

Computer Network and Distributed Systems

Spread Spectrum

Why Spread Spectrum ?

❑ In wired communication a narrowband signal is good enough to carry users data.

❑ But in wireless or mobile communication air being the medium of transmission, user data gets corrupt due to interference.

Spread Spectrum address this issue

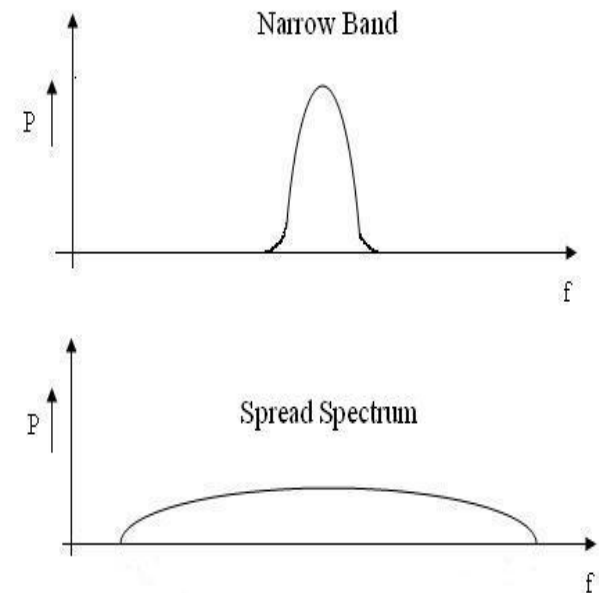
Principles of Spread Spectrum

❑ The basic principle of Spread Spectrum is to spread the bandwidth of the information signal to a its several orders of magnitude **greater** than the minimum required signal bandwidth

❑ It converts the information signal as noise like signal. And noise like signals are less susceptible to narrowband interference

❑ Spread signals are much wider, they are transmitted in lesser power density.

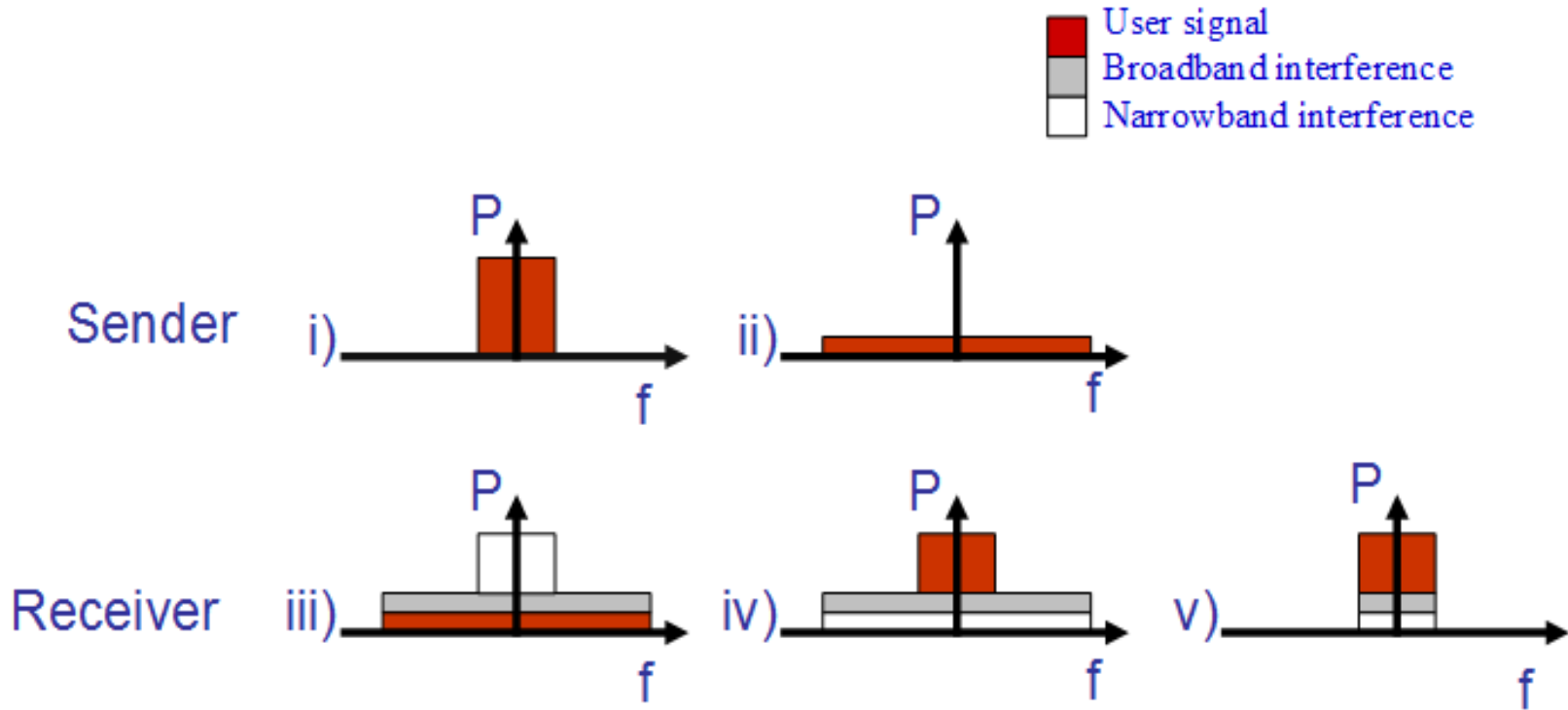
❑ Lesser power means lesser energy, lesser interference and greater capacity



