

## Parcial 2

1)

linea	nombre del estado	estado/guardas	aclaracion
-	$\sigma_0$	$x \mapsto 20, r \mapsto 21, y \mapsto 22$	Estado inicial
$\ell_1$	$\sigma_1$	$x \mapsto 20, r \mapsto 0, y \mapsto 1$	
$\ell_2$		<i>True</i>	
$\ell_3$	$\sigma_2$	$x \mapsto 20, r \mapsto 1, y \mapsto 2$	
$\ell_2$		<i>True</i>	
$\ell_3$	$\sigma_3$	$x \mapsto 20, r \mapsto 2, y \mapsto 4$	
$\ell_2$		<i>True</i>	
$\ell_3$	$\sigma_4$	$x \mapsto 20, r \mapsto 3, y \mapsto 8$	
$\ell_2$		<i>True</i>	
$\ell_3$	$\sigma_5$	$x \mapsto 20, r \mapsto 4, y \mapsto 16$	
$\ell_2$		<i>True</i>	
$\ell_3$	$\sigma_6$	$x \mapsto 20, r \mapsto 5, y \mapsto 32$	
$\ell_2$		<i>False</i>	
$\ell_5$	$\sigma_7$	$x \mapsto 20, r \mapsto 5, y \mapsto 12$	Estado final

2)

$$\begin{aligned}
 & \langle \Sigma i : 0 \leq i < N/2 : A.(2 * i) \rangle \\
 & \equiv \{A=[1,2,3,4,5,6]\} \\
 & \langle \Sigma i : 0 \leq i < 3 : A.(2 * i) \rangle \\
 & \equiv \{\text{Evaluo rango}\} \\
 & \langle \Sigma i : i \in \{0, 1, 2\} : A.(2 * i) \rangle \\
 & \equiv \{\text{Evaluo termino con el rango}\} \\
 & A.(2 * 0) + A.(2 * 1) + A.(2 * 2) \\
 & \equiv \{\text{Aritmetica}\} \\
 & A.(0) + A.(2) + A.(4)
 \end{aligned}$$

3)

$$\begin{aligned}
& (n \bmod 2 \equiv 0 \vee n \bmod 3 \equiv 0) \\
& \wedge (n \bmod 2 \equiv 0 \Rightarrow wp.(r, n := r + 1, n + 1).(r = \langle N \ i : 0 \leq i < n : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle)) \\
& \wedge (n \bmod 3 \equiv 0 \Rightarrow wp.(r, n := r + 1, n + 1).(r = \langle N \ i : 0 \leq i < n : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle)) \\
& \equiv \{\text{Suponemos antecedente y demostramos consecuente}\} \\
& (n \bmod 2 \equiv 0 \vee n \bmod 3 \equiv 0) \\
& \wedge wp.(r, n := r + 1, n + 1).(r = \langle N \ i : 0 \leq i < n : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle) \\
& \wedge wp.(r, n := r + 1, n + 1).(r = \langle N \ i : 0 \leq i < n : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle) \\
& \equiv \{\text{wp asignacion}\} \\
& (n \bmod 2 \equiv 0 \vee n \bmod 3 \equiv 0) \\
& \wedge (r + 1 = \langle N \ i : 0 \leq i < n + 1 : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle) \\
& \wedge (r + 1 = \langle N \ i : 0 \leq i < n + 1 : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle) \\
& \equiv \{\text{Logica}\} \\
& r = \langle N \ i : 0 \leq i < n : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle
\end{aligned}$$

d)

$$\begin{aligned}
2 &= \langle N \ i : 0 \leq i < 3 : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle \\
&\equiv \{\text{Evaluó rango}\} \\
2 &= \langle N \ i : i \in \{0, 1, 2\} : i \bmod 2 \equiv 0 \vee i \bmod 3 \equiv 0 \rangle \\
&\equiv \{\text{Evaluó termino en el rango}\} \\
2 &= (0 \bmod 2 \equiv 0 \vee 0 \bmod 3 \equiv 0) + (1 \bmod 2 \equiv 0 \vee 1 \bmod 3 \equiv 0) + (2 \bmod 2 \equiv 0 \vee 2 \bmod 3 \equiv 0) \\
&\equiv \{\text{Evaluó}\} \\
2 &= 1 + 0 + (True \vee 2 \bmod 3 \equiv 0) \\
&\equiv \{\text{Logica}\} \\
2 &= 1 + 0 + 1 \\
&\equiv \{\text{Aritmetica}\} \\
2 &= 2 \\
&\equiv \{\text{True}\}
\end{aligned}$$

4)

### Encontrar Invariante

$$I : 0 \leq n \leq N-1 \wedge u = \text{sum}.A.0.n \wedge r = \langle \exists i : 0 \leq i < n : \text{sum}.A.0.i = 3 \wedge A.i = 3 \rangle$$

## Inicializacion

$\{N > 0\}$   
 $r, u, n := False, 0, 0;$   
 $\{0 \leq n \leq N - 1 \wedge u = sum.A.0.n \wedge r = \langle \exists i : 0 \leq i < n : sum.A.0.i = 3 \wedge A.i = 3 \rangle\}$

## Finalizacion

$0 \leq n \leq N - 1 \wedge u = sum.A.0.n \wedge r = \langle \exists i : 0 \leq i < n : sum.A.0.i = 3 \wedge A.i = 3 \rangle \wedge \neg B$   
 $\Rightarrow r = \langle \exists i : 0 \leq i < N : sum.A.0.i = 3 \wedge A.i = 3 \rangle$