韩子坚

华中师范大学计算机学院

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- 环境说明
- 2 前期准备
- 3 安装运行
- 4 其他尝试
- 5 对话实例

环境说明 •0000

环境说明 00000

硬件环境

环境说明

- **CPU**: AMD EPYC 7742 64-Core Processor * 2
- **GPU**: NVIDIA A100-SXM4-80GB * 8
- **RAM**: 1.5 TB



环境说明

环境说明 00000

- 软件环境



环境说明 0000

软件环境

- **OS**: Ubuntu 22.04.5 LTS
- Container: Docker 26.1.3
- LLM Backend: Ollama 0.5.7 (Python Version)
- LLM Chat UI: Open WebUI 0.5.7 (Docker Version)



- 前期准备

安装 Ollama + Open Webui

- 前期准备

安装 Docker

前期准备 0.000000000000



前期准备 0000000000000

- 如果是一个空白机器, 并且具有 Docker 权限, 推荐直接使用 Docker 版本的 Ollama + Open Webui.
- 有 sudo 权限的用户, 推荐使用 1panel 的安装脚本。
- 安装 1panel 的同时会自动安装 Docker + Docker Compose。
- 中国大陆境内机器使用 1panel 的安装脚本会自动使用镜像地址安装,免去网络 的烦恼。

curl -sSL https://resource.fit2cloud.com/1panel/package/ quick_start.sh -o quick_start.sh && sudo bash quick_start.sh



sudo sh get-docker.sh

本地部署 DeepSeek-R1-671B-q4 K M 完全指南

安装 Docker (续)

• 如果不喜欢使用 1panel 或没有面板安装权限, 可直接安装 Docker。

```
export DOWNLOAD URL="https://mirrors.tuna.tsinghua.edu.cn/docker-
   ce" # 非中国大陆区域可取消这行
curl -fsSL https://get.docker.com -o get-docker.sh
```



前期准备

安装 Ollama + Open Webui

前期准备 0000000000000





安装 Ollama + Open Webui

前期准备

• 如果已有 Ollama, 必须升级到最新版本, 否则可能不支持 Deepseek R1 架构, 升级脚本与安装脚本是相同的。

curl -fsSL https://ollama.com/install.sh | sh



安装 Open Webui

前期准备 00000000000000

- Open Webui 不建议使用 pip 安装, 升级 pip 安装的 Open Webui 出现过数据丢失 的情况, 建议使用 Docker 安装。
- 如果服务器本身处在 Docker 容器内,或者无法处理后面的网络问题,也可以使 用 pip 安装,建议使用 uv。

安装 Open Webui - uv pip 方式

前期准备 00000000000000

- 适用干服务器本身在 Docker 容器内或无法解决网络问题的场景。
- 推荐使用 uv pip 安装。

uv 不会继承 pip.conf 以及其他 pip 中设置镜像地址的方法, 需要另外设置 uv 镜像 (非中国大陆地区可以跳过)

```
# vim ~/.config/uv/uv.toml
[[index]]
url = "https://pypi.tuna.tsinghua.edu.cn/simple"
default = true
```

安装 Open Webui - uv pip 方式(续)

前期准备 00000000000000

```
# 安装 uv
pip install uv
# 创建虚拟环境
uv venv --python python3.11
#激活虚拟环境
source .venv/bin/activate
# 安装 Open Webui
uv pip install open-webui
```

安装 Open Webui - Docker 方式

• 使用 Docker 方式安装 Open Webui。

```
docker run -d -p 8080:8080 -e HF_ENDPOINT=https://hf-mirror.com/
--add-host=host.docker.internal:host-gateway -v open-webui:/app
/backend/data --name open-webui --restart always swr.cn-north
-4.myhuaweicloud.com/ddn-k8s/ghcr.io/open-webui/open-webui:
latest
```

注意: Open Webui 的 Docker image 不在 Docker 官方 hub.docker.com 中,而是在ghcr.io 中,中国大陆境内无法访问 hub.docker.com,寻常 Docker 镜像站都是镜像于官方 hub.docker.com,可以访问渡渡鸟镜像站: https://docker.aityp.com/查询 Open Webui 的镜像,也可以使用南京大学的 ghcr 同步站: ghcr.nju.edu.cn(实测限速)



Docker 网络配置

前期准备 0000000000000000

- 桥接模式下,容器访问宿主机网络可以使用 --add-host=host.docker.internal:host-gateway.
- 但这样并不总是可行,例如机器网络比较复杂或者 Docker 版本比较低等等原 因,可以这样做(前提是已经设置OLLAMA HOST=0.0.0.0):



