

Tutorial-3

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10

1) Pseudo code for linear search

```
for (i = 0 to n)
{
    if (arr[i] == key)
        printf "Element found"
}
```

2)

Recursive

void insertion (int arr[], int n)

```
{
    if (n <= 1)
        return;
    insertion (arr, n-1);
    int num = arr[n-1];
    int j = n-2;
    while (j >= 0 && arr[j] > num)
    {
        arr[j+1] = arr[j];
        j--;
    }
    arr[j+1] = num;
}
```

iterative

```
for (i = 1 to n)
    key = A[i]
    j = i-1
    while (j >= 0 && A[j] > key)
    {
        A[j+1] = A[j]
        j = j-1;
    }
```

A[j+1] = key;

insertion sort is online sorting because it doesn't know the whole input, more input can be inserted while the insertion sorting is running.

(2)

3) Complexity of different sorting algorithms

Name	Best case	worst case	Average
Selecting sorting	$O(n^2)$	$O(n^2)$	$O(n^2)$
Bubble sorting	$O(n)$	$O(n^2)$	$O(n^2)$
Insertion sorting	$O(n)$	$O(n^2)$	$O(n^2)$
Heap	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$
Quick	$O(n \log n)$	$O(n^2)$	$O(n \log n)$
merge	$O(n \log n)$	$O(n \log n)$	$O(n \log n)$

4) Inplace sorting

Bubble
 Selecting Selection
 Insertion
 quick
 heap

Stable sorting

Merge
 Bubble
 Insertion

Online sorting

Insertion

5) Iterative

```

int b-search (int arr[], int l, int r, int key)
{
    while (l <= r) {
        int m = ((l+r)/2);
        if (arr[m] == key)
            return m;
        else if (key < arr[m])
            r = m-1;
        else {
            l = m+1;
        }
    }
    return -1;
}

```

Time complexity $\rightarrow O(n)$

(3)

Recursive

```

int b_search(int arr[], int l, int r, int key)
{
    while (l <= r) {
        int m = (l + r) / 2;
        if (key == arr[m])
            return m;
        else if (key < arr[m])
            return b_search(arr, l, mid-1, key);
        else
            return b_search(arr, mid+1, r, key);
    }
    return -1;
}

```

Time complexity $\rightarrow \boxed{O(\log n)}$

$$\begin{aligned}
 6) \quad T(n) &= T(n/2) + 1 \quad \text{--- (1)} \\
 T(n/2) &= T(n/4) + 1 \quad \text{--- (2)} \\
 T(n/4) &= T(n/8) + 1 \quad \text{--- (3)}
 \end{aligned}$$

$$\begin{aligned}
 T(n) &= T(n/2) + 1 \\
 &= T(n/4) + 2 \\
 &= T(n/8) + 3 \\
 &= T\left(\frac{n}{2^k}\right) + k
 \end{aligned}$$

$$\text{let } 2^k = n$$

$$k = \log n$$

$$T(n) = T\left(\frac{n}{n}\right) + \log n$$

$$T(n) = T(1) + \log n$$

Time complexity $\Rightarrow \boxed{O(\log n)}$

```

7) for (i=0; i < n; i++)
    { for (int j=0; j < n; j++)
        { if (arr[i] + arr[j] == k)
            printf("%d %d", i, j); }
    }

```

8) Quick sort is fastest general-purpose sort. In most practical situation quick sort is the method of choice if stability is important and space is available, merge sort might be best.

9) Inversions in array:-

- A pair $(A[i], A[j])$ is said to be inversion if $A[i] > A[j]$
- $i < j$
- Total no of inversions in given array are 31 using merge sort.

10) Worst case ($O(n^2)$):-

- when the pivot element is an extreme (smallest/largest) element. This happens when input array is sorted or reverse sorted and either first or last element is selected as pivot.

Best case ($O(n \log n)$):-

- The best case occurs when we will select pivot element as a mean element.

11) Merge sort:-

$$\begin{aligned}
 & \text{Best case} \rightarrow T(n) = 2T(n/2) + O(n) \\
 & \text{Worst case} \rightarrow T(n) = 2T(n/2) + O(n)
 \end{aligned}
 \left. \vphantom{\begin{aligned} \text{Best case} \\ \text{Worst case} \end{aligned}} \right\} O(n \log n)$$

Quick sort:-

$$\begin{aligned}
 & \text{Best case} \rightarrow T(n) = 2T(n/2) + O(n) \rightarrow O(n \log n) \\
 & \text{Worst case} \rightarrow T(n) = T(n-1) + O(n) \rightarrow O(n^2)
 \end{aligned}$$

⑤

- In quick sort, array of elements are divided into 2 parts repeated until it is not possible to divide further.
- In merge sort, the elements are split into 2 subarray $(n/2)$ again and again until only 1 element is left.

```
12) for (int i=0; i < n-1; i++)  
{  
    int min = i;  
    for (int j = i+1; j < n; j++)  
        if (a[min] > a[j])  
            min = j;  
  
    int key = a[min];  
    while (min > i)  
    {  
        a[min] = a[min-1];  
        min--;  
    }  
    a[i] = key;  
}
```

- 13) A better version of bubble sort, known as modified bubble sort, includes a flag that is set if an exchange is made after an entire pass over. If no exchange is made then it should be called the array is already sorted because no 2 elements need to be switched.

```
void bubble (int arr[], int n)  
{  
    for (int i=0; i < n; i++)  
    {  
        swaps = 0;  
        for (int j=0; j < n-i-1; j++)  
        {  
            if (arr[j] > arr[j+1])  
            {  
                int t = arr[j];  
                arr[j] = arr[j+1];  
                arr[j+1] = t;  
                swap++;  
            }  
        }  
        if (swap == 0)  
            break;  
    }  
}
```

}}}