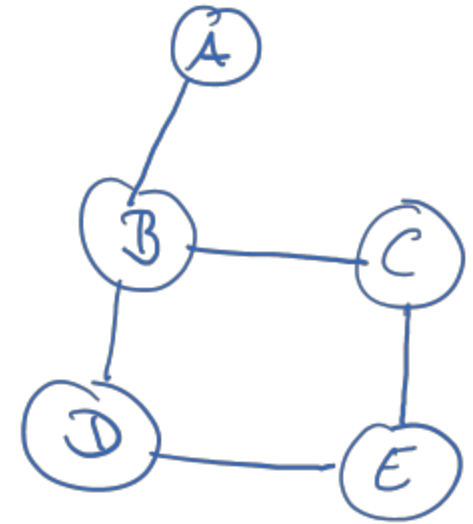


Graphs

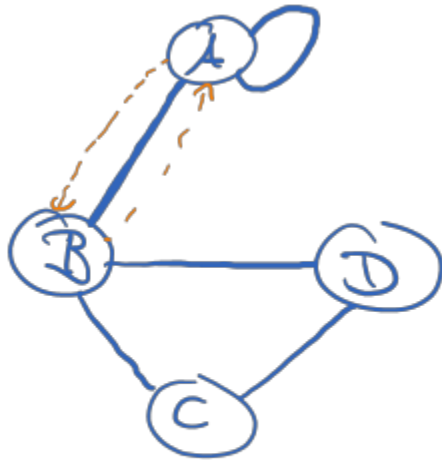
What is a Graph?

- Graphs consist of
 - Set of nodes (or vertices)
 - Set of links (or edges)
- Two nodes connected with direct edge are
 - Neighbors or adjacent nodes

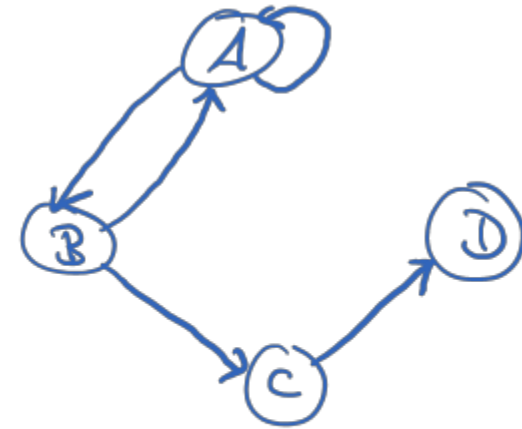


Two Types of Graphs

undirected



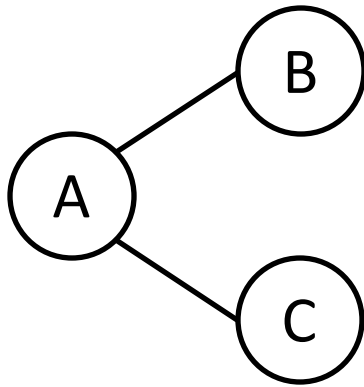
directed



Source \longrightarrow target

Degree of a Node

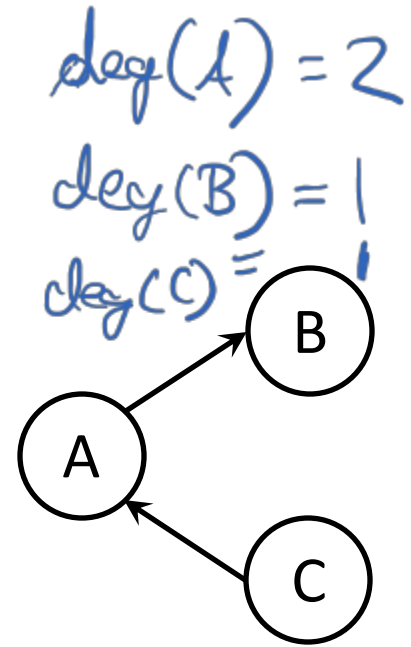
The number of edges connected to a node.



