

# **CS666-ASSIGNMENT - 4**

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**Group Details**: Group No. 3

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**Python Tool Version used:**  
**3.9.12(Spyder-Anaconda)**

## **Question:**

In this assignment, you have to implement a Differential Fault Attack on AES. You would be supplied with two pairs of faulty and correct ciphertext and using that you need to recover the first column (first 32 bits) of the round 10 key. The two pairs of faulty and correct ciphertext for each group are as given below:

## Group 3:

Correct Ciphertext1:

0x317982fa5666677f86b021f313e21725

Correct Ciphertext2:

0xeeade76ae853ccea45dddf257c63c0

Faulty Ciphertext1:

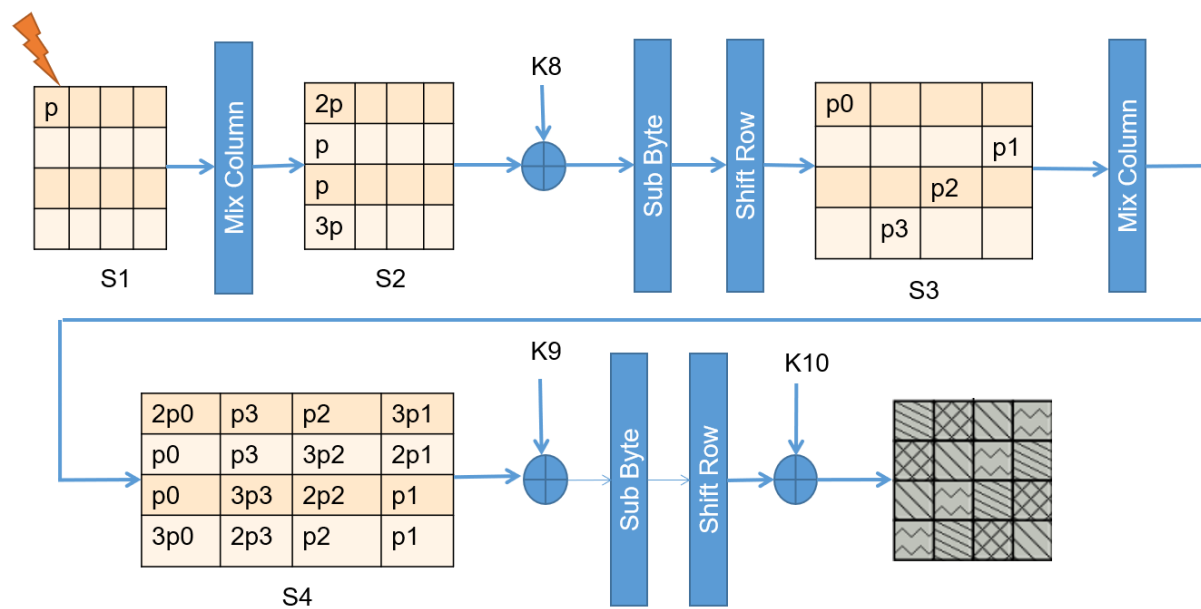
0x77f3dac61758cdb2cd9ab5d532d4ec8d

Faulty Ciphertext2:

0x0c1531c070f9303ef23fca1f5bab1007

## Logic:

The faults are assumed to be introduced before the MixColumn of round 8. So, the propagation of fault across various subsequent rounds of AES is shown in the image below.



The above image shows the propagation of fault when the fault is introduced in the first byte of the 8<sup>th</sup> round AES.

So, the idea is to introduce 4 different faults at 4 different locations to get 16 equations. Solution to these 16 equations will give the entire 128-bit key of the 10<sup>th</sup> round of AES.

Notations:

C - Correct Ciphertext

C\* - Faulty Ciphertext

f - Fault

S<sup>-1</sup> - Inverse Sub-box

So, firstly introducing fault at (0,0) before the 8<sup>th</sup> round MixColumn operation, this fault propagates in (0,0) (1,3) (2,2) (3,1) after MixColumn operation and AddRoundKey of 9<sup>th</sup> round, these 4 faults further propagate all over the AES box after MixColumn and AddRoundKey of 9<sup>th</sup> round, so we will have 16 equations out of which four will correspond to 1 fault, such for 4 faults we have 16 equations

