

Introduction to Linear Algebra

5th Edition

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Summary

Six great theorem in LA:

Dimension Th. All bases for a vec sp. have the same number ofvecs.

Counting Th. Dimension of row sp. + dimension of nullspace = numbers of cols.

Rank Th. Dimension of row sp. = dimension of col sp. This is the rank.

Fundamental Th. The row sp. and nullspace of A are otho. complements in \mathbb{R}^n .

SVD There are ortho. bases (v 's and u 's for r and c sp.) so that $Av_i = \sigma_i u_i$.

Spectral Th. If $A^T = A$ there are ortho. q 's so that $Aq_i = \lambda q_i$ and $A = Q\Lambda Q^T$.

LA in a nutshell: nonsingular singular

A is inv. rows/cols are ind.

A is not inv. rows/cols are dependent.

$\det A$ is not zero.

$\det A$ is zero.

$Ax = 0$ has one solution $x = 0$.

$Ax = 0$ has inf. many solutions.

$Ax = b$ has one solution $x = A^{-1}b$.

$Ax = b$ has no solution or inf. many.

A has n pivots. full rank $r = n$.

A has $r < n$ pivots. rank $r < n$.

rref $R = I$. row/col sp. is all of \mathbb{R}^n .

rref R has at least one zero row. row/col sp.

A 's all eigenvalues are nonzero.

zero is an eigenvalue of A . $\dim \leq n$.

$A^T A$ is sym. pos. has n (pos.) singular values.

$A^T A$ is only semipos. has $n < n$ singular values.

Matrix factorizations:

$$A = LU, A = LDU, PA = LU, EA = R, A = QR, A = X \Lambda X^{-1}, S = Q \Lambda Q^T$$

$$(S = LDL^T)$$

$$A = U \Sigma V^T = C^T C, A^T = V \Sigma^T U^T, A = Q S, A = B J B^{-1}$$

$$A = U \Lambda U^{-1}, A = Q T Q^H$$

$$(U \Lambda U^H)$$

Alphabet

A any matrix

L lower tri. matrix

U upper tri. matrix . left singular matrix

E elimination matrix (. eye matrix)

P permutation matrix . pascal matrix . projection matrix

I identity matrix

D diagonal matrix

D_n (\bar{D}_n / D_n) lower/upper difference matrix

R ref matrix . upper tri. matrix . reflection matrix . rotation matrix
(reduced row echelon form)

Q orthogonal matrix

C cofactor matrix (companion matrix) . correlation matrix

S symmetric matrix

X eigenvector matrix

Λ eigenvalue matrix

V right singular matrix

Σ singular value matrix

J Jordan matrix

B basis matrix

O.A.I & O.A.IIA 2

1. Introduction to Vectors

1.1 vectors and linear combinations

lin. comb.: $c\vec{v} + d\vec{w}$ fill a plane? solution: if right side is on the plane.

1.2 lengths and dot products

dot p: $\vec{v} \cdot \vec{w}$ perpendicular: $\vec{v} \cdot \vec{w} = 0$, length $\|\vec{v}\| = \sqrt{\vec{v} \cdot \vec{v}}$.

all vectors have $|\vec{v} \cdot \vec{w}| \leq \|\vec{v}\| \|\vec{w}\|$ angle $\theta = \cos^{-1} \frac{\vec{v} \cdot \vec{w}}{\|\vec{v}\| \|\vec{w}\|}$.

1.3 matrices

$A\vec{x}$: comb of the cols. components of $A\vec{x}$ are the dots of A's rows

equation: $A\vec{x} = \vec{b}$ solution could write $\vec{x} = A^{-1}\vec{b}$ and \vec{x} .

2. Solving Linear Equations

2.1 vectors and linear combinations $A\vec{x} = \vec{b}$

col picture: comb of n cols of A produces the vec \vec{b} .

row picture: m eqs from m rows give m planes meeting at \vec{x} .

$\vec{b} = 0$, at least one possibility is $\vec{x} = 0$.

2.2 the idea of elimination

$(\begin{array}{|ccc|} \hline a_{11} & & \\ \hline \downarrow & \downarrow & \downarrow \\ a_{n1} & & \end{array})$ pivot a_{ii} , multiplier $l_{ij} = \frac{a_{ji}}{a_{ii}}$ and then subtract it from j.

elimination (Gauss-) $A \rightarrow U$ elimination doesn't change

solutions.

$A\vec{x} = \vec{b}$: forward elimination + back substitution

elimination breaks down if zero appears in pivot.

exchanging two eqs may solve it. (zero is not allowed in pivot ← nonsingular)

when breakdown is permanent, $A\vec{x} = \vec{b}$ has no solution or inf. many.

$$\Leftrightarrow A^2 = I \cdot A \cdot O \quad A^3 = I \cdot A \cdot O$$

幂等矩阵的性质:

($\frac{1}{A} = A$) $\cdot A^2 = O$: $C(A)$ is contained in N_A
so $r \leq n - r$. (结论)

$\cdot A^2 = A$: if idem. sym. mt, we call proj.

(idempotent mt.) if idem. mt. full rank, its

$\text{tr}(A) = \text{rank}(A)$ $\otimes A \sim B$, B idem.
 $\det(A) = 0$ or 1. $A \sim (I_r, 0)$

$(A^2 = A \text{ odd idem. } \lambda = 0, \pm 1)$

vice versa!



• 引入 LA?

1. 线性空间 $\vec{x} = \alpha\vec{u} + \beta\vec{v} + \gamma\vec{w}$

$$\vec{x} = \begin{pmatrix} \alpha \\ \beta \\ \gamma \end{pmatrix} \quad \vec{u} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \quad \text{向量}$$

基底, 线性变换与变基 $A\vec{x} = \vec{b}$ inv. 范数 $\lambda = 0, 1$.

2. 分块矩阵. ($\frac{A}{B}$) 多向度 AB 矩阵乘法 一对 \downarrow $(A^2 = A \text{ odd idem. } \lambda = 0, \pm 1)$

• (课) 求行列式 直算 / 余子式 / 降阶 Big Formular Δ 递推 Δ 全排列

消元 \rightarrow 上下三角

大分拆 Big Partition Δ 剪子阵 (Δ) $= \begin{pmatrix} 0 \\ 1 \end{pmatrix}$

/ 加边 $\left| \begin{array}{cccc} & & 0 & 0 \\ & \ddots & & 0 \\ 0 & & \ddots & 0 \\ & & & 1 \end{array} \right|$ shear matrix

特殊阵

Vandermonde $V_n = \prod_{i>j} (x_i - x_j)$

cycle (n为偶数)

triangular (band) triD det: $D_n = aD_{n-1} - bcD_{n-2}$

(R, i.e. more properties.)

求逆

直算 / A^{-1}

解方程 $AA^{-1} = I$

初等变换 / 加边 $E \leftrightarrow E^T (L)$

(消元) $(AI) \leftrightarrow (IA^{-1})$ — G-J elimination

分块 $\begin{pmatrix} A & C \\ 0 & B \end{pmatrix}$ Δ Sherman-Morrison公式: $(A + uv^T)^{-1} = A^{-1} - \frac{A^T uv^T A}{1 + v^T A^T u}$

(见2.4)

等式变形 $A^n = I \cdot f(A)$

esp. $(I + uv^T)^{-1} = I - \frac{uv^T}{1 + v^T u}$

• 左行右列

行 $\overset{T}{\leftrightarrow}$ 列互证

• 矩阵等式 (张量代数) 证明

• 非方阵不同时存在左右逆: $A mxn$
(广义逆)

left inv., right inv. $n \times m$

comb 出 $(\vdots) \cdots (\vdots) \Rightarrow \text{rank } A \geq n, \text{rank } A \geq m$

且 $\text{rank } A \leq n, m \Rightarrow n = m$.

• 左逆 = 右逆

$$\text{left } \left\{ \begin{array}{l} A\vec{x} = \vec{b} \Rightarrow \vec{x} = A_1^{-1}\vec{b} \text{ 唯一解} \\ A\vec{x} = (AA_1^{-1})\vec{b} = \vec{b} \end{array} \right.$$

$$A\vec{b} \Rightarrow AA_1^{-1} = I, A_1^{-1} = A_r^{-1}$$

$$\text{inv. uniqueness: } A_1^{-1} = A_r^{-1}, I = A_1^{-1}AA_2^{-1} = A_2^{-1}$$

$$\text{right } \left\{ \begin{array}{l} \vec{x}_0 = A_r^{-1}\vec{b} \text{ 为一解} \\ A_r^{-1}A\vec{x}_0 = \vec{x}_0 \end{array} \right.$$

$$(或 T) \quad A_r^{-1}A = I, A_r^{-1} = A_1^{-1}$$

如采用 P,
则无法有
 $A\vec{x} = \vec{b}$

2.3 elimination using matrices

views of matrix multiplications : $A = (\bar{a}_1, \dots, \bar{a}_n)$, $A\bar{x} = x_1\bar{a}_1 + \dots + x_n\bar{a}_n$

$E(A\bar{b}) = (EA, E\bar{b})$ augmented matrix $AB = (A\bar{b}_1, \dots, A\bar{b}_n)$

elimination matrix / elementary transform (ET) : $E_{ij} = \begin{pmatrix} 1 & & \\ & \ddots & \\ & & 1 \\ -i & & 0 \end{pmatrix}$

permutation matrix $P_{ij} = \begin{pmatrix} 0 & \dots & 0 & 1 \\ \vdots & \ddots & \vdots & \vdots \\ 0 & \dots & 0 & 0 \end{pmatrix}$ (may used) \rightarrow could see from ① multiply / act vecs.

one matrix both steps : $PE\dots = E$ \rightarrow do same step onto I. rows/cols
times then. to create a $U = EA$

2.4 rules for matrix operations

$(m \times n)(n \times p) = (m \times p)$ mnp separated muls → A rows · B cols (mp entries)

$(AB)C = A(BC)$. $AB \neq BA$ usually → A cols · B rows (n adds)

block multiplication, if their shapes permit. → A rows · B (vec-mt)

block elimination. Schur complement: → A · B cols (int. vec)

so Ax → B rows recomb (comb)
 \rightarrow recomb of A cols → A cols recomb (comb)

$\rightarrow x$ after a linear transform (bases (dual) transform)

2.5 inverse matrices

invertible: $AA^{-1} = A^{-1}A = I$. (means $n \times n$) test invertibility:

① A must have n (nonzero) pivots.

(not sq. mt do not have inv., for transpass \rightarrow log:)

③ ② det A must not be zero.
over diff. IR ($n \leftrightarrow m$) inv. ④ $AX = 0, \bar{x} = 0$ must be the only solution.

if all invertible, $(AB)^{-1} = B^{-1}A^{-1}$

Gauss-Jordan elimination:

$O_{(G-J)}$

$\approx n^3$

③ $A\bar{x} = \bar{b}, \bar{x} = A^{-1}\bar{b}$ be the only solution, not none or inf.)

(for means back elimination, the same essence)

④ row/cols of A are independent. = n

$$\square \quad \begin{pmatrix} -\frac{1}{2}, 0 \\ 0, \frac{1}{2} \end{pmatrix}^{-1} = \begin{pmatrix} \frac{1}{2}, 0 \\ 0, \frac{1}{2} \end{pmatrix}$$

so: how you view
a trans?

$\circ \bar{x} = A^{-1}\bar{b}$'s a math trans!

* elim' is no interception (coupling, rolling)

1. acts on row or col? / how it produced?

elim is rolling.

2. view as E or directly E^{-1} ?

3. start from tip or top? one can be rolling.

to get an inv., must go through rolling once.
(turn an unrolling way back)

矩阵相乘 ($A = B \cdot \text{inv}(B)$) 与 行/列空间的基 (combs)

- useful justify dependent/uninv.: $r/c s$ adds to zero, or one has combs of others.

blockize: better
sq. in diag.

矩阵收集 (combs):

$$\text{adjacency matrix } S = \begin{pmatrix} 0 & 1 & \dots \\ 1 & 0 & \dots \\ \vdots & \vdots & \ddots \end{pmatrix}$$

(无向图) 为 sym. $S^2 \dots S^n$: length n, node i $\rightarrow j$

$$\text{sum matrix } S = \begin{pmatrix} n & 0 \\ 0 & 1 \end{pmatrix}$$

$$\text{difference matrix } \bar{D} = \begin{pmatrix} 0 & 1 & 0 & \dots \\ 1 & 0 & -1 & \dots \\ 0 & -1 & 0 & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}, D = \begin{pmatrix} 1 & 0 & \dots \\ 0 & 1 & -1 & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}, D^{-1} = \begin{pmatrix} 1 & 0 & \dots \\ 0 & 1 & 1 & \dots \\ \vdots & \vdots & \vdots & \ddots \end{pmatrix}$$

n .th diff matrix D_n

$$D_n^{-1} = S^n$$

- cycle matrix (zero-cycle matrix means adds equal to zero,

$$C = \begin{pmatrix} 0 & 1 & \dots \\ 1 & 0 & \dots \\ \vdots & \vdots & \ddots \end{pmatrix} \text{ it's uninv.}$$

