

11.1 Groups and Coding

D: \oplus : 加和模2

D: $e: B^m \rightarrow B^n$ ^{字空间 \rightarrow 码字空间} encoding Function, m, n 为长度

D: weight 权值 $\rightarrow |x|$ 中1的个数

D: Hamming distance: $|x_1 \oplus x_2|$ $s(x, y)$
 \downarrow
 the weight of $x_1 \oplus x_2$ 加和模2.

I Theorem: ① $s(x, y) = s(y, x)$ ② $s(x, y) \geq 0$ ③ $s(x, y) = 0 \iff x = y$

④ ? 走神了.

D: Minimum distance of an encoding function e :

$$\min \{d(e(x_1), e(x_2)) \mid x_1, x_2 \in B^m\}.$$

II Theorem: $e(\varphi) = \dots$
 \downarrow \downarrow
 字空间 码字空间
 $\underbrace{\text{最小距离}}_{\geq k+1}$
 \downarrow
 error

III Theorem:

$e: B^m \rightarrow B^n$ be a group code,
 the min distance of e is the
 min weight of a nonzero code.

D: group code $e(B^m) = \{e(b) \mid e(b) \in B^n\} = \text{Ran}(e)$

$e: B^m \rightarrow B^n$ is a subgroup of B^n .

(group: $B^3 \rightarrow B^6$)
 封闭、有e、有 x^{-1} 、可结合??

D: definition I.II... Theorem.

D: Boolean \sim of matrix

$$A \odot B = \dots \vee a_i b_i$$

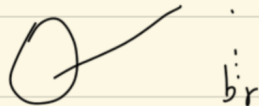
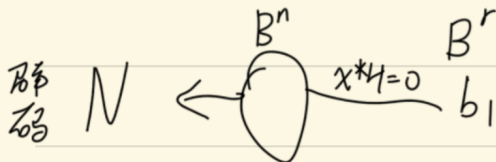
V Theorem:

奇偶校验矩阵

$m < n$, nonnegative, $r = n - m$, H be an $m \times n$ Boolean matrix

$$f_H: B^n \rightarrow B^r : \underline{f_H(x) = x * H, x \in B^n}$$

is a homomorphism from Group B^n to B^r



$$x_1, x_2 \in B^n, \text{ 则 } f_H(x_1 \oplus x_2) = f_H(x_1) \oplus f_H(x_2)$$

Corollary:

$N = \{x \in B^n \mid x * H = 0\}$ is a normal subgroup of B^n .

证: ① $x_1, \dots, x_n \in N$ 非空

② 单位元 0

③ 逆元 x^{-1}

④ $aH = Ha$

\downarrow
 $H \leq G$ 是 normal group, $\left\{ \begin{array}{l} \forall g \in G, h \in H, ghg^{-1} \in H \\ gH = Hg \end{array} \right.$

$$e_H: e^m \rightarrow e^n$$

$$b = b_1 b_2 \dots b_m, \quad x = e_H(b) = b_1 b_2 \dots b_m x_1 x_2 \dots x_r$$

$$x_i = b_1 \underline{h_{1i}} + b_2 \underline{h_{2i}} + \dots + b_m \cdot h_{mi}$$

$$\vdots$$

第 i 列

$$x_r$$

① m 的字节空间 ② $n=1$, 补3位.

$$\begin{array}{c} \text{11} \\ \text{2} \end{array} \quad \begin{array}{c} \text{00} \\ \text{01} \\ \text{10} \\ \text{11} \end{array}$$

$b_1 b_2$

11.2. Decoding and Error Correction. $\leftarrow \text{译码} \rightarrow d_{min} \rightarrow s$
 decoding F 去校验 ???

I. $e: (m, n)$ encoding F, d is a maximum likelihood decoding F ,
 then (e, d) can correct K or fewer errors

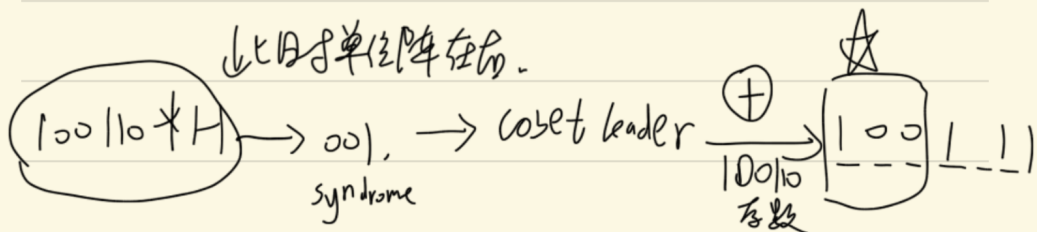
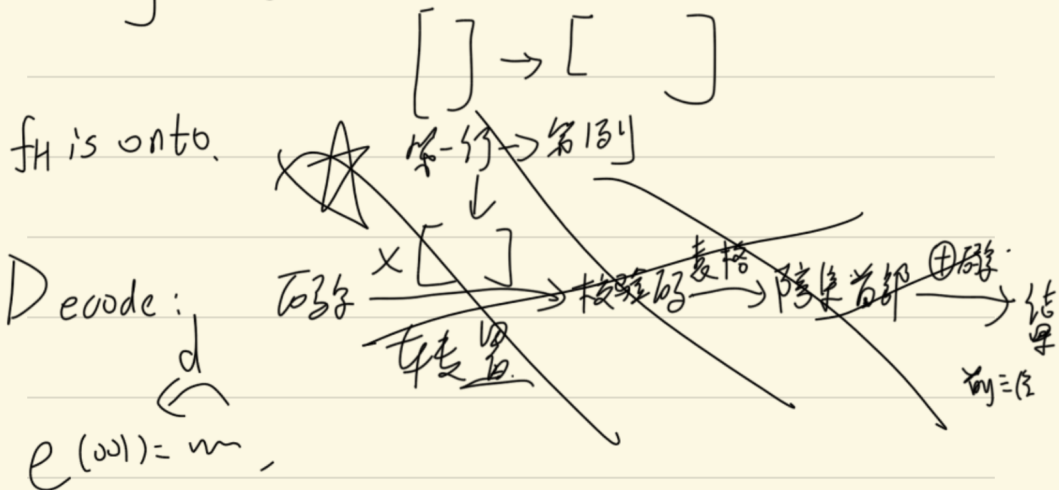
\updownarrow
 the minimum distance of e is $2k+1$

$$e^3 \rightarrow e^8 \quad 3 \geq 2k+1$$

II. ?

coset header: coset 中 weight 最小的元素.

Decoding table



(a)
 $e(00) = 0000$

$$e(01) = 0101$$

$$e(10) = 1001$$

$$e(11) = 1100$$

(b).