

Where did my RAM go? Using algebraic cryptanalysis in practice

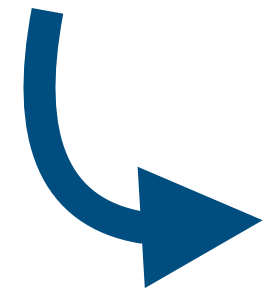
Lars Ran

Monika Trimoska

Summer school on RWC and privacy
July 1, Dubrovnik, Croatia

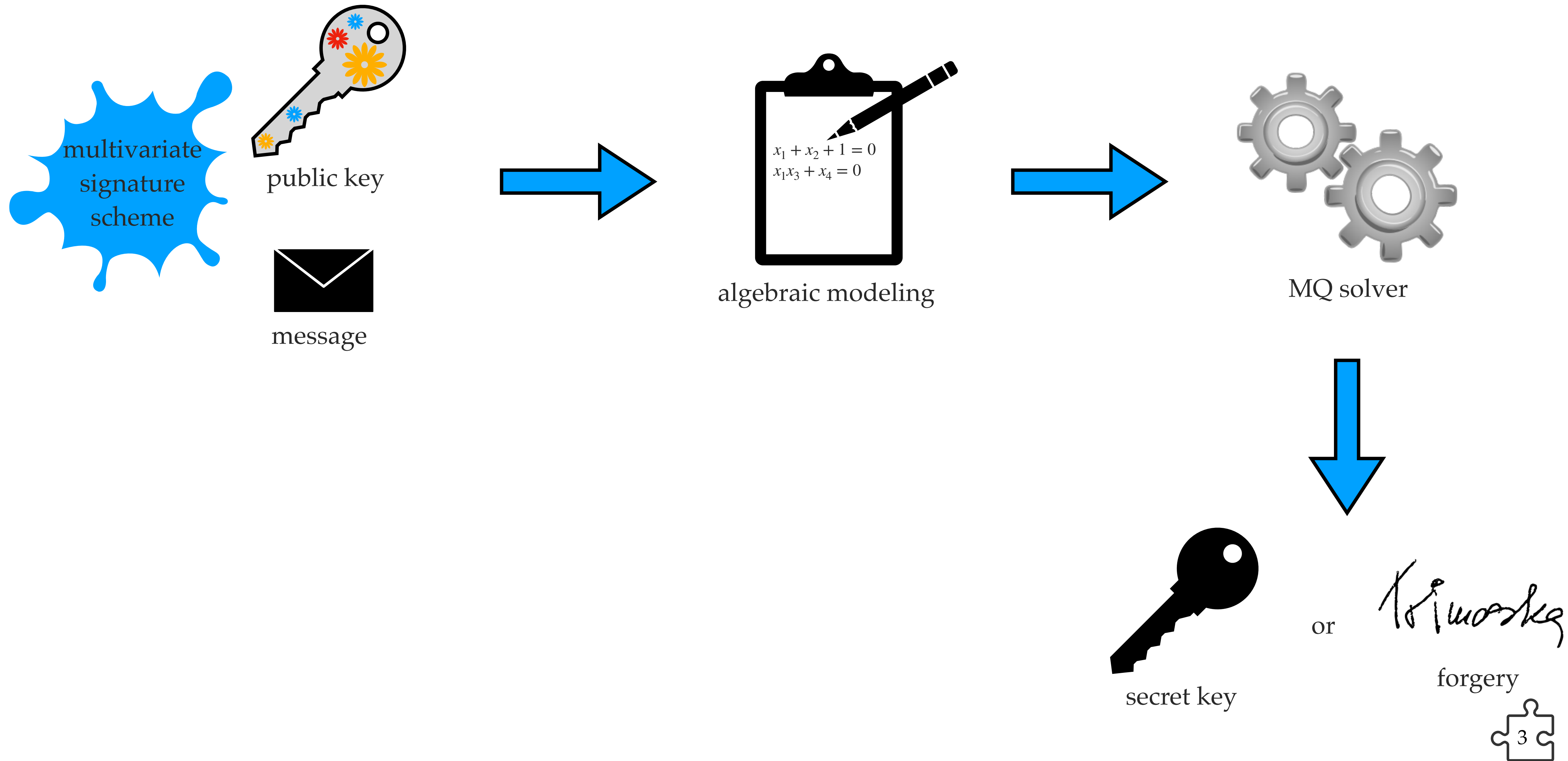
TU/e

Algebraic cryptanalysis

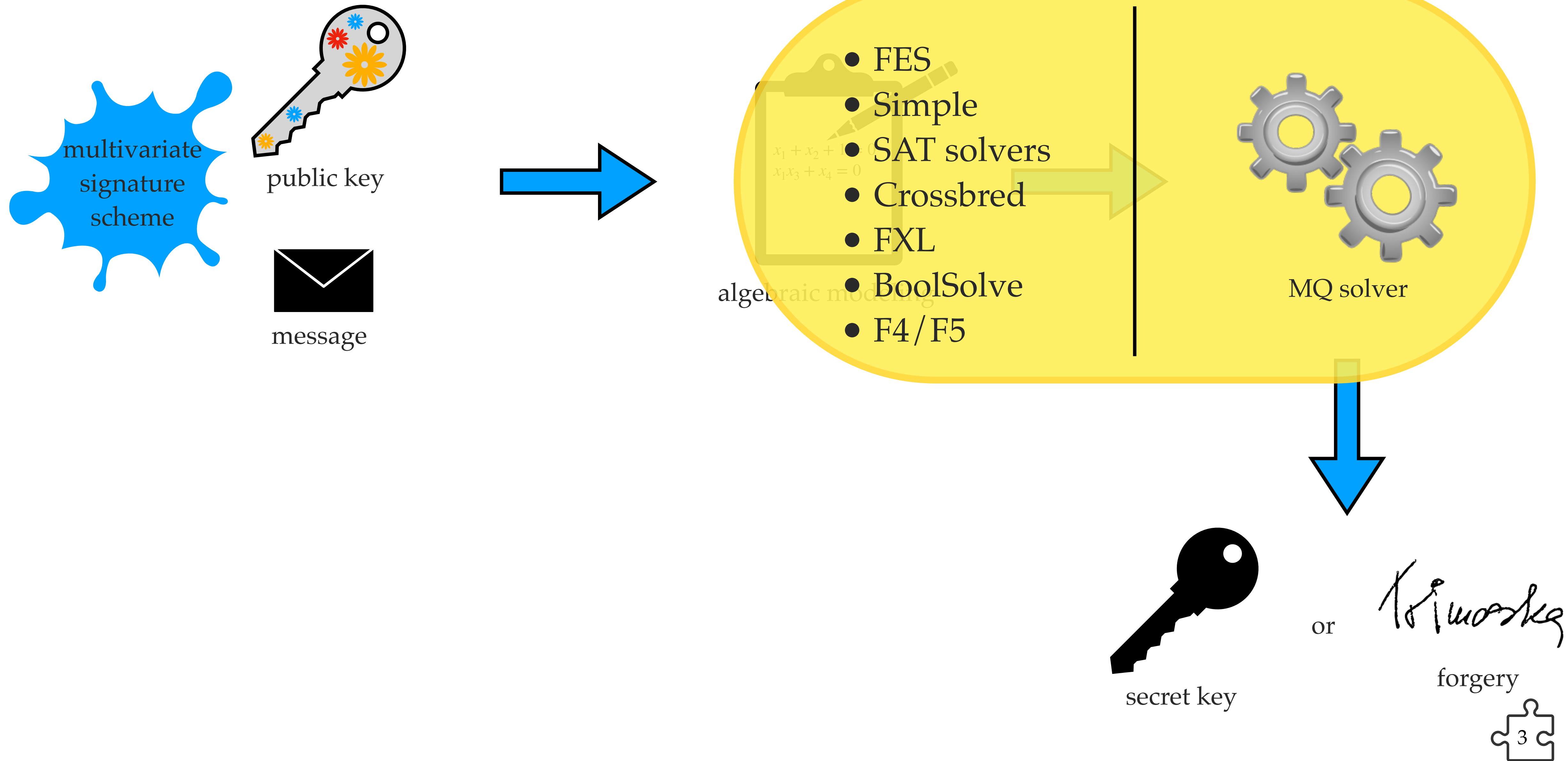


A type of cryptanalytic methods where the problem of finding the secret key (or any attack goal) is **reduced** to the problem of finding a solution to a **nonlinear multivariate polynomial system of equations**.

