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AlphaDeepChess: motor de ajedrez basado en  
podas alpha-beta  
AlphaDeepChess: chess engine based on  
alpha-beta pruning

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Trabajo de Fin de Grado  
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Grado en Ingeniería de Computadores

Grado en Desarrollo de Videojuegos

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# Dedication

*To our younger selves, for knowing the art of  
chess*



# Acknowledgments

To our family members for their support and for taking us to chess tournaments to compete.





# Abstract

## AlphaDeepChess: chess engine based on alpha-beta pruning

Chess engines have played a fundamental role in the advancement of artificial intelligence applied to chess since the mid-20th century. Pioneers such as Alan Turing and Claude Shannon established the theoretical principles that laid the foundation for this field. Building upon these foundations, the evolution of hardware and the refinement of search techniques have enabled significant advancements, such as alpha-beta pruning, an optimization of the minimax algorithm that drastically reduces the number of nodes evaluated in the game tree. Today, Stockfish, the most powerful and open-source chess engine, continues to rely on alpha-beta pruning but also incorporates deep learning techniques and neural networks.

The goal of this project is to develop a chess engine capable of competing against both other engines and human players, using alpha-beta pruning as its core. Additionally, we will analyze the impact of other classical algorithmic techniques such as transposition tables, iterative deepening, and a move generator based on magic bitboards.

The chess engine has finally been uploaded to Lichess platform, where AlphaDeepChess achieved an ELO rating of 1900 while running on a Raspberry Pi 5 equipped with a 2TB transposition table.

## Keywords

chess, chess engine, alpha-beta pruning, iterative deepening, quiescence search, move ordering, transposition table, zobrist hashing, pext instruction, magic bitboards



# Resumen

## AlphaDeepChess: motor de ajedrez basado en podas alpha-beta

Los motores de ajedrez han desempeñado un papel fundamental en el avance de la inteligencia artificial aplicada al ajedrez desde mediados del siglo XX. Pioneros como Alan Turing y Claude Shannon establecieron los principios teóricos que sentaron las bases de este campo. Sobre estos cimientos, la evolución del hardware y el perfeccionamiento de técnicas de búsqueda permitieron importantes avances, como la poda alfa-beta, una optimización del algoritmo minimax que reduce drásticamente el número de nodos evaluados en el árbol de juego. Hoy en día, Stockfish, el motor de ajedrez más potente y de código abierto, sigue basándose en técnicas algorítmicas clásicas, pero también incorpora deep learning y redes neuronales.

El objetivo de este proyecto es desarrollar un motor de ajedrez capaz de competir tanto contra otros motores como contra jugadores humanos, utilizando la poda alfa-beta como núcleo del algoritmo. Además, se analizará el impacto de otras técnicas algorítmicas clásicas, como las tablas de transposición, la búsqueda en profundidad iterativa y un generador de movimientos basado en bitboards mágicos.

Finalmente, el motor de ajedrez ha sido subido a la plataforma de Lichess, donde AlphaDeepChess ha alcanzado una puntuación ELO de 1900, ejecutándose en una Raspberry Pi 5 con una tabla de transposiciones de 2TB.

## Palabras clave

ajedrez, motor de ajedrez, poda alfa-beta, búsqueda en profundidad iterativa, búsqueda quiescente, ordenación de movimientos, tabla de transposiciones, zobrist hashing, instrucción pext, bitboards mágicos



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# Chapter 1

## Introduction

*“The most powerful weapon in chess is to have the next move”*

— David Bronstein

Chess, one of the oldest strategy games in human history, has long been a domain for both intellectual competition and computational research. The pursuit of creating a machine that could compete with the best human players, chess Grandmasters, was present. It was only a matter of time before computation surpassed human capabilities.

In 1997, the chess engine **Deep Blue** made history by defeating the reigning world champion at the time, Garry Kasparov, marking the first time a computer had defeated a sitting world chess champion.

Today, we find ourselves in an era where chess engines have reached unprecedented strength. This has been achieved through a combination of classical techniques like alpha-beta pruning, and modern advancements such as deep learning and neural networks.

### 1.1. Objectives

The objectives of this project are the following:

- Develop a chess engine based on alpha-beta pruning that follows the UCI protocol (Stefan-Meyer Kahlen (2004)). The engine will be a console application capable of playing chess against humans or other engines, as well as analyzing and evaluating positions to determine the best legal move.
- Implement various optimization techniques, including move ordering, quiescence search, iterative deepening, transposition tables, multithreading, and a move generator based on magic bitboards.
- Measure the impact of these optimization techniques and profile the engine to identify performance bottlenecks.

- Upload the engine to `lichess.org` and compete against other chess engines.

## 1.2. Work plan

The project will be divided into several phases, each focusing on specific aspects of the engine's development. The timeline for each phase is as follows:

1. Research phase and basic implementation: understand the fundamentals of alpha-beta pruning with minimax and position evaluation. Familiarize with the UCI (Universal Chess Interface) and implement the move generator with its specific exceptions and rules.
2. Optimization: implement quiescence search and iterative deepening to improve pruning effectiveness.
3. Optimization: improve search efficiency using transposition tables and Zobrist hashing.
4. Optimization: implement multithreading to enable parallel search.
5. Profiling: use a profiler to identify performance bottlenecks and optimize critical sections of the code.
6. Testing: use Stockfish to compare efficiency generating tournaments between chess engines and estimate the performance of the engine. Also, compare different versions of the engine to evaluate the impact of optimizations.
7. Analyze the results and write the final report.

In the following Section 1.3, we will talk about the basic concepts of chess, but if you already have the knowledge we recommend you to advance directly to the next chapter 2.

## 1.3. Basic concepts

Chess is a board game where two players who take white pieces and black pieces respectively compete to first checkmate the opponent. Checkmate occurs when the king is under threat of capture (known as check) by a piece or pieces of the enemy, and there is no legal way to escape or remove the threat.

What about a chess engine? A chess engine consists of a software program that analyzes chess positions and returns optimal moves depending on its configuration. In order to help users to use these engines, chess community agreed on creating an open communication protocol called **Universal Chess Interface** or commonly referred to as UCI, that provides the interaction with chess engines through user interfaces.

A chess game takes place on a chessboard with specific rules governing the movement and interaction of the pieces. This section introduces the fundamental concepts necessary to understand how chess is played.

### 1.3.1. Chessboard

A chessboard is a game board of 64 squares arranged in 8 rows and 8 columns. To refer to each of the squares we mostly use **algebraic notation** using the numbers from 1 to 8 and the letters from “a” to “h”. There are also other notations like descriptive notation (now obsolete) or ICCF numeric notation due to chess pieces have different abbreviations depending on the language.

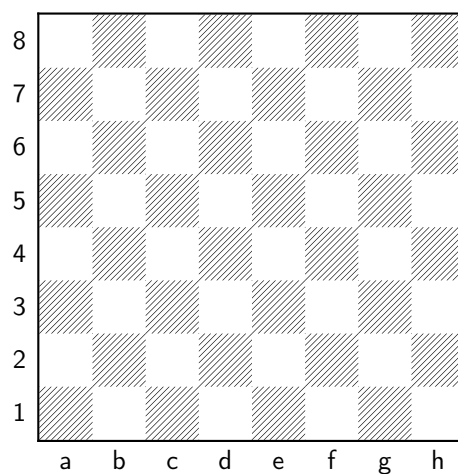


Figure 1.1: Empty chessboard.

For example,  $g5$  refers to the following square:

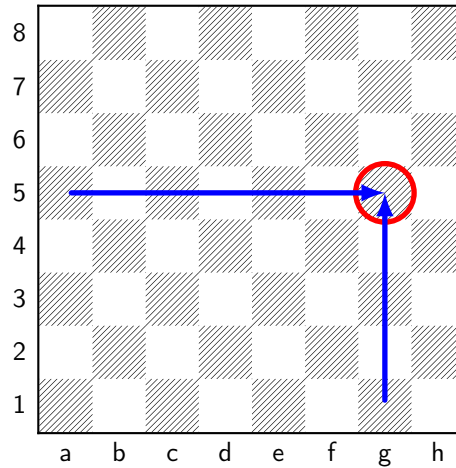


Figure 1.2: Example: square  $g5$  highlighted and arrows pointing to it.

It is important to know that when placing a chessboard in the correct orientation, there should always be a white square in the bottom-right corner or a black square in the bottom-left corner.

### 1.3.2. Chess pieces

There are 6 types of chess pieces: king, queen, rook, bishop, knight and pawn, and each side has 16 pieces:

































Piece	White Pieces	Black Pieces	Number of Pieces
King			1
Queen			1
Rook	 	 	2
Bishop	 	 	2
Knight	 	 	2
Pawn	       	       	8

Table 1.1: Number of chess pieces by type and color.

The starting position of the chess pieces on a chessboard is the following:

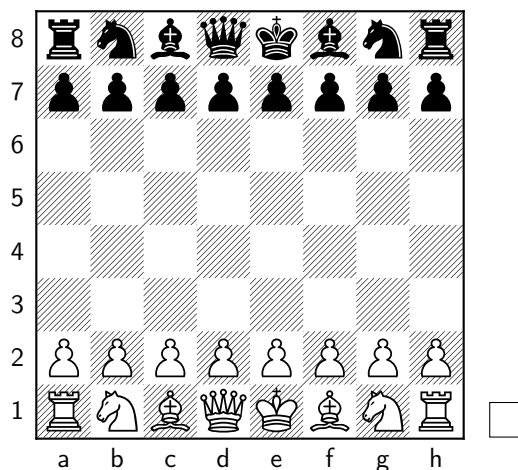


Figure 1.3: Starting position.

The smaller white square next to the board indicates which side is to move in the current position. If the square is white, it means it is white's turn to move; if the square is black, it means it is black's turn to move. This visual indicator helps clarify which player has the next move in the game. Notice that the queen and king are placed in the center columns. The queen is placed on a square of its color, while the king is placed on the remaining central column. The rest of the pieces are positioned symmetrically, as shown in Figure 1.3. This means that the chessboard is divided into two sides relative to the positions of the king and queen at the start of the game:

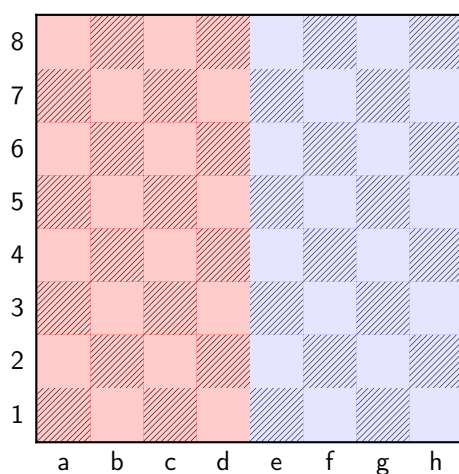


Figure 1.4: King's side (blue) and Queen's side (red).

### 1.3.3. Movement of the pieces

#### 1.3.3.1. Pawn

The pawn can move one square forward, but it can only capture pieces one square diagonally. On its first move, the pawn has the option to move two squares forward. If a pawn reaches the last row of the opponent's side, it promotes to any other piece (except for a king). Promotion is a term to indicate the mandatory replacement of a pawn with another piece, usually providing a significant advantage to the player who promotes.

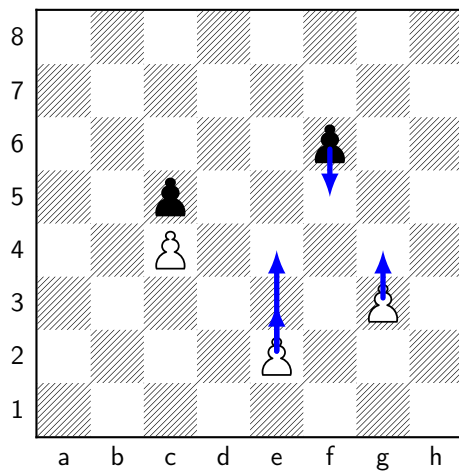


Figure 1.5: Pawn's movement.

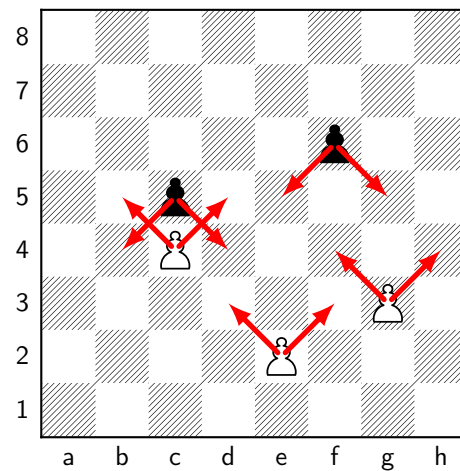


Figure 1.6: Pawn attack.

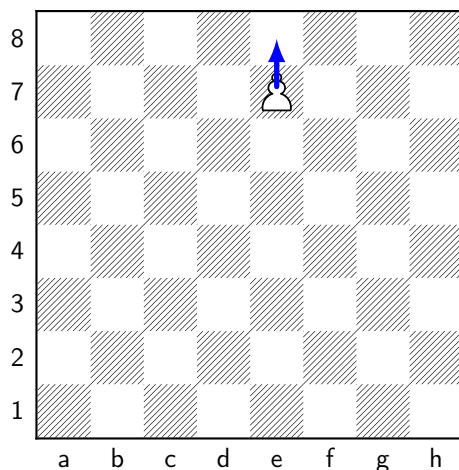


Figure 1.7: Promotion.

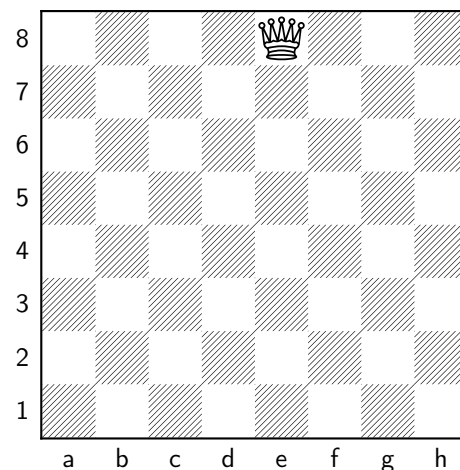


Figure 1.8: Pawn promotes to queen.

There is a specific capture movement which is **en passant**. This move allows a pawn that has moved two squares forward from its starting position to be captured by an



opponent's pawn as if it had only moved one square. The capturing pawn must be on an adjacent file and can only capture the en passant pawn immediately after it moves.

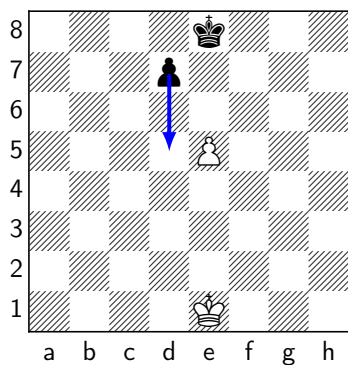


Figure 1.9: En passant (1).

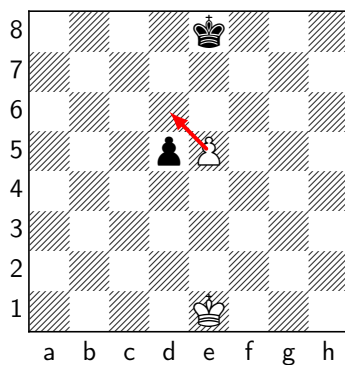


Figure 1.10: En passant (2).

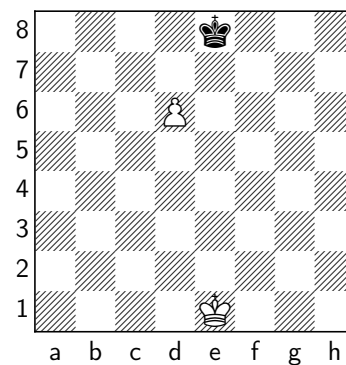


Figure 1.11: En passant (3).

