

ECE 252: Intro to Computer Engineering

Week 12 Discussion



Attendance via TopHat

Course code: **265393**

Attendance code:

2017



ECE Discovery Panel Series

Get to know a technical area:
Applied Electromagnetics and Acoustics

Application Areas -- Advanced Electives -- Job Opportunities



Acoustic sensing, microwave ablation, electronic jamming, wireless power transfer







Tuesday, Nov 19 , 12:20-12:50pm **1413 EH (Cheney Room)** Grab the latest ECE

Come for the insights, stay for the pizza!



Subroutines

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- Why have subroutines?
 - The same sequence of instructions needs to be used repeatedly in a program
 - Avoid copy/paste errors
 - Make it easier to modify/maintain
 - Create a library of functions that can be used by you and other programmers
 - Reusability!
 - And who wants to look at a massive program and try to figure out why it doesn't work?
 - Yuck!
- Subroutines support the development of modular software



Calling a Subroutine Attendance code: 2017

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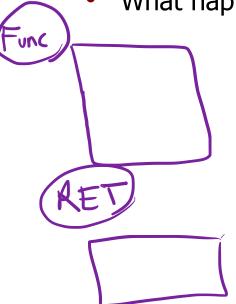
- JSR label ; Jump to SubRoutine
 - $R7 \leftarrow PC^+, PC \leftarrow PC^+ + SEXT(PCoffset11)$
 - Range is approximately +/-1024 words

- What if the subroutine is further away?
- JSRR BaseR; Jump to SubRoutine (Register)
 - temp \leftarrow PC⁺, PC \leftarrow BaseR, R7 \leftarrow temp
 - JSRR R7 will work as expected
 - In reality, we can update R7 and PC simultaneously and there is no need for a temporary register



Returning from a Subroutine

- To return from a subroutine, the subroutine code must use the value in R7
 - Execute the RET instruction, which does PC ← R7
 - Programming tip: Put the **RET** instruction at the end of the subroutine right away!
 - What happens if you forget **RET**?



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Returning from a Subroutine

- All subroutine calls will modify R7!
 - If a subroutine calls a subroutine, it needs to save R7, call the subroutine, and then restore R7

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Returning from a Subroutine

 What happens if subroutine A calls subroutine B, but A does not save and restore R7?

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Well-Behaved Subroutines

- If a subroutine uses a register (and changes its value), it could cause problems for the caller
- There are two approaches to prevent subroutines from corrupting the caller's state
 - Have the <u>caller</u> save all registers in use so they will not be corrupted by the subroutine (**caller-save**)
 - Have the <u>subroutine</u> save all registers it uses, so the caller can assume no registers (other than **R7**) are corrupted by the subroutine (**callee-save**)
- Callee-save is most commonly used
 - In LC-3 (and others), caller must still save/restore R7



Subroutine Parameters and Results

- To send information to a subroutine, the caller can use registers or memory
 - The subroutine code must assume that the caller put the information where it was expected to be
 - The subroutine code <u>must</u> be well-documented as to parameter locations and constraints!
 - The subroutine CANNOT possibly know whether or not the caller put the correct parameters in

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Subroutine Parameters and Results

- The results from a subroutine can be returned in registers or memory
 - Common error: If returning a result in a register, be sure that the register is <u>NOT</u> saved and restored by the subroutine (or you'll lose your result!)



Using Return Values

- If result is returned in memory, you need to read from memory to see the result
- If the result is returned in a register...
 - You don't have to do anything special before using that register in any instruction that <u>explicitly</u> uses a register
 - What about branches? Can you assume a return value set the condition codes?
 - What if the subroutine saves/restores registers?
 - What if the subroutine is changed later?
 - Good (safe) practice is to explicitly set condition codes based on the returned value

branch based on value returned in R0

JSR mysub
ADD R0, R0, #0
BRz DoThing



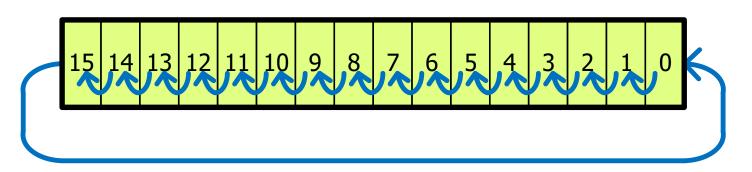
How do you get good at this?