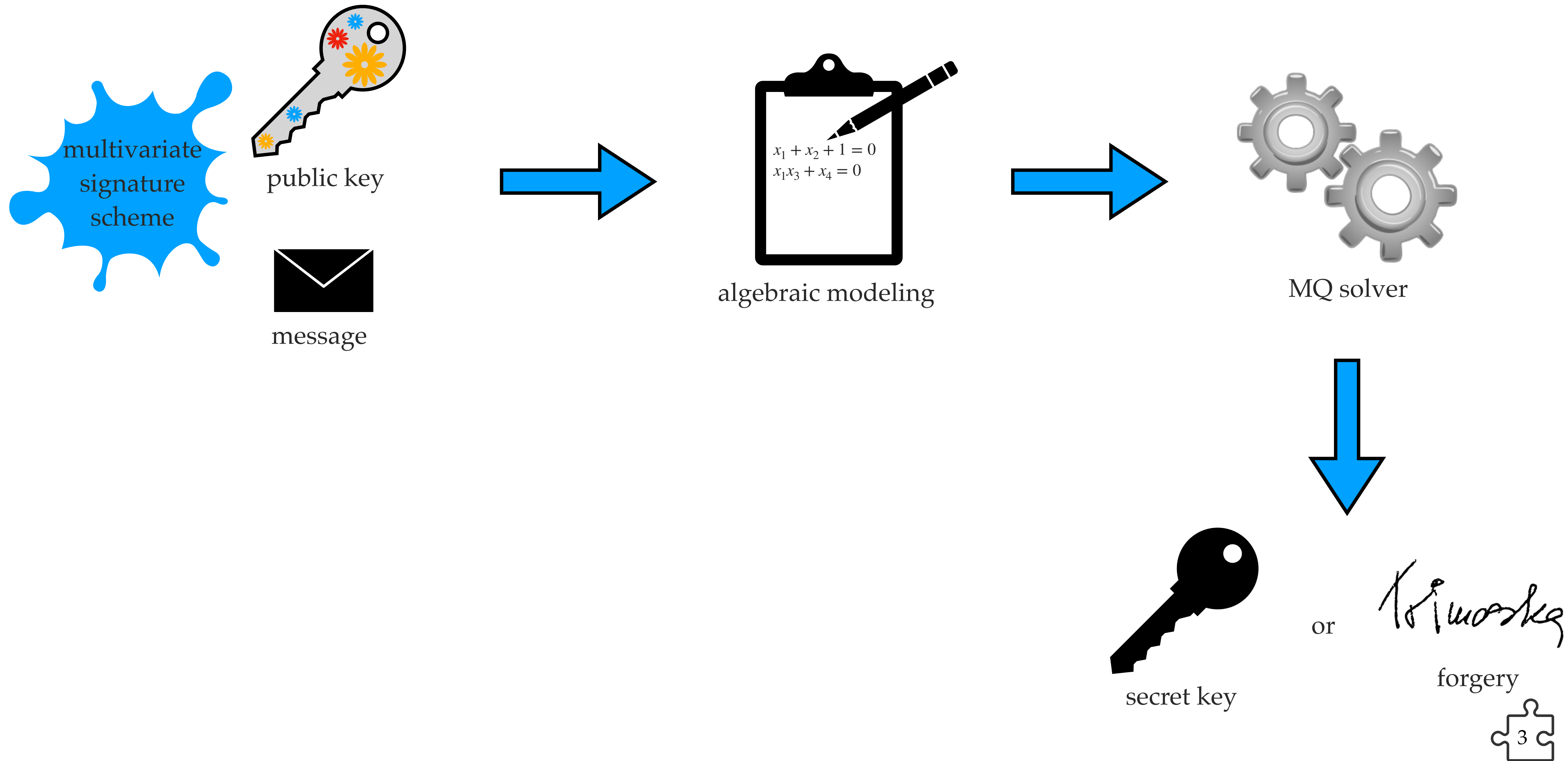
Algebraic cryptanalysis



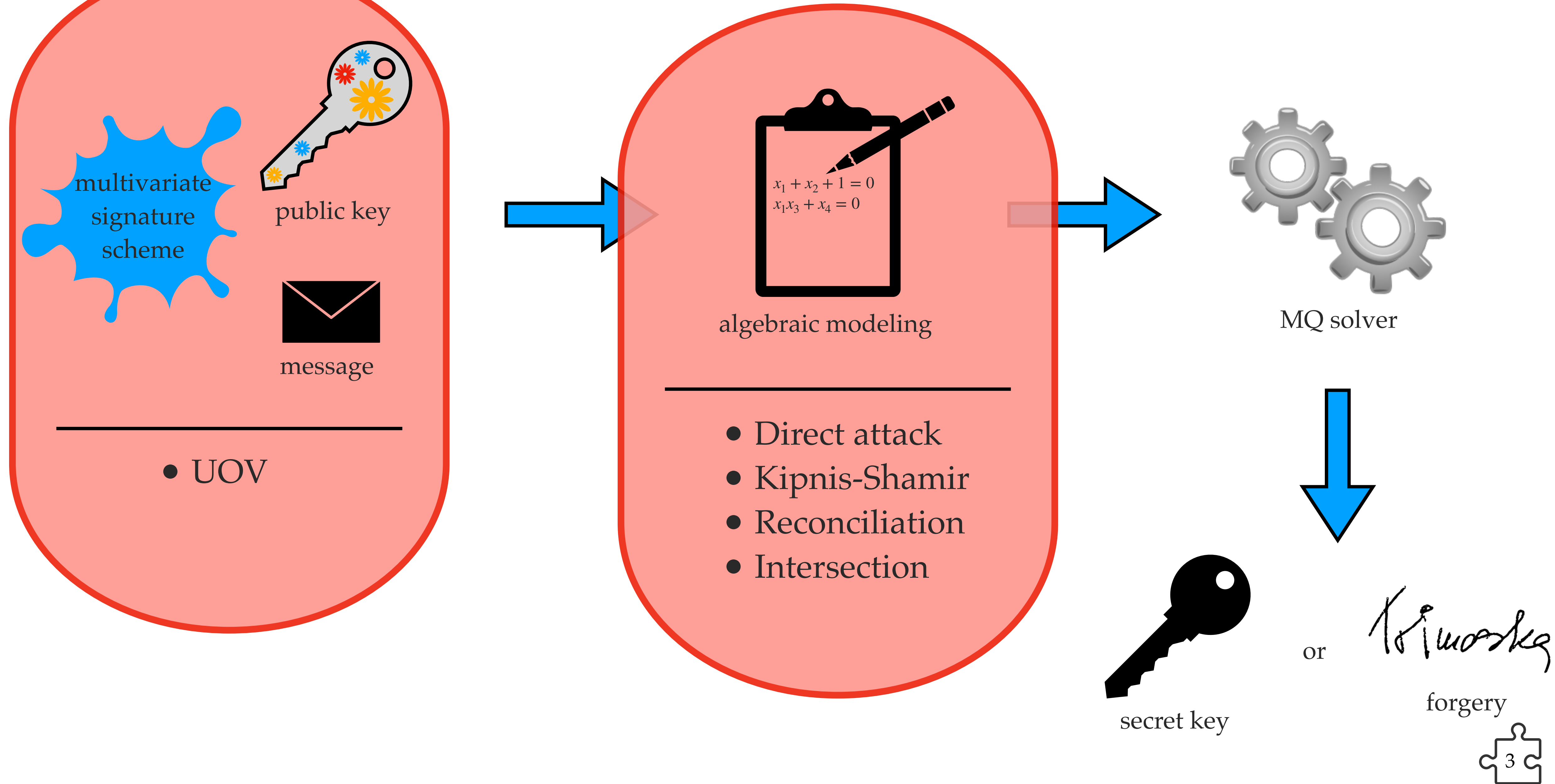
Algebraic cryptanalysis



Algebraic cryptanalysis



Algebraic cryptanalysis



The MQ problem (recall)

The MQ problem

Given m multivariate quadratic polynomials f_1, \dots, f_m of n variables over a finite field \mathbb{F}_q , find a tuple $\mathbf{x} = (x_1, \dots, x_n)$ in \mathbb{F}_q^n , such that $f_1(\mathbf{x}) = \dots = f_m(\mathbf{x}) = 0$.

Example.

$$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$$

$$f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$$

$$f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$$

$$f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$$

$$f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$$

$$f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$$

Algebraic cryptanalysis : modelisation example

Example.

Given matrices $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2 \in \mathcal{M}_{n,n}(\mathbb{F}_q)$ (the space of matrices over \mathbb{F}_q of size $n \times n$), find $\mathbf{A}, \mathbf{B} \in \text{GL}_n(\mathbb{F}_q)$ (the space of invertible matrices over \mathbb{F}_q of size $n \times n$), such that

$$\mathbf{D}_1 = \mathbf{A} \mathbf{C}_1 \mathbf{B}$$

$$\mathbf{D}_2 = \mathbf{A} \mathbf{C}_2 \mathbf{B}$$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A}\mathbf{C}_1\mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A}\mathbf{C}_2\mathbf{B}$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

\mathbf{A}

$a_{1,1}$	$a_{1,2}$	$a_{1,3}$	$a_{1,4}$	$a_{1,5}$
$a_{2,1}$	$a_{2,2}$	$a_{2,3}$	$a_{2,4}$	$a_{2,5}$
$a_{3,1}$	$a_{3,2}$	$a_{3,3}$	$a_{3,4}$	$a_{3,5}$
$a_{4,1}$	$a_{4,2}$	$a_{4,3}$	$a_{4,4}$	$a_{4,5}$
$a_{5,1}$	$a_{5,2}$	$a_{5,3}$	$a_{5,4}$	$a_{5,5}$

\mathbf{C}_1

$c_{1,1}$	$c_{1,2}$	$c_{1,3}$	$c_{1,4}$	$c_{1,5}$
$c_{2,1}$	$c_{2,2}$	$c_{2,3}$	$c_{2,4}$	$c_{2,5}$
$c_{3,1}$	$c_{3,2}$	$c_{3,3}$	$c_{3,4}$	$c_{3,5}$
$c_{4,1}$	$c_{4,2}$	$c_{4,3}$	$c_{4,4}$	$c_{4,5}$
$c_{5,1}$	$c_{5,2}$	$c_{5,3}$	$c_{5,4}$	$c_{5,5}$

\mathbf{B}

$b_{1,1}$	$b_{1,2}$	$b_{1,3}$	$b_{1,4}$	$b_{1,5}$
$b_{2,1}$	$b_{2,2}$	$b_{2,3}$	$b_{2,4}$	$b_{2,5}$
$b_{3,1}$	$b_{3,2}$	$b_{3,3}$	$b_{3,4}$	$b_{3,5}$
$b_{4,1}$	$b_{4,2}$	$b_{4,3}$	$b_{4,4}$	$b_{4,5}$
$b_{5,1}$	$b_{5,2}$	$b_{5,3}$	$b_{5,4}$	$b_{5,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A} \mathbf{C}_1 \mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A} \mathbf{C}_2 \mathbf{B}$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

$\mathbf{A} \mathbf{C}_1 \mathbf{B}$

$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A} \mathbf{C}_1 \mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A} \mathbf{C}_2 \mathbf{B}$

$$d_{1,1} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1} = 0,$$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

$\mathbf{A} \mathbf{C}_1 \mathbf{B}$

$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A} \mathbf{C}_1 \mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A} \mathbf{C}_2 \mathbf{B}$

$$d_{1,1} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1} = 0,$$

$$d_{2,1} - \sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1} = 0,$$

$$\mathbf{A} \mathbf{C}_1 \mathbf{B}$$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A}\mathbf{C}_1\mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A}\mathbf{C}_2\mathbf{B}$

$$d_{1,1} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1} = 0, \quad d_{1,2} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2} = 0, \dots$$

$$d_{2,1} - \sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1} = 0,$$

$\mathbf{A}\mathbf{C}_1\mathbf{B}$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{D}_1 = \mathbf{A} \mathbf{C}_1 \mathbf{B}$
 $\mathbf{D}_2 = \mathbf{A} \mathbf{C}_2 \mathbf{B}$

$$d_{1,1} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1} = 0, \quad d_{1,2} - \sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2} = 0, \dots$$

$$d_{2,1} - \sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1} = 0, \quad d_{l,p} - \sum_{i=1}^n \sum_{j=1}^n a_{l,j} c_{j,i} b_{i,p} = 0$$

$\mathbf{A} \mathbf{C}_1 \mathbf{B}$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

=

$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{1,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{2,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{3,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{4,j} c_{j,i} b_{i,5}$
$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,1}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,2}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,3}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,4}$	$\sum_{i=1}^n \sum_{j=1}^n a_{5,j} c_{j,i} b_{i,5}$

Algebraic cryptanalysis : modelisation example



A motivating example: a better idea for modelisation.

Given matrices $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2 \in \mathcal{M}_{n,n}(\mathbb{F}_q)$ (the space of matrices over \mathbb{F}_q of size $n \times n$), find $\mathbf{A}, \mathbf{B} \in \text{GL}_n(\mathbb{F}_q)$ (the space of invertible matrices over \mathbb{F}_q of size $n \times n$), such that

$$\begin{aligned}\mathbf{A}^{-1}\mathbf{D}_1 &= \mathbf{C}_1\mathbf{B} \\ \mathbf{A}^{-1}\mathbf{D}_2 &= \mathbf{C}_2\mathbf{B}\end{aligned}$$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{A}^{-1}\mathbf{D}_1 = \mathbf{C}_1\mathbf{B}$
 $\mathbf{A}^{-1}\mathbf{D}_2 = \mathbf{C}_2\mathbf{B}$

\mathbf{A}^{-1}

$a_{1,1}$	$a_{1,2}$	$a_{1,3}$	$a_{1,4}$	$a_{1,5}$
$a_{2,1}$	$a_{2,2}$	$a_{2,3}$	$a_{2,4}$	$a_{2,5}$
$a_{3,1}$	$a_{3,2}$	$a_{3,3}$	$a_{3,4}$	$a_{3,5}$
$a_{4,1}$	$a_{4,2}$	$a_{4,3}$	$a_{4,4}$	$a_{4,5}$
$a_{5,1}$	$a_{5,2}$	$a_{5,3}$	$a_{5,4}$	$a_{5,5}$

\mathbf{D}_1

$d_{1,1}$	$d_{1,2}$	$d_{1,3}$	$d_{1,4}$	$d_{1,5}$
$d_{2,1}$	$d_{2,2}$	$d_{2,3}$	$d_{2,4}$	$d_{2,5}$
$d_{3,1}$	$d_{3,2}$	$d_{3,3}$	$d_{3,4}$	$d_{3,5}$
$d_{4,1}$	$d_{4,2}$	$d_{4,3}$	$d_{4,4}$	$d_{4,5}$
$d_{5,1}$	$d_{5,2}$	$d_{5,3}$	$d_{5,4}$	$d_{5,5}$