### 1.3.3.2. Rook

The rook can move any number of squares horizontally or vertically. It can also capture pieces in the same way.

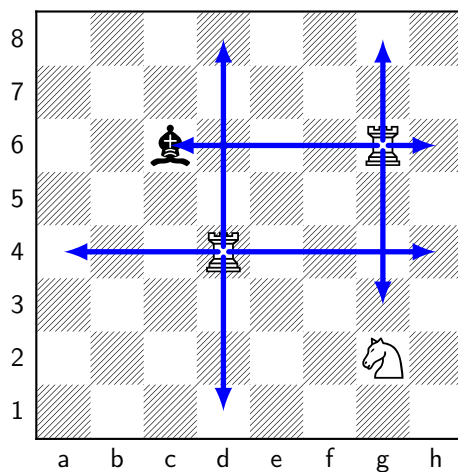


Figure 1.12: Rook's movement.

### 1.3.3.3. Knight

The knight moves in an L-shape: two squares in one direction and then one square perpendicular to that direction. The knight can jump over other pieces, making it a unique piece in terms of movement. It can also capture pieces in the same way.

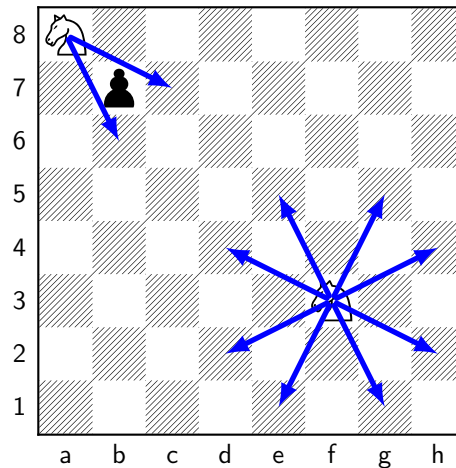


Figure 1.13: Knight's movement.

#### 1.3.3.4. Bishop

The bishop can move any number of squares diagonally. It can also capture pieces in the same way. Considering that each side has two bishops, one bishop moves on light squares and the other on dark squares.

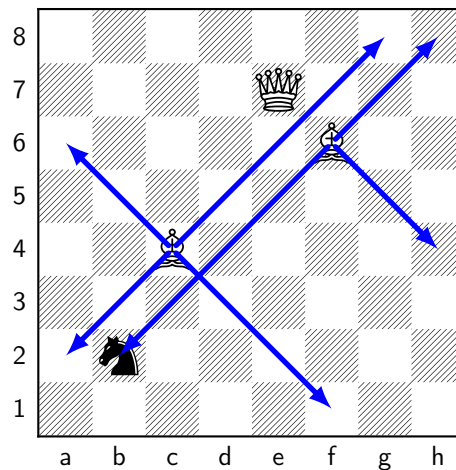


Figure 1.14: Bishop's movement.

#### 1.3.3.5. King

The king can move one square in any direction: horizontally, vertically, or diagonally. However, the king cannot move to a square that is under attack by an opponent's piece. The king can also capture pieces in the same way. The king is a crucial piece in chess, as the game ends when one player checkmates the opponent's king.

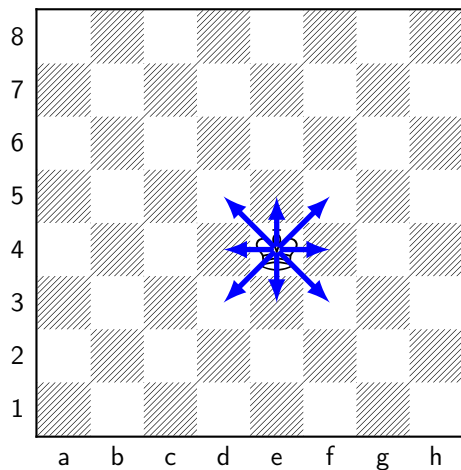


Figure 1.15: King's movement.

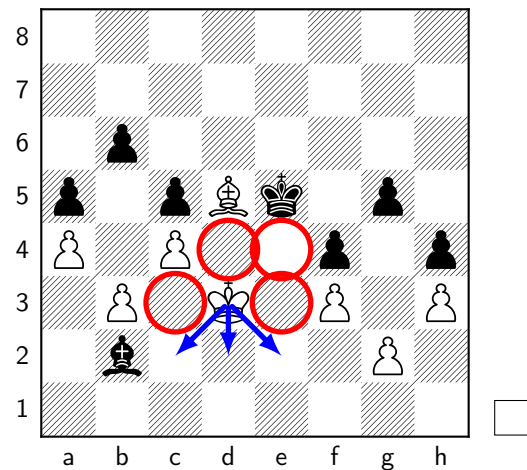


Figure 1.16: White King's movement in a game.

In Figure 1.16, the white king cannot move to  $e4$  because the black king is attacking that square.

**Castling** is a special move which involves moving the king two squares towards a rook and moving the rook to the square next to the king. Castling has specific conditions which are:

- Neither the king nor the rook involved in castling must have moved previously.
- There must be no pieces between the king and the rook.
- The king cannot be in check, move through a square under attack, or end up in check.

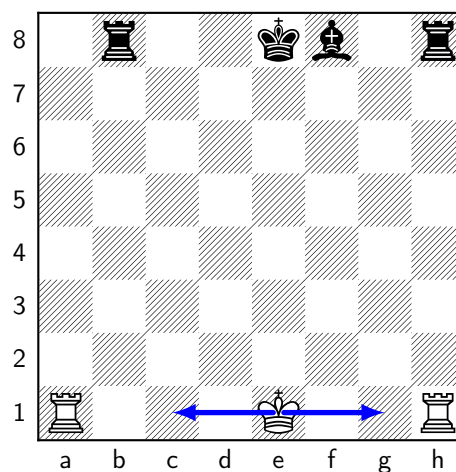


Figure 1.17: Castling

In Figure 1.17, the white king can castle on either the king's side or the queen's side as long as the rooks have not been moved from their starting position, but the black king cannot castle because there is a bishop on  $f8$  interfering with the movement and the rook on the queen's side has been moved to  $b8$ .

### 1.3.3.6. Queen

The queen can move any number of squares in any direction: horizontally, vertically, or diagonally. It can also capture pieces in the same way.

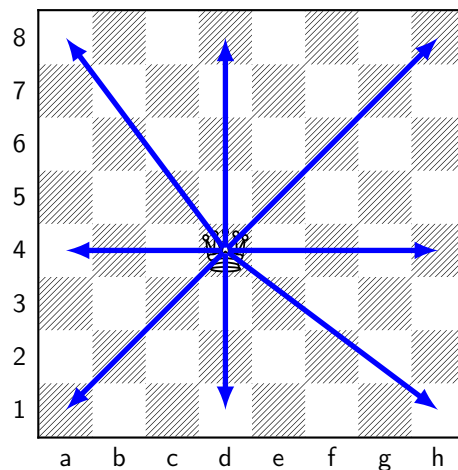


Figure 1.18: Queen's movement.

## 1.3.4. Rules

The rules of chess follow the official regulations established by FIDE Fédération Internationale des Échecs (2023). The objective of each player is to checkmate the opponent's king, meaning the king is under attack and cannot escape.

In every game, white starts first, and the possible results of each game can be win for white, win for black or draw. A draw or tie could be caused by different conditions:

1. Stalemate: the player whose turn it is to move has no legal moves, and their king is not in check.
2. Insufficient material: neither player has enough pieces to checkmate. Those cases are king vs king, king and bishop vs king, king and knight vs king, and king and bishop vs king and bishop with the bishops on the same color.
3. Threefold repetition: it occurs when same position happens three times during the game, with the same player to move and the same possible moves (including castling and en passant).
4. Fifty-move rule: if 50 consecutive moves are made by both players without a pawn move or a capture, the game can be declared a draw.

5. Mutual agreement: both players can agree to a draw at any point during the game.
6. Dead position: a position where no legal moves can be made, and the game cannot continue. This includes cases like king vs king, king and knight vs king, or king and bishop vs king.

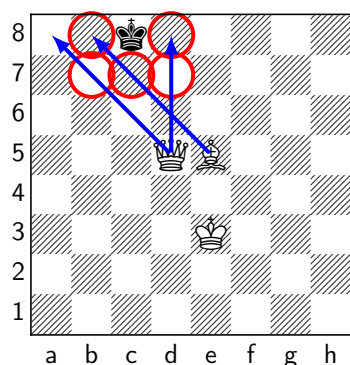


Figure 1.19: Stalemate.

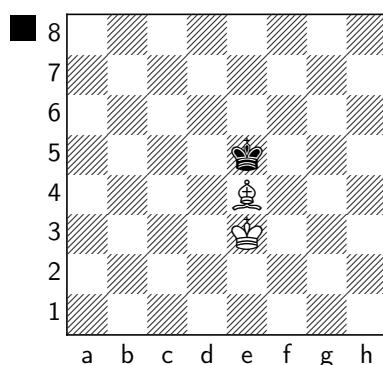


Figure 1.20: Insufficient material.

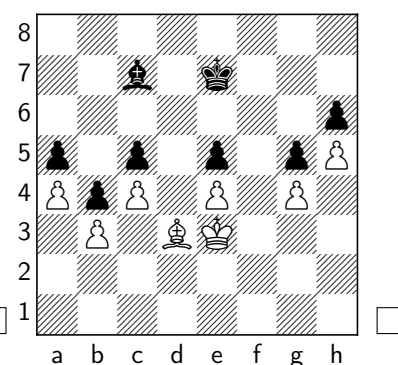


Figure 1.21: Dead position.

Players can also resign at any time, conceding victory to the opponent. Also, if a player runs out of time in a timed game, they lose unless the opponent does not have enough material to checkmate, in which case the game is drawn.

### 1.3.5. Notation

Notation is important in chess to record moves and analyze games.

#### 1.3.5.1. Algebraic notation

In addition to the algebraic notation of the squares in Section 1.3.1, each piece is identified by an uppercase letter, which may vary across different languages:

Piece	English Notation	Spanish Notation
Pawn	<i>P</i>	<i>P</i> (peón)
Rook	<i>R</i>	<i>T</i> (torre)
Knight	<i>N</i>	<i>C</i> (caballo)
Bishop	<i>B</i>	<i>A</i> (alfil)
Queen	<i>Q</i>	<i>D</i> (dama)
King	<i>K</i>	<i>R</i> (rey)

Table 1.2: Chess piece notation in English and Spanish.

**Normal moves (not captures nor promoting)** are written using the piece uppercase letter plus the coordinate of destination. In the case of pawns, it can be written only with the coordinate of destination:

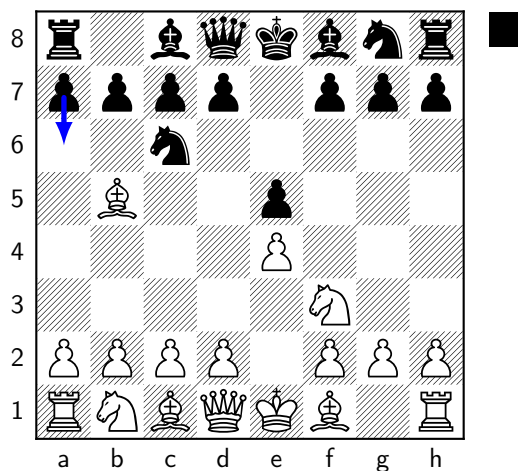


Figure 1.22: Pawn goes to a6.

In Figure 1.22, the pawn's movement is written as *Pa6* or directly as *a6*.

**Captures** are written with an "x" between the piece uppercase letter and coordinate of destination or the captured piece coordinate. In the case of pawns, it can be written with the column letter of the pawn that captures the piece. Also, if two pieces of the same type can capture the same piece, the piece's column or row letter is added to indicate which piece is moving:

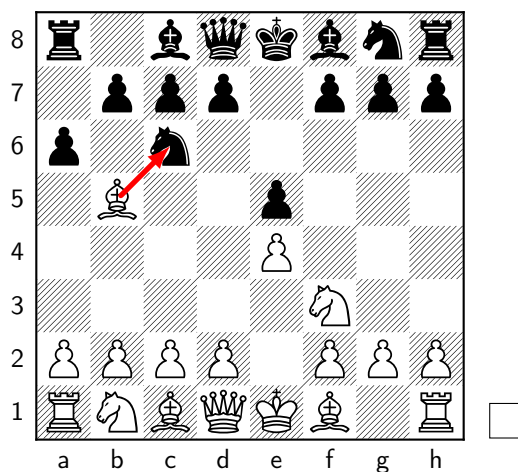


Figure 1.23: Bishop captures knight.

In Figure 1.23, the white bishop capturing the black knight is written as *Bxc6*. If it

were black's turn, the pawn on *a6* could capture the white bishop, and it would be written as *Pxb5* or simply *axb5*, indicating the pawn's column.

**Pawn promotion** is written as the pawn's movement to the last row, followed by the piece to which it is promoted:

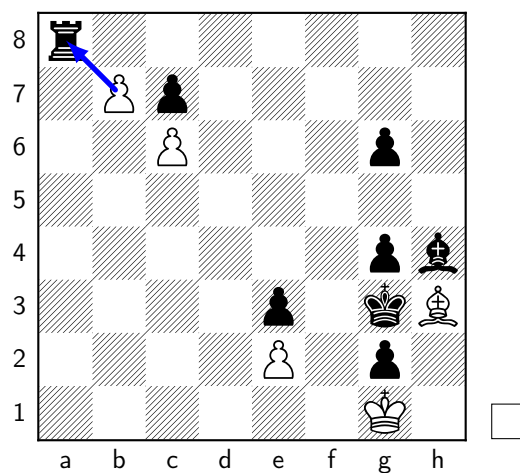


Figure 1.24: Pawn captures rook.

In Figure 1.24, white pawn capturing and promoting in *a8* to a queen is written as *bx a8Q* or *bx a8 = Q*.

**Castling** depending on whether it is on the king's side or the queen's side, it is written as *0 - 0* and *0 - 0 - 0*, respectively.

**Check and checkmate** are written by adding a *+* sign for check or *++* for checkmate, respectively.

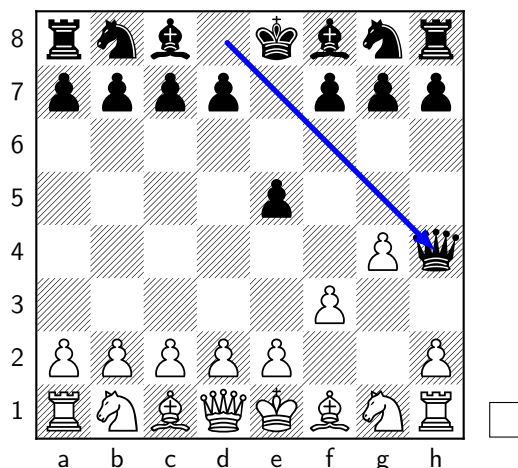


Figure 1.25: Black queen checkmates.

In Figure 1.25, black queen movement checkmates and it is written as  $Dh4++$ .

**The end of game notation** indicates the result of the game. It is typically written as:

- $1-0$ : White wins.
- $0-1$ : Black wins.
- $1/2-1/2$ : The game ends in a draw.

### 1.3.5.2. Forsyth–Edwards Notation (FEN)

This is a notation that describes a specific position on a chessboard. It includes 6 fields separated by spaces: the piece placement, whose turn it is to move, castling availability, en passant target square, halfmove clock, and fullmove number. For example, the FEN for the starting position is:

```
rnbqkbnr/pppppppp/8/8/8/8/PPPPPPPP/RNBQKBNR w KQkq - 0 1
```

Keep in mind this notation is important for the engine to understand the position of the pieces on the board.

