## Competitive Programming Team Notebook

## BUBT\_ModZero

## Team Members:

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```
/************************************/
// Double Hashing
// 1. Modular Exponentian Needed
// 2. Init must be call and set the maximum length of the string
// 3. If sub string hash required then Compute hash have to call
// 4. If prefix and suffix hash required ComputePreAndSufHash have to call
struct ModularExponentiation {
   template <typename T> T Pow(T b, T p) {
       T res = 1;
       while (p > 0) {
          if (p % 2 == 1) res = res * b;
          b = b * b;
          p /= 2;
       }
       return res;
   template <typename T> T Mod(T a, T m) {
       return (((a % m) + m) % m);
   }
   template <typename T> T BigMod(T b, T p, T m) {
       T res = 1;
       if (b > m) b %= m;
       while (p) {
          if (p % 2 == 1) res = res * b % m;
          b = b * b % m;
          p /= 2;
       }
       return res;
   template <typename T> T ModInv(T b,T m) {
       return BigMod(b , m - 2 , m);
   }
};
struct DoubleHashing {
   long long base[2] = {1949313259, 1997293877};
   long long mod[2] = \{2091573227, 2117566807\};
   vector <long long> pow[2] , inv[2];
   vector <long long> prehash[2] , sufhash[2];
   int maxN , flag = 0 , len;
   void Init(int n) {
       maxN = n + 2;
       for (int i = 0; i < 2; i++) {
          pow[i].resize(maxN);
          inv[i].resize(maxN);
       }
       Generate();
   void Generate() {
       ModularExponentiation Ex;
       for (int j = 0; j < 2; j++) {
          pow[j][0] = 1;
          inv[j][0] = 1;
          long long minv = Ex.ModInv(base[j] ,mod[j]);
          for (int i = 1; i < maxN; i++) {
              pow[j][i] = pow[j][i - 1] * base[j] % mod[j];
              inv[j][i] = inv[j][i - 1] * minv % mod[j];
          }
       }
```

```
}
long long GetHash(string &s) {
    long long hash_val[2] = {0 , 0};
    int n = s.size();
    for (int j = 0; j < 2; j++) {
        for (int i = 0; i < n; i++) {
            hash_val[j] = (hash_val[j] + s[i] * pow[j][i]) % mod[j];
    }
    return (hash val[0] << 32LL) | hash val[1];</pre>
void ComputeHash(string &s) {
   flag = 1;
    len = s.size();
    for (int j = 0; j < 2; j++) prehash[j].resize(maxN);
    for (int j = 0; j < 2; j++) prehash[j][0] = 0;
    for (int j = 0; j < 2; j++) {
        for (int i = 0; i < len; i++) {
            prehash[j][i + 1] = (prehash[j][i] + pow[j][i] * s[i]) % mod[j];
        }
    }
}
long long GetSubstrHash(int 1 , int r) {
    if (!flag) { cout << "ComputeHash\n"; return -1;}</pre>
    long long hash_val[2];
    for (int j = 0; j < 2; j++)
        hash\_val[j] = (prehash[j][r + 1] - prehash[j][l]) * inv[j][l] % mod[j];
    for (int j = 0; j < 2; j++) if (hash val[j] < 0) hash val[j] += mod[j];
    return (hash_val[0] << 32) | hash_val[1];</pre>
void ComputePreAndSufHash(string &s) {
    flag = 1;
    len = s.size();
    for (int j = 0; j < 2; j++) {
        prehash[j].resize(maxN);
        sufhash[j].resize(maxN);
    for (int j = 0; j < 2; j++) prehash[j][0] = sufhash[j][0] = 0;
    for (int j = 0; j < 2; j++) {
        for (int i = 0; i < len; i++) {
            prehash[j][i + 1] = (prehash[j][i] + pow[j][i] * s[i]) % mod[j];
            sufhash[j][i + 1] = (sufhash[j][i] + pow[j][len - i + 1] * s[i]) % mod[j];
        }
    }
long long GetPrefixHash(int 1 , int r) {
    return GetSubstrHash(1 , r);
long long GetSuffixHash(int 1 , int r) {
    if (!flag) { cout << "ComputePreAndSufHash\n"; return -1;}</pre>
    long long hash_val[2];
    for (int j = 0; j < 2; j++)
        hash_val[j] = (sufhash[j][r + 1] - sufhash[j][l]) * inv[j][len - r + 1] % mod[j];
    for (int j = 0; j < 2; j++) if (hash_val[j] < 0) hash_val[j] += mod[j];
    return (hash_val[0] << 32) | hash_val[1];</pre>
bool IsPallindrome(int 1 , int r) {
    return (GetPrefixHash(1 , r) == GetSuffixHash(1 , r));
vector <int> RabinKarp(string &txt , string &ptrn) {
```

```
long long ptrn_hash = GetHash(ptrn);
       vector <int> occurences;
       int txtlen = txt.size();
       int ptrnlen = ptrn.size();
       for (int i = 0; i < txtlen - ptrnlen + 1; i++) {
           long long cur_hash = GetSubstrHash(i , ((i + ptrnlen) - 1));
           // pattern match...
           if (cur_hash == ptrn_hash)
              occurences.emplace back(i + 1);
       }
       return occurences;
   }
};
                   // Building Prefix array
vector <int> BuildPrefixArray(string pattern) {
    vector <int> pfix(pattern.length());
    pfix[0] = 0;
    for(int i = 1 ,j = 0 ; i < pattern.length() ; ) {
       if(pattern[i] == pattern[j])
           pfix[i++] = ++j ;
       else {
           if(j == 0) pfix[i++] = 0;
           else j = pfix[j - 1];
   }
   return pfix;
int Kmp(string text , string pattern) {
   vector <int> pfix = BuildPrefixArray(pattern) ;
    int cnt = 0;
    for(int i = 0 , j = 0 ; i < (int)text.length() && j < (int)pattern.length() ; ) {</pre>
       if(text[i] == pattern[j]) {
           i++ ;
           j++ ;
       }
       else {
           if(j == 0) i++ ;
           else j = pfix[j - 1];
       if(j == (int)pattern.length()) {
           cnt++; // Number of occurances..
           j = pfix[j - 1];
       }
    }
    return cnt;
}
Here we discuss two problems at once.
   Given a string s of length n.
   In the first variation of the problem we want to count the number of appearances of each
prefix s[0...i] in the same string.
    In the second variation of the problem another string t is given and we want to count the
number of appearances of each prefix s[0...i] in t.
Let us compute the prefix function for s.
Using the last value of it we define the value k=n-\pi[n-1].
We will show, that if k divides n, then k will be the answer,
```

ComputeHash(txt);

