

Implementing an Efficient Shuffle Operator for Streaming Database Systems

Bachelor Thesis

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Problem Setting

Streaming Shuffle Simulation:

- 1. **Tuple Generation:** Randomly generated tuples with 32-bit keys and optional data fields.
- 2. **Data Shuffle:** Tuples stored in partition buckets using slotted pages.
- 3. Storing on Slotted Pages: Thread-local vs. shared (locking/lock-free) write-out strategies.

Header		Slot 1		Slot 2	Slot 3
Slot 4	Slot 5	Slots	grov	v o	
← Data grow			Data 5		
Data 4			Data 3		
Data 2			Data 1		

Figure: Slotted Page grow visualization

Key Contribution: Efficient, multithreaded shuffle operator implementations.



Naive approach: OnDemand

OnDemand:

• Tuples are directly written to the partition buckets.

Problems:

- Each tuples causes a write-out to a shared slotted page.
- Very high contention on partition buckets.



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Optimized approach: Smb

Software Managed Buffers (SMBs):

- Cacheline-sized, thread-local buffers for each partition.
- Flush partition when buffer is full.

Problems:

High contention on partition buckets.



Histogram-based approach: Radix

Materialize large chunk of incoming tuples:

- Assign each thread a memory region within this chunk.
- Memory region can be reused for next chunk.

Local Histogram:

Create thread-local histogram (frequency map) for each partition.

Global Histogram:

- Merge local histograms into a global histogram.
- Pre-allocate slotted pages based on global histogram.
- Assign exclusive write-out locations based on local histograms.

Problems:

Synchronization overhead during histogram merging and page allocation.



Histogram-based approach: Hybrid

Materialize small chunk of incoming tuples:

- Assign each thread a memory region within this chunk.
- Memory region can be reused for next chunk.

Local Histogram:

Create thread-local histogram (frequency map) for each partition.

Ad-hoc page allocation:

- Assign exclusive write-out locations based on local histograms.
- Allocate slotted page for local histogram if not already allocated.

Problems:

Synchronization overhead during page allocation.



Reducing contention: CmpProcessingUnits

Processing Units:

- Partition threads into *Processing Units*
- Within each Processing Unit:
 - Each thread is assigned an exclusive partition range
 - No overlap between partition ranges
 - Only a single thread writes to a partition

Problems:

Each tuple must be processed by all threads of a Processing Unit



Avoiding contention: LocalPagesAndMerge

Thread-local Pages:

- Each thread has its own slotted pages
- Avoids synchronization

Merging Phase:

- All non-full pages have to be merged
- We assign each thread a group of partitions to merge
- Each thread merges the pages of its assigned partitions without synchronization

Problems:

Huge initial memory consumption



Evaluation Die beiden Grafiken werden noch aktualisert (9 statt 12 Impl)

