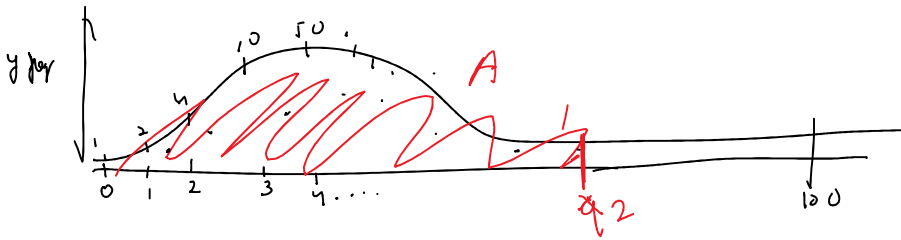
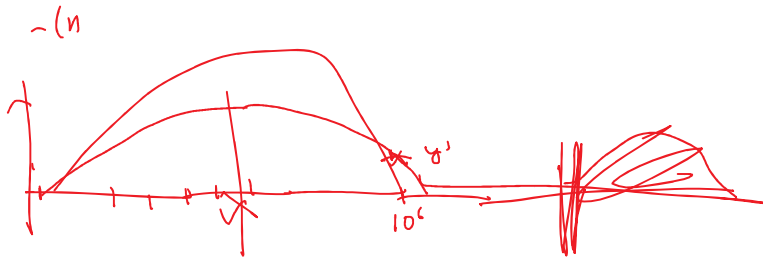


$$f(n) = 2 + 3n \rightarrow O(n)$$



A
Total ans



$$K + C \cdot \log n = f(n)$$

$$\rightarrow O(n)$$

$$f(n) = 1 + 2(1 + \log_2 n)$$

$f(n) = 1 + 2n \rightarrow O(n)$

```

while (i <= n) {
    i += 1;
}
            
```

$f(n) = 1 + 2 \cdot \log n$

```

while (i <= n) {
    i = i + 2;
}
            
```

```

while ( ) {
    i = i + 2;
    i = i + 3;
}
            
```

$i = n$

```

w( ) {
    i = i * 2;
}
            
```

```

w( ) {
    i = i * 3;
}
            
```

```

w( ) {
    i = i * 2;
    i = i * 3;
}
            
```

l	i
1	1
2	K
3	K^2
l	K^{l-1}

$$K^{l-1} = n$$

```

1 = 1 + 2
while ( ) {
  i = i + k
}

```

$$K^{l-1} C = n$$

$$l-1 = \log_k n + 1$$

$$f(n) = 1 + 2(1 + \log_k n)$$

$$O(\log_k n)$$

$$\log_a C$$

$$\frac{a}{u} = 0$$

```

while (n > 0) {
  n = n / 2;
}

```

$$n = 2^{l-1}$$

$$\log_2 n = l$$

l	n
1	n
2	$n/2$
3	$n/2^2$
4	$n/2^3$
l	$n/2^{l-1} \Rightarrow 0 = 1$

```

for ( i = 0; ( i^2 <= n; i++ ) {
  syso(i);
}

```

$$\sqrt{n} = n^{1/2}$$

$$n^{-1} = \frac{1}{n}$$

$$(\sqrt{n} + 1) \cdot C_1 + C_2 = f(n)$$

$$\begin{array}{c}
 i=1 \\
 j=1 \text{ to } 1 \\
 k \rightarrow 1000 \\
 1k
 \end{array}
 \left|
 \begin{array}{c}
 2 \\
 j=1 \\
 1000 \\
 2k
 \end{array}
 \right|
 \left|
 \begin{array}{c}
 3 \\
 1 \quad 2 \quad 3 \\
 1000 \quad 1000 \quad 1000 \\
 3k
 \end{array}
 \right|
 \left|
 \begin{array}{c}
 4 \\
 \dots \\
 n \cdot 1k
 \end{array}
 \right|
 \dots
 \left|
 \begin{array}{c}
 n
 \end{array}
 \right|$$