```
otherwise there doesn't exists an effective compression and the answer is \ensuremath{\text{n.}}
```

```
vector<int> ans(n + 1);
for (int i = 0; i < n; i++)
   ans[pi[i]]++;
for (int i = n-1; i > 0; i--)
   ans[pi[i-1]] += ans[i];
for (int i = 0; i <= n; i++)
   ans[i]++;
template<typename T, bool maximum_mode = false>
struct RangeMinQuery {
   int n = 0;
   vector<vector<T>> sptab;
   static int LargestBit(int x) {
       return 31 - __builtin_clz(x);
   }
   static T Better(T a, T b) {
       return maximum_mode? max(a , b): min(a, b);
   }
   // O(NlogN)
   void Build(const vector<T> &values) {
       n = int(values.size());
       int levels = LargestBit(n) + 1;
       sptab.resize(levels);
       for (int k = 0; k < levels; k++) sptab[k].resize(n - (1 << k) + 1);
       if (n > 0) sptab[0] = values;
       for (int k = 1; k < levels; k++) {
          for (int i = 0; i <= n - (1 << k); i++) {
              sptab[k][i] = Better(sptab[k - 1][i], sptab[k - 1][i + (1 << (k - 1))]);
          }
       }
   }
   // 0(1)
   int Query(int a, int b) const {
       assert(0 <= a && a < b && b <= n);
       int level = LargestBit(b - a);
       return Better(sptab[level][a], sptab[level][b - (1 << level)]);</pre>
   }
};
RangeMinQuery<int>rmq;
// Complexity O(N * LogN)
struct SuffixArray{
   string s;
   int n;
   vector <int> sar , lcp, rank; // suffix array, lcp , rank
   SuffixArray() {}
   SuffixArray(string _s) {
       Init(_s);
   }
   void Init(string s) {
       s = _s;
       s += "$";
       n = s.size();
       sar.resize(n);
       rank.resize(n);
       lcp.resize(n);
       BuildSuffixArray();
       BuildLCPArray();
```

```
rmq.Build(lcp);
    }
    // O(NLogN)
    void BuildSuffixArray() {
        vector <int> cnt(256, 0), pos(256);
        for (int i = 0; i < n; i++) cnt[s[i]]++;
        pos[0] = 0;
        for (int i = 1; i < 256; i++) pos[i] = pos[i - 1] + cnt[i - 1];
        for (int i = 0; i < n; i++) {
            sar[pos[s[i]]] = i;
            pos[s[i]]++;
        rank[sar[0]] = 0;
        for (int i = 1; i < n; i++) {
            rank[sar[i]] = rank[sar[i - 1]];
            if (s[sar[i]] != s[sar[i - 1]]) rank[sar[i]]++;
        }
        pos.resize(n);
        for (int k = 0; (1 << k) < n; k++) {
            for (int i = 0; i < n; i++) sar[i] = (sar[i] - (1 <math><< k) + n) % n;
            cnt.assign(n, 0);
            for (int x : rank) cnt[x]++;
            pos[0] = 0;
            for (int i = 1; i < n; i++) pos[i] = pos[i - 1] + cnt[i - 1];
            vector <int> sar_new(n);
            for (int x : sar) {
                int i = rank[x];
                sar new[pos[i]] = x;
                pos[i]++;
            }
            sar = sar_new;
            vector <int> rank_new(n);
            rank_new[sar[0]] = 0;
            for (int i = 1; i < n; i++) {
                rank_new[sar[i]] = rank_new[sar[i - 1]];
                pair <int, int> prev = \{rank[sar[i - 1]], rank[(sar[i - 1] + (1 << k)) \%\}
n]};
                pair <int, int> cur = {rank[sar[i]] , rank[(sar[i] + (1 << k)) % n]};</pre>
                if (prev != cur) rank_new[sar[i]]++;
            }
            rank = rank_new;
        }
    }
    // algorithm of Kasai, Arimura, Arikawa, Lee and Park.
    // Complexity O(n)
    // Longest Common prefix of adjacent suffixes
    void BuildLCPArray() {
        int k = 0;
        for (int i = 0; i < n - 1; i++) {
            int pi = rank[i];
            int j = sar[pi - 1];
            while (s[i + k] == s[j + k]) k++;
            lcp[pi] = k;
            k = max(k - 1, 0);
        }
    }
    // 0(1)
    int GetLcpFromRanks(int a, int b) const {
        if (a == b) return n - sar[a];
        if (a > b) swap(a, b);
```

```
return rmq.Query(a + 1, b + 1);
    }
    // 0(1)
    int GetLcp(int a, int b) const {
        if (a >= n \mid | b >= n) return 0;
        if (a == b) return n - a;
        return GetLcpFromRanks(rank[a], rank[b]);
    }
    // Compares the substrings starting at `a` and `b` up to `length`
    // 0(1)
    int Compare(int a, int b, int length = -1) const {
        if (length < 0) length = n;</pre>
        if (a == b) return 0;
        int common = GetLcp(a, b);
        if (common >= length) return 0;
        if (a + common >= n \mid | b + common >= n) {
            return a + common >= n ? -1 : 1;
        }
        return s[a + common] < s[b + common] ? -1 : (s[a + common] == s[b + common] ? 0 :
1);
    }
    // O(N)
    long long NumOfUniqueSubStr() {
        long long sum = 0;
        for (int i = 1; i < lcp.size(); i++) sum += lcp[i];
        long long sz = n - 1; // actual size of string
        long long totalSubStrings = (sz * (sz + 1)) / 2LL;
        return totalSubStrings - sum;
    }
    // 0(|t|*LogN)
    bool IsSubstr(string &t) {
        int tlen = t.size();
        if (tlen > n) return false;
        int low = LowerBound(t);
        if (low < n && n - sar[low] >= tlen && s.compare(sar[low] , tlen, t) == 0) return
true;
        return false;
    }
    // O(|t|*LogN)
    int LowerBound(string &t) {
        int low = 0 ,high = n ,tlen = t.size();
        while (low < high) {</pre>
            int mid = (low + high) >> 1;
            if (s.compare(sar[mid], tlen, t) < 0)</pre>
                low = mid + 1;
            else
                high = mid;
        return low;
    }
    // 0(|t|*LogN)
    int UpperBound(string &t) {
        int low = 0 ,high = n ,tlen = t.size();
        while (low < high) {
            int mid = (low + high) >> 1;
            if (s.compare(sar[mid], tlen, t) <= 0)</pre>
                low = mid + 1;
            else
                high = mid;
        }
```

