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## 89. Gray Code

The gray code is a binary numeral system where two successive values differ in only one bit.

Given a non-negative integer  $n$  representing the total number of bits in the code, print the sequence of gray code. A gray code sequence must begin with 0.

For example, given  $n = 2$ , return  $[0, 1, 3, 2]$ . Its gray code sequence is:

```
00 - 0
01 - 1
11 - 3
10 - 2
```

### Note:

For a given  $n$ , a gray code sequence is not uniquely defined.

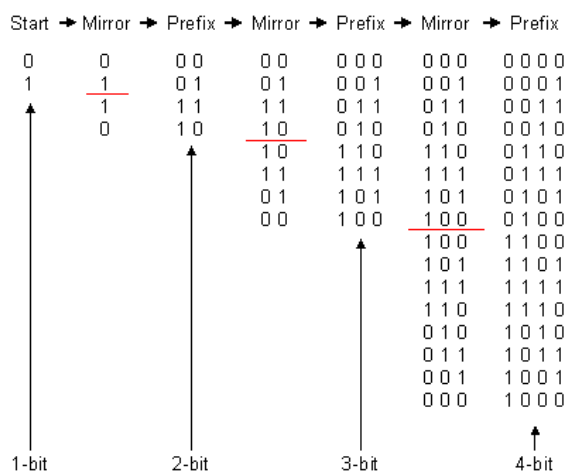
For example,  $[0, 2, 3, 1]$  is also a valid gray code sequence according to the above definition.

For now, the judge is able to judge based on one instance of gray code sequence. Sorry about that.

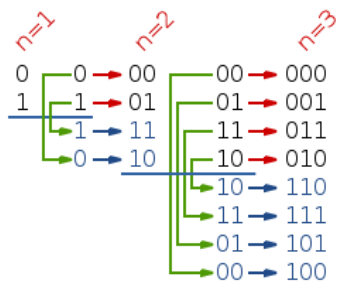
**Tags: Backtracking**

```
class Solution {
public:
    vector<int> grayCode(int n) {
        vector<int> re(1,0);
        int max_value = ((unsigned)0 - 1) >> (sizeof(int) * 8 - n);
        for (int i = 1; i <= max_value; i++) {
            re.push_back(i^(i >> 1));
        }
        return re;
    }
};
```

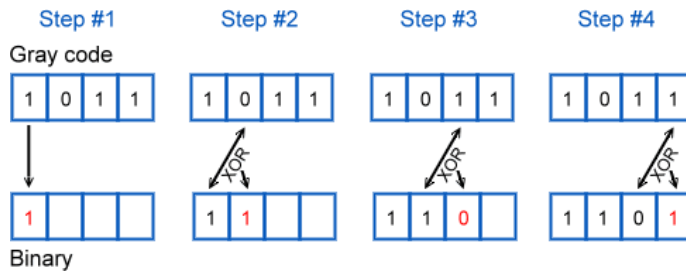
思路:



- 除了最高位，格雷码的位元完全上下对称。如果我们生成了 $n$ 位的格雷码，可将其顺序反转，然后在两个列表的前面分别补上0和1，就可生成 $n+1$ 为的格雷码。(可用bitset实现)



1. 自然数 $n$ 对应的二进制数为 $B$ ,其格雷码为  $G(i) = B(i+1) \text{ xor } B(i)$  。



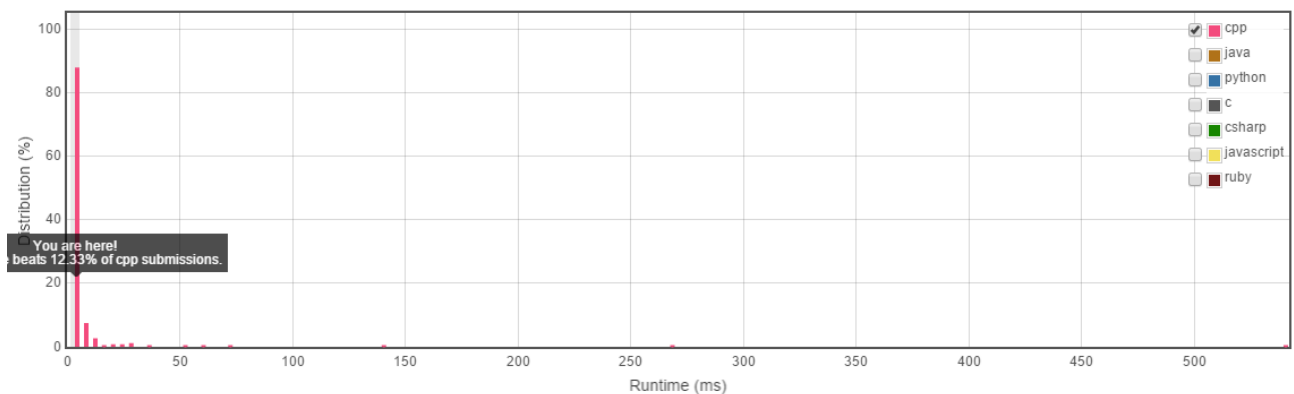
12 / 12 test cases passed.

