

A **transformation sequence** from word `beginWord` to word `endWord` using a dictionary `wordList` is a sequence of words `beginWord` ->  $s_1$  ->  $s_2$  -> ... ->  $s_k$  such that:

- Every adjacent pair of words differs by a single letter.
  - Every  $s_i$  for  $1 \leq i \leq k$  is in wordList. Note that beginWord does not need to be in wordList.
  - $s_k == \text{endWord}$

Given two words, beginWord and endWord, and a dictionary wordList, return the *number of words in the shortest transformation sequence* from beginWord to endWord, or 0 if no such sequence exists.

### Example 1:

**Input:** beginWord = "hit", endWord = "cog", wordList = ["hot", "dot", "dog", "lot", "log", "cog"]

**Output:** 5

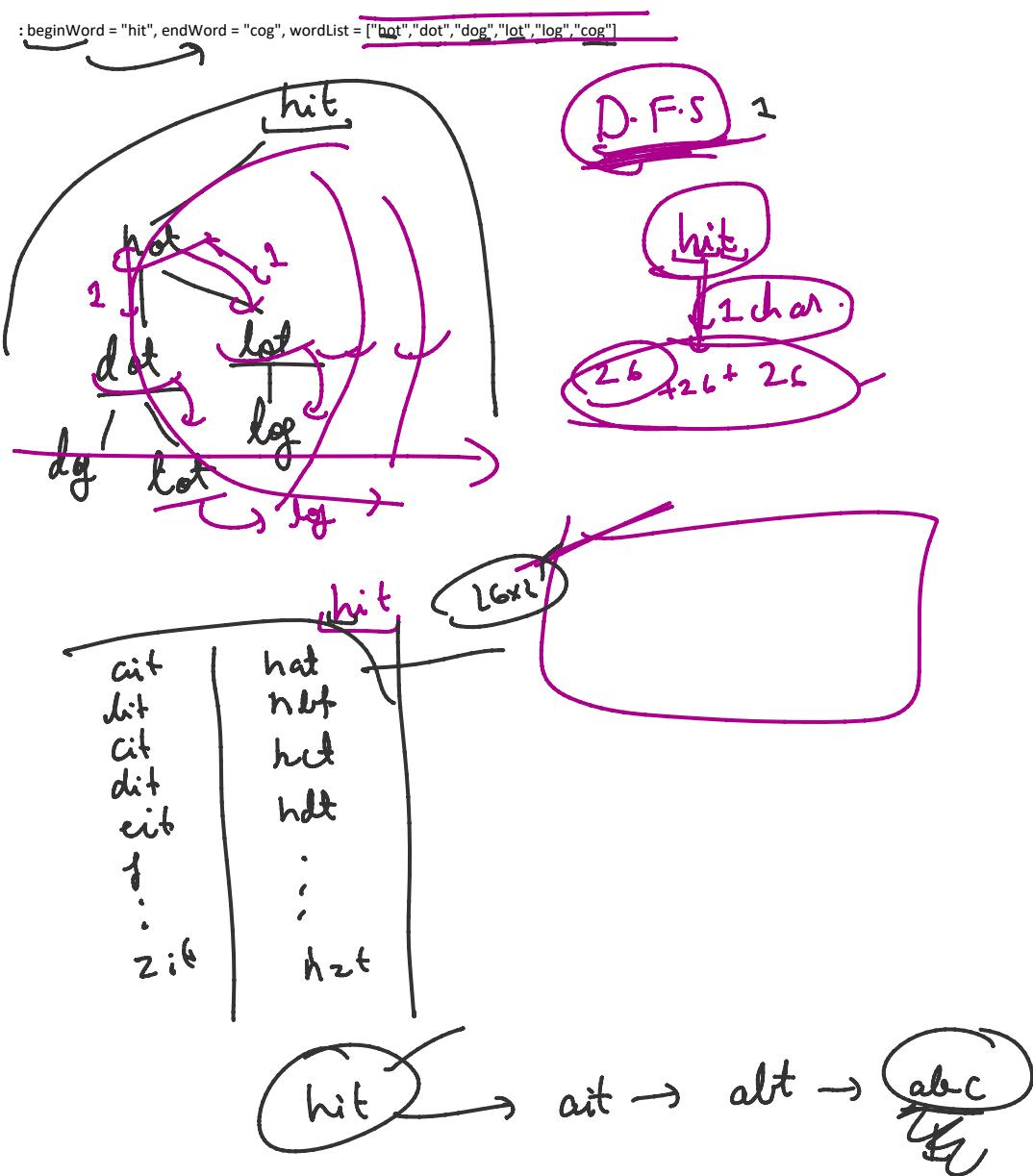
**Explanation:** One shortest transformation sequence is "hit" -> "hot" -> "dot" -> "dog" -> cog", which is 5 words long.

