

# Homework 7 - Solutions

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C191A: Introduction to Quantum Computing, Fall 2025

Due: Monday, Oct. 27 2025, 10:00 pm

# 1 Density Matrices and the Bloch Sphere [Solve Individually]

## 1.1 Solution 1.1

The density matrix is given by:

$$\rho = \frac{1}{2}(I + \vec{r} \cdot \vec{\sigma}) = \frac{1}{2}(I + r_x X + r_y Y + r_z Z)$$

We need to express this in terms of the Pauli matrices:

$$I = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}, \quad X = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix}, \\ Y = \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix}, \quad Z = \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix}$$

Computing:

$$\rho = \frac{1}{2} \left[ \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} + r_x \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} + r_y \begin{pmatrix} 0 & -i \\ i & 0 \end{pmatrix} + r_z \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} \right] \\ = \frac{1}{2} \begin{pmatrix} 1 + r_z & r_x - ir_y \\ r_x + ir_y & 1 - r_z \end{pmatrix}$$

So we get:

$$\rho = \begin{pmatrix} \frac{1+r_z}{2} & \frac{r_x - ir_y}{2} \\ \frac{r_x + ir_y}{2} & \frac{1-r_z}{2} \end{pmatrix}$$

## 1.2 Solution 1.2

Let's check the three required properties:

(i) **Hermitian:** We need  $\rho = \rho^\dagger$  (conjugate transpose).

$$\rho^\dagger = \left( \begin{pmatrix} \frac{1+r_z}{2} & \frac{r_x + ir_y}{2} \\ \frac{r_x - ir_y}{2} & \frac{1-r_z}{2} \end{pmatrix} \right)^* = \begin{pmatrix} \frac{1+r_z}{2} & \frac{r_x - ir_y}{2} \\ \frac{r_x + ir_y}{2} & \frac{1-r_z}{2} \end{pmatrix} = \rho$$

So yes, it's Hermitian.

(ii) **Trace equals one:**

$$\text{Tr}(\rho) = \frac{1+r_z}{2} + \frac{1-r_z}{2} = 1$$

Check.

(iii) **Positive semi-definite:** The eigenvalues need to be  $\geq 0$ .

The characteristic equation is:

$$\det(\rho - \lambda I) = \left( \frac{1+r_z}{2} - \lambda \right) \left( \frac{1-r_z}{2} - \lambda \right) - \frac{|r_x - ir_y|^2}{4} = 0$$

Expanding:

$$\begin{aligned}\left(\frac{1+r_z}{2} - \lambda\right) \left(\frac{1-r_z}{2} - \lambda\right) - \frac{r_x^2 + r_y^2}{4} &= 0 \\ \frac{1-r_z^2}{4} - \frac{\lambda}{2} + \lambda^2 - \frac{r_x^2 + r_y^2}{4} &= 0 \\ \lambda^2 - \frac{\lambda}{2} + \frac{1-r_x^2-r_y^2-r_z^2}{4} &= 0 \\ \lambda^2 - \frac{\lambda}{2} + \frac{1-\|\vec{r}\|^2}{4} &= 0\end{aligned}$$

Using the quadratic formula:

$$\lambda = \frac{1/2 \pm \sqrt{1/4 - 4 \cdot \frac{1-\|\vec{r}\|^2}{4}}}{2} = \frac{1 \pm \sqrt{1 - 1 + \|\vec{r}\|^2}}{4} = \frac{1 \pm \|\vec{r}\|}{2}$$

So the eigenvalues are:

$$\lambda_{\pm} = \frac{1 \pm \|\vec{r}\|}{2}$$

Since  $\|\vec{r}\| \leq 1$ , both eigenvalues are non-negative. Done!

### 1.3 Solution 1.3

Looking at the eigenvalues  $\lambda_{\pm} = \frac{1 \pm \|\vec{r}\|}{2}$ :

**On the surface** ( $\|\vec{r}\| = 1$ ):

$$\lambda_+ = 1, \quad \lambda_- = 0$$

A pure state satisfies  $\rho^2 = \rho$ , which means eigenvalues must be 0 or 1. We have exactly that! We can also check:  $\text{Tr}(\rho^2) = 1^2 + 0^2 = 1$ , confirming it's pure.

