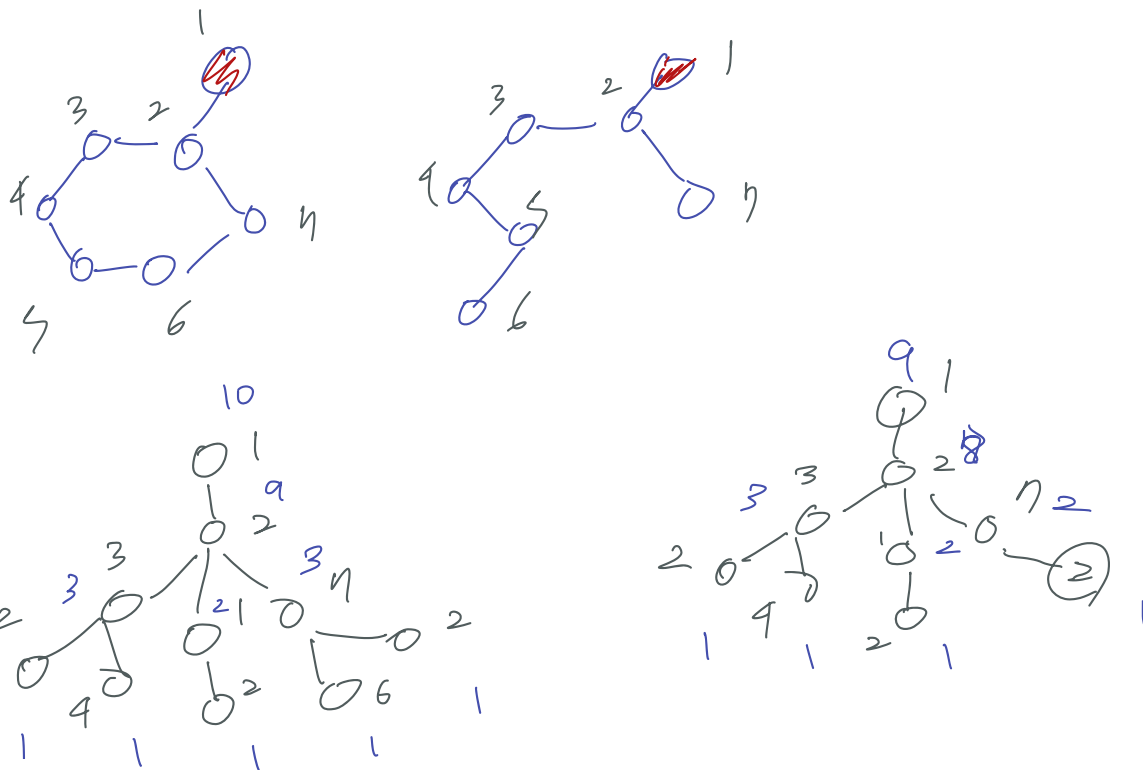


1-a)



⇒ 2 layer 이면 node 1에 attached 노드 9개 존재 embedding 됨

1-b)

$$m(h_v^k) = \begin{cases} 1, & \text{if } h_v^k \neq 1 \\ 0, & \text{otherwise} \end{cases}$$

$$h_{N(v)}^{k+1} = \max_{n \in N(v)} (m(h_n^k), 0)$$

$$h_v^{k+1} = \max (m(h_v^k), h_{N(v)}^{k+1})$$

2-a)

$$i) \quad -(p^2 + p^2 - 2p + 1) = -2p^2 + 2p = \text{GINI}$$

$$\frac{d^2 \text{GINI}}{dp^2} = \frac{d}{dp} (-4p + 2) = -4 < 0 \Rightarrow \text{concave/}$$

$$ii) \text{ Entropy} \quad p \log_2 \left(\frac{1}{p} \right) + (1-p) \log_2 \left(\frac{1}{1-p} \right) \\ = -p \log_2 p - (1-p) \log_2 (1-p)$$