$$\deg(A) = 2$$

$$\deg(B) = 1$$

$$\deg(C) = 1$$



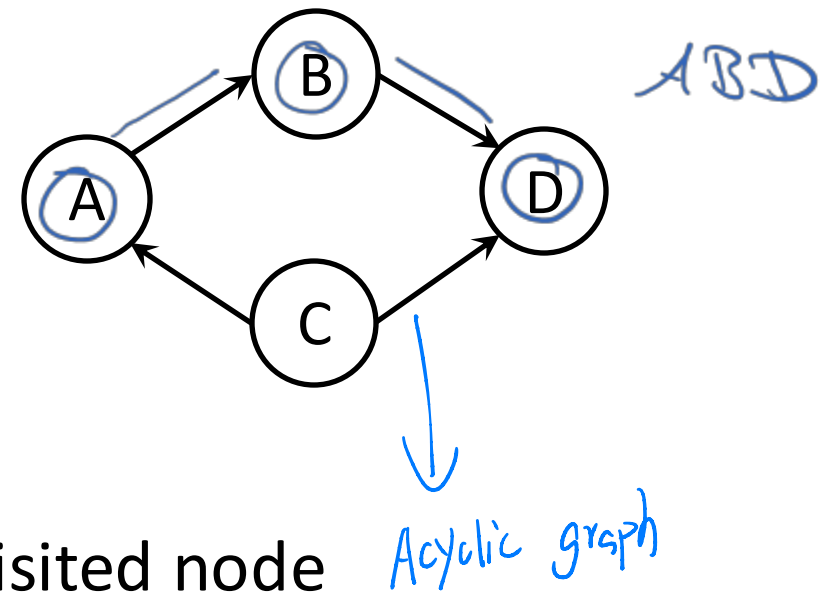
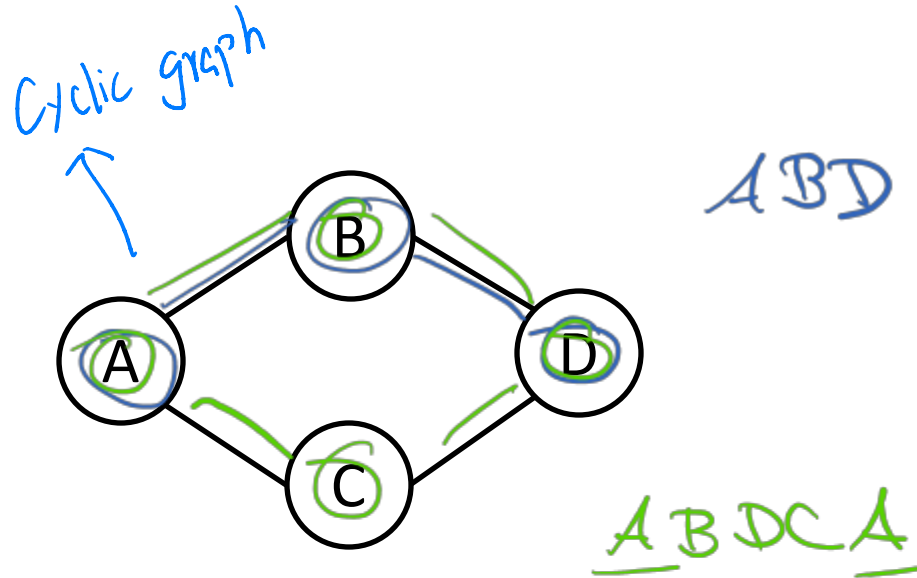
$$\deg(A) = 2$$

$$\deg(B) = 1$$

$$\deg(C) = 1$$

$\text{indeg}(A) = 1$	$\text{outdeg}(A) = 1$
$\text{indeg}(B) = 0$	$\text{outdeg}(B) = 0$
$\text{indeg}(C) = 1$	$\text{outdeg}(C) = 0$

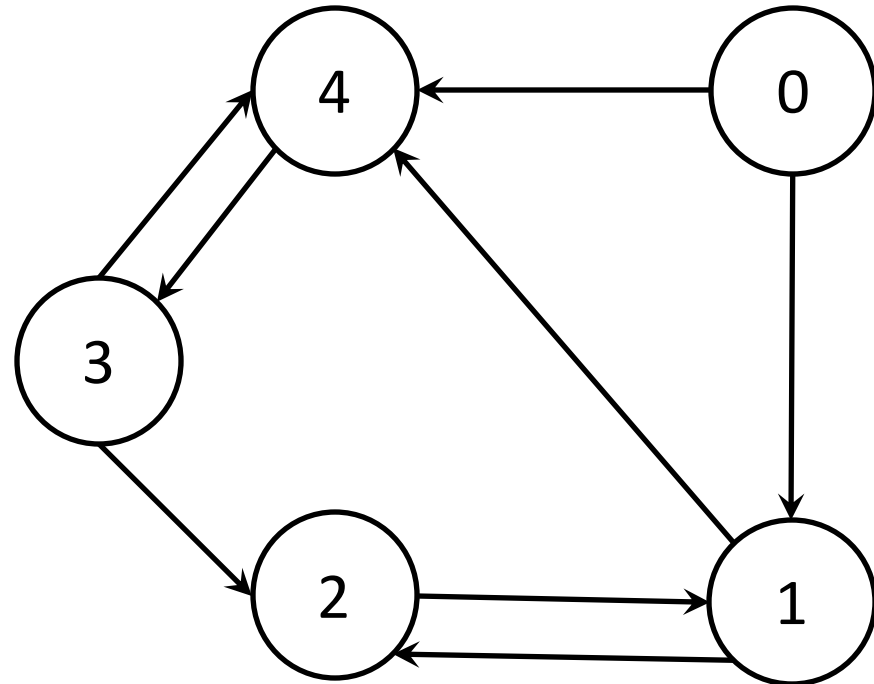
Paths in a Graph



- Cycle: A path that returns to a previously visited node
- Cyclic graph: contains at least one cycle
- Acyclic graph: contains no cycle

