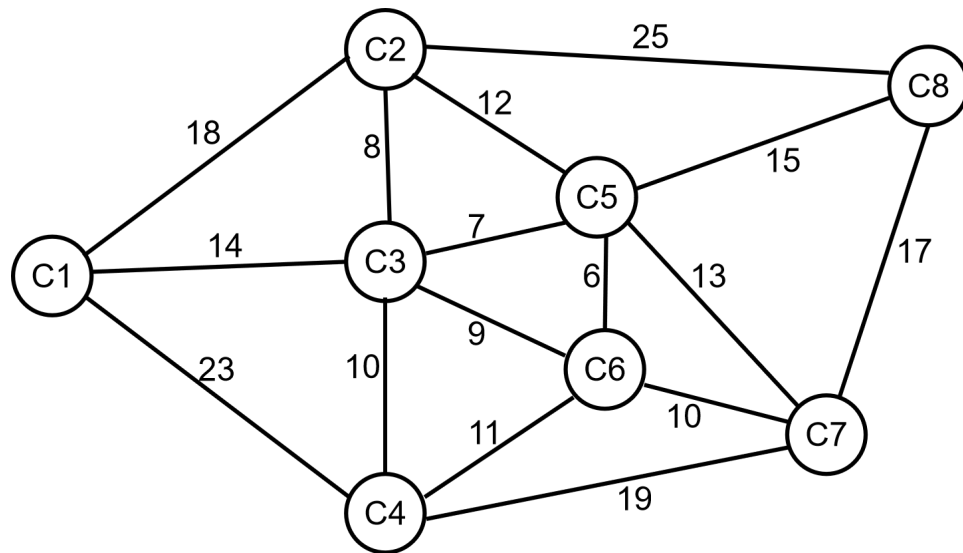


Graph

- Graph - shortest paths
 - directed graph $G=(V,E)$
 - Dijkstra* algorithm - single-pair shortest path
 - Dijkstra* algorithm - **all-pairs shortest paths**

	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0

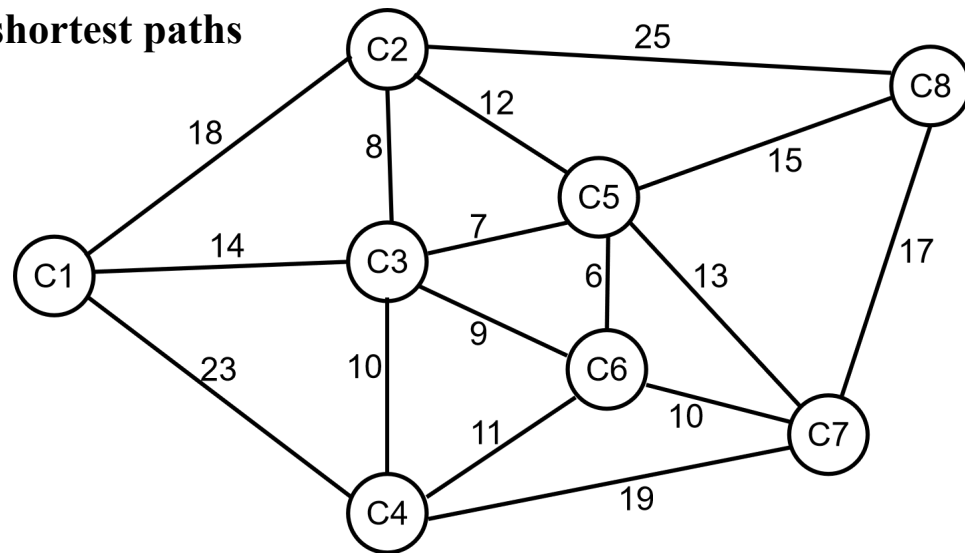
solution directly based on the *Dijkstra* algorithm
 FOR i IN {C1, ..., C8}
 Dijkstra(i) until all other vertices are visited



Graph

- **Graph - shortest paths**
 - directed graph $G=(V,E)$
 - *Dijkstra* algorithm - single-pair shortest path
 - *Dijkstra* algorithm - all-pairs shortest paths
 - *Floyd* algorithm - **all-pairs shortest paths**
 - *dynamic programming*

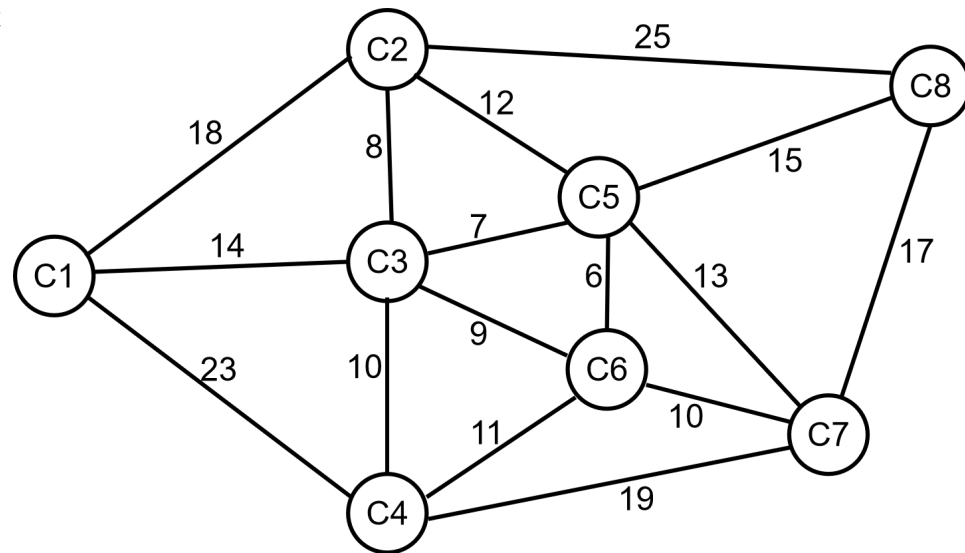
	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0



Graph



- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - best (k+1)-path from s to t is either case 1) or case 2)
 - 1) the best k-path from s to k followed by the best k-path from k to t
 - 2) the best k-path from s to t



Floyd algorithm

INIT: Set the 0-path (direct edges among vertices)

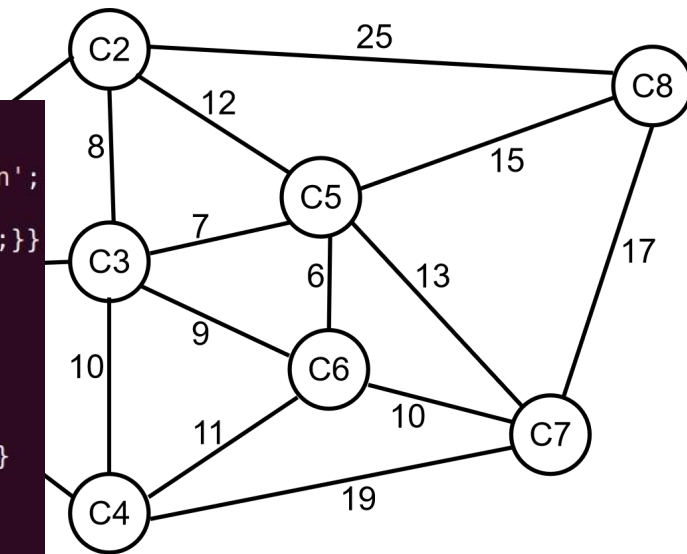
LOOP: k FROM 1 TO n

 Compute the k-path according to the (k-1)-path

Graph

- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - best (k+1)-path from s to t is either case 1) or case 2)
 - 1) the best k-path from s to k followed by the best k-path from k to t
 - 2) the best k-path from s to t

```
// Floyd algorithm - all-pairs shortest paths
void showFloyd(int *d[],int n){
    std::cout<<'\\t';for(int j=0;j<n;j++) std::cout<<j<<'\\t'; std::cout<<'\\n';
    for(int i=0;i<n;i++){std::cout<<i<<'\\t';
        for(int j=0;j<n;j++) std::cout<<d[i][j]<<'\\t'; std::cout<<'\\n';}}
void Floyd(LGraph* g,int *d[]){int n=g->num(),i,j,k;
    for(i=0;i<n;i++) for(j=0;j<n;j++)
        if(g->wgt(i,j)<0) d[i][j]=D_INF; else d[i][j]=g->wgt(i,j);
    for(i=0;i<n;i++) d[i][i]=0;
    for(k=0;k<n;k++){std::cout<<k<<"-path ==> \\n";showFloyd(d,n);
        for(i=0;i<n;i++) for(j=0;j<n;j++)
            if(d[i][j]>(d[i][k]+d[k][j])) d[i][j]=d[i][k]+d[k][j];}
    std::cout<<k<<"-path ==> \\n";showFloyd(d,n);
}
// END Floyd algorithm - all-pairs shortest paths
```



Graph

- Graph - Floyd algorithm - all-pairs shortest paths

0-path =>

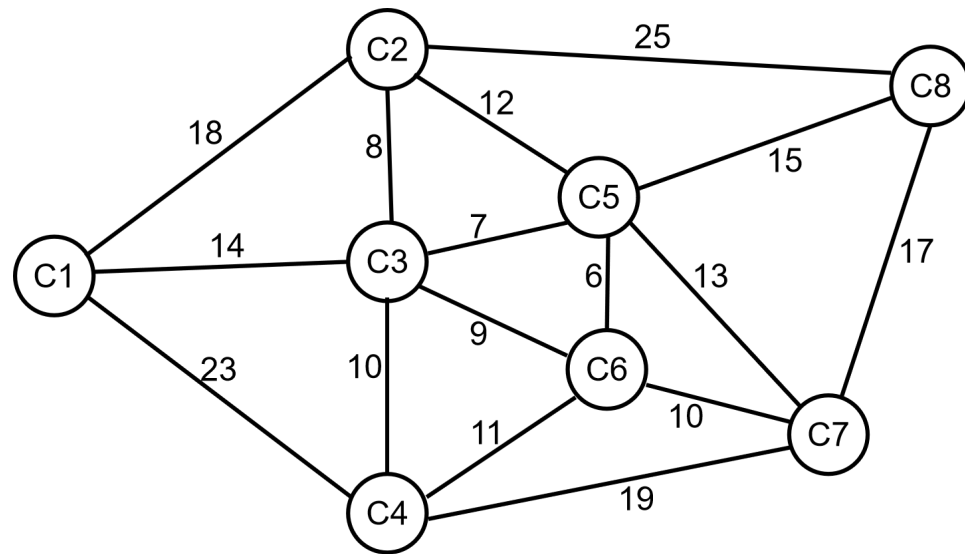
	0	1	2	3	4	5	6	7
0	0	18	14	23	10000	10000	10000	10000
1	18	0	8	10000	12	10000	10000	25
2	14	8	0	10	7	9	10000	10000
3	23	10000	10	0	10000	11	19	10000
4	10000	12	7	10000	0	6	13	15
5	10000	10000	9	11	6	0	10	10000
6	10000	10000	10000	19	13	10	0	17
7	10000	25	10000	10000	15	10000	17	0

1-path =>

	0	1	2	3	4	5	6	7
0	0	18	14	23	10000	10000	10000	10000
1	18	0	8	41	12	10000	10000	25
2	14	8	0	10	7	9	10000	10000
3	23	41	10	0	10000	11	19	10000
4	10000	12	7	10000	0	6	13	15
5	10000	10000	9	11	6	0	10	10000
6	10000	10000	10000	19	13	10	0	17
7	10000	25	10000	10000	15	10000	17	0

2-path =>

	0	1	2	3	4	5	6	7
0	0	18	14	23	30	10000	10000	43
1	18	0	8	41	12	10000	10000	25
2	14	8	0	10	7	9	10000	33
3	23	41	10	0	53	11	19	66
4	30	12	7	53	0	6	13	15
5	10000	10000	9	11	6	0	10	10000
6	10000	10000	10000	19	13	10	0	17
7	43	25	33	66	15	10000	17	0



Graph

- Graph - Floyd algorithm - all-pairs shortest paths

3-path =>

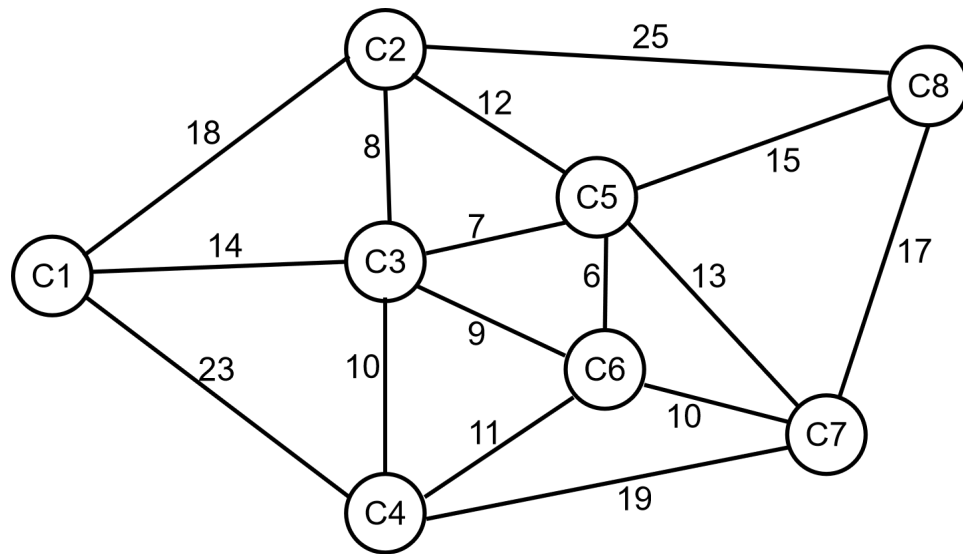
	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	10000	43
1	18	0	8	18	12	17	10000	25
2	14	8	0	10	7	9	10000	33
3	23	18	10	0	17	11	19	43
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	42
6	10000	10000	10000	19	13	10	0	17
7	43	25	33	43	15	42	17	0

