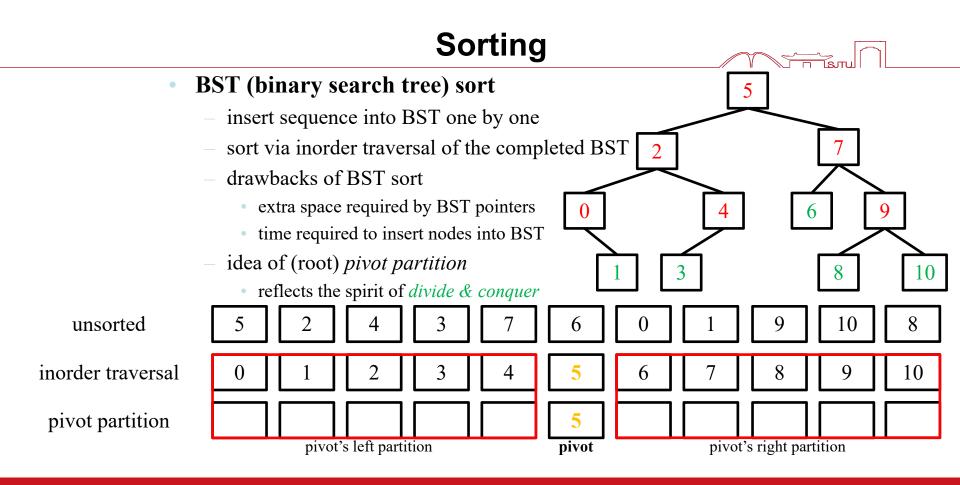




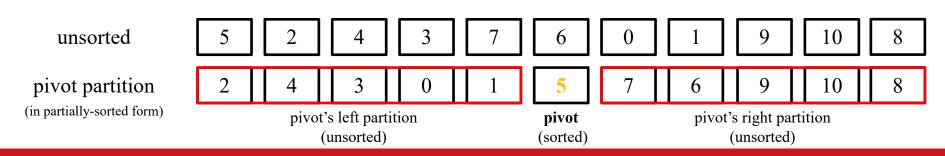
- BST (binary search tree) sort
 - insert sequence into BST one by one
 - sort via inorder traversal of the completed BST
 - drawbacks of BST sort
 - extra space required by BST pointers
 - time required to insert (implying to create as well) nodes into BST
 - even aided by the *freelist* mechanism, creation needs time after all







- Quicksort
 - divide & conquer
 - divide into sub-tasks that are much easier and can be "merged" efficiently
 - pivot partition in partially-sorted form (PSF)
 - put sequence elements that are prior to the pivot to the left partition
 - put sequence elements to which the pivot is prior to the right partition
 - pivot is **sorted** namely its ordinal position is determined immediately
 - left & right partitions are **unsorted** and hence can be formed conveniently & quickly
 - unsorted left & right partitions are sorted via further recursive pivot partition in PSF





Quicksort

- divide & conquer
- pivot partition in partially-sorted form (PSF)

unsorted

pivot partition

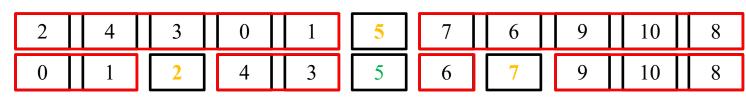
(in partially-sorted form)

5	2	4	3	7	6	0	1	9	10	8
2	4	3	0	1	5	7	6	9	10	8
0	1	2	4	3	5	6	7	9	10	8



Quicksort - divide & conquer

- pivot partition in partially-sorted form (PSF)
 - put sequence elements that are prior to the pivot to the left partition
 - put sequence elements to which the pivot is prior to the right partition
 - pivot is **sorted** namely its ordinal position is determined immediately
 - left & right partitions are **unsorted** and hence can be formed conveniently & quickly
 - unsorted left & right partitions are sorted via further **recursive pivot partition in PSF** pivot partition in PSF is what makes a difference between quicksort & BST sort
 - enable quicksort to be effective implemented on a single fixed array
 - exempt from dynamic maintenance of BST nodes (one-by-one insertion)
 - exempt from establishing unnecesary relations among sequence elements (extra pointers)



Quicksort - divide & conquer

- pivot partition in partially-sorted form (PSF)
 - put sequence elements that are prior to the pivot to the left partition
 - put sequence elements to which the pivot is prior to the right partition
 - pivot is **sorted** namely its ordinal position is determined immediately
 - left & right partitions are **unsorted** and hence can be formed conveniently & quickly
 - unsorted left & right partitions are sorted via further recursive pivot partition in PSF
- pivot partition in PSF is what makes a difference between quicksort & BST sort
- complexity
 - worst case: O(n²)
 - best case: O(n log n) $c(n) = 2c(\frac{n}{2}) + O(n)$
 - average case: O(n log n) $c(n) = \frac{1}{n-1} \sum_{k=1}^{n-1} [c(k) + c(n-k)] + O(n) = \frac{2}{n-1} \sum_{k=1}^{n-1} c(k) + O(n)$