Docker 网络配置 - 访问宿主机网络(续)

查看容器所属网络的网关地址

```
| docker inspect -f '{{range .NetworkSettings.Networks}}{{.Gateway }}{{end}}' <容器名称或ID>
```

2 # 例如: 172.17.0.0/

容器内部测试访问

```
1 curl http://172.17.0.0/: 宿 主 机 端 口
```

• 注意: 但是这样也未必可联通, 很有可能会被防火墙阻挡, 需要添加放行规则



Docker 网络配置 - 防火墙放行

前期准备

添加防火墙放行规则(示例)

sudo ufw allow proto tcp from 172.17.0.0/16 to any port 11434 sudo ufw reload



Docker 网络配置 - 容器互联

• 如果 Ollama 和 Open Webui 分属不同 Docker 容器,可以将它们添加到同一个网络,通过容器名:端口互联。

创建 Docker 网络并连接容器

```
docker network create llm
```

- 2 | docker network connect llm < ollama_id >
 - |docker network connect llm < openwebui_id >

- 3 安装运行 使用 Open Webui 对话

- 3 安装运行 运行 Ollama 使用 Open Webui 对话

运行 Ollama

启动 Ollama 服务

```
export OLLAMA MODELS=/usr/share/ollama/.ollama/models
export OLLAMA HOST=0.0.0.0
export OLLAMA SCHED SPREAD=1
export CUDA VISIBLE DEVICES=0,1,2,3,4,5,6,7
nohup sh -c 'ollama serve' > ollama.log 2>&1 &
```

• 注意:

- OLLAMA HOST 必须设置为 0.0.0.0,否则后续 Open Webui 无法访问 Ollama。
- OLLAMA SCHED SPREAD 必须设置为 1. 否则无法多卡运行单一模型。



拉取 DeepSeek R1

使用 Ollama 拉取 DeepSeek R1 模型

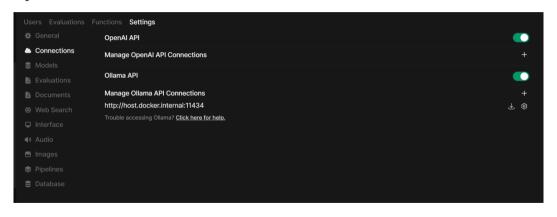
ollama pull deepseek-r1:671b # 默认就是 q4_K_M 量化



- 1 环境说明
- 2 前期准备
- 3 安装运行运行 Ollama使用 Open Webui 对话
- 4 其他尝证
- 5 对话实例

|使用 Open Webui 对话

Open Webui 后端地址设置



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Open Webui 配置详解

- Open Webui 配置分三层:模型级,账户级,聊天级。
- 优先级: 模型级 > 账户级 > 单次聊天级。
- **单次聊天级 (Per-Chat)**: 聊天右侧 Chat Controls 中设定,仅对当前会话生效,不能覆盖模型预设。
- 模型级 (Per-Model): Admin Panel-Settings-Models 中设定,适用于所有使用该模型的聊天,优先级最高。
- 账户级 (Per-Account): Settings-Advanced Parameters 中设定,会被模型级配置覆盖。

Open Webui 调整参数

- 8卡 A100 在设定 Context Length 和 Max Tokens 都为 8k 后可正常推理, 每张卡显存占用约为 50-60GB.
- 默认情况下,5min 内如果没有新的对话,模型会从GPU中 offload,重新加载至显存需要很长时间。可以在Settings-Advanced Parameters 中设定 Keep Alive参数,设为-1则用不卸载。
- 如果显存不够,可通过设置 num_gpu 参数来灵活调整 CPU/GPU 推理比例。
 例如设为 60 则 60 层加载到 GPU、剩余在 CPU 推理。
- 完成上述配置后,即可在 Open Webui 中与 DeepSeek R1 模型进行对话。



- 4 其他尝试

使用 Ollama 并不是最优选择, 无论是速度还是显存占用都是不佳的。

OLLAMA_SCHED_SPREAD=1 本来也是不建议开启的, Ollama 只能用单一的 gguf 格式 的模型参数进行推理, 但是对于 671 B 这么大的模型来说, 全部加载在同一张 GPU 上是不太现实的,只能分担在多张 GPU 上, Ollama 对于一个模型在多卡中推理优 化较差,这样做增加了 GPU 之间通信的成本,即使所有参数全部在 GPU 中推理, 速度也只有 10-40 tokens/s。