4-path =>

	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	42	43
1	18	0	8	18	12	17	37	25
2	14	8	0	10	7	9	29	33
3	23	18	10	0	17	11	19	43
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	42
6	42	37	29	19	13	10	0	17
7	43	25	33	43	15	42	17	0

5-path =>

	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	34	36
1	18	0	8	18	12	17	25	25
2	14	8	0	10	7	9	20	22
3	23	18	10	0	17	11	19	32
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	21
6	34	25	20	19	13	10	0	17
7	36	25	22	32	15	21	17	0



Graph

- Graph - Floyd algorithm - all-pairs shortest paths

6-path =>

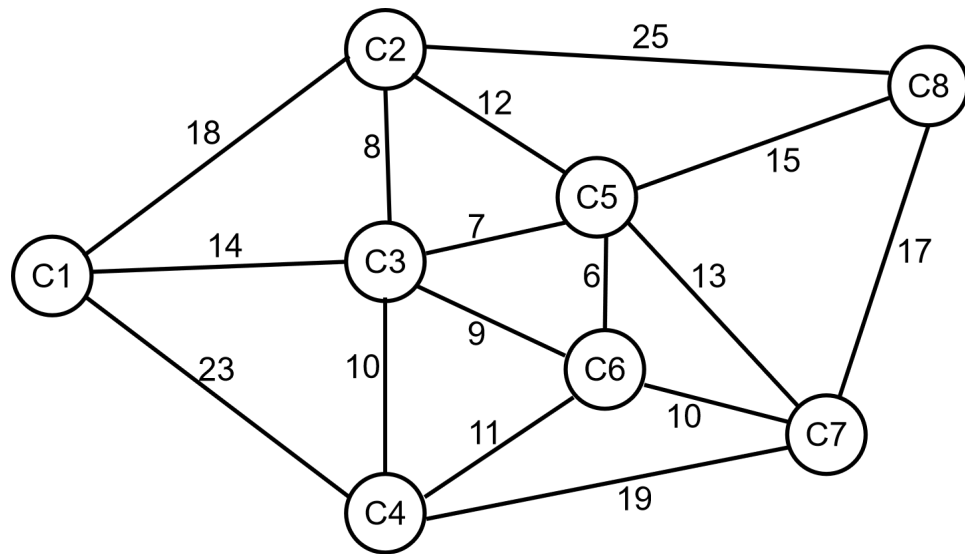
	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	33	36
1	18	0	8	18	12	17	25	25
2	14	8	0	10	7	9	19	22
3	23	18	10	0	17	11	19	32
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	21
6	33	25	19	19	13	10	0	17
7	36	25	22	32	15	21	17	0

7-path =>

	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	33	36
1	18	0	8	18	12	17	25	25
2	14	8	0	10	7	9	19	22
3	23	18	10	0	17	11	19	32
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	21
6	33	25	19	19	13	10	0	17
7	36	25	22	32	15	21	17	0

8-path =>

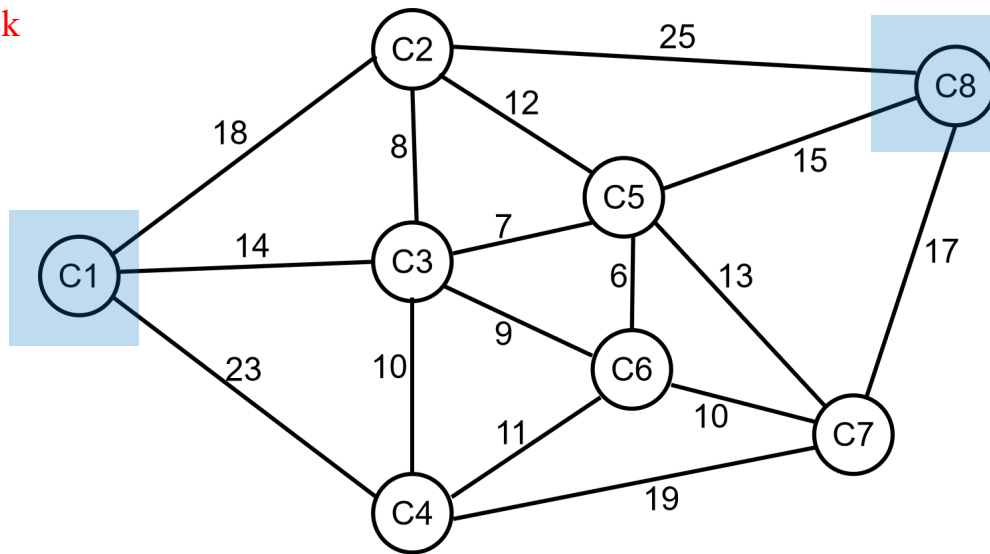
	0	1	2	3	4	5	6	7
0	0	18	14	23	21	23	33	36
1	18	0	8	18	12	17	25	25
2	14	8	0	10	7	9	19	22
3	23	18	10	0	17	11	19	32
4	21	12	7	17	0	6	13	15
5	23	17	9	11	6	0	10	21
6	33	25	19	19	13	10	0	17
7	36	25	22	32	15	21	17	0



Graph

- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - **D.P.** $(k+1)\text{-path}^{\text{best}}[s \text{ to } t] \leq k\text{-path}^{\text{best}}[s \text{ to } k] + k\text{-path}^{\text{best}}[k \text{ to } t]$ or $k\text{-path}^{\text{best}}[s \text{ to } t]$
 - retrieve the min-path
 - no preceding vertex to track
 - how ?

	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0



Graph

- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - **D.P.** $(k+1)\text{-path}^{\text{best}}[s \text{ to } t] \leq k\text{-path}^{\text{best}}[s \text{ to } k] + k\text{-path}^{\text{best}}[k \text{ to } t]$ or $k\text{-path}^{\text{best}}[s \text{ to } t]$
 - retrieve the min-path
 - check the distance consistency

	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0

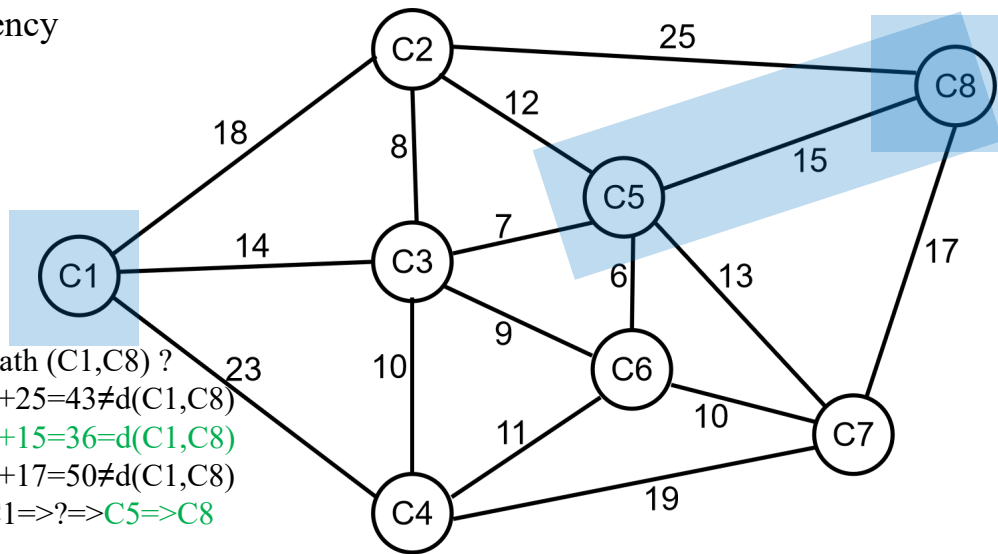
e.g. $d(C1, C8) = 36$, min-path $(C1, C8)$?

$d(C1, C2) + w(C2, C8) = 18 + 25 = 43 \neq d(C1, C8)$

$d(C1, C5) + w(C5, C8) = 21 + 15 = 36 = d(C1, C8)$

$d(C1, C7) + w(C7, C8) = 33 + 17 = 50 \neq d(C1, C8)$

thus min-path $(C1, C8)$: $C1 \Rightarrow ? \Rightarrow C5 \Rightarrow C8$

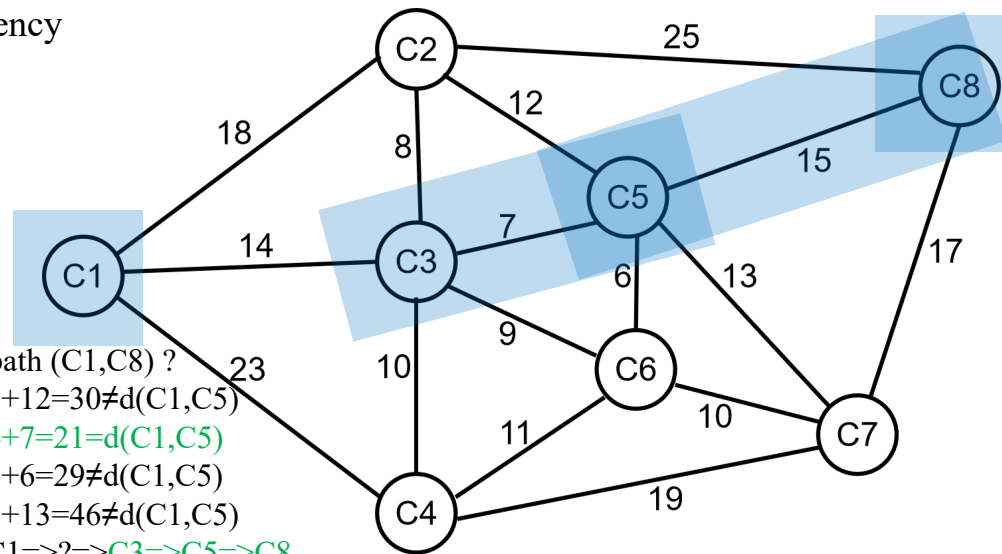


Graph

- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - **D.P.** $(k+1)\text{-path}^{\text{best}}[s \text{ to } t] \leq k\text{-path}^{\text{best}}[s \text{ to } k] + k\text{-path}^{\text{best}}[k \text{ to } t]$ or $k\text{-path}^{\text{best}}[s \text{ to } t]$
 - retrieve the min-path
 - check the distance consistency

	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0

e.g. $d(C1, C8) = 36$, min-path $(C1, C8)$?
 $d(C1, C2) + w(C2, C5) = 18 + 12 = 30 \neq d(C1, C5)$
 $d(C1, C3) + w(C3, C5) = 14 + 7 = 21 = d(C1, C5)$
 $d(C1, C6) + w(C6, C5) = 23 + 6 = 29 \neq d(C1, C5)$
 $d(C1, C7) + w(C7, C5) = 33 + 13 = 46 \neq d(C1, C5)$
 thus min-path $(C1, C8)$: $C1 \Rightarrow ? \Rightarrow C3 \Rightarrow C5 \Rightarrow C8$



Graph

- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - **D.P.** $(k+1)\text{-path}^{\text{best}}[s \text{ to } t] \leq k\text{-path}^{\text{best}}[s \text{ to } k] + k\text{-path}^{\text{best}}[k \text{ to } t]$ or $k\text{-path}^{\text{best}}[s \text{ to } t]$
 - retrieve the min-path
 - check the distance consistency

	C1	C2	C3	C4	C5	C6	C7	C8
C1	0	18	14	23	21	23	33	36
C2	18	0	8	18	12	17	25	25
C3	14	8	0	10	7	9	19	22
C4	23	18	10	0	17	11	19	32
C5	21	12	7	17	0	6	13	15
C6	23	17	9	11	6	0	10	21
C7	33	25	19	19	13	10	0	17
C8	36	25	22	32	15	21	17	0

e.g. $d(C1, C8) = 36$, min-path $(C1, C8)$?

