



int xc → 4 b l

-2*10^9 to 2*10^9

123456789
123443210

9999999
9000000

v—

Inverse of a number

2 3 5 1 4
5 4 3 2 1

=====

3 1 4 5 2
5 4 3 2 1

6 at pos 3?!

600

7 at pos 5

70000

9 at pos 1

9

90

9*10^(pos-1)

13356789
Sum(9*10^0 Ans = ans +sum(digit*mult);
8*10^1
7*10^2
6*10^3
5*10^4

```
int num = 2789;  
int mult = 1;  
int ans = 0;  
while (num > 0) {
```

```

int num = 2789;
int mult = 1;
int ans = 0;
while (num > 0) {
    int digit = 9;
    if (9 - digit < digit) {
        if (0 < 9) {
            digit = 0;
        }
        num = 278;
        ans = 0+0*1;
    }
    mult = 1* 10;
}
System.out.println(ans);

```

```

while (num > 0) {
    int digit = 8;
    if (1 < 8) {
        digit =1;
    }
    num = 27;
    ans = 0 +1 *10;
    Ans =10
    mult = 100;
}

```

```

while (num > 0) {
    int digit = 7;
    if (2 < 7) {
        digit =2;
    }
    num =2;
    ans = 10+ 2 *100;
    Ans = 210
    mult = 1000* 10;
}

```

```

while (num > 0) {
    int digit =2;
    if (9 - 2 < 2) {
        digit = 9-digit;
    }
    num = 0;
    ans = 210+ 2* 1000;
    mult = mult * 10;
}
2210

```

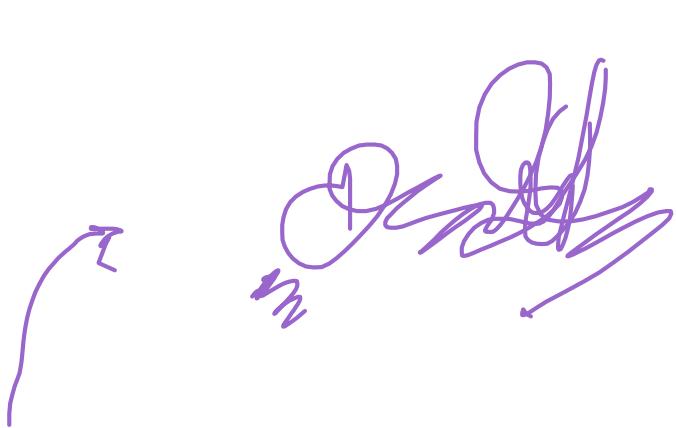
```
R total_star_in_row(tst)
1 5
2 5
3 5
4 5
5 5
```

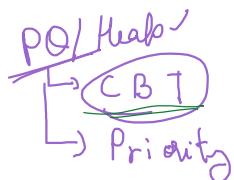
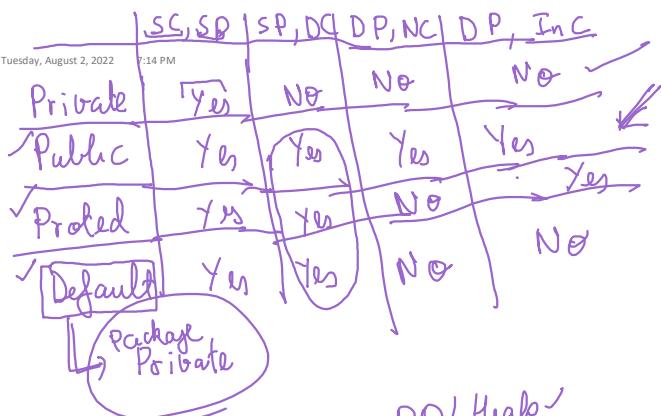
```
R tst
1 1
2 2
3 3
4 4
5 5
```

```
R tst
1 5
2 4
3 3
4 2
5 1
```

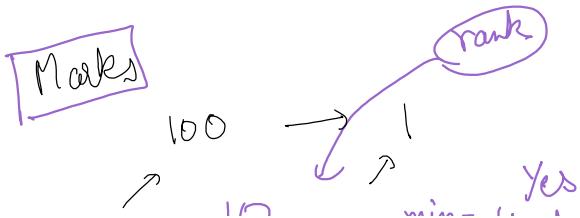
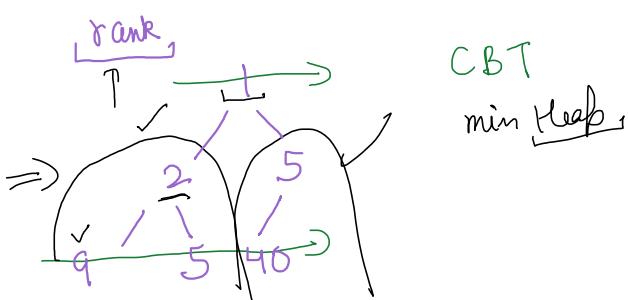
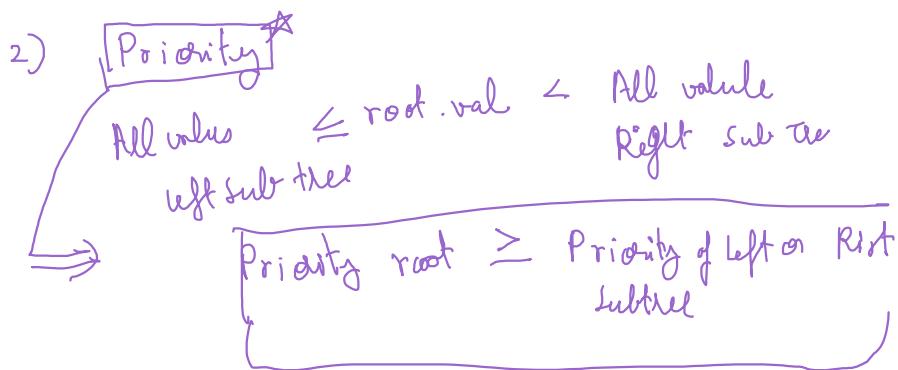
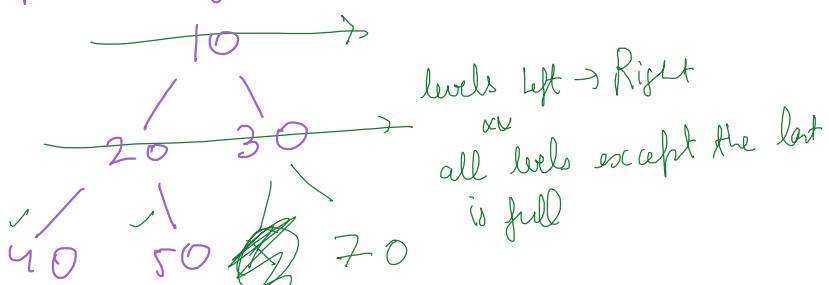
R , t_sp,t_st
1 4 1
2 3 2
3 2 3
4 1 4
5 0 5

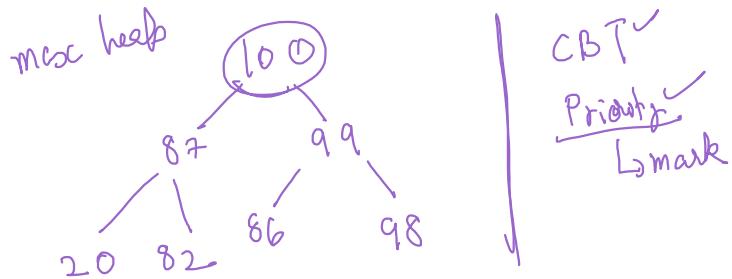
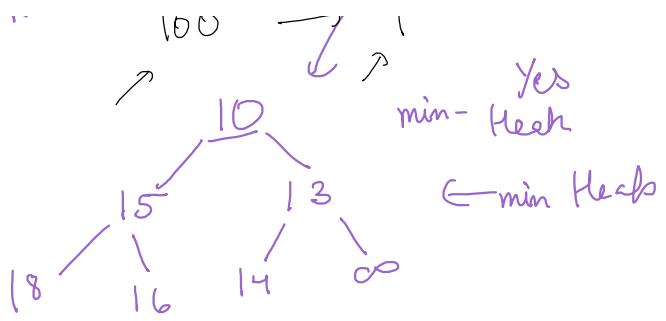
*
- - - * *
- - - * * *
- - - * * * *
- - - * * * *





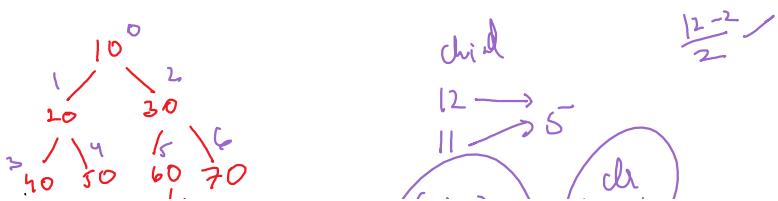
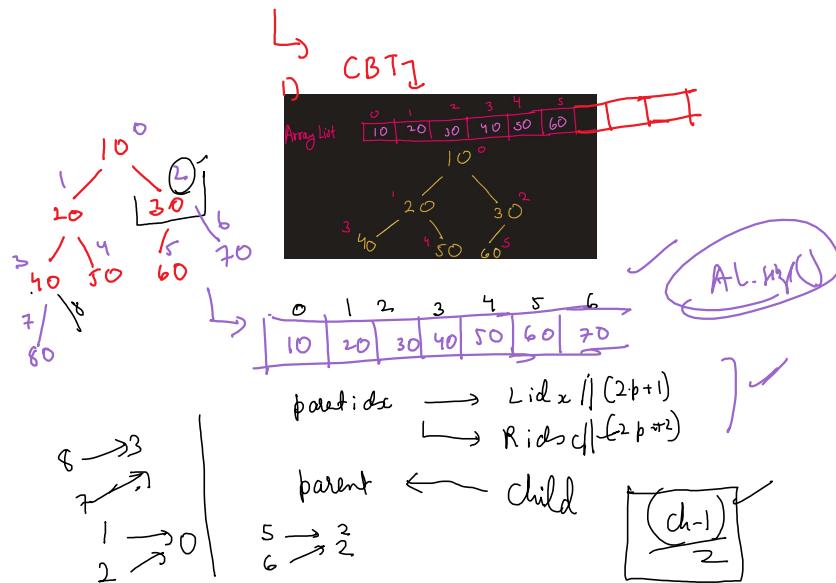
Complete Binary Tree

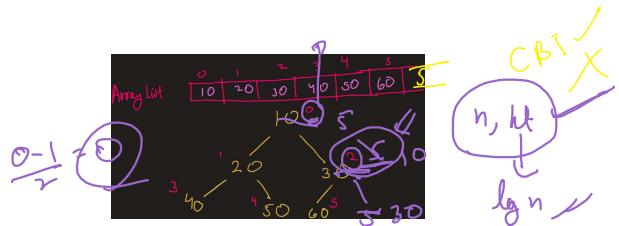
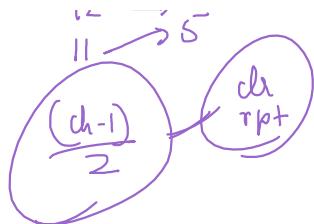
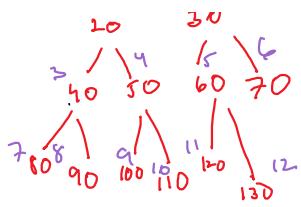




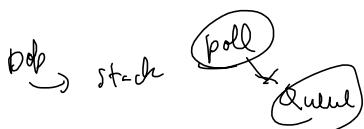
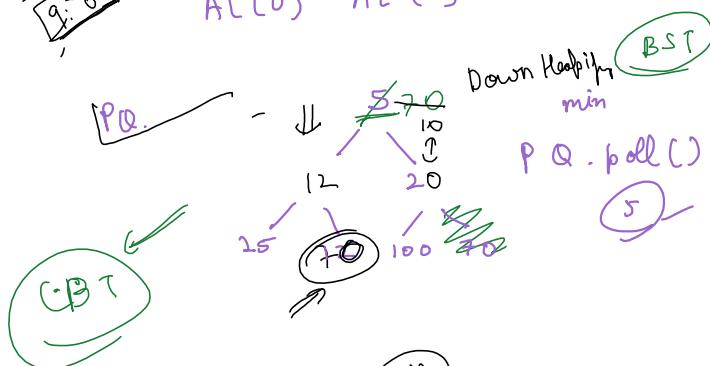
	Sorted array (increasing order)	Unsorted Array	Heap
Add	$O(n)$, find correct index and shift elements	$O(1)$, last place	$O(\log n)$
Delete Min element	shifting, $O(n)$	shifting $O(n)$	$O(\log n)$
Get Min element	$O(1)$	$O(n)$	$O(1)$

Heap





AL[0] AL[0]



Q largest K elements

$K=3$

$n \cdot \log K$

M1 → 2 loops $n \times K$

M2 → $n \log n + K$

M3 → $\boxed{n \log n + K \log n}$

Queue

[20, 30, 60, 50, 10, 55, 57, 40]

{ 20, 30, 60, 50, 10, 55, 57, 40 }

+

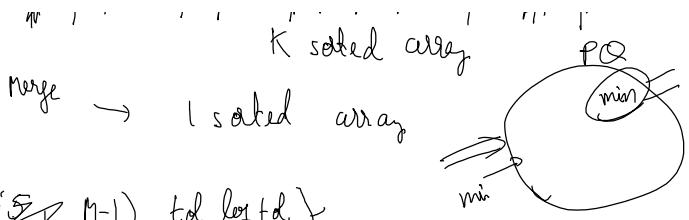
Q[[10, 20, 30], [15, 35], [5, 11, 25], [17, 22, 40]]

K sorted array

Merge

...

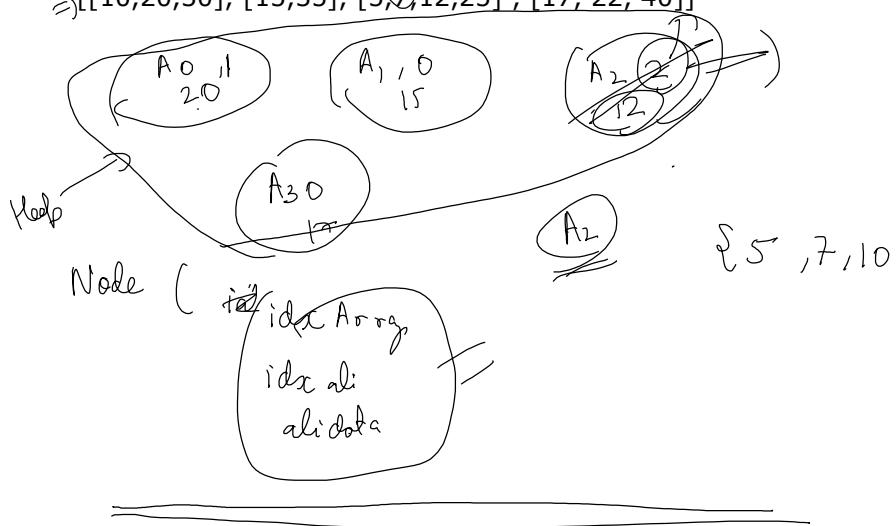
PQ min



~~M-1) tot log td~~

M-2) {5, 7, 10, 12, 15, 17, 20, 22, 25,
30, 35, 40}
K. total

→ [[10, 20, 30], [15, 35], [5, 12, 25], [17, 22, 40]]



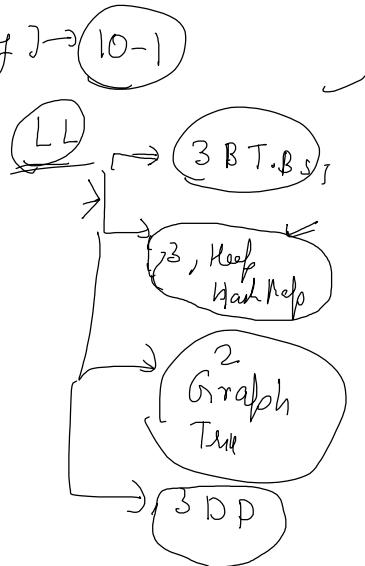
Hash Map

May → 3

Weedy → 2 ✓

May
Weedy → 2 class

[June + July + Aug]



Hash Map

Arrays

A.L.

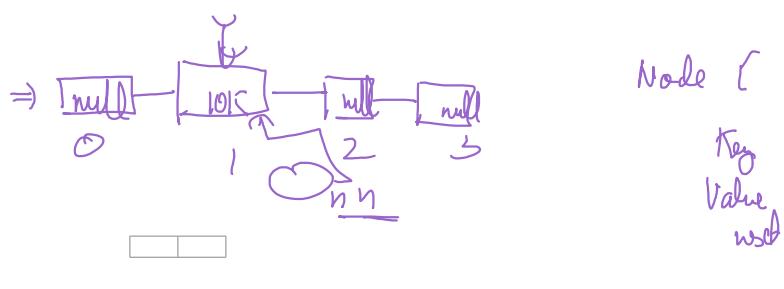
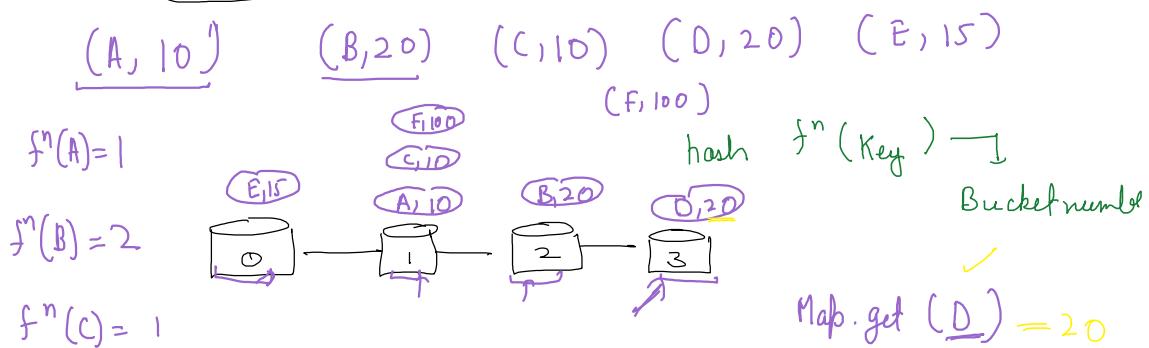
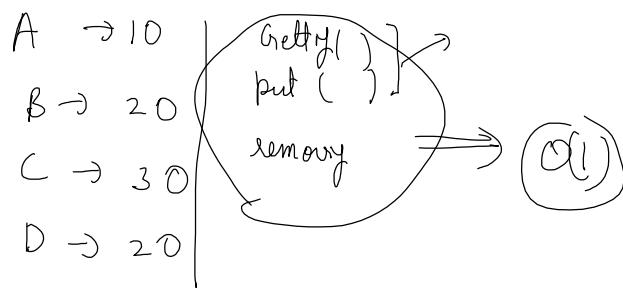
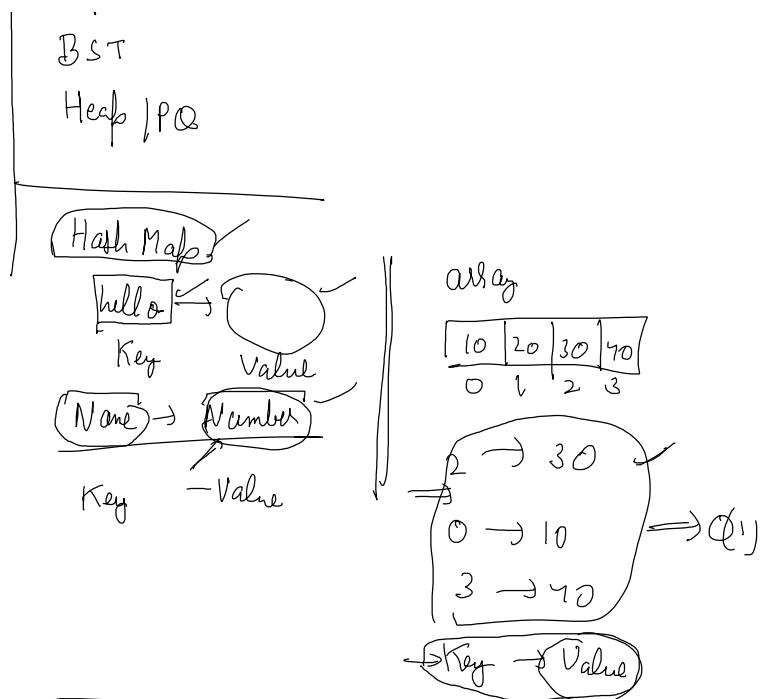
C.L.L.

|
BT
BST
Heap

Groups
A.L.
Stack
Queue

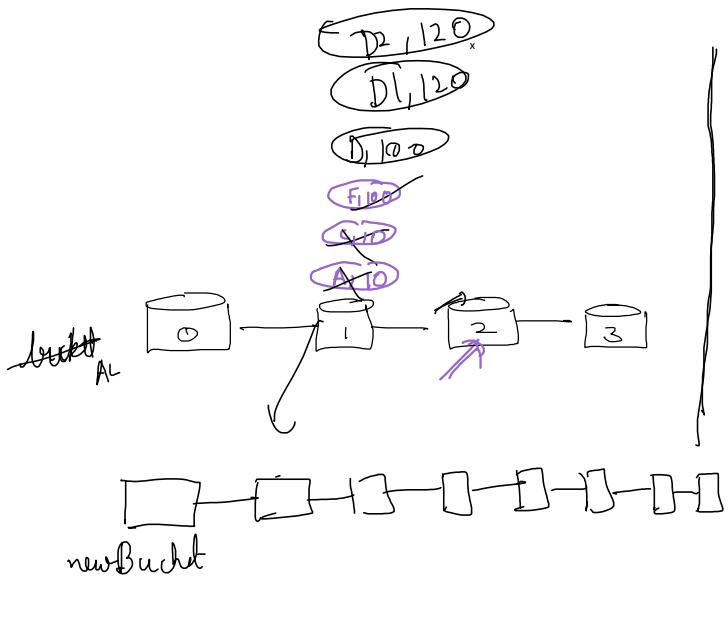
$L \cdot L \rightarrow i^{\text{th}} \rightarrow O(1)$
 $\rightarrow O(n)$

$\text{Key} \rightarrow 0 \text{ to } \text{len}-1$





WSD



Add $\rightarrow O(1)$ time
put $\rightarrow O(\text{size of bucket})$

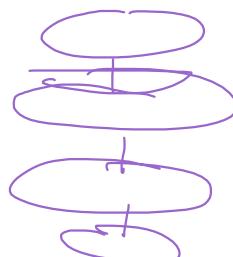
Avg bucket size = $\frac{\text{total pair}}{\text{total bucket}}$
load factor
thresh olding

load factor ≥ 2
rehash() double buckets

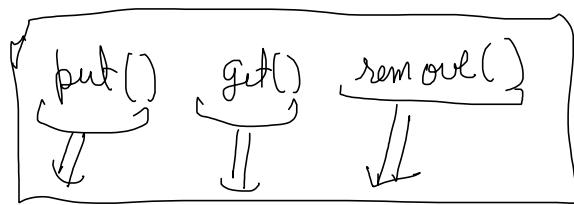
Q Java Hash Map solves collision ??

hash fⁿ \rightarrow $O(\text{no of bucket - 1})$

Chaining each bucket
is a Linked List *



Initial 16
thresh old
0.75



[load factor] \rightarrow avg

$O(n)$

Heap ✓
HashMap ✓

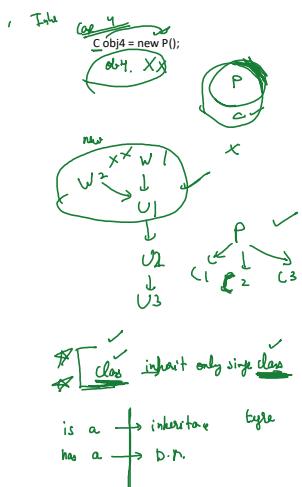
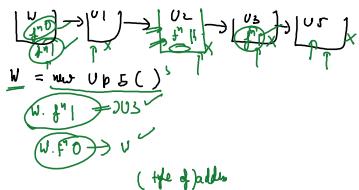
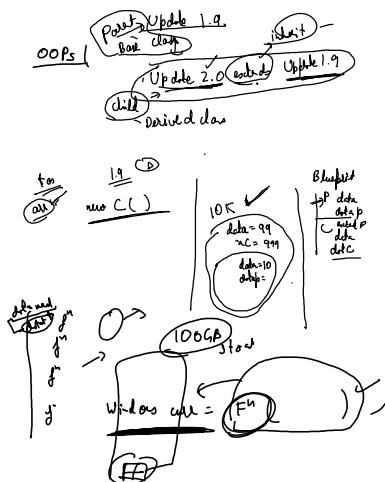
(ht)

7

(32 pointer)

[6,4,6,2,5,4,8,3,7]

From: <<https://leetcode.com/problems/container-with-most-water/>>



Value Class

Car is a Value

H is a Mammal

H has a Mammal

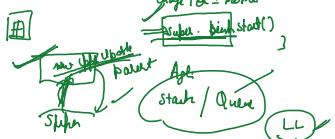
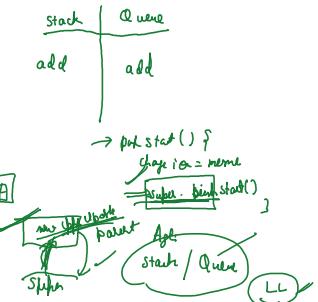
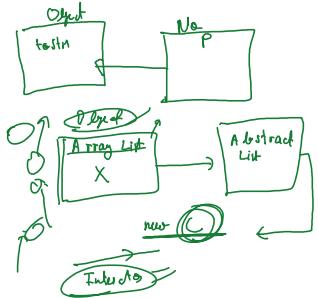
Mammal $\xrightarrow{\text{is a}}$ Living Being

Object

\downarrow
 $\text{new } p$

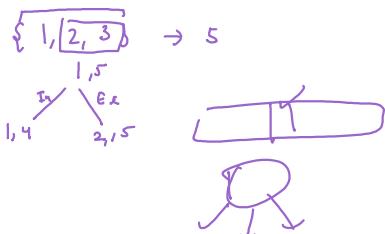
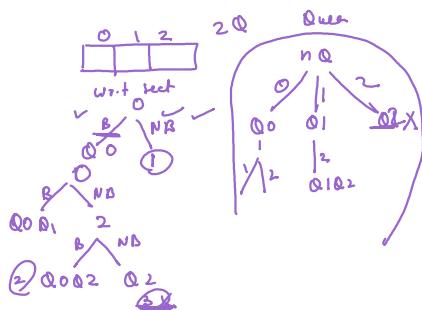
D.LN (" r. ")

\downarrow
 new P
 \downarrow
 Println(x) {
 Value of (x) \rightarrow ② tostring()

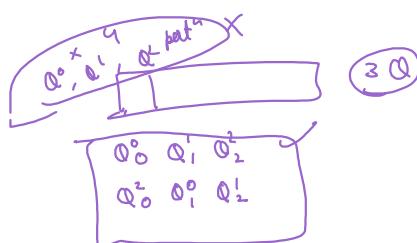


Geim Den ommeint"