本地部署 DeenSeek-R1-671B-q4 K M 完全指南

- vllm 可能是比 Ollama 更优的选择。
 - 原生支持 safetensor 格式。
 - 但是 671B fp8 太大了, 8 卡 a100 也装不下。
- 考虑使用量化模型以降低显存需求。
 - vllm 官 四中有一句话, Please note that GGUF support in vLLM is highly experimental and under-optimized at the moment, it might be incompatible with other features.
 - 尝试 1.58 bit 动态量化版本报错 (ValueError: GGUF model with architecture deepseek2 is not supported yet.)。
 - INT4 W4A16 量化可能是可行的方案 (未尝试)。



- 拉取 1.58 bit 动态量化模型时. hf mirror 速度慢且易断连。
- modelscope 限速在 15MB/s 左右。
- 对于 671B 这么大的模型,可以使用 snapshot download 下载部分文件夹。

```
from huggingface_hub import snapshot_download
snapshot_download(repo_id='unsloth/DeepSeek-R1-GGUF'.
   allow patterns='DeepSeek-R1-UD-IQ1 S/*', cache dir='./')
 modelscope
from modelscope.hub.snapshot_download import snapshot_download
model_dir = snapshot_download(repo_id='unsloth/DeepSeek-R1-GGUF',
   allow_patterns='DeepSeek-R1-UD-IQ1_S/*',cache_dir='./')
```

6

- 1 玩诗兴明
- 2 前期准备
- 3 安装运行
- 4 其他尝证
- 5 对话实例

2024 AIME I Problems/Problem 4

2025 考研数学-T18

- 环境说明
- 2 前期准备
- 3 安装运行
- 4 其他尝试
- 5 对话实例

2024 AIME I Problems/Problem 4

2025 考研数学一T18



2024 AIME I Problem 4

grand prize if all four of her numbers were the randomly chosen numbers. The probability of her winning the

Ol deepseek-r1:671b

"Okay, so Jen plays this lottery by picking 4 distinct numbers from the set S = (1, 2, ..., 10). Then, 4 numbers are randomly chosen from S as well. She wins a prize if at least two of her numbers match the drawn numbers, and she gets the grand prize if all four match. The question is asking for the probability that Jen wins the grand prize given that she won at least some prize (i.e., at least two matches). They want this probability expressed as a reduced fraction m/n, and then we need to find m + n.

First, let me make sure I understand conditional probability here. So, we're looking for P(Grand Prize | At Least Two Matches), which by definition is equal to the probability of both events happening divided by the probability of the conditioning event (wholing at least a prize). But since if you win the grand prize, you have automatically won at least two matches (in fact, all four), so this simplifies to PIGrand Prize) / PIALL east Two Matches)

1. Calculate the total number of possible draws: Since they are choosing 4 numbers from 10, that's C(10,4). Let me confirm what notation we use here. I think it's combination, right? So C(n,k) = nl / (k)(n - k))).

So total possibilities for the draw is C(10.4), which should be 210.

2. Then compute P(Grand Prize): That's just 1 divided by this number because there's only one winning combination that matches all four of her numbers. So probability is 1 / 210.

Wait but hold on: Wait. Jen picks her own four numbers first. But the problem says "4 numbers are randomly chosen from S," Does she pick hers first? Yes, and then if they match... so actually, the total number of possible draws is C(10.4) = 210. The probability that all four of her numbers were drawn would be how many combinations contain exactly those four numbers: which is just 1 (assuming Jen's numbers are fixed). So yeah, P(Grand Prize) = 1 / 210.

3. Then compute PIAt Least Two Matches): This is the sum over k=2 to 4 of C(4,k)*C(6,4 -k), divided by total number of possible draws C(10.4). Because she has 4 numbers picked: if we want exactly k matches, that's choosing k correct numbers from her 4 and (4 - k) incorrect numbers from the remaining 6. Then sum this for k=2 to 4.

