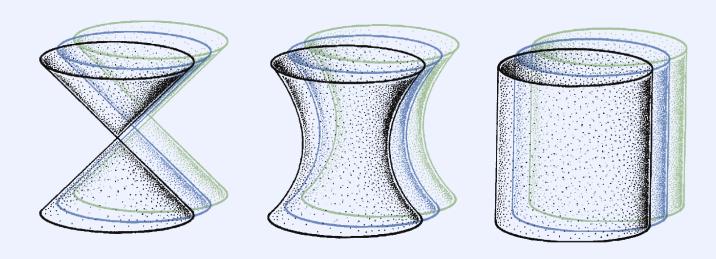
# Geometric Measure Theory

An Introduction to the Theory of Inconvenient Surfaces

Saturday, December 10, 2022



### Course Website

largoscv.github.io

### Element.io LaTeX Integration

5:47PM

Let  $\Sigma$  be a smooth oriented surface in  ${\bf R}^3$  with boundary  $\partial \Sigma$ . If a vector field  ${\bf F}(x,y,z)=(F_x(x,y,z),F_y(x,y,z),F_z(x,y,z))$  is defined and has continuous first order partial derivatives in a region containing  $\Sigma$ , then

$$\iint_{\Sigma} (
abla imes F) \cdot d^2 \Sigma = \oint_{\partial \Sigma} \mathbf{F} \cdot d\mathbf{\Gamma}$$

With the above notation, if  ${f F}$  is any smooth vector field on  ${f R}^3$ , then

$$\oint_{\partial \Sigma} \mathbf{F} \cdot d\mathbf{\Gamma} = \iint_{\Sigma} 
abla imes \mathbf{F} \cdot d^2 \Sigma$$
 (edited)

@ 27182818284tropy

$$\forall X \ (\forall \varnothing \ (\forall t \ (t \not\in \varnothing) \Rightarrow \varnothing \not\in X) \Rightarrow \exists f \ (\forall e \ (e \in f \Rightarrow \exists a \ (a \in X \land \exists b \ (\exists edited))) \}$$

#### To Enable LaTeX Rendering:

- 1. Navigate to:
  - Linux:cd ~/.config/Element/
  - Windows:C:/Users/user/Appdata/Roaming/Element
  - MacOS:~/Library/Application Support/Element
- 2. Add/replace config.json with contents in

https://pastebin.com/8wV9nGR3

3. Restart Element (Task Manager). Then go to All Settings > Labs > "Render LaTeX Maths"

#### **Basic Outline**

**GMT Part 1:** 10 December 2022 to 21 January 2023

- 6 week course in general measure theory a la Carathéorory
- Concepts from geometric measure theory & applications (minimal surfaces) taught in lecture
- 6 homework assignments, first HW due 17 December (3:30pm EST)
- Pace: One section of Federer a week (~10-15 pgs/wk)

Saturday

10
December

- Introduction to GMT
- Relations & Orders
- Transfinite Induction
- Ordinals and Cardinals

Wednesday

14

December

Study Session

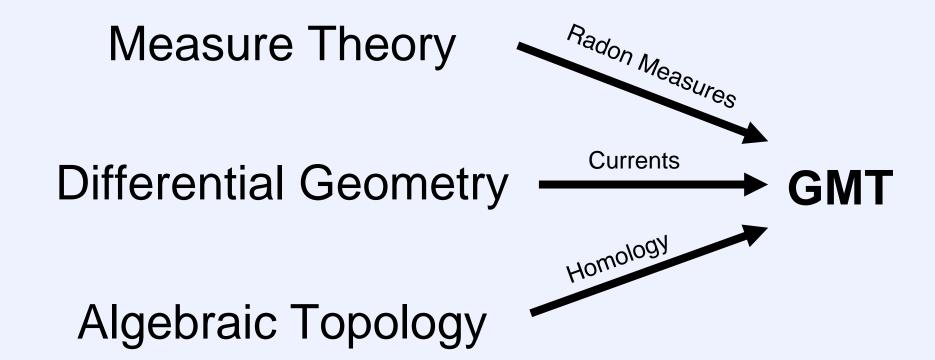
Saturday

**17** 

December

- Hausdorff Measure
- Densities
- Approximate limits
- Approximate continuity

## **Geometric Measure Theory**



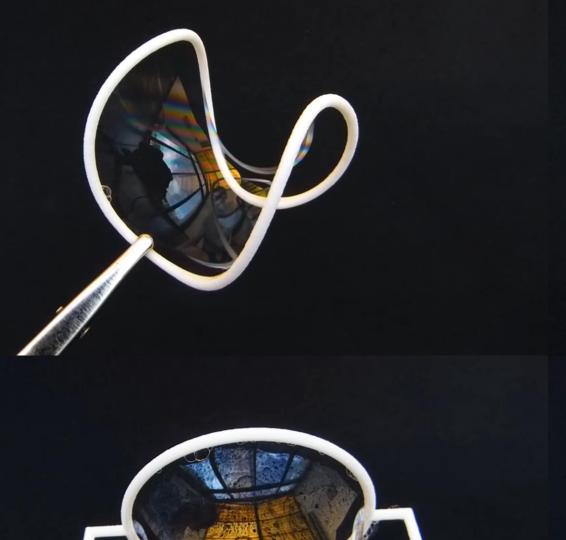
Joseph-Louis Lagrange (1736-1813)