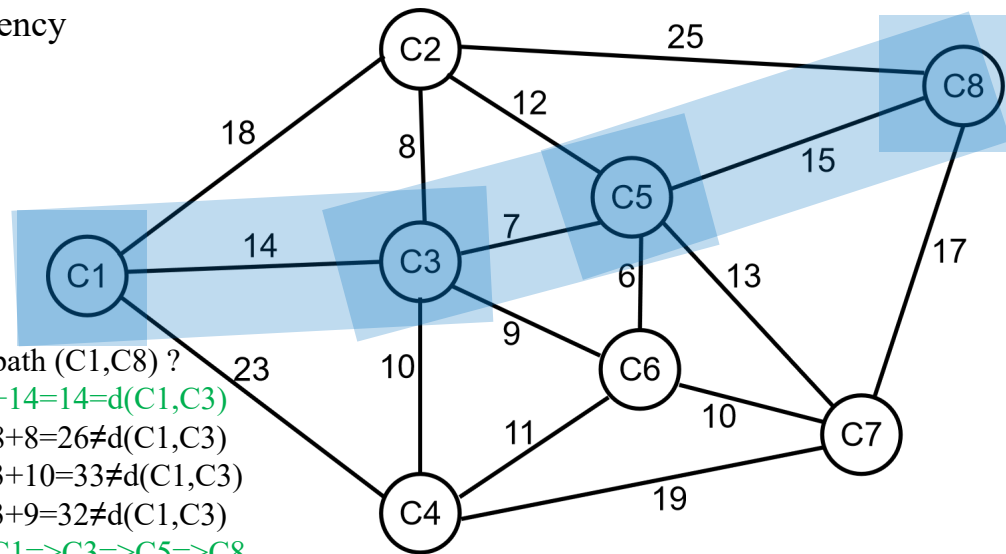
$d(C1, C1) + w(C1, C3) = 0 + 14 = 14 = d(C1, C3)$

$d(C1, C2) + w(C2, C3) = 18 + 8 = 26 \neq d(C1, C3)$

$d(C1, C4) + w(C4, C3) = 23 + 10 = 33 \neq d(C1, C3)$

$d(C1, C6) + w(C6, C3) = 23 + 9 = 32 \neq d(C1, C3)$

thus min-path $(C1, C8)$: $C1 \Rightarrow C3 \Rightarrow C5 \Rightarrow C8$



Graph



- **Graph - Floyd algorithm** - all-pairs shortest paths
 - k-path : intermediate vertices (except the two ends) all have indices less than k
 - **D.P.** $(k+1)\text{-path}^{\text{best}}[s \text{ to } t] \leq k\text{-path}^{\text{best}}[s \text{ to } k] + k\text{-path}^{\text{best}}[k \text{ to } t]$ or $k\text{-path}^{\text{best}}[s \text{ to } t]$
 - retrieve the min-path - check the distance consistency
 - complexity
 - adjacency list based *Dijkstra* with heap based MDFO: $O(|V| (|V|+|E|) \log |V|)$
 - suitable to sparse graphs
 - adjacency matrix based *Dijkstra*: $O(|V| |V|^2) = O(|V|^3)$
 - *Floyd*: $O(|V|^3)$
 - suitable to dense graphs (in terms of not only efficiency but also implementation simplicity)

```
void Floyd(LGraph* g,int *d[]){int n=g->num(),i,j,k;
    for(i=0;i<n;i++) for(j=0;j<n;j++)
        if(g->wgt(i,j)<0) d[i][j]=D_INF; else d[i][j]=g->wgt(i,j);
    for(i=0;i<n;i++) d[i][i]=0;
    for(k=0;k<n;k++){std::cout<<k<<"-path => \n";showFloyd(d,n);
        for(i=0;i<n;i++) for(j=0;j<n;j++)
            if(d[i][j]>(d[i][k]+d[k][j])) d[i][j]=d[i][k]+d[k][j];
        std::cout<<k<<"-path => \n";showFloyd(d,n);
    }
}
```



THANK YOU



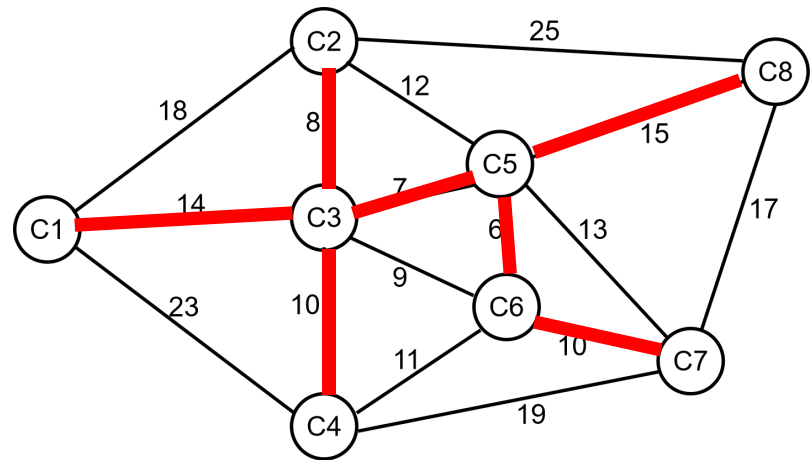
上海交通大學
SHANGHAI JIAO TONG UNIVERSITY

Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - minimum-cost means having the minimum sum of all edge weights of the sub-graph
 - such sub-graph must be a *tree* (so why called MS“T”)
 - a cycle must have a redundant edge that can be removed without violating sub-graph connectivity

MST = {C1C3,C2C3,C3C4,C3C5,C5C6,C5C8,C6C7}

minimum-cost |MST| = $14+8+10+7+6+15+10 = 70$

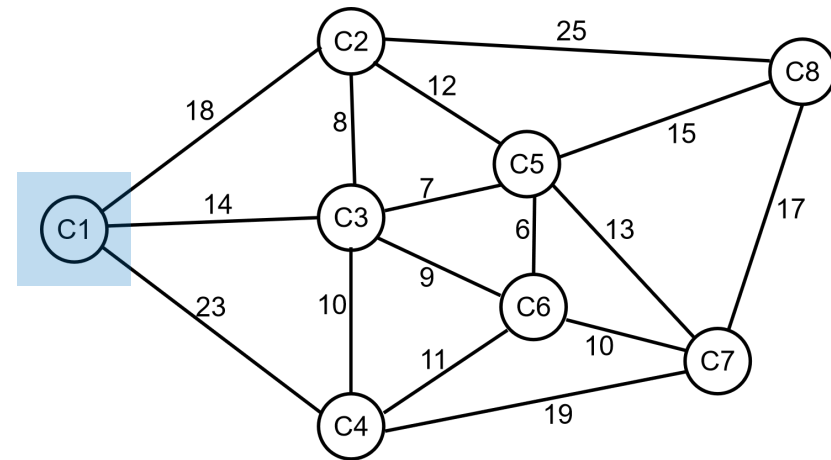
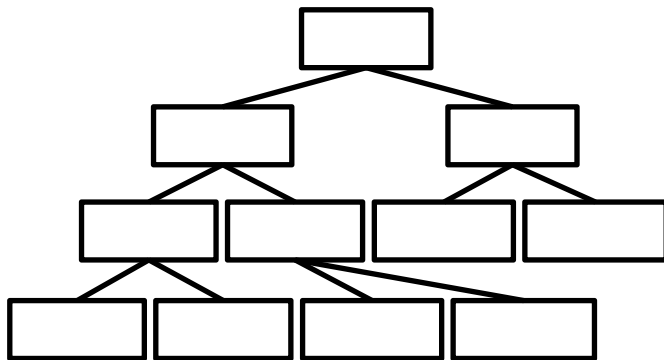


Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

$V = \{C1\}$

$E = \{\}$

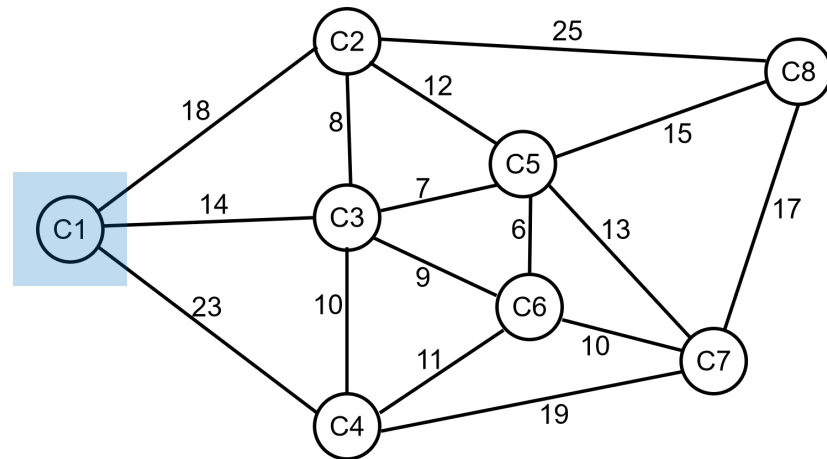
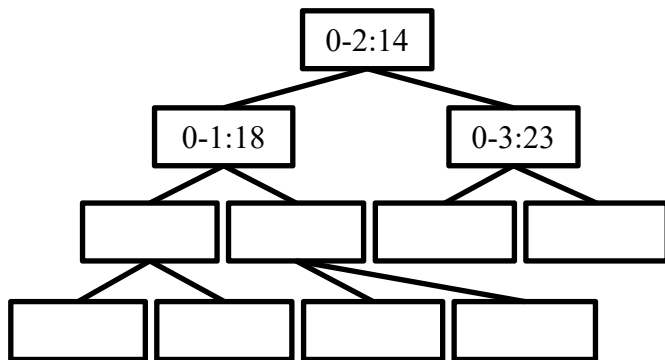


Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
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$E = \{\}$

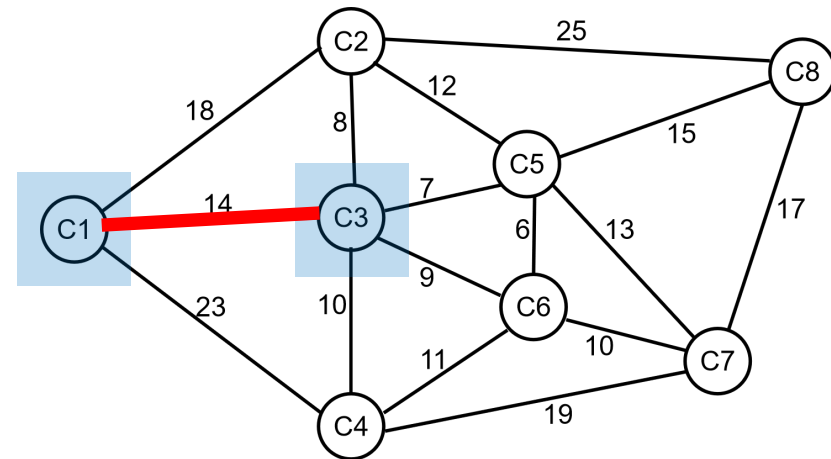
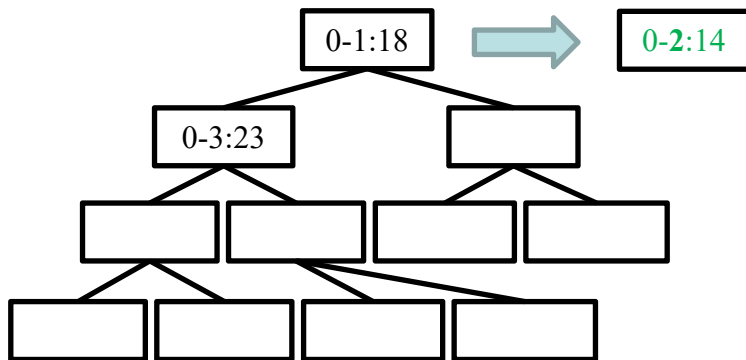


Graph

- Minimum-cost spanning tree (MST)
 - connected & undirected graph $G=(V,E)$
 - connected sub-graph* that has the minimum cost
 - Prim algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

$V = \{C1, C3\}$

$E = \{C1C3\}$

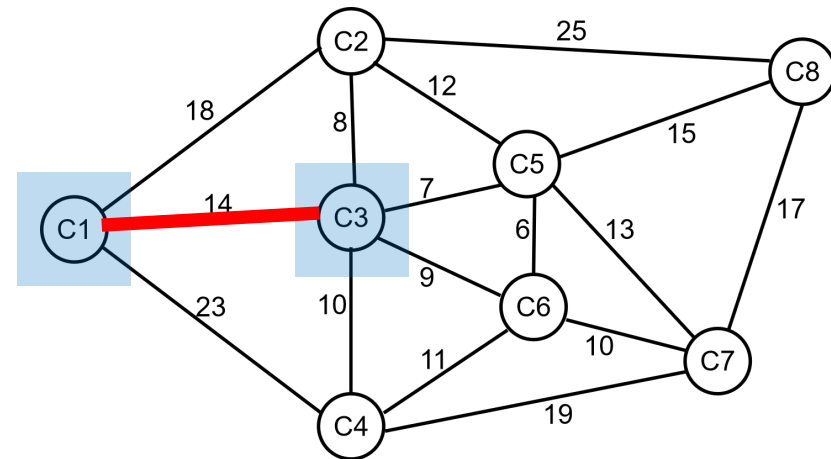
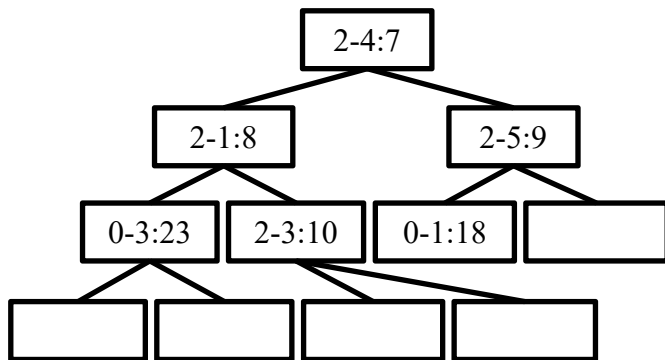


Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

$V = \{C1, C3\}$

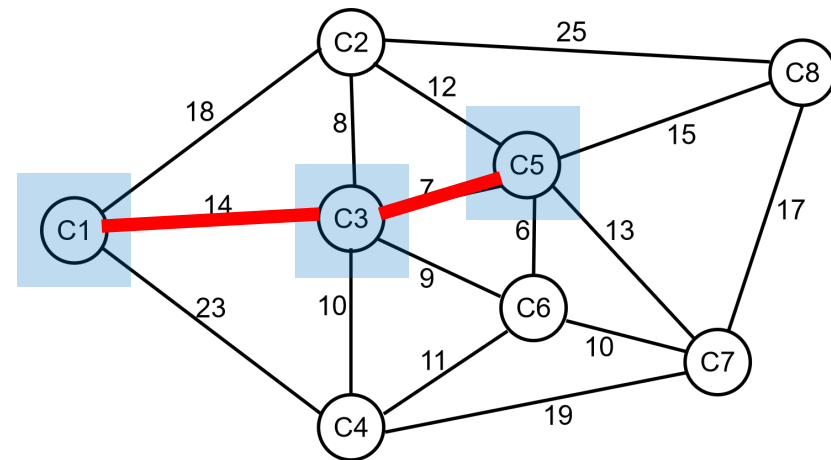
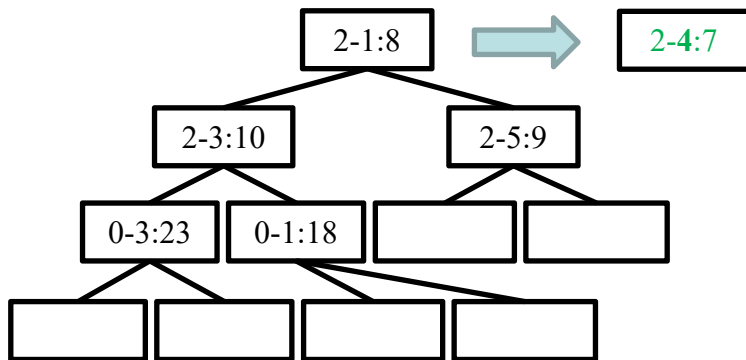
$E = \{C1C3\}$



Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

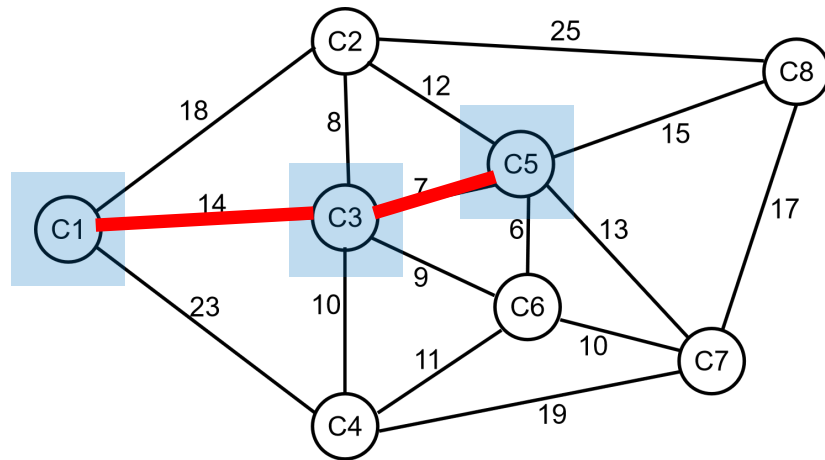
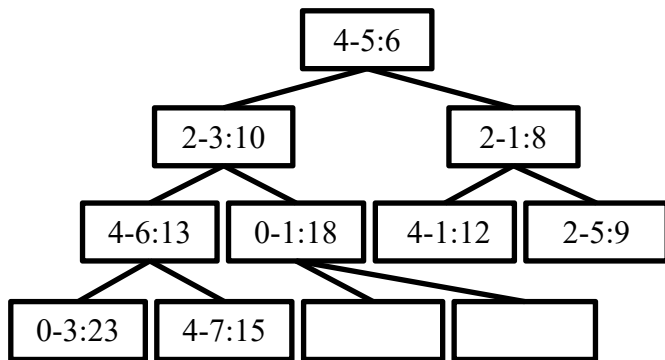
$V = \{C1, C3, C5\}$ —
 $E = \{C1C3, C3C5\}$



Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

$V = \{C1, C3, C5\}$
 $E = \{C1C3, C3C5\}$

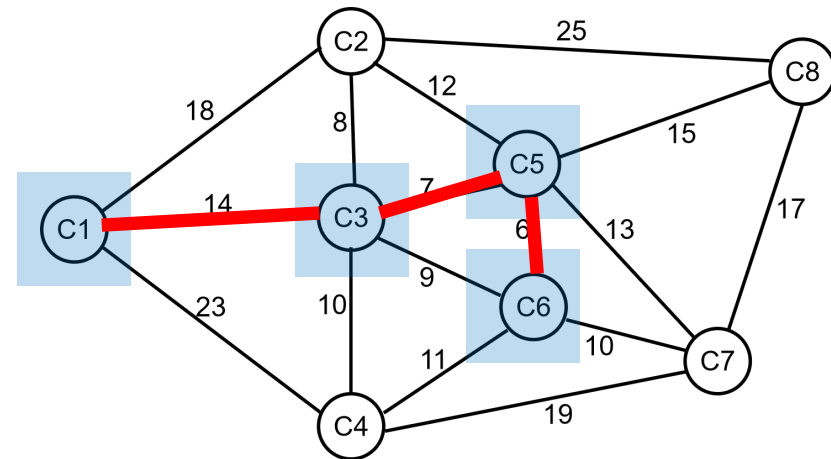
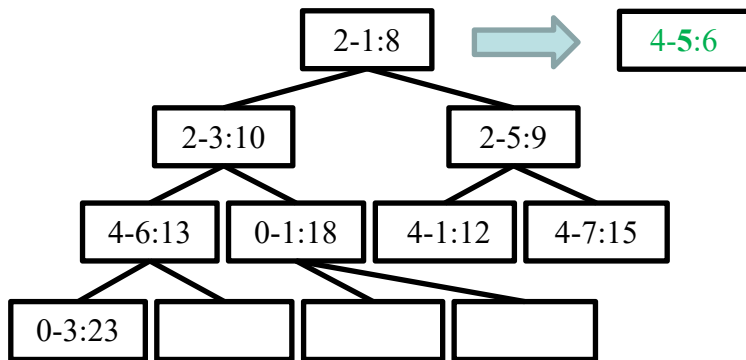


Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm
 - incrementally incorporate the closest vertex until all vertices are incorporated

$V = \{C1, C3, C5, C6\}$

$E = \{C1C3, C3C5, C5C6\}$

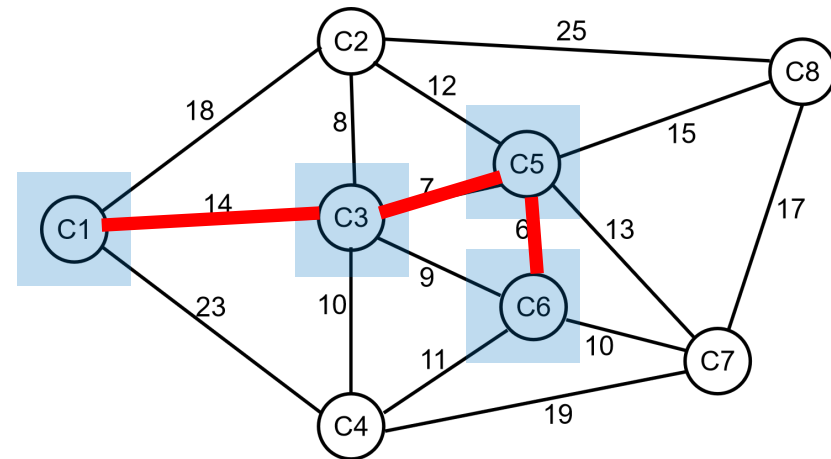
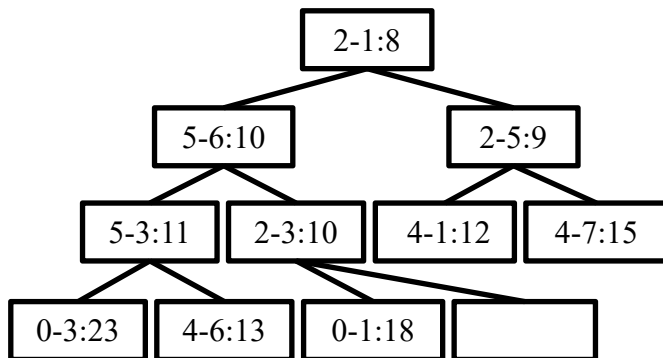


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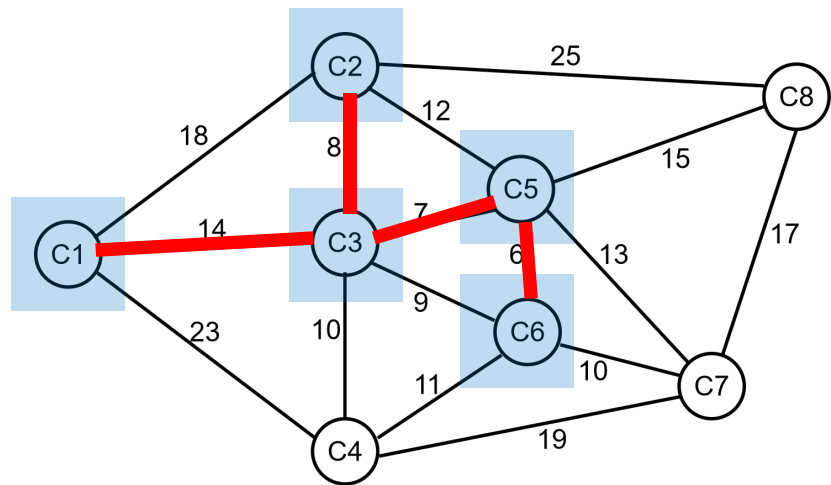
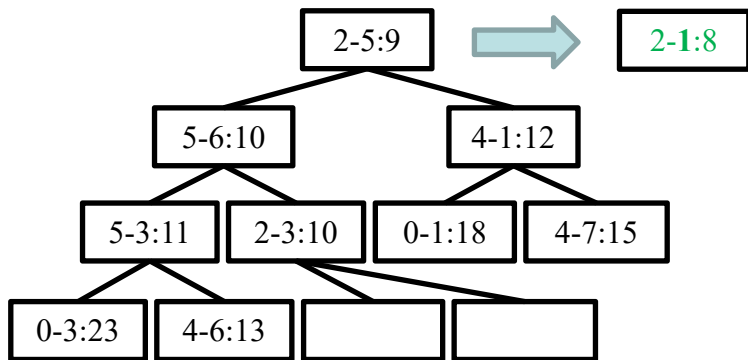


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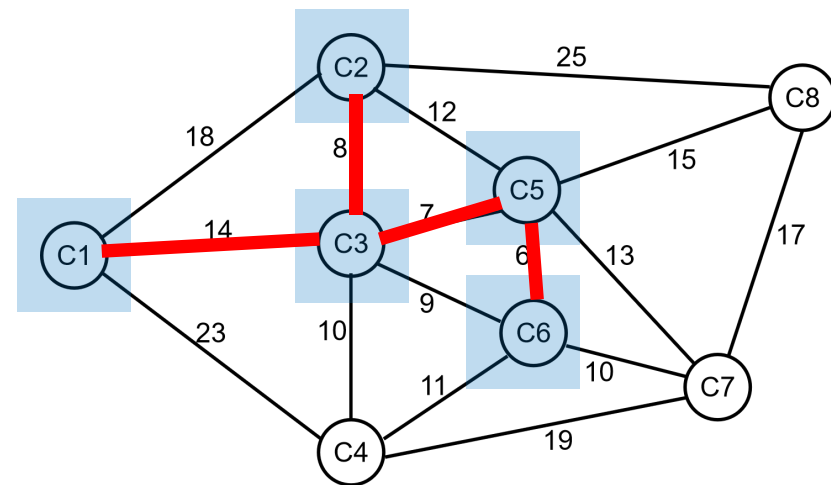
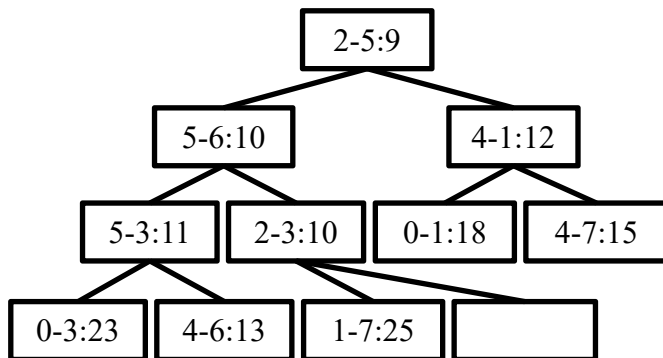


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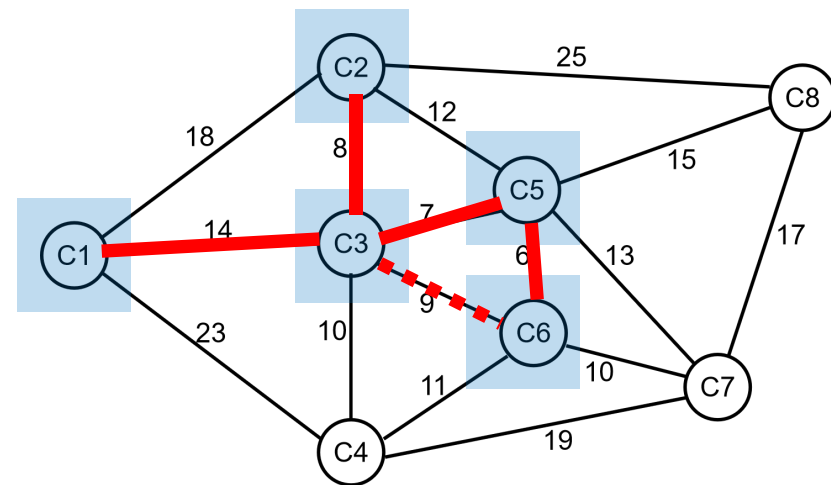
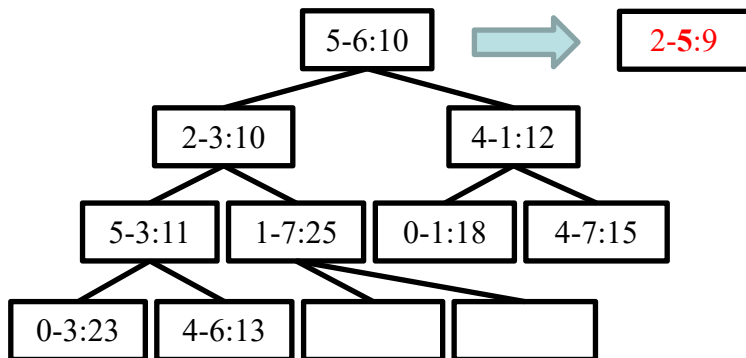


