

# H 335: Computer Structures

## Spring 2016 – Homework 2

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**Due:** 2-23-2016

1. **Prove, using induction, that for radix-r, the largest number that can be represented with  $N$  digits is  $r^N - 1$ :**

$$(d_{N-1}, \dots, d_0)_r = \sum_{i=0}^{N-1} d_i \times r^i, d_i \in \{0..r-1\}$$

**Base case:**  $N = 1$

$$\sum_{i=0}^{N-1} d_i \times r^i = r^N - 1$$

$$\sum_{i=0}^0 0 \times r^0 = r^0 - 1$$

$$0 = 0$$

**Inductive hypothesis:**  $N = k$

$$\sum_{i=0}^{k-1} d_i \times r^i = r^k - 1$$

**Inductive step:**  $N = k + 1$

$$\sum_{i=0}^k d_i \times r^i = r^{k+1} - 1$$

$$(\sum_{i=0}^1 d_i \times r^i) + (\sum_{i=0}^{k-1} d_i \times r^i) = r^{k+1} - 1$$

$$(\sum_{i=0}^1 d_i \times r^i) + (r^k - 1) = r^{k+1} - 1$$

$$(\sum_{i=0}^1 d_i \times r^i) + r^k = r^{k+1}$$

$$(\sum_{i=0}^1 d_i \times r^i) = r^{k+1} - r^k$$

$$(\sum_{i=0}^1 d_i \times r^i) = (r - 1)r^k$$

2. **Prove that for radix-r addition, the carry bits are always 0 or 1:**

The carry bit for any size radix will always be either 0 or 1. This is because when adding radix number we simply add the two first digits and move on. A single digit number when added to another single digit number will never give a number greater than or equal to 20. Therefore, every carry can only be 1 or zero to represent carrying a ten or a radix-r where r is the given radix system.

$$9 + 9 = 18$$

$18 < 20 \therefore$  carry would only be a 1.

3. **Given the formal definition, derive the minimum and maximum two's complement numbers that can be represented in  $N$  bits:**

$$(b_N - 1, \dots, b_0)_2 = -b_{N-1} \times 2^{N-1} + \sum_{i=0}^{N-2} b_i \times 2^i$$

When finding the largest two's complement number given  $N$  bits:  $2^{N-1} - 1$

$$2^{8-1} - 1 = 2^7 - 1 = 128 - 1 = +127$$

$$\text{largest} = +127 = 0111\ 1111$$

When finding the smallest two's complement number given  $N$  bits:  $-2^{N-1}$

$$-2^{8-1} = -2^7 = -128$$

$$\text{smallest} = -128 = 1000\ 0000$$

4. **For a number  $B$  with magnitude less than  $2^{N-2}$ , show that if  $B$  is represented by 2's complement number with  $N$  bits  $b_{N-1}..b_0$  then  $-(b_N - 1, \dots, b_0)_2 = (b_N - 1, \dots, b_0)_2 + 1$**

$$-(b_N - 1, \dots, b_0)_2 = \overline{(b_N - 1, \dots, b_0)}_2 + 1$$

$$-(0010)_2 = (\overline{0010})_2 + 1$$

$$lhs = -(0111)_2 = (1111)_2$$

$$rhs = (\overline{0111})_2 + 1 = (1000)_2 + 1 = (1001)_2$$

$$lhs = rhs$$

5. **Prove the "sign-extension" is value preserving:**

Sign-extension is value preserving because when using the two's complement, the most significant bit is the value needed to show if a value is positive or negative. That is why the greatest value cannot be used with an  $N$ -bit number. When using two's complement, the greatest value in an  $N$ -bit number must be  $(2^{N-1} - 1)$  because if we used the most significant bit as a value instead of a marker for positive or negative, we would not be able to know if the value was positive or negative. This was shown in question 3.

