

Lecture 14:

Compiler Correctness

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Compilers

This lecture

- How can we ensure that a compiler is **correct**?
- Compiler correctness is important because compilers are **prevalent** and **trusted**.

Coursework assessment

- List of language features to support is on Github.
- Compilers will be assessed solely on whether the generated assembly produces the correct output.
- Compilers will only be tested on valid inputs.
- I don't expect your compiler to work every input I've given you; this list is meant as an exhaustive upper bound for everyone!

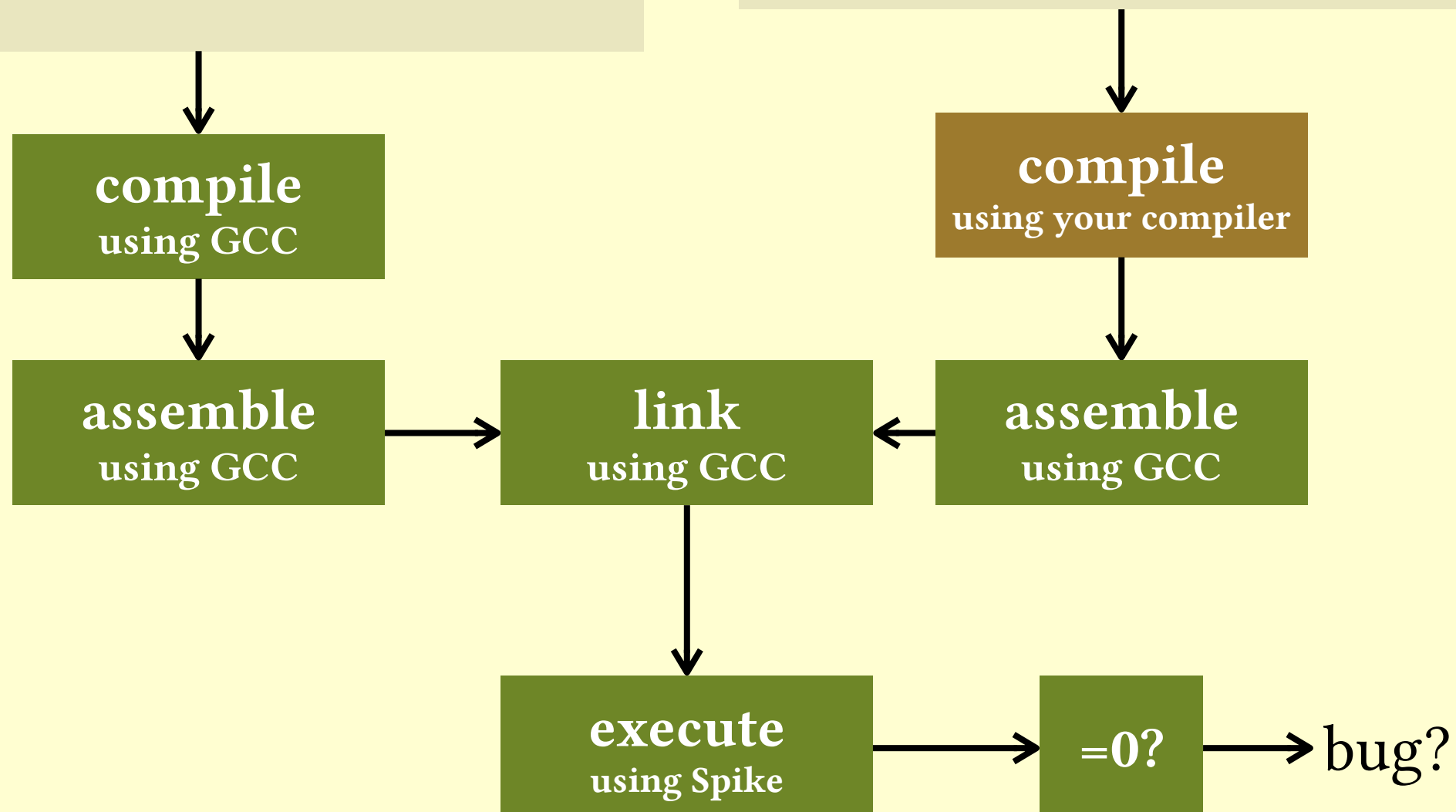
Coursework assessment

test_ADD0_driver.c

```
int f(int x, int y);  
int main() {  
    return !(40==f(30,10));  
}
```

