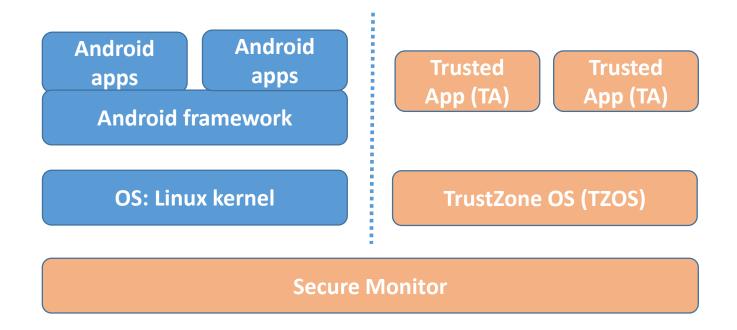
# PartEmu: Enabling Dynamic Analysis of Real-World TrustZone Software Using Emulation

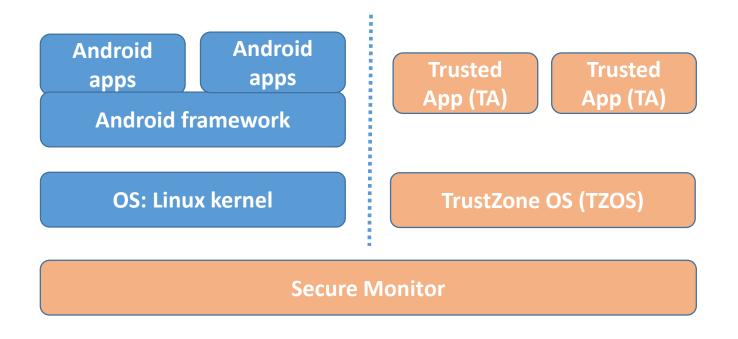
Lee Harrison, Hayawardh Vijayakumar, Michael Grace

Samsung Knox, Samsung Research America

Rohan Padhye, Koushik Sen

EECS Department, University of California, Berkeley





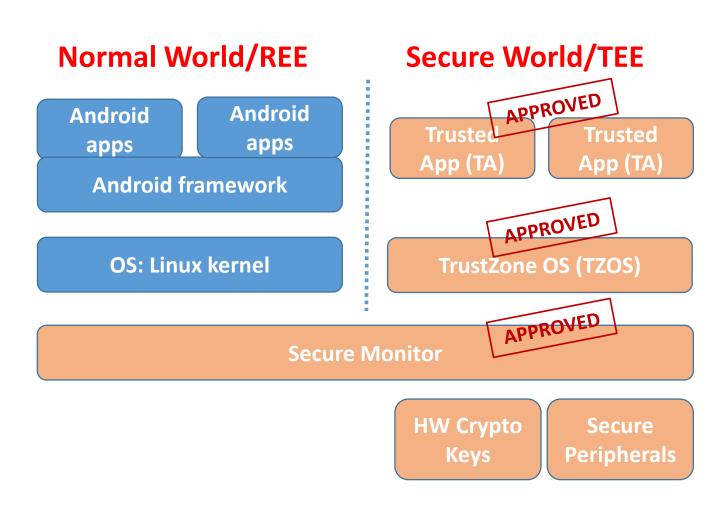
- Separate software stack
  - Trusted applications (TAs)
  - TrustZone OS (TZOS)

# Android apps apps Android framework OS: Linux kernel Secure World/TEE Trusted App (TA) Trusted App (TA) TrustZone OS (TZOS) Secure World/TEE

- Separate software stack
  - Trusted applications (TAs)
  - TrustZone OS (TZOS)
  - TEE/REE

#### **Secure World/TEE** Normal World/REE **Android Android Trusted** Trusted apps apps App (TA) **Android framework OS: Linux kernel** TrustZone OS (TZOS) **Secure Monitor HW Crypto** Secure **Peripherals** Keys

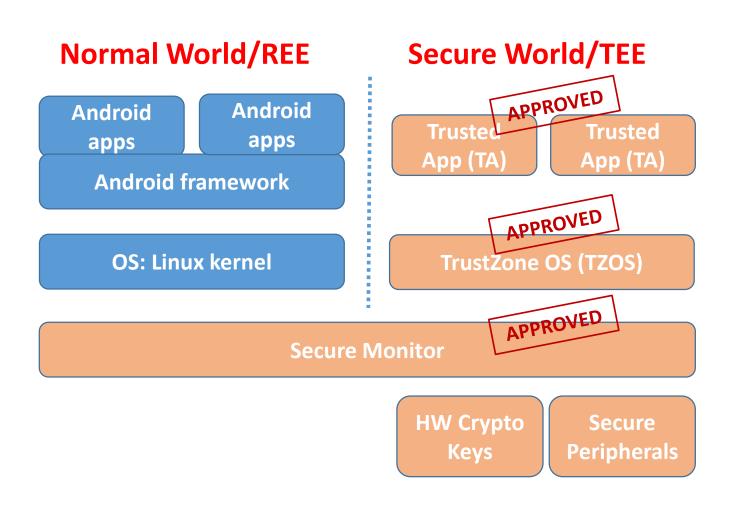
- Separate software stack
  - Trusted applications (TAs)
  - TrustZone OS (TZOS)
  - TEE/REE
- Basis for security: Has access to hardware keys

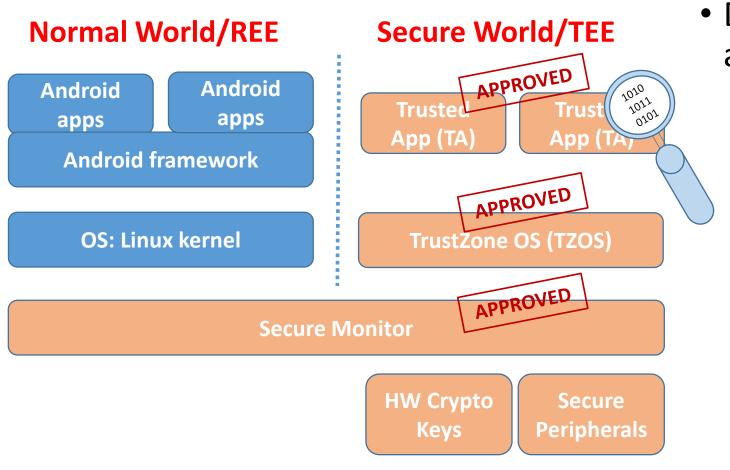


- Separate software stack
  - Trusted applications (TAs)
  - TrustZone OS (TZOS)
  - TEE/REE
- Basis for security: Has access to hardware keys
- Access to TZ locked down: Only signed software can run

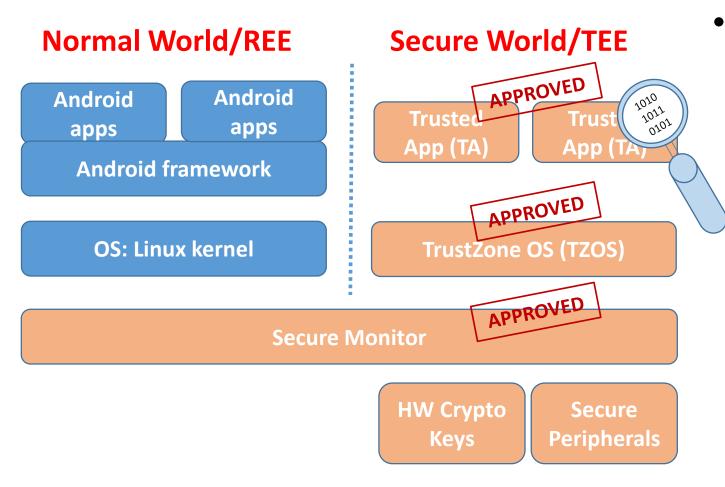
Approach

Results: What did we learn?