```
return low;
   }
   // O(|t|*LogN)
   int SubstrOccurences(string &t) {
      int tlen = t.size();
      int low = LowerBound(t) , up = UpperBound(t);
       if (low == up || s.compare(sar[low] , tlen , t) != 0)                       return 0;
      return (up - low);
   }
   // Longest common substring of two string
   // O(NLogN)
   string LongestCommonSubstring(string &a, string &b) {
      int alen = a.size();
      Init(a + "&" + b);
      int tot = sar.size();
      bool color[tot];
      for (int i = 0; i < tot; i++) {
          if (sar[i] < alen) color[i] = 1;</pre>
          else color[i] = 0;
      int mx = -1, mxid = -1;
      for (int i = 1; i < tot; i++) {
          if (color[i] != color[i - 1] && mx < lcp[i]) {
             mx = lcp[i];
             mxid = sar[i];
          }
       if (mxid != -1) return s.substr(mxid , mx);
      return "";
   void ShowSuffixArray() {
      cout << "Rank LCP
                        SA
                            Suffixes\n";
      for (int i = 0; i < n; i++) {
          cout << rank[i] << " " << lcp[i] << " " << sar[i] << " " <<
s.substr(sar[i] , n - sar[i]) << "\n";
      }
   }
};
/*.....*/
struct node{
   bool endmark;
   node *next[27];
   node() {
      endmark = false ;
      for(int i = 0; i < 26; i++)
          next[i] = NULL ;
   }
};
node *root = new node();
void Insert(string str , int len) {
   node *curr = root;
   for(int i = 0 ; i < len ; i++ ) {
      int id = str[i] - 'a';
      if(curr->next[id] == NULL)
          curr->next[id] = new node();
      curr = curr->next[id] ;
   }
   curr->endmark = true ;
```

```
bool Search(string str , int len) {
   node *curr = root ;
   for(int i = 0; i < len; i++) {
       int id = str[i] - 'a';
       if(curr->next[id] == NULL)
           return false;
       curr = curr->next[id] ;
    return curr->endmark;
}
/*******************Manacher's Algorithm********************/
/*
   Algorithm:
   1. Finds all sub-palindromes in O(N) , Applications: 1. Finds Longest Palindromes O(N)
*/
/*
   1. string is 0 based indexed of length n.
   2. the manacher will make a string of length 2 * n + 1.
       ex: "abba" will be "#a#b#b#a#".
    3. Odd length palindromes:
       for all i = 0 to i < n
           i is a center and stored maxlen palindrome in 2i + 1 th index in P.
   4. Even length palindromes:
       there are n - 1 centered positions.
       for all i = 0 to i < n - 1
           max length palindrome centered at i found in 2i + 2 th index in P.
*/
vector<int> Manacher_Odd_Length(string T) {
    int n = T.size();
   vector <int> P(n);
   int C = 0, R = -1, rad;
   for (int i = 0; i < n; ++i) {
       if (i <= R) {
           rad = min(P[2 * C - i], R - i);
       } else {
           rad = 0;
       // Try to extend
       while (i + rad < n \&\& i - rad >= 0 \&\& T[i - rad] == T[i + rad]) {
           rad++;
       }
       P[i] = rad;
       if (i + rad - 1 > R) {
           C = i;
           R = i + rad - 1;
       }
   for (int i = 0; i < (int)P.size(); i ++) {
       P[i]--;
    }
   return P;
vector <int> Manacher(const string &s) {
   string t;
   for(auto c: s) {
       t += string("#") + c;
   }
```

```
return Manacher_Odd_Length(t + "#");
}
string LongestPalindrome(const string &s) {
   vector <int> P = Manacher(s);
   int mxlen = 0, startpos = -1;
   for (int i = 0; i < (int)s.size(); i++) {</pre>
       if (mxlen < P[2 * i + 1]) {
          mxlen = P[2 * i + 1];
          startpos = i - P[2 * i + 1] / 2;
       }
   }
   for (int i = 0; i < (int)s.size() - 1; i++) {
       if (mxlen < P[2 * i + 2]) {
          mxlen = P[2 * i + 2];
          startpos = i - P[2 * i + 2] / 2 + 1;
       }
   }
   return s.substr(startpos, mxlen);
2147483647
#define INF9
#define INF18
                  9223372036854775806
template <typename T> struct SegmentTree {
   vector <T> seg;
   vector <T> lazy;
   vector <T> ar;
   int type , up;
   SegmentTree() {
       type = 0;
       up = 0;
   SegmentTree(int tp , int u) {
       type = tp;
       up = u;
   void Init(int N) {
       seg.assign(N << 2, 0);</pre>
       lazy.assign(N << 2 , 0);</pre>
   void Init(vector <T> &s) {
       Init(s.size() + 1);
       ar = s;
   void PushDown(int cur , int left , int right) {
       if (type == 0) {
          if (up == 1) seg[cur] += (right - left + 1) * lazy[cur];
          else seg[cur] = (right - left + 1) * lazy[cur];
       } else {
          if (up == 1) seg[cur] += lazy[cur];
          else seg[cur] = lazy[cur];
       }
       if (left != right) {
          if (up == 1) {
              lazy[cur << 1] += lazy[cur];</pre>
              lazy[cur << 1 | 1] += lazy[cur];</pre>
          } else {
              lazy[cur << 1] = lazy[cur];</pre>
```

