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CSCI 230

README

After testing linear, quadratic, and pseudo-random probing on a hash table of size 5, it was determined that the cost of each function was constant when the position after the first probe was the EMPTYKEY or null. If the key needed to be probed more to find the next empty index, then the cost increased to linear. During experimentation, I noticed that some values could not be inserted or removed with the quadratic and pseudo-random probing functions because the function would keep alternating between 2 of the same indices. Even if there was an empty or null position, the function would never locate it, thus never inserting/removing the value. This would cause the cost of the function to increase. The above graph is not an accurate representation of what the plot should be based off Figure 9.10 in the HW6 pdf. One reason the above graph is different is because for some insert and remove methods, the timestamp never increased, making the time 0. This in turn, would cause the trend line to be flat(constant).