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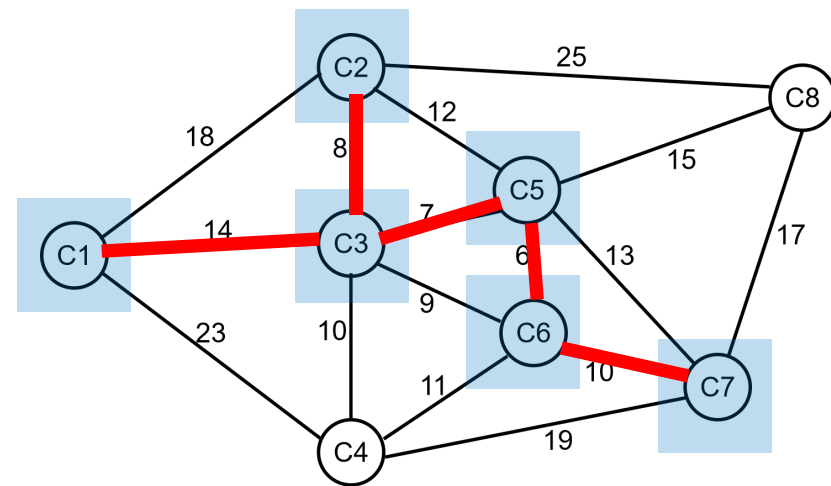
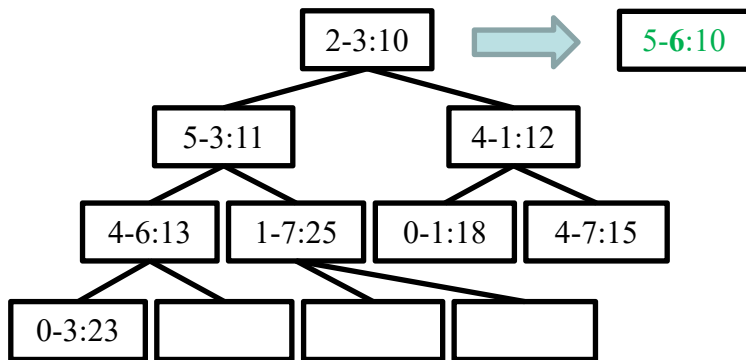


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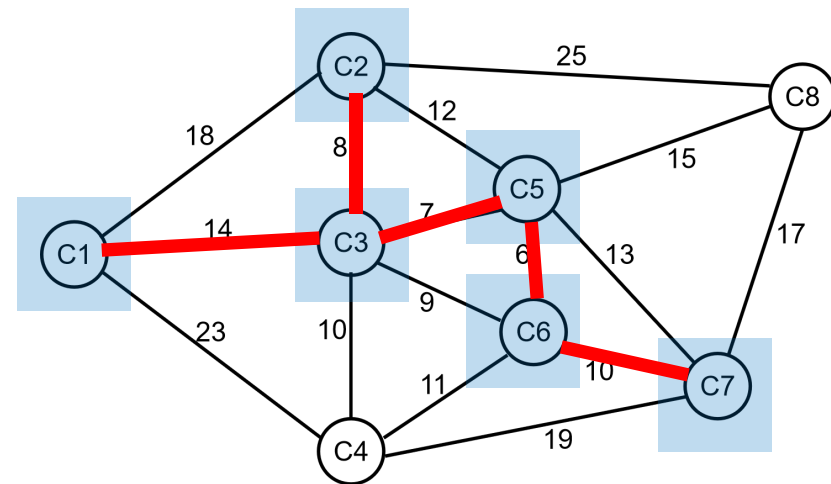
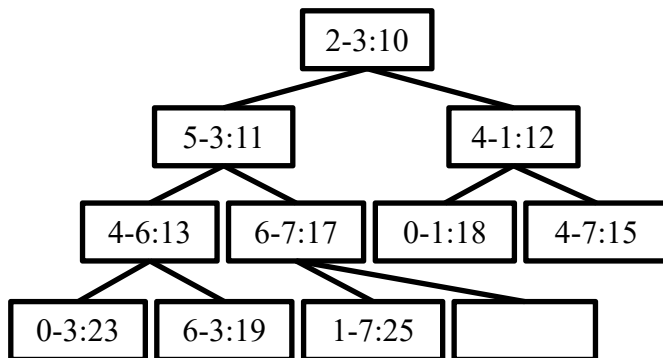


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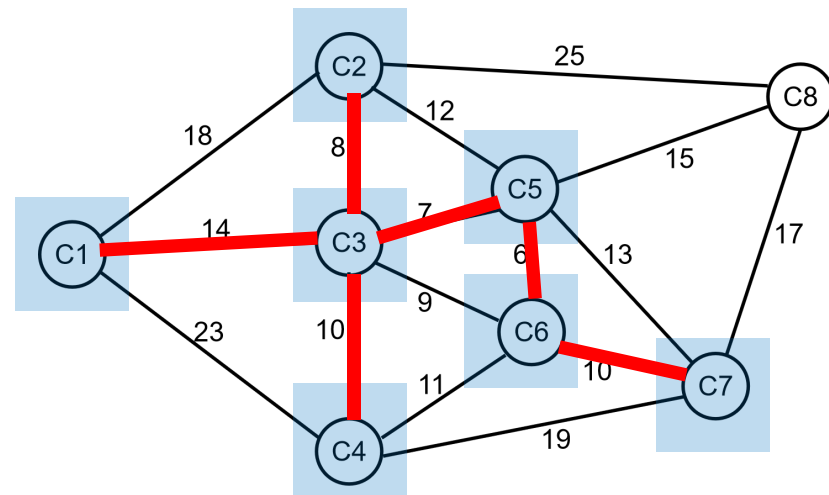
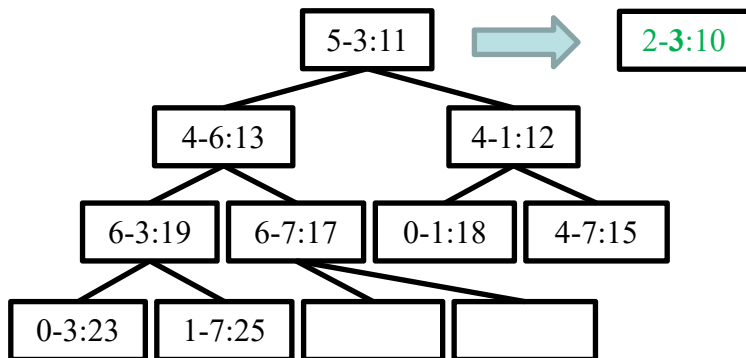


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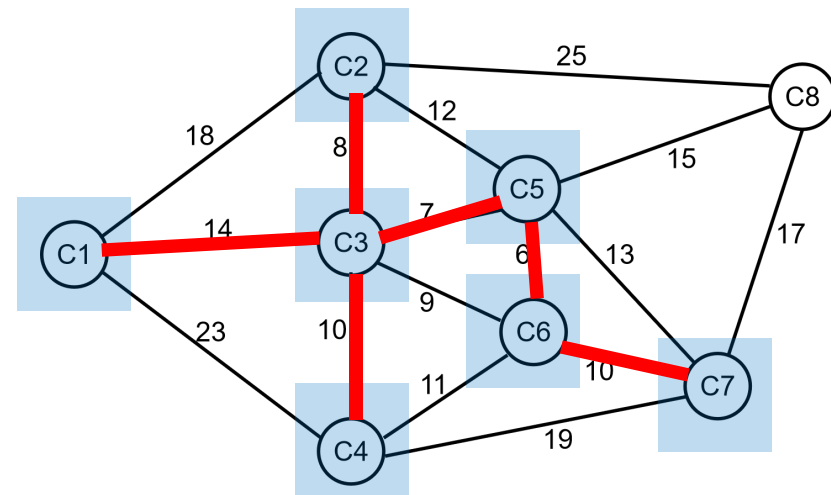
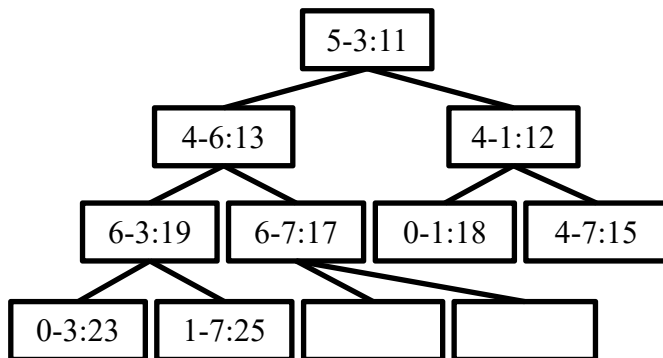


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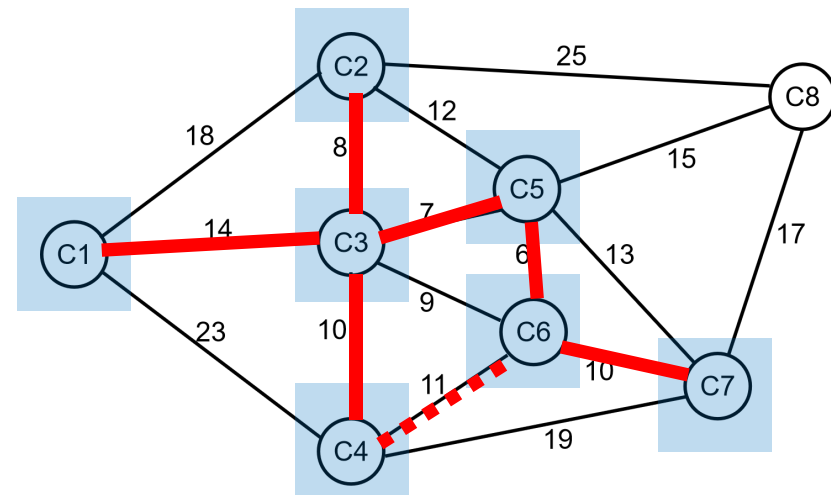
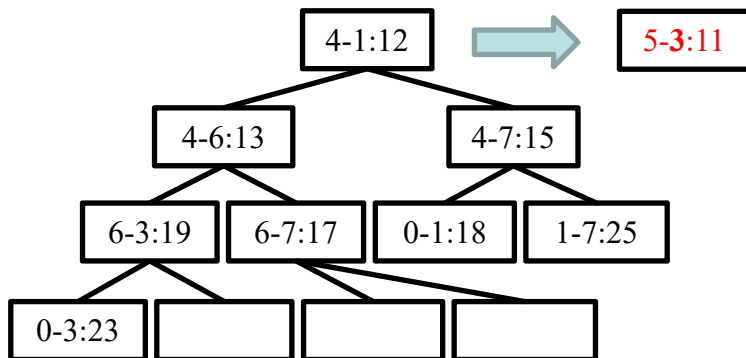


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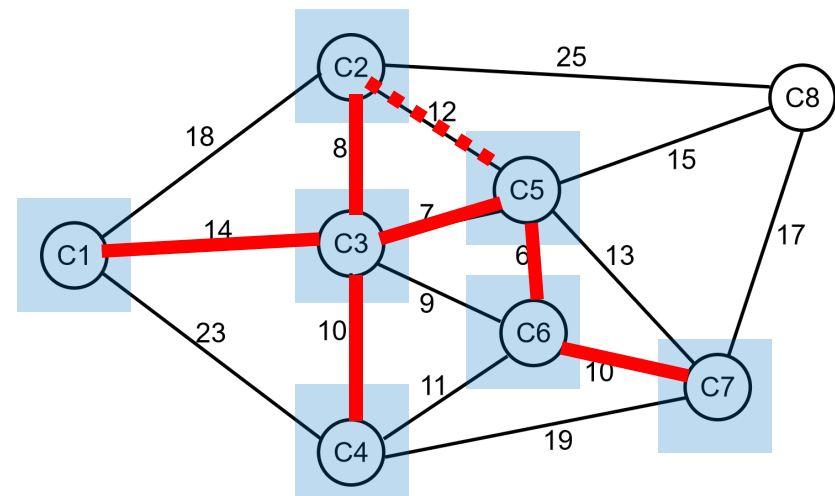
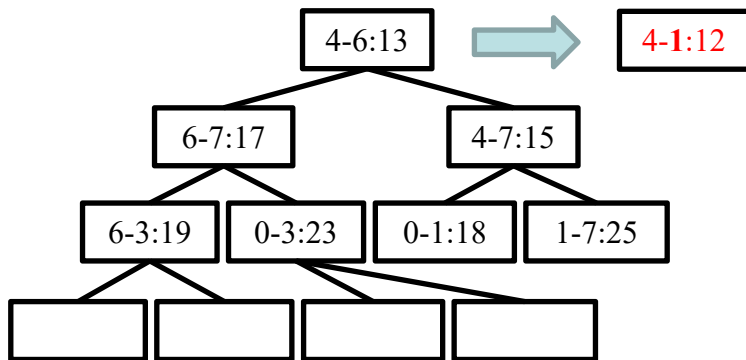


