

# 实验报告（Josephus 问题）

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时间：2023 年 9 月 25 日 上海

## 源码与结果

### 源码

见附录A

### 结果

n	Result
10	4
11	6
12	8
13	10
14	12
15	14
16	0
17	2
18	4
19	6
20	8

```
Insert number of players (n): 10
In the case of 10 players, the winner is 4
> ./testfile
Insert number of players (n): 11
In the case of 11 players, the winner is 6
> ./testfile
Insert number of players (n): 12
In the case of 12 players, the winner is 8
> ./testfile
Insert number of players (n): 13
In the case of 13 players, the winner is 10
> ./testfile
Insert number of players (n): 14
In the case of 14 players, the winner is 12
> ./testfile
Insert number of players (n): 15
In the case of 15 players, the winner is 14
> ./testfile
Insert number of players (n): 16
In the case of 16 players, the winner is 0
> ./testfile
Insert number of players (n): 17
In the case of 17 players, the winner is 2
> ./testfile
Insert number of players (n): 18
In the case of 18 players, the winner is 4
> ./testfile
Insert number of players (n): 19
In the case of 19 players, the winner is 6
> ./testfile
Insert number of players (n): 20
In the case of 20 players, the winner is 8
```

## 算法理解

### 定义

- 所有的玩家编号成一有序序列：  $L_n = \llbracket 0, n - 1 \rrbracket = [0, 1, \dots, n - 1]$
- 场上存活的玩家（无序）集合为：  $S$ ，定义  $S$  中最大、最小编号分别为  $\bar{s}, s$

- 在所有报数之后，最后的玩家编号为  $w$ ，即我们所求；定义函数  $f: \mathbb{R} \rightarrow \mathbb{R}$ ，输入总人数  $n$ ，输出最终玩家：

$$f: n \mapsto w$$

- 在一次报数后，“退出”的人编号为  $x$ ；
- 在一次报数、 $x$  退出后，待抽签序列为  $P_{n-1} = [x+1, \dots, \bar{s}, \underline{s}, \dots, x-1]$ .

- | 这么定义是为了符合报数规则

- 对  $P$  序列进行不断操作，定义函数  $g: \mathbb{R} \rightarrow \mathbb{R}$ ，为从拥有  $n$  个元素的  $P$  开始得到最终玩家胜利编号  $w$  的映射

$$g: n \mapsto w$$

- 设  $L_1, L_2$  为两个包含参数个数相同，且参数两两不同的两个序列，定义  $p: \mathbb{R} \rightarrow \mathbb{R}$  为两者编码之间一一映射关系。举例：假设在  $L_1, L_2$  两次编码中，在  $a, b$  实际位置上两序列中的编码如下图所示：

实际位置	$a$	...	$b$
$L_1$	3	...	1
$L_2$	7	...	0

则， $p(3) = 7, p(1) = 0$

## 对从零开始标号序列操作

### 1. 起始数列

$$L_n = [0, 1, \dots, n-1], \quad f(n) = w$$

- 报数后求出被移除的数的表达式： $x = (2-1)\%n = 1\%n$ ，一定要记得减 1，这是我们编码决定的。
- 移除后，待报数序列如下图所示。易得，分别对  $L_n$  和  $P_{n-1}$  进行操作，两者得到的最终胜者应该相同（为  $w$ ）：

$$P_{n-1} = [x+1, \dots, n-1, 0, 1, \dots, x-1], \quad g(n-1) = w$$

- 已知以下两个序列之间呈一一映射关系，我们假设该关系（从  $P_{n-1}$  到  $L_{n-1}$ ）为  $p: \mathbb{R} \rightarrow \mathbb{R}$ 。重要的结论：**胜者的实际位置没有变化，变化的是其位置上的编码**。假设  $L_{n-1}$  得到的胜者为  $w': f(n-1) = w'$ ，在  $P_{n-1}$  中的胜者即为  $p^{-1}(w')$ 。

$$\begin{array}{c} P_{n-1}: [x+1, \dots, n-1, 0, 1, \dots, x-1] \\ \downarrow p \\ L_{n-1}: [0, 1, \dots, x-1, x+1, \dots, n-1] \end{array}$$

### 5. 建立关系

$$\begin{cases} f(n) = g(n-1) = w \\ g(n-1) = p^{-1}(f(n-1)) = p^{-1}(w') \end{cases} \implies \boxed{f(n) = p^{-1}(f(n-1))}$$

- 求函数  $p^{-1}$ ，两者成一次函数关系：

$$p^{-1}(t) = \begin{cases} t + (x+1), & 0 < t + (x+1) < n \\ t + (x+1) - n, & t + (x+1) \geq n \end{cases}$$