$$\begin{aligned} \frac{d^2 \text{Entropy}}{dp^2} &= \frac{d}{dp} \left(\frac{d \text{Entropy}}{dp} \right) = \frac{d}{dp} \left(-\log_2 p - \frac{1}{\ln 2} + \log_2 (1-p) - (1-p) \cdot \frac{1}{\ln 2} \cdot \frac{-1}{1-p} \right) \\ &= \frac{d}{dp} \left(-\log_2 p - \frac{1}{\ln 2} + \log_2 (1-p) + \frac{1}{\ln 2} \right) \\ &= \frac{d}{dp} \left(\log_2 (1-p) - \log_2 p \right) \\ &= \frac{1}{\ln 2} \cdot \frac{-1}{1-p} - \frac{1}{\ln 2} \cdot \frac{1}{p} < 0 \end{aligned}$$

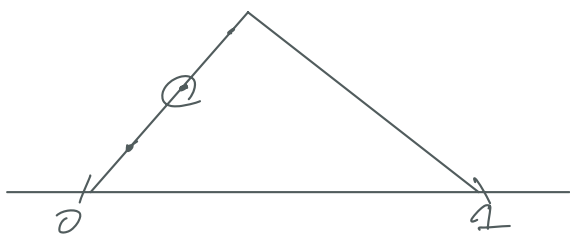
iii)

$$p < 0.5$$

$$1 - \max(p, 1-p) = p$$

$$p \geq 0.5$$

$$1 - \max(p, 1-p) = 1-p$$



$$x = 0.1 \quad y = 0.5, \quad z = 0.3$$

$$z = 0.3 \quad f(z) = 0.3$$

$$x = 0.1 \quad f(x) = 0.1$$

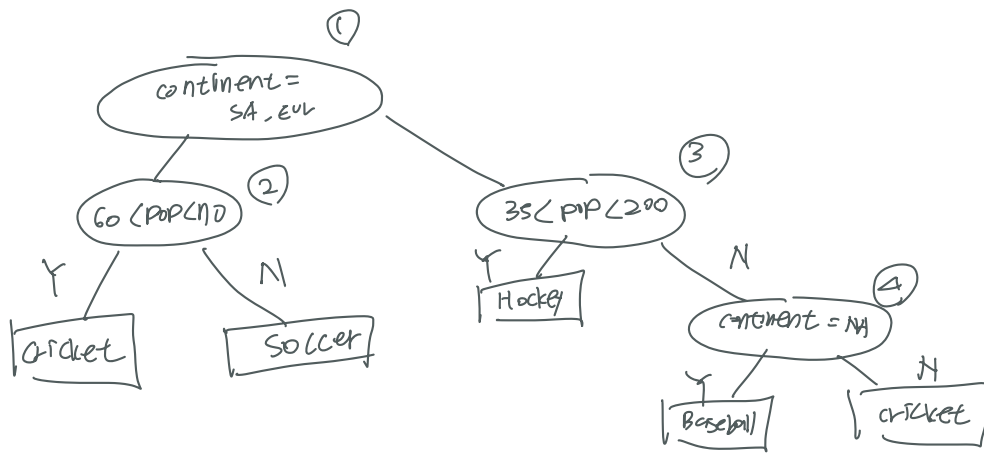
$$y = 0.5 \quad f(y) = 0.5$$

$$f(z) = 0.3$$

$$\frac{y-z}{y-x} f(x) + \frac{z-x}{y-x} f(y) = \frac{0.2}{0.4} \cdot 0.1 + \frac{0.2}{0.4} \cdot 0.5 = \frac{1}{2} \cdot 0.6 = 0.3$$

Same

2-b)



① - GINI Index

$$P_{\text{Soccer}} = \frac{5}{12}, P_{\text{Cricket}} = \frac{3}{12}, P_{\text{Hockey}} = \frac{2}{12}, P_{\text{Baseball}} = \frac{2}{12}$$

$$1 - \left(\frac{25}{144} + \frac{9}{144} + \frac{4}{144} + \frac{4}{144} \right) = 1 - \frac{42}{144} = \frac{102}{144}$$

① - accuracy

$$1 - \max \left(\frac{5}{12}, \frac{3}{12}, \frac{2}{12}, \frac{2}{12} \right) = \frac{7}{12}$$

