



# CS 4400

# Computer Systems

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## LECTURE 8

Array allocation and access

# Arrays in C

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- Array declaration `T A[N];`
  - allocates a contiguous region of  $L \cdot N$  bytes,  $L$  is the size of  $T$
  - introduces  $A$  as a constant pointer to the beginning of the array
- Let  $x_A$  be address stored in  $A$ , element  $i$  is stored at  $x_A + L \cdot i$ .

# Arrays in C

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- With IA32's flexible addressing modes, translation to assembly code is straightforward.

- Suppose:

```
int E[ ]
```

```
%edx = &E
```

```
%ecx = i
```

```
movl (%edx,%ecx,4),%eax
```

```
results in %eax = E[i]
```

- Optimizing compilers are particularly good at simplifying address computations, which may make assembly code hard to read

# Pointer Arithmetic

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- Computed value is scaled according to size of data type.
  - for `int* p`, expression `p+k` has value  $x_p + 4 \cdot k$
  - for `char* str`, what is the value of expression `str+j`?
- Array subscripting operation can be applied to array names and other pointers.
  - `A[i]` equivalent to `*(A+i)`

# Pointer Arithmetic

---

- Examples
  - `%edx`: address of `E`
  - `%ecx`: value of `i`
  - `%eax`: result

`E[ 2 ]`

`E+i-1`

`* ( &E[ i ] + i )`

# Pointer Arithmetic

---

- Examples
  - `%edx`: address of `E`
  - `%ecx`: value of `i`
  - `%eax`: result

`E[2]`                      `movl 8(%edx), %eax`

`E+i-1`

`*(&E[i]+i)`

# Pointer Arithmetic

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- Examples
  - `%edx`: address of `E`
  - `%ecx`: value of `i`
  - `%eax`: result

`E[2]`                      `movl 8(%edx), %eax`

`E+i-1`                    `leal -4(%edx, %ecx, 4), %eax`

`*(&E[i]+i)`

# Pointer Arithmetic

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- Examples
  - `%edx`: address of `E`
  - `%ecx`: value of `i`
  - `%eax`: result

`E[2]`                      `movl 8(%edx), %eax`

`E+i-1`                    `leal -4(%edx,%ecx,4), %eax`

`*(&E[i]+i)`              `leal (%edx,%ecx,4), %eax`  
                             `leal (%eax,%ecx,4), %eax`  
   or  
                             `leal (%edx,%ecx,8), %eax`



# A Few Pointer Arithmetic Rules

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- OK to add an integer to a pointer
- OK to subtract an integer from a pointer
- OK to subtract a pointer from a pointer
  - But only if both pointers have same type and point into the same array
  - Otherwise, behavior is undefined
  - Result is an integer
- NOT OK to add a pointer to a pointer

# Exercise: Pointer Arithmetic

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- For: short `S[ ]` and index `i`
- Let: `%edx = &S` and `%ecx = i`
- Put a pointer result in `%eax`
- Put short result in `%ax`

	Type	Value	Assembly Code
<code>S+1</code>			
<code>S[3]</code>			
<code>&amp;S[i]</code>			
<code>S[4*i+1]</code>			
<code>S+i-5</code>			

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- Put short result in `%ax`

	Type	Value	Assembly Code
<code>S+1</code>	<code>short*</code>	<code>xS+2</code>	<code>leal 2(%edx),%eax</code>
<code>S[3]</code>			
<code>&amp;S[i]</code>			
<code>S[4*i+1]</code>			
<code>S+i-5</code>			

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<code>S[3]</code>	<code>short</code>	<code>*(xS+6)</code>	<code>movw 6(%edx),%ax</code>
<code>&amp;S[i]</code>			
<code>S[4*i+1]</code>			
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# Exercise: Pointer Arithmetic

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<code>&amp;S[i]</code>	<code>short*</code>	<code>xS+2*i</code>	<code>leal (%edx,%ecx,2),%eax</code>
<code>S[4*i+1]</code>			
<code>S+i-5</code>			

# Exercise: Pointer Arithmetic

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<code>&amp;S[i]</code>	<code>short*</code>	<code>xS+2*i</code>	<code>leal (%edx,%ecx,2),%eax</code>
<code>S[4*i+1]</code>	<code>short</code>	<code>*(xS+8*i+2)</code>	<code>movw 2(%edx,%ecx,8),%ax</code>
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# Exercise: Pointer Arithmetic

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- For: short `S[ ]` and index `i`
- Let: `%edx = &S` and `%ecx = i`
- Put a pointer result in `%eax`
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<code>S+1</code>	<code>short*</code>	<code>xS+2</code>	<code>leal 2(%edx),%eax</code>
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<code>&amp;S[i]</code>	<code>short*</code>	<code>xS+2*i</code>	<code>leal (%edx,%ecx,2),%eax</code>
<code>S[4*i+1]</code>	<code>short</code>	<code>*(xS+8*i+2)</code>	<code>movw 2(%edx,%ecx,8),%ax</code>
<code>S+i-5</code>	<code>short*</code>	<code>xS+2*i-10</code>	<code>leal 10(%edx,%ecx,2),%eax</code>

# Question

---

Suppose we have declared `int arr[N]`.

Are the following is equivalent to the reference `arr[i]`?

A. `*(arr + 4 * i)`

B. `*(&arr[0] + i)`

C. `*((int*)((char*)arr + 4 * i))`



# Question

---

Suppose we have declared `char* arr[N]`.