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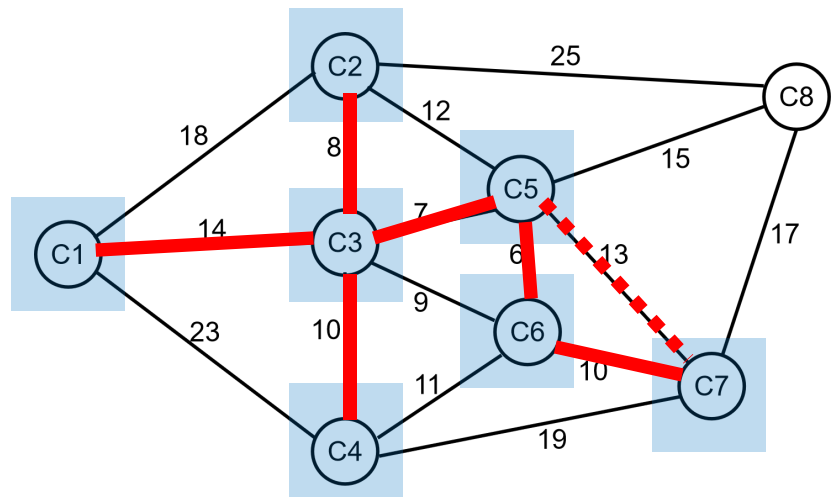
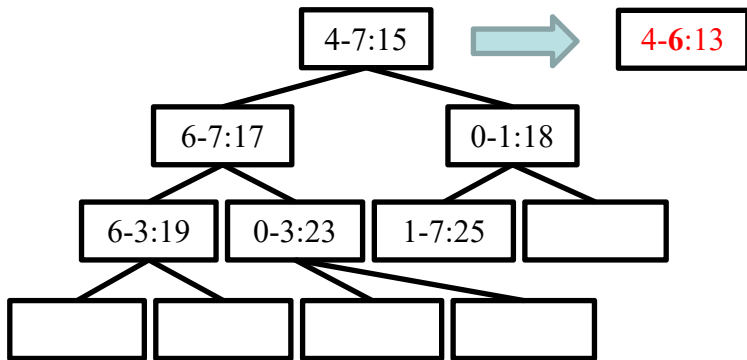


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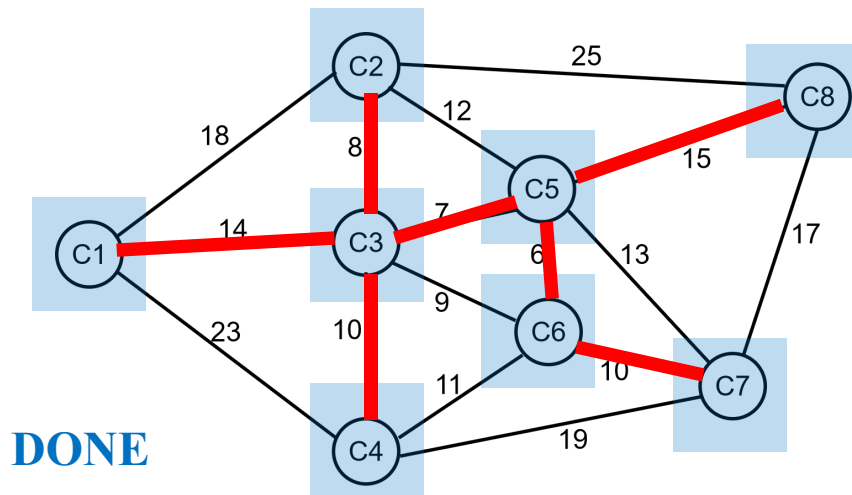
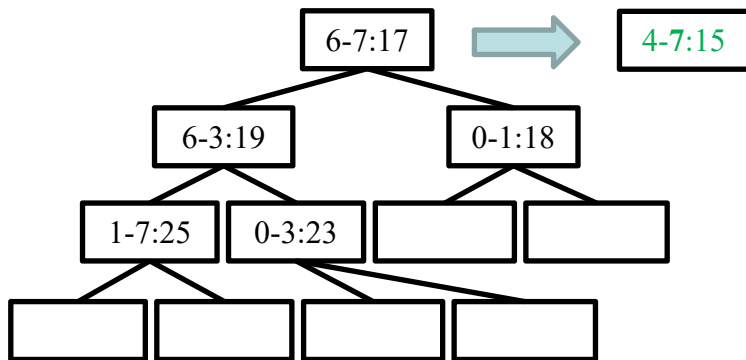


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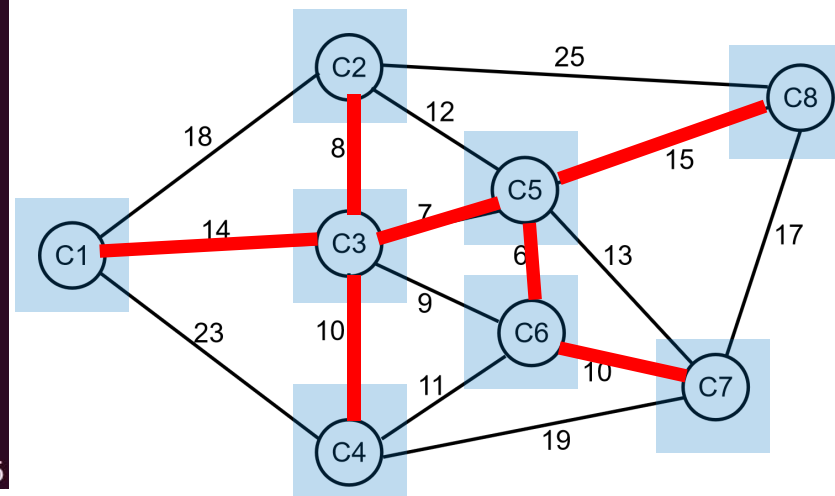
$E = \{C1C3, C3C5, C5C6, C3C2, C6C7, C3C4, C5C8\}$



Graph

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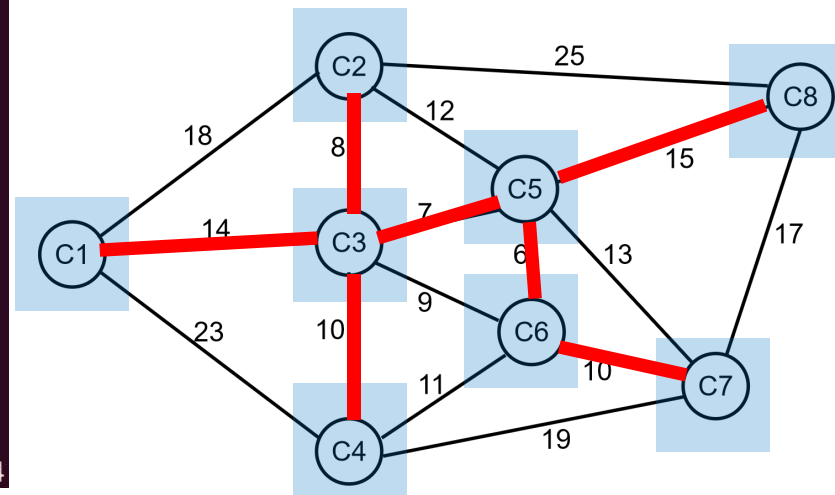
```
Prim MST from 0 =>
add out-MST edges from 0 : =>0-1:18=>0-2:14=>0-3:23
remove min-edge 0-2:14=>2 incorporated into MST
add out-MST edges from 2 : =>2-1:8=>2-3:10=>2-4:7=>2-5:9
remove min-edge 2-4:7=>4 incorporated into MST
add out-MST edges from 4 : =>4-1:12=>4-5:6=>4-6:13=>4-7:15
remove min-edge 4-5:6=>5 incorporated into MST
add out-MST edges from 5 : =>5-3:11=>5-6:10
remove min-edge 2-1:8=>1 incorporated into MST
add out-MST edges from 1 : =>1-7:25
remove min-edge 2-5:9=>5 already in MST
remove min-edge 5-6:10=>6 incorporated into MST
add out-MST edges from 6 : =>6-3:19=>6-7:17
remove min-edge 2-3:10=>3 incorporated into MST
add out-MST edges from 3 :
remove min-edge 5-3:11=>3 already in MST
remove min-edge 4-1:12=>1 already in MST
remove min-edge 4-6:13=>6 already in MST
remove min-edge 4-7:15=>7 incorporated into MST
Prim MST from 0 =>| 0-2:14 2-4:7 4-5:6 2-1:8 5-6:10 2-3:10 4-7:15
```



Graph

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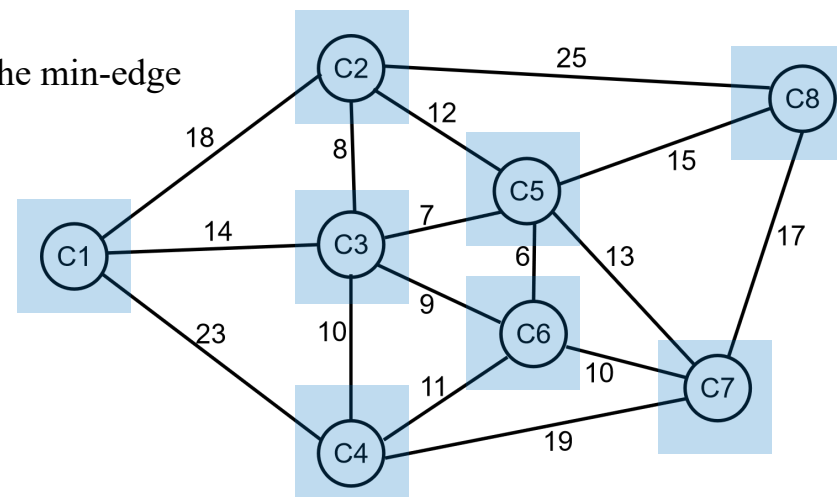
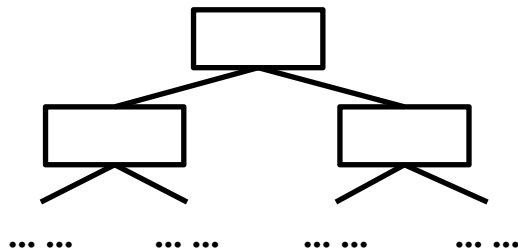
```
Prim MST from 7 =>
add out-MST edges from 7 : =>7-1:25=>7-4:15=>7-6:17
remove min-edge 7-4:15=>4 incorporated into MST
add out-MST edges from 4 : =>4-1:12=>4-2:7=>4-5:6=>4-6:13
remove min-edge 4-5:6=>5 incorporated into MST
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add out-MST edges from 2 : =>2-0:14=>2-1:8=>2-3:10
remove min-edge 2-1:8=>1 incorporated into MST
add out-MST edges from 1 : =>1-0:18
remove min-edge 5-2:9=>2 already in MST
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add out-MST edges from 3 : =>3-0:23=>3-6:19
remove min-edge 5-6:10=>6 incorporated into MST
add out-MST edges from 6 :
remove min-edge 5-3:11=>3 already in MST
remove min-edge 4-1:12=>1 already in MST
remove min-edge 4-6:13=>6 already in MST
remove min-edge 2-0:14=>0 incorporated into MST
Prim MST from 7 =>| 7-4:15 4-5:6 4-2:7 2-1:8 2-3:10 5-6:10 2-0:14
```



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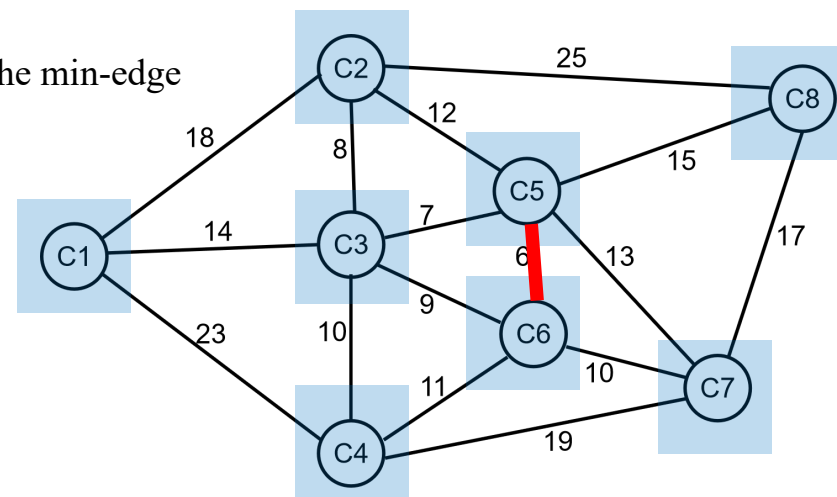
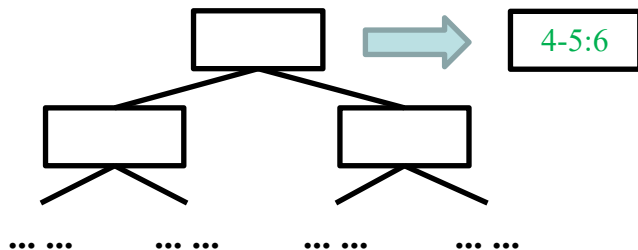
{C1} {C2} {C3} {C4} {C5} {C6} {C7} {C8}
E = {}