### Example 2:

**Input:** beginWord = "hit", endWord = "cog", wordList = ["hot", "dot", "dog", "lot", "log"]

**Output:** 0

**Explanation:** The endWord "cog" is not in wordList, therefore there is no valid transformation sequence.



There is an **undirected** graph with  $n$  nodes, where each node is numbered between 0 and  $n - 1$ . You are given a 2D array  $\text{graph}$ , where  $\text{graph}[u]$  is an array of nodes that node  $u$  is adjacent to. More formally, for each  $v$  in  $\text{graph}[u]$ , there is an undirected edge between node  $u$  and node  $v$ . The graph has the following properties:

- There are no self-edges ( $\text{graph}[u]$  does not contain  $u$ ).
- There are no parallel edges ( $\text{graph}[u]$  does not contain duplicate  $v$ ).
- If  $v$  is in  $\text{graph}[u]$ , then  $u$  is in  $\text{graph}[v]$  (the graph is undirected).
- The graph may not be connected, meaning there may be two nodes

A graph is **bipartite** if the nodes can be partitioned into two independent sets such that every edge connects a node in one set to a node in the other. Return true if and only if it is **bipartite**.

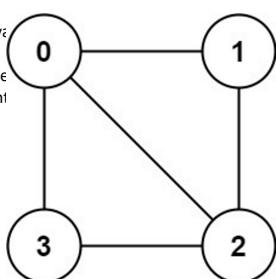
#### Example 1:

A graph is **bipartite** if the nodes can be partitioned into two independent sets  $A$  and  $B$  such that **every** edge in the graph connects a node in set  $A$  and a node in set  $B$ .

**Input:**  $\text{graph} = [[1,2,3],[0,2],[0,1,3],[0,2]]$

**Output:** false

**Explanation:** There is no way to partition the nodes into two independent sets such that every edge connects a node in one and a node in the other.



between them.  
the graph connects a node in set A and a node in set B.