Quicksort - divide & conquer

- pivot partition in partially-sorted form (PSF)
 - put sequence elements that are prior to the pivot to the left partition
 - put sequence elements to which the pivot is prior to the right partition
 - pivot is **sorted** namely its ordinal position is determined immediately
 - left & right partitions are **unsorted** and hence can be formed conveniently & quickly
 - unsorted left & right partitions are sorted via further recursive pivot partition in PSF
- pivot partition in PSF is what makes a difference between quicksort & BST sort
- complexity worst case O(n²); best & average cases O(n log n)
- find pivot
 - to select first or last tends to cause unbalanced partition (for nearly sorted or reverse sorted)
 - better to select a random element or expediently array middle (good enough & efficient)
 - even better to select "median of (random) three" or expediently median of first, middle, last



- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T,class P> void quicksort(T s[],int iL,int iR,int d){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```

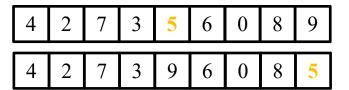




Quicksort - divide & conquer

b

- pivot partition in partially-sorted form (PSF)
- partition trick



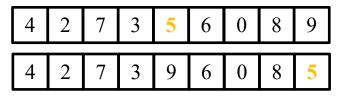
a

```
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template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s.iP.iR-1):iP=iR-1:int a=iL-1.b=iP:
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
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                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



++a b

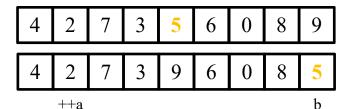
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template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
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        swap(s.iP.iR-1):iP=iR-1:int a=iL-1.b=iP:
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
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        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
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template <class T,class P> void quicksort(T s[],int n){
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        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```



insertionsort<T,P>(s,n);}
// END quick sort: divide & conquer



- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick

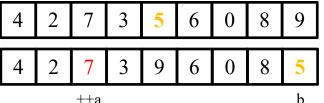


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        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s.iP.iR-1):iP=iR-1:int a=iL-1.b=iP:
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
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                swap(s,a,b);}while(a<b);</pre>
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        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



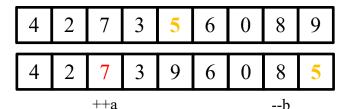
```
b swap(s,iP, do{while((
```

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        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
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        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s.iP.iR-1):iP=iR-1:int a=iL-1.b=iP:
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
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        // take advantage of the best-case performance of insertion sort by a
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        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```



insertionsort<T,P>(s,n);}
END quick sort: divide & conquer

- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick

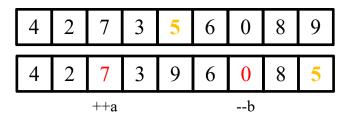


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        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
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        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s.iP.iR-1):iP=iR-1:int a=iL-1.b=iP:
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
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        quicksort<T,P>(s,0,n);
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```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick

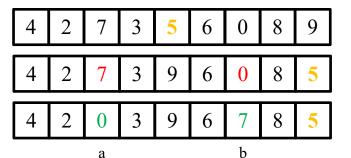


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template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
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        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
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        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
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        // take advantage of the best-case performance of insertion sort by a
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        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



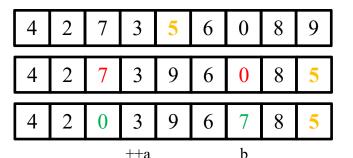
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template <class T,class P> inline int findpivot(T s[],int iL,int iR){
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        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```



// END quick sort: divide & conquer



- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick

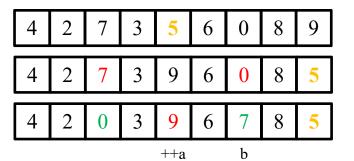


```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick

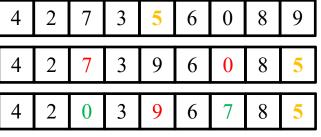


```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



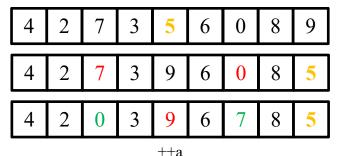
```
++a --b
```

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



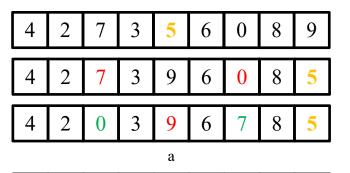
```
--b
```

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T, class P> void quicksort(T s[], int iL, int iR, int d){ // [iL, iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
 // END quick sort: divide & conquer
```





- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick



a

```
5 0 / 8 9
```

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T.class P> void quicksort(T s[].int iL.int iR.int d){ // [iL.iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```



