# 2014-05-19.sagews

#### May 19, 2014

## Contents

1	Mat	th 480b Sage Course	1
	1.1	Linear Algebra, part 2	1
	1.2	May 19, 2014	1
	1.3	Vector spaces	1
	1.4	Linear algebra over finite fields (very important for coding theory)	2
	1.5	Remarks about asymptotically fast algorithms	4

## 1 Math 480b Sage Course

## 1.1 Linear Algebra, part 2

#### 1.2 May 19, 2014

Screencast: REMIND ME. And watch to see if it crashes.

- Questions
- Homework:
  - hw7, etc., collected this morning, and re-destributed for grading
  - hw8 assigned (should find it in your project). This is the last homework assignment.
- Topic: Exact linear algebra, part 2
  - vector spaces
  - linear algebra over finite fields and coding theory
  - remarks about asymptotically fast algorithms
- Wednesday and Friday: Graph theory, Group Theory

#### 1.3 Vector spaces

```
# The vector space of all 3-tuples of rational numbers (i.e., vectors in \
   3-space with tail at the origin)
V = QQ^3
Vector space of dimension 3 over Rational Field
# These arise natural as spans, kernels (=nullspaces), etc.
m = matrix(QQ, 2,3, [2,3,5, 7,-4,0]); m
[2 3 5]
[7-40]
# Compute the vector space of vector x such that m*x = 0
V = m.right_kernel(); V
Vector space of degree 3 and dimension 1 over Rational Field
Basis matrix:
    1 7/4 -29/20]
[
<class 'sage.modules.free_module.FreeModule_submodule_field_with_category'>
V.dimension()
V.basis()
(1, 7/4, -29/20)
1
# compute another 1-dimensional vector space
m = matrix(QQ, 2,3, [1,2,3,4,5,6]); m
W = m.right_kernel(); W
[1 2 3]
[4 5 6]
Vector space of degree 3 and dimension 1 over Rational Field
Basis matrix:
[ 1 -2 1]
V.intersection(W)
Vector space of degree 3 and dimension 0 over Rational Field
Basis matrix:
Г٦
Vector space of degree 3 and dimension 2 over Rational Field
Basis matrix:
   1
          0 -23/75]
1 -49/75]
```

### 1.4 Linear algebra over finite fields (very important for coding theory)

```
# define a finite field
F = GF(7)
F
list(F) # the elements of F
Finite Field of size 7
[0, 1, 2, 3, 4, 5, 6]
# define a matrix and vector over F
m = matrix(F, 3,3, [2,3,5,
                             7,-4,0,
                                        2,-5,1]); m
v = vector(F, [10,5,2]); v
# notice how 7 == 0 below, since we are working in F.
[2 3 5]
[0 3 0]
[2 2 1]
(3, 5, 2)
# solve system
x = m.solve_right(v); x
(0, 4, 1)
m * x
(3, 5, 2)
```

In fact, Sage has extensive coding theory functionality. (See http://www.sagemath.org/doc/reference/coding/sage/coding/code\_constructions.html and http://www.sagemath.org/doc/reference/coding/index.html)

- A code is a subspace of a finite dimensional vector space.
- One encodes messages as elements of this subspace.
- When a message is corrupted (say one bit flipped) it becomes something not in the subspace.
- Decoding involves finding the closest vector in the subspace to what you get.

```
C = codes.HammingCode(3,GF(2)); C
Linear code of length 7, dimension 4 over Finite Field of size 2

C.basis()
[(1, 0, 0, 0, 0, 1, 1), (0, 1, 0, 0, 1, 0, 1), (0, 0, 1, 0, 1, 1, 0), (0, 0, 0, 1, 1, 1, 1)]

span(C.basis())
Vector space of degree 7 and dimension 4 over Finite Field of size 2
Basis matrix:
[1 0 0 0 0 1 1]
[0 1 0 0 1 0 1]
[0 0 1 0 1 1 0]
[0 0 0 0 1 1 1 1]
```

```
len(C)
16
for v in C:
    print v
(0, 0, 0, 0, 0, 0, 0)
(1, 0, 0, 0, 0, 1, 1)
(0, 1, 0, 0, 1, 0, 1)
(1, 1, 0, 0, 1, 1, 0)
(0, 0, 1, 0, 1, 1, 0)
(1, 0, 1, 0, 1, 0, 1)
(0, 1, 1, 0, 0, 1, 1)
(1, 1, 1, 0, 0, 0, 0)
(0, 0, 0, 1, 1, 1, 1)
(1, 0, 0, 1, 1, 0, 0)
(0, 1, 0, 1, 0, 1, 0)
(1, 1, 0, 1, 0, 0, 1)
(0, 0, 1, 1, 0, 0, 1)
(1, 0, 1, 1, 0, 1, 0)
(0, 1, 1, 1, 1, 0, 0)
(1, 1, 1, 1, 1, 1, 1)
# a corrupted message
corrupted_message = [1, 1, 0, 0, 0, 1, 1]
# check this out:
C.decode(corrupted_message)
(1, 0, 0, 0, 0, 1, 1)
```