- The only way to learn programming...
 is to DO programming!
- The more you do, the easier it gets... because you know
 - ...what the instructions do
 - ...how to do stuff



Subroutine Practice 1

- Write a subroutine ROL1 that rotates an operand left by one bit
 - Parameters: R0 is value
 - Returns: R0 is rotated result
- Subroutine should not corrupt any caller registers



- Rotate left is a left shift except LSb←MSb instead of 0
 - Left shift (value) = 2(value) = (value) + (value)
 - Then, if original MSb was 1, add 1 to complete rotate



Rotate Left Examples

Rotate Left by 1

Before	<mark>0</mark> 011	1100	0101	1010
After	0111	1000	1011	0100

Rotate Left by 1

Before	<mark>1</mark> 010 01	01 1100	0011
After	0100 10°	11 1000	011 <mark>1</mark>

Rotate Left by 2

Before	<mark>01</mark> 10	0010	1001	1111
After	1000	1010	0111	11 <mark>01</mark>

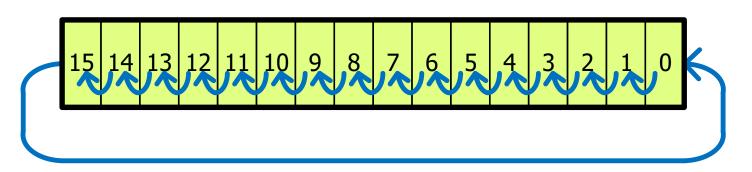
Rotate Left by 4

Before	0110	0010	1001	1111
After	0010	1001	1111	0110



Subroutine Practice 1

- Write a subroutine ROL1 that rotates an operand left by one bit
 - Parameters: R0 is value
 - Returns: R0 is rotated result
- Subroutine should not corrupt any caller registers

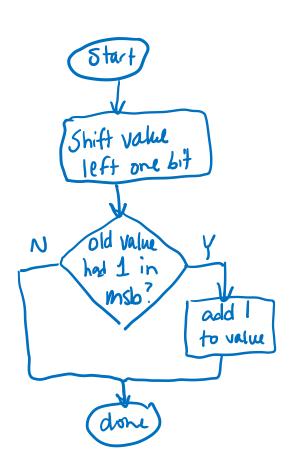


- Rotate left is a left shift except LSb←MSb instead of 0
 - Left shift (value) = 2(value) = (value) + (value)
 - Then, if original MSb was 1, add 1 to complete rotate



ROL1 rotates left by 1 bit

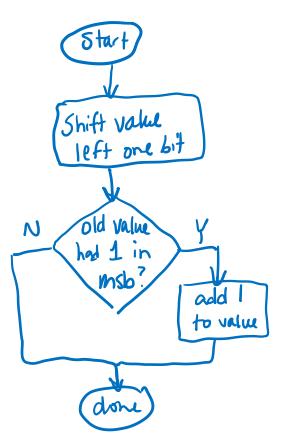
Parameters: **R0** is value Returns: **R0** is rotated result

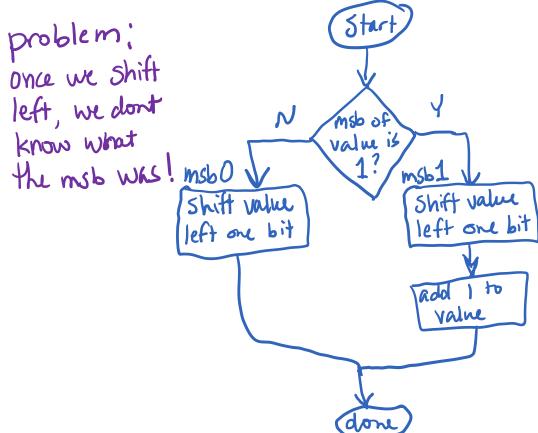




ROL1 rotates left by 1 bit

Parameters: **R0** is value Returns: **R0** is rotated result

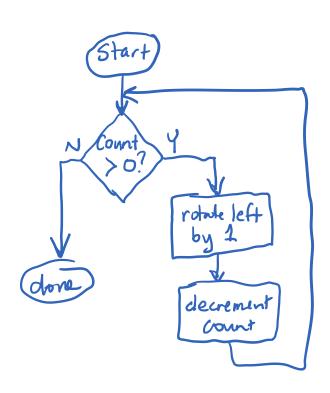






ROL rotates left by N bits

Parameters: **R0** is value, **R1** is count (N) Returns: **R0** is rotated result





ROR rotates right by N bits

Parameters: **R0** is value, **R1** is count (N) Returns: **R0** is rotated result



Wrapping Up

- Up Next:
 - I/O Concepts
 - LC-3 I/O
- Remember your videos and reading
 - Including the video quiz!

Questions?