化简得到：

$$p^{-1}(t) = (t + x + 1) \% n$$

7. 代入式子：

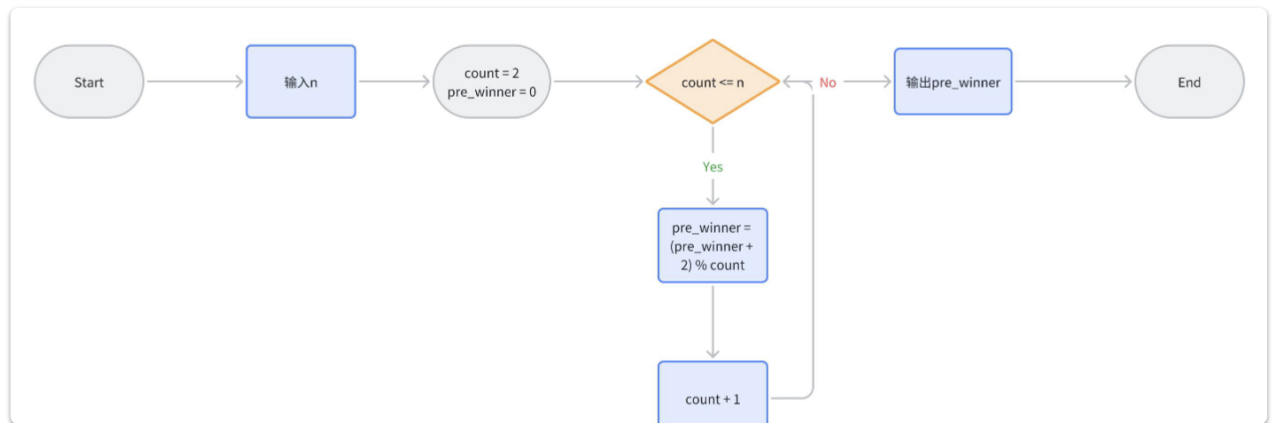
$$\begin{aligned} f(n) &= p^{-1}(f(n-1)) \\ &= (f(n-1) + x + 1) \% n \\ &= (f(n-1) + 1 \% n + 1) \% n \\ &= (f(n-1) + 2) \% n \end{aligned}$$

结论：对  $L_n$  进行操作，最后剩余的玩家通过以下式子得到：

$$f(n) = (f(n-1) + 2) \% n$$

## 对任意序列任意操作

1. 当  $n = 1$  时，易得  $f(n)|_{n=1} = 1$ .
2. 当  $n > 1$  时，求带有  $n$  玩家的问题，通过对上式不断递推得到。  
所得结论即流程图所示图：



## 附录

### A：源码

```

/* Homework : Josephus 问题
 * Author : Brandon Lin 林楠
 * Date : 25 September, 2023 */

#include <stdio.h>

// Declaration
int find_winner ( int player_amounts );

int main(){
    // Definitions (without initialization)
    int player_amounts; /* 玩家数量 */
    int pre_winner; /* 赢家 */

    // Get the amount of players
    printf("Insert number of players (n): ");
    scanf("%d", &player_amounts);

    // Find the final winner
    pre_winner = find_winner(player_amounts);
  
```

```

    // Print the result
    printf("In the case of %2d players, the winner is %d\n", player_amounts,
pre_winner);
    return 0;
}

int
find_winner ( int player_amounts )
{
    // Definitions
    int pre_winner = 0;
    int count;
    const int step = 2;

    // Run the game
    for ( count = step; count <= player_amounts; ++count ) {
        pre_winner = ( pre_winner + step ) % count;
    }

    return pre_winner;
}

```

## B : 数组实现

时间复杂度明显高于（效率更差）A 实现方式。  
所以！数学理论才是王道。

```

/* Homework : Josephus 问题
 * Author : Brandon Lin 林楠
 * Date : 25 September, 2023 */

#include <stdio.h>

// Declaration
int find_winner(int player_amounts);

int main() {
    int player_amounts;
    int pre_winner;

    // 获取玩家数量
    printf("Insert number of players (n): ");
    scanf("%d", &player_amounts);

    // 找到最终赢家
    pre_winner = find_winner(player_amounts);

    // 打印最终结果
    printf("In the case of %2d players, the winner is %d\n", player_amounts,
pre_winner);
    return 0;
}

int find_winner(int player_amounts) {

```

```

// 创建玩家列表
int player_arr[player_amounts];
int exist_players = player_amounts;
for (int i = 0; i < player_amounts; i++) {
    player_arr[i] = i;
}

int current_player_index = -1; // 一定一定要从 -1 开始!!!!!! 血的教训啊啊啊啊啊

// 运行游戏
while (exist_players > 1) {
    // 寻找下一个活着的玩家
    int count = 0;
    while (count < 2) {
        current_player_index = (current_player_index + 1) % player_amounts;
        if (player_arr[current_player_index] != -1) {
            count++;
        }
    }

    printf("current_player_index %d.\n", current_player_index);
    // 淘汰当前玩家
    player_arr[current_player_index] = -1;
    exist_players--;

}

// 返回赢家的编号
for (int i = 0; i < player_amounts; i++) {
    if (player_arr[i] != -1) {
        return player_arr[i];
    }
}

return 0; // 只是不让编译器爆出警告.....
}

```