#### 1.5 Remarks about asymptotically fast algorithms

- All the problems I showed you above are trivial and you could do them by hand.
- One of the key things that distinguishes Sage from certain other famous (or not) programs is that it implements many asymptotically fast algorithms for exact linear algebra, i.e., these algorithms work even if the matrices are a bit bigger. (Tell Alan Steels store about him getting money from Knuth for proving with Magma in the 90sthat asymptotically fast algorithms are practical, which Knuth said in his book they arent.)
- Some examples to get a sense of speed and capabilities.

```
m = random_matrix(ZZ, 100)
m[0] # Oth row of our 100x100 matrix
(4, 0, -2, -3, -2, 0, -3, 0, -6, -1, -1, 0, -1, -1, 1, 7, 1, -1, 1, -1, -1, -1, 2, -2, 0, 1, 0, 0, -1, -1, 1, -1, -10, -1, -1, 0, 0, -1, 6, -1, 1, 1, -16, 0, 2, 0, 0, -1, -6, 3, -6, 33, 0, -2, -1, -1, 10, 1, -1, -7, 1, -3, 5, -1, -1, 58, -1, -71, 1, -53, 1, 1, 3, -1, -1, 0, 1, -2, 1, 2, -3, 7, -1, 1, 1, 0, 2, -2, 0, -3, 0, -2, 0, 0, 4, 26, -1, -1, 7, 3)
# LIE!!!
%timeit m.det()
625 loops, best of 3: 208 ns per loop
```

Note, the above 208ns is a very misleading. The reason is because m.det() caches the result of the computation.

And the time it command takes the best of 3 the first time is long, and the others are short.

You can use m.\_clear\_cache() to delete everything from this cache.

```
# very fast
%time m.det()
```

 $749963331861405888818547429151919773237672392143832019032854540604538053353704554622059203\\302906647014042800194447461657550709663669037247751657423676451189168714574332072236654488\\74535011757119353$ 

CPU time: 0.00 s, Wall time: 0.00 s

```
%timeit m._clear_cache(); m.det()
5 loops, best of 3: 98.8 ms per loop
```

```
m = random_matrix(ZZ, 200)
%time m.det()
```

 $-64669184758437239663183673979917197112110496549640403795057967741844317344472132082700872\\154668833160552494368323839338066085287883221575176676490342951367404623495306652944781582\\441536472093846800309570810483796344473084572291617332841990161013300998240186468559391490\\470919565641814961196573721530953737756400690804642916690448036268023345060055365765180543\\024624602123866399730144748758038656931789371038406200969185401458605233665580015630409753\\748368355891777932$ 

CPU time: 1.21 s, Wall time: 0.00 s

```
# PARI -- an open source "competitor" -- which doesn't implement \
    asymptotically
# fast algorithms... takes 61 seconds on what takes Sage only 1.2 seconds\
    :

m = random_matrix(ZZ,200)
g = gp(m)
%time g.matdet()
```

 $135753046785592272670949639113731138010628571292147640483630705584920731036615594516655463\\648920594266652841613244812506025860156753659593320638910235683149694055820945847286579164\\088583533400876342839197653862920340029840314309077747640520241247155187555406364215325475\\278287831411236815157067998590441380705417248522641661888528610264660692779031684195568888\\798547548520635907458061399083242472195280650808745547355233910181385442740889768590555152\\3414618216712$ 

```
CPU time: 61.81 s, Wall time: 0.00 s
```

```
m = random_matrix(ZZ, 300)
%time m.det()
```

 $-36656715590627309259273556481881789462901939923531360793867281610385107201681220482880222\\393761888024820796864311562604558417184724840961390842317156721562411764807649790447353028\\762036659749422497603071523179929464858478336910202053252562638678128526896500157785165054\\293794181261236340877929324049832895881162922644941697319222838224274528581864774816760098\\633846426568893761513261653040173779579974072207726365822070522566752960386653985032850957\\285364543777278122604573855524765351327787341793324699648149366240468175888574163954601787$ 

067935414108698721339178390706013060867742054317774668049577948380584476918673603456117521789370051146973348953237116581299139570448159964895572091946444810118180145708719947061161201369437878805232850725058812193721000

CPU time: 3.52 s, Wall time: 1.92 s

```
m = random_matrix(ZZ, 500)
%time m.det()
```

124423606683294696600781583201186670364081090679629871550747263911593699563304600113975106 400944403310607026290606788346046531219242161553754683944465019655784506677458738897853766 91833236832957627046

CPU time: 10.63 s, Wall time: 3.56 s

Depending on time, say something about how m.det is so frickin fast it uses a whole bunch of surprising tricks.

- Strassen: matrix multiplication done by decomposition matrix into blocks and doing 7 multiplies instead of 8
- Multimodular: working modulo prime powers, and using the Chinese Remainder theorem.
- Cramers Rule: Solving a random linear system and looking at the denominator of the resulting vector, then fixing it.