- Pascal matrix / binomial coeff. matrix

$$P_L = \begin{pmatrix} 1 & 1 & 1 & \dots \\ 1 & 2 & 1 & \dots \\ 1 & 3 & 3 & 1 & \dots \\ 1 & 4 & 6 & 4 & \dots \\ 1 & 5 & 10 & 10 & 5 & \dots \\ 1 & 6 & 15 & 20 & 15 & 6 & \dots \\ 1 & 7 & 21 & 35 & 35 & 21 & 7 & \dots \\ 1 & 8 & 28 & 56 & 70 & 56 & 28 & 8 & \dots \\ 1 & 9 & 36 & 84 & 126 & 126 & 84 & 36 & 9 & \dots \\ 1 & 10 & 45 & 120 & 210 & 252 & 210 & 120 & 45 & 10 & \dots \\ 1 & 11 & 55 & 165 & 330 & 462 & 462 & 330 & 165 & 55 & 11 & \dots \\ 1 & 12 & 66 & 220 & 495 & 792 & 924 & 792 & 495 & 220 & 66 & 12 & \dots \\ 1 & 13 & 78 & 286 & 715 & 1287 & 1716 & 1716 & 1287 & 715 & 286 & 78 & 13 & \dots \\ 1 & 14 & 91 & 364 & 1001 & 2002 & 3003 & 3432 & 3003 & 2002 & 1001 & 364 & 91 & 14 & \dots \\ 1 & 15 & 105 & 455 & 1365 & 3465 & 6435 & 9450 & 9450 & 6435 & 3465 & 1365 & 455 & 105 & \dots \\ 1 & 16 & 120 & 560 & 1701 & 4804 & 12650 & 24310 & 3150 & 24310 & 12650 & 4804 & 1701 & 560 & 120 & \dots \\ 1 & 17 & 136 & 680 & 2380 & 7315 & 19360 & 45580 & 7315 & 19360 & 45580 & 7315 & 19360 & 2380 & 680 & 136 & \dots \\ 1 & 18 & 153 & 816 & 3060 & 10200 & 27300 & 64350 & 12650 & 64350 & 12650 & 27300 & 10200 & 3060 & 816 & 153 & \dots \\ 1 & 19 & 171 & 960 & 3720 & 13200 & 34320 & 84700 & 17160 & 84700 & 17160 & 34320 & 13200 & 3720 & 960 & 171 & \dots \\ 1 & 20 & 190 & 1120 & 4800 & 16800 & 48000 & 114400 & 21000 & 114400 & 21000 & 48000 & 16800 & 4800 & 1120 & 190 & \dots \\ 1 & 21 & 210 & 1300 & 6300 & 21000 & 67200 & 168000 & 33600 & 168000 & 33600 & 67200 & 21000 & 6300 & 1300 & 210 & \dots \\ 1 & 22 & 231 & 1500 & 7800 & 25200 & 84000 & 210000 & 44800 & 210000 & 44800 & 84000 & 25200 & 7800 & 1500 & 231 & \dots \\ 1 & 23 & 253 & 1700 & 9300 & 30600 & 102000 & 252000 & 50400 & 252000 & 50400 & 102000 & 30600 & 9300 & 1700 & 253 & \dots \\ 1 & 24 & 276 & 1900 & 11000 & 37200 & 120000 & 312000 & 62400 & 312000 & 62400 & 120000 & 37200 & 11000 & 1900 & 276 & \dots \\ 1 & 25 & 300 & 2100 & 13000 & 45000 & 150000 & 400000 & 80000 & 400000 & 80000 & 150000 & 45000 & 13000 & 2100 & 300 & \dots \\ 1 & 26 & 325 & 2300 & 15000 & 54000 & 180000 & 504000 & 120000 & 504000 & 120000 & 180000 & 54000 & 15000 & 2300 & 325 & \dots \\ 1 & 27 & 351 & 2500 & 17000 & 63000 & 210000 & 616000 & 168000 & 616000 & 168000 & 210000 & 63000 & 17000 & 2500 & 351 & \dots \\ 1 & 28 & 378 & 2700 & 19000 & 72000 & 240000 & 720000 & 192000 & 720000 & 192000 & 240000 & 72000 & 19000 & 2700 & 378 & \dots \\ 1 & 29 & 406 & 2900 & 21000 & 81000 & 270000 & 840000 & 216000 & 840000 & 216000 & 270000 & 81000 & 21000 & 2900 & 406 & \dots \\ 1 & 30 & 435 & 3100 & 23000 & 90000 & 300000 & 960000 & 240000 & 960000 & 240000 & 300000 & 90000 & 23000 & 3100 & 435 & \dots \\ 1 & 31 & 465 & 3300 & 25000 & 100000 & 330000 & 1080000 & 288000 & 1080000 & 288000 & 330000 & 100000 & 25000 & 3300 & 465 & \dots \\ 1 & 32 & 500 & 3500 & 27000 & 110000 & 360000 & 1200000 & 312000 & 1200000 & 312000 & 360000 & 110000 & 27000 & 3500 & 500 & \dots \\ 1 & 33 & 536 & 3700 & 29000 & 120000 & 400000 & 1360000 & 408000 & 1360000 & 408000 & 400000 & 120000 & 29000 & 3700 & 536 & \dots \\ 1 & 34 & 574 & 3900 & 31000 & 130000 & 440000 & 1520000 & 456000 & 1520000 & 456000 & 440000 & 130000 & 31000 & 3900 & 574 & \dots \\ 1 & 35 & 613 & 4100 & 33000 & 140000 & 480000 & 1680000 & 512000 & 1680000 & 512000 & 480000 & 140000 & 33000 & 3700 & 613 & \dots \\ 1 & 36 & 653 & 4300 & 35000 & 150000 & 520000 & 1840000 & 544000 & 1840000 & 544000 & 520000 & 150000 & 35000 & 3500 & 653 & \dots \\ 1 & 37 & 694 & 4500 & 37000 & 160000 & 560000 & 1920000 & 568000 & 1920000 & 568000 & 560000 & 160000 & 37000 & 3500 & 694 & \dots \\ 1 & 38 & 736 & 4700 & 39000 & 170000 & 600000 & 1980000 & 584000 & 1980000 & 584000 & 580000 & 170000 & 39000 & 3500 & 736 & \dots \\ 1 & 39 & 779 & 4900 & 41000 & 180000 & 640000 & 2040000 & 616000 & 2040000 & 616000 & 580000 & 180000 & 41000 & 3500 & 779 & \dots \\ 1 & 40 & 823 & 5100 & 43000 & 190000 & 680000 & 2100000 & 648000 & 2100000 & 648000 & 580000 & 190000 & 43000 & 3500 & 823 & \dots \\ 1 & 41 & 868 & 5300 & 45000 & 200000 & 720000 & 2160000 & 680000 & 2160000 & 680000 & 580000 & 190000 & 45000 & 3500 & 868 & \dots \\ 1 & 42 & 914 & 5500 & 47000 & 210000 & 760000 & 2220000 & 712000 & 2220000 & 712000 & 580000 & 190000 & 47000 & 3500 & 914 & \dots \\ 1 & 43 & 961 & 5700 & 49000 & 220000 & 800000 & 2280000 & 744000 & 2280000 & 744000 & 580000 & 190000 & 49000 & 3500 & 961 & \dots \\ 1 & 44 & 1009 & 5900 & 51000 & 230000 & 840000 & 2340000 & 776000 & 2340000 & 776000 & 580000 & 190000 & 51000 & 3500 & 1009 & \dots \\ 1 & 45 & 1058 & 6100 & 53000 & 240000 & 880000 & 2400000 & 808000 & 2400000 & 808000 & 580000 & 190000 & 53000 & 3500 & 1058 & \dots \\ 1 & 46 & 1108 & 6300 & 55000 & 250000 & 920000 & 2460000 & 840000 & 2460000 & 840000 & 580000 & 190000 & 55000 & 3500 & 1108 & \dots \\ 1 & 47 & 1160 & 6500 & 57000 & 260000 & 960000 & 2520000 & 872000 & 2520000 & 872000 & 580000 & 190000 & 57000 & 3500 & 1160 & \dots \\ 1 & 48 & 1213 & 6700 & 59000 & 270000 & 1000000 & 2580000 & 904000 & 2580000 & 904000 & 580000 & 190000 & 59000 & 3500 & 1213 & \dots \\ 1 & 49 & 1267 & 6900 & 61000 & 280000 & 1040000 & 2640000 & 936000 & 2640000 & 936000 & 580000 & 190000 & 61000 & 3500 & 1267 & \dots \\ 1 & 50 & 1323 & 7100 & 63000 & 290000 & 1080000 & 2700000 & 968000 & 2700000 & 968000 & 580000 & 190000 & 63000 & 3500 & 1323 & \dots \\ 1 & 51 & 1380 & 7300 & 65000 & 300000 & 1120000 & 2760000 & 1000000 & 2760000 & 1000000 & 580000 & 190000 & 65000 & 3500 & 1380 & \dots \\ 1 & 52 & 1440 & 7500 & 67000 & 310000 & 1160000 & 2820000 & 1032000 & 2820000 & 1032000 & 580000 & 190000 & 67000 & 3500 & 1440 & \dots \\ 1 & 53 & 1501 & 7700 & 69000 & 320000 & 1200000 & 2880000 & 1064000 & 2880000 & 1064000 & 580000 & 190000 & 69000 & 3500 & 1501 & \dots \\ 1 & 54 & 1564 & 7900 & 71000 & 330000 & 1240000 & 2940000 & 1096000 & 2940000 & 1096000 & 580000 & 190000 & 71000 & 3500 & 1564 & \dots \\ 1 & 55 & 1630 & 8100 & 73000 & 340000 & 1280000 & 3000000 & 1128000 & 3000000 & 1128000 & 580000 & 190000 & 73000 & 3500 & 1630 & \dots \\ 1 & 56 & 1700 & 8300 & 75000 & 350000 & 1320000 & 3060000 & 1160000 & 3060000 & 1160000 & 580000 & 190000 & 75000 & 3500 & 1700 & \dots \\ 1 & 57 & 1772 & 8500 & 77000 & 360000 & 1360000 & 3120000 & 1192000 & 3120000 & 1192000 & 580000 & 190000 & 77000 & 3500 & 1772 & \dots \\ 1 & 58 & 1846 & 8700 & 79000 & 370000 & 1400000 & 3180000 & 1224000 & 3180000 & 1224000 & 580000 & 190000 & 79000 & 3500 & 1846 & \dots \\ 1 & 59 & 1922 & 8900 & 81000 & 380000 & 1440000 & 3240000 & 1256000 & 3240000 & 1256000 & 580000 & 190000 & 81000 & 3500 & 1922 & \dots \\ 1 & 60 & 1999 & 9100 & 83000 & 390000 & 1480000 & 3300000 & 1288000 & 3300000 & 1288000 & 580000 & 190000 & 83000 & 3500 & 1999 & \dots \\ 1 & 61 & 2078 & 9300 & 85000 & 400000 & 1520000 & 3360000 & 1320000 & 3360000 & 1320000 & 580000 & 190000 & 85000 & 3500 & 2078 & \dots \\ 1 & 62 & 2158 & 9500 & 87000 & 410000 & 1560000 & 3420000 & 1352000 & 3420000 & 1352000 & 580000 & 190000 & 87000 & 3500 & 2158 & \dots \\ 1 & 63 & 2240 & 9700 & 89000 & 420000 & 1600000 & 3480000 & 1384000 & 3480000 & 1384000 & 580000 & 190000 & 89000 & 3500 & 2240 & \dots \\ 1 & 64 & 2323 & 9900 & 91000 & 430000 & 1640000 & 3540000 & 1416000 & 3540000 & 1416000 & 580000 & 190000 & 91000 & 3500 & 2323 & \dots \\ 1 & 65 & 2408 & 10100 & 93000 & 440000 & 1680000 & 3600000 & 1448000 & 3600000 & 1448000 & 580000 & 190000 & 93000 & 3500 & 2408 & \dots \\ 1 & 66 & 2494 & 10300 & 95000 & 450000 & 1720000 & 3660000 & 1480000 & 3660000 & 1480000 & 580000 & 190000 & 95000 & 3500 & 2494 & \dots \\ 1 & 67 & 2582 & 10500 & 97000 & 460000 & 1760000 & 3720000 & 1512000 & 3720000 & 1512000 & 580000 & 190000 & 97000 & 3500 & 2582 & \dots \\ 1 & 68 & 2672 & 10700 & 99000 & 470000 & 1800000 & 3780000 & 1544000 & 3780000 & 1544000 & 580000 & 190000 & 99000 & 3500 & 2672 & \dots \\ 1 & 69 & 2763 & 10900 & 101000 & 480000 & 1840000 & 3840000 & 1576000 & 3840000 & 1576000 & 580000 & 190000 & 101000 & 3500 & 2763 & \dots \\ 1 & 70 & 2856 & 11100 & 103000 & 490000 & 1880000 & 3900000 & 1608000 & 3900000 & 1608000 & 580000 & 190000 & 103000 & 3500 & 2856 & \dots \\ 1 & 71 & 2951 & 11300 & 105000 & 500000 & 1920000 & 3960000 & 1640000 & 3960000 & 1640000 & 580000 & 190000 & 105000 & 3500 & 2951 & \dots \\ 1 & 72 & 3048 & 11500 & 107000 & 510000 & 1960000 & 4020000 & 1672000 & 4020000 & 1672000 & 580000 & 190000 & 107000 & 3500 & 3048 & \dots \\ 1 & 73 & 3147 & 11700 & 109000 & 520000 & 2000000 & 4080000 & 1704000 & 4080000 & 1704000 & 580000 & 190000 & 109000 & 3500 & 3147 & \dots \\ 1 & 74 & 3248 & 11900 & 111000 & 530000 & 2040000 & 4140000 & 1736000 & 4140000 & 1736000 & 580000 & 190000 & 111000 & 3500 & 3248 & \dots \\ 1 & 75 & 3351 & 12100 & 113000 & 540000 & 2080000 & 4200000 & 1768000 & 4200000 & 1768000 & 580000 & 190000 & 113000 & 3500 & 3351 & \dots \\ 1 & 76 & 3456 & 12300 & 115000 & 550000 & 2120000 & 4260000 & 1800000 & 4260000 & 1800000 & 580000 & 190000 & 115000 & 3500 & 3456 & \dots \\ 1 & 77 & 3563 & 12500 & 117000 & 560000 & 2160000 & 4320000 & 1832000 & 4320000 & 1832000 & 580000 & 190000 & 117000 & 3500 & 3563 & \dots \\ 1 & 78 & 3672 & 12700 & 119000 & 570000 & 2200000 & 4380000 & 1864000 & 4380000 & 1864000 & 580000 & 190000 & 119000 & 3500 & 3672 & \dots \\ 1 & 79 & 3783 & 12900 & 121000 & 580000 & 2240000 & 4440000 & 1896000 & 4440000 & 1896000 & 580000 & 190000 & 121000 & 3500 & 3783 & \dots \\ 1 & 80 & 3896 & 13100 & 123000 & 590000 & 2280000 & 4500000 & 1928000 & 4500000 & 1928000 & 580000 & 190000 & 123000 & 3500 & 3896 & \dots \\ 1 & 81 & 4011 & 13300 & 125000 & 600000 & 2320000 & 4560000 & 19$$

LU LDU
(triangular / Doolittle / Crout decomposition)

2.6 elimination = factorization: $A = LU$

the whole forward elimination process (with no row exchanges) is inverted by L , which still lower triangular, every multiplier l_{ij} stands explicitly at r_{icj} .

and to unitize, pops D (a diag) out.

uniqueness: $L_1 D_1 U_1 = L_2 D_2 U_2$

solving a tri system: $O(\frac{n^2}{2})$ (lu itself isn't unique) leads to $L_1^{-1} L_2 D_2 = D_1 U_1 U_2^{-1}$
(backward)

elimination / solving $A\vec{x} = \vec{b}$: $O(\frac{n^3}{3})$ $L = U$, must be diag!
 $D = A$ so $L_1^{-1} L_2 = I$, $L_1 = L_2 \dots$

o again: echelon! no interact.

($\xrightarrow{\text{row3}} \text{row3 of } U = \text{row3 of } A - l_{31} \text{row1 of } U - l_{32} \text{row2 of } U$)

that's $\text{row3 of } A = (l_{31} \ l_{32} \ 1) \cdot \begin{pmatrix} \text{row1 of } U \\ \text{row2 of } U \\ \text{row3 of } U \end{pmatrix}$
 U' acts (r/c?)

the algorithm of $A\vec{x} = \vec{b}$: 1. factor into $A = LU$ (muler stored in L) $\frac{n^3}{3} \ n n^2$
2. solve $Lc = b$ ($b \xrightarrow{L^{-1}} c$) n^2 $2n^2$
 $Ux = c$ (backward substitution)

sparse: faster. for a random/full mt. $n=1000$ on a PC takes 1 sec.

$\Delta \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 2 & 1 \\ 0 & 0 & 1 & 2 \\ 0 & 0 & 0 & 1 \end{pmatrix} = \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix} \begin{pmatrix} 1 & 1 & 0 & 0 \\ 0 & 1 & 1 & 0 \\ 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & 1 \end{pmatrix}$ the zeros at start (r or c) then n^3 increases
will still be reserved in L or U . the band width.

□ Cholesky factorization (sq root factorization) $A = LL^T$, when A is sym. and positive.
(lemma: B is triang., $A = BB^T$ is sym.pos. for $x^T Ax = (xB)^T (xB) > 0$)

$A = LDU = A^T = UTDL^T \rightarrow L = U^T$ (B's row axes)

compare: $LU \quad n^2 + cn - 1 + \dots \sim \frac{n^3}{3}$ $A = (L\sqrt{D})(L\sqrt{D})^T$ (lu chooses the largest pivot in col in algorithm.)

Chol $\frac{n^3}{2} + \frac{(n-1)^2}{2} + \dots \sim \frac{n^3}{6}$ half ← sym.'s credit.

algo: $A = (A_{11} \ A_1^T \ A_1) = (L_{11} \ L_1^T \ L_1) (L_{11} \ L_1^T \ L_1^T) \sim$ to reduce error.
(recurs) since $LUX = b$ not stable

$L' L'^T = A' - L_1 L_1^T$ — recursion. (big swallow small)

$O(\frac{n^2}{2})$ (end: $\sqrt{A_{11}}$, that's $\sqrt{A_{nn}}$)

L sym.

• 短阵收集仓 (2):

- Toeplitz matrix $T = \begin{pmatrix} t_0 & t_1 & \dots & t_{n-1} \\ t_1 & t_0 & \dots & t_{n-2} \\ \vdots & \vdots & \ddots & \vdots \\ t_{n-1} & t_{n-2} & \dots & t_0 \end{pmatrix}$ persymmetric: $A^T = JAJ$ (由 π)
- reversal matrix $J = \begin{pmatrix} 0 & 1 & \dots & 1 \\ 1 & 0 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 1 & 0 & \dots & 0 \end{pmatrix}$ persym. $^{-1}$ is also persym.

cyclic matrix is a esp. Toeplitz matrix.

□ full-rank decomposition / CR factorization

满秩分解

- $A = CR$ C contains r independent cols of A . (a base of $\mathcal{C}(A)$)

$R = (I F)P$, $A = (C CF)P$ F expresses the remaining $n-r$ dependent cols

(P permutes them according to A)

well, the rows of R can be a base of row space of A . ($\mathcal{R}(A)$) (expressed by C)

row R is independent, for it contains I .

In space view, A :

$$\begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \\ 7 & 8 & 9 \end{pmatrix} = \begin{pmatrix} 1 & 2 \\ 4 & 5 \\ 7 & 8 \end{pmatrix} \begin{pmatrix} 1 & 0 & -1 \\ 0 & 1 & 2 \end{pmatrix} = \begin{pmatrix} 1 & 2 \\ 4 & 5 \\ 7 & 8 \end{pmatrix} \begin{pmatrix} 1 & 2 \\ 4 & 5 \end{pmatrix}^{-1} \begin{pmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{pmatrix} \quad \mathbb{R}^n \rightarrow \mathbb{R}^r \xrightarrow{W^{-1}} \mathbb{R}^r \rightarrow \mathbb{R}^m$$

$R(A)$ $\mathcal{C}(A)$

- $A = CW^{-1}B$ B also contains r actual rows of A .

$$\text{if } P=I, A = \begin{pmatrix} W & H \\ J & K \end{pmatrix} = \begin{pmatrix} W \\ J \end{pmatrix} W^{-1} (W H) \quad (\text{so } R = W^{-1}B = (I, W^{-1}H), B = WR)$$

W is $m \times r$, where B meets C at $r \times r$ positions (base), which called 'r-c intersection'.

(for a sq. and inv. A , it reduced to $A = WW^{-1}W$ or $A = C = B = W$)

(considering P is nowhere any diff, you can AP rather than A , or $CW^{-1}B$ not influence)

. Uniqueness: C is one basis you choose, then R is unique, but W, B aren't unique.

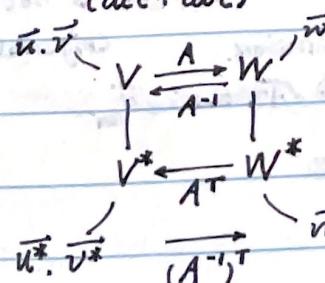
This is consistent with the uniqueness of the ref R_0 , which has

$$\square' \text{ dual } \vec{w} = a\vec{u} + b\vec{v} \quad \begin{pmatrix} a \\ b \end{pmatrix} = \begin{pmatrix} u_1 & v_1 \\ u_2 & v_2 \end{pmatrix}^{-1} \begin{pmatrix} w_1 \\ w_2 \end{pmatrix} = \begin{pmatrix} u^* \cdot \vec{w} \\ v^* \cdot \vec{w} \end{pmatrix} = \begin{pmatrix} u^* \\ v^* \end{pmatrix} \vec{w} \quad \text{non-uniqueness steps from } A \text{ to } R_0$$

dual. ortho. basis (mutual) $u^*(\vec{u}) = 1, u^*(\vec{v}) = 0, v^*(\vec{u}) = 0, v^*(\vec{v}) = 1$

linear functional ~ row vec. $u^*(\vec{u}) = 0, v^*(\vec{v}) = 1, v^*(\vec{w}) = b$

(act: dot)



$$\Rightarrow \vec{u}^* = \begin{pmatrix} \frac{v_2}{\langle u, v \rangle} \\ -\frac{v_1}{\langle u, v \rangle} \end{pmatrix}^T \quad \vec{v}^* = \begin{pmatrix} -\frac{u_2}{\langle u, v \rangle} \\ \frac{u_1}{\langle u, v \rangle} \end{pmatrix}^T$$

$\cdot (\varphi u, v) = (u, \varphi^{\text{adj}} v)$ $\varphi: V \rightarrow W$
on the dual basis, $\varphi \cdot \varphi^{\text{adj}} = \text{id}_V$ $\varphi^{\text{adj}} \cdot \varphi = \text{id}_W$

induced mt. are transpose.

pull-back:
 $A: V \rightarrow W$, for W .
 $(A^T f)(x) = f(A(x))$

2.7 transposes and permutations

$$(Ax)^T = x^T A^T, (AB)^T = B^T A^T, (A^{-1})^T = (A^T)^{-1} \text{ note as } A^{-T}$$

inner product. $x \cdot y = x^T y = y^T x$ the idea behind A^T is $Ax \cdot y = x \cdot A^T y$

symmetric matrix $S^T = S$. antisymmetric matrix $S^T = -S$ (skew-sym.)

and $A^T A$ always sym. orthogonal matrix $Q^T = Q^{-1}$ ($QQ^T = I$)
(and I_n is $S = LDL^T$)

it's standard cunit/orthonormal

cols of Q are ortho. unit vectors. (ortho axes)

permu. mt. P puts x_1, \dots, x_n in new order, and $P^T = P^{-1}$, and can be combed by
($n!$ kinds, half odd and half even) several P_{ij} .

3. Vector Spaces and Subspaces

$$\begin{pmatrix} 0 & 1 \\ 1 & 0 \\ 0 & 0 \end{pmatrix} \Phi \begin{matrix} 3 \\ 1 \\ 2 \\ 2 \\ 3 \\ 1 \end{matrix} \quad (P_{ij} = P_{ij}^T = P_{ij}^{-1})$$

3.1 spaces of vectors Vector Space $S \leftarrow$ rec can be matrices even funcs of x .

\mathbb{R}^n contains all real col.vec with n components. if $v, w \in S$, every comb $cv + dw$

Subspace of \mathbb{R}^n is a vecspace inside \mathbb{R}^n , closure to adds and must be in S .

(\mathbb{Z} : one-point space consists of $x=0$.) must contain zero.

col space of A contains all combs of A 's cols, which is a subspace of \mathbb{R}^n .

col space contains all vecs Ax , so $Ax = b$ is solvable when b is in $C(A)$

we call it's spanned by A 's cols. $C(A)$: Im A (image) $N(A)$: Ker A (kernel)

3.2 the nullspace of A : solving $Ax = 0$ and $Rx = 0$

null space $N(A)$ (a subspace in \mathbb{R}^n) contains all solution x to $Ax = 0$, and it must contain

elimination from A to U to R does not change the nullspace $N(A) = N(U) = N(R)$

$R = \text{ref}(A)$ means reduced row echelon form, has all pivots = 1, with zeros above and

numbers of pivots = numbers of nonzero rows in R = rank R (also rank of cols.)

every matrix with $m < n$ has nonzero

There are $n-r$ free cols. their combs is the

solutions in its N .

so we have $\dim C(A) + \dim N(A) = n$.