6. **Convert the following decimal numbers to binary:**

(a) **107 = 110 1011**

Repeated Division (mod 2)							
Quotients	107	53	26	13	6	3	1
Remainders	1	1	0	1	0	1	1

$$(107)_{10} = (1101011)_2$$

(b) **2312 = 1001 0000 1000**

Repeated Division (mod 2)												
Quotients	2312	1156	578	289	144	72	36	18	9	4	2	1
Remainders	0	0	0	1	0	0	0	0	1	0	0	1

$(2312)_{10} = (100100001000)_2$

(c) **31333 = 111 1010 0110 0101**

Repeated Division (mod 2)														
Quotients	31333	15666	7833	3916	1958	979	489	244	122	61	30	15	7	3
Remainders	1	0	1	0	0	1	1	0	0	1	0	1	1	1

$(31333)_{10} = (111101001100101)_2$

(d) **97 = 110 0001**

Repeated Division (mod 2)							
Quotients	97	48	24	12	6	3	1
Remainders	1	0	0	0	0	1	1

$(97)_{10} = (1100001)_2$

7. Perform the following subtraction operations using complements:

(a) **103 - 92 = 11**

0: $(10 - 1) - 0 = 9 - 0 = 9$	$\boxed{\cancel{1}}$	$\boxed{0}$	$\boxed{1}$	
9: $(10 - 1) - 9 = 9 - 9 = 0$		1	0	3
2: $(10 - 1) - 2 = 9 - 2 = 7$		9	0	7
	+			1
			1	1

(b) **1027 - 11 = 1016**

0: $(10 - 1) - 0 = 9 - 0 = 9$	$\boxed{\cancel{1}}$	$\boxed{1}$	$\boxed{1}$	$\boxed{1}$	
0: $(10 - 1) - 0 = 9 - 0 = 9$		1	0	2	7
1: $(10 - 1) - 1 = 9 - 1 = 8$		9	9	8	8
1: $(10 - 1) - 1 = 9 - 1 = 8$					1
	+				6
		1	0	1	

(c) **129 - 33 = 96**

0: $(10 - 1) - 0 = 9 - 0 = 9$	$\boxed{\cancel{1}}$	$\boxed{0}$	$\boxed{1}$	
3: $(10 - 1) - 3 = 9 - 3 = 6$		1	2	9
3: $(10 - 1) - 3 = 9 - 3 = 6$		9	6	6
	+			1
			9	6

(d) **2222 - 222 = 2000**

0: $(10 - 1) - 0 = 9 - 0 = 9$
2: $(10 - 1) - 2 = 9 - 2 = 7$
2: $(10 - 1) - 2 = 9 - 2 = 7$
2: $(10 - 1) - 2 = 9 - 2 = 7$

$\overline{1}$	1	1	1	
	2	2	2	2
	9	7	7	7
+				1
	2	0	0	0

8. Convert the following decimal numbers to 8-bit two's complement (show your work)

(a) **-91 = 1101 1011**

Repeated Division (mod 2)							
Quotients	91	45	22	11	5	2	1
Remainders	1	1	0	1	1	0	1

$(-91)_{10} = (1101\ 1011)_2$

The most significant bit will be the bit to show if the number is positive or negative. Since the decimal number will be an 8-bit two's complement, the most significant bit will be in the last position on the left and for -91 the most significant bit will be 1 to represent a negative.

(b) **-96 = 1110 0000**

Repeated Division (mod 2)							
Quotients	96	48	24	12	6	3	1
Remainders	0	0	0	0	0	1	1

$(-96)_{10} = (1110\ 0000)_2$

(c) **-126 = 1111 1110**

Repeated Division (mod 2)							
Quotients	126	63	31	15	7	3	1
Remainders	0	1	1	1	1	1	1

$(-126)_{10} = (1111\ 1110)_2$

(d) **101 = 0110 0101**

Repeated Division (mod 2)							
Quotients	101	50	25	12	6	3	1
Remainders	1	0	1	0	0	1	1

$(101)_{10} = (0110\ 0101)_2$

Now that the decimal is positive, the most significant bit will be a 0. 1 is for negative and 0 is for positive.

(e) **78 = 0100 1110**

Repeated Division (mod 2)							
Quotients	78	39	19	9	4	2	1
Remainders	0	1	1	1	0	0	1

$(78)_{10} = (0100\ 1110)_2$