Let me verify:

For exactly 2 matches: C(4,2)*C(6,2), Because choose 2 of Jen's numbers and 2 others Exactly 3 matches: C(4,3)*C(6,1) Evently 4 matches: CI4 4)*C(6.0) = 1

So total favorable for at least two is sum from kii 2 to 4 of the above. Let me compute each

First, calculate combinations

For exactly 2: C(4,2)=6



C(6,2)=15 Multiply: 90

| 2024 AIME I Problem 4 (续)

Exactly 3: C(4.3) = 4C(6.1) = 6 Multiply:24 Evantly 4: So total is 90 + 24 + 1 = 115. Therefore PIALLeast Two Matches)=115 / C(10.4) But C(10.4) is 210. So that's 115/210. Let me check if this fraction reduces: both numerator and denominator divided by 52 Wait 215 + 5x23x5; wait no Wait 105 divides into 21, right? But first let's see GCD of 114 (wait no. 90 +24 is 14(?) No. The total was 115. So yes: So 115 and 210: Divide numerator and denominator by 57 Let's check if both are divisible. 115 +5=23; 210+5=42. Yes! So 2 divides into 230 / 4207 Wait, sorry, simplifying 115/210 divided by GCD(115.210). The factors of 115: 1,5,23,115 Factors of 210:1,2,3,5,6,7,10,..., etc. Common factor is 5. So divide numerator and denominator by 5: 23 /42. So 23/42 Therefore PIAt Least Two!=23/42? Wait wait but hold on: Wait let's check appin: CI4 21 C/6 2J =615=90 C(4.3) C/6 f/=4 6=24 C/4 4)*C/6 0)=1 Total is 115 Total possible draws are C(10.4)=210. So 115/210 simplifies to dividing numerator and denominator by 5:115 +5 =23: 210 +5=42. So yes, 23/42 =0.547. So B(Grand Prize | At Least Two) is 11 (210) divided by [23 JA2]. Which equals (1 (210) (42 (23) s (42) (23023). Simplify: 42 and 210 share a factor of 4221 at's see Weit, actually Let me compute numerator/denominator step-wise (1/210) + (23/42) = (1/210) * (42/23) Simplify this fraction. So the 420 in denominator cancels with the 41? Wait no. wait: The 21 in 120? Let's think about it. Multiplying numerator: 1+42=42

Then remaining: 750?

No. let me do straight division:

No. For example.

How many time does 7 go into 7007:100+7=700. Then remainder 5. So no: so GCD(7,805)=7.

Walt actually: Let's compute 735 +7=105 (since 70×10 +35). But then 705 =7+100 +35. Therefore 700+0 and 5?

Next digit5 bring down; makes35, 7 divides into35 five times. So total is115. Therefore, 805=7.115. Thus GCD/7805/=7.

First divide 80 by7 which gives 11 with remainder3, bring down the five: 30 +5=35. Then 355? Wait no:

Recause 800/77 Not exact. Wait sorry, check if 8 times is:

Wait wait 805 divided by 7 is as follows:

Start from left digit (8), 7 goes into 8 once, write 1. Remainder 1.
Bring next digit 0: making into 7 into 10 once, remainder 3.

2024 AIME I Problem 4 (续)

Divide numerator and denominator by 6? Let's see if that works. 42 +6=7 4830 +6=805 Then we have 7 /805 Check GCD of 7 and 805: Since 5 divides into 800 so 80 is divisible by 5 but 8 +0+5=13 not sure. Wait 85 divided by 752 Wait no But 710 over here? No well 807 But ofter divide 4830 +6 +805 So 7 and 805 Since 735 is 7+105: so 705+700 +5, but 704/. Alternatively, check if 785 divided by 7: 800 +7=114.28, which doesn't fit into 784 (7×112). But wait 80(5) =70×11.5? Not integer. 805 divide by 7: 700=100 2: 105=157 so total 115, so yes! Wait 705 divided by 7? No sorry, in denominator after dividing 4830/6 =805. So check if 5 divides into that? 8+0+5=13 not divisible by 3. But does 800 +5 is 80+10 +5 which ends with a 5 so yes Wait no: Wait 70 over here, sorry: So after dividing numerator and denominator by 6 we get 7 /805 Then check if GCD(7805): Since 700 =100 ×7: then 705=700+5. So 8007 No. Hold on: Compute 8 times 10 is 80. Wait 805 divided by7: Divide 70 into each part: 735 +7=105





2024 AIME I Problem 4 (续)

Simplify fraction: [7/805] = [1 /115]. Because divide numerator and denominator by?.

Thus the probability of winning grand prize given that she won a prize is Y₁₀. So m=1, n=115, which are coprime.

But wait this contradicts something? Let me double-check steps because I might have messed up.

Wait:

Probability (Grand Prize)/ Probability/At Least Two Matches) = [1/210] / [23/42]. Which is equal to (1+42)/(210 *23) = 42/(4830). Simplify that. Then divide numerator and denominator by 6:7/805, then check if OCD/7,805): As above, since 8 times? Wall 1.

Wait wait, in previous thought, I realized 805 divided by7=115. So 4×100 is...?

