Project2

Disk Based B+Tree



B+Tree

- B+Tree soruce code (originated http://www.amittai.com/prose/bpt.c)
 - Download the basic source code uploaded in piazza (bpt.zip)
 - Unzip the file.

```
[[hyeongwon@dev project]$ ls
bpt.zip
[[hyeongwon@dev project]$ unzip bpt.zip
Archive: bpt.zip
    creating: bpt/include/
    inflating: bpt/include/bpt.h
    creating: bpt/lib/
    inflating: bpt/Makefile
    creating: bpt/src/
    inflating: bpt/src/
    inflating: bpt/src/bpt.c
    inflating: bpt/src/main.c
[[hyeongwon@dev project]$ ls
bpt bpt.zip
```



B+Tree

- B+Tree soruce code (originated http://www.amittai.com/prose/bpt.c)
 - Compile the source file using 'Makefile'
 - If you don't have make util, type 'sudo apt-get install make'
 - Don't change the makefile unless you add another source file.
 - You will see the executable file (main) like below.

```
[hyeongwon@dev project]$ cd bpt
[hyeongwon@dev bpt]$ ls
include lib Makefile src
[hyeongwon@dev bpt]$ make
gcc -g -fPIC -I include/ -c -o src/main.o src/main.c
gcc -g -fPIC -I include/ -o src/bpt.o -c src/bpt.c
make static_library
make[1]: Entering directory `/home/hyeongwon/workspace/db.class/project/bpt'
ar cr lib/libbpt.a src/bpt.o
make[1]: Leaving directory `/home/hyeongwon/workspace/db.class/project/bpt'
gcc -g -fPIC -I include/ -o main src/main.o -L lib/ -lbpt
[hyeongwon@dev bpt]$ ls
include lib main Makefile src
```



Basic trial

• Execute it and try some basic commands.

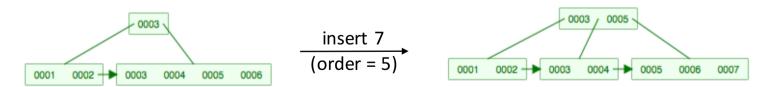


Basic trial

• You can adjust order by giving argument. (see the usage() functions)

```
$ ./main 5
...
3 | insert key 7 after inserting
1 2 | 3 4 5 6 | key 1 to 6 (sequentially)
3 5 |
1 2 | 3 4 | 5 6 7 |
>
```

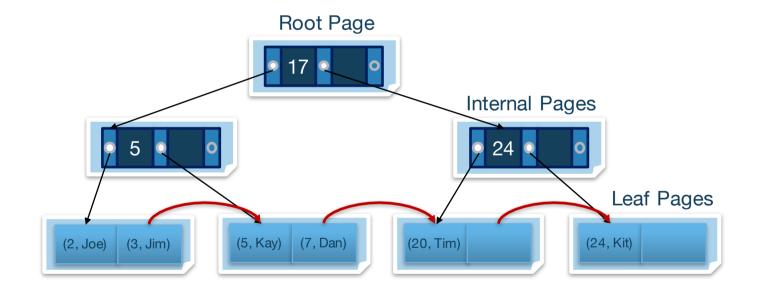
- You'd better understand the code fully before implementing the project.
- You can get some help from https://www.cs.usfca.edu/~galles/visualization/BPlusTree.html





Disk-based B+tree

- Note that current design only considers in-memory b+tree.
- Our goal is to implement disk-based b+ tree. (like below example)





- Implement 4 commands : open / insert / find / delete
- > There should be an appropriate **data file** in your system (you can call it a very simple database), maintaining disk-based b+ tree after serving those commands.

1. open <pathname>

- Open existing data file using 'pathname' or create one if not existed.
- All other 3 commands below should be handled after open data file.

2. insert <key> <value>

- Insert input 'key/value' (record) to data file at the right place.
- Same key should not be inserted (no duplicate).

A "record" means a <key/value> pair

3. find <key>

• Find the record containing input 'key' and return matching 'value' or 'null' if not found.

4. delete <key>

Find the matching record and delete it if found.