(no pivot in) (the comb to produce

itself. or $x_j = 1$)

(for R (pivot rows and cols) contains I , $\mathcal{C} \rightarrow \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$)

complete solution of $Ax = 0$ or

$N(A)$ ($Rx = 0$)

The rank of A is the true size of A or a linear system.

'special solution'

A and U and R have r ind. rows/cols.

□ full-rank factor: $A = (\text{pivot cols of } A)(\text{first } r \text{ rows of } A)$ $m \times n = (m \times r)(r \times n)$

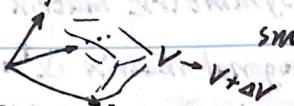
$S =$
- free colvec
with $x_j = 1$ and
zeros.

• 几个想问 LA?

1. basic: 基、变换: 特征向量、行列式: Cramer's rule 4. SVD (if so)

2. det and tr: $\det = \prod \lambda_i$, $\text{tr} = \sum \lambda_i$, $\text{tr} = \det'$

① $y' = Ay$, $y = e^{tA} y_0$ in small period, trace is the velocity of volume change.

$|e^{tA}| = 1 + t \cdot \text{tr} A + O(t^2)$  small changes mainly in the edges, on its own direction.

② intrinsic (coordinate free):

mt. is depicted wholly by traces, $\text{tr}(A) + \dots + \text{tr}(A^n) = \lambda_1^n + \dots + \lambda_n^n$.

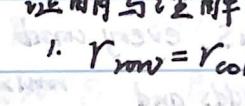
③ fix-point: $\text{tr} \phi = \sum_i \langle \phi | e_i \rangle, e_i \rangle$ (like 'how much remained?')

esp. $A^2 = A$. the rank of $A = A$'s dim of fix-point spanned space = $\text{tr} A$ for full fix, $\langle \phi, e_i \rangle = 1$.

3. transpose: $A^T x$: x 向列向量基上投影 (实为点乘), 注意与逆区分.

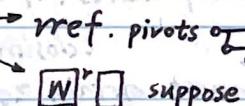
用矩阵解释: ① $\langle x, Ay \rangle = \langle A^T x, y \rangle$ ② $(AB)^T = B^T A^T$ ③ $Q^T = Q^{-1}$, $Q^T Q = I$ (not unit: D^2)

• 几个证明与理解:

1. $r_{\text{row}} = r_{\text{col}}$  full-rank decomp.

④ $(A^{-1})^T = (A^T)^{-1}$

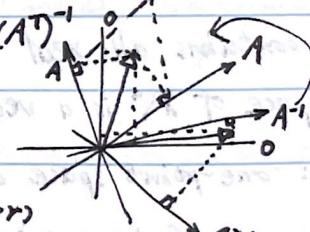
2. $|A| = |A^T|$

 rref. pivots

suppose $r_{\text{row}} < r_{\text{col}}$.

then the remaining cols ($>r$)

could be expressed by W (rank = r)



geometry analogy

LU/rref

SVD

理解

全排列

\vec{x}

$A = C \cdot B$

$m \times n \quad m \times r \quad r \times n$

且 r_{col} 组成 C , 用 B 表达.

反过来 r 行被 C 组成 A ,

$r_{\text{row}}(A) \leq c = r_{\text{col}}(A)$.

同理 $r_{\text{col}}(A) \leq r_{\text{row}}(A)$, 两数相等!

since $\langle \vec{x}_1, A \vec{x}_2 \rangle = \langle A \vec{x}_1, \vec{x}_2 \rangle$, if \vec{x}_2 is eigen direction, then $\vec{x}_1 \cdot \vec{x}_2 = 0 \Rightarrow A \vec{x}_1 \cdot A \vec{x}_2 = \lambda^2 \vec{x}_1 \cdot \vec{x}_2 = 0$.

so A maintains the ortho of eigen direction. sym. mt's eigen directions are ortho basis,

(\vec{x}_1 also is) its diagonalize just rotate the axes.
~~won't keep the origin direc.~~

(if $\lambda_1 = \lambda_2$, it's true that the whole plane is eigen.)

so $S = Q A Q^T$ orthonormal

another algebraic proof (之后)
and 2 for λ unequal.

• 更多关于 R 的:

R could be write as $\begin{pmatrix} I & F \\ 0 & 0 \end{pmatrix} P$.

'row-col reduced form' $R' = (\text{rref}(R^T))^T = \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix}$

• $A^T A$ always has the same N with A , ($A^T A x = 0$ 与 $A x = 0$ 为同解方程组)

proof: ① $A x_0 = 0 \Rightarrow A^T(A x_0) = 0$

② x_0 是正交性或写出 R 的形式.

$A^T A x_0 = 0 \Rightarrow x_0^T A^T A x_0 = 0, (A x_0)^T (A x_0) = 0, A x_0 = 0$.

$(C x_0) \perp N(A^T A)$,

► Generally, $r(C A^T A) = r(A)$ ($= r(A A^T)$) ($= r(I) = r(\Sigma)$)

so $A^T A x = 0$ only when $A x = 0$

3.3 the complete solution to $Ax = b$

- again: If a row/col contains a pivot, it's not a comb of previous rows/cols, otherwise is. $\mathcal{C}(A)$ tells us which row/col are combs of earlier ones; R tells us what those combs are and the special solution to $Ax=0$. R is always the same (determined by A).

Three fundamental subspaces: $\mathcal{C}(A)$ — choose the pivot cols of A as a basis.

Counting Theorem: $\mathcal{R}(A)$ — choose the nonzero rows of R as a basis.

$r = r_{\text{pivot/ind. vars}} + n - r_{\text{free vars}}$ $\mathcal{N}(A)$ — choose the special solution to $Rx=0$. ($Ax=0$)

(AI) $\rightarrow (R, E)$ will virtually tell you everything about A , when A is sq. and inv., R is I and E is A^{-1} .

pivot vars are determined after free vars are chosen.

Complete solution to $Ax = b$: $x = (\text{one particular solution } x_p) + (\text{any } x_n \text{ in nullspace})$
 $(A\bar{b}) \rightarrow (R\bar{d})$ (thus comb of special/solution basis)

$Ax = b / Rx = d$ is solvable only when all zero rows of R have zeros in \bar{d} .

full row rank $r = m$ when its col space $\mathcal{C}(A)$ is \mathbb{R}^m , $Ax = b$ is always solvable, but may many.

full col rank $r = n$ when its nullspace $\mathcal{N}(A) = \mathbb{Z}$, no free vars, one solution or none.

When is solvable, one particular x_p can be all free vars equal to zero, pivot vars from \bar{d} .

To conclude, there are four cases:

ranks	$r = m = n$	$r = m < n$	$r = n < m$	$r < m, r < n$
types of R	(I)	(I F)	(I) (0)	(I F) (0 0)
conclusion	A is inv. every $Ax = b$ is solvable,	$Ax = b$ has but inf.	$Ax = b$ has 0 or ∞ solutions	0 or 0 solutions

- if $Ax = b$ and $Cx = b$ have same solutions for every b , that means $A = C$.
(proof: let $x = (1, 0, 0, \dots)$ to get $\text{colvec } 1_A = C$)

3.4 independence, basis and dimension

linear independence: $x_1\vec{v}_1 + x_2\vec{v}_2 + \dots + x_n\vec{v}_n = 0$ only happens when all x 's are zero.

Any set of n vecs in \mathbb{R}^m must be dependent if $n > m$.

basis: (linearly) independent vectors that span the space, like pivot cols of A for $\mathcal{C}(A)$.
every vector in the space is a unique comb of the basis vecs. (empty set for \mathbb{Z})

dimension: All bases for a space have the same number of vecs called dim, like r_A for $\mathcal{C}(A)$.

Basis of \vec{v}_i 's from basis of \vec{w}_i 's when the change of basis matrix is inv. $V = WB$

dim of outputs + dim of nullspace = dim of inputs. $\Delta \dim(V+W) + \dim(V \cap W) = \dim V + \dim W$

- $\vec{x} = (x_1, x_2, \dots, x_n)$ 的全排列 ($n!$ 个) 3张成的 S , $\dim S$ 有 4 种: ① 0 ② 1 ③ $n-1$ ④ n
(④: find a vec that dot to zero / always perpendicular to \vec{x})

- AB 's rows are combs of B 's rows. AB 's cols are combs of A 's cols.

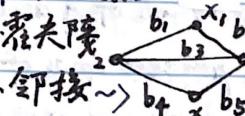
(so its row/col space contain in B/A 's) $\text{rank}(AB) \leq \text{rank}(A) \cdot \text{rank}(B)$. (这是矩阵法的第)

If multiply by an inv. mt, the rank will not change, because rank can't jump back

so if $AB=0$, we have $C_{(B)} \subset N(A) / C(A^T) \subset N(B^T)$, $r_A + r_B \leq n$. (mul an inverse mt, simply through acts.)

- Kirchhoff's Law . graphs and incidence matrix:

<参考 基于霍夫曼
关联矩阵、邻接~>



$$Ax = b, A = \begin{pmatrix} -1 & 1 & 1 \\ -1 & 1 & 1 \\ -1 & 1 & 1 \\ -1 & 1 & 1 \end{pmatrix}$$

since $\vec{b} = 0 \Rightarrow x_1 = x_2 = x_3 = x_4$
thus $\vec{x} = (c, c, c, c)$ in $N(A)$,
the rank must be 3.

since $b_1 - b_2 - b_3$ is a loop, rows 1, 2, 3 of A

are not lin. ind. but $(1, -1, 1, 0, 0)$ will be in A 's left nullspace.

(two ind. loops, $m-r=5-3=2$)

- Kirchhoff's Voltage Law $Ax = b$

Kirchhoff's Current Law $A^T y = 0$ —— 'balance equation': most important eq. in applied maths.

- $A = uv^T$'s four subspaces are ~~v^\perp~~ ~~u^\perp~~ . If B have those same subspaces, then $B = \lambda I$.

Generally, mt. A and B have same four subspaces means their rref equals.

(proof: same row space is the key. every row in rref(A) must be comb of rref(B), but notice I in these rref, it means each row should be equal)

- More Advanced Level: Matrix Spaces and Function Spaces

- dim of the subspace of sym. mts. is $\frac{1}{2}n^2 + \frac{1}{2}n$.

- second derivative equations: solution space $y'' = -y$ has two basis funcs: $\sin x$ and $\cos x$.

but $y'' = 2$ don't form a subspace.

a particular solution $y_p = x^2$. then the complete solution is $y = x^2 + cx + d$.

- Ex. $A = \begin{pmatrix} 1 & 0 & -1 \\ -1 & 1 & 0 \\ 0 & -1 & 1 \end{pmatrix}$, notice $A \begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 0 \\ 0 \end{pmatrix}$. The counting theorem writes $6+3=9$.

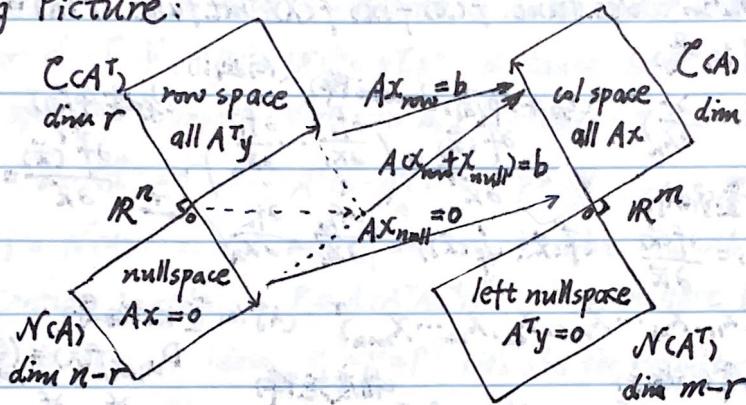
$AX=B, X \in 3 \times 3$ mt. space.

$\mathcal{C} \mathcal{N}$

- Fredholm's Alternative: Exactly one of these problems has a solution $\begin{cases} Ax=b \\ ATy=0 \text{ with } y \neq 0 \end{cases}$

3.5 dimensions of the four subspaces

The Big Picture:



1. A has the same row space as R .
same dim r and same basis.

2. The col space of A has dim r .
col rank equals row rank. (Ranking
 $\triangleright C(A) \neq C(R)$! but dim ✓. Theorem)

(for same combs of col = 0 for A and R ,
thus $Ax = 0 \Leftrightarrow Rx = 0$) ($R_{(A)} = R_{(R)}$)

3. A has same nullspace as R .
same dim $n-r$ and same basis.
(again, doesn't change solutions)

4. The left nullspace of A has dim $m-r$.
(Counting Theorem or $R^T y = 0 / y^T R = 0$,
 $R = EA$)

Fundamental Theorem of Linear Algebra:

Part 1. The col space and row space both have dim r ;
The nullspaces have dim $n-r$ and $m-r$.

rank-one matrix: $A = uv^T = \text{col times row}$. $C(A)$ has basis u , $C(AT)$ has basis v .

rank-one's sum: $EA = R$, $A = E^{-1}R = CR$ Ex. $A = \begin{pmatrix} 1 & 0 & 3 \\ 1 & 1 & 7 \\ 4 & 2 & 20 \end{pmatrix} = \begin{pmatrix} 1 & 0 & 3 \\ 1 & 1 & 7 \\ 4 & 2 & 1 \end{pmatrix} \begin{pmatrix} 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} = CR$.

(pre-read: full-rank decomp.) Every rank r mt. is a sum of r rank-one mts.
o left eigenvcs.

rank-one expansion)

$$A = (\vec{u}_1 \ \vec{u}_2 \ \vec{u}_3) \begin{pmatrix} \vec{v}_1^T \\ \vec{v}_2^T \\ \vdots \\ \vec{v}_r^T \\ \text{zero rows} \end{pmatrix} = \vec{u}_1 \vec{v}_1^T + \vec{u}_2 \vec{v}_2^T + \dots + \vec{u}_r \vec{v}_r^T = \text{rank } 1 + \text{rank } 1$$

4. Orthogonality

4.1 orthogonality of the four subspaces

orthogonal vectors $v^T w = 0$ and have

$$\|v\|^2 + \|w\|^2 = \|v+w\|^2$$

The orthogonal complement of a subspace V contains orthogonal spaces $v^T w = 0$ for every vector that is perpendicular to V , denoted by V^\perp .
all v in V and all w in W .
Orthogonality is impossible when

Part 2. $N(A)$ is the ortho complement of $C(AT)$; (in R^n)

$$\dim V + \dim W > \dim \text{(whole space)}$$

$N(AT)$ is the ortho complement of $C(A)$, (in R^m), since it only can be zero

$$(V+W)$$

to be contained in both ortho spaces.)

This means every x can be split into $x_r + x_n$. When $v^T v = 0 \Rightarrow v = 0$, a row of A can't be in $N(A)$.

A multiplies, nullspace component x_n goes to zero $Ax_n = 0$.

rowspace component x_r goes to colspace $Ax_r = Ax$.

Every b in colspace comes from one and only one vec x_r in rowspace.

There is an $r \times r$ inv. mt. hiding inside A , if throw away two nullspaces.

From the rowspace to the colspace, A is inv. ('pseudoinverse', see 7.4)

n vecs

o' the definition of basis has two properties, but one implies the other: ind. $\Leftrightarrow \text{span } R^n$

更多关于 tri. mat 的: L/U 的 inv 仍为 L/U; $L_1 L_2 / U_1 U_2$ 积仍为 L/U ; $LX = L'/UX = U'$ 则
(且对角线 $d = d_1 d_2$) X 一定由为 L/U .

口 matrix derivative

矩阵求导

· 定义概念 vec. func. $\vec{f}(x)$ $\vec{f}(\bar{x})$ $\vec{f}(X)$ mt. func. $F(X) = (F_{ij}(x_{mn}))$

(标函:)

$$\text{向元: } \frac{\partial f(\bar{x})}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial f}{\partial x_1} \\ \vdots \\ \frac{\partial f}{\partial x_n} \end{pmatrix} \xrightarrow{I} \frac{\partial f(\bar{x})}{\partial \bar{x}^T} = \left(\frac{\partial f}{\partial x_1} \cdots \frac{\partial f}{\partial x_n} \right)^T \quad (\text{分子布局, Jacobi 阵}) \quad (\text{向量})$$

$$\frac{\partial f(\bar{x})}{\partial \bar{x}^T} = \begin{pmatrix} \frac{\partial f_1}{\partial x_1} & \cdots & \frac{\partial f_1}{\partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial f_m}{\partial x_1} & \cdots & \frac{\partial f_m}{\partial x_n} \end{pmatrix} \quad (\text{分母布局})$$

$$\frac{\partial \bar{x}^T}{\partial \bar{x}} = \begin{pmatrix} \frac{\partial x_1}{\partial x_1} & \cdots & \frac{\partial x_1}{\partial x_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial x_n}{\partial x_1} & \cdots & \frac{\partial x_n}{\partial x_n} \end{pmatrix} \quad (\text{梯度/列向量})$$

$$\text{向元: } D_{\bar{x}} f(\bar{x}) = \frac{\partial f(\bar{x})}{\partial \bar{x}^T} \leftrightarrow \nabla_{\bar{x}} f(\bar{x}) = \frac{\partial f(\bar{x})}{\partial \bar{x}} \quad (\vec{f}(\bar{x}), J_f(\bar{x}))$$

矩阵: $\text{vec}(X)$: (列堆栈) 向量化 $= (x_{11}, x_{21} \cdots x_{m1}, x_{12}, x_{22} \cdots x_{mn})^T$ (行向量化)

$$(\text{Jacobi 阵}) \quad D_X f(X) = \begin{pmatrix} \frac{\partial f}{\partial x_{11}} & \cdots & \frac{\partial f}{\partial x_{1n}} \\ \vdots & \ddots & \vdots \\ \frac{\partial f}{\partial x_{m1}} & \cdots & \frac{\partial f}{\partial x_{mn}} \end{pmatrix} \quad (\text{列向量化})$$

$$\nabla_{\text{vec}} f(X) = \frac{\partial f(X)}{\partial \text{vec}} = \left(\frac{\partial f}{\partial x_{11}} \cdots \frac{\partial f}{\partial x_{mn}} \right)^T \quad (\text{梯度矩阵})$$

$$\nabla_X f(X) = \frac{\partial f}{\partial X} = \begin{pmatrix} \frac{\partial f}{\partial x_{11}} & \cdots & \frac{\partial f}{\partial x_{1n}} \\ \vdots & \ddots & \vdots \\ \frac{\partial f}{\partial x_{m1}} & \cdots & \frac{\partial f}{\partial x_{mn}} \end{pmatrix}$$

$$\text{同理矩阵也有 } D_X F(X) = \frac{\partial \text{vec}(F(X))}{\partial \text{vec} X}, \nabla_X F(X) = \frac{\partial \text{vec}^T(F(X))}{\partial \text{vec} X} \quad \Delta \text{ Transpose 的 T 位置可以变动.}$$