Feature of Spread Spectrum

- ❑ Frequency range of the transmitted signal is deliberately varied, resulting in a much greater bandwidth than the signal would have if its frequency were not varied. Variation is done according to spreading sequences
- ❑ Power of the spread signal is the same as of the narrowband signal, resulting in a lower power spectral density (due to the higher bandwidth)
- ❑ Spread spectrum is very bandwidth **inefficient** for a single user
- ❑ But many users can **simultaneously** use the same bandwidth without significantly interfering with one another

Effects of spreading and interference



Classification

Direct Sequence Spread Spectrum (DSSS)

- Signal is spread before modulating it onto a carrier
- Each symbol is represented by a number of **chips**, resulting in a signal with a larger range of frequencies but with redundant data
- Spread signal is modulated onto a carrier signal with fixed carrier frequency

Frequency Hopping Spread Spectrum (FHSS)

- Signal is spread during modulation
- Carrier frequency is rapidly changed according to a hopping sequence

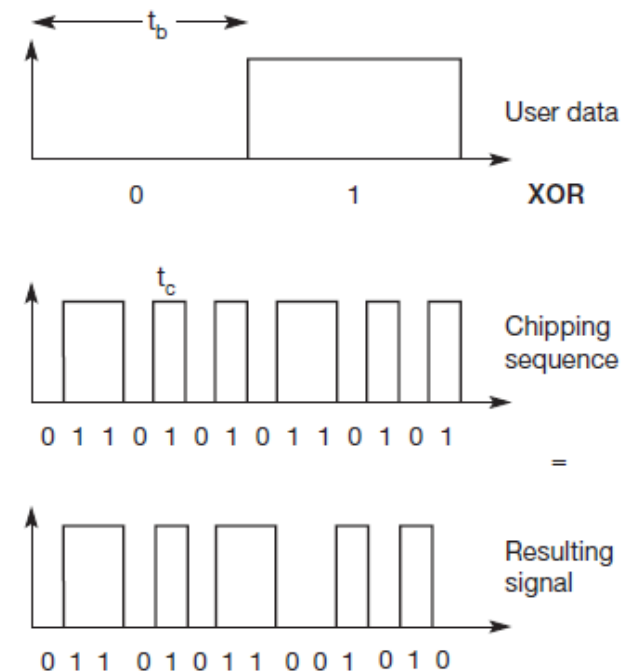
Direct Sequence Spread Spectrum (DSSS)

❑ DSSS systems take a user bit stream and perform an (XOR) with a so-called chipping sequence (the code) and transmit

❑ While each user bit has a duration t_b , the chipping sequence consists of smaller pulses, called chips, with a duration t_c .

❑ The **spreading factor $s = t_b/t_c$** determines the bandwidth of the resulting signal.

❑ If the original signal needs a bandwidth w , the resulting signal needs $s * w$ after spreading



Direct Sequence Spread Spectrum (DSSS) Contd

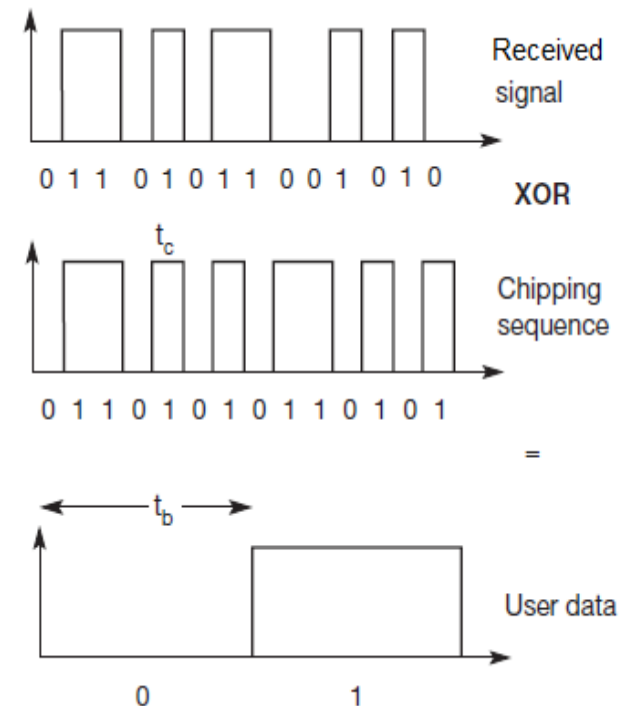
• **Spreading Factor/Gain (Process Gain) =**
Chip Rate / Data Rate