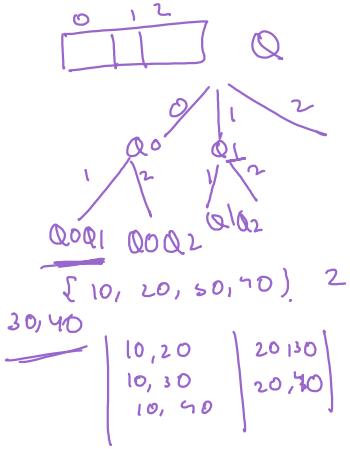
{ 1, 2, 3 } 5
|||| Team



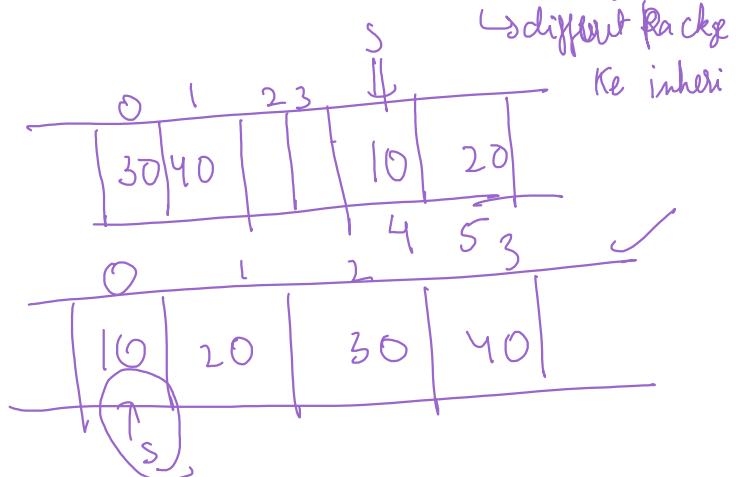
3 Q



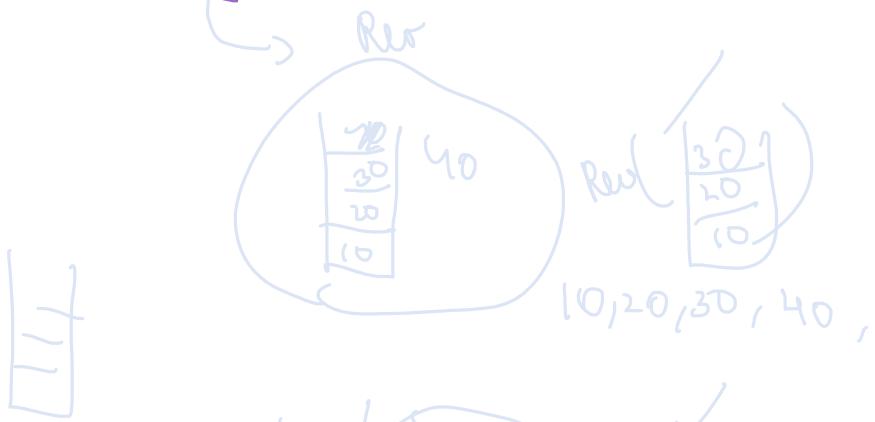
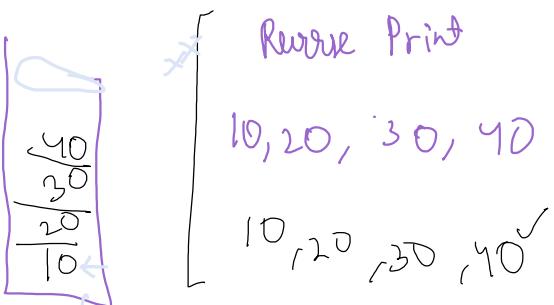
0 1 2
Q



protected
↳ same
package



Q



```

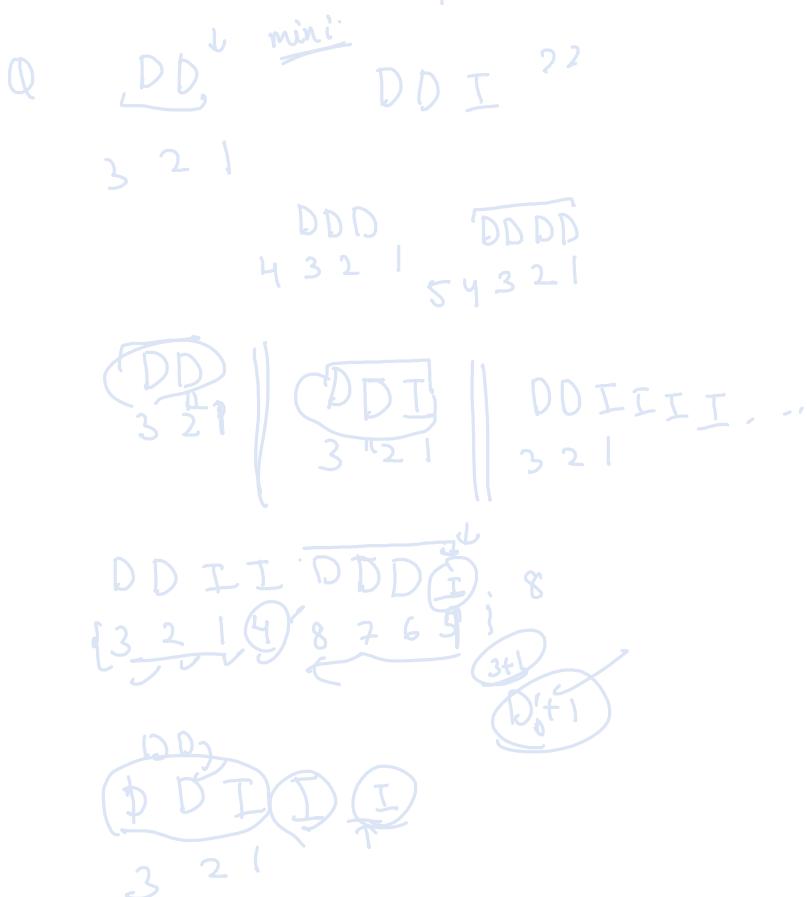
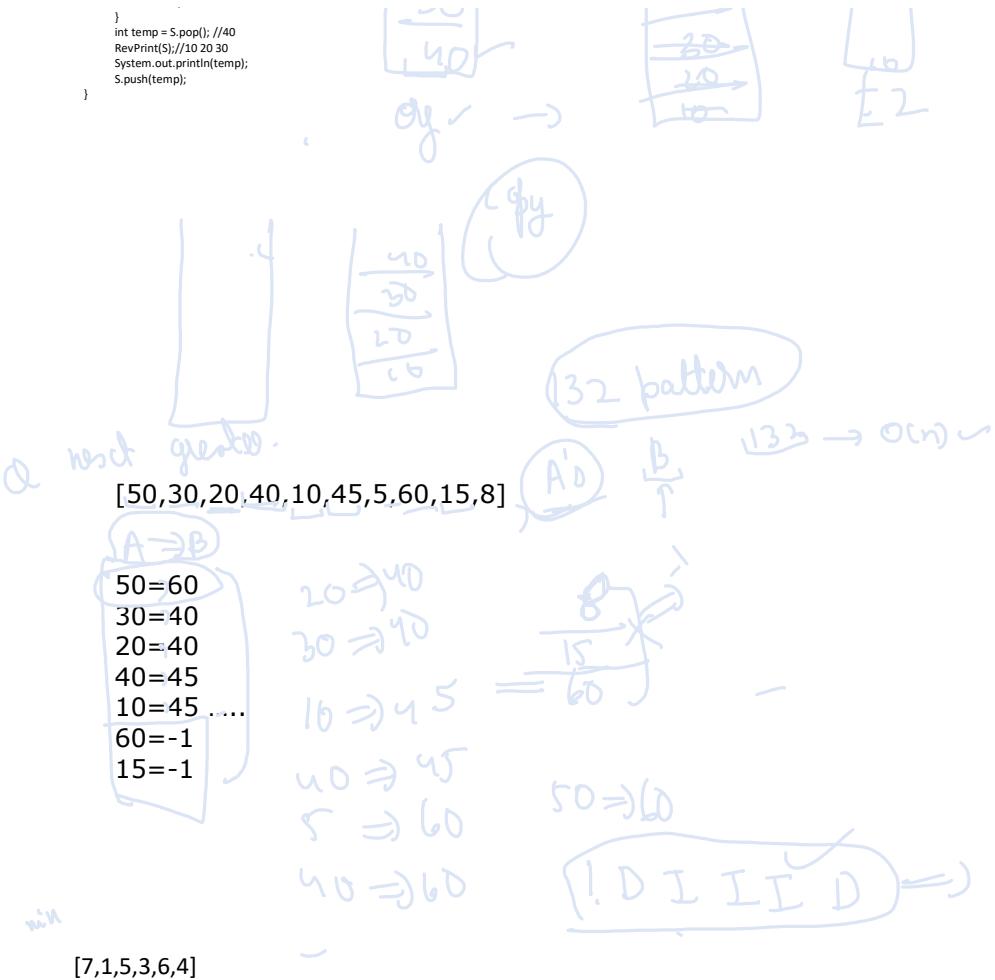
public static void RevPrint(Stack<Integer> S) {
    if(S.isEmpty()) {
        return;
    }
    int temp = S.pop(); // 40
    RevPrint(S); // 10 20 30
    System.out.println(temp);
    S.push(temp);
}
  
```

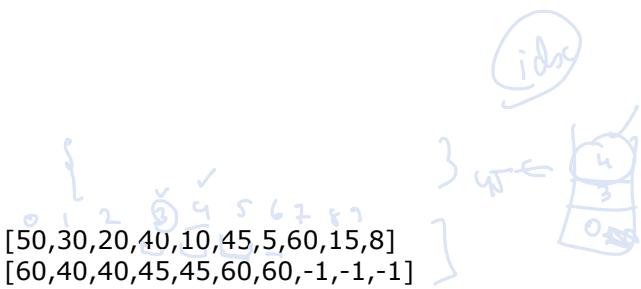


```

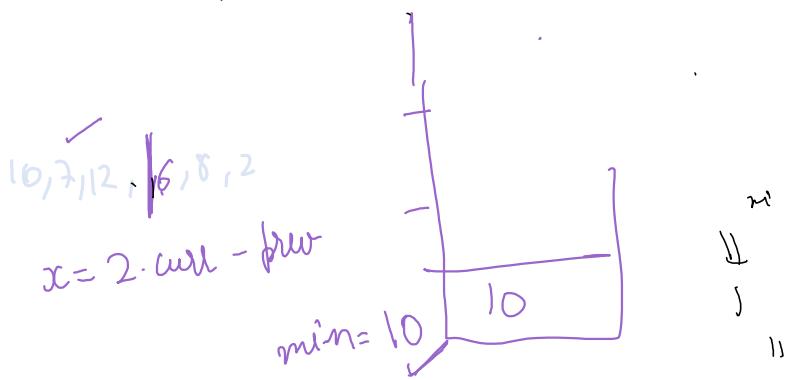
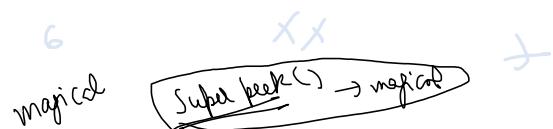
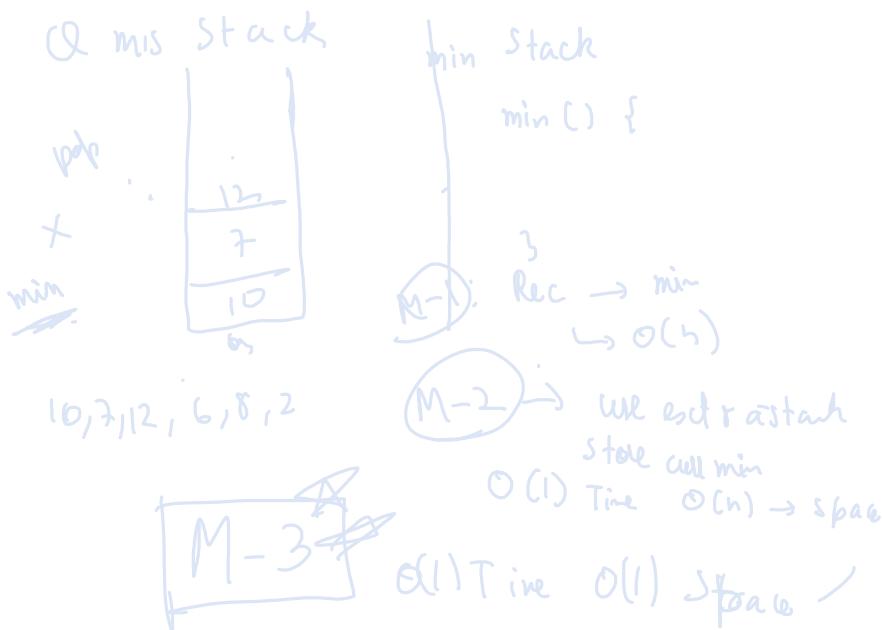
    }
    int temp = S.pop(); //40
    RevPrint(S); //10 20 30
    System.out.println(temp);
    S.push(temp);
}

```





do this in $O(n)$



$$pw = 2 \cdot curl - x$$

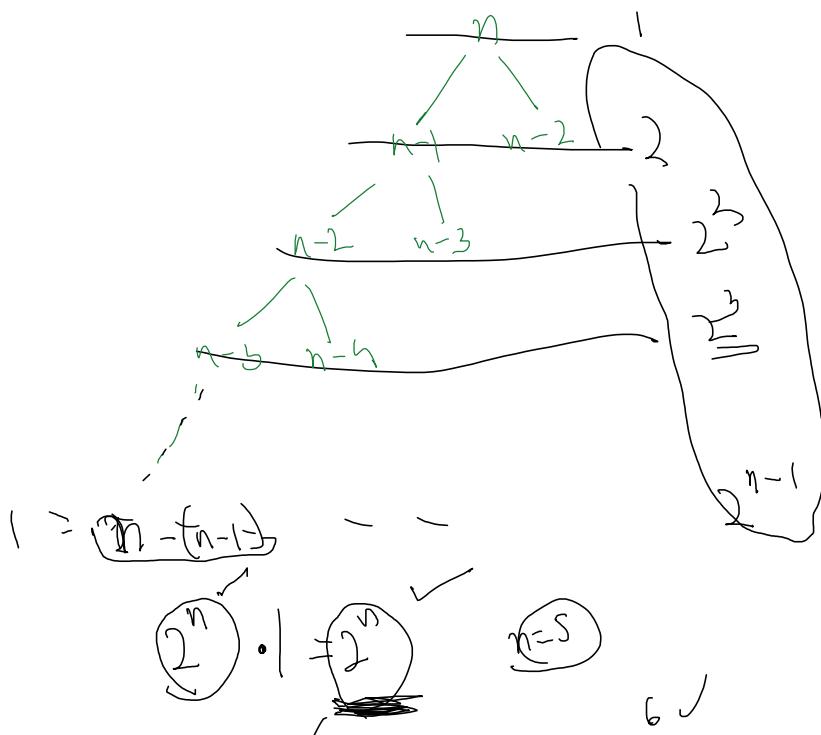
$$2 \cdot 7 - 4 \\ 10$$

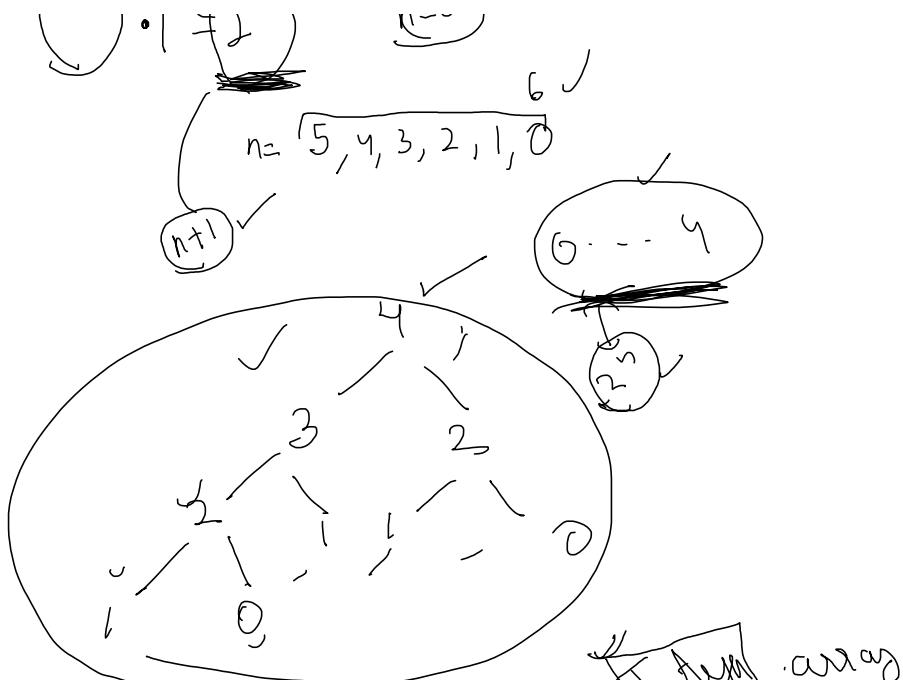
$$6 \times 2 - 5 = 12 - 5 = 7$$

0, 1, 1, 2, 3, 5, 8,
 0 1 2 3 4 5 6

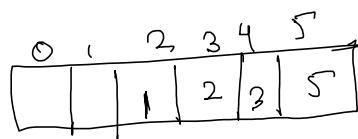
$\text{BP} \Rightarrow \underbrace{\text{Fib}(n)}$
 $\text{SP} \Rightarrow \underbrace{\text{Fib}(n-1)}$
 $\text{SP}_2 = \text{Fib}(n-2)$
 return $\text{SP}_1 + \text{SP}_2,$

$$\begin{aligned} T(n) &= \underbrace{T(n-1)}_{=} + \underbrace{T(n-2)}_{=} + 0 \\ &= \dots T(n-1) + T(n-1) + 1 \\ \boxed{T(n) = 2T(n-1) + 1} &\quad \downarrow \\ &\quad 2^n \end{aligned}$$





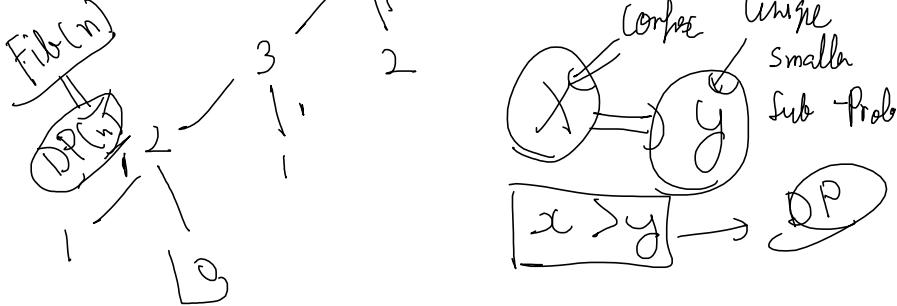
```
public static int Fibo(int n, int[] dp) {
    if (n <= 1) {
        return n;
    }
    BP : Fibo(n)
    SP : Fibo(n-1)
    if (dp[n] != 0) {
        return dp[n];
    }
    int sp1 = Fibo(n - 1);
    int sp2 = Fibo(n - 2);
    dp[n] = sp1 + sp2;
}
```



Integers .array

n

$O(n)$



$O(n)$

Solve ??

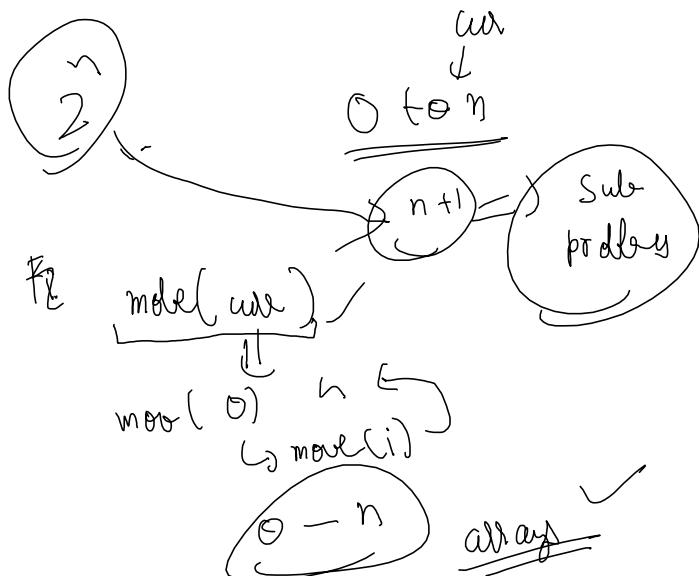
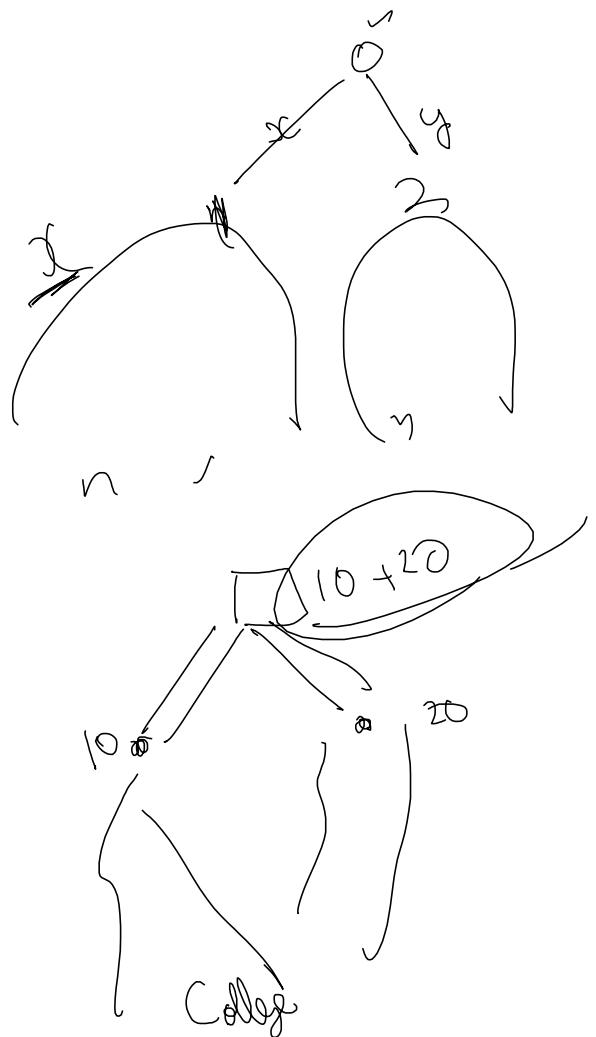
Biggest (smallest) fibo

$O(1)$

for (n = 0, n = nth; n++)

[]

1



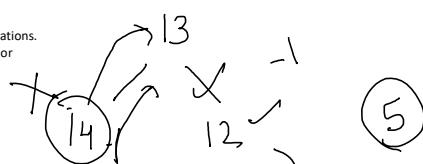
Given a positive number N your task is to bring this number to 1 by performing only a set of operations.
The operations can be either dividing the number by 2 only if the number is even or you can add or subtract 1 only if the number is odd.

More Precisely:

1) $N=N/2$ (if N is even)

2) $N=N-1$ (if N is odd)

Your task is to minimize these number of operations.



SOURCE & ONLY IF THE NUMBER IS EVEN:

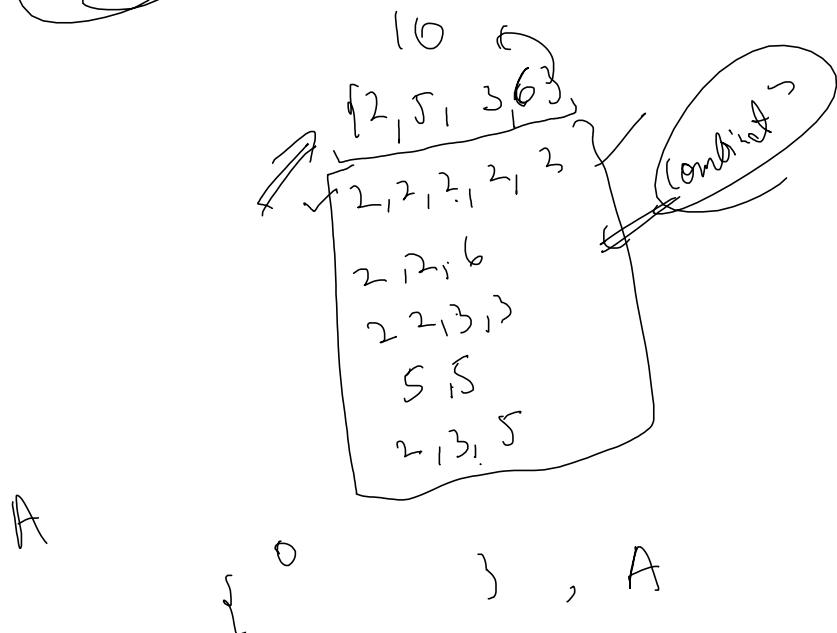
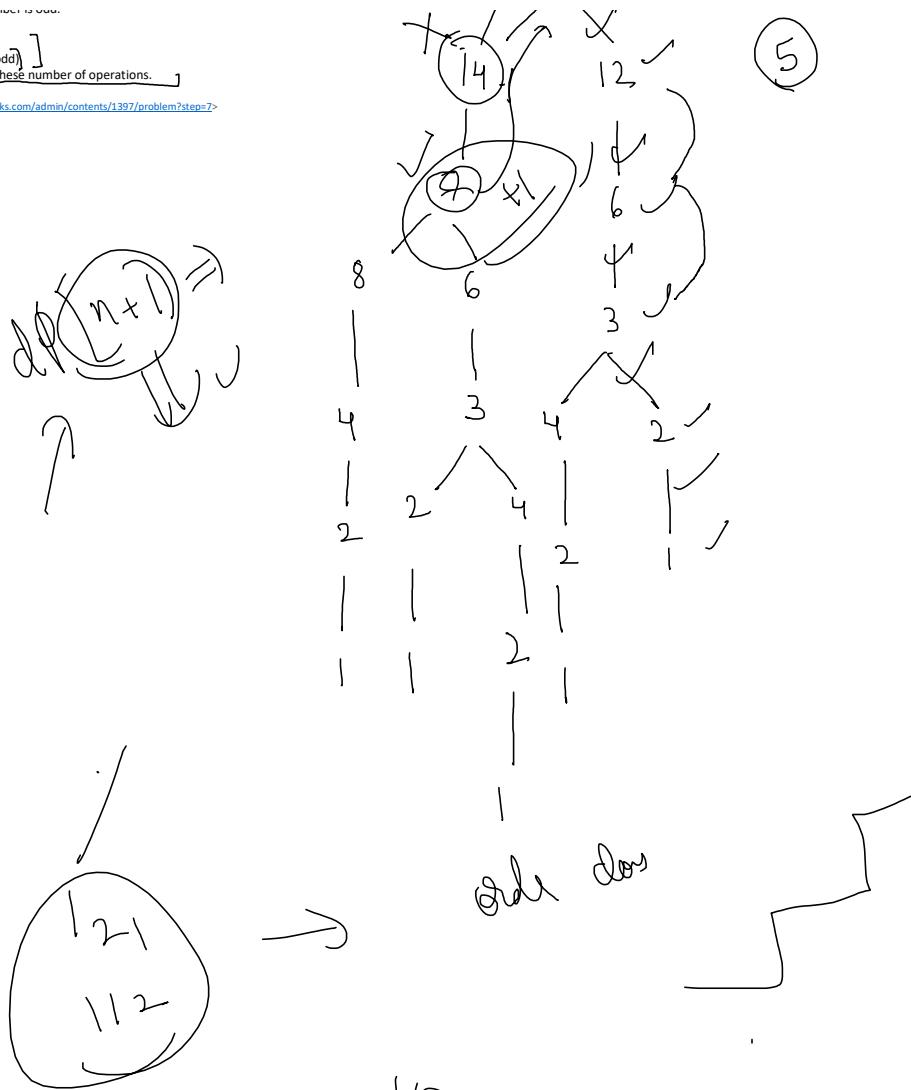
More Precisely:

1) $N=N/2$ (If N is even)

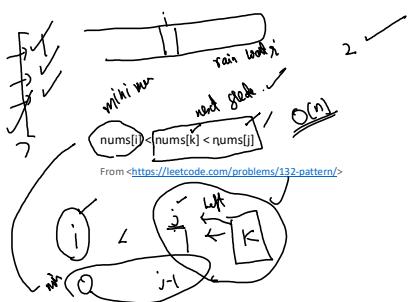
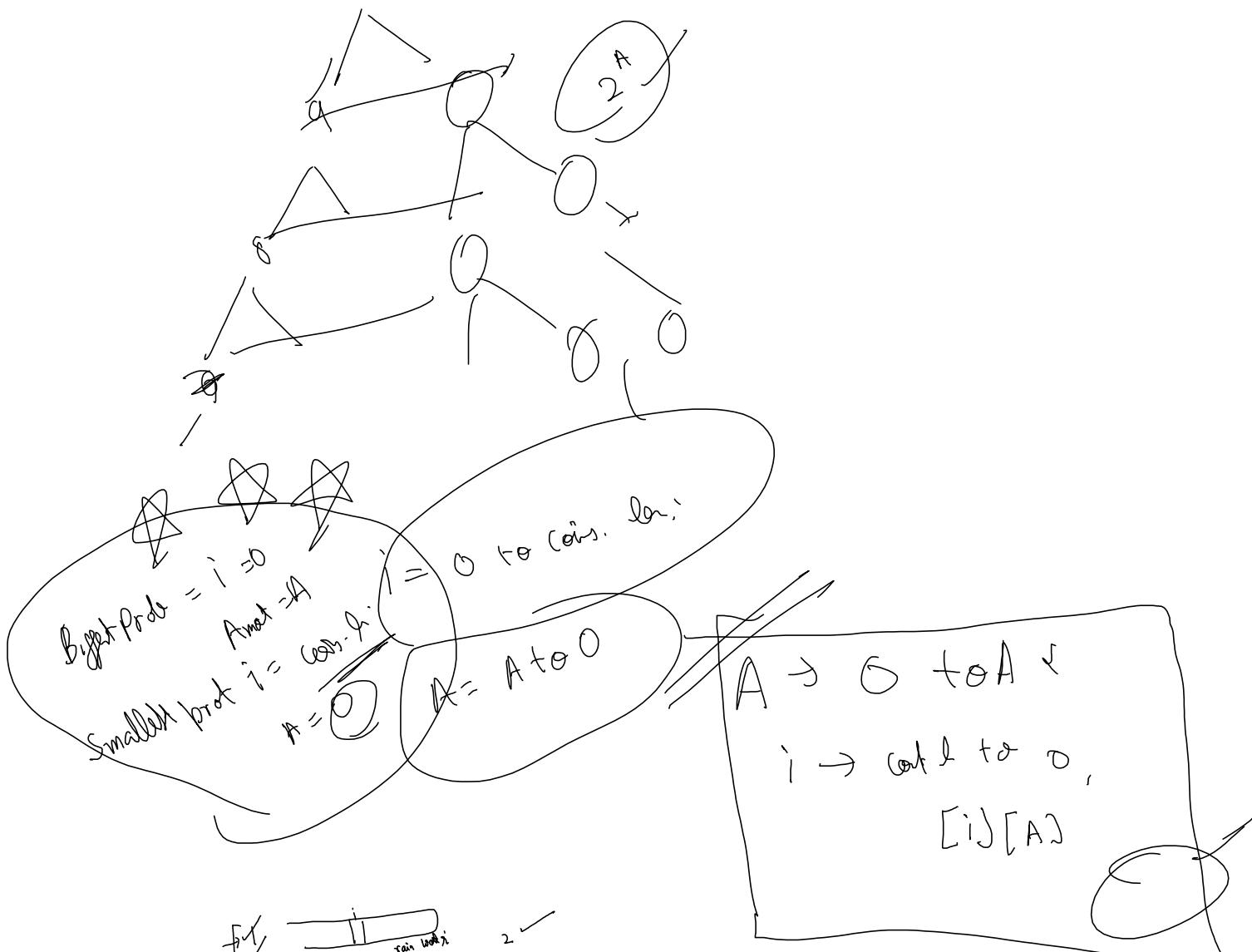
2) $N=N-1$ (If N is odd)

Your task is to minimize these number of operations.

