## **Exhaustive Algorithm Solution Pseudocode and Time Analysis**

```
function crane unloading exhaustive(setting):
  assert setting.rows() > 0 —---> 1 tu
  assert setting.columns() > 0 —----> 1tu
  max steps = setting.rows() + setting.columns() - 2 ----->1tu
  assert max steps < 64 —----> 1 tu
best = initialize best path(setting) —----> 1tu
for steps = 0 to max steps:
  for bits = 0 to pow(2, steps) - 1:
     nextStep = initialize path with setting(setting) —----> 1tu
     valid = true —----> 1tu
     for k = 0 to steps - 1:
       bit = (bits >> k) & 1
       if bit == 1 -----> 1tu
         step_direction = go right ----->1tu
       else
         step direction = go down—---->1tu
       if nextStep.is step valid(step direction) —---->1tu
         nextStep.add step(step direction) ----->1tu
       else
         valid = false -----> 1tu
         break
     endfor
    if valid and nextStep.total cranes() > best.total cranes() —----> 1tu
       best = nextStep \longrightarrow 1tu
   endfor
endfor
return best
      max pow(2,s) s-1
5 + \sum
           \sum 4 \sum 6
```

Since there's a power function, then the Exhaustive Algorithm will end up having an exponential time complexity.

## **Dynamic Algorithm Solution Pseudocode and Time Analysis**

function crane\_unloading\_dyn\_prog(setting):

```
assert setting.rows() \geq 0 —---> 1 tu
assert setting.columns() > 0 —----> 1 tu
cell type = optional < path> -----> 1 tu
A = create 2d vector(setting.rows(), setting.columns()) —----> r * c
A[0][0] = \text{create new path(setting)} \longrightarrow 1 \text{ tu}
assert A[0][0].has value() —----> 1 tu
for r = 0 to setting.rows() - 1
  for c = 0 to setting.columns() - 1
     if setting.get(r, c) == CELL BUILDING \longrightarrow 1 tu
       A[r][c].reset() \longrightarrow 1 tu
       continue
     from above = null —----> 1 tu
     from left = null —----> 1 tu
     if r > 0 and A[r - 1][c].has value()
       from above = A[r - 1][c] —---> 1 tu
       if from above is step valid(go down)
          from above.add step(go down) -----> 1 tu
    if c > 0 and A[r][c - 1].has value() —----> 1 tu
       from left = A[r][c - 1] —----> 1 tu
       if from left.is step valid(go right) —----> 1 tu
          from left.add step(go right) -----> 1 tu
     if from above has value and from left has value
       if from above.total cranes() > from left.total cranes() —----> 1 tu
          A[r][c] = \text{from above} \longrightarrow 1 \text{ tu}
       else
          A[r][c] = from left \longrightarrow 1 tu
     else if from left has value
       A[r][c] = from left \longrightarrow 1 tu
     else if from above has value
```

```
A[r][c] = \text{from\_above} -----> 1 \text{ tu}
\text{endfor}
\text{endfor}
\text{best} = A[0][0] -----> 1 \text{ tu}
\text{for } r = 0 \text{ to setting.rows()} - 1
\text{for } c = 0 \text{ to setting.columns()} - 1
\text{if } A[r][c].\text{has\_value()} \text{ and } A[r][c].\text{total\_cranes()} > \text{best.total\_cranes()} -----> 1 \text{ tu or}
r^*c
\text{best} = A[r][c] -----> 1 \text{ tu}
\text{endfor}
\text{endfor}
\text{assert best.has\_value()}
\text{return best}
r - 1c - 1
\sum_{r=0}^{\infty} \sum_{c=0}^{\infty} 17 + 1 + \sum_{r=0}^{\infty} \sum_{c=0}^{\infty} 2
```

Since the Dynamic Algorithm appears to have a r\*c, it would indicate that it has a polynomial time complexity which would mean it is faster than the exhaustive algorithm.

## Questions

- 1. Is there a noticeable difference in the performance of the two algorithms? Which is faster, and by how much? Does this surprise you?
  - a. After running the code, it seems like that Dynamic Algorithm is the faster one of the two algorithms by a lot. This does not surprise me because of the nature of how Exhaustive search has to explore all the potential paths to decide which is the best. Exhaustive searches typically have exponential time complexities while Dynamic Algorithms can have polynomial time complexities.
- 2. Are your empirical analyses consistent with your mathematical analyses? Justify your answer.
  - a. An empirical analysis should be consistent with a mathematical analysis, depending on a problem's time complexities. Since an expected time complexity for the Dynamic is a polynomial and exhaustive is exponential, if they were to change, then it would not be consistent with a mathematical analysis.
- 3. Is this evidence consistent or inconsistent with hypothesis 1? Justify your answer.
  - a. The hypothesis: "This experiment will test the following hypothesis:

    Polynomial-time dynamic programming algorithms are more efficient than
    exponential-time exhaustive search algorithms that solve the same problem."

- b. Based on this hypothesis, the evidence would be considered consistent. This is because when running the code, while both are able to return 4 cranes, the elapsed time for exhaustive optimization is around 0.122773 and dynamic is around 0.0003671. While these numbers do change, the dynamic algorithm demonstrates a faster time every time.
- 4. Is this evidence consistent or inconsistent with hypothesis 2? Justify your answer.
  - a. Could not find hypothesis 2.