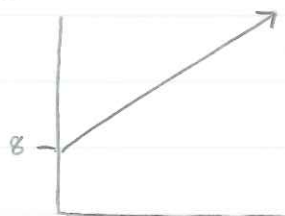


HW 2

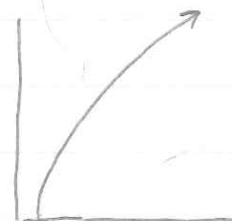
Nick Palumbo

3.1

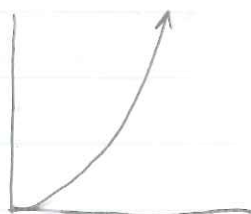
$8n$



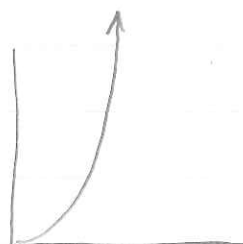
$4n \log n$



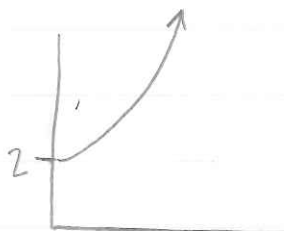
$2n^2$



n^3



2^n



3.2

$$A = 8n \log n$$

$$8n \log n = 2n^2$$

$$B = 2n^2$$

$$4n \log n = n^2$$

$$4 \log n = n$$

$$4 = n / \log n$$

$$16 / \log 216 = 16 / 4 = 4$$

$$n_0 = 17$$

3.8

$$2^{10}, 2^{\log n}, 3n + 100 \log n, 4n, n \log n, 4n \log n + 2n, n^2 + 10n, n^2 + 10n, n^3, 2^n$$

3.10

If $d(n)$ is $O(f(n))$, then

$$d(n) \leq c_1 f(n) \quad n \geq n_1$$

If $e(n)$ is $O(g(n))$, then

$$e(n) \leq (2g(n)) \quad n \geq n_2$$

$$\begin{aligned}\text{Thus } d(n)e(n) &\leq c_1 f(n)e(n) \\ c_1 f(n)e(n) &\leq c_1 f(n)c_2 g(n)\end{aligned}$$

$$\begin{aligned}\text{Thus } d(n)e(n) &\leq c_1 c_2 f(n)g(n) \\ d(n)e(n) &\text{ is } O(f(n) \cdot g(n)) \quad n \geq \max[n_1, n_2]\end{aligned}$$

$$\begin{aligned}3.14 \quad 2 \max(f(n), g(n)) &\geq (g(n) + f(n)) \\ \max(f(n), g(n)) &\geq \frac{1}{2} (g(n) + f(n)) \\ n &\geq n_0\end{aligned}$$

$$\begin{aligned}\text{Thus } \max\{f(n), g(n)\} &= O(g(n) + f(n)) \\ \text{where } c &= \frac{1}{2}\end{aligned}$$

$$\begin{aligned}3.17 \quad (n+1)^5 &= n^5 + 5n^4 + n^3 + 10n^2 + 5n + 1 \\ n^5 + 5n^4 + 0n^3 + 10n^2 + 5n + 1n^0 &\leq n^5 + 5n^5 + 0n^5 + 10n^5 \\ &\quad + 5n^5 + 1n^5\end{aligned}$$

$$\begin{aligned}\text{Thus } 1n^5 + 5n^4 + 10n^2 + 5n + 1 &\leq 22n^5 \\ c = 22 \quad \text{for any } n &\geq 1\end{aligned}$$

$$\begin{aligned}3.18 \quad 2^{n+1} &\in O(2^n) \\ 2^{n+1} &= 2^n \cdot 2^1 = 2 \cdot 2^n \\ \text{Thus } c &= 2 \quad n_0 = 1\end{aligned}$$

3.20 show that n^2 is $\Omega(n \log n)$

$$n^2 \in \Omega(n \log n)$$

$$n \quad n^2 \quad n \log n \quad n_0 = 1$$

1 1 0 $C=1$ Therefore n^2 is $(n \log n)$

2 4 2

3 9 4.74

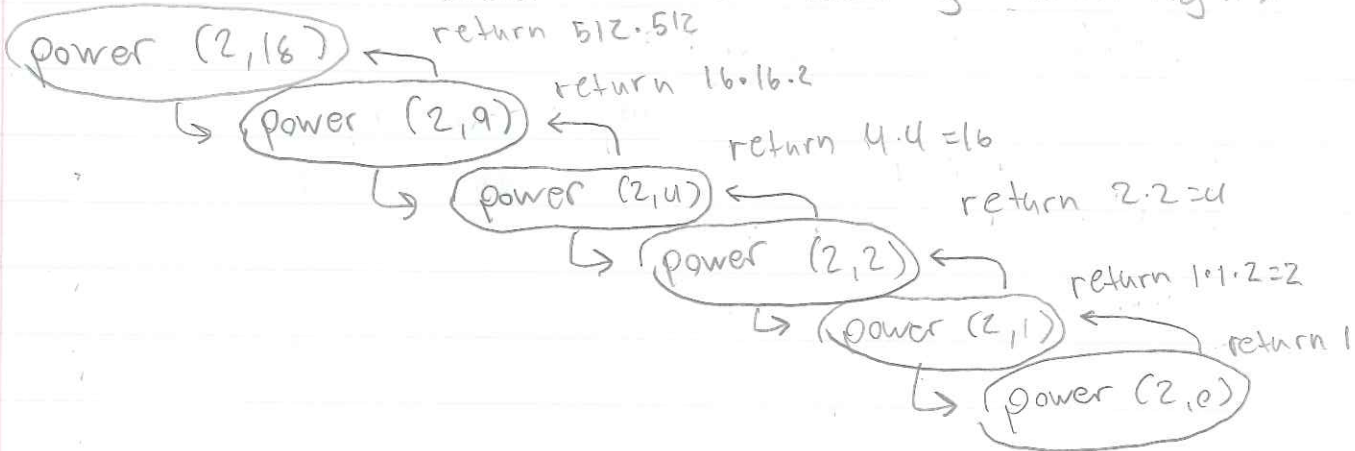
4 16 8

5 32 11.5

3.27 There are two nested loops in an outer loop so $O(n^3)$

3.35 Combine all 3 sets and sort. That takes $O(n \log n)$. Then traverse the list and find 3 matching elements consecutively. That's $O(n)$. n traversal would be canceled because it is a lower order term, leaving $O(n \log n)$.

4.3



```

4.7 def strToNum(strNum):
    if len(strNum) == 1:
        return int(strNum)
    else:
        new_digit = int(num[len(num)-1])
        return new_digit + 10 * (strToNum(strNum[:len(num)-1]))

```

The base case is a string of length 1.
 This should be the first number of
 the string (1 3531).

The recursive case casts the last number
 of the string. It returns the number plus 10×4
 recursive call on the string except the last
 number.

$$\begin{aligned}
 &1 + 10 \cdot (\text{strNum}(1353)) \\
 &1 + 10 (3 + 10 (\text{strNum}(135))) \\
 &1 + 10 (3 + 10 (5 + 10 (\text{strNum}(1311))) \\
 &1 + 10 (3 + 10 (5 + 10 (3 + \text{strNum}(\text{strNum}(1)))))) \\
 &1 + 10 (3 + 10 (5 + 10 (3 + 10 \cdot 1))) \\
 &13531
 \end{aligned}$$