No need to confuse here.

So the fraction after dividing by6 was7 /805. Then recognize that both are divisible by7: so divide numerator and

Numerator becomes1 (since7 +7 =1), Denominator becomes 115 (805 +7=115). So yes, simplifies to1/115.

Therefore m/n is/is which is reduced as coprime because gcd(\115)=1. Therefore answer is1 +115=116.
But wall the problem says "the probability of her winding the grand prize given that she won a prize isn'n
wheremandane relatively prime positive integers. Findin =n."

Hence m =1,n=115? That leads to sum 116?

Wait but let me check if my calculations are correct again because sometimes these combinatorial problems can be

First, total number of possible draws: C(10,4)=210. Correct.

Number of grand prize winners: only one set (assuming Jen's numbers are fixed). So probability is 1/210. That seems right.

Then the number of ways to have at least two matches: Compute for k=2: C(4.2) C(6.2) = 615±90

k=3:C(4,3)C(6,1)=4+6=24,

k=4:1. Total favorable is115.

Therefore probability of at least two matches:115/210 simplifies to23/42 (divided by5)

Then conditional probability [1 /210] divided by[23/42].

Which is equal to (1 * 42)/(210 *23) *42/(4830). Simplify, divide numerator and denominator by6 gives7/805. Then divide by7:1/115.

So yeah m=1,n=115 sum116. Seems correct.

But let me check alternative approach using conditional probability via counting number of favorable outcomes over possible given the condition.

Alternative Approach:

Given that she already has won a prize (i.e., at least two matches), we need to compute how many such possible draws satisfy this and then among those, how many have all four matched.

satisfy this and then among those, now many have all four matched.

Total possibilities where she wins a prize:115 as computed. The grand prize is 1 of these outcomes because there's only

| 2024 AIME I Problem 4 (续完)

one way for all her numbers to be drawn. Hence the conditional probability would simply be 1/115 which matches previous result. So here we can see directly that once you condition on having at least two, how many possible draws are in that pool:115. The grand prize is1 of them, so indeed it's Vis. Therefore m=1 n=115: sum116. Hence the answer should helhoved/1161 Final Answer The value of m + n is \boxed(116)). Jen enters a lottery by cicking 4 distinct numbers from the set $S = \{1, 2, \dots, 10\}$. The probability of her winning the grand prize given that she won a prize (at least two matches) needs to be found. First, we calculate the total number of possible draws: The probability of winning the grand prize is: combinations for exactly 2, 3, and 4 matches: • For exactly 2 matches: $\binom{4}{5}\binom{6}{5} = 6 \times 15 = 90$ For exactly 3 matches: $\binom{4}{5}\binom{6}{5} = 4 \times 6 = 24$ 90 + 24 + 1 = 115The probability of winning at least a prize is: Simplifying this fraction by dividing both the numerator and denominator by their greatest common divisor (5): Using conditional probability, we find the probability of winning the grand prize given that she won at least a prize: Here, m=1 and n=115, which are relatively prime. Thus, m+n=1+115=116The final answer is:

输出情况

Listing 1: 2024 AIME I Problem 4 生成情况

```
1    response_token/s: 14.55
2    prompt_token/s: 54.51
3    total_duration: 302493858721
4    load_duration: 88603384116
5    prompt_eval_count: 124
6    prompt_eval_duration: 2275000000
7    eval_count: 3079
8    eval_duration: 211613000000
9    approximate_total: 0h5m2s
```



