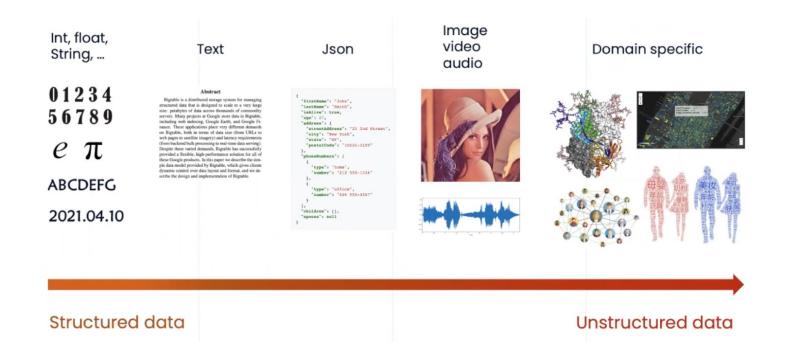
# Milvus: A Purpose-Built Vector Data Management System

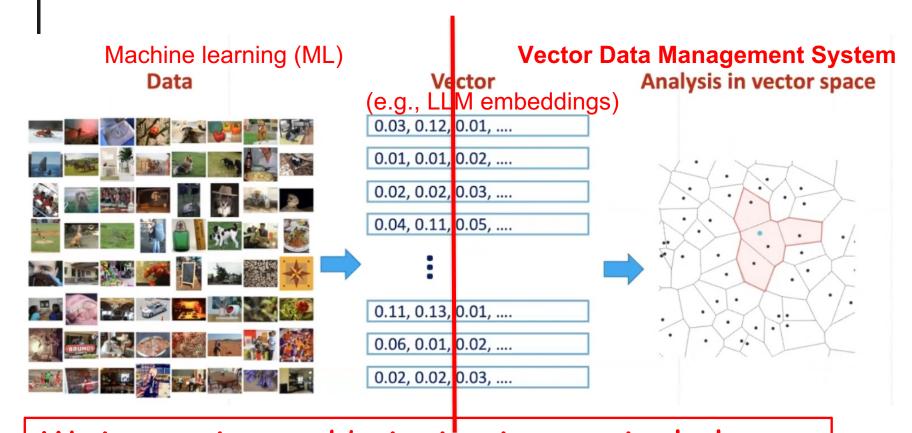
Jianguo Wang etc.
SIGMOD 2021

# Background: Unstructured Data are Dominating



IDC: 80% data will be unstructured by 2025

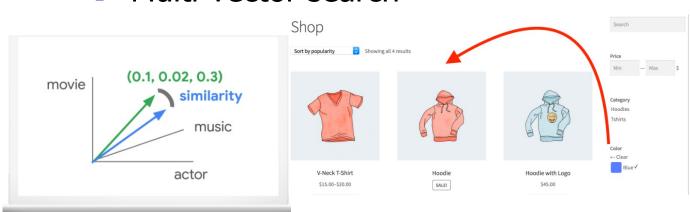
# Background: Vector Analysis Work Flow

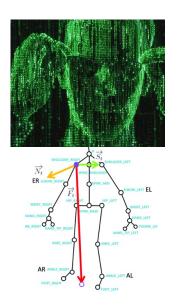


ML is soaring and bringing increasingly large volume of vectors, and how about the Vector Data Management System?

#### Requirements and Challenges for Scalable Vector Data Management System

- Fast query processing and dynamic vector manipulation
- Advanced query semantics:
  - Vector similarity search
  - Attribute filtering
  - Multi-vector search





However, existing work fail for poor performance and limited functionalities...

#### **Limitations of Existing Work**

	Billion-Scale Data	Dynamic Data	GPU	Attribute Filtering	Multi-Vector Query	Distributed System
Facebook Faiss [3, 35]	✓	Х	✓	Х	Х	Х
Microsoft SPTAG [14]	✓	X	X	X	X	X
ElasticSearch [2]	×	✓	X	✓	×	✓
Jingdong Vearch [4, 39]	×	✓	✓	✓	×	✓
Alibaba AnalyticDB-V [65]	✓	✓	X	✓	×	✓
Alibaba PASE (PostgreSQL) [68]	×	✓	X	✓	×	X
Milvus (this paper)	✓	✓	✓	✓	✓	✓

- Existing algorithms (e.g., [14]) don't support large data storage at scale or dynamic data with fast updates.
- Existing systems (e.g., [3]) simply extend relational DB and fail to fully optimize query processing on vector data

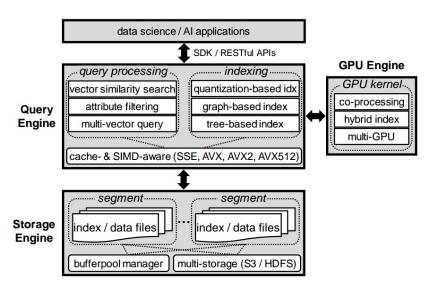
Both lack the support for more advanced query processing beyond simple vector similarity search.