(★ There are 2 ways to represent the graph)

① Adjacency Matrix Example 1



rows:
source

→

→

→

→

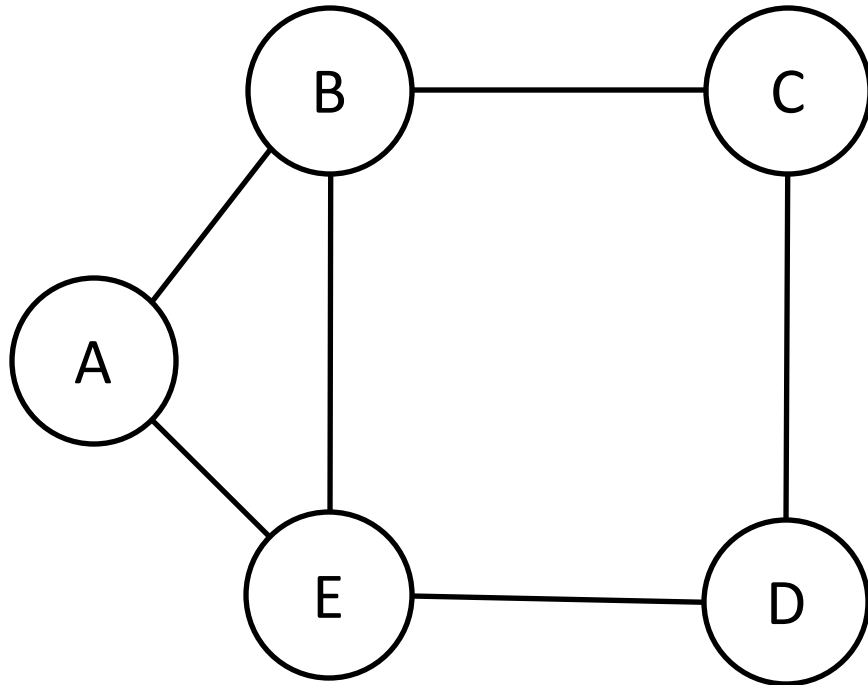
→

	0	1	2	3	4
0	F	T	F	F	T
1	F	F	T	F	T
2	F	T	F	F	F
3	F	F	T	F	T
4	F	F	F	T	F

Usually represented as 2-D array in Java

column = target

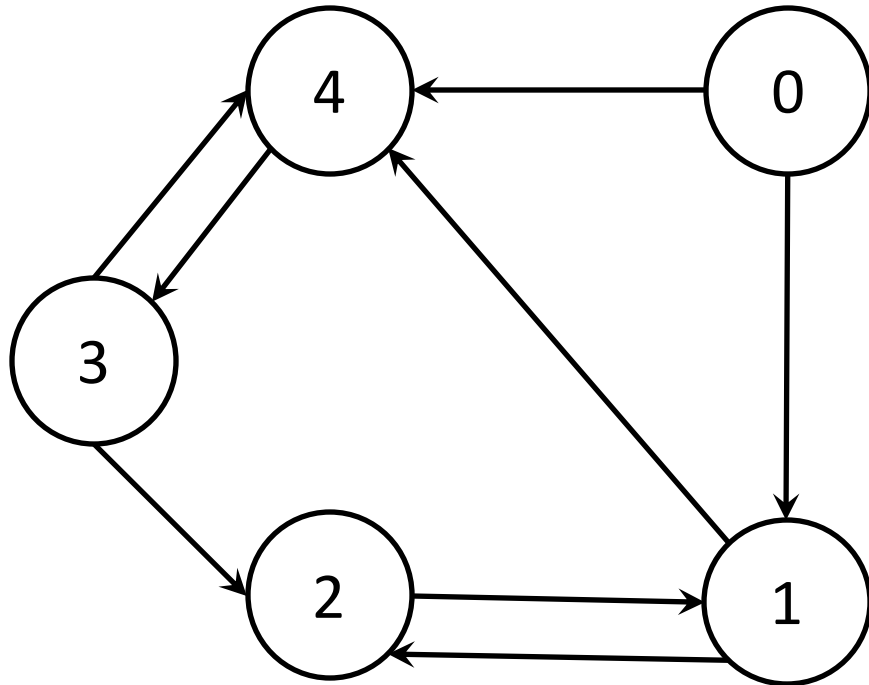
Adjacency Matrix Example 2



		A	B	C	D	E
→	A	F	T	F	F	T
→	B	T	F	T	F	T
→	C	F	T	F	T	F
→	D	F	F	T	F	T
→	E	T	T	F	T	F

undirected: symmetric at diagonal

② Adjacency List Example 1

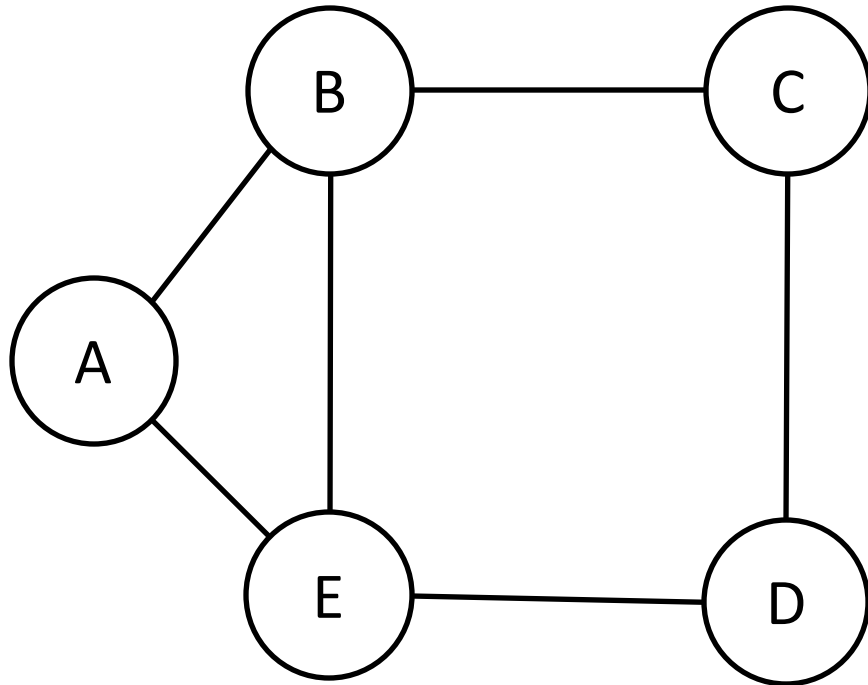


(★ List of neighbors of the selected node)

→	0:	1, 4
→	1:	2, 4
→	2:	1
→	3:	2, 4
→	4:	3

(Here, it will have 5 lists for each of the node)

Adjacency List Example 2



→ **A:** B, E
→ **B:** A, C, E
→ **C:** B, D
→ **D:** C, E
→ **E:** A, B, D

Implementation of Graphs

public class Graph <T> {

node type

boolean[][] adjacencyMatrix;
Map<T,Integer> nodeIndices;

OR

Map<T,Graphnode<T>> vertexTable;
protected class Graphnode<T> {
 protected T data;
 protected List<Graphnode<T>> adjacencyList;
}