➤ Your library (libbpt.a) should provide those API services.

1. int open (char *pathname);

- Open existing data file using 'pathname' or create one if not existed.
- If success, return 0. Otherwise, return non-zero value.

2. int insert (int64_t key, char * value);

- Insert input 'key/value' (record) to data file at the right place.
- If success, return 0. Otherwise, return non-zero value.

char * find (int64_t key);

- Find the record containing input 'key'.
- If found matching 'key', return matched 'value' string. Otherwise, return NULL.

int delete (int64_t key);

- Find the matching record and delete it if found.
- If success, return 0. Otherwise, return non-zero value.



- ➤ All update operation (insert/delete) should be applied to your data file **as an operation unit.** That means one update operation should change the data file layout correctly.
- ➤ Note that your code must be worked as other students' data file. That means, your code should handle open(), insert(), find() and delete() API with other students' data file as well.
- >So follow the data file layout described from next slides.

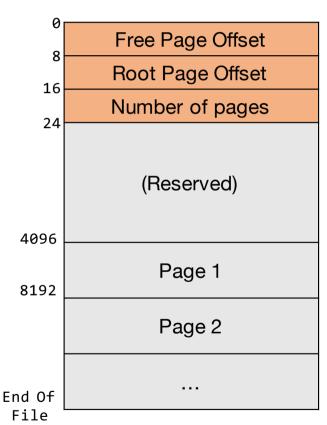


- ➤ We fixed the on-disk page size with **4096** Bytes.
- > We fixed the record (key + value) size with 128 (8 + 120) Bytes.
 - type : key => integer & value => string
- > There are 4 types of page. (detail next slides..)
 - **1. Header page** (special, containing metadata)
 - 2. Free page (maintained by free page list)
 - **3. Leaf page** (containing records)
 - Internal page (indexing internal/leaf page)



Header Page (Special)

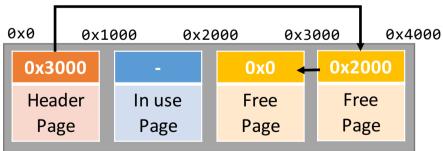
- ➤ Header page is the **first page (offset 0-4095)** of a data file, and contains metadata.
- ➤ When we open the data file at first, initializing disk-based b+tree should be done using this header page.
- Free page offset: [0-7]
 - points the first free page (head of free page list)
 - 0, if there is no free page left.
- Root page offset: [8-15]
 - pointing the root page within the data file.
- Number of pages: [16-23]
 - how many pages exist in this data file now.





Free Page

- ➤ In previous slide, header page contains the position of the first free page.
- Free pages are linked and allocation is managed by the free page list.
- Next free page offset: [0-7]
 - points the next free page.
 - 0, if end of the free page list.



Data file example



4096

8 Next Free Page Offset

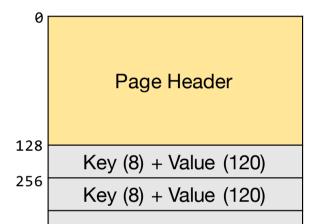
(Not used)



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Leaf Page

- ➤ Leaf page contains the **key/value records**.
- ➤ Keys are sorted in the page.
- ➤ One record size is 128 bytes and we contain maximum 31 records per one data page.
- First 128 bytes will be used as a page header for other types of pages. (see next slides)
- ➤ Branch factor (order) = 32



Key (8) + Value (120)

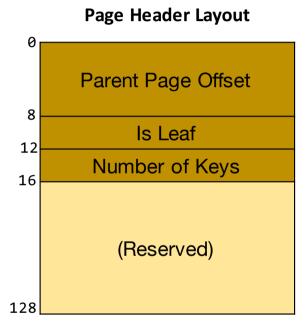
4096

Leaf Page Layout



Page Header

- ➤ Internal/Leaf page have first 128 bytes as a page header.
- Leaf/Internal page should contain those data (see the *node* structure in include/bpt.h)
 - Parent page offset [0-7]: If internal/leaf page, this field points the position of parent page.
 - Is Leaf [8-11]: 0 is internal page, 1 is leaf page.
 - Number of keys [12-15]: the number of keys within this page.

