

# Implementing an Efficient Shuffle Operator for Streaming Database Systems

**Bachelor Thesis** 

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Garching, 11.03.2025
Technical University of Munich



### **Problem Setting**

#### **Streaming Shuffle Simulation:**

- 1. **Tuple Generation:** Randomly generated tuples with 32-bit keys and optional data fields.
- 2. **Data Shuffle:** Tuples stored in partition buckets using slotted pages.
- 3. Storing on Slotted Pages: Thread-local vs. shared (locking/lock-free) write-out strategies.

Header		Slot 1		Slot 2	Slot 3
Slot 4	Slot 5	Slots	grov	v  o	
← Data grow			Data 5		
Data 4			Data 3		
Data 2			Data 1		

Figure: Slotted Page grow visualization

**Key Contribution:** Efficient, multithreaded shuffle operator implementations.



# Naive approach: OnDemand

#### **OnDemand:**

• Tuples are directly written to the partition buckets.

#### **Problems:**

- Each tuples causes a write-out to a shared slotted page.
- Very high contention on partition buckets.



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# Optimized approach: Smb

#### **Software Managed Buffers (SMBs):**

- Cacheline-sized, thread-local buffers for each partition.
- Flush partition when buffer is full.

#### **Problems:**

High contention on partition buckets.



# Histogram-based approach: Radix

#### **Materialize large chunk of incoming tuples:**

- Assign each thread a memory region within this chunk.
- Memory region can be reused for next chunk.

#### **Local Histogram:**

Create thread-local histogram (frequency map) for each partition.

#### **Global Histogram:**

- Merge local histograms into a global histogram.
- Pre-allocate slotted pages based on global histogram.
- Assign exclusive write-out locations based on local histograms.

#### **Problems:**

Synchronization overhead during histogram merging and page allocation.



# Histogram-based approach: Hybrid

#### Materialize small chunk of incoming tuples:

- Assign each thread a memory region within this chunk.
- Memory region can be reused for next chunk.

#### **Local Histogram:**

Create thread-local histogram (frequency map) for each partition.

#### Ad-hoc page allocation:

- Assign exclusive write-out locations based on local histograms.
- Allocate slotted page for local histogram if not already allocated.

#### **Problems:**

Synchronization overhead during page allocation.



# Reducing contention: CmpProcessingUnits

#### **Processing Units:**

- Partition threads into *Processing Units*
- Within each Processing Unit:
  - Each thread is assigned an exclusive partition range
  - No overlap between partition ranges
  - Only a single thread writes to a partition

#### **Problems:**

Each tuple must be processed by all threads of a Processing Unit



# Avoiding contention: LocalPagesAndMerge

#### **Thread-local Pages:**

- Each thread has its own slotted pages
- Avoids synchronization

#### **Merging Phase:**

- All non-full pages have to be merged
- We assign each thread a group of partitions to merge
- Each thread merges the pages of its assigned partitions without synchronization