test_ADD0.c

```
int f(int a, int b) {  
    return a+b;  
}
```



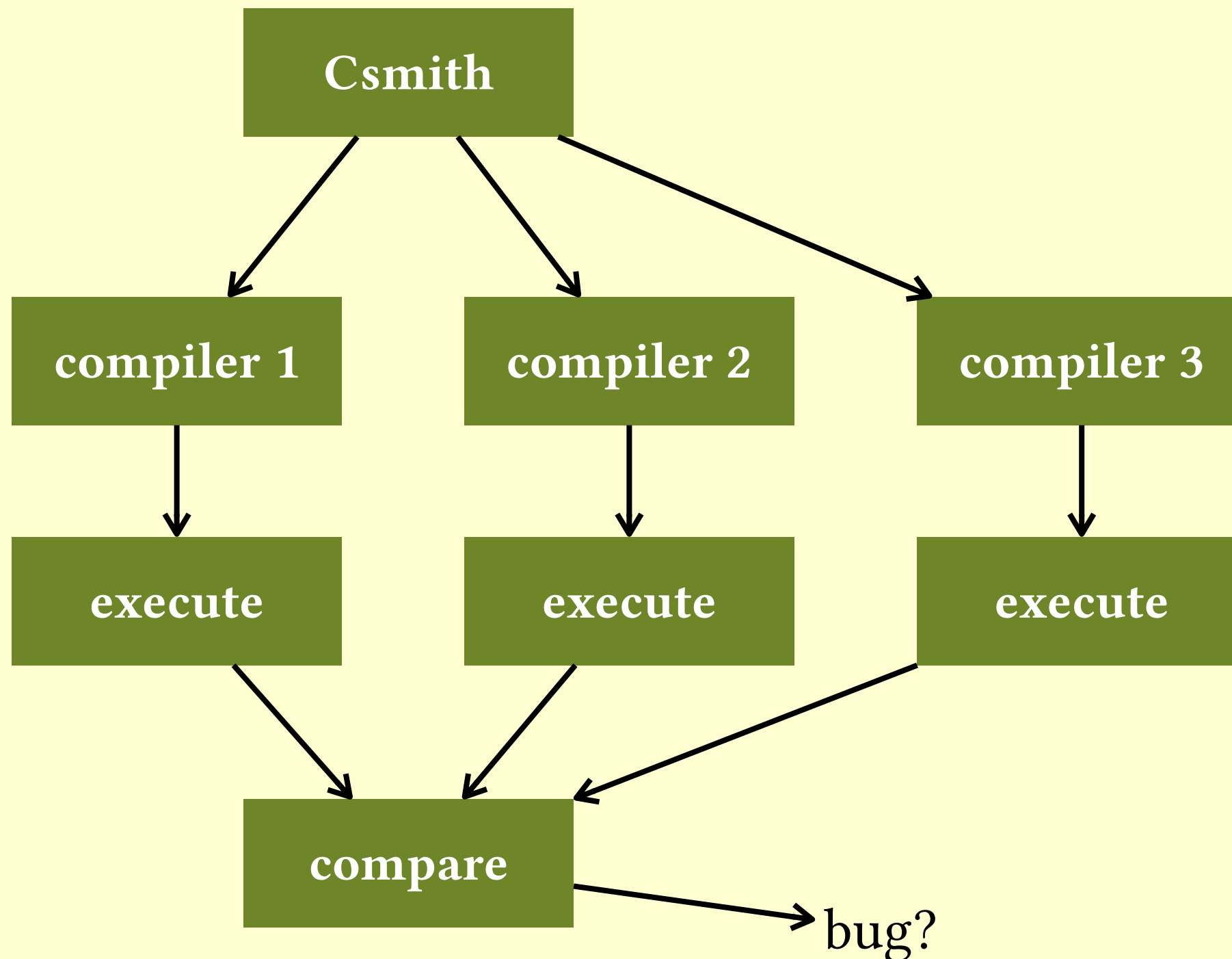
Random testing

- **Csmith** is a tool from the University of Utah that generates random C programs.

