

Chapter 10, section 6

- **R-10.6**

Which of the hash table collision-handling schemes could tolerate a load factor above 1 and which could not?

could handle: Separate Chaining

could not handle: Open Addressing (e.g linear probing, quadratic probing, double hashing)

- **R-10.9**

Draw the 11-entry hash table that results from using the hash function, $h(i) = (3i+5) \bmod 11$, to hash the keys 12, 44, 13, 88, 23, 94, 11, 39, 20, 16, and 5, assuming collisions are handled by chaining.

Index	0	1	2	3	4	5	6	7	8	9	10
Buckets	13	94 39				44 88 11			12 23	16 5	20

- **R-10.10**

What is the result of the previous exercise, assuming collisions are handled by linear probing?

0	1	2	3	4	5	6	7	8	9	10
13	94	39	16	5	44	88	11	12	23	20

- **R-10.12**

What is the result of Exercise R-10.9 when collisions are handled by double hashing using the secondary hash function $h(k) = 7 - (k \bmod 7)$?

0	1	2	3	4	5	6	7	8	9	10
13	94	5	39	88	44	23	11	12	16	20

- **R-10.14**

Show the result of rehashing the hash table shown in Figure 10.6 into a table of size 19 using the new hash function $h(k) = 3k \bmod 17$.

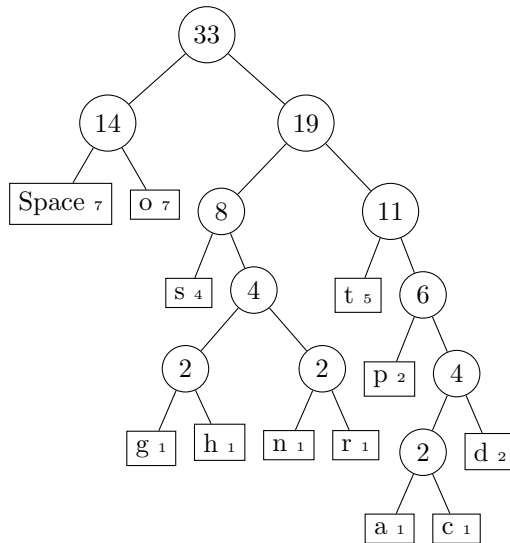
0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18
		12	18	41		36	25		54			38	10		90	28		

Chapter 13, section 6

• R-13.11

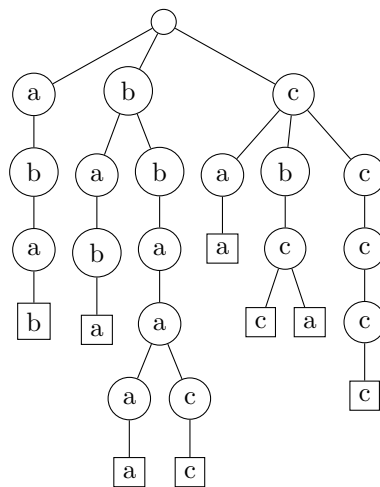
Draw the frequency array and Huffman tree for the following string:
"dogs do not spot hot pots or cats"

Character		a	c	d	g	h	n	o	p	r	s	t	Total
Frequency		7	1	1	2	1	1	7	2	1	4	5	33



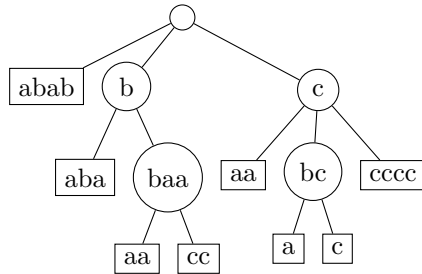
• R-13.12

Draw a standard trie for the following set of strings:
{ abab, baba, ccccc, bbaaaa, caa, bbaacc, cbcc, cbca }



• **R-13.13**

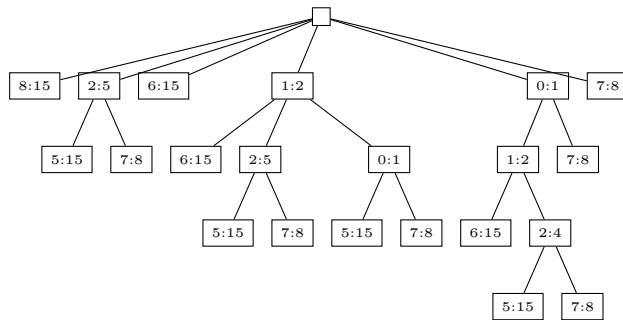
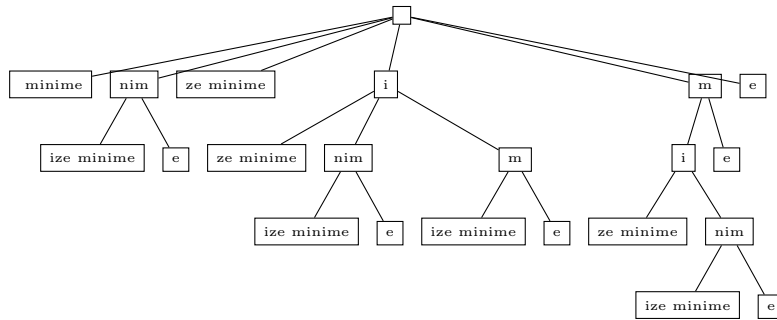
Draw a compressed trie for the strings given in the previous problem.



• **R-13.14**

Draw the compact representation of the suffix trie for the string: "minimize minimize"

0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
m	i	n	i	m	i	z	e		m	i	n	i	m	e



• **C-13.43**

Give an efficient algorithm for deleting a string from a standard trie and analyze its running time.

Input a Trie and Search for the desired string to delete. If the string reaches a leaf then begin to delete each node upwards until hitting a node with more than one child. A node cannot be deleted if it does not reach

a leaf which is a termination node because that means it is only a partial string. The running time for deletion of a string would be $O(l)$ where l is the length of the string to be deleted.

Chapter 15, section 5

- **R-15.4**

For what values of d is the tree T of the previous exercise an order- d B-tree?

The value for an order- d B-tree is 8. A B-Tree of order- d is an (a,b) with $a = \lceil d/2 \rceil$ and $b = d$. For the previous question the (a,b) tree was equal to $a = 4$ and $b = 8$. Therefore $d = 8$ for the order- d tree.

- **R-15.8**

Draw the result of inserting, into an initially empty order-7 B-tree, entries with keys $(4,40,23,50,11,34,62,78,66,22,90,59,25,72,64,77,39,12)$, in this order.