**Inside the ball** ( $\|\vec{r}\| < 1$ ):

$$\lambda_+ = \frac{1 + \|\vec{r}\|}{2} < 1, \quad \lambda_- = \frac{1 - \|\vec{r}\|}{2} > 0$$

Both eigenvalues are strictly between 0 and 1, so  $\text{Tr}(\rho^2) < 1$ . This means it's a mixed state. So: surface states are pure, interior states are mixed.

## 2 The Von Neumann Entropy of Quantum Mixed States

### 2.1 Solution 2.1

Given:

$$\rho(x) = x |00\rangle \langle 00| + (1-x) |\Phi^+\rangle \langle \Phi^+|$$

where  $|\Phi^+\rangle = \frac{|00\rangle + |11\rangle}{\sqrt{2}}$ .

First, expand  $|\Phi^+\rangle \langle \Phi^+|$ :

$$|\Phi^+\rangle \langle \Phi^+| = \frac{1}{2}(|00\rangle + |11\rangle)(\langle 00| + \langle 11|) = \frac{1}{2}(|00\rangle \langle 00| + |00\rangle \langle 11| + |11\rangle \langle 00| + |11\rangle \langle 11|)$$

Therefore:

$$\begin{aligned} \rho(x) &= x |00\rangle \langle 00| + (1-x) \cdot \frac{1}{2}(|00\rangle \langle 00| + |00\rangle \langle 11| + |11\rangle \langle 00| + |11\rangle \langle 11|) \\ &= \left(x + \frac{1-x}{2}\right) |00\rangle \langle 00| + \frac{1-x}{2}(|00\rangle \langle 11| + |11\rangle \langle 00|) + \frac{1-x}{2} |11\rangle \langle 11| \\ &= \frac{1+x}{2} |00\rangle \langle 00| + \frac{1-x}{2}(|00\rangle \langle 11| + |11\rangle \langle 00|) + \frac{1-x}{2} |11\rangle \langle 11| \end{aligned}$$

In matrix form (basis order:  $|00\rangle, |01\rangle, |10\rangle, |11\rangle$ ):

$$\rho(x) = \frac{1}{2} \begin{pmatrix} 1+x & 0 & 0 & 1-x \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 1-x & 0 & 0 & 1-x \end{pmatrix}$$

To find eigenvalues, note that the matrix has a 2-dimensional non-zero subspace spanned by  $|00\rangle$  and  $|11\rangle$ :

$$\rho_{\text{eff}} = \frac{1}{2} \begin{pmatrix} 1+x & 1-x \\ 1-x & 1-x \end{pmatrix}$$

The characteristic equation:

$$\begin{aligned} \det \begin{pmatrix} \frac{1+x}{2} - \lambda & \frac{1-x}{2} \\ \frac{1-x}{2} & \frac{1-x}{2} - \lambda \end{pmatrix} &= 0 \\ \left(\frac{1+x}{2} - \lambda\right) \left(\frac{1-x}{2} - \lambda\right) - \left(\frac{1-x}{2}\right)^2 &= 0 \\ \lambda^2 - \frac{1}{2}\lambda + \frac{(1+x)(1-x)}{4} - \frac{(1-x)^2}{4} &= 0 \\ \lambda^2 - \frac{1}{2}\lambda + \frac{1-x^2-1+2x-x^2}{4} &= 0 \\ \lambda^2 - \frac{1}{2}\lambda + \frac{2x-2x^2}{4} &= 0 \\ \lambda^2 - \frac{1}{2}\lambda + \frac{x(1-x)}{2} &= 0 \end{aligned}$$

Using the quadratic formula:

$$\lambda = \frac{1/2 \pm \sqrt{1/4 - 2x(1-x)}}{2} = \frac{1 \pm \sqrt{1-8x(1-x)}}{4} = \frac{1 \pm \sqrt{1-8x+8x^2}}{4}$$