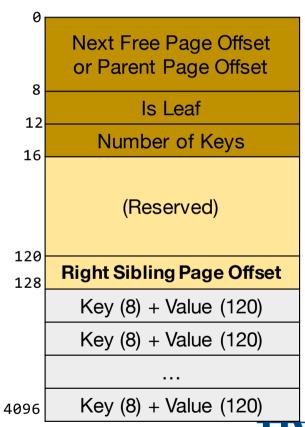




Leaf Page (Cont.)

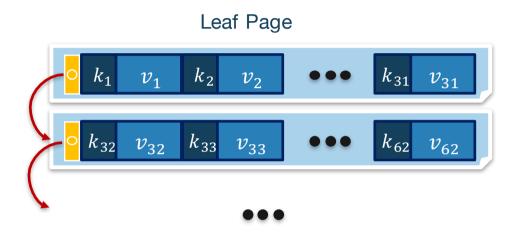
- ➤ We can say that the order of leaf page in disk-based b+tree is 32, but there is a minor problem.
- There should be one more offset added to store right sibling page offset for leaf page. (see the comments of *node* structure in include/bpt.h)
- ➤ So we define one special offset at the end of page header.
- ➤ If rightmost leaf page, right sibling page offset field is 0.

Leaf Page Layout



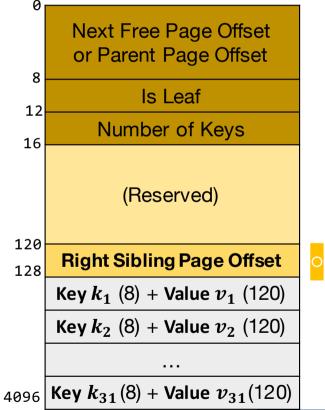
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Leaf Page (Cont.)





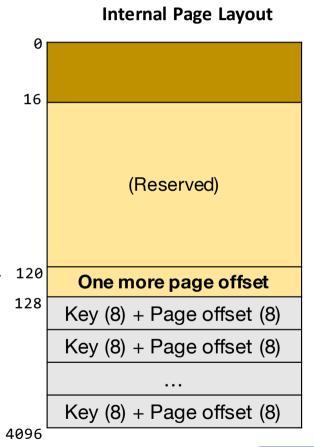
Leaf Page Layout



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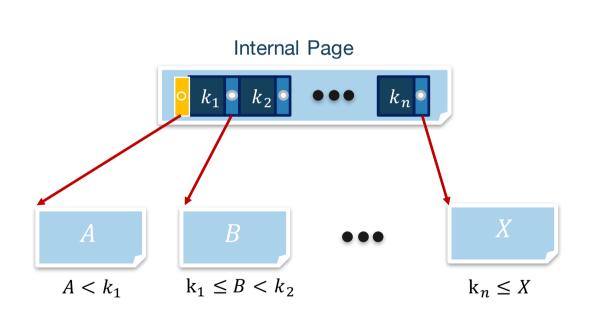
Internal Page

- Internal page is similar to leaf page, but instead of containing 120 bytes of values, it contains 8 bytes of another page (internal or leaf) offset.
- Internal page also needs one more page offset to interpret key ranges and we use the field which is specially defined in the leaf page for indicating right sibling.
- ➤ Branch factor (order) = 249
 - Internal page can have <u>maximum 248 entries</u>, because 'key + page offset' (8+8 bytes) can cover up to whole page (except page header) with the number of 248.
 - (4096-128)/(8+8)=248

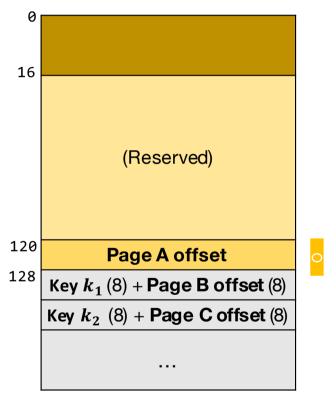




Internal Page



Internal Page Layout



4096



Disk-based B+tree Example

