初等数论

吴伊涛

Lecture 9. Congruence

,	_			
(1	1	1)
	10	1	1	/
/	0		/	

We already use the fails "odd + odd = even, odd x odd = odd..."
many limes. We can summarize these fails as the following:

	l odd	even			odd	even
odd	even	odd		odd	odd	even
even	odd	even	westabilg count	even	even	even

These seems really elementary, can we get anything interesting faits from these?

Let's see ...

First, let's generalize the above table:

Clearly, $\forall n \in \mathbb{R}$, the remainder of n divides by 3 should be one of 0, 1, 2. i.e. n is of 3k, 3k+1, 3k+2 type integer.

Thus we have the following table:

	I TEM (telet) " Stembling to			Theorem	By Fry Somental Theorem				
J	3k	3k+1	3k+2		3k	3k+1	3k+2		
3k		6k+1	6k+2	3k	9k²	9k2+3k	9k2+6k		
	bk+1	6k+2	6k+3	3k+1	9k²+3k	9k+6k+1	9k2+9k+2		
	bk+2	6k+3	6k+4	3k+2	9k+6k	9k2+9k+2	9k2+12k+4		

Now, if we just record the remainder of the number divided by 3, the above two tables turns like:

		, x xxxx	"uaria =			X	
1	0	woll of 3	2	J. 35 M.A.	0	Nel con	2
0	0	1	2	0	0	0	0
1	1	2	0	1	0	1	2
2	2	A 90 0	1	2	0	2	1100
					1		

We can find, from the above multiplication lable, that "a 3k+2 type integer n, if we fattorize n as $n=a\cdot b$, then a or b is also a 3k+2 type integer". Thus we get the following facts:

Positive

Prop. A 3k+2 type integer always has a 3k+2 type prime factor.

Proof: Clearly, all primes except 3 are of 3k+1 type or 3k+2 type.

If P., Pz, ... Pe are all of 3k+1-lype integer, then by the muttiplication table above, p. p. p. is also a 3k+1 lype integer. A contradiction. So, at least one of $p_1, p_2, \dots p_L$ is a 3k+2 type integer.

We can deduce the following consequence:

Theorem. There are infinitely many 3k+2 type primes.

Proof: Assume not, let $p_1 = 2$, $p_2 = 5$, ..., p_e are all 3k+2 type primes.

We construct $N = 3p_2p_3 \cdots p_l + 2$.

As N is a 3k+2 type positive integer, by the above Prop, N has a 3k+2 type prime factor 9. So 9 is one of 9, 9, 9, 9.

If $g = p_1 = 2$, then $g \mid N-2 = 3 \cdot p_2 \cdot p_3 \cdots p_L$, which is an odd number, a contradiction.

If $g = p_i$, $i \in \{2, 3, \dots, l\}$, then $g \mid 3p_2 \cdots p_l$, thus $g \mid N-3p_2 \cdots p_l = 2$, also a contradiction.

So There are infinitely many 3k+2 type primes.

Exercise:

I. Prove that there are infinitely many 4k+3 type primes. (Hint: Prove that a 4k+3 type positive integer always has a 4k+3 type prime factor).

Example Solve the Diophantine equation

$$x^{2022} + 115x^2 + 2333 = 0$$

Solution:

Assume x_0 is an integral solution of the Diophantine equation $x^{2022} + 115x^2 + 2333 = 0$.

If x_0 is an odd number, then x_0^{2022} , $115x_0^2$ both are odd numbers, thus $x_0^{2022} + 115x_0^2 + 2333$ is a sum of three odd numbers, which is still an odd number, thus unequal to 0. A contradiction.

If x_0 is an even number, then x_0^{202} , $115x_0^2$ both are even numbers, thus $x_0^{2022} + 115x_0^2 + 2333$ is a sum of an odd number and two even numbers, which is an odd number, thus unequal to 0. A contradiction.

