

## Problem 1

1.1

$$\text{Size} = \text{Header} + \text{name} + \text{student} + \text{birthday} + \text{state} \\ = 4 + 1 + 1 + 8 + 1 = 15 \text{ B}$$

1.2.

$$4 + 1 + 1 + 8 + 1 = 15 \text{ B}$$

↓ word-alignment

$$(4 + 1 + 1 + 2) + (8) + (1 + 7) = 24 \text{ B}$$

At page level:

$$\text{slot count} + \text{free space pointer} = 8 + 8 = 16 \text{ B}$$

For each record:

$$\text{record pointer} + \text{record length} = 8 + 8 = 16 \text{ B}$$

Totally we have:

$$16 + (24 + 16)x = 16 + 40x \quad (x \in \mathbb{N})$$

$$\therefore 16 + 40x \leq 2048 \quad (x \in \mathbb{N})$$

$$\therefore x \leq 50.8 \quad (x \in \mathbb{N})$$

$$\therefore \text{max } x = 50$$

$$\therefore \text{maximum number} = 50$$

1.3. reorder:

$$\text{birthday} + \text{Record Header} + \text{name} + \text{student} + \text{state} + \text{padding}$$

$$= (8) + (4 + 1 + (1 + 1) + 1) = 16 \text{ B}$$

$$\therefore 16 + (16 + 16)x \leq 2048$$

$$x \leq 63.5$$

$$\therefore \text{max } x = 63$$

Problem 2.

2.1

PSM: for size: 100 pages  
for credits: 100 pages  
 $\therefore$  total: 200 pages

NSM: read both size and credits: all tuples are needed  
 $\therefore$  worst cases  
 $\therefore$  500 pages

2.2.

DSM: min: 1 for id 3 for <sup>name</sup>instructor = 4 pages

max: 100 for id 2x3 for <sup>name</sup><sub>Size</sub>instructor = 106 pages

NSM: min 1 for all = 1 pages

max: worst cases: 500 pages.

## Problem 3: Bloom Filter:

3.1:

insert 7:

$$h_1(7) = 7 \quad 0.7 \bmod 10 =$$

$$h_2(7) = 0$$

set 0th and 7th to 1:

$$[1 \ 0 \ 0 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0]$$

insert 27

$$h_1(27) = 7$$

$$h_2(27) = 2$$

set 7th and 2nd to 1:

$$[1 \ 0 \ 1 \ 0 \ 0 \ 0 \ 1 \ 0 \ 0]$$

3.2

Because: Hash functions can collide: different element may map to same bits  
When number of elements increases, more bits are set to 1. Causing bit overlaps.

3.3.  
the false positive probability  $p$  for a Bloom Filter is approximately

$$p \approx (1 - e^{-\frac{kn}{m}})^k$$

where  $m$  is the size of the bit array  
 $n$  is the number of elements inserted  
 $k$  is the number of hash functions

We choose  $k = \frac{m}{n} \ln 2$

$$\Rightarrow p \approx \left(\frac{1}{2}\right)^k < 0.01$$

$$\Rightarrow k > \log_2(100) \approx 6.64 \quad k \in \mathbb{N}^+$$

$$\therefore \text{set } k=7$$

$$\therefore m = \frac{kn}{\ln 2} = \frac{7000}{\ln 2} \approx 10100$$

$\therefore$  we can choose  $m \approx 10000$  and  $k=7$