- **Quicksort divide & conquer**
- pivot partition in partially-sorted form (PSF)
- partition trick
- do nothing for small

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
       if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<I,P>(s,iL,iR);else iP=findpivot<I>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T,class P> void quicksort(T s[],int iL,int iR,int d){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```



- Quicksort divide & conquer
 - pivot partition in partially-sorted form (PSF)
 - partition trick
 - do nothing for small
 - final insertion sort

```
// quick sort: divide & conquer
#define QUICKSORT SMALL T 3
#define M OF 3 0
template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
template <class T,class P> inline int findpivot(T s[],int iL,int iR){
        int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
        else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T) return; // do nothing for small sub-sequences
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a);quicksort<T,P>(s,a+1,iR);}
template <class T,class P> void quicksort(T s[],int iL,int iR,int d){ // [iL,iR)
        if((iR-iL)<=QUICKSORT SMALL T || d<=0) return;</pre>
        int iP;if(M OF 3) iP=findpivot<T,P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
        swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
        do\{while((a < b)\&P::p(s[++a],s[iP]));while((a < b)\&P::p(s[iP],s[--b]));
                swap(s,a,b);}while(a<b);</pre>
        swap(s,a,iP);quicksort<T,P>(s,iL,a,d-1);quicksort<T,P>(s,a+1,iR,d-1);}
template <class T,class P> void quicksort(T s[],int n){
        quicksort<T,P>(s,0,n);
        // take advantage of the best-case performance of insertion sort by a
        // single final call to insertion sort to process the entire array
        insertionsort<T,P>(s,n);}
  END quick sort: divide & conquer
```



unsorted sequence: 42 92 96 79 93 4 85 66 68 76 74 63 39 17 71 3

quick sort: 3 17 39 4 42 79 85 66 68 76 74 63 96 92 71 93

quick sort: 3 4 39 17 42 71 63 66 68 76 74 79 96 92 93 85 quick sort: 3 4 39 17 42 68 63 66 71 76 74 79 85 92 93 96

DEMO : quick sort =>

find pivot (median of three) => 42



cout<<"DEMO : quick sort =>\n";

Quicksort - divide & conquer

- pivot partition in partially-sorted form (PSF)
- partition trick
- do nothing for small
- final insertion sort

```
// quick sort: divide & conquer
                                                                  #define QUICKSORT SMALL T 3
                                                                  #define M OF 3 0
                                                                  template <class T> inline int findpivot(T s[],int iL,int iR){return (iL+iR-1)/2;}
                                                                  template <class T,class P> inline int findpivot(T s[],int iL,int iR){
                                                                          int iM=(iL+iR-1)/2; if(P::p(s[iL],s[iM])==P::p(s[iM],s[iR-1])) return iM;
                                                                          else if(P::p(s[iM],s[iL])==P::p(s[iL],s[iR-1])) return iL;else return iR-1;}
                                                                  template <class T,class P> void quicksort(T s[],int iL,int iR){ // [iL,iR)
                                                                          if((iR-iL)<=OUICKSORT SMALL T) return; // do nothing for small sub-sequences
                                                                          int iP;if(M OF 3) iP=findpivot<T.P>(s,iL,iR);else iP=findpivot<T>(s,iL,iR);
                                                                          swap(s,iP,iR-1);iP=iR-1;int a=iL-1,b=iP;
                                                                          do{while((a<b)&P::p(s[++a],s[iP])); while((a<b)&P::p(s[iP],s[--b]));</pre>
                                                                                  swap(s,a,b);}while(a<b);</pre>
                                                                          swap(s.a.iP):guicksort<T.P>(s.iL.a):guicksort<T.P>(s.a+1.iR):
unsorted sequence: machine intelligence system automation program technique computer data
```

```
DEMO : quick sort =>
unsorted sequence: 42 92 96 79 93 4 85 66 68 76 74 63 39 17 71 3
find pivot (simply middle) => 66
quick sort: 42 17 39 63 3 4 66 93 68 76 74 79 96 92 71 85
quick sort: 3 17 4 39 42 63 66 71 68 76 74 79 96 92 93 85
quick sort: 3 17 4 39 42 63 66 68 74 76 71 79 85 92 93 96
quick sort (with final insertionsort): 3 4 17 39 42 63 66 68 71 74 76 79 85 92 93 96
unsorted sequence: machine intelligence system automation program technique computer da
quick sort: automation computer data intelligence machine program system technique
```