### 1.3.5.3. Portable Game Notation (PGN)

This notation is mostly used for keeping information about the game and a header section with metadata: the name of the event, site, date of play, round, color and name of each player, and result. For example, the PGN for a game could look like this:



Listing 1.1: Example of a PGN file

```
[Event "XX_Gran_Torneo_Internacional_Aficionado"]
[Site "?"]
[Date "2024.06.25"]
[Round "1"]
[White "Tejedor_Barber,_Lorenzo"]
[Black "Giron_Herranz,_Juan"]
[Result "0-1"]
1. e4 c6 2. d4 d5 3. exd5 cxd5 4. Bd3 Nc6 5. c3 Nf6 6. Bf4 Bg4
   7. Qb3 Qd7 8. h3 Bh5 9. Nd2 e6 10. Ngf3 a6 11. O-O Be7 12.
   Rfe1 O-O 13. Re3 b5 14. Ne5 Qb7 15. Nxc6 Qxc6 16. Nf3 Nd7
   17. a3 Bg6 18. Bxg6 hxg6 19. Rae1 a5 20. Qd1 b4 21. axb4
   axb4 22. h4 bxc3 23. bxc3 Ra3 24. Qd3 Rc8 25. Rc1 Bb4 26.
   h5 gxh5 27. Ng5 Nf6 28. Be5 Rxc3 29. Rxc3 Qxc3 30. Qe2 Qc1+
   31. Kh2 Bd2 32. Bxf6 Bxe3 33. fxe3 gxf6 34. Nh3 Qb1 35.
   Nf4 Rc1 36. Nxh5 Rh1+ 37. Kg3
```



# Chapter 2

## State of the art

In this chapter, we will explore the fundamental concepts and techniques on which our chess engine is based. This includes board representation, move generation, game trees, etc. Each section will provide an overview of the concepts and techniques used by our engine, and additional tools.

### 2.1. Board representation

The chessboard is where the game takes place and which serves as the foundation for all operations. In order to store a position setting with its pieces and other additional information like the side to move (the colour that has the turn of movement), the castling rights (the possibilities for each side to castle both short and long, explained in Section 1.3.5.1) or the fifty-move rule counter (explained in Section 1.3.4 Item 4), we can encounter different types of representations: piece centric, square centric and hybrid solutions. We chose to use bitboards as the primary representation, complemented by a piece list to store the piece on each square (piece centric representation). Additionally, the game state is stored in a bit field.

**Bitboard** also known as a bitset or bitmap, is a 64-bit data structure that efficiently represents a chessboard.

Bitboards can be used to represent different aspects of the board:

- All pieces: a bit is set to 1 for every square occupied by a piece, regardless of its type or color.
- Pieces by color: a bit is set to 1 for every square occupied by a piece of a specific color.
- Specific piece types: a bit is set to 1 for every square occupied by a specific type of piece.

This structure is highly efficient for operations such as move generation and attack detection, as bitwise operations (e.g., AND, OR, XOR, shifts) can be used to manipulate and query the board state quickly.

**Bit field** in contrast to a bitboard, uses a fixed number of bits within an integer to store multiple small values or flags.

For example, using a 64-bit bitfield, we can allocate 6 bits to store the number of pieces on the board, as  $2^6 = 64$  possibilities. This means the bits would occupy an interval from  $X$  to  $X + 5$  inclusive.

## 2.2. Move generation

An essential part of any chess engine is move generation. It involves generating all possible legal moves from a given position ensuring chess rules.

There are two types of move generation:

- Pseudolegal move generation: generates all moves without considering whether the king is left in check after the move. It requires additional filtering to remove illegal moves which is more computationally expensive than legal move generation.
- Legal move generation: generates only moves that are valid according to the chess rules, ensuring that the king is not left in check. The more accurate, the computationally more expensive it is.

We have preferred to use legal move generation because it ensures that the generated moves are correct.

Additionally, we have chosen to implement **magic bitboards** for this type of move generator, particularly for sliding pieces such as rooks, bishops, and queens. Magic bitboards use precomputed attack tables and bitwise operations. This approach significantly reduces the computational cost of move generation, enabling the engine to explore deeper levels of the game tree while maintaining accuracy and performance.

## 2.3. Game trees

Sequential games, such as chess or tic-tac-toe, where players take turns alternately, unlike simultaneous games, can be represented in a game tree or graph. In this representation, the root node is the main position from which we look for the best move, and each subsequent node is a possible option or game state, forming a tree-like structure. This tree has a height or depth that refers to the number of levels or layers in the tree, starting from the root node (the initial game state) and extending to the leaf nodes.

The depth of a chess game tree is important because it determines the extent to which it will be analysed and evaluated. A depth of 1 represents all possible moves for the current player or side to move, while a depth of 2 includes the opponent's responses to those moves. As the depth increases, the tree grows exponentially, making it computationally expensive to explore all possible states.

## 2.4. Search algorithms

There are different approaches to find the best move from a position. Some of these search algorithms are: Depth-First Search (DFS), Best-First Search (not to be confused with Breadth-First Search or BFS but they are related) and Parallel Search.

Note that these search algorithms are the foundation of more advanced and practical algorithms used today. However, explaining them is essential to understand the underlying principles.

**Depth-First Search** refers to the process of exploring each branch of a tree or graph to its deepest level before backtracking. Unfortunately, in chess, this cannot be possible because the number of possible moves grows exponentially with the depth of the search tree, leading to the so-called combinatorial explosion. To address this, depth-first search is often combined with techniques like alpha-beta pruning (discussed below) to reduce the number of nodes evaluated, making the search more efficient while still exploring the tree deeply. The following pseudocode illustrates the working of the DFS algorithm:

Listing 2.1: Pseudocode of the Depth-First Search algorithm.

```

1 Procedure DepthFirstSearch(Graph G, Node v):
2   Mark v as visited
3   For each neighbor w of v in G.adjacentEdges(v):
4     If w is not visited:
5       Recursively call DFS(G, w)

```

DFS visits nodes by marking them as visited (line 2) and recursively explores all adjacent nodes until no unvisited nodes remain (lines 3 to 5). It has a worst-case performance of  $O(|V| + |E|)$  and worst-case space complexity of  $O(|V|)$ , with  $|V|$  = number of nodes and  $|E|$  = number of edges.

**Best-First Search** refers to the way of exploring the most promising nodes first. It is similar to a breadth-first search but prioritizes some nodes before others. They typically require significant memory resources, as they must store a search space (the collection of all potential solutions in search algorithms) that grows exponentially.

Listing 2.2: Pseudocode of the Best-First Search algorithm.

```

1 Procedure BestFirstSearch(Graph G, Node start, Node goal):
2   Create an empty priority queue PQ
3   Add start to PQ with priority 0
4   Mark start as visited
5
6   While PQ is not empty:
7     Node current = Remove the node with the highest
7       priority from PQ
8     If current is the goal:
9       Return the path to the goal

```

```
10 |
11 |     For each neighbor w of current in G.adjacentEdges(
12 |         current):
13 |         If w is not visited:
14 |             Calculate priority for w (e.g., using a
15 |                 heuristic)
16 |             Add w to PQ with the calculated priority
17 |             Mark w as visited
```

In this case, the priority queue contains nodes along with their associated priorities, which are determined by a heuristic function.

**Parallel Search** refers to multithreaded search, a technique used to accelerate search processes by leveraging multiple processors.

Next, we will explore some of the most used search algorithms in chess engines.

### 2.4.1. Minimax algorithm

The **minimax** algorithm is a decision making algorithm that follows DFS principles. It is based on the assumption that both players play optimally, with one player (the maximizer) trying to maximize his score and the other player (the minimizer) trying to minimize his score. It explores the game tree to evaluate all possible moves and determines the best move for the current player.

Listing 2.3: Pseudocode of the Minimax algorithm.

```

1 Procedure Minimax(Node position , Integer depth , Boolean
  maximizingPlayer):
2   If depth == 0 or position is a terminal node:
3     Return the evaluation of the position
4
5   If maximizingPlayer:
6     Integer maxEval = -Infinity
7     For each child of position:
8       Integer eval = Minimax(child , depth - 1 , False)
9       maxEval = max(maxEval, eval)
10    Return maxEval
11  Else: // minimizingPlayer
12    Integer minEval = +Infinity
13    For each child of position:
14      Integer eval = Minimax(child , depth - 1 , True)
15      minEval = min(minEval, eval)
16    Return minEval

```

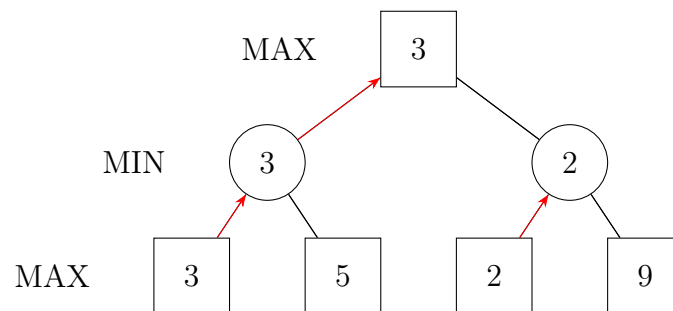


Figure 2.1: Example of minimax.

In this example, white is represented by square nodes and black by circle nodes. Note that this example is a binary tree, but there might be more moves or nodes in a real scenario. Each of them wants to maximize or minimize their respective final value in each position. For the leftmost pair of leaf nodes with values of 3 and 5, 3 is chosen because black tries to get the lowest score between them. Then, the other pair of leaf nodes with values of 2 and 9, 2 is chosen for the same reason. Lastly, at the root node, white selects 3 as the maximum number between 3 and 2.

### 2.4.2. Alpha-beta pruning

Alpha-beta pruning is an optimization of minimax that reduces significantly the number of evaluated nodes in the game tree. It uses two values, alpha and beta, to discard branches that cannot influence the final decision, improving the efficiency.

Listing 2.4: Pseudocode of the Alpha-Beta Pruning algorithm.

```

1 Procedure AlphaBeta(Node position , Integer depth , Integer alpha
  , Integer beta , Boolean maximizingPlayer):
2   If depth == 0 or position is a terminal node:
3     Return the evaluation of the position
4   If maximizingPlayer:
5     Integer maxEval = -Infinity
6     For each child of position:
7       Integer eval = AlphaBeta(child , depth - 1 , alpha ,
        beta , False)
8       maxEval = max(maxEval , eval)
9       alpha = max(alpha , eval)
10      If beta <= alpha:
11        Break // Beta cutoff
12    Return maxEval
13  Else: // minimizingPlayer
14    Integer minEval = +Infinity
15    For each child of position:
16      Integer eval = AlphaBeta(child , depth - 1 , alpha ,
        beta , True)
17      minEval = min(minEval , eval)
18      beta = min(beta , eval)
19      If beta <= alpha:
20        Break // Alpha cutoff
21    Return minEval

```

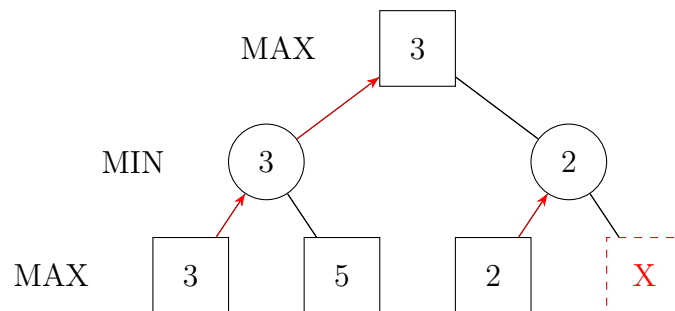
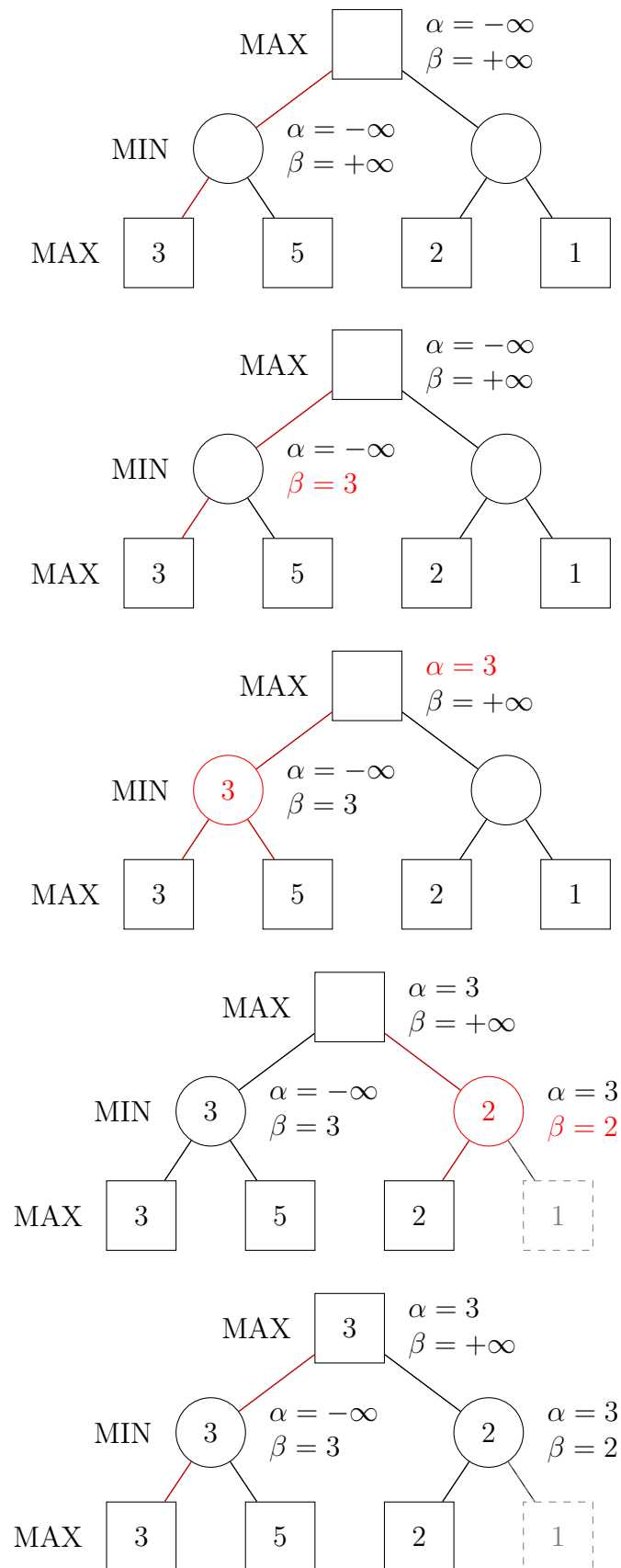


Figure 2.2: Example of alpha-beta pruning.

In this other example, the red dashed node is pruned because it cannot influence the final decision independently of its value. If its value is less than or equal to 2, it will never improve the previously analyzed value of 3. On the other hand, if its value is greater than 2, black will still choose 2 to minimize the score. Another formal way to explain this is by using *alpha* and *beta* values:

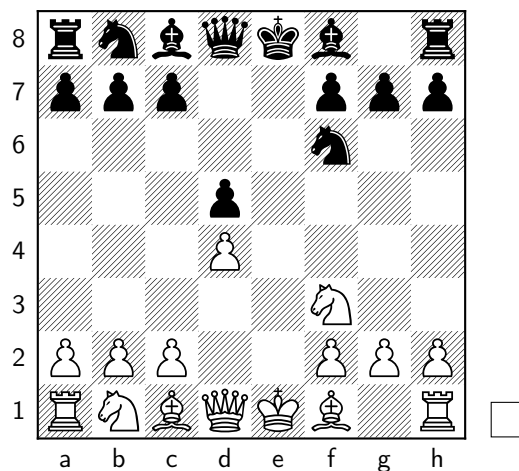


Figure 2.3: Example of alpha-beta pruning with  $\alpha$  and  $\beta$  values.