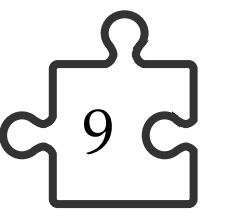
=

\mathbf{C}_1

$c_{1,1}$	$c_{1,2}$	$c_{1,3}$	$c_{1,4}$	$c_{1,5}$
$c_{2,1}$	$c_{2,2}$	$c_{2,3}$	$c_{2,4}$	$c_{2,5}$
$c_{3,1}$	$c_{3,2}$	$c_{3,3}$	$c_{3,4}$	$c_{3,5}$
$c_{4,1}$	$c_{4,2}$	$c_{4,3}$	$c_{4,4}$	$c_{4,5}$
$c_{5,1}$	$c_{5,2}$	$c_{5,3}$	$c_{5,4}$	$c_{5,5}$

\mathbf{B}

$b_{1,1}$	$b_{1,2}$	$b_{1,3}$	$b_{1,4}$	$b_{1,5}$
$b_{2,1}$	$b_{2,2}$	$b_{2,3}$	$b_{2,4}$	$b_{2,5}$
$b_{3,1}$	$b_{3,2}$	$b_{3,3}$	$b_{3,4}$	$b_{3,5}$
$b_{4,1}$	$b_{4,2}$	$b_{4,3}$	$b_{4,4}$	$b_{4,5}$
$b_{5,1}$	$b_{5,2}$	$b_{5,3}$	$b_{5,4}$	$b_{5,5}$



Algebraic cryptanalysis : modelisation example

Example. $\mathbf{A}^{-1}\mathbf{D}_1 = \mathbf{C}_1\mathbf{B}$
 $\mathbf{A}^{-1}\mathbf{D}_2 = \mathbf{C}_2\mathbf{B}$

$$\mathbf{A}^{-1}\mathbf{D}_1$$

$\sum_{i=1}^n a_{1,i}d_{i,1}$	$\sum_{i=1}^n a_{1,i}d_{i,2}$	$\sum_{i=1}^n a_{1,i}d_{i,3}$	$\sum_{i=1}^n a_{1,i}d_{i,4}$	$\sum_{i=1}^n a_{1,i}d_{i,5}$
$\sum_{i=1}^n a_{2,i}d_{i,1}$	$\sum_{i=1}^n a_{2,i}d_{i,2}$	$\sum_{i=1}^n a_{2,i}d_{i,3}$	$\sum_{i=1}^n a_{2,i}d_{i,4}$	$\sum_{i=1}^n a_{2,i}d_{i,5}$
$\sum_{i=1}^n a_{3,i}d_{i,1}$	$\sum_{i=1}^n a_{3,i}d_{i,2}$	$\sum_{i=1}^n a_{3,i}d_{i,3}$	$\sum_{i=1}^n a_{3,i}d_{i,4}$	$\sum_{i=1}^n a_{3,i}d_{i,5}$
$\sum_{i=1}^n a_{4,i}d_{i,1}$	$\sum_{i=1}^n a_{4,i}d_{i,2}$	$\sum_{i=1}^n a_{4,i}d_{i,3}$	$\sum_{i=1}^n a_{4,i}d_{i,4}$	$\sum_{i=1}^n a_{4,i}d_{i,5}$
$\sum_{i=1}^n a_{5,i}d_{i,1}$	$\sum_{i=1}^n a_{5,i}d_{i,2}$	$\sum_{i=1}^n a_{5,i}d_{i,3}$	$\sum_{i=1}^n a_{5,i}d_{i,4}$	$\sum_{i=1}^n a_{5,i}d_{i,5}$

=

$$\mathbf{C}_1\mathbf{B}$$

$\sum_{i=1}^n c_{1,i}b_{i,1}$	$\sum_{i=1}^n c_{1,i}b_{i,2}$	$\sum_{i=1}^n c_{1,i}b_{i,3}$	$\sum_{i=1}^n c_{1,i}b_{i,4}$	$\sum_{i=1}^n c_{1,i}b_{i,5}$
$\sum_{i=1}^n c_{2,i}b_{i,1}$	$\sum_{i=1}^n c_{2,i}b_{i,2}$	$\sum_{i=1}^n c_{2,i}b_{i,3}$	$\sum_{i=1}^n c_{2,i}b_{i,4}$	$\sum_{i=1}^n c_{2,i}b_{i,5}$
$\sum_{i=1}^n c_{3,i}b_{i,1}$	$\sum_{i=1}^n c_{3,i}b_{i,2}$	$\sum_{i=1}^n c_{3,i}b_{i,3}$	$\sum_{i=1}^n c_{3,i}b_{i,4}$	$\sum_{i=1}^n c_{3,i}b_{i,5}$
$\sum_{i=1}^n c_{4,i}b_{i,1}$	$\sum_{i=1}^n c_{4,i}b_{i,2}$	$\sum_{i=1}^n c_{4,i}b_{i,3}$	$\sum_{i=1}^n c_{4,i}b_{i,4}$	$\sum_{i=1}^n c_{4,i}b_{i,5}$
$\sum_{i=1}^n c_{5,i}b_{i,1}$	$\sum_{i=1}^n c_{5,i}b_{i,2}$	$\sum_{i=1}^n c_{5,i}b_{i,3}$	$\sum_{i=1}^n c_{5,i}b_{i,4}$	$\sum_{i=1}^n c_{5,i}b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{A}^{-1}\mathbf{D}_1 = \mathbf{C}_1\mathbf{B}$
 $\mathbf{A}^{-1}\mathbf{D}_2 = \mathbf{C}_2\mathbf{B}$

$$\sum_{i=1}^n a_{1,i}d_{i,1} - \sum_{i=1}^n c_{1,i}b_{i,1} = 0,$$

$\mathbf{A}^{-1}\mathbf{D}_1$

$\sum_{i=1}^n a_{1,i}d_{i,1}$	$\sum_{i=1}^n a_{1,i}d_{i,2}$	$\sum_{i=1}^n a_{1,i}d_{i,3}$	$\sum_{i=1}^n a_{1,i}d_{i,4}$	$\sum_{i=1}^n a_{1,i}d_{i,5}$
$\sum_{i=1}^n a_{2,i}d_{i,1}$	$\sum_{i=1}^n a_{2,i}d_{i,2}$	$\sum_{i=1}^n a_{2,i}d_{i,3}$	$\sum_{i=1}^n a_{2,i}d_{i,4}$	$\sum_{i=1}^n a_{2,i}d_{i,5}$
$\sum_{i=1}^n a_{3,i}d_{i,1}$	$\sum_{i=1}^n a_{3,i}d_{i,2}$	$\sum_{i=1}^n a_{3,i}d_{i,3}$	$\sum_{i=1}^n a_{3,i}d_{i,4}$	$\sum_{i=1}^n a_{3,i}d_{i,5}$
$\sum_{i=1}^n a_{4,i}d_{i,1}$	$\sum_{i=1}^n a_{4,i}d_{i,2}$	$\sum_{i=1}^n a_{4,i}d_{i,3}$	$\sum_{i=1}^n a_{4,i}d_{i,4}$	$\sum_{i=1}^n a_{4,i}d_{i,5}$
$\sum_{i=1}^n a_{5,i}d_{i,1}$	$\sum_{i=1}^n a_{5,i}d_{i,2}$	$\sum_{i=1}^n a_{5,i}d_{i,3}$	$\sum_{i=1}^n a_{5,i}d_{i,4}$	$\sum_{i=1}^n a_{5,i}d_{i,5}$

=

$\mathbf{C}_1\mathbf{B}$

$\sum_{i=1}^n c_{1,i}b_{i,1}$	$\sum_{i=1}^n c_{1,i}b_{i,2}$	$\sum_{i=1}^n c_{1,i}b_{i,3}$	$\sum_{i=1}^n c_{1,i}b_{i,4}$	$\sum_{i=1}^n c_{1,i}b_{i,5}$
$\sum_{i=1}^n c_{2,i}b_{i,1}$	$\sum_{i=1}^n c_{2,i}b_{i,2}$	$\sum_{i=1}^n c_{2,i}b_{i,3}$	$\sum_{i=1}^n c_{2,i}b_{i,4}$	$\sum_{i=1}^n c_{2,i}b_{i,5}$
$\sum_{i=1}^n c_{3,i}b_{i,1}$	$\sum_{i=1}^n c_{3,i}b_{i,2}$	$\sum_{i=1}^n c_{3,i}b_{i,3}$	$\sum_{i=1}^n c_{3,i}b_{i,4}$	$\sum_{i=1}^n c_{3,i}b_{i,5}$
$\sum_{i=1}^n c_{4,i}b_{i,1}$	$\sum_{i=1}^n c_{4,i}b_{i,2}$	$\sum_{i=1}^n c_{4,i}b_{i,3}$	$\sum_{i=1}^n c_{4,i}b_{i,4}$	$\sum_{i=1}^n c_{4,i}b_{i,5}$
$\sum_{i=1}^n c_{5,i}b_{i,1}$	$\sum_{i=1}^n c_{5,i}b_{i,2}$	$\sum_{i=1}^n c_{5,i}b_{i,3}$	$\sum_{i=1}^n c_{5,i}b_{i,4}$	$\sum_{i=1}^n c_{5,i}b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{A}^{-1}\mathbf{D}_1 = \mathbf{C}_1\mathbf{B}$
 $\mathbf{A}^{-1}\mathbf{D}_2 = \mathbf{C}_2\mathbf{B}$

$$\mathbf{A}^{-1}\mathbf{D}_1$$

$\sum_{i=1}^n a_{1,i}d_{i,1}$	$\sum_{i=1}^n a_{1,i}d_{i,2}$	$\sum_{i=1}^n a_{1,i}d_{i,3}$	$\sum_{i=1}^n a_{1,i}d_{i,4}$	$\sum_{i=1}^n a_{1,i}d_{i,5}$
$\sum_{i=1}^n a_{2,i}d_{i,1}$	$\sum_{i=1}^n a_{2,i}d_{i,2}$	$\sum_{i=1}^n a_{2,i}d_{i,3}$	$\sum_{i=1}^n a_{2,i}d_{i,4}$	$\sum_{i=1}^n a_{2,i}d_{i,5}$
$\sum_{i=1}^n a_{3,i}d_{i,1}$	$\sum_{i=1}^n a_{3,i}d_{i,2}$	$\sum_{i=1}^n a_{3,i}d_{i,3}$	$\sum_{i=1}^n a_{3,i}d_{i,4}$	$\sum_{i=1}^n a_{3,i}d_{i,5}$
$\sum_{i=1}^n a_{4,i}d_{i,1}$	$\sum_{i=1}^n a_{4,i}d_{i,2}$	$\sum_{i=1}^n a_{4,i}d_{i,3}$	$\sum_{i=1}^n a_{4,i}d_{i,4}$	$\sum_{i=1}^n a_{4,i}d_{i,5}$
$\sum_{i=1}^n a_{5,i}d_{i,1}$	$\sum_{i=1}^n a_{5,i}d_{i,2}$	$\sum_{i=1}^n a_{5,i}d_{i,3}$	$\sum_{i=1}^n a_{5,i}d_{i,4}$	$\sum_{i=1}^n a_{5,i}d_{i,5}$

=

$$\mathbf{C}_1\mathbf{B}$$

$\sum_{i=1}^n c_{1,i}b_{i,1}$	$\sum_{i=1}^n c_{1,i}b_{i,2}$	$\sum_{i=1}^n c_{1,i}b_{i,3}$	$\sum_{i=1}^n c_{1,i}b_{i,4}$	$\sum_{i=1}^n c_{1,i}b_{i,5}$
$\sum_{i=1}^n c_{2,i}b_{i,1}$	$\sum_{i=1}^n c_{2,i}b_{i,2}$	$\sum_{i=1}^n c_{2,i}b_{i,3}$	$\sum_{i=1}^n c_{2,i}b_{i,4}$	$\sum_{i=1}^n c_{2,i}b_{i,5}$
$\sum_{i=1}^n c_{3,i}b_{i,1}$	$\sum_{i=1}^n c_{3,i}b_{i,2}$	$\sum_{i=1}^n c_{3,i}b_{i,3}$	$\sum_{i=1}^n c_{3,i}b_{i,4}$	$\sum_{i=1}^n c_{3,i}b_{i,5}$
$\sum_{i=1}^n c_{4,i}b_{i,1}$	$\sum_{i=1}^n c_{4,i}b_{i,2}$	$\sum_{i=1}^n c_{4,i}b_{i,3}$	$\sum_{i=1}^n c_{4,i}b_{i,4}$	$\sum_{i=1}^n c_{4,i}b_{i,5}$
$\sum_{i=1}^n c_{5,i}b_{i,1}$	$\sum_{i=1}^n c_{5,i}b_{i,2}$	$\sum_{i=1}^n c_{5,i}b_{i,3}$	$\sum_{i=1}^n c_{5,i}b_{i,4}$	$\sum_{i=1}^n c_{5,i}b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example.

