

Figure 1: Caption

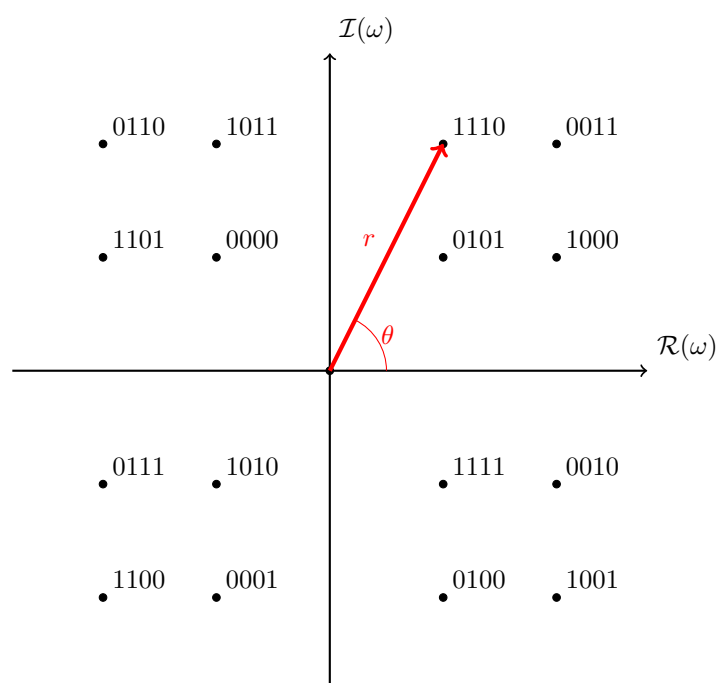
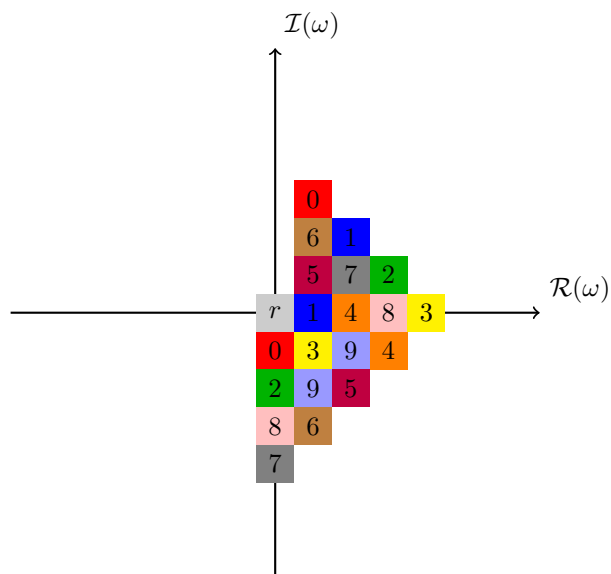


Figure 2: Encoding 4-bit words into frequencies.



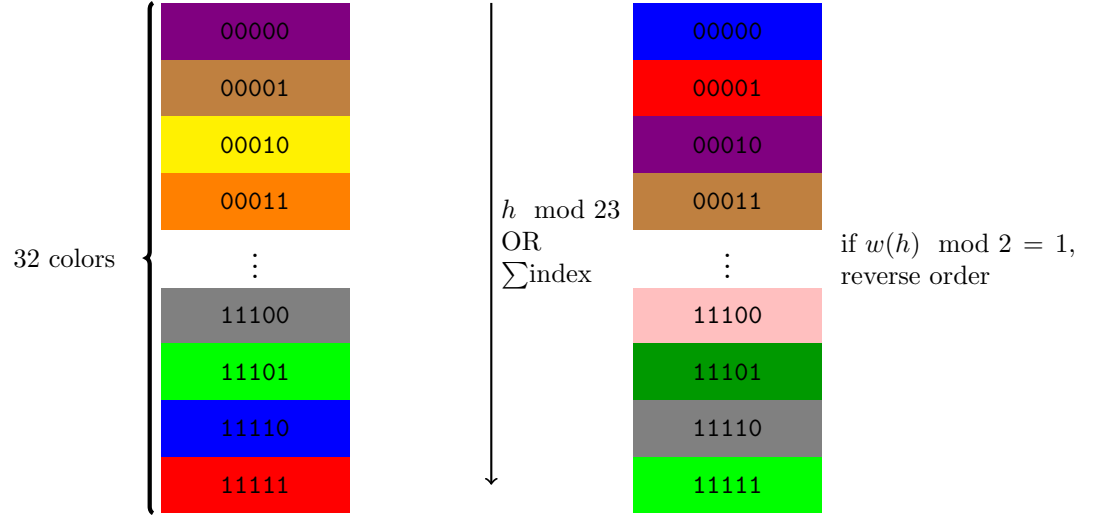
$$\mathcal{B}_m = I(f_m > 0.5), \quad m \in \{1, 2, 3\}$$

$$\mathcal{P}(\{\mathcal{B}_1, \mathcal{B}_2, \mathcal{B}_3\}) \rightarrow \{Color\#1, Color\#2, \dots, Color\#2^m\}$$