Thus the Diophantine equation $x^{2022} + 115x^2 + 2333 = 0$ is unsolvable.

· Congruence (in Z).

(23)

Dog: Let a, b E Z, m E Z⁺.

We say that $\underline{\alpha}$ is congruent to \underline{b} modulo \underline{m} if $\underline{m} | \underline{b} - \underline{a}$. We denote it by $\underline{a} \equiv \underline{b} \pmod{m}$.

F.g. 3=18 (mod S), 19=11 (mod 2), 2020=4 (mod 9).

Prop. Congruence modulo m is an equivalence relation on the set of integers &, i.e.

- 1) $\alpha \equiv \alpha \pmod{m}$, $\forall \alpha \in \mathbb{Z}$,
- 2) $a \equiv b \pmod{m}$ implies $b \equiv a \pmod{m}$, $\forall a, b \in \mathbb{Z}$,
- 3) $\alpha \equiv b \pmod{m}$ and $b \equiv C \pmod{m}$ implies $\alpha \equiv C \pmod{m}$ $\forall a, b, c \in \mathbb{R}$.

Proof:

- 1) a-a=0 and $m \mid 0$.
- 2) If m | b-a, then m | a-b, i.e. $b \equiv a \pmod{m}$
- 3) If m|b-a, m|c-b. Then m|c-a=(c-b)+(b-a), i.e. $a\equiv C \pmod{m}$.

Prop. 1) If $a \equiv b \pmod{m}$, $c \equiv d \pmod{m}$, then a+c = b+d (mod m) + g am gad a tol pa

- A-d 2) If $A \equiv b \pmod{m}$, $C \equiv d \pmod{m}$ then $a-c = b-d \pmod{m}$.
 - 3) If $a \equiv b \pmod{m}$, $t \in A$, ud t = d = bthen $ta \equiv tb \pmod{m}$. (2 bom) $81 \equiv 8$
 - 4) If $a \equiv b \pmod{m}$, $C \equiv d \pmod{m}$ Then $a \cdot c \equiv b \cdot d \pmod{m}$.

Proof:

1) 2) 3) leave to the reader. (1) ham) 0 = 0 (1)

4). If $\Delta \equiv b \pmod{m}$, $C \equiv d \pmod{m}$.

Thus $bd-ac = (b-a) \cdot d + a \cdot (d-c)$ is a multiple of m,

: ac = bd (mod m).

Cor. Let f(x) ∈ Z[x] be a polynomial with integral coefficients. If $a \equiv b \pmod{m}$, then $f(a) \equiv f(b) \pmod{m}$.

Given f(X) E Z[X], Then Y KEZ, Cor. $f(k) = f(0) \pmod{2}$ or $f(k) = f(1) \pmod{2}$

Example Solve the Diophantine equation

$$x^{2022} + 115x^2 + 2333 = 0$$

Solution:

Denote the Diophantine equation $x^{2022} + 115x^2 + 2333 = 0$ by f(x) = 0.

Assume x_0 is an integral solution of the Diophantine equation f(x) = 0.

Clearly, $x_0 \equiv 0$ or $1 \pmod{2}$.

If $x_0 \equiv 0 \pmod{2}$, then $f(x_0) \equiv f(0) \pmod{2}$, so $f(x_0) \equiv 2333 \equiv 1 \pmod{2}$, a contradiction to $f(x_0) \equiv 0 \pmod{2}$.

If $x_0 \equiv 1 \pmod{2}$, then $f(x_0) \equiv f(1) \pmod{2}$, so $f(x_0) \equiv 1 + 115 + 2333 \equiv 1 \pmod{2}$, a contradiction to $f(x_0) \equiv 0 \pmod{2}$.

Thus the Diophantine equation $x^{2022} + 115x^2 + 2333 = 0$ is unsolvable.