· 向元标函: $\frac{\partial c}{\partial \bar{x}} = 0$, $\frac{\partial f(\bar{x})g(\bar{x})}{\partial \bar{x}} = \frac{\partial f(\bar{x})}{\partial \bar{x}} g(\bar{x}) + f(\bar{x}) \frac{\partial g(\bar{x})}{\partial \bar{x}}$ 等其他老规律, 常用新规律:

$$1. \frac{\partial (a^T \bar{x})}{\partial \bar{x}} = \frac{\partial (a^T x)}{\partial x} = a \quad 2. \frac{\partial (x^T x)}{\partial \bar{x}} = 2x \quad 3. \frac{\partial (x^T a)}{\partial \bar{x}} = Ax + A^T x \quad 4. \frac{\partial (a^T x x^T b)}{\partial \bar{x}} = (ab^T + ba^T)x$$

$$\text{矩阵标函: } 5. \frac{\partial (a^T X b)}{\partial X} = ab^T, \frac{\partial (a^T X^T b)}{\partial X} = ba^T \quad 7. \frac{\partial (a^T X X^T b)}{\partial X} = (ab^T + ba^T)X, \frac{\partial (a^T X^T X b)}{\partial X} = X(ab^T + ba^T)$$

$$\cdot 维元标函: df(X) = \frac{\partial f}{\partial x_{11}} dx_{11} + \cdots + \frac{\partial f}{\partial x_{mn}} dx_{mn} = \text{tr} \left(\frac{\partial f(X)}{\partial X^T} dX \right) \quad \text{常用新规律:}$$

$$1. d(AXB) = AdXB \quad 2. d|X| = |X| \text{tr}(X^{-1}dX) = \text{tr}(X^*dX) \quad 3. dX^{-1} = -X^{-1}dXX^{-1} \quad (XX^{-1} = I \text{ 伟微分})$$

$$\Delta df(X) = d(\text{tr} f(X)) = \text{tr}(df(X)) \quad (\text{标量进不变})$$

$$\cdot 例 1: 1. \frac{\partial (a^T X X^T b)}{\partial X} : \quad d(a^T X X^T b) = \text{tr}(a^T d(X X^T) b) = \text{tr}(a^T d(X X^T b)) + \text{tr}(a^T X d(X^T b)) \quad (d(X^T) = (dX)^T)$$

$$\text{故 } \frac{\partial (a^T X X^T b)}{\partial X^T} = X(ba^T + ab^T), \quad \text{tr}(AB) = \text{tr}(BA) \Rightarrow \text{tr}(X^T b a^T dX) + \text{tr}(X^T a b^T dX)$$

$$\frac{\partial (a^T X X^T b)}{\partial X} = (ab^T + ba^T)X \quad 2. \frac{\partial \text{tr}(X^T X)}{\partial X} = 2X : \quad d(\text{tr}(X^T X)) = \text{tr}(dX^T X) + \text{tr}(X^T dX) = 2\text{tr}(X^T dX) \quad \text{故 } \frac{\partial \text{tr}(X^T X)}{\partial X^T} = 2X^T, \frac{\partial \text{tr}(X^T X)}{\partial X} = 2X$$

$$3. \frac{\partial \log |X|}{\partial X} = X^{-T} : \quad d(\ln |X|) = \text{tr} \left(\frac{1}{|X|} d|X| \right) = \text{tr}(X^{-1} dX) \quad 4. \frac{\partial \text{tr}(X+A)^{-1}}{\partial X} = -(X+A)^{-2} \quad \text{故 } \frac{\partial \text{tr}(X+A)^{-1}}{\partial X^T} = 2X^T, \frac{\partial \text{tr}(X^T+A)^{-1}}{\partial X} = 2X$$

$$5. \frac{\partial |X|^3}{\partial X} = 3|X|^3 X^{-T} \quad 6. 概率分布: X \sim N(\mu, \Sigma) \quad f(x) = (2\pi)^{-\frac{p}{2}} |\Sigma|^{-\frac{1}{2}} e^{-\frac{1}{2} (x-\mu)^T \Sigma^{-1} (x-\mu)}$$

$$\text{极大似然} \quad \ln L(\mu, \Sigma) = \sum_i \ln f(x_i) = -\frac{p}{2} n \ln 2\pi - \frac{1}{2} n \ln |\Sigma| - \frac{1}{2} \sum_i (x_i - \mu)^T \Sigma^{-1} (x_i - \mu)$$

$$\frac{\partial \ln L(\mu, \Sigma)}{\partial \mu} = \sum_i (x_i - \mu) = 0$$

$$\left\{ \begin{array}{l} \frac{\partial \ln L(\mu, \Sigma)}{\partial \Sigma} = 0 \\ \text{(后)} \end{array} \right.$$

4.2 projections

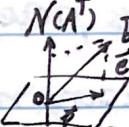
(in dim.)

The projection of \vec{b} onto a subspace S is the closest vector \vec{p} in S with $\vec{b} - \vec{p}$ is ortho to S .

$$\text{error: } \vec{e} = \vec{b} - \vec{p}, \|\vec{p}\|^2 + \|\vec{e}\|^2 = \|\vec{b}\|^2, \text{ distance: } \|\vec{e}\|$$

A 's cols span S , suppose they are ind. (otherwise use those r ind. cols)

$$\vec{a}_i \cdot (\vec{b} - A\vec{x}) = 0 \rightarrow A^T(\vec{b} - A\vec{x}) = 0, p = A\vec{x}. \quad A^T b = A^T A \vec{x}$$



$$N(A^T A) = N(A): A^T A \text{ is inv. if and only if } A \text{ has ind. cols. so } p = A(A^T A)^{-1} A^T b.$$

The projection matrix is $P = A(A^T A)^{-1} A^T$ Δ we have $P_c A P_R = A (= P_c A = A P_R)$. $P_{c/R}$: proj. to col/row space.

proj. mat is sym. sq. idem. $P^2 = P = P^T$ (it's also the definition of proj. mat.) $\text{rank}(P) = \text{rank}(A) = r$

$$(\text{for a vec. } \vec{z} = \frac{a^T b}{a^T a}, P = \frac{aa^T}{a^T a} \text{ a rank-one mat.}) \text{ (When } A \text{ is inv., } P \text{ is } I_r \text{)} \quad (P \sim (I_r)_0)$$

4.3 least squares approximations

$$A = \begin{pmatrix} 1 & t_1 \\ 1 & t_2 \\ \vdots & \vdots \\ 1 & t_m \end{pmatrix} \quad x = \begin{pmatrix} c \\ D \end{pmatrix} \quad b = \begin{pmatrix} b_1 \\ \vdots \\ b_m \end{pmatrix} \quad (b = C + Dt, m \geq 2, n=2)$$

There are no solutions to $Ax = b$, we splitting up $b = p + e$.

\hat{x} is the least square solution, for the square $P = A\hat{x}$

length of $Ax - b$ is minimized: $E = \|Ax - b\|^2 \text{ min. } (= \sum e_i^2)$

By geometry: the nearest point is the projection. solve $A^T A \hat{x} = A^T b$.

By algebra/calculus: the partial derivatives of E are zero. $\frac{\partial E}{\partial C} = 0, \frac{\partial E}{\partial D} = 0 \text{ or } \frac{\partial E}{\partial x} = 0 \text{ as a whole, we get } (Ax - b)^T A = 0.$

$$A^T A = \begin{pmatrix} m & \sum t_i \\ \sum t_i & \sum t_i^2 \end{pmatrix} \quad A^T b = \begin{pmatrix} \sum b_i \\ \sum t_i b_i \end{pmatrix}, \quad \hat{x} = \begin{pmatrix} c \\ D \end{pmatrix} = (A^T A)^{-1} A^T b$$

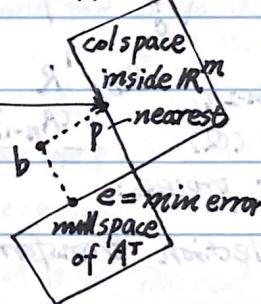
If $\sum t_i = 0$ then A has ortho cols and $A^T A$ is diag. It's worth shifting times t_i by subtracting the average $\bar{t} = \frac{\sum t_i}{m}$, $T = t - \bar{t}$ or make cols ortho. in advance by Gram-Schmidt. $b = C + Dt - \bar{t}$

Like fitting a straight line, the unknowns $C + Dt + Et^2$ still appear linearly in fitting a

(if A has dependent cols, we have many solution lines, Then $\hat{x} = (C, D, E)$, nothing diff. parabola.)

pseudoinv. of A will choose the shortest solution to $A\hat{x} = p$.

when A is ind. cols the pseudoinv. is usual left inv. $L = (A^T A)^{-1} A^T$)



Δ since first col of A is $(1, 1, \dots, 1)$, $\sum e_i = 0$ or it mean will go to b mean. ($\bar{b} = C + Dt$)

$(A \text{ is Vandermonde mt.})$
 $(\text{or fitting a plane, } C + Dx + Ey, \text{ and so on.})$

- $A = QR = QHR'$ \leftarrow diag R all 1's. with heights h_i in
 h_i tells the height col i above the plane of H)
 $\text{cols } 1 \text{ to } i-1$.
 $O_{(G-S)} \approx n^3$

ways of QR decomp.: Householder decomp. and Givens decomp.

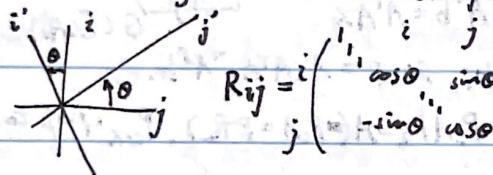
QR 分解之高斯-施密特法分解、吉文斯分解

QR's uniqueness: $Q_1 R_1 = Q_2 R_2$

Givens process: (A has ind. cols!)

$Q_1 = Q_2 R_2 R_1^{-1} = Q_2 U$ leads to
 $Q_1^T Q_1 = I = U^T U$ must be diag!

elementary rotation transform:



$$R_{ij} = \begin{pmatrix} 1 & 0 & i \\ 0 & \cos\theta & \sin\theta \\ 0 & -\sin\theta & \cos\theta \end{pmatrix} \quad (\text{left mult})$$

since we make R positive diag, so $U = I$, $R_1 = R_2 \dots$

choose $\cos\theta = \frac{x_i}{\sqrt{x_i^2 + x_j^2}}$ rotate x in $i-j$ plane
 notice that R_{ij} is orthogonal.

$$A_1 = R_{1n} R_{m-1} \cdots R_{12} A = \begin{pmatrix} \|a_1\| & & & \\ 0 & & & \\ 0 & & \ddots & \\ \vdots & & & 0 \end{pmatrix} \quad \text{then rotate } a_2. \quad A_2 = R_{2n} R_{2n-1} \cdots R_{23} A_1 = \begin{pmatrix} \|a_1\| & a_2 \\ 0 & \|a_2 - a_1\| \dots \\ \vdots & 0 \\ 0 & \dots \end{pmatrix}$$

(since A ind. cols, a_2 can't all be zeros.)

$$\text{so } A = (R_{n-1} R_{n-2} \cdots R_{12})^{-1} R \quad O(\text{Givens}) \approx \frac{4}{3} n^3 \quad \text{it's better for sparse matrix.}$$

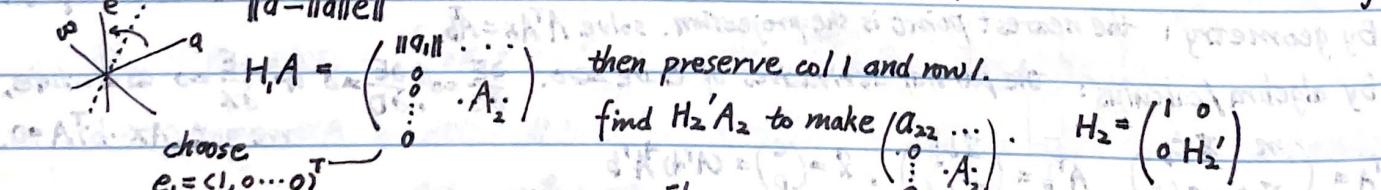
(Q) $\quad (A_{n-1})$ ortho. mts' product is also ortho.

Householder process:

elementary reflection transform: $H = I - 2ww^T$ w is normalized unitvec.
 notice that H is orthogonal and sym.

$$\text{or } H = Q^{-1} \begin{pmatrix} 1 & 0 \\ 0 & -1 \end{pmatrix} Q \quad (IH) = -I, |R_{ij}| = 1$$

choose $w = \frac{a - \|a\|e}{\|a - \|a\|e\|}$ to reflect a to e (Q^T) where $Qw = (1, 0)$ $(H = 2P - I \text{ ano. way})$



$$e = (1, 0 \dots 0)^T$$

$$H_A = \begin{pmatrix} \|a\| & & & \\ 0 & & & \\ 0 & & \ddots & \\ 0 & & & A_{22} \end{pmatrix} \quad \text{then preserve col 1 and row 1.}$$

find $H_2^T A_2$ to make $(A_{22} \dots)$.

$$H_2 = \begin{pmatrix} 1 & 0 \\ 0 & H_2^T \end{pmatrix}$$

$$\text{so } A = (H_{n-1} \cdots H_1)^{-1} R \quad O(\text{House}) \approx \frac{2}{3} n^3$$

(Q) $\quad (A_{n-1})$

To conclude: Gram-Smidt ✓ small mt.; generate q in every step

X not for sparse mt.; not stable; save all mt. thus cost large space.

Householder

✓ suitable for dense mt.; not calculate Q explicitly (if only need R)

X not generate q 's until finished; not for sparse mt.

Givens

✓ suitable for sparse mt.

X not generate q 's until finished; not for dense mt.

a rotation mt. = two reflection mts: $R_{ij} = H_1 H_2$

$$H_1 = I - 2w_1 w_1^T, w_1 = (0 \dots \sin\frac{\theta}{4}, 0 \dots 0, \cos\frac{\theta}{4}, 0 \dots 0)^T; H_2 = I - 2w_2 w_2^T, w_2 = (0 \dots 0, \sin\frac{3}{4}\theta, 0 \dots 0, \cos\frac{3}{4}\theta, 0 \dots 0)^T$$

4.4 orthonormal bases and Gram-Schmidt

not require sq.!

A matrix Q with orthonormal cols satisfies $Q^T Q = I$. It leaves lengths unchanged: $\|Qx\| = \|x\|$, and also preserves dot.p. $(Qx)^T (Qy) = x^T y$. Some Q 's examples: Rotation, Permutation, Reflection ($Q = I - 2uu^T$)

The least squares solution to $Qx = b$ is $\hat{x} = Q^T b$, the proj: $p = Q Q^T b = Pb$. thus \hat{u} is the minor (If Q is sq. /orthogonal/, it spans all space so $P = I$, $Q^T = Q^{-1}$) dot.p of b with \hat{u} normal with $\|\hat{u}\|^2 = 1$

Gram-Schmidt process:

q_1, q_2, \dots, q_n , then use Q comb/proj.

The one-dim proj's are uncoupled.

Subtract from every new vec its proj's in the directions already set. $b = \sum q_i (q_i^T b)$
 $\text{first } C = c - \frac{A^T c}{A^T A} A - \frac{B^T c}{B^T B} B / \|c\|$ At the end, divide the ortho. vecs by their length.
 $B = b - \frac{A^T b}{A^T A} A / \|B\|$

factorization: $A = QR$ (orthogonality decomp.) $(q_1, q_2, \dots, q_n) \begin{pmatrix} q_1^T a_1 & q_1^T a_2 & \dots & q_1^T a_n \\ q_2^T a_1 & q_2^T a_2 & \dots & q_2^T a_n \\ \vdots & \vdots & \ddots & \vdots \\ q_n^T a_1 & q_n^T a_2 & \dots & q_n^T a_n \end{pmatrix}$

$a_1, \dots, a_n \rightarrow q_1, \dots, q_n$ later q 's are ortho. to earlier a 's (and q 's.) $R = Q^T A$ to recognize the Any $m \times n$ ind. cols A can be factored. It's useful for least squares: (R 's diag is $\|A\|, \|B\|, \dots$) upper tri.

$A^T A = R^T R$. we solve $R \hat{x} = Q^T b$ by back substitution ($\hat{x} = R^{-1} Q^T b$). The real cost is G-S: $O(mn^2)$.

(modified G-S: more stable than original which subtracts all proj's at once. it subtracts one proj at a time as in

$$C^*, C: C^* = a_3 - (q_1^T a_3) q_1, C = C^* - (q_2^T C^*) q_2, q_3 = \frac{c}{\|c\|}.$$

5. Determinants

5.1 the properties of determinants

$|\det A| = \text{volume of box whose sides are } A$'s rows/cols. (sign is the direction! \downarrow)

① the det of $I_{n \times n}$ is 1. ② the det changes sign when two rows exchanged.

③ the det is a linear func of each row separately. $\begin{vmatrix} a+ta' & b+tb' \\ c & d \end{vmatrix} = \begin{vmatrix} a & b \\ c & d \end{vmatrix} + t \begin{vmatrix} a' & b' \\ c & d \end{vmatrix}$

R) then there are some important corollaries: ④ if two rows of A are equal, then $\det A = 0$.

⑤ subtracting a multiple of one row from ano. row leaves A unchanged. $\begin{vmatrix} a & b \\ c-ta & d-tb \end{vmatrix} = \begin{vmatrix} a & b \\ c & d \end{vmatrix}$

⑥ a mt. with a row of zeros has $\det = 0$. ⑦ if A is triangular then $\det A = a_{11} \cdots a_{nn}$ ← diag.

⑧ singular A 's $\det = 0$, inv. A 's $\det \neq 0$. ⑨ the det AB is $\det A$ times $\det B$. esp. $\det A^{-1} = \frac{1}{\det A}$.

⑩ the transpose A^T has the same det as A . (use $PA = LU$, $A^T P^T = U^T L^T$ and ⑨ proof) (proof: $\frac{|AB|}{|B|}$ has the same three properties that define $|A|$ (⑨⑩⑪), and the case $|B| = 0$ is easy for AB is singular.)

△ try to image those in geometry!

- o if A is singular, $AC^T = 0$, then each col of C^T is in the nullspace of A .
 ($A_{n \times n}$ of rank $n-2$ or smaller, all cofactors are zero and we only find $\vec{x} = 0$.)
- o $AC^T = (\det A) I \rightarrow (\det A)^{n-1} = \det C$. so if you know all cofactors of A (with position) you could get A .
- Cauchy-Binet formula : 总积-比内公式

It gives the det of a sq. mt. AB (esp. $A^T A$) when the factors A, B are rectangular.