From <<https://hack.codingblocks.com/admin/contents/1397/problem?step=7>>

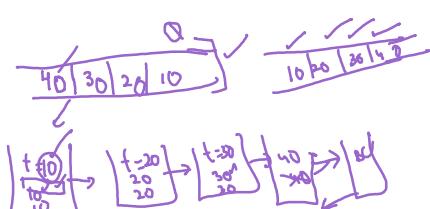


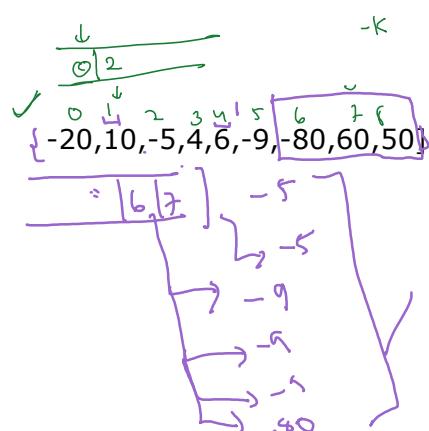
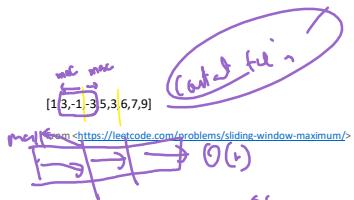
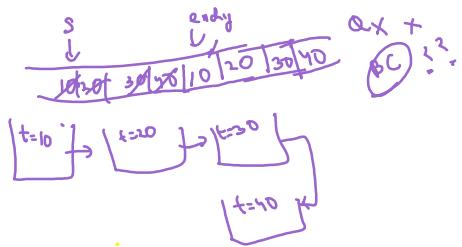
$$\{ \underline{1}, \underline{2}, \underline{3} \} \xrightarrow{10}$$



10 20 30 0(10) 50
 Restore Display 50, 40, 30, 20, 10 ✓
 SP → 50, 40, 30, 20, 10 ✓

0. call (10)





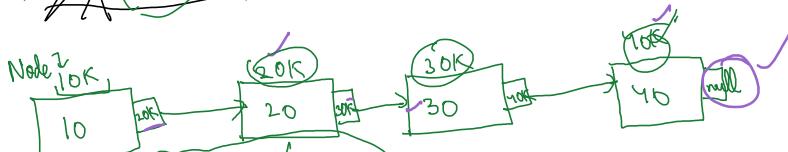
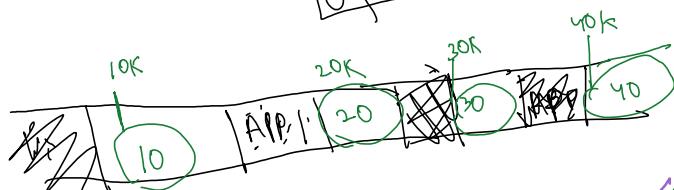
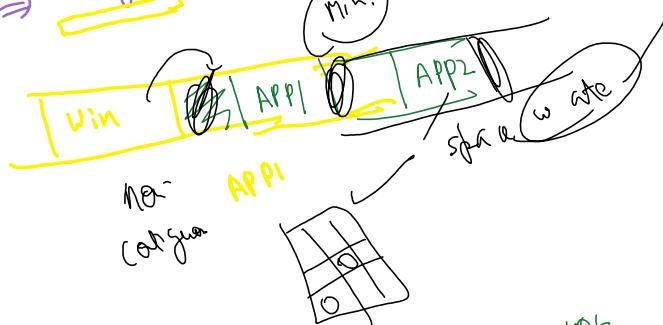
Data Structure

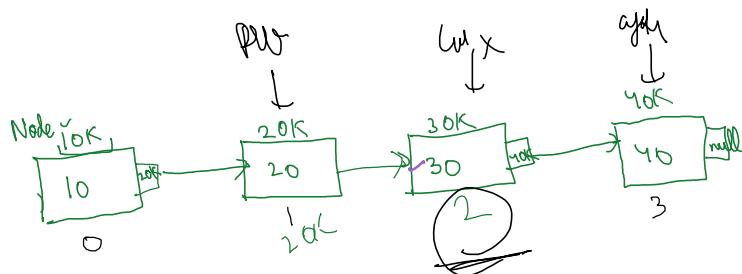
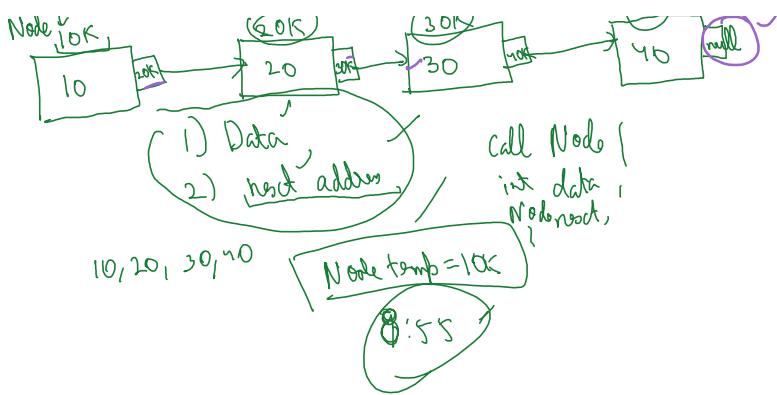
A.L \rightarrow array

Stack

Que

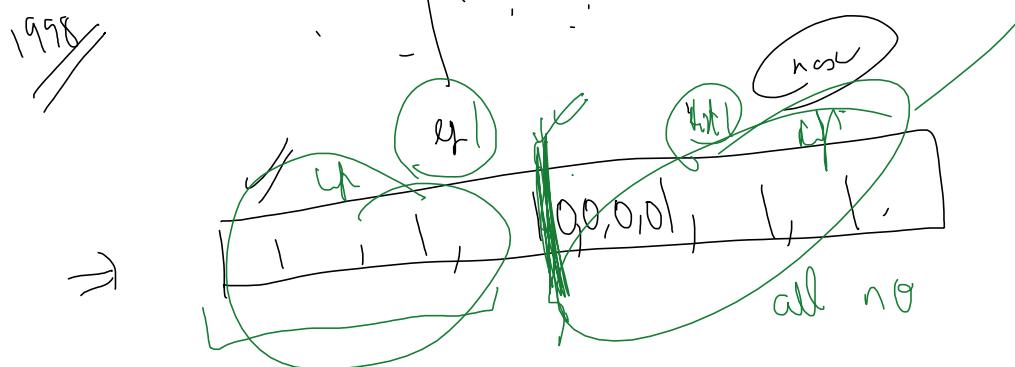
Q \Rightarrow Continuous \Rightarrow it's $O(1)$





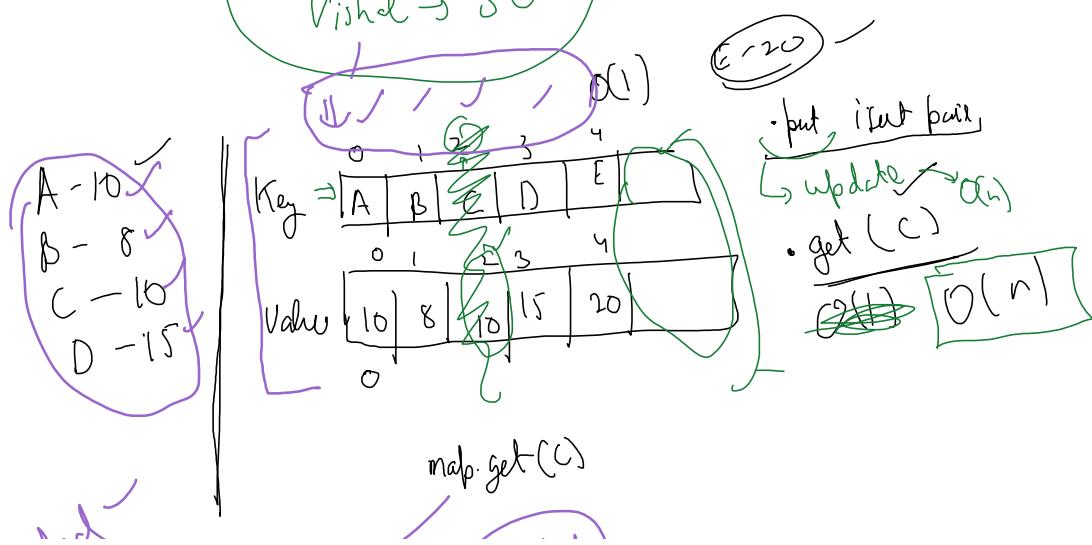
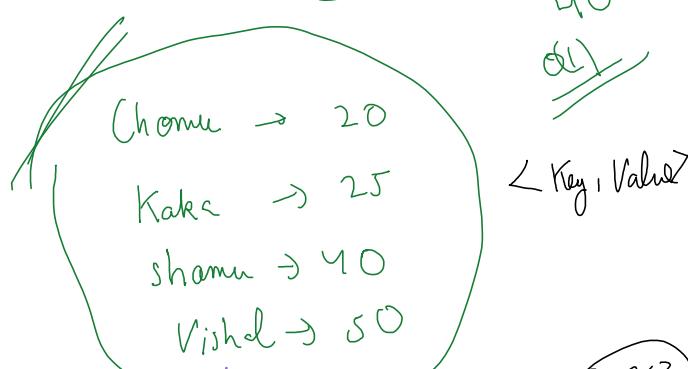
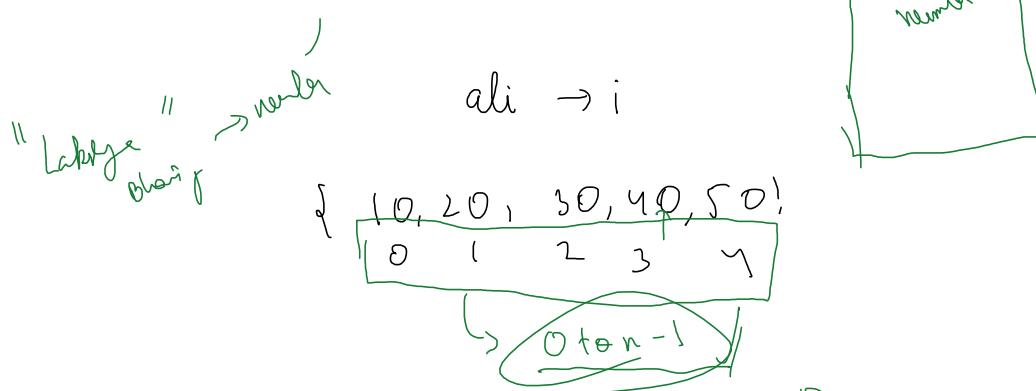
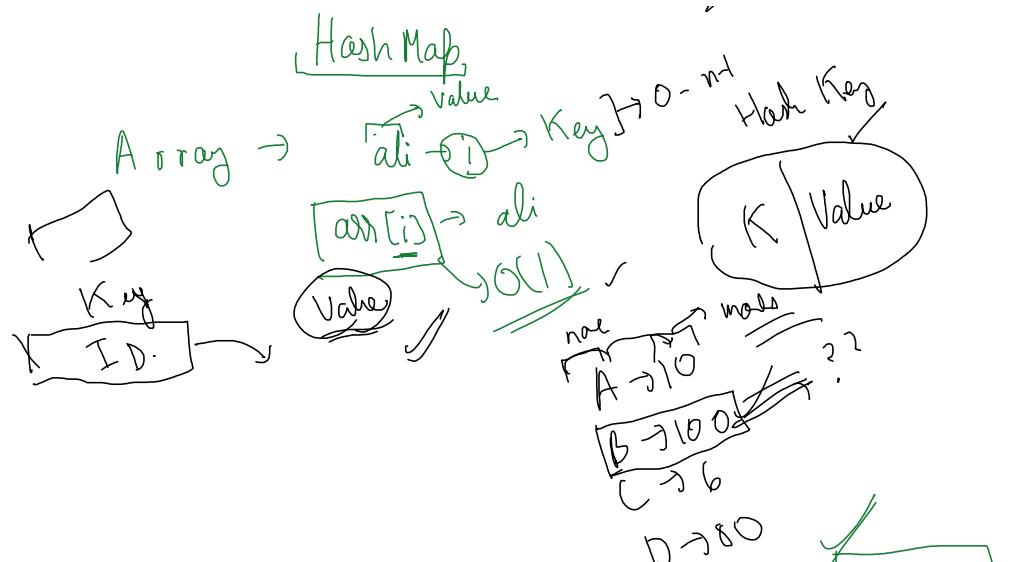
```
public void addAt(int idx, int all) {
    Node prev = getNode(idx - 1);
    Node nn = new Node(all);
    prev.next = nn;
    nn.next = prev.next;
}
```

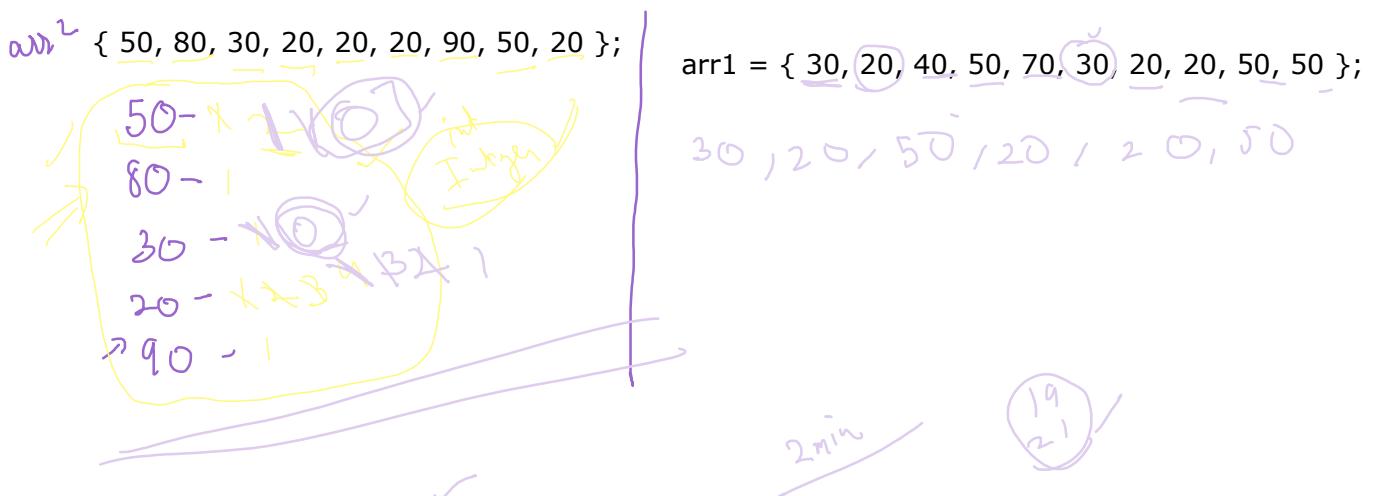
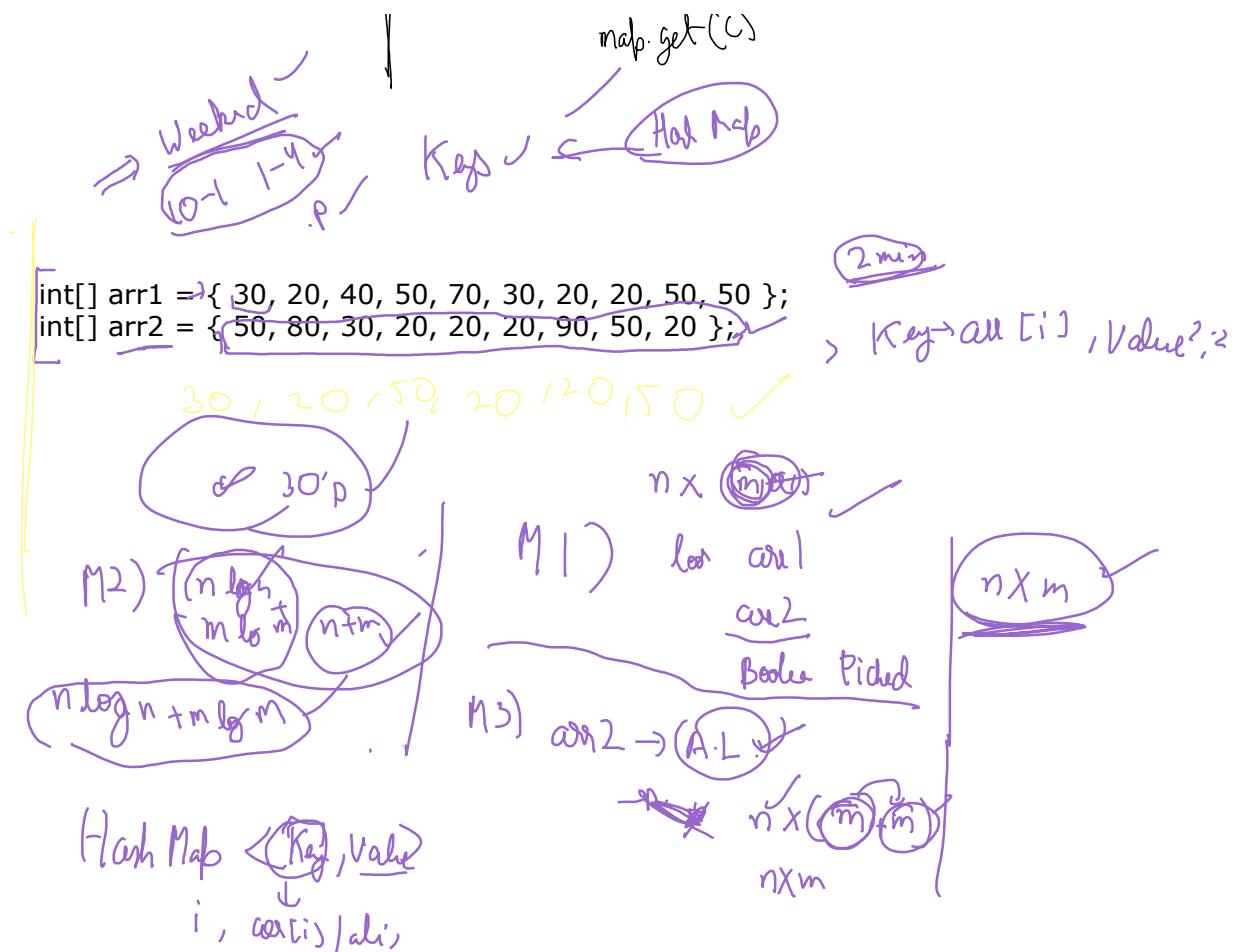
prev. next = curr



HashMap
value

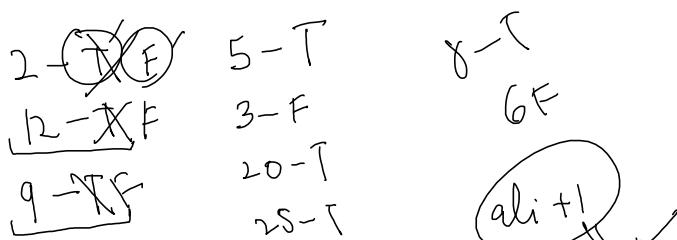
~ O~nd, 1 Key.





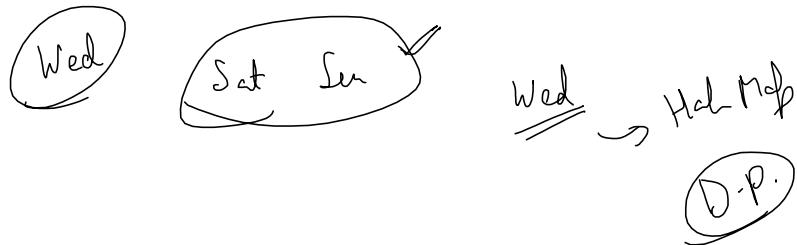
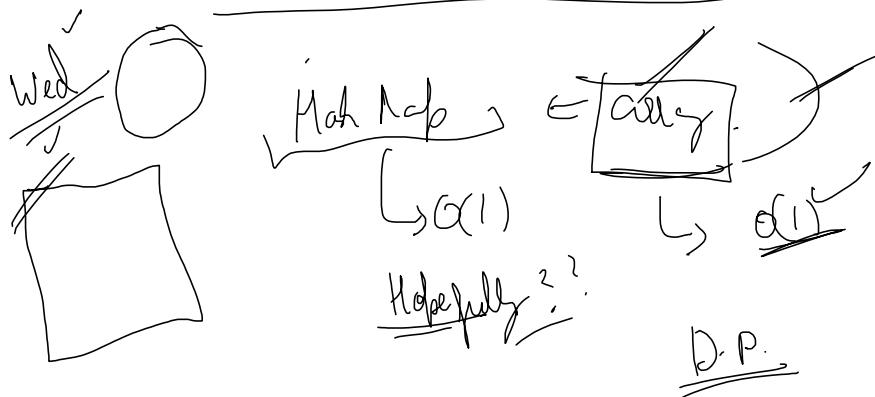
- Q Given an array of Integers in no particular order. Write a Program to find the longest possible sequence of consecutive numbers using the numbers from the array.

