 i1-len inv-mempool-info-def invariant.inv-def)
  have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va p) !
?i2))
    using i1-len i2-len j1-len inv-mempool from-l-suc
    by simp

```

```

have from-l-gt0:0 ≤ from-l Va t using a4 a3 by linarith
{ assume a03:?i2 = ?k
  { assume a04:j ≥ ?j2 ∧ j < ?j2+4
    { assume a05:j = ?j2
      then have get-bit-s x p' ?i2 j = ALLOCATING using a00 a01 a02 a03
        get-bit-x-l1-b4[OF - from-l-gt0 a11 i2-len, of ?j2] a8 zero-lt-len-levels
        mp-alloc-stm4-lvl-len[OF a1 a8] len-eq mp-alloc-stm4-froml[OF a13]
        from-l-gt0 a13 j2-len
        by (meson Suc-lessD mult.commute)
      then have get-bit-s x p' ?i2 j ≠ DIVIDED by auto
    }
    moreover {
      assume a05:j ≥ ?j2 + 1 ∧ j < ?j2+4
      then have get-bit-s x p' ?i2 j = FREE using a00 a01 a02 a03
        get-bit-x-l1-b41[OF - from-l-gt0 a10[simplified from-l a13[THEN sym]]
a11 i2-len, of ?j2] a8 zero-lt-len-levels
        mp-alloc-stm4-lvl-len[OF a1 a8] len-eq mp-alloc-stm4-froml[OF a13]
        from-l-gt0 a13 j2-len a11 mp-alloc-stm4-pre-precond-f-bn
        by (smt One-nat-def add.commute add-Suc-shift length-list-update-n
list-updates-n-eq
        numeral-2-eq-2 numeral-3-eq-3 numeral-Bit0 plus-1-eq-Suc)
      then have get-bit-s x p' ?i2 j ≠ DIVIDED by auto
    } ultimately have get-bit-s x p' ?i2 j ≠ DIVIDED using a04 by fastforce
  }
  moreover {
    assume ¬(j ≥ ?j2 ∧ j < ?j2+4)
    moreover have eq-get-bit-i-j:get-bit-s x p ?i2 j = get-bit-s Va p ?i2 j
      using a03 from-l-suc same-bit-mp-alloc-x-va[OF
a10[simplified a13[simplified mp-alloc-stm4-froml[OF a13], THEN sym]]
a11, of ?i2 j]
      from-l-gt0 calculation
      by force
    ultimately have get-bit-s x p' ?i2 j ≠ DIVIDED
      using a0 a02 a03 unfolding inv-def inv-bitmapn-def Let-def
      by (metis One-nat-def a01 a8 a10 a11 a1 length-list-update-n
mp-alloc-stm4-inv-bits-len mp-alloc-stm4-lvl-len)
    } ultimately have get-bit-s x p' ?k j ≠ DIVIDED using a03 by auto
  }
  moreover {
    assume ?i2 ≠ ?k
    moreover have ?i2 < ?k
      using calculation a00 a02 a8 i2-len mem-pools mp-alloc-stm4-lvl-len by auto
    then have ?i1 ≠ ?k
      by linarith
    ultimately have eq-get-bit-i-j:get-bit-s x p ?k j = get-bit-s Va p ?k j
      using from-l-suc same-bit-mp-alloc-x-va[OF
a10[simplified a13[simplified mp-alloc-stm4-froml[OF a13], THEN sym]]
a11, of ?i2 j]
      from-l-gt0

```

```

    by (metis a10 a13 from-l same-bit-mp-alloc-stm4-pre-precond-f1)
  then have get-bit-s x p' ?k j ≠ DIVIDED
    using a0 a02 unfolding inv-def inv-bitmapn-def Let-def
    by (metis One-nat-def a01 a8 a10 a11 a1 length-list-update-n
        mp-alloc-stm4-inv-bits-len mp-alloc-stm4-lvl-len)
  } ultimately have get-bit-s x p' ?k j ≠ DIVIDED by auto
} ultimately have get-bit-s x p' ?k j ≠ DIVIDED by auto
}
then show ∀ p ∈ mem-pools x.
  ∀ i < length (bits (levels (mem-pool-info x p) ! (length (levels (mem-pool-info
x p)) - Suc NULL))).
    get-bit-s x p (length (levels (mem-pool-info x p)) - Suc NULL) i ≠
DIVIDED by auto

```

qed

**lemma** mp-alloc-stm4-inv-bitmap4free:

```

  assumes
    a0:inv Va and
    a1:freeing-node Va t = None and
    a2:p ∈ mem-pools Va and
    a3:∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info
Va p)) div 4 ^ ii and
    a4:length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) and
    a5:alloc-l Va t < int (n-levels (mem-pool-info Va p)) and
    a6:¬ free-l Va t < OK and
    a7:free-l Va t ≤ from-l Va t and
    a8:allocating-node Va t =
      Some (pool = p, level = nat (from-l Va t),
        block = block-num (mem-pool-info Va p)
          (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t)))
          (lsizes Va t ! nat (from-l Va t)),
        data = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div
4 ^ nat (from-l Va t))) and
    a9:n = block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
      (lsizes Va t ! nat (from-l Va t)) ∨
    max-sz (mem-pool-info Va p) div 4 ^ nat (from-l Va t) = NULL and
    a10:block-num (mem-pool-info Va p)
      (buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4 ^ nat
(from-l Va t)))
      (lsizes Va t ! nat (from-l Va t))
    < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) and
    a11:from-l Va t < alloc-l Va t and
    a12:blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p)
div 4 ^ nat (from-l Va t)) and
    a13:(x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable and

```

```

a14:  $\forall pa. pa \neq p \longrightarrow mem\text{-}pool\text{-}info\ x\ pa = mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ pa$  and
a15:  $\forall jj. jj \neq nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1) \longrightarrow$ 
 $levels\ (mem\text{-}pool\text{-}info\ x\ p)!\ jj = levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)!\ jj$  and
a16:  $bits\ (levels\ (mem\text{-}pool\text{-}info\ x\ p)!\ nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)) =$ 
 $list\text{-}updates\text{-}n$ 
 $(bits\ (levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)!\$ 
 $nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)))$ 
 $(Suc\ (bn\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t * 4))\ 3\ FREE$  and
a17:  $free\text{-}list\ (levels\ (mem\text{-}pool\text{-}info\ x\ p)!\ nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)) =$ 
 $inserts$ 
 $(map\ (\lambda ii. lsize\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t!\$ 
 $nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1) *$ 
 $ii +$ 
 $blk\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t)$ 
 $[Suc\ NULL..<4])$ 
 $(free\text{-}list$ 
 $(levels\ (mem\text{-}pool\text{-}info\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ p)!\$ 
 $nat\ (from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)\ t + 1)))$  and
a18:  $lsize\ x = lsize\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)$  and
a19:  $from\text{-}l\ x = from\text{-}l\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)$  and
a20:  $freeing\text{-}node\ x = freeing\text{-}node\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)$  and
a21:  $allocating\text{-}node\ x = allocating\text{-}node\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)$  and
a22:  $alloc\text{-}l\ Va\ t = int\ (length\ (lsize\ Va\ t)) - 1 \wedge length\ (lsize\ Va\ t) = n\text{-}levels$ 
 $(mem\text{-}pool\text{-}info\ Va\ p) \vee$ 
 $alloc\text{-}l\ Va\ t = int\ (length\ (lsize\ Va\ t)) - 2 \wedge lsize\ Va\ t!\ nat\ (alloc\text{-}l\ Va\ t + 1)$ 
 $< sz$  and
a23:  $blk\ x = blk\ (mp\text{-}alloc\text{-}stm4\text{-}pre\text{-}precond\text{-}f\ Va\ t\ p)$ 
shows  $inv\text{-}bitmap\text{-}not4free\ x$ 
proof–
{ fix  $p'\ i\ j$ 
assume  $a00: p' \in mem\text{-}pools\ x$  and
 $a01: i < length\ (levels\ (mem\text{-}pool\text{-}info\ x\ p'))$  and
 $a02: j < length\ (bits\ (levels\ (mem\text{-}pool\text{-}info\ x\ p')!\ i))$ 
{ assume  $a03: 0 < i$  and
 $a04: get\text{-}bit\text{-}s\ x\ p'\ i\ (Suc\ (Suc\ (j\ div\ 4 * 4))) = FREE$  and
 $a05: get\text{-}bit\text{-}s\ x\ p'\ i\ (Suc\ (j\ div\ 4 * 4)) = FREE$  and
 $a06: get\text{-}bit\text{-}s\ x\ p'\ i\ (j\ div\ 4 * 4) = FREE$ 
{ assume  $p' \neq p$ 
then have  $get\text{-}bit\text{-}s\ x\ p'\ i\ (j\ div\ 4 * 4 + 3) \neq FREE$ 
using  $a00\ a01\ a0\ a8$  using  $a00\ a01\ a0\ a8\ a9\ mp\text{-}alloc\text{-}stm4\text{-}mempools$ 
 $mp\text{-}alloc\text{-}stm4\text{-}pres\text{-}mpinfo$ 
unfolding  $inv\text{-}bitmap\text{-}not4free\text{-}def\ inv\text{-}def$ 
by  $(metis\ a02\ a03\ a04\ a05\ a06\ a13\ a14\ add.\text{commute}$ 
 $add\text{-}2\text{-}eq\text{-}Suc'\ partner\text{-}bits\text{-}def\ plus\text{-}1\text{-}eq\text{-}Suc)$ 
} note  $not\text{-}p = this$ 

```

```

moreover{
  assume a07:p' = p
  let ?i1=(nat (from-l Va t)) and
    ?j1=(block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))) and
    ?i2 = (nat (from-l Va t + 1)) and
    ?j2 = (block-num (mem-pool-info Va p) (blk Va t) (lsizes Va t ! nat (from-l
Va t))*4)
  have from-l-gt0:(nat (from-l Va t + 1)) > 0
    using a6 a7 by linarith
  have zero-lt-len-levels:0 < length (levels (mem-pool-info x p))
    using a0 a2 mp-alloc-stm4-lvl-len unfolding inv-mempool-info-def inv-def
    using a01 a07 gr-implies-not0 by blast
  then have len-eq:length (bits (levels (mem-pool-info x p) ! 0)) =
    length (bits (levels (mem-pool-info Va p) ! 0))
    using a13 mp-alloc-stm4-inv-bits-len
    unfolding gvars-conf-stable-def gvars-conf-def
    by fastforce
  have mem-pools:mem-pools x = mem-pools Va using mp-alloc-stm4-mempools[OF
a13] by auto
    have inv-mempool-info-mp Va p
      using a2 mem-pools a0 unfolding inv-def inv-mempool-info-def Let-def
by auto
    note inv-mempool=this[simplified Let-def]
    have from-l:from-l x = from-l Va
      using mp-alloc-stm4-froml[OF a19] by auto
    have from-l-suc:nat (from-l Va t + 1) = nat(from-l Va t) + 1
      using from-l-gt0 by auto
    have i1-len:?i1 < length (levels (mem-pool-info Va p))
      using a2 a11 a5 from-l-gt0 a0 unfolding inv-def inv-mempool-info-def
Let-def
      by auto
    have i2-len:?i2 < length (levels (mem-pool-info Va p))
      using a0 a5 a2 a11 from-l-gt0 unfolding inv-def inv-mempool-info-def
Let-def
      by auto
    have j1-len:?j1 < length (bits (levels (mem-pool-info Va p) ! ?i1))
      using assms(11) assms(13) i1-len inv-mempool by presburger
    have j2-len:Suc (Suc (Suc ?j2)) < length (bits (levels (mem-pool-info Va
p) ! ?i2))
      using i1-len i2-len j1-len inv-mempool from-l-suc
      by simp

    have from-l-gt0:0 ≤ from-l Va t
      using a6 a7 by linarith
    { assume a08:i ≠ ?i1 ∧ i ≠ ?i2
      then have eq-get-bit-i-j:get-bit-s x p i (j div 4 * 4 + 3) = get-bit-s Va
p i (j div 4 * 4 + 3)
        using same-bit-mp-alloc-x-va

```



```

      [OF a15[simplified a19[simplified mp-alloc-stm4-froml[OF a19],
        THEN sym]] a16, of i (j div 4 * 4 + 3)]
    from-l-gt0 by auto
  moreover have eq-get-bit-i-j:get-bit-s x p i (j div 4 * 4) = get-bit-s Va
p i (j div 4 * 4)
    using same-bit-mp-alloc-x-va
      [OF a15[simplified a19[simplified mp-alloc-stm4-froml[OF a19],
        THEN sym]] a16, of i (j div 4 * 4)] a08
    from-l-gt0 by auto
  moreover have eq-get-bit-i-j:get-bit-s x p i (Suc (j div 4 * 4)) = get-bit-s
Va p i (Suc (j div 4 * 4))
    using same-bit-mp-alloc-x-va
      [OF a15[simplified a19[simplified mp-alloc-stm4-froml[OF a19],
        THEN sym]] a16, of i (Suc (j div 4 * 4))] a08
    from-l-gt0 by auto
  moreover have eq-get-bit-i-j:get-bit-s x p i (Suc (Suc (j div 4 * 4))) =
get-bit-s Va p i (Suc (Suc (j div 4 * 4)))
    using same-bit-mp-alloc-x-va
      [OF a15[simplified a19[simplified mp-alloc-stm4-froml[OF a19],
        THEN sym]] a16, of i (Suc (Suc (j div 4 * 4)))] a08
    from-l-gt0 by auto
  ultimately have get-bit-s x p' i (j div 4 * 4 + 3) ≠ FREE
using a07 a03 a04 a05 a06 a01 a02 a0 a13 a15 a16 a2 mp-alloc-stm4-inv-bits-len
mp-alloc-stm4-lvl-len
  unfolding inv-bitmap-not4free-def inv-def Let-def partner-bits-def
by (metis add.commute add-2-eq-Suc' length-list-update-n plus-1-eq-Suc)

}
moreover { assume i=?i1
  then have get-bit-s x p' i (j div 4 * 4 + 3) ≠ FREE
    using not-p a0 a02 a03 a04 a05 a06 a15 a19 a2 from-l from-l-gt0
from-l-suc i1-len j1-len
      mp-alloc-stm4-inv-bits-len mp-alloc-stm4-pre-precond-f-bitmap-not-free
      same-bit-mp-alloc-stm4-pre-precond-f1
    unfolding inv-bitmap-not4free-def invariant.inv-def partner-bits-def
  by (smt add-2-eq-Suc' add-eq-self-zero le-zero-eq nat-int-add not-one-le-zero

      plus-1-eq-Suc )
}note i1= this
moreover {
  assume a08:i=?i2
  { assume j ≥ ?j2 ∧ j ≤ ?j2 + 3
    then have j = ?j2 ∨ j = ?j2 + 1 ∨ j = ?j2 + 2 ∨ j = ?j2 + 3
      by auto
    then have j div 4 * 4 = ?j2 by auto
    moreover have get-bit-s x p' i ?j2 = ALLOCATING
      using get-bit-x-l1-b4[OF - from-l-gt0 a16 i2-len ] a08 a07 mult.commute
j2-len
    by (metis Suc-lessD)
  }
}

```