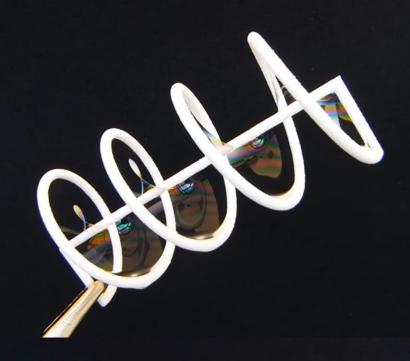
• Given a smooth closed curve  $\gamma$  in  $\mathbb{R}^3$ , does there exist a surface S of least area such that

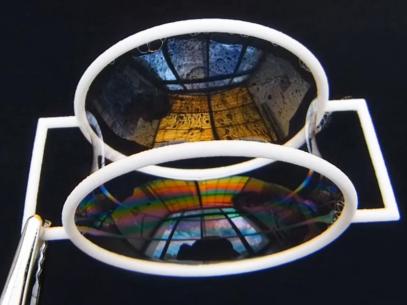
$$\partial S = \operatorname{graph}(\gamma)$$
?

- ❖ Such minimal surfaces may be modeled using soap films, as done by Joseph Plateau (1801-1883). Hence, this problem is known as **Plateau's Problem**.
- ❖ Solved by Douglas & Rado (1930) under topological restrictions, Federer & Fleming (1960) used the theory of currents to solve the problem without topological restrictions (orientable version).









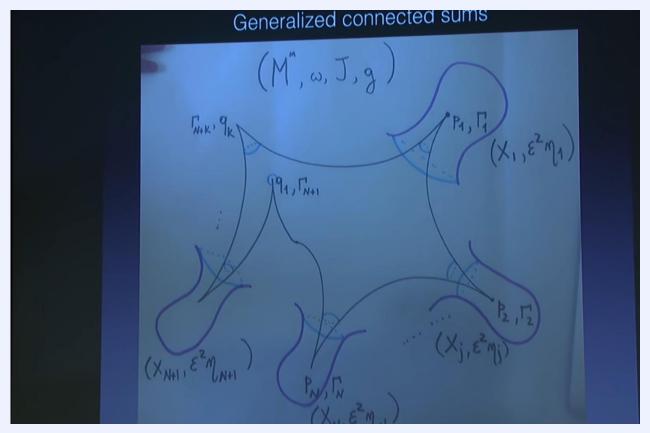


#### **Plateau's Problem in Higher Dimensions**

❖ Any solution of Plateau's Problem in R³ is a "convenient surface" (no corners, tangent plane that exists everywhere).

 $\bullet$  In higher dimensions, i.e.,  $\mathbb{R}^n$ , there are solutions of Plateau's Problems that have singularities (points with no tangent space)

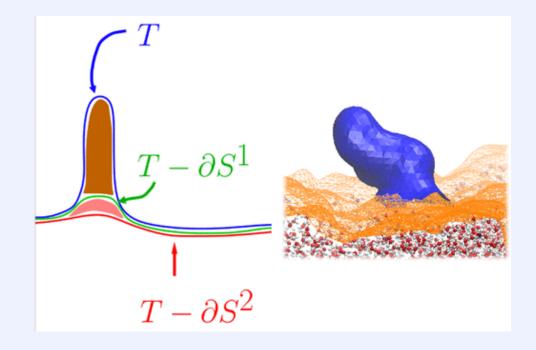
- $\bullet$  To study Plateau's Problem in  $\mathbb{R}^n$ , we need:
  - Theory of Caccioppoli Sets (De Giorgi)
  - Theory of Rectifiable Currents (Federer & Fleming)



Claudio Arezzo, 2019

### **GMT Sans Minimal Surfaces**

- Partial Differential Equations (Evans)
- Several Complex Variables (Krantz)
- Calculus of Variations
- Differential & Riemannian Geometry



- Continuum Mechanics, Cauchy's Stress Theory (Falach, 2013)
- Analysis of Soft Matter Surfaces (Alvarado et. al., 2020)

## **Background on ZFC Set Theory**

### Binary Relations

Given sets X and Y, the Cartesian product  $X \times Y$  is the set

$$\{(x, y): x \in X \text{ and } y \in Y\}$$

A **binary relation** over X and Y is a subset  $R \subseteq X \times Y$ .

If  $(x, y) \in R$ , we may write xRy.

#### Let *A* be a non-empty set and $\leq$ , $\prec$ be binary relations on *A*.

The relation ≤ is...