$$1k(1 + 2 + 3 + \dots + n) \Rightarrow 1000 \cdot \frac{n \cdot (n+1)}{2}$$

$$O(n^2) \checkmark$$

$$\begin{array}{c}
 i=1 \\
 j=1 \\
 n/2
 \end{array}
 \left|
 \begin{array}{c}
 i=2 \\
 j=1 \text{ to } 2^2 \\
 n/2
 \end{array}
 \right|
 \left|
 \begin{array}{c}
 i=3 \\
 j=1 \text{ to } 3^2 \\
 n/2
 \end{array}
 \right|
 \left|
 \begin{array}{c}
 i=4 \\
 4^2 \cdot n/2
 \end{array}
 \right|
 \left|
 \begin{array}{c}
 i=n \\
 n^2 \cdot n/2
 \end{array}
 \right|$$

$$n/2 + 4 \cdot n/2 + 3^2 \cdot n/2 + \dots + n^2 \cdot n/2$$

$$n/2 (1 + 2^2 + 3^2 + \dots + n^2)$$

$$O(n^4)$$

$$n/2 \cdot \frac{n \cdot (n+1) (2n+1)}{6}$$

$$\left(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n} \right)$$

$$\begin{array}{c}
 i=1 \quad 2 \quad 3 \quad \dots \quad n \\
 n + n/2 + n/4 + \dots + n/n
 \end{array}$$

$$\log n$$

$$n \left(1 + \frac{1}{2} + \frac{1}{3} + \frac{1}{4} + \dots + \frac{1}{n} \right) = n \log n$$

$$\int \frac{1}{x}$$



$$\frac{n}{k}$$

$$O(1) \xrightarrow{O(\log n)} O(\sqrt{n}), O(n), O(n \log n)$$

$$O(n^2), O(n^2 \log n), O(n^3), O(c^n), O(n!), O(n^n)$$

~~$$O(n^n)$$~~

$$c^n < \underline{n!}, n^n$$

$$(c \cdot c \cdot c \cdot \dots \cdot c) \quad c^{k-1}$$

$$(1 \cdot 2 \cdot 3 \cdot \dots \cdot (i-1) \cdot (i+1) \cdot (i+2) \cdot \dots \cdot n)$$

$$1) \rightarrow \frac{1}{4} \cdot 4 \cdot 4 \cdot 4$$

More subtle recursion seen

Q Max Sub array sum
 $n \leq 10^5$

$$O(n^3)$$

$$10^{15} / 10^9 = 10^6$$

$$O(n^2)$$

$$O(10^{10})$$

$$O(n)$$

$$10^5$$

$$10^{10} / 10^9 \approx 10$$

	W	Best
Bubble	$O(n^2)$	$\Omega(n^2) \rightarrow \Omega(n)$
Selec	$O(n^2)$	$\Omega(n^2)$
Ins	$O(n^2)$	$\Omega(n)$

Given : $1 * x^n + 2 * x^{n-1} + 3 * x^{n-2} \dots + nx$;
 solve, return number
 $n=3$;
 $x=7$;

$$1 \cdot (7)^3 + 2 \cdot (7)^2 + 3 \cdot 7^1$$

$$Q \rightarrow 1 \cdot x^n + 2 \cdot x^{n-1} + 3 \cdot x^{n-2} + 4 \cdot x^{n-3} \dots + nx^1$$

$$n \rightarrow 3, x = 7$$

$$O(n)$$

$$7^3 + 2 \cdot 7^2 + 3 \cdot 7^1$$

$$343 + 2 \times 49 + 21$$

$$98$$

$$462$$

$$n \rightarrow 2$$

$$x \rightarrow 1$$

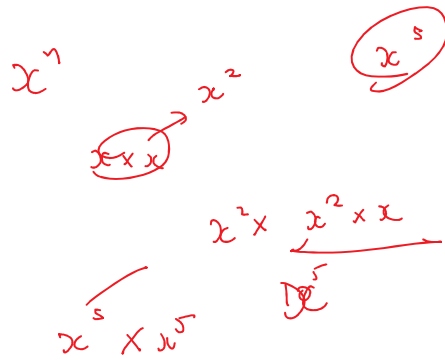
$$1 \cdot (1)^2 + 2 \cdot (1)$$

$$= 3$$

$$x^n$$

$$x$$

$$x^s$$



(Key, Value) →

unique

array → 0

$n-1$

A → 10

B → 20

C → 15

D → 10

E → 20

Search → $O(n)$, Key & Value

Key →	A	B	C	D	E
Value	10	20	15	10	20

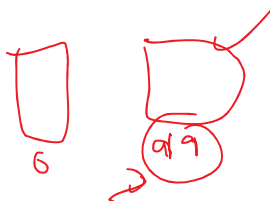
Add → $O(1)$
 Remove → $O(n)$
 Update → $O(n)$