#### **Problems:**

Huge initial memory consumption



# Evaluation Die beiden Grafiken werden noch aktualisert (9 statt 12 Impl)

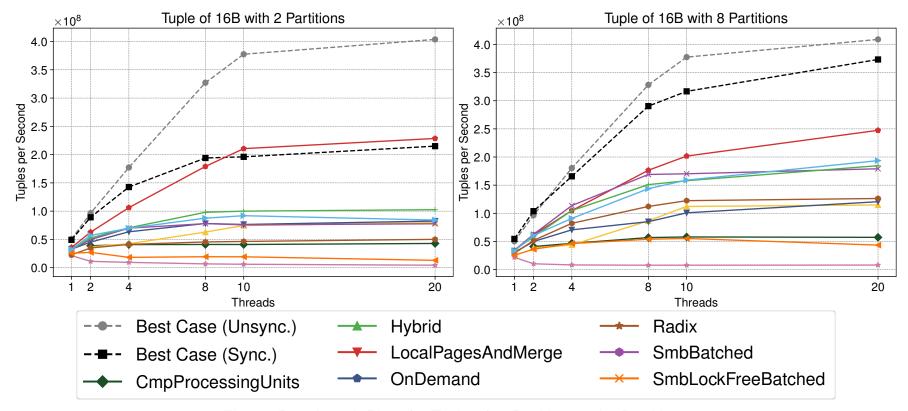


Figure: Benchmark Plots for Tuple of 16B with 2 and 8 Partitions



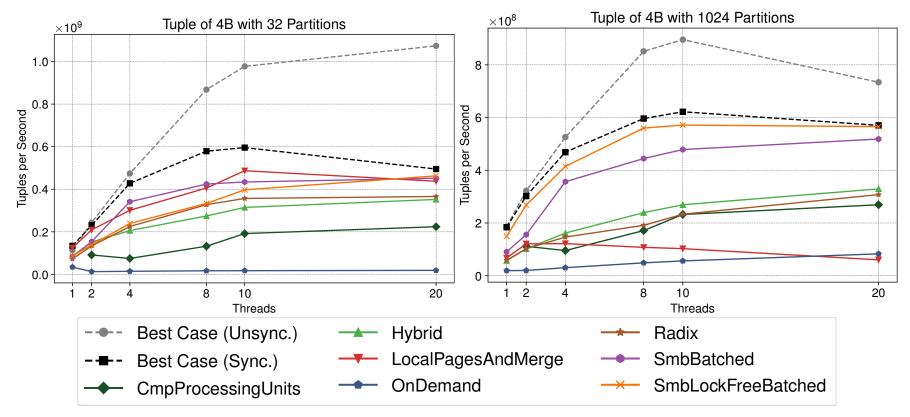


Figure: Benchmark Plots for Tuple of 4B with 32 and 1024 Partitions



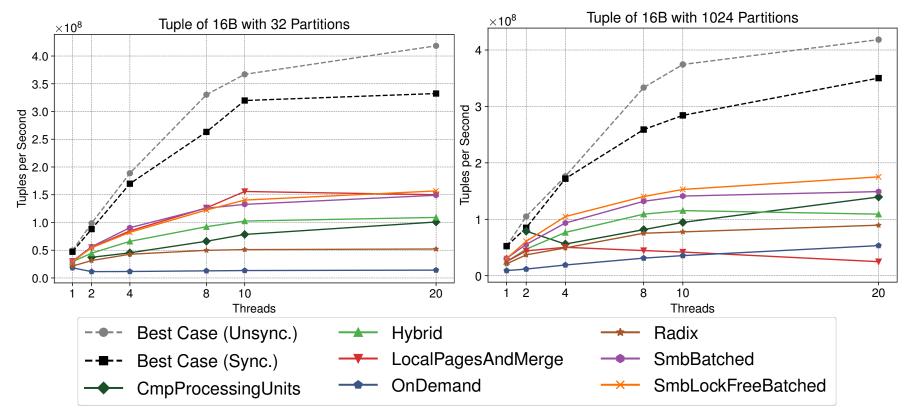


Figure: Benchmark Plots for Tuple of 16B with 32 and 1024 Partitions



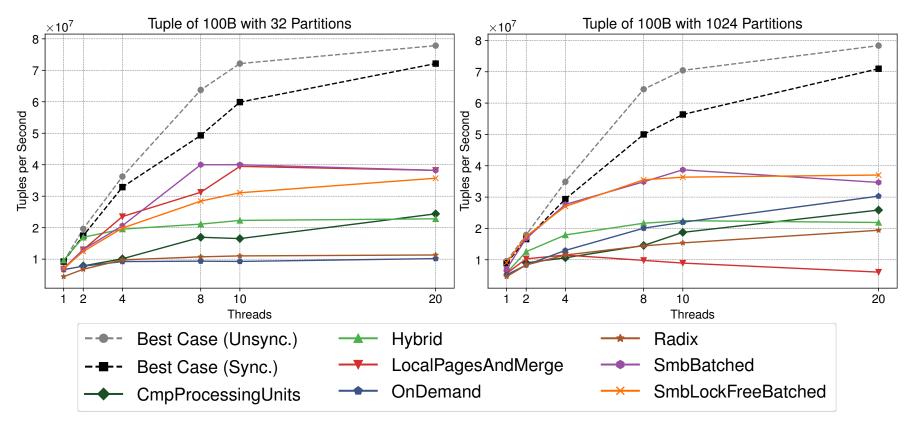


Figure: Benchmark Plots for Tuple of 100B with 32 and 1024 Partitions



### Peak Heap Memory

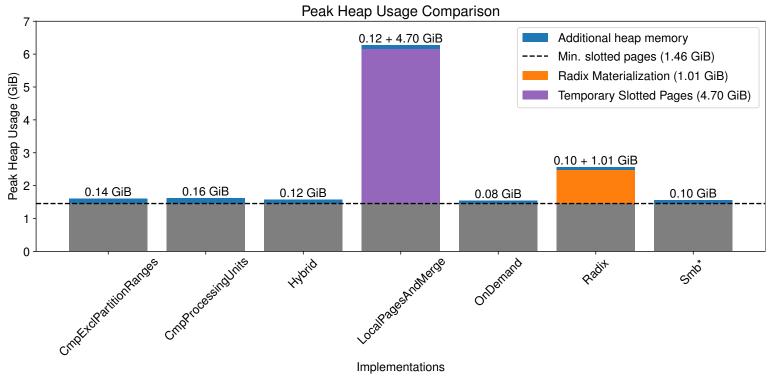


