## **Probabilistic Paths**

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A probabilistic path is similar to an existential path, but for reasoning about uncertainty. For example, the following two functions  $f_0$  and  $f_1$  have the same existential path:

$$f_0 := \(x : nat) = if (x \% 2) == 0 \{ 1 \} else \{ 0 \}$$
  
 $f_1 := \(x : nat) = if (x \% 10) == 0 \{ 1 \} else \{ 0 \}$   
 $\exists f_0 <=> \exists f_1 <=> \(x : nat) = x == 0 \lor x == 1$ 

Yet, the function  $f_1$  returns 1 five times as often compared to  $f_1$ :

```
f_0: 10101010101010...

f_1: 100000000100...
```

So, their probabilistic paths are different:

```
\exists_p f_0 := \setminus (\_: nat) = if x == 0 \lor x == 1 \{ 1/2 \} else \{ 0 \}

\exists_p f_1 := \setminus (x : nat) = if x == 0 \{ 9/10 \} else if x == 1 \{ 1/10 \} else \{ 0 \}
```

All boolean functions has an assigned probability of returning `true`, by using a higher order `if` function:

$$P(if(A, B)(x)) = P(x) \cdot P(A) + (1-P(x)) \cdot P(B)$$
$$if := \langle (a, b) = \langle (c) = if \ c \ \{ a \} \ else \ \{ b \} \}$$

Just like the existential path can not inherit knowledge about the input, the probability path uses a uniform random distribution for the input. This means the probability of input P(x) gets replaced by 1/2 after deriving the probability formula for a boolean function. The probabilistic path of every boolean function can be constructed by the probability function:

```
probx := \(k : real \land [prob] true) = \(x : bool) = if x \{ k \} else \{ 1 - k \}
prob := \(x : real) = x >= 0 \land x <= 1
```

For example, for the function `and`:

$$k = P(and(x_0, x_1)) = P(x_0) \cdot P(x_1) = 1/2 \cdot 1/2 = 1/4$$
  
 $\exists_p and \le probx(1/4)$ 

In probability theory there are conditional probabilities. The equivalent concept for probabilistic paths is constructed by adding sub-type constraints to the input before taking the probabilistic path:

$$\exists_{p}$$
 and  $\{(= true), _{} \} <=> \exists_{p}$  id  $<=> probx(1/2)$ 

Translating back and forth between the two notations is straight forward:

```
P(and(a, b) | a) = (\exists_p \text{and}\{(= \text{true}), \_\})(\text{true})
P(\(\sigma \text{and}\{(= \text{true}), \_\})(false)
```

In proofs, inject the sub-type constraints so you get only free inputs and then take the probabilistic path:

```
P(and(a, b) \mid a) = (\exists_p id)(true)
```

```
\begin{split} & P(\text{and}(a,b) \mid a) \\ & P(\text{and}(a:(=\text{true}),b)) \\ & P(\text{and}\{(=\text{true}),\_\}(a,b)) \\ & (\exists_p \text{and}\{(=\text{true}),\_\})(\text{true}) \\ & \quad \text{and}\{(=\text{true}),\_\} <=> \text{id} \\ & (\exists_p \text{id})(\text{true}) \end{split}
```

## $P(\neg and(a, b) \mid a) = (\exists_p id)(false)$

```
\begin{split} &P(\neg and(a,b) \mid a) \\ &P(\neg and(a:(=true),b)) \\ &P(\neg and\{(=true),\_\}(a,b)) \\ &\quad not \cdot and <=> nand \\ &P(nand\{(=true),\_\}(a,b)) \\ &(\exists_p nand\{(=true),\_\})(true) \\ &\quad \textit{Here I do a trick by inverting the output.} \\ &(\exists_p and\{(=true),\_\})(false) \\ &\quad and\{(=true),\_\} <=> id \\ &(\exists_p id)(false) \end{split}
```