Yes, they can.



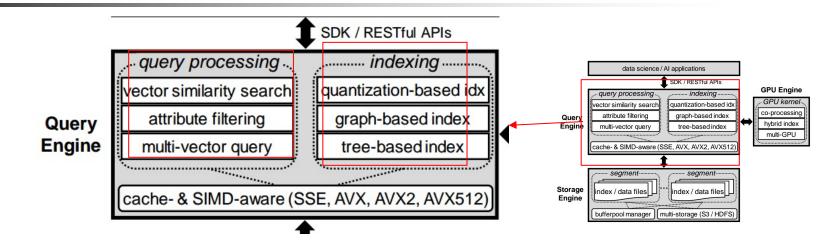
#### Milvus Overview



#### All-in-one solution

- Functional complete
   system + optimized algos
- Hardware conscious design on both CPU and GPU
- EasySDK in Python/Java/Go/C++ and RESTful APIs
- Three major components
  - Query Engine
  - GPU Engine
  - Storage Engine

#### The Query Engine

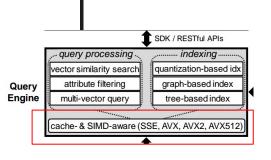


- Support vector(s) and vector(s)+non-vector attributes, therefore 3 primitive advanced queries\*
- Cover common similarity and index functions
- Adopts LSM-tree [47] for efficient insertions and deletions
  - First in MemTable
  - Next flush into disk by rules (e.g., by size or time)

\*Recall: Existing vectorDB only support similarity search

47: Chen Luo and Michael J. Carey. LSM-based Storage Techniques: A Survey.

## The Query Engine



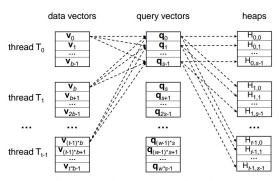
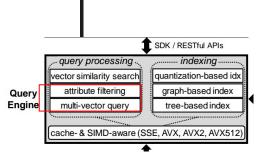


Figure 3: Cache-aware design in Milvus

- Cache-aware and SIMD optimizations
  - Cache-aware
    - Reuse accessed data for L3 opt.
      - Hence reduce L3 miss
    - Fine-grained parallelism of coredatablock assigns
      - Outperform OpenMP
  - SIMD
    - AVX512 support and automatic instruction choose in <u>runtime</u>

## The Query Engine



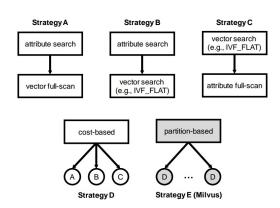
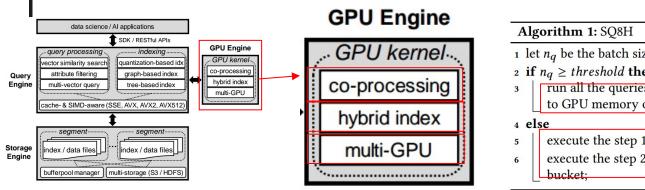


Figure 4: Different strategies for attribute filtering

- Further opt. on <u>attribute filtering</u> and <u>multi-vector query\*</u>
  - Four adopted + one newly proposed attribute filtering
    - Partition-based is up to 13.7x faster than existing algos
  - Two new approaches instead of naïve solution for multivector query
    - i.e., vector fusion and iterative merging

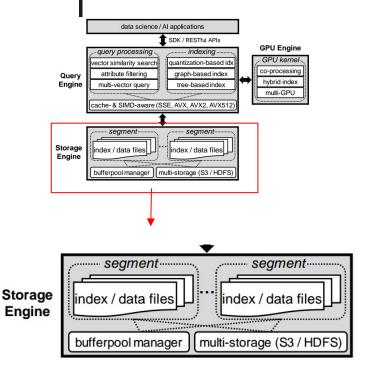
#### The GPU Engine



# Algorithm 1: SQ8H 1 let $n_q$ be the batch size; 2 if $n_q \ge threshold$ then 3 run all the queries entirely in GPU (load multiple buckets to GPU memory on the fly); 4 else 5 execute the step 1 of SQ8 in GPU: finding $n_{probe}$ buckets; 6 execute the step 2 of SQ8 in CPU: scanning every relevant

- Optimized SQ8 algo [3] for CPU-GPU co-processing
  - Communication opt. by multi-bucket copying
  - Computation opt. by fine-grained decomposition
- Optimized GPU kernel to supported larger query within GPU memory
  - <u>Progressively, round-by-round</u> flushing results, instead of loading the entire and flushing only once
- Multi-GPU support by adjusting device number in runtime