Input : [2, 12, 9, 16, 10, 5, 3, 20, 25, 11, 1, 8, 6]
 Output : [8, 9, 10, 11, 12]



$$\begin{array}{lll}
 \overbrace{\quad}^9 \rightarrow F & 20-1 & \text{ali} + 1 \\
 16-1 & 25-1 & \downarrow \\
 10 \checkmark & 11-F & F \\
 1-\checkmark & &
 \end{array}$$

Intya Boolean
ali ← is the starting point?



$$B^P \rightarrow Fib(n)$$

$$SP \rightarrow Fib(n^{-1})$$

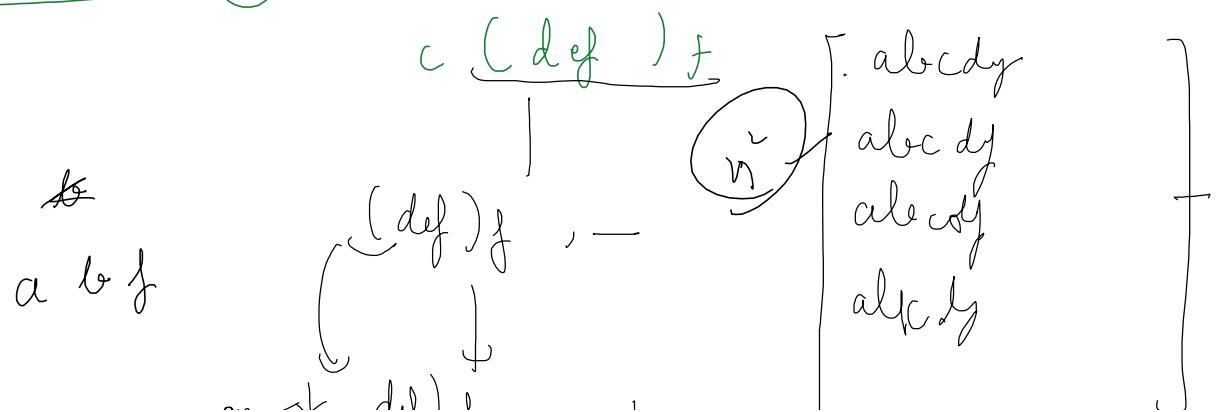
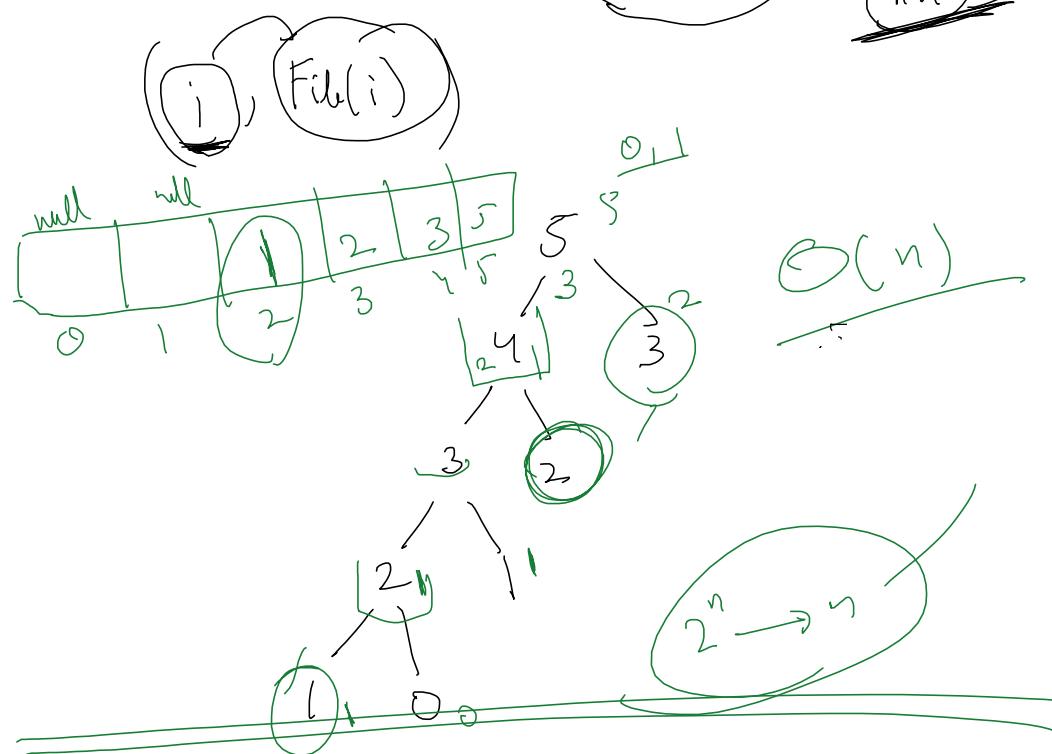
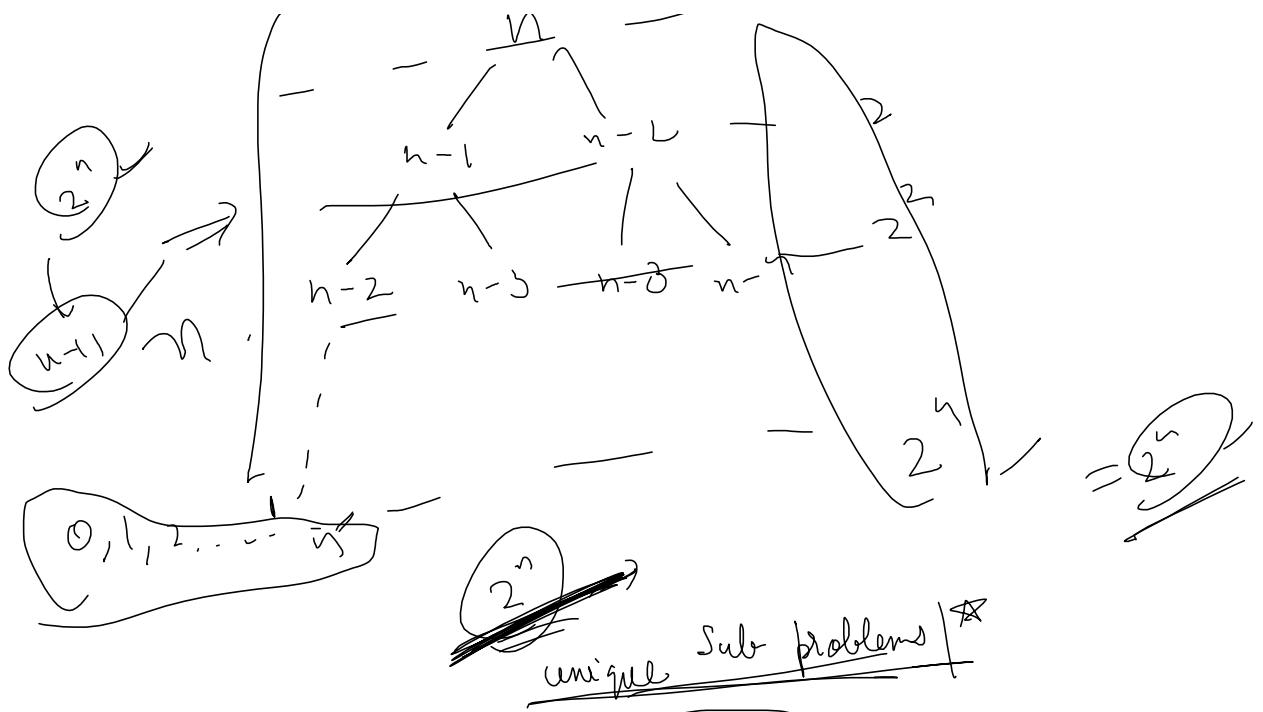
Fib(n^{-1})

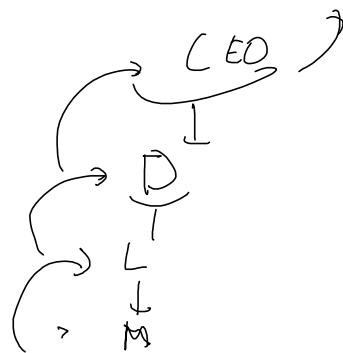
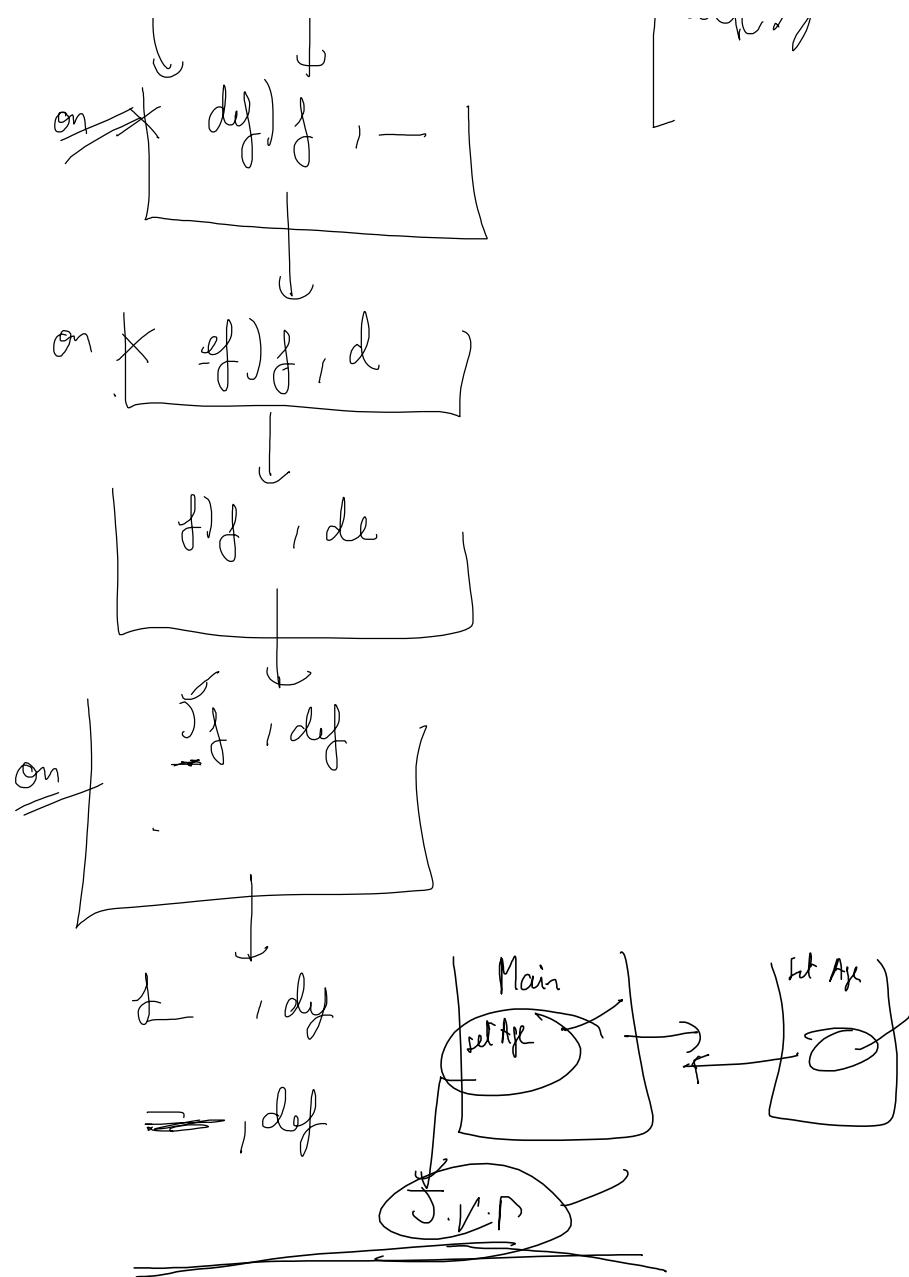
$$T(n) = T(n-1) + T(n-2)$$

$$T(n) = 2T(n-1) + 1$$

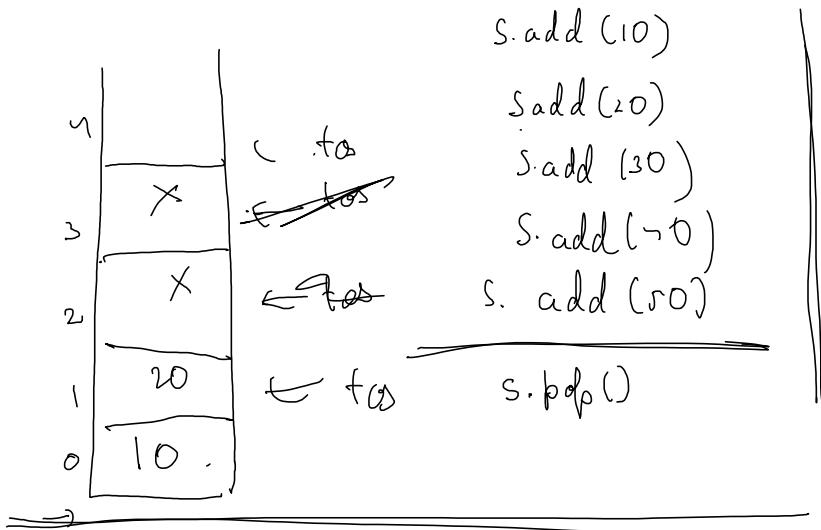
2 ↙



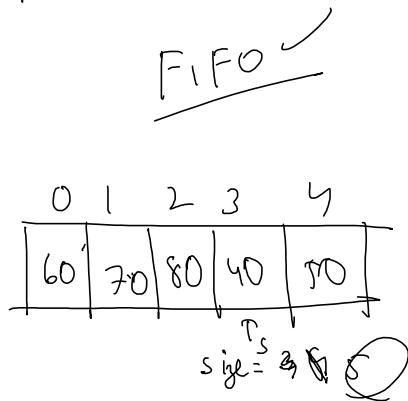
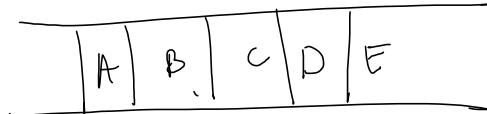




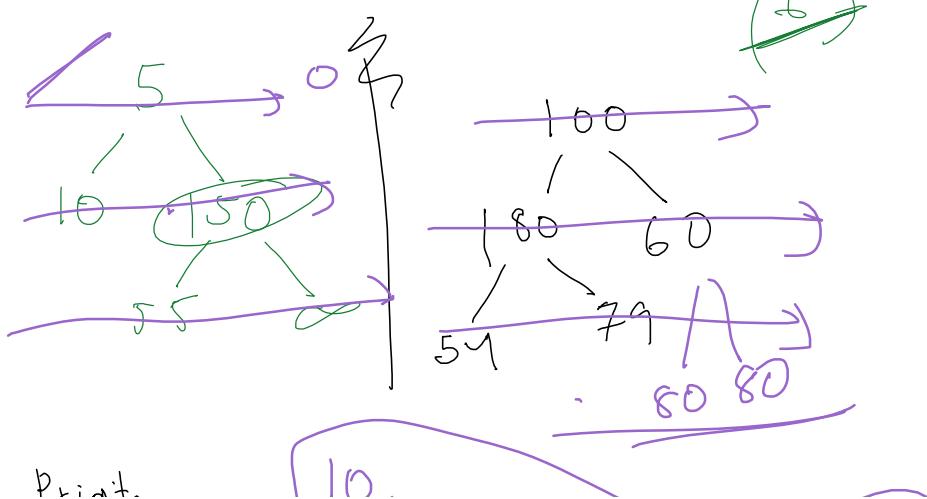
S.add(10)

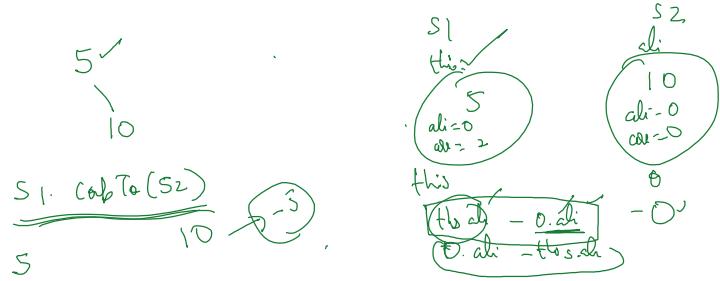


Queue $\xrightarrow{\text{add}}$ $\xleftarrow{\text{poll}}$

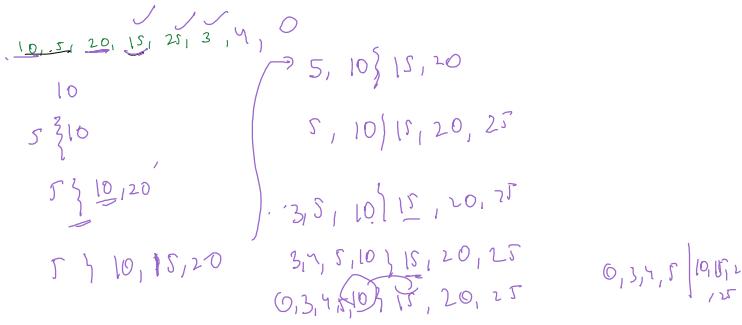
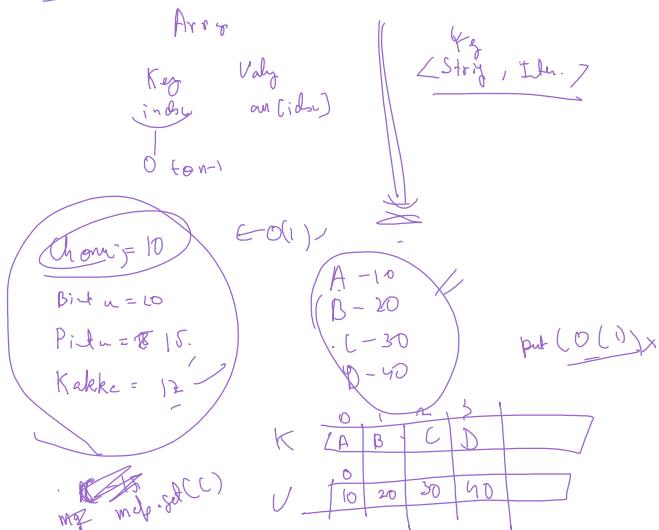


- Q.add(10)
Q.add(20)
Q.add(30)
Q.add(40)
Q.add(50)
Q.poll() \rightarrow 10
Q.poll() \rightarrow 20
Q.poll() \rightarrow 30

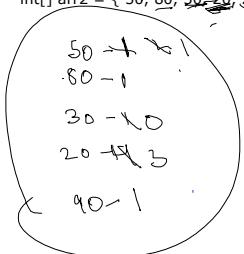
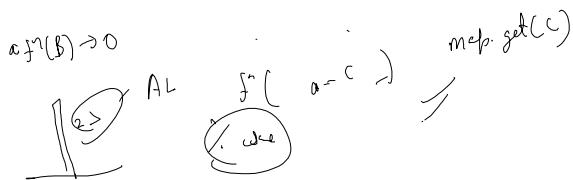
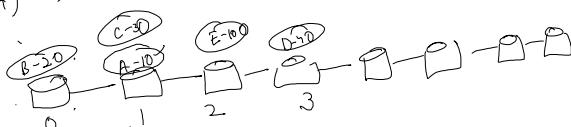




Rummy median
 of a street.
 $10, 15, 20, 15, 25, 30, 40$
 $10 \rightarrow 10$
 $5, 10 \rightarrow 7$
 $5, 10, 20 \rightarrow 10$
 $5, 10, 15, 20 \rightarrow 12$
 $5, 10, 15, 20, 25 \rightarrow 15$
 $5, 10, 15, 20, 25, 30, 40 \rightarrow 17$
 $5, 10, 15, 20, 25, 30, 40 \rightarrow 2$

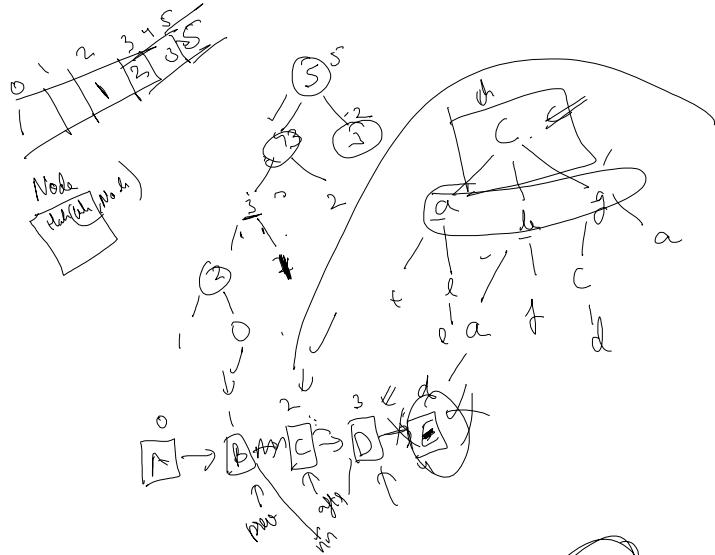
Hash Map

```
int[] arr1 = { 30, 20, 40, 50, 70, 30, 20, 20, 50, 50 };
int[] arr2 = { 50, 80, 30, 20, 20, 20, 90, 50, 20 };
```

 $f^n(\text{Key})$ $f^n(A) \Rightarrow 1$ $\text{Avg Bucket size falls}$ 

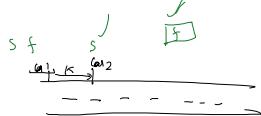
$$fib(n) = fib(n-1) + fib(n-2)$$

$$T(n) = T(n-1) + T(n-2) + 1$$



$O(2V)$ without

$M \times M \rightarrow O(n^2)$
 $M \rightarrow O(n^2 \text{ stem})$
 $O(n^2)$

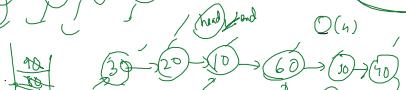


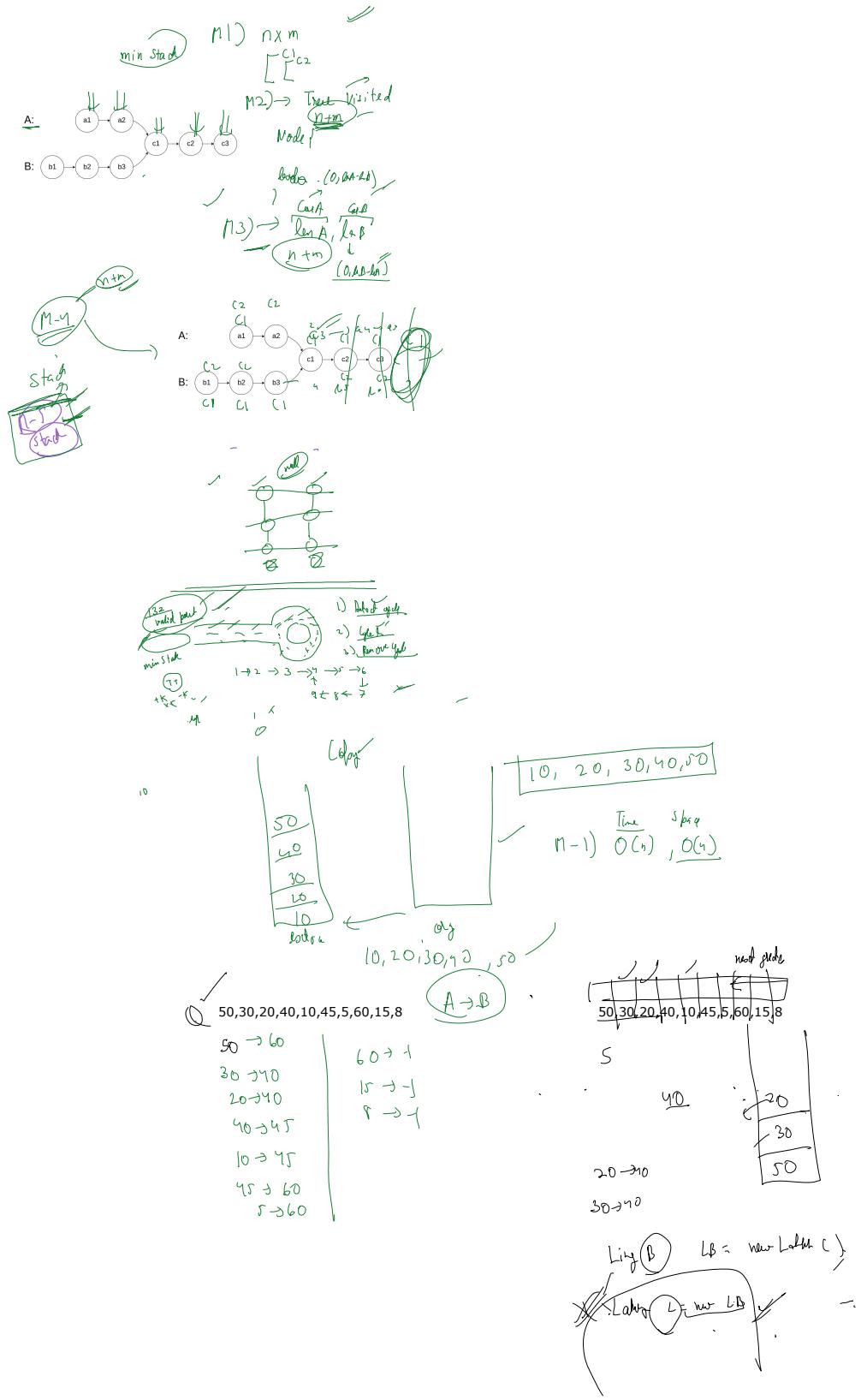
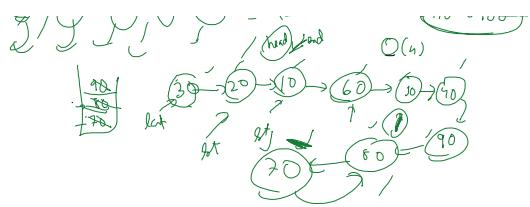
$K=2$
 $10 \rightarrow 20 \rightarrow 30$
 $O(2)$

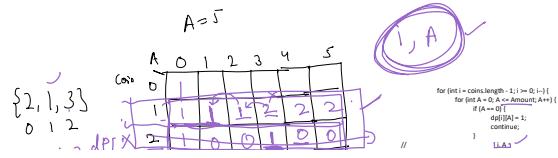
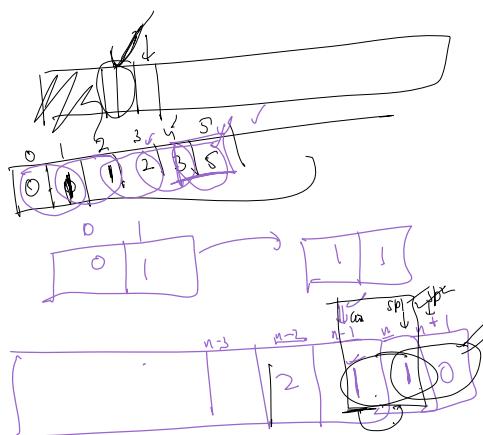
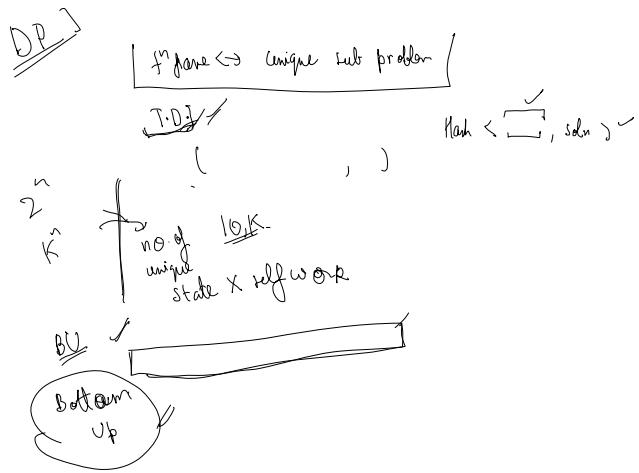
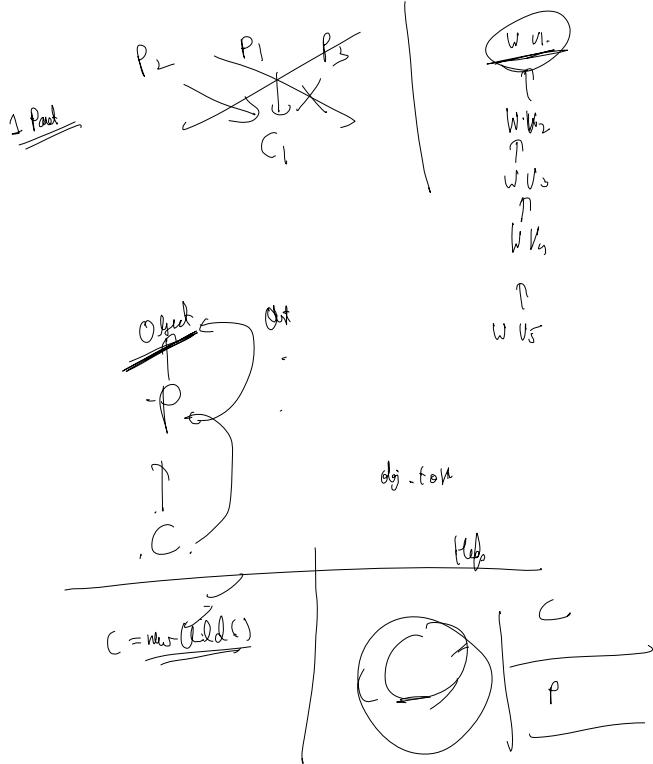
$K=1 \rightarrow 50$
 $K=2 \rightarrow 40$
 $K=3 \rightarrow 30$
 $O(1)$