Runtime: 4 ms

Status: Accepted

Submitted: 8 minutes ago

Accepted Solutions Runtime Distribution



## 121. Best Time to Buy and Sell Stock

Say you have an array for which the  $i$ th element is the price of a given stock on day  $i$ .

If you were only permitted to complete at most one transaction (ie, buy one and sell one share of the stock), design an algorithm to find the maximum profit.

**Tags:** Array, Dynamic Programming

```
class Solution {
public:
    int maxProfit(vector<int>& prices) {
        if (prices.size() <= 1) return 0;
        int low = prices[0];
        int re = 0;
        for (int i = 1; i < prices.size(); i++) {
            low = min(low, prices[i]);
            re = max(re, prices[i] - low);
        }
        return re;
    }
};
```

思路:

设  $dp[i]$  为到第  $i$  天为止的最大收益，则显然有  $dp[i+1] = \max(dp[i], \text{prices}[i+1] - \text{lowest})$ 。

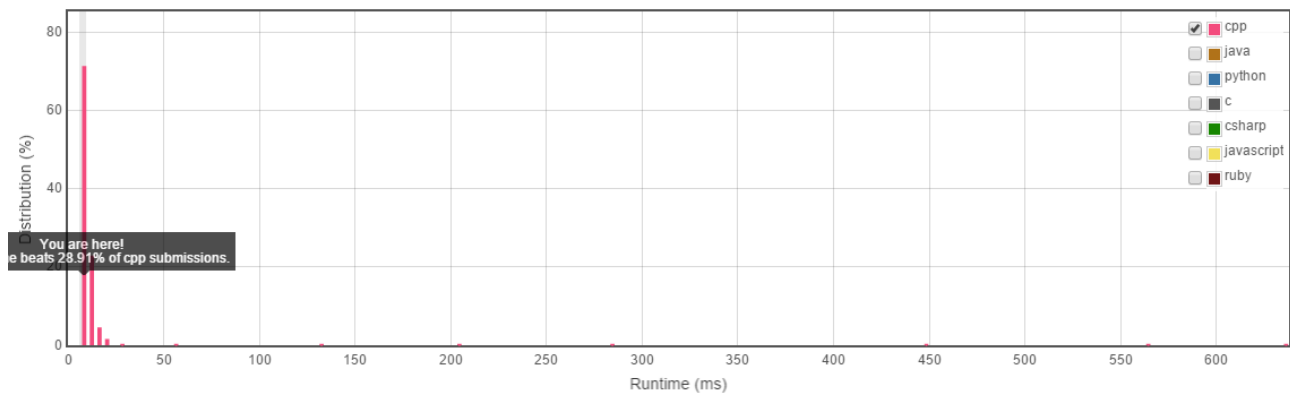
198 / 198 test cases passed.

Runtime: 8 ms

Status: Accepted

Submitted: 21 hours, 51 minutes ago

Accepted Solutions Runtime Distribution



## 123. Best Time to Buy and Sell Stock III

Say you have an array for which the  $i$ th element is the price of a given stock on day  $i$ .

Design an algorithm to find the maximum profit. You may complete at most two transactions.

**Note:** You may not engage in multiple transactions at the same time (ie, you must sell the stock before you buy again).

**Tags:** Array, Dynamic Programming

```
class Solution {
public:
    int maxProfit(vector<int>& prices) {
        if (prices.size() <= 1) return 0;
        int i;
        int re = 0;
        int *forward_dp = new int[prices.size()];
        int *back_dp = new int[prices.size()];
        forward_dp[0] = 0;
        back_dp[prices.size() - 1] = 0;

        int low = prices.front();
        for (i = 1; i < prices.size(); i++) {
            low = min(low, prices[i]);
            forward_dp[i] = max(forward_dp[i - 1], prices[i] - low);
        }

        int high = prices.back();
        for (i = prices.size() - 2; i >= 0; i--) {
            high = max(high, prices[i]);
            back_dp[i] = max(back_dp[i + 1], high - prices[i]);
        }

        for (i = 0; i < prices.size(); i++) {
            re = max(re, forward_dp[i] + back_dp[i]);
        }
        return re;
    }
};
```

思路：

其最多允许进行两次交易，还包括进行一次交易或者不交易，后者也可以看作进行两次交易，如在当天进行一次或两次买入和卖出。

若发生了两次交易，则第二次买入需在第一次卖出以后进行，以 `forward_dp[i]` 表示第0天到第i天的最大收益，`back_dp[i]` 表示第i天到最后一天的最大收益，显然有 `dp[i] = forward_dp[i] + back_dp[i]`，其中 `dp[i]` 代表总收益。

由121题可知，`forward_dp[i+1] = max(forward_dp[i], prices[i+1] - lowest)`，则 `back_dp[i] = max(back_dp[i+1], highest - prices[i])`。二者得到以后，即可从前向后扫描，计算 `dp[i]`，找到最大值即可。

198 / 198 test cases passed.

Runtime: 12 ms

Status: Accepted

Submitted: 21 hours, 46 minutes ago

#### Accepted Solutions Runtime Distribution