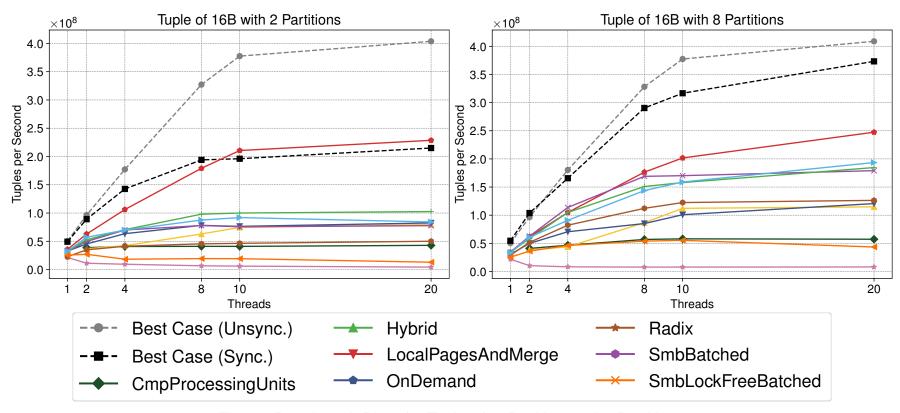


Figure: Benchmark Plots for Tuple of 16B with 2 and 8 Partitions



Evaluation

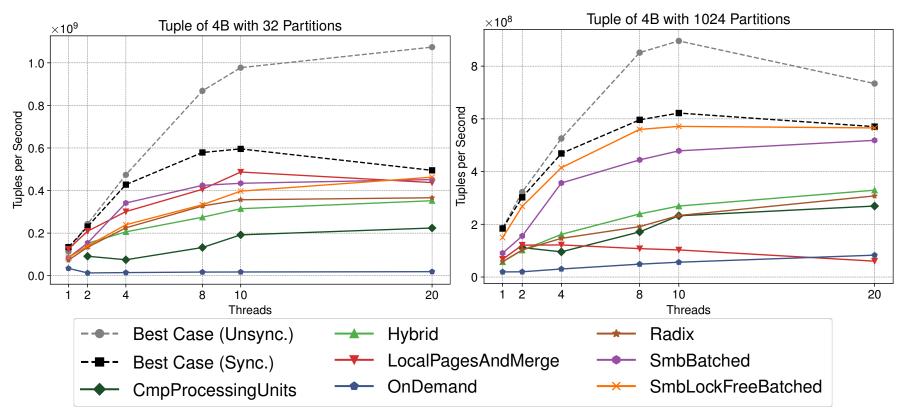


Figure: Benchmark Plots for Tuple of 4B with 32 and 1024 Partitions



Evaluation

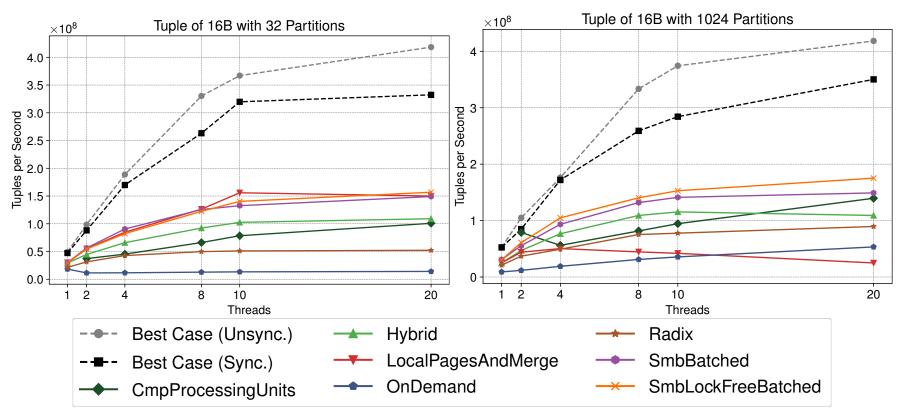


Figure: Benchmark Plots for Tuple of 16B with 32 and 1024 Partitions



Evaluation

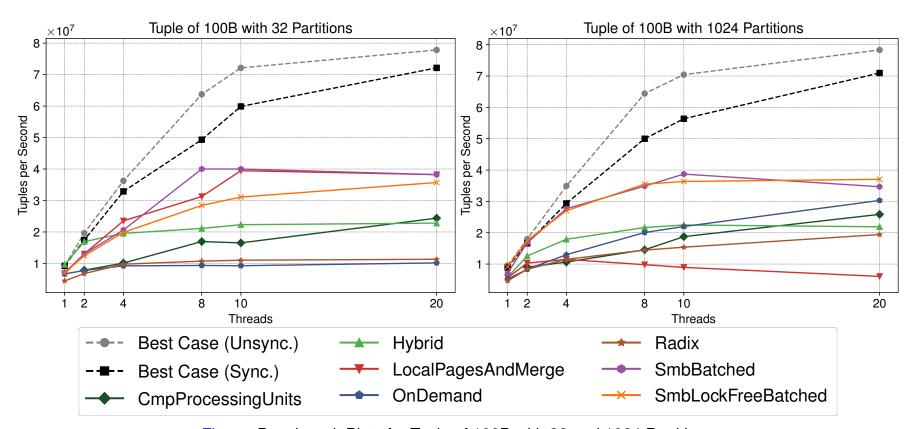


Figure: Benchmark Plots for Tuple of 100B with 32 and 1024 Partitions



Peak Heap Memory

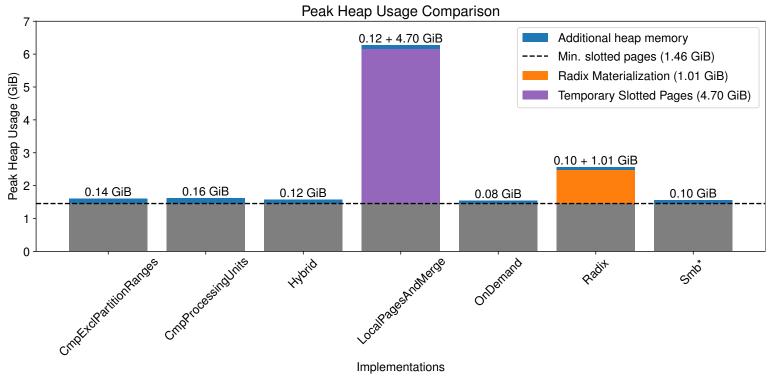


Figure: Peak Heap Usage when using 32 Partitions, 40 Threads and 67.2 Mio. 16B Tuples (1 GiB)



Comparison with Apache Flink

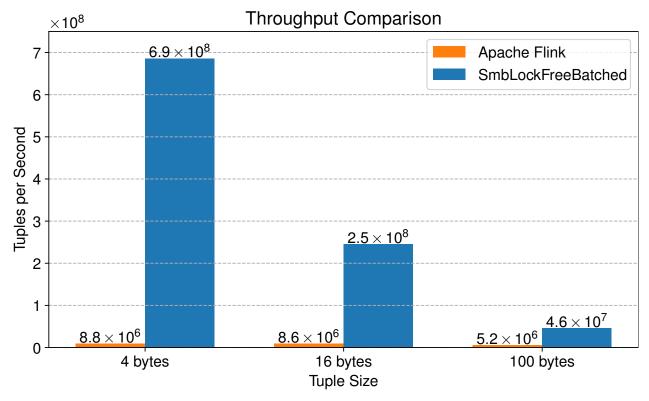


Figure: Tuples per Second Comparison when using 1024 Partitions



TODO: 2 und 8 Partitionen auf 128 Threads Server



Conclusion

- Thread-local slotted pages are optimal for low partition counts (<32).
- SMB-based methods scale best for high partition counts (>32).

Future Work

- Further reducing contention on machines with 20+ cores.
- Evaluate the implementations in a real-world streaming system.