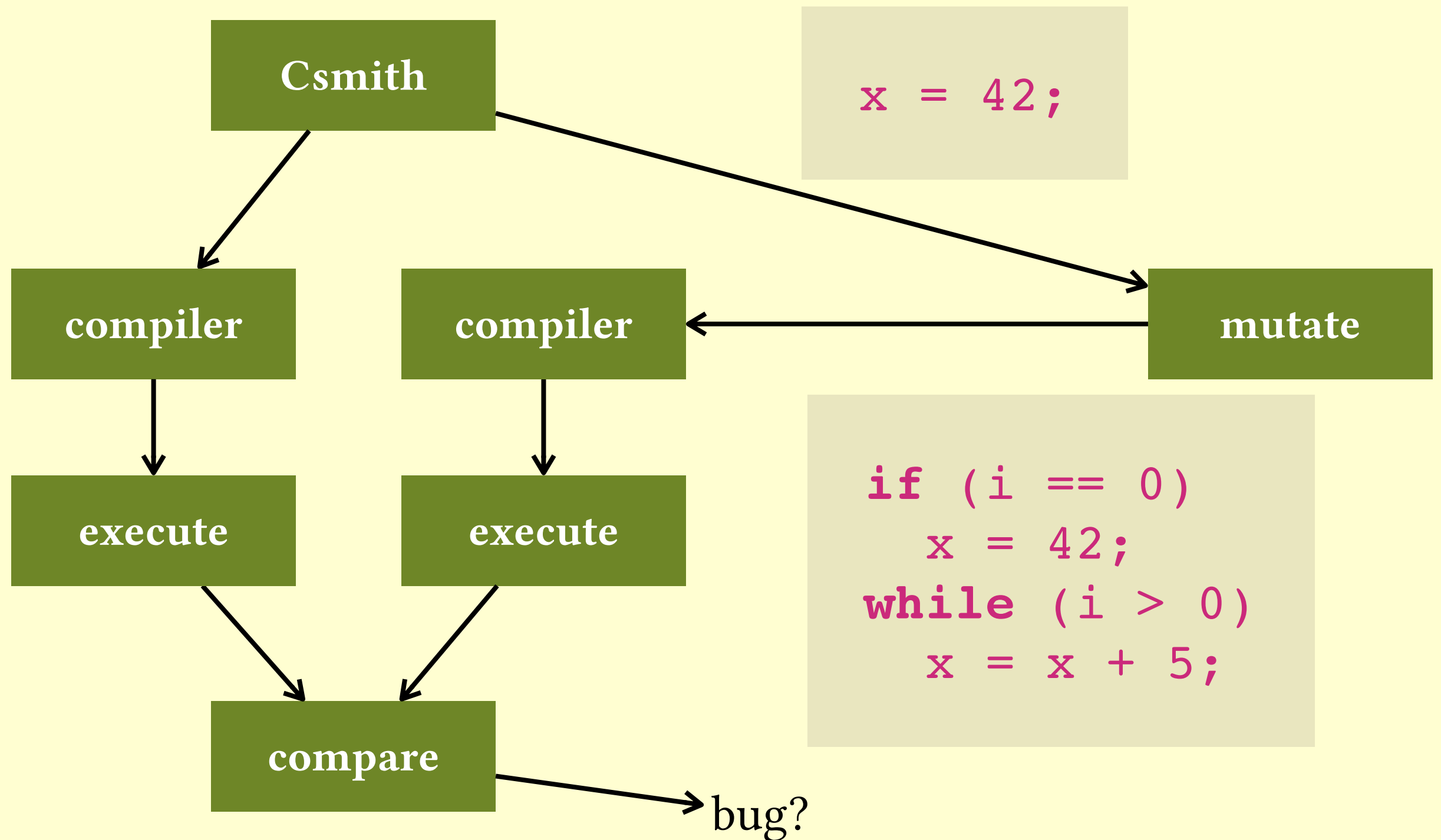


- **Question.** How do we know if the random program has been compiled correctly?
- **Solution 1.** Random differential testing.
- **Solution 2.** EMI testing.

Random differential testing



EMI testing



Problem: non-determinism

- C99 has 52 **unspecified** behaviours (e.g. the order in which operands are evaluated) and 114 **implementation-defined** behaviours (e.g. the size of an **int**). Generated programs must produce output that does not depend on how a compiler implements these behaviours.
- C99 also has 191 **undefined** behaviours (e.g. **x = *NULL**, or overflowing a **signed int**). Generated programs must not contain undefined behaviours at all.

```
foo(*p);
```



```
if (p!=NULL)  
    foo(*p);
```



Features of Csmith

- **Included:** global variables, local variables, if-then-else statements, for loops, break/continue statements, goto statements, function calls, signed and unsigned integers of various sizes, arithmetic and logical operations, structs, arrays, pointers.
- **Excluded:** strings, malloc, floating point types, unions, recursion, function pointers.

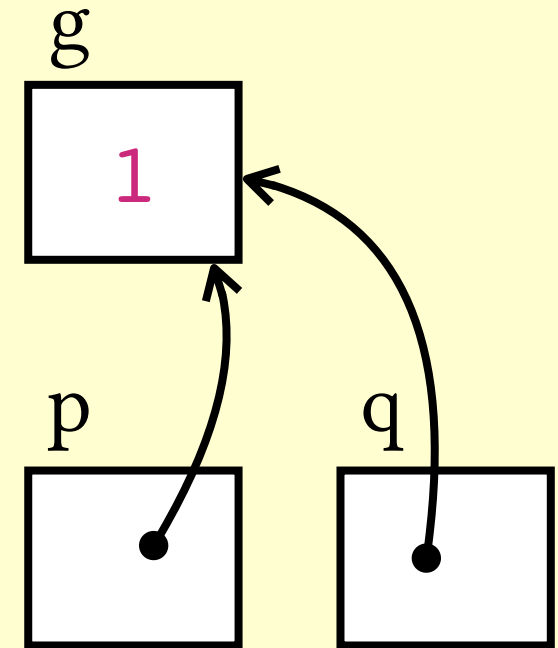
Example bug

```
int foo () {  
    signed char x = 1;  
    unsigned char y = 255;  
    return x > y;  
}
```

- When operands have different types, both are promoted to **int**. So this function should return **0**.
- However, when compiled with an old version of GCC, this function returns **1**.