**Reflexive** if  $(\forall a \in A)(a \leq a)$ 

**Transitive** if  $(\forall a, b, c \in A)$ ,

$$(a \le b) \land (b \le c) \Rightarrow (a < c)$$

Weakly antisymmetric if  $(\forall a, b \in A)$ ,

$$((a \le b) \land (b \le a) \Rightarrow (a = b)$$

A **pre-order** on *A* satisfies the first two properties. A **partial order** satisfies all three.

The relation  $\prec$  is...

**Irreflexive** if  $(\forall a \in A)(\neg(a \prec a))$ 

**Antisymmetric** if  $(\forall a, b \in A)$ ,

$$(a < b) \Rightarrow (\neg(b < a))$$

A **strict partial order** on *A* is an irreflexive, transitive, and antisymmetric relation.

### Examples

#### **Total Orders**

A partial order  $\leq$  on a set A is called **total** (or a linear order) if it satisfies the additional property

$$(\forall a, b \in A)(a \leq b \lor b \leq a)$$

A strict partial order  $\prec$  on a set A is a (strict) total order if the associated partial ordering  $\leq$  is total.

#### Well Orders

A **well-order** on *A* is a total order  $\prec$  on *A* such that every non-empty subset of *A* has a  $\prec$ -least element. That is,

$$(\forall B \subseteq A \text{ nonempty})(\exists x \in B)(\forall b \in B)$$

### Well-Ordering Theorem:

Every set can be equipped with a well-order.

(Equivalent to the Axiom of Choice)

### Initial Segments

Given a set A equipped with a well order  $\prec$ , a set  $I \subseteq A$  is...

 $\Leftrightarrow$  An **initial segment** of A if

$$(\forall i \in I)(\forall a \in A)(a < i \Rightarrow a \in I)$$

**\Leftrightarrow** A **proper initial segment** if *I* is an initial segment and  $I \neq A$ .

### Induction on Well Orderings

**Theorem.** Given a set A equipped with a strict well order  $\prec$ , and  $\Psi(x)$  a property defined for all  $x \in A$ . If for all  $a \in A$ , we have that

$$(\forall b < a) \big( \Psi(b) \big) \Rightarrow \Psi(a)$$

then  $(\forall a \in A)(\Psi(a))$  holds true.

Proof. Straightforward.

### The Regularity Axiom:

$$(\forall x)[x \neq \emptyset \Rightarrow (\exists y \in x)(x \cap y = \emptyset)]$$

(Every non-empty set contains an element that is disjoint from it)

#### **Ordinals**

A set  $\alpha$  is called an **ordinal** if

1.  $\alpha$  is transitive, i.e.,  $(\forall \beta \in \alpha)(\beta \subseteq \alpha)$ , and

2. 
$$(\forall \beta, \gamma \in \alpha)(\beta = \gamma \lor \beta \in \gamma \lor \gamma \in \beta)$$

**Theorem**. A set  $\alpha$  is an ordinal if and only if  $\alpha$  is a transitive set and  $(\alpha, \in)$  constitutes a well-ordering.

*Proof.* Regularity axiom.

### Properties of Ordinals I

### Properties of Ordinals II

For ordinals  $\alpha$ ,  $\beta$ , we write  $\alpha < \beta$  whenever  $\alpha \in \beta$ .

#### Natural numbers

Naturals are finite ordinals defined by the recursive rule

$$0 = \emptyset, \qquad n+1 = n \cup \{n\}$$

The first infinite ordinal is called  $\omega_0$  and it is

$$\omega_0 = \{0, 1, 2, 3, 4, \dots\}$$

The first uncountable ordinal is called  $\omega_1$ , and the first uncountable ordinal which is not equinumerous with  $\omega_1$  is called  $\omega_2$ .

In terms of cardinality, we call those sets  $\aleph_0$ ,  $\aleph_1$ , and  $\aleph_2$  respectively.

#### Ordinal Arithmetic

The **successor ordinal**  $\beta$  of an ordinal  $\alpha$  is given by

$$\beta = \alpha \cup \{\alpha\} = \alpha + 1$$

An ordinal  $\beta$  is called a **limit ordinal** if  $\beta \neq 0$  and  $\beta$  is not a successor.

## Example

#### Transfinite Induction

**Theorem.** Let  $\Psi(x)$  be a property defined for all ordinals  $\alpha$ . If we have that, for every ordinal  $\alpha$ ,

 $\Psi(\beta)$  is true for all  $\beta < \alpha \Rightarrow \Psi(\alpha)$  is true

Then Ψ holds for all ordinals.

### How to Induct Transfinitely

To prove a property  $\Psi$  on ordinals:

1. Base case: Prove that  $\Psi(0)$  holds.

- 2. Successor case: Prove that for any successor ordinal  $\alpha + 1$ ,  $\Psi(\alpha) \Rightarrow \Psi(\alpha + 1)$
- 3. Limit case: Prove that for any limit ordinal  $\beta$ ,  $\Psi(\beta)$  follows from  $\Psi(\alpha)$  for all  $\alpha < \beta$ .