$$\begin{aligned} \mathbf{A}^{-1}\mathbf{D}_1 &= \mathbf{C}_1\mathbf{B} \\ \mathbf{A}^{-1}\mathbf{D}_2 &= \mathbf{C}_2\mathbf{B} \end{aligned}$$

$$\mathbf{A}^{-1}\mathbf{D}_1$$

$\sum_{i=1}^n a_{1,i}d_{i,1}$	$\sum_{i=1}^n a_{1,i}d_{i,2}$	$\sum_{i=1}^n a_{1,i}d_{i,3}$	$\sum_{i=1}^n a_{1,i}d_{i,4}$	$\sum_{i=1}^n a_{1,i}d_{i,5}$
$\sum_{i=1}^n a_{2,i}d_{i,1}$	$\sum_{i=1}^n a_{2,i}d_{i,2}$	$\sum_{i=1}^n a_{2,i}d_{i,3}$	$\sum_{i=1}^n a_{2,i}d_{i,4}$	$\sum_{i=1}^n a_{2,i}d_{i,5}$
$\sum_{i=1}^n a_{3,i}d_{i,1}$	$\sum_{i=1}^n a_{3,i}d_{i,2}$	$\sum_{i=1}^n a_{3,i}d_{i,3}$	$\sum_{i=1}^n a_{3,i}d_{i,4}$	$\sum_{i=1}^n a_{3,i}d_{i,5}$
$\sum_{i=1}^n a_{4,i}d_{i,1}$	$\sum_{i=1}^n a_{4,i}d_{i,2}$	$\sum_{i=1}^n a_{4,i}d_{i,3}$	$\sum_{i=1}^n a_{4,i}d_{i,4}$	$\sum_{i=1}^n a_{4,i}d_{i,5}$
$\sum_{i=1}^n a_{5,i}d_{i,1}$	$\sum_{i=1}^n a_{5,i}d_{i,2}$	$\sum_{i=1}^n a_{5,i}d_{i,3}$	$\sum_{i=1}^n a_{5,i}d_{i,4}$	$\sum_{i=1}^n a_{5,i}d_{i,5}$

=

$$\mathbf{C}_1\mathbf{B}$$

$$\sum_{i=1}^n a_{1,i}d_{i,1}-\sum_{i=1}^n c_{1,i}b_{i,1}=0, \quad \sum_{i=1}^n a_{1,i}d_{i,2}-\sum_{i=1}^n c_{1,i}b_{i,2}=0,\dots$$

$$\sum_{i=1}^n a_{2,i}d_{i,1}-\sum_{i=1}^n c_{2,i}b_{i,1}=0,$$

C_1
 B

$\sum_{i=1}^n c_{1,i}b_{i,1}$	$\sum_{i=1}^n c_{1,i}b_{i,2}$	$\sum_{i=1}^n c_{1,i}b_{i,3}$	$\sum_{i=1}^n c_{1,i}b_{i,4}$	$\sum_{i=1}^n c_{1,i}b_{i,5}$
$\sum_{i=1}^n c_{2,i}b_{i,1}$	$\sum_{i=1}^n c_{2,i}b_{i,2}$	$\sum_{i=1}^n c_{2,i}b_{i,3}$	$\sum_{i=1}^n c_{2,i}b_{i,4}$	$\sum_{i=1}^n c_{2,i}b_{i,5}$
$\sum_{i=1}^n c_{3,i}b_{i,1}$	$\sum_{i=1}^n c_{3,i}b_{i,2}$	$\sum_{i=1}^n c_{3,i}b_{i,3}$	$\sum_{i=1}^n c_{3,i}b_{i,4}$	$\sum_{i=1}^n c_{3,i}b_{i,5}$
$\sum_{i=1}^n c_{4,i}b_{i,1}$	$\sum_{i=1}^n c_{4,i}b_{i,2}$	$\sum_{i=1}^n c_{4,i}b_{i,3}$	$\sum_{i=1}^n c_{4,i}b_{i,4}$	$\sum_{i=1}^n c_{4,i}b_{i,5}$
$\sum_{i=1}^n c_{5,i}b_{i,1}$	$\sum_{i=1}^n c_{5,i}b_{i,2}$	$\sum_{i=1}^n c_{5,i}b_{i,3}$	$\sum_{i=1}^n c_{5,i}b_{i,4}$	$\sum_{i=1}^n c_{5,i}b_{i,5}$

Algebraic cryptanalysis : modelisation example

Example. $\mathbf{A}^{-1}\mathbf{D}_1 = \mathbf{C}_1\mathbf{B}$
 $\mathbf{A}^{-1}\mathbf{D}_2 = \mathbf{C}_2\mathbf{B}$

$$\mathbf{A}^{-1}\mathbf{D}_1$$

$\sum_{i=1}^n a_{1,i}d_{i,1}$	$\sum_{i=1}^n a_{1,i}d_{i,2}$	$\sum_{i=1}^n a_{1,i}d_{i,3}$	$\sum_{i=1}^n a_{1,i}d_{i,4}$	$\sum_{i=1}^n a_{1,i}d_{i,5}$
$\sum_{i=1}^n a_{2,i}d_{i,1}$	$\sum_{i=1}^n a_{2,i}d_{i,2}$	$\sum_{i=1}^n a_{2,i}d_{i,3}$	$\sum_{i=1}^n a_{2,i}d_{i,4}$	$\sum_{i=1}^n a_{2,i}d_{i,5}$
$\sum_{i=1}^n a_{3,i}d_{i,1}$	$\sum_{i=1}^n a_{3,i}d_{i,2}$	$\sum_{i=1}^n a_{3,i}d_{i,3}$	$\sum_{i=1}^n a_{3,i}d_{i,4}$	$\sum_{i=1}^n a_{3,i}d_{i,5}$
$\sum_{i=1}^n a_{4,i}d_{i,1}$	$\sum_{i=1}^n a_{4,i}d_{i,2}$	$\sum_{i=1}^n a_{4,i}d_{i,3}$	$\sum_{i=1}^n a_{4,i}d_{i,4}$	$\sum_{i=1}^n a_{4,i}d_{i,5}$
$\sum_{i=1}^n a_{5,i}d_{i,1}$	$\sum_{i=1}^n a_{5,i}d_{i,2}$	$\sum_{i=1}^n a_{5,i}d_{i,3}$	$\sum_{i=1}^n a_{5,i}d_{i,4}$	$\sum_{i=1}^n a_{5,i}d_{i,5}$

=

$$\mathbf{C}_1\mathbf{B}$$

$$\sum_{i=1}^n a_{1,i}d_{i,1} - \sum_{i=1}^n c_{1,i}b_{i,1} = 0, \quad \sum_{i=1}^n a_{1,i}d_{i,2} - \sum_{i=1}^n c_{1,i}b_{i,2} = 0, \dots$$

$$\sum_{i=1}^n a_{2,i}d_{i,1} - \sum_{i=1}^n c_{2,i}b_{i,1} = 0, \quad \sum_{i=1}^n a_{l,i}d_{i,p} - \sum_{i=1}^n c_{l,i}b_{i,p} = 0$$

$\sum_{i=1}^n c_{1,i}b_{i,1}$	$\sum_{i=1}^n c_{1,i}b_{i,2}$	$\sum_{i=1}^n c_{1,i}b_{i,3}$	$\sum_{i=1}^n c_{1,i}b_{i,4}$	$\sum_{i=1}^n c_{1,i}b_{i,5}$
$\sum_{i=1}^n c_{2,i}b_{i,1}$	$\sum_{i=1}^n c_{2,i}b_{i,2}$	$\sum_{i=1}^n c_{2,i}b_{i,3}$	$\sum_{i=1}^n c_{2,i}b_{i,4}$	$\sum_{i=1}^n c_{2,i}b_{i,5}$
$\sum_{i=1}^n c_{3,i}b_{i,1}$	$\sum_{i=1}^n c_{3,i}b_{i,2}$	$\sum_{i=1}^n c_{3,i}b_{i,3}$	$\sum_{i=1}^n c_{3,i}b_{i,4}$	$\sum_{i=1}^n c_{3,i}b_{i,5}$
$\sum_{i=1}^n c_{4,i}b_{i,1}$	$\sum_{i=1}^n c_{4,i}b_{i,2}$	$\sum_{i=1}^n c_{4,i}b_{i,3}$	$\sum_{i=1}^n c_{4,i}b_{i,4}$	$\sum_{i=1}^n c_{4,i}b_{i,5}$
$\sum_{i=1}^n c_{5,i}b_{i,1}$	$\sum_{i=1}^n c_{5,i}b_{i,2}$	$\sum_{i=1}^n c_{5,i}b_{i,3}$	$\sum_{i=1}^n c_{5,i}b_{i,4}$	$\sum_{i=1}^n c_{5,i}b_{i,5}$

Algebraic cryptanalysis : modelisation example

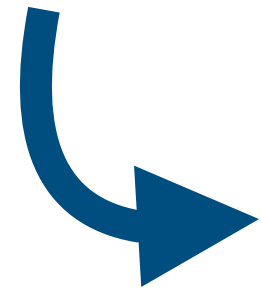


A motivating example: a better idea for modelisation.

Given matrices $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2 \in \mathcal{M}_{n,n}(\mathbb{F}_q)$ (the space of matrices over \mathbb{F}_q of size $n \times n$), find $\mathbf{A}, \mathbf{B} \in \text{GL}_n(\mathbb{F}_q)$ (the space of invertible matrices over \mathbb{F}_q of size $n \times n$), such that

$$\begin{aligned}\mathbf{A}^{-1}\mathbf{D}_1 &= \mathbf{C}_1\mathbf{B} \\ \mathbf{A}^{-1}\mathbf{D}_2 &= \mathbf{C}_2\mathbf{B}\end{aligned}$$

Algebraic cryptanalysis : modelisation example



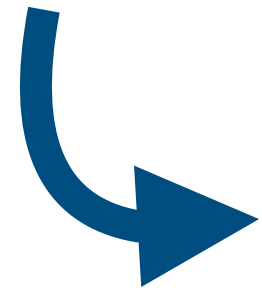
A motivating example: a better idea for modelisation.

Given matrices $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2 \in \mathcal{M}_{n,n}(\mathbb{F}_q)$ (the space of matrices over \mathbb{F}_q of size $n \times n$), find $\mathbf{A}, \mathbf{B} \in \text{GL}_n(\mathbb{F}_q)$ (the space of invertible matrices over \mathbb{F}_q of size $n \times n$), such that

$$\begin{aligned}\mathbf{A}^{-1}\mathbf{D}_1 &= \mathbf{C}_1\mathbf{B} \\ \mathbf{A}^{-1}\mathbf{D}_2 &= \mathbf{C}_2\mathbf{B}\end{aligned}$$

→ Results in a **linear** system with the same number of variables and equations.