$A_{s \times n}, B_{n \times s}$: $s > n$ leads to $\det AB = 0$;

$s \leq n$ then $\det AB = \text{sum of } (s \times s \text{ dets in } A)(s \times s \text{ dets in } B)$.

$$\left(\frac{\underline{a}}{\underline{d}}, \frac{\underline{b}}{\underline{e}}, \frac{\underline{c}}{\underline{f}} \right) \left(\begin{matrix} \underline{g} & \underline{j} \\ \underline{h} & \underline{k} \end{matrix} \right) \cdot \left(\frac{\underline{a}}{\underline{d}}, \frac{\underline{b}}{\underline{e}}, \frac{\underline{c}}{\underline{f}} \right) \left(\begin{matrix} \underline{g} & \underline{j} \\ \underline{i} & \underline{l} \end{matrix} \right) \cdot \left(\frac{\underline{a}}{\underline{d}}, \frac{\underline{b}}{\underline{e}}, \frac{\underline{c}}{\underline{f}} \right) \left(\begin{matrix} \underline{g} & \underline{j} \\ \underline{h} & \underline{k} \end{matrix} \right) \quad \begin{array}{l} (s \text{ 's cols of } A \text{ correspond with } s \text{ 's rows of } B) \\ (\text{calculate the } 2 \times 2 \text{ dets underlined}) \end{array}$$

o Jacobi formula: The cofactor formula can be generalized. (or 'Laplace expansion')

For $n \times n$ mt., we can choose $k \times k$ det times $(n-k) \times (n-k)$ det. Then add them with signs.

The sign depends on the permutation of $a_1, \dots, a_k, b_1, \dots, b_{n-k}$ (a and b is the column you choose)

$$\begin{vmatrix} 2 & -1 & 0 & 0 & 0 \\ 1 & 2 & -1 & 0 & 0 \\ 0 & 1 & 2 & -1 & 0 \\ 0 & 0 & 1 & 2 & 1 \\ 0 & 0 & 0 & 0 & 2 \end{vmatrix} = \begin{vmatrix} 2 & -1 \\ -1 & 2 \end{vmatrix} \cdot \begin{vmatrix} 2 & -1 & 0 \\ -1 & 2 & -1 \\ 0 & -1 & 2 \end{vmatrix} + \begin{vmatrix} 2 & 0 \\ -1 & -1 \end{vmatrix} \cdot \begin{vmatrix} -1 & 1 & 0 \\ 0 & 2 & 1 \\ 0 & -1 & 2 \end{vmatrix} + P(1, 2, 3, 4, 5) = 6$$

o [rank-one mt.] $A = uv^T$ has an eigenvector u . its $\lambda = v^T u$. Other eigenvalues are zeros.

更多关于
permuation mt. P have $P^k = I \Rightarrow \lambda^k = 1$ Then we find P 's eigens. (maybe complex)

$\sum_i \lambda_i^k = \text{tr } A^k$. (every perm. mt. leaves $\vec{x} = (1, 1, \dots, 1)^T$ unchanged so $\lambda = 1$)

rotation mt. can do the same. $\lambda^{\frac{2\pi}{\theta}} = 1$ we find $e^{\pm i\theta}$ (for higher dim, $\lambda = 1, 1, \dots, e^{\pm i\theta}, \dots$)

o When a mt. is diagl, the number of nonzero eigenvalues is its rank. $\sim \begin{pmatrix} 1_r \\ 0 \end{pmatrix}$

o $A+B, AB$ might not have $\lambda = \lambda_a + \lambda_b, \lambda_a \lambda_b$! It only happens when A, B have same eigenvcs.

o esp. $B = I$, $A+cI$ can have $\lambda = \lambda_a + c$ for A 's eigenvcs must in I 's eigen.

(when λ has multiplicity, the whole plane spanned by its eigenvess are eigen comb) for I it's what space)

o no interconnection between invertibility ... diagonalizability!

$$\lambda = 0?$$

$$X \exists ?$$

5.2 permutations and cofactors

The determinant of any $n \times n$ matrix can be found in three ways: (more: Big Partition...)

1. pivot formula: $PA = LU \rightarrow \det A = \pm \det U = \pm$ (product of pivots)

(we don't need row exchanges when all the upper left submts. have $\det A_k \neq 0$. The k th pivot is

2. big formula: $\det A = \text{sum over all } n! \text{ col permutations } P = (\alpha, \beta, \dots, w) \quad d_k = \frac{\det A_k}{\det A_{k,k}} \quad = \sum \det P \cdot a_{1\alpha} a_{2\beta} \cdots a_{n w}$ choose one entry from every row and col.

$$|a_{1\alpha} a_{2\beta} \cdots a_{n w}| = a_{11} a_{22} a_{33} | \cdot \cdot \cdot | + a_{12} a_{23} a_{31} | \cdot \cdot \cdot | + a_{13} a_{21} a_{32} | \cdot \cdot \cdot | + a_{11} a_{23} a_{32} | \cdot \cdot \cdot | + a_{12} a_{21} a_{33} | \cdot \cdot \cdot | + a_{13} a_{22} a_{31} | \cdot \cdot \cdot |$$

$n!$ ways to order, half odd and half even.

3. cofactor formula: $\det A = a_{11} C_{11} + \cdots + a_{in} C_{in}$. each cofactor C_{ij} includes its correct sign:

since a_{ij} removes row i and col j , we can linearly expand \det by r_i/c_j . $C_{ij} = (-1)^{i+j} M_{ij}$

$$\begin{aligned} |a_{11} \cdots a_{33}| &= |a_{11} a_{22} a_{33} \cdots| + |a_{12} a_{23} a_{31} \cdots| + |a_{13} a_{21} a_{32} \cdots| \quad \text{to permute } a_{ij} \text{ to } a_{11} \text{ need } (-1)^{i+j-2} \text{ reverse,} \\ &\quad \text{then } M_{ij} \text{ 's det permutes as itself.} \\ &= a_{11}(a_{22}a_{33} - a_{32}a_{23}) + a_{12}(a_{23}a_{31} - a_{31}a_{23}) + a_{13}(a_{21}a_{32} - a_{32}a_{21}) \quad (\det A^T = \det A, \text{ so did col } j.) \end{aligned}$$

5.3 Cramer's rule, inverse, and volumes

Cramer's rule solves $Ax = b$. we construct $A \begin{pmatrix} x_1 & 0 & 0 \\ x_2 & 1 & 0 \\ x_3 & 0 & 1 \end{pmatrix} = \begin{pmatrix} b & a_2 & a_3 \end{pmatrix} = B_1$. $x_1 = \frac{\det B_1}{\det A}$.

then $x_n = \frac{\det B_n}{\det A}$, where mt. B_n has the n th col of replaced by b .

Then we find the cols of A^{-1} by solving $AA^{-1} = I$. $B_n = \begin{pmatrix} \dots & 0 & \dots \end{pmatrix}$. so

$$\begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} x_1 \\ x_2 \\ x_3 \end{pmatrix} = \begin{pmatrix} 1 \\ 0 \\ 0 \end{pmatrix} \quad \begin{pmatrix} a & b \\ c & d \end{pmatrix} \begin{pmatrix} y_1 \\ y_2 \\ y_3 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \\ 0 \end{pmatrix} \quad (A^{-1})_{ij} = \frac{C_{ji}}{\det A} \quad \text{and } A^{-1} = \frac{C^T}{\det A}.$$

A direct proof of $AC^T = (\det A)I$:

(C_{ij} go into the cofactor matrix C)

The cofactor formula yields $\det A$ on the diagonal, multiplying cofactors from diff rows

is the cofactor formula for a new matrix when the second row is copied into first row,

so yields zero.

The triangle with corners $(x_1, y_1), (x_2, y_2), (x_3, y_3)$ has area $= \frac{\det}{2} = \frac{1}{2} \begin{vmatrix} x_1 & y_1 & 1 \\ x_2 & y_2 & 1 \\ x_3 & y_3 & 1 \end{vmatrix}$,

We can proof that the parallelogram's area when $(x_3, y_3) = (0, 0)$, area $= \frac{1}{2} \begin{vmatrix} x_1 & y_1 \\ x_2 & y_2 \end{vmatrix}$.

or the box's volume obeys the determinant properties ①②③.

cross p.: $\vec{u} \times \vec{v} = \begin{vmatrix} i & j & k \\ u_1 & u_2 & u_3 \\ v_1 & v_2 & v_3 \end{vmatrix}$, is perpendicular to \vec{u} and \vec{v} .

triple p.: $(\vec{u} \times \vec{v}) \cdot \vec{w} = \begin{vmatrix} u_1 & u_2 & u_3 \\ v_1 & v_2 & v_3 \\ w_1 & w_2 & w_3 \end{vmatrix}$ = determinant = volume.

(antisym.)

- Heisenberg's uncertainty principle: (inf. mt.) position mt. P , momentum mt. Q . (sym.)

$$QP - PQ = I \quad (\text{To have } Px=0 \text{ and } Qx=0 \text{ would require } x=0)$$

$$x^T x = x^T Q P x - x^T P Q x \leq 2 \|Px\| \|Qx\| \quad \text{Explain last step by using Schwarz ineq.}$$

$$\text{So } \frac{\|Px\| \|Qx\|}{\|x\|^2} \geq \frac{1}{2}. \quad (\text{Impossible to get position and momentum error all very small.})$$

$$\|uv\| \leq \|u\| \|v\|.$$

(proof Rjá)

(commutable)

- A and B share the same n ind. eigenvects if and only if $AB = BA$.

- If A is $m \times n$ and B is $n \times m$, then AB and BA have some nonzero eigenvalues.

(proof: $\begin{pmatrix} I & -A \\ B & 0 \end{pmatrix} \begin{pmatrix} AB & 0 \\ B & I \end{pmatrix} \begin{pmatrix} I & A \\ 0 & I \end{pmatrix} = \begin{pmatrix} 0 & 0 \\ B & BA \end{pmatrix} \begin{pmatrix} m \times m & m \times n \\ n \times m & n \times n \end{pmatrix}$)

(and $|n-m|$ zeros.)

$$\text{so } E = \begin{pmatrix} AB & 0 \\ B & 0 \end{pmatrix} \sim F = \begin{pmatrix} 0 & 0 \\ B & BA \end{pmatrix}, \text{ they have same } m+n \text{ values: } F \text{ has } n \text{ eigenvalues of } BA + m \text{ zero}$$

- Suppose A_1, A_2 $n \times n$ inv., then $A_1 A_2$ is similar to $A_2 A_1$, same eigenvalues. (B choose A_2)

- Gershgorin circles: Every λ is in the circle around one or more diagonal entries a_{ii} :

(reason: $A - \lambda I$ is singular, it cannot be diag dominant.)

$$|\lambda_{ii} - \lambda| \leq R_i = \sum_{j \neq i} |a_{ij}|$$

- eigenvalues of A equals eigenvalues of A^T . (since $\det(A - \lambda I) = \det(A^T - \lambda I)$)

The eigenvectors of A and A^T from different λ 's are perpendicular. (esp. sym. mt's since $A\vec{x}_1 \cdot \vec{x}_2 = \vec{x}_1 \cdot A^T \vec{x}_2 \Rightarrow \lambda_1 \vec{x}_1 \cdot \vec{x}_2 = \lambda_2 \vec{x}_1 \cdot \vec{x}_2$, $\lambda_1 \neq \lambda_2$ leads to $\vec{x}_1 \perp \vec{x}_2$.)

eigenvecs are ortho.)

(or $A = X \Lambda X^{-1}$, $A^T = X^{-T} \Lambda X^T$, the eigenvecs of A^T are X^{-T} , $X^{-T} X = I$ leads to $x_i \cdot x_j$ ($i \neq j$) = 0.)

- 矩阵收集包(3): 左 eigenvectors of A ($y^T A = \lambda y^T$). rank-one's sum: $A = \lambda_1 x_1 y_1^T + \dots + \lambda_n x_n y_n^T$.

• Hadamard matrix $H_2 = \begin{pmatrix} 1 & 1 \\ 1 & -1 \end{pmatrix}$ $H_{ij} = \pm 1$ and orthogonal. (so n cannot be odd) $H = H^{-1}$

(Sylvester theorem) $H v \rightarrow \begin{pmatrix} H & H \\ H & -H \end{pmatrix} v$ Hadamard Theorem:

(Hadamard conjecture: $n \mid \det H_n v$)

$A = (a_{ij})_{n \times n}$ with $-1 \leq a_{ij} \leq 1$,

then $|\det A| \leq n^{\frac{n}{2}}$. equality achieves when A is H ($a_{ij} = \pm 1$ and rows are ortho.)

$$|\det H| = |\det H_{n-1}| + |\det H_{n-2}| \quad (\text{Fibonacci } F_{n+2})$$

• Markov matrix largest eigenvalue $\lambda_m = 1$. this eigenvect is the steady state. (A^∞)

Mar. mt. each column adds up to 1. ('possibility transfer')

(proof: since $\lambda_A = \lambda_{A^T}$, thus λ can be 1;

and Gershgorin circle tells $\lambda - a_{ii} \leq \sum_{j \neq i} a_{ji}, \lambda \leq 1$)

Chapter 2: Eigenvectors and Eigenvalues
 (Eigenvalues and Eigenvectors)

6. Eigenvectors and Eigenvectors

6.1 Introduction to eigenvectors

An eigen vector \vec{x} lies along the same line as $Ax : Ax = \lambda x$. The eigenvalue is λ .

Then $(A - \lambda I)x = 0$ and $A - \lambda I$ is singular and $\det(A - \lambda I) = 0$. So 1. Compute the det of $A - \lambda I$.

2. Find the n roots of this polynomial of degree n . 3. For each λ , solve $(A - \lambda I)x = 0$ to find an

The eigenvalues of A^2 and A^{-1} are λ^2 and λ^{-1} , with the same eigenvectors.

(Vieta's formula or geometry) The sum of the λ 's equals the sum down the main diag of A or called the trace. $\sum_{i=1}^n \lambda_i = a_{11} + a_{22} + \dots + a_{nn} = \text{tr}(A)$. The products of the λ 's equals the det of A .

Special mts. (properties) lead to special eigenvalues and eigenvectors: $\prod_{i=1}^n \lambda_i = \det(A)$.

Proj. P have $\lambda = 1$ and 0. Reflection R have $\lambda = 1$ and -1. Rotation have $\lambda = e^{i\theta}$ and $e^{-i\theta}$.

Singular mt. has $\lambda = 0$. Triangular mt. has λ 's on its diagonal. Identity mt. has all $n \lambda = 1$.

Sym. mt. can be compared to a real number, antisym. mt. can be compared to an imaginary

Orthogonal mt. corresponds to a complex number with $|\lambda| = 1$. For these three special mts., their eigenvecs are perpendicular/orthonormal.

6.2 Diagonalizing a matrix

$Ax_i = \lambda x_i$ are the cols of $AX = X\Lambda$. The eigenvalue matrix Λ is diagonal and correspond with the order of x_i in X .

n ind. eigenvecs diagonalize $A : A = X\Lambda X^{-1}$. $\Lambda = x^T A x$ (also diagonalize all powers $A^k = X\Lambda^k X^{-1}$)

No equal eigenvalues : ind. x from diff. λ , so eigenvecs are ind. X is inv. and A is

Equal eigenvalues : A might have too few ind. eigenvecs. X^{-1} fails.

we have $GM \leq AM$ (geometric multiplicity \leq algebraic multiplicity)

when $GM < AM$, A is not diagl. $\dim(\mathcal{N}(A - \lambda I))$ roots of $\det(A - \lambda I) = 0$ multiply A several times

Similar matrices: (only I itself) (because maybe two cases during $A - \lambda I$) - multiply λ_j to

all $A = BCB^{-1}$, $A \sim C$, similar I). $\begin{vmatrix} 0 & 0 & 0 \\ 0 & 0 & 0 \\ 0 & 0 & 0 \end{vmatrix}$ pivots here. $\therefore \begin{vmatrix} 1 & (1-\lambda_1)^2 & \dots \\ 0 & 1 & \dots \\ 0 & 0 & 1 \end{vmatrix} \Rightarrow \prod_j (\lambda_1 - \lambda_j)$

they all share the eigenvalues of C . (black is zero when λ_i) compressed so $c_1 = 0$. $c_1 x_i = 0$ in one pivot. all other $c_j = 0$, ind.

Classical examples:

① Markov $\begin{pmatrix} .8 & .3 \\ .2 & .7 \end{pmatrix} \xrightarrow{\sim} \begin{pmatrix} .6 & .4 \\ .4 & .4 \end{pmatrix} \oplus I, \begin{pmatrix} 1 & 1 \\ 0 & 1 \end{pmatrix}$ ③ $A^2 = A \cdot \frac{xy^T}{y^T x}$. ④ Fibnums. $\vec{U} = A \vec{U}_K$, $(U_K = A^K U_0)$

(Jordan form) ($\lambda = 1, 0$) ($\lambda = 1$)

$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$ (difference eq.)

(ano2. way): $Sx = \lambda x$, $Sy = \beta y$. its colspace $C \perp$ nullspace N ,
 $S - \beta I$ is sym. so $x \perp y$ ($\lambda_1 \neq \lambda_2$);

o' we can call R is S (sym.)'s sq.root
if $S = R^T R$. so from diaglize we

□ 对称矩阵的证明

a proof of sym. mts has ortho. eigenvecs in algebraic way:

$$S = X \Lambda X^{-1}, S^T = X^{-T} \Lambda X^T \Rightarrow X^T X \Lambda = \Lambda X^T X \text{ or } \Lambda X^T X \text{ is sym.}$$

find R can be $= \frac{1}{\sqrt{\lambda}} Q$
(not unique)
 R to

we know if A, B sym. (AB) sym. needs $AB = BA$. so Λ and $X^T X$ commutable and have same eigenv.

since Λ is diag., $X^T X$ must be like I or Λ , it's Q . (ano. way: $\lambda_i x^T y = (Sx)^T y = x^T S y = \lambda_j x^T y$)

② above has a premise that sym. mts can be diagonalized: $(x: \lambda_1, y: \lambda_2)$ so $x^T y = 0$ ($\lambda_1 \neq \lambda_2$),

[Schur's Theorem]: Every sq. A factors into $Q T Q^{-1}$ where T is upper tri. and $Q^T = Q^{-1}$.

If A has real evalues then Q and T can be chosen real: $Q^T Q = I$.

(informal: if S has When $S^T = S$, T must be diagonal: $S = Q \Lambda Q^T$.

repeated λ 's, slightly change S in diag($c, 2c, \dots, nc$). with $c \rightarrow 0$, orthonormal evecs remained.)

③ real eigenvalues: $Sx = \lambda x \rightarrow S\bar{x} = \bar{\lambda}\bar{x}, \bar{x}^T S = \bar{x}^T \bar{\lambda}$.

$$\downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow \quad \quad \quad \downarrow$$

$$x^T S x = x^T \lambda x \quad \quad \quad \bar{x}^T S x = \bar{x}^T \bar{\lambda} x \quad \quad \quad |A_{\text{anti}}| = \text{positive or zero.}$$

(all mts have $TT^T = T\bar{\lambda}$)

④ same signs with pivots: $\bar{x}^T x$ is the squared length > 0 . (for antisym., $\lambda = -\bar{\lambda}$, ↑ it's pure imag.)

$$\text{tri. fac. } S = LDL^T = \begin{pmatrix} 1 & & \\ 0 & \ddots & \\ & & 1 \end{pmatrix} \begin{pmatrix} d & & \\ & \ddots & \\ & & d \end{pmatrix} \begin{pmatrix} 1 & & \\ 0 & \ddots & \\ & & 1 \end{pmatrix} \quad \text{so } \lambda = \bar{\lambda}, \text{ it's real}$$

when L moves to I ($0 \rightarrow 0$), IDI^T has eigenvalues = pivots d 's. (evalue → pivot)

But to change sign, a real evalue would cross zero → matrix singular at a moment X

• 算零阵 nilpotent matrix: the following propositions are equivalent: no sign change in moving

① A is nilpotent ($\exists m$, s.t. $A^m = 0$) ② $A^n = 0 \leftarrow A^{r(A)+1} = 0$ ③ all λ 's of $A = 0$

④ (from ③) $\det(A + kI) = k^n$, $\det(kA + I) = 1$ ⑤ $\forall k \text{ tr}(A^k) = 0$ (o [Cayley-Hamilton's Theorem]:

$\forall k$

($\lambda = -1 \nmid A^2 = I$: nothing new. $\lambda = \pm i$, esp. sym.+ortho.)

eigenpolynomial: $\det(\lambda I - A) = f(\lambda)$
for nilp. A it's λ^n ,
then we have $f(A) = 0$.)