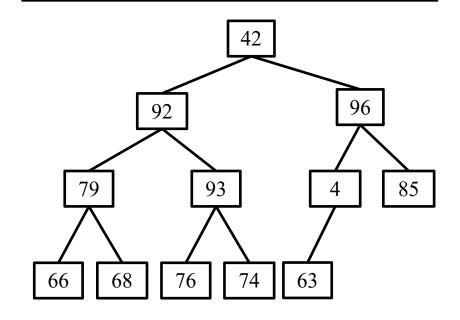
quick sort (with final insertionsort): 3 4 17 39 42 63 66 68 71 74 76 79 85 92 93 96

quick sort: automation computer data intelligence machine program system technique

```
cp<int>(iA,iTab,ni);cp<cptr>(cA,cTab,nc);
cout<<"unsorted sequence: ";show<int>(iA,ni);
int iP;if(M OF 3) iP=findpivot<int,IntPriorMin>(iA,0,ni);else iP=findpivot<int>(iA,0,ni);
if(M OF 3) cout<<"find pivot (median of three) => "<<iA[iP]<<endl;</pre>
else cout<<"find pivot (simply middle) => "<<iA[iP]<<endl;</pre>
cout<<"quick sort: ";quicksort<int,IntPriorMin>(iA,0,ni,1);show<int>(iA,ni);
cout<<"quick sort: ";cp<int>(iA,iTab,ni);quicksort<int,IntPriorMin>(iA,0,ni,2);show<int>(iA,ni);
cout<<"quick sort: ";cp<int>(iA,iTab,ni);quicksort<int,IntPriorMin>(iA,0,ni);show<int>(iA,ni);
cout<<"quick sort (with final insertionsort): ";quicksort<int,IntPriorMin>(iA,ni);show<int>(iA.ni);
cout<<"unsorted sequence: ";show<cptr>(cA,nc);
cout<<"guick sort: ":guicksort<cptr.CharsPriorMin>(cA.nc):show<cptr>(cA.nc):
```



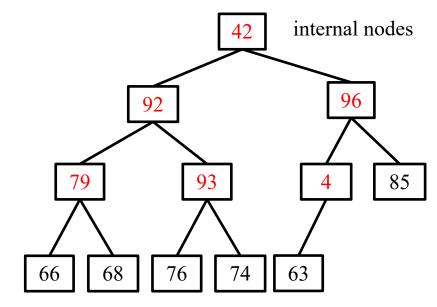
- Heap sort
 - batch initialization
 - iterative removeroot





Heap sort

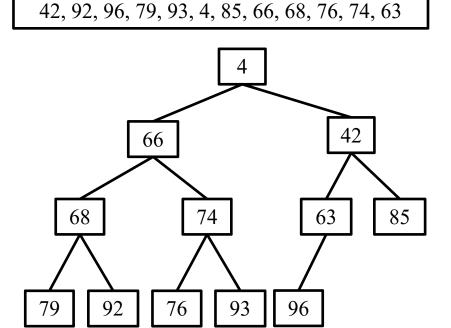
- batch initialization
 - backward iterated siftdown
- iterative removeroot





Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

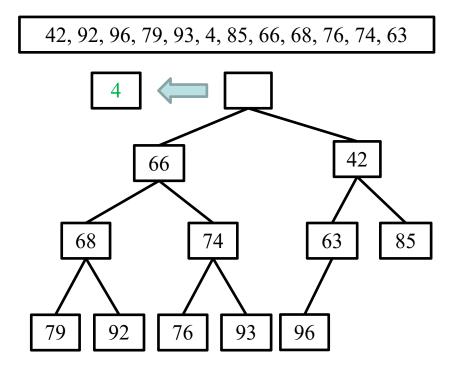




Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

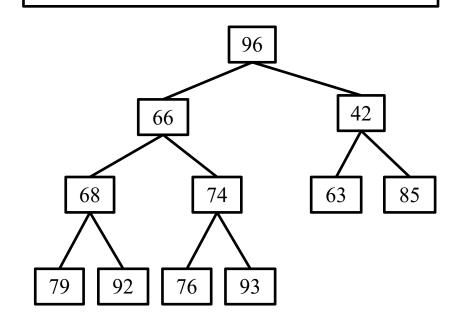




Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

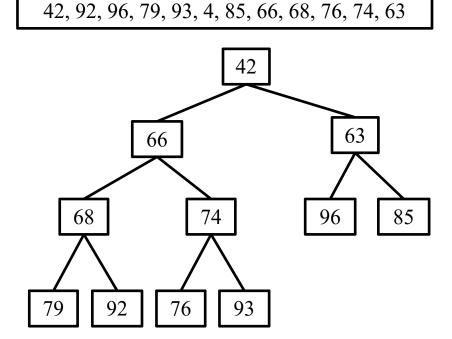




Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot



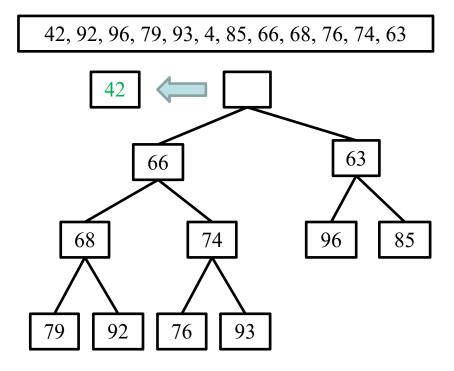


42

Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

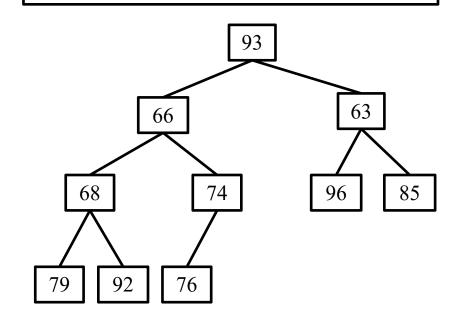




42

Sorting

- Heap sort
 - batch initialization
 - backward iterated siftdown
 - complexity: O(n)
 - iterative removeroot

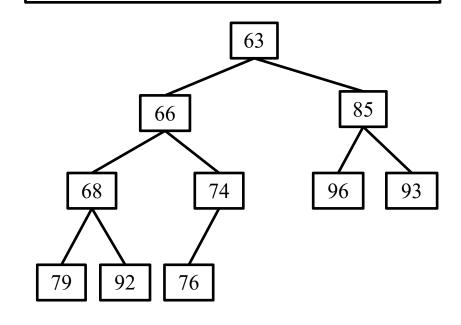




42

Sorting

- Heap sort
 - batch initialization
 - backward iterated siftdown
 - complexity: O(n)
 - iterative removeroot





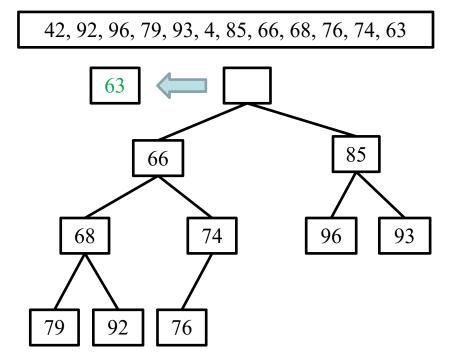
42

63

Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





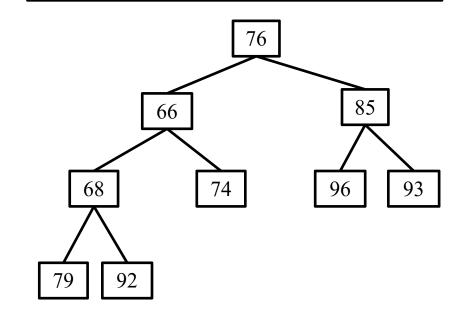
42

63

Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





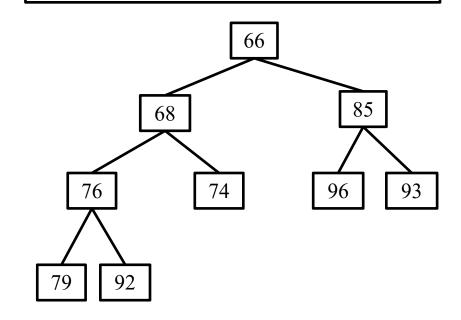
42

63

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot



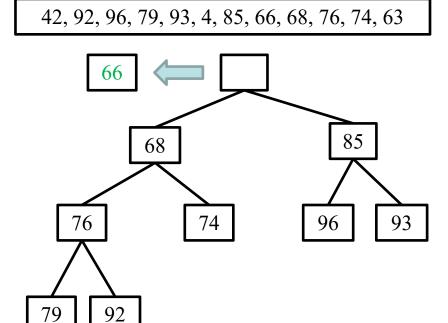


42

6366

Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





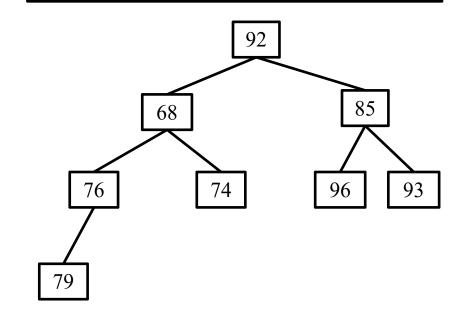
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6366

Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





42

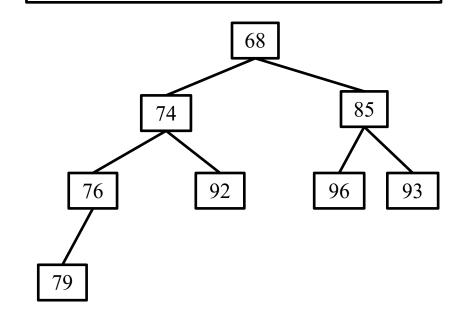
63

66

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





42

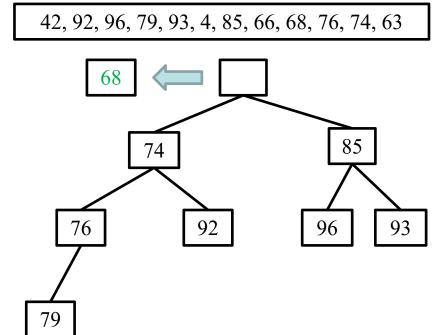
66

63

68

Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





4263

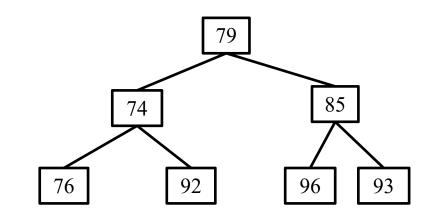
66

68

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





42 63

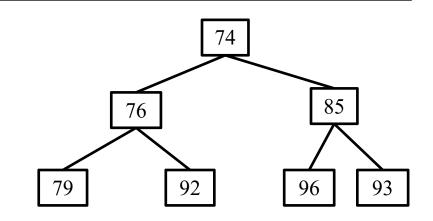
66

68

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





42

63

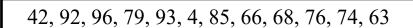
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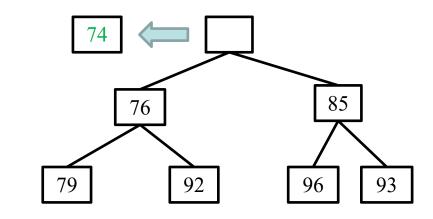
68

74

Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot







42

63

66

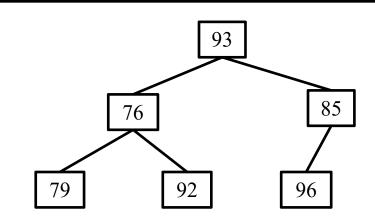
68

74

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





42

6366

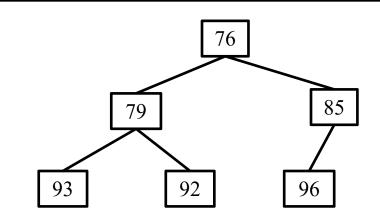
68

74

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





63

66

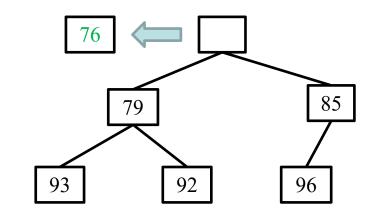
6874

76

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot





63

42

66

68

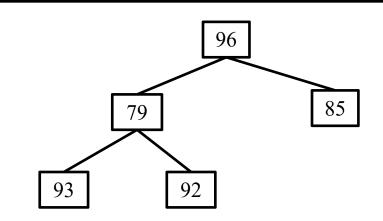
74

76

Sorting

• Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot



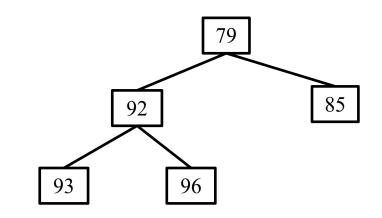


76

Sorting

Heap sort

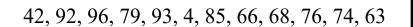
- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

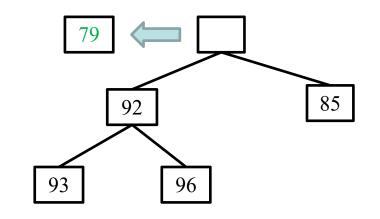




Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot



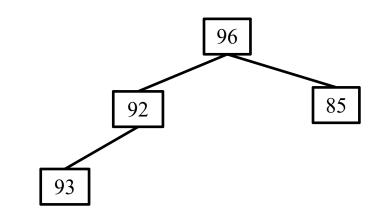




Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

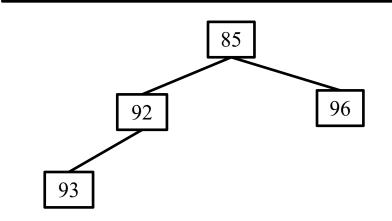




Sorting

Heap sort

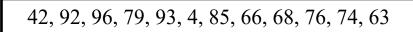
- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

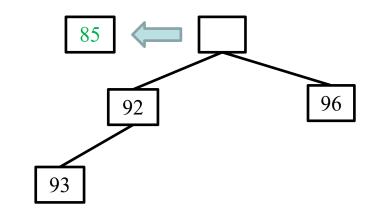




Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot



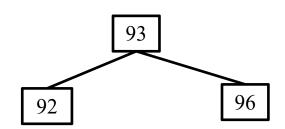




Sorting

Heap sort

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

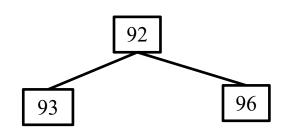




Sorting

Heap sort

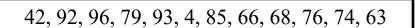
- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