```
lazy[cur << 1 | 1] = lazy[cur];</pre>
            }
        lazy[cur] = 0;
    T Merge(T x , T y) {
        if (type == 0) return x + y;
        if (type == 1) return max(x , y);
        if (type == 2) return min(x , y);
    void Build(int cur , int left , int right) {
        lazy[cur] = 0;
        if (left == right) {
            seg[cur] = ar[left];
            return;
        }
        int mid = (left + right) >> 1;
        Build(cur << 1 , left , mid);</pre>
        Build(cur << 1 | 1 , mid + 1 , right);
        seg[cur] = Merge(seg[cur << 1] , seg[cur << 1 | 1]);</pre>
    void Update(int cur , int left , int right , int pos , T val) {
        Update(cur , left , right , pos , pos , val);
    void Update(int cur , int left , int right , int l , int r , T val) {
        if (lazy[cur] != 0) PushDown(cur , left , right);
        if (1 > right || r < left) return;</pre>
        if (left >= 1 && right <= r) {
            if (up == 0) lazy[cur] = val;
            else lazy[cur] += val;
            PushDown(cur , left , right);
            return;
        }
        int mid = (left + right) >> 1;
        Update(cur << 1 , left , mid , l , r , val);</pre>
        Update(cur << 1 | 1 , mid + 1 , right , l , r , val);</pre>
        seg[cur] = Merge(seg[cur << 1] , seg[cur << 1 | 1]);
    T Query(int cur , int left , int right , int l , int r) {
        if (1 > right || r < left) {
            if (type == 0) return 0;
            if (type == 1) return -INF18;
            if (type == 2) return INF18;
        if (lazy[cur] != 0) PushDown(cur , left , right);
        if (left >= 1 && right <= r) return seg[cur];
        int mid = (left + right) >> 1;
        T p1 = Query(cur << 1 , left , mid , l , r);
        T p2 = Query(cur << 1 | 1 , mid + 1 , right , l , r);
        return Merge(p1 , p2);
    }
//for sum = 0, max = 1, min = 2, for assignment update send 0 or 1 for increment.
SegmentTree <long long> tr(0 , 0);
/***********************************/
const int mxN = 1e3;
int a[mxN + 1][mxN + 1];
int t[mxN << 2][mxN << 2];
int m;
```

**}**;

```
void build_y(int vx, int lx, int rx, int vy, int ly, int ry) {
    if (ly == ry) {
        if (lx == rx)
            t[vx][vy] = a[lx][ly];
        else
            t[vx][vy] = t[vx*2][vy] + t[vx*2+1][vy];
    } else {
        int my = (ly + ry) / 2;
        build y(vx, 1x, rx, vy*2, 1y, my);
        build_y(vx, lx, rx, vy*2+1, my+1, ry);
        t[vx][vy] = t[vx][vy*2] + t[vx][vy*2+1];
    }
}
void build_x(int vx, int lx, int rx) {
    if (lx != rx) {
        int mx = (lx + rx) / 2;
        build_x(vx*2, lx, mx);
        build_x(vx*2+1, mx+1, rx);
    build_y(vx, lx, rx, 1, 0, m-1);
int sum_y(int vx, int vy, int tly, int try_, int ly, int ry) {
    if (ly > ry)
        return 0;
    if (ly == tly && try_ == ry)
        return t[vx][vy];
    int tmy = (tly + try) / 2;
    return sum_y(vx, vy*2, tly, tmy, ly, min(ry, tmy))
         + sum_y(vx, vy*2+1, tmy+1, try_, max(ly, tmy+1), ry);
int sum_x(int vx, int tlx, int trx, int lx, int rx, int ly, int ry) {
    if (lx > rx)
        return 0;
    if (1x == t1x \&\& trx == rx)
        return sum_y(vx, 1, 0, m-1, ly, ry);
    int tmx = (tlx + trx) / 2;
    return sum_x(vx*2, tlx, tmx, lx, min(rx, tmx), ly, ry)
         + sum_x(vx*2+1, tmx+1, trx, max(lx, tmx+1), rx, ly, ry);
void update_y(int vx, int lx, int rx, int vy, int ly, int ry, int x, int y, int new_val) {
    if (ly == ry) {
        if (lx == rx)
            t[vx][vy] ^= 1;
        else
            t[vx][vy] = t[vx*2][vy] + t[vx*2+1][vy];
    } else {
        int my = (ly + ry) / 2;
        if (y \le my)
            update_y(vx, lx, rx, vy*2, ly, my, x, y, new_val);
        else
            update_y(vx, lx, rx, vy*2+1, my+1, ry, x, y, new_val);
        t[vx][vy] = t[vx][vy*2] + t[vx][vy*2+1];
    }
void update_x(int vx, int lx, int rx, int x, int y, int new_val) {
    if (lx != rx) {
        int mx = (lx + rx) / 2;
        if (x <= mx)
            update_x(vx*2, lx, mx, x, y, new_val);
```