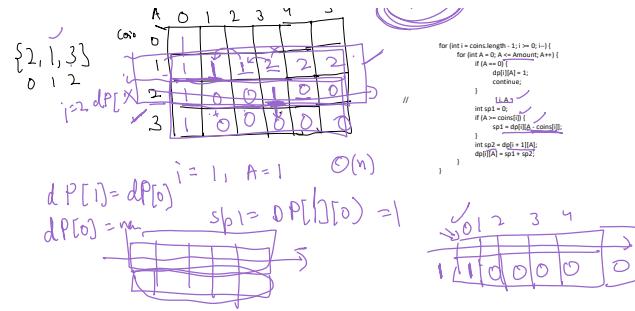


$k=3$
 $10 \rightarrow 20 \rightarrow 30 \rightarrow 40 \rightarrow 50 \rightarrow 60 \rightarrow 70 \rightarrow 80 \rightarrow 90 \rightarrow 100 \rightarrow 110$
 $30 \rightarrow 20 \rightarrow 10 \rightarrow 60 \rightarrow 50 \rightarrow 40 \rightarrow 90 \rightarrow 80 \rightarrow 70 \rightarrow 110 \rightarrow 100$









```

public static int minDistance_BFS(String word1, String word2) {
    int[][] dp = new int[word1.length() + 1][word2.length() + 1];
    for (int i = 0; i < word1.length(); i++) {
        for (int j = 0; j < word2.length(); j++) {
            if (i == 0 & j == 0) {
                dp[0][0] = 0;
            } else if (i == 0) {
                dp[0][j] = j;
            } else if (j == 0) {
                dp[i][0] = i;
            } else {
                dp[i][j] = dp[i - 1][j - 1];
            }
        }
    }
    for (int i = 1; i < word1.length(); i++) {
        for (int j = 1; j < word2.length(); j++) {
            if (word1.charAt(i) == word2.charAt(j)) {
                dp[i][j] = dp[i - 1][j - 1];
            } else {
                dp[i][j] = Math.min(dp[i - 1][j], dp[i][j - 1], dp[i - 1][j - 1]) + 1;
            }
        }
    }
    return dp[word1.length()][word2.length()];
}

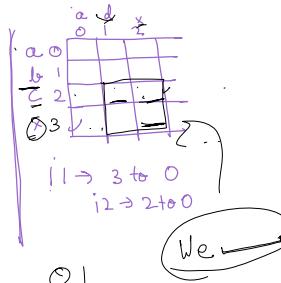
```

```

public static int minDistance_BFS(String word1, String word2) {
    int[][] dp = new int[word1.length() + 1][word2.length() + 1];
    for (int i = 0; i < word1.length(); i++) {
        for (int j = 0; j < word2.length(); j++) {
            if (i == 0 & j == 0) {
                dp[0][0] = 0;
            } else if (i == 0) {
                dp[0][j] = j;
            } else if (j == 0) {
                dp[i][0] = i;
            } else {
                dp[i][j] = dp[i - 1][j - 1];
            }
        }
    }
    for (int i = 1; i < word1.length(); i++) {
        for (int j = 1; j < word2.length(); j++) {
            if (word1.charAt(i) == word2.charAt(j)) {
                dp[i][j] = dp[i - 1][j - 1];
            } else {
                dp[i][j] = Math.max(dp[i - 1][j], dp[i][j - 1]);
            }
        }
    }
    return dp[word1.length()][word2.length()];
}

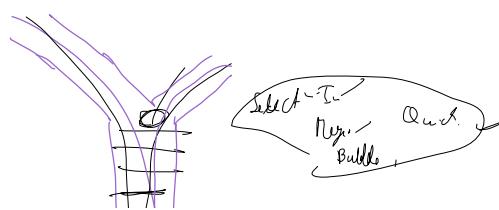
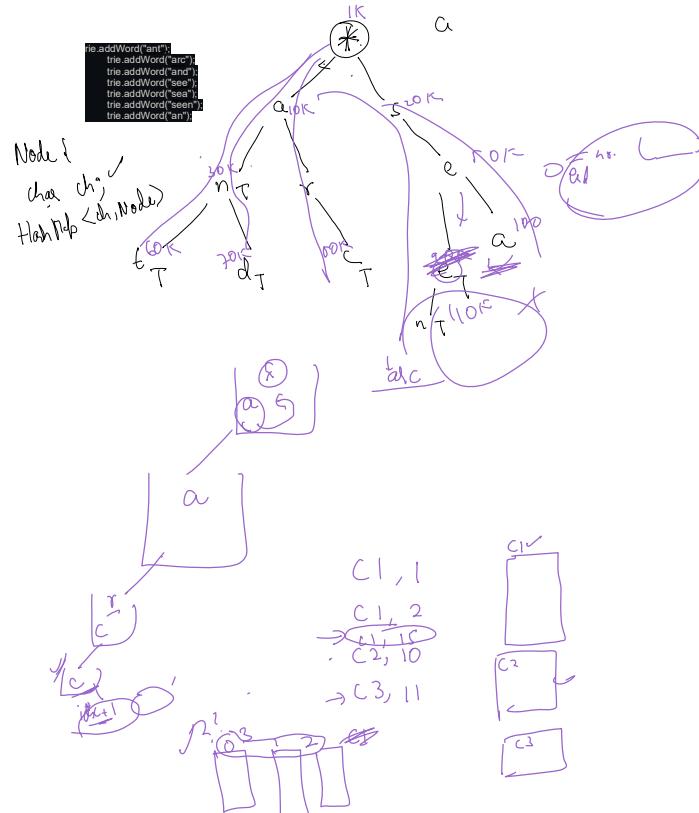
```

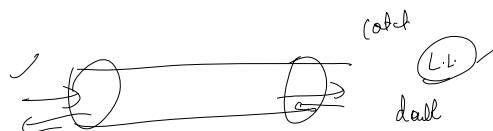
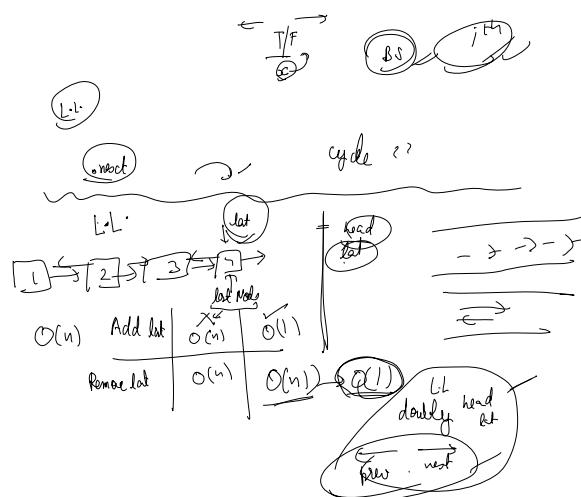
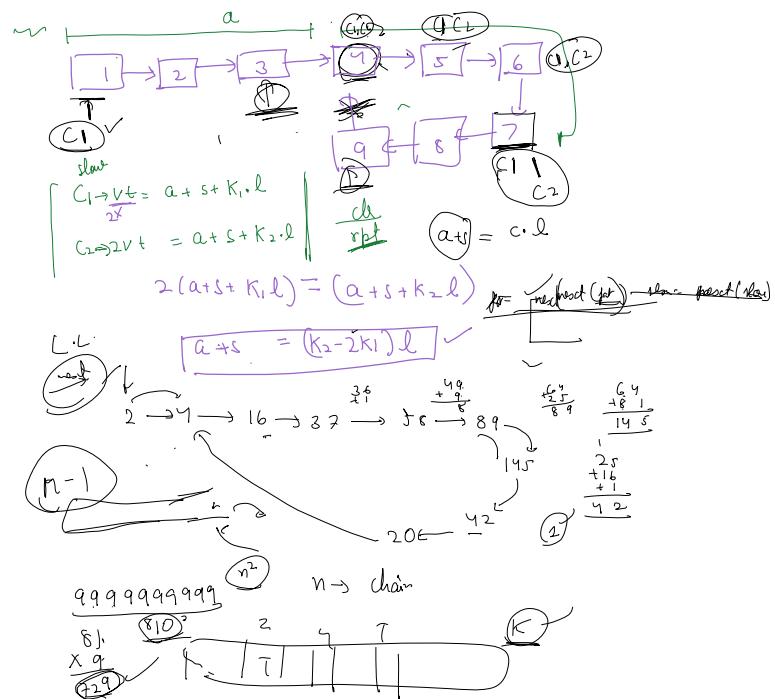
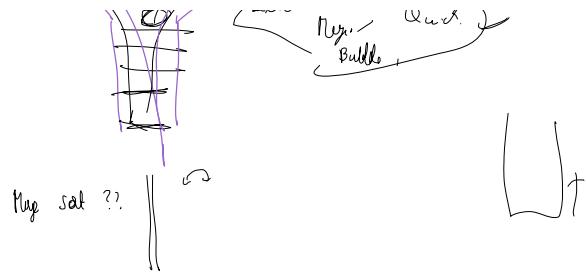
w1 "abc"
w2 "ad"

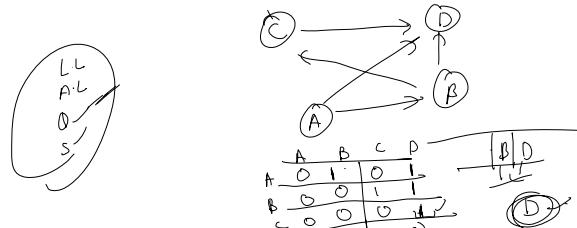


①

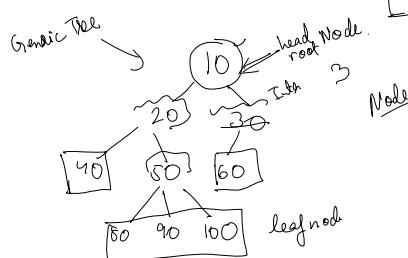
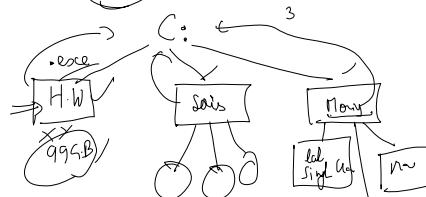
We



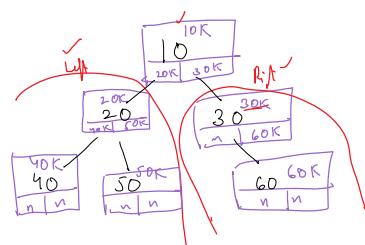
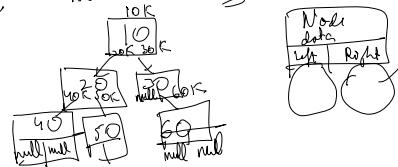




Tree
L-L array \times root node head node

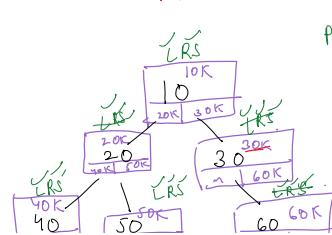


B: $L \rightarrow$ Node maximum \geq

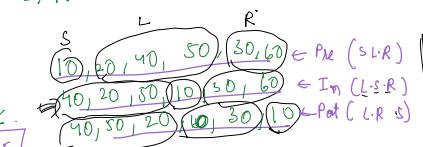


B : Print (10K)
SP : Print (20K)
Print (30K)
Symbol (10K)

10 F
null null

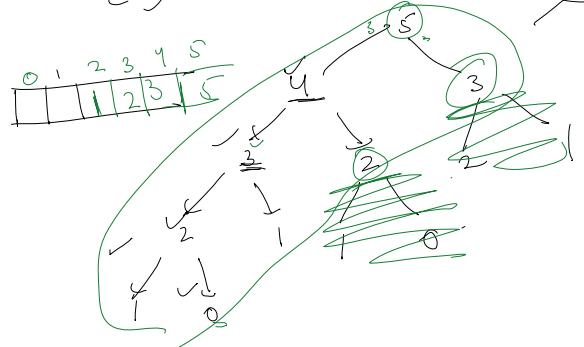
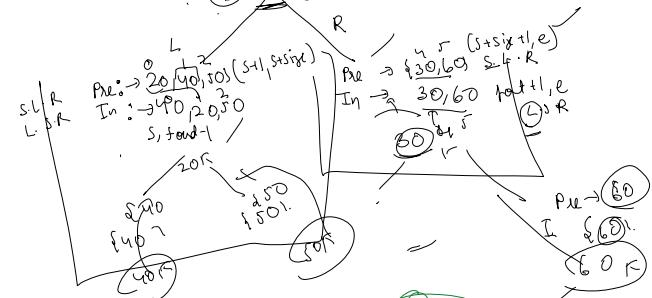
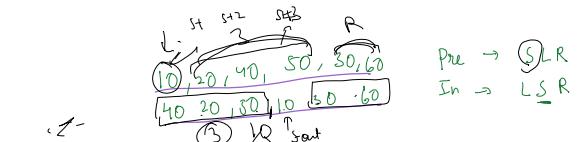
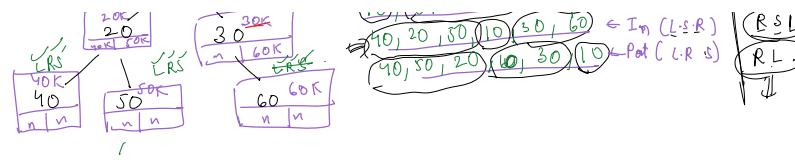


Print (n) \leftarrow S.L.R Pre order BT



translates

mn
spod
→ in
→ Pλν ✓

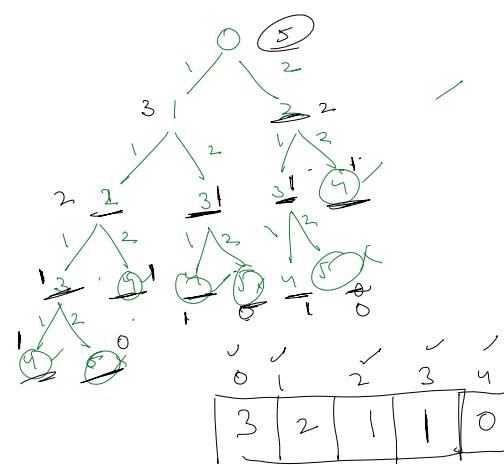


$\checkmark \text{ Fibo}(n) \rightarrow \text{Seq}$

$\text{dp}[i]$
Fibonacci(i)

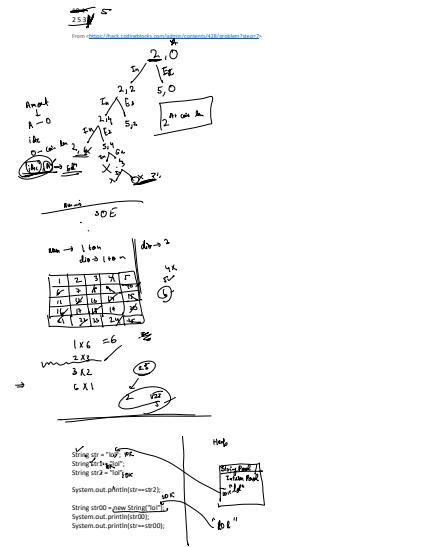
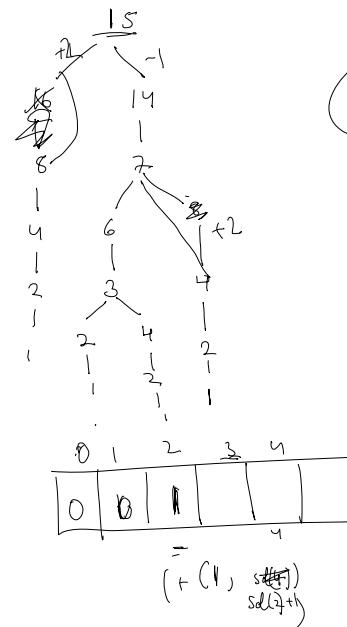
$\text{BPC}(n)$

$\text{sp} \quad \text{sp}_1(n-1)$
 $\text{sp}_2(n-2)$

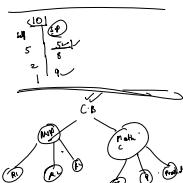


$\lambda \nu$

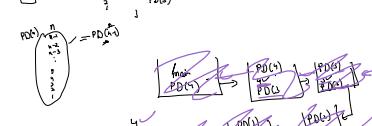
```
]] dp = new int[n+2];
[n]=1;
{int curr = n-1; curr >= 0; curr--) {
    imib(curr)
    int sp1 = dp[curr + 1];
    int sp2 = dp[curr + 2];
    dp[curr] = sp1 * sp2;
}
return dp[0];
```



- 1) Biggs Problem.
 - 2) Smaller Probleme
 - 3) Präsentation & PPT
 - 4) Biggs Sol" unters.



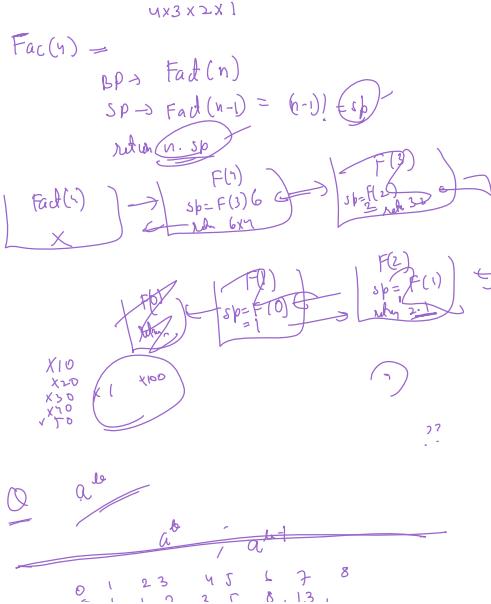
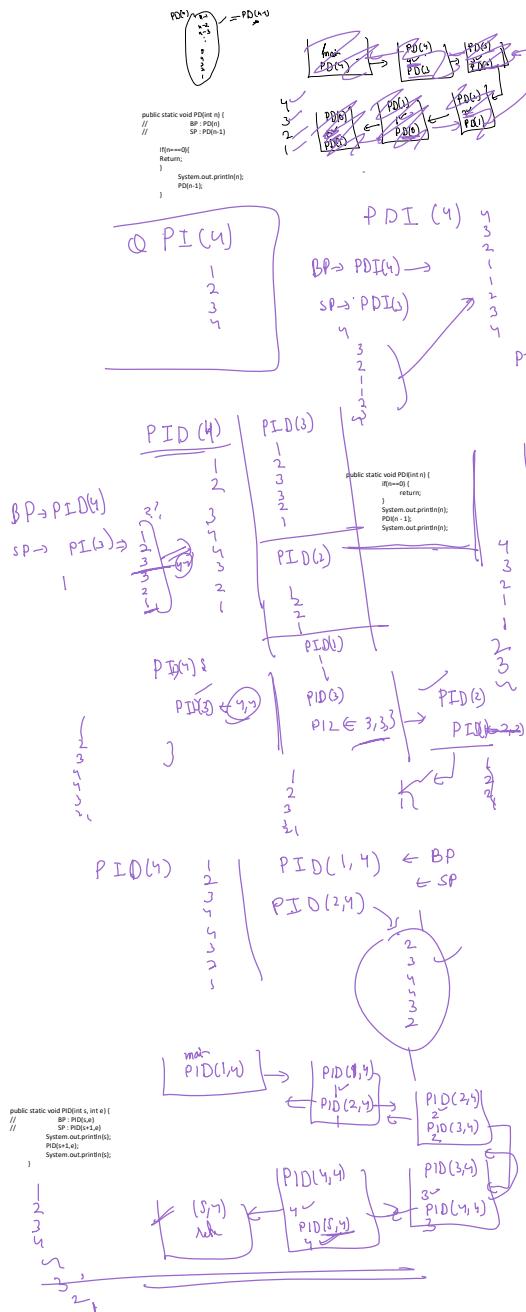
$$\begin{array}{c} \textcircled{8} \quad P D(4) \rightarrow \begin{matrix} 4 \\ 3 \\ 2 \\ 1 \end{matrix} \quad P D(5) \begin{matrix} 5 \\ 4 \\ 3 \\ 2 \\ 1 \end{matrix} \quad P D(6) \begin{matrix} 6 \\ 5 \\ 4 \\ 3 \\ 2 \\ 1 \end{matrix} \\ P D(3) \begin{matrix} 3 \\ 2 \\ 1 \end{matrix} \quad P D(2) \begin{matrix} 2 \\ 1 \end{matrix} \quad P D(1) \begin{matrix} 1 \\ \end{matrix} \\ P D(1) \begin{matrix} 1 \\ \end{matrix} \end{array}$$

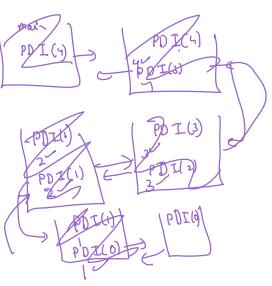
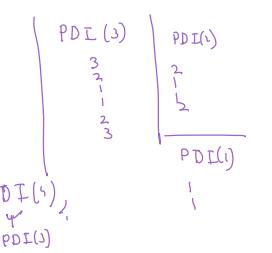


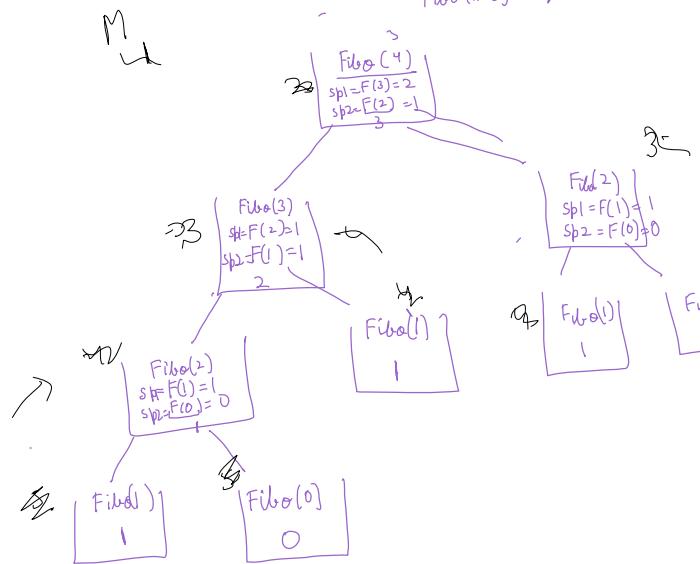
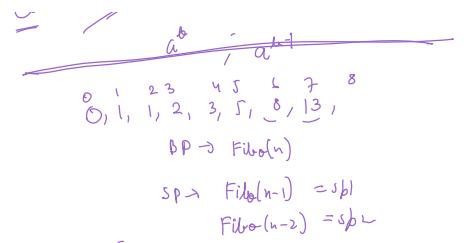
17
16

+

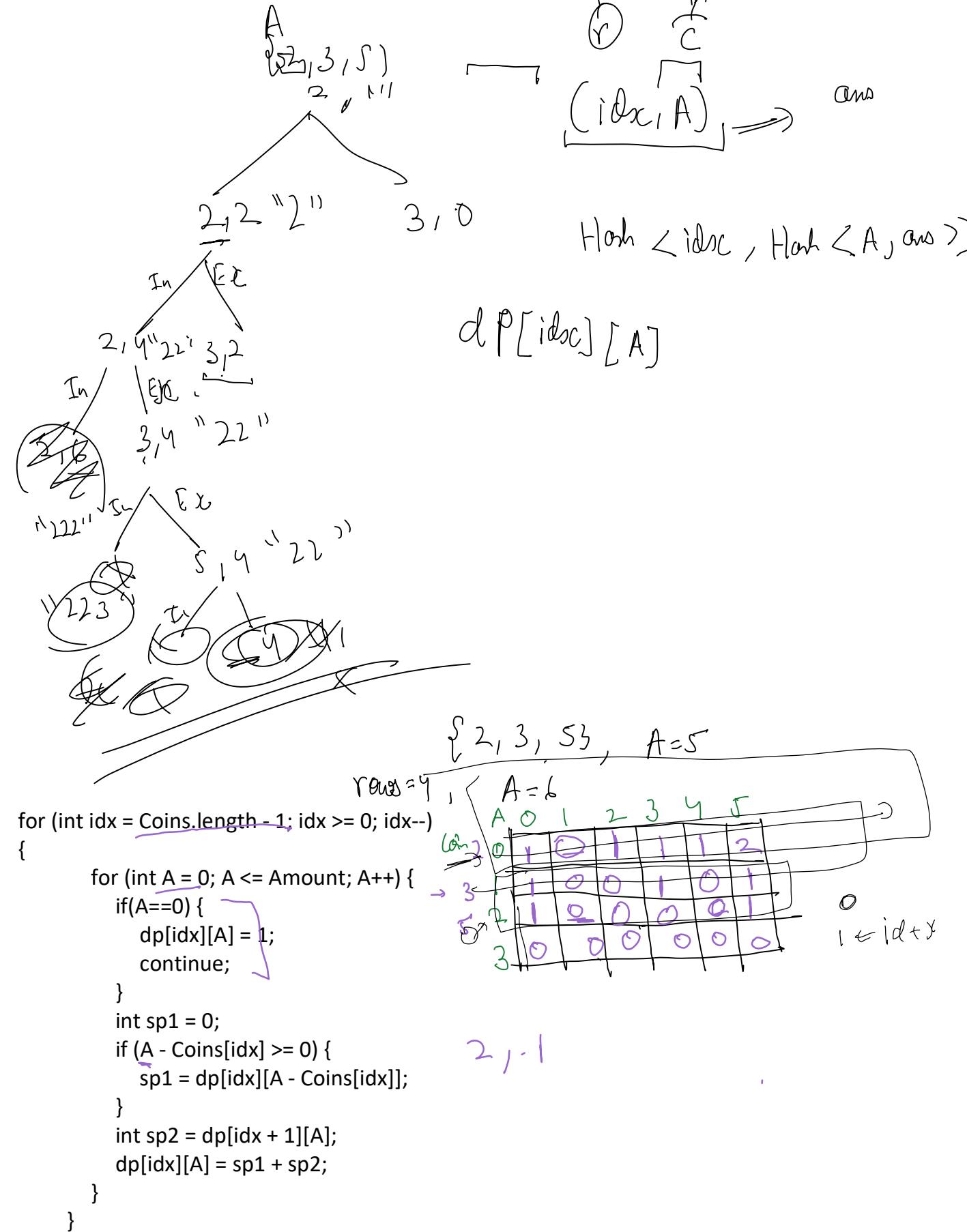
7







$b_0(b)$

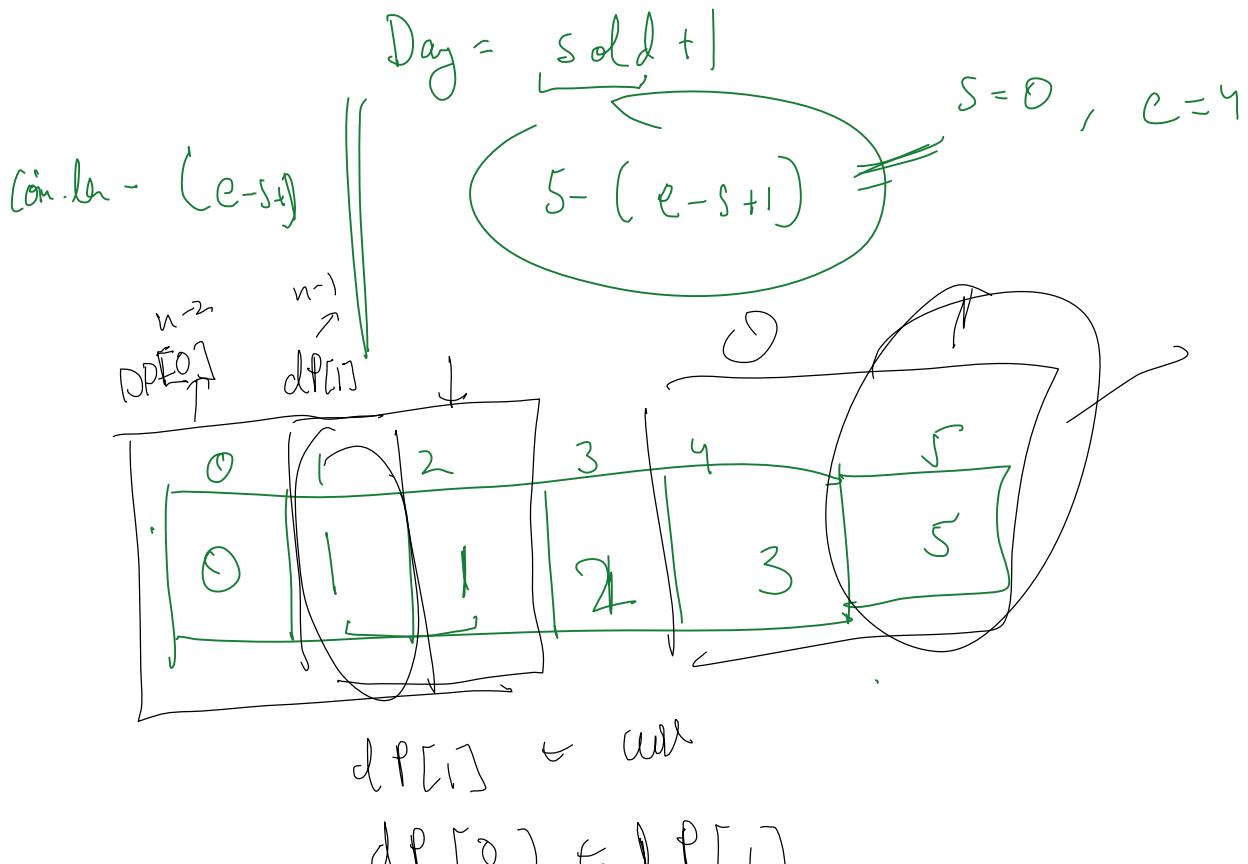


$B \rightarrow \{10, 20, 5, 15\}$
 $D-2d \quad 20, 40, 10, 50$
 $D-1d \quad 30, 60, 15, 45$

$D1 = 2, 3, 5, 1, 4$	Sell bottle 2
$D2 = 4, 6, 10, 2, 8$	Sell bottle 6
$D3 = 6, 9, 15, 3, 12$	Sell bottle 12
$D4 = 8, 12, 20, 4, 16$	Sell bottle 4
$D4 = 8, 15, 25, 4, 16$	