- 1 环境说明
- 2 前期准备
- 3 安装运行
- 4 其他尝试
- 5 对话实例

2024 AIME I Problems/Problem 4

2025 考研数学一 T18

2025 考研数学一T18



是一个常数。因此: $\frac{\partial^2 \alpha}{\partial x^2} = (1/y) + \frac{d}{dx} [f'(u)] = (1/y) / f''(u) \frac{\partial u}{\partial x} = (1/y) (f''(u) + 1/y) = f''(u) / (y^2)$ 接下来算混合导数ð²g/(ðxðy), 也就是先对x求导, 再对y求导: 或者也可以先对v求导、但按题目给出的顺序是先x后v。 原来的表达式是 $\partial/\partial y$ [$\partial g/\partial x$],已经知道 $\partial g/\partial x = f(u)/y$ 。现在对这个结果关于y求导数 首先写出表达式: a/av [f'(u) / v] 应用乘积法则或者商数法则。 图导数为 [d/dv (f'(u))] * (1/v) + f'(u)* (-1/v2) 计算第一部分: d/dy(f'(u))=f"(u)* ðu/ðy = f"(u)*(-x / y^2) 所以整体的混合导数量: $[f''(u)(-x/y^2)](1/y) + f'(u)(-1/y^2)$ $= -x/(y^2)f''(u) - f'(u)/y^2$ 接下来计算82g/8y2: 先求一阶号数 ðg/ðy = f'(u)*(-x / y*2) 再对v来导: $d/dv [f'(u)(-x/v^2)] = (-x)[d/dv (f'(u)/v^2)]$ 田亜印法制、得到: $(-x) [d/dy(f'(u)) * 1/y^2 + f'(u)*d/dy(1/y^2)]$ 首先计算d/dy(f'(u))=f'(u)*(-x / y^2) 然后 d/dy(1/y^2)= -2/y3 代入得到: $(-x)[(f''(u)(-x/v^2)) * 1/v^2 + f'(u)*(-2/v^2)]$ $= (-x) [(-x)f^{-}(u)/y^{4} + (-2 f'(u))/y^{2}]$ 展开后得 $(-x)(-x f'(u)/y^4) + (-x)(-2 f(u)/y^3)$ $= x^2 f'(u)/(y^5) + 2x f'(u)/y^4$ 不过可能我的计算有些错误,需要重新仔细算一遍:

首先写成: -x f'(u) y^{-2} 所以导数为 d/dy [-x f'(u) y*{-2}] 应用乘积法则,这里的变量是y,而x被视为常数。因此, $= -x [d/dy(f'(u)) * y^{-2} + f'(u)*d/dy(y^{-2})]$ 首先计算d/dy(f'(u))=f'(u) * du/dy = f''(u)*(-x/y2) ifii d/dy(y^{-2})= -2 y^{-3} = -x*(-x f"(u))/y4 - (-x)* (2 f'(u)) / y^3 = x2 f"(u)/(y^4 * y??)不对,这里的分母是四次吗? 仔细看看: 第一项里的(-x)乘上作的导数是-x[(f"*(-x/y^2))*y^(-2)]? 是的。所以第一个项的系数应该是: -f"* (-x / y2)* y^{-2}= f" *x/(y4) 然后事以-x的活、即整个是 -xf"(x)/y^4)= -x^2 f"(u)/(y^4)? 不,这里我需要仔细代入步程; 原式: -x [(-x f'(u))/ y2 * (1/y2) + f'(u)*(-2/y3)] 也就是将表达式代入后变成 $= -x*[(-x f^*(u))/(y^4) - 2 f^*(u)/y^2]$ 然后分配重法律: 世級提. $= (x^2 f''(u))/ y^4 + (2x f'(u))/(y^3)$ EECLA20/Av2 = [v2 f=(u)]/(v4) + [2v f*(u)]/(v2) ◆ロト ◆部 → ◆注 → 注 り へ ○

现在将这些代入原偏微分方程 左边是 $x^2 \partial^2 g/\partial x^2 + xy * \partial^2 g/(\partial x \partial y) + y^2 \partial^2 g/\partial y^2$ 分别计算每项。 首先,第一项: x²乘以ð²g/ðx²= x²*(f''(u)/y²) 即 (x2 / y2)f"(u) 第二项是xy乘以混合导数 之前求的混合导数是 -x/(v3) f"(u) - f'(u)/(v*2). 因此。 xy * [-x/(y3)f"(u)-f'(u)/y2] $= xy*(-x f'(u)/y^3) + xy*(-f'(u)/y^2)$ = (- x^2 v / v^3) f"(u)) - (xy f'(u))/ v^2 $= (-x^2/y^2)f''(u) - (x/y)f'(u)$ 第三项是y2乘以ð2g/ðy2、图。 y2 * [x2 f"(u)/y4 + 2x f'(u)/(y8)] $= v^{2+}(x^2 f''(u))/v^4 + v^{2+}(2x f'(u))/v^3)$ (P.85) = x2 f'(u) / y^2 + 2x f'(u)/ y 现在将三项相加: 左边 = [第一项] + [第二项] + [第三项] $(x^2/y^2)f'' + (-x^2/y^2)f'' - (x/y)f' + x^2/y^2f'' + 2x/(y) f'$ 先看前两项的第一部分: (x²/y²)f" - (x²/y²)f", 加上第三项的x²/y² f" 所以第一部分应该是: $(1)x^2/y^2$ f" $-1x^2/y^2$ f" $+1x^2/y^2$ f" = $x^2/(y^2)$ f" 然后是第二部分: - (x/y) f' + 2x/y f'= (+x/y)f'