$$\begin{aligned}\{Color\#1, Color\#2, \dots, Color\#2^m\} = & \{f_1 \leq 0.5 \wedge f_2 \leq 0.5 \wedge f_3 \leq 0.5, \\ & f_1 \leq 0.5 \wedge f_2 \leq 0.5 \wedge f_3 > 0.5, \\ & \dots, \\ & f_1 > 0.5 \wedge f_2 > 0.5 \wedge f_3 > 0.5\end{aligned}$$

$$h = \underbrace{\underline{1100}}_{f_1} | \underbrace{\underline{1000}}_{f_2} | \underbrace{\underline{0110}}_{f_3} | \underline{1110} | \underline{0101} | \underline{1011} | \underline{0100} | \underline{1111} | \underline{0110} | \underline{0010} | \underline{1001} | \underline{0110} \dots \underline{0010} | \underline{1001} | \underline{0101} | \underline{1001} | \underline{0100}$$

$$p_i = \begin{cases} b_{0,i} \oplus b_{1,i} \oplus b_{2,i} \oplus \cdots \oplus b_{9,i} = \bigoplus_{k=0}^9 b_{k,i} & 0 \leq i < 12 \\ h_{120+i-12} \oplus h_{120+i-8} & 12 \leq i < 16 \end{cases}$$



$$\sum \text{index} = \sum_{k=0}^{128} I(b_k = 1)k \mod 31$$

Proof that  $\tilde{x} \neq x \mod 23$ , where  $\tilde{x}$  is  $x$  with 2 bits of the same group changed: Let  $\ell$  and  $\ell + 12m$  the indices of the 2 flipped bits, with  $0 < m \leq 10$ . Ring of  $\mathbb{Z}_p$  is an integral domain.

- Case  $b_\ell = 1, b_{\ell+12m} = 0$ :

$$\begin{aligned} x &= \tilde{x} = x - 2^\ell + 2^{\ell+12m} \mod 23 \\ \implies x &= x + 2^\ell (2^{12m} - 1) \mod 23 \\ \implies 0 &= 2^\ell (2^{12m} - 1) \mod 23 \\ \implies 0 &= 2^{12m} - 1 \mod 23 \\ \implies 1 &= 2^{12m} \mod 23 \end{aligned}$$

which has no solution for  $1 < m \leq 10$ .

- Case  $b_\ell = 0, b_{\ell+12m} = 1$ :

$$\begin{aligned} x &= \tilde{x} = x - 2^\ell + 2^{\ell+12m} \mod 23 \\ \implies x &= x + 2^\ell (1 - 2^{12m}) \mod 23 \\ \implies 0 &= 2^\ell (1 - 2^{12m}) \mod 23 \end{aligned}$$

which has no solution for  $1 < m \leq 10$ , same as the previous case.

- Case  $b_\ell = 1, b_{\ell+12m} = 1$ :

$$\begin{aligned}
x &= \tilde{x} = x + 2^\ell + 2^{\ell+12m} \pmod{23} \\
\implies x &= x + 2^\ell (1 + 2^{12m}) \pmod{23} \\
\implies 0 &= 2^\ell (1 + 2^{12m}) \pmod{23} \\
\implies 0 &= 1 + 2^{12m} \pmod{23} \\
\implies -1 &= 2^{12m} \pmod{23}
\end{aligned}$$

which has no solution for  $1 < m \leq 10$ .

- Case  $b_\ell = 0, b_{\ell+12m} = 0$ :

$$\begin{aligned}
x &= \tilde{x} = x - 2^\ell - 2^{\ell+12m} \pmod{23} \\
\implies x &= x + 2^\ell (-1 - 2^{12m}) \pmod{23} \\
\implies 0 &= 2^\ell (-1 - 2^{12m}) \pmod{23}
\end{aligned}$$

which has no solution for  $1 < m \leq 10$ , same as the previous case

Attacks if shift is  $\pmod{23}$ :

- Adv flips  $n$  bits of the same index, keeping same modulo 23 :
  - Effect : the same palette is used. If  $n$  is even and  $n > 2$ , then all colors are exactly the same as intra group parity is the same
  - "Impossible" for odd  $n$  as palette is flipped
  - Rare (I guess) for  $n = 4$ : maximum 8 if the ten bits are e.g. 1101011100
- Adv flips  $n$  bits, keeping same modulo 23 :
  - Effect : the same palette is used. The colors corresponding to the  $n$  indices are changed
  - "Impossible" for odd  $n$  as palette is flipped

$$\begin{aligned}
x + 2^k - 2^\ell &= x \pmod{p} \\
2^k - 2^\ell &= 0 \pmod{p} \\
2^\ell (2^{k-\ell} - 1) &= 0 \pmod{p} \\
2^{k-\ell} &= 1 \pmod{p} \\
k - \ell &= m \cdot |2| \pmod{p}, m \in \mathbb{Z}
\end{aligned}$$

Problem :  $11 \equiv -1 \pmod{12} \implies$  flipping  $b_i$  and  $b_{i+11}$  might only change 1 color.

Attacks if shift is  $\sum index$ :

Lots (sum of 4 indices that are for the same function divide 31)