Another example bug

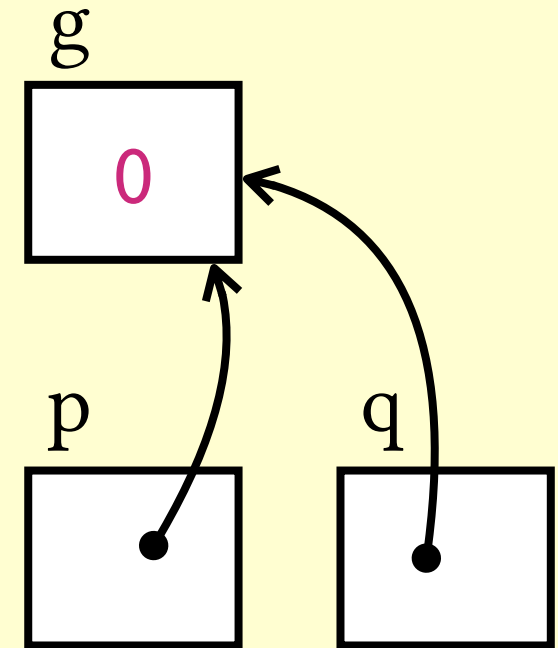
```
int g[1];  
int *p = &g[0];  
int *q = &g[0];  
  
int foo () {  
    g[0] = 1;  
    *p = 0;  
    *p = *q;  
    return g[0];  
}
```



Another example bug

```
int g[1];
int *p = &g[0];
int *q = &g[0];



int foo () {
    g[0] = 1;
    *p = 0;
    *p = *q;
    return g[0];
}
```



- Should return 0, but an old version of GCC did faulty dead-code elimination (removing `*p = 0`) because it didn't realise that `p` and `q` are aliases, and thus returned 1.

