Programming Languages and Environments (Lecture 3)

LEI - Licenciatura em Engenharia Informática

João Costa Seco (joao.seco@fct.unl.pt)



Syllabus

- Name declarations
- Function declaration, with and without parameters.
- Evaluation of expression by substitution.
- Functions as values.
- Function partial applications.

Declarations

Variable declarations

Global declarations

Local declarations

```
let x = 42 in (string_of_int x)^": the ultimate question of life, the universe, and everything"

\[ \sqrt{13} \sqrt{0.0s} \]

\[ \cdots \]

- : string = "42: the ultimate question of life, the universe, and everything"
```

Scoping

- A declaration of a name (x) is limited to the body of the declaration (e2).
- It is not, for instance, visible in the expression that defines its value (e1).

let
$$x = e1$$
 in $e2$

let
$$y = let y = 1 in y + 1 in let y = y + 2 in y + 2$$

 Declarations follow the principle of name irrelevance, meaning that the chosen names should not affect the evaluation of an expression.

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 in $e2$

let
$$y =$$
 let $y = 1$ in $y + 1$ in let $y = y + 2$ in $y + 2$

 Declarations follow the principle of name irrelevance, meaning that the chosen names should not affect the evaluation of an expression.

```
let x =
 let y = 1 in
 let z = 2 in
 Y + Z
let w = 3+x in
```

```
let x =
 let y = 1 in
 let z = 2 in
let w = 3+x in
```

```
let x =
 let y = 1 in
 let z = 2 in
let w = 3+x in
```

```
let x =
 let y = 1 in
 let z = 2 in
 Y + Z
1n
let w = 3+x in
```

```
let x =
 let y = 1 in let z = 2 in
let w = 3+x in
```

```
let x =
 let y = 1 in
 let z = 2 in
 Y + Z
1
let w = 3+x in
```

```
let z = 2 in 1 + z
let w = 3+x in
```

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$$let x =$$

```
1 + 2
in
let w = 3+x in
w + x
```

```
3
in
let w = 3+x in
w + x
```

let
$$w = 3+3$$
 in $w + 3$

let
$$w = 6$$
 in $w = 3$



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Functions

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.

let
$$f x = e1$$
 in $e2$

let f
$$x = x + 1$$
 in f $(1 + 1)$

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.
- Functions with "no parameters", have a parameter of type unit.

let
$$x = 1$$
 in let $f() = 1 + x$ in $f()$

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.
- Declarations with parameters is a syntactic alternative to using functions as values (the arrow type is composed by two characters ->).

let
$$f = fun x \rightarrow x + 1 in f (1 + 1)$$

Definition and evaluation of functions

- The application of functions can be defined by the substitution of the parameter by the value of the argument.
- OCaml implements *call-by-value* evaluation strategy, meaning that the arguments are evaluated before expanding the body of the function.

(fun
$$x \rightarrow x + 1$$
) (1 + 1)
(fun $x \rightarrow x + 1$) 2
2 + 1
3

Recursive definitions (scoping)

• The declaration of a name (x) is visible in the body of the declaration (e2) and in the body of the declaration (e1).

let rec x = e1 in e2

```
(* [fact x] computes the factorial of x
Requires: [x >= 0] *)
```

let rec fact x = if x = 0 then 1 else x * fact (x - 1)

Mutually recursive declarations (scoping)

The declaration of a name (x) is visible in the body of the declaration (e2) and
in the body of the declaration (e1).

let rec x = e1 in e2

Mutually recursive declarations in C

• Declare a function without defining it.

```
bool odd(int x);
bool even(int x) {
  if(x == 0)
    return false;
  } else if( x == 1 ) {
    return false;
  } else {
    return odd(x-1);
```

```
bool odd(int x) {
  if( x == 0 ) {
    return false;
  } else if( x == 1 ) {
    return true;
  } else {
    return even(x-1);
  }
}
```

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.

let
$$f x y = x + y in f 1 1$$

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.

let
$$f = fun x y \rightarrow x + y in f 1 1$$

- The name declaration (f) is limited to the body (e2).
- The name (f) is not visible in the expression that defines the value (e1).
- Parameters are listed in the declaration.

let
$$f = fun x \rightarrow fun y \rightarrow x + y in f 1 1$$

Partial evaluation of functions

- A function with multiple parameters is essentially the composition of multiple functions.
- Parameters can be instantiated one at a time, resulting in partial applications until the evaluation is complete.

```
let add x y = x + y
 ✓ 0.0s
val add : int \rightarrow int = <fun>
   add 2 3
-: int = 5
   let add1 = add 1
val add1 : int \rightarrow int = <fun>
```

Summary

- Name declarations
- Function declaration, with and without parameters.
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Programming Languages and Environments (Lecture 4)

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