

11.1 Groups and Coding

D: \oplus : 加和模2

D: $e: B^m \rightarrow B^n$ $\xrightarrow{\text{字空间}} \text{码字空间}$ encoding Function, m, n 为长度

D: weight 权值 $\rightarrow |x|_1$ 中的个数

D: Hamming distance: $|x_1 \oplus x_2| = s(x, y)$

↓
the weight of $x_1 \oplus x_2$ 加和模2.

I Theorem: ① $s(x, y) = s(y, x)$ ② $s(x, y) \geq 0$ ③ $s(x, y) = 0 \Leftrightarrow x = y$

④ ? 走神了.

D: Minimum distance of an encoding function e :

$$\min \{d(e(x), e(y)) \mid x, y \in B^m\}.$$

II Theorem: $e(\varphi) = \dots$
 \downarrow \downarrow
 字空间 码字空间
 $\xrightarrow{\text{最短距离}} \geq k+1$
 $\xrightarrow{\text{error}}$

III Theorem:

$e: B^m \rightarrow B^n$ be a group code,
 the min distance of e is the
 min weight of a nonzero code.

D: group code $e(B^m) = \{e(b) \mid e(b) \in B^n\} = \text{Ran}(e)$

$e: B^m \rightarrow B^n$ is a subgroup of B^n .

(group: ① $B^3 \rightarrow B^6$)
 对于封闭、有0, 有 x^{-1} , 可结合???

D: definition I, II ... Theorem.

D: Boolean \wedge of matrix

$$A \odot B = \dots \vee a \wedge b$$

Theorem:

奇偶校验矩阵

$m < n$, nonnegative, $r = n - m$, H be an $m \times n$ Boolean matrix

$$f_H: B^n \rightarrow B^r : f_H(x) = x * H, x \in B^n$$

is a homomorphism from Group B^n to B^r

$$\begin{array}{ccc} B^n & & B^r \\ \text{是} & \xleftarrow{x * H = 0} & \text{是} \\ \text{ZG} & N & b_1 \\ \text{2}^{n/r} \text{块} & \text{O} & b_r \\ & \vdots & \vdots \end{array}$$

$$\begin{array}{ccc} 2^{n/r} \text{块} & \text{O} & b_r \\ & \vdots & \vdots \end{array}$$

$$\begin{array}{ccc} & \text{O} & \cdot \\ & \vdots & \vdots \end{array}$$

$$x_1, x_2 \in B^n, \text{ 则 } f_H(x_1 \oplus x_2) = f_H(x_1) \oplus f_H(x_2)$$

Corollary:

$N = \{x \in B^n | x * H = 0\}$ is a normal subgroup of B^n .

- 证: ① 封闭 (Group)
② 单位元 0
③ 逆元 x'
④ $aH = Ha$

$$H \leq G \text{ 是 normal group, } \forall g \in G, h \in H, ghg^{-1} \in H / gH = Hg$$

$e_H: e^m \rightarrow e^n$

$b = b_1 b_2 \dots b_m, \quad x = e_H(b) = b_1 b_2 \dots b_m x_1 x_2 \dots x_r$

$$x_i = \underbrace{b_1 h_{1i}} + \underbrace{b_2 h_{2i}} + \dots + \underbrace{b_m h_{mi}}$$

$\vdots \qquad \text{第 } i \text{ 列}$

$$x_r$$

① m 的子空间 ② $n=5, \underline{\text{补3维}}$.

$\begin{smallmatrix} 1 & 0 \\ 0 & 1 \\ \vdots & \vdots \end{smallmatrix}$

b_1, b_2

II.2. Decoding and Error Correction. ↗ $\text{VBB} \rightarrow d_{\min} \xrightarrow{d} e$
 decoding F 檢驗 ???

I. e: (m, n) encoding F , d is a maximum likelihood decoding F ,
 then (e, d) can correct K or fewer errors



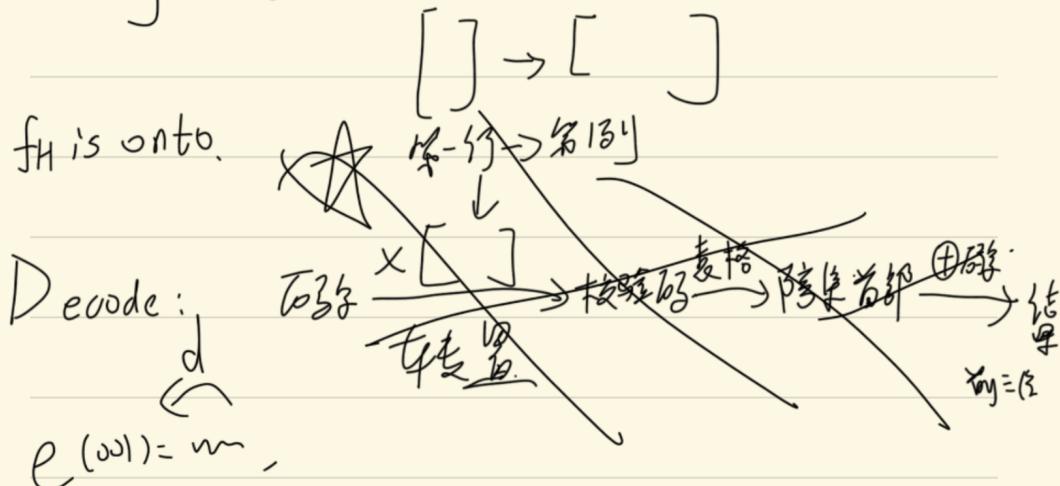
the minimum distance of e is $2k+1$

$$e^3 \rightarrow e^8 \quad 3 \geq 2k+1$$

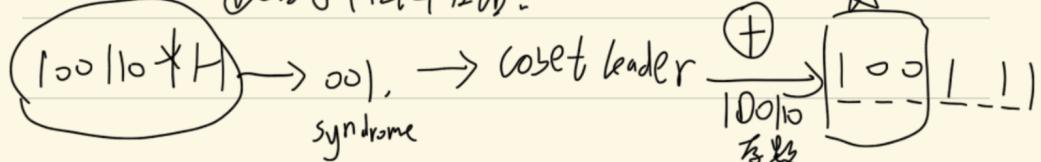
II. ?

coset header: coset 中 weight 最小的元素.

Decoding table



↓ 以对角线阵在场.



$$(a) \quad e(00) = 00000$$

$$e(01) = 01011$$

$$e(10) = 10011$$

$$e(11) = 11000$$

(b).