## Assignment 6 Computer Science Fall 2017 B461

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## Answers

1. (a)

 $\pi_{sid,bookno}(\sigma_{Buys,bookno\neq Buys2,bookno\wedge Student.name='Eric'\wedge Buys,bookno\neq' 2010'}(Student\times Buys\times Buys2))$ 

(b) i. Original Query

ii. Convert from using SOME to an exists statement

iii. Push down student s relation into subquery as s1.

iv. The inner query can be translated as:

```
\pi_{b1.price}(\sigma_{b1.price} > 50 \land s1.sid = t.sid \land t.bookno = b1.bookno}(T \times B1 \times S1))
```

v. Finally, we perform the semi-join with the inner query and the student and book relation:

```
\pi_{b.bookno,b.title}(S \times B) \ltimes \pi_{b1.price}(\sigma_{b1,price > 50 \land s1.sid = t.sid \land t.bookno = b1.bookno}(T \times B1 \times S1))
```

(c) i. Original Query

```
SELECT b.bookno
FROM book b
WHERE b.bookno IN
(SELECT b1.bookno FROM book b1 WHERE b1.price > 50)
UNION
(SELECT c.bookno FROM cites c);
```

ii. RA representation of inner query (simple translation)

$$\pi_{bookno}(\sigma_{price>50}(book)) \cup \pi_{bookno}(cites)$$

iii. Since the IN predicate is equivalent to saying that there exists one bookno for which a bookno in the inner query matches, we can use the semi join

$$\pi_{bookno}(book) \ltimes \pi_{bookno}(\sigma_{price > 50}(book)) \cup \pi_{bookno}(cites)$$

(d) i. Original Query

ii. First, push down book b relation into subquery as book b2.

iii. Now, we can properly translate the inner query as:

$$\pi_{b1.bookno}(\sigma_{b1.price>b2.price}(B1 \times B2))$$

iv. To preserve the semantics of the original outer query, we need to perform an anti-semi join:

$$\pi_{b.bookno}(\sigma_{b.price})=80(B \ltimes \pi_{b1.bookno}(\sigma_{b1.price}) + 2.price(B1 \times B2))))$$

(e) i. Original query:

ii. First, push down parameterized values in the inner-most query. We will push down the student s relation as s1.

```
SELECT s.sid
FROM Student s
WHERE EXISTS(SELECT 1
FROM Book b
WHERE b.price > 50 AND
b.bookno IN (SELECT t.bookno
FROM Buys t, Student s1
WHERE s1.sid = t.sid AND
s1.sname = 'Eric'))
```

iii. Next, convert the first inner subquery IN expression into an EXISTS statement. We can do this by pushing the book relation completely into the inner most subquery and adding its conditions

iv. We can translate the inner query as:

```
\pi_{t.bookno}(\sigma_{s1.sid=t.sid} \land b.price > 50 \land b.bookno=t.bookno \land s1.sname='Eric'(S1 \times B \times T))
```

v. Finally, we can translate our outer exists expression with a semijoin and project the student sid.

```
\pi_{s.sid}(S \ltimes \pi_{t.bookno}(\sigma_{s1.sid=t.sid \land b.price>50 \land b.bookno=t.bookno \land s1.sname='Eric'}(S1 \times B \times T)))
```

(f) i. Original query:

```
SELECT s1.sid, s2.sid
FROM student s1, student s2
WHERE s1.sid <> s2.sid AND
NOT EXISTS(SELECT 1
FROM Buys t1
WHERE t1.sid = s1.sid AND
t1.bookno NOT IN (SELECT t2.bookno
FROM Buys t2
WHERE t2.sid = s2.sid));
```

ii. For the deepest query, we need to translate the NOT IN statement to NOT EXISTS.

iii. We now need to push down our relations recursively from the top and middle queries to the deepest query.

iv. Now that we've pushed down the upper query's relations into deeper queries, we can translate the deepest query as:

```
\varepsilon = \pi_{t2.sid,t1.bookno,t2.bookno,s1.sid}(\sigma_{t2.sid=s2.sid \land t1.bookno \neq t2.bookno}(T1 \times T2 \times S1 \times S2))
```

v. Our middle query then can be translated as:

$$\tau = \pi_{s1.sid,s2.sid,t.bookno}(\sigma_{t.sid=s1.sid}(T \times S1 \times S2)\overline{\ltimes}\varepsilon)$$

vi. Finally, we can semijoin the top level query in an expression like so:

$$\pi_{s1.sid.s2.sid}(\sigma_{s1.sid\neq s2.sid}(S1 \times S2) \overline{\ltimes} \tau)$$

2.