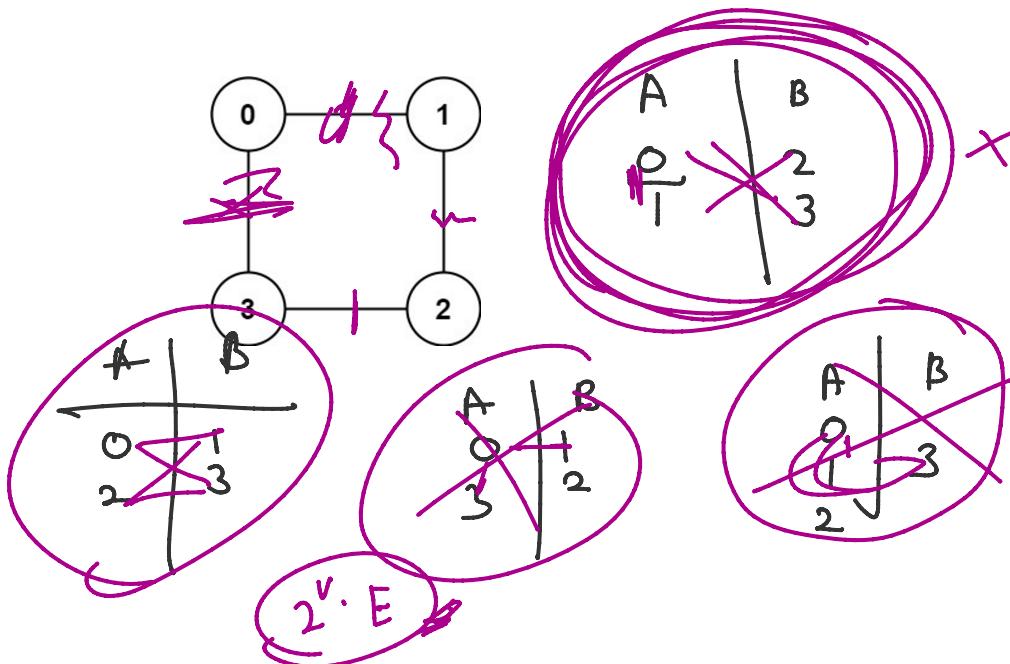
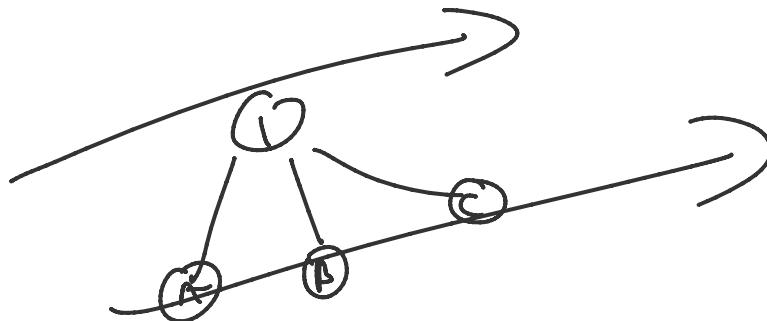
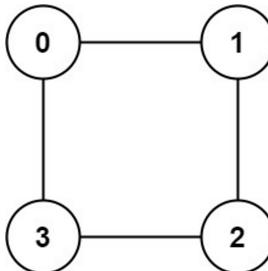
#### Example 2:

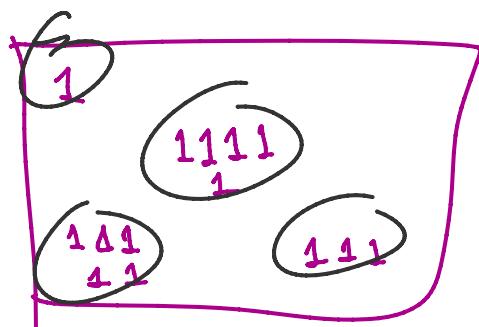
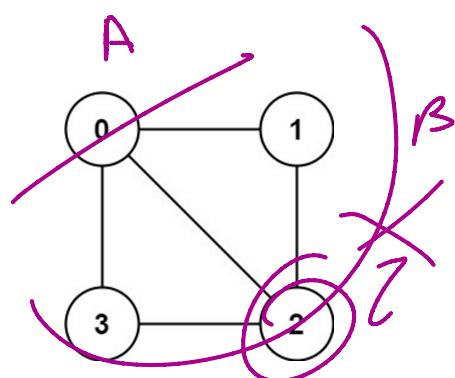
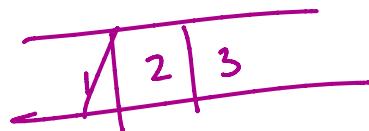
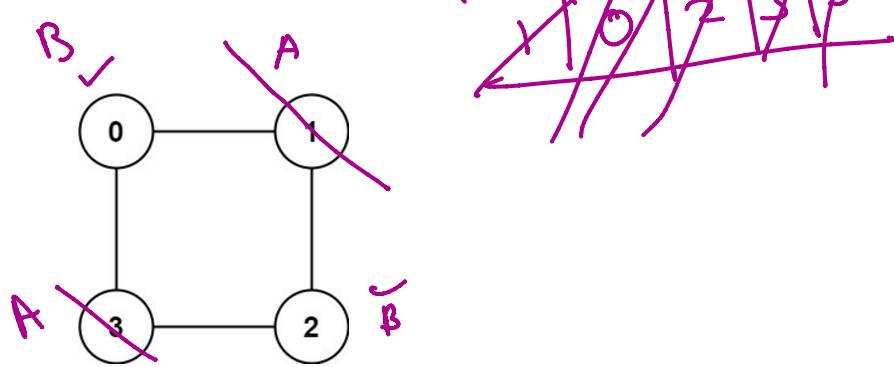
**Input:**  $\text{graph} = [[1,3],[0,2],[1,3],[0,2]]$