### 2.4.2.1. Alpha-Beta Enhancements

The alpha-beta algorithm has been further improved over time with various enhancements to increase the overall efficiency. Some of these are: transposition tables, iterative deepening, aspiration windows, quiescence search, move ordering. . .

Take into consideration that many positions can be reached in different ways. This is formally known as transpositions. Just like in dynamic programming, the evaluations of different positions are stored in a structure, the **transposition tables**, to avoid repeating the process of searching and evaluating, which improves efficiency. Take the following example:



French Defense:  
1. e4 e6 2. d4 d5 3. exd5  
exd5 4. Nf3 Nf6

Petrov Defense:  
1. e4 e5 2. Nf3 Nf6 3. Nxe5  
d6 4. Nf3 Nxe4 5. d3 Nf6 6.  
d4 d5

Figure 2.4: Example of transposition.

Both games reach to the same position, although they involve a different number of moves.

In order to store these different positions and access them in a map or dictionary structure, there is a need for a unique and efficient way to index positions: **Zobrist hashing**. Zobrist hashing maps a large number of possible positions to a fixed-size hash value, which can lead to collisions, as different positions may produce the same hash. To handle these collisions, storing additional information like the depth is used to verify the correctness of the entry. In some cases, overwriting the older or less relevant entries can be also useful.

As it is mentioned in <https://www.chessprogramming.org/Search>, «*Depth-first algorithms are generally embedded inside an iterative deepening framework for time control and move ordering issues.*». **Iterative deepening** refers to the combination of DFS with limited depth searches. It performs successive searches by increasing the depth limit at each iteration, allowing you to obtain partial results quickly and improve accuracy over time.

An important concept related to iterative deepening is the use of **aspiration windows**. Their main objective is to reduce the search space by narrowing the search

bounds. In other words, by adjusting *alpha* and *beta* values in each iteration of the iterative deepening. If the value of the evaluation in the iteration falls outside this range or window, a re-search is performed with a wider window to ensure accuracy.

Another critical concept is **quiescence search**. This is a search technique used to address the horizon effect at the end of the search. Simply stopping the search at a fixed or desired depth and evaluating the position can lead to inaccuracies, as critical tactical moves, such as captures, are often overlooked.

*Consider the situation where the last move you consider is QxP. If you stop there and evaluate, you might think that you have won a pawn. But what if you were to search one move deeper and find that the next move is PxQ? You didn't win a pawn, you actually lost a queen. Hence the need to make sure that you are evaluating only quiescent (quiet) positions.<sup>1</sup>*

Finally, alpha-beta algorithm could not perform well without **move ordering**. It is important to ensure that best moves are searched first and to reduce the search space of the game tree. Some of the techniques for move ordering are: Most Valuable victim - Least Valuable Aggressor (MVV-LVA) for captures and killer moves for non-captures.

**MVV-LVA** is a heuristic that prioritizes capturing moves by evaluating the value of the piece being captured (the victim) and the value of the piece performing the capture (the aggressor). The goal is to maximize the gain while minimizing the risk. The following example reflects this:

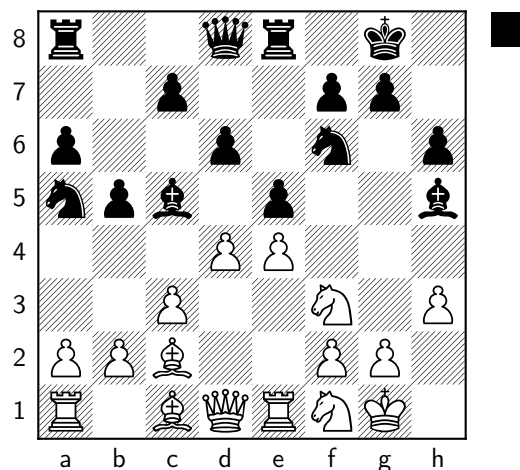


Figure 2.5: Example of MVV-LVA.

In this position, it is black's turn to play after white has moved *d4*, black has the option to capture the pawn on *d4* with the pawn on *e5* or with the bishop on *c5*.

<sup>1</sup>[https://www.chessprogramming.org/Quiescence\\_Search](https://www.chessprogramming.org/Quiescence_Search)

Between the two capturing movements, the best option is to take  $d4$ 's pawn with  $e5$ 's pawn because the pawn has less value than the bishop. Then, after  $exd4$ , white can re-capture with the pawn on  $c3$  or the knight on  $f3$ . If black had taken the pawn with the bishop, white would have won a bishop for a pawn. This simple heuristic that orders what is best in capturing movements for each side can efficiently evaluate tactical exchanges and focus on moves that are more likely to yield a favorable outcome.

**Killer moves** is a heuristic that considers moves that produced cutoffs or pruning while searching. When the engine encounters a similar position at the same depth later in the search, it will prioritize the last move that caused a cutoff, potentially leading to faster pruning.

**Young Brothers Wait Concept** is a parallel search algorithm designed to optimize the distribution of work among multiple threads. This is particularly effective in alpha-beta pruning, where the search tree is explored selectively. It is divided into two phases: the principal variation move and the wait concept. The principal variation is searched sequentially by the main thread which ensures that the most promising move is evaluated first. Then, once the first move is evaluated, the remaining moves are distributed among multiple threads for parallel evaluation.

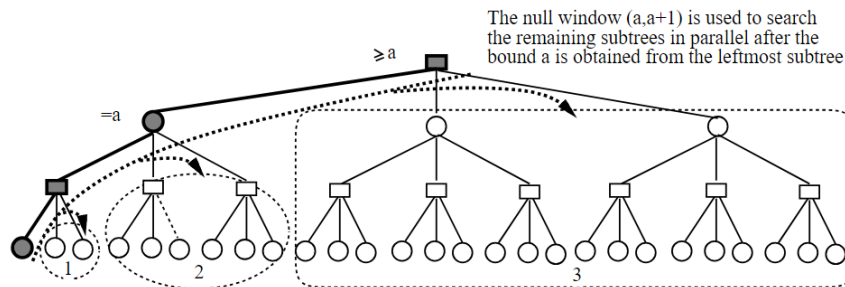


Figure 2.6: Principal variation splitting. Gao and Marsland (1996)

## 2.5. Evaluation

For each position, a numerical value is assigned representing how favorable the position is for one side: positive (+) for white and negative (-) for black. These symmetric values were first formulated by Shannon (1950) and which guide the engine in selecting the best moves during the search process. Generally, this value is expressed in centipawns (cp) that represents one hundredth of a pawn's value. There are two ways to implement it:

- Traditional hand-crafted evaluation
- Multi-layer neuronal networks

The use of neural networks is beyond the scope of the project, so traditional evaluation has been implemented and discussed below.

When teaching new chess players how to evaluate their position, assigning a simple value to each piece on board is an effective approach.

Piece	Value
Pawn	100
Knight	320
Bishop	330
Rook	500
Queen	950
King	$\infty$

Table 2.1: Standard values assigned to chess pieces in centipawns.

Over the time, there have been different point values for each piece. A table can be found in [https://www.chessprogramming.org/Point\\_Value](https://www.chessprogramming.org/Point_Value) in *Basic values* section.

By summing up the piece values for each side, we end up the concept of material. For example, if white has 1 knight and 1 bishop, while black has 2 bishops, the material calculation is as follows:

$$(+)(1 \times \text{knightValue} + 1 \times \text{bishopValue}) - (2 \times \text{bishopValue}) \quad (2.1)$$

Substituting the standard piece values in centipawns (cp):

$$(1 \times 320 + 1 \times 330) - (2 \times 330) = -10 \text{ cp} \quad (2.2)$$

This result indicates that black has a slight material advantage of 10 centipawns.

There are other material considerations like *bonus for the bishop pair* depending on the position that can be advantageous to control squares of different color, or insufficient material, which refers to situations where neither side has enough pieces to deliver checkmate. This is also mentioned in Section 1.3.4.

However, the higher the level one aims to reach in chess, the greater the need to evaluate other aspects. For instance, there are squares in which the pieces have less value because of their activity.

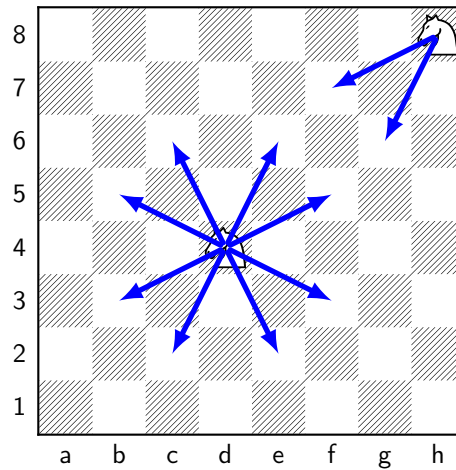


Figure 2.7: Knight's movement on corner square and center square.

A knight on a corner has only 2 moves, while one in the center can reach up to 8.

A way to add this to the evaluation is to have precomputed piece-square tables like the following one for the bishop:

-20	-10	-10	-10	-10	-10	-10	-20
-10	0	0	0	0	0	0	-10
-10	0	5	10	10	5	0	-10
-10	5	5	10	10	5	5	-10
-10	0	10	10	10	10	0	-10
-10	10	10	10	10	10	10	-10
-10	5	0	0	0	0	5	-10
-20	-10	-10	-10	-10	-10	-10	-20

Figure 2.8: Chessboard with precomputed piece-square values for the bishop.

It is important to consider the game phase to select these piece-square tables, especially in the endgame phase. These phases are:

- Opening: where players develop their minor pieces (knights and bishops), control the center of the board and try to ensure king safety.
- Middlegame: where players try to create tactical opportunities and attack the opponent's position once they have developed their pieces and secured their king.
- Endgame: where players have fewer pieces on the board and the focus is on promoting pawns and achieving checkmate.

In the endgame phase, when there are no queens, king activity becomes extremely significant in supporting the advancement of the pawns.

In relation to the above and Section 2.2, the mobility score measures the number of legal moves a piece can make from its current position. As it is mentioned in <https://www.chessprogramming.org/Mobility>,

*...the more choices you have at your disposal, the stronger your position.*

This score is calculated as the number of legal moves adjusted for certain factors like blocked moves by friendly pieces (pieces of the same color as the current piece) and enemy attacks (the squares controlled by the opponent).

During opening and middlegame, king safety is an important factor to consider. When the king has castled, it is crucial to maintain the pawns nearby to shield it from attacks. Generally, it is preferable to keep the pawns unmoved or advanced by only one square. This is called king's shield and a bonus is awarded for each pawn that still has not moved from the castling area.

On the contrary, king safety penalization can also mean to punish the player based on the number and type of enemy attacks targeting the area around the king.

All these ensures that positions with exposed kings are evaluated as less favorable. If there are open (no pawns of either color) or semi-open (no pawns of one color but at least one pawn of the opposite color) files near the king, the position is penalized further, as these files provide attacking opportunities for enemy rooks and queens.

## 2.6. How can we determine the strength of our engine?

This can be answered by playing against other engines and analyzing the results. The most common way to do this is by using the Computer Chess Rating Lists which ranks chess engines based on their performance in various tournaments and matches. By the time being, we have chosen to compare different versions of the engine with Stockfish, currently ranked as the number one on the list. Continue reading to learn about the used tools or directly read about the work behind it in Section 3.4.

### 2.6.1. Profiler

To analyze the performance of our chess engine and identify potential bottlenecks, we used the `perf` tool available on Linux systems. `perf` provides robust profiling capabilities by recording CPU events, sampling function execution, and collecting stack traces.

Our profiling goal is to identify which parts of the code consume the most execution time.

We run the engine under **perf** using the following commands:

Listing 2.5: Profiling AlphaDeepChess with perf

```
# Record performance data with function stack traces
sudo perf record -g ./build/release/AlphaDeepChess

# Display interactive report
sudo perf report -g --no-children
```

After recording, **perf report** opens an interactive terminal interface where functions are sorted by CPU overhead.

### 2.6.2. CustomTkinter

CustomTkinter is a modern UI-library for Python that extends from Tkinter module and considers both OOP (Object Oriented Programming) and simple implementations.

Although UCI implements a command that draws the current position through standard output, making moves and showing the evaluation is somewhat a time-consuming task when debugging and testing while programming. We resorted to using an interface to help us do this job and a really fast solution was to use Python to make a GUI. In this case, one of the most used UI libraries was CustomTkinter and it was used to build a friendly interface from scratch for bridging between executable and command sending.

### 2.6.3. Cutechess

Cutechess is an open-source tool designed to perform automated games between chess engines. It is widely used in the chess programming community to test and compare engines, evaluate their performance, and analyze games.

It provides both command-line interface (CLI) and a graphic user interface (GUI), with cross-platform compatibility for Windows, macOS, and Linux. For our purposes, we utilized the CLI version to automate the tests with Python scripts and commands, integrating it into a CI/CD workflow.

### 2.6.4. Stockfish

Stockfish is also an open-source tool and command-line program and chess engine with which we will compare with our engine. It is available for multiple platforms (Windows, macOS, Linux, Android, and iOS).

It is necessary to highlight that it provides multiple versions optimized for different hardware structures. These versions leverage specific CPU instruction sets to improve performance. For instance, the AVX2 version is recommended for most users



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with Intel processors from 2013 onwards or AMD processors from 2015 onwards, as it utilizes advanced vectorization instructions.

### **2.6.5. GitHub Actions and workflows**

GitHub Actions is a CI/CD tool integrated into GitHub that allows developers to automate tasks such as building, testing, and deploying code. Workflows are defined in YAML files and specify the tasks to be executed, the jobs or events that trigger them, and the environment in which they run.

In this project, since it is public in a GitHub repository, GitHub Actions was used to automate the testing and evaluation of the chess engine using the last two mentioned tools. A workflow was configured to compile the engine, run automated games using CUTEchess between different versions of the engine or our engine versus Stockfish.



# Chapter 3

## Work description

We designed a simple structure to organize all the scripts. As we are implementing different versions of the engine, there are specific modules that are compulsory to have well distinguished and possibly selected between them. Those modules are presented in the following Section 3.1.

Next, the key point is how to represent each required structure, including the board, chess pieces, precomputed tables, transposition tables, etc., as discussed in Section 3.2.

### 3.1. Modules

Each module is responsible for a specific aspect of the chess engine's functionality. The good part about this modular design is that it ensures clarity, maintainability, and the ability to test and improve individual components independently even more so when it comes to development with more people.

#### 3.1.1. Board

This module handles the representation of the chessboard, as previously mentioned in Section 2.1, and the state of the game. It includes:

- Representation of the chessboard with its pieces using bitboards for efficiency.
- Functions to make and unmake moves, including special moves like castling and en passant.
- Updating game state variables, such as castling rights or the half-move counter.

This module is vital, as it provides the data structures and operations required by other modules.

### 3.1.2. Evaluation

This module returns the final evaluation, a positive, 0, or negative number, depending on the current position.

#### 3.1.2.1. Move Generator

This module generates the list of legal moves generated depending on the board while updating relevant information such as king's danger masks, pinned pieces or checking pieces, and calculating each piece moves. All chess rules are taken into consideration.

### 3.1.3. Move ordering

This module uses a heuristic to return a prioritized list of moves to improve the efficiency of the search.

### 3.1.4. Search

This module must return a set of search results, each containing the best move and the achieved evaluation at a specific depth.

### 3.1.5. UCI

This module is responsible for processing commands to invoke each module, following UCI standards.

## 3.2. Code implementation

Although the previous algorithms were presented in pseudocode, from this point onward, all code implementations will be written in C++.