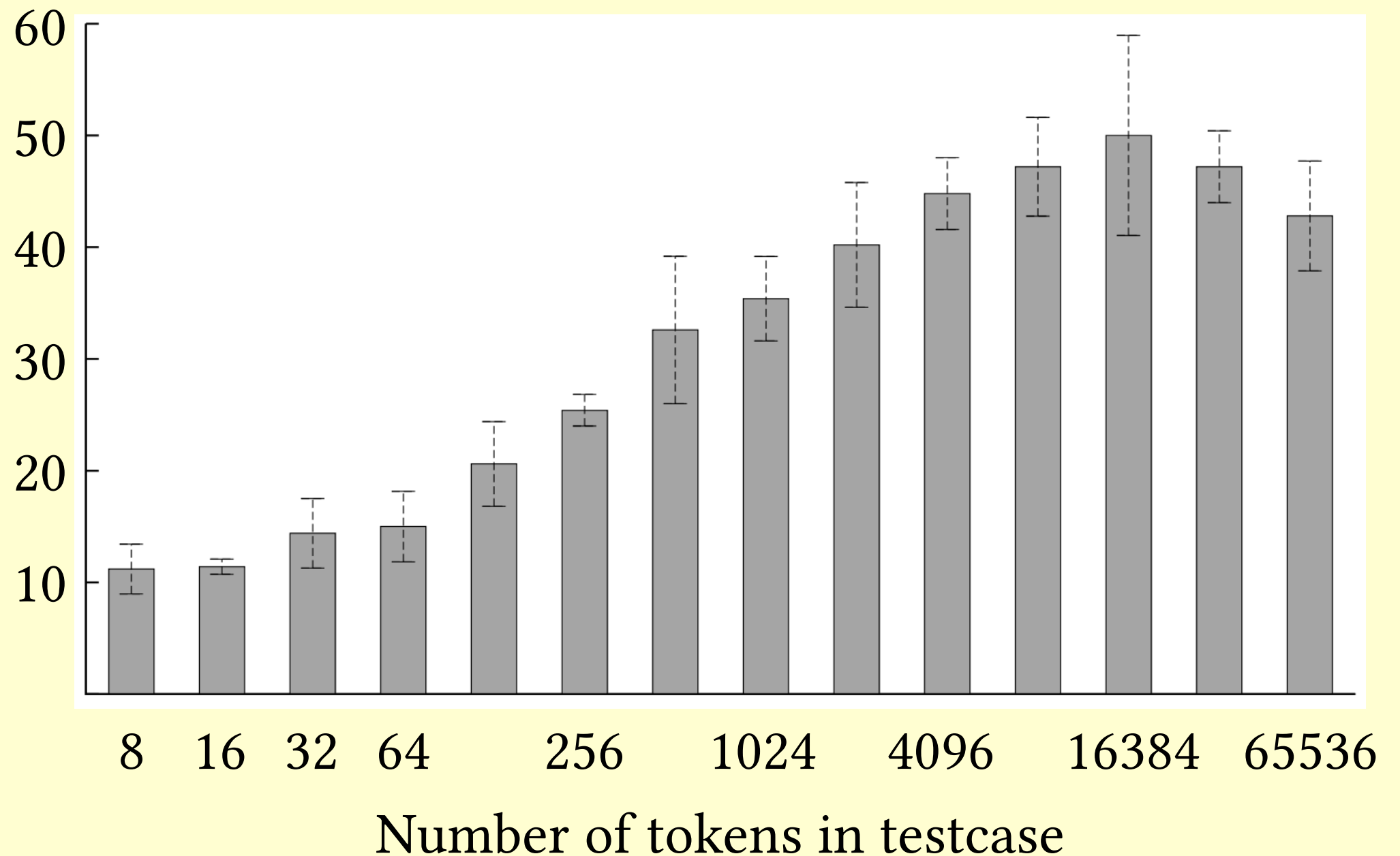
Bugs found



		
crashes	50	136
wrong code bugs	29	66

How large should the random programs be?

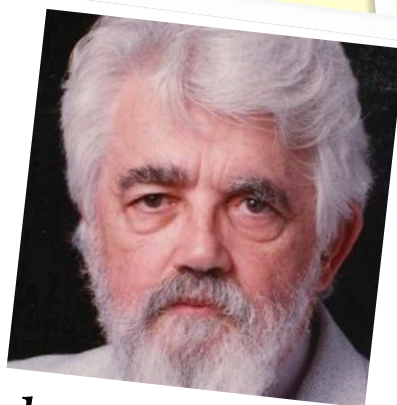
Number of
bugs found
in 24 hours



Can we do better?

Verified compilation

- Can we prove (mathematically) that our compiler will never generate wrong code?



John McCarthy
1927–2011

CORRECTNESS OF A COMPILER FOR ARITHMETIC EXPRESSIONS*

JOHN McCARTHY and JAMES PAINTER

1967

1 Introduction

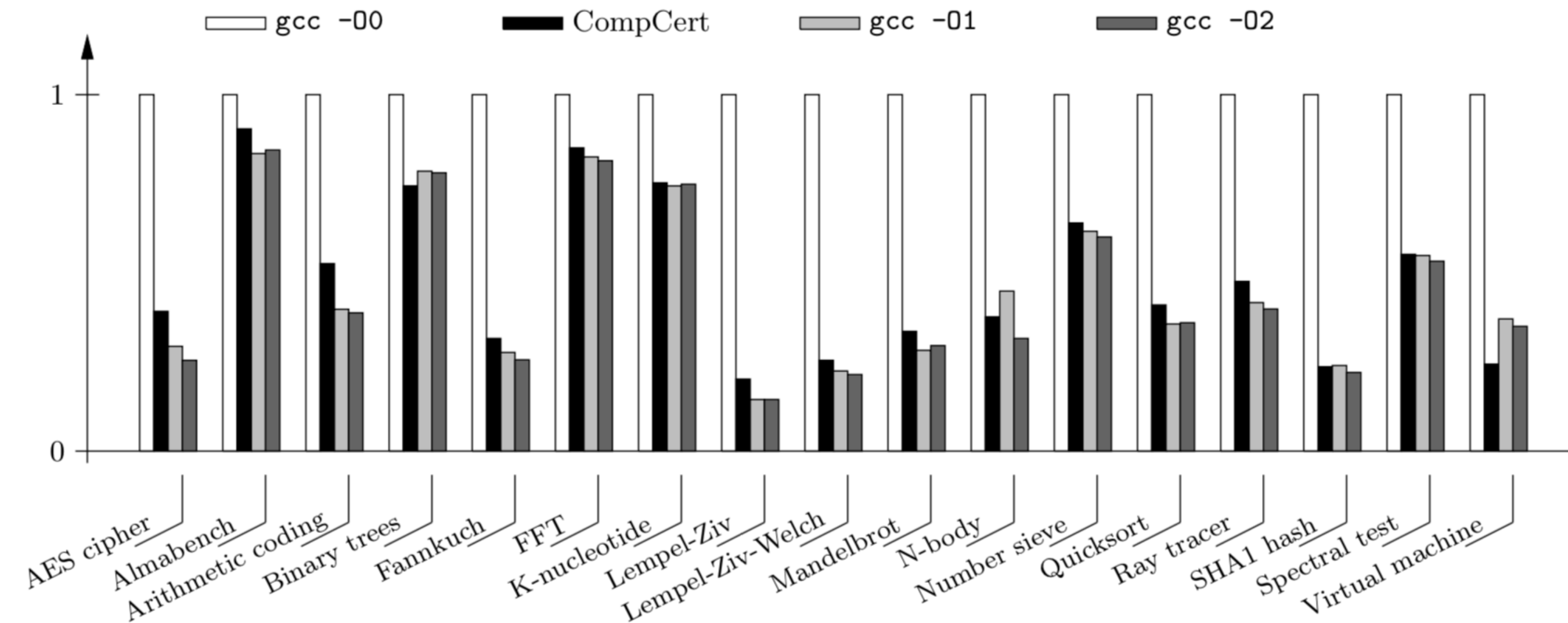
This paper contains a proof of the correctness of a simple compiling algorithm for compiling arithmetic expressions into machine language.
The definition of correctness, the formalism used to express the description of the compiler, and the methods of proof are