**Output:** true

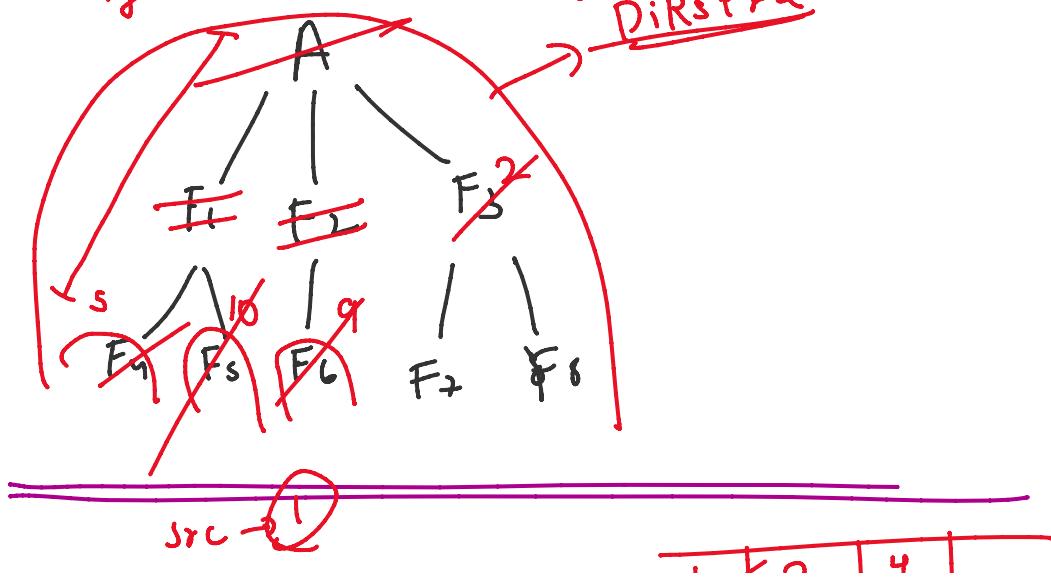
**Explanation:** We can partition the nodes into two sets:  $\{0, 2\}$  and  $\{1, 3\}$ .