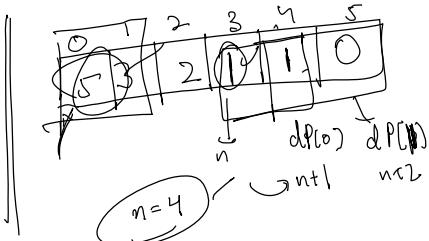
$s=1, e=4$

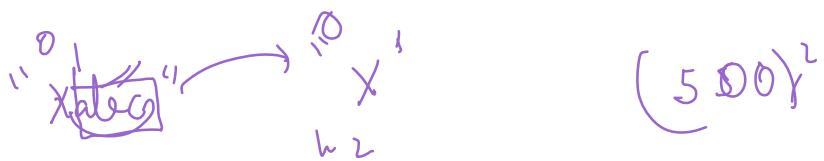
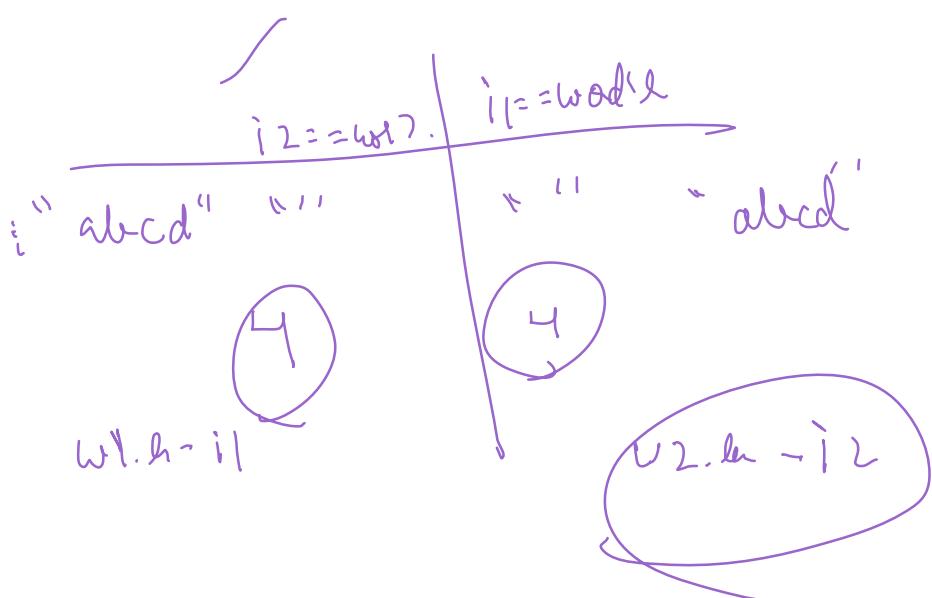
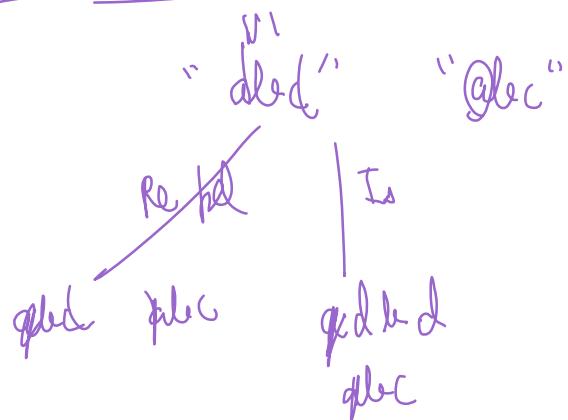
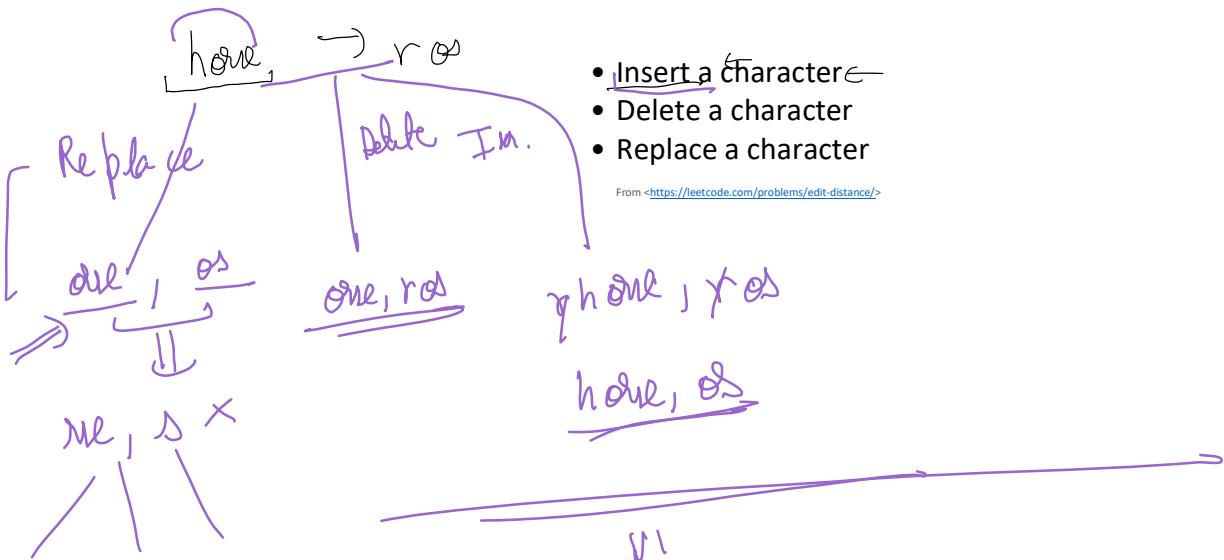
$s=0, e=3$



```

public static int climb_BU(int n) {
    int[] dp = new int[n+2];
    dp[0]=1;
    for (int curr = n-1; curr >= 0; curr--) {
        climb(curr)
            int sp1 = dp[curr + 1];
            int sp2 = dp[curr + 2];
            dp[curr] = sp1 + sp2;
    }
    return dp[0];
}
    
```





$w_1 \downarrow$
 $(1,1)$
 4-1
 5×10^2
 25×10^3
 $\cancel{250,000}$

```

for (int i1 = w1.length(); i1 >= 0; i1--) {
    for (int i2 = w2.length(); i2 >= 0; i2--) {
        solve(i1, i2)
        if (i1 == w1.length() || i2 == w2.length()) {
            int ans1 = w1.length() - i1;
            int ans2 = w2.length() - i2;
            dp[i1][i2] = Math.max(ans2, ans1);
            continue;
        }
        if (w1.charAt(i1) == w2.charAt(i2)) {
            int ans = solve(w1, i1 + 1, w2, i2 + 1);
            dp[i1][i2] = ans;
        } else {
            int rep = 1 + dp[i1 + 1][i2 + 1];
            int del = 1 + dp[i1 + 1][i2];
            int ins = 1 + dp[i1][i2 + 1];
            dp[i1][i2] = Math.min(rep, Math.min(del, ins));
        }
    }
}

```

BU SE
 SE
 HR-1
 HR-2
 Rec
 Sat
 4-1

smallt
 Rec
 del
 ins
 rep

(idk A)
 (idx A)

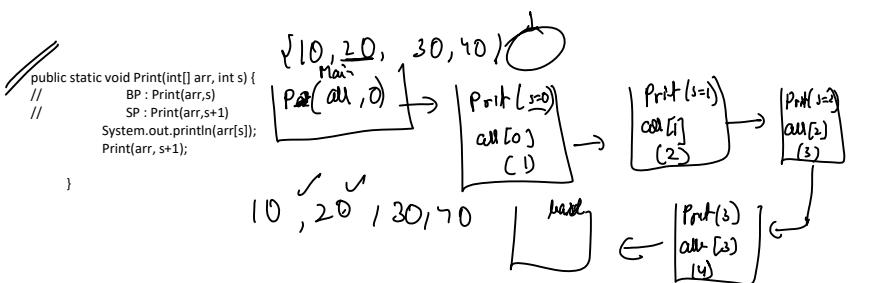
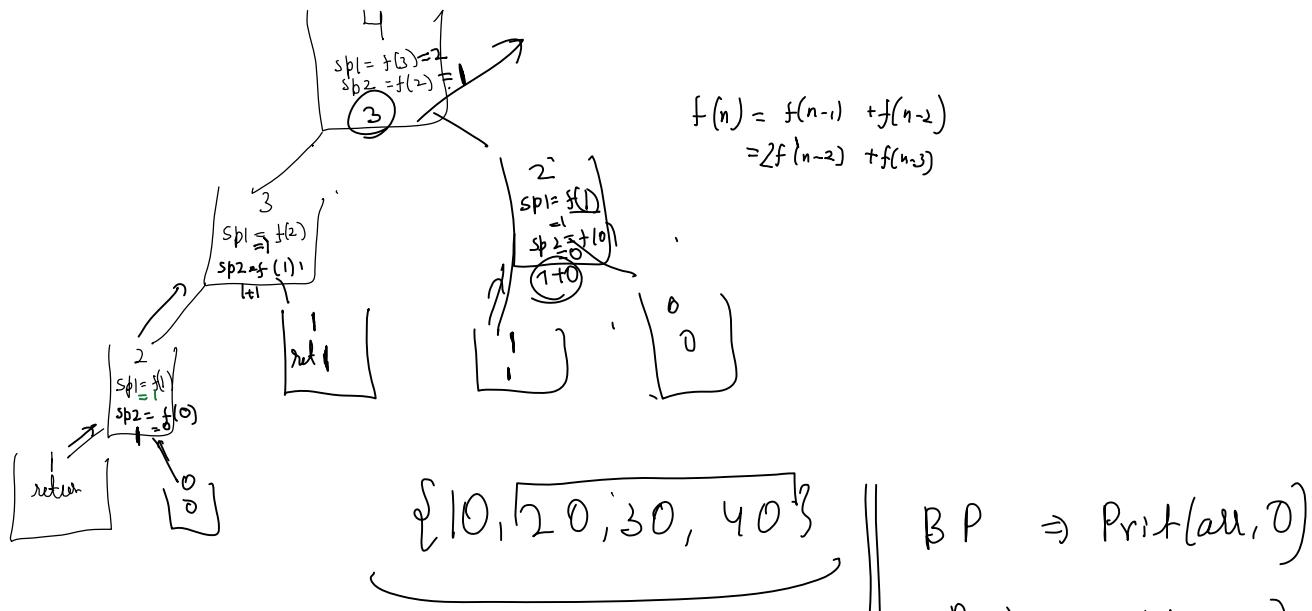
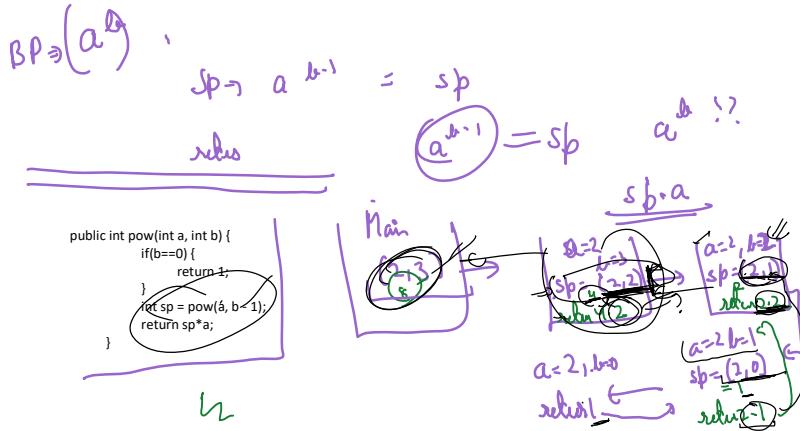
$3/\text{Mount}$
 $20^{\circ}/\text{air}$
 30°
 10°
 21

$\text{DP} + \text{Rec}$

$\text{PID}(y)$
 $\frac{1}{2}$
 $\frac{2}{3}$
 $\frac{3}{4}$
 $\frac{4}{5}$
 $\frac{5}{6}$
 $\frac{6}{7}$
 $\frac{7}{8}$
 $\frac{8}{9}$
 $\frac{9}{10}$

$\$P \rightarrow \text{PID}(1,y)$
 $\$P \rightarrow \text{PID}(2,y)$
 $\text{PID} \leftarrow \frac{\text{Fib}(1)}{a}$

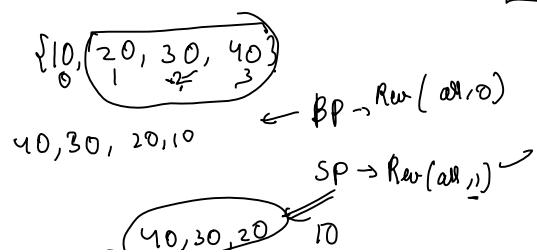
$M(1,2)$
 $\text{public static void PID(int s, int e) {$
 $\quad \text{if}(s > e) {$
 $\quad \quad \text{return};$
 $\quad }$
 $\quad \text{System.out.println(s);}$
 $\quad \text{PID}(s+1, e);$
 $\quad \text{System.out.println(s);}$
 $\quad }$
 $\Rightarrow 1$
 $\Rightarrow 2$
 $\Rightarrow 1$
 $\Rightarrow 2$
 $\Rightarrow 3$
 $\Rightarrow 4$
 $\Rightarrow 5$
 $\Rightarrow 6$
 $\Rightarrow 7$
 $\Rightarrow 8$
 $\Rightarrow 9$
 $\Rightarrow 10$

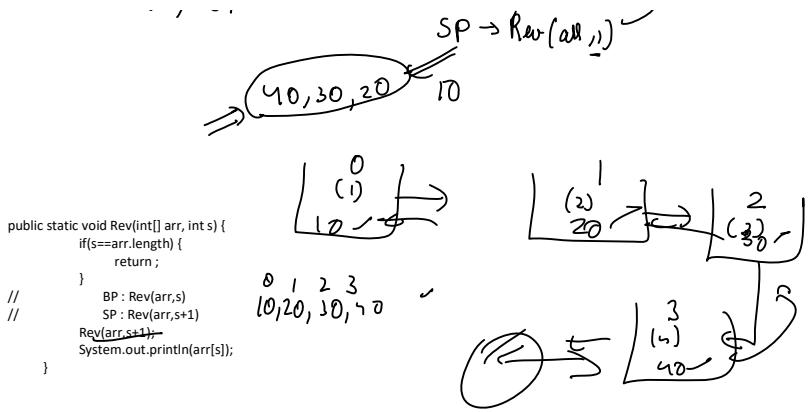


P (all, 0) {

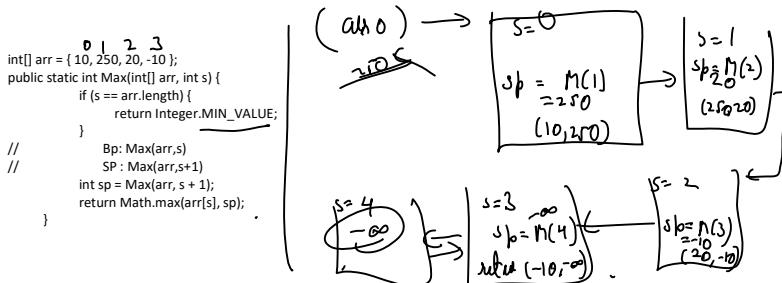
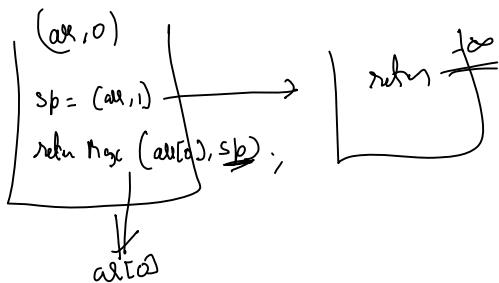
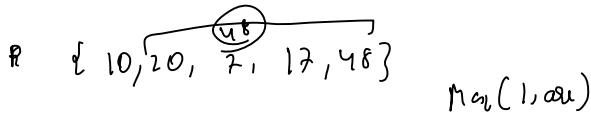
\downarrow print(arr[0])

 P(all, 1)

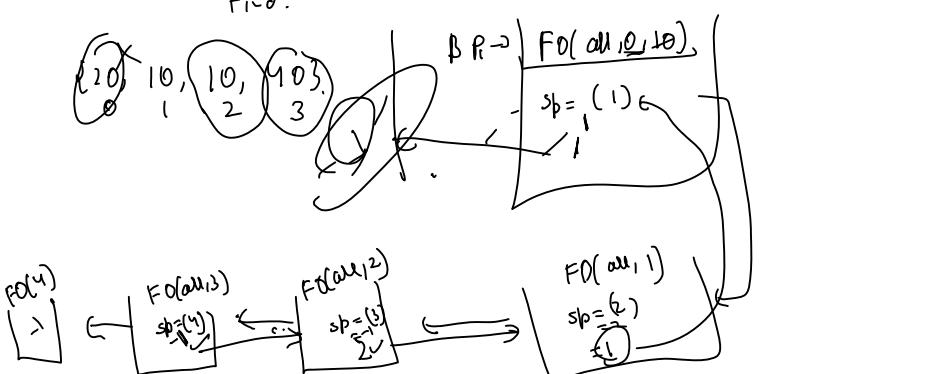




Mac



Find.



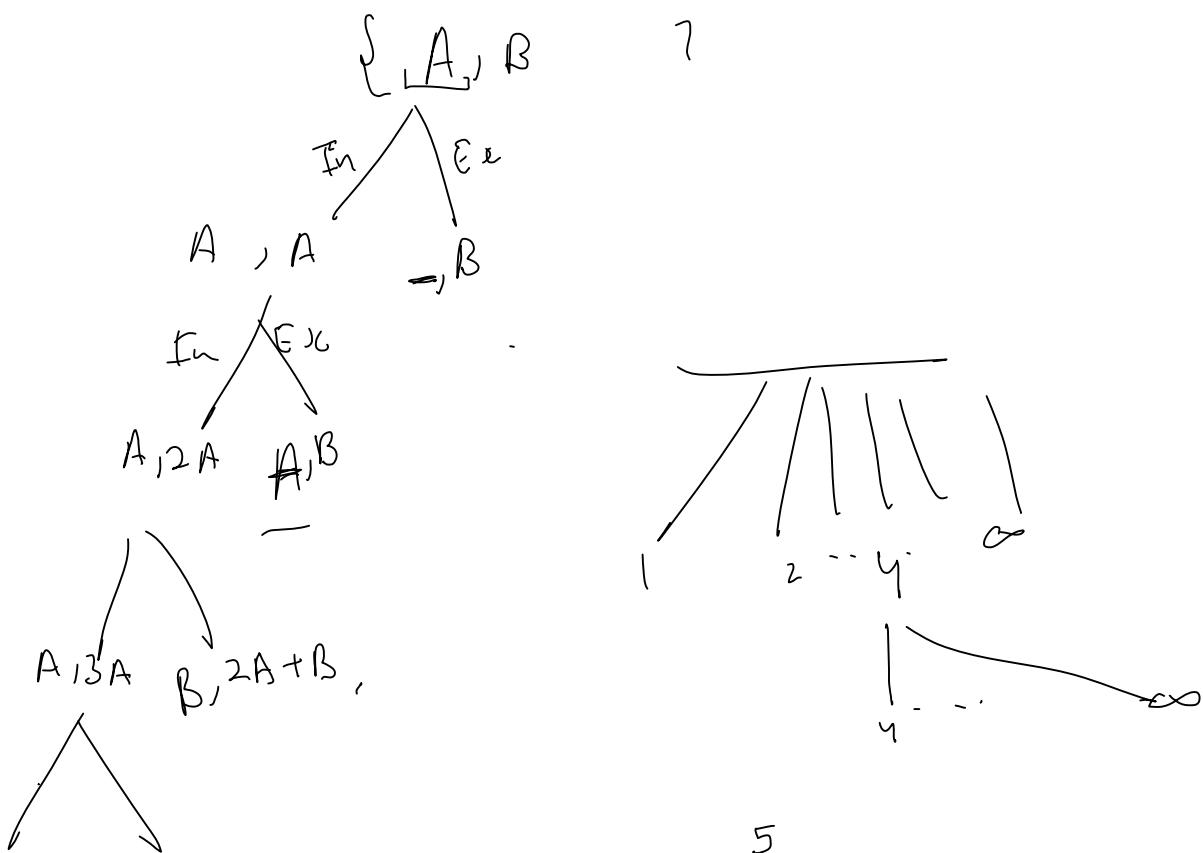
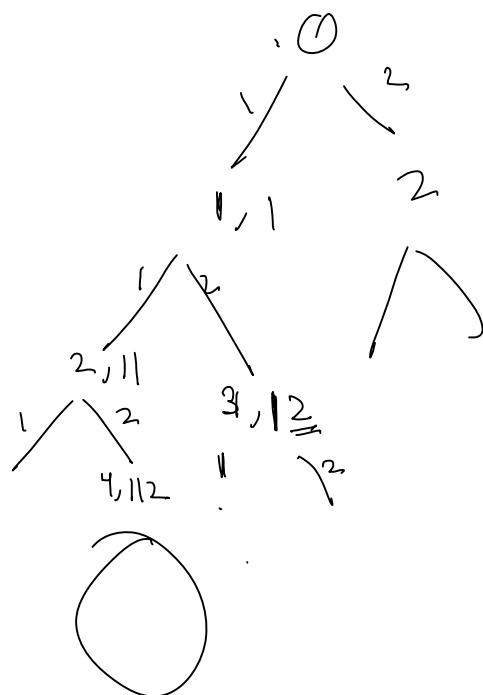
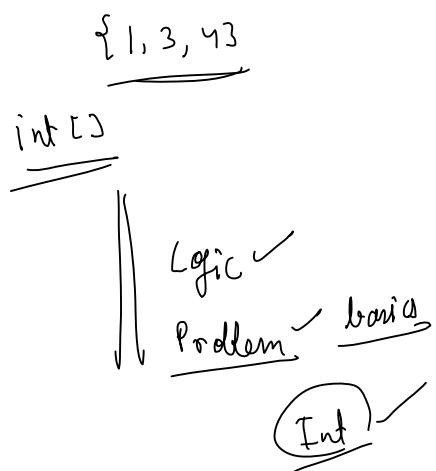
$\{ 15, 10, 20, 10, 30 \}$

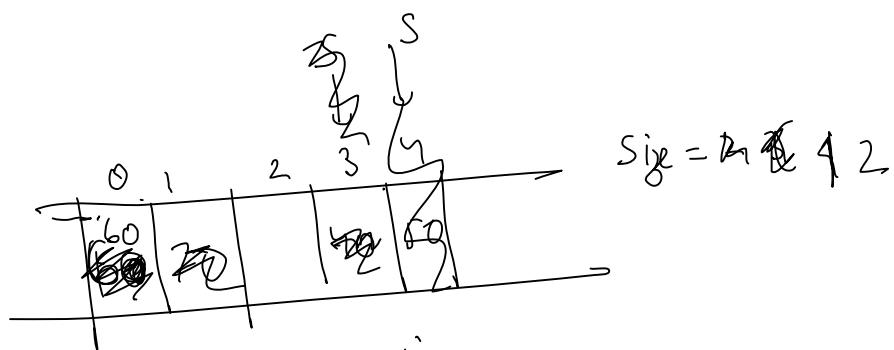
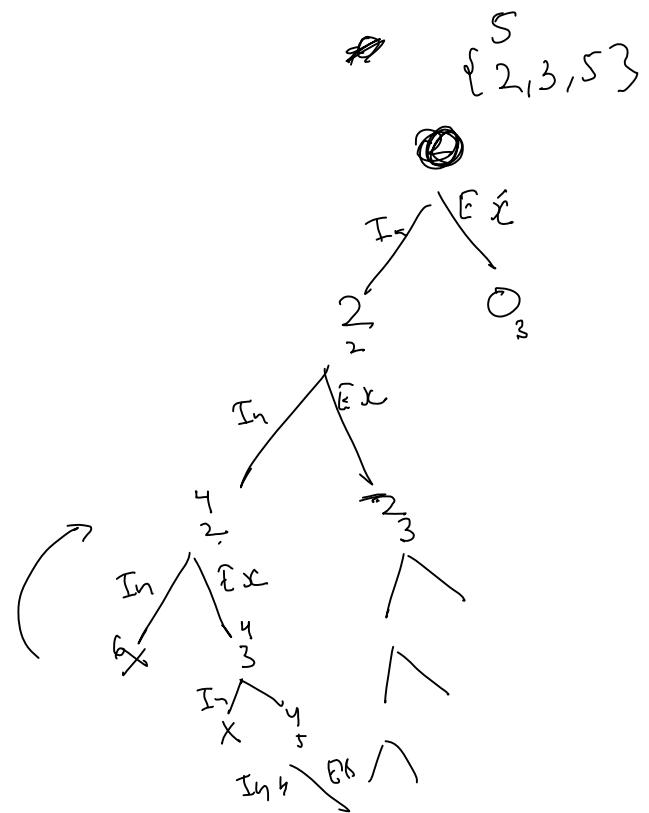
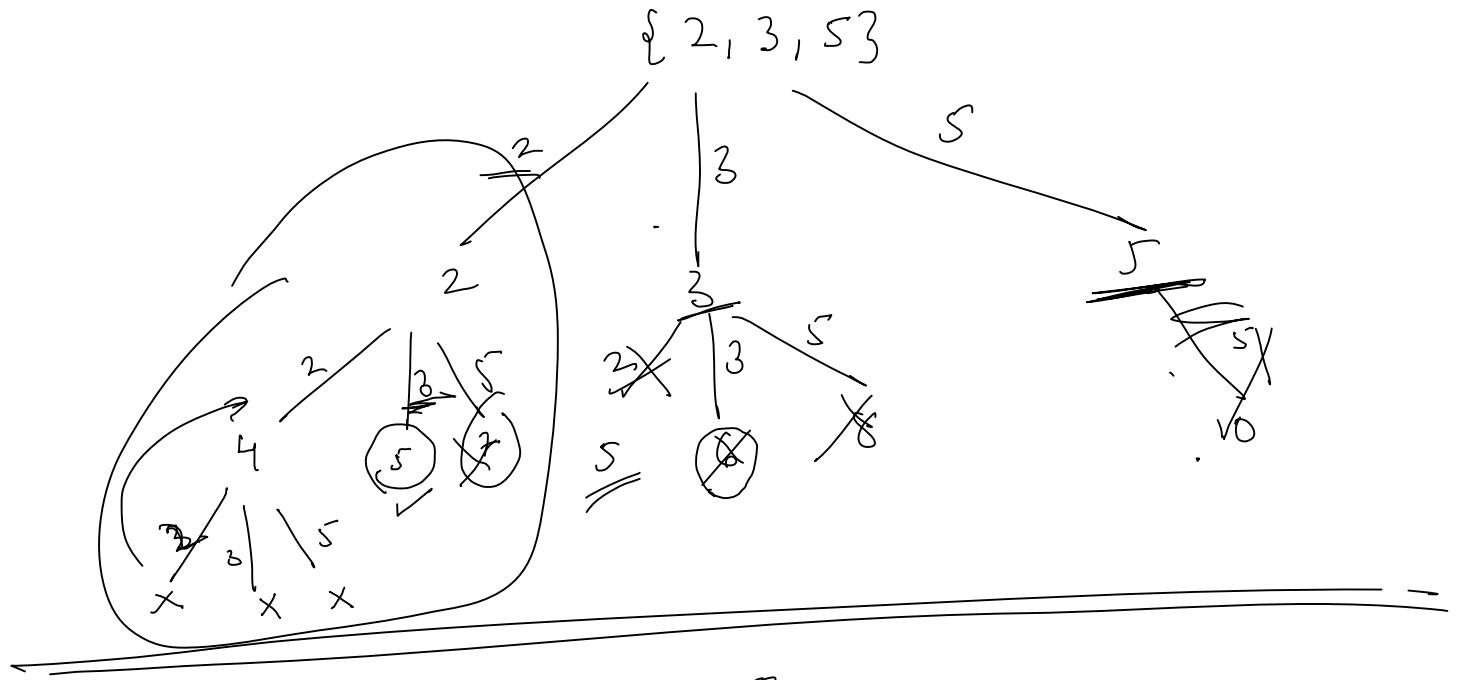
Q { 15, 10, 11, 10, 10, 150 } ||