CompCert

- A verified C compiler
- Development started around 2005; first release was in 2008.
- Programmed (and proved correct) in Coq.
- Performance is comparable to `gcc -O1`.

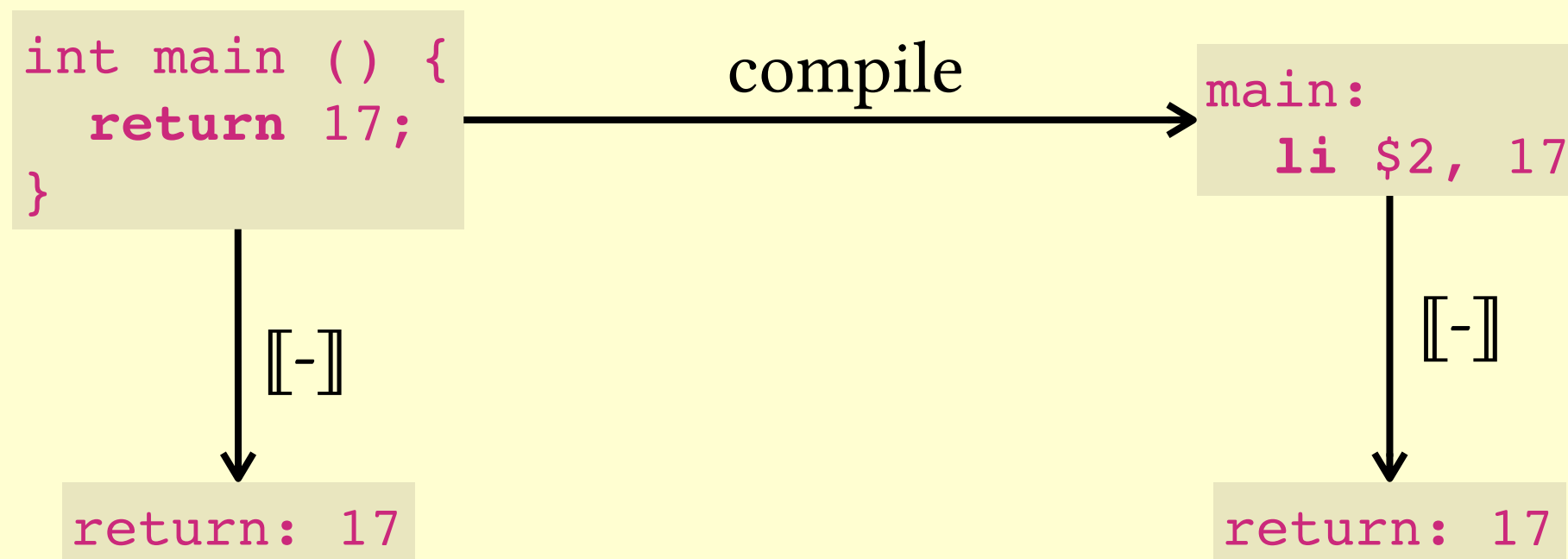


CompCert performance



CompCert correctness

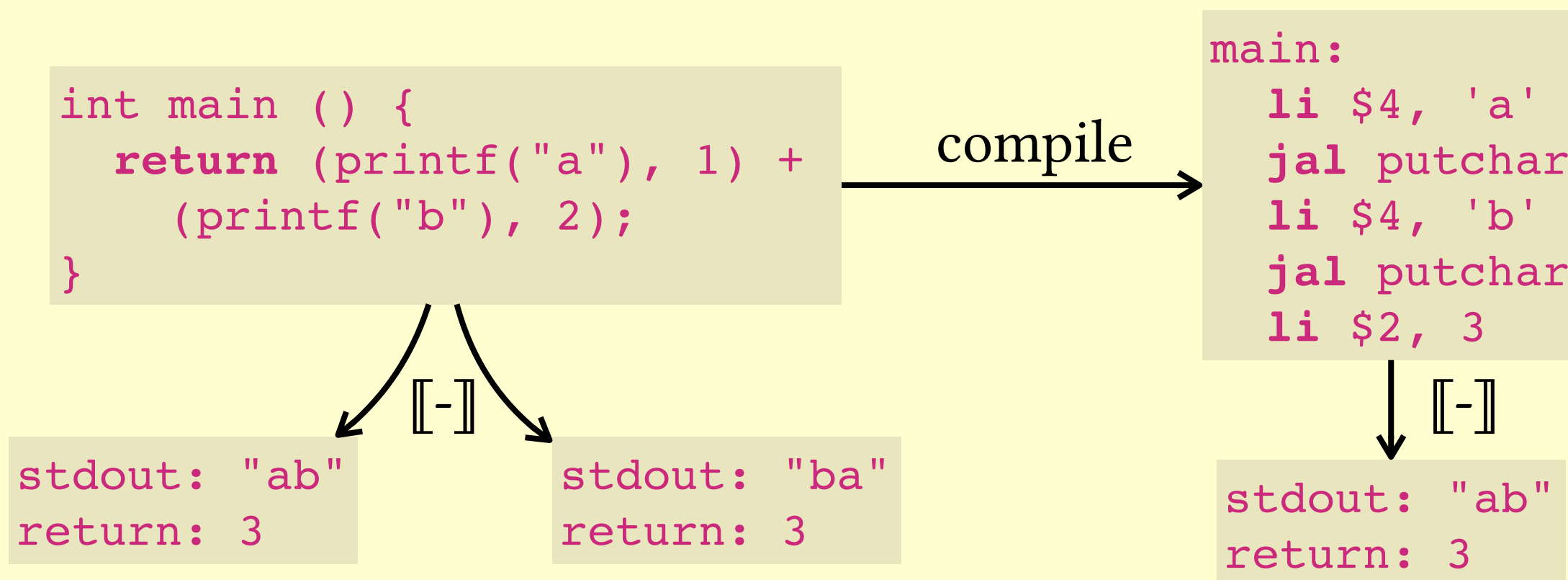
- **Main theorem:**
For any C program P,