• 更多关于矩阵指數: e^{At_1} times e^{At_2} must be $e^{A(t_1+t_2)}$, but $e^A e^B e^C e^D$ can be all diff!

• when A is tri., X, X^{-1}, e^{At} are also tri. (since $A - \lambda I$: $\boxed{\begin{matrix} 0 & & \\ & \ddots & \\ & & 0 \end{matrix}}$, X is tri.) ($e^A \cdot e^A = e^{2A}$)

• repeated λ : diaglize is impossible so compute e^{At} directly:

$$A = \begin{pmatrix} 0 & 1 \\ -1 & 2 \end{pmatrix}, \lambda = 1, 1, \quad e^{At} = e^{It} e^{(A-I)t} = e^t (I + (A-I)t) \quad \text{some mt.'s series stops early.}$$

$(A-I)^2$ is zero. (o)

• antisym. $A = \begin{pmatrix} 0 & 1 \\ -1 & 0 \end{pmatrix}$, then $e^{At} = \begin{pmatrix} \cos t & \sin t \\ -\sin t & \cos t \end{pmatrix}$, its eigenvalues are e^{it}, e^{-it} . so we find t here.

$(A^2 = I)$

• extension: Cosine form: solve $\frac{d^2 u}{dt^2} = -A^2 u$ $\cos At = I - \frac{1}{2} A^2 + \frac{1}{4!} A^4 + \dots$

Short form $u_{\text{cos}} = \cos(At) u_{\text{cos}}$ / specific form $u_{\text{cos}} = c_1 x_1 + c_2 x_2$ which corresponds.

$$u_{\text{cos}} = c_1 \cos \lambda_1 t x_1 + c_2 \cos \lambda_2 t x_2$$

6.3 systems of differential equations

- $\frac{du}{dt} = Au$. if $Ax = \lambda x$ then $u_{ct} = e^{\lambda t} x$ will solve it. each λ and x gives a ind. solution.
- start from u_{co} , it can be written as comb $c_1 x_1 + \dots + c_n x_n$ or $Xc = u_{co}$, $c = X^{-1} u_{co}$.
- then each eigenvect is muled by $e^{\lambda_i t}$, the solution is $u_{ct} = c_1 e^{\lambda_1 t} x_1 + \dots + c_n e^{\lambda_n t} x_n$.
- (at past $\frac{d}{dt} \begin{pmatrix} y \\ z \end{pmatrix} = \begin{pmatrix} 0 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} y \\ z \end{pmatrix}$ you may use $\frac{d}{dt} (y \pm z) = \pm (y \pm z)$ instead of $u_{ct} = e^{At} \begin{pmatrix} 1 \\ 1 \end{pmatrix}$, $u_{2t} = e^{-At} \begin{pmatrix} 1 \\ -1 \end{pmatrix}$)
- Matrix exponential form: $e^{At} = I + At + \frac{1}{2!}(At)^2 + \dots \quad .(e^{At})' = Ae^{At}$.
2. e^{At} 's eigenvalues are $e^{\lambda t}$. (sub $A = X\Lambda X^{-1}$ into exponential, 3. e^{At} has inverse e^{-At} . When A is antisym, ($A^T = -A$) e^{At} is orthogonal.)
 $u_{ct} = \sum_i c_i e^{\lambda_i t} x_i = X e^{At} X^{-1} u_{co} = e^{At} X e^{At} X^{-1} u_{co}$.
- (we call A is stable when $e^{At} \rightarrow 0$ thus all eigenvalues of A have real part < 0 . imaginary part gives oscillation.)

(Not included: if two λ 's are equal, with only one eigenvect, ano. solution is needed: $t e^{\lambda t} x$)

6.4 symmetric matrices

- Every sym. matrix S can be diagonalized, and has n real eigenvalues λ_i and $\det > 0$.
- (Principle Axis / Spectral Theorem) $S = Q \Lambda Q^T$ n orthonormal eigenvectors q_1, \dots, q_n .
- The number of positive eivals of S = The number of positive pivots of S (proof $\lambda_i = q_i^T S q_i$)
- (Antisym. matrix $A^T = -A$ have imag. λ 's and ortho. complex 'sign's matched')

For real matrices (nonsym.), complex λ 's and x 's come in conjugate pairs: $Ax = \lambda x$

(every ortho. Q have $|Q| = 1$ since then)

$$A\bar{x} = \bar{\lambda}\bar{x}$$

6.5 positive definite matrices

- positive definite (sym.) S : all eivals $> 0 \leftrightarrow$ all pivots $> 0 \leftrightarrow$ all upper left dets > 0 .
4. (Energy-based def/test) $x^T S x > 0$ for all vecs $x \neq 0$. (so obviously S pos. Tpos. $\rightarrow S + T$ pos.)
5. A col ind. (to rule out zero), $S = A^T A$, ($A = \text{chol}(S) = L \sqrt{D} / A = Q \sqrt{\Lambda} Q^T$ etc.) ($x^T S x = \|Ax\|^2 > 0$)
- positive semidefinite: allows $\lambda = 0$, pivot = 0, det = 0. $x^T S x = 0$ ($x \in N(S)$). A has dep. cols. $= N(A)$
- geometry: The graph of $x^T S x = 1$ is an ellipse: $(x, y) Q \Lambda Q^T \begin{pmatrix} x \\ y \end{pmatrix} = (X, Y) \Lambda \begin{pmatrix} X \\ Y \end{pmatrix} = \lambda_1 X^2 + \lambda_2 Y^2 = 1$.
- y. The axes point along eigenvects of S , half-lengths are $\frac{1}{\sqrt{\lambda}}$ ($X = q_1^T x, Y = q_2^T x$) (standard form)
- (if one λ is neg. ellipse changes to hyperbola.)

o' \$S \text{ pos. } T \text{ pos.} \rightarrow ST \text{ might not sym.}\$
 but pos. evalues: \$STx = \lambda x\$.

$$\lambda = \frac{(Tx)^T STx}{\lambda x^T Tx} > 0$$

(Normal includes sym. antisym. ortho.)

- \$N\$ has \$n\$ orthonormal eigenvectors (\$N = Q\Lambda Q^H\$) if and only if \$N\$ is normal. (\$N^H N = N N^H\$)
 (proof: ① \$N = Q\Lambda Q^H\$ sub. ② Schur T: \$N = Q T Q^H\$ and \$T\$ is \$\Lambda\$.)

o' sym. \$S\$. ① no diag entry can be larger than \$\lambda_{\max}\$ (use \$S = \lambda_i q_i q_i^T\$)

② diag entries fall in between the \$\lambda\$'s. (\$S_{2x2} = \frac{\det(A-\lambda I)}{\lambda^2 - \lambda_1 \lambda_2}\$)

□ Congruence . similarity — same transform on diff. basis.

合同矩阵 congruence — same bilinear / (0,2) tensor / metric on diff. basis.

\$x^T Ax\$ \$T = T_{ij} e_i e_j^T\$, if \$e_i = R_{ij} e'_j\$ then \$T'_{ij} = T_{kl} R_{ki} R_{jl}\$, mt. form: \$T' = R^T T R\$.

\$x^T \underbrace{C^T A C}_B x\$ some quadratic surface on diff. basis: \$\oplus \times \oplus \oplus\$

In real, congruence \$\leftrightarrow\$ equal inertia index. keep the pos./neg. inertia index.

□ Jordan canonical form
 In complex, congruence \$\leftrightarrow\$ equal rank.

约当标准形 (观后) . sq. \$A \sim J = \begin{pmatrix} J_1 & & \\ & \ddots & \\ & & J_m \end{pmatrix}\$, \$J_i = \begin{pmatrix} \lambda_i & & \\ & \ddots & \\ & & \lambda_i \end{pmatrix}\$ (可考若 \$J_i\$: Jordan pat.)

< 可考若不变因子和各阶 \$gm = \text{number of Jordan blocks} \leq cm = \text{scale/cols of Jordan blocks}\$.

余子式的方法替代 G.M. AM>

. \$A = X J X^{-1}\$. \$AX = XJ\$ solve \$x_i\$ one by one: \$(\begin{matrix} x_1 & x_2 & x_3 & \dots \\ \lambda & \lambda & \lambda & \dots \\ & & & \ddots \end{matrix}) (A - \lambda I)x_1 = 0\$ (call \$\lambda_i\$ adds)
 (each Jord. block)

. a proof of nilpot. \$A\$ has \$A^{r(A+1)} = 0\$:

since \$\lambda=0\$, \$A \sim J = \begin{pmatrix} J_{(0,k_1)} & & \\ & \ddots & \\ & & J_{(0,k_g)} \end{pmatrix}\$

\$r(A) = r(J) \geq r(J_{(0,k_i)}) = k_i - 1\$ so \$k_i \leq r(A) + 1\$,

$J_{(0,k_i)}^{k_i} = 0 \rightarrow A^{r(A)+1} = 0$

' a more classical way of Jordan: \$(A - \lambda I)x_j = x_{j-1}

① Primary Decomp. \$f(\lambda) = (\lambda - \lambda_1)^{m_1} \cdots (\lambda - \lambda_m)^{m_m}\$

惟素分解 \$V_i = \ker(A - \lambda_i I)^{m_i}\$ then \$V = \bigoplus V_i\$ (根子

② Cyclic Decomp. \$V_i = \bigoplus_{j=1}^l \text{span}\{N^j w_j : 0 \leq j \leq l\}\$.

弱环解 (t = dim(ker \$N\$), \$N = A - \lambda I\$) (强循环子空间)

< 可考若 Young 图 > (t = dim(ker \$N\$), \$N = A - \lambda I\$) (强循环子空间)

◦ rank-one decompr. / Spectral decompr.: a synthesis concept of decompr. \$A = \sum_i \lambda_i G_i = \sum_i \lambda_i u_i v_i^T\$

秩一分解

谱分解 (\$G_i G_j = \delta_{ij} G_i\$ if \$S\$)

\$G_i\$ is idempot. \$\lambda_i\$ is the eigenvalue. (\$G_i = q_i q_i^T\$ or \$u_i v_i^T\$)

I generalize it: all \$n\$ rank-one mts \$u_i v_i^T\$ adds up to \$A\$. Obviously it's not unique.

◦ a rank \$r\$ mt. can be expanded by \$r\$ rank-one mts and at least \$r\$. (\$r = r-k+k\$, gen.)

proof: \$A\$ can be elementary transformed to \$P \begin{pmatrix} I_r & 0 \\ 0 & 0 \end{pmatrix} P^{-1}\$ so \$A = \sum_{i=1}^r p_i \cdot I \cdot q_i^T\$.

and since \$r(A+B) \leq r(A) + r(B)\$, at least should \$r\$ rank-one mts.

□ Table of Eigens [Matrix/\$\lambda/x\$] P362

特殊特征值向量表

◦ \$Q\$ has all singular values = 1 as well as \$I\$.

◦ a proof of commutable mts have common eigenvecs: esp. normal mt \$A\$ have same \$A\$ has a \$\lambda\$ and its eigensubspace \$V_\lambda\$. \$\nu \in V_\lambda\$, so \$V_\lambda\$ also a invariant subspace of \$B\$. There must has diff. order!

\$AB\nu = BA\nu = \lambda B\nu\$

be a \$\mu\$ that s.t. \$B\nu = \mu \nu\$ with \$A\nu = \lambda \nu\$. ('a paradox')

Then \$\nu\$ is the common eivec.

o' 补充说明: when $q_1 \cdots q_n$ ortho., we have $q_1 q_1^T + \cdots + q_n q_n^T = I$, vice versa. (?)

o $Q_1 A Q_2^T$ has the same σ 's as A and $Q S Q^T$ has the same λ 's as S .

Computing of singular: $Q^T S Q \rightarrow$ triang. $Q_1 A Q_2^T \rightarrow$ bidiag.

o $\sigma_{\max} \geq |\lambda|_{\max}$; $\sigma_{\min} \leq |\lambda|_{\min}$. (ex. if) $\|Ax\| \leq \sigma_1 \|x\|$; $\|Ax\| \geq \sigma_{\min} \|x\|$ or
 $= |\lambda| \|x\|$ inverse to A^{-1} : $\frac{1}{\sigma_{\min}} \geq \frac{1}{|\lambda|_{\min}}$

□ Norm of matrices: 向量 p 范数 $\|x\|_p = (\sum_j |x_j|^p)^{\frac{1}{p}}$ ($1 \leq p$) · $\frac{1}{p}$ l-norm, 2-norm, ∞ -norm
 矩阵范数 范数满足 $\|x\| \geq 0$ ② $\|\alpha x\| = |\alpha| \|x\|$ ③ $\|x+y\| \leq \|x\| + \|y\|$ (Euclidean) ($\max_j |x_j|$)

· inducing norm/operator norm: $\|A\|_p = \max_{x \neq 0} \frac{\|Ax\|_p}{\|x\|_p}$. 特别 $\|A\|_1 = \max_j \sum_i |a_{ij}|$ (max col sum)

Frobenius norm (F-norm): $\|A\|_F = \sqrt{\sum_{i,j} |a_{ij}|^2} = \sqrt{\text{tr}(A^T A)}$; $\|A\|_2 = \sqrt{\lambda_{\max}(A^T A)} = \sigma_1$ (Eig. / spectrum)

· theorem 1: $\rho(A) \leq \|A\|$ (ρ is spectrum radius, $= \sqrt{\sum_i \sigma_i^2}$) $\|A\|_{\infty} = \max_j \sum_i |a_{ij}| = \|A^T\|_1$,

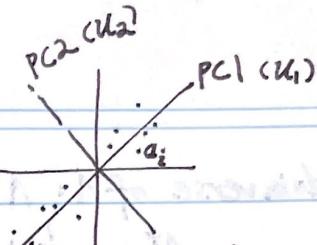
$\rho(A) = \max_i |\lambda_i|$ theorem 2 (Gelfand formula):

eq. when $\|\cdot\|$ is 2-norm and A is sym.

$$\rho(A) = \lim_{k \rightarrow \infty} \|A^k\|^{\frac{1}{k}}$$

o rank-one mt. $A = \sigma u v^T$ (unit vec u and v), $A^+ = \frac{v u^T}{\sigma}$. $AA^+ = u u^T$, $A^+ A = v v^T$

o' SVD of A is equivalent to the diagonalization of sym. $M = \begin{pmatrix} 0 & A^T \\ A & 0 \end{pmatrix}$, with eigenv = (u, v) .



7.3 principal component analysis (PCA by the SVD)

Data comes in a matrix: n samples and m measurements per sample.

Substracting the mean of each row to center A. Then SVD finds the (\vec{u}, \vec{v})

(sample) covariance matrix: $S = \frac{AA^T}{n-1}$ 'eigen-' or 'comb-' that contains the most info. or explains the most variance.

Then we find λ and σ , u and the leading direction in the scatter plot.

Total Variance $T = s_1^2 + \dots + s_m^2 = \sigma_1^2 + \dots + \sigma_m^2 = \text{tr}(S)$. PC1 accounts for $\frac{\sigma_1^2}{T}$ of total var.

1) PCA = perpendicular least squares (orthogonal regression): (not OLS / vertical least squares!)

$$\sum_{j=1}^n \|a_j\|^2 = \|A^T u_1\|^2 + \|A^T u_2\|^2 = \sum_{j=1}^n |a_j^T u_1|^2 + \sum_{j=1}^n |a_j^T u_2|^2. \quad \text{When } u_1^T A A^T u_1 \text{ maximized, } A^T A x = A^T b \quad \text{the squared distances sum minimize.}$$

(Correlation Matrix, $C = \frac{DATAD}{n-1} = DSD$. (and u_2 the second direction in ortho) measurements may have diff units and original scaling is not meaningful.

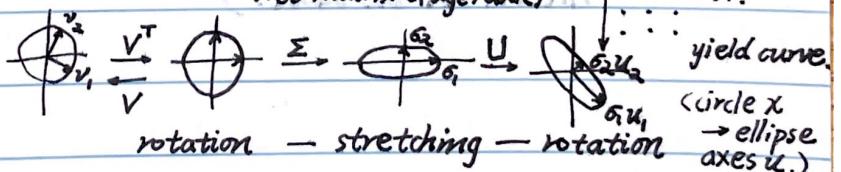
DA makes row have same length $\sqrt{n-1}$, or makes S has 1's on diag by dividing s_1, \dots, s_n .

App. Examples: ① Genetic variation in Europe. ② Eigenfaces. ③ Model order reduction, POD.

(When A is centered and $m > n$, σ 's have at most $n-1$ nonzeros.)

④ Web matrix (PageRank) ⑤ $T^T \dots$ ir.

7.4 the geometry of the SVD



The norm $\|Ax\| = \max \frac{\|Ax\|}{\|x\|} = \sigma_1$.

rotation — stretching — rotation

tri. ineq. $\|A+B\| \leq \|A\| + \|B\|$, product ineq. $\|AB\| \leq \|A\| \|B\|$ (proof: $\|Ax\| \leq \|A\| \|x\|$, times x)

Eckart - Young (Mirsky) Theorem: The closest rank k mt to A is $A_k = \sigma_1 u_1 v_1^T + \dots + \sigma_k u_k v_k^T$. ($\|A-B\| \geq \|A-A_k\| = \sigma_{k+1}$, $\forall B$ of $\text{rc}(B)=k$.)

Polar Decomposition: like $e^{i\theta} \cdot r$, $A = Q S$ where Q is orthogonal and S is sym. pos. semidefinite. rotation-stretching $Q = UV^T$, $S = V \Sigma V^T$. if A is inv., S is pos. definite.

$Q = UV^T$ is the nearest ortho. mt to A. (or S is the sqrt of $A^T A = V \Sigma^2 V^T$) (also $A = kQ$ with $k = U \Sigma U^T$ and Q the same.) ($\|Q-A\|_{\min}$)

A_0 is the nearest sim. mt to A by changing σ_{\min} to zero.

$n \times m$

pseudoinverse of A : $A^+ = V\Sigma^+U^T$ ($v_1 \dots v_n$) $\begin{pmatrix} \sigma_1^{-1} \\ \vdots \\ \sigma_r^{-1} \end{pmatrix}$ ($u_1 \dots u_m$)^T

$A^+ u_i = \frac{1}{\sigma_i} v_i$ ($i \leq r$) and $A^+ u_i = 0$ ($i > r$) $\xrightarrow{n \times n \quad n \times m \quad m \times m$ always exists. $(CA^+) = R(CA)$, $R(A^+) = C(A)$.