Remark

- In general, we can prove: Given an integral coefficients polynomial f(x), if the constant term of f(x) (i.e. f(0)) and the sum of all coefficients of f(x) (i.e. f(1)) both are odd numbers, then f(x) = 0 is unsolvable.
- Try to solve the Diophantine equation

$$x^{2022} - 2022x^{2020} + 1234 = 0$$

Example Solve the Diophantine equation

$$x^2 + y^2 = 3z^2$$

Solution:

Assume (a, b, c) is a primitive integer solution of the Diophantine equation $x^2 + y^2 = 3z^2$, i.e. gcd(a, b, c) = 1.

Since $a^2 + b^2 \equiv 3c^2 \equiv 0 \pmod{3}$, it then follows that $a^2 \equiv 0 \pmod{3}$ and $b^2 \equiv 0 \pmod{3}$, which leads to 3|a| and 3|b|.

Apply a = 3u, b = 3v to the equation $a^2 + b^2 = 3c^2$, it then follows that 3|c.

So 3|gcd(a, b, c) = 1, a contradiction.

So the Diophantine equation $x^2 + y^2 = 3z^2$ has no primitive integer solutions. The only possible solution is (0,0,0).

By checking, we find (0,0,0) is indeed a solution. Thus the Diophantine equation $x^2 + y^2 = 3z^2$ has a unique solution (0,0,0).

. Repeated squaring method

(mod 2000)

Given a, b, m & Zt, we want to calculate the remainder of ab divided by m.

We introduce a notation first:

Def: $a \pmod{m} := \text{the remainder of } a \text{ divided by } m$.

E.g. 11 (mod 8) = 3, 30 (mod 11) = 8, 2020 (mod 9) = 4 8 (mod S) = 4.

E.g. Calculate 1148 (mod 2020).

Clearly, it is cumbersome to calculate 1148 out and then divide it by 2020. We introduce the following "Repeated squaring method"

Sol:

 $11 \pmod{2020} = 11, \quad 11^2 \pmod{2020} = 11^2 = 121$ 114 (mod 2020) = 1212 (mod 2020) = 501 118 (mod 2020) = 5012 (mod 2020) = 521 1116 (mod 2020) = 5212 (mod 2020) = 761 1132 (mod 2020) = 7612 (mod 2020) = 1401.

Now, 48 = 32+16.

 $(mod 2020) = 11^{32} \cdot 11^{16} (mod 2020) = (13^{32} (mod 2020) \cdot 11^{16} (mod 2020))$ = 1401 × 761 (mod 2020) = 1621.

The above method is called "repeated squaring method". The (128) principle is:

 $ab \pmod{m} = (a(mod m) \cdot b(mod m)) \pmod{m}$

which clearly from $a \equiv x \pmod{b}$, $b \equiv y \pmod{m} \Longrightarrow ab \equiv xy \pmod{m}$ (page 126, 4)).

E.g. 119 × 748 (mod 10) = (119 (mod 10) × 748 (mod 10) (mod 10) hom) 060= 9x8 (mod 10) = 2.

Now, let's state the "repeated squaring method"

Aim: Calculate ab (mod m)

Step 1: Convert b to binary representation:

doublement b= (Ck Ck-195). C, Col210+ att suborter all about

Step 2: Calculate $\alpha^2 \pmod{m}$, $i=0,1,2,\cdots k$ Note that $a^{2^{j+1}} \pmod{m} = (a^{2^{j}} \pmod{m})^2 \pmod{m}$ So, just "repealed square".

Multiply those a (mod m) where $C_i = 1$. (Actually more delicate if m >> 1).

In fail, $b = C_R \cdot 2^k + C_{R-1} \cdot 2^{k-1} + \cdots + C_1 \cdot 2 + C_0$.

 $(acachem) = (a^{2k})^{Ck} \pmod{m} \cdot (a^{2k})^{Ck-1} \pmod{m} \cdot (a^{2k})^{Ck-1} \pmod{m}$ (mod m)

(Note that if $C_{ni}=0$, $(\alpha^{2i})^{C_i}$ (mod m)=1. if Ci = 1, Q2i Ci (mod m) = Q2i (mod m). E.g. Calculate 999 (mod 1001).