```

ultimately have get-bit-s  $x \ p' \ i \ (j \text{ div } 4 * 4 + 3) \neq \text{FREE}$  using a06
by auto
}
moreover {
  assume  $\neg(j \geq ?j2 \wedge j \leq ?j2 + 3)$ 
  then have  $j < ?j2 \vee j > ?j2 + 3$ 
  by auto
  moreover { assume  $j < ?j2$ 
    then have  $j \text{ div } 4 * 4 + 3 < ?j2$ 
    by presburger
    moreover have get-bit-s  $x \ p \ i \ (j \text{ div } 4 * 4) = \text{get-bit-s } Va \ p \ i \ (j \text{ div } 4 * 4)$ 
      using same-bit-mp-alloc-x-va[OF a15[simplified from-l a19[THEN sym]] a16, of i (j div 4 * 4)]
      a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation
      a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc
      by (auto simp add: a16)
    moreover have get-bit-s  $x \ p \ i \ (j \text{ div } 4 * 4 + 1) = \text{get-bit-s } Va \ p \ i \ (j \text{ div } 4 * 4 + 1)$ 
      using same-bit-mp-alloc-x-va[OF a15[simplified from-l a19[THEN sym]] a16, of i (j div 4 * 4 + 1)]
      a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation
      a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc
      by (auto simp add: a16)
    moreover have get-bit-s  $x \ p \ i \ (j \text{ div } 4 * 4 + 2) = \text{get-bit-s } Va \ p \ i \ (j \text{ div } 4 * 4 + 2)$ 
      using same-bit-mp-alloc-x-va[OF a15[simplified from-l a19[THEN sym]] a16, of i (j div 4 * 4 + 2)]
      a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation
      a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc
      by (auto simp add: a16)
    moreover have get-bit-s  $x \ p \ i \ (j \text{ div } 4 * 4 + 3) = \text{get-bit-s } Va \ p \ i \ (j \text{ div } 4 * 4 + 3)$ 
      using same-bit-mp-alloc-x-va[OF a15[simplified from-l a19[THEN sym]] a16, of i (j div 4 * 4 + 3)]
      a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation
      a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc
      by (auto simp add: a16)
    ultimately have get-bit-s  $x \ p' \ i \ (j \text{ div } 4 * 4 + 3) \neq \text{FREE}$ 
    using same-bit-mp-alloc-x-va[OF - a16] a15 a01 a02 a03 a04 a05 a06 a07 a08 a00
      a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len i1
    unfolding inv-def inv-bitmap-not4free-def partner-bits-def
    by (smt add commute add-2-eq-Suc' length-list-update-n plus-1-eq-Suc)
  }
}
moreover {
  assume  $j > ?j2 + 3$ 
  then have  $j \text{ div } 4 * 4 > ?j2 + 3$ 

```

by *presburger*  
 moreover have  $\text{get-bit-s } x \ p \ i \ (j \ \text{div} \ 4 * 4) = \text{get-bit-s } Va \ p \ i \ (j \ \text{div} \ 4 * 4)$   
 using *same-bit-mp-alloc-x-va*[*OF* *a15*[*simplified from-l a19*[*THEN sym*]] *a16*, of *i* (*j div 4 \* 4*)]  
     *a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation*  
     *a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc*  
 by (*auto simp add: a16*)  
 moreover have  $\text{get-bit-s } x \ p \ i \ (j \ \text{div} \ 4 * 4 + 1) = \text{get-bit-s } Va \ p \ i \ (j \ \text{div} \ 4 * 4 + 1)$   
 using *same-bit-mp-alloc-x-va*[*OF* *a15*[*simplified from-l a19*[*THEN sym*]] *a16*, of *i* (*j div 4 \* 4 + 1*)]  
     *a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation*  
     *a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc*  
 by (*auto simp add: a16*)  
 moreover have  $\text{get-bit-s } x \ p \ i \ (j \ \text{div} \ 4 * 4 + 2) = \text{get-bit-s } Va \ p \ i \ (j \ \text{div} \ 4 * 4 + 2)$   
 using *same-bit-mp-alloc-x-va*[*OF* *a15*[*simplified from-l a19*[*THEN sym*]] *a16*, of *i* (*j div 4 \* 4 + 2*)]  
     *a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation*  
     *a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc*  
 by (*auto simp add: a16*)  
 moreover have  $\text{get-bit-s } x \ p \ i \ (j \ \text{div} \ 4 * 4 + 3) = \text{get-bit-s } Va \ p \ i \ (j \ \text{div} \ 4 * 4 + 3)$   
 using *same-bit-mp-alloc-x-va*[*OF* *a15*[*simplified from-l a19*[*THEN sym*]] *a16*, of *i* (*j div 4 \* 4 + 3*)]  
     *a01 a02 a03 a04 a05 a06 a07 a08 a00 calculation*  
     *a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len from-l-suc*  
 by (*auto simp add: a16*)  
 ultimately have  $\text{get-bit-s } x \ p' \ i \ (j \ \text{div} \ 4 * 4 + 3) \neq \text{FREE}$   
 using *same-bit-mp-alloc-x-va*[*OF* - *a16*] *a15 a01 a02 a03 a04 a05 a06 a07 a08 a00*  
     *a0 a16 a19 a2 from-l i2-len mp-alloc-stm4-inv-bits-len i1*  
 unfolding *inv-def inv-bitmap-not4free-def partner-bits-def*  
 by (*smt add.commute add-2-eq-Suc' length-list-update-n plus-1-eq-Suc*)  
 }  
 ultimately have  $\text{get-bit-s } x \ p' \ i \ (j \ \text{div} \ 4 * 4 + 3) \neq \text{FREE}$  by *auto*  
 } ultimately have  $\text{get-bit-s } x \ p' \ i \ (j \ \text{div} \ 4 * 4 + 3) \neq \text{FREE}$  by *auto*  
 } ultimately have  $\text{get-bit-s } x \ p' \ i \ (j \ \text{div} \ 4 * 4 + 3) \neq \text{FREE}$  by *auto*  
 } ultimately have  $\text{get-bit-s } x \ p' \ i \ (j \ \text{div} \ 4 * 4 + 3) \neq \text{FREE}$  by *auto*  
 }  
 } then show *inv-bitmap-not4free x*  
   unfolding *inv-bitmap-not4free-def Let-def partner-bits-def*  
   by *auto*  
 qed

**lemma** *mp-alloc-stm4-whlpst-in-post-inv:*  
*inv Va*  $\implies$   
*freeing-node Va t = None*  $\implies$

$p \in \text{mem-pools } Va \implies$   
 $ETIMEOUT \leq \text{timeout} \implies$   
 $\text{timeout} = ETIMEOUT \longrightarrow \text{tmout } Va \ t = ETIMEOUT \implies$   
 $\neg \text{rf } Va \ t \implies$   
 $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \wedge ii \implies$   
 $\text{length } (\text{lsizes } Va \ t) \leq \text{n-levels } (\text{mem-pool-info } Va \ p) \implies$   
 $\text{alloc-l } Va \ t < \text{int } (\text{n-levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $\neg \text{free-l } Va \ t < \text{OK} \implies$   
 $\text{NULL} < \text{buf } (\text{mem-pool-info } Va \ p) \vee \text{NULL} < n \wedge \text{NULL} < \text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) \implies$   
 $\text{free-l } Va \ t \leq \text{from-l } Va \ t \implies$   
 $\text{allocating-node } Va \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } Va \ t),$   
 $\text{block} = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))) \implies$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)),$   
 $\text{data} = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t))) \implies$   
 $n = \text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))) \implies$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)) \vee$   
 $\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t) = \text{NULL} \implies$   
 $\text{block-num } (\text{mem-pool-info } Va \ p)$   
 $(\text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)))) \implies$   
 $(\text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t))$   
 $< \text{n-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \implies$   
 $\text{from-l } Va \ t < \text{alloc-l } Va \ t \implies$   
 $\text{cur } Va = \text{Some } t \implies$   
 $n < \text{n-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \implies$   
 $\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) \implies$   
 $\text{mempoolalloc-ret } Va \ t = \text{None} \implies$   
 $\forall ii \leq \text{nat } (\text{alloc-l } Va \ t). \text{sz} \leq \text{lsizes } Va \ t \ ! \ ii \implies$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 1 \wedge \text{length } (\text{lsizes } Va \ t) = \text{n-levels } (\text{mem-pool-info } Va \ p) \vee$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t \ ! \ \text{nat } (\text{alloc-l } Va \ t + 1)$   
 $< \text{sz} \implies$   
 $i \ x \ t = 4 \implies$   
 $\text{cur } x = \text{cur } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \implies$   
 $\text{tick } x = \text{tick } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \implies$   
 $\text{thd-state } x = \text{thd-state } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \implies$   
 $(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \implies$   
 $\forall pa. pa \neq p \longrightarrow \text{mem-pool-info } x \ pa = \text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ pa \implies$   
 $\text{wait-q } (\text{mem-pool-info } x \ p) = \text{wait-q } (\text{mem-pool-info } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ p)$

$Va\ t\ p)\ p) \implies$   
 $\forall t'.\ t' \neq t \implies lvars-nochange\ t'\ x\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $\forall jj.\ jj \neq nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1) \implies$   
 $levels\ (mem-pool-info\ x\ p)!\ jj = levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)!\ jj \implies$   
 $bits\ (levels\ (mem-pool-info\ x\ p)!\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)) =$   
 $list-updates-n$   
 $(bits\ (levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)!\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)))$   
 $(Suc\ (bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t * 4))\ 3\ FREE \implies$   
 $free-list\ (levels\ (mem-pool-info\ x\ p)!\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1)) =$   
 $inserts$   
 $(map\ (\lambda ii.\ lsizes\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t!\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1) * ii + blk\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t)$   
 $[Suc\ NULL..<4])$   
 $(free-list$   
 $(levels\ (mem-pool-info\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ p)!\ nat\ (from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t + 1))) \implies$   
 $j\ x = j\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $ret\ x = ret\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $end\ x = end\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $rf\ x = rf\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $tmout\ x = tmout\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $lsizes\ x = lsizes\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $alloc-l\ x = alloc-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $free-l\ x = free-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $from-l\ x = from-l\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $blk\ x = blk\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $nodev\ x = nodev\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $bn\ x = bn\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $alloc-lsize-r\ x = alloc-lsize-r\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $lvl\ x = lvl\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $bb\ x = bb\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $block-pt\ x = block-pt\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $th\ x = th\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $need-resched\ x = need-resched\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $mempoolalloc-ret\ x = mempoolalloc-ret\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $freeing-node\ x = freeing-node\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies$   
 $allocating-node\ x = allocating-node\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \implies inv$   
 $x$

$\text{apply}(simp\ add:inv-def)$   
 $\text{apply}(rule\ conjI)\ \text{apply}(simp\ add:inv-cur-def)\ \text{apply}\ clarify\ using\ mp-alloc-stm4-inv-cur$   
 $\text{apply}\metis$   
 $\text{apply}(rule\ conjI)\ using\ mp-alloc-stm4-inv-thd-state\ \text{apply}\metis$   
 $\text{apply}(rule\ conjI)\ using\ mp-alloc-stm4-inv-mempool-info\ \text{apply}\metis$

**apply**(rule conjI) **using** mp-alloc-stm4-inv-bitmap-freelist **apply** blast  
**apply**(rule conjI) **using** mp-alloc-stm4-inv-bitmap **unfolding** inv-def **apply** blast  
**apply**(rule conjI) **using** mp-alloc-stm4-inv-aux-vars **unfolding** inv-def **apply** blast  
**apply**(rule conjI) **using** mp-alloc-stm4-inv-bitmap0 **unfolding** inv-def **apply** blast  
**apply** (rule conjI) **using** mp-alloc-stm4-inv-bitmapn **unfolding** inv-def **apply** blast  
**using** mp-alloc-stm4-inv-bitmap4free **unfolding** inv-def **by** blast

**lemma** mp-alloc-stm4-whlpst-in-post-h1:

$p \in \text{mem-pools } Va \implies$   
 $\text{inv } Va \implies$   
 $\text{alloc-l } Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $\text{from-l } Va \ t < \text{alloc-l } Va \ t \implies$   
 $\neg \text{free-l } Va \ t < 0 \implies$   
 $\text{free-l } Va \ t \leq \text{from-l } Va \ t \implies$   
 $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \text{ div } 4 \ ^\wedge ii \implies$   
 $\text{length } (\text{lsizes } Va \ t) \leq n\text{-levels } (\text{mem-pool-info } Va \ p) \implies$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 1 \wedge \text{length } (\text{lsizes } Va \ t) = n\text{-levels } (\text{mem-pool-info } Va \ p) \vee$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t \ ! \ \text{nat } (\text{alloc-l } Va \ t + 1) < \text{sz} \implies$   
 $\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \text{ div } 4 \ ^\wedge \text{nat } (\text{from-l } Va \ t)) \implies$   
 $(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \implies$   
 $\text{allocating-node } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1),$   
 $\text{block} = \text{block-num } (\text{mem-pool-info } x \ p) (\text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$

$(\text{lsizes } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t \ !$   
 $\text{nat } (\text{from-l } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t + 1)),$   
 $\text{data} = \text{blk } (\text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \ t)$   
**apply**(simp add:mp-alloc-stm4-pre-precond-f-def block-num-def)  
**apply**(rule subst[**where** s=buf (mem-pool-info Va p) **and** t=buf (mem-pool-info x p)])  
**apply**(simp add:gvars-conf-stable-def gvars-conf-def set-bit-def)

**apply**(rule subst[**where** s=n \* (max-sz (mem-pool-info Va p) div 4 ^ nat (from-l Va t)) **and** t=buf (mem-pool-info Va p) + n \* (max-sz (mem-pool-info Va p) div 4 ^ nat (from-l Va t)) -  
 $\text{buf } (\text{mem-pool-info } Va \ p)]])$   
**apply** arith

**apply**(subgoal-tac  $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = (\text{max-sz } (\text{mem-pool-info } Va \ p))$ )

$Va\ p)) \text{ div } 4 \wedge ii)$   
**prefer 2 using** *inv-maxsz-align4*[of *Va*] **apply** *metis*

**apply**(rule *subst*[**where**  $s=lsizes\ Va\ t ! nat\ (from-l\ Va\ t) \text{ div } 4$  **and**  $t=lsizes\ Va\ t ! nat\ (from-l\ Va\ t + 1)$ ])  
**apply** (*smt div-mult-self1-is-m mp-alloc-stm4-blockfit-help4 nat-less-iff*  
*semiring-normalization-rules(7) zero-less-numeral*)  
**by** (*smt div-eq-0-iff m-mod-div mod-mult-self2-is-0 mp-alloc-stm4-blockfit-help4*  
*nat-less-iff nonzero-mult-div-cancel-right semiring-normalization-rules(7)*)

**lemma** *mp-alloc-stm4-whlpst-in-post-h2*:  
 $p \in mem-pools\ Va \implies$   
 $inv\ Va \implies$   
 $alloc-l\ Va\ t < int\ (n-levels\ (mem-pool-info\ Va\ p)) \implies$   
 $from-l\ Va\ t < alloc-l\ Va\ t \implies$   
 $\neg free-l\ Va\ t < 0 \implies$   
 $free-l\ Va\ t \leq from-l\ Va\ t \implies$   
 $\forall ii < length\ (lsizes\ Va\ t). lsizes\ Va\ t ! ii = ALIGN4\ (max-sz\ (mem-pool-info\ Va\ p)) \text{ div } 4 \wedge ii \implies$   
 $length\ (lsizes\ Va\ t) \leq n-levels\ (mem-pool-info\ Va\ p) \implies$   
 $alloc-l\ Va\ t = int\ (length\ (lsizes\ Va\ t)) - 1 \wedge length\ (lsizes\ Va\ t) = n-levels\ (mem-pool-info\ Va\ p) \vee$   
 $alloc-l\ Va\ t = int\ (length\ (lsizes\ Va\ t)) - 2 \wedge lsizes\ Va\ t ! nat\ (alloc-l\ Va\ t + 1) < sz \implies$   
 $blk\ Va\ t = buf\ (mem-pool-info\ Va\ p) + n * (max-sz\ (mem-pool-info\ Va\ p) \text{ div } 4 \wedge nat\ (from-l\ Va\ t)) \implies$   
 $(x, mp-alloc-stm4-pre-precond-f\ Va\ t\ p) \in gvars-conf-stable \implies$   
 $data\ (the\ (allocating-node\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t)) =$   
 $buf\ (mem-pool-info\ x\ p) +$   
 $block\ (the\ (allocating-node\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t)) *$   
 $(max-sz\ (mem-pool-info\ x\ p) \text{ div } 4 \wedge level\ (the\ (allocating-node\ (mp-alloc-stm4-pre-precond-f\ Va\ t\ p)\ t)))$   
**apply**(*simp add:mp-alloc-stm4-pre-precond-f-def block-num-def*)

**apply**(rule *subst*[**where**  $s=buf\ (mem-pool-info\ Va\ p)$  **and**  $t=buf\ (mem-pool-info\ x\ p)$ ])  
**apply**(*simp add:gvars-conf-stable-def gvars-conf-def set-bit-def*)

**apply**(rule *subst*[**where**  $s=n * (max-sz\ (mem-pool-info\ Va\ p) \text{ div } 4 \wedge nat\ (from-l\ Va\ t))$  **and**  $t=buf\ (mem-pool-info\ Va\ p) + n * (max-sz\ (mem-pool-info\ Va\ p) \text{ div } 4 \wedge nat\ (from-l\ Va\ t)) -$   
 $buf\ (mem-pool-info\ Va\ p)$ ])  
**apply** *arith*

**apply**(*subgoal-tac*  $\forall ii < length\ (lsizes\ Va\ t). lsizes\ Va\ t ! ii = (max-sz\ (mem-pool-info\ Va\ p)) \text{ div } 4 \wedge ii)$   
**prefer 2 using** *inv-maxsz-align4*[of *Va*] **apply** *metis*

**apply**(rule subst[**where**  $s = \text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t) \ \text{div } 4$  **and**  $t = \text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t + 1)$ ])  
**apply** (smt div-mult-self1-is-m mp-alloc-stm4-blockfit-help4 nat-less-iff  
semiring-normalization-rules(7) zero-less-numeral)  
**apply**(rule subst[**where**  $s = \text{max-sz } (\text{mem-pool-info } Va \ p)$  **and**  $t = \text{max-sz } (\text{mem-pool-info } x \ p)$ ])  
**apply**(simp add:gvars-conf-stable-def gvars-conf-def set-bit-def)  
**apply**(rule subst[**where**  $s = \text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge \text{nat } (\text{from-l } Va \ t) \ \text{div } 4$   
**and**  $t = \text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge \text{nat } (\text{from-l } Va \ t + 1)$ ])  
**apply** (metis inv-maxsz-align4 mp-alloc-stm4-blockfit-help4 nonzero-mult-div-cancel-right  
zero-neq-numeral)  
**apply**(rule subst[**where**  $s = \text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge \text{nat } (\text{from-l } Va \ t)$   
**and**  $t = \text{lsizes } Va \ t \ ! \ \text{nat } (\text{from-l } Va \ t)$ ])  
**apply** (smt nat-less-iff)  
**apply**(subgoal-tac  $\exists m > 0. \text{max-sz } (\text{mem-pool-info } Va \ p) = (4 * m) * (4 \wedge n\text{-levels } (\text{mem-pool-info } Va \ p))$ )  
**prefer** 2 **apply**(simp add:inv-def inv-mempool-info-def Let-def) **apply** metis  
**by** (smt add-left-cancel inv-maxsz-align4 mp-alloc-stm4-blockfit-help4 mult.assoc  
mult-is-0 nonzero-mult-div-cancel-left semiring-normalization-rules(7))

**lemma** mp-alloc-stm4-whlpst-in-post-h3-1:  
 $\text{from-l } Va \ t \geq 0 \implies n < n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \implies$   
 $4 * n < n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t + 1)$   
**by** (smt mult.assoc Divides.div-mult2-eq Suc-nat-eq-nat-zadd1 div-eq-0-iff  
div-mult-mult1-if gr-implies-not0 mult commute mult-eq-0-iff power-Suc  
semiring-normalization-rules(7) zero-less-numeral)

**lemma** mp-alloc-stm4-whlpst-in-post-h3:  
 $p \in \text{mem-pools } Va \implies$   
 $\text{inv } Va \implies$   
 $\text{alloc-l } Va \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } Va \ p)) \implies$   
 $\text{from-l } Va \ t < \text{alloc-l } Va \ t \implies$   
 $\neg \text{free-l } Va \ t < 0 \implies$   
 $\text{free-l } Va \ t \leq \text{from-l } Va \ t \implies$   
 $\forall ii < \text{length } (\text{lsizes } Va \ t). \text{lsizes } Va \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } Va \ p)) \ \text{div } 4 \wedge ii \implies$   
 $\text{length } (\text{lsizes } Va \ t) \leq n\text{-levels } (\text{mem-pool-info } Va \ p) \implies$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 1 \wedge \text{length } (\text{lsizes } Va \ t) = n\text{-levels } (\text{mem-pool-info } Va \ p) \vee$   
 $\text{alloc-l } Va \ t = \text{int } (\text{length } (\text{lsizes } Va \ t)) - 2 \wedge \text{lsizes } Va \ t \ ! \ \text{nat } (\text{alloc-l } Va \ t + 1) < \text{sz} \implies$   
 $n < n\text{-max } (\text{mem-pool-info } Va \ p) * 4 \wedge \text{nat } (\text{from-l } Va \ t) \implies$   
 $(x, \text{mp-alloc-stm4-pre-precond-f } Va \ t \ p) \in \text{gvars-conf-stable} \implies$   
 $\text{blk } Va \ t = \text{buf } (\text{mem-pool-info } Va \ p) + n * (\text{max-sz } (\text{mem-pool-info } Va \ p) \ \text{div } 4 \wedge \text{nat } (\text{from-l } Va \ t)) \implies$



```

    block (the (allocating-node (mp-alloc-stm4-pre-precond-f Va t p) t))
      < n-max (mem-pool-info x p) * 4 ^ level (the (allocating-node (mp-alloc-stm4-pre-precond-f
        Va t p) t))
  apply(simp add:mp-alloc-stm4-pre-precond-f-def block-num-def)
  apply(rule subst[where s=max-sz (mem-pool-info Va p) div 4 ^ nat (from-l Va
    t)
      and t=lsizes Va t ! nat (from-l Va t)])
    using inv-maxsz-align4 apply auto[1]
  apply(rule subst[where s=n-max (mem-pool-info Va p) and t=n-max (mem-pool-info
    x p)])
    apply(simp add:mp-alloc-stm4-pre-precond-f-def set-bit-def gvars-conf-stable-def
      gvars-conf-def)
  apply(subgoal-tac  $\exists m > 0. \text{max-sz (mem-pool-info Va p) = (4 * m) * (4 ^ n-levels (mem-pool-info Va p))}$ )
  prefer 2 apply(simp add:inv-def inv-mempool-info-def Let-def) apply metis
  apply(subgoal-tac nat (from-l Va t) < n-levels (mem-pool-info Va p)) prefer 2
  apply linarith

  apply(rule subst[where s=n and t=(n * (max-sz (mem-pool-info Va p) div 4 ^
    nat (from-l Va t)) div
      (max-sz (mem-pool-info Va p) div 4 ^ nat (from-l Va t)))]])
  apply (simp add: mp-alloc-stm3-lm2-inv-1-2)
  apply clarsimp

  apply(rule mp-alloc-stm4-whlpst-in-post-h3-1)
  apply arith apply blast

done

lemma mp-alloc-stm4-whlpst-in-post-h4:
  p ∈ mem-pools Va ⇒
  inv Va ⇒
  alloc-l Va t < int (n-levels (mem-pool-info Va p)) ⇒
  from-l Va t < alloc-l Va t ⇒
  ¬ free-l Va t < 0 ⇒
  free-l Va t ≤ from-l Va t ⇒
  ∀ ii < length (lsizes Va t). lsizes Va t ! ii = ALIGN4 (max-sz (mem-pool-info Va
    p)) div 4 ^ ii ⇒
  length (lsizes Va t) ≤ n-levels (mem-pool-info Va p) ⇒
  alloc-l Va t = int (length (lsizes Va t)) - 1 ∧ length (lsizes Va t) = n-levels
    (mem-pool-info Va p) ∨
  alloc-l Va t = int (length (lsizes Va t)) - 2 ∧ lsizes Va t ! nat (alloc-l Va t +
    1) < sz ⇒
  n < n-max (mem-pool-info Va p) * 4 ^ nat (from-l Va t) ⇒
  blk Va t = buf (mem-pool-info Va p) + n * (max-sz (mem-pool-info Va p) div 4
    ^ nat (from-l Va t)) ⇒
  (x, mp-alloc-stm4-pre-precond-f Va t p) ∈ gvars-conf-stable ⇒
  (∃ n < n-max (mem-pool-info x p) * 4 ^ nat (from-l (mp-alloc-stm4-pre-precond-f
    Va t p) t + 1)).

```

```

      blk (mp-alloc-stm4-pre-precond-f Va t p) t =
      buf (mem-pool-info x p) +
      n * (max-sz (mem-pool-info x p) div 4 ^ nat (from-l (mp-alloc-stm4-pre-precond-f
      Va t p) t + 1)))
apply(rule subst[where s=n-max (mem-pool-info Va p) and t=n-max (mem-pool-info
x p)])
apply(simp add:gvars-conf-stable-def gvars-conf-def mp-alloc-stm4-pre-precond-f-def
set-bit-def)
apply(rule subst[where s=buf (mem-pool-info Va p) and t=buf (mem-pool-info
x p)])
apply(simp add:gvars-conf-stable-def gvars-conf-def mp-alloc-stm4-pre-precond-f-def
set-bit-def)
apply(rule subst[where s=from-l Va and t=from-l (mp-alloc-stm4-pre-precond-f
Va t p)])
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(rule subst[where s=blk Va and t=blk (mp-alloc-stm4-pre-precond-f Va t
p)])
apply(simp add:mp-alloc-stm4-pre-precond-f-def)
apply(rule subst[where s=max-sz (mem-pool-info Va p) and t=max-sz (mem-pool-info
x p)])
apply(simp add:gvars-conf-stable-def gvars-conf-def mp-alloc-stm4-pre-precond-f-def
set-bit-def)
apply(rule exI[where x=4 * n])
by (smt inv-maxsz-align4 mp-alloc-stm4-blockfit-help4 mp-alloc-stm4-whlpst-in-post-h3-1

      mult.assoc semiring-normalization-rules(7))

```

**lemma** mp-alloc-stm4-whlpst-in-post:

```

Va ∈ mp-alloc-precond2-1-1-loopinv-0 t p sz timeout ∩ {cur = Some t} ⇒
  mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va t p) t p ∩ {i t ≥ 4}
  ⊆ { (Pair Va) ∈ Mem-pool-alloc-guar t } ∩ mp-alloc-precond2-1-1-loopinv-1 t p
sz timeout
apply clarsimp
apply(rule conjI)
apply(simp add:Mem-pool-alloc-guar-def) apply clarsimp
apply(rule conjI)
apply(simp add:gvars-conf-stable-def gvars-conf-def)
apply(rule conjI) using mp-alloc-stm4-mempools2 apply metis
apply clarify
apply(rule conjI) using mp-alloc-stm4-inv-mif-buf apply metis
apply(rule conjI) using mp-alloc-stm4-inv-mif-mxsz apply metis
apply(rule conjI) using mp-alloc-stm4-inv-mif-nmax apply metis
apply(rule conjI) using mp-alloc-stm4-inv-mif-nlvs apply metis
apply(rule conjI) using mp-alloc-stm4-inv-mif-len apply metis
apply clarify using mp-alloc-stm4-inv-bits-len apply metis
apply(rule conjI)
using mp-alloc-stm4-whlpst-in-post-inv[of Va t p timeout - sz -] apply auto[1]
apply(rule conjI)

```

```

apply clarsimp
apply(subgoal-tac lvars-nochange t' x (mp-alloc-stm4-pre-precond-f Va t p))
  prefer 2 apply metis
apply(subgoal-tac lvars-nochange t' Va (mp-alloc-stm4-pre-precond-f Va t p))
  prefer 2 using mp-alloc-stm4-pre-precond-f-lvars-nochange[of - t Va p] apply
metis
  using lvars-nochange-trans[of - Va mp-alloc-stm4-pre-precond-f Va t p -]
    lvars-nochange-sym apply metis
using mp-alloc-stm4-pre-precond-f-tick apply metis

apply(rule conjI)
apply clarsimp
using mp-alloc-stm4-whlpst-in-post-inv[of Va t p timeout - sz -] apply auto[1]

apply(rule conjI)
apply clarsimp
using mp-alloc-stm4-pre-precond-f-def-frnode apply metis

apply(rule conjI)
apply clarsimp
using mp-alloc-stm4-pre-precond-f-mpls apply metis

apply(rule conjI)
apply clarsimp
apply(rule conjI) apply clarsimp
  apply(subgoal-tac rf Va t) prefer 2 using mp-alloc-stm4-pre-precond-f-rf ap-
ply metis
  apply fast
apply(rule conjI) apply clarsimp
  using mp-alloc-stm4-pre-precond-f-ret apply metis
apply clarsimp using mp-alloc-stm4-pre-precond-f-tmout apply metis

apply(rule conjI)
apply clarsimp
apply(subgoal-tac rf Va t) prefer 2 using mp-alloc-stm4-pre-precond-f-rf apply
metis
  apply fast

apply(rule conjI)
apply clarsimp
apply(rule conjI)
apply clarsimp
apply(subgoal-tac max-sz (mem-pool-info x p) = max-sz (mem-pool-info Va p))
  prefer 2 apply(subgoal-tac max-sz (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p)
    = max-sz (mem-pool-info Va p))
  prefer 2 using mp-alloc-stm4-pre-maxsz apply metis
apply(simp add:gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac lsizes (mp-alloc-stm4-pre-precond-f Va t p) t = lsizes Va t)

```

```

    prefer 2 using mp-alloc-stm4-pre-precond-f-lsz apply metis
    apply metis
    apply(subgoal-tac n-levels (mem-pool-info x p) = n-levels (mem-pool-info Va p))
    prefer 2 apply(subgoal-tac n-levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p)
    = n-levels (mem-pool-info Va p))
    prefer 2 using mp-alloc-stm4-inv-mif-nlvl apply metis
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
    apply(rule conjI)
    apply(subgoal-tac lsizes (mp-alloc-stm4-pre-precond-f Va t p) t = lsizes Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-lsz apply metis
    apply metis
    apply(rule conjI)
    apply(subgoal-tac alloc-l (mp-alloc-stm4-pre-precond-f Va t p) t = alloc-l Va t)

    prefer 2 using mp-alloc-stm4-pre-precond-f-allocl apply metis
    apply metis
    apply(rule conjI)
    apply(rule subst[where t= free-l (mp-alloc-stm4-pre-precond-f Va t p) t and
s= free-l Va t])
    using mp-alloc-stm4-pre-precond-f-freel apply metis
    apply linarith
    apply(rule conjI)
    apply(subgoal-tac alloc-l (mp-alloc-stm4-pre-precond-f Va t p) t = alloc-l Va t)

    prefer 2 using mp-alloc-stm4-pre-precond-f-allocl apply metis
    apply(rule subst[where t= free-l (mp-alloc-stm4-pre-precond-f Va t p) t and
s= free-l Va t])
    using mp-alloc-stm4-pre-precond-f-freel apply metis
    apply linarith

    apply(rule disjI2)
    apply(rule subst[where s=alloc-l Va and t=alloc-l (mp-alloc-stm4-pre-precond-f
Va t p)])
    using mp-alloc-stm4-pre-precond-f-allocl apply metis
    apply(rule subst[where s=lsizes Va and t=lsizes (mp-alloc-stm4-pre-precond-f
Va t p)])
    using mp-alloc-stm4-pre-precond-f-lsz apply metis
    apply(rule conjI) apply linarith
    apply(rule conjI) apply blast
    apply(subgoal-tac n-levels (mem-pool-info x p) = n-levels (mem-pool-info Va p))
    prefer 2 apply(subgoal-tac n-levels (mem-pool-info (mp-alloc-stm4-pre-precond-f
Va t p) p)
    = n-levels (mem-pool-info Va p))
    prefer 2 using mp-alloc-stm4-inv-mif-nlvl apply metis
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
    apply metis

    apply(rule conjI)

```

```

apply clarsimp
apply(subgoal-tac alloc-l (mp-alloc-stm4-pre-precond-f Va t p) t = alloc-l Va t)
  prefer 2 using mp-alloc-stm4-pre-precond-f-allocl apply metis
apply arith

apply(rule conjI)
  apply clarsimp
  apply(subgoal-tac free-l (mp-alloc-stm4-pre-precond-f Va t p) t = free-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-freel apply metis
  apply arith

apply(rule conjI)
  apply clarsimp
  apply(subgoal-tac blk Va t > 0) prefer 2
    apply(simp add:inv-def inv-mempool-info-def)
  apply(subgoal-tac blk (mp-alloc-stm4-pre-precond-f Va t p) t = blk Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-blk apply metis
  apply arith

apply(rule conjI)
  apply clarsimp
  apply(subgoal-tac alloc-l (mp-alloc-stm4-pre-precond-f Va t p) t = alloc-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-allocl apply metis
  apply(subgoal-tac from-l (mp-alloc-stm4-pre-precond-f Va t p) t = from-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-froml apply metis
  apply arith

apply clarsimp
apply(rule conjI)
  apply(subgoal-tac alloc-l (mp-alloc-stm4-pre-precond-f Va t p) t = alloc-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-allocl apply metis
  apply(subgoal-tac from-l (mp-alloc-stm4-pre-precond-f Va t p) t = from-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-froml apply metis
  apply arith

apply(rule conjI)
  apply(subgoal-tac from-l (mp-alloc-stm4-pre-precond-f Va t p) t = from-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-froml apply metis
  apply(subgoal-tac free-l (mp-alloc-stm4-pre-precond-f Va t p) t = free-l Va t)
    prefer 2 using mp-alloc-stm4-pre-precond-f-freel apply metis
  apply arith

apply(rule conjI)
  using mp-alloc-stm4-whlpst-in-post-h1 apply blast

apply(rule conjI)
  using mp-alloc-stm4-whlpst-in-post-h2 apply blast

apply(rule conjI)

```

```

using mp-alloc-stm4-whlpst-in-post-h3 apply blast

using mp-alloc-stm4-whlpst-in-post-h4 apply blast
done

lemma thd-state (mp-alloc-stm4-pre-precond-f Va t p) = thd-state Va
  by(simp add:mp-alloc-stm4-pre-precond-f-def)

lemma thd-state (mp-alloc-stm4-pre-precond-f Va t p) = thd-state Va
  by(simp add:mp-alloc-stm4-pre-precond-f-def)

lemma  $\forall p \in \text{mem-pools } Va. \text{wait-q (mem-pool-info Va p) = wait-q (mem-pool-info (mp-alloc-stm4-pre-precond-f Va t p1) p)}$ 
  apply clarify
  apply(simp add:mp-alloc-stm4-pre-precond-f-def)
  apply(simp add: set-bit-def)
done

term mp-alloc-precond2-1-1-loopinv-0 t p sz timeout
term mp-alloc-precond2-1-1-loopinv-1 t p sz timeout

lemma mp-alloc-stm4-lm1-1:
  Va  $\in$  mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap \{\text{'cur} = \text{Some } t\} \implies$ 
   $\Gamma \vdash_I \text{Some ( 'bn := 'bn (t := block-num ( 'mem-pool-info p) ( 'blk t) ( ('lsizes t)!(nat ( 'from-l t)))));}$ 
  'mem-pool-info := set-bit-divide 'mem-pool-info p (nat ( 'from-l t)) ( 'bn t);;
  'mem-pool-info := set-bit-allocating 'mem-pool-info p (nat ( 'from-l t + 1))
  (4 * 'bn t);;
  'allocating-node := 'allocating-node (t := Some (pool = p, level = nat ( 'from-l t + 1),
  block = 4 * 'bn t, data = 'blk t ));;
  FOR 'i := 'i (t := Suc 0);
  'i t < 4;
  'i := 'i (t := Suc ( 'i t)) DO
  'lbn := 'lbn (t := 4 * 'bn t + 'i t);;
  'lsz := 'lsz (t := ( 'lsizes t) ! (nat ( 'from-l t + 1)));;
  'block2 := 'block2(t := 'lsz t * 'i t + 'blk t);;
  'mem-pool-info := set-bit-free 'mem-pool-info p (nat ( 'from-l t + 1)) ( 'lbn t);;

  IF block-fits ( 'mem-pool-info p) ( 'block2 t) ( 'lsz t) THEN
    'mem-pool-info := 'mem-pool-info (p :=
      append-free-list ( 'mem-pool-info p) (nat ( 'from-l t + 1)) ( 'block2 t) )
  FI
  ROF satp [{Va}, {(s, t). s = t}, UNIV,
   $\{\text{'(Pair Va) } \in \text{Mem-pool-alloc-guar } t\} \cap \text{mp-alloc-precond2-1-1-loopinv-1 } t$ 
  p sz timeout]

```

```

apply(rule Seq[where mid={mp-alloc-stm4-pre-precond4
                        (mp-alloc-stm4-pre-precond3
                        (mp-alloc-stm4-pre-precond2
                        (mp-alloc-stm4-pre-precond1 Va t p) t p) t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-pre-precond3
                        (mp-alloc-stm4-pre-precond2
                        (mp-alloc-stm4-pre-precond1 Va t p) t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-pre-precond2
                        (mp-alloc-stm4-pre-precond1 Va t p) t p}])
apply(rule Seq[where mid={mp-alloc-stm4-pre-precond1 Va t p}])

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Seq[where mid={mp-alloc-stm4-pre-precond-f Va t p}])
apply(rule Basic)
apply(simp add:mp-alloc-stm4-pre-precond-f-def) apply simp apply(simp add:stable-def)
apply(simp add:stable-def)

apply(rule Conseq[where pre={mp-alloc-stm4-pre-precond-f Va t p}
                and pre'=mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f Va
t p) t p
                and rely={s, t}.s = t} and rely'={s, t}.s = t} and guar=UNIV
and guar'=UNIV
                and post'=mp-alloc-stm4-loopinv (mp-alloc-stm4-pre-precond-f
Va t p) t p  $\cap$  {i t  $\geq$  4})
                using mp-alloc-stm4-pre-precond-f-in-mp-alloc-stm4-loopinv apply auto[1]
                apply simp apply simp using mp-alloc-stm4-whlpst-in-post[of Va t p timeout
sz] apply argo
                using mp-alloc-stm4-while[of Va t p timeout sz] apply fastforce
done

term mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap$  {cur = Some t}
term mp-alloc-precond2-1-1-loopinv-1 t p sz timeout
term { (Pair Va)  $\in$  Mem-pool-alloc-guar t }  $\cap$  mp-alloc-precond2-1-1-loopinv-1 t
p sz timeout

lemma mp-alloc-stm4-lm1:

```

$mp\_alloc\_precond2-1-1-loopinv-0 \ t \ p \ sz \ timeout \cap \{\cur = Some \ t\} \cap \{Va\} = \{Va\} \implies$   
 $\Gamma \vdash_I Some \ ( \ 'bn := 'bn(t := block\_num \ ( \ 'mem\_pool\_info \ p) \ ( \ blk \ t) \ ( \ lsizes \ t ! nat \ ( \ 'from-l \ t))) ; ;$   
 $\quad \quad \quad \ 'mem\_pool\_info := set\_bit\_divide \ 'mem\_pool\_info \ p \ (nat \ ( \ 'from-l \ t)) \ ( \ 'bn \ t) ; ;$   
 $\quad \quad \quad \ 'mem\_pool\_info := set\_bit\_allocating \ 'mem\_pool\_info \ p \ (nat \ ( \ 'from-l \ t + 1))$   
 $\quad \quad \quad \ (4 * \ 'bn \ t) ; ;$   
 $\quad \quad \quad \ 'allocating\_node := 'allocating\_node(t \mapsto \langle pool = p, level = nat \ ( \ 'from-l \ t + 1), block = 4 * \ 'bn \ t, data = \ 'blk \ t \rangle) ; ;$   
 $\quad \quad \quad \ ( \ 'i := \ 'i(t := Suc \ NULL) ; ;$   
 $\quad \quad \quad \ WHILE \ 'i \ t < 4$   
 $\quad \quad \quad \ DO \ 'lbn := \ 'lbn(t := 4 * \ 'bn \ t + \ 'i \ t) ; ; \ 'lsz := \ 'lsz(t := \ 'lsizes \ t ! nat$   
 $\quad \quad \quad \ ( \ 'from-l \ t + 1)) ; ;$   
 $\quad \quad \quad \ 'block2 := \ 'block2(t := \ 'lsz \ t * \ 'i \ t + \ 'blk \ t) ; ;$   
 $\quad \quad \quad \ 'mem\_pool\_info := set\_bit\_free \ 'mem\_pool\_info \ p \ (nat \ ( \ 'from-l \ t + 1))$   
 $\quad \quad \quad \ ( \ 'lbn \ t) ; ;$   
 $\quad \quad \quad \ IF \ block\_fits \ ( \ 'mem\_pool\_info \ p) \ ( \ 'block2 \ t)$   
 $\quad \quad \quad \ \ ( \ 'lsz \ t) \ THEN \ 'mem\_pool\_info := \ 'mem\_pool\_info$   
 $\quad \quad \quad \ \ \ \ \ \ (p := append\_free\_list \ ( \ 'mem\_pool\_info \ p) \ (nat \ ( \ 'from-l \ t +$   
 $\quad \quad \quad \ 1)) \ ( \ 'block2 \ t)) \ FI ; ;$   
 $\quad \quad \quad \ \ 'i := \ 'i(t := Suc \ ( \ 'i \ t))$   
 $\quad \quad \quad \ OD)) \ sat_p \ [mp\_alloc\_precond2-1-1-loopinv-0 \ t \ p \ sz \ timeout \cap \{\cur = Some$   
 $\quad \quad \quad \ t\} \cap \{Va\},$   
 $\quad \quad \quad \ \{(s, t). s = t\}, \ UNIV,$   
 $\quad \quad \quad \ \{\cur(Pair \ Va) \in Mem\_pool\_alloc\_guar \ t\} \cap mp\_alloc\_precond2-1-1-loopinv-1$   
 $\quad \quad \quad \ t \ p \ sz \ timeout]$   
 $\quad \quad \quad \ apply(rule \ subst[where \ t=mp\_alloc\_precond2-1-1-loopinv-0 \ t \ p \ sz \ timeout \cap$   
 $\quad \quad \quad \ \{\cur = Some \ t\} \cap \{Va\} \ and \ s=\{Va\}])$   
 $\quad \quad \quad \ apply \ metis$   
 $\quad \quad \quad \ apply(subgoal\_tac \ Va \in mp\_alloc\_precond2-1-1-loopinv-0 \ t \ p \ sz \ timeout \cap \{\cur$   
 $\quad \quad \quad \ = Some \ t\})$   
 $\quad \quad \quad \ prefer \ 2 \ apply \ auto[1]$   
 $\quad \quad \quad \ using \ mp\_alloc\_stm4-lm1-1 \ apply \ meson$   
 $\quad \quad \quad \ done$

**term**  $mp\_alloc\_precond2-1-1-loopinv \ t \ p \ sz \ timeout$

**term**  $mp\_alloc\_precond2-1-2 \ t \ p \ sz \ timeout$

**lemma**  $mp\_alloc\_stm4-lm:$

$\Gamma \vdash_I Some \ ( \ WHILE \ 'from-l \ t < \ 'alloc-l \ t \ DO$   
 $\quad \quad \quad \ (* \ == \ start: \ blk = break\_block(p, blk, from-l, lsizes); *)$   
 $\quad \quad \quad \ (t \blacktriangleright \ ATOMIC$   
 $\quad \quad \quad \ \ 'bn := \ 'bn \ (t := block\_num \ ( \ 'mem\_pool\_info \ p) \ ( \ blk \ t) \ (( \ 'lsizes \ t)!(nat$   
 $\quad \quad \quad \ ( \ 'from-l \ t)))) ; ;$   
 $\quad \quad \quad \ \ 'mem\_pool\_info := set\_bit\_divide \ 'mem\_pool\_info \ p \ (nat \ ( \ 'from-l \ t)) \ ( \ 'bn \ t) ; ;$   
 $\quad \quad \quad \ \ 'mem\_pool\_info := set\_bit\_allocating \ 'mem\_pool\_info \ p \ (nat \ ( \ 'from-l \ t + 1))$   
 $\quad \quad \quad \ \ (4 * \ 'bn \ t) ; ;$



```

    (* set the allocating node info of the thread *)
    'allocating-node := 'allocating-node (t := Some (pool = p, level = nat
('from-l t + 1),
    block = 4 * 'bn t, data = 'blk t ));

FOR 'i := 'i (t := 1);
    'i t < 4;
    'i := 'i (t := 'i t + 1) DO
    'lbn := 'lbn (t := 4 * 'bn t + 'i t);
    'lsz := 'lsz (t := ('lsizes t) ! (nat ('from-l t + 1)));
    'block2 := 'block2(t := 'lsz t * 'i t + 'blk t);

    (* set-free-bit(p, l + 1, lbn); *)
    'mem-pool-info := set-bit-free 'mem-pool-info p (nat ('from-l t + 1)) ('lbn
t));

IF block-fits ('mem-pool-info p) ('block2 t) ('lsz t) THEN

    (* sys-dlist-append(&p->levels[l + 1].free-list, block2); *)
    'mem-pool-info := 'mem-pool-info (p :=
    append-free-list ('mem-pool-info p) (nat ('from-l t + 1)) ('block2
t) )