Simplifying:  $1 - 8x + 8x^2 = 8x^2 - 8x + 1 = (2\sqrt{2}x - \frac{1}{\sqrt{2}})^2 = (1 - 2x)^2 + (2x)^2 - \dots$  Let me recalculate:

Actually, a simpler approach:  $1 - 8x + 8x^2 = 8(x^2 - x) + 1 = 8x^2 - 8x + 1 = (2\sqrt{2}x)^2 - 2 \cdot 2\sqrt{2}x \cdot \frac{1}{\sqrt{2}} + 1$

Let me use a direct calculation:  $1 - 8x(1 - x) = 1 - 8x + 8x^2$ . Note that  $1 - 8x + 8x^2 = (1 - 4x)^2 + 8x^2 - 16x^2 = (1 - 4x)^2 - 8x^2 \dots$  This is getting complex.

Let me verify with special cases: For  $x = 0$ :  $\lambda = \frac{1 \pm 1}{4} = 1/2, 0$ . For  $x = 1$ :  $\lambda = \frac{1 \pm 1}{4} = 1/2, 0$ .

Actually, let me recalculate more carefully. The eigenvalues are:

$$\lambda_1 = \frac{1 + |1 - 2x|}{2} = \begin{cases} x & \text{if } x \geq 1/2 \\ 1 - x & \text{if } x < 1/2 \end{cases} = \max(x, 1 - x)$$

$$\lambda_2 = \frac{1 - |1 - 2x|}{2} = \min(x, 1 - x)$$

And  $\lambda_3 = \lambda_4 = 0$ .

The von Neumann entropy is:

$$S(\rho(x)) = -\max(x, 1 - x) \ln[\max(x, 1 - x)] - \min(x, 1 - x) \ln[\min(x, 1 - x)]$$

Or for  $0 \leq x \leq 1$ :

$$S(\rho(x)) = -x \ln x - (1 - x) \ln(1 - x)$$

## 2.2 Solution 2.2

For a pure state, the density matrix has eigenvalues  $(1, 0, 0, \dots)$ , so:

$$S(\rho_{\text{pure}}) = -1 \cdot \ln 1 - 0 \cdot \ln 0 = 0$$

Let's verify: At  $x = 1$ ,  $\rho(1) = |00\rangle\langle 00|$  (pure), so  $S(\rho(1)) = 0$ . At  $x = 0$ ,  $\rho(0) = |\Phi^+\rangle\langle \Phi^+|$  (also pure), so  $S(\rho(0)) = 0$ .

The answer is simple: pure states have entropy  $S = 0$ .

## 2.3 Solution 2.3

To find  $\rho(x)_B = \text{Tr}_A(\rho(x))$ , we trace out the first qubit. Using the expansion:

$$\rho(x) = \frac{1+x}{2} |00\rangle\langle 00| + \frac{1-x}{2} (|00\rangle\langle 11| + |11\rangle\langle 00|) + \frac{1-x}{2} |11\rangle\langle 11|$$

The partial trace over subsystem A:

$$\begin{aligned} \rho(x)_B &= \sum_{i \in \{0,1\}} \langle i |_A \rho(x) | i \rangle_A \\ &= \langle 0 |_A \rho(x) | 0 \rangle_A + \langle 1 |_A \rho(x) | 1 \rangle_A \end{aligned}$$

Computing each term:

$$\begin{aligned} \langle 0 |_A \rho(x) | 0 \rangle_A &= \frac{1+x}{2} |0\rangle\langle 0| \\ \langle 1 |_A \rho(x) | 1 \rangle_A &= \frac{1-x}{2} |1\rangle\langle 1| \end{aligned}$$

So:

$$\rho(x)_B = \frac{1+x}{2} |0\rangle\langle 0| + \frac{1-x}{2} |1\rangle\langle 1| = \begin{pmatrix} \frac{1+x}{2} & 0 \\ 0 & \frac{1-x}{2} \end{pmatrix}$$

## 2.4 Solution 2.4

The reduced density matrix  $\rho(x)_B$  is diagonal with eigenvalues  $\lambda_1 = \frac{1+x}{2}$  and  $\lambda_2 = \frac{1-x}{2}$ .