### 3.2.1. Data representation

#### 3.2.1.1. Square

There are 64 possible squares on a chessboard so the first thought is to use a 6-bit structure which seem like a perfect match ( $2^6 = 64$ ). However, modern processors are optimized to work with data sizes aligned to multiples of 8 bits (1 byte). Then, using 6 bits would require packing the data into more complex structures, which could introduce additional overhead in terms of bit manipulation and memory access. We preferred clarity and performance over micro-optimizations that complicate readability so we used a `uint8_t` to describe a square.

Moreover, masks are extremely useful for efficiently identifying and manipulating the squares on a chessboard using bitwise operations. Some of these masks are defined

as constants in the code and others are calculated during compilation time. For example, they can be used to identify the column of a square or simply placing a new piece on board.

Listing 3.1: Square class.

```
class Square {
    uint8_t sq_value;

    constexpr bool is_valid() const {
        return sq_value < 64U;
    }

    constexpr uint64_t mask() const {
        return is_valid() ? 1ULL << sq_value : 0ULL;
    }

    // Calculating the column of a square (sq_value % 8)
    constexpr int col() const {
        return is_valid() ? sq_value & 7U : COL_INVALID;
    }
}

// Placing a piece by setting the bit corresponding to the
// square to 1
const uint64_t mask = square.mask();
bitboard_all |= mask;
```

Just to clarify, `sq_value` must be a value between 0 and 63, inclusive, for the 64 squares on a chessboard. To calculate the column of a square, an AND operation is applied: `sq_value & 7U` extracts the 3 least significant bits, which correspond to the column of the square. Columns are enumerated from 0 (A) to 7 (H).

### 3.2.1.2. Piece and PieceType

They are simply enumerations where each piece and piece type correspond to an integer number to improve code readability.

Listing 3.2: Piece and PieceType enumerators.

```
enum class Piece : int
{
    W_PAWN = 0,
    W_KNIGHT = 1,
    W_BISHOP = 2,
    ...
    B_QUEEN = 10,
    B_KING = 11,
    EMPTY = 12,
    NUM_PIECES = 13
};
```

```
enum class PieceType : int
{
    PAWN = 0,
    KNIGHT = 1,
    BISHOP = 2,
    ROOK = 3,
    QUEEN = 4,
    KING = 5,
    EMPTY = 6,
    NUM_PIECES = 7
};
```

### 3.2.1.3. Move and MoveType

There are four types of moves: normal, promotion, en passant, and castling. These are represented using an integer-based enumeration:

Listing 3.3: MoveType enumerator.

```
enum class MoveType
{
    NORMAL = 0,
    PROMOTION = 1,
    EN_PASSANT = 2,
    CASTLING = 3
};
```

Meanwhile, moves are represented as a combination of two squares (the origin square and the destination square), 2 bits for the promotion piece, and 2 bits for the move type, all encoded in a `uint16_t`. In this case, each square is represented using 6 bits, resulting in a total of 16 bits ( $6 + 6 + 2 + 2 = 16$  bits). Each bit field has a unique mask and its specific shift, which must remain unchanged throughout the development.

### 3.2.1.4. Game State

The game state must store important information during the game, including the Zobrist hash key of the current position, the number of moves, the en passant square, the castling rights for each side and color, the side to move, the last captured piece, the fifty-move rule counter, the number of pieces, and a flag indicating whether the attacks are updated.

There are two `uint64_t` variables: one for the Zobrist hash key and the other for the remaining bit fields:

Listing 3.4: GameState class.

```
class GameState {
    uint64_t zobrist_key;
```

```

// 50 : attacks_updated : 1 if updated, 0 if not
// 43-49 : num_pieces : 0 to 64 pieces
// 35-42 : fifty_move_rule_counter : if counter gets to 100
// then game is a draw.
// 32-34 : last_captured_piece : PieceType::Empty if last
// move was not a capture.
// 31 : side_to_move : 0 if white, 1 if black.
// 30 : castle_king_white : 1 if available, 0 if not.
// 29 : castle_queen_white : 1 if available, 0 if not.
// 28 : castle_king_black : 1 if available, 0 if not.
// 27 : castle_queen_black : 1 if available, 0 if not.
// 26-20 : en_passant_square : 0-63 if available, >=64 if not
// available
// 19-0 : move_number : 0-1048575 number of moves in the game
.
uint64_t state_register;
}

```

### 3.2.1.5. Board

For the chessboard, as previously mentioned, bitboards are represented as 64-bit structures using `uint64_t`. Since the sign is not relevant, an unsigned type is used.

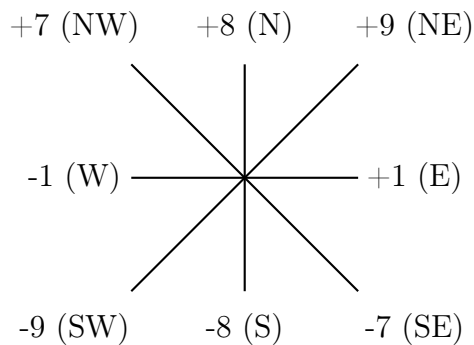
The bitboard representation follows the Little-Endian Rank-File Mapping convention also called as LERF. In this mapping, each bit in the 64-bit integer corresponds to a square on the chessboard where the least significant bit (0) is square A1, and the most significant bit (63) is square H8.

8	56	57	58	59	60	61	62	63
7	48	49	50	51	52	53	54	55
6	40	41	42	43	44	45	46	47
5	32	33	34	35	36	37	38	39
4	24	25	26	27	28	29	30	31
3	16	17	18	19	20	21	22	23
2	8	9	10	11	12	13	14	15
1	0	1	2	3	4	5	6	7
	A	B	C	D	E	F	G	H

Figure 3.1: Little-Endian Rank-File Mapping with Coordinates.

There are bitboards for all pieces (`bitboard_all`), for each piece color (`bitboard_color[0]` and `bitboard_color[1]`), and for each piece type. (`bitboard_piece[Piece]` like `bitboard_piece[Piece::W_QUEEN]`)

To identify ray directions on the board, we used the compass rose:



This means that, to get the numerical value that identifies the square to the north-east of a given square, you only need to add 9. For example, given the square  $f6$  (45), the north-east square  $g7$  has a value of 54 ( $45 + 9 = 54$ ). It is really effective for sliding pieces to calculate their attacks.

#### 3.2.1.6. Transposition table

The transposition table contains a list of entries. These entries are defined as a storage of information about a specific chess position, including its Zobrist key, evaluation score, best move, node type, and search depth.

The use of this will be discussed in Section 3.3.1.

#### 3.2.1.7. History

It contains a circular array (`uint64_t positions[HISTORY_MAX_SIZE]` with a position index that runs circularly `int next_position_index`) to store the hashes of previous board positions, ensuring efficient memory usage and quick access. This is essential for detecting threefold repetitions. Additionally, it provides functions to add new positions, remove the most recent ones, and clear the history when needed. The array is implemented with a fixed size (`HISTORY_MAX_SIZE`) of 128, defined as a power of two, to optimize indexing operations.

The detection of threefold repetition is performed by comparing the hash of the current position with the hashes stored in the history. If the same hash appears at least twice, it indicates that the position has been repeated three times.

#### 3.2.1.8. Precomputed Data

Precomputed data refers to information that is calculated and stored in advance to accelerate runtime operations. It reduces the need for repetitive calculations during gameplay. This approach allows the engine to focus on more computationally intensive tasks, such as search and evaluation.

The following types of precomputed data are used in AlphaDeepChess:

- Piece-Square Tables: These tables assign a static value to each piece on every square of the board, reflecting its positional strength. For example, pawns are



more valuable in the center of the board, while knights are stronger when placed near the middle. These are used for evaluation.

Listing 3.5: Initialization of knight piece-square table

```
static constexpr int KNIGHT_PIECE_SQUARE_TABLE[64] = {
    -50, -40, -30, -30, -30, -30, -40, -50,
    -40, -20, 0, 0, 0, 0, -20, -40,
    -30, 0, 10, 15, 15, 10, 0, -30,
    -30, 5, 15, 20, 20, 15, 5, -30,
    -30, 0, 15, 20, 20, 15, 0, -30,
    -30, 5, 10, 15, 15, 10, 5, -30,
    -40, -20, 0, 5, 5, 0, -20, -40,
    -50, -40, -30, -30, -30, -30, -40, -50
};
```

- Piece Attack Tables: Precomputed bitboards that represent the attack patterns of each piece type from every square on the board. These masks are used to quickly determine legal moves and attacked squares and they are also used in evaluation.

Listing 3.6: Initialization of bishop attack tables

```
constexpr uint64_t calculate_bishop_moves(Square square,
    uint64_t blockers)
{
    assert(square.is_valid());

    uint64_t moves_mask = 0ULL;

    const Direction dirs[4] = {NORTH_EAST, NORTH_WEST,
        SOUTH_EAST, SOUTH_WEST};

    for (const Direction dir : dirs) {

        Square aux_sq = square;
        aux_sq.to_direction(dir);

        while (aux_sq.is_valid()) {

            moves_mask |= aux_sq.mask();

            if (blockers & aux_sq.mask()) {
                break; // if there is a blocker in the
                        // square we stop in this direction
            }
            else {
                aux_sq.to_direction(dir);
            }
        }
    }
}
```

```

    return moves_mask;
}

```

- Ray Directions: Precomputed arrays that store the sliding attack directions for each square. These are used to efficiently calculate moves for sliding pieces during move generation.

Listing 3.7: Initialization of directions masks

```

static constexpr const std::array<std::array<uint64_t, 64>,
64> init_direction_mask()
{
    std::array<std::array<uint64_t, 64>, 64>
        PRECOMPUTED_DIR_MASK = {};

    for (Square sq1 = Square::A1; sq1.is_valid(); sq1++) {
        for (Square sq2 = Square::A1; sq2.is_valid(); sq2++) {
            if (sq1.row() == sq2.row()) {
                PRECOMPUTED_DIR_MASK[sq1][sq2] =
                    get_row_mask(sq1.row());
            }
            else if (sq1.col() == sq2.col()) {
                PRECOMPUTED_DIR_MASK[sq1][sq2] =
                    get_col_mask(sq1.col());
            }
            else if (squares_in_same_diagonal(sq1, sq2)) {
                PRECOMPUTED_DIR_MASK[sq1][sq2] =
                    get_diagonal_mask(sq1.diagonal());
            }
            else if (squares_in_same_antidiagonal(sq1, sq2)) {
                PRECOMPUTED_DIR_MASK[sq1][sq2] =
                    get_antidiagonal_mask(sq1.antidiagonal());
                ;
            }
            else {
                PRECOMPUTED_DIR_MASK[sq1][sq2] = 0ULL;
            }
        }
    }

    return PRECOMPUTED_DIR_MASK;
}

```

- Zobrist Hash Seeds: Randomly generated 64-bit numbers used for Zobrist hashing. Each piece-square combination, castling right, en passant square, and side to move is assigned a unique random value.

Listing 3.8: Initialization of seeds

```

static void init_random_numbers_only_once()
{
    static bool initialized = false;

    if (initialized) {
        return;
    }

#ifdef NDEBUG
    const uint64_t SEED = 123456789ULL;
    std::mt19937_64 gen(SEED);    // Use constant seed in
    debug mode
#else
    std::random_device rd;
    std::mt19937_64 gen(rd());    // Use random seed in
    release mode
#endif

    std::uniform_int_distribution<uint64_t> dis;

    for (Square sq = Square::A1; sq.is_valid(); sq++) {
        for (int p = 0; p < 12; p++) {
            square_piece_seed[sq.value()][p] = dis(gen);
        }
    }

    for (auto& elem : en_passant_seed) {
        elem = dis(gen);
    }

    king_white_castle_seed = dis(gen);
    queen_white_castle_seed = dis(gen);
    king_black_castle_seed = dis(gen);
    queen_black_castle_seed = dis(gen);

    black_to_move_seed = dis(gen);

    initialized = true;
}

```

### 3.2.2. Evaluation

In addition to the basic information that is mentioned in Section 2.5, in order to calculate the game phase, we can perform a dynamic evaluation (tampering evaluation): evaluate twice, once as if we were in the middlegame and once as if we were in the endgame. The final evaluation will be the sum of both, but each with a different weight. To do this, we calculate the middlegame percentage (24 pieces means 100 %

middlegame) and the endgame percentage (0 pieces = 100 % endgame).

$$\text{Evaluation} = \text{middlegame} \% \cdot \text{eval\_middlegame} + \text{endgame} \% \cdot \text{eval\_endgame}$$

Listing 3.9: Counting pieces and depending on game phase evaluation.

```

int evaluate_position(Board& board)
{
    board.update_attacks_bb();

    int middlegame_eval = 0;
    int endgame_eval = 0;

    const int middlegame_percentage =
        calculate_middlegame_percentage(board);
    const int endgame_percentage = 24 -
        calculate_middlegame_percentage(board);

    uint64_t pieces = board.get_bitboard_all();

    while (pieces) {

        const Square square(pop_lsb(pieces));
        const Piece piece = board.get_piece(square);
        const ChessColor color = get_color(piece);
        const int piece_raw_value = raw_value(piece);
        const int bonus_middlegame = PrecomputedEvalData::
            get_piece_square_table<PST_TYPE_MIDDLEGAME>(piece,
                square);
        const int bonus_endgame = PrecomputedEvalData::
            get_piece_square_table<PST_TYPE_ENDGAME>(piece, square
            );

        const int piece_value_middlegame = piece_raw_value +
            bonus_middlegame;
        const int piece_value_endgame = piece_raw_value +
            bonus_endgame;

        middlegame_eval += is_white(color) ?
            piece_value_middlegame : -piece_value_middlegame;
        endgame_eval += is_white(color) ? piece_value_endgame : -
            piece_value_endgame;
    }
    const int blended_eval = (middlegame_eval *
        middlegame_percentage + endgame_eval * endgame_percentage)
        / 24;

    return blended_eval;
}

```

---

}

---

Listing 3.10: King safety, piece mobility, depending on game phase evaluation.

```

int evaluate_position(Board& board)
{
    board.update_attacks_bb();

    int middlegame_eval = 0;
    int endgame_eval = 0;

    const int middlegame_percentage =
        calculate_middlegame_percentage(board);
    const int endgame_percentage = 24 -
        calculate_middlegame_percentage(board);

    const Square white_king_sq = lsb(board.get_bitboard_piece(
        Piece::W_KING));
    const Square black_king_sq = lsb(board.get_bitboard_piece(
        Piece::B_KING));

    uint64_t pieces = board.get_bitboard_all();

    while (pieces) {

        const Square square(pop_lsb(pieces));
        const Piece piece = board.get_piece(square);
        const ChessColor color = get_color(piece);
        const int piece_raw_value = raw_value(piece);
        const int bonus_middlegame = PrecomputedEvalData::
            get_piece_square_table<PST_TYPE_MIDDLEGAME>(piece,
            square);
        const int bonus_endgame = PrecomputedEvalData::
            get_piece_square_table<PST_TYPE_ENDGAME>(piece, square
            );
        const int mobility = mobility_piece_score(square, piece,
            board);

        const int piece_value_middlegame = piece_raw_value +
            bonus_middlegame + mobility;
        const int piece_value_endgame = piece_raw_value +
            bonus_endgame + mobility;

        middlegame_eval += is_white(color) ?
            piece_value_middlegame : -piece_value_middlegame;
        endgame_eval += is_white(color) ? piece_value_endgame : -
            piece_value_endgame;
    }

    const int king_safety_penalty_white =

```

```

        king_safety_penalization<ChessColor::WHITE>(white_king_sq,
            board);
    const int king_safety_penalty_black =
        king_safety_penalization<ChessColor::BLACK>(black_king_sq,
            board);
    const int king_safety_penalty = king_safety_penalty_black -
        king_safety_penalty_white;

    const int king_shield_bonus_white = king_shield<ChessColor::
        WHITE>(white_king_sq, board);
    const int king_shield_bonus_black = king_shield<ChessColor::
        BLACK>(black_king_sq, board);
    const int king_shield_bonus = king_shield_bonus_white -
        king_shield_bonus_black;

    middlegame_eval += king_shield_bonus + king_safety_penalty;

    const int blended_eval = (middlegame_eval *
        middlegame_percentage + endgame_eval * endgame_percentage)
        / 24;

    return blended_eval;
}

```

Listing 3.11: Different pawn square values depending on game phase.

```

// Opening and Middlegame pawn square values
static constexpr int PAWN_PIECE_SQUARE_TABLE[64] = {
    0, 0, 0, 0, 0, 0, 0, 0,
    50, 50, 50, 50, 50, 50, 50, 50,
    10, 10, 20, 30, 30, 20, 10, 10,
    5, 5, 10, 25, 25, 10, 5, 5,
    0, 0, 0, 20, 20, 0, -5, 0,
    5, -5, -10, 0, 0, -10, -5, 5,
    5, 10, 10, -20, -20, 10, 10, 5,
    0, 0, 0, 0, 0, 0, 0, 0
};

// Endgame pawn square values
static constexpr int PAWN_ENDGAME_SQUARE_TABLE[64] = {
    0, 0, 0, 0, 0, 0, 0, 0,
    80, 80, 80, 80, 80, 80, 80, 80,
    50, 50, 50, 50, 50, 50, 50, 50,
    30, 30, 30, 30, 30, 30, 30, 30,
    20, 20, 20, 20, 20, 20, 20, 20,
    10, 10, 10, 10, 10, 10, 10, 10,
    10, 10, 10, 10, 10, 10, 10, 10,
    0, 0, 0, 0, 0, 0, 0, 0
};

```

### 3.2.3. Move generator

Calculating the legal moves in a chess position is a more difficult and tedious task than it might seem, mainly due to the unintuitive rules of en passant and castling, and it is also difficult to restrict the moves of pinned pieces. Jones (2023)

To create an efficient move generator, we will be using bitboards. The steps to generate legal movements efficiently are the following:

- Calculate the bitboard of attacked squares by the waiting side.
- Calculate bitboard of pinned pieces.
- Generate the legal moves of each piece of the side whose turn is, knowing that the king cannot move to any attacked square and that pinned pieces can only move in the direction of the pin.
- Generate the special legal moves like en passant and castling.