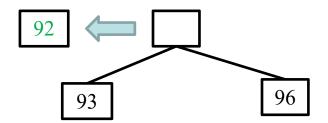




Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

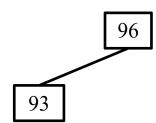






Sorting

- Heap sort
 - batch initialization
 - backward iterated siftdown
 - complexity: O(n)
 - iterative removeroot

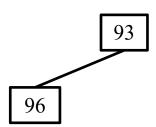




Sorting

Heap sort

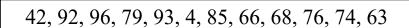
- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot

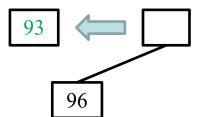




Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot







Sorting

- Heap sort
 - batch initialization
 - backward iterated siftdown
 - complexity: O(n)
 - iterative removeroot

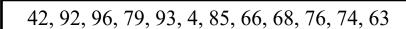
42, 92, 96, 79, 93, 4, 85, 66, 68, 76, 74, 63

96



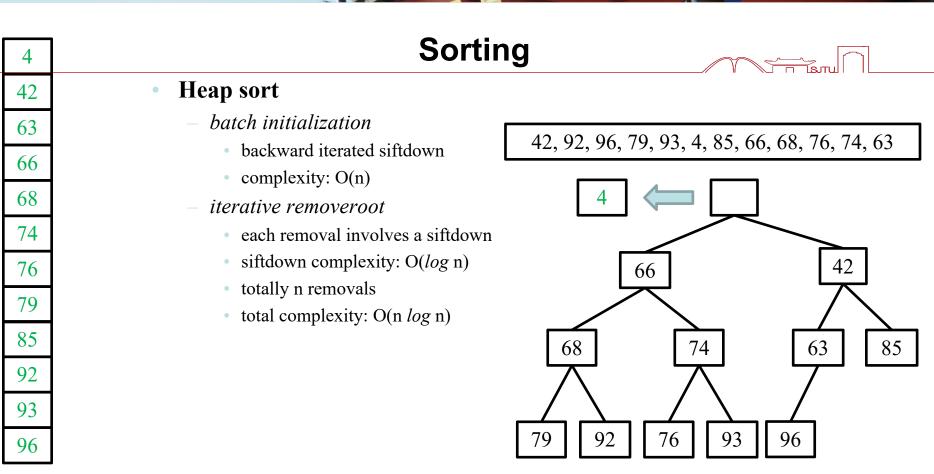
Sorting

- batch initialization
 - backward iterated siftdown
 - complexity: O(n)
- iterative removeroot











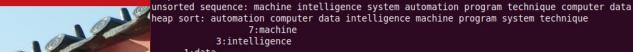
- Heap sort
 - batch initialization
 - iterative removeroot

```
heap sort
                           template <class T,class P> void heapsort(T s[],int n){
                                    Heap<T,P> h(n,s,n);
                                    for(int i=0;i<n;i++) s[i]=h.removeroot();}</pre>
                              END heap sort
cout<<"DEMO : heap sort =>\n";
```

```
cp<int>(iA,iTab,ni);cp<cptr>(cA,cTab,nc);
          cout<<"unsorted sequence: ";show<int>(iA,ni);
          cout<<"heap sort: ";heapsort<int,IntPriorMin>(iA,ni);show<int>(iA,ni);
          cout<<"unsorted sequence: ";show<cptr>(cA,nc);
          cout<<"heap sort: ";heapsort<cptr,CharsPriorMin>(cA,nc);show<cptr>(cA,nc);
          cp<cptr>(cA,cTab,nc);Heap<cptr,CharsPriorMin> h(nc,cA,nc);h.S();
          for(int i=0;i<nc;i++){cA[i]=h.removeroot();cout<<"heap sort: ";show<cptr>(cA,i+1);h.S();}
DEMO : heap sort =>
```

```
unsorted sequence: 42 92 96 79 93 4 85 66 68 76 74 63 39 17 71 3
heap sort: 3 4 17 39 42 63 66 68 71 74 76 79 85 92 93 96
unsorted sequence: machine intelligence system automation program technique computer data
heap sort: automation computer data intelligence machine program system technique
```