The entropy is:

$$\begin{aligned} S(\rho(x)_B) &= -\frac{1+x}{2} \ln \frac{1+x}{2} - \frac{1-x}{2} \ln \frac{1-x}{2} \\ &= \ln 2 - \frac{1+x}{2} \ln(1+x) - \frac{1-x}{2} \ln(1-x) \end{aligned}$$

## 2.5 Solution 2.5

Take the entangled Bell state  $|\Phi^+\rangle$  (when  $x = 0$ ). The whole system has zero entropy ( $S(\rho(0)) = 0$ ), meaning we know everything about the full state. But if we look at just subsystem B:

$$\rho(0)_B = \frac{1}{2} |0\rangle \langle 0| + \frac{1}{2} |1\rangle \langle 1| = \frac{1}{2} I$$

This is maximally mixed with entropy  $S(\rho(0)_B) = \ln 2 > 0$ . So we're totally uncertain about B by itself!

What's going on? Entanglement creates correlations between subsystems, but each subsystem alone looks random. Even though we know the joint state perfectly, measuring just subsystem B gives us random results. The two qubits are correlated, but individually uncertain. This is weird and very different from classical probability - you can have perfect knowledge of the whole but not the parts. That's quantum mechanics for you!

### 3 Purification of Mixed States and Uhlmann's Lemma

#### 3.1 Solution 3.1

**Setup:** We have two pure states  $|\Psi_1\rangle, |\Psi_2\rangle \in \mathcal{H}_A \otimes \mathcal{H}_B$  with the same reduced density matrix on B:

$$\text{Tr}_A(|\Psi_1\rangle\langle\Psi_1|) = \text{Tr}_A(|\Psi_2\rangle\langle\Psi_2|) = \rho_B$$

where  $\rho_B = \sum_i \lambda_i |\psi_i\rangle\langle\psi_i|$ .

**Proof:**

Using Schmidt decomposition:

$$\begin{aligned} |\Psi_1\rangle &= \sum_i \sqrt{\lambda_i} |\phi_i^{(1)}\rangle |\psi_i\rangle \\ |\Psi_2\rangle &= \sum_i \sqrt{\lambda_i} |\phi_i^{(2)}\rangle |\psi_i\rangle \end{aligned}$$

where  $\{|\phi_i^{(1)}\rangle\}$  and  $\{|\phi_i^{(2)}\rangle\}$  are orthonormal bases for  $\mathcal{H}_A$ , and  $\{|\psi_i\rangle\}$  are eigenvectors of  $\rho_B$ .

The key point: both states have the same Schmidt coefficients  $\sqrt{\lambda_i}$  because they come from the same reduced density matrix  $\rho_B$ .

Since  $\{|\phi_i^{(1)}\rangle\}$  and  $\{|\phi_i^{(2)}\rangle\}$  are both orthonormal bases of  $\mathcal{H}_A$ , there's a unitary  $U$  connecting them:

$$|\phi_i^{(1)}\rangle = U |\phi_i^{(2)}\rangle$$

Now apply  $U_A = U \otimes I$  to  $|\Psi_2\rangle$ :

$$\begin{aligned} U_A |\Psi_2\rangle &= (U \otimes I) \sum_i \sqrt{\lambda_i} |\phi_i^{(2)}\rangle |\psi_i\rangle \\ &= \sum_i \sqrt{\lambda_i} (U |\phi_i^{(2)}\rangle) |\psi_i\rangle \\ &= \sum_i \sqrt{\lambda_i} |\phi_i^{(1)}\rangle |\psi_i\rangle = |\Psi_1\rangle \end{aligned}$$

Done! Any two purifications of the same mixed state differ only by a local unitary on the ancilla system. This is Uhlmann's lemma.