② - GINI Index

$$P_{\text{Soccer}} = \frac{5}{6}, P_{\text{Cricket}} = \frac{1}{6}$$

$$1 - \left(\frac{25}{36} + \frac{1}{36} \right) = 1 - \frac{26}{36} = \frac{10}{36} = \frac{5}{18}$$

② - accuracy

$$1 - \max \left(\frac{25}{36}, \frac{1}{36} \right) = \frac{11}{36}$$

③ - GINI Index

$$P_{\text{Hockey}} = \frac{2}{6}, P_{\text{Baseball}} = \frac{2}{6}, P_{\text{Cricket}} = \frac{2}{6}$$

$$1 - \left(\frac{4}{36} + \frac{4}{36} + \frac{4}{36} \right) = 1 - \frac{12}{36} = \frac{24}{36} = \frac{2}{3}$$

③ - accuracy

$$1 - \frac{2}{6} = \frac{2}{3}$$

④ - GINI Index

$$P_{\text{baseball}} = \frac{2}{4}$$

$$P_{\text{cricket}} = \frac{2}{4}$$

$$1 - \left(\frac{1}{4} + \frac{1}{4} \right) = \frac{1}{2}$$

④ - accuracy

$$1 - \frac{1}{2} = \frac{1}{2}$$

3-a)

i) $(1 - e^{-km/n})^k$

$k=3, n=8 \text{ billion}, m=1 \text{ billion}$

$(1 - e^{-3 \cdot \frac{1}{8}})^3 = 0.030$

$k=4, n=8 \text{ billion}, m=1 \text{ billion}$

$(1 - e^{-4 \cdot \frac{1}{8}})^4 = 0.024$

ii) $1 - \frac{k}{n}$

$(1 - \frac{k}{n})^m$

$1 - (1 - \frac{k}{n})^m$

$= 1 - (1 - \frac{k}{n})^{m \cdot \frac{n}{k} \cdot \frac{k}{n}}$

$= 1 - (1 - \frac{1}{x})^{m \cdot x \cdot \frac{k}{n}}$

$= 1 - e^{-m \cdot \frac{k}{n}}$ for each array

$\frac{k}{n} \approx \frac{1}{x}$
 $x = \frac{n}{k}$

for every k array

$(1 - e^{-m \cdot \frac{k}{n}})^k$

$\frac{d}{dx} (f(x))^x$

iii) optimal value of $k = \frac{n}{m} \ln 2$

$(1 - e^{-m \cdot \frac{k}{n}})^k \cdot \ln$

3-b)

i)

a) $h(x) = 2x+1 \mod 32$

elem	hash	bit	tail length
3	7	00111	0
1	3	00011	0
4	9	01001	0
1	3	00011	0
5	11	01011	0
9	19	10011	0
2	5	00101	0
6	13	01101	0
5	11	01011	0

\Rightarrow estimated distance elem
 $\Rightarrow 2^0 = 1$

b) $h(x) = 3x + 17 \pmod{32}$

elem	hash	bit	tail length
3	16	10000	4
1	10	01010	1
4	19	10011	0
1	10	01010	1
5	22	10110	1
9	2	00010	1
2	13	01101	0
6	25	11001	0
5	22	10110	1

\Rightarrow number of distinct elements
 $\Rightarrow 2^4 = 16$

c) $h(x) = 4x \pmod{32}$

elem	hash	bit	tail length
3	12	01100	2
1	4	00100	2
4	16	10000	4
1	4	00100	2
5	20	10100	2
9	4	00100	2
2	8	01000	3
6	24	11000	3
5	20	10100	2

\Rightarrow number of distinct elements
 $\Rightarrow 2^4 = 16$

i)

hash function $a \times b \pmod{2^k}$ 일때, 이 식은 k -bit length로 사용된다.
 전체 element를 생성할 수 있을 정도의 k bit가 있어야 한다. 31 예제를 보았을 때 $k=5$ 는

충분하다.

이러한 hash function의 over or under specification에 따라 fluctuation이 심해진다.
 일정한 적절한 hash function을 사용해야 한다.

또한 각 hash function의 범위를 평균을 내서 distinct element