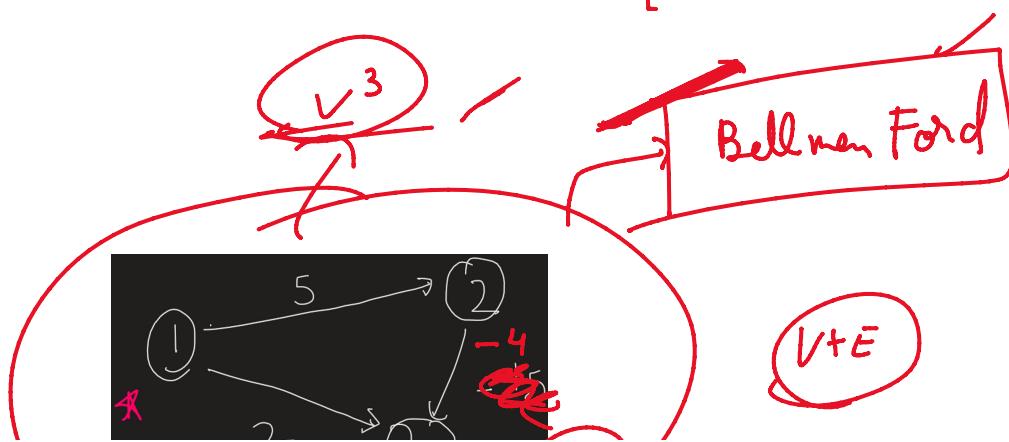
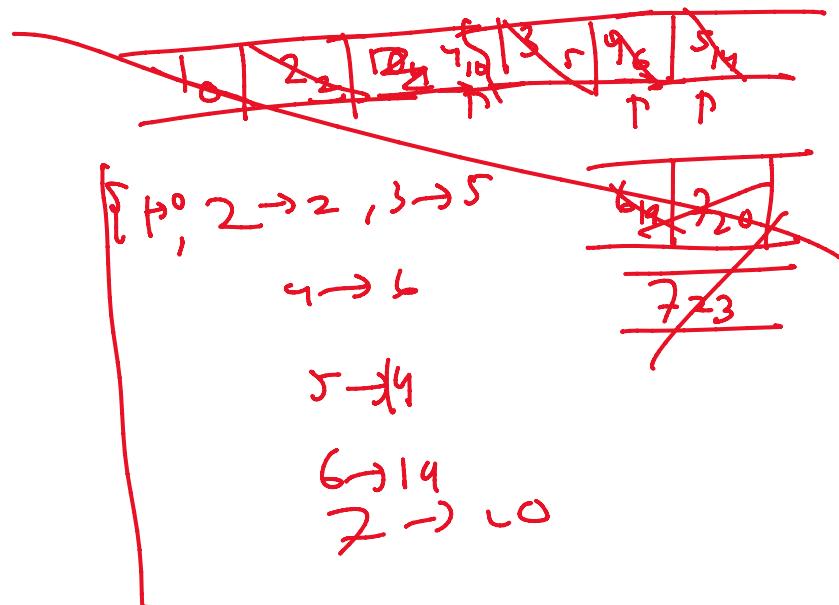
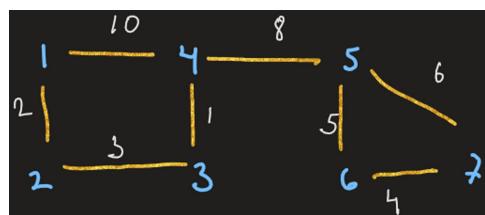
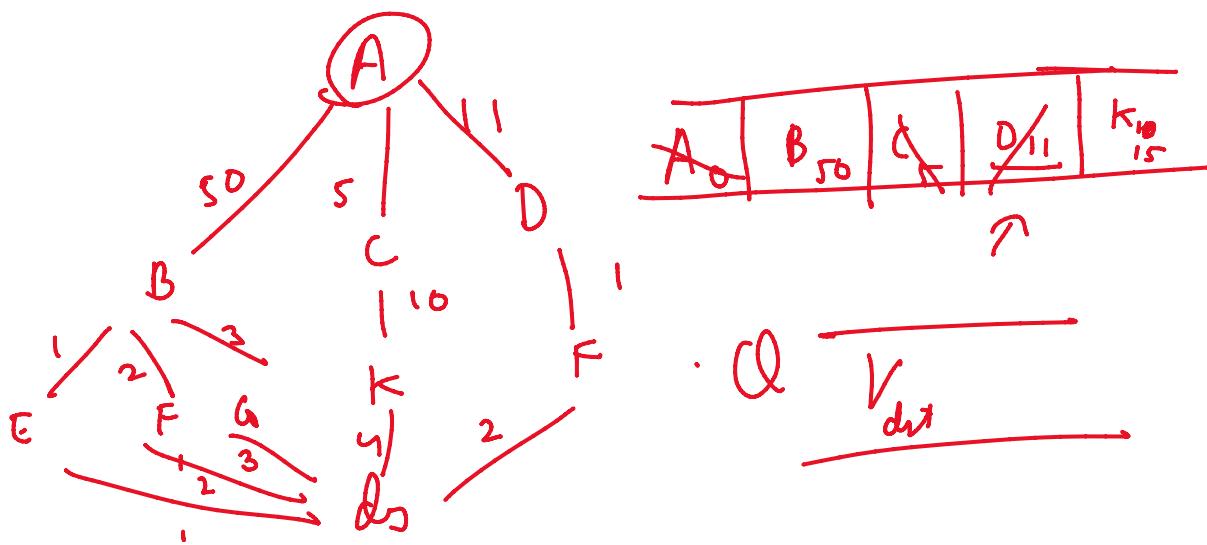
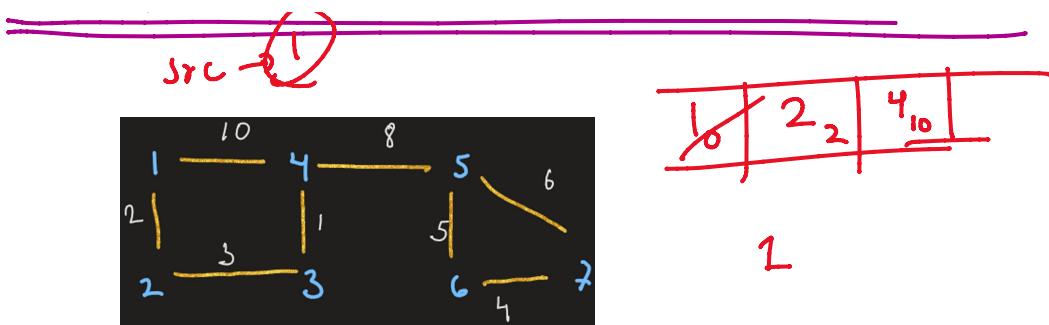
From <<https://leetcode.com/problems/is-graph-bipartite/>>