 $\llbracket P \rrbracket = \llbracket \text{compile}(P) \rrbracket$.



CompCert correctness

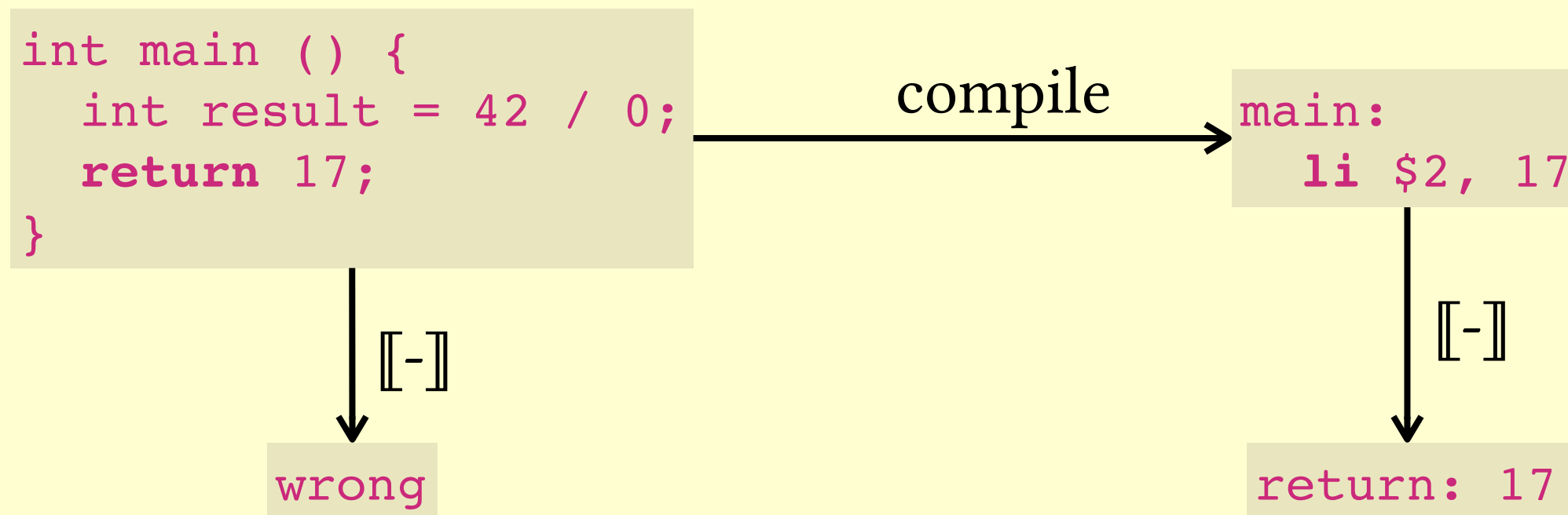
- **Main theorem:**

For any C program P ,
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CompCert correctness

- **Main theorem:**
For any C_{light} program P ,
 $\llbracket P \rrbracket = \llbracket \text{compile}(P) \rrbracket$.



