## Path Semantics in Concept Grammar Language

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## Abstract:

In this paper I formalize some aspects of deriving normal programming from path semantics using the Concept Grammar Language<sup>[1]</sup>.

Path semantics is a formal language which concept grammar is constructed from symbols `F` and `X` according to an axiom of equality:

$$\frac{F_0 X_0, F_1 X_1, F_0 = F_1, F_0 > X_0}{X_0 = X_1}$$

- `F=F`, `X=X` are weak equalities and `F>X` is a weak partial order
- "weak" means encoded and interpreted in some external language<sup>[1]</sup>
- Indices are local and chosen for a single sentence
- Matching is decided against some syntactic structure

Using the advanced index notation of Concept Grammar Language the axiom looks like this:

$$\{F_i \; X_i, \; F_j \; X_j, \; F_i \!\!=\!\! F_j, \; F_i \!\!>\!\! X_i\} \; \rightarrow \\ \{X_i \!\!=\!\! X_j\}$$
 
$$\}^{ij}$$

Normal programming languages do have path semantics in the strict sense, but not directly according to the axiom of equality. When somebody says "using path semantics" they refer to extending a language with "paths":

[F] 
$$X_1 \rightarrow (F(X_0) \rightarrow X_1)$$
 (path)  
 $F_0([F_1] X) \rightarrow [F_2] X$  (path between functions)

## References

[1] Concept Grammar Language (Sven Nilsen, 2016)