### 3.2.4. Move ordering

Order the legal moves from most to least likely to be the best move in the position. The sooner we explore the best move, the more branches of the tree will be pruned. To do this, we use the MVV-LVA heuristic (most valuable victim, least valuable aggressor). We give higher scores to capturing a low-value piece over a higher-value piece. Capturing a queen with a pawn scores highly. We also give a bonus to piece promotions.

#### 3.2.4.1. Killer moves

A killer move is a non-capturing (quiet) move that previously caused a beta-cutoff during the search in a sibling node or any other branch at the same depth in the game tree. These moves are often strong candidates, as they have previously led to pruning in similar positions. Promoting them early in the move ordering increases the chances of early cutoffs, which improves search efficiency.

To take advantage of this heuristic we store up to two killer moves for each search depth. During move ordering, these killer moves are given a bonus score, allowing them to be explored before other quiet moves.

### 3.2.5. The Core: Search Algorithm

The core of the chess engine is its search algorithm, in this case alpha-beta pruning algorithm. The entire game tree is generated up to a selected maximum depth. At each node, the next player evaluates the position and during execution, the values of alpha and beta are updated. Pruning is performed when a branch of the tree

is detected as irrelevant because the evaluation being examined is worse than the current value of alpha or beta for MAX or MIN, respectively.

Therefore, the following events happen at each node of the tree:

- Check if we are at an end node because of a checkmate, a draw by triple repetition, the 50-move rule, or because we have reached the maximum selected depth.
- Evaluate the position: a positive value (+) means that White has an advantage, and a negative value (-) means that Black has an advantage. A limit is set that represents mate in one; we have arbitrarily chosen 3,200,000.
- Generate legal moves: create a list of every possible legal move in the position.
- Order the legal moves: from greatest to least intuition of being the best move for the position. The sooner we explore the best move, the more branches of the tree will be pruned.
- Explore each of the legal moves from the position in order, update the evaluation, the value of alpha and beta, and check if we can perform pruning.

Listing 3.12: Alpha-Beta search.

```

template<SearchType searchType>
static int alpha_beta_search(std::atomic<bool>& stop, int depth,
    int ply, int alpha, int beta, SearchContext& context)
{
    // ...

    MoveList moves;
    bool isCheck;
    generate_legal_moves<ALL_MOVES>(moves, board, &isCheck);
    // Verify ending position: checkmate, stalemate, fifty-move
    // rule counter...

    int final_node_evaluation = MAXIMIZING_WHITE ? -INF_EVAL : +
        INF_EVAL;
    const GameState game_state = board.state();

    order_moves(moves, board, ply);

    for (int i = 0; i < moves.size(); i++) {
        // ...
        constexpr SearchType nextSearchType = MAXIMIZING_WHITE ?
            MINIMIZE_BLACK : MAXIMIZE_WHITE;

        board.make_move(moves[i]);
    }
}

```



```

    int eval = alpha_beta_search<nextSearchType>(stop , depth
        - 1, ply + 1, alpha , beta , context);
    board.unmake_move(moves[i] , game_state);
    History::pop_position();

    if constexpr (MAXIMIZING_WHITE) {
        // ...

        final_node_evaluation = std::max(
            final_node_evaluation , eval);
        alpha = std::max(alpha , eval);

        if (final_node_evaluation >= beta) {
            // ...
            break;    // beta cutoff
        }
    }
    else if (MINIMIZING_BLACK) {
        // ...

        final_node_evaluation = std::min(
            final_node_evaluation , eval);
        beta = std::min(beta , eval);

        if (final_node_evaluation <= alpha) {
            // ...
            break;    // alpha cutoff
        }
    }
}

return final_node_evaluation;
}

```

#### 3.2.5.1. Search: iterative deepening

At what depth do we decide to search? Actually, the simplest thing is to perform an infinite search, first searching at depth 1, then 2, then 3... to infinity. The engine will update the evaluation and the best move for the position in each iteration. Simply by signaling *stop* the search will stop.

Listing 3.13: Iterative deepening.

```

static void iterative_deepening(std::atomic<bool>& stop ,
    SearchResults& results , int max_depth , SearchContext& context)
{
    // ...

    int alpha = -INF_EVAL;

```

```

int beta = +INF_EVAL;

for (int depth = 1; depth <= max_depth; depth++) {
    context.bestMoveInIteration = Move::null();
    context.bestEvalInIteration = is_white(side_to_move) ? -
        INF_EVAL : +INF_EVAL;

    // ...

    is_white(side_to_move) ? alpha_beta_search<MAXIMIZE_WHITE>
        (stop, depth, 0, alpha, beta, context)
        : alpha_beta_search<MINIMIZE_BLACK>
        (stop, depth, 0, alpha, beta, context);

    // ...

    context.bestMoveFound = context.bestMoveInIteration;
    context.bestEvalFound = context.bestEvalInIteration;

    assert(context.bestMoveFound.is_valid());

    insert_new_result(results, depth, context.bestEvalFound,
        context.bestMoveFound);

    if (abs(context.bestEvalFound) > MATE_THRESHOLD) {
        break; // We found a checkmate, we stop because we
               cant find a shorter checkMate
    }
}

```

### 3.2.5.2. Search: aspiration window

Aspiration window's main objective is to reduce the number of nodes to explore by simply restricting the range of alpha and beta values (window). A search is performed with this narrow window. If the position evaluation falls within the window, it is accepted as valid and additional node exploration is avoided. However, if the evaluation is outside the limits of the window (when a fail-low or fail-high occurs), the window is expanded to the extreme values (-INF and +INF) and a new search is performed to obtain an accurate evaluation:

Listing 3.14: Aspiration window.

```

static void iterative_deepening(std::atomic<bool>& stop,
    SearchResults& results, int max_depth, SearchContext& context)
{
    // ...

```

```

int alpha = -INF_EVAL;
int beta = +INF_EVAL;

for (int depth = 1; depth <= max_depth; depth++) {
    context.bestMoveInIteration = Move::null();
    context.bestEvalInIteration = is_white(side_to_move) ? -
        INF_EVAL : +INF_EVAL;

    // if it is not the first iteration, adjust alpha and
    // beta
    if (depth > 1) {
        alpha = context.bestEvalFound - ASPIRATION_MARGIN;
        beta  = context.bestEvalFound + ASPIRATION_MARGIN;
    }

    is_white(side_to_move) ? alpha_beta_search<MAXIMIZE_WHITE>
        (stop, depth, 0, alpha, beta, context)
        : alpha_beta_search<MINIMIZE_BLACK>
        (stop, depth, 0, alpha, beta, context);

    // if the evaluation is out of bounds of the window, redo
    // the search with -INF_EVAL and +INF_EVAL
    if (context.bestEvalInIteration <= alpha || context.
        bestEvalInIteration >= beta) {
        alpha = -INF_EVAL;
        beta  = +INF_EVAL;

        is_white(side_to_move) ? alpha_beta_search<
            MAXIMIZE_WHITE>(stop, depth, 0, alpha, beta,
                context)
            : alpha_beta_search<
                MINIMIZE_BLACK>(stop, depth
                    , 0, alpha, beta, context);
    }

    // ...
}

```

### 3.2.5.3. Search: horizon effect, quiescence search

What happens if, upon reaching maximum depth, we evaluate the position in the middle of a piece exchange? For example, if a queen captures a pawn. It will seem like we have won a pawn, but on the next move, another pawn captures the queen, and now we lose a queen. This is known as the horizon effect. To avoid this, when we reach the end of the tree at maximum depth, we must extend the search to include only piece captures. This is known as quiescence search.

The purpose of this search is only to stop the search and evaluate quiet positions, where there is no capture or tactical movement.

Listing 3.15: Quiescence search implementation.

```

template<SearchType searchType>
static int quiescence_search(bool& stop, int ply, int alpha, int
    beta, SearchContext& context)
{
    constexpr bool MAXIMIZING_WHITE = searchType ==
        MAXIMIZE_WHITE;
    constexpr bool MINIMIZING_BLACK = searchType ==
        MINIMIZE_BLACK;

    Board& board = context.board;
    const uint64_t zobrist_key = board.state().get_zobrist_key();

    const uint8_t fifty_move_rule_counter = board.state().
        fifty_move_rule_counter();
    const bool fifty_move_rule_draw = fifty_move_rule_counter >=
        100U;

    if (History::threefold_repetition_detected(
        fifty_move_rule_counter) || fifty_move_rule_draw) {
        return 0;
    }

    int static_evaluation = evaluate_position(board);

    if (ply >= MAX_PLY) {
        return static_evaluation;
    }

    if constexpr (MAXIMIZING_WHITE) {
        if (static_evaluation >= beta) {
            return beta;    // beta cutoff
        }
        alpha = std::max(alpha, static_evaluation);
    }
    else if constexpr (MINIMIZING_BLACK) {
        if (static_evaluation <= alpha) {
            return alpha;    // Alpha cutoff
        }
        beta = std::min(beta, static_evaluation);
    }

    MoveList capture_moves;
    generate_legal_moves<ONLY_CAPTURES>(capture_moves, board);

    if (capture_moves.size() == 0) {

```

```

        return static_evaluation;    // No captures: return static
                                evaluation
    }

    order_moves(capture_moves, board, ply);

    const GameState game_state = board.state();
    int final_node_evaluation = static_evaluation;

    for (int i = 0; i < capture_moves.size(); i++) {

        if (stop) {
            return 0;
        }

        constexpr SearchType nextSearchType = MAXIMIZING_WHITE ?
            MINIMIZE_BLACK : MAXIMIZE_WHITE;

        History::push_position(zobrist_key);
        board.make_move(capture_moves[i]);
        int eval = quiescence_search<nextSearchType>(stop, ply +
            1, alpha, beta, context);
        board.unmake_move(capture_moves[i], game_state);
        History::pop_position();

        if constexpr (MAXIMIZING_WHITE) {
            final_node_evaluation = std::max(
                final_node_evaluation, eval);
            alpha = std::max(alpha, eval);

            if (final_node_evaluation >= beta) {
                break;    // Beta cutoff
            }
        }
        else if constexpr (MINIMIZING_BLACK) {
            final_node_evaluation = std::min(
                final_node_evaluation, eval);
            beta = std::min(beta, eval);

            if (final_node_evaluation <= alpha) {
                break;    // Alpha cutoff
            }
        }
    }

    return final_node_evaluation;
}

```

### 3.2.6. UCI

Universal Chess Interface specifications are independent of the operating system. To ensure the synchronization of the engine with the GUI, a `isready` command is sent and engine should respond with `readyok`.

The engine must be capable of continuously processing input from standard input, even during evaluation. Moreover, if an unknown command is received, it should be ignored.

The move format is in long algebraic notation which means sending two squares coordinates like `e2e4` or `b1c3` independently of the type of piece because the engine must be the one checking that the movement is legal.

Some of the most important commands are the following:

- `position [fen <fenstring>| startpos | actualpos] moves <move1>...<movei>` : sets the current position of the board to the FEN string or make the list of moves from starting position or current position.
- `go`: starts evaluating the current position. Some important subparameters are:
  - `depth <x>`: specifies the number of `x` plies to search.
  - `movetime <x>`: specifies the number of `x` seconds to search.
- `stop`: stops evaluation if it is running.

## 3.3. Improvements

Some improvements and new structures were later added for different versions: transposition tables, Zobrist hashing, table entries, PEXT instructions, search multithread and search reductions that will be discussed and analysed below.

### 3.3.1. Transposition Table

The basic implementation of the chess engine generates a large amount of redundant calculations due to transpositions: situations in which the same board position is reached through different sequences of moves in the game tree. Figure 3.2 illustrates a position that can arise through multiple move orders. Where the white king could go to the g3 square from multiple paths.

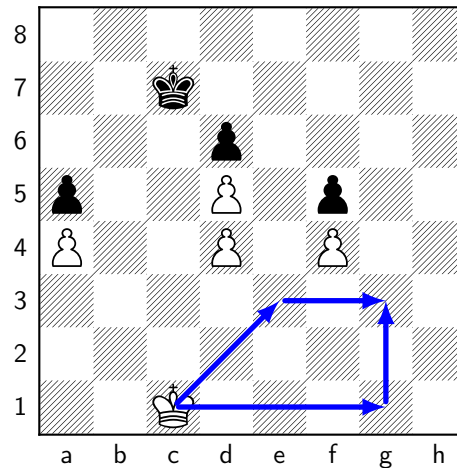


Figure 3.2: Lasker-Reichhelm Position, transposition example

Taking advantage of the concept of dynamic programming, we are going to create a look-up table of chess positions and its evaluation. So if we encounter the same position again, the evaluation is already precalculated. However, we ask ourselves the following question: how much space does the look-up table take up if there are an astronomical amount of chess positions? What we can do is assign a hash to each position and make the table index the last bits of the hash. The larger the table, the less likely access collisions will be. We also want a hash that is fast to calculate and has collision-reducing properties; for this, we will use the Zobrist hashing technique in the following subsection.

#### 3.3.1.1. Zobrist Hashing

Zobrist Hashing is a technique to transform a board position of arbitrary size into a number of a set length, with an equal distribution over all possible numbers invented by Albert Zobrist. (Zobrist (1970))

To generate a 64-bit hash for a position, the following steps are followed:

- There are 12 different types of chess pieces. For each of the 64 squares on the board, we generate 12 random 64-bit integers. That is, each piece-square combination is assigned a unique random value. This initialization step is performed only once when the program starts.
- The hash value for a given position is computed by performing the XOR operation between the hash accumulator and the random value corresponding to each piece on its square.
- In addition to the pieces, we also include:
  - A random value for the side to move.

