Programming Languages and Environments (Lecture 15)

LEI - Licenciatura em Engenharia Informática

João Costa Seco (joao.seco@fct.unl.pt)



Syllabus

- Module system
- Functors
- Concurrency
 - Promises

Module System

- Name spaces
 - Groups of statements (usually) related isolated from other groups through the qualification of the names in a module.
 - Allows the reuse of the same names in different contexts without collisions.
 - Packages and classes in Java, files/modules in C, structures/modules in OCaml
- Abstraction
 - Allows you to selectively hide/reveal information (information hiding)
 - Code isolation, better development and maintenance, ownership, etc.
- Code reuse
 - Reuse without copy, modularity, (cf. inheritance in Java)
- (In OCaml) Module parameterisation
 - The Functors in OCaml are like functions from modules to modules (cf. traits in Scala)

Modules in OCaml

- The modules are defined by structures (struct)
- The types for the modules are signatures (sig)
- Default type definitions are public (type)
- Name implementations are private (val)

```
module MyModule = struct
  type primary_color = Red | Green | Blue

let inc x = x + 1
  let dec x = x - 1
end
```

```
(* Inferred signature *)
module MyModule :
    sig
      type primary_color = Red | Green | Blue
    val inc : int -> int
    val dec : int -> int
    end
```

utop

Namespaces

- The names declared in a module can be used in a qualified way (with the module name and a period: List.fold_right)
- Or you can use the open directive to expand the names of the module used in the client module.
- The StdLib module is always open.

Module MyList

```
module MyList = struct
  type 'a list = Nil | Cons of 'a * 'a list
  let empty = Nil
  let rec length = function
      Nil -> 0
      Cons (\_, xs) \rightarrow 1 + length xs
  let insert x xs = Cons (x, xs)
  let head = function
    Nil -> None
      Cons (x, _) \rightarrow Some x
  let tail = function
     Nil -> None
     Cons (_, xs) -> Some xs
end
```

```
module MyList :
    sig
    type 'a list = Nil | Cons of 'a * 'a list
    val empty : 'a list
    val length : 'a list -> int
    val insert : 'a -> 'a list -> 'a list
    val head : 'a list -> 'a option
    val tail : 'a list -> 'a list option
    end
```

Java vs. Ocaml

Java
s = new List();
s.insert(1);
OCaml
let s = MyList.empty;;
let s' = MyList.insert 6 s;;

Name abstraction

- The types of the modules also allow you hide the definition of the types
- A subscription can have multiple compatible implementations (opaque)

```
module type Stack = sig
  type 'a stack
  val empty : 'a stack
  val is_empty : 'a stack -> bool
  val push : 'a -> 'a stack -> 'a stack
  val top : 'a stack -> 'a option
  val pop : 'a stack -> 'a stack option
end
```

```
module MyListStack : Stack = struct
  type 'a stack = 'a MyList.list
  let empty = MyList.empty
  let is_empty = MyList.is_empty
  let push x s = MyList.insert x s
  let top = MyList.head
  let pop = MyList.tail
end;;
```

Namespaces (again)

```
type 'a stack = 'a MyList.list
let empty = MyList.empty
let is_empty = MyList.is_empty
let push x s = MyList.insert x s
let top = MyList.head
let pop = MyList.tail
end;;
```

```
module MyListStack : Stack = struct
  open MyList

type 'a stack = 'a list
  let empty = empty
  let is_empty = is_empty
  let push x s = insert x s
  let top = head
  let pop = tail
end;;
```

Module ListStackCachedSize

Implement and test!

```
module ListStackCachedSize : Stack = struct
  type 'a stack = 'a list * int

let empty = ([], 0)
let is_empty s =
   match s with
  | ([], _) -> true
  | _ -> false
let push x s = (x::(fst s), (snd s)+1)
let top s = match s with
  | ([], _) -> None
  | (x :: xs, _) -> Some x
let pop s = match s with
  | ([], _) -> None
  | (x :: xs, n) -> Some (xs, n-1)
let size s = snd s
end
```

Module Counter with Refs!