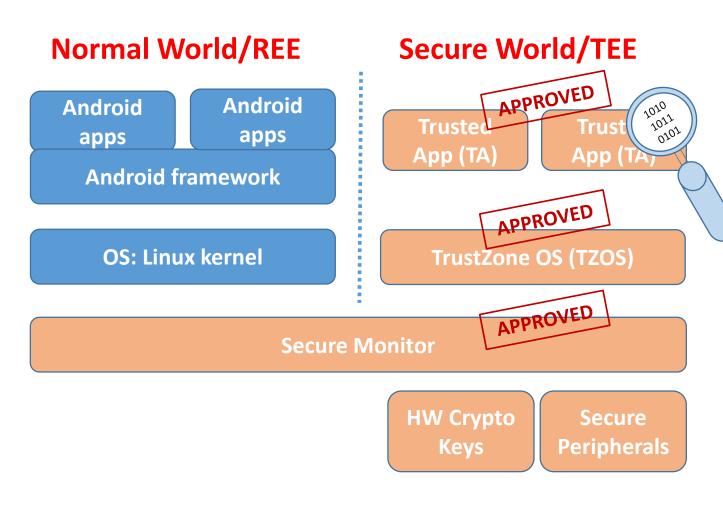




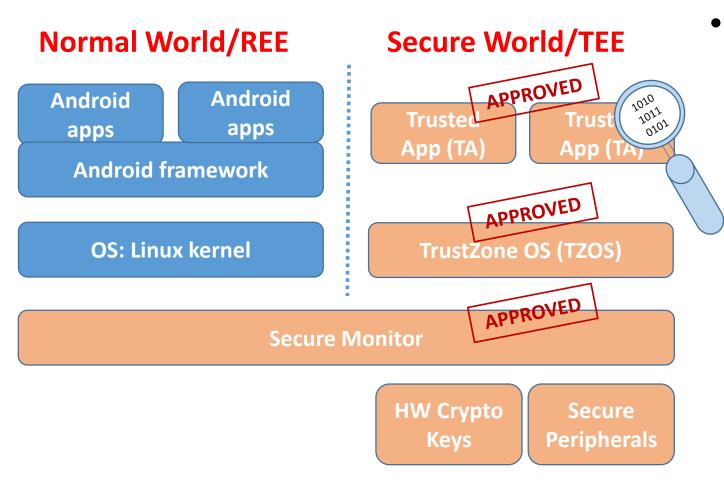
 Dynamic analysis needs ability to monitor target



- Dynamic analysis needs ability to monitor target
  - Debugging needs memory/registers
  - Feedback-driven fuzz testing – needs list of basic blocks covered



- Dynamic analysis needs ability to monitor target
  - Debugging needs memory/registers
  - Feedback-driven fuzz testing – needs list of basic blocks covered
- However, cannot instrument TZ software or monitor TZ memory due to signing!



- Prior dynamic analysis approaches limited!
  - TA/TZOS binary reverse engineering
  - Fuzz testing without feedback

# Solution: Dynamic Analysis By Emulation

We build an emulator that runs real-world TZOSes and TAs

# Solution: Dynamic Analysis By Emulation

- We build an emulator that runs real-world TZOSes and TAs
- Emulation enables dynamic analysis
  - Allows introspection and monitoring of TZ execution

# Solution: Dynamic Analysis By Emulation

- We build an emulator that runs real-world TZOSes and TAs
- Emulation enables dynamic analysis
  - Allows introspection and monitoring of TZ execution
- We support four widely-used real-world TZOSes:
  - Qualcomm's QSEE
  - Trustonic's Kinibi
  - Samsung's TEEGRIS
  - Linaro's OP-TEE

Approach: How did we run TZ in an emulator?

Results: What did we learn?

**Android Apps** 

**Android FW** 

**TEE Userspace** 

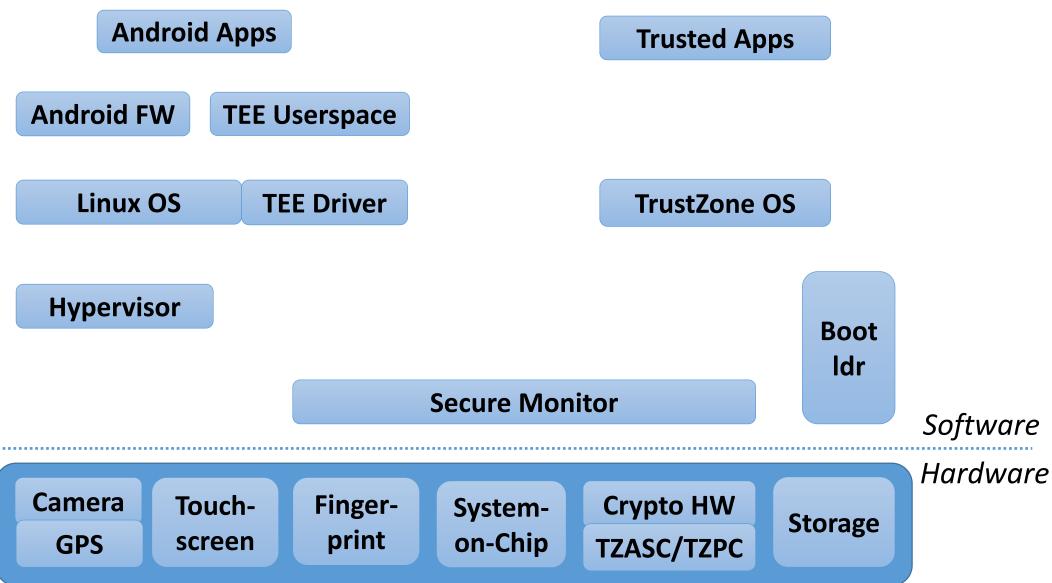
**Linux OS** 

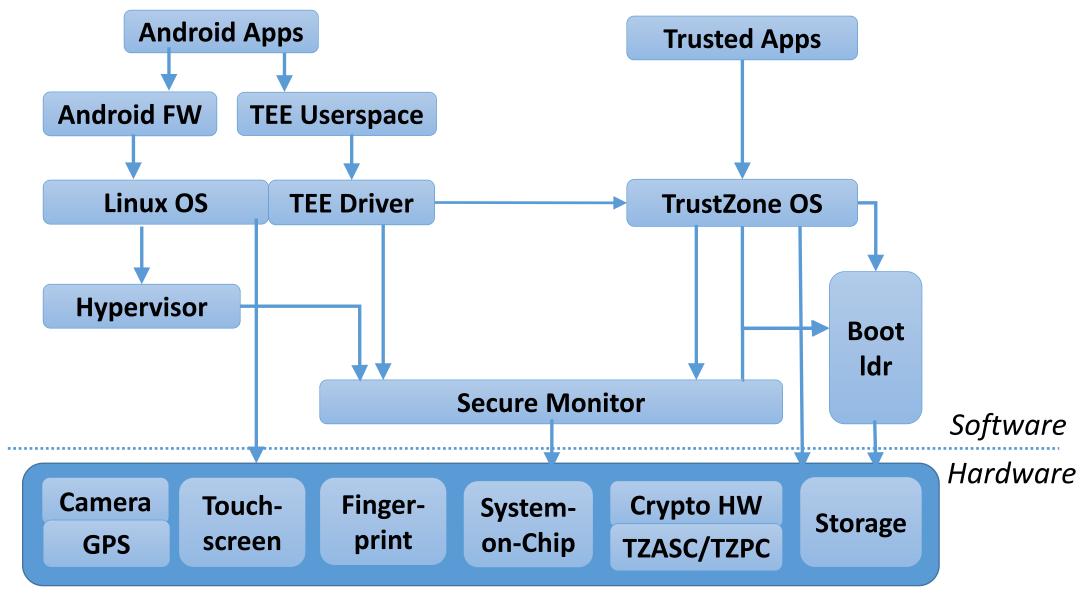
**TEE Driver** 

**Hypervisor** 

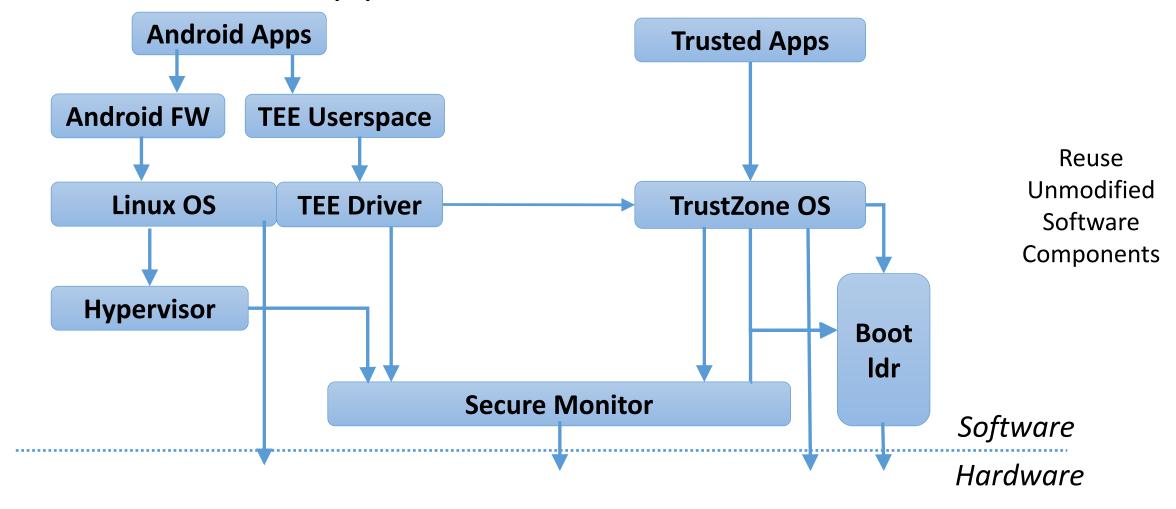
Software

**Android Apps Trusted Apps Android FW TEE Userspace Linux OS TEE Driver TrustZone OS Hypervisor Boot** ldr **Secure Monitor** 