    FI
    ROF

END);
(t ► 'from-l := 'from-l(t := 'from-l t + 1))
OD) satp [mp-alloc-precond2-1-1-loopinv t p sz timeout, Mem-pool-alloc-rely t,
Mem-pool-alloc-guar t,
    mp-alloc-precond2-1-2 t p sz timeout]
apply(rule While)
using mp-alloc-precond2-1-1-loopinv-stb apply simp
apply(rule Int-greatest) apply(rule Int-greatest) apply(rule Int-greatest)
apply(rule Int-greatest) apply(rule Int-greatest) apply(rule Int-greatest)
apply(rule Int-greatest) apply(rule Int-greatest)
apply auto[1] apply auto[1] apply auto[1] apply auto[1] apply auto[1]
apply auto[1] apply auto[1] apply auto[1] apply auto[1] apply clarify apply auto[1] apply
auto[1]
apply(rule subst[where t=⌈'from-l t ≤ 'alloc-l t ∧ 'allocating-node t =
    Some (pool = p, level = nat ('from-l t), block = block-num ('mem-pool-info
p) ('blk t) ('lsizes t ! nat ('from-l t)),
    data = 'blk t)⌋ and s=⌈'from-l t ≤ 'alloc-l t⌋ ∩ ⌈'allocating-node t =
    Some (pool = p, level = nat ('from-l t), block = block-num ('mem-pool-info
p) ('blk t) ('lsizes t ! nat ('from-l t)),
    data = 'blk t)⌋]) apply auto[1]

using mp-alloc-precond2-1-2-stb apply simp

```

```

apply(rule Seq[where mid=mp-alloc-precond2-1-1-loopinv-1 t p sz timeout])

apply(unfold stm-def)[1]
apply(rule Await)
  using mp-alloc-precond2-1-1-loopinv-0-stb apply auto[1]
  using mp-alloc-precond2-1-1-loopinv-1-stb apply simp
  apply clarify
  apply(rule Await)
    using stable-id2 apply fast using stable-id2 apply fast
    apply clarify
    apply(case-tac V = Va) prefer 2 apply simp using Emptyprecond apply
auto[1]
    apply simp
    apply(case-tac mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap \llbracket 'cur =$ 
Some t  $\rrbracket \cap \{Va\} = \{\}$ )
      using Emptyprecond[of -  $\{(s, t). s = t\}$  UNIV ] apply auto[1]
      apply(subgoal-tac mp-alloc-precond2-1-1-loopinv-0 t p sz timeout  $\cap \llbracket 'cur$ 
= Some t  $\rrbracket \cap \{Va\} = \{Va\}$ )
        prefer 2 using int1-eq[where P=mp-alloc-precond2-1-1-loopinv-0 t p sz
timeout  $\cap \llbracket 'cur =$  Some t  $\rrbracket$ ] apply meson
        using mp-alloc-stm4-lm1[of t p timeout sz] apply auto[1]

apply(unfold stm-def)
apply(rule Await)
  using mp-alloc-precond2-1-1-loopinv-1-stb apply simp
  using mp-alloc-precond2-1-1-loopinv-stb apply auto[1]
  apply clarify
  apply(rule Basic)
    apply(case-tac mp-alloc-precond2-1-1-loopinv-1 t p sz timeout  $\cap \llbracket 'cur =$ 
Some t  $\rrbracket \cap \{V\} = \{\}$ )
      apply auto[1]
      apply(subgoal-tac mp-alloc-precond2-1-1-loopinv-1 t p sz timeout  $\cap \llbracket 'cur$ 
= Some t  $\rrbracket \cap \{V\} = \{V\}$ )
        prefer 2 using int1-eq[where P=mp-alloc-precond2-1-1-loopinv-1 t p sz
timeout  $\cap \llbracket 'cur =$  Some t  $\rrbracket$ ] apply meson
        apply simp
        apply(rule conjI) apply(simp add:Mem-pool-alloc-guar-def) apply(rule
disjI1)
          apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
          apply(rule conjI) apply(subgoal-tac (V, V( $\llbracket from-l := (from-l V)(t :=$ 
from-l V t + 1  $\rrbracket$ )) $\in$ lvars-nochange1-4all)
            using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
              apply(simp add:lvars-nochange-def)
              apply(rule conjI) apply(subgoal-tac (V, V( $\llbracket from-l := (from-l V)(t :=$ 
from-l V t + 1  $\rrbracket$ )) $\in$ lvars-nochange1-4all)
                using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)

```

```

apply(rule conjI) apply auto[1]
apply(rule conjI) apply (metis less-minus-one-simps(1))
apply(rule conjI) apply smt
apply (metis (no-types, hide-lams) Mem-block.simps(2) Mem-block.simps(3)
Mem-block.simps(4) option.sel)

```

```

apply simp using stable-id2 apply blast using stable-id2 apply blast

```

```

apply(simp add:Mem-pool-alloc-guar-def)
done

```

## 10.8 stm5

**lemma** mp-alloc-stm5-lm-1-inv-mempool-info:

```

free-l V t ≤ alloc-l V t ⇒
  alloc-l V t < int (n-levels (mem-pool-info V p)) ⇒
  p ∈ mem-pools V ⇒
  inv-mempool-info V ⇒
  ¬ free-l V t < OK ⇒
  NULL < blk V t ⇒
  inv-mempool-info
  (V (mem-pool-info := (mem-pool-info V)
    (p := mem-pool-info V p
      (levels := levels (mem-pool-info V p)
        [nat (alloc-l V t) := (levels (mem-pool-info V p) ! nat (alloc-l V t))
          (bits := bits (levels (mem-pool-info V p) ! nat (alloc-l V t))
            [(blk V t - buf (mem-pool-info V p)) div lsize V t ! nat (alloc-l
V t) := ALLOCATED])))),
    allocating-node := (allocating-node V)(t := None)))
apply(simp add:inv-mempool-info-def)
apply(rule conjI) apply metis
apply(rule conjI) apply metis
apply(rule conjI) apply metis
apply(rule conjI) apply metis
apply clarify apply(rename-tac ii) apply(subgoal-tac length (bits (levels (mem-pool-info
V p)
  [nat (alloc-l V t) := (levels (mem-pool-info V p) ! nat (alloc-l V t))
    (bits := bits (levels (mem-pool-info V p) ! nat (alloc-l V t))
      [(blk V t - buf (mem-pool-info V p)) div lsize V t ! nat (alloc-l
V t) := ALLOCATED])) !
    ii)=length (bits (levels (mem-pool-info V p) ! ii)))
  prefer 2 apply(case-tac ii = nat (alloc-l V t)) apply force apply force
apply metis
done

```

**lemma** mp-alloc-stm5-lm-1-inv-bitmap-h1:

```

allocating-node V t =
  Some (pool = p, level = nat (alloc-l V t), block = (blk V t - buf (mem-pool-info

```

$V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t), \text{ data} = \text{blk } V\ t) \implies$   
 $\forall t\ n. \text{ allocating-node } V\ t = \text{Some } n \longrightarrow \text{get-bit-s } V\ (\text{pool } n) (\text{level } n) (\text{block } n)$   
 $= \text{ALLOCATING} \implies$   
 $\text{get-bit-s } V\ p\ (\text{nat } (\text{alloc-l } V\ t)) ((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes}$   
 $V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t)) = \text{ALLOCATING}$   
**by** *fastforce*

**lemma** *mp-alloc-stm5-lm-1-inv-bitmap-freelist*:

$\text{allocating-node } V\ t =$   
 $\text{Some } (\downarrow \text{pool} = p, \text{level} = \text{nat } (\text{alloc-l } V\ t), \text{block} = (\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t)) \implies$   
 $\text{alloc-l } V\ t < \text{int } (n\text{-levels } (\text{mem-pool-info } V\ p)) \implies$   
 $p \in \text{mem-pools } V \implies$   
 $\text{inv-mempool-info } V \wedge \text{inv-aux-vars } V \wedge \text{inv-bitmap-freelist } V \implies$   
 $\text{inv-bitmap-freelist}$   
 $(V(\downarrow \text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V\ p$   
 $(\downarrow \text{levels} := \text{levels } (\text{mem-pool-info } V\ p)$   
 $[\text{nat } (\text{alloc-l } V\ t) := (\text{levels } (\text{mem-pool-info } V\ p) \ ! \ \text{nat } (\text{alloc-l } V\ t))$   
 $(\downarrow \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V\ p) \ ! \ \text{nat } (\text{alloc-l } V\ t))$   
 $[(\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t) := \text{ALLOCATED}])])])$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*rule subst[where s=inv-bitmap-freelist*  
 $(V(\downarrow \text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V\ p$   
 $(\downarrow \text{levels} := \text{levels } (\text{mem-pool-info } V\ p)$   
 $[\text{nat } (\text{alloc-l } V\ t) := (\text{levels } (\text{mem-pool-info } V\ p) \ ! \ \text{nat } (\text{alloc-l } V\ t))$   
 $(\downarrow \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V\ p) \ ! \ \text{nat } (\text{alloc-l } V\ t))$   
 $[(\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t) := \text{ALLOCATED}])])])$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*simp add:inv-bitmap-freelist-def*)  
**apply**(*rule subst[where s=inv-bitmap-freelist (set-bit-s V p (nat (alloc-l V t))*  
 $((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t)) \text{ ALLO-}$   
 $\text{CATED}])$   
**apply**(*unfold set-bit-s-def set-bit-def*)[1] **apply** *blast*  
**apply**(*subgoal-tac get-bit-s V p (nat (alloc-l V t))*  
 $((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t)) = \text{ALLOCATING})$  **prefer** 2  
**apply**(*subgoal-tac  $\forall t\ n. \text{ allocating-node } V\ t = \text{Some } n \longrightarrow \text{get-bit-s } V\ (\text{pool } n)$*   
 $(\text{level } n) (\text{block } n) = \text{ALLOCATING})$  **prefer** 2  
**apply**(*simp add:inv-aux-vars-def Let-def*)  
**using** *mp-alloc-stm5-lm-1-inv-bitmap-h1* **apply** *blast*

**using** *inv-bitmap-freelist-presv-setbit-notfree*[*of p V ALLOCATED nat (alloc-l V t)*  
 $(\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } l\text{sizes } V\ t \ ! \ \text{nat } (\text{alloc-l } V\ t)]$   
**apply** *fastforce*

done

**lemma** *mp-alloc-stm5-lm-1-inv-bitmap*:

*allocating-node*  $V\ t =$   
 $\text{Some } (\downarrow \text{pool} = p, \text{level} = \text{nat } (\text{alloc-l } V\ t), \text{block} = (\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t), \text{data} = \text{blk } V\ t) \implies$   
 $p \in \text{mem-pools } V \implies$   
 $\text{inv-bitmap } V \wedge \text{inv-aux-vars } V \implies$   
 $\text{inv-bitmap}$   
 $(V(\downarrow \text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V\ p$   
 $(\downarrow \text{levels} := \text{levels } (\text{mem-pool-info } V\ p)$   
 $[\text{nat } (\text{alloc-l } V\ t) := (\text{levels } (\text{mem-pool-info } V\ p) ! \text{nat } (\text{alloc-l } V\ t))$   
 $(\downarrow \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V\ p) ! \text{nat } (\text{alloc-l } V\ t))$   
 $[(\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t) := \text{ALLOCATED}])])])$ ,  
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$ )  
**apply**(*rule subst*[**where**  $s = \text{inv-bitmap}$   
 $(V(\downarrow \text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V\ p$   
 $(\downarrow \text{levels} := \text{levels } (\text{mem-pool-info } V\ p)$   
 $[\text{nat } (\text{alloc-l } V\ t) := (\text{levels } (\text{mem-pool-info } V\ p) ! \text{nat } (\text{alloc-l } V\ t))$   
 $(\downarrow \text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V\ p) ! \text{nat } (\text{alloc-l } V\ t))$   
 $[(\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t) := \text{ALLOCATED}])])])])$ )  
**apply**(*simp add:inv-bitmap-def*)  
**apply**(*rule subst*[**where**  $s = \text{inv-bitmap } (\text{set-bit-s } V\ p\ (\text{nat } (\text{alloc-l } V\ t))$   
 $((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t)) \text{ ALLOCATED}])$ )  
**apply**(*unfold set-bit-s-def set-bit-def*)[1] **apply** *blast*  
**apply**(*subgoal-tac get-bit-s*  $V\ p\ (\text{nat } (\text{alloc-l } V\ t))$   
 $((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t)) = \text{ALLOCATING})$  **prefer** 2  
**apply**(*subgoal-tac*  $\forall t\ n. \text{allocating-node } V\ t = \text{Some } n \longrightarrow \text{get-bit-s } V\ (\text{pool } n)$   
 $(\text{level } n) (\text{block } n) = \text{ALLOCATING})$  **prefer** 2  
**apply**(*simp add:inv-aux-vars-def Let-def*)  
**using** *mp-alloc-stm5-lm-1-inv-bitmap-h1* **apply** *blast*  
**using** *inv-bitmap-presv-setbit*[*of*  $V\ p\ \text{nat } (\text{alloc-l } V\ t) (\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t)$   
 $\text{ALLOCATED set-bit-s } V\ p\ (\text{nat } (\text{alloc-l } V\ t)) ((\text{blk } V\ t - \text{buf } (\text{mem-pool-info } V\ p)) \text{ div } \text{lsizes } V\ t ! \text{nat } (\text{alloc-l } V\ t)) \text{ ALLOCATED}$ ]  
**apply** *blast*  
done

**lemma** *mp-alloc-stm5-lm-1-inv-aux-vars*:

$(\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l } V \ t) < n\text{-max}$   
 $(\text{mem-pool-info } V \ p) * 4 \wedge \text{nat } (\text{alloc-l } V \ t) \implies$   
 $\text{blk } V \ t =$   
 $\text{buf } (\text{mem-pool-info } V \ p) +$   
 $(\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l } V \ t) * (\text{max-sz}$   
 $(\text{mem-pool-info } V \ p) \text{ div } 4 \wedge \text{nat } (\text{alloc-l } V \ t)) \implies$   
 $0 < \text{blk } V \ t \implies$   
 $\text{allocating-node } V \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat } (\text{alloc-l } V \ t), \text{block} = (\text{blk } V \ t - \text{buf } (\text{mem-pool-info}$   
 $V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l } V \ t),$   
 $\text{data} = \text{blk } V \ t) \implies$   
 $\text{alloc-l } V \ t < \text{int } (n\text{-levels } (\text{mem-pool-info } V \ p)) \implies$   
 $p \in \text{mem-pools } V \implies$   
 $\text{inv-mempool-info } V \wedge \text{inv-aux-vars } V \implies$   
 $\forall ii < \text{length } (\text{lsizes } V \ t). \text{lsizes } V \ t \ ! \ ii = \text{ALIGN4 } (\text{max-sz } (\text{mem-pool-info } V \ p))$   
 $\text{div } 4 \wedge ii \implies$   
 $\text{inv-aux-vars}$   
 $(V(\text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V \ p$   
 $(\text{levels} := \text{levels } (\text{mem-pool-info } V \ p)$   
 $[\text{nat } (\text{alloc-l } V \ t) := (\text{levels } (\text{mem-pool-info } V \ p) \ ! \ \text{nat } (\text{alloc-l } V \ t))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \ p) \ ! \ \text{nat } (\text{alloc-l } V \ t))$   
 $[(\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l}$   
 $V \ t) := \text{ALLOCATED}]))))$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*unfold inv-aux-vars-def*)  
**apply**(*subgoal-tac get-bit-s*  $V \ p \ (\text{nat } (\text{alloc-l } V \ t))$   
 $((\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l}$   
 $V \ t)) = \text{ALLOCATING})$  **prefer** 2  
**using** *mp-alloc-stm5-lm-1-inv-bitmap-h1* **apply** *presburger*  
  
**apply**(*subgoal-tac mem-block-addr-valid*  $V \ (\text{pool} = p, \text{level} = \text{nat } (\text{alloc-l } V \ t),$   
 $\text{block} = (\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat}$   
 $(\text{alloc-l } V \ t), \text{data} = \text{blk } V \ t))$ ) **prefer** 2  
**apply**(*simp add:mem-block-addr-valid-def*)  
  
**apply**(*rule conjI*)  
**apply** *clarify*  
**apply**(*subgoal-tac freeing-node*  $V \ ta = \text{Some } n$ ) **prefer** 2 **apply** *force*  
  
**apply**(*subgoal-tac*  $\neg(\text{pool } n = p \wedge \text{level } n = \text{nat } (\text{alloc-l } V \ t)$   
 $\wedge \text{block } n = (\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l}$   
 $V \ t))$ )  
**prefer** 2 **apply** *metis*  
**apply**(*subgoal-tac get-bit-s*  $V \ (\text{pool } n) (\text{level } n) (\text{block } n) = \text{FREEING})$  **prefer** 2  
**apply** *presburger*  
**apply**(*subgoal-tac get-bit-s*  $V \ (\text{pool } n) (\text{level } n) (\text{block } n) = \text{get-bit-s}$   
 $(V(\text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V \ p$

$\llbracket levels := levels \ (mem\text{-}pool\text{-}info \ V \ p)$   
 $\llbracket nat \ (alloc\text{-}l \ V \ t) := (levels \ (mem\text{-}pool\text{-}info \ V \ p) ! nat \ (alloc\text{-}l$   
 $V \ t))$   
 $\llbracket bits := bits \ (levels \ (mem\text{-}pool\text{-}info \ V \ p) ! nat \ (alloc\text{-}l \ V \ t))$   
 $\llbracket (blk \ V \ t - buf \ (mem\text{-}pool\text{-}info \ V \ p)) \div lsizes \ V \ t ! nat$   
 $(alloc\text{-}l \ V \ t) := ALLOCATED \rrbracket \rrbracket \rrbracket$ ,  
 $allocating\text{-}node := (allocating\text{-}node \ V)(t := None) \rrbracket \rrbracket (pool \ n) \ (level$   
 $n) \ (block \ n))$  **prefer** 2  
**apply**(*case-tac* *pool*  $n \neq p$ ) **apply** *force*  
**apply**(*case-tac* *level*  $n \neq nat \ (alloc\text{-}l \ V \ t)$ ) **apply** *force*  
**apply**(*case-tac* *block*  $n \neq (blk \ V \ t - buf \ (mem\text{-}pool\text{-}info \ V \ p)) \div lsizes \ V \ t !$   
 $nat \ (alloc\text{-}l \ V \ t)$ )  
**apply**(*case-tac* *level*  $n \geq length \ (levels \ (mem\text{-}pool\text{-}info \ V \ (pool \ n)))$ ) **apply**  
*fastforce*  
**apply** *force* **apply** *blast*  
**apply** *argo*

**apply**(*rule* *conjI*)  
**apply** *clarify*  
**apply**(*subgoal-tac*  $\exists ta. \text{freeing-node} \ V \ ta = Some \ n$ ) **prefer** 2  
**apply**(*subgoal-tac* *get-bit-s*  $V \ (pool \ n) \ (level \ n) \ (block \ n) = FREEING$ ) **prefer**  
2  
**apply**(*case-tac* *pool*  $n \neq p$ ) **apply** *force*  
**apply**(*case-tac* *level*  $n \neq nat \ (alloc\text{-}l \ V \ t)$ ) **apply** *force*  
**apply**(*case-tac* *block*  $n \neq (blk \ V \ t - buf \ (mem\text{-}pool\text{-}info \ V \ p)) \div lsizes \ V \ t !$   
 $nat \ (alloc\text{-}l \ V \ t)$ )  
**apply**(*case-tac* *level*  $n \geq length \ (levels \ (mem\text{-}pool\text{-}info \ V \ (pool \ n)))$ )  
**apply** *fastforce* **apply** *force*  
**apply**(*case-tac* *level*  $n \geq length \ (levels \ (mem\text{-}pool\text{-}info \ V \ (pool \ n)))$ )  
**apply** *fastforce*  
**apply**(*case-tac* *block*  $n \geq length \ (bits \ (levels \ (mem\text{-}pool\text{-}info \ V \ p) ! nat \ (alloc\text{-}l$   
 $V \ t)))$ )  
**apply** *fastforce* **apply** *fastforce*  
**apply**(*subgoal-tac* *mem-block-addr-valid*  $V \ n$ ) **prefer** 2  
**apply**(*simp* *add:mem-block-addr-valid-def*)  
**apply** *blast*  
**apply** *force*

**apply**(*rule* *conjI*)  
**apply** *clarify*  
**apply**(*subgoal-tac*  $t \neq ta$ ) **prefer** 2 **apply** *fastforce*  
**apply**(*subgoal-tac* *allocating-node*  $V \ ta = Some \ n$ ) **prefer** 2 **apply** *force*

**apply**(*subgoal-tac*  $\neg(pool \ n = p \wedge level \ n = nat \ (alloc\text{-}l \ V \ t)$   
 $\wedge block \ n = (blk \ V \ t - buf \ (mem\text{-}pool\text{-}info \ V \ p)) \div lsizes \ V \ t ! nat \ (alloc\text{-}l$   
 $V \ t)))$   
**prefer** 2 **apply** (*metis* *Mem-block.select-convs*(1) *Mem-block.select-convs*(2) *Mem-block.select-convs*(3))  
**apply**(*subgoal-tac* *get-bit-s*  $V \ (pool \ n) \ (level \ n) \ (block \ n) = ALLOCATING$ ) **pre-**  
**fer** 2 **apply** *presburger*

**apply**(*subgoal-tac* *get-bit-s* *V* (*pool* *n*) (*level* *n*) (*block* *n*) = *get-bit-s*  
 ( *V* (*mem-pool-info* := (*mem-pool-info* *V*)

(*p* := *mem-pool-info* *V* *p*  
 (*levels* := *levels* (*mem-pool-info* *V* *p*)

[*nat* (*alloc-l* *V* *t*) := (*levels* (*mem-pool-info* *V* *p*) ! *nat* (*alloc-l*  
*V* *t*))

(*bits* := *bits* (*levels* (*mem-pool-info* *V* *p*) ! *nat* (*alloc-l* *V* *t*))

[(*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* ! *nat*  
 (*alloc-l* *V* *t*) := *ALLOCATED*]))]),

*allocating-node* := (*allocating-node* *V*)(*t* := *None*)) (*pool* *n*) (*level*  
*n*) (*block* *n*)) **prefer** 2

**apply**(*case-tac* *pool* *n* ≠ *p*) **apply** *force*  
**apply**(*case-tac* *level* *n* ≠ *nat* (*alloc-l* *V* *t*)) **apply** *force*  
**apply**(*case-tac* *block* *n* ≠ (*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* !

*nat* (*alloc-l* *V* *t*))

**apply**(*case-tac* *level* *n* ≥ *length* (*levels* (*mem-pool-info* *V* (*pool* *n*)))) **apply**  
*fastforce*

**apply** *force* **apply** *blast*

**apply** *argo*

**apply**(*rule* *conjI*)

**apply** *clarify*

**apply**(*subgoal-tac* *nat* (*alloc-l* *V* *t*) < *length* (*levels* (*mem-pool-info* *V* *p*))) **prefer**  
2

**apply**(*simp* *add:inv-mempool-info-def* *Let-def*)

**apply** (*metis* *int-nat-eq* *of-nat-0-less-iff* *of-nat-less-imp-less*)

**apply**(*subgoal-tac* (*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* ! *nat* (*alloc-l*  
*V* *t*)

< *length* (*bits* (*levels* (*mem-pool-info* *V* *p*) ! *nat* (*alloc-l* *V* *t*))))

**prefer** 2

**apply**(*simp* *add:inv-mempool-info-def* *Let-def*)

**apply**(*subgoal-tac* ¬(*pool* *n* = *p* ∧ *level* *n* = *nat* (*alloc-l* *V* *t*)

∧ *block* *n* = (*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* ! *nat* (*alloc-l*  
*V* *t*)))

**prefer** 2

**apply**(*case-tac* *pool* *n* ≠ *p*) **apply** *fastforce*  
**apply**(*case-tac* *level* *n* ≠ *nat* (*alloc-l* *V* *t*)) **apply** *fastforce*  
**apply**(*case-tac* *block* *n* ≠ (*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* !

*nat* (*alloc-l* *V* *t*)) **apply** *fastforce*

**apply** *simp*

**apply**(*subgoal-tac* ∃ *ta*. *ta* ≠ *t* ∧ *allocating-node* *V* *ta* = *Some* *n*) **prefer** 2

**apply**(*subgoal-tac* *get-bit-s* *V* (*pool* *n*) (*level* *n*) (*block* *n*) = *ALLOCATING*)

**prefer** 2

**apply**(*case-tac* *pool* *n* ≠ *p*) **apply** *force*  
**apply**(*case-tac* *level* *n* ≠ *nat* (*alloc-l* *V* *t*)) **apply** *force*  
**apply**(*case-tac* *block* *n* ≠ (*blk* *V* *t* - *buf* (*mem-pool-info* *V* *p*)) *div* *lsizes* *V* *t* !

*nat* (*alloc-l* *V* *t*))

**apply**(*case-tac* *level* *n* ≥ *length* (*levels* (*mem-pool-info* *V* (*pool* *n*))))