If A^- exists, A^+ is the same as A^- in that case $m=n=r$ and $V\Sigma^{-1}U^T$.

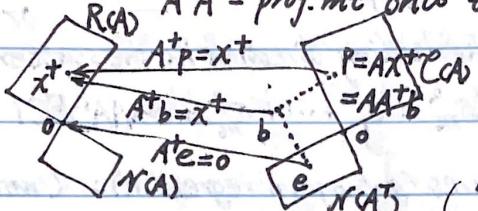
since $\Sigma^+\Sigma = (I_n)$, we get

$$AA^{-1} = A^{-1}A = I$$

(proj. itself)

AA^+ = proj. mt onto the col space of A .

A^+A = proj. mt onto the row space of A .



concisely:



$N =$

A^+ can get Ax^+ in col space back to $A^+Ax^+ = x^+$ in row space

$$(A^+A = (I_n)) \xrightarrow{n \times n \quad \text{nullspace}} \sum u_i u_i^T$$

$$(AA^+ = (I_m)) \xrightarrow{m \times m \quad \text{left nullspace}} \sum u_i u_i^T$$

least squares with dependent cols: $A^T A x = A^T b$ has inf. many solutions when A is singular. solution: $x^+ = A^+ b$ which is the shortest ($\|x^+\|_{\min}$, no nullspace part) solution to $A^T A \hat{x} = A^T b$ and $A \hat{x} = p$.

If A is full col rank ($r=n$) then it has left inv. $L = (A^T A)^{-1} A^T$, $LA = I$. $A^+ = L$ in this case.

If A is full row rank ($r=m$) then it has right inv. $R = A^T (A A^T)^{-1}$, $AR = I$. $A^+ = R$ in this case.

(proof: LA and AL are projs onto row space (\mathbb{R}^n) and col space (subset of \mathbb{R}^m , the origin P); AR and RA the same.)

8. Linear Transformations

8.1 the idea of a linear transformation

A transformation assigns an output $T(v)$ to each input vector $v \in V$.

Linearity requires $T(cv + dw) = cT(v) + dT(w)$. Note $T(0) = 0$, so affine trans. isn't linear.

Lines \rightarrow Lines, Triangles \rightarrow Triangles. Comb \rightarrow Comb, so basis is important.
 $(T(v) = Av + u_0, \text{ex.})$

$A = \begin{pmatrix} 0 & 1 & 0 \\ 0 & 2 & 0 \\ 0 & 0 & \dots & n \end{pmatrix}$ = mt. form of the derivative $T = \frac{d}{dx}$, with basis $v_i = x^{i-1}$. nullspace a line in func. sp.

and integral give the pseudoinv. $T^+ = \int_0^x dx$. (mt. form $A^+ = \begin{pmatrix} 0 & \dots & 0 \\ 1 & 0 & 0 \\ 0 & \dots & 1 \end{pmatrix}$) $A^+A = \begin{pmatrix} 0 & I \\ 0 & 0 \end{pmatrix}$, $AA^+ = I$.

The product of two trans. $(ST)(v) = S(T(v))$ still linear. Inverse trans. T^{-1} brings every vec

$T: V \rightarrow W$ Range of T = set of outputs. Kernel of T = set of all inputs for which $T(v) = 0$.
Ex. $T(v) = Av$ (V Image) ($C(A)$) ($N(A)$)

8.2 the matrix of a linear transformation

Every lin. trans. from V to W can be converted to a matrix. This mt. depends on bases.

Cols of 1-to-n of the matrix will contain those outputs $T(v_1)$ to $T(v_n)$.

$T(v_i) = \text{comb of output basis vectors} = a_{1i}w_1 + \dots + a_{ni}w_i$.

For every $v \in V$ that $v = c_1v_1 + \dots + c_nv_n$, linearity gives $T(v) = c_1T(v_1) + \dots + c_nT(v_n)$, as Ac.

Multiplication corresponds. $TS: U \rightarrow V \rightarrow W$. $AB: (m \times n)(n \times p) = (m \times p)$.

Change of basis: $Vc = V'c'$. (same vec written in diff basis) (B contains new basis vecs b in cols
written in the standard/old basis.)
1. $V'B = V$, $B = V'^{-1}V$; $c' = Bc$.
2. $A' = B_{\text{out}}^{-1}AB_{\text{in}}$ (of course it's inv.)

8.3 the search for a good basis

vec. sp. 1. $B_{\text{in}} = B_{\text{out}} = \text{eigen vec mt. } X$, $X^{-1}AX = \text{eigenvalues in } \Lambda$. (sq. A and has n ind. evecs.)

2. $B_{\text{in}} = V$ and $B_{\text{out}} = U$: singular vecs of A. $U^{-1}AV = \text{diag } \Sigma$. — 'isometric' (def: are ortho.)
 $\Sigma = Q_1^{-1}AQ_2$ when Q_1, Q_2

3. $B_{\text{in}} = B_{\text{out}} = \text{generalized eigenvects of } A$. $B^{-1}AB = \text{Jordan form } J$. (sq. A but only s < n ind. vcs.)

$J = \begin{pmatrix} \lambda_1 & & & \\ & \ddots & & \\ & & \lambda_s & \end{pmatrix}$, Jordan block: $J_i = \begin{pmatrix} \lambda_i & & & \\ & \ddots & & \\ & & \lambda_i & \end{pmatrix}$, $(A - \lambda_i I) b_{j+1} = b_j$ (at one block)
 $((A - \lambda_i I) x_{j+1} = x_j \text{ instead of } 0, x_j = \text{col } 0, 0, \dots, 0)$.

Matrices are similar if they share the same Jordan form, not otherwise. and $BX_j = b_j$.)

In app. we take J^k and e^{Jt} (involves $e^{\lambda t}$ times powers 1, t, \dots, t^{s-1})

(Jordan form is unstable!)

矩阵收敛性(四):

Hilbert matrix $H = \begin{pmatrix} 1 & \frac{1}{2} & \frac{1}{3} & \frac{1}{4} \\ \frac{1}{2} & \frac{1}{3} & \frac{1}{4} & \dots \\ \frac{1}{3} & \frac{1}{4} & \dots \\ \vdots & \vdots & \ddots \end{pmatrix}$, $H_{ij} = \frac{1}{i+j-1}$ if use powers $1, x, x^2, x^3, \dots$ as a func basis,
 $\int_0^1 x^{i+j-2} dx$. $B^T B$ is not ortho. and equals H : bad!
 H is ill-conditioned.

FFT 形象化补充:

逆归 $F_n = \begin{pmatrix} 1 & 1 & 1 & \dots & 1 \\ 1 & w & w^2 & \dots & w^{n-1} \\ 1 & w^2 & w^4 & \dots & w^{2(n-1)} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & w^{n-1} & w^{2(n-1)} & \dots & w^{(n-1)^2} \end{pmatrix}$

$w^n = 1$ $w^2 = -1$

proof of Schur Theorem: $U^H A U = T$. so (I).
(Schur decomp.) $(A \in \mathbb{C}^n, U \text{ unitary}, T \text{ tri.})$ (I).

induction: find A 's λ_i , a eigenvector e_i , expand it to a basis.

$$U = (e_1, e_2, \dots, e_n) \quad U^H A U = (\lambda_1 \dots 0 \ A_1) \quad \text{then do the same thing on } A_1, \text{ till } A_{n-1}: \quad \sum_{i,j} |a_{ij}|^2$$

$$\text{finally } (U_1 U_2 \dots U_{n-1})^H A (U_1 U_2 \dots U_{n-1}) = (\lambda_1 \dots \lambda_{n-1}) \quad U_2^H A_1 U_2 = (\lambda_2 \dots 0 \ A_2), \quad U_2 = (1 \ 0 \ \dots \ 0 \ U_2). \quad \text{tr}(AA^T)$$

Schur inequality: $\sum_i |\lambda_i|^2 \leq \sum_{i,j} |a_{ij}|^2$. (proof: $\sum_i |\lambda_i|^2 = \sum_i |t_{ii}|^2 \leq \sum_i |t_{ii}|^2 + \sum_i |t_{ij}|^2 = \text{tr}(TT^T)$)

change in A^{-1} from a change in A : Sherman - Woodbury - Morrison formula:

$$1. M = I - uv^T, M^{-1} = I + \frac{uv^T}{1-v^Tu} \quad (\text{rank-one change}) \quad Ax = b \rightarrow My = b: (M = A - uv^T)$$

$$2. M = A - uv^T, M^{-1} = A^{-1} + \frac{A^{-1}uv^TA^{-1}}{1-v^TA^{-1}u} \quad \text{① solve } A\mathbf{g} = u, \text{ compute } C = 1 - v^T\mathbf{g}$$

$$3. M = I - UV, M^{-1} = I_n + U(I_m - VU)^{-1}V \quad \text{② } c \neq 0 \text{ then } y = x + \frac{v^Tx}{c} \mathbf{g}.$$

$$4. (\text{matrix inversion lemma}) M = A - UW^{-1}V, M^{-1} = A^{-1} + A^{-1}U(W - VA^{-1}U)^{-1}VA^{-1}$$

Hamilton - Cayley Theorem. $f(\lambda) = \det(\lambda I - A)$ then we have $f(A) = 0$.

corollary: $\cdot A^K$ can be simplified by $f(A) = 0$ (downgrade) \uparrow
diagonalizable mts. are dense.

$\cdot A^{-1}$ can be expanded by a $g(A)$ ($f(A)$ times A^{-1} to get)

(DFT)

4. $B_{in} = B_{out} = \text{Fourier matrix } F$, Fx is a Discrete Fourier Transformation of x .

the eigenvector mt. F diagonalizes the permutation mt. P : ex. $\begin{pmatrix} 1 & 1 & 1 \\ 1 & -1 & -1 \\ -1 & 1 & -1 \end{pmatrix}$ $A = \begin{pmatrix} 1 & 1 & 1 \\ 1 & i & -1 \\ -1 & -i & 1 \end{pmatrix}$ $F = \begin{pmatrix} 1 & 1 & 1 \\ 1 & i & -1 \\ -1 & -i & 1 \end{pmatrix}$

F is sym. ortho? Circulant matrix: $C = \begin{pmatrix} c_0 & c_1 & c_2 & c_3 \\ c_3 & c_0 & c_1 & c_2 \\ c_2 & c_3 & c_0 & c_1 \\ c_1 & c_2 & c_3 & c_0 \end{pmatrix}$ eigenvalues $c_0 + c_1\lambda + c_2\lambda^2 + c_3\lambda^3$
(unitary)

func.sp. its four eigenvalues are given by $F_C = \begin{pmatrix} 1 & 1 & 1 & 1 \\ 1 & i & -1 & -i \\ -1 & -i & 1 & i \\ -i & 1 & -1 & -i \end{pmatrix}$ $\lambda = \pm 1, \pm i$

5. The Fourier basis. $1, \sin x, \cos x, \sin 2x, \cos 2x \dots$ periodic and orthogonal, so Fourier coeff.

6. The Legendre basis. $1, x, x^2 - \frac{1}{3}, x^3 - \frac{3}{5}x \dots$ ex. $a_i = \frac{\int f(x) \cos ix dx}{\int \cos ix \cos ix dx}$.

come from applying the G-S idea to orthogonalize $1, x, x^2, x^3 \dots$ ex. $\frac{(x^3, x)}{(x, x)} = \frac{3}{5}$, so $x^3 - \frac{3}{5}x \perp 1, x, x^2$

7. The Chebyshev basis. $1, x, 2x^2 - 1, 4x^3 - 3x \dots$ $\cos \theta \rightarrow \cos \theta$ (remind even \perp odd (x))

9. Complex Vectors and Matrices

9.1 complex numbers complex plane. polar form. n th roots of 1. $|z|^2 = r^2 = 3\bar{z}$.
 $w = e^{\frac{2\pi i}{n}}$. $\frac{z}{3} = \frac{\bar{z}}{|z|^2}, \bar{z}_1 = \bar{z}_2$.

9.2 hermitian and unitary matrices

Real versus Complex in a nutshell: When you transpose vector or mt A , take the complex conjugate too.

$\bar{z}^T = z^H$ called Hermitian or adjoint. (For truly useful $u^H u = \|u\|^2$ not $u^T u$.)

dot.p / inner.p: $u^H v$ the order is now important $u^H v \neq v^H u$. $(Au)^H v = u^H (A^H v)$

$$(AB)^H = B^H A^H$$

$S = S^H$ called Hermitian matrix. $z^H S z$ is real for $v z$; all eigenvalues of S is real;

$Q^H Q = I$, $Q^H = Q^{-1}$ (sq) called unitary. eigenvcs from diff eigenvalues are ortho. $q_1^H q_2 = 0$

$$\|Qz\| = \|z\| \text{ and then } Qz = \lambda z \text{ leads to } |\lambda| = 1.$$

a good row space is no longer $C(A^T)$ but $C(A^H)$.

9.3 the Fast Fourier Transform (FFT)

$$F_n^H = \bar{F}_n. F_n \bar{F}_n = nI, F_n^{-1} = \frac{1}{n} \bar{F}_n.$$

frequency space \xrightarrow{F} physical space $\xleftarrow{F^{-1}}$ $y_j =$ the Fourier series $\sum_k c_k e^{ikx_j}$.
those n points x (angles) equally spaced around.

(with $w = \bar{w} = e^{-\frac{2\pi i}{n}}$, or
permute row $i \leftrightarrow N-i$)

a entry in row j col k is w^{jk} , zeroth row and col contains all $w^{0k} = 1$.

recursion: $F_n \rightarrow F_{\frac{n}{2}}$ even-odd perm. Δ Fr.mt. is the Vandermonde mt. for interpolation at n .

$$F_n = \begin{pmatrix} I_{\frac{n}{2}} & D_{\frac{n}{2}} \\ I_{\frac{n}{2}} & -D_{\frac{n}{2}} \end{pmatrix} \begin{pmatrix} F_{\frac{n}{2}} & \\ & I_{\frac{n}{2}} \end{pmatrix} \begin{pmatrix} 1 & 1 & \dots \\ \dots & 1 & \dots \\ & & 1 \end{pmatrix}$$

$$y_j = \sum_{k=0}^{n-1} w^{jk} c_k = \sum_{k=0}^{\frac{n}{2}-1} w^{jk} c_{2k} + \sum_{k=0}^{\frac{n}{2}-1} w^{j(2k+1)} c_{2k+1}$$

$$(P(z) = f)$$

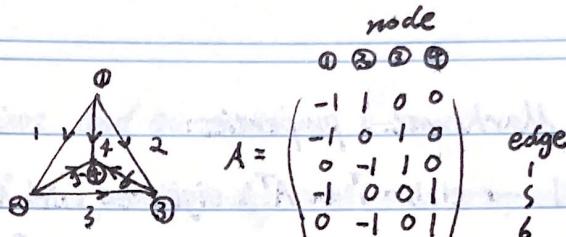
$$(w' = w^2) = \sum_{k=0}^{\frac{n}{2}-1} w^{jk} c_{2k} + \sum_{k=0}^{\frac{n}{2}-1} w^{j(2k+1)} c_{2k+1} \cdot w^j = y'_j, \text{ even} + w^j y'_j, \text{ odd}$$

$$D = \text{diag}(1, w, \dots, w^{\frac{n}{2}-1})$$

$O(n^2) \rightarrow O(\frac{1}{2}n \log_2 n)$ $\xrightarrow{\text{bit-reversed order}}$ $\underbrace{0 \ 1 \ 0 \ 1}_{C's \text{ final permutation}}$

10. Applications

10.1 graphs and networks



The incidence matrix A comes from a connected graph with n nodes and m edges.

Elimination reduces every graph to a tree. Rows' dependency \leftrightarrow edges form a loop.

$N(A)$: The constant vec ($c, c \dots, c$) make up the nullspace, $\dim = 1$.

△ Use orthogonality

CATs: The edges of any tree give r ind. rows, $r = n-1$.

to judge a ver whether

, in a space of 1.

(CA): Ax gives voltage diffs. Voltage Law: The components of Ax add to zero in a space of $n-1$. dim = $n-1$

NCAT_j : Current Law: $A^T y = 0 = (\text{flowin}) - (\text{flowout})$ is solved by loop currents around all loops.

since Ohm's Law $y = -CAx$. C is the conductance mt. (diag) | F 1 : 5

if all $c=1$, we get the graph's Laplacian matrix $A^T A$. (diag) Euler's Formula
 $A^T A x = f$ (With A and C : Nodes - Edges + Loops = 1.
 we call it a network)

10.2 matrices in engineering

$A^T C A x = f$ (we call it a Network.)
 ~~αx~~ (batteries / current sources ..) csmodels

$\frac{d}{dx} (C(x) \frac{d}{dx}) = f(x)$, with boundary conditions : $x = \frac{1}{n+1}$. AT^T is not inv. One node has to be grinded.

divide the bar into n pieces of $5x$; $U(50) = 0$ and $(U(5)) = 0$ or Thus remove one row and col.

replace $\frac{d}{dx}$ by A and $-\frac{d}{dx}$ by A^T ; (they include $\frac{d}{dx}$)
 end conditions are $u_n = 0$ and ($u_n = 0$ or $y_n = 0$);

thus corresponds to C, thus: $f = A^T y$, $y = Ce$, $e = Au$ give $A^T C A u = f$.

and $\lambda = \mu$. This is a contradiction. \square

We have three choices in replacing $\frac{dx}{dx}$: $\frac{dx}{dx}$, $\frac{dx}{dx}$, $\frac{2dx}{dx}$. "3
 $(\text{forward diff})^T = - \text{backward diff}$. forward backward centered free

fixed-fixed $\begin{pmatrix} -1 \\ 1 \\ -1 \end{pmatrix}$ $\begin{pmatrix} -1 \\ 1 \\ 1 \end{pmatrix}$, fixed-free $\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$ $\begin{pmatrix} 1 \\ -1 \\ 1 \end{pmatrix}$, free-free $\begin{pmatrix} 1 \\ 1 \\ 1 \end{pmatrix}$ $\begin{pmatrix} -1 \\ -1 \\ 1 \end{pmatrix}$.

$$K's \text{ properties: } \begin{pmatrix} 2 & -1 \\ -1 & 2 \end{pmatrix} \quad \text{circular} \quad \begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix} \quad \begin{pmatrix} 2 & -1 \\ 1 & 2 \end{pmatrix} \quad \begin{matrix} \text{sin. semipos.} \\ a\lambda = d = 0, x_N \end{matrix} \quad \begin{pmatrix} 1 & -1 \\ -1 & 2 \end{pmatrix}$$

④ sym. ⑤ tridiag. $\begin{pmatrix} 1 & -1 \\ -1 & 1 \end{pmatrix}$ $\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$, thus
 ⑥ inv. (6) $\begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$ $\begin{pmatrix} 1 & -1 & 0 \\ 0 & 1 & -1 \\ 0 & 0 & 1 \end{pmatrix}$
 f have $f + f + f = 0$

$$\oplus \begin{pmatrix} k^{-1} & -2 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix} \text{ also } X_{\mathcal{M}} = \begin{pmatrix} 2 & -1 & 1 \\ -1 & 2 & -1 \end{pmatrix} \text{ have } j_1, j_2, j_3$$