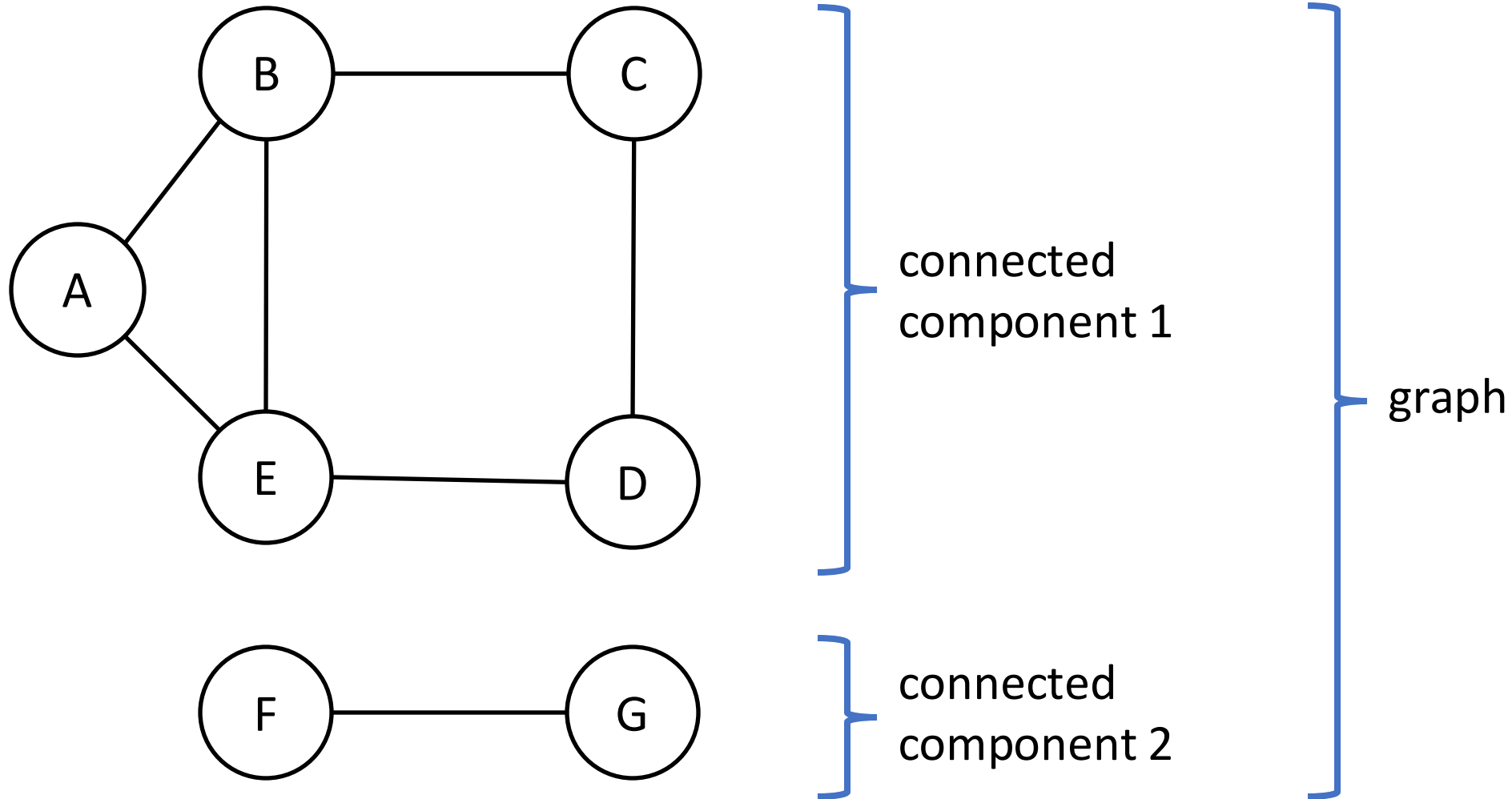
}

Graph Traversals

Graph Traversals

- ✧ • Goal: Traverse graph by crossing edges to visit each node exactly once
- Points to consider
 - ➔ • need to pick starting node
 - ➔ • nodes might be unreachable from starting node
 - ➔ • cycles can lead to infinite loops

Unreachable Nodes: Connected Components



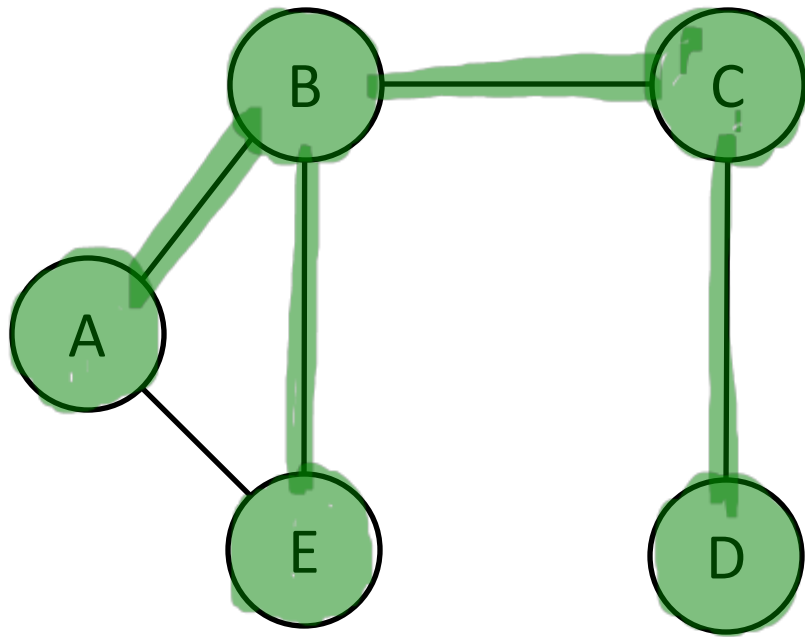
Detecting Cycles

- Detect cycles to avoid visiting nodes multiple times

- ★ Strategy: check if node is unvisited before visiting it
 - need to keep track of visited nodes:
 - either with Boolean field in node type, or
 - additional data structure to keep track of visited nodes