Has

Search → Key $O(1)$, Update $O(1)$

Add → $O(1)$

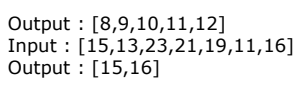
Remove → $O(1)$

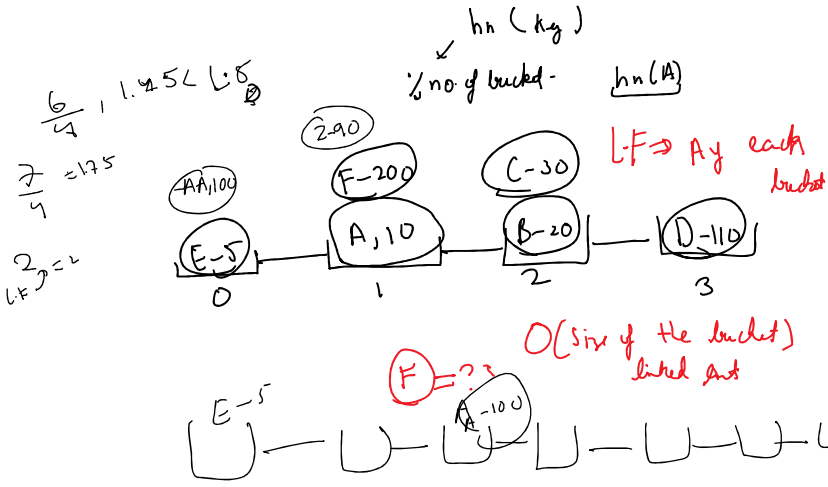
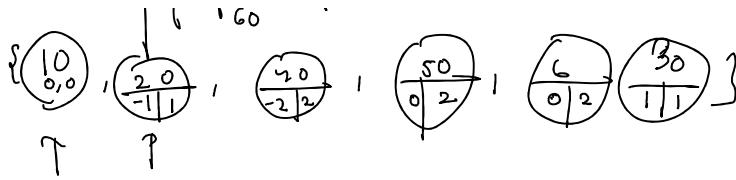


int[] arr1 = {30, 20, 40, 50, 70, 20, 20, 20, 50, 50};
 int[] arr2 = {30, 80, 20, 20, 20, 90, 50, 20};
 30, 20, 50, 20, 20, 50
 30, 20, 50, 20, 20, 50

Q number of string concatenation for string s to be a subsequence of string t

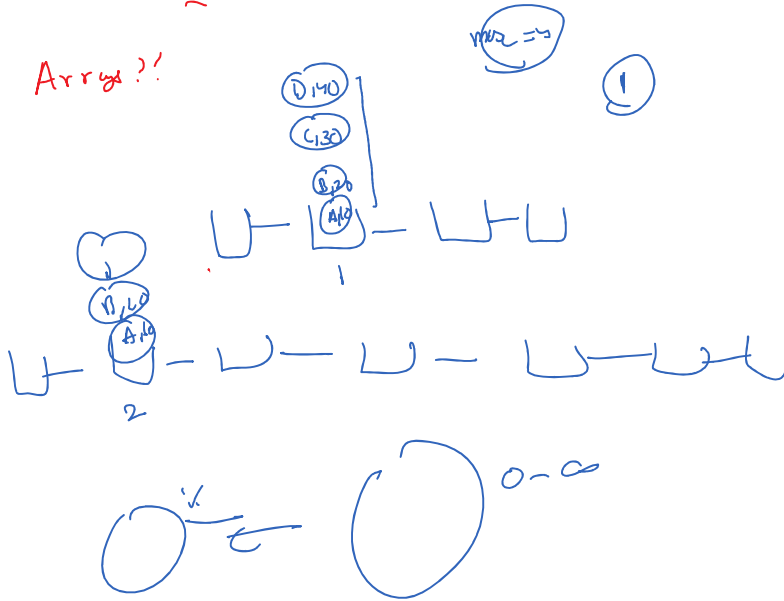
From
<https://docs.google.com/spreadsheets/d/1Ww2mE71X2rte9kVV3cu74wBapoch18P-y18kcn0OM4/edit#gid=1635212751>





any size of each bucket \Rightarrow L.F??
 collision resolve?? \rightarrow chaining.
 default bucket \rightarrow 16

Array??