Which of the following correctly puts `arr[i]` in `%eax`? (Suppose that `arr` in `%edx` and `i` in `%eax`.)

- A. `leal (%edx,%eax), %eax`
- B. `leal (%edx,%eax,4), %eax`
- C. `movl (%edx,%eax), %eax`
- D. `movl (%edx,%eax,4), %eax`
- E. none of the above

# Arrays and Loops

---

- Array references in loops often have very regular patterns.

```
for(i = 0, val = 0; i < 5; i++)  
    val = (10 * val) + x[i];
```

- For efficiency, optimizing compilers exploit these patterns.

```
xorl %eax,%eax           ;val=0  
leal 16(%ecx),%ebx       ;xend=x+4  
.L12:  
    leal (%eax,%eax,4),%edx ;compute 5*val  
    movl (%ecx),%eax       ;compute *x  
    leal (%eax,%edx,2),%eax ;compute *x+2*(5*val)  
    addl $4,%ecx           ;x++  
    cmpl %ebx,%ecx        ;compare x-xend  
    jbe .L12              ;if x<=xend, goto loop
```

- Uses pointer arithmetic instead of loop index `i`.

```
int* xend = x + 4;  
do {  
    val = (10 * val) + *x;  
} while(++x <= xend);
```

# Multi-Dimensional Arrays

---

- The same principles hold for arrays of arrays.
  - `int A[4][3];` is an array of four 3-integer arrays (“rows”)
  - arrays are linearized in memory in row-major order
- `A[i][j]` is at memory address  $\text{xA} + \text{L} \cdot (\text{C} \cdot i + j)$ .
- Example (`%eax`: address of `A`, `%edx`: value of `i`, `%ecx`: value of `j`)

```
sall $2,%ecx ; j*4
```

```
leal (%edx,%edx,2),%edx ; i*3
```

```
leal (%ecx,%edx,4),%edx ; j*4 + i*12
```

```
movl (%eax,%edx),%eax ; read A[i][j]
```

```

#define N 16
typedef int fix_matrix[N][N];    fixed-sized array

int fix_prod(fix_matrix A, fix_matrix B, int i, int
k) {
    int j, result;

    for(j = 0, result = 0; j < N; j++)
        result += A[i][j] * B[j][k];

    return result;
}

```

```

int fix_prod(fix_matrix A, fix_matrix B,
              int i, int k) {

    int *Aptr, *Bptr, cnt, result;
    Aptr = &A[i][0];
    Bptr = &B[0][k];
    cnt = N-1;
    result = 0;

    do {
        result += (*Aptr) * (*Bptr);
        Aptr++;
        Bptr += N;
        cnt--;
    } while(cnt >= 0);

    return result;
}

```

*compiler optimizations*

Aptr is in %edx  
 Bptr is in %ecx  
 result is in %esi  
 cnt is in %ebx

```

.L23:
    movl (%edx),%eax
    imull (%ecx),%eax
    addl %eax,%esi
    addl $64,%ecx
    addl $4,%edx
    decl %ebx
    jns .L23

```

# Exercise: Nested Arrays

```
#define M ??  
#define N ??  
  
int mat1[M][N];  
int mat2[N][M];  
  
int sum_element(int i, int j) {  
    return mat1[i][j] + mat2[j][i];  
}
```

```
movl 8(%ebp),%ecx  
movl 12(%ebp),%eax  
leal 0(,%eax,4),%ebx  
leal 0(,%ecx,8),%edx  
subl %ecx,%edx  
addl %ebx,%eax  
sall $2,%eax  
movl mat2(%eax,%ecx,4),%eax  
addl mat1(%ebx,%edx,4),%eax
```

# Question

---

Will the following compile successfully?

```
#define N 100
int foo(int arr[][N], int i, int j) {
    return arr[i][j];
}
```

# Dynamic Memory Allocation

---

- For allocation of memory at run time, library routine `malloc` is used.
  - arguments specify number of bytes to be allocated
  - return value is a pointer to the allocated memory or `NULL`
- `malloc` allocates one contiguous block (of specified size).

```
NODE* head = malloc(sizeof(NODE)); // implicit  
head->next = malloc(sizeof(NODE)); // cast
```

- To release dynamically-allocated memory, the library routine `free` is used.
    - argument is the pointer to the block of memory to be released
- ```
free(ptr);
```

# Question

---

Suppose we have

```
short* arr = malloc(user_input*sizeof(short));
```

Which of the following references the second element?

A. `arr[1]`

B. `*(arr+1)`

C. `*(arr+2)`

D. exactly 2 of the above

E. all of A-C

F. none of A-C



# Question

---

Suppose we have

```
short* matrix = malloc(N*N*sizeof(short));
```

Which of the following references the element in the second row and second column?

A. `arr[1]`

B. `arr[N]`

C. `arr[N+1]`

D. `arr[1][1]`

E. none of the above

# Exercise: Compiler Optimizations

```
#define N 16
typedef int fix_matrix[N][N];

void fix_set_diag(fix_matrix A,
                 int val) {
    int i;
    for (i = 0; i < N; i++)
        A[i][i] = val;
}
```

```
movl 8(%ebp),%ecx
movl 12(%ebp),%edx
movl $0,%eax
.L14:
movl %edx, (%ecx,%eax)
addl $68,%eax
cmpl $1088,%eax
jne .L14
```

- Write a function `fix_set_diag_opt` that uses optimizations similar to those in the assembly code. Do not assume that `N` is 16.