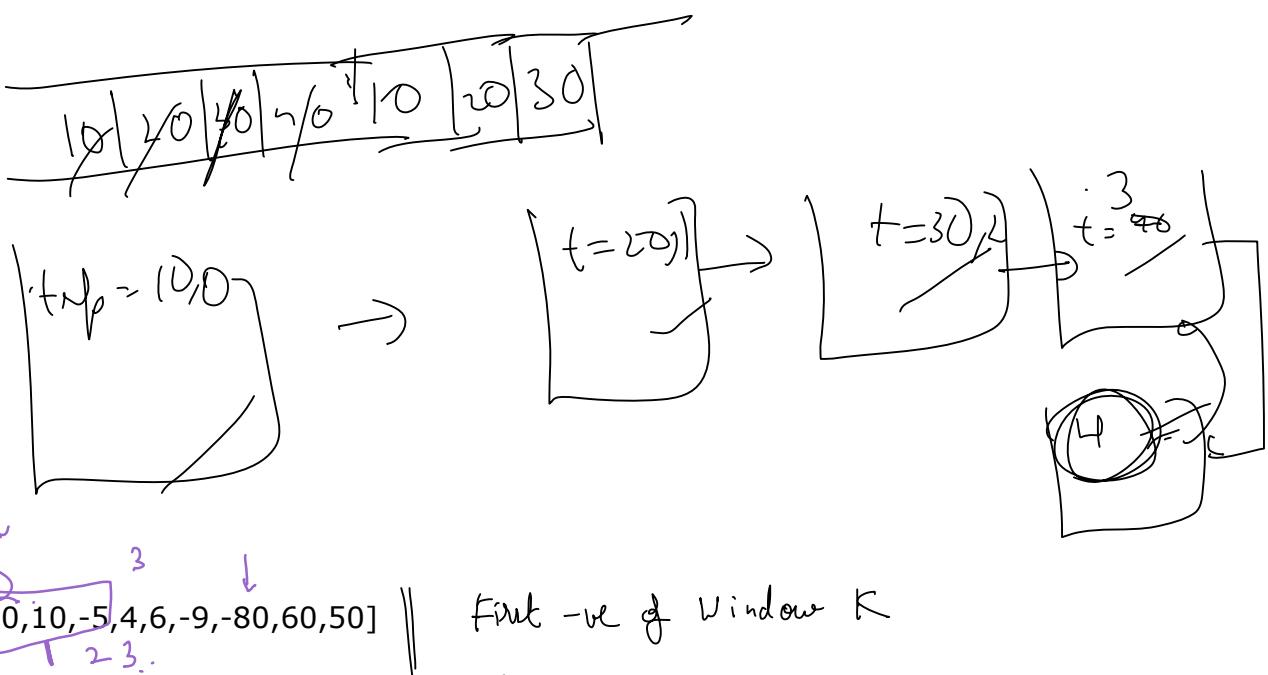
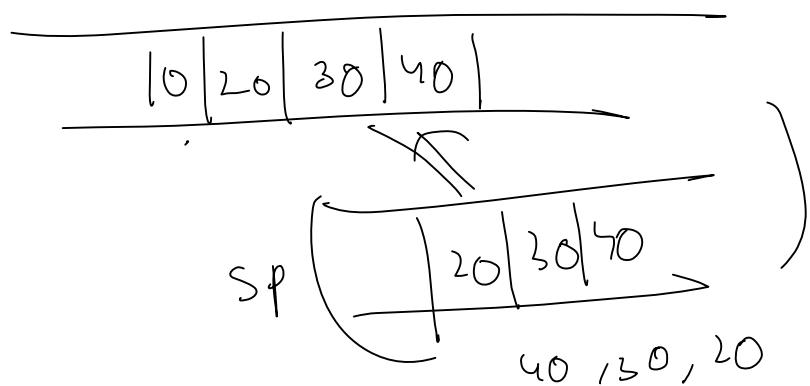
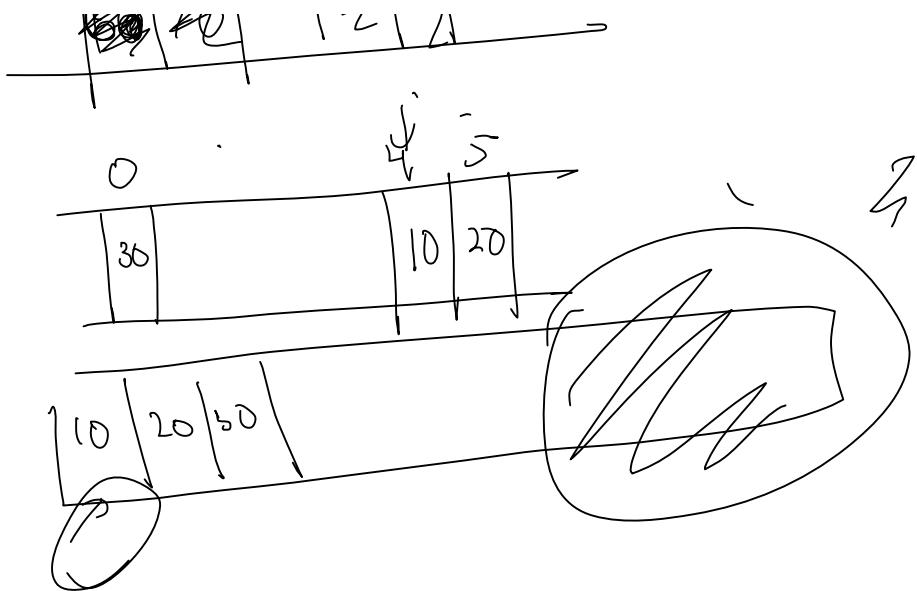
$BP \rightarrow LO(arr, 0, d)$

$SP \rightarrow LO(arr, 1, d)$

Q q 18, 10, 11, 10, 10, 15 of || v ③

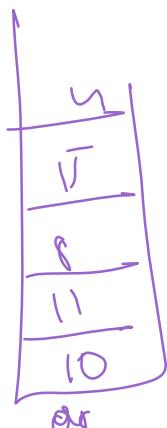






Q $[-20, 10, -5, 4, 6, -9, -80, 60, 50]$ || First -ve of Window K
 $K=3$

Q Min Stack



getMin()

$$\min = 4$$

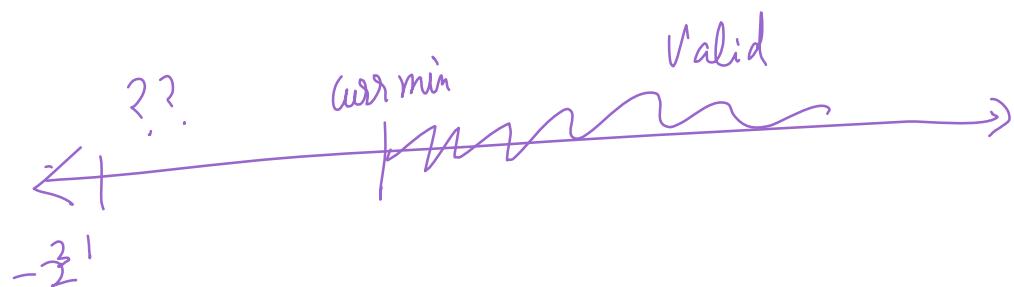
M1) $n \times K$

M) $n + K$

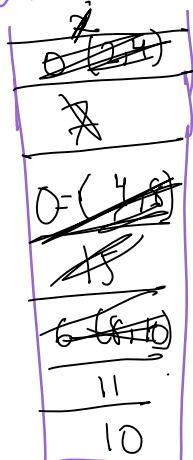
M1) $O(n)$ $O(n)$
Rec

M2) $O(1)$ $O(n)$
T S

{10, 11, 8, 15, 14, 13}



{10, 11, 8, 15, 14, 13}

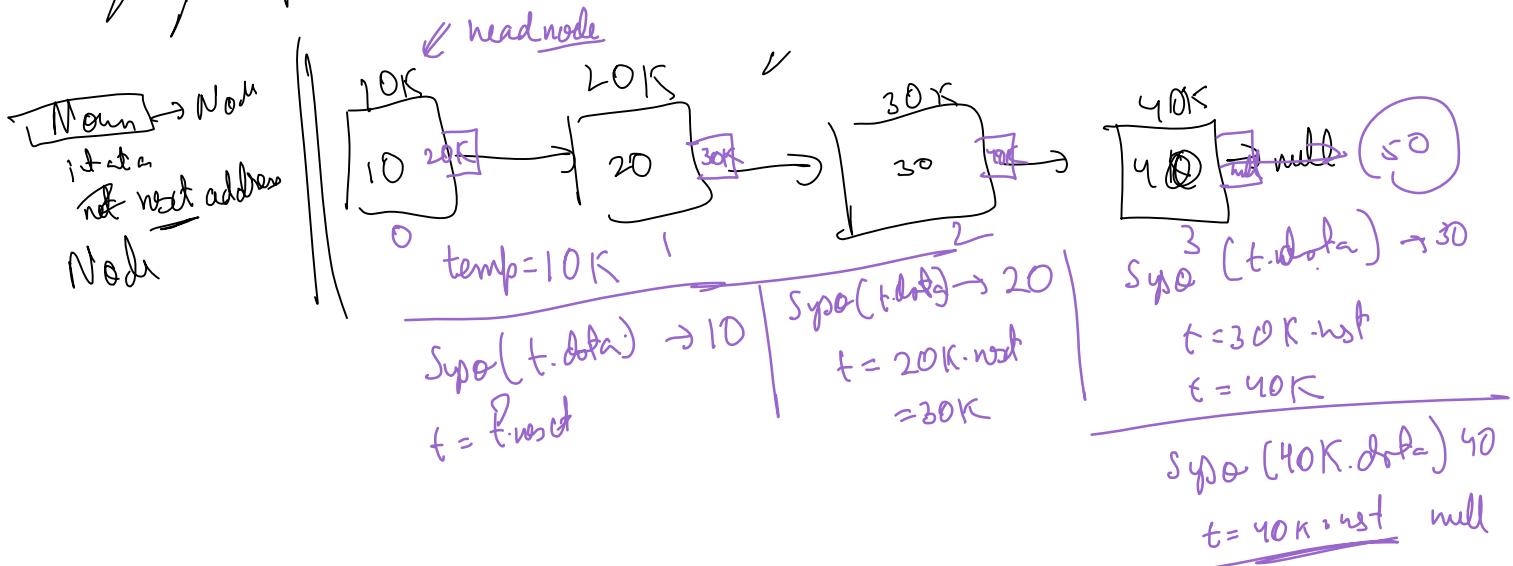
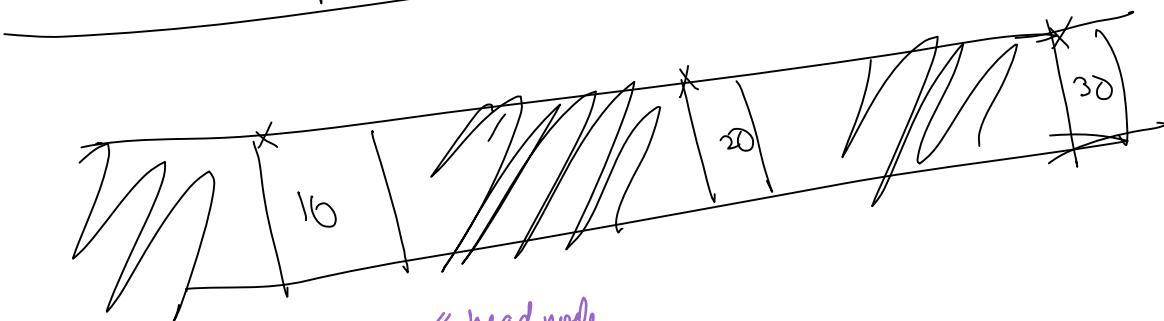
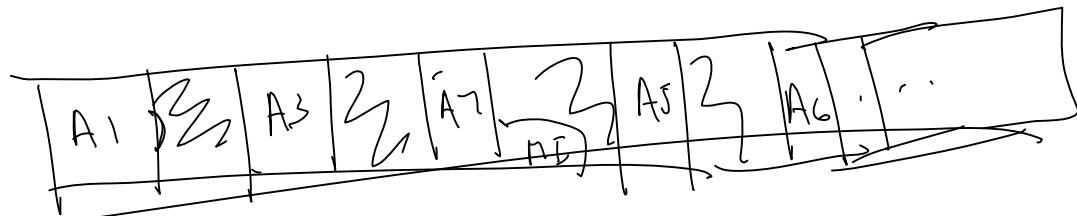
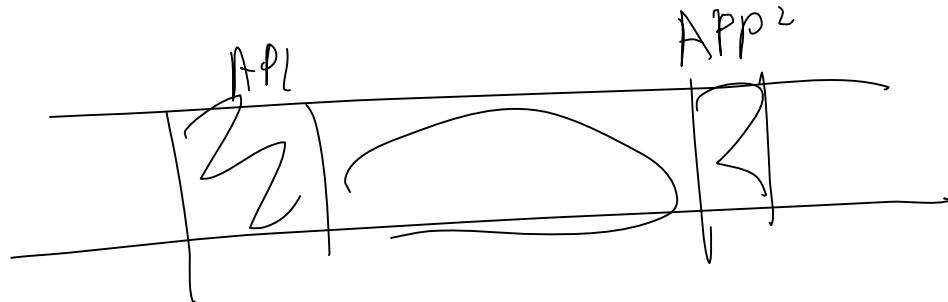
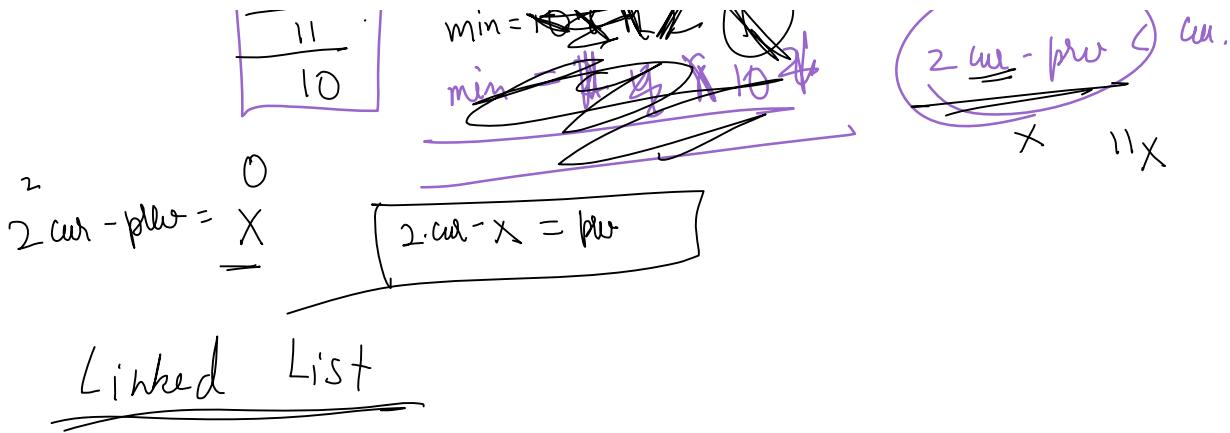


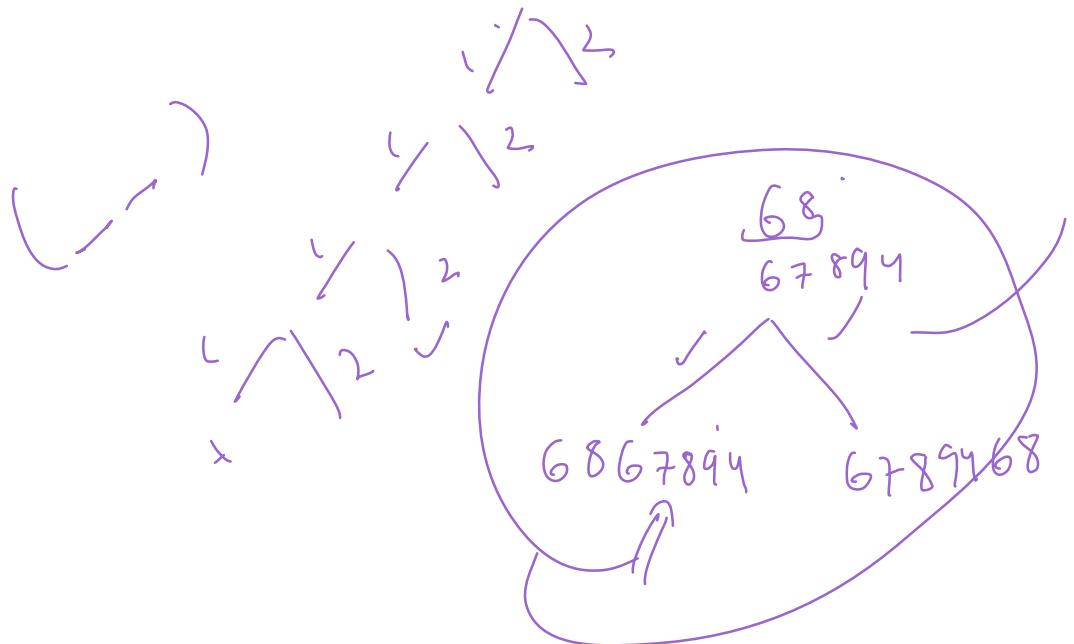
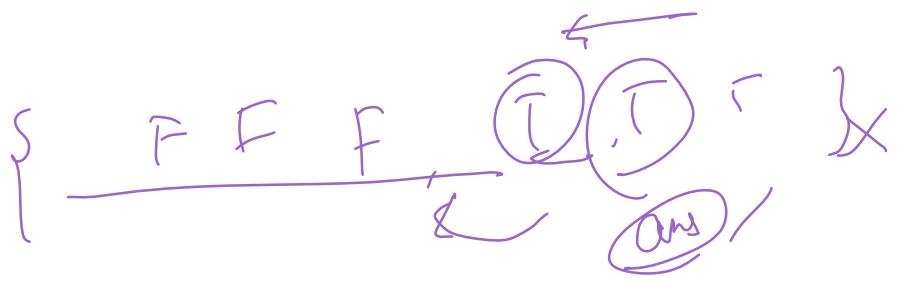
$f(x)(curr, prev min) \leftarrow curr.$

curr < prev

~~min = 8~~
~~min = 10~~

curr - prev < 0
2 curr - prev < curr.

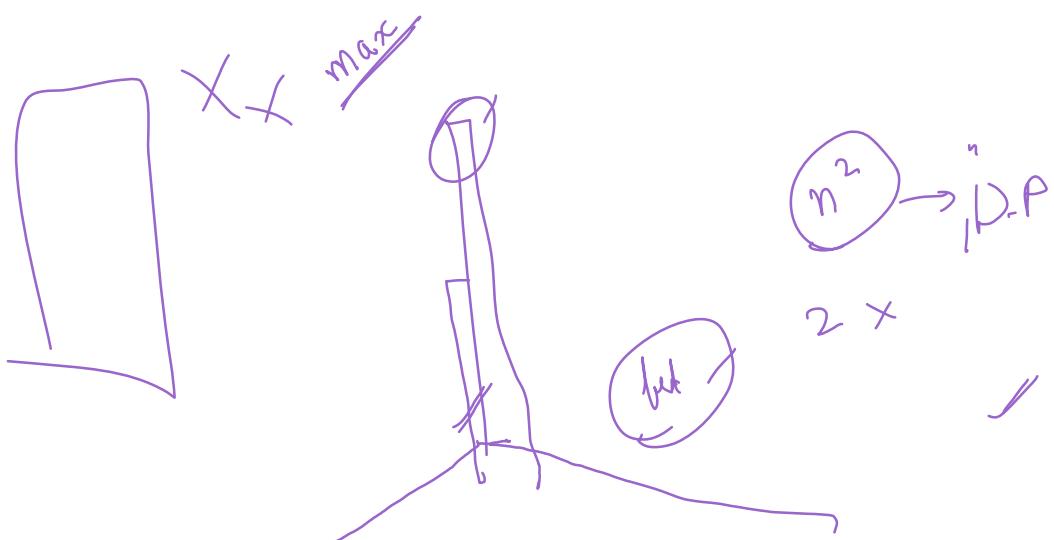


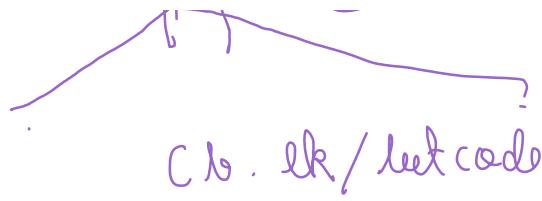


Q $S_1 = "abcabc"$
 $S_2 = "aabab"$

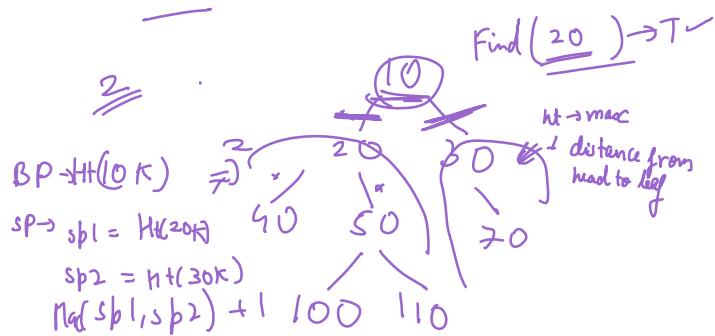
|| subset

b aab

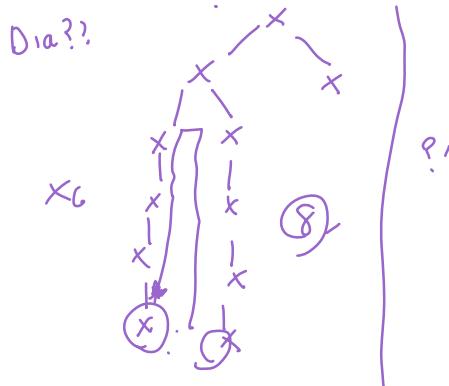
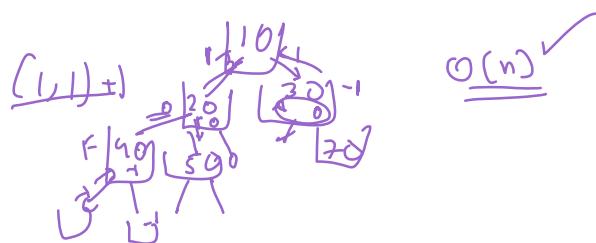
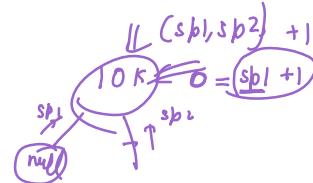


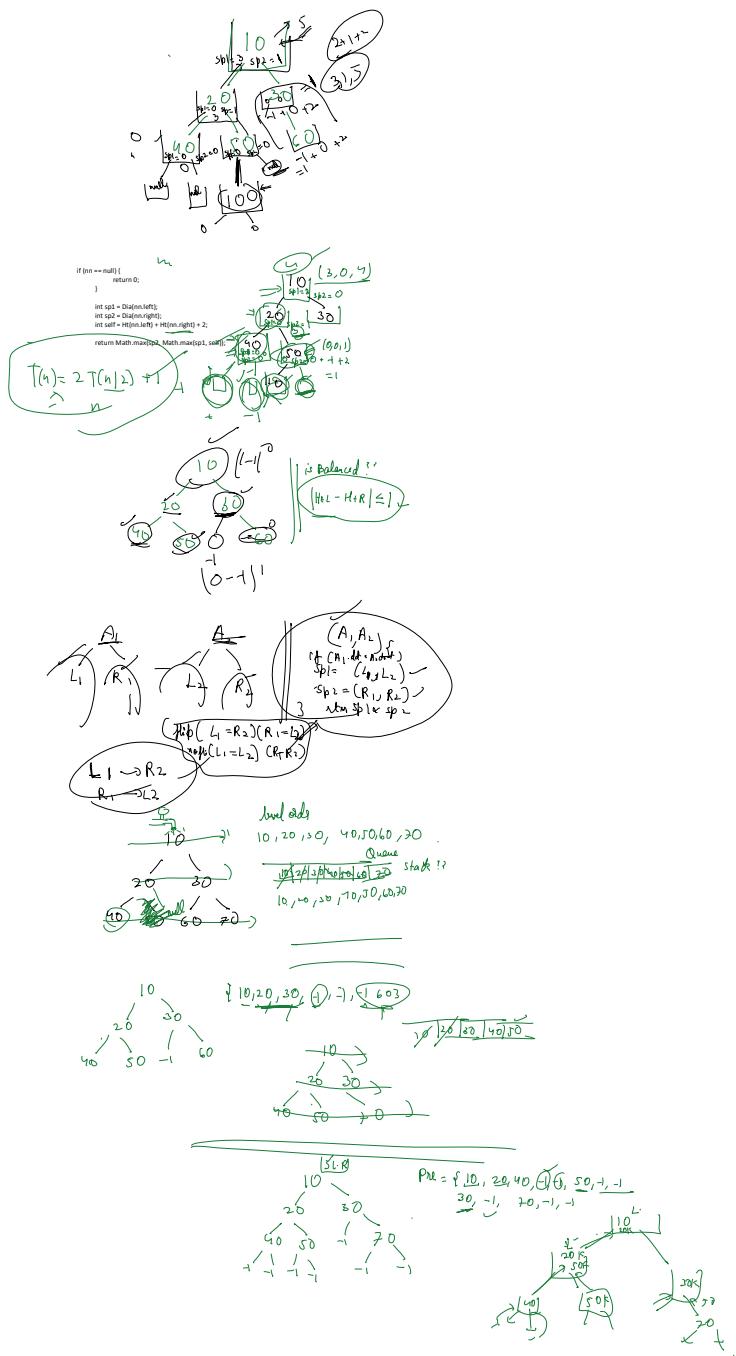
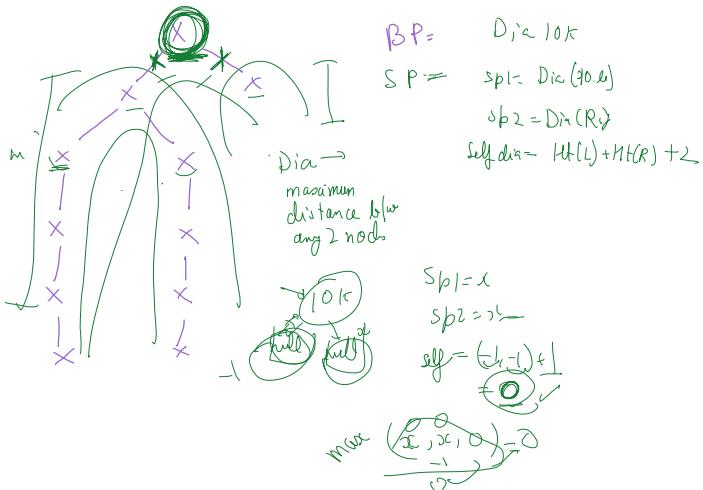


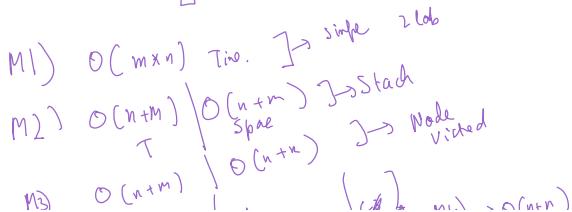
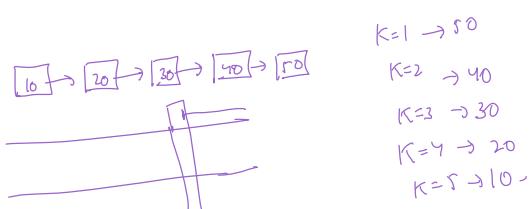
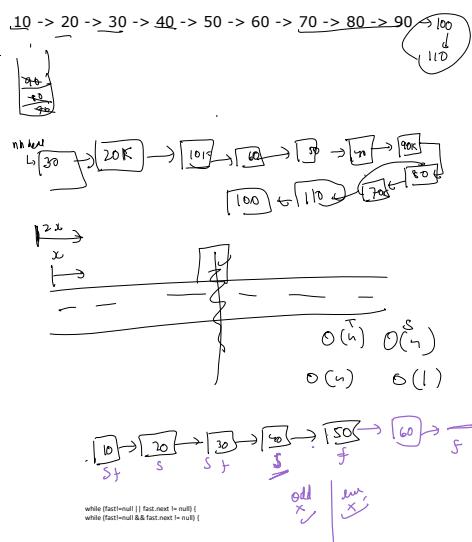
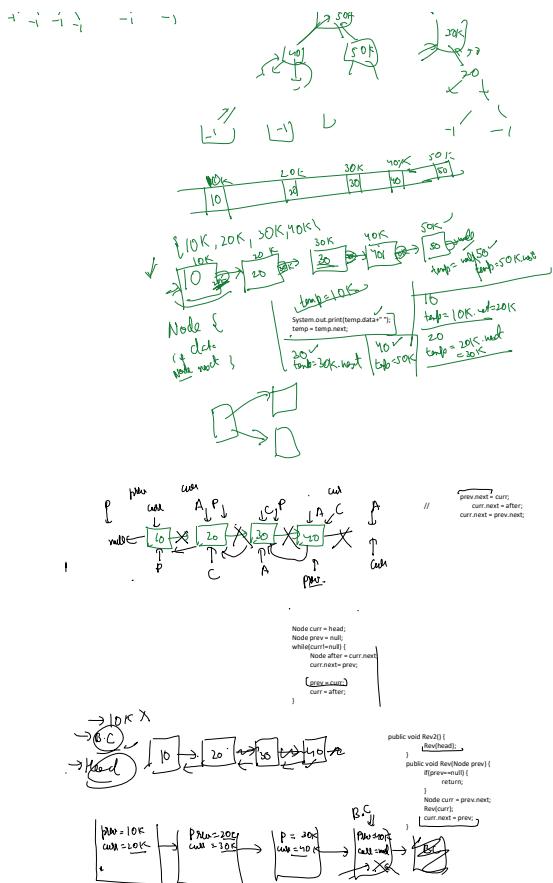
8
12 15 10 11 5 6 2 3