- One random value to account for the possibility of an en passant capture.
  - One random value per possible castling move.
- These random values are useful so that even slightly different positions produce very different hash values. This greatly reduces the chance of collisions.
  - The XOR operation is used not only because it is computationally inexpensive, but also because it is reversible. This means that when a move is made or undone, we can update the hash incrementally by applying XOR only to the affected squares, without needing to recompute the entire hash.

Listing 3.16: Code example to calculate zobrist hash of position

```

/**
 * @brief Calculate zobrist hash of position
 *
 * @param[in] position board containing the chess position
 *
 * @return the hash key of the chess position
 */
static uint64_t hash(const Board& position)
{
    // initialize the random numbers only once
    init_random_numbers_only_once();

    uint64_t hash = 0ULL;

    for (Square square = Square::A1; square.is_valid(); square++)
    {
        const Piece piece = position.get_piece(square);
        if (piece != Piece::EMPTY) {
            hash ^= get_seed(square, piece);
        }
    }
    const GameState state = position.state();

    const Square eps_square = state.en_passant_square();
    if (eps_square.is_valid()) {
        hash ^= get_en_passant_seed(eps_square.col());
    }

    if (state.castle_king_white()) {
        hash ^= get_king_white_castle_seed();
    }
    if (state.castle_queen_white()) {
        hash ^= get_queen_white_castle_seed();
    }
    if (state.castle_king_black()) {

```



```

        hash ^= get_king_black_castle_seed();
    }
    if (state.castle_queen_black()) {
        hash ^= get_queen_black_castle_seed();
    }

    if (state.side_to_move() == ChessColor::BLACK) {
        hash ^= get_black_to_move_seed();
    }

    return hash;
}

```

### 3.3.1.2. Table Entry

Each entry in the transposition table stores the following information:

- **Zobrist Hash:** The full 64-bit hash of the position. This is used to verify that the entry corresponds to the current position and to detect possible index collisions in the table.
- **Evaluation:** The numerical evaluation of the position, as computed by the evaluation function.
- **Depth:** The depth at which the evaluation was calculated. A deeper search could potentially yield a more accurate evaluation, so this value helps determine whether a new evaluation should overwrite the existing one.
- **Node Type:** Indicates the type of node stored:
  - **EXACT** the evaluation is precise for this position.
  - **UPPERBOUND** the evaluation is an upper bound, typically resulting from an alpha cutoff.
  - **LOWERBOUND** the evaluation is a lower bound, typically resulting from a beta cutoff.
  - **FAILED** entry is empty or with invalid information.

Listing 3.17: Code example transposition table entry

```

class TranspositionTable::Entry
{
    uint64_t key; // full Zobrist key of the chess position

    int evaluation; // Score of the chess position

    Move move; // Best move found for the position.

    NodeType type; // FAILED, EXACT, LOWERBOUND, UPPERBOUND
}

```

```

    int8_t depth; // At which the position was evaluated
    ...
}

```

### 3.3.1.3. Collisions

As discussed earlier, index collisions in the transposition table are handled by verifying the full Zobrist hash stored in the entry. However, it is still theoretically possible for a full hash collision to occur, that is two different positions producing the same hash.

This scenario is extremely rare. With 64-bit hashes, there are  $2^{64}$  possible unique values, which is more than sufficient for practical purposes. In the unlikely event of a true hash collision, it could result in an incorrect evaluation being reused for a different position.

### 3.3.1.4. Analysis

To evaluate the improvement introduced by the transposition table, we conducted a 100-game tournament against the basic version of the engine. We selected 50 random starting positions from an opening book and played each position twice, alternating colors to ensure fairness. Each bot has 4 seconds to think per move.



Figure 3.3: 64MB Transposition Table bot vs basic bot

We see a substantial improvement by adding the transposition table with 46 wins versus 32 losses.

## 3.3.2. Move generator with Magic Bitboards and PEXT instructions

To identify potential performance bottlenecks, we performed profiling on the engine.

Samples: 15K of event 'cycles:P', Event count (approx.): 15313528435				
Overhead	Command	Shared Object	Symbol	
+ 36.07%	AlphaDeepChess	AlphaDeepChess	[.] generate_legal_moves(MoveList&, Board const&, bool*, bool*)	
+ 19.30%	AlphaDeepChess	AlphaDeepChess	[.] calculate_moves_in_direction(Square, Direction, MoveGeneratorInfo&) [clone .isra.0]	
+ 16.63%	AlphaDeepChess	AlphaDeepChess	[.] evaluate_position(Board const&)	
+ 16.23%	AlphaDeepChess	AlphaDeepChess	[.] update_danger_in_direction(Square, Direction, MoveGeneratorInfo&) [clone .isra.0]	
+ 1.24%	AlphaDeepChess	AlphaDeepChess	[.] calculate_king_moves(Square, MoveGeneratorInfo&) [clone .isra.0]	
0.96%	AlphaDeepChess	AlphaDeepChess	[.] quiescence_maximize_white(Board&, int, int, int)	
0.81%	AlphaDeepChess	AlphaDeepChess	[.] Board::make_move(Move) [clone .isra.0]	
0.74%	AlphaDeepChess	AlphaDeepChess	[.] order_moves(MoveList&, Board const&)	
0.73%	AlphaDeepChess	AlphaDeepChess	[.] quiescence_minimize_black(Board&, int, int, int)	
0.60%	AlphaDeepChess	AlphaDeepChess	[.] Board::put_piece(Piece, Square) [clone .isra.0]	
0.59%	AlphaDeepChess	libc.so.6	[.] memset_avx2_unaligned_erms	

Profiling results show that most part of the total execution time is spent in the generate legal moves function. Therefore, optimizing this component is expected to lead to significant performance improvements.

### 3.3.2.1. Magic bitboards

We can create a look up table of all the rook and bishop moves for each square on the board and for each combination of pieces that blocks the path of the slider piece (blockers bitboard). Basically we need a hash table to store rook and bishop moves indexed by square and bitboard of blockers. The problem is that this table could be very big. Kannan (2007)

Magic bitboards technique used to reduce the size of the look up table. We cut off unnecessary information in the blockers bitboard, excluding the board borders and the squares outside its attack pattern.

A **magic number** is a multiplier to the bitboard of blockers with the following properties:

- Preserves relevant blocker information: The nearest blockers along a piece's movement direction are preserved. *Example:* Consider a rook with two pawns in its path:

Rook  $\rightarrow$   $\rightarrow$   $\rightarrow$  [Pawn1] [Pawn2]

In this case, only 'Pawn1' blocks the rook's movement, while 'Pawn2' is irrelevant.

- Compresses the blocker bitboard, pushing the important bits near the most significant bit.
- The final multiplication must produce a unique index for each possible blocker configuration. The way to ensure the uniqueness is by brute force testing.

As illustrated in Figure 3.4, we aim to compute the legal moves of the white rook in the given position. In practice, the only pieces that truly block the rook's path are those marked with a red circle.

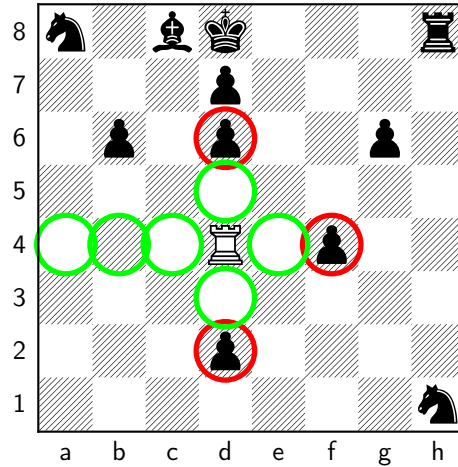


Figure 3.4: Initial chess position with white rook and blockers

First, we mask out all pieces outside the rook's attack pattern or on the board borders, as shown in Figure 3.5.

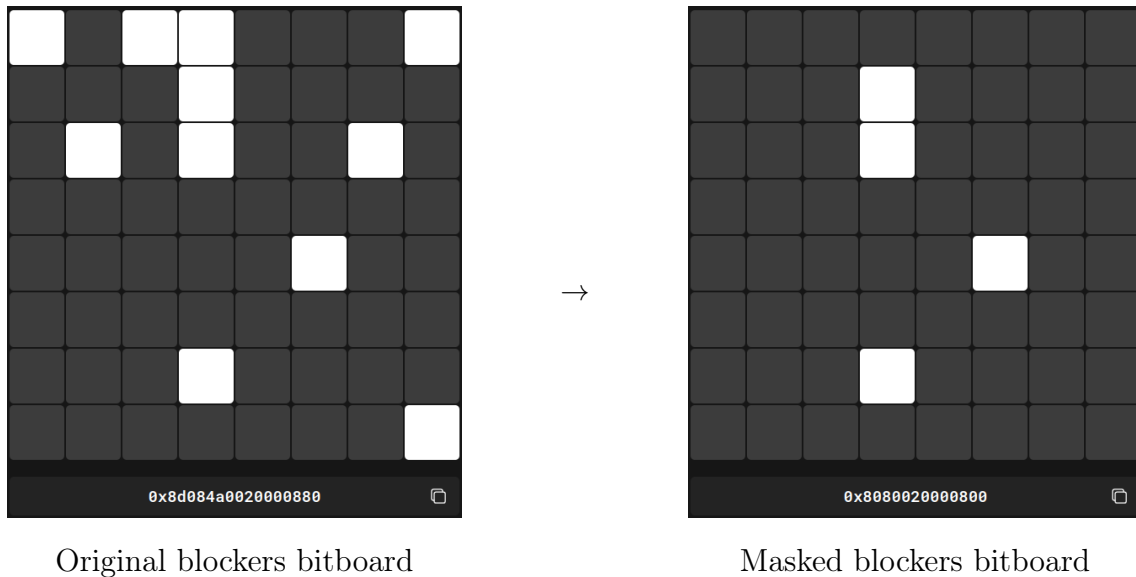


Figure 3.5: Pre-processing of the blockers bitboard

As illustrated in Figure 3.6, the masked blockers bitboard is then multiplied by the magic number. The result retains only the three relevant pawns that obstruct the rook's movement, pushing them toward the most significant bits.

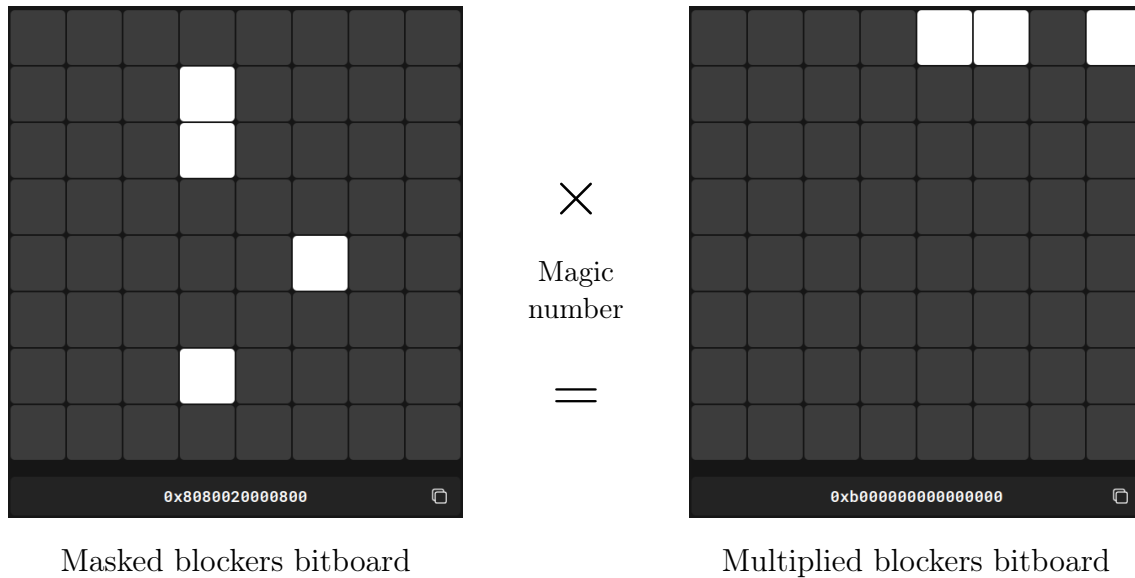


Figure 3.6: Multiplication by magic number to produce an index

Next, we compress the index toward the least significant bits by shifting right by `64-relevant_squares`. The number of relevant squares varies per board square; Listing 3.18 shows this for the rook:

Listing 3.18: Number of relevant squares for each rook square

```

/**
 * @brief ROOK_RELEVANT_SQUARES
 *
 * Number of squares where the rook could move
 * from each square minus the board border squares.
 */
static constexpr int ROOK_RELEVANT_SQUARES[NUM_SQUARES] =
{
    12, 11, 11, 11, 11, 11, 11, 12,
    11, 10, 10, 10, 10, 10, 10, 11,
    11, 10, 10, 10, 10, 10, 10, 11,
    11, 10, 10, 10, 10, 10, 10, 11,
    11, 10, 10, 10, 10, 10, 10, 11,
    11, 10, 10, 10, 10, 10, 10, 11,
    11, 10, 10, 10, 10, 10, 10, 11,
    12, 11, 11, 11, 11, 11, 11, 12
};

```

The final index is thus computed as:

```
index = (bitboard_of_blockers  $\times$  magic_number)  $\gg$  (64 - relevant_squares).
```

### 3.3.2.2. PEXT instruction

The **PEXT** (Parallel Bits Extract) instruction—available on modern x86\_64 CPUs—extracts bits from a source operand according to a mask and packs them into the lower bits of the destination operand. Hilewitz and Lee (2006) It is ideally suited for computing our table index.

Figure 3.7 illustrates how **PEXT** works: it selects specific bits from register **r2**, as specified by the mask in **r3**, and packs the result into the lower bits of the destination register **r1**.

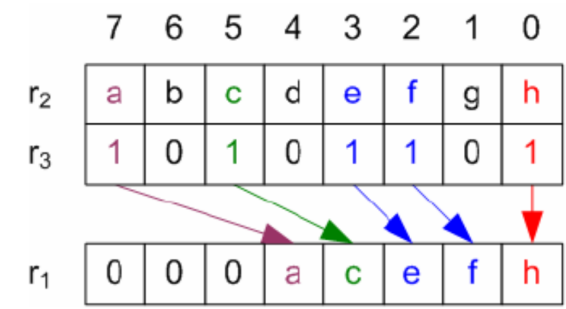


Figure 3.7: Example of the **PEXT** instruction: extracting bits from **r2** using **r3** as a mask, and storing the result in **r1**. Hilewitz and Lee (2006)

For our previous example (see Figure 3.4), we only need the full bitboard of blockers and the rook’s attack pattern (excluding the borders to reduce space), as illustrated in Figure 3.8.

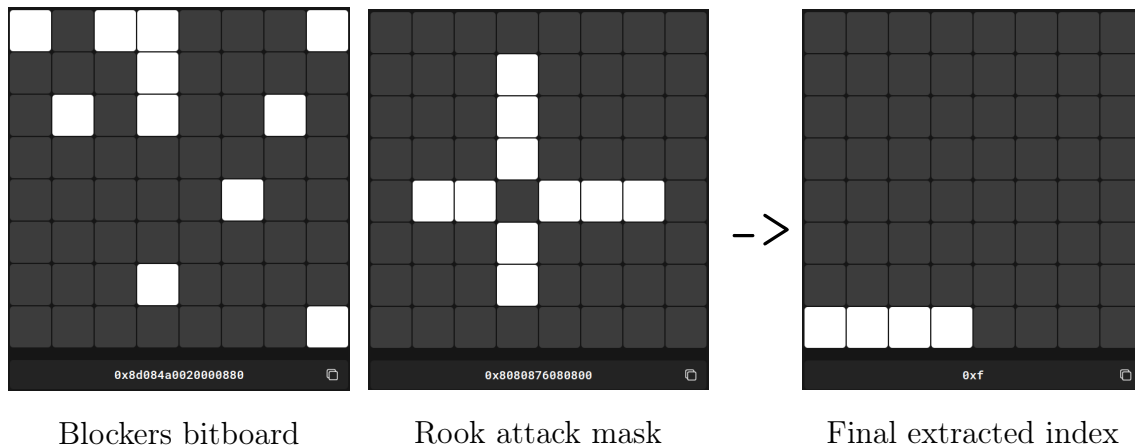


Figure 3.8: index extraction with Pext example

The final index used to access the lookup table is calculated using the **pext** instruction as follows:

```
index = _pext_u64(blockers, attack_pattern).
```

To maintain compatibility and performance across different hardware platforms, we provide two implementations:

- If PEXT support is detected at compile time, the engine uses it to compute the index directly.
- Otherwise, the engine falls back to the Magic Bitboards approach using multiplication and bit shifts.