Algebraic cryptanalysis : modelisation example



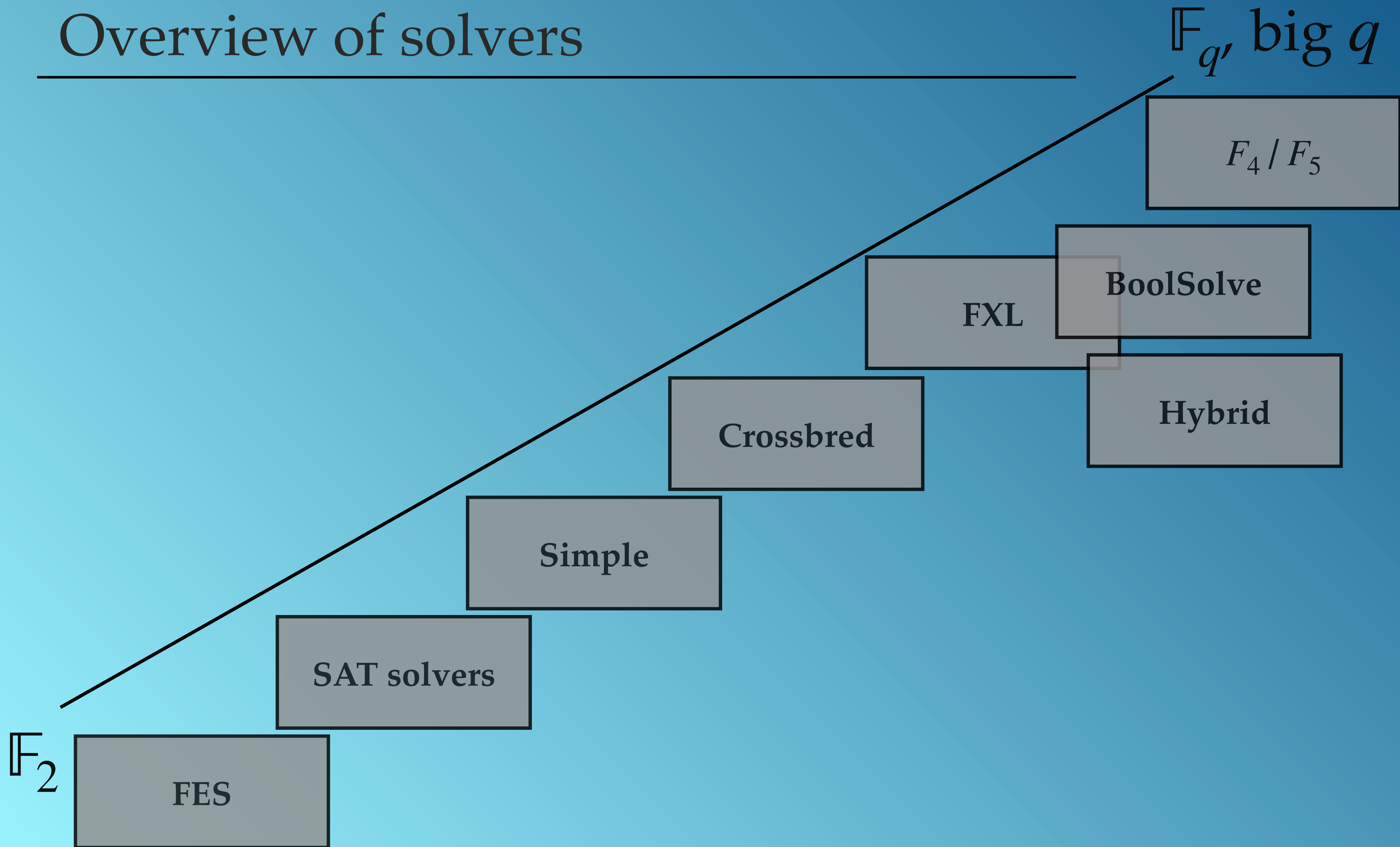
A motivating example: a better idea for modelisation.

Given matrices $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2 \in \mathcal{M}_{n,n}(\mathbb{F}_q)$ (the space of matrices over \mathbb{F}_q of size $n \times n$), find $\mathbf{A}, \mathbf{B} \in \text{GL}_n(\mathbb{F}_q)$ (the space of invertible matrices over \mathbb{F}_q of size $n \times n$), such that

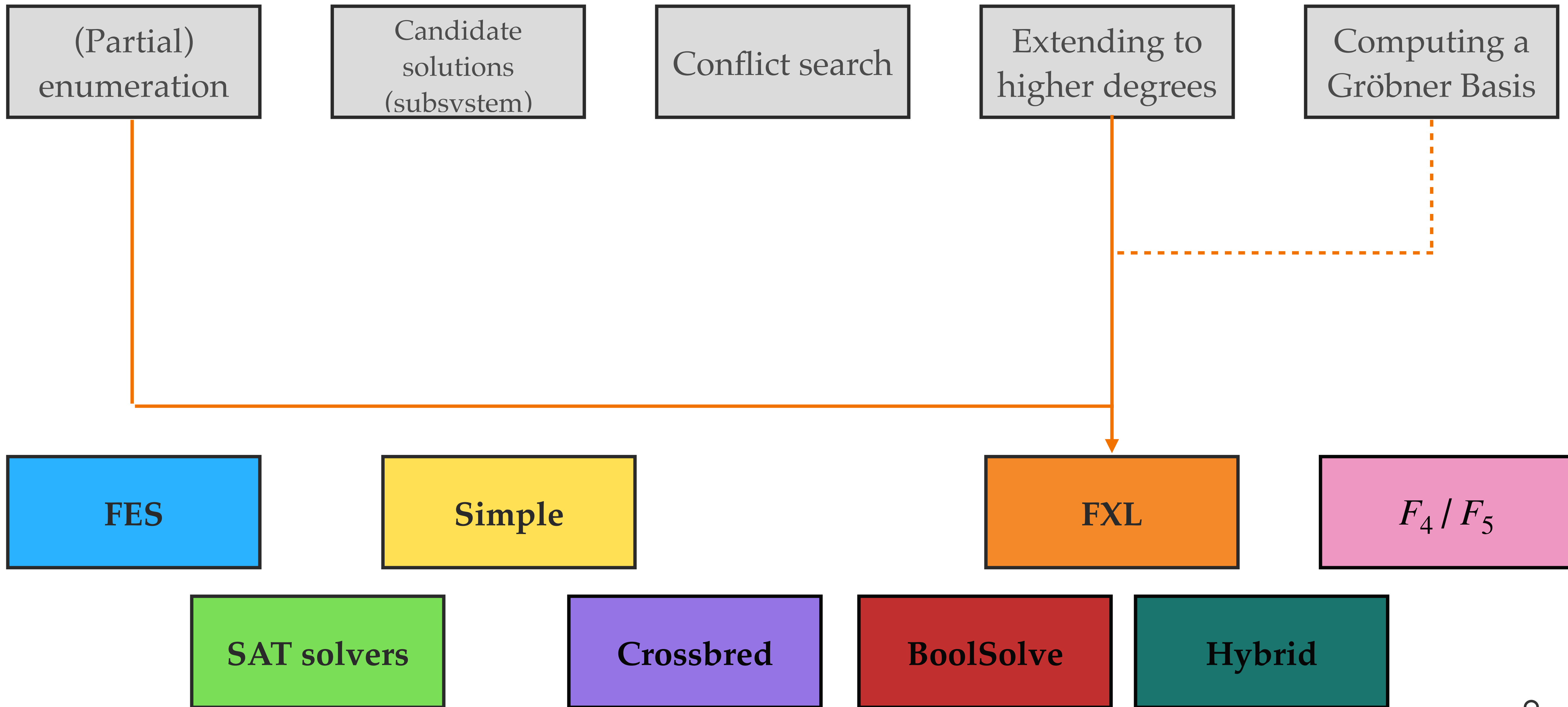
$$\begin{aligned}\mathbf{A}^{-1}\mathbf{D}_1 &= \mathbf{C}_1\mathbf{B} \\ \mathbf{A}^{-1}\mathbf{D}_2 &= \mathbf{C}_2\mathbf{B}\end{aligned}$$

- Results in a **linear** system with the same number of variables and equations.
- If $\mathbf{C}_1, \mathbf{C}_2, \mathbf{D}_1, \mathbf{D}_2$ are all full rank, we should have a unique solution.
- We can easily recover \mathbf{A} from \mathbf{A}^{-1} .

Overview of solvers



Summary



Summary

(Partial)
enumeration

Candidate
solutions
(subsystem)