```
else
          update_x(vx*2+1, mx+1, rx, x, y, new_val);
   update_y(vx, lx, rx, 1, 0, m-1, x, y, new_val);
}
int main() {
   int n, q; cin >> n >> q;
   m = n;
   for (int i = 0; i < n; i++) {
       for (int j = 0; j < n; j++) {
          char ch; cin >> ch;
          a[i][j] = (ch == '*')? 1 : 0;
       }
   }
   build_x(1 , 0 , m - 1);
   while (q--) {
       int typ; cin >> typ;
       if (typ == 2) {
          int y1, x1, y2, x2;
          cin >> y1 >> x1 >> y2 >> x2;
          y1--, x1--,y2--,x2--;
          cout << sum_x(1 , 0 , m - 1, y1 , y2 , x1 , x2) << '\n';
       } else {
          int y, x; cin >> y >> x;
          y--,x--;
          update_x(1 , 0 , m - 1, y , x , 1);
       }
   }
   return 0;
// 1. assign max size with constructor
// 2. each test case call makeset
struct DSU {
   vector <int> parent;
   vector <int> siz;
   DSU(int mxN) {
      mxN <<= 1;
      parent.resize(mxN + 1);
      siz.resize(mxN + 1);
   void Makeset(int n) {
      for (int i = 1; i <= n; i++) {
          parent[i] = n + i;
          parent[n + i] = n + i;
          siz[n + i] = 1;
       }
   int Find(int u) {
       if (parent[u] == u) return u;
       return parent[u] = Find(parent[u]);
   void Union(int u , int v) {
      u = Find(u);
      v = Find(v);
       if (u != v) {
          if (siz[u] < siz[v]) swap(u , v);</pre>
          parent[v] = u;
          siz[u] += siz[v];
       }
```

```
bool SameSet(int u , int v) {
       return (Find(u) == Find(v));
   void MoveUtoSetV(int u , int v) {
       if (SameSet(u , v)) return;
       int x = Find(u);
       int y = Find(v);
       siz[x]--;
       siz[y]++;
       parent[u] = y;
   }
   int Size(int u) {
       return siz[Find(u)];
   }
};
1. All nodes are number from 0 to n - 1 */
   2. Assign the graph by Init(graph) or simply Init(total nodes) and
       call AddEdge(u , v) for all the edges */
      Must be Take the node value from input directly or use the
      TakeNodeVal(nodeval) to assigning the node value */
   4. Call Build() to construct hld and segment tree */
/*
      simply use the path query by query(u , v) and update(pos , val)
/* 6. use optimized segment tree sometimes it cz TLE
/*******************Segment tree*************************/
//for sum = 0, max = 1, min = 2, for assignment update send 0 or 1 for increment.
SegmentTree <long long> T(0 , 0);
struct HeavyLightDecompose {
   vector <vector <int> > g ; // graph
   vector <long long> node_val;
   int N, root = 0;
   vector <int> depth , parent , sub;
   // HLD staffs
   int chain_no, indx;
   vector <int> chain_head , chain_ind;
   vector <int> node_serial , serial_node;
   vector <long long> segarr; // tree on linear format
   void Init(int n) {
      N = n;
       g.assign(N , {});
      node_val.clear();
       segarr.resize(N);
       depth.resize(N);
       parent.resize(N);
       sub.resize(N);
       chain_head.assign(N, -1);
       chain_ind.resize(N);
       node_serial.resize(N);
       serial node.resize(N);
       return;
   void Init(const vector <vector<int>> &_g) {
       Init(_g.size());
       g = g;
       return;
   void AddEdge(int u , int v) {
```

```
g[u].push_back(v);
    g[v].push_back(u);
    return;
void TakeNodeVal(const vector <long long> &_node_val) {
    node_val = _node_val;
}
void Build() {
    Dfs(root);
    chain no = 0, indx = 0;
    HLD(0);
    T.Init(segarr);
    T.Build(1 , 0 , N - 1);
}
void Dfs(int u, int par = -1) {
    sub[u] = 1;
    if (par == -1) {
        depth[u] = 0;
        parent[u] = -1;
    for (int v : g[u]) {
        if (v == par) continue;
        parent[v] = u;
        depth[v] = depth[u] + 1;
        Dfs(v, u);
        sub[u] += sub[v];
    }
    return;
void HLD(int u , int par = -1) {
    if (chain_head[chain_no] == -1) chain_head[chain_no] = u;
    chain_ind[u] = chain_no;
    node_serial[u] = indx;
    serial_node[indx] = u;
    segarr[indx] = node_val[u]; // tree flatting..
    indx++;
    int heavychild = -1 , heavysize = 0;
    for (int v : g[u]) {
        if (v == par) continue;
        if (sub[v] > heavysize) {
            heavysize = sub[v];
            heavychild = v;
        }
    }
    if (heavychild != -1) HLD(heavychild , u);
    for (int v : g[u]) {
        if (v != par && v != heavychild) {
            chain_no++;
            HLD(v, u);
        }
    }
    return;
void Update(int p , int val) {
    T.Update(1 , 0 , N - 1, node_serial[p] , val);
    node_val[p] = val;
long long Query(int u , int v) {
    long long ans = 0;
```

```
for ( ; chain_ind[u] != chain_ind[v] ; v = parent[chain_head[chain_ind[v]]]) {
          if (depth[chain_head[chain_ind[u]]] > depth[chain_head[chain_ind[v]]])
              swap( u , v );
          ans += T.Query(1 , 0 , N - 1 , node_serial[chain_head[chain_ind[v]]] ,
node_serial[v]);
       if (depth[u] > depth[v])
          swap(u, v);
       ans += T.Query(1 , 0 , N - 1 , node_serial[u] , node_serial[v]);
       return ans;
   }
} hd;
/*********************************
/*********************************
/*.....*/
/st Given an array of length n and q querys of range 1 , r . Find the number of unique
elements in the given range */
const int BLOCK = 555;
const int mxN = 100000;
struct query {
   int l , r , idx;
};
query Q[mxN + 5];
int ar[mxN + 5], ans[mxN + 5];
int freq[mxN + 5];
int cnt = 0;
bool Cmp(query &a , query &b) {
   if (a.l / BLOCK != b.l / BLOCK)
       return a.1 / BLOCK < b.1 / BLOCK;
   return a.r < b.r;
}
void Add(int pos) {
   freq[ar[pos]]++;
   if (freq[ar[pos]] == 1) cnt++;
void Remove(int pos) {
   freq[ar[pos]]--;
   if (freq[ar[pos]] == 0) cnt--;
void Input_Query(int q) {
   for (int i = 0; i < q; i++) {
       cin >> Q[i].1 >> Q[i].r;
       Q[i].idx = i;
       Q[i].1--; Q[i].r--;
   }
void MosAlgo(int q) {
   Input_Query(q);
   sort(Q, Q + q, Cmp);
   int ML = 0 , MR = -1;
   for (int i = 0; i < q; i++) {
       int L = Q[i].1;
       int R = Q[i].r;
       while (ML > L) Add(--ML);
       while (MR < R) Add(++MR);
       while (ML < L) Remove(ML++);
       while (MR > R) Remove(MR--);
       ans[Q[i].idx] = cnt;
```