Heap sort

- batch initialization
- iterative removeroot

```
heap sort
template <class T,class P> void heapsort(T s[],int n){
           Heap<T,P> h(n,s,n);
           for(int i=0;i<n;i++) s[i]=h.removeroot();}</pre>
    END heap sort
cout<<"DEM0 : heap sort =>\n";
cp<int>(iA,iTab,ni);cp<cptr>(cA,cTab,nc);
cout<<"unsorted sequence: ";show<int>(iA,ni);
cout<<"heap sort: ";heapsort<int,IntPriorMin>(iA,ni);show<int>(iA,ni);
cout<<"unsorted sequence: ";show<cptr>(cA,nc);
cout<<"heap sort: ";heapsort<cptr,CharsPriorMin>(cA,nc);show<cptr>(cA,nc);
cp<cptr>(cA,cTab,nc);Heap<cptr,CharsPriorMin> h(nc,cA,nc);h.S();
for(int i=0;i<nc;i++){cA[i]=h.removeroot();cout<<"heap sort: ";show<cptr>(cA,i+1);h.S();}
```

```
2:computer
            6:system
heap sort: automation
            3:intelligence
      1:data
            4:program
0:computer
            5:technique
      2:machine
            6:system
heap sort: automation computer
            3:system
      1:intelligence
            4:program
0:data
            5:technique
      2:machine
heap sort: automation computer data
            3:system
      1:program
            4:technique
0:intelligence
      2:machine
heap sort: automation computer data intelligence
            3:system
      1:program
0:machine
      2:technique
heap sort: automation computer data intelligence machine
      1:system
```

7:machine 3:intelligence

4:program

5:technique

1:data

0:automation

0:program

0:svstem

0:technique

Heap is empty!

2:technique

1:technique

heap sort: automation computer data intelligence machine program

heap sort: automation computer data intelligence machine program system

heap sort: automation computer data intelligence machine program system technique



Sorting complexity

insertion sort	$W:O(n^2)$	$A:O(n^2)$	B:O(n)
----------------------------------	------------	------------	--------

- bubble sort
$$W:O(n^2)$$
 A: $O(n^2)$ B: $O(n^2)$

- selection sort
$$W:O(n^2)$$
 A:O(n²) B:O(n²)

- shell sort (depending on *shell sequence*)
$$O(n (log n)^2) \sim O(n^{1.5})$$

mergesort	$W:O(n \log n)$	A:O(n log n)	B:O(n log n)
		(()	\ (3 /

- BST sort
$$W:O(n^2)$$
 A:O(n log n) B:O(n log n)

- quicksort
$$W:O(n^2)$$
 A:O(n log n) B:O(n log n)

- heap sort
$$W:O(n \log n)$$
 A: $O(n \log n)$ B: $O(n \log n)$



Sorting complexity

insertion sort	$W:O(n^2)$	$A:O(n^2)$	B:O(n)
----------------------------------	------------	------------	--------

- bubble sort
$$W:O(n^2)$$
 A:O(n²) B:O(n²)

- selection sort
$$W:O(n^2)$$
 A:O(n²) B:O(n²)

- shell sort (depending on *shell sequence*) O(n
$$(log n)^2$$
) ~ O(n^{1.5})

mergesort	$W:O(n \log n)$	A:O(n log n)	$B:O(n \log n)$
11101 20001 t	W.O(II 102 II)		$\mathbf{D} \cdot \mathbf{O} \in \mathbf{H}$

- BST sort
$$W:O(n^2)$$
 A: $O(n log n)$ B: $O(n log n)$

- quicksort
$$W:O(n^2)$$
 A:O(n log n) B:O(n log n)

- heap sort $W:O(n \log n)$ A: $O(n \log n)$ B: $O(n \log n)$

Reflection

- why is quicksort called "quick"? why is not mergesort or heap sort so called?
- BST sort seems not so bad, why is BST sort hardly used?



Sorting complexity

- insertion sort $W:O(n^2)$ A:O(n²) B:O(n)
- bubble sort $W:O(n^2)$ A:O(n²) B:O(n²)
- selection sort $W:O(n^2)$ A:O(n²) B:O(n²)
- shell sort (depending on *shell sequence*) $O(n (log n)^2) \sim O(n^{1.5})$
- mergesort $W:O(n \log n)$ A:O(n log n) B:O(n log n)
- BST sort $W:O(n^2)$ A:O(n log n) B:O(n log n)
- quicksort $W:O(n^2)$ A:O(n log n) B:O(n log n)
- heap sort $W:O(n \log n)$ A:O(n log n) B:O(n log n)

Reflection

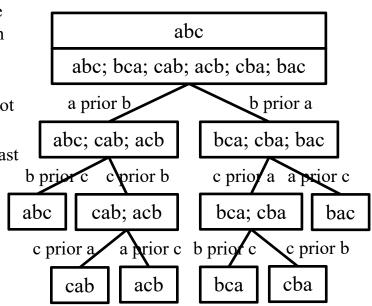
- why is quicksort called "quick"? why is not mergesort or heap sort so called?
- BST sort seems not so bad, why is BST sort hardly used?

not only complexity level matters, but also does concrete implementation (computation & memory)



Lower bounds for sorting

- worst-case complexity: \ge O(n *log* n)
- perspective of *decision tree*
 - each comparison (& potential swap) can be regarded as decision to distinguish between two branches of possible permutations
 - complexity of sorting a permutation equals the depth of decision that leads from the root to the permutation (leaf)
 - the decision tree must have a depth of at least log(n!) to accommodate all n! possible permutations
 - $log(n!) \approx n log n$





THANK YOU

