

# Inverse Zero Order Existential Paths

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The first order existential path of `f` is written  $\exists f$ . Assume that you know  $\exists f$ , what should you call `f` without referring to it by name? I call this the *zero order existential path*. This terminology is useful when you want to express relationships between functions through their existential paths.

An *inverse zero order existential path* is a function `h` that outputs only values that are not output by the zero order existential path `f`:

$$\exists h \iff \neg \exists f$$

The negation symbol  $\neg$  is interpreted literally, by flipping `true` to `false` and `false` to `true`. If `h` is the inverse zero order existential path of `f`, then `f` is the inverse zero order existential path of `h`:

$$\neg \exists h \iff \exists f$$

The second order existential path of both `f` and `h` must be  $\text{true}_1$ , since `f` halts and it must have some values it does not output, such that `h` can output them instead.

$$\begin{aligned}\exists \exists f &\iff \text{true}_1 \\ \exists \exists h &\iff \text{true}_1\end{aligned}$$

Otherwise, `h` would never halt, since there are values that are not returned, but none that are returned.

For every halting function `f`, one can prove the existence of the inverse zero existential path by showing that  $\exists \exists f \iff \text{true}_1$ . To do this it is sufficient to find two objects `A` and `B` for which:

$$\begin{aligned}\exists f(A) &= \text{true} \\ \exists f(B) &= \text{false}\end{aligned}$$

Therefore, because the requirement of two such objects existing, the smallest output type of `f` is a boolean type:

$$f : A \rightarrow \text{bool} \quad (\text{smallest type})$$

Notice that  $\exists h$  is mentioned in the existential path relationship, but never `h` itself. This means that one is allowed to change the input type of `h` arbitrary, as long the returned values are those not returned by `f`:

$$\begin{aligned}f &: A \rightarrow C \\ h &: B \rightarrow C\end{aligned}$$

If both `f` and `h` have an inverse, their domains are smaller than their output type. This is because there must be a one-to-one correspondence between input and output, and the presence of an inverse zero order existential path means the output type contains a value that is not returned by either function.