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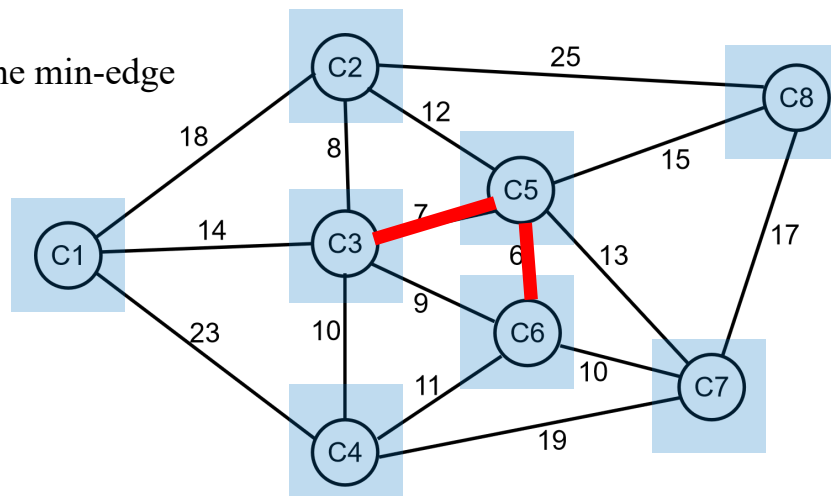
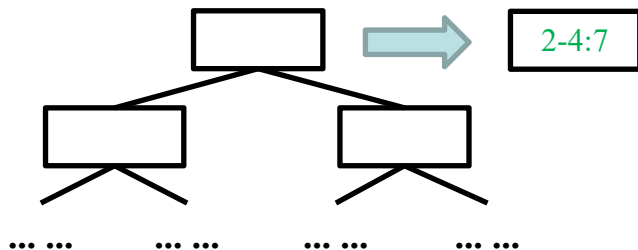
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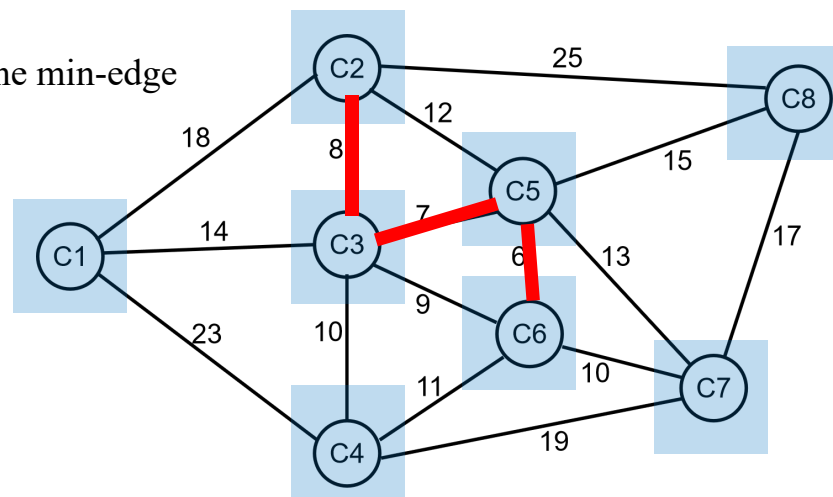
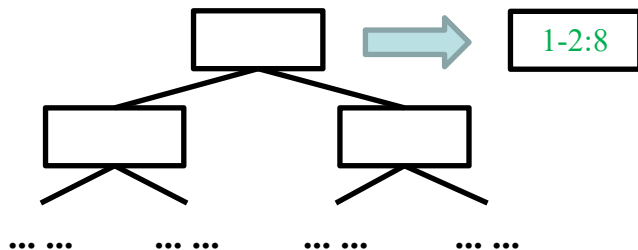
$\{C1\} \{C2\} \{C3, C5, C6\} \{C4\} \{C7\} \{C8\}$
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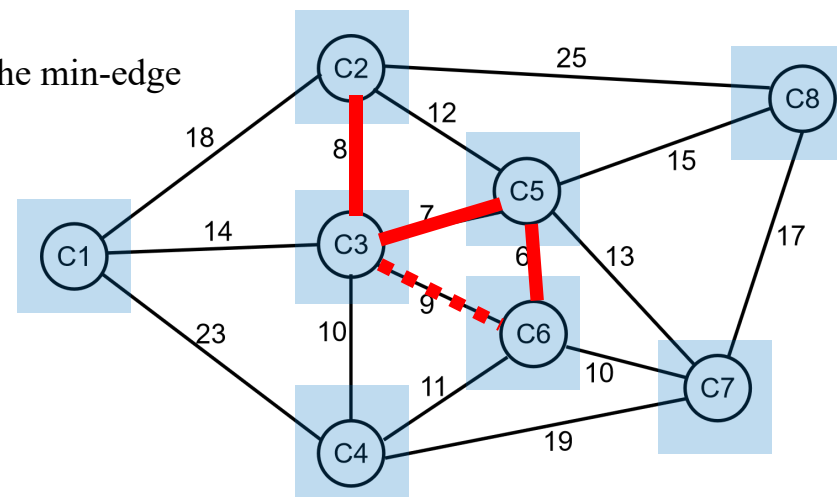
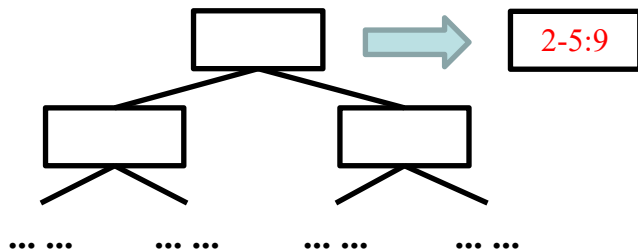
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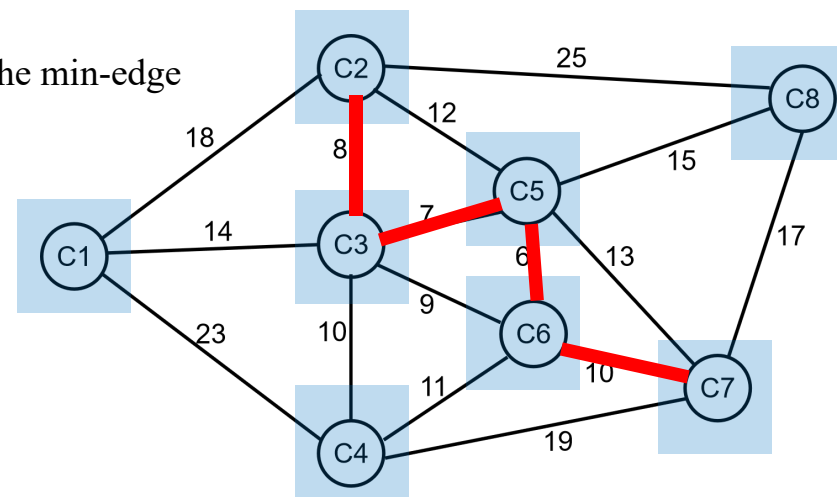
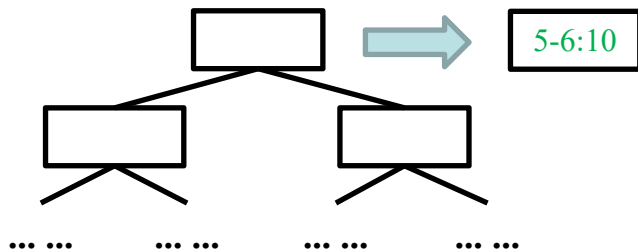
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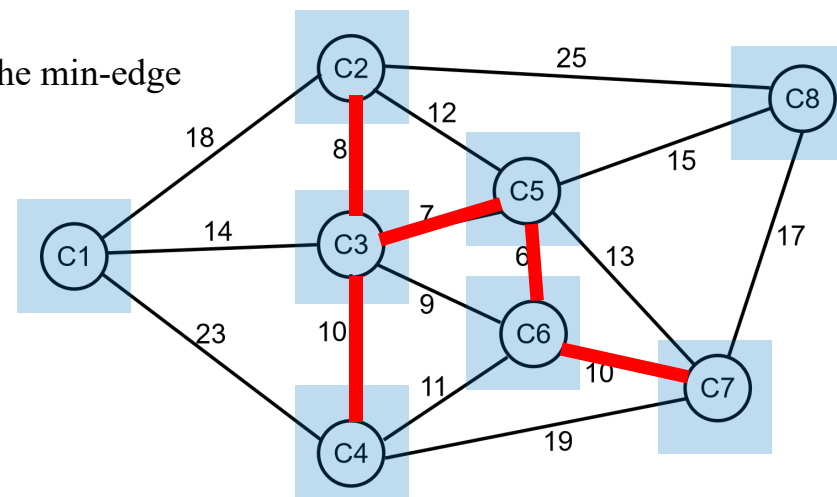
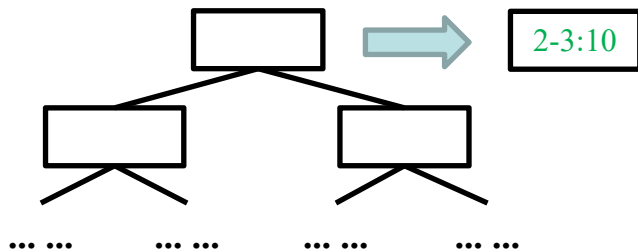
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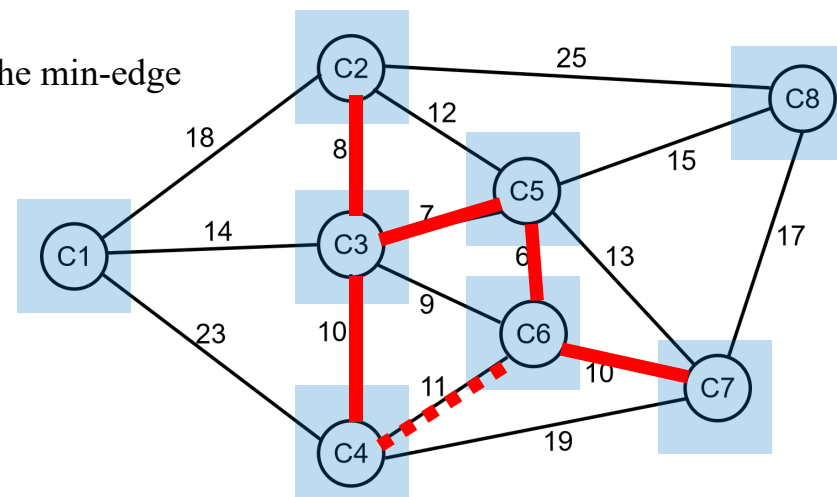
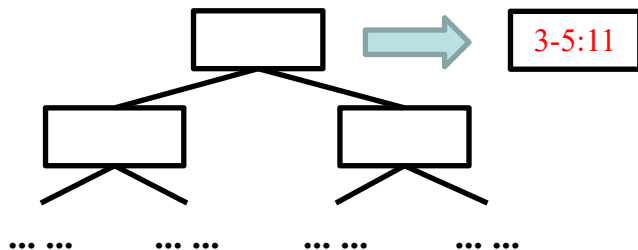
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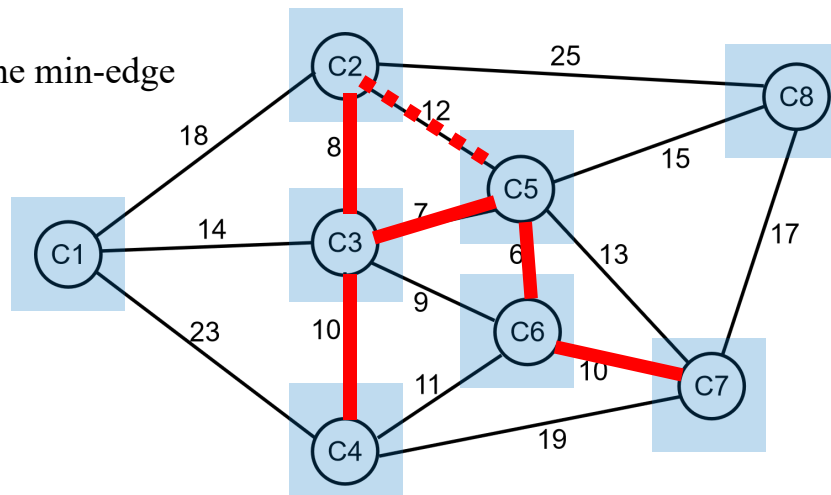
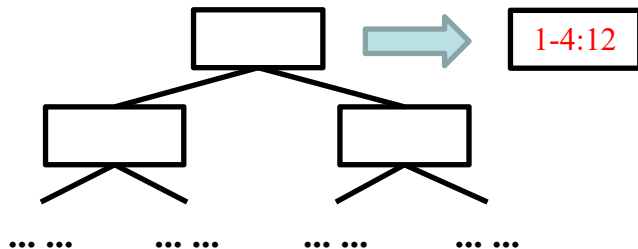
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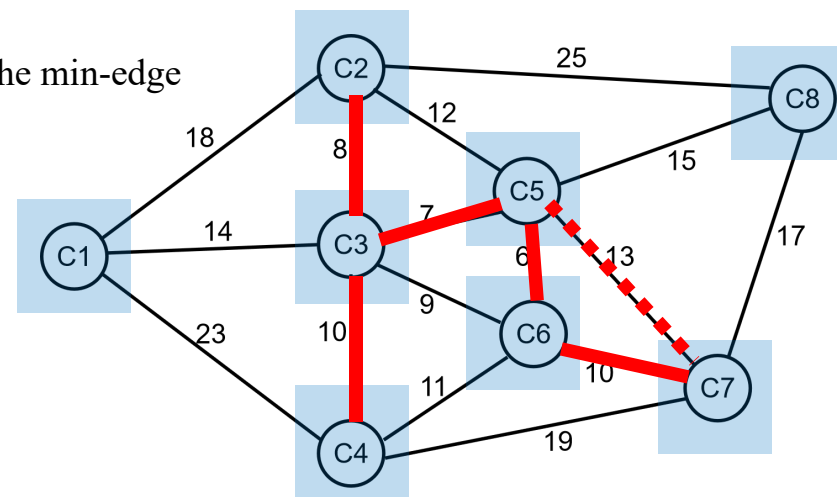
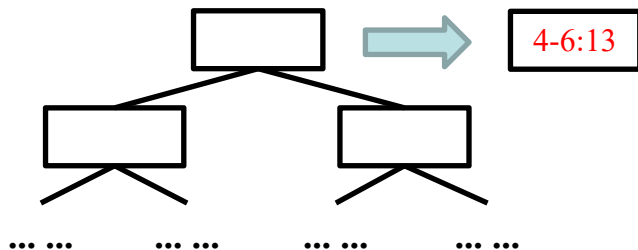
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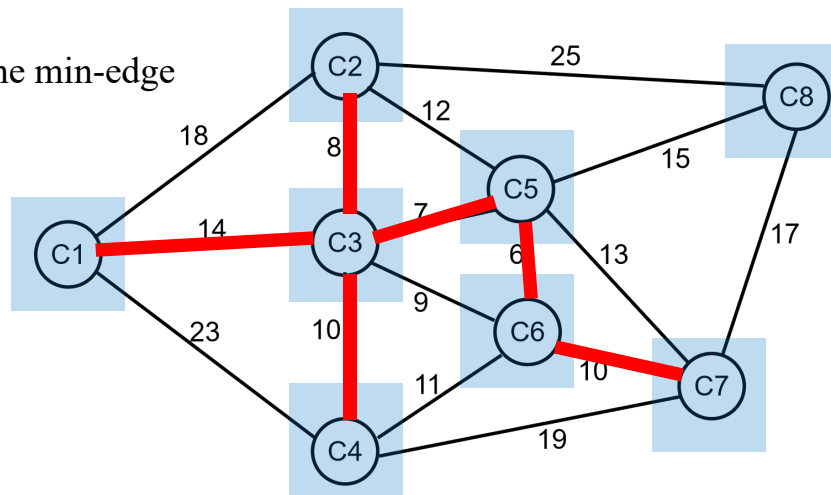
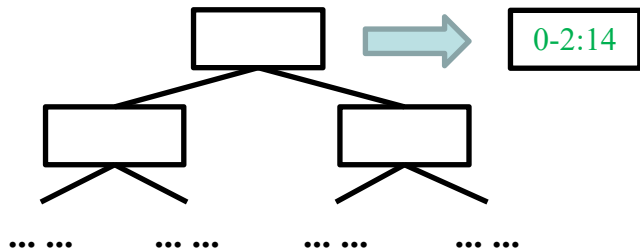
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 $E = \{C5C6, C3C5, C2C3, C6C7, C3C4, C1C3\}$