#### The Storage Engine



- Support various storage for different data structures and storage hardware
  - Vector storage (mem)
    - Continuous row for single vector
    - Columnar fashion for multiple vectors
  - Attribute storage (mem)
    - <key,value> pairs for each attribute column
    - Further optimize by building skip pointers[16]
  - Buffer pool to bridge mem and disk
  - Various file system support for in-disk and in-network

#### **Implementation Optimizations**

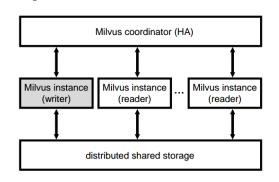
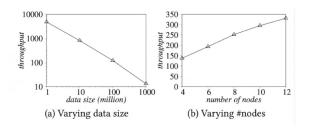


Figure 5: Milvus distributed system



- Asynchronous processing
  - To minimize foreground processing (e.g., receiving request)
  - Write-ahead logs + operating at back gound
- Snapshot isolation
  - To make sure of the R/W consistency
  - Using LSM structure[47]
- Distributed system
  - For scalability and availability
  - Shared-storage and separated compute/storage
  - Stateless computing instances
    - Single writer multiple reader as Milvus is <u>read-heavy</u>

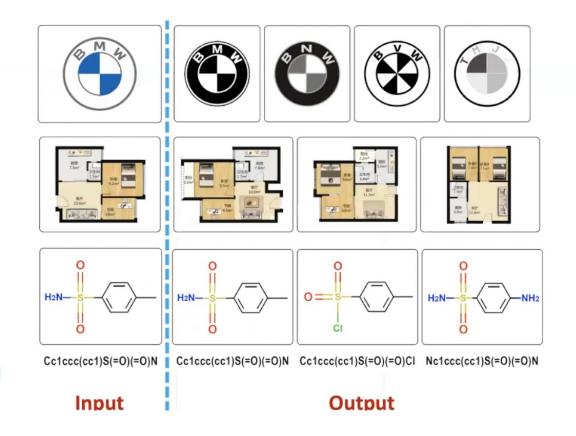
#### **Application Examples**

#### Similar trademark search

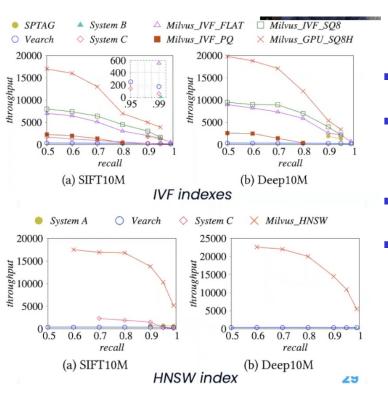
#### Floor plan search

#### Chemical Structure Search

Video search, chatbots, recommendation engine, biological multi-factor authentication...



## Outperform SOTA Significantly



Platform: Alibaba cloud

Dataset: SIFT1B and Deep1B

Both are 1 billion, 128-dimension

Open source SOTA: SPTAG, Vearch

Commercial system: System A, B, C

#### Conclusion and Future Work

 Milvus satisfies the requirements of versatile applications by heuristic optimization in system and algorithm

- SIMD support
- Cache-aware optimization
- GPU accelerations
   LSM structure
- Snapshot isolation
   Distributed processing
- Hybrid search algorithm
- Cost base optimization
- Vector fusion
- Iterative merging
- Future work can be following:
  - Cloud native architecture
  - FPGA acceleration

#### What Can We Learn and Do?

- This is an extremely complete system research.
  - Inclusively enough for wide community use System:
  - SIMD support
- Cache-aware optimization
- GPU accelerations
   LSM structure

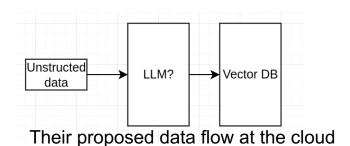
- Snapshot isolation
   Distributed processing

#### Algorithm:

- Hybrid search algorithm
- Cost base optimization
- Vector fusion
- Iterative merging

But it is not a good idea to apply Milvus on Edges...

#### What if On Edge?



People flow in a building

Control the aircons on-demand

Legend Data Data source sink

Nector DB LLM?

Control the aircons energy constrained edge device

One possible edge application

- Our toy edge application
  - <u>Continuously</u> tracing people flow and control the aircons on demand
  - VectorDB bridges two subsystems on a device
- Problem 1: The indexing and storage overhead is still too heavy for constrained edges
  - Is there any cure by using approximate computation, or even new indexing ways?
- Problem 2: Two subsystems may raise high-speed and concurrent write (streams)
  - Will stateless, single writer still work well?