```
}
}
int main() {
   FasterI0
   int n , q ;
   cin >> n;
   for(int i = 0; i < n; i++)
       cin >> ar[i] ;
   cin >> q;
   MosAlgo(q);
   for(int i = 0; i < q; i++) {
       cout << ans[i] << "\n" ;
   return 0;
              ********************************
// 1 based index...
const int N = 3e5, M = N;
const int MAX = 1e6;
int a[N];
struct wavelet_tree {
   int lo, hi;
   wavelet_tree *1, *r;
   vector <int> b;
   vector <int> c; // c holds the prefix sum of elements
   //nos are in range [x,y]
   //array indices are [from, to)
   wavelet_tree(int *from, int *to, int x, int y) {
       lo = x, hi = y;
       if( from >= to) return;
       if( hi == lo ) {
          b.reserve(to - from + 1);
          b.push_back(0);
          c.reserve(to - from + 1);
          c.push_back(0);
          for(auto it = from; it != to; it++) {
              b.push_back(b.back() + 1);
              c.push_back(c.back() + *it);
          }
          return ;
       int mid = (lo + hi) / 2;
       auto f = [mid](int x) {
          return x <= mid;
       b.reserve(to - from + 1);
       b.push back(0);
       c.reserve(to - from + 1);
       c.push_back(0);
       for(auto it = from; it != to; it++) {
          b.push_back(b.back() + f(*it));
          c.push_back(c.back() + *it);
       }
       //stable_partition the lamda function
       auto pivot = stable_partition(from, to, f);
      1 = new wavelet_tree(from, pivot, lo, mid);
       r = new wavelet_tree(pivot, to, mid + 1, hi);
   // swap a[i] with a[i+1] , if a[i]!=a[i+1] call swapadjacent(i)
```

```
void swapadjacent(int i) {
        if(lo == hi) return ;
        b[i] = b[i - 1] + b[i + 1] - b[i];
        c[i] = c[i - 1] + c[i + 1] - c[i];
        if(b[i+1] - b[i] == b[i] - b[i-1]) {
            if(b[i] - b[i - 1])
                return this->l->swapadjacent(b[i]);
            else
                return this->r->swapadjacent(i - b[i]);
        return;
    //kth smallest element in [l, r]
    int kth(int 1, int r, int k) {
        if(1 > r) return 0;
        if(lo == hi) return lo;
        int inLeft = b[r] - b[l - 1];
        int lb = b[1 - 1]; //amt of nos in first (1-1) nos that go in left
        int rb = b[r]; //amt of nos in first (r) nos that go in left
        if(k <= inLeft) return this->l->kth(lb + 1, rb, k);
        return this->r->kth(1 - lb, r - rb, k - inLeft);
    }
    //count of nos in [l, r] Less than or equal to k
    int LTE(int 1, int r, int k) {
        if(1 > r \mid \mid k < lo) return 0;
        if(hi \leftarrow k) return r - l + 1;
        int lb = b[1-1], rb = b[r];
        return this->l->LTE(lb + 1, rb, k) + this->r->LTE(l - lb, r - rb, k);
    }
    //count of nos in [1, r] equal to k
    int count(int 1, int r, int k) {
        if(l > r || k < lo || k > hi) return 0;
        if(lo == hi) return r - l + 1;
        int lb = b[1 - 1], rb = b[r], mid = (lo + hi) / 2;
        if(k <= mid) return this->l->count(lb + 1, rb, k);
        return this->r->count(1 - lb, r - rb, k);
    //sum of nos in [l ,r] less than or equal to k
    int sumk(int 1, int r, int k) {
        if(l > r \mid \mid k < lo) return 0;
        if(hi <= k) return c[r] - c[l-1];
        int 1b = b[1 - 1], rb = b[r];
        return this->l->sumk(lb + 1, rb, k) + this->r->sumk(l - lb, r - rb, k);
    ~wavelet_tree() {
       delete 1;
        delete r;
    }
int main() {
    int n ; cin >> n;
    for(int i = 1; i <= n; i++) {
       cin >> a[i];
    }
    // wavelet_tree T(array start address, array end address, min element, max element);
   wavelet_tree T(a + 1, a + n + 1, 1, MAX);
    int q; cin >> q;
   while(q--) {
        int x; cin >> x;
        int k, l, r;
```

};

```
if(x == 0) {
           //kth smallest element in range [1 , r]
           cin >> 1 >> r >> k;
           cout << "Kth smallest: ";</pre>
           cout << T.kth(l, r, k) << endl;</pre>
       if(x == 1) {
           //Number of Elements less than or equal to K in range [l , r]
           cin >> 1 >> r >> k;
           cout << "LTE: ";</pre>
           cout << T.LTE(1, r, k) << endl;</pre>
       if(x == 2) {
           //count occurence of K in [1, r]
           cin >> 1 >> r >> k;
           cout << "Occurence of K: ";</pre>
           cout << T.count(1, r, k) << endl;</pre>
       if(x == 3) {
           //sum of elements less than or equal to K in [1, r]
           cin >> 1 >> r >> k;
           cout << "Sum: ";</pre>
           cout << T.sumk(1, r, k) << endl;</pre>
       if(x == 4) {
           int pos ; cin >> pos ;
           if(a[pos] != a[pos + 1])
               T.swapadjacent(pos);
       }
   }
   return 0;
    /****************************/owest Common Anceston********************/
/* 1.
       All nodes are number 0 to n - 1 */
       simply Init(total nodes) and call AddEdge(u , v) for all the edges */
/* 2.
       Call Build() to run dfs and build the sparse table */
/* 3.
struct LowestCommonAncestor {
   int N, root = 0, po;
   vector <vector <int> > g;
   vector <vector <int> > sptab;
   vector <int> depth;
   vector <int> parent;
   void Init(int _n) {
       N = _n;
       po = log2((N)) + 1;
       g.assign(N, {});
       depth.resize(N);
       parent.resize(N);
       sptab.assign(N, {});
   void AddEdge(int u , int v) {
       g[u].push_back(v);
       g[v].push_back(u);
   void Dfs(int u , int par = -1) {
       if(par == -1) {
```