\checkmark R is full (not sparse) with all entries $\begin{pmatrix} -1 & 2 & -1 \\ -1 & -1 & 2 \end{pmatrix}$ > 0 , in nullspace.

o some proof of Markovitz's properties:

A 's cols adds to 1 $\rightarrow (1, 1, \dots)^T$ an A^T 's eigenvect with $\lambda = 1$, also A 's λ . \rightarrow other eigenvects will perpendicular to A^T 's $(1, 1, \dots)^T$, which means $\xrightarrow{\text{so all } u_0 \text{ would finally go to } x_1}$

\rightarrow we need to proof that all other $|\lambda| < 1$:

□ Perron-Frobenius Theorem:

$$\Phi P(A) < 1$$

• spectrum radius $P(A) = \max |\lambda(A)|$, here are some equal conditions: ② $\lim_{k \rightarrow \infty} A^k = 0$

$$(\approx 2\text{-norm}) \quad (\text{Gelfand Formula: } P(A) = \lim_{k \rightarrow \infty} \|A^k\|^{\frac{1}{k}} \leq \|A\|) \quad \text{③ } \sum_{k=0}^{\infty} A^k \text{ converges. } (= (I-A)^{-1})$$

• $(A \geq 0)$ specr's monotocity: $B \in \mathbb{C}$ s.t. $|b_{ij}| < a_{ij} \forall i, j$. then $P(B) \leq P(BI) \leq P(A)$. Neumann Series, specr with sums: $\min_{\text{row/col}} a_{ij} \leq P(A) \leq \max_{\text{row/col}} a_{ij}$. Collatz-Wielandt Formula: with $P(A) < 1$ it can converge: $\min_{1 \leq i \leq n} \frac{(Ax)_i}{x_i} \leq P(A) \leq \max_{1 \leq i \leq n} \frac{(Ax)_i}{x_i}$ ($= 1$ when x is eivec.)

then 1. $P(A) > 0$ and has $CW = AW = I$. 2. all nonneg. eigenvect correspond to eigenvalue

3. \exists only one $v \in \mathbb{R}^n$ s.t. $v > 0$, $\sum_i v_i = 1$ and $Av = P(A)v$; (left eivec) $v^T w = 1$ and $w^T A = P(A)w$ (all other λ 's vec has entries < 0)

4. Perron Proj.: $\lim_{k \rightarrow \infty} \left(\frac{A}{P(A)} \right)^k = vw^T$ (reducible: $\exists P$ s.t. $P^T AP = \begin{pmatrix} B & C \\ 0 & D \end{pmatrix}$, (or $A^T w = P(A)w$)

(blockized upper tri.) then A 's graph X strongly connected)

o sensitivity of eigenvalues: A makes small change ΔA , which results in $\Delta \lambda$:

$$\Delta \lambda = (X + \Delta X)^{-1}(A + \Delta A)(X + \Delta X) - X^{-1}AX = (X^{-1} - X^{-1}\Delta X X^{-1})A(X + \Delta X) + X^{-1}\Delta AX - X^{-1}AX$$

$$\text{we use } X^{-1}A = \lambda X^{-1} \rightarrow = \lambda X^{-1}\Delta X - X^{-1}\Delta X \lambda + X^{-1}\Delta AX. \quad (\text{omit the second term})$$

now we'll show that $\lambda X^{-1}\Delta X$ and $X^{-1}\Delta X \lambda$ cancel on the diag: $\lambda \in X^{-1}\Delta X$ with $(X^{-1}\Delta X)\lambda$
so $\Delta \lambda = X^{-1}\Delta AX$,

o proof of $(I-A)^{-1}$: $P(A) < 1$ vice above; $P(A) = 1$ means a $\lambda = 0$ or $\Delta \lambda = y^T \Delta A x$.

$\lambda = 0$, singular; $P(A) > 1$ means a $\lambda = \frac{1}{1-p} < 0$, since its eigenvect $u > 0$, $(Ax = \lambda x, A^T y = \lambda y)$

o matrix series: $\sum_i c_i B^i$ has convergence radius R , then shows there must be entries c_i ($(I-A)^{-1} u = \lambda u < 0$)

$(A \in \mathbb{C}^{n \times n}) \sum_i c_i A^i$: $\Phi P(A) < R$, absolute converge; $\Phi P(A) > R$, diverge.

10.3 Markov matrices, population, and economics

Markov mt.: every entry of A is positive $a_{ij} > 0$ ($A > 0$); every cols of A adds to 1.

Then $\lambda_1 = 1$ is larger than any other values, its vec x_1 is the steady state: $u_k = x_1 + \sum_{i=2}^n c_i (\lambda_i)^k x_i$

• Perron-Frobenius Theorem (PFA): for $A > 0$, All numbers in $Ax = \lambda_{\max} x$ are positive.

Leslie mt.: population: $\begin{pmatrix} F_1 & F_2 & F_3 \\ P_1 & & \\ & P_2 & \end{pmatrix}$ with $\lambda_{\max} \geq 1$ maybe. F 's and P 's change in αA makes λ_{\max} change.

Consumption mt.: $p - Ap = y$, $p = (I - A)^{-1} y$ Neumann/geometric Series:
(Leontief) product. demand. $(I - A)^{-1} = I + A + A^2 + \dots$

it converges if $P(A) < 1$; $P(A) = 1$ makes $(I - A^{-1})$ fail to exist; $P(A) > 1$ then $(I - A^{-1})$ has neg. entries.
(cannot meet demand)

10.4 linear programming



(A has $m < n$ usually.)

Primal problem: minimize $c^T x$ with $Ax = b$ ($AX \geq b$) and $x \geq 0$

Dual problem: maximize $y^T b$ with $A^T y \leq c$ and $y \geq 0$

$x \geq 0$ and the slack $s = c - A^T y \geq 0$ gave $x^T s \geq 0$ which means $y^T b \leq x^T c$, when $y^T b$ max. equality means $x^T s = 0$ thus either $x_j = 0$ or $s_j = 0$, i.e. y solves m eqs. $A^T y = c$ in $c^T x$ min.

geometry: a feasible set ('triangle'), maybe unbounded. the m components that are nonzero in x .

The optimal solution will be one of corners. (m nonzeros and $n-m$ zeros)

① list them and compare. ② the simplex method: start from one corner,

enter one variable $0 \rightarrow 1$, others have to adjust to keep $Ax = b$.

iii) ③ the interior point method: compute $c \cdot x$ and choose that gives the most neg. change.

remove the constraints $x_j \geq 0$ by then decide how much it can enter and drop one var. to 0.

minimize $c^T x - \theta \sum \log x_i$ with $Ax = b$. repeat until all changes are pos.

(move to neighbor)

Lagrange $L(x, y, \theta) = c^T x - \theta \sum \log x_i - y^T (Ax - b)$

$\frac{\partial L}{\partial x_j} = c_j - \frac{\theta}{x_j} - (A^T y)_j = 0$ which is $x_j c_j - \theta = 0 \rightarrow 0$. actually solved by Newton's method (iteration)

10.5 Fourier series: linear algebra for functions

Hilbert space: if and only if their length are finite: $\|v\|^2 = v_1^2 + v_2^2 + v_3^2 + \dots$ (Schwarz ineq. ✓ $|v \cdot w| \leq \|v\| \|w\|$)

$f(x) = a_0 + a_1 \cos x + b_1 \sin x + a_2 \cos 2x + b_2 \sin 2x + \dots$ (orthon. means $\int_0^{2\pi} f(x) g(x) dx = 0$)

$\|f\|^2 = 2\pi a_0^2 + \pi (a_1^2 + b_1^2 + a_2^2 + b_2^2 + \dots)$ basis. $\frac{1}{\sqrt{2\pi}}, \frac{\cos x}{\sqrt{\pi}}, \frac{\sin x}{\sqrt{\pi}}, \dots$)

since basis are ortho. $a_n = \frac{1}{\pi} \int_0^{2\pi} f(x) \cos nx dx$, $b_n = \frac{1}{\pi} \int_0^{2\pi} f(x) \sin nx dx$ ($a_0 = \frac{1}{2\pi} \int_0^{2\pi} f(x) dx = \bar{f(x)}$)

(orthon.)

ex. $f(x)$,

10.6 computer graphics

$$(x, y, z, 1) \quad (x, y, z)$$

we use homogeneous coordinates. note to separate 'point' with 'vector'.

Here are some operations: translation $T = \begin{pmatrix} 1 & & & \\ & cx, y, z, 1 \end{pmatrix}$ (like to rescale the space by $\frac{1}{c_i}$)
 rescaling $S = \begin{pmatrix} c_1 & & & \\ & c_2 & & \\ & & c_3 & \\ & & & 1 \end{pmatrix}$; rotation $R = \begin{pmatrix} 1 & & & \\ & \cos\theta & -\sin\theta & 0 \\ & \sin\theta & \cos\theta & 0 \\ & 0 & 0 & 1 \end{pmatrix}$: (add two points like to find middle point.)
 any affine matrix $A = \begin{pmatrix} T_{11}, 0, 0, 0 \\ T_{12}, 0, 0, 0 \\ T_{13}, 0, 0, 0 \\ T_{14}, 0, 0, 1 \end{pmatrix}$ projection $\begin{pmatrix} I - uu^T & 0 \\ 0 & 1 \end{pmatrix}$ (minor $I - uu^T$) sometimes there are more steps: ex. proj. to a flat.
 $(\frac{I}{-v_0}, 1)(I - nn^T)(\frac{I}{v_0}, 1)$

10.7 linear algebra for cryptography

modular arithmetic: Inversion of every $y \in \mathbb{Z}/p\mathbb{Z}$ will be possible if and only if p is prime.

Hill cipher: $y, 2y, \dots, py$ have diff reminders, there must be 1. (con-n)y $\equiv 0 \pmod{p}$, since m-n or y < p, cannot.)
 divide message into blocks $\vec{v}_1, \vec{v}_2, \dots$ of size n.

multiply each by encryption mt. $E \pmod{p}$. decryption mt. will be E^{-1} (p is prime and this exists.)
 finite field $F_p = \{0, 1, 2, \dots, p-1\}$, or p^k members. define its +, \times rules. (p's factors won't have inverse, $\equiv 1$)

F_2^3 have 8 vecs. 29 mts. over F_2 (3×3), maybe sin. (ex. mod) (It can be difficult when p is large or done many times.)

11. Numerical Linear Algebra

speed, accuracy, stability

11.1 Gaussian elimination in practice

roundoff error \rightarrow partial pivoting: choose the largest num in row k or below, exchange for pos. definite mt. (may also exchange cols or rescale)
 exchanges aren't required. (no improve)

operation counts of full mt.: Gaussian elim. has two advantages over $A^{-1}b$:

orthogonalization: $\frac{2}{3}n^3$ (Householder) $\Theta \frac{1}{3}n^3$ mul-sub compared to n^3 in A^{-1} (next n^2 equal)
 (Givens) (Givens) $\Theta 2wn + 2wn$ if A is banded so are L, U, but A^{-1} is full.

Euler angles: $\Omega_{32}\Omega_{31}\Omega_{21}A = R$, where A is ortho so $R = I$! $A = \Omega_{32}^{-1}\Omega_{31}^{-1}\Omega_{21}^{-1}$

□ LA in probability and statistics:

线代在概率统计中(简)

Monte Carlo method: $E(x) \leftarrow \frac{\sum x_i}{N}, \sigma \sim \frac{1}{\sqrt{N}}$

multilevel MC: 2-level $E(x) = \frac{1}{N} \sum_{i=1}^N y(b_i) + \frac{1}{N^*} \sum_{i=1}^{N^*} (x(b_i) - y(b_i))$

(for same accuracy, since $x-y$ has a smaller σ^* , N^* can be smaller than N)

fixed cost $NC + N^*C^* = T$, optimal ratio $\frac{N^*}{N} = \frac{\sigma^*}{\sigma} \sqrt{\frac{C}{C^*}}$

3-level $E(x) = \frac{\sum x}{N} + \frac{\sum (y-x)}{N^*} + \frac{\sum (x-y)}{N^{**}}$

\dots and optimize those N .
($\sum \sum$ or $\int dx$)

Covariance matrix: $V = E((x-\bar{x})(x-\bar{x})^T)$

correlation mt.

$$u^T V u = E(\|u^T(x-\bar{x})\|^2) \geq 0.$$

(sim. $|V|=0$)

$R = \begin{pmatrix} 1 & p_{xy} \\ p_{yx} & 1 \end{pmatrix}$ and we have $\sigma_i^2 \sigma_j^2 \geq \sigma_{ij}^2$ transform for $Z = AX$,

$p_{xy} = \frac{\rho_{xy}}{\sigma_x \sigma_y} \in [-1, 1]$

$D = \text{diag}(\frac{1}{\sigma_i})$

$$V_Z = A V_X A^T$$

(multivariate) Gaussian $p(x) = \frac{1}{(\sqrt{2\pi})^m |V|} e^{-\frac{(x-m)^T V^{-1} (x-m)}{2}}$, m = means of x .

we can diagonalize the cov. mt $V^{-1} = Q \Lambda^{-1} Q^T$, $Y = Q^T X$, $X = Q Y$, which means finding

verify: $\int p(x) dx = \int e^{-\frac{Y^T \Lambda^{-1} Y}{2}} dY = \prod_i \sqrt{2\pi \lambda_i} = (\sqrt{2\pi})^m |V|$ ($dX = dY$) corrs. that are ind. (experi.)

$$\int x p(x) dx = \int x p(Y) dY + m \int p(Y) dY = m.$$

$$\int (x-m) p(x) (x-m)^T dx = Q \int Y Y^T e^{-\frac{Y^T \Lambda^{-1} Y}{2}} dY \cdot Q^T = Q \Lambda^{-1} Q^T = V$$

weighted least squares:

$$\frac{1}{\sqrt{2\pi}^m |V|} \quad (\hookrightarrow Z = \frac{1}{\sqrt{\Lambda}} Y, dY = \sqrt{\Lambda} dZ)$$

previously we choose \hat{x} minimize $\|b-Ax\|^2$, get $\int Z Z^T e^{-\frac{Z^T Z}{2}} dZ = \frac{1}{\sqrt{2\pi}^m |V|}$

now we minimize $E = (b-Ax)^T V^{-1} (b-Ax)$, $A^T A \hat{x} = A^T b$. \leftarrow normal dis. $\sigma^2 = 1, m=0$.

'whitening noise' — $A^T V^{-1} A x = A^T V^{-1} b$ (if b 's errors are not ind. or variances not equal)

(means $\frac{b_i}{\sigma_i}$ to get $N(0, 1)$, then $V^{-\frac{1}{2}} A x = V^{-\frac{1}{2}} b$)

substitute by $A \rightarrow V^{-\frac{1}{2}} A$ and $b \rightarrow V^{-\frac{1}{2}} b$)

variance of \hat{x} : since $\hat{x} = Lb = (A^T V^{-1} A)^{-1} A^T V^{-1} b$ use trans. formula to get $V(\hat{x}) = (A^T V^{-1} A)^{-1}$

recursive least squares: static / dynamic: Kalman filter:

update: $(A_0) \hat{x}_0 = b_0$ (without $x_k = F_{k-1} x_{k-1}$, C_k changing state here)

$$(A_0^T A_0^T) (V_0^{-1}) (A_0) \hat{x}_1 = (A_0^T A_0^T) (V_0^{-1}) (b_0)$$

$$\rightarrow \hat{x}_1 = \hat{x}_0 + K_1 (b_1 - A_0 \hat{x}_0)$$

$$\text{cov. mt. } W_1^{-1} = W_0^{-1} + A_0^T V_0^{-1} A_0 \quad (\text{the same form})$$

$$\text{Kalman gain mt. } K_1 = W_0 A_0^T V_0^{-1}$$

11.2 norms and condition numbers

we choose ℓ^2 norm here. sym. $A : \|A\| = |\lambda_{\max}(A)| (= \rho(A))$.

sensitivity to error: sym. or unsym. $A : \|A\| = \sqrt{\lambda_{\max}(A^T A)} (= \sigma_{\max}(A))$.

$$Ax = b \quad \rightarrow \quad A(x + \alpha x) = b + \alpha b \quad \frac{\|\alpha x\|}{\|x\|} \leq c \frac{\|ab\|}{\|b\|} \quad \text{condition number } c = \|A\| \|A^{-1}\|$$

$$\Downarrow \quad (A + \alpha A)(x + \alpha x) = b \quad \frac{\|\alpha x\|}{\|x\|} \approx \frac{\|\alpha x\|}{\|x + \alpha x\|} \leq c \frac{\|\alpha A\|}{\|A\|} \quad (\text{for pos. mt. } c = \frac{\lambda_{\max}}{\lambda_{\min}})$$

equal when b along largest eigen, αb along smallest eigen, amplified by c .

11.3 iterative methods and preconditioners

) split A into $S - T$. iteration $Sx_{k+1} = Tx_k + b$ — two goals: speed per step and fast converg.

1. Jacobi method: keep the diag(A) as S , off-diag part is T .

$P < 1 \leftarrow$
(diag or not.)

2. Gauss-Seidel method: keep the lower tri. part as S . error eq. $e_{k+1} = S^{-1} T e_k$

& SOR (successive overrelaxation): $\omega Ax = wb$ (cut storage and usually speed up iter.) S has diag(A) and below diag of ωA . $e_i = x_\infty - x_i$

3. Elimination (exact LU) (triang. choose ω to make the $P(S^T T)$ small asp. very fast.)

& incomplete LU: ('fill-in'. sparse mt. that has nonzeros far from diag.)

$L_U_0 x_{k+1} = (L_U_0 - A)x_k + b$ set small nonzeros to 0 in L.U.

4. multigrid 5. conjugate gradients & preconditioned CG.

& GMRES, MINRES. $\begin{matrix} \text{Asym. H} \\ \text{Krylov} \end{matrix}$ $\begin{matrix} \Delta \text{tridiag.} \\ A\Omega = QH \end{matrix}$

for eigens: 6. power method & inverse power method: \langle Arnoldi iter. ; Lanczos iter. \rangle

$u_k = A^k u_0$ the largest eigenvalue dominates. (u_k gradually points towards x_i);

or by A^{-1} , solve $Au_{k+1} = u_k$ (by saving L.U), λ_{\min} dominates. For speed we shifted:

8. QR method: $A = QR$, $A_i = R_i Q_i = Q_i^{-1} A_i Q_i$ thus same $A - \lambda^* I$ (λ^* can choose a to make λ_{\min} Rayleigh quo. $\frac{x^T A x}{x^T x}$)

$A_1 = Q_1 R_1$, $A_2 = R_1 Q_2$, ... eigenvalues eigenvalue with A even smaller.

two extra ideas: begin to show up in diag of A_n . (last entry $(0 \dots 0 \lambda_n)$), then

① shift. $A_n - c_n I$ into $Q_n R_n$, $A_{n+1} = R_n Q_n + c_n I$ back. (c_n choose to near an unknown remove this r.c.)

② elimination. obtain zeros before. EAE' eigenvalue through diag of A_n . (or Givens) (move) leave nonzeros along subdiag, which QR factor drops from (Hessenberg mt.) $O(n^3) \rightarrow O(n^2)$.