CompCert correctness

- **Main theorem:**
For any C_{light} program P ,
if $\llbracket P \rrbracket \neq \text{wrong}$ then
 $\llbracket P \rrbracket = \llbracket \text{compile}(P) \rrbracket$.

```
int main () {  
    return 1 * 1 * ... * 1;  
}
```

compile

Compilation error!

CompCert correctness

- **Main theorem:**



For any C_{light} program P ,

if $\llbracket P \rrbracket \neq \text{wrong}$ and the compiler finishes then
 $\llbracket P \rrbracket = \llbracket \text{compile}(P) \rrbracket$.

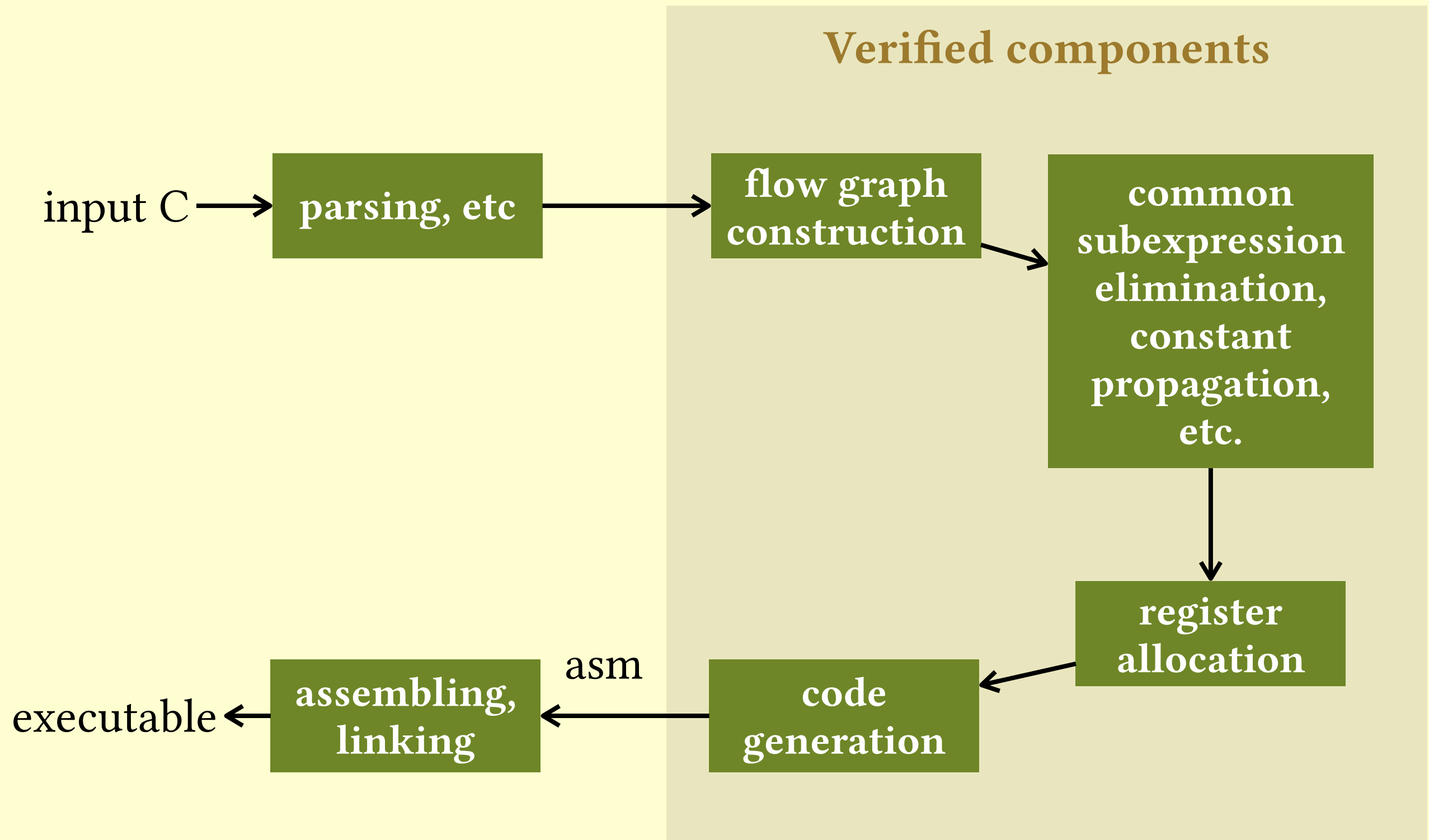


<https://compcert.org>

CompCert vs. Csmith

			CompCert
crashes	50	136	~5
wrong code bugs	29	66	~8

CompCert structure



Summary

- Compiler correctness is crucial because compilers are **prevalent** and **trusted**.
- Csmith generates random, deterministic C programs, which can be used for **differential testing** or **EMI testing** of compilers.
- CompCert is an optimising C compiler that is **proven** correct.