```
depth[u] = 0;
        parent[u] = -1;
   for(int v : g[u]) {
        if (v == par) continue;
        parent[v] = u;
        depth[v] = depth[u] + 1;
        Dfs(v, u);
    }
void SparceTable() {
    for(int i = 0; i < N; i++) sptab[i][0] = parent[i];</pre>
    for(int j = 1; (1 << j) < N; j++) {
        for(int i = 0; i < N; i++) {
            if(sptab[i][j - 1] != -1) {
                sptab[i][j] = sptab[sptab[i][j - 1]][j - 1];
            }
        }
    }
}
void Build() {
   for(int i = 0; i < N; i++) {
        for(int j = 0; j \le po; j++) {
            sptab[i].push_back(-1);
        }
   Dfs(root);
   SparceTable();
}
int Lca(int u , int v) {
    if(depth[u] < depth[v]) swap(u , v);</pre>
    for(log = 1 ; (1 << log) <= depth[u] ; log++); log--;</pre>
    for(int i = log ; i >= 0 ; i--) {
        if(depth[u] - (1 << i) >= depth[v]) {
            u = sptab[u][i];
        }
    if(u == v) return u;
    for(int i = log ; i >= 0 ; i--) {
        if(sptab[u][i] != -1 && sptab[u][i] != sptab[v][i]) {
            u = sptab[u][i];
            v = sptab[v][i];
        }
    }
    return parent[u];
int KthAncestor(int u, int k) {
    int log;
    for(log = 1; (1 << log) <= depth[u]; log++); log--;
    for(int i = log; i >= 0; i--) {
        if(k - (1 << i) >= 0) {
            u = sptab[u][i];
            k = (1 << i);
        }
    }
    return u;
int Getdist(int u , int v) {
    return (depth[u] + depth[v] - (2 * (depth[Lca(u , v)])));
```

```
}
   bool IsAnsector(int u , int v) {
      int cur = Lca(u , v);
      if(cur == u) return 1;
      return 0;
} lca;
struct StronglyConnectedComponent {
   vector <vector <int>> g , gr;
   vector <bool> vis;
   vector <int> order, sccid;
   int nodes , edges , scc = 0;
   void Init(int _nodes) {
      nodes = _nodes;
      edges = 0;
      scc = 0;
      g.clear();
      gr.clear();
      g.assign(nodes + 1 , {});
      gr.assign(nodes + 1, {});
      order.clear();
      vis.assign(nodes + 1, 0);
      sccid.resize(nodes + 1);
   void AddEdge(int u , int v) {
      if (u == -1 || v == -1) return;
      g[u].push_back(v);
      gr[v].push_back(u);
      edges++;
   }
   void Init() {
      for (int i = 0; i <= nodes; i++) vis[i] = 0;
   void Dfs1(int u) {
      if (vis[u]) return;
      vis[u] = 1;
      for (int v : g[u]) Dfs1(v);
      order.push_back(u);
   }
   void TopSort() {
      Init();
      for (int i = 0; i < nodes; i++) {
          if (!vis[i]) Dfs1(i);
      }
      reverse(order.begin() , order.end());
   void Dfs2(int u, int id) {
      if (vis[u]) return;
      vis[u] = 1;
      sccid[u] = id;
      for (int v : gr[u]) Dfs2(v, id);
   void KosarajuSCC() {
      TopSort();
      Init();
      for (int i : order) {
          if (!vis[i]) {
```

```
Dfs2(i , scc);
              scc++;
           }
       }
   }
   int GetAns() {
       vector <int> here(nodes + 1);
       for (int i = 0; i < scc; i++ ) {
           here[i] = 1;
       for (int i = 0; i < nodes; i++) {
           for (int j : g[i]) {
              if (sccid[i] != sccid[j])
                  here[sccid[j]] = 0;
           }
       }
       int ans = 0;
       for (int i = 0; i < scc; i++) ans += here[i];
       return ans;
   }
};
/*.....finding strongly connected component............
/* Given a graph , Find if there is one or more connected component. print 1 for 1 or 0 for
more.. */
// complexity : O(V + E)
#define mxN 2000
vector \langle int \rangle G[mxN + 10] , GR[mxN + 10] ;
vector <int> order , component ;
bool vis[mxN + 10];
void init(int nodes) {
   for(int i = 0 ; i <= nodes ; i++)
       vis[i] = 0;
}
void dfs1(int u) {
   vis[u] = 1;
   for(int v : G[u]) {
       if(!vis[v])
           dfs1(v);
   order.push_back(u);
}
void reverse_edges(int nodes , int edges) {
   for(int i = 1; i <= nodes; i++) {
       for(int j : G[i]) {
           GR[j].push_back(i);
       }
   }
}
void dfs2(int u) {
   vis[u] = 1;
   component.push_back(u);
   for(int v : GR[u]) {
       if(!vis[v])
           dfs2(v);
   }
}
void KosarajuSCC(int nodes ,int edges) {
```

```
init(nodes);
   for(int i = 1; i <= nodes; i++) {
       if(!vis[i])
           dfs1(i);
    }
   reverse_edges(nodes , edges) ;
    init(nodes);
   int cnt = 0;
   for(int i = order.size() - 1; i >= 0; i--) {
       int v = order[i];
       if(!vis[v]) {
           dfs2(v);
           cnt++; // Component containg here all the nodes which are in a SCC..
           component.clear();
       }
   if(cnt == 1)
       cout << "1\n" ;
   else
       cout << "0\n" ;
}
int main() {
   int nodes , edges ;
   while(cin >> nodes >> edges) {
       if(nodes == edges && nodes == 0)
           break;
       for(int i = 0; i < edges; i++) {
           int u , v , w ;
           cin >> u >> v >> w ;
           if(w == 1) {
              G[u].push_back(v);
           } else {
              G[u].push_back(v);
              G[v].push_back(u);
           }
       KosarajuSCC(nodes , edges) ;
       graph_clear(nodes);
   return 0;
/************************************/hortest Path Faster Algorithm**********************/
#define INF 1e9
vector < pair <int , int> > G[1010] ;
int cost[1010] ; // containing shortest path costs..
int vis[1010];
// The chinese algorithm..
void ShortestPathFasterAlgorithm(int nodes) {
   for(int i = 0; i < 1010; i++) {
       cost[i] = INF ;
       vis[i] = 0;
   cost[0] = 0; // setting the source cost as 0
   queue <int> q;
   q.push(0);
   while(!q.empty()) {
       int u = q.front();
       q.pop();
       vis[u] = 1;
```