Implement and test!

```
module type Counter = sig
  type t
  (** [create v] makes a new counter the initial value [v]. *)
  val create : int -> t

  (** [inc c] increments the counter by 1. *)
  val inc : t -> unit

  (** [dec c] decrements the counter by 1. *)
  val dec : t -> unit

  (** [get c] returns the current value. *)
  val get : t -> int

  (** [reset c] sets the counter to zero. *)
  val reset : t -> unit
end
```

```
module CounterRef : Counter = struct
  type t = int ref

let create v = ref v

let inc c = c := !c + 1

let dec c = c := !c - 1

let get c = !c

let reset c = c := 0

end
```

Types and names

The specialisation of module types can be done with an adaptation module.

```
module IntStack = (struct
  (* 1. Build on a generic "ListStack" module
        by fixing its element type to int. *)
  type stack = int MyListStack.stack
  (* 2. Re-export the operations from ListStack *)
  let empty = MyListStack.empty
  let push = MyListStack.push
  let pop = MvListStack.pop
  let top = MyListStack.top
end : sia
  (* 3. Users of IntStack only see the abstract type [stack]
        and the four operations with the following types *)
  type stack
  val empty : stack
  val push : int -> stack -> stack
  val pop : stack -> stack option
  val top : stack -> int option
end)
```

```
module IntStack :
    sig
     type stack
    val empty : stack
    val push : int -> stack -> stack
    val pop : stack -> int option
    val top : stack -> int option
    end
```

Modules and files

- The organisation in files separates the structure (struct) from the signature (sig)
- Files MyList.ml, Stack.mli, MyStackList.ml

```
module · MyList · = · struct

· type · 'a · list · = · Nil · | · Cons · of · 'a · * · 'a · list

· let · empty · = · Nil

· let · is_empty · = · function
· let · rec · length · = · function
· let · insert · x · xs · = · Cons · (x, · xs)

· let · insert · x · xs · = · Cons · (x, · xs)

· let · head · = · function
· let · head · function
· let · head · = · function
· let · head · = · function
```

```
LAP-2025 > Stack.mli > ...

module type Stack = sig

type 'a stack

val empty : 'a stack

val is_empty : 'a stack -> bool

val push · · · : 'a -> 'a stack -> 'a stack

val top · · · : 'a stack -> 'a option

val pop · · · : 'a stack -> 'a stack option

end
```

```
s > LAP 2024-12 > MyStack.ml > ...

type 'a stack = 'a MyList.list
'a
let empty = MyList.empty
'a -> 'b -> 'c
let push x xs = MyList.insert
'a
let pop = MyList.tail
'a
let top = MyList.head
```

Modules and Functors (module functions for modules)

```
module type X = sig
  val x : int
end

module IncX (M : X) = struct
  let x = M.x + 1
end
```

```
module type X = sig val x : int end
module IncX : functor (M : X) -> sig val x : int end
```

Modules and Functors (module functions for modules)

```
module type X = sig
 type t
end
module Stack = struct
 module Make (M : X) = struct
                                                       module Stack:
   type stack = M.t list
                                                         siq
   let empty = []
                                                           module Make:
   let push x xs = x :: xs
                                                             functor (M : X) ->
   let pop = function
                                                               siq
      [] -> None
                                                                type stack = M.t list
       :: xs -> Some xs
                                                                val empty : 'a list
   let top = function
                                                                val push : 'a -> 'a list -> 'a list
       [] -> None
                                                                val pop : 'a list -> 'a list option
      x :: -> Some x
                                                                val top: 'a list -> 'a option
 end
                                                               end
end
                                                         end
module IntStack = Stack.Make (struct type t = int end)
let s = IntStack.empty
let = assert (IntStack.top s = None)
let = assert (IntStack.top (IntStack.push 1 s) = Some 1)
let = assert (IntStack.pop (IntStack.push 1 s) = Some IntStack.empty)
```

Modules and Functors (module functions for modules)

```
module Pair = struct
  type t = int * string

let compare (x1, y1) (x2, y2) =
    if x1 < x2 then -1
    else if x1 = x2 && y1 < y2 then -1
    else if x1 = x2 && y1 = y2 then 0
    else 1
end

module Str = Set.Make(Pair)</pre>
```

Concurrency

Concurrency

- Capacity to perform several calculations that overlap in time, and not require them to be done in sequence.
- Graphical interfaces; to ensure quick responses to user actions.
- Spreadsheets; to recalculate all calculations without interruptions.
- Web Browser; to load and show pages incrementally.
- Servers; to serve several customers without making them wait.
- How?
- Interlacing: switching between the various active computes quickly.
- Parallelism: using several physical processors (multi-core).
- Preemptive / Collaborative

Non-determinism

- A program that can have different results each time it runs is a nondeterministic program.
- The introduction of competition causes non-determinism, by not fixing the order in which operations are executed.
- When two imperative programs share state variables, the possibilities of changing that same state are indeterminate. Therefore, it is not easy to predict the exact behaviour of the programs.
- Interactions/interferences between programs are benign or malignant (race conditions)
- Functional languages make it easier to reason about programs because a pure expression always denotes the same value.