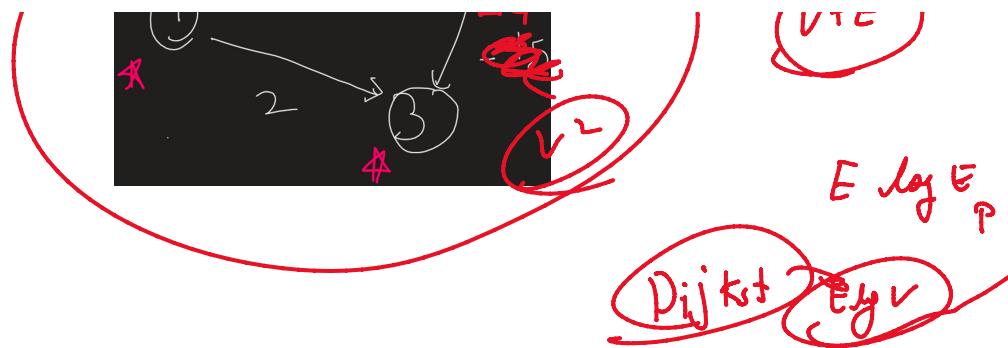




Single source shortest Algo :- Dijkstra







$$Fib(0) \rightarrow 0$$

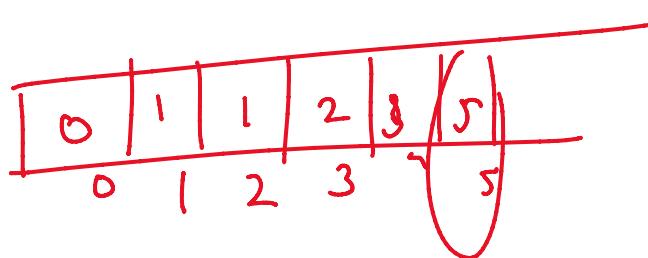
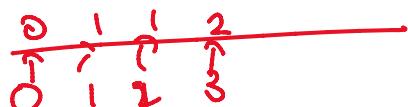
$$Fib(1) \rightarrow 1$$