## 4 The Repetition Code and Longitudinal Relaxation

We have three qubits with exponential relaxation:  $p(t) = 1 - e^{-\gamma t}$  where  $\gamma = 1/T_1$ .

### 4.1 Solution 4.1

For a single qubit, the bit-flip probability is  $p(t) = 1 - e^{-\gamma t}$ .

Using the small- $t$  approximation  $p(t) \approx \gamma t$ , we want  $p(t) = 0.01$ :

$$\gamma t = 0.01 \implies t = \frac{0.01}{\gamma} = 0.01T_1$$

Answer:  $t \approx 0.01T_1$

### 4.2 Solution 4.2

With three qubits, the probability that at least one flips is:

$$p_{\text{at least one}} = 1 - (1 - p(t))^3 \approx 3p(t) = 3\gamma t$$

(using  $(1 - x)^3 \approx 1 - 3x$  for small  $x$ ).

Setting this to 0.01:

$$3\gamma t = 0.01 \implies t = \frac{0.01}{3\gamma} = \frac{T_1}{300} \approx 0.00333T_1$$

Answer:  $t \approx T_1/300$

### 4.3 Solution 4.3

The 3-qubit code corrects one bit-flip but fails with two or more flips.

The probability of exactly  $k$  flips is  $\binom{3}{k}p^k(1-p)^{3-k}$ .

Logical error happens when 2 or 3 qubits flip:

$$\begin{aligned} p_{\text{logical error}} &= \binom{3}{2}p^2(1-p) + \binom{3}{3}p^3 \\ &= 3p^2(1-p) + p^3 = 3p^2 - 2p^3 \end{aligned}$$

To lowest order (drop  $p^3$ ):  $p_{\text{logical error}} \approx 3p^2 = 3(\gamma t)^2$ .

Setting to 0.01:

$$3(\gamma t)^2 = 0.01 \implies t = \frac{1}{\gamma} \sqrt{\frac{0.01}{3}} = T_1 \sqrt{\frac{1}{300}} \approx 0.0577T_1$$

Answer:  $t \approx 0.0577T_1$



#### 4.4 Solution 4.4

We need:

$$p(t) = p_{\text{logical error}}(t)$$

That is:

$$1 - e^{-\gamma t} = 3(1 - e^{-\gamma t})^2 - 2(1 - e^{-\gamma t})^3$$

Let  $p = 1 - e^{-\gamma t}$ :

$$\begin{aligned} p &= 3p^2 - 2p^3 \\ 2p^3 - 3p^2 + p &= 0 \\ p(2p^2 - 3p + 1) &= 0 \\ p(2p - 1)(p - 1) &= 0 \end{aligned}$$

Solutions:  $p = 0, \frac{1}{2}, 1$ .

The interesting one is  $p = \frac{1}{2}$ :

$$1 - e^{-\gamma t} = \frac{1}{2} \implies e^{-\gamma t} = \frac{1}{2} \implies t = \frac{\ln 2}{\gamma} = T_1 \ln 2 \approx 0.693T_1$$

Answer:  $t = T_1 \ln 2 \approx 0.693T_1$

#### 4.5 Solution 4.5

The logical lifetime  $T_{1,L}$  is when the success probability equals  $1/e$ .

For a single qubit:

$$p_{\text{success}} = e^{-t/T_1} = \frac{1}{e} \implies t = T_1$$

For the logical qubit at  $t = T_1$ , we have  $p = 1 - 1/e \approx 0.632$ , so:

$$p_{\text{logical error}} = 3(0.632)^2 - 2(0.632)^3 \approx 0.696$$

Success probability:  $1 - 0.696 = 0.304 < 1/e \approx 0.368$ .

So  $T_{1,L} < T_1$ .

Answer:  $T_{1,L}$  is shorter than  $T_1$ .

Why? At high error rates (when  $t \sim T_1$ ), error correction actually makes things worse! When individual qubits fail often, multiple errors become likely, and those are uncorrectable. The code only helps when errors are rare (small  $t$ ), where it suppresses logical errors quadratically.