

Computer and Network Security: Random Number Generators

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Usage

- Secure communication requires high quality (pseudo) random numbers
 - One time pads; stream ciphers (symmetric key)
 - RSA public key (generate and test for primality)
 - Salts in hashing
 - Session keys; nonces during handshaking
- Note: Security requirements not the same across all use cases

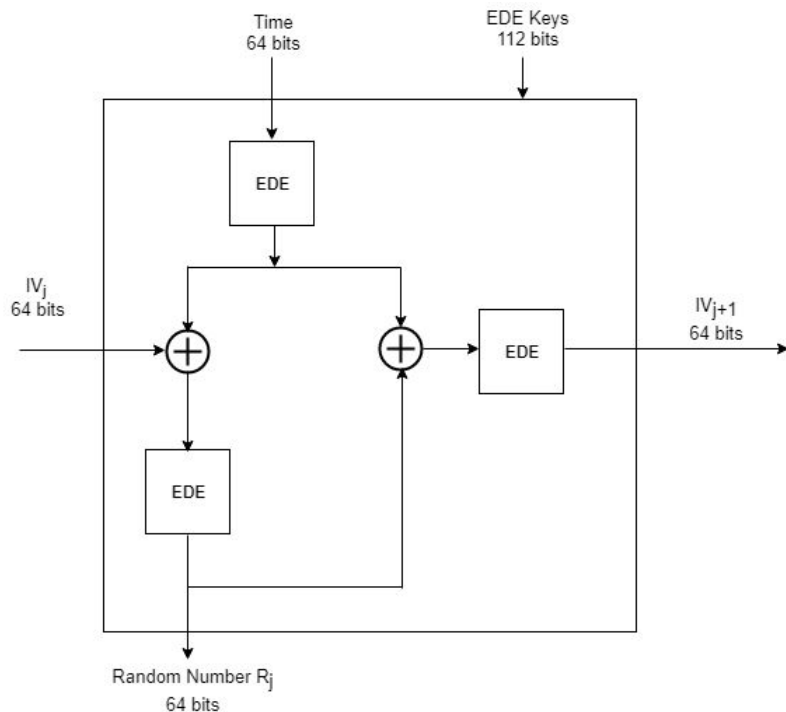
Truly Random Requirements

- Uniform Distribution: all numbers in given range occur equally often
- Independence: Cannot predict future numbers, given the past
- True randomness can come from nature (thermal noise, background radiation from space)
 - Leveraging it is slow and expensive
 - Not easy to convert unpredictability to uniform distribution

Pseudo Random Number Generators (PRNG)

- Algorithmically generated random numbers (approximation of truly random)
- Linear Congruential Generators
$$X_{n+1} = (a * X_n + c) \bmod m; X_0 \text{ is the seed}$$
 - Certain choices of a , c and m can produce uniform distribution
- Not secure since attacker needs to be able to predict just m , a and c
- Better Variant: Take current clock time modulo m as a new seed; change seed after some numbers

Cryptographically Secure PRNG (CSPRNG)



Based on Cryptographic Primitives

- Many designs in literature
- Example: Two keys, two inputs and 3 EDE blocks
 - EDE: encrypt/decrypt/encrypt (2-key DES); can be replaced by AES
 - Next random number dictated by IV_{j+1} and one more independent pseudorandom input (time)
 - Cannot be predicted from R_j
 - Between two consecutive random numbers 3 EDE encryptions (nine DES encryptions)
 - IV_{j+1} cannot be predicted from IV_j; 2 EDE encryptions
- Much slower to generate pseudo-random numbers

True Random Number Generators (TRNG)

- Relies on analog phenomena (entropy source) that change with time in unpredictable ways and generates truly random stream
 - Entropy: extent of uncertainty with a random process; high ☐ good
 - E.g. amplify thermal noise in resistors and then digitize; Very power intensive
 - TRNG cannot generate at the rate applications need
- Leverage TRNG in CSPRNG. How?
 - Use TRNG as seeds to CSPRNGs

- Hardware: e.g. Intel digital random number generator (DRNG)
 - Uses invertors and thermal noise
- Software:
 - Timing between key presses
 - Mouse interrupt and position
 - Interrupt timings (disk better than timers)
 - Finishing time of block device requests
 - Few hundred bits of entropy per sec

Randomness in Linux

- Stores random data (see software) in a random pool
- Uses `/dev/random` (blocks) and `/dev/urandom` (non-blocking) to turn randomness into pseudo random numbers
- Any time random number generated ☐ entropy of randomness pool decreased
 - Entropy reaches zero `/dev/random` blocks
 - `/dev/urandom` will use data in pool as seed and generate as many pseudo random numbers as needed

References

- <https://hackaday.com/2017/11/02/what-is-entropy-and-how-do-i-get-more-of-it/>
- <https://people.eecs.berkeley.edu/~daw/rnd/linux-rand>
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