Sol: Since $999 \equiv -2 \pmod{1001}$, $999^{1000} \equiv (-2)^{1000} \equiv 2^{1000} \pmod{1001}$ $\therefore 999^{1000} \pmod{1001} = 2^{1000} \pmod{1001}$.

Now, Convert 1000 lo binary representation, we get: $|000 = (|||||0||000)_2 = 2^3 + 2^5 + 2^6 + 2^7 + 2^8 + 2^9$ Then, let's calculate 2 (mod [001], i=0,1,2,...,9. $2^{2^{\circ}} \pmod{|00|} = 2$. $2^{2^{\circ}} \pmod{|00|} = 2^{2^{\circ}} \pmod{|00|} = 4$ $2^{2^{2}} \pmod{\lfloor 001 \rfloor} = 4^{2} \pmod{\lfloor 001 \rfloor} = 16, \quad 2^{2^{3}} \pmod{\lfloor 001 \rfloor} = 16^{2} \pmod{\lfloor 001 \rfloor} = 256$ 24 (mod [001) = 2562 (mod [001] = 65536 (mod [001] = 471. 25 (mod [001) = 4712 (mod [001) = 221841 (mod [001) = 620. 226 (mod [001) = 6202 (mod [001) = 384400 (mod [001) = 16 227 (mod [001) = 162 (mod [001] = 256 228 (mod 1001) = 2562 (mod 1001) = 471 29 (mod [001) = 4712 (mod [001) = 620.

Thus $2^{1000} (\text{mod } 1001) = 2^{2^3} \cdot 2^{2^5} \cdot 2^{2^6} \cdot 2^{2^7} \cdot 2^{2^8} \cdot 2^{2^9} (\text{mod } 1001)$ $= 256 \times 620 \times 16 \times 256 \times 471 \times 620 (\text{mod } 1001)$ = 562.

So 999 (mod 1001) = 562.

Rmk 1: We will introduce some simplifications of the above method. (Bo)

Anyway, the "repeated squaring method" shows that

Ob (mod m) is easy to get: roughly O(log2b) times of

two integers' (smaller than m) multiplications.

Rmk 2: Note that 2²ⁱ (mod 1001) is a cyclic sequence: 2, 4, 16, 256, 471, 620, 16, 256, 471, 620, We will explain these later.

Remark

• The Repeated Squaring Method is efficient.

Example

Alice is using her computer to calculate

 $1234567890987654321^{2021202220232024} (mod 31415926535897932626) \\$

If the computer use 1 second to calculate the product of two 20-digits number, and 1 second to calculate the remainder of a 40-digits number divided by a 20-digits number, then is it possible to get the result in 5 minutes? Give your reason and estimate a bound of time.

Solution

Let's use "repeated squaring method".

First, write out the binary presentation of b=2021202220232024, it's just divide b by 2 repeatedly, thus takes almost no time.

Now, suppose we already know $x^{2^{i}}(modm)$ (where x=1234567890987654321, m=31415926535897932626), which should be a 20-digit(or less) number, then we square it to get a 40-digit(or less) number and then divide the result by m, all these cost at most 2 seconds, we then get $x^{2^{i+1}}(modm)$, thus by induction, we could get $x^{2^{i}}(modm)$ for all $i \leq k$ in 2k seconds.

Now multiply all $x^{2^i}(modm)$ if the *i*-th digit in the binary presentation is 1 in the same way as above: multiply the first 2 numbers and then divide it by m, and then multiply the result with the 3rd number and then divide it by m.....

Thus the total time is at most 4k where $k = [log_2b] = 50$, so we can finish the computation in 5 minutes. In fact, an upper bound of the time is 200 seconds.