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因此总和左边整体是
(x2 / y^2 )f"(u) + (x/y) f'(u)
根据贾月的等式。这个等于1. 所以有:
(x^2/y^2)f''(u)+(x/y)f'(u)=1
现在需要将这个方程用变量替换为关于u的常数分方程。因为u = x / y . 我们可以尝试将自变量转换为u。
◇u = x/v、服么可以将所有的变量转化为以u的形式表达。例如,在左边的项中:
(x^2 / y^2)f''(u) + (x/y)f'(u)=1
注意到左边可以写成 u*2 f"(u)+ u f(u), 因为 u= x/v, 所以 x = uv,
代入得到:
(u y)^2 / y^2 * f"(u) + (uy)/y *f'(u)
tb就是.
u^2 f''(u) + u f'(u)=1
因此, 方程变为: u2 f"(u)+u f(u)=1.
这是一个关于自变量为u的常微分方程。现在需要解这个ODE来找到f(u).
首先改写方程为:
或者等价地:
u^2 f'(u) + u f'(u) -1=0
可以尝试将变量替换,令v=f(u),则方程变成关于v的方程:
III: u2 dv/du + u v =1.
或者除以u.
得到: u (dv/du )+v = 1/u.
不过这可能更直观地用另一种变量替换。例如、观察左边是否为某个导数的形式。让我们重新整理方程:
将原式写为:
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u^2 f" + u f' = 1 可以视为政拉-柯西方程的一种类型,通常这类方程的解可以通过假设解的形式为f(u)=u^k来求得齐次解,但在这里非齐次项是常数 吗? 不对, 右边是1, 或者更准确地说, 原方程为 $u^2 d^2 f/du^2 + u df/du = 1$ 不过现在华处理这个ODE ◇p=f(u)、则方程变为 这是一个一阶线性常微分方程关于p(u). 标准形式是: p' + (1/u) p = 1/(u^2) 这里除以2: $p' + (1/u)p = 1/u^2$ 积分因子 u(u)=exp(f (1/u) du)= exp(in u)=u. 两边面以前公园子 $\mu(u)\rho' + \mu(u)(1/u) p = \mu(u)/(u^2)$ $d/du [\mu(u) * p] = 1/(u)$ 所以左边是 d/du [u*p] 、右边为1/u。 对两端积分: [d(up)=[(1/u)du+C 得到: up = In |u| + C. 解得p= f'(u)= (in u + C)/u. 然后求f(u)需要再次积分: f(u)= [[(ln u +C)/u] du 令t=In u, 那么dt=(1/u)du, 所以积分变为: [t + C dt = 1/2 t2 + Ct + D