Conflict search

Extending to
higher degrees

Computing a
Gröbner Basis

FES

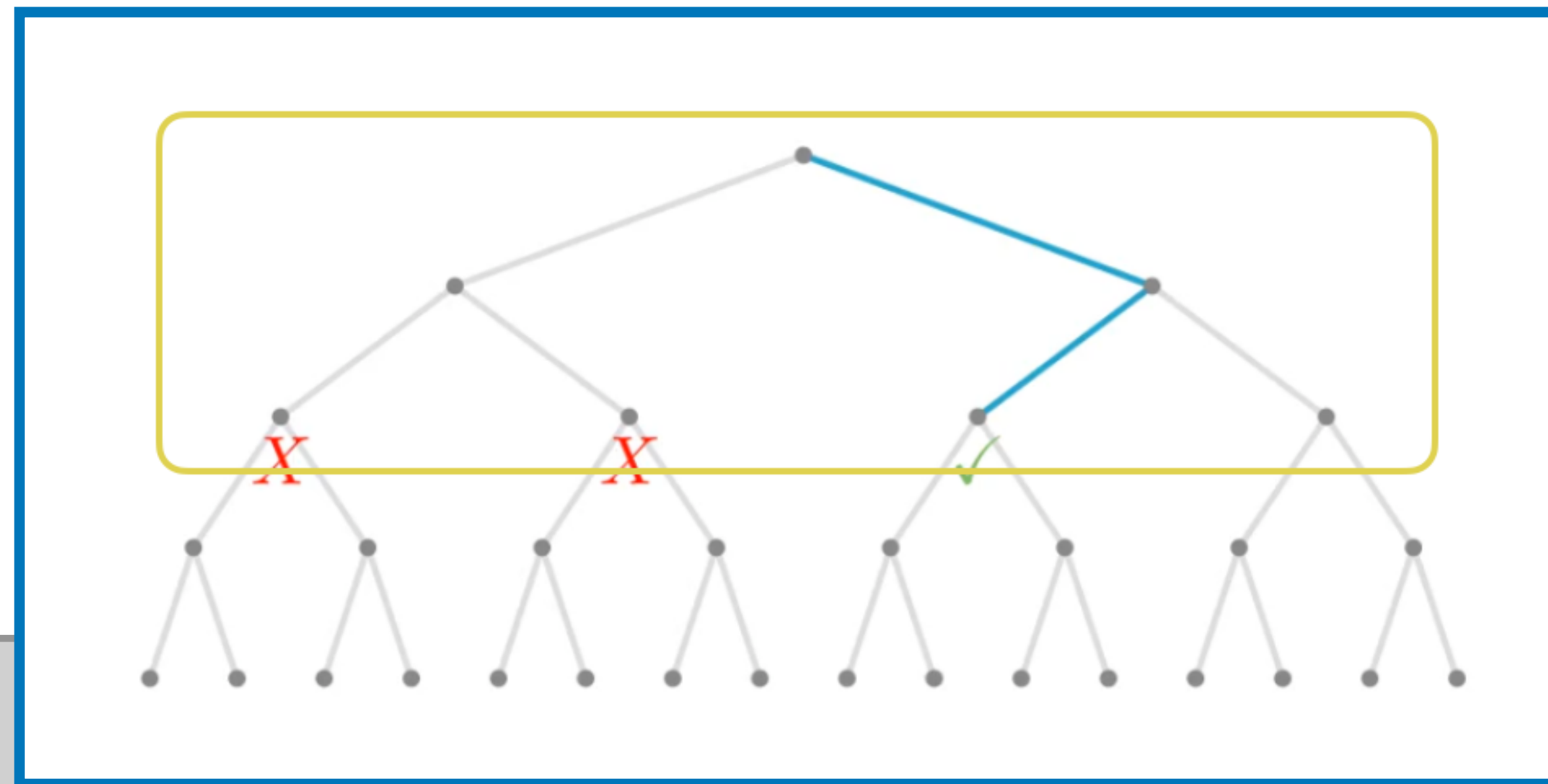
SAT solvers

Crossbred

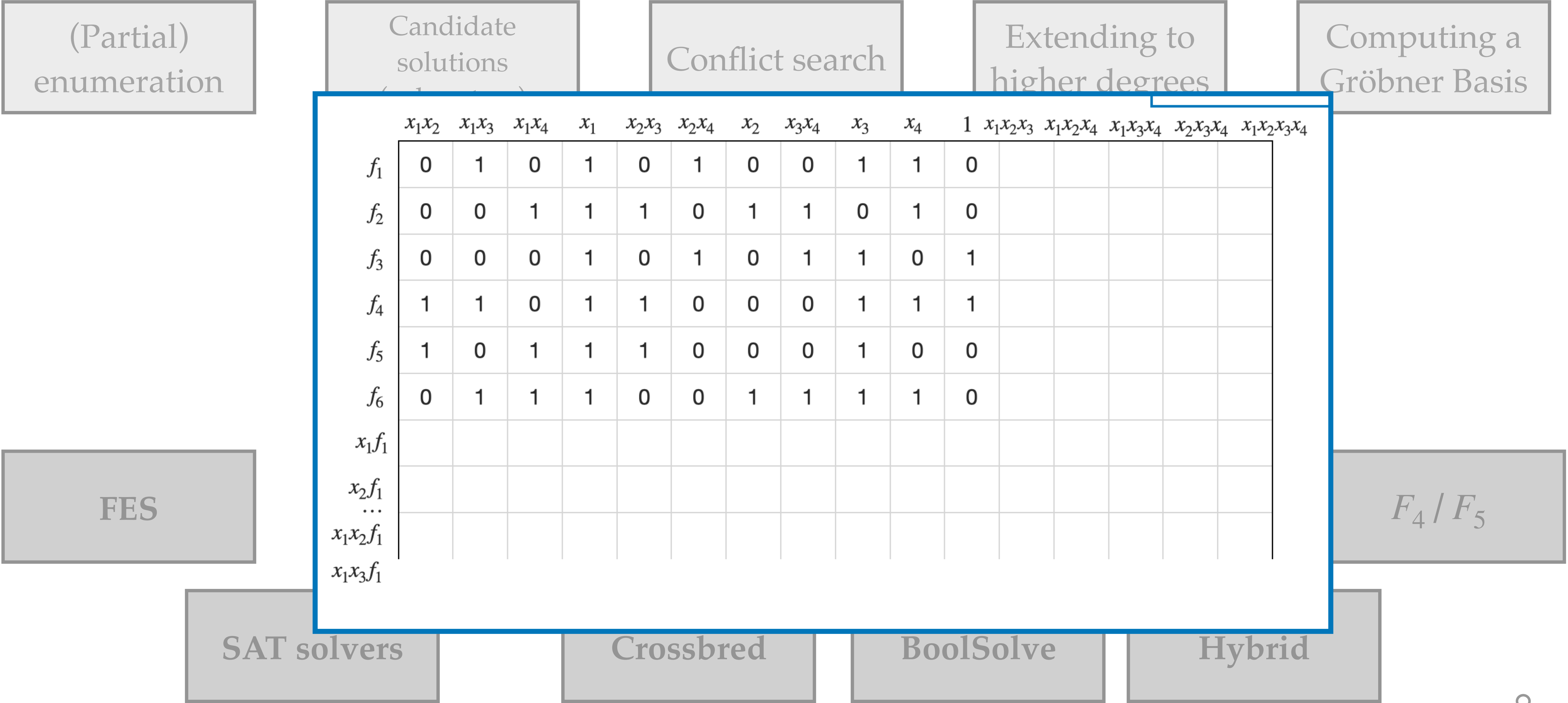
BoolSolve

Hybrid

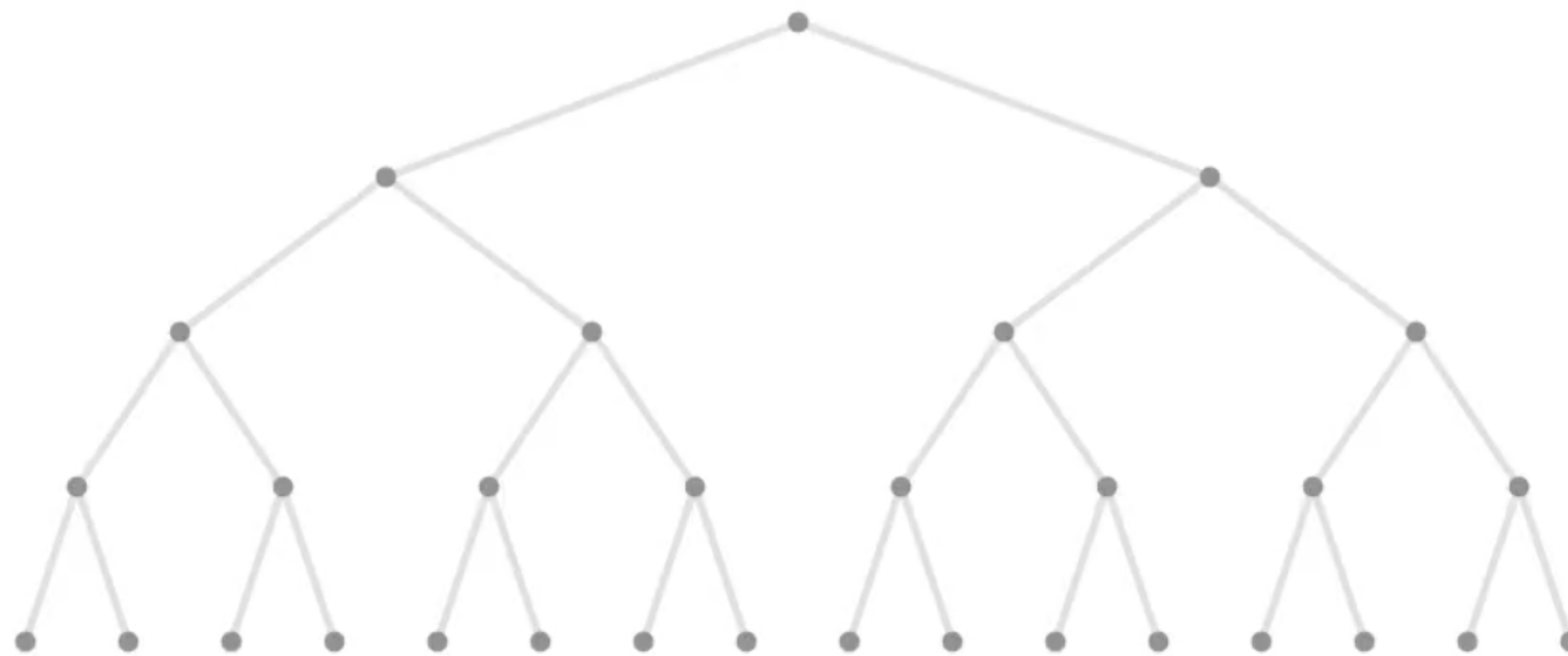
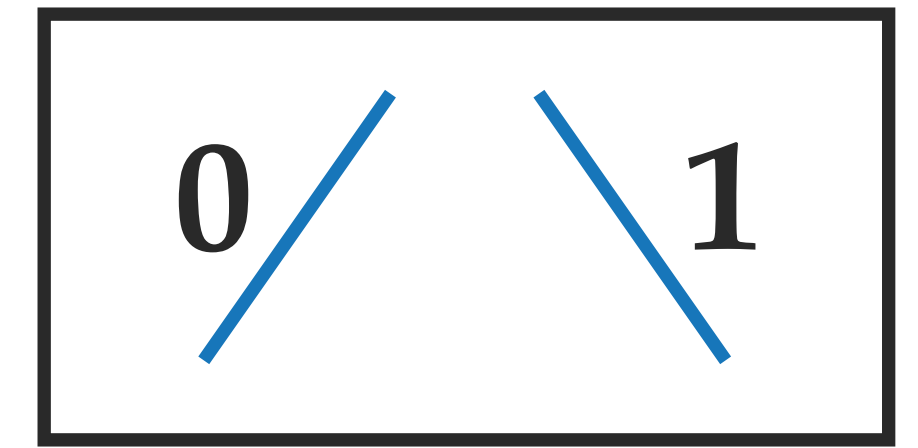
F_4 / F_5



Summary



Exhaustive Search



$$x_1 \cdot x_2 + x_1 \cdot x_3 + x_3 \cdot x_4 + x_3 = 0$$

$$x_2 \cdot x_3 + x_2 \cdot x_4 + x_1 + x_2 + 1 = 0$$

$$x_1 \cdot x_2 + x_2 \cdot x_3 + x_2 \cdot x_4 + x_1 + x_4 = 0$$

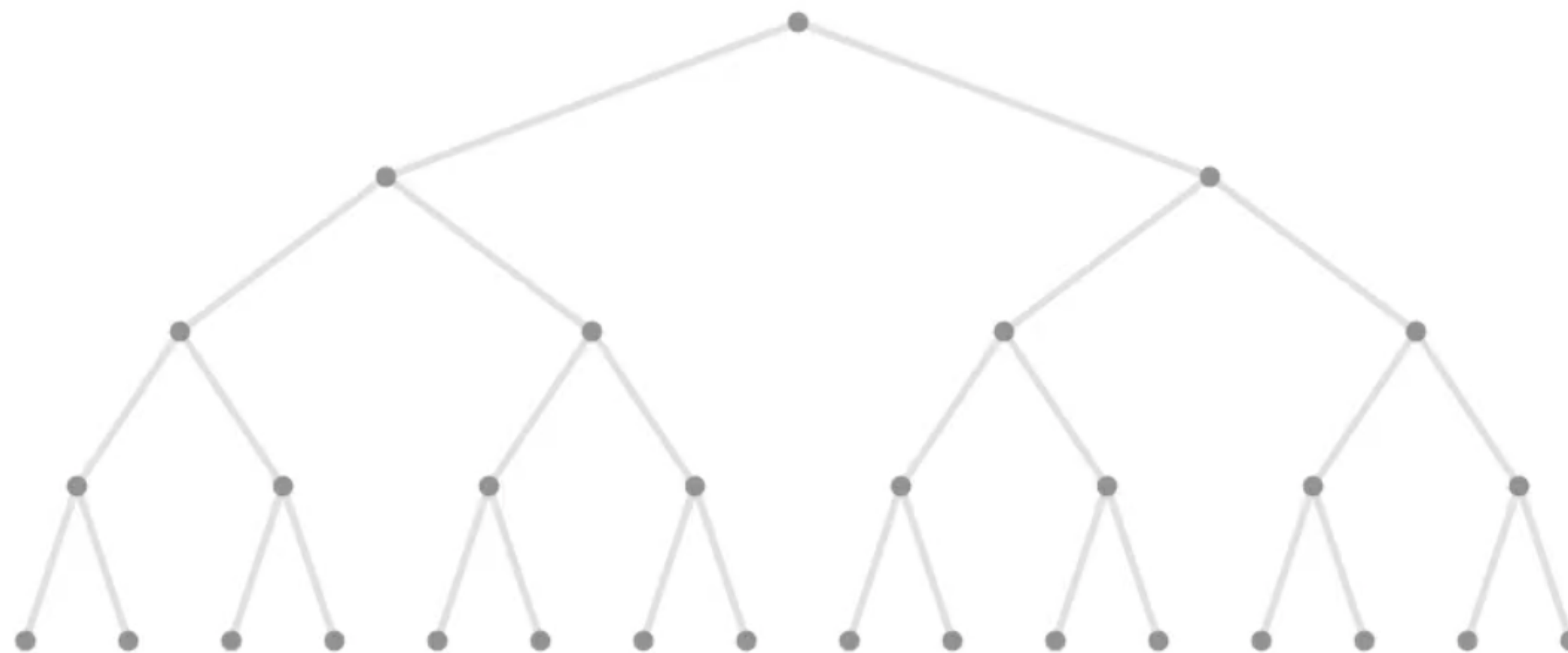
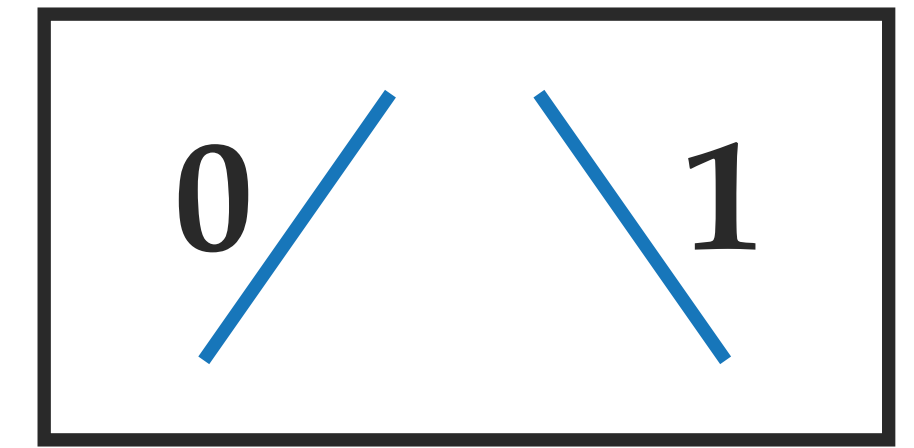
$$x_1 \cdot x_4 + x_2 \cdot x_3 + x_2 + x_3 + x_4 = 0$$