1	2	3	4	5
6	7	8	9	10
11	12	13	14	15
16	17	18	19	20
21	22	23	24	25

Num = 2	Num = 3	Num = 4	Num = n
Div = 1, 2	Div = 1, 2, 3	Div = 1, 2, 3, 4	Div = 1 to n

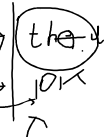
$N \Rightarrow n/2$
 $N \Rightarrow N/3$

$n/2 + n/3 + n/5 + n/7 + n/11 + \dots + n/n$
 $N \log(\log(n))$

F	F	F	F	F
F	F	F	F	F
F	F	F	F	F
F	F	F	F	F

1, 2, 3, 4, 5, ^

200 x 500 x 4 x 2
8,00,00,00 byte
8 MB

V1 = 
V2
V3

V999

String	n	Operation
"0"	1	1+1
"01"	2	2+1
"012"	3	3+1

String str1 = "Hello"; // Literal
String str2 = "Hello";
String str3 = "Hello";

H
I
P
ok
Hello

	String	StringBuilder
Create	String str = "hello"; String str = new String("hello")	String sb = new StringBuilder("hello");
Print	Sysout(str)	Sysout(sb)
Char at	str.charAt(index)	str.charAt(index)
SubString	str.substring(index);	str.substring(index); // return String
Add Insert Update Remove	X	sb.append(""); sb.insert(index, ""); char or string sb.setCharAt(index, ""); sb.deleteCharAt(index); sb.delete(index, excluding) sb.toString(); // return to string

1) → B.P.

4) → Self-w
Bi soln

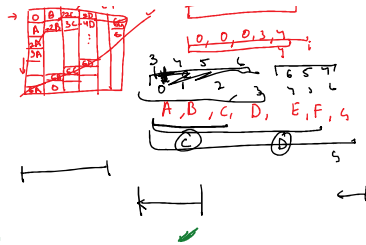
4

4
3
2
1

PD(4)

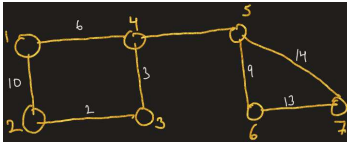
PD(3)

3
2
1

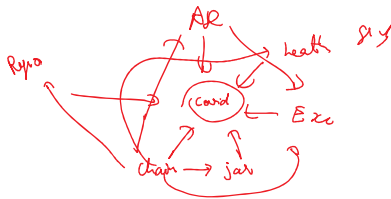


$\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 & 6 \\ A & B & C & D & E & F & G \end{matrix}$
 $\{3, 6, 7, 2, 5, 2\}$ -000
 $\begin{matrix} 0 & 1 & 2 & 3 & 4 & 5 \\ 0 & 1 & 2 & 3 & 4 \end{matrix}$
 -0

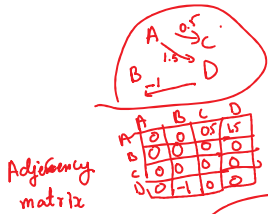
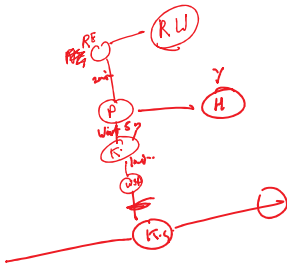
$\rightarrow \{A, B, C, D\}$



8 edges
7 \rightarrow vertex / node

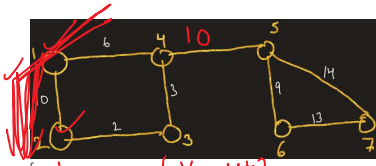


$\begin{matrix} \uparrow & \uparrow & \uparrow & \uparrow \\ \downarrow & \downarrow & \downarrow & \downarrow \end{matrix}$ $\begin{matrix} \hookrightarrow \text{visited} \hookrightarrow \text{directed} \\ \hookrightarrow \text{non-visited} \hookrightarrow \text{undirected} \end{matrix}$