```

    apply fastforce apply force
  apply(case-tac level n ≥ length (levels (mem-pool-info V (pool n))))
    apply fastforce
  apply(case-tac block n ≥ length (bits (levels (mem-pool-info V p) ! nat (alloc-l
V t))))
    apply fastforce apply fastforce
  apply(subgoal-tac mem-block-addr-valid V n) prefer 2
  apply(simp add:mem-block-addr-valid-def)
  apply (metis Mem-block.select-convs(1) Mem-block.select-convs(2) Mem-block.select-convs(3)
option.sel)

apply auto[1]

apply(rule conjI)
apply clarify
apply auto[1]

apply(rule conjI)
apply clarify
apply(subgoal-tac allocating-node V t1 = Some n1) prefer 2
  apply(case-tac t = t1) apply force apply force
apply(subgoal-tac allocating-node V t2 = Some n1) prefer 2
  apply(case-tac t = t2) apply force apply force
apply metis

apply clarify
apply(subgoal-tac allocating-node V t1 = Some n1) prefer 2
  apply(case-tac t = t1) apply force apply force
apply(subgoal-tac freeing-node V t2 = Some n1) prefer 2 apply force
apply metis
done

lemma mp-alloc-stm5-lm-1-inv-bitmap0:
p ∈ mem-pools V ⇒
  inv-mempool-info V ∧ inv-bitmap0 V ⇒
  inv-bitmap0
  (V (mem-pool-info := (mem-pool-info V)
    (p := mem-pool-info V p
      (levels := levels (mem-pool-info V p)
        [nat (alloc-l V t) := (levels (mem-pool-info V p) ! nat (alloc-l V t))
        (bits := bits (levels (mem-pool-info V p) ! nat (alloc-l V t))
        [(blk V t - buf (mem-pool-info V p)) div lsize V t ! nat (alloc-l
V t) := ALLOCATED]]]]),
    allocating-node := (allocating-node V)(t := None)))
  apply(simp add:inv-bitmap0-def Let-def)
  apply clarsimp
  apply(subgoal-tac length (levels (mem-pool-info V p)) > 0) prefer 2
    apply(simp add:inv-mempool-info-def Let-def) apply fastforce

```

**apply**(*case-tac* *nat* (*alloc-l* *V t*) = 0)  
**apply**(*case-tac* *i* = (*blk* *V t* - *buf* (*mem-pool-info* *V p*)) *div* *lsizes* *V t* ! *nat* (*alloc-l* *V t*))  
**apply** *fastforce* **apply** *force*  
**by** *fastforce*

**lemma** *mp-alloc-stm5-lm-1-inv-bitmapn*:  
 $p \in \text{mem-pools } V \implies$   
 $\text{inv-mempool-info } V \wedge \text{inv-bitmapn } V \implies$   
 $\text{inv-bitmapn}$   
 $(V(\text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V p$   
 $(\text{levels} := \text{levels } (\text{mem-pool-info } V p)$   
 $[\text{nat } (\text{alloc-l } V t) := (\text{levels } (\text{mem-pool-info } V p) ! \text{nat } (\text{alloc-l } V t))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V p) ! \text{nat } (\text{alloc-l } V t))$   
 $[(\text{blk } V t - \text{buf } (\text{mem-pool-info } V p)) \text{div } \text{lsizes } V t ! \text{nat } (\text{alloc-l}$   
 $V t) := \text{ALLOCATED}]))],$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*simp* *add:inv-bitmapn-def* *Let-def*)  
**apply** *clarsimp*  
**apply**(*subgoal-tac* *length* (*levels* (*mem-pool-info* *V p*)) > 0) **prefer** 2  
**apply**(*simp* *add:inv-mempool-info-def* *Let-def*) **apply** *fastforce*

**apply**(*case-tac* *nat* (*alloc-l* *V t*) = *length* (*levels* (*mem-pool-info* *V p*)) - *Suc* 0)  
**apply**(*case-tac* *i* = (*blk* *V t* - *buf* (*mem-pool-info* *V p*)) *div* *lsizes* *V t* ! *nat* (*alloc-l* *V t*))  
**apply** *fastforce* **apply** *force*  
**by** *fastforce*

**lemma** *mp-alloc-stm5-lm-1-inv-bitmap-not4free*:  
 $(\text{blk } V t - \text{buf } (\text{mem-pool-info } V p)) \text{div } \text{lsizes } V t ! \text{nat } (\text{alloc-l } V t) < n\text{-max}$   
 $(\text{mem-pool-info } V p) * 4 ^ \text{nat } (\text{alloc-l } V t) \implies$   
 $\text{alloc-l } V t < \text{int } (n\text{-levels } (\text{mem-pool-info } V p)) \implies$   
 $p \in \text{mem-pools } V \implies$   
 $\text{inv-mempool-info } V \wedge \text{inv-bitmap-not4free } V \implies$   
 $\text{inv-bitmap-not4free}$   
 $(V(\text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V p$   
 $(\text{levels} := \text{levels } (\text{mem-pool-info } V p)$   
 $[\text{nat } (\text{alloc-l } V t) := (\text{levels } (\text{mem-pool-info } V p) ! \text{nat } (\text{alloc-l } V t))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V p) ! \text{nat } (\text{alloc-l } V t))$   
 $[(\text{blk } V t - \text{buf } (\text{mem-pool-info } V p)) \text{div } \text{lsizes } V t ! \text{nat } (\text{alloc-l}$   
 $V t) := \text{ALLOCATED}]))],$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*subgoal-tac* *length* (*levels* (*mem-pool-info* *V p*)) > 0) **prefer** 2  
**apply**(*simp* *add:inv-mempool-info-def* *Let-def*) **apply** *fastforce*  
**apply**(*subgoal-tac* *nat* (*alloc-l* *V t*) < *length* (*levels* (*mem-pool-info* *V p*))) **prefer** 2  
2

**apply**(*simp add:inv-mempool-info-def Let-def*)  
**apply** (*metis int-nat-eq of-nat-0-less-iff of-nat-less-imp-less*)

**apply**(*simp add:inv-bitmap-not4free-def partner-bits-def Let-def*)  
**apply** *clarsimp*

**apply**(*subgoal-tac (blk V t - buf (mem-pool-info V p)) div lsizes V t ! nat (alloc-l V t)*  
 $< \text{length (bits (levels (mem-pool-info V p) ! nat (alloc-l V t)))}$ )

**prefer** 2  
**apply**(*simp add:inv-mempool-info-def Let-def*)

**apply**(*case-tac nat (alloc-l V t) < length (levels (mem-pool-info V p))*)  
**apply**(*case-tac i = nat (alloc-l V t)*)  
**apply**(*case-tac (blk V t - buf (mem-pool-info V p)) div lsizes V t ! nat (alloc-l V t) = j div 4 \* 4*)  
**apply** *fastforce*  
**apply**(*case-tac (blk V t - buf (mem-pool-info V p)) div lsizes V t ! nat (alloc-l V t) = Suc (j div 4 \* 4)*)  
**apply** *fastforce*  
**apply**(*case-tac (blk V t - buf (mem-pool-info V p)) div lsizes V t ! nat (alloc-l V t) = Suc (Suc (j div 4 \* 4))*)  
**apply** *fastforce*  
**apply**(*case-tac (blk V t - buf (mem-pool-info V p)) div lsizes V t ! nat (alloc-l V t) = j div 4 \* 4 + 3*)  
**apply** *fastforce*  
**apply** *fastforce*  
**apply** *force*  
**by** *blast*

**lemma** *mp-alloc-stm5-lm-1-inv*:  
 $(\text{blk } V \ t - \text{buf (mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t ! \text{ nat (alloc-l } V \ t) < n\text{-max}$   
 $(\text{mem-pool-info } V \ p) * 4 \wedge \text{ nat (alloc-l } V \ t) \implies$   
 $\text{blk } V \ t = \text{buf (mem-pool-info } V \ p) + (\text{blk } V \ t - \text{buf (mem-pool-info } V \ p)) \text{ div}$   
 $\text{lsizes } V \ t ! \text{ nat (alloc-l } V \ t) *$   
 $(\text{max-sz (mem-pool-info } V \ p) \text{ div } 4 \wedge \text{ nat (alloc-l } V \ t)) \implies$   
 $\text{allocating-node } V \ t =$   
 $\text{Some } (\text{pool} = p, \text{level} = \text{nat (alloc-l } V \ t), \text{block} = (\text{blk } V \ t - \text{buf (mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t ! \text{ nat (alloc-l } V \ t),$   
 $\text{data} = \text{blk } V \ t) \implies$   
 $\text{free-l } V \ t \leq \text{alloc-l } V \ t \implies$   
 $\text{alloc-l } V \ t < \text{int (n-levels (mem-pool-info } V \ p)) \implies$   
 $\text{length (lsizes } V \ t) \leq \text{n-levels (mem-pool-info } V \ p) \implies$   
 $p \in \text{mem-pools } V \implies$   
 $\text{inv } V \implies$   
 $\forall ii < \text{length (lsizes } V \ t). \text{lsizes } V \ t ! ii = \text{ALIGN4 (max-sz (mem-pool-info } V \ p)) \text{ div } 4 \wedge ii \implies$   
 $\neg \text{free-l } V \ t < \text{OK} \implies$   
 $\text{NULL} < \text{blk } V \ t \implies$

$\forall ii \leq \text{nat } (\text{alloc-l } V \ t). \ sz \leq \text{lsizes } V \ t \ ! \ ii \implies$   
 $\text{alloc-l } V \ t = \text{int } (\text{length } (\text{lsizes } V \ t)) - 1 \wedge \text{length } (\text{lsizes } V \ t) = n\text{-levels}$   
 $(\text{mem-pool-info } V \ p) \vee$   
 $\text{alloc-l } V \ t = \text{int } (\text{length } (\text{lsizes } V \ t)) - 2 \wedge \text{lsizes } V \ t \ ! \ \text{nat } (\text{int } (\text{length } (\text{lsizes}$   
 $V \ t)) - 1) < sz \implies$   
 $\text{inv } (V \ ( \text{mem-pool-info} := (\text{mem-pool-info } V)$   
 $(p := \text{mem-pool-info } V \ p$   
 $(\text{levels} := \text{levels } (\text{mem-pool-info } V \ p)$   
 $[\text{nat } (\text{alloc-l } V \ t) := (\text{levels } (\text{mem-pool-info } V \ p) \ ! \ \text{nat } (\text{alloc-l } V \ t))$   
 $(\text{bits} := \text{bits } (\text{levels } (\text{mem-pool-info } V \ p) \ ! \ \text{nat } (\text{alloc-l } V \ t))$   
 $[(\text{blk } V \ t - \text{buf } (\text{mem-pool-info } V \ p)) \text{ div } \text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l}$   
 $V \ t) := \text{ALLOCATED}])])],$   
 $\text{allocating-node} := (\text{allocating-node } V)(t := \text{None}))$   
**apply**(*simp add:inv-def*)  
**apply**(*rule conjI*) **apply**(*simp add:inv-cur-def*)  
**apply**(*rule conjI*) **apply**(*simp add:inv-thd-waitq-def*)  
**apply**(*rule conjI*) **apply** *metis* **apply** *metis*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-mempool-info* **apply** *blast*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-bitmap-freelist* **apply** *blast*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-bitmap* **apply** *blast*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-aux-vars* **apply** *blast*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-bitmap0* **apply** *blast*  
**apply**(*rule conjI*) **using** *mp-alloc-stm5-lm-1-inv-bitmapn* **apply** *blast*  
**using** *mp-alloc-stm5-lm-1-inv-bitmap-not4free* **apply** *blast*  
**done**

**term** *mp-alloc-precond2-1-2* *t p sz timeout*  $\cap \{ \text{'cur} = \text{Some } t \}$

**lemma** *mp-alloc-stm5-lm-1*:

$\text{mp-alloc-precond2-1-2 } t \ p \ sz \ \text{timeout} \cap \{ \text{'cur} = \text{Some } t \} \cap \{ V \} \neq \{ \} \implies$   
 $\Gamma \vdash_I \text{Some } ( \text{'mem-pool-info} :=$   
 $\text{set-bit-alloc } \text{'mem-pool-info } p \ (\text{nat } ( \text{'alloc-l } t)) \ (\text{block-num } ( \text{'mem-pool-info}$   
 $p) \ (\text{'blk } t) \ (\text{'lsizes } t \ ! \ \text{nat } ( \text{'alloc-l } t))) ;;$   
 $\text{'allocating-node} := \text{'allocating-node } (t := \text{None}) )$   
 $\text{sat}_p [\text{mp-alloc-precond2-1-2 } t \ p \ sz \ \text{timeout} \cap \{ \text{'cur} = \text{Some } t \} \cap \{ V \},$   
 $\{ (s, t). s = t \}, \text{UNIV}, \{ \text{'(Pair } V) \in \text{Mem-pool-alloc-guar } t \} \cap \text{mp-alloc-precond2-1-3}$   
 $t \ p \ sz \ \text{timeout}]$   
**apply**(*subgoal-tac mp-alloc-precond2-1-2* *t p sz timeout*  $\cap \{ \text{'cur} = \text{Some } t \} \cap$   
 $\{ V \} = \{ V \}$ )  
**prefer** 2 **using** *int1-eq*[**where**  $P = \text{mp-alloc-precond2-1-2 } t \ p \ sz \ \text{timeout} \cap \{ \text{'cur}$   
 $= \text{Some } t \}$ ] **apply** *meson*

**apply** *simp*

**apply**(*rule Seq*[**where**  $\text{mid} = \{ V \ ( \text{mem-pool-info} := \text{set-bit-alloc } (\text{mem-pool-info}$   
 $V) \ p \ (\text{nat } (\text{alloc-l } V \ t))$   
 $(\text{block-num } ((\text{mem-pool-info } V) \ p) \ (\text{blk } V \ t) \ (\text{lsizes } V \ t \ ! \ \text{nat } (\text{alloc-l } V$   
 $t))) \} \}$ ])

```

apply(rule Basic)
apply simp apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(rule Basic)
apply clarsimp apply(simp add: set-bit-def block-num-def)
apply(rule conjI)
apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
apply(rule conjI)
apply(simp add:gvars-conf-stable-def gvars-conf-def) apply clarsimp
apply(case-tac i = nat (alloc-l V t)) apply(case-tac i < length (levels
(mem-pool-info V p)))
apply auto[1] apply auto[1] apply auto[1]
apply(rule conjI) using mp-alloc-stm5-lm-1-inv apply clarsimp
apply(simp add:lvars-nochange-def)
apply(rule conjI) using mp-alloc-stm5-lm-1-inv apply clarsimp
apply(case-tac alloc-l V t = int (length (lsizes V t)) - 1  $\wedge$  length (lsizes V
t) = n-levels (mem-pool-info V p))
apply simp apply simp

apply simp apply(simp add:stable-def) using stable-id2 apply metis
done

lemma mp-alloc-stm5-lm:
   $\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{'mem-pool-info} := \text{set-bit-alloc 'mem-pool-info } p \text{ (nat ('alloc-l } t))$ 
     $(\text{block-num ('mem-pool-info } p) \text{ ('blk } t) ((\text{'lsizes } t)!(\text{nat}$ 
     $(\text{'alloc-l } t)))));;$ 
     $\text{'allocating-node} := \text{'allocating-node } (t := \text{None})$ 
    )  $\text{sat}_p [\text{mp-alloc-precond2-1-2 } t \text{ } p \text{ sz timeout, Mem-pool-alloc-rely } t, \text{Mem-pool-alloc-guar}$ 
     $t,$ 
     $\text{mp-alloc-precond2-1-3 } t \text{ } p \text{ sz timeout}]$ 
apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond2-1-2-stb apply auto[1]
using mp-alloc-precond2-1-3-stb apply auto[1]

apply clarify
apply(case-tac mp-alloc-precond2-1-2 t p sz timeout  $\cap \{\text{'cur} = \text{Some } t\} \cap \{V\}$ 
=  $\{\}$ )
apply simp using Emptyprecond apply metis
using mp-alloc-stm5-lm-1[of t p timeout sz] apply clarsimp
done

term mp-alloc-precond2-1-2 t p sz timeout
term mp-alloc-precond2-1-3 t p sz timeout

```

## 10.9 stm6

**lemma** mp-alloc-stm6-lm:

```

Γ ⊢I Some (t ► 'mempoolalloc-ret := 'mempoolalloc-ret (t :=
  Some (pool = p, level = nat ('alloc-l t),
    block = block-num ('mem-pool-info p) ('blk t) (('lsizes t)!(nat ('alloc-l
t))),
  data = 'blk t )))
satp [mp-alloc-precond2-1-3 t p sz timeout, Mem-pool-alloc-rely t, Mem-pool-alloc-guar
t,
  mp-alloc-precond2-1-4 t p sz timeout]
apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond2-1-3-stb apply simp
using mp-alloc-precond2-1-4-stb apply simp
apply clarify
apply(rule Basic)
apply clarsimp
apply(rule conjI)
apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
apply(rule conjI)
apply(simp add:gvars-conf-stable-def gvars-conf-def)
apply(rule conjI)
apply(subgoal-tac (V, V (pool := mempoolalloc-ret V (t ↦
  (pool = p, level = nat (alloc-l V t), block = block-num (mem-pool-info V
p) (blk V t) (lsizes V t ! nat (alloc-l V t)),
  data = blk V t)))) ∈ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply(simp add:lvars-nochange-def)
apply(rule conjI)
apply(subgoal-tac (V, V (pool := mempoolalloc-ret V (t ↦
  (pool = p, level = nat (alloc-l V t), block = block-num (mem-pool-info V
p) (blk V t) (lsizes V t ! nat (alloc-l V t)),
  data = blk V t)))) ∈ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply(simp add:alloc-memblk-valid-def)
apply(rule conjI)
apply (smt int-nat-eq inv-maxsz-align4 less-imp-le-nat not-less-of-nat-less-iff)

apply clarify
apply(subgoal-tac ¬(alloc-l V t = int (length (lsizes V t)) - 1 ∧ length
(lsizes V t) = n-levels (mem-pool-info V p)))
prefer 2 apply auto[1]
apply simp apply (smt Suc-nat-eq-nat-zadd1 inv-maxsz-align4 lessI nat-int
power-Suc)

apply simp using stable-id2 apply metis using stable-id2 apply metis
done

```

## 10.10 stm7

**abbreviation** *mp-alloc-stm7-precond1*  $Va \equiv Va(\text{thd-state} := (\text{thd-state } Va)(\text{the } (cur \text{ } Va) := \text{BLOCKED}))$

**abbreviation** *mp-alloc-stm7-precond3*  $Va \text{ } t \text{ } p \equiv$   
 $Va(\text{mem-pool-info} := (\text{mem-pool-info } Va)(p := (\text{mem-pool-info } Va \text{ } p)(\text{wait-q} :=$   
 $(\text{wait-q } (\text{mem-pool-info } Va \text{ } p)) @ [the (cur \text{ } Va)]))$

**lemma** *mp-alloc-stm7-lm-2-1*:  $(\lambda a. \text{if } a = p \text{ then } \text{mem-pool-info } Va \text{ } p(\text{wait-q} :=$   
 $\text{wait-q } (\text{mem-pool-info } Va \text{ } p) @ [t])$   
 $\text{else } \text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED})))$   
 $a) \text{ } x$   
 $= (\lambda a. \text{if } a = p \text{ then } \text{mem-pool-info } Va \text{ } p(\text{wait-q} := \text{wait-q } (\text{mem-pool-info } Va \text{ } p) @ [t])$   
 $\text{else } \text{mem-pool-info } Va \text{ } a) \text{ } x$   
**apply** (*case-tac*  $x = p$ )  
**apply** *auto*  
**done**

**lemma** *mp-alloc-stm7-lm-2-2*:  
 $cur \text{ } Va = \text{Some } t \implies$   
 $(\lambda a. \text{if } a = p$   
 $\text{then } \text{mem-pool-info } Va \text{ } p(\text{wait-q} := \text{wait-q } (\text{mem-pool-info } Va \text{ } p) @ [the (cur$   
 $(Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))))$   
 $\text{else } \text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))) a)$   
 $=$   
 $(\lambda a. \text{if } a = p \text{ then } \text{mem-pool-info } Va \text{ } p(\text{wait-q} := \text{wait-q } (\text{mem-pool-info } Va \text{ } p)$   
 $@ [t]) \text{ else } \text{mem-pool-info } Va \text{ } a)$   
**using** *mp-alloc-stm7-lm-2-1* **by** *auto*

**lemma** *mp-alloc-stm7-lm-2*:  
 $cur \text{ } Va = \text{Some } t \implies$   
 $(\lambda a. \text{if } a = p \text{ then}$   
 $\text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))) p$   
 $(\text{wait-q} := \text{wait-q } (\text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t$   
 $:= \text{BLOCKED}))) p)$   
 $@ [the (cur (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))))]$   
 $\text{else } \text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))) a)$   
 $=$   
 $(\text{mem-pool-info } Va)(p := \text{mem-pool-info } Va \text{ } p(\text{wait-q} := \text{wait-q } (\text{mem-pool-info } Va \text{ } p) @ [t]))$   
**apply** (*rule subst* [**where**  $t = \text{mem-pool-info } (Va(\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}))) p$   
**and**  $s = \text{mem-pool-info } Va \text{ } p]$ ) **apply** *simp*  
**apply** (*simp add:fun-upd-def*)  
**using** *mp-alloc-stm7-lm-2-2* **apply** *auto*  
**done**

**lemma** *mp-alloc-stm7-swap-iffbody-inv*:

$p \in \text{mem-pools } Va \implies$

$\text{inv } Va \implies$

$\text{cur } Va = \text{Some } t \implies$

$(\text{if } ta = t \text{ then } \text{BLOCKED} \text{ else } \text{thd-state } Va \text{ } ta) = \text{READY} \implies$

$\text{inv } (\text{mp-alloc-stm7-precond3 } Va \text{ } t \text{ } p$

$\quad (\text{cur} := \text{Some } (\text{SOME } ta. \text{ } ta \neq t \wedge (ta \neq t \longrightarrow \text{thd-state } Va \text{ } ta =$

$\text{READY}))),$

$\text{thd-state} :=$

$\quad \lambda x. \text{ if } x = (\text{SOME } ta. \text{ } ta \neq t \wedge (ta \neq t \longrightarrow \text{thd-state } Va \text{ } ta =$

$\text{READY})) \text{ then } \text{RUNNING}$

$\quad \text{else } \text{thd-state}$

$\quad \quad (Va (\text{thd-state} := (\text{thd-state } Va)(t := \text{BLOCKED}),$

$\quad \quad \quad \text{mem-pool-info} := (\text{mem-pool-info } Va)$

$\quad \quad \quad (p := \text{mem-pool-info } Va \text{ } p (\text{wait-q} := \text{wait-q } (\text{mem-pool-info}$

$Va \text{ } p) @ [t])),$

$\quad \quad \quad \text{cur} := \text{Some } (\text{SOME } ta. \text{ } (ta = t \longrightarrow \text{BLOCKED} =$

$\text{READY}) \wedge (ta \neq t \longrightarrow \text{thd-state } Va \text{ } ta = \text{READY})))$

$\quad \quad \quad x))$

**apply**(*subgoal-tac* *thd-state* *Va* *t* = *RUNNING*)

**prefer** 2 **apply**(*simp* *add:inv-def* *inv-cur-def*) **apply** *auto*[1]

**apply**(*subgoal-tac* *ta*  $\neq t \wedge \text{thd-state } Va \text{ } ta = \text{READY}$ )

**prefer** 2 **apply** *auto*[1] **using** *Thread-State-Type.distinct*(3) **apply** *presburger*

**apply**(*subgoal-tac* (*SOME* *ta*. *ta*  $\neq t \wedge (ta \neq t \longrightarrow \text{thd-state } Va \text{ } ta = \text{READY})$ )

$\neq t$ )

**prefer** 2 **using** *exE-some*[**where** *P*= $\lambda tb. tb \neq t \wedge (tb \neq t \longrightarrow \text{thd-state } Va \text{ } tb = \text{READY})$ ]

**and** *c*=*SOME* *tb*. *tb*  $\neq t \wedge (tb \neq t \longrightarrow \text{thd-state } Va \text{ } tb = \text{READY})$ ]

**apply** *auto*[1]

**apply**(*subgoal-tac* *thd-state* *Va* (*SOME* *ta*. *ta*  $\neq t \wedge (ta \neq t \longrightarrow \text{thd-state } Va \text{ } ta = \text{READY})$ ) = *READY*)

**prefer** 2 **using** *exE-some*[**where** *P*= $\lambda tb. tb \neq t \wedge (tb \neq t \longrightarrow \text{thd-state } Va \text{ } tb = \text{READY})$ ]

**and** *c*=*SOME* *tb*. *tb*  $\neq t \wedge (tb \neq t \longrightarrow \text{thd-state } Va \text{ } tb = \text{READY})$ ]

**apply** *auto*[1]

**apply**(*simp* *add:inv-def*)

**apply**(*rule* *conjI*) **apply**(*simp* *add:inv-cur-def*) **apply** *auto*[1]

**apply**(*rule* *conjI*) **apply**(*simp* *add:inv-thd-waitq-def*)

**apply**(*rule* *conjI*) **apply** *auto*[1]

**apply**(*rule* *conjI*) **apply** *auto*[1]

**apply**(*rule* *conjI*) **apply** (*metis* (*no-types*, *lifting*) *Thread-State-Type.distinct*(5)

*diff-is-0-eq'*

$\text{less-Suc-eq less-Suc-eq-le nth-Cons-0 nth-append nth-mem}$ )

**apply** *auto*[1]

**apply**(*rule* *conjI*) **apply**(*simp* *add:inv-mempool-info-def*) **apply** *meson*

**apply**(*rule* *conjI*) **apply**(*simp* *add:inv-bitmap-freelist-def*) **apply** *meson*

**apply**(*rule* *conjI*) **apply**(*simp* *add:inv-bitmap-def*) **apply**(*simp* *add:Let-def*)



```

    apply(rule conjI) apply(simp add: inv-aux-vars-def mem-block-addr-valid-def)
  apply meson
    apply(rule conjI) apply(simp add: inv-bitmap0-def)
    apply(rule conjI) apply(simp add: inv-bitmapn-def)
    apply(simp add: inv-bitmap-not4free-def partner-bits-def) apply
meson
done

```

**lemma** *mp-alloc-stm7-swap-elsebody-inv*:

```

  p ∈ mem-pools Va ⇒
    inv Va ⇒
    cur Va = Some t ⇒
    (if ta = t then BLOCKED else thd-state Va ta) ≠ READY ⇒
    inv (cur-update Map.empty
      (Va (thd-state := (thd-state Va)(t := BLOCKED),
        mem-pool-info := (mem-pool-info Va)(p := mem-pool-info Va p (wait-q
:= wait-q (mem-pool-info Va p) @ [t])))
    apply(subgoal-tac thd-state Va t = RUNNING)
    prefer 2 apply(simp add: inv-def inv-cur-def) apply auto[1]

```

```

    apply(simp add: inv-def)
    apply(rule conjI) apply(simp add: inv-cur-def) apply auto[1]
    apply(rule conjI) apply(simp add: inv-thd-waitq-def)
    apply(rule conjI) apply auto[1]
    apply(rule conjI) apply (metis Thread-State-Type.distinct(6) diff-is-0-eq'
less-Suc-eq
      less-Suc-eq-le nth-Cons-0 nth-append nth-mem)
    apply (metis (no-types, lifting) Thread-State-Type.distinct(5))
    apply(rule conjI) apply(simp add: inv-mempool-info-def) apply meson
    apply(rule conjI) apply(simp add: inv-bitmap-freelist-def) apply meson
    apply(rule conjI) apply(simp add: inv-bitmap-def) apply(simp add: Let-def)
    apply(rule conjI) apply(simp add: inv-aux-vars-def mem-block-addr-valid-def)
  apply meson
    apply(rule conjI) apply(simp add: inv-bitmap0-def)
    apply(rule conjI) apply(simp add: inv-bitmapn-def)
    apply(simp add: inv-bitmap-not4free-def partner-bits-def) apply
meson
done

```

**lemma** *mp-alloc-stm7-lm-1*:

```

  mp-alloc-precond1-8-2-2 t p sz timeout ∩ {cur = Some t} ∩ {V} ∩ UNIV ∩
  {Va} ≠ {} ⇒
  Γ ⊢I Some (thd-state := thd-state(the cur := BLOCKED));
    mem-pool-info := mem-pool-info(p := mem-pool-info p (wait-q := wait-q
(mem-pool-info p) @ [the cur] ));
    swap )
  satp [mp-alloc-precond1-8-2-2 t p sz timeout ∩ {cur = Some t} ∩ {V} ∩ UNIV

```

$\cap \{Va\},$   
 $\{(s, t). s = t\}, UNIV, \llbracket '(Pair\ Va) \in UNIV \rrbracket \cap (\llbracket '(Pair\ V) \in Mem\text{-}pool\text{-}alloc\text{-}guar$   
 $t \rrbracket \cap$   
 $(mp\text{-}alloc\text{-}precond1\text{-}8\text{-}2\text{-}2\ t\ p\ sz\ timeout))]$   
**apply**(*subgoal-tac*  $V = Va$ )  
**prefer 2 apply simp**  
**apply**(*subgoal-tac*  $mp\text{-}alloc\text{-}precond1\text{-}8\text{-}2\text{-}2\ t\ p\ sz\ timeout \cap \llbracket 'cur = Some\ t \rrbracket \cap$   
 $\{V\} \cap UNIV \cap \{Va\} = \{Va\}$ )  
**prefer 2 apply auto[1]**  
**apply**(*rule subst*[**where**  $t = mp\text{-}alloc\text{-}precond1\text{-}8\text{-}2\text{-}2\ t\ p\ sz\ timeout \cap \llbracket 'cur =$   
 $Some\ t \rrbracket \cap \{V\} \cap UNIV \cap \{Va\}$  **and**  $s = \{V\}$ ])  
**apply simp**  
**apply clarsimp**  
  
**apply**(*rule Seq*[**where**  $mid = \{mp\text{-}alloc\text{-}stm7\text{-}precond3\ (mp\text{-}alloc\text{-}stm7\text{-}precond1$   
 $Va)\ t\ p\}$ ])  
**apply**(*rule Seq*[**where**  $mid = \{mp\text{-}alloc\text{-}stm7\text{-}precond1\ Va\}$ ])  
  
**apply**(*rule Basic*)  
**apply**(*simp add:fun-upd-def*)  
**apply simp apply**(*simp add:stable-def*) **apply**(*simp add:stable-def*)  
  
**apply**(*rule Basic*)  
**apply simp using**  $mp\text{-}alloc\text{-}stm7\text{-}lm\text{-}2[of\ Va\ t\ p]$  **apply metis**  
**apply simp apply**(*simp add:stable-def*) **apply**(*simp add:stable-def*)  
  
**apply**(*simp add:swap-def*)  
**apply**(*rule Cond*)  
**apply**(*simp add:stable-def*)  
  
  
**apply**(*case-tac*  $\{Va \llbracket thd\text{-}state := (thd\text{-}state\ Va)(t := BLOCKED),$   
 $mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info\ Va)$   
 $(p := mem\text{-}pool\text{-}info\ Va\ p \llbracket wait\text{-}q := wait\text{-}q\ (mem\text{-}pool\text{-}info$   
 $Va\ p) \text{ @ } [t] \rrbracket \rrbracket \rrbracket \} \cap$   
 $\{\exists t. 'thd\text{-}state\ t = READY\} = \{\}$ )  
**apply simp using Emptyprecond apply metis**  
**apply**(*rule subst*[**where**  $t = \{Va \llbracket thd\text{-}state := (thd\text{-}state\ Va)(t := BLOCKED),$   
 $mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info\ Va)$   
 $(p := mem\text{-}pool\text{-}info\ Va\ p \llbracket wait\text{-}q := wait\text{-}q\ (mem\text{-}pool\text{-}info$   
 $Va\ p) \text{ @ } [t] \rrbracket \rrbracket \rrbracket \} \cap$   
 $\{\exists t. 'thd\text{-}state\ t = READY\}$  **and**  $s = \{Va \llbracket thd\text{-}state :=$   
 $(thd\text{-}state\ Va)(t := BLOCKED),$   
 $mem\text{-}pool\text{-}info := (mem\text{-}pool\text{-}info\ Va)$   
 $(p := mem\text{-}pool\text{-}info\ Va\ p \llbracket wait\text{-}q := wait\text{-}q\ (mem\text{-}pool\text{-}info$   
 $Va\ p) \text{ @ } [t] \rrbracket \rrbracket \rrbracket \}$ )  
**apply simp**  
**apply**(*rule Seq*[**where**  $mid = \{let\ V = mp\text{-}alloc\text{-}stm7\text{-}precond3\ (mp\text{-}alloc\text{-}stm7\text{-}precond1$   
 $Va)\ t\ p\ in$

$$V(\lfloor cur := Some (SOME t. (thd-state V) t = READY) \rfloor \rfloor))$$

```

apply(rule Basic)
apply auto[1] apply simp apply(simp add:stable-def) apply(simp add:stable-def)

```

```

apply(rule Basic)
  apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
    using mp-alloc-stm7-swap-ifbody-inv apply auto[1]
    apply(simp add:lvars-nochange-def)
    using mp-alloc-stm7-swap-ifbody-inv apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
    using mp-alloc-stm7-swap-ifbody-inv apply auto[1]
    apply(simp add:lvars-nochange-def)
    using mp-alloc-stm7-swap-ifbody-inv apply auto[1]

```

```

apply simp apply(simp add:stable-def) using stable-id2 apply metis

```

```

apply(rule Basic)
  apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
    using mp-alloc-stm7-swap-elsebody-inv apply auto[1]
    apply(simp add:lvars-nochange-def)
using mp-alloc-stm7-swap-elsebody-inv apply auto[1]
apply(simp add:Mem-pool-alloc-guar-def)
  apply(rule disjI1)
  apply(rule conjI)
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
  apply(rule conjI)
    using mp-alloc-stm7-swap-elsebody-inv apply auto[1]
    apply(simp add:lvars-nochange-def)
    using mp-alloc-stm7-swap-elsebody-inv apply auto[1]

```

```

apply simp apply(simp add:stable-def) using stable-id2 apply metis
apply simp

```

done

**lemma** *mp-alloc-stm7-lm*:

$\Gamma \vdash_I \text{Some } (t \blacktriangleright \text{ATOMIC})$   
 $\text{'thd-state} := \text{'thd-state}(\text{the 'cur} := \text{BLOCKED});;$   
 $\text{'mem-pool-info} := \text{'mem-pool-info}(p := \text{'mem-pool-info } p \langle \text{wait-q} := \text{wait-q}$   
 $(\text{'mem-pool-info } p) @ [\text{the 'cur}] \rangle);;$   
 $\text{swap}$   
 $\text{END}) \text{ sat}_p [\text{mp-alloc-precond1-8-2-2 } t \text{ } p \text{ sz timeout, Mem-pool-alloc-rely } t,$   
 $\text{Mem-pool-alloc-guar } t,$   
 $\text{mp-alloc-precond1-8-2-2 } t \text{ } p \text{ sz timeout}]$   
**apply**(*simp add:stm-def*)  
**apply**(*rule Await*)  
**using** *mp-alloc-precond1-8-2-2-stb* **apply** *simp*  
**using** *mp-alloc-precond1-8-2-2-stb* **apply** *simp*  
**apply** *clarify*  
**apply**(*rule Await*)  
**using** *stable-id2* **apply** *metis*  
**using** *stable-id2* **apply** *metis*  
**apply** *clarify*  
**apply**(*case-tac mp-alloc-precond1-8-2-2 t p sz timeout  $\cap$*   
 $\{\text{'cur} = \text{Some } t\} \cap \{V\} \cap \text{UNIV} \cap \{Va\} = \{\}$   
**using** *Emptyprecond* **apply** *metis*  
**using** *mp-alloc-stm7-lm-1* **apply** *meson*

done

**term** *mp-alloc-precond1-8-2-2 t p sz timeout*

## 10.11 final proof

**lemma** *mp-alloc-stm8-guar*:

$\text{cur } V = \text{Some } t \implies \text{inv } V \implies V \langle \text{rf} := (\text{rf } V)(t := \text{True}) \rangle \in \{\text{'(Pair } V) \in$   
 $\text{Mem-pool-alloc-guar } t\}$   
**apply** *auto* **apply**(*simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def*  
 $\text{lvars-nochange-def}$ )  
**apply**(*rule disjI1*)  
**apply**(*subgoal-tac (V, V \langle \text{rf} := (\text{rf } V)(t := \text{True}) \rangle) \in \text{lvars-nochange1-4all}*)  
**using** *glnochange-inv0* **apply** *auto[1]* **apply**(*simp add:lvars-nochange1-4all-def*  
 $\text{lvars-nochange1-def}$ )  
done

**lemma** *mp-alloc-stm9-guar*:

$\text{cur } V = \text{Some } t \implies \text{inv } V \implies V \langle \text{ret} := (\text{ret } V)(t := \text{ETIMEOUT}) \rangle \in \{\text{'(Pair}$   
 $V) \in \text{Mem-pool-alloc-guar } t\}$   
**apply** *auto* **apply**(*simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def*  
 $\text{lvars-nochange-def}$ )  
**apply**(*rule disjI1*)

```

apply(subgoal-tac (V, V( $\lfloor$ ret := (ret V)(t := ETIMEOUT) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
done

```

**lemma** Mempool-alloc-satRG:  $\Gamma$  (Mem-pool-alloc t p sz timeout)  $\vdash$  Mem-pool-alloc-RGCond t p sz timeout