Depth First Traversal

Assume: all vertices are marked unvisited at start



adj. matrix

for each call to DFT:

V for loop iter.

$$\Rightarrow \underline{V * V}$$

adj. list

across V calls to DFT:

E for-loop iter.

$$\Rightarrow \underline{V + E}$$

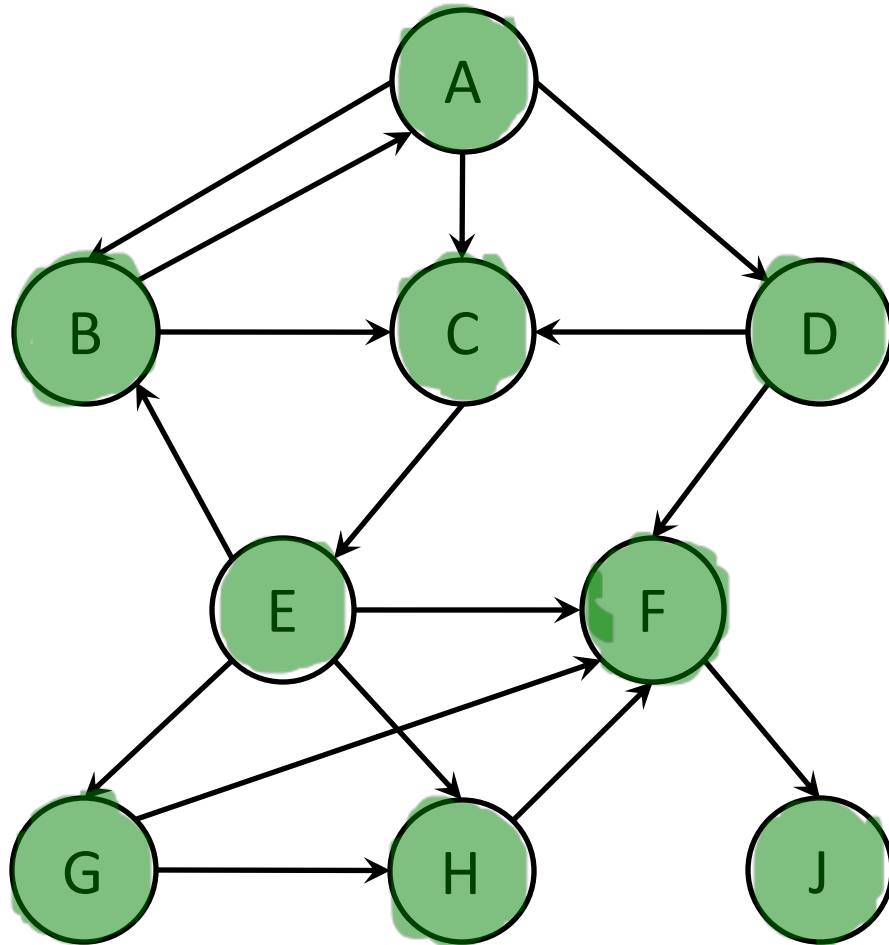
Starting node
DFT(v): \leftarrow V calls

\rightarrow mark v as visited

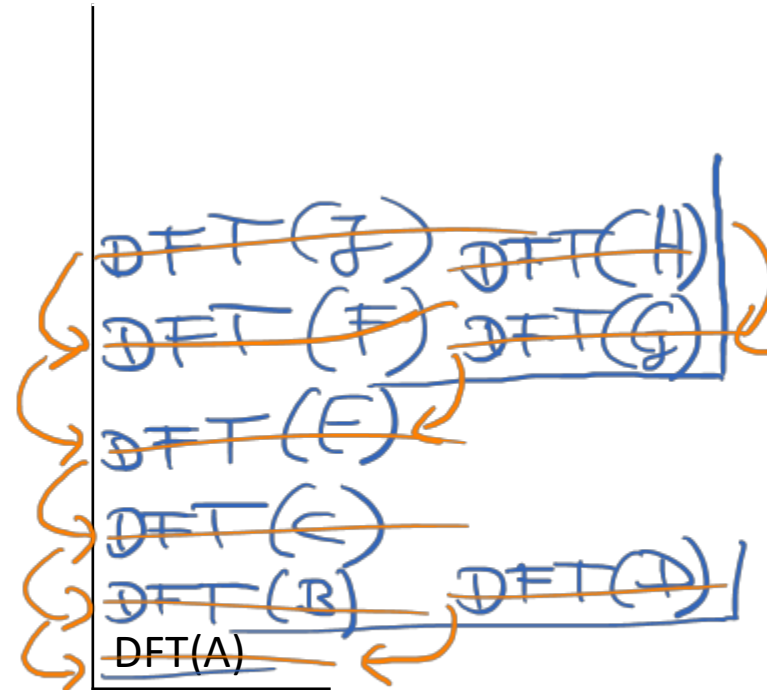
\rightarrow for each unvisited neighbor u of v:

DFT(u)

Depth First Example



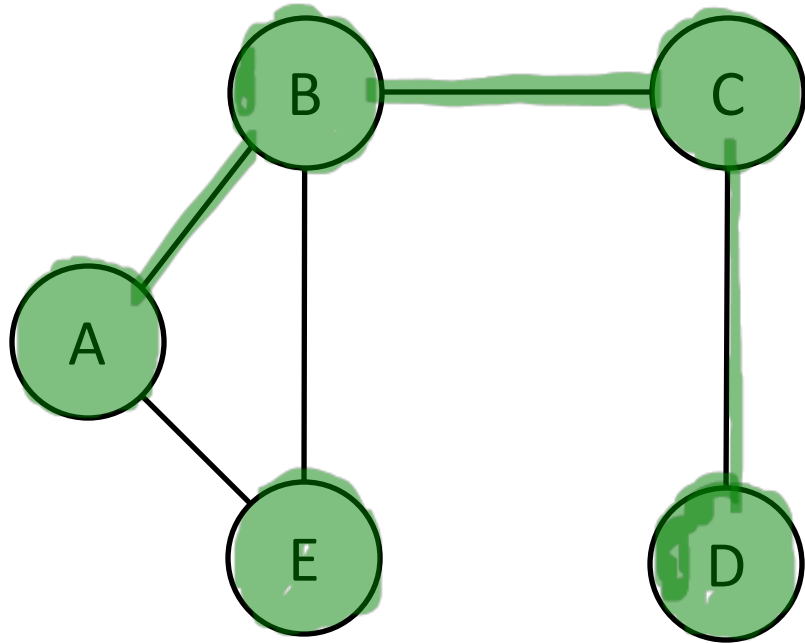
starting node: A



visit sequence: A, B, C, E, F, I, G, H, D

Breadth First Traversal

Assume: all vertices are marked unvisited at start



BFT(v):

→ q = new Queue()

→ mark v as visited

→ q.enqueue(v)

→ while (!q.isEmpty()):

→ c = q.dequeue()

→ for each unvisited neighbor u of c:

→ mark u as visited

→ q.enqueue(u)

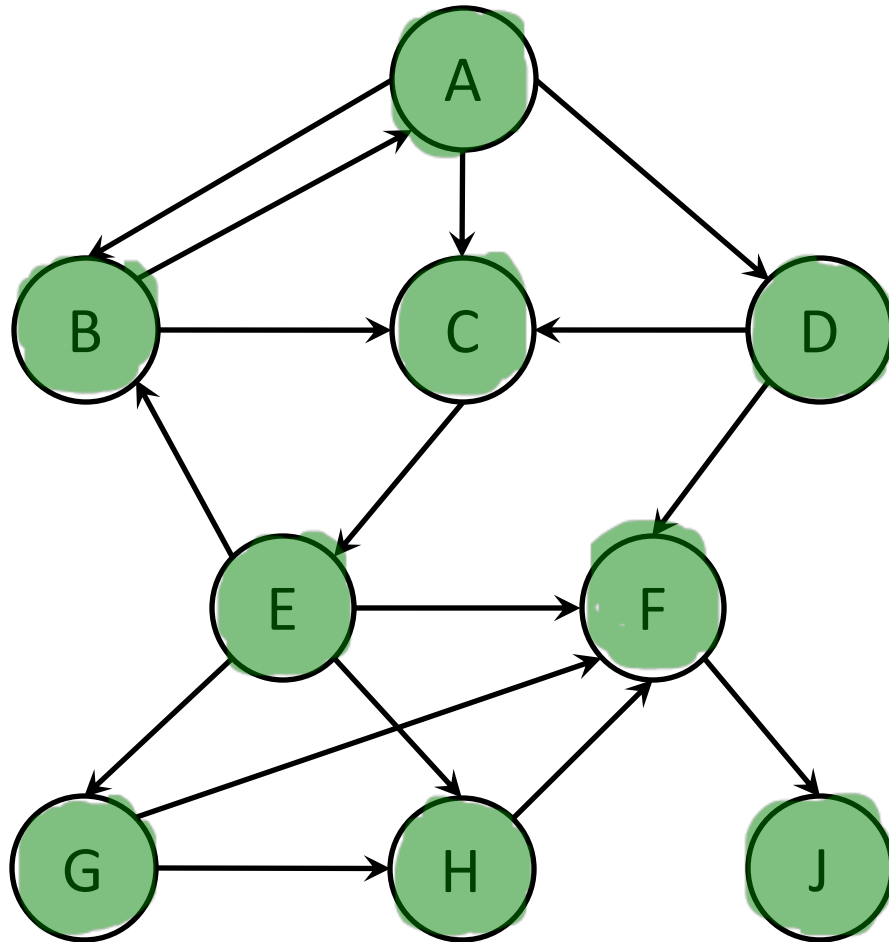
adj. matrix
for each while
iter: V for-loop iter
→ $V * V$