Non-determinism

x = 0

Thread A

A1: READ(x)

A2: WRITE(x, x+1)

Thread B

B1: READ(x)

B2: WRITE(x, x+1)

$$x = ?$$

Promises or Futures

 They represent computations that have not yet ended but that will denote a value somewhere in a future (deferred) instant.

Usually associated with a concurrent computing to the main thread.

```
async function getNames() {
  const response = await fetch('https://server.com/users')
  const data = await response.json()
  return data.map(user ⇒ user.name)
}
```

 Async (Jane Street) and Lwt (Ocsigen) are two popular libraries for implementing asynchronous computing in OCaml.

Lwt-style promises

They are data abstractions for an asynchronous computing model.

- Promises are references, their value can change.
- When it is created, it does not contain anything.
- A promise can be fulfilled and fulfilled for a value
- A promise can be rejected (filled by an exception)
- In both cases it is said to be resolved.

Signature of the Promise Module

```
module type PROMISE = sig
 type 'a state = Pending | Fulfilled of 'a | Rejected of exn
 type 'a promise
 type 'a resolver
  (** [make ()] is a new promise and resolver. The promise is pending. *)
 val make : unit -> 'a promise * 'a resolver
  (** [return x] is a new promise that is already fulfilled with value [x]. *)
 val return : 'a -> 'a promise
  (** [state p] is the state of the promise *)
 val state : 'a promise -> 'a state
  (** [fulfill r x] fulfills the promise [p] associated with [r] with value [x], meaning that
      [state p] will become [Fulfilled x]. Requires: [p] is pending. *)
 val fulfill : 'a resolver -> 'a -> unit
  (** [reject r x] rejects the promise [p] associated with [r] with exception [x], meaning that
      [state p] will become [Rejected x].
                                           Requires: [p] is pending. *)
 val reject : 'a resolver -> exn -> unit
end
```

Implementation of the Promise Module

```
module Promise : PROMISE = struct
        type 'a state =
            Pending
            Fulfilled of 'a
            Rejected of exn
        type 'a promise = 'a state ref
        type 'a resolver = 'a promise
        (** [write once p s] changes the state of [p] to be [s]. If [p] and [s] are both pending,
            that has no effect. Raises: [Invalid arg] if the state of [p] is not pending. *)
        let write once p s =
          if !p = Pending then p := s else invalid arg "cannot write twice"
        let make () = let p = ref Pending in (p, p)
        let return x = ref (Fulfilled x)
        let state p = !p
        let fulfill r x = write once r (Fulfilled x)
        let reject r x = write_once r (Rejected x)
      end
Linguagens e ambientes de Programação, NUVA FUI, © 2024, João Costa
```

Use of Promises

```
(** Lets see how to use promises *)
let compute fact sync n : int Promise.promise =
 let p, r = Promise.make () in
 begin
   try
     Promise.fulfill r result (* settle with the computed value *)
   with exn ->
     Promise reject r exn (* an error, which never happens here*)
 end;
let check name p =
 match Promise.state p with
  | Pending ->
     Printf.printf "%s is still computing...\n" name
 | Fulfilled v ->
     Printf.printf "%s = %d\n" name v
 | Rejected exn ->
     Printf.printf "%s failed\n" name
let p5 = compute_fact_sync 5 in check "5!" p5
```

Use of Promises in LWT

```
(* Pure-function *)
let rec fact n =
  if n \le 1 then 1 else n * fact (n - 1)
(* Initialize the preemptive pool with default bounds (min=0, max=4) *)
let () = Lwt preemptive.simple init ()
(* Offload [fact n] onto Lwt's preemptive pool *)
let compute_fact n : int Lwt.t =
 Lwt preemptive.detach fact n
let () =
 let work = [
   compute_fact 20 >>= fun r -> Lwt_io.printf "20! = %d\n" r;
   compute_fact 25 >>= fun r -> Lwt_io.printf "25! = %d\n" r;
  l in
 Lwt_main.run (Lwt.join work)
                                                             In utop:
                                                             #require "lwt.unix";;
```

#open Lwt;;

#open Lwt.Infix;;