### 3.3.2.3. Analysis

To evaluate the improvement in the move generator, we conducted the same 100 game match vs the basic bot version.



Figure 3.9: Move generator with PEXT instructions bot vs basic bot

Huge improvement with 64 wins versus 22 losses.

### 3.3.3. Evaluation with King Safety and piece mobility

It is often beneficial to evaluate additional aspects of a position beyond simply counting material. We introduce the following positional evaluation parameters:

1. King Shield Bonus: The king is typically safer when protected by friendly pawns in front of it. We assign a bonus in the evaluation score for each allied pawn positioned directly in front of the king.
2. King Safety Penalty: For each square within a  $3 \times 3$  area surrounding the king that is attacked by enemy pieces, we apply a penalty to reflect increased vulnerability.
3. Piece Mobility: Greater piece mobility is generally indicative of a stronger position. Each piece receives a bonus for every available move to a square that is not attacked by enemy pawns.

#### 3.3.3.1. Analysis

To evaluate the improvement in the new evaluation, we conducted the same 100 game match vs the basic bot version.



Figure 3.10: King Safety and Piece mobility evaluation bot vs basic bot

The results are slightly worse compared to the match using the material-only evaluation, 3.9 with 8 more losses than before. This may be due to the increased computational cost of evaluating these additional parameters. Furthermore, although these are abstract concepts commonly used by humans to assess positions, the engine may struggle to find a clear correlation between them and actual positional strength.

### 3.3.4. Search Multithread

This version of search follows YBWC mentioned in Section 2.4.2.1. It searches the first move sequentially after ordering the moves and launches the remaining moves in parallel using a thread. If it turns out that the first move was the best and pruning was applied, directly return the final node evaluation without the additional thread.

```

template<SearchType searchType>
static int alpha_beta_search(std::atomic<bool>& stop, int depth,
    int ply, int alpha, int beta, SearchContext& context, Board
    board)
{
    // ...
    // Check if it is in transposition table
    // If it is, return previously calculated evaluation

    generate_legal_moves<ALL_MOVES>(moves, board, &isCheck);
    // ...
    order_moves(moves, board, ply);
    // ...

    constexpr SearchType nextSearchType = MAXIMIZING_WHITE ?
        MINIMIZE_BLACK : MAXIMIZE_WHITE;

    // Search the first move sequentially
    board.make_move(moves[0]);
    int eval = alpha_beta_search<nextSearchType>(stop, depth - 1,
        ply + 1, alpha, beta, context, board);
    board.unmake_move(moves[0], game_state);
    // ...

    // Alpha-Beta pruning with first move
    if (/*pruned*/) goto end_search;

    // Search remaining moves in parallel
    thread = std::thread([&]() {
        for (int i = 1; i < moves.size(); i++) {
            // ...
            // Alpha-Beta pruning with the remaining moves
            if (/*pruned*/) break;
        }
    });
}

```



```

    if (thread.joinable()) {
        thread.join();
    }

end_search:
    // Store entry in transposition table

    return final_node_evaluation;
}

```

#### 3.3.4.1. Analysis

To evaluate the improvement in the new evaluation, we conducted the same 100 game match vs the basic bot version.



Figure 3.11: Multithread Search bot vs basic bot

The results are slightly worse compared to the match using the material-only evaluation, with 8 more losses than before. This may be due to the increased computational cost of evaluating these additional parameters. Furthermore, although these are abstract concepts commonly used by humans to assess positions, the engine may struggle to find a clear correlation between them and actual positional strength.

### 3.3.5. Late Move Reductions

We experiment with the use of late move pruning, under the assumption that if our move ordering is good, the best move in the position should be among the first explored moves. We reduce the depth by one unit starting from the tenth movement.

#### 3.3.5.1. Analysis

To evaluate the improvement in the new evaluation, we conducted the same 100 game match vs the basic bot version.

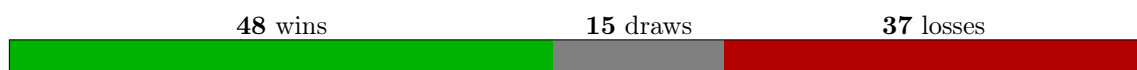


Figure 3.12: Reductions Search bot vs basic bot

The results are worse, with 15 more losses than the version without this aggressive pruning. 3.9 This could be because our move ordering is not that strong, and the best move in the position sometimes is in the last positions.

## 3.4. Additional tools and work

As it was mentioned in Section 2.6, some additional work apart from the engine implementation was done to ensure quality of the final product. This section includes some important specifications and features for the interactive board visualizer and testing workflow.

### 3.4.1. Interactive board visualizer using CustomTkinter

The board is implemented as a grid of squares, with each square capable of displaying a piece image. Piece images are loaded dynamically from the `assets` directory, which contains PNG files for each piece.

The visualizer employs an event-driven architecture, where user actions, such as clicks, trigger events handled by the `EventManager` class. When the Python script is executed, it ensures that an instance of the engine in release version executable exists, starts a new subprocess with the engine, and launches a window displaying the chessboard and interface.

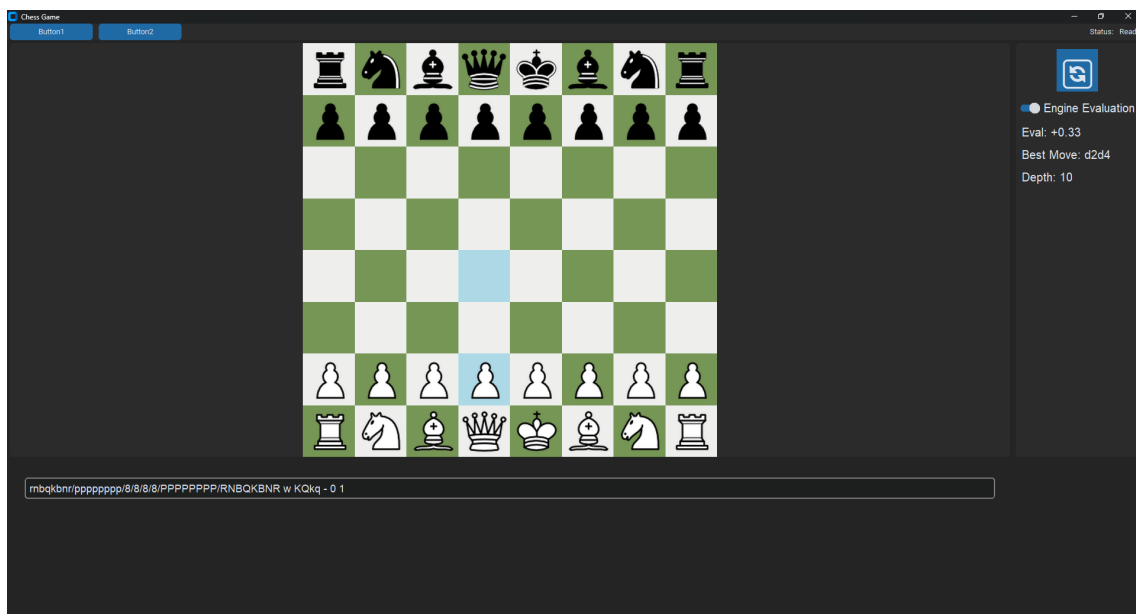


Figure 3.13: AlphaDeepChess GUI

On the right panel, there is a switch button to activate and deactivate the evaluation and suggested move in an specific depth. Also, current position's fen on the bottom panel can be modified by exchanging it to a new one.

### 3.4.2. Testing engine strength with Cutechess, Stockfish, and GitHub Actions

Cutechess and Stockfish can be downloaded from their respective webpages. Then, a JSON file must be provided to Cutechess to configure the available engines. An example of this `engines.json` file could be:

```
[
  {
    "name": "AlphaDeepChess",
    "command": "../build/release/AlphaDeepChess",
    "protocol": "uci",
    "options": {}
  },
  {
    "name": "Stockfish",
    "command": "./stockfish/stockfish/stockfish-windows-x86-64",
    "protocol": "uci",
    "options": {
      "UCI_LimitStrength": "true",
      "UCI_Elo": "1500"
    }
  }
]
```

By simply calling the Cutechess CLI with the names of the engines and additional specifications such as search time (`-st`), the number of games (`-games`), or depth (`-depth`), the matches can be executed.

Finally, automating this process of selecting the versions of the engine and Stockfish we created a workflow using a YAML file.

This workflow is triggered manually. It includes the following steps:

- Setup environment: Install dependencies such as Python packages, Cutechess and Stockfish.
- Build the engine: Compile the chess engine from source code.
- Run tests: Execute previously designed tests to validate the functionality of the engine using Cutechess CLI.
- Export generated data from Cutechess for later analysis.

This generated data are the results of the run games that includes:

- A `results.epd` with all the final positions in FEN format.

- A `results.log` with Cutechess additional detailed information like the beginning and end of a game, the summary of the game, and ELO stadistics.
- A `results.pgn` with all the played moves and games.

# Chapter 4

## Conclusions and Future Work

AlphaDeepChess has proven to be an effective implementation of classical techniques focused on alpha-beta pruning. Despite the rise of neural network-based engines, this project demonstrates that well-optimized traditional approaches remain competitive.

The engine achieved an ELO score of 1900 on Lichess, running on a Raspberry Pi 5 with a 2TB transposition table, confirming its efficiency even on limited hardware. Link to the lichess engine profile: <https://lichess.org/@/AlphaDeepChess>

Incorporating techniques such as transposition tables, iterative deepening, move ordering heuristics, and magic bitboards contributed significantly to the engine's performance. The use of Cutechess testing and comparisons with Stockfish allowed this impact to be measured.

The next steps to be implemented would be the application of neural networks (NNUE) which, although intended for CPUs, could be thought of as a streamlined evaluation with GPUs as performed by Leela Chess Zero.



# Personal contributions

## Juan Girón Herranz

- Contributed to the alpha-beta pruning algorithm. Taking part in the research and implementation of the iterative deepening and the algorithm core.
- Designed and implemented The bitboard-based move generator, then optimized the calculation of the slider pieces moves with magic bitboards technique and the PEXT hardware instruction. In addition to the integration of the move generator in the search algorithm.
- Involved in the implementation of the board data structure, with emphasis on the game state bit field design.
- Design and Developed the move data structure, which was also optimized as a bit field to reduce space consumption.
- Designed and implemented auxiliary data structures for rows, columns, diagonals, and directions. These structures played a key role in simplifying and optimizing bitboard masking operations, enabling more efficient move generation and attack pattern calculations.
- Research, designed and implemented the transposition table using Zobrist hashing, with total integration in the alpha-beta search.
- Introduced MVV-LVA (Most Valuable Victim, Least Valuable Aggressor), then research and enhanced the algorithm with killer-move heuristics.
- Research and developed the quiescence search enhancement to avoid the horizon effect in the alpha-beta pruning.
- Developed and test the search with late move reductions technique.
- Developed the tampering evaluation, adjusting the weight for the middlegame and endgame evaluations, also contributed to optimizations in the implementation for king safety and piece mobility.

- Created the algorithm to detect threefold repetition using the game's position history and its implementation in the search.
- Implemented part of the UCI command parsing and communication for engine integration.
- Conducted multiple 100-game matches using CuteChess between engine versions to measure the impact of each optimization.
- Creation of hundreds of unit tests, which have been a fundamental part of finding bugs and ensuring code quality. In addition to Perft testing of the move generator to ensure correctness of the chess engine.
- Contributed to the development of a helper GUI in Python to facilitate interactive testing of the engine, with support for the UCI protocol.
- Using Linux's `perf` tool, analyze the CPU overhead of the different parts of the chess engine.
- Compiled and deployed the engine on a Raspberry Pi 5, configuring it as a Lichess.org bot. Running under limited hardware resources, the engine achieved competitive ELO ratings while demonstrating our code's efficiency and portability.

## Yi Wang Qiu

- Responsible for the architectural design and full implementation of the alpha-beta pruning algorithm, established as the foundational search technique of the engine. This algorithm was enhanced through iterative refinement, such as aspiration windows, and theoretical benchmarking, enabling effective traversal of the game tree while significantly reducing the computational overhead associated with brute-force minimax strategies.
- Developed an optimized multithreaded search version incorporating the Young Brothers Wait Concept (YBWC), a parallelization paradigm specifically tailored for game tree evaluation. This technique allows the principal variation to be explored first in a sequential way, deferring sibling node evaluations to parallel workers only after the most promising path has been examined. The principal variation refers to the first move after generated legal moves were ordered.
- Engineered the parsing and command interpretation system compliant with the Universal Chess Interface (UCI) protocol. This subsystem ensures seamless bidirectional communication between the chess engine and external graphical user interfaces, testing suites, and benchmarking frameworks.
- Designed and implemented the core engine abstractions, `Square` and `Board` classes, which supports the representation and manipulation of chess positions. These classes encapsulate critical logic such as coordinate translation



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or position translation from FEN, piece tracking, castling rights, and en passant possibilities, all integrated with a bitboard backend. This design allows high-level readability while preserving low-level computational performance.

- Constructed the internal board representation model using 64-bit bitboards. This representation supports highly efficient binary operations such as masking, shifting, and logical conjunctions to simulate piece movement and board updates. Bitboards were used extensively to implement both legal move generation and tactical evaluation routines, leading to a compact and performant engine state.
- Developed a modular and extensible evaluation system capable of quantifying chess positions through multiple heuristic lenses. The implemented strategies range from basic material balance (expressed in centipawns) to more sophisticated models that incorporate positional features such as game phase, piece activity, mobility scoring, and king vulnerability or safety. These heuristics were designed to be dynamically weighted depending on the stage of the game (opening, middlegame, or endgame).
- Integrated and calibrated precomputed positional data structures, including piece-square tables to accelerate the static evaluation of positions.
- Designed and authored a suite of automated Python scripts for orchestrating engine versus engine tournaments and performance benchmarking using Cutchess CLI. These scripts included configurable match parameters like search time, depth of search, number of games, book of openings or initial positions. They were essential in enabling the reproducibility of experiments, comparison of successive versions of the engine, and quantification of the impact of algorithmic refinements.
- Established a robust continuous integration and delivery (CI/CD) pipeline using GitHub Actions. This infrastructure automated the build, deployment, and testing stages of the engine. Used the above python scripts to automate the tournaments in an independent machine to avoid wasting time and computation capacity while still developing.
- Contributed to the frontend layer of the project by prototyping a graphical user interface (GUI) in Python, designed to allow interactive execution of the engine in a visual environment. The GUI included subprocess communication features, move display, and optional positional evaluations. Although later iterations focused on headless execution, this interface was key during early debugging and demonstration phases.
- Authored detailed and structured online documentation describing the engine's internal architecture, modular hierarchy, function-level responsibilities, and usage guidelines. The documentation was designed not only as an educational resource for future contributors, but also as a formal exposition of the system's logic for academic evaluation purposes like this exact document. It includes

illustrative diagrams or graphs, code references, and configuration examples to support transparency and reproducibility.

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