Figure: Peak Heap Usage when using 32 Partitions, 40 Threads and 67.2 Mio. 16B Tuples (1 GiB)



# Comparison with Apache Flink

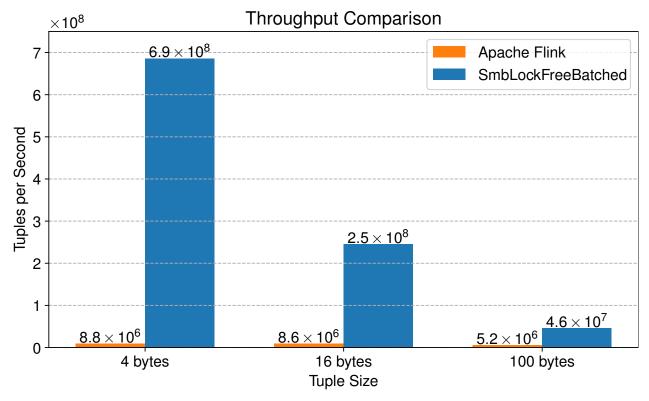


Figure: Tuples per Second Comparison when using 1024 Partitions



AMD EPYC 7713, 2x64 threads, 2 NUMA nodes, 1 TB memory

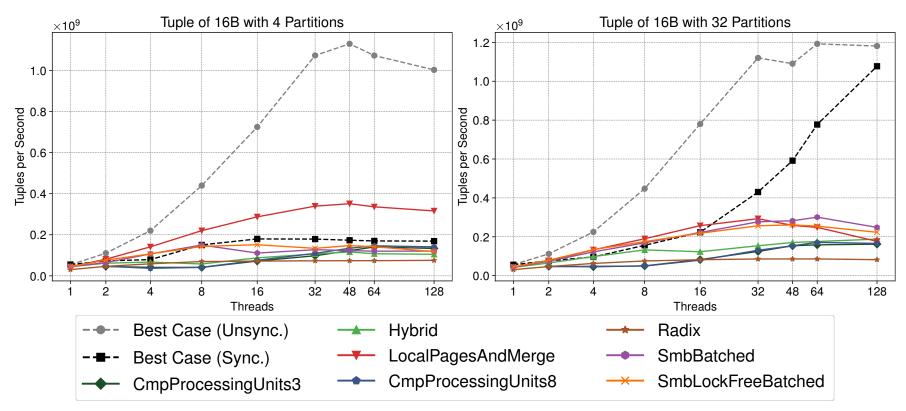


Figure: Benchmark Plots for Tuple of 16B with 32 and 1024 Partitions



### Conclusion

- Thread-local slotted pages are optimal for low partition counts (<32).
- SMB-based methods scale best for high partition counts (>32).

### **Future Work**

- Further reducing contention on machines with 20+ cores.
- Evaluate the implementations in a real-world streaming system.
- Investigate the impact of data skew.