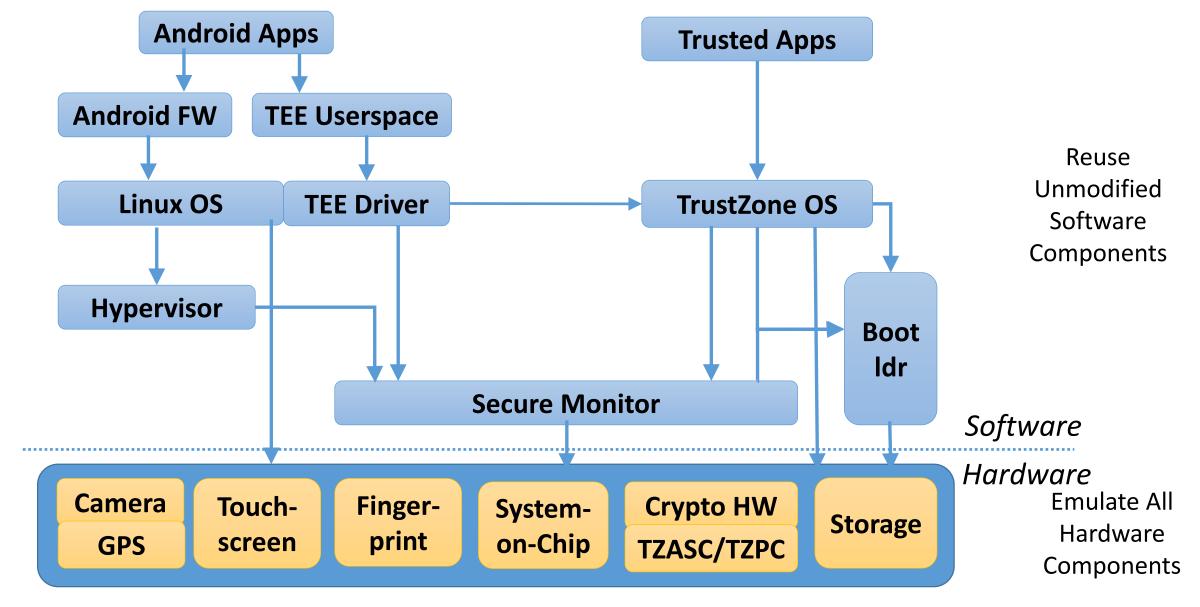




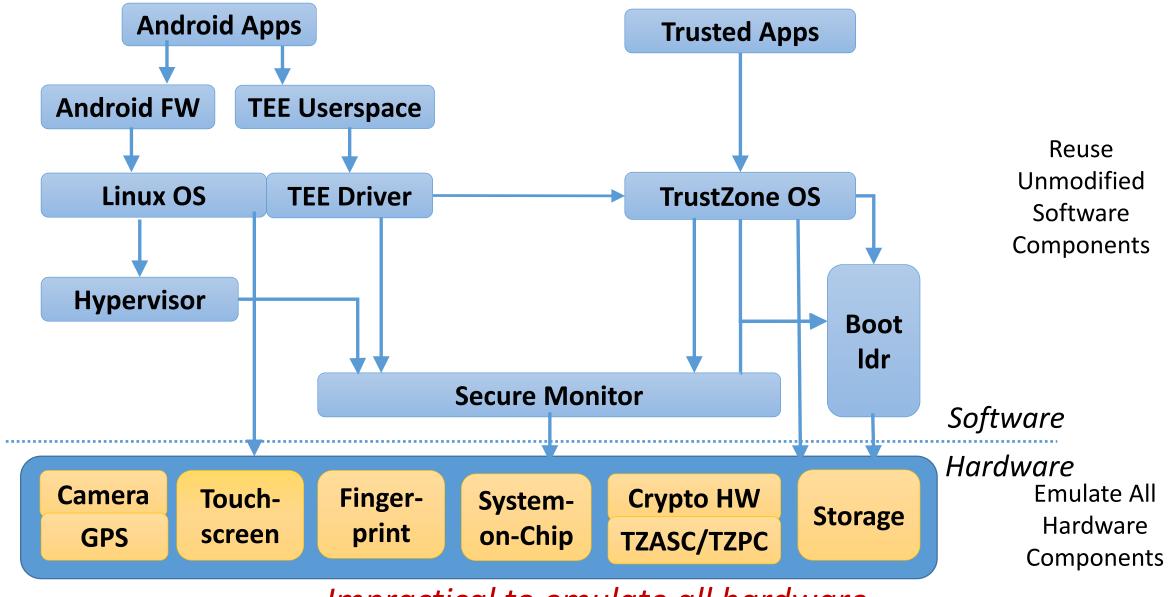
# Traditional Approach: Emulate all HW



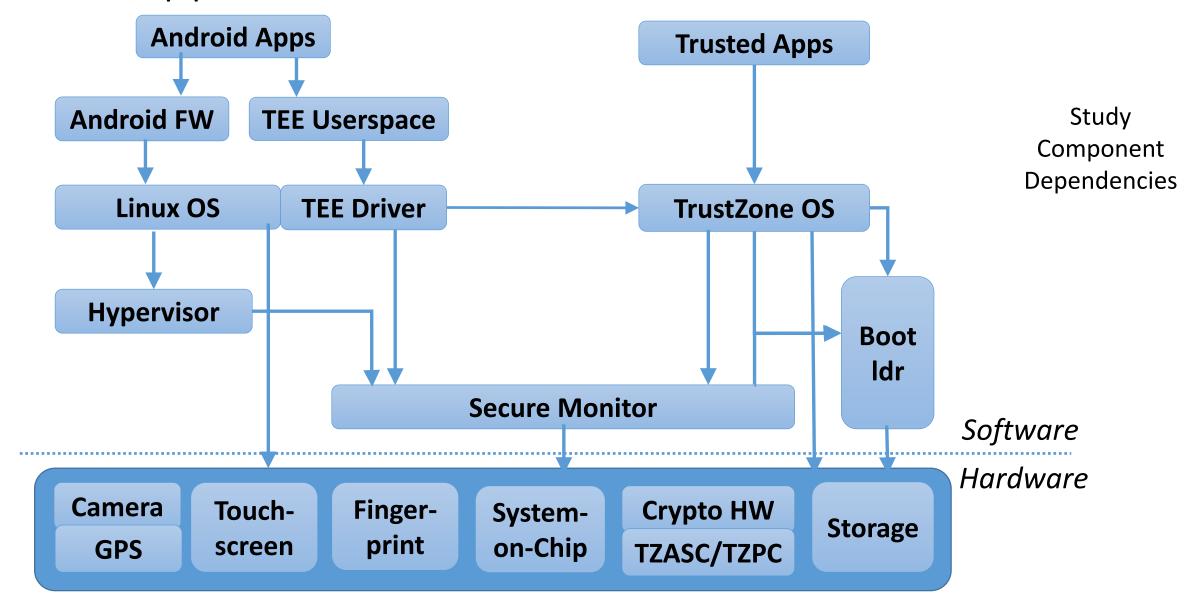
# Traditional Approach: Emulate all HW

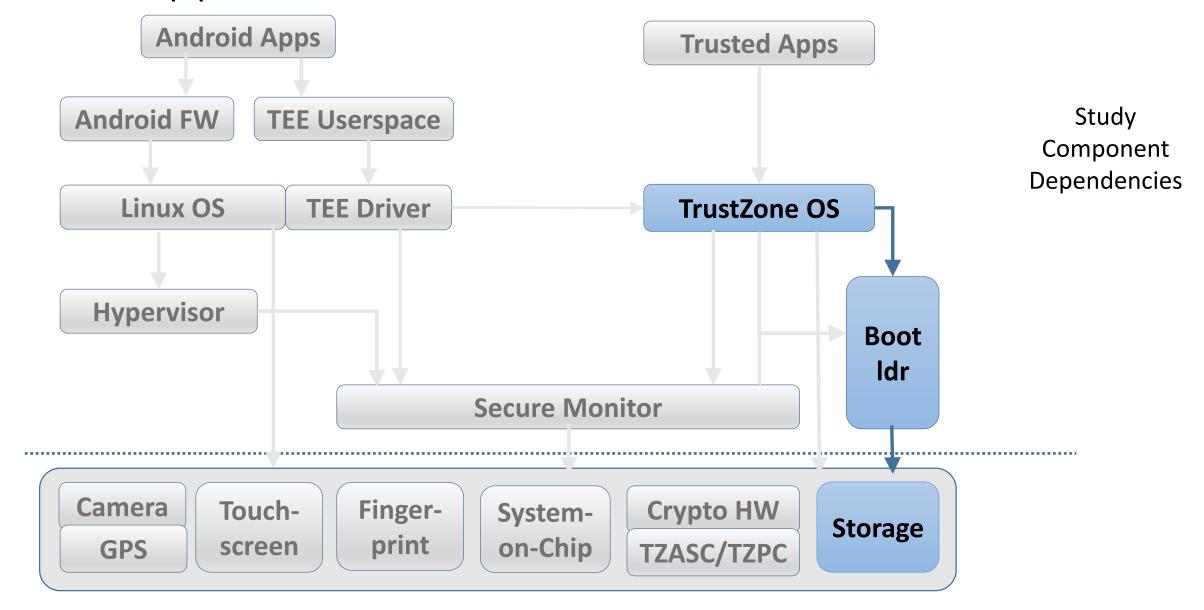


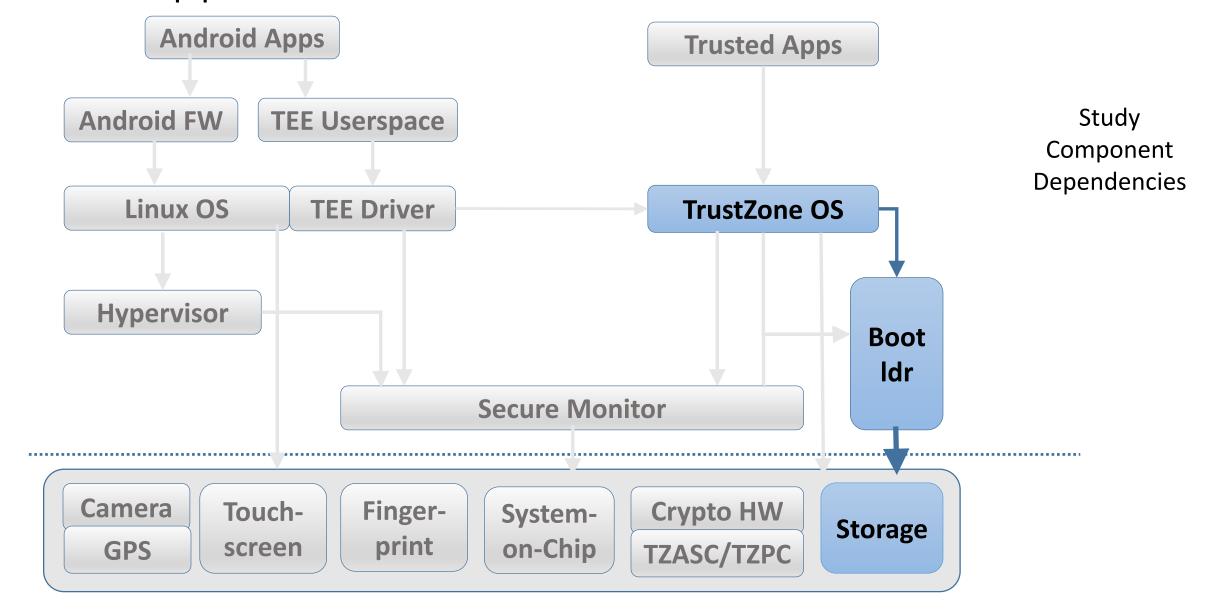
# Traditional Approach: Emulate all HW

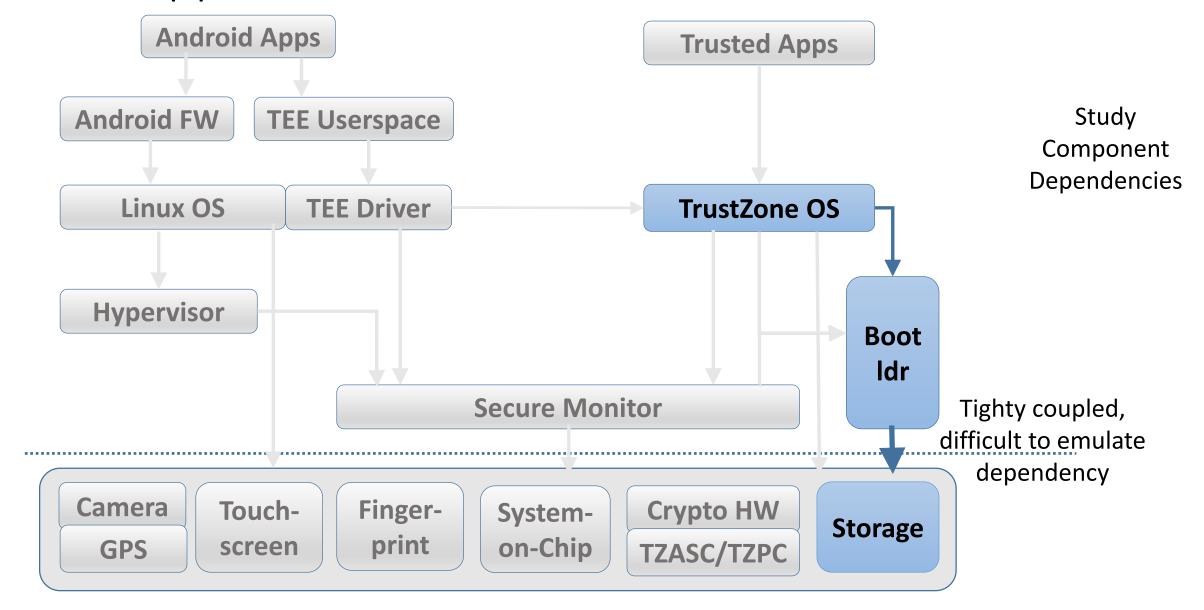


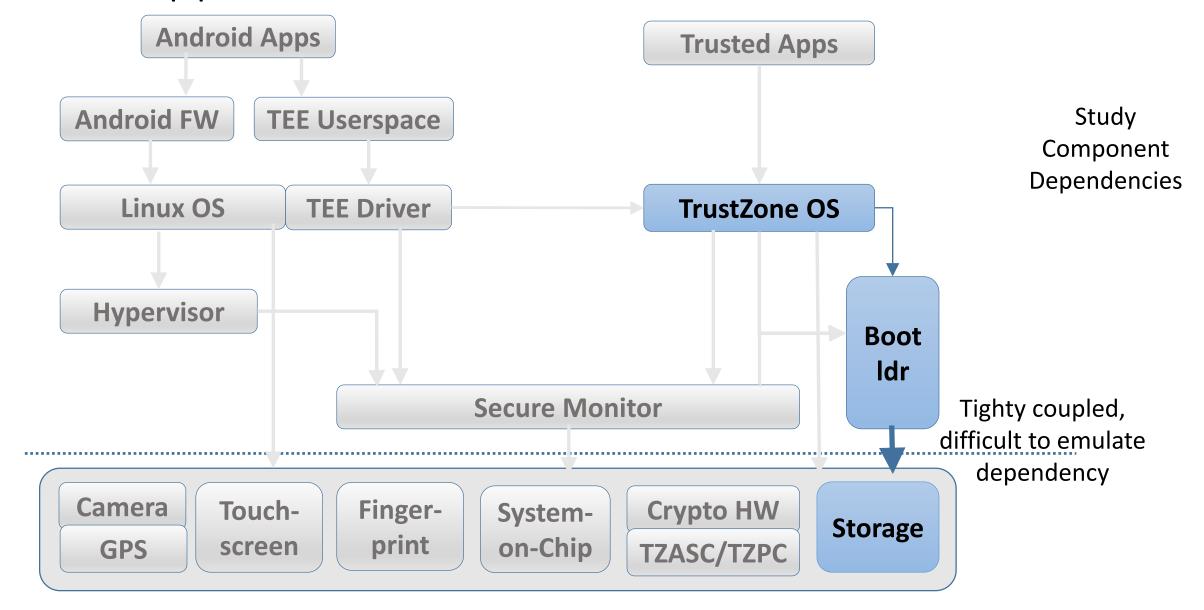
Impractical to emulate all hardware

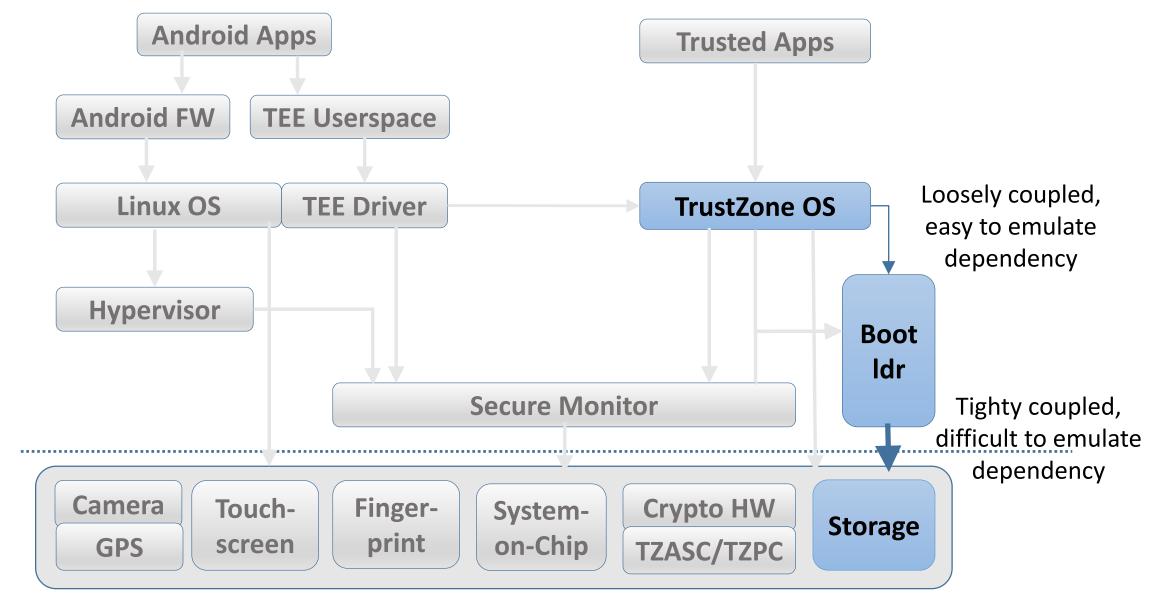