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f(u) = \frac{1}{2} (\ln u)^2 + C \ln u + D.
加在黑海确党党数仓和自
接下来应用题目绘出的条件。然而题目中的条件是当某个条件下的g(x v)等于1.以及偏导数的值在点处的值为2/x,可能这里需要注
意条件的转换问题.
但题目中给出的具体条件是什么? 再仔细看一下:
题目说: "目 g(x,y) = 1.\$ \frac{1.\$ \frac{1}{y}}{y} = \frac{1.\$ \frac{1}{y}}{y}
这可能是在某个特定的曲线上的条件吗?通常这种条件下,可能需要更明确的条件。但原文中是"且"连接的这两个条件,所以可能
是指对于所有满足q(x,y)=1的点处,同时有\partial q/\partial x=2/x?
或者可能是存在某些矛盾的条件,需要澄清。
不过另一个可能的理解是,在某个具体的位置上的值,比如当u=x/y取特定值时f(u)的值以及其导数。例如可能在点(x,y)满足某
种关系时,如可能是在x=v的情况下?但题目并没有明确指出条件应用的目体位置,所以可能需要更仔细的分析。
或老这些条件是对于某些目体的食物而言的?
这里可能存在翻译或题度的理解问题,所以我需要再分析一下:
由于q(x,y)=f(x/y),而当它满足"q(x,y)=1"和在该点的导数。这里的条件可能对应于某个特定的点(x0,y0),或者在某种条件下的
函数值,但由于题目没有给出具体的点,所以可能需要进一步推断。
或者这实际上是初始条件应用于当u=x/y=1的时候?比如在u=18ff(u)的值以及导数值?
这里有可能存在误解。因为如果q(x, v)=1、也就是f(x/v)=1、假设这可能对接于某个结合位置的情况。例如当x=v号的情形。因此
Bfu=x/y=1.
而同时在该点处(可能是在x=v的情况下)。 求意的偏导数8g/8x的值是2/x,这时需要将条件转化为关于f(u)的条件:
例如在u=18t.可能有某些约束条件。让我们尝试这样处理。
假设当u=1即x/y=1也就是y=x时的位置处,应用给定的条件:q(x,x)=f(1)=1,
同时在该点处的偏导数ag/ax的值是2/x.
现在计算在点 (x,x) 处的偏导数
因为对于一般的(x,y), 有\partial \alpha/\partial x = f'(u) * 1/v。
所以在点(x,x), y=x。所以这里的导数是f'(1)/x。顾日中说这个值等于2/x。
ED: f'(1)/x=2/x \Rightarrow f'(1)=2.
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另外在u=1时, g(x, x)=f(1)=1. 因此得到:
当u=18t.
f(1) =1
现在问到BQH的通报: f(u)=½ (in u)*2 + C in u + D.
我们需要应用这两个条件来确定常数C和D。注意这里是在u=1处的条件和导数。因此需要代入u=1:
首先计算f(1):
f(1) = 1/2*(0)*2 + C*0 + D = D = 1 → D=1.
接下来求号: f'(u)= (in u)/u*(1/u)? 不, 让我重新算一下号数:
f(u) = (1/2)(\ln u)^2 + C \ln u + D
First f'(u) = (1/2) \mathcal{L}(\ln u) / I/u) + C (1/u) = (\ln u) / u + C / u
在u=1时, 代入得: f(1) = 0 + C*1 → C=2 (因为题目条件给出f(1)=2)
BECLARAS:
f(u)=(1/2)(in u)^2 +2 in u +1
不过需要检查这个结果是否符合原方程以及导数的计算。
现在综化入觉数分方程验证是否正确。例如、在ODE中: u2f"+uf=1.
先算二阶号数:
f'(u) = (\ln u)/u + 2/u = (\ln u + 2)/u
然后作(u)的导数为:
使用商数法則: [d/du (ln u +2)] * (1/u ) - (ln u +2)*(1 /u² )
RO.
[ (1/u) 1 / (1/u ) - (In u +2)/u^2
= 1/u<sup>2</sup> - (in u +2)/u<sup>2</sup> = [1 - in u -2]/u<sup>2</sup> = (-1 - in u)/u<sup>2</sup>.
代入原ODE左边:
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= (-1 - in u)+ (in u + 2 )= [-1 - in u] + [in u +2 ] = (-1+2)+( -in u +in u) )=1.
   满足等式右边为1、正确、因此解是正确的。
   另外检查初始条件:
   ≌iu=18t.
   f(1)=1/2*(0)^2 + 2*0 +1=1.
   f'(1)=(0 +2)/1 =2.符合题目要求的结果。
   所以最终得到函数是 f(u) = (1/2)(ln u)^2 +2 ln u +1.
   可能需要再整理一下:
   或者写成: f(u)= \frac{1}{2}{\lin u}^2 + 2\lin u + 1
   不过可能需要检查是否存在其他可能的解,或者在求解ODE过程中是否有遗漏的情况。但根据我们的步骤和验证过程是正确的。
   因此最终的答案应该是 f(u) = (1/2)(ln u)*2 + 2 ln u +1.
   当然也可以把它写成更简洁的形式:
   f(u)= \frac(1)(2)(\ln u)^2 + 2\ln u + 1
   这个表达式满足所有的条件。同时代入原偏微分方程之后也成立。因此是正确的解。
   器室
x^2 \frac{\partial^2 g}{\partial x^2} + xy \frac{\partial^2 g}{\partial x \partial y} + y^2 \frac{\partial^2 g}{\partial y^2} = 1,
首先, 令 u=\frac{\pi}{n}, 则 g(x,y)=f(u)。计算各阶偏导数:
  1. 一脸偏导数:
                                           \frac{\partial g}{\partial x} = f'(u) \cdot \frac{1}{u}, \quad \frac{\partial g}{\partial u} = f'(u) \cdot \left(-\frac{x}{u^2}\right)
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2025 考研数学一T19(续完)



输出情况

Listing 2: 2025 考研数学一 T18 生成情况

```
1    response_token/s: 13.96
2    prompt_token/s: 60.25
3    total_duration: 419088681937
4    load_duration: 74808920509
5    prompt_eval_count: 144
6    prompt_eval_duration: 2390000000
7    eval_count: 4774
8    eval_duration: 341887000000
9    approximate_total: 0h6m59s
```



Thank you!