For fault at (0,0) we get following 4 equations:

$$2p0 = S^{-1} (C_{0,0} \oplus K_{0,0}) \oplus S^{-1} (C^*_{0,0} \oplus K_{0,0})$$

$$p0 = S^{-1} (C_{1,3} \oplus K_{1,3}) \oplus S^{-1} (C^*_{1,3} \oplus K_{1,3})$$

$$p0 = S^{-1} (C_{2,2} \oplus K_{2,2}) \oplus S^{-1} (C^*_{2,2} \oplus K_{2,2})$$

$$3p0 = S^{-1} (C_{3,1} \oplus K_{3,1}) \oplus S^{-1} (C^*_{3,1} \oplus K_{3,1})$$

For fault at (3,1) we get following 4 equations:

$$p3 = S^{-1} (C_{0,1} \oplus K_{0,1}) \oplus S^{-1} (C^*_{0,1} \oplus K_{0,1})$$

$$p3 = S^{-1} (C_{1,0} \oplus K_{1,0}) \oplus S^{-1} (C^*_{1,0} \oplus K_{1,0})$$

$$3p3 = S^{-1} (C_{2,3} \oplus K_{2,3}) \oplus S^{-1} (C^*_{2,3} \oplus K_{2,3})$$

$$2p3 = S^{-1} (C_{3,2} \oplus K_{3,2}) \oplus S^{-1} (C^*_{3,2} \oplus K_{3,2})$$

For fault at (2,2) we get following 4 equations:

$$p2 = S^{-1} (C_{0,2} \oplus K_{0,2}) \oplus S^{-1} (C^*_{0,2} \oplus K_{0,2})$$

$$3p2 = S^{-1} (C_{1,1} \oplus K_{1,1}) \oplus S^{-1} (C^*_{1,1} \oplus K_{1,1})$$

$$2p2 = S^{-1} (C_{2,0} \oplus K_{2,0}) \oplus S^{-1} (C^*_{2,0} \oplus K_{2,0})$$

$$p2 = S^{-1} (C_{3,3} \oplus K_{3,3}) \oplus S^{-1} (C^*_{3,3} \oplus K_{3,3})$$

For fault at (1,3) we get following 4 equations:

$$3p1 = S^{-1} (C_{0,3} \oplus K_{0,3}) \oplus S^{-1} (C^*_{0,3} \oplus K_{0,3})$$

$$2p1 = S^{-1} (C_{1,2} \oplus K_{1,2}) \oplus S^{-1} (C^*_{1,2} \oplus K_{1,2})$$

$$p1 = S^{-1} (C_{2,1} \oplus K_{2,1}) \oplus S^{-1} (C^*_{2,1} \oplus K_{2,1})$$

$$p1 = S^{-1} (C_{3,0} \oplus K_{3,0}) \oplus S^{-1} (C^*_{3,0} \oplus K_{3,0})$$

Now, the overall approach is to solve the above 16 sets of equations to get the complete 128-bits of round 10 key out of which the first 32-bits required can be easily calculated by solving for the guessed key and fault, obtain the probable list of keys for both faulty ciphertexts and having a intersection of both those lists will give the correct key.

### **Answer:**

The total 128-bits of 10<sup>th</sup> round key found is =

**0x ade208dd 380df400 7b37a576 e918ed66**

Hence, the first 32-bits of 10<sup>th</sup> round key =

**0x ade208dd**

## Output:

```
----final key---  
OrderedDict([(0, '0xad'), (1, '0xe2'), (2, '0x8'), (3, '0xdd'), (4,  
'0x38'), (5, '0xd'), (6, '0xf4'), (7, '0x0'), (8, '0x7b'), (9, '0x37'),  
(10, '0xa5'), (11, '0x76'), (12, '0xe9'), (13, '0x18'), (14, '0xed'),  
(15, '0x66')])
```

Website Referred to check whether the 10<sup>th</sup> round key generated is correct or not by using the Master Key, **(0x 02a21a3ccdd2223ff75128421d1af222)** which was provided:

<https://www.cryptool.org/en/cto/aes-step-by-step>

Key													^
02a21a3c cdd2223f f7512842 1d1af222													
Expanded Key													^
02a21a3c	cdd2223f	f7512842	1d1af222	a12b8998	6cf9aba7	9ba883e5	86b271c7	94884fdc	f871e47b	63d9679e	e56b1659	efcf8405	
17be607e	746707e0	910c11b9	194dd284	0ef3b2fa	7a94b51a	eb98a4a3	4f04d86d	41f76a97	3b63df8d	d0fb7b2e	6025e91d	21d2838a	
1ab15c07	ca4a2729	f6e94c69	d73bcfe3	cd8a93e4	07c0b4cd	cc64f1ac	1b5f3e4f	d6d5adab	d1151966	8eb0c292	95effcdd	433a5176	
922f4810	ade208dd	380df400	7b37a576	e918ed66									