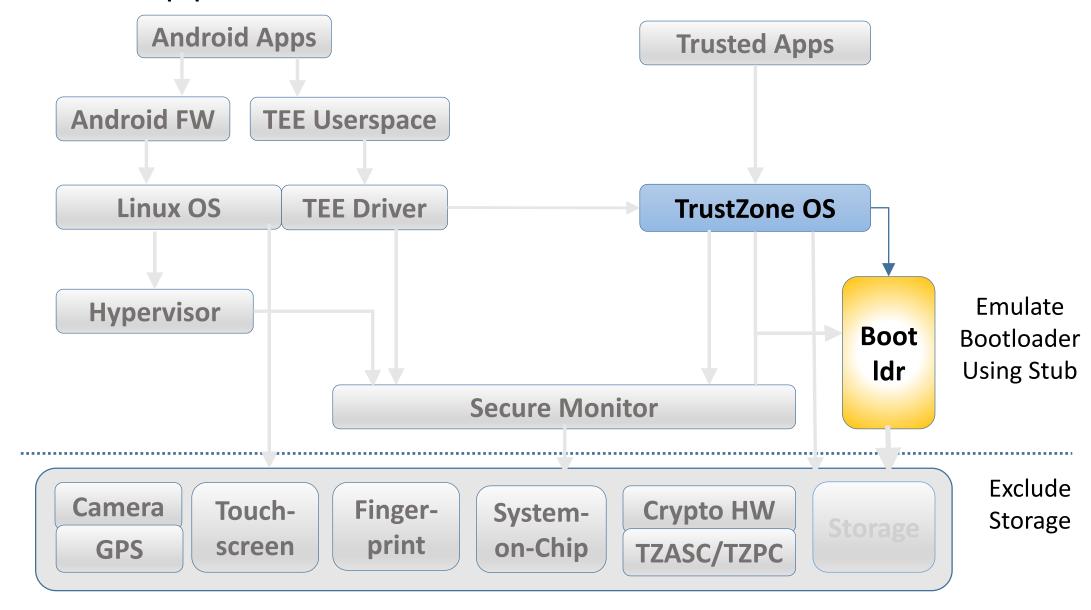


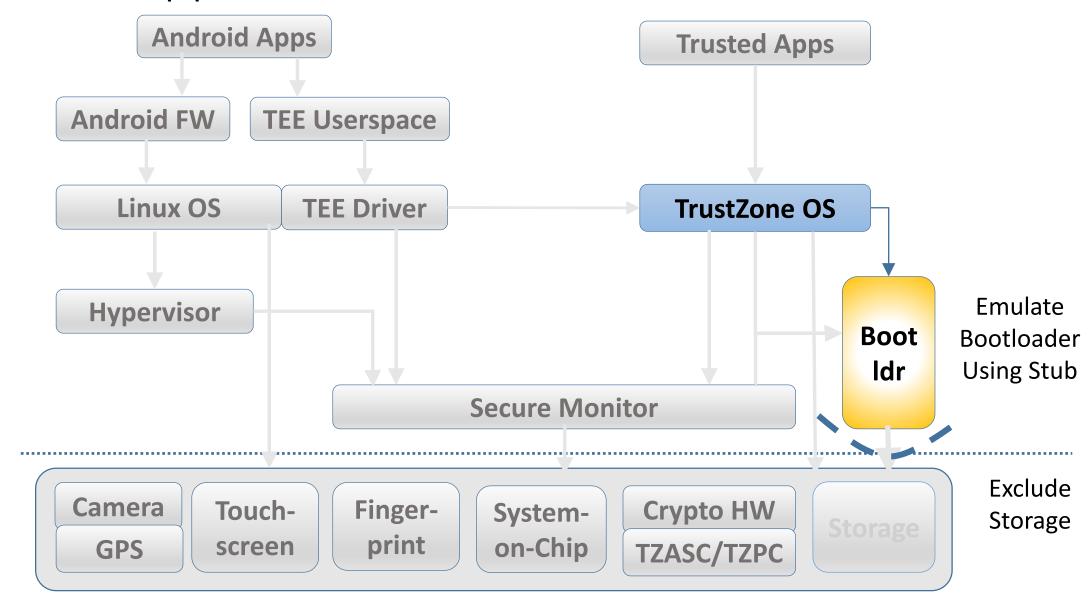


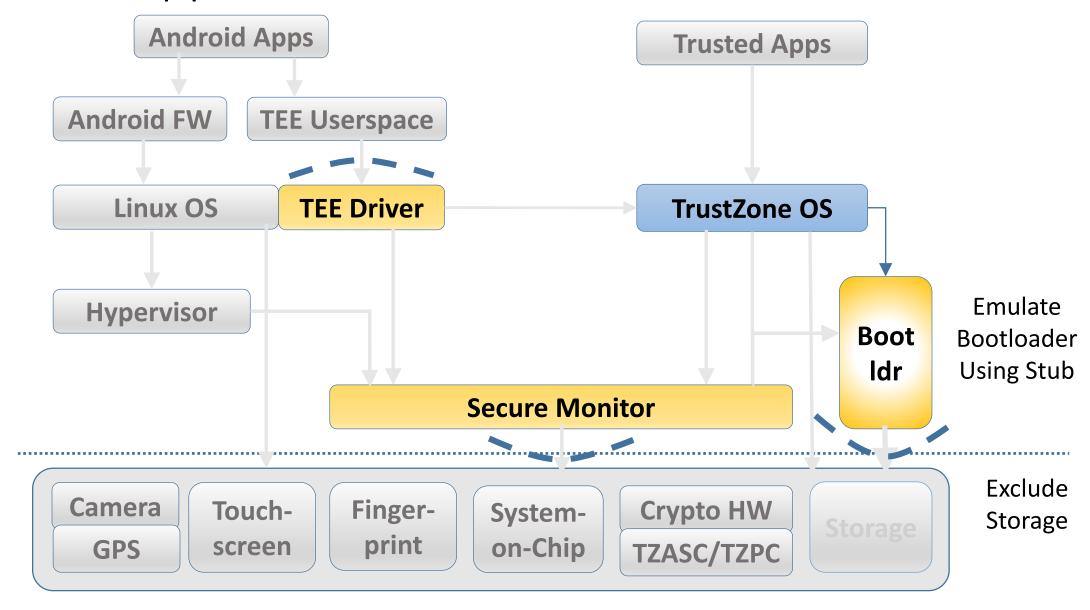












# Emulation Effort Feasible Using Patterns

#### Emulation Effort Feasible Using Patterns

• Patterns to Emulate Hardware (MMIO Loads and Stores)

```
# Constant read (CONSTANT_READ_REG)
v = read(CONSTANT_READ_REG);
if (v != VALID VALUE)
    fail();
# Read-write (READ WRITE REG)
write (READ WRITE \overline{REG}, \overline{v1});
v2 = read(READ_WRITE_REG);
if (v2 != v1)
    fail();
# Increment (INCR REG)
v = read(INCR REG);
if (read(INCR REG) < v)
    fail();
# Poll (POLL REG)
while (read(POLL_REG) != READY);
```

```
# Random (RAND REG)
v1 = read(RAND REG)
v2 = read(RAND REG)
if (v1 == v2)
   fail();
# Shadow (SHADOW REG1, SHADOW REG2)
# Commit (COMMIT REG)
# Target (TARGET REG1, TARGET REG2)
write(SHADOW_REG1, v1)
write (SHADOW REG2, v2)
write (COMMIT REG, COMMIT VALUE)
v3 = read(TARGET REG1)
v4 = read(TARGET REG2)
if ((v1 != v3) \text{ or } (v2 != v4))
    fail();
```