```
for(int i = 0; i < G[u].size(); i++) {
          int v = G[u][i].first;
          int wv = G[u][i].second ;
          if(cost[u] + wv < cost[v]) {</pre>
             cost[v] = cost[u] + wv;
             if(!vis[v]) {
                 q.push(v);
                 vis[v] = 1;
             }
          }
      }
   }
}
// If there exist a negative weight cycle returning true
bool NegativeCycle(int nodes) {
   for(int u = 0; u < nodes; u++) {
      for(int i = 0; i < G[u].size(); i++) {
          int v = G[u][i].first;
          int wv = G[u][i].second ;
          if(cost[v] > cost[u] + wv)
             return 1;
      }
   }
   return 0;
         ***********************************
typedef long long 11;
typedef unsigned long long ull;
typedef long double ld;
支持更高的快速幂操作
expower.mod_pow(a,b,mod);
struct Expower {
   ull Mod_mul(ull a, ull b, ull M) {
      11 \text{ ret} = a * b - M * ull(ld(a) * ld(b) / ld(M));
      return ret + M * (ret < 0) - M * (ret >= (11)M);
   ull Mod_pow(ull b, ull e, ull mod) {
      ull ans = 1;
      for (; e; b = Mod_mul(b, b, mod), e /= 2)
          if (e & 1) ans = Mod_mul(ans, b, mod);
      return ans % mod;
   }
} Expower;
struct BigPrime {
   Miller-Rubin 素数判别
   is prime(n);
   bool IsPrime(ull n) {
      if (n < 2 | | n % 6 % 4 != 1)
          return n - 2 < 2;
      ull A[] = {2, 325, 9375, 28178, 450775, 9780504, 1795265022};
      ull s = __builtin_ctzll(n - 1), d = n >> s;
      for (auto a : A) {
          ull p = Expower.Mod_pow(a, d, n), i = s;
```

```
while (p != 1 && p != n - 1 && a % n && i--)
              p = Expower.Mod_mul(p, p, n);
           if (p != n - 1 && i != s)
              return 0;
       }
       return 1;
   }
   /*
   素因数分解
   ret=factorization(n);
   */
   ull Pollard(ull n) {
       auto f = [n](ull x) {
           return ( Expower.Mod_mul(x, x, n) + 1) % n;
       };
       if (!( n & 1))
           return 2;
       for (ull i = 2; i++) {
           ull x = i, y = f(x), p;
           while ((p = \_gcd(n + y - x, n)) == 1)
              x = f(x), y = f(f(y));
           if (p != n) return p;
       }
   }
   vector <ull> Factorization(ull n){
       if (n == 1) return {};
       if (IsPrime(n)) return {n};
       ull x = Pollard(n);
       auto 1 = Factorization(x), r = Factorization(n/x);
       1.insert(1.end(), begin(r), end(r));
       return 1;
   }
};
int main() {
   FasterI0
   BigPrime ob;
   int tc; cin >> tc;
   while (tc--) {
       ull n; cin >> n;
       if (ob.IsPrime(n)) {
           cout << "YES\n";</pre>
       } else {
          cout << "NO\n";</pre>
       }
   }
   return 0;
    ***********************************
/ // Build Complexity : O(NlogN)
void SieveOfEulersPhi() {
   for (int i = 0; i <= mxN; i++) phi[i] = 0;
   phi[1] = 1;
   for (int i = 2; i <= mxN; i++) {
       if (phi[i] == 0) {
           phi[i] = i - 1;
           for (int j = i + i; j \leftarrow mxN; j += i) {
              if (phi[j] == 0) phi[j] = j;
              phi[j] -= phi[j] / i;
           }
```

```
}
   }
}
// if n <= 1e7, always works with O(1) per query
// if n >= 1e7, most often works with O(logN) per query but sometimes goes O(sqrtN) in very
rare.
long long Phi(long long n) {
   if (n <= mxN) return phi[n];</pre>
   long long coprime = n;
   for (int i = 0; i < tot_primes; i++) {</pre>
       if (n \le mxN \&\& isp[n] == 0) break;
       long long x = prime[i];
       if (x * x > n) break;
       if (n \% x == 0) {
           while (n \% x == 0) n /= x;
           coprime -= coprime / x;
       }
   }
   if (n != 1) coprime -= coprime / n;
   return coprime;
}
/****************************
// Sum of Number of divisors in range 1 to N .
// Complexity O(sqrt(N))
int SNOD( int n ) {
int res = 0;
                int u =
sqrt(n);
   for ( int i = 1; i <= u; i++ ) {
res += ( n / i ) - i; //Step 1
   }
   res *= 2; //Step 2
res += u; //Step 3
return res;
// sum of coprimes of n int
sumofcoprimesN(int n){
int x = phi(n);
                   int ans
= (x * n) / 2;
                   return
ans;
/*********************************
// Sum of Number of divisors in range 1 to N .
// Complexity O(sqrt(N))
int SNOD( int n ) {
   int res = 0;
   int u = sqrt(n);
   for ( int i = 1; i <= u; i++ ) {
       res += ( n / i ) - i; //Step 1
   res *= 2; //Step 2
   res += u; //Step 3
   return res;
}
// sum of coprimes of n
int sumofcoprimesN(int n){
   int x = phi(n);
    int ans = (x * n) / 2;
   return ans;
}
```