```

apply(simp add: Mem-pool-alloc-RGCond-def getrgformula-def)
apply(simp add: Evt-sat-RG-def Pref-def Postf-def Relyf-def Guarf-def)
apply(unfold Mem-pool-alloc-def)
apply(rule BasicEvt)
apply(unfold body-def guard-def snd-conv fst-conv)

apply(rule Seq[where mid=mp-alloc-precond7 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond6 t p timeout])
apply(rule Seq[where mid=mp-alloc-precond5 t p timeout])
apply(rule Seq[where mid=mp-alloc-precond4 t p timeout])
apply(rule Seq[where mid=mp-alloc-precond3 t p timeout])
apply(rule Seq[where mid=mp-alloc-precond2 t p timeout])

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond1-stb apply auto[1]
  using mp-alloc-precond2-stb apply simp
apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1 t p timeout  $\cap$   $\{\text{'cur} = \text{Some } t\} \cap \{V\} =$ 
{}))
    apply auto[1] apply simp
    apply(rule conjI)
      apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
      apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
      apply(rule conjI)
      apply(subgoal-tac (V, V( $\lfloor$ tmout := (tmout V)(t := timeout) $\rfloor$ )) $\in$ lvars-nochange1-4all)
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
        apply(simp add:lvars-nochange-def)
        apply(subgoal-tac (V, V( $\lfloor$ tmout := (tmout V)(t := timeout) $\rfloor$ )) $\in$ lvars-nochange1-4all)
        using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
          apply(simp add:stable-def)+

```

```

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond2-stb apply simp
  using mp-alloc-precond3-stb apply simp

```

```

apply(rule allI)
apply(rule Basic)
apply(case-tac mp-alloc-precond2 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} =$ 
{ })
  apply auto[1] apply simp
  apply(rule conjI)
    apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
    apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
    apply(rule conjI)
    apply(subgoal-tac (V, V( $\llbracket endt := (endt\ V)(t := NULL) \rrbracket$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
      apply(simp add:lvars-nochange-def)
      apply(subgoal-tac (V, V( $\llbracket endt := (endt\ V)(t := NULL) \rrbracket$ )) $\in$ lvars-nochange1-4all)
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
        apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(unfold stm-def)[1]
apply(rule Await)
using mp-alloc-precond3-stb apply simp
using mp-alloc-precond4-stb apply simp
apply clarify
apply(rule Cond)
  apply(simp add:stable-def)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond3 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} =$ 
{ })
    apply auto[1] apply auto[1]
    apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]
    apply(simp add:gvars-conf-stable-def gvars-conf-def)
    apply(subgoal-tac (V, V( $\llbracket endt := (endt\ V)(t := tick\ V + nat\ (tmout\ V$ 
t)) $\rrbracket$ ) $\in$ lvars-nochange1-4all)
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
        apply(simp add:lvars-nochange-def)
        apply(subgoal-tac (V, V( $\llbracket endt := (endt\ V)(t := tick\ V + nat\ (tmout\ V$ 
t)) $\rrbracket$ ) $\in$ lvars-nochange1-4all)
          using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
            apply simp apply(simp add:stable-def) apply(simp add:stable-def)

apply(unfold Skip-def)[1]
apply(rule Basic)
apply(case-tac mp-alloc-precond3 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} \cap$ 
 $-\llbracket OK < timeout \rrbracket = \{\}$ )
  apply auto[1] apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def)+

```

```

    apply(simp add:stable-def)+

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond4-stb apply simp
  using mp-alloc-precond5-stb apply simp
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond2 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} =$ 
{ })
    apply auto[1] apply simp
    apply(rule conjI)
      apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
      apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
      apply(rule conjI)
      apply(subgoal-tac (V, V  $\llbracket mempoolalloc-ret := (mempoolalloc-ret\ V)(t :=$ 
None  $\rrbracket) \in lvars-nochange1-4all$ )
        using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
          apply(simp add:lvars-nochange-def)
          apply(subgoal-tac (V, V  $\llbracket mempoolalloc-ret := (mempoolalloc-ret\ V)(t :=$ 
None  $\rrbracket) \in lvars-nochange1-4all$ )
            using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
              apply(simp add:stable-def)+

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond5-stb apply simp
  using mp-alloc-precond6-stb apply simp
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond5 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} =$ 
{ })
    apply auto[1] apply simp
    apply(rule conjI)
      apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
      apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
      apply(rule conjI)
      apply(subgoal-tac (V, V  $\llbracket ret := (ret\ V)(t := ESIZEERR) \rrbracket) \in lvars-nochange1-4all$ )
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
        apply(simp add:lvars-nochange-def)
        apply(subgoal-tac (V, V  $\llbracket ret := (ret\ V)(t := ESIZEERR) \rrbracket) \in lvars-nochange1-4all$ )
          using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
            apply(simp add:stable-def)+

```

```

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond6-stb apply simp
  using mp-alloc-precond7-stb apply simp
apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond6 t p timeout  $\cap \llbracket 'cur = Some\ t \rrbracket \cap \{V\} =$ 
{ })
    apply auto[1] apply simp
    apply(rule conjI)
      apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
      apply(rule conjI) apply(simp add:gvars-conf-stable-def gvars-conf-def)
      apply(rule conjI)
      apply(subgoal-tac (V, V( $\llbracket rf := (rf\ V)(t := False) \rrbracket$ )) $\in$ lvars-nochange1-4all)
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
        apply(simp add:lvars-nochange-def)
        apply(subgoal-tac (V, V( $\llbracket rf := (rf\ V)(t := False) \rrbracket$ )) $\in$ lvars-nochange1-4all)
        using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
          apply(simp add:stable-def)+

apply(rule While)
  using mp-alloc-precond7-stb apply simp
  apply(simp add:Mem-pool-alloc-post-def) apply auto[1]
  using mp-alloc-post-stb apply simp

prefer 2 apply(simp add:Mem-pool-alloc-guar-def)
prefer 2 apply (simp add: mem-pool-alloc-pre-stb)
prefer 2 apply(simp add:Mem-pool-alloc-guar-def)

apply(rule Seq[where mid=mp-alloc-precond1-8 t p sz timeout])
apply(rule Seq2[where mida=mp-alloc-precond1-7 t p sz timeout and midb=mp-alloc-precond1-70
t p sz timeout])

apply(rule Seq[where mid=mp-alloc-precond1-6 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond1-5 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond1-4 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond1-3 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond1-2 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond1-1 t p sz timeout])

apply(simp add:stm-def)
apply(rule Await)

```



```

using mp-alloc-precond1-0-stb apply simp
using mp-alloc-precond1-1-stb apply simp
apply(rule allI)
apply(rule Basic)
  apply(case-tac mp-alloc-precond1-0 t p sz timeout  $\cap \{\text{'cur} = \text{Some } t\} \cap$ 
{ V } = { })
  apply auto[1] apply clarify
    apply(rule IntI) apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac ( V, V( $\text{blk} := (\text{blk } V)(t := \text{NULL})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac ( V, V( $\text{blk} := (\text{blk } V)(t := \text{NULL})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:alloc-memblk-valid-def)
  apply(subgoal-tac ( V, V( $\text{blk} := (\text{blk } V)(t := \text{NULL})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply(simp add:stable-def)+

apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond1-1-stb apply simp
using mp-alloc-precond1-2-stb apply simp
apply(rule allI)
apply(rule Basic)
  apply(case-tac mp-alloc-precond1-1 t p sz timeout  $\cap \{\text{'cur} = \text{Some } t\} \cap$ 
{ V } = { })
  apply auto[1] apply clarify
    apply(rule IntI) apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac ( V, V( $\text{alloc-lsize-r} := (\text{alloc-lsize-r } V)(t := \text{False})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac ( V, V( $\text{alloc-lsize-r} := (\text{alloc-lsize-r } V)(t := \text{False})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)

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    apply(simp add:alloc-memblk-valid-def)
  apply(subgoal-tac (V, V(⟦alloc-lsize-r := (alloc-lsize-r V)(t := False)⟧) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:stable-def))+

  apply(simp add:stm-def)
  apply(rule Await)
  using mp-alloc-precond1-2-stb apply simp
  using mp-alloc-precond1-3-stb apply simp
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1-2 t p sz timeout ∩ ⟦'cur = Some t⟧ ∩
{V} = {})
  apply auto[1] apply clarify
  apply(rule IntI) apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

    gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac (V, V(⟦alloc-l := (alloc-l V)(t := ETIMEOUT)⟧) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

    gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac (V, V(⟦alloc-l := (alloc-l V)(t := ETIMEOUT)⟧) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:alloc-memblk-valid-def)
  apply(subgoal-tac (V, V(⟦alloc-l := (alloc-l V)(t := ETIMEOUT)⟧) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:stable-def))+

  apply(simp add:stm-def)
  apply(rule Await)
  using mp-alloc-precond1-3-stb apply simp
  using mp-alloc-precond1-4-stb apply simp
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1-3 t p sz timeout ∩ ⟦'cur = Some t⟧ ∩
{V} = {})
  apply auto[1] apply clarify
  apply(rule IntI) apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

    gvars-conf-stable-def gvars-conf-def)

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```

apply(subgoal-tac (V, V( $\lfloor$ free-l := (free-l V)(t := ETIMEOUT) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
apply(subgoal-tac (V, V( $\lfloor$ free-l := (free-l V)(t := ETIMEOUT) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
  apply(simp add:alloc-memblk-valid-def)
apply(subgoal-tac (V, V( $\lfloor$ free-l := (free-l V)(t := ETIMEOUT) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply(simp add:stable-def)+

apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond1-4-stb apply simp
using mp-alloc-precond1-5-stb apply simp
apply(rule allI)
apply(rule Basic)
  apply(case-tac mp-alloc-precond1-4 t p sz timeout  $\cap$   $\{ \text{'cur} = \text{Some } t \}$   $\cap$ 
{ V } = { })
  apply auto[1] apply clarify
  apply(rule IntI) apply auto[1]
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac (V, V( $\lfloor$ lsizes := (lsizes V)(t := [ALIGN4 (max-sz
(mem-pool-info V p)) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac (V, V( $\lfloor$ lsizes := (lsizes V)(t := [ALIGN4 (max-sz
(mem-pool-info V p)) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
apply(simp add:alloc-memblk-valid-def)
  apply(subgoal-tac (V, V( $\lfloor$ lsizes := (lsizes V)(t := [ALIGN4 (max-sz
(mem-pool-info V p)) $\rfloor$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply(simp add:stable-def)+

apply(simp add:stm-def)

```

```

apply(rule Await)
  using mp-alloc-precond1-5-stb apply simp
  using mp-alloc-precond1-6-stb apply simp
  apply(rule allI)
  apply(rule Basic)
    apply(case-tac mp-alloc-precond1-5 t p sz timeout  $\cap \{\cur = \text{Some } t\} \cap$ 
 $\{V\} = \{\}$ )
    apply auto[1] apply clarify
    apply(rule IntI) apply auto[1]
    apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
    apply(subgoal-tac ( $V, V(i := (i \ V)(t := 0)) \in \text{lvars-nochange1-4all}$ )
      using glnochange-inv0 apply auto[1])
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
    apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
    apply(subgoal-tac ( $V, V(i := (i \ V)(t := 0)) \in \text{lvars-nochange1-4all}$ )
      using glnochange-inv0 apply auto[1])
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def lvars-nochange-def)
    apply(simp add:alloc-memblk-valid-def)
    apply(rule conjI)
    apply(subgoal-tac ( $V, V(i := (i \ V)(t := 0)) \in \text{lvars-nochange1-4all}$ )
      using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(simp add:inv-def inv-mempool-info-def) apply (meson Suc-leI)
    apply simp apply(simp add:stable-def) using stable-id2 apply auto[1]

using lsize-loop-stm[of t p sz timeout] apply clarsimp

using precnd17-bl-170 apply simp

apply(rule Cond)
  using mp-alloc-precond1-70-stb apply simp

  apply(simp add:stm-def)
  apply(rule Await)
    using mp-alloc-precond1-70-1-stb apply simp
    using mp-alloc-precond1-8-stb apply auto[1]

    apply(rule allI)
    apply(rule Basic)
    apply(case-tac mp-alloc-precond1-70-1 t p sz timeout  $\cap \{\cur = \text{Some } t\} \cap$ 
 $\{V\} = \{\}$ )

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```

    apply auto[1] apply clarify
    apply(rule IntI) apply simp
    apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def
lvars-nochange-def)
    apply(rule disjI1)
    apply(subgoal-tac (V, V( $\lfloor$ ret := (ret V)(t := ESIZEERR) $\rfloor$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply(simp add:inv-def inv-mempool-info-def)
apply simp
    apply(rule IntI) prefer 2 apply simp
    apply(subgoal-tac (V, V( $\lfloor$ ret := (ret V)(t := ESIZEERR) $\rfloor$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply simp using stable-id2 apply auto[1] using stable-id2 apply auto[1]

apply(rule Cond)
using mp-alloc-precond1-70-2-stb apply simp

apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond1-70-2-1-stb apply simp
using mp-alloc-precond1-8-stb apply auto[1]

apply(rule allI)
apply(rule Basic)
    apply(case-tac mp-alloc-precond1-70-2-1 t p sz timeout  $\cap$   $\lfloor$ 'cur = Some
t $\rfloor$   $\cap$  {V} = {})
    apply auto[1] apply clarify
    apply(rule IntI) apply simp
    apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def
lvars-nochange-def)
    apply(rule disjI1)
    apply(subgoal-tac (V, V( $\lfloor$ ret := (ret V)(t := ENOMEM) $\rfloor$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply(simp add:inv-def inv-mempool-info-def)
apply simp
    apply(rule IntI) prefer 2 apply simp
    apply(subgoal-tac (V, V( $\lfloor$ ret := (ret V)(t := ENOMEM) $\rfloor$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply simp using stable-id2 apply auto[1] using stable-id2 apply auto[1]

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apply(rule Seq[where mid=mp-alloc-precond2-1 t p sz timeout])

using mp-alloc-stm3-lm apply simp

apply(rule Cond)
  using mp-alloc-precond2-1-stb apply simp

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond2-1-0-stb apply simp
  using mp-alloc-precond1-8-stb apply auto[1]

apply(rule allI)
apply(rule Basic)
apply(case-tac mp-alloc-precond2-1-0 t p sz timeout  $\cap \{\text{'cur} = \text{Some } t\}$ 
 $\cap \{V\} = \{\}$ )
  apply auto[1] apply clarify
  apply(rule IntI) apply simp
    apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def
gvars-conf-def lvars-nochange-def)
    apply(rule disjI1)
    apply(subgoal-tac (V, V( $\text{ret} := (\text{ret } V)(t := \text{EAGAIN})$ )) $\in$ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply(simp add:inv-def inv-mempool-info-def)
apply simp
  apply(rule IntI) prefer 2 apply simp
  apply(subgoal-tac (V, V( $\text{ret} := (\text{ret } V)(t := \text{EAGAIN})$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply simp using stable-id2 apply auto[1] using stable-id2 apply
auto[1]

apply(rule Seq[where mid=mp-alloc-precond2-1-4 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond2-1-3 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond2-1-2 t p sz timeout])
apply(rule Seq[where mid=mp-alloc-precond2-1-1-loopinv t p sz timeout])

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond2-1-1-stb apply simp
  using mp-alloc-precond2-1-1-loopinv-stb apply simp

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```

    apply(rule allI)
    apply(rule Basic)
    apply(case-tac mp-alloc-precond2-1-1 t p sz timeout  $\cap \{\}'cur = Some\ t\}$ 
 $\cap \{V\} = \{\}$ )
    apply auto[1] apply clarify
    apply(rule IntI) apply simp
    apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def
gvars-conf-def lvars-nochange-def)
    apply(rule disjI1)
    apply(subgoal-tac (V, V( $\lfloor from-l := (from-l\ V)(t := free-l\ V$ 
t) $\rfloor$ ) $\in lvars-nochange1-4all$ )
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply(rule IntI) prefer 2
    apply(case-tac i V t = 0) apply(simp add:inv-def inv-mempool-info-def)
apply simp
    apply(rule IntI) prefer 2 apply simp
    apply(subgoal-tac (V, V( $\lfloor from-l := (from-l\ V)(t := free-l\ V\ t)\rfloor$ ) $\in lvars-nochange1-4all$ )
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
    apply simp using stable-id2 apply auto[1] using stable-id2 apply
auto[1]

```

```

using mp-alloc-stm4-lm apply simp

```

```

using mp-alloc-stm5-lm apply simp

```

```

using mp-alloc-stm6-lm apply simp

```

```

    apply(simp add:stm-def)
    apply(rule Await)
    using mp-alloc-precond2-1-4-stb apply simp
    using mp-alloc-precond1-8-stb apply auto[1]

    apply(rule allI)
    apply(rule Basic)
    apply(case-tac mp-alloc-precond2-1-4 t p sz timeout  $\cap \{\}'cur = Some\ t\}$ 
 $\cap \{V\} = \{\}$ )
    apply auto[1] apply clarify
    apply(rule IntI) apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def
gvars-conf-def lvars-nochange-def)
    apply(rule disjI1)
    apply(subgoal-tac (V, V( $\lfloor ret := (ret\ V)(t := OK)\rfloor$ ) $\in lvars-nochange1-4all$ )
    using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def

```

```

lvars-nochange1-def)
  apply(rule IntI) apply(rule IntI) apply(rule IntI) apply(rule IntI)
  apply(subgoal-tac (V, V( $\text{ret} := (\text{ret } V)(t := \text{OK})$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
  apply auto[1] apply auto[1] apply auto[1]
  apply auto[1] apply(simp add:alloc-memblk-valid-def) apply auto[1]
  apply(simp add:alloc-memblk-valid-def) apply auto[1]
  apply simp using stable-id2 apply auto[1] using stable-id2 apply
auto[1]

  apply(simp add:Mem-pool-alloc-guar-def)
  apply(simp add:Mem-pool-alloc-guar-def)
  apply(simp add:Mem-pool-alloc-guar-def)

apply(rule Cond)
using mp-alloc-precond1-8-stb apply simp

apply(rule Seq[where mid=mp-alloc-precond1-8-1-1 t p sz timeout])

apply(simp add:stm-def)
apply(rule Await)
  using mp-alloc-precond1-8-1-stb apply auto[1]
  using mp-alloc-precond1-8-1-1-stb apply auto[1]
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1-8-1 t p sz timeout  $\cap$   $\{ \text{'cur} = \text{Some } t \}$ 
 $\cap \{ V \} = \{ \}$ )
  apply auto[1] apply clarify
  apply(rule IntI)
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def) apply(rule disjI1)
  apply(subgoal-tac (V, V( $\text{rf} := (\text{rf } V)(t := \text{True})$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
  apply(rule IntI) apply(rule IntI) apply(rule IntI)
  apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def lvars-nochange-def

      gvars-conf-stable-def gvars-conf-def)
  apply(subgoal-tac (V, V( $\text{rf} := (\text{rf } V)(t := \text{True})$ )) $\in$ lvars-nochange1-4all)
  using glnochange-inv0 apply auto[1]
  apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
  apply simp
  apply(simp add:alloc-memblk-valid-def) apply simp

```



```

apply(simp add:stable-def)+

apply(rule Cond)
using mp-alloc-precond1-8-1-1-stb apply auto[1]

apply(simp add:stm-def)
apply(rule Await)
using mp-alloc-precond1-8-1-2-stb apply auto[1]
using mp-alloc-precond7-stb apply auto[1]
apply(rule allI)
apply(rule Basic)
apply(case-tac mp-alloc-precond1-8-1-2 t p sz timeout  $\cap \{ \text{cur} = \text{Some } t \} \cap \{ V \} = \{ \}$ )
apply auto[1] apply auto[1]
apply(simp add:Mem-pool-alloc-guar-def gvars-conf-stable-def gvars-conf-def
lvars-nochange-def)
apply(rule disjI1)
apply(subgoal-tac (V, V( $\text{ret} := (\text{ret } V)(t := \text{ENOMEM})$ )) $\in$ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply(subgoal-tac (V, V( $\text{ret} := (\text{ret } V)(t := \text{ENOMEM})$ )) $\in$ lvars-nochange1-4all)
using glnochange-inv0 apply auto[1] apply(simp add:lvars-nochange1-4all-def
lvars-nochange1-def)
apply simp
apply(simp add:stable-def) apply auto[1]
apply(simp add:stable-def)

apply(simp add:Skip-def)
apply(rule Basic)
apply auto[1] apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]

using mp-alloc-precond1-8-1-3-stb apply auto[1]
using mp-alloc-precond7-stb apply auto[1]

apply(simp add:Mem-pool-alloc-guar-def)

apply(rule Cond)
using mp-alloc-precond1-8-2-stb apply simp

apply(simp add:Skip-def)
apply(rule Basic) apply auto[1]
apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]
using mp-alloc-precond1-8-2-1-stb apply simp
using mp-alloc-precond7-stb apply simp

```