Patterns to Emulate Hardware (MMIO Loads and Stores)

```
# Constant read (CONSTANT_READ_REG)
v = read(CONSTANT_READ_REG);
if (v != VALID VALUE)
    fail();
# Read-write (READ WRITE REG)
write (READ WRITE REG, v1);
v2 = read(READ_WRITE_REG);
if (v2 != v1)
    fail();
# Increment (INCR REG)
v = read(INCR REG);
if (read(INCR REG) < v)
    fail();
# Poll (POLL REG)
while (read(POLL_REG) != READY);
```

```
# Random (RAND REG)
v1 = read(RAND REG)
v2 = read(RAND REG)
if (v1 == v2)
   fail();
# Shadow (SHADOW REG1, SHADOW REG2)
# Commit (COMMIT REG)
# Target (TARGET REG1, TARGET REG2)
write(SHADOW_REG1, v1)
write (SHADOW REG2, v2)
write (COMMIT REG, COMMIT VALUE)
v3 = read(TARGET REG1)
v4 = read(TARGET REG2)
if ((v1 != v3) \text{ or } (v2 != v4))
   fail();
```

Patterns to Emulate Hardware (MMIO Loads and Stores)

```
# Constant read (CONSTANT READ REG)
v = read(CONSTANT READ REG);
if (v != VALID VALUE)
   fail();
# Read-write (READ WRITE REG)
write (READ WRITE REG, v1);
v2 = read(READ_WRITE_REG);
if (v2 != v1)
   fail();
# Increment (INCR REG)
v = read(INCR REG);
if (read(INCR REG) < v)
    fail();
# Poll (POLL REG)
while (read(POLL_REG) != READY);
```

```
# Random (RAND REG)
v1 = read(RAND REG)
v2 = read(RAND REG)
if (v1 == v2)
   fail();
# Shadow (SHADOW REG1, SHADOW REG2)
# Commit (COMMIT REG)
# Target (TARGET REG1, TARGET REG2)
write (SHADOW REG1, v1)
write (SHADOW REG2, v2)
write (COMMIT REG, COMMIT VALUE)
v3 = read(TARGET REG1)
v4 = read(TARGET REG2)
if ((v1 != v3) \text{ or } (v2 != v4))
   fail();
```

Patterns to Emulate Software APIs

Category	Difficulty	K	Q	T	0			
Emulated Boot Information Structure								
Constants	Low	13	8	2	3			
Any value	Low	1	3	0	0			
Simple value	Low	2	1	14	2			
Complex values	High	2	1 <sup>[note a]</sup>	0	0			
Total	-	18	13	16	5			
Emulated Secure Monitor Calls <sup>[note b]</sup>								
Return simple value	Low	0	-	3	-			
Return constant	Low	1	-	5	-			
Store/retrieve values	Low	1	-	2	-			
Control transfer	High	3	-	2	-			
Total	-	5	-	12	-			

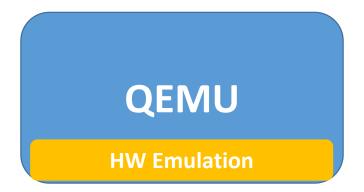
Patterns to Emulate Software APIs

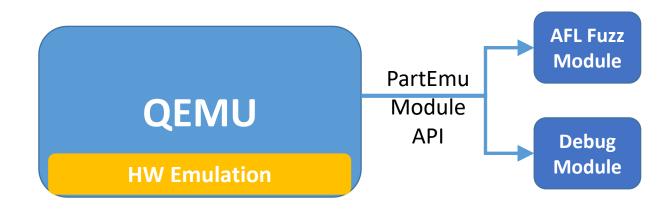
Category	Difficulty	K	Q	T	O	
Emulated Boot Information Structure						
Constants	Low	13	8	2	3	
Any value	Low	1	3	0	0	
Simple value	Low	2	1	14	2	
Complex values	High	2	1[note a]	0	0	
Total	-	18	13	16	5	
Emulated	Secure Monit	or Ca	lls <sup>[note b]</sup>			
Return simple value	Low	0	-	3	-	
Return constant	Low	1	-	5	-	
Store/retrieve values	Low	1	-	2	-	
Control transfer	High	3	-	2	-	
Total	-	5	-	12	-	
						-

Patterns to Emulate Software APIs

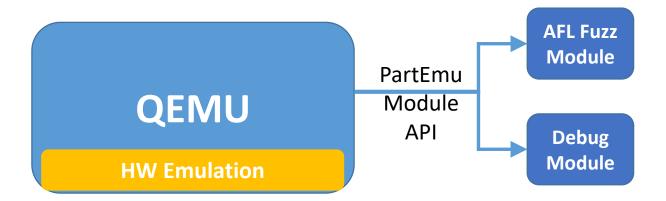
Category	Difficulty	K	Q	T	O			
Emulated Boot Information Structure								
Constants	Low	13	8	2	3			
Any value	Low	1	3	0	0			
Simple value	Low	2	1	14	2			
Complex values	High	2	1[note a]	0	0			
Total	-	18	13	16	5			
Emulated Secure Monitor Calls <sup>[note b]</sup>								
Return simple value	Low	0	-	3	-			
Return constant	Low	1	-	5	-			
Store/retrieve values	Low	1	-	2	_			
Control transfer	High	3	-	2	-			
Total	-	5	-	12	-			

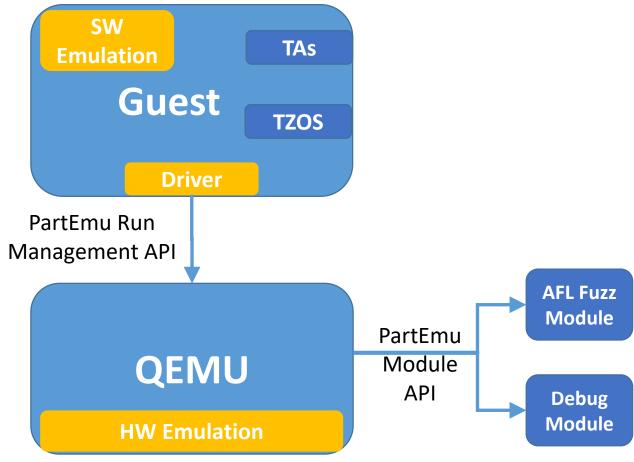












Problem: Dynamic analysis of TZ is hard!

Approach: How did we run TZ in an emulator?

Results: What did we learn?

16 Firmware Images

16 Firmware Images

12 Smartphone / loT vendors

16 Firmware Images

12 Smartphone / loT vendors

196 Unique TAs

16 Firmware Images

12 Smartphone / loT vendors

196 Unique TAs

AFL Crashed
48 TAs

16 Firmware Images

12 Smartphone / loT vendors

**196** Unique TAs

AFL Crashed 48 TAs

Found TZ-specific coding anti-patterns that led to crashes

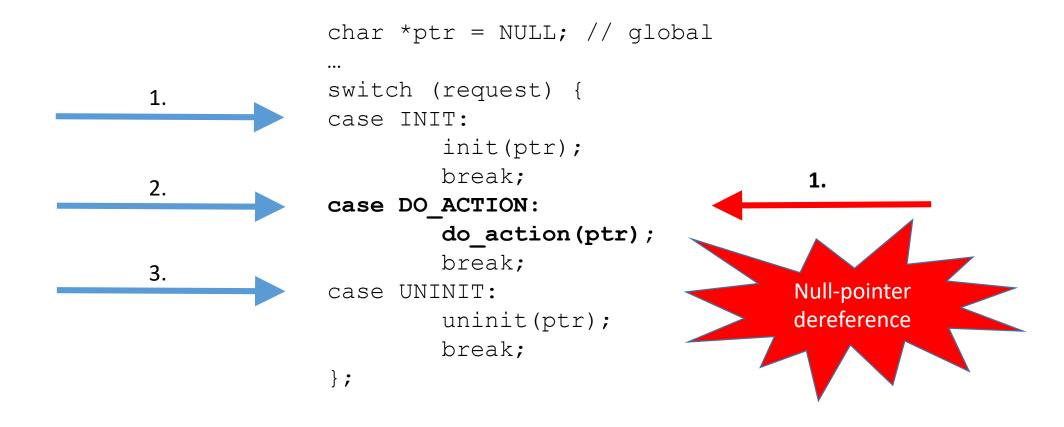
```
char *ptr = NULL; // global
switch (request) {
case INIT:
        init(ptr);
        break;
case DO ACTION:
        do action(ptr);
        break;
case UNINIT:
        uninit(ptr);
        break;
};
```

```
char *ptr = NULL; // global
switch (request) {
case INIT:
        init(ptr);
        break;
case DO ACTION:
        do action(ptr);
        break;
case UNINIT:
        uninit(ptr);
        break;
};
```