adj. list
across V while iter:
 E for-loop iter
⇒ $V + E$

starting node

Iterations

Breadth First Example



starting node: A

queue:



visit sequence: A, B, C, D, E, F, G, H, J

Time Complexity

V : # of nodes in graph

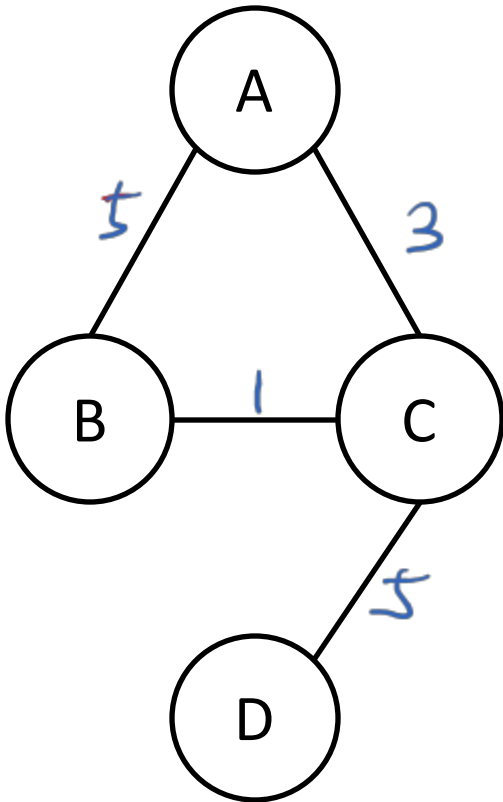
E : # of edges in graph

Depth First		Breadth First	
adjacency matrix	adjacency list	<u>adjacency matrix</u>	adjacency list
$O(V^2)$	$O(V+E)$	$O(V^2)$	$O(V+E)$

More Graph Terminology

Weighted Graphs

Weighted edges assign a cost or weight to each edge.



```
double[][] adjacencyMatrix;
```

```
List<Graphnode<T>> adjacencyList;  
List<Double> edgeWeights;
```

Cost of a Path

- For weighted graphs: sum of edge weights on path

$$A \xrightarrow{3} B \xrightarrow{4} C$$
$$\text{cost}(ABC) = 3 + 4 = 7$$

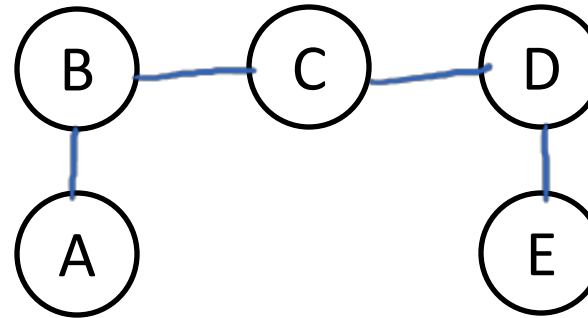
- For unweighted graphs: length of path (# of edges)

$$\text{cost}(CDE) = 2$$

Connected Graph

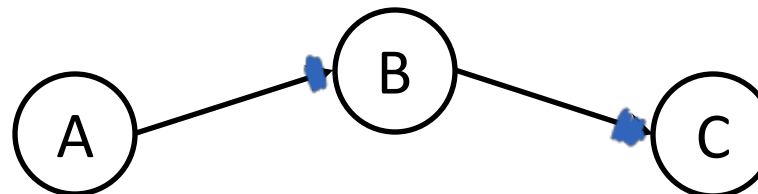
(It means you can arrive at any node from any other node)

- In a connected graph, a path exists between every pair of nodes.



- For directed graphs:

- ★ • strongly connected: a path exists between every pair of nodes with edge directions respected
- ★ • weakly connected: a path exists between every pair of nodes with edge directions ignored



Subgraphs

- G' is a subgraph of G if:
 - The set of nodes of G' is a subset of the nodes of G , and
 - The set of edges of G' is a subset of the edges of G .