Binary search tree

Exhaustive Search



Worst-case complexity: $\mathcal{O}(2^n)$



Binary search tree

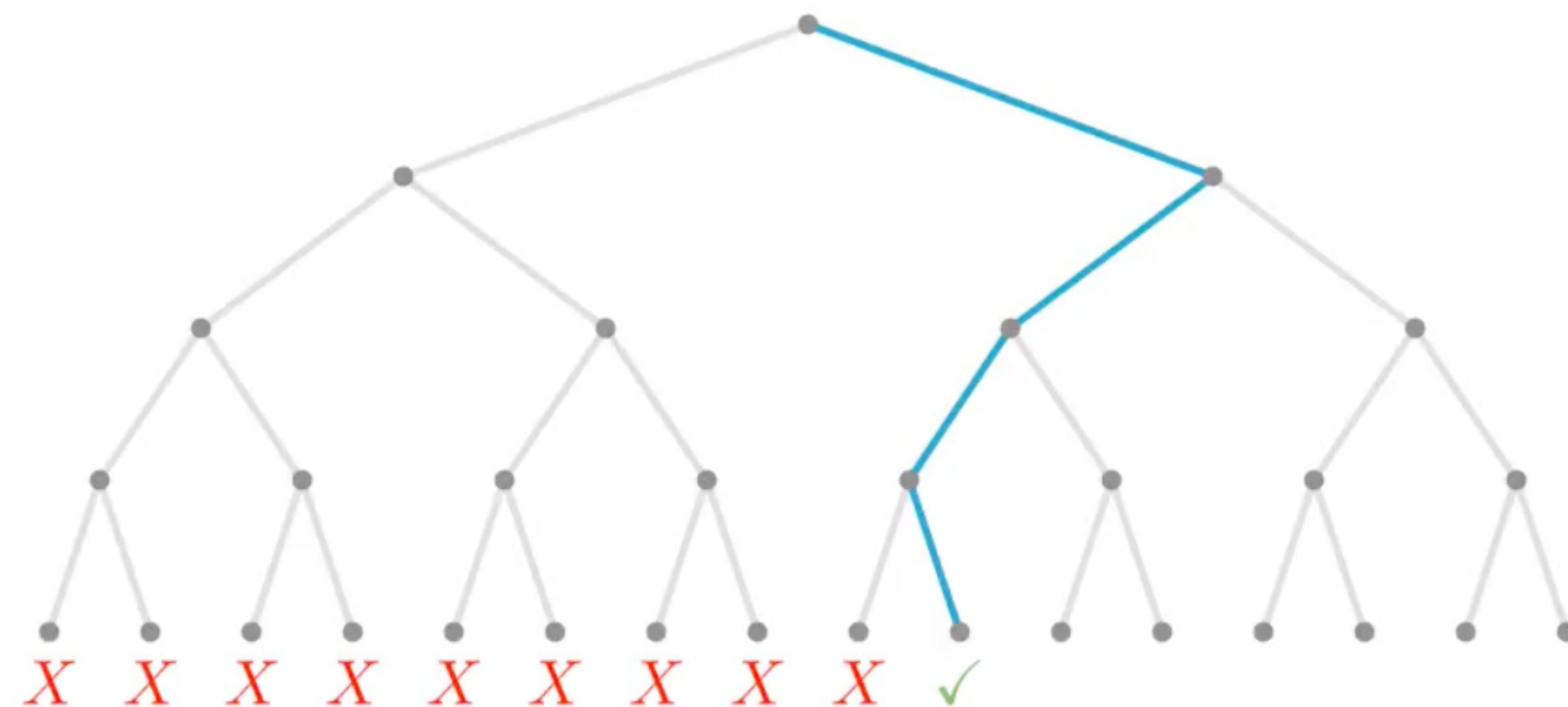
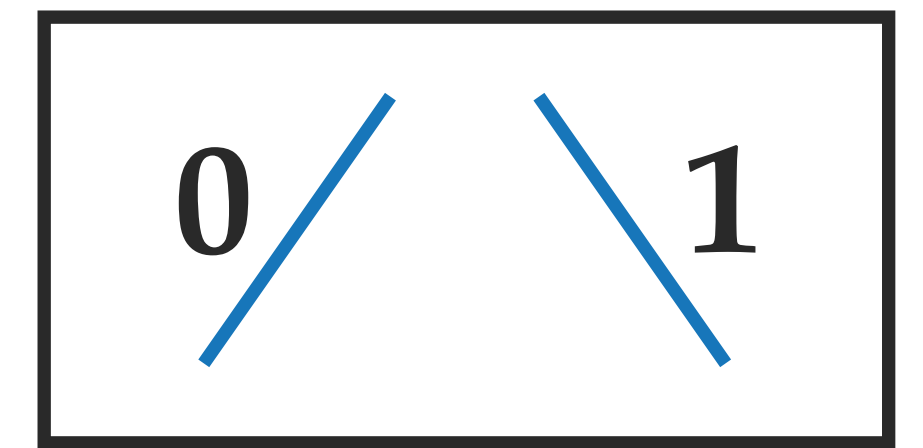
$$x_1 \cdot x_2 + x_1 \cdot x_3 + x_3 \cdot x_4 + x_3 = 0$$

$$x_2 \cdot x_3 + x_2 \cdot x_4 + x_1 + x_2 + 1 = 0$$

$$x_1 \cdot x_2 + x_2 \cdot x_3 + x_2 \cdot x_4 + x_1 + x_4 = 0$$

$$x_1 \cdot x_4 + x_2 \cdot x_3 + x_2 + x_3 + x_4 = 0$$

Exhaustive Search

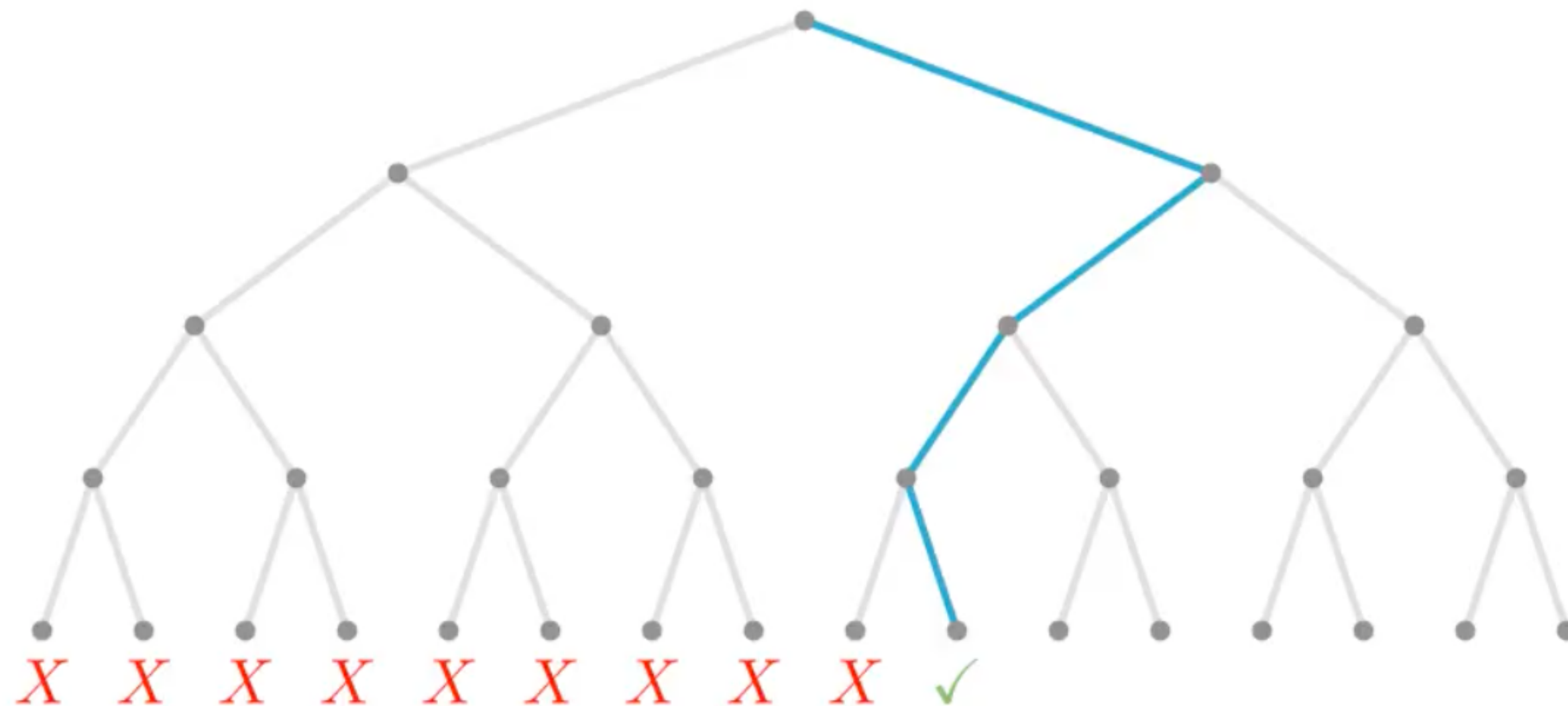
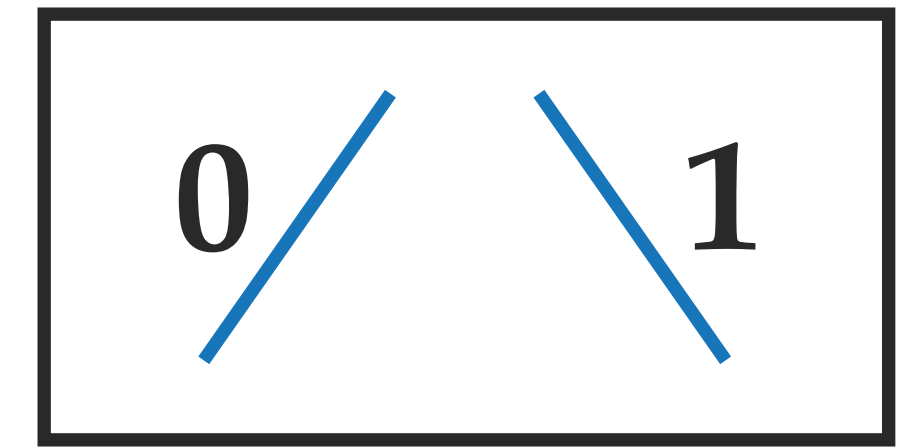


Binary search tree

$$\begin{aligned} 1 \cdot 0 + 1 \cdot 0 + 0 \cdot 1 + 0 &= 0 \\ 0 \cdot 0 + 0 \cdot 1 + 1 + 0 + 1 &= 0 \\ 1 \cdot 0 + 0 \cdot 0 + 0 \cdot 1 + 1 + 1 &= 0 \\ 1 \cdot 1 + 0 \cdot 0 + 0 + 0 + 1 &= 0 \end{aligned}$$

Exhaustive Search

Worst-case complexity: $\mathcal{O}(2^n)$



Binary search tree

$$\begin{aligned} 1 \cdot 0 + 1 \cdot 0 + 0 \cdot 1 + 0 &= 0 \\ 0 \cdot 0 + 0 \cdot 1 + 1 + 0 + 1 &= 0 \\ 1 \cdot 0 + 0 \cdot 0 + 0 \cdot 1 + 1 + 1 &= 0 \\ 1 \cdot 1 + 0 \cdot 0 + 0 + 0 + 1 &= 0 \end{aligned}$$

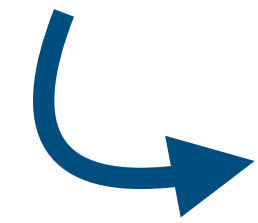
Macaulay matrix

Linearisation

Linear systems are easy to solve, nonlinear systems are hard.