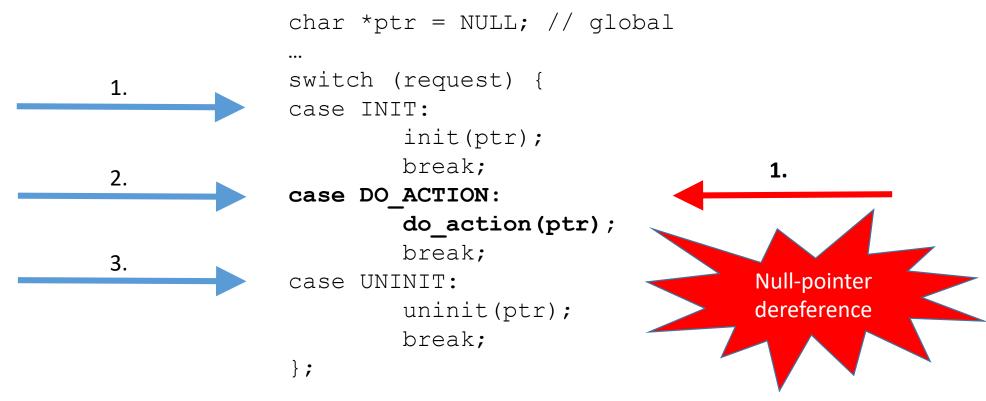
```
char *ptr = NULL; // global
switch (request) {
case INIT:
        init(ptr);
        break;
case DO ACTION:
        do action(ptr);
        break;
case UNINIT:
        uninit(ptr);
        break;
};
```

```
char *ptr = NULL; // global
switch (request) {
case INIT:
        init(ptr);
        break;
case DO ACTION:
        do action(ptr);
        break;
case UNINIT:
        uninit(ptr);
        break;
};
```

```
char *ptr = NULL; // global
switch (request) {
case INIT:
        init(ptr);
        break;
                                    1.
case DO ACTION:
        do action(ptr);
        break;
case UNINIT:
        uninit(ptr);
        break;
};
```



• TAs split work into small units → receive a sequence of requests



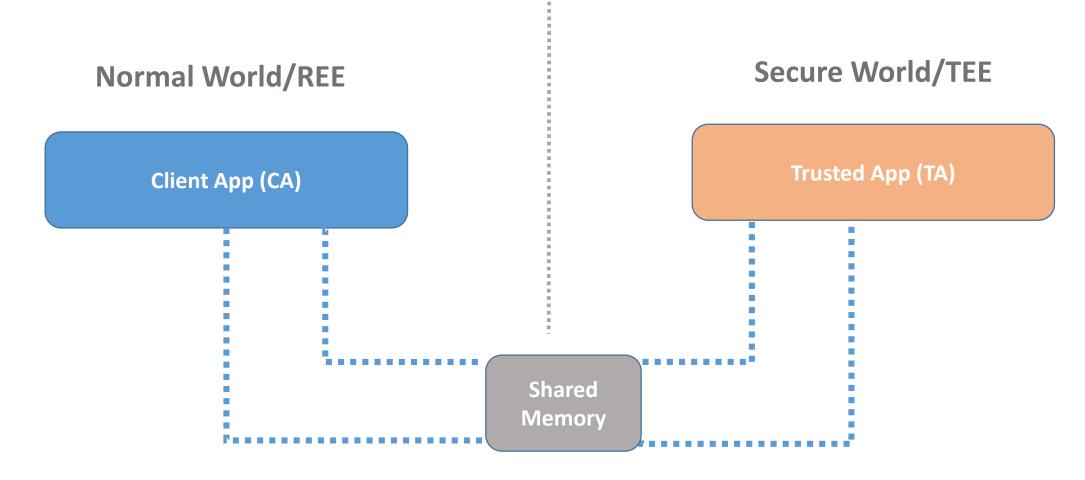
TA should properly handle any sequence of requests from CA

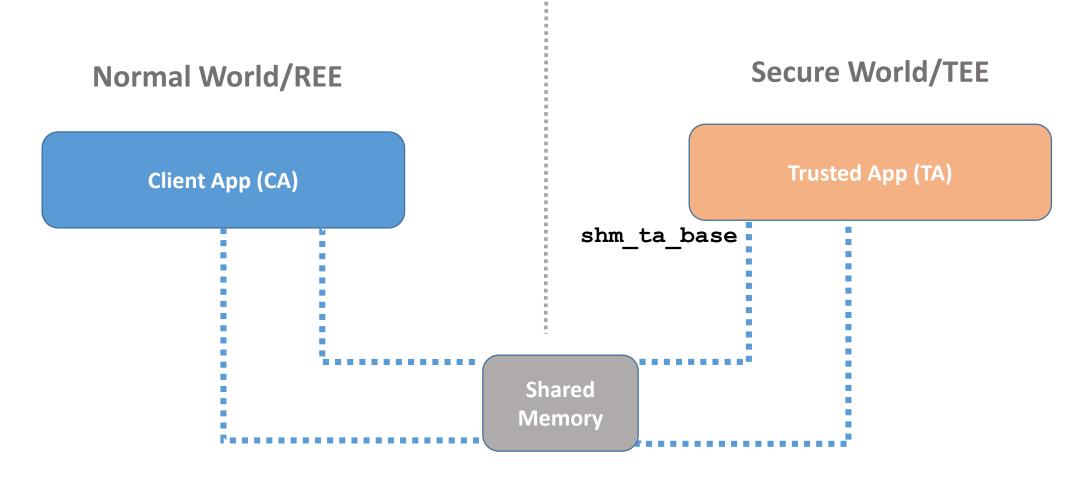
Normal World/REE

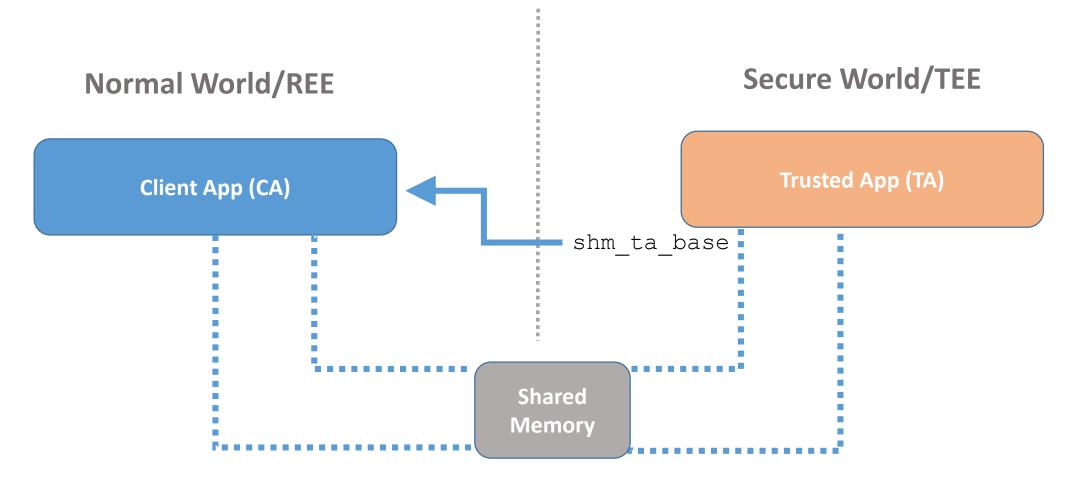
Client App (CA)