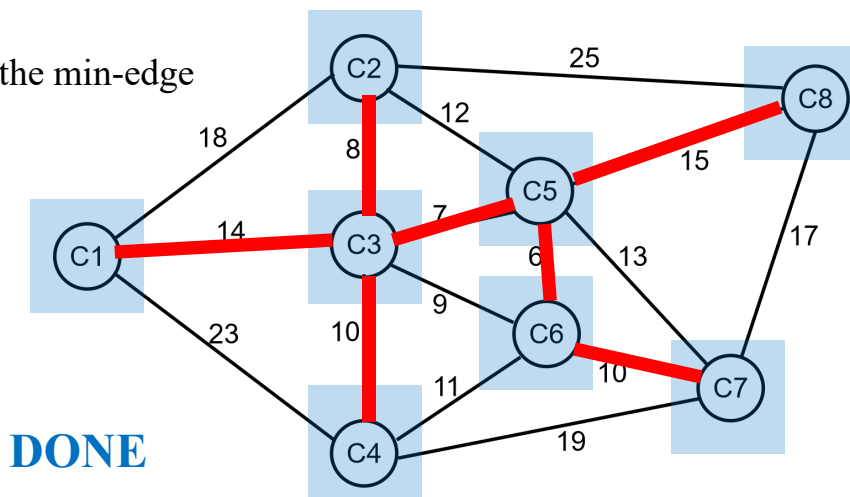
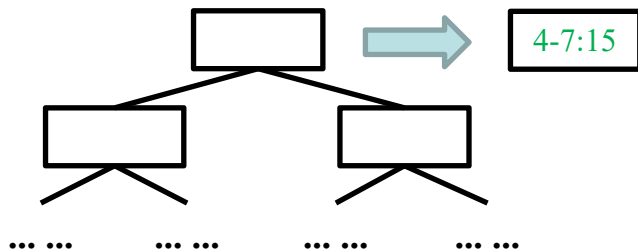


Graph

- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm - incrementally incorporate the closest vertex
 - *Kruskal* algorithm
 - organize graph edges via heap
 - incrementally merge sub-graphs via the min-edge

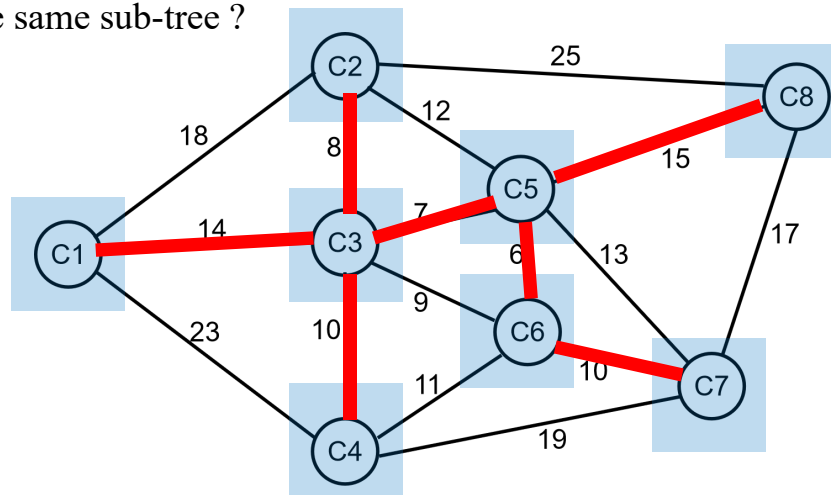
$\{C1, C2, C3, C4, C5, C6, C7, C8\}$

$E = \{C5C6, C3C5, C2C3, C6C7, C3C4, C1C3, C5C8\}$



Graph

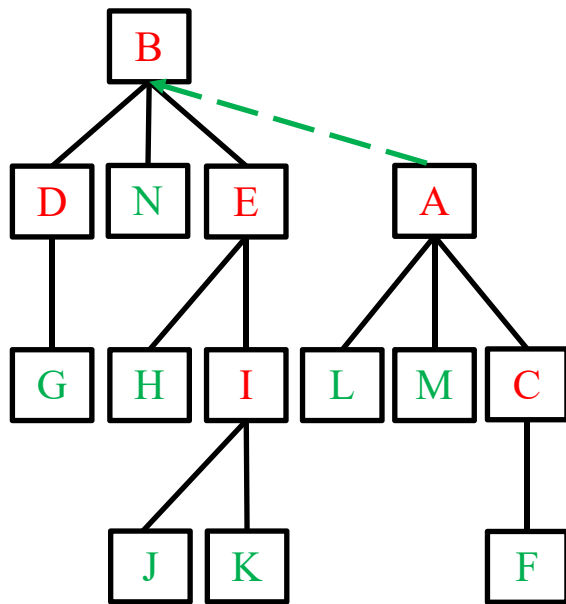
- **Minimum-cost spanning tree (MST)**
 - connected & undirected graph $G=(V,E)$
 - *connected sub-graph* that has the minimum cost
 - *Prim* algorithm - incrementally incorporate the closest vertex
 - *Kruskal* algorithm - incrementally merge sub-graphs (sub-trees) via the min-edge
 - how to check if two vertices are in the same sub-tree ?



Graph

REVIEW

- General tree
 - parent pointer implementation



```
template <class T> class PPNode{ // parent pointer GT node abstract class
private: T e; int n; // node's element; number of nodes of the node's tree
        PPNode *p; // node's parent
public: PPNode(){p=NULL;n=1;} ~PPNode(){} PPNode(const T& ei){e=ei;p=NULL;n=1;}
        const T& getE() const{return e;} void setE(const T& ei){e=ei;}
        int getN() const{return n;} void setN(int ni){n=ni;}
        inline PPNode* getP() const{return p;} void setP(PPNode* g){p=g;}
        static PPNode* find(PPNode* g){while(g->getP()!=NULL) g=g->getP();return g;}
        static void ppunion(PPNode* a,PPNode* b){a=find(a);b=find(b);if(a==b) return;
                if(a->getN()<=b->getN()){a->setP(b);b->setN(b->getN()+a->getN());}
                else{b->setP(a);a->setN(a->getN()+b->getN());}}
};

// ostream overloading, so that 'cout<<{PPNode<T> object}' can have meaning
template <typename T> std::ostream& operator<<(std::ostream& out,PPNode<T>* b){
        out<<b->getE()<<':'<<b->getN(); return out;}

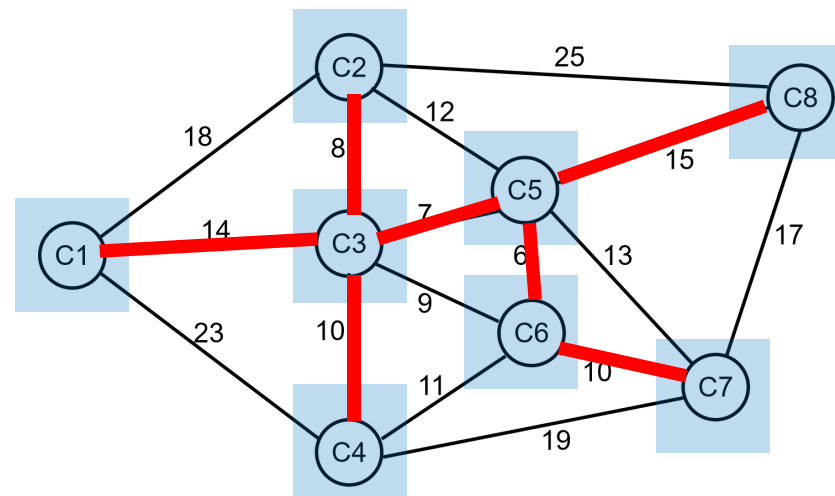
PPNode<char>* p1,*p2;PPNode<char>* pp[26]={new PPNode<char>('A'),
new PPNode<char>('B'), new PPNode<char>('C'), new PPNode<char>('D'), new PPNode<char>('E'), new PPNode<char>('F'),
new PPNode<char>('G'), new PPNode<char>('H'), new PPNode<char>('I'), new PPNode<char>('J'), new PPNode<char>('K'),
new PPNode<char>('L'), new PPNode<char>('M'), new PPNode<char>('N'), new PPNode<char>('O'), new PPNode<char>('P'),
new PPNode<char>('Q'), new PPNode<char>('R'), new PPNode<char>('S'), new PPNode<char>('T'), new PPNode<char>('U'),
new PPNode<char>('V'), new PPNode<char>('W'), new PPNode<char>('X'), new PPNode<char>('Y'), new PPNode<char>('Z')};
while(true){int i=rand()%26,j=rand()%26;p1=PPNode<char>::find(pp[i]);p2=PPNode<char>::find(pp[j]);
        cout<<"merge "<<pp[i]<<" & "<<pp[j]<<" : "<<pp[i]<<" => "<<p1<<" | "<<pp[j]<<" => "<<p2<<" : ";
        if(p1==p2){cout<<p1<<"=="<<p2<<" unnecessary to merge!\n";}
        else{PPNode<char>::ppunion(p1,p2);
                if(p1->getN()<p2->getN()){cout<<" merge into "<<p2<<endl;if(26==p2->getN()) break;}
                else{cout<<" merge into "<<p1<<endl;if(26==p1->getN()) break;}}}
```

Graph

- **Minimum-cost spanning tree (MST)**
 - *Kruskal* algorithm - incrementally merge sub-graphs (sub-trees) via the min-edge

```

Kruskal MST =>
remove min-edge 4-5:6=>vertices 4 & 5 connected (merged)
remove min-edge 2-4:7=>vertices 2 & 4 connected (merged)
remove min-edge 1-2:8=>vertices 1 & 2 connected (merged)
remove min-edge 2-5:9=>vertices 2 & 5 in the same sub-tree
remove min-edge 2-3:10=>vertices 2 & 3 connected (merged)
remove min-edge 5-6:10=>vertices 5 & 6 connected (merged)
remove min-edge 3-5:11=>vertices 3 & 5 in the same sub-tree
remove min-edge 1-4:12=>vertices 1 & 4 in the same sub-tree
remove min-edge 4-6:13=>vertices 4 & 6 in the same sub-tree
remove min-edge 0-2:14=>vertices 0 & 2 connected (merged)
remove min-edge 4-7:15=>vertices 4 & 7 connected (merged)
Kruskal MST =>| 4-5:6 2-4:7 1-2:8 2-3:10 5-6:10 0-2:14 4-7:15
  
```





THANK YOU



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