Linearisation

Linear systems are easy to solve, nonlinear systems are hard.



Linearisation: for each nonlinear monomial, replace all of its occurrences by a new variable.

Linearisation

Linear systems are easy to solve, nonlinear systems are hard.

↪ **Linearisation:** for each nonlinear monomial, replace all of its occurrences by a new variable.

Example.

$$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$$

$$f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$$

$$f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$$

$$f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$$

$$f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$$

$$f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$$



$$f_1 : y_2 + y_5 + x_1 + x_3 + x_4 = 0$$

$$f_2 : y_4 + y_3 + y_6 + x_1 + x_2 + x_4 = 0$$

$$f_3 : y_5 + y_6 + x_1 + x_3 + 1 = 0$$

$$f_4 : y_1 + y_2 + y_4 + x_3 + x_4 + 1 = 0$$

$$f_5 : y_1 + y_4 + y_3 + x_3 = 0$$

$$f_6 : y_2 + y_3 + y_6 + x_1 + x_2 + x_3 + x_4 = 0$$

Linearisation

Linear systems are easy to solve, nonlinear systems are hard.

↪ **Linearisation:** for each nonlinear monomial, replace all of its occurrences by a new variable.

Example.

$$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$$

$$f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$$

$$f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$$

$$f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$$

$$f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$$

$$f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$$



$$f_1 : y_2 + y_5 + x_1 + x_3 + x_4 = 0$$

$$f_2 : y_4 + y_3 + y_6 + x_1 + x_2 + x_4 = 0$$

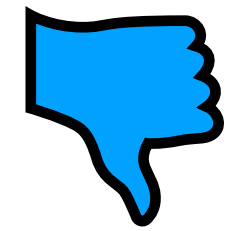
$$f_3 : y_5 + y_6 + x_1 + x_3 + 1 = 0$$

$$f_4 : y_1 + y_2 + y_4 + x_3 + x_4 + 1 = 0$$

$$f_5 : y_1 + y_4 + y_3 + x_3 = 0$$

$$f_6 : y_2 + y_3 + y_6 + x_1 + x_2 + x_3 + x_4 = 0$$

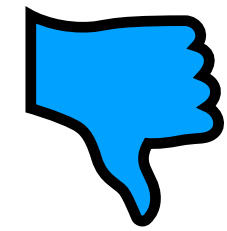
Linearisation



Linearisation adds solutions: a *random* quadratic system of m equations in n variables, when $n = m$, is expected to have one solution (probability is $\sim \frac{1}{q}$ for systems over \mathbb{F}_q). The corresponding linearised system has a solution space of dimension $\binom{n+1}{2} - m$.

$\uparrow \binom{n}{2}$ quadratic plus n linear monomials

Linearisation



Linearisation adds solutions: a *random* quadratic system of m equations in n variables, when $n = m$, is expected to have one solution (probability is $\sim \frac{1}{q}$ for systems over \mathbb{F}_q). The corresponding linearised system has a solution space of dimension $\binom{n+1}{2} - m$.

 $\binom{n}{2}$ quadratic plus n linear monomials



Loss of information: e.g. assignment $x_1 = 1; x_2 = 0; y_1 = 1$; is part of a valid solution to the linearised system, but $x_1 x_2 \neq y_1$.

Macaulay matrix

Equations
↓

Monomials
→

	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1
f_1											
f_2											
f_3											
f_4											
f_5											
f_6											

$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$ $f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$ $f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$ $f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$ $f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$ $f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$

Macaulay matrix

Equations
↓

	Monomials →										
	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1
f_1	0	1	0	1	0	1	0	0	1	1	0
f_2	0	0	1	1	1	0	1	1	0	1	0
f_3	0	0	0	1	0	1	0	1	1	0	1
f_4	1	1	0	1	1	0	0	0	1	1	1
f_5	1	0	1	1	1	0	0	0	1	0	0
f_6	0	1	1	1	0	0	1	1	1	1	0

$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$
 $f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$
 $f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$
 $f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$
 $f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$
 $f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$



Gröbner basis algorithms

[Buchberger, 1965]

[Lazard, 1983]

F_4/F_5 [Faugère, 1999/2002]

(XL [Courtois, Klimov, Patarin, Shamir, 2000])

Gröbner basis algorithms (intuition)

*We are essentially describing the XL algorithm.

Gröbner basis algorithms (intuition)

*We are essentially describing the XL algorithm.

	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1
f_1	0	1	0	1	0	1	0	0	1	1	0
f_2	0	0	1	1	1	0	1	1	0	1	0
f_3	0	0	0	1	0	1	0	1	1	0	1
f_4	1	1	0	1	1	0	0	0	1	1	1
f_5	1	0	1	1	1	0	0	0	1	0	0
f_6	0	1	1	1	0	0	1	1	1	1	0

Gröbner basis algorithms (intuition)

*We are essentially describing the XL algorithm.

$$f_1 : x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0$$
$$f_2 : x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0$$
$$f_3 : x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0$$
$$f_4 : x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0$$
$$f_5 : x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0$$
$$f_6 : x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0$$

	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1
f_1	0	1	0	1	0	1	0	0	1	1	0
f_2	0	0	1	1	1	0	1	1	0	1	0
f_3	0	0	0	1	0	1	0	1	1	0	1
f_4	1	1	0	1	1	0	0	0	1	1	1
f_5	1	0	1	1	1	0	0	0	1	0	0
f_6	0	1	1	1	0	0	1	1	1	1	0

Gröbner basis algorithms (intuition)

*We are essentially describing the XL algorithm.

$D = 3$

$$\begin{aligned} f_1 &: x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0 \\ f_2 &: x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0 \\ f_3 &: x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0 \\ f_4 &: x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0 \\ f_5 &: x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0 \\ f_6 &: x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0 \end{aligned}$$

	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1	$x_1x_2x_3$	$x_1x_2x_4$	$x_1x_3x_4$	$x_2x_3x_4$
f_1	0	1	0	1	0	1	0	0	1	1	0				
f_2	0	0	1	1	1	0	1	1	0	1	0				
f_3	0	0	0	1	0	1	0	1	1	0	1				
f_4	1	1	0	1	1	0	0	0	1	1	1				
f_5	1	0	1	1	1	0	0	0	1	0	0				
f_6	0	1	1	1	0	0	1	1	1	1	0				
x_1f_1															
x_2f_1															
...															

Gröbner basis algorithms (intuition)

*We are essentially describing the XL algorithm.

$D = 4$

$$\begin{aligned} f_1 &: x_1x_3 + x_2x_4 + x_1 + x_3 + x_4 = 0 \\ f_2 &: x_2x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_4 = 0 \\ f_3 &: x_2x_4 + x_3x_4 + x_1 + x_3 + 1 = 0 \\ f_4 &: x_1x_2 + x_1x_3 + x_2x_3 + x_3 + x_4 + 1 = 0 \\ f_5 &: x_1x_2 + x_2x_3 + x_1x_4 + x_3 = 0 \\ f_6 &: x_1x_3 + x_1x_4 + x_3x_4 + x_1 + x_2 + x_3 + x_4 = 0 \end{aligned}$$

	x_1x_2	x_1x_3	x_1x_4	x_1	x_2x_3	x_2x_4	x_2	x_3x_4	x_3	x_4	1	$x_1x_2x_3$	$x_1x_2x_4$	$x_1x_3x_4$	$x_2x_3x_4$	$x_1x_2x_3x_4$
f_1	0	1	0	1	0	1	0	0	1	1	0					
f_2	0	0	1	1	1	0	1	1	0	1	0					
f_3	0	0	0	1	0	1	0	1	1	0	1					
f_4	1	1	0	1	1	0	0	0	1	1	1					
f_5	1	0	1	1	1	0	0	0	1	0	0					
f_6	0	1	1	1	0	0	1	1	1	1	0					
x_1f_1																
x_2f_1																
...																
$x_1x_2f_1$																
$x_1x_3f_1$																

XL / Gröbner basis algorithms: complexity

XL / Gröbner basis algorithms: complexity

$$\mathcal{O} \left(m D_{reg} \binom{n + D_{reg} - 1}{D_{reg}}^{\omega} \right)$$

XL / Gröbner basis algorithms: complexity

$$\mathcal{O} \left(m D_{reg} \binom{n + D_{reg} - 1}{D_{reg}}^{\omega} \right)$$

D_{reg} : degree of regularity



the power of the first non-positive coefficient in the expansion of

$$\frac{(1 - t^2)^m}{(1 - t)^n}$$

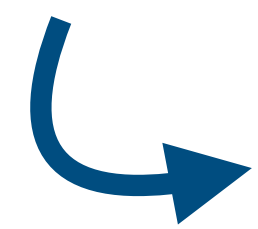
XL / Gröbner basis algorithms: complexity

```
▷ m=8
  n=7
  R.<t> = PowerSeriesRing(ZZ)
  hs = ((1-t^2)^m) / (1-t)^n
  print(hs)
[3] ✓ 0.0s
... 1 + 7*t + 20*t^2 + 28*t^3 + 14*t^4 - 14*t^5 - 28*t^6 - 20*t^7 - 7*t^8 - t^9 + 0(t^20)
```

The number of monomials (columns) **minus** linearly independent equations (rows) at **degree $D = 4$** is 14.

XL / Gröbner basis algorithms: complexity

D_{reg} : degree of regularity



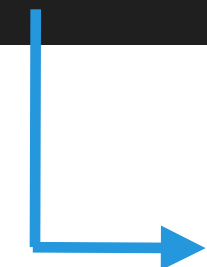
the power of the first non-positive coefficient in the expansion of

$$\frac{(1 - t^2)^m}{(1 - t)^n}$$

```

m=8
n=7
R.<t> = PowerSeriesRing(ZZ)
hs = ((1-t^2)^(m)) / (1-t)^(n)
print(hs)
[3]  ✓  0.0s
...  1 + 7*t + 20*t^2 + 28*t^3 + 14*t^4 - 14*t^5 - 28*t^6 - 20*t^7 - 7*t^8 - t^9 + 0(t^20)

```



The number of monomials (columns) **minus** linearly independent equations (rows) at **degree $D = 4$** is 14.

Summary