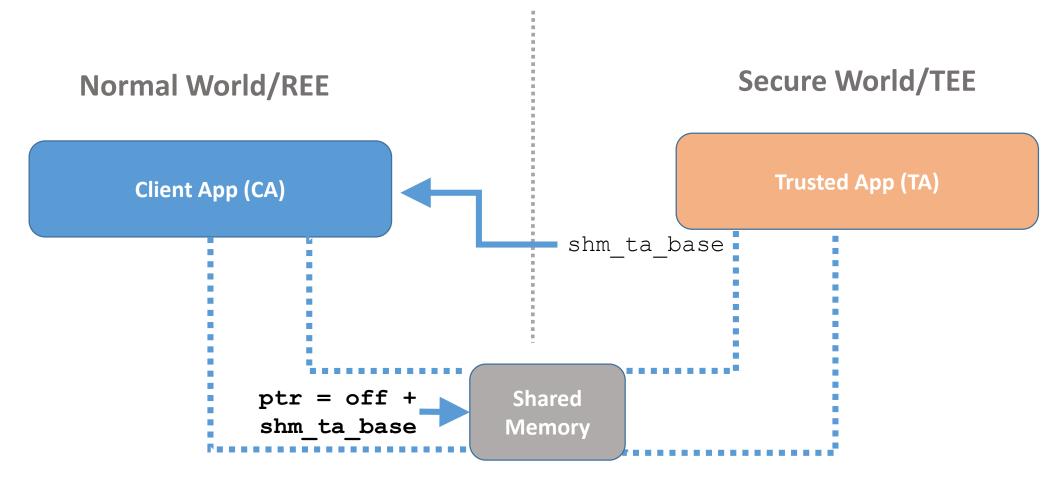
**Secure World/TEE** 

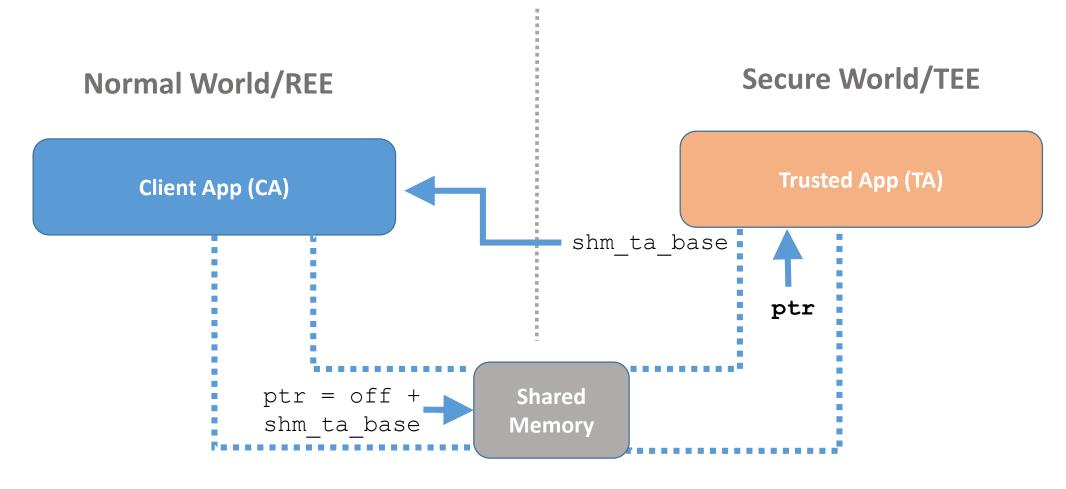
**Trusted App (TA)** 

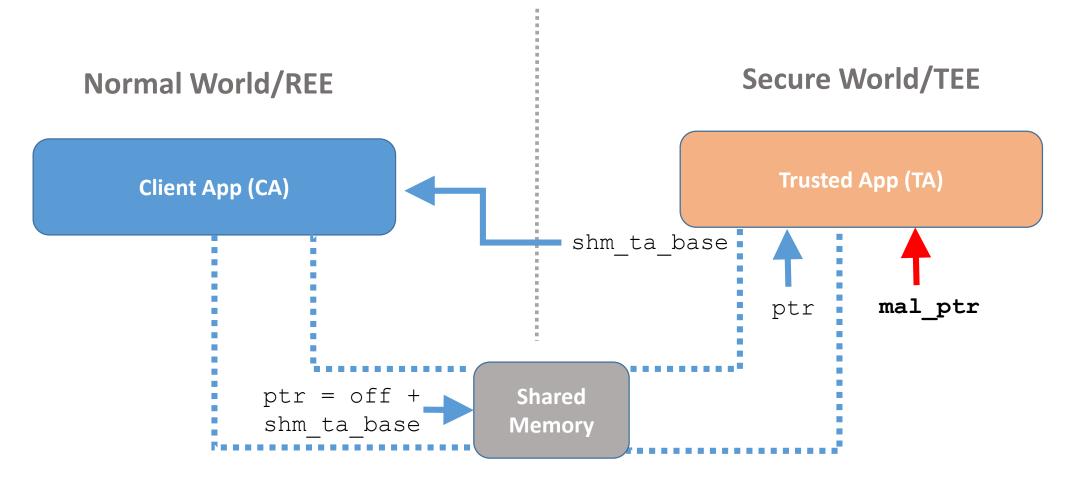


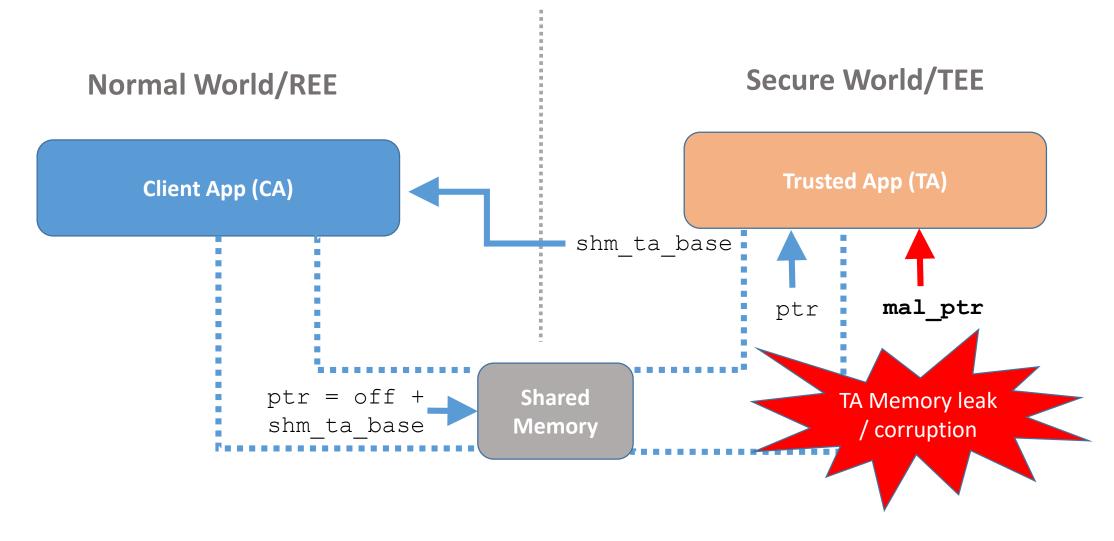


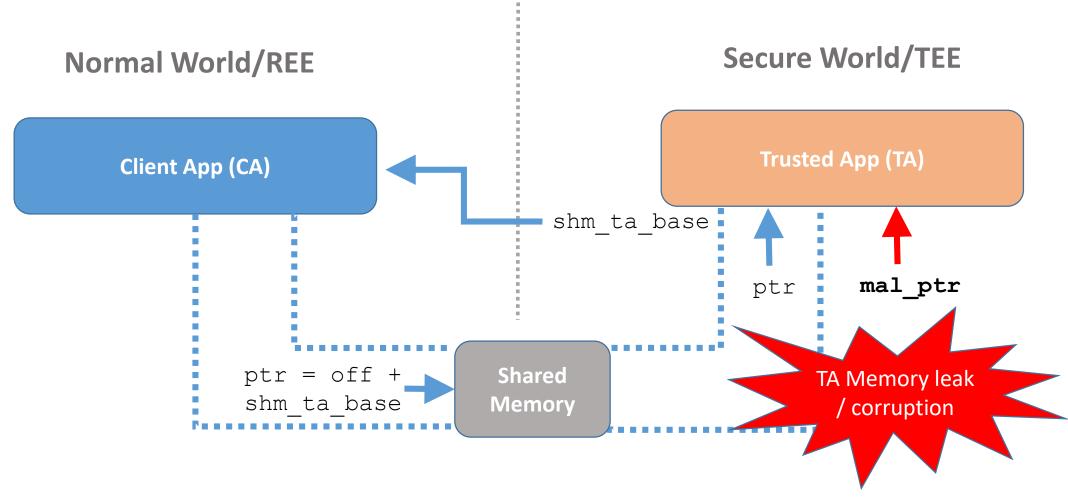












TA should check that CA-supplied pointers point to shared memory

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

```
TEE Result TA InvokeCommandEntryPoint(void *session, uint32 t cmd,
      uint32 t paramTypes, TEE Params params[4])
       // Use params[0] as a buffer
                                                              paramTypes(0) =
      request ptr = (struct request struct *) params[0];
                                                                     TEEC MEMREF;
       switch (request ptr->request) {
                                                              paramTypes(0) =
                                                                     TEEC VALUE;
                                                                   TA Memory leak
                                                                     / corruption
```

- GlobalPlatform TEE API allows 4 parameters in TA calls
  - Each parameter can be either a value or a pointer to a buffer

```
TEE Result TA InvokeCommandEntryPoint(void *session, uint32 t cmd,
      uint32 t paramTypes, TEE Params params[4])
                                                             paramTypes(0) =
      // Use params[0] as a buffer
      request ptr = (struct request struct *) params[0];
       switch (request ptr->request) {
                                                             paramTypes(0) =
                                                                    TEEC VALUE;
                                                                  TA Memory leak
                                                                   / corruption
 TA should check CA-supplied parameter types
```

 We showed that it is practically feasible to run real-world TZOSes and TAs in an emulator

 We showed that it is practically feasible to run real-world TZOSes and TAs in an emulator

 Large-scale fuzz testing of TAs using the emulator found several previously unknown vulnerabilities in TAs with high reproducibility

 We showed that it is practically feasible to run real-world TZOSes and TAs in an emulator

 Large-scale fuzz testing of TAs using the emulator found several previously unknown vulnerabilities in TAs with high reproducibility

- We identified vulnerability patterns unique to TA development
  - Pointing to the need for TZ-specific developer education

 We showed that it is practically feasible to run real-world TZOSes and TAs in an emulator

 Large-scale fuzz testing of TAs using the emulator found several previously unknown vulnerabilities in TAs with high reproducibility

- We identified vulnerability patterns unique to TA development
  - Pointing to the need for TZ-specific developer education

### Thank you!