$$Fib(1) \rightarrow 1$$

$$\vdots \quad (3) \rightarrow 2$$

$$\quad \quad \quad \downarrow \rightarrow 3$$

$$Fib(5) \rightarrow 5$$



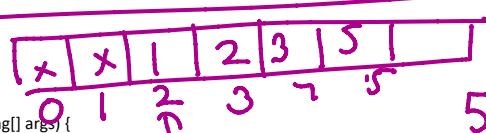
T.D

```

public static void main(String[] args) {
    int n = 5;
    System.out.println(Fibo(n, new Integer[n + 1]));
}

public static int Fibo(int n, Integer[] dp) {
    if (n <= 1) {
        return n;
    }
    if (dp[n] != null) {
        return dp[n];
    }
    int sp1 = Fibo(n - 1, dp);
    int sp2 = Fibo(n - 2, dp);
    Fibo(n) = > index?!
    dp[n] = sp1 + sp2;
    return sp1 + sp2;
}

```



Top-Down / Memoized  
Rec      Memorizing  
+ Space

1) Recursive

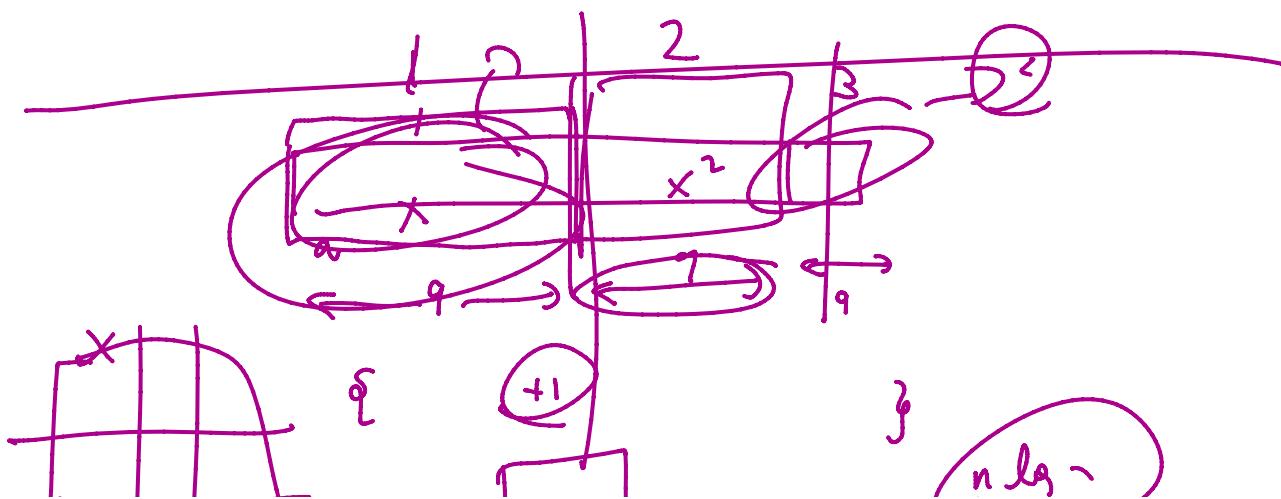
2) D.S  $\rightarrow$  Sol<sup>m</sup> ?? ] Confuse

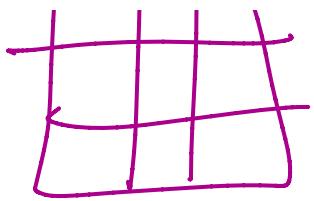
- 1) Recursive
- 2) D.S  $\rightarrow$  Sol<sup>n</sup> ?? } Confuse
- 3) D.S use
- 4) D.S ka size

- BU  $\leftarrow$  TD  $\rightarrow$  D.S means
- 1) D.S  
smallest to Biggest Problem  $\rightarrow$  loop
  - 2) Recur cur sol<sup>n</sup>  $\leftarrow$  T.D. copy  
 $\rightarrow$  recursive calls  
D.S ki form.

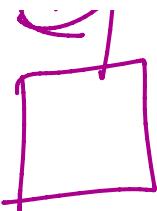
3) Base Case

- 4) Biggest Problem
- 5) D.S ka size





2



3



9