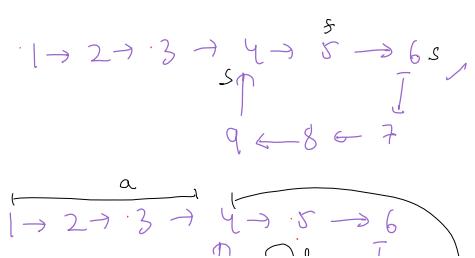
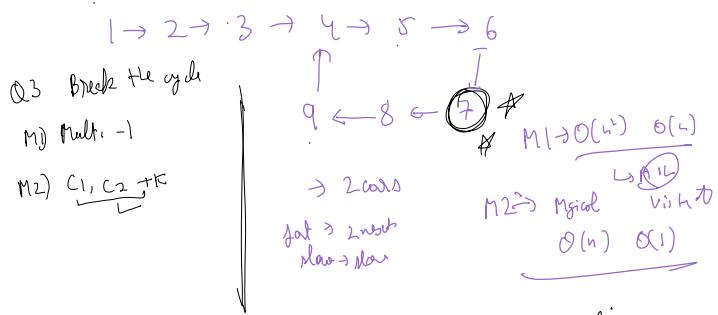
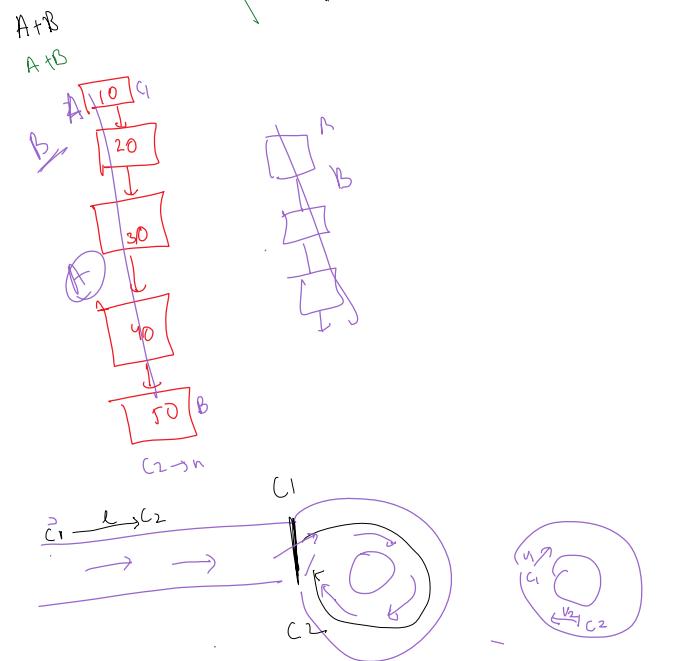
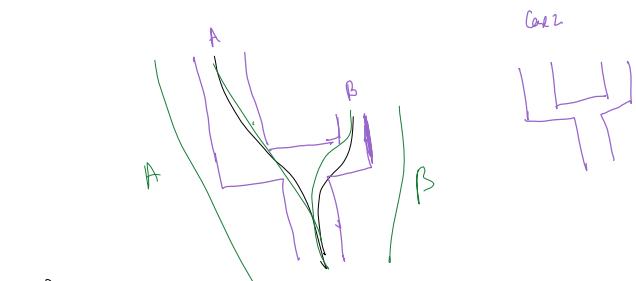
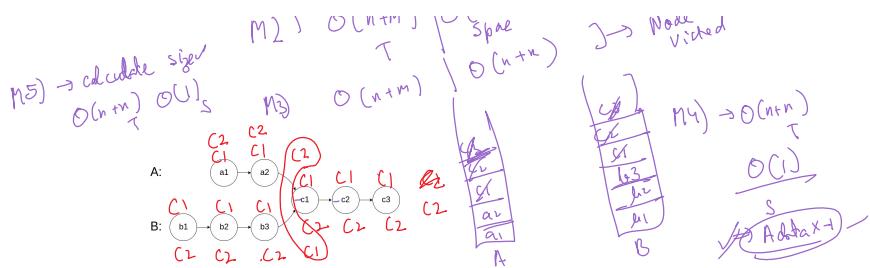


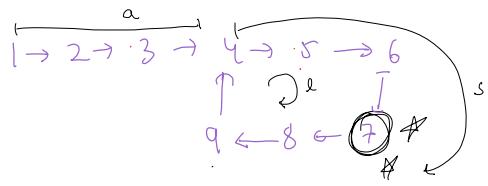
```
private int Ht(Node nn) {
    int sp1 = Ht(nn.left);
    int sp2 = Ht(nn.right);
    return Math.max(sp1, sp2)+1;
}
```



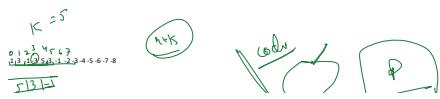
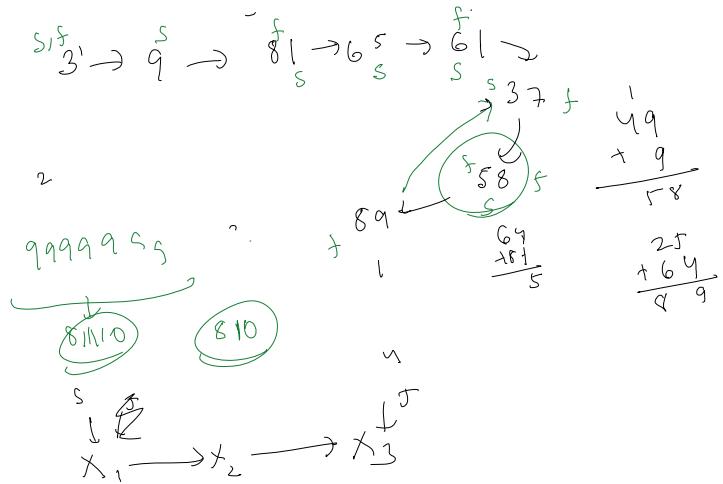
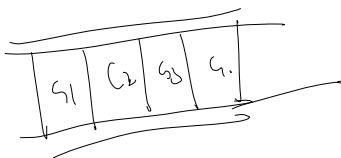
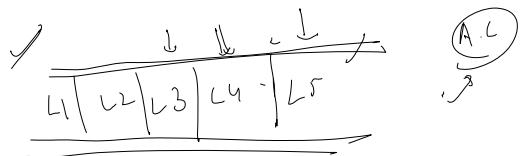
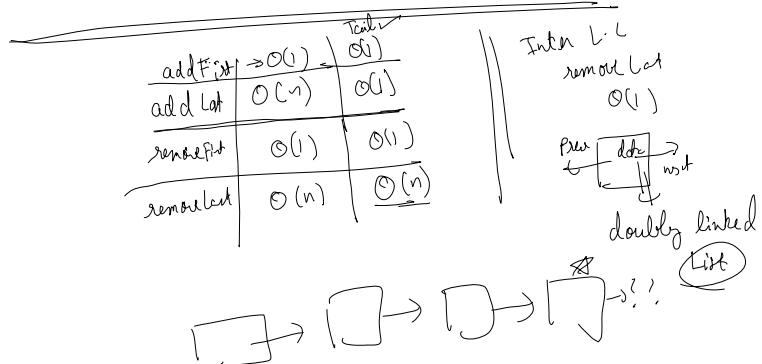


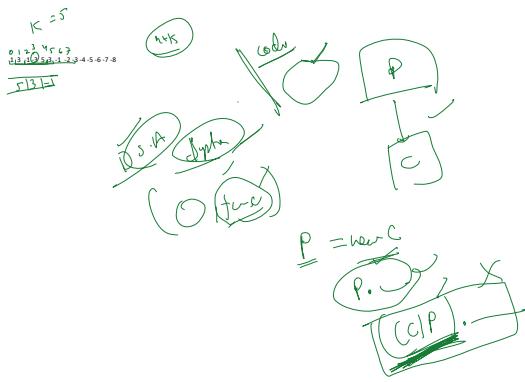




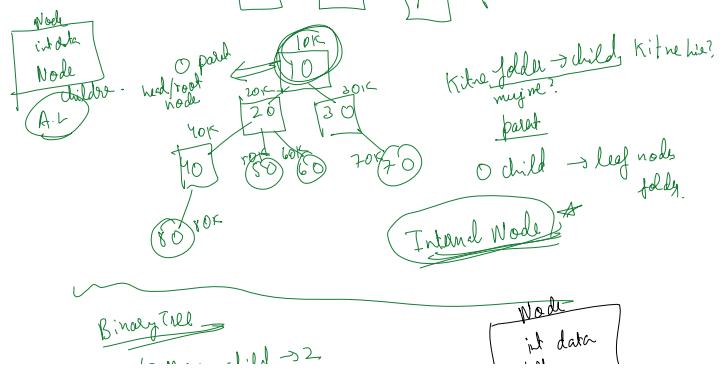
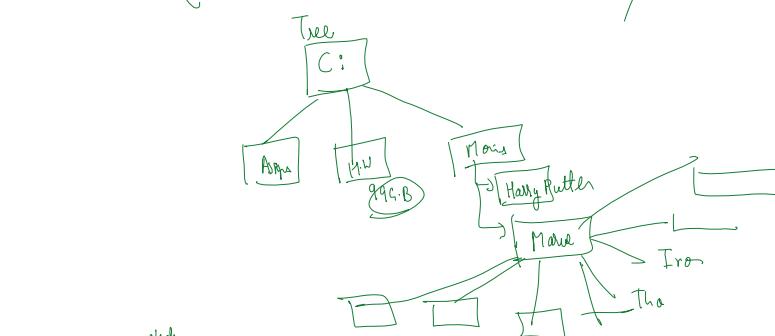
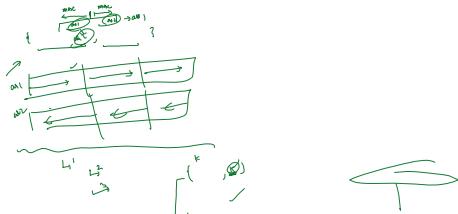


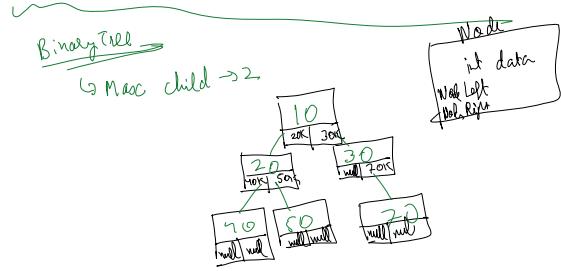
$$\begin{aligned}
 & 2 \times (v \cdot t = a + s + c_1 l) \Rightarrow 2v \cdot t = 2a + 2s + 2c_1 l \\
 & 2v \cdot t = a + s + c_2 l \quad 2v \cdot t = a + s + c_2 l \\
 & \therefore a + s + (c_1 - c_2)l \\
 & \boxed{a + s = (c_2 - 2c_1) \cdot l}
 \end{aligned}$$



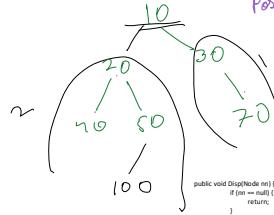


[1,3,1,-3,5,3,6,7]
From <https://www.geeksforgeeks.org/level-order-traversal-binary-tree/>





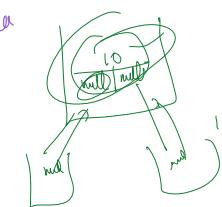
HT



Pre → S. L. R ⇒ 10, 20, 40, 50, 10, 30, 70
 In → L. S. R ⇒ 40, 20, 50, 10, 30, 70
 Post → L. R. S ⇒ 40, 50, 20, 70, 30, 10

S. R. L ⇒ 10, 30, 70, 20, 50, 40
 R. S. L ⇒ 70, 30, 10, 50, 20, 40
 R. L. S ⇒ 70, 30, 50, 40, 20, 10

S, L. R. Pre order



```

public void DisplNode(n) {
    // se n == null
    Disp(n.left);
    Disp(n.right);
    System.out.println(n.data);
}
  
```

Pre → S. L. R ⇒ 10, 20, 40, 50, 10, 30, 70

In → L. S. R ⇒ 40, 20, 50, 10, 30, 70

BP (S.L.R)
 Pre ⇒ ①(20, 40, 50) ②(30, 70)
 In ⇒ ③(40, 20, 50) ④(10, 30, 70)

SP) L.R. max
 Pre = {40, 20, 50}
 In = {40, 20, 50}

is = 4, j = 1

found = 1, ie

is = 4, ie = 5

pre = 5

is = 4, ie = 5

pre = 5

is = 4, ie = 5

pre = 5

is = 4, ie = 5

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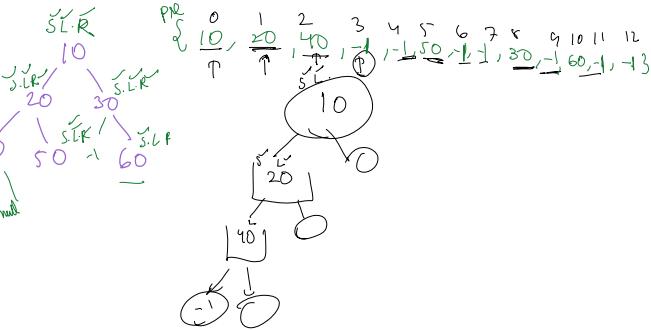
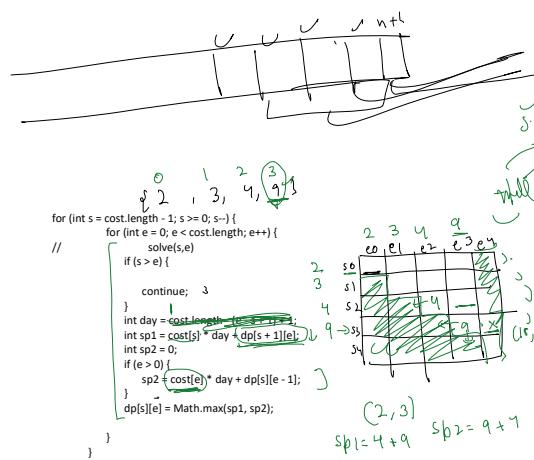
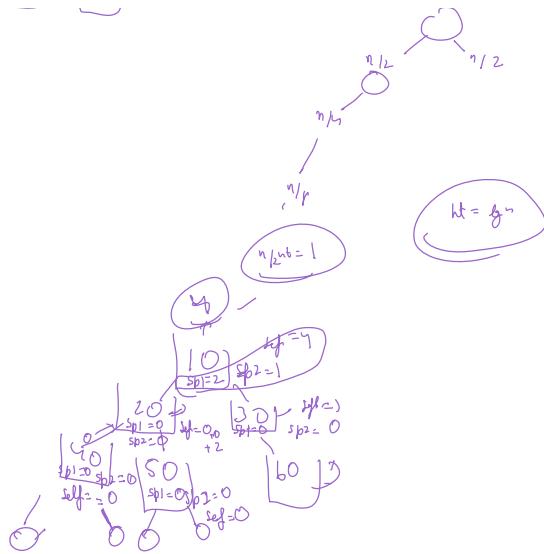
is = 4, ie = 5

$\hat{b}_i^{\dagger} \hat{b}_i$

ψ_i^k

\hat{b}_i^{\dagger}
 \hat{b}_i

$T(n) =$



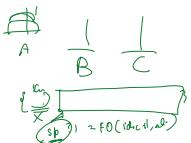
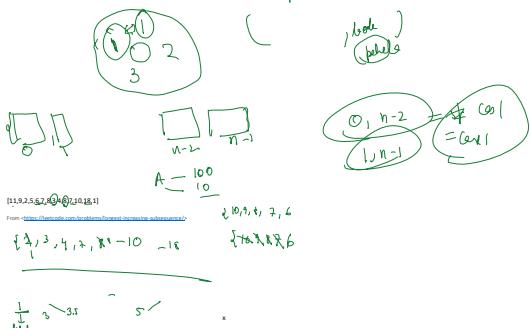
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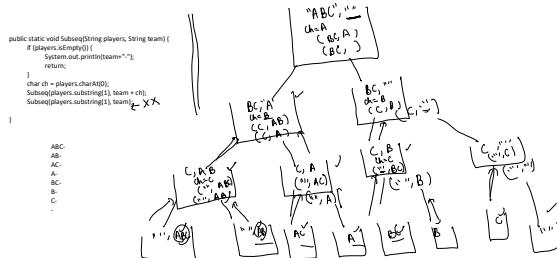
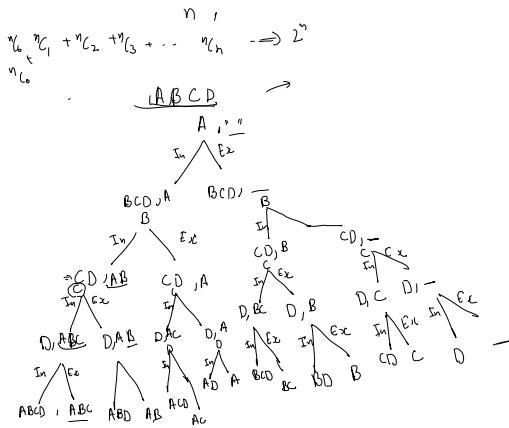
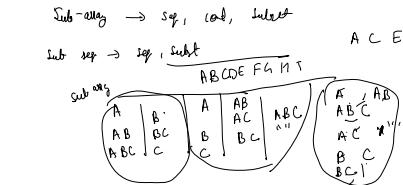
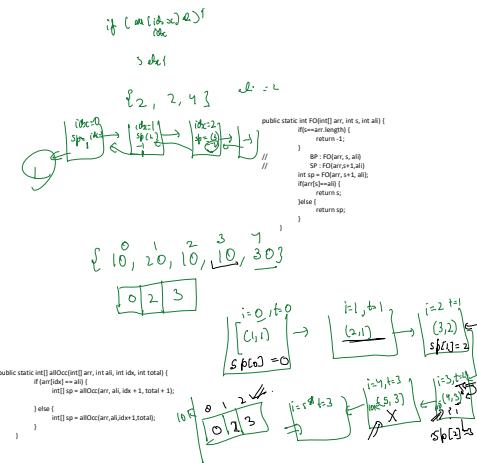
{ $3, 1, 1, 2$ } 5 $1+1=2$

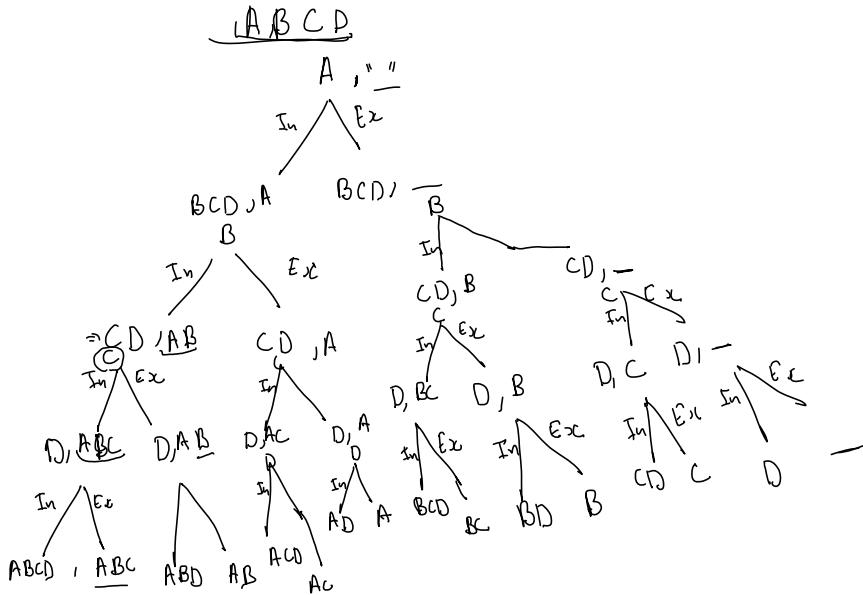
$$sp_2 \rightarrow 0 \quad sp_1 \rightarrow 1 \quad sp_1 = (1)$$

num = [1,2,3,1]

From <https://leetcode.com/problems/coin-change/>

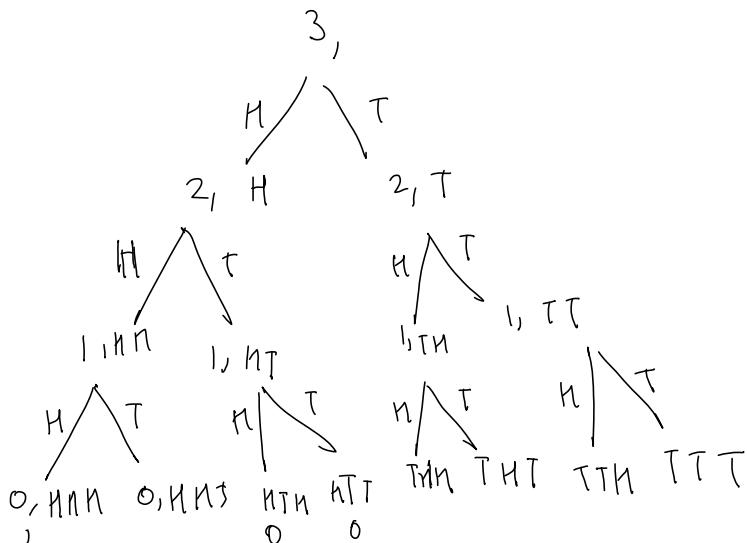
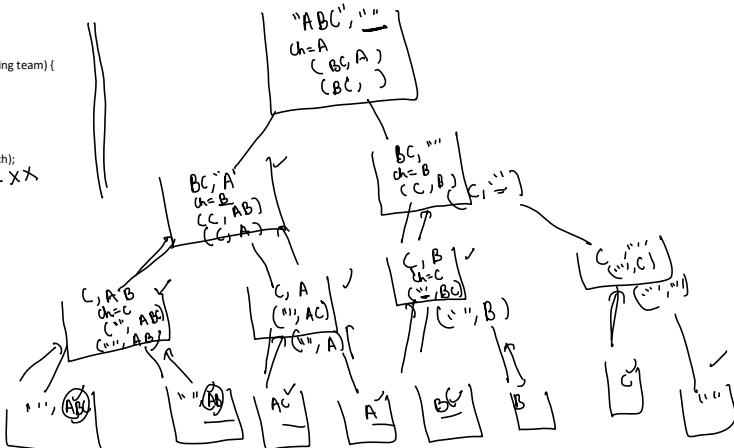






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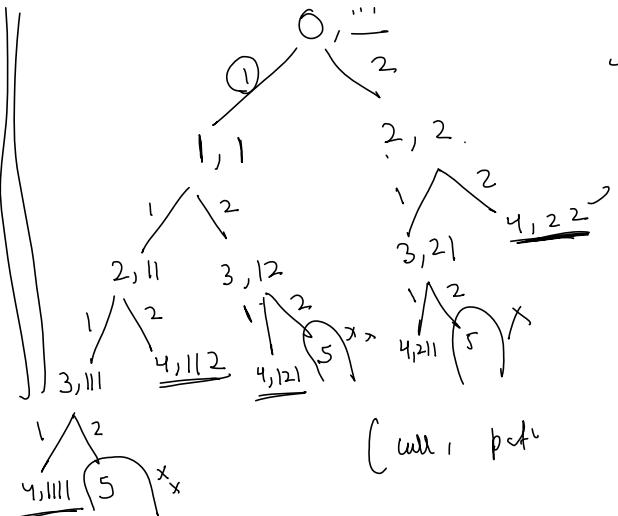
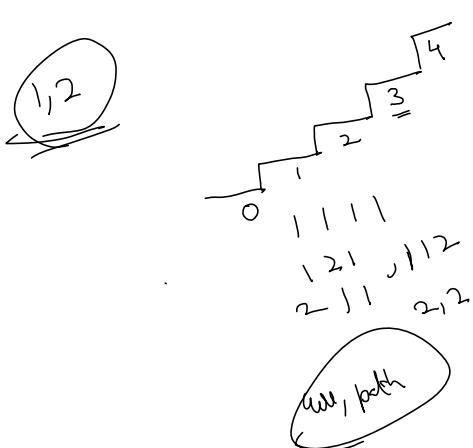
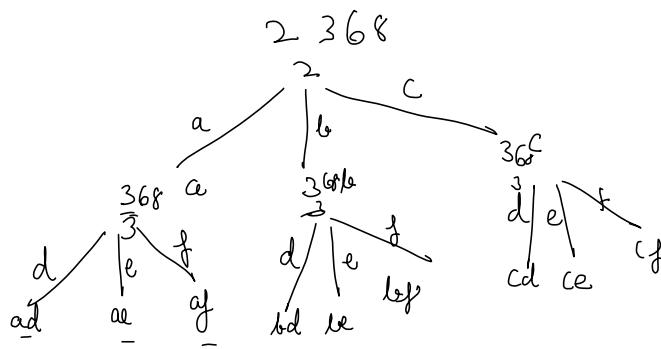
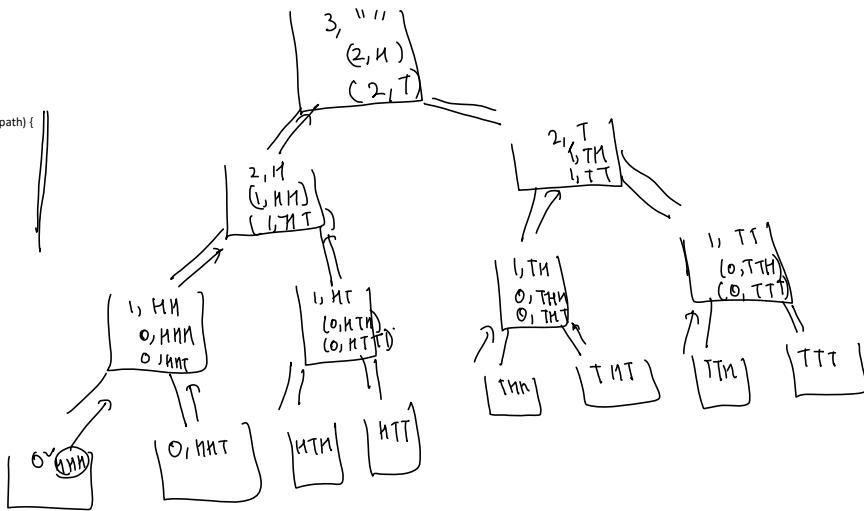
public static void Subseq(String players, String team) {
    if (players.isEmpty()) {
        System.out.println(team + ".");
        return;
    }
    char ch = players.charAt(0);
    Subseq(players.substring(1), team + ch);
    Subseq(players.substring(1), team);
}
    
```



```

public static void CoinToss(int n, String path) {
    if(n==0) {
        System.out.println(path);
        return;
    }
    CoinToss(n-1, path+'H');
    CoinToss(n-1, path+'T');
}

```



```

public static void Climb(int curr, String path, int dest) {
    if(curr==dest) {
        System.out.println(path);
        return;
    }
    if(curr>dest) {
        return;
    }
    Climb(curr+1, path+1, dest);
    Climb(curr+2, path+2, dest);
}

```