$N \times N^2$
 $V \rightarrow V_{i+1}$
 $V \cdot (V-1)$

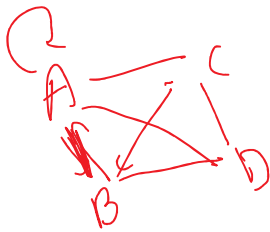
$A \rightarrow B?$ $\{A \rightarrow \{C, D\}, B \rightarrow \{E\}, C \rightarrow \{F\}, D \rightarrow \{B\}\}$ \rightarrow Adjacency matrix
 $\rightarrow O(N)$
 $A \rightarrow B$



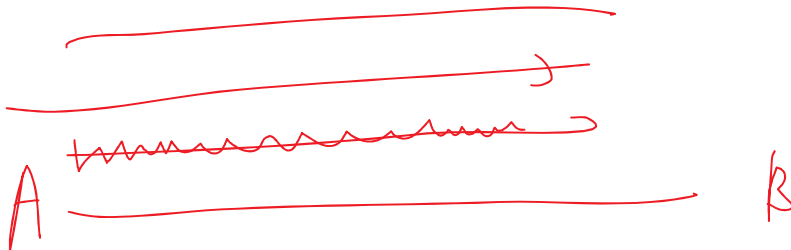
- $V_1 \rightarrow (V_2, w_t)$
 $1 \rightarrow \{(2, 10), (4, 6)\}$
 $2 \rightarrow \{(1, 10), (3, 2)\}$
 $3 \rightarrow \{(4, 3), (2, 2)\}$
 $4 \rightarrow \{(1, 6), (3, 3), (5, 10)\}$

- $5 \rightarrow \{(4, 10), (6, 9), (7, 14)\}$
 $6 \rightarrow \{(5, 9), (7, 13)\}$
 $7 \rightarrow \{(6, 13), (5, 14)\}$

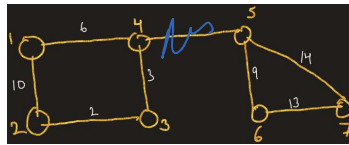
$\langle V_1, N_{V_1} \rangle$
 \downarrow
 $\langle V_2, w_t \rangle$



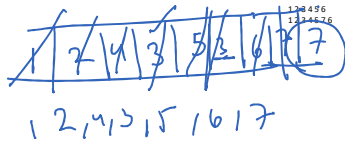
$$\begin{array}{r}
 4C2 \\
 \hline
 6 \quad 4P_2 \\
 4 \cdot 3 \\
 \hline
 12
 \end{array}$$



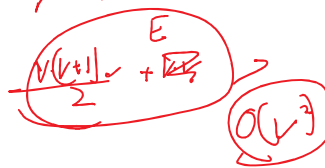
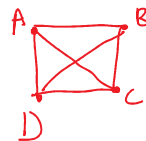
C.B



connected \rightarrow 1 component



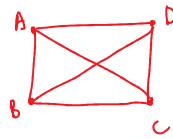
B.F.T traversal



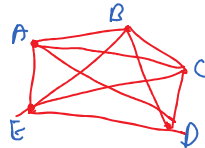
```

public boolean hasPath2(int src, int dest) {
    boolean ans = false;
    HashSet<Integer> Visited = new HashSet<>();
    Queue<Integer> Q = new LinkedList<>();
    Q.add(src);
    while (!Q.isEmpty()) {
        int curr = Q.poll();
        if (curr == dest) {
            ans = true;
        }
        if (Visited.contains(curr)) {
            System.out.println("Cycle exist");
            continue;
        }
        Visited.add(curr);
        for (Integer nbr : adjList.get(curr)) {
            if (!Visited.contains(nbr)) {
                Q.add(nbr);
            }
        }
    }
    return ans;
}

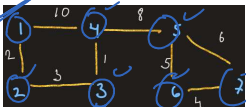
```



$O(E+V)$ DFI
 $O(E+V)$



Min spanning tree
Single src shortest Path
Tree



Fileged search
Union Find