```

apply(rule Seq[where mid=mp-alloc-precond1-8-2-2 t p sz timeout])

using mp-alloc-stm7-lm apply simp

apply(rule Cond)
  using mp-alloc-precond1-8-2-2-stb apply auto[1]
  apply(rule Seq[where mid=mp-alloc-precond1-8-2-4 t p sz timeout])

apply(unfold stm-def)[1]
apply(rule Await)
  using mp-alloc-precond1-8-2-3-stb mp-pred1823-eq apply auto[1]
  using mp-alloc-precond1-8-2-4-stb apply blast
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1-8-2-3 t p sz timeout  $\cap \llbracket 'cur = Some$ 
 $t \rrbracket \cap \{V\} = \{\}$ )
    apply auto[1] apply auto[1]
    apply(simp add:Mem-pool-alloc-guar-def) apply(rule disjI1)
    apply(simp add:Mem-pool-alloc-guar-def lvars-nochange1-def
lvars-nochange-def
gvars-conf-stable-def gvars-conf-def)
      apply(subgoal-tac (V, V( $\llbracket tmout := (tmout V)(t := int (endt V t) -$ 
 $int (tick V) \rrbracket \rrbracket \in lvars-nochange1-4all$ ))
        using glnochange-inv0 apply auto[1]
        apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
          apply(subgoal-tac (V, V( $\llbracket tmout := (tmout V)(t := int (endt V t) -$ 
 $int (tick V) \rrbracket \rrbracket \in lvars-nochange1-4all$ ))
            using glnochange-inv0 apply auto[1]
            apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def) apply auto[1]
            apply(simp add:stable-def) apply auto[1] apply(simp add:stable-def)
            apply auto[1]

apply(rule Cond)
  using mp-alloc-precond1-8-2-4-stb apply blast

apply(rule Seq[where mid=mp-alloc-precond1-8-2-5 t p sz timeout])

apply(unfold stm-def)[1]
apply(rule Await)
  using mp-alloc-precond1-8-2-40-stb apply blast
  using mp-alloc-precond1-8-2-5-stb apply blast
  apply(rule allI)
  apply(rule Basic)
  apply(case-tac mp-alloc-precond1-8-2-40 t p sz timeout  $\cap \llbracket 'cur =$ 
 $Some t \rrbracket \cap \{V\} = \{\}$ )

```

```

    apply auto[1] apply auto[1]
    using mp-alloc-stm8-guar apply simp
    apply(subgoal-tac (V, V(|rf := (rf V)(t := True)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
    using mp-alloc-stm8-guar apply simp
    apply(subgoal-tac (V, V(|rf := (rf V)(t := True)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
    using mp-alloc-stm8-guar apply simp
    apply simp apply(simp add:stable-def) apply auto[1] apply(simp
add:stable-def) apply auto[1]

    apply(unfold stm-def)[1]
    apply(rule Await)
    using mp-alloc-precond1-8-2-5-stb apply blast
    using mp-alloc-precond7-stb apply blast
    apply(rule allI)
    apply(rule Basic)
    apply(case-tac mp-alloc-precond1-8-2-5 t p sz timeout ∩ {cur =
Some t} ∩ {V} = {}))
    apply auto[1] apply auto[1]
    using mp-alloc-stm9-guar apply simp
    apply(subgoal-tac (V, V(|ret := (ret V)(t := ETIMEOUT)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
    using mp-alloc-stm9-guar apply simp
    apply(subgoal-tac (V, V(|ret := (ret V)(t := ETIMEOUT)|)) ∈ lvars-nochange1-4all)
    using glnochange-inv0 apply auto[1]
    apply(simp add:lvars-nochange1-4all-def lvars-nochange1-def
lvars-nochange-def)
    apply simp apply(simp add:stable-def) apply auto[1] apply(simp
add:stable-def)

    apply(unfold Skip-def)[1]
    apply(rule Basic)
    apply auto[1]
    prefer 2 using mp-alloc-precond1-8-2-41-stb apply fast
    prefer 2 using mp-alloc-precond7-stb apply blast
    apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]

    apply(simp add:Mem-pool-alloc-guar-def)

    apply(unfold Skip-def)[1]
    apply(rule Basic)

```

```

    apply auto[1]
    apply(simp add:Mem-pool-alloc-guar-def) apply auto[1]
    using mp-alloc-precond1-8-2-20-stb apply fast
    using mp-alloc-precond7-stb apply blast

  apply(simp add:Mem-pool-alloc-guar-def)+
done

end

theory memory-manage-sys
imports rg-cond func-cor-other func-cor-mempoolfree func-cor-mempoolalloc
begin

```

## 11 formal specification of Zephyr memory management

**definition** *Mem-pool-alloc-RGF* :: *Thread*  $\Rightarrow$  *mempool-ref*  $\Rightarrow$  *nat*  $\Rightarrow$  *int*  $\Rightarrow$  (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-e*  
**where** *Mem-pool-alloc-RGF* *t p sz timeout*  
 $\equiv$  (*Mem-pool-alloc* *t p sz timeout*, *Mem-pool-alloc-RGCond* *t p sz timeout*)

**definition** *Mem-pool-free-RGF* :: *Thread*  $\Rightarrow$  *Mem-block*  $\Rightarrow$  (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-e*  
**where** *Mem-pool-free-RGF* *t b*  $\equiv$  (*Mem-pool-free* *t b*, *Mem-pool-free-RGCond* *t b*)

**definition** *Schedule-RGF* :: *Thread*  $\Rightarrow$  (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-e*  
**where** *Schedule-RGF* *t*  $\equiv$  (*Schedule* *t*, *Schedule-RGCond* *t*)

**definition** *Tick-RGF* :: (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-e*  
**where** *Tick-RGF*  $\equiv$  (*Tick*, *Tick-RGCond*)

**definition** *Thread-RGF* :: *Thread*  $\Rightarrow$  (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-es*  
**where** *Thread-RGF* *t*  $\equiv$  (*rgf-EvtSys* (( $\bigcup$  (*p*, *sz*, *timeout*).{*Mem-pool-alloc-RGF* *t p sz timeout*}  $\cup$   
 $(\bigcup b. \{ \text{Mem-pool-free-RGF } t \ b \} )$ ),  
 $RG[(\text{Mem-pool-free-pre } t \cap \text{Mem-pool-alloc-pre } t),$   
 $(\text{Mem-pool-free-rely } t \cap \text{Mem-pool-alloc-rely } t),$   
 $(\text{Mem-pool-free-guar } t \cup \text{Mem-pool-alloc-guar } t),$   
 $(\text{Mem-pool-free-post } t \cup (\bigcup (p, sz, timeout). \text{Mem-pool-alloc-post } t \ p \ sz$   
 $timeout))]]$ )

**definition** *Scheduler-RGF* :: (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-es*  
**where** *Scheduler-RGF*  $\equiv$  (*rgf-EvtSys* ( $\bigcup t. \{ \text{Schedule-RGF } t \}$ ),  
 $RG[\{s. \text{inv } s\}, \text{Schedule-rely}, \text{Schedule-guar}, \{s. \text{inv } s\}]$ )

**definition** *Timer-RGF* :: (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-es*  
**where** *Timer-RGF*  $\equiv$  (*rgf-EvtSys* {*Tick-RGF*},  
*RG*[{*True*}, *Tick-rely*, *Tick-guar*, {*True*}])

**definition** *Memory-manage-system-Spec* :: (*EventLabel*, *Core*, *State*, *State com option*) *rgformula-par*  
**where** *Memory-manage-system-Spec* *k*  $\equiv$   
*case k of* ( $\mathcal{T}$  *t*)  $\Rightarrow$  *Thread-RGF t*  
 $\mid \mathcal{S} \Rightarrow$  *Scheduler-RGF*  
 $\mid$  *Timer*  $\Rightarrow$  *Timer-RGF*

## 12 functional correctness of the whole specification

**definition** *sys-rely*  $\equiv$  {}

**definition** *sys-guar*  $\equiv$  *Tick-guar*  $\cup$  *Schedule-guar*  $\cup$  ( $\bigcup t. (\text{Mem-pool-free-guar } t \cup \text{Mem-pool-alloc-guar } t)$ )

**lemma** *scheduler-esys-sat*:  $\Gamma \vdash \text{fst } (\text{Memory-manage-system-Spec } \mathcal{S})$   
 $\text{sat}_s [\text{Pre}_{es} (\text{Memory-manage-system-Spec } \mathcal{S}),$   
 $\text{Rely}_{es} (\text{Memory-manage-system-Spec } \mathcal{S}),$   
 $\text{Guar}_{es} (\text{Memory-manage-system-Spec } \mathcal{S}),$   
 $\text{Post}_{es} (\text{Memory-manage-system-Spec } \mathcal{S})]$   
**apply**(*simp add:Memory-manage-system-Spec-def Scheduler-RGF-def Schedule-RGF-def*)  
**apply**(*rule EvtSys-h*)  
**apply** *auto*[1] **apply**(*simp add:E<sub>e</sub>-def Pre<sub>e</sub>-def Rely<sub>e</sub>-def Guar<sub>e</sub>-def Post<sub>e</sub>-def*)  
**using** *Schedule-satRG* **apply**(*simp add:Schedule-RGCond-def Evt-sat-RG-def*  
*Pre<sub>f</sub>-def Rely<sub>f</sub>-def Guar<sub>f</sub>-def Post<sub>f</sub>-def*)  
**apply** *fast*  
**apply**(*simp add:Pre<sub>es</sub>-def Pre<sub>e</sub>-def Schedule-RGCond-def*)  
**apply**(*simp add:Rely<sub>es</sub>-def Rely<sub>e</sub>-def Schedule-RGCond-def*)  
**apply**(*simp add:Guar<sub>es</sub>-def Guar<sub>e</sub>-def Schedule-RGCond-def*)  
**apply**(*simp add:Post<sub>es</sub>-def Post<sub>e</sub>-def Schedule-RGCond-def*)  
**apply**(*simp add:Post<sub>e</sub>-def Pre<sub>e</sub>-def Schedule-RGCond-def getrgformula-def*)  
**apply**(*simp add:Pre<sub>es</sub>-def Rely<sub>es</sub>-def getrgformula-def*)  
**using** *stable-inv-sched-rely* **apply**(*simp add:stable-def*)  
**apply**(*simp add:Guar<sub>es</sub>-def getrgformula-def Schedule-guar-def*)  
**done**

**lemma** *thread-esys-sat*:  $\Gamma \vdash \text{fst } (\text{Memory-manage-system-Spec } (\mathcal{T} \ x))$   
 $\text{sat}_s [\text{Pre}_{es} (\text{Memory-manage-system-Spec } (\mathcal{T} \ x)),$   
 $\text{Rely}_{es} (\text{Memory-manage-system-Spec } (\mathcal{T} \ x)),$   
 $\text{Guar}_{es} (\text{Memory-manage-system-Spec } (\mathcal{T} \ x)),$   
 $\text{Post}_{es} (\text{Memory-manage-system-Spec } (\mathcal{T} \ x))]$

```

apply(simp add:Memory-manage-system-Spec-def Thread-RGF-def Mem-pool-alloc-RGF-def
Mem-pool-free-RGF-def)
apply(rule EvtSys-h)
  apply auto[1]
    apply(simp add:Ee-def Pree-def Relye-def Guare-def Poste-def getrgformula-def
Mem-pool-alloc-RGCond-def)
      using Mempool-alloc-satRG apply(simp add:Evt-sat-RG-def Mem-pool-alloc-RGCond-def

getrgformula-def Pref-def Relyf-def Guarf-def
Postf-def) apply fast
        apply(simp add:Ee-def Pree-def Relye-def Guare-def Poste-def getrgformula-def
Mem-pool-free-RGCond-def)
          using Mempool-free-satRG apply(simp add:Evt-sat-RG-def Mem-pool-free-RGCond-def

getrgformula-def Pref-def Relyf-def Guarf-def
Postf-def) apply fast
            apply auto[1]
              apply(simp add:Prees-def Pree-def getrgformula-def Mem-pool-alloc-RGCond-def)
              apply(simp add:Prees-def Pree-def getrgformula-def Mem-pool-free-RGCond-def)
              apply auto[1]
                apply(simp add:Relyes-def Relye-def getrgformula-def Mem-pool-alloc-RGCond-def)
                apply(simp add:Relyes-def Relye-def getrgformula-def Mem-pool-free-RGCond-def)
                apply auto[1]
                  apply(simp add:Guares-def Guare-def getrgformula-def Mem-pool-alloc-RGCond-def)
                  apply(simp add:Guares-def Guare-def getrgformula-def Mem-pool-free-RGCond-def)
                  apply auto[1]
                    apply(simp add:Postes-def Poste-def getrgformula-def Mem-pool-alloc-RGCond-def)
apply auto[1]
  apply(simp add:Postes-def Poste-def getrgformula-def Mem-pool-free-RGCond-def)
  apply auto[1]
    apply(simp add:Poste-def Pree-def Mem-pool-alloc-RGCond-def getrgformula-def
Mem-pool-alloc-post-def)
      apply(simp add:Poste-def Pree-def Mem-pool-alloc-RGCond-def Mem-pool-free-RGCond-def

getrgformula-def Mem-pool-alloc-post-def)
        apply(simp add:Poste-def Pree-def Mem-pool-alloc-RGCond-def Mem-pool-free-RGCond-def

getrgformula-def Mem-pool-free-post-def)
          apply(simp add:Poste-def Pree-def Mem-pool-free-RGCond-def getrgformula-def
Mem-pool-free-post-def)
            apply(simp add:Prees-def Relyes-def getrgformula-def Mem-pool-free-rely-def Mem-pool-alloc-rely-def)

defer 1
  apply(simp add:Guares-def getrgformula-def Mem-pool-free-guar-def)
  using mem-pool-free-pre-stb apply(simp add:Mem-pool-free-rely-def)
done

lemma timer-esys-sat:  $\Gamma \vdash \text{fst } (\text{Memory-manage-system-Spec Timer})$ 
  sats [Prees (Memory-manage-system-Spec Timer),

```

```

    Relyes (Memory-manage-system-Spec Timer),
    Guares (Memory-manage-system-Spec Timer),
    Postes (Memory-manage-system-Spec Timer)]
apply(simp add:Memory-manage-system-Spec-def Timer-RGF-def Tick-RGF-def)
apply(rule EvtSys-h)
    apply auto[1] apply(simp add:Ee-def Pree-def Relye-def Guare-def Poste-def)
    using Tick-satRG apply(simp add:Tick-RGCond-def Evt-sat-RG-def Pref-def
Relyf-def Guarf-def Postf-def)
    apply fast
    apply(simp add:Prees-def Pree-def Tick-RGCond-def)
    apply(simp add:Relyes-def Relye-def Tick-RGCond-def)
    apply(simp add:Guares-def Guare-def Tick-RGCond-def)
    apply(simp add:Postes-def Poste-def Tick-RGCond-def)
    apply(simp add:Poste-def Pree-def Tick-RGCond-def getrgformula-def)
    apply(simp add:Prees-def Relyes-def getrgformula-def)
    using stable-inv-sched-rely apply(simp add:stable-def)
    apply(simp add:Guares-def getrgformula-def Tick-guar-def)
done

```

**lemma** esys-sat:  $\Gamma \vdash \text{fst} \text{ (Memory-manage-system-Spec } k)$

```

    sats [Prees (Memory-manage-system-Spec k),
    Relyes (Memory-manage-system-Spec k),
    Guares (Memory-manage-system-Spec k),
    Postes (Memory-manage-system-Spec k)]
    apply(induct k)
    using scheduler-esys-sat apply fast
    using thread-esys-sat apply fast
    using timer-esys-sat apply fast
done

```

**lemma** s0-esys-pre:  $\{s0\} \subseteq \text{Pre}_{es} \text{ (Memory-manage-system-Spec } k)$

```

apply(induct k)
    apply(simp add:Memory-manage-system-Spec-def Prees-def Scheduler-RGF-def
getrgformula-def)
    using s0-inv apply fast
    apply(simp add:Memory-manage-system-Spec-def Prees-def Thread-RGF-def getrgformula-def)
    using s0-inv s0a4 s0a10 apply auto[1]
    apply(simp add:Memory-manage-system-Spec-def Prees-def Timer-RGF-def getrgformula-def)
done

```

**lemma** alloc-free-eq-guar:  $\text{Mem-pool-free-guar } x = \text{Mem-pool-alloc-guar } x$

```

    by(simp add:Mem-pool-free-guar-def Mem-pool-alloc-guar-def)

```

**lemma** alloc-free-eq-rely:  $\text{Mem-pool-free-rely } x = \text{Mem-pool-alloc-rely } x$

```

    by(simp add:Mem-pool-free-rely-def Mem-pool-alloc-rely-def)

```

**lemma** esys-guar-in-other:

```

     $jj \neq k \longrightarrow \text{Guar}_{es} \text{ (Memory-manage-system-Spec } jj) \subseteq \text{Rely}_{es} \text{ (Memory-manage-system-Spec } k)$ 

```

```

apply auto
apply(induct jj)
  apply(induct k)
    apply simp
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Scheduler-RGF-def
Thread-RGF-def getrgformula-def)
    using schedguar-in-allocrely apply(simp add: Mem-pool-free-rely-def Mem-pool-alloc-rely-def)
apply auto[1]
  apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Scheduler-RGF-def
Timer-RGF-def getrgformula-def)
  using schedguar-in-tickrely apply auto[1]
  apply(induct k)
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Scheduler-RGF-def
Thread-RGF-def getrgformula-def)
    apply auto[1]
    using allocguar-in-schedrely alloc-free-eq-guar apply fast
    using allocguar-in-schedrely apply fast
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Thread-RGF-def
getrgformula-def)
    apply auto[1]
    using allocguar-in-allocrely alloc-free-eq-guar alloc-free-eq-rely apply fast+
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Timer-RGF-def
Thread-RGF-def getrgformula-def)
    apply auto[1]
    using allocguar-in-tickrely alloc-free-eq-guar alloc-free-eq-rely apply fast+
  apply(induct k)
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Scheduler-RGF-def
Timer-RGF-def getrgformula-def)
    using tickguar-in-schedrely apply fast
    apply(simp add: Guares-def Relyes-def Memory-manage-system-Spec-def Thread-RGF-def
Timer-RGF-def getrgformula-def)
    apply auto[1]
    using tickguar-in-allocrely alloc-free-eq-guar alloc-free-eq-rely apply fast+
done

```

```

lemma esys-guar-in-sys: Guares (Memory-manage-system-Spec k) ⊆ sys-guar
apply(induct k)
  apply(simp add: Guares-def Memory-manage-system-Spec-def Scheduler-RGF-def
getrgformula-def sys-guar-def) apply auto[1]
  apply(simp add: Guares-def Memory-manage-system-Spec-def Thread-RGF-def
getrgformula-def sys-guar-def) apply auto[1]
  apply(simp add: Guares-def Memory-manage-system-Spec-def Timer-RGF-def
getrgformula-def sys-guar-def) apply auto[1]
done

```

```

lemma mem-sys-sat:  $\Gamma \vdash \text{Memory-manage-system-Spec SAT } [\{s0\}, \text{sys-rely}, \text{sys-guar},$ 
UNIV]
apply(rule ParallelESys[of  $\Gamma$  Memory-manage-system-Spec  $\{s0\}$  sys-rely sys-guar
UNIV])

```



```

apply clarify using esys-sat apply fast
using s0-esys-pre apply fast
apply(simp add:sys-rely-def)
using esys-guar-in-other apply fast
using esys-guar-in-sys apply fast
apply simp
done

end

```

```

theory memory-management-inv
  imports memory-manage-sys
begin

```

## 13 invariant verification

```

theorem invariant-presv-pares  $\Gamma$  inv (paresys-spec Memory-manage-system-Spec)
  {s0} sys-rely
  apply(rule invariant-theorem[where  $G = \text{sys-guar}$  and  $pst = UNIV$ ])
    using mem-sys-sat apply fast
    apply(simp add:sys-rely-def stable-def)
    apply(simp add:sys-guar-def)
    apply(rule stable-un-R) apply(rule stable-un-R)
      using tick-guar-stb-inv apply(simp add:stable-def)
      using sched-guar-stb-inv apply(simp add:stable-def)
    apply(rule stable-un-S) apply clarify apply(rule stable-un-R)
      using alloc-guar-stb-inv alloc-free-eq-guar apply(simp add:stable-def)
      using alloc-guar-stb-inv apply(simp add:stable-def)
    using s0-inv apply simp
  done

end

```