❑ In every system, the chip rate is fixed, hence data rate varies based on the size of spreading code

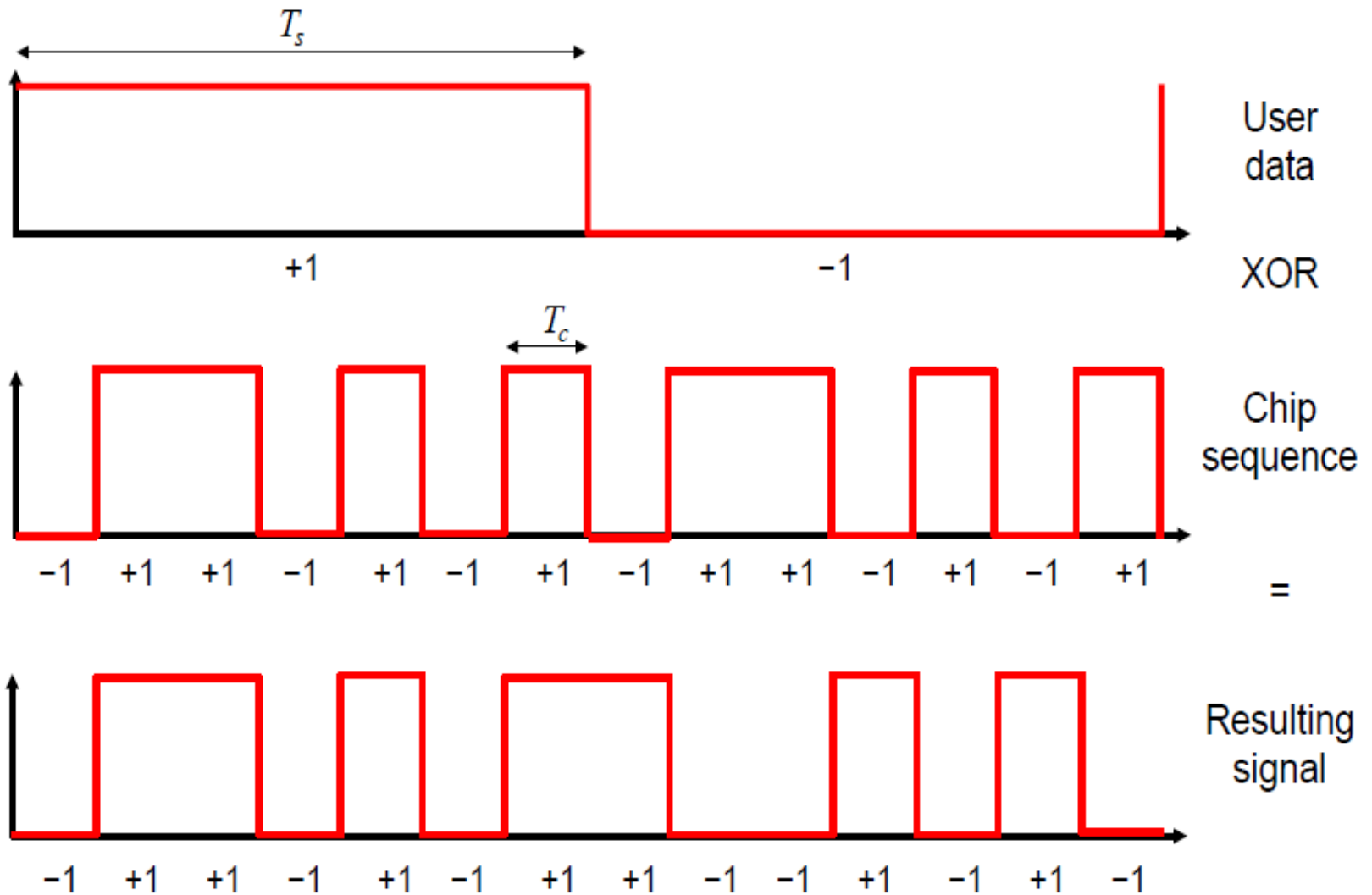
❑ Such that higher the size of the code used for spreading, lower the data rate achieved.

❑ At receiver, it takes the received chip stream and perform an (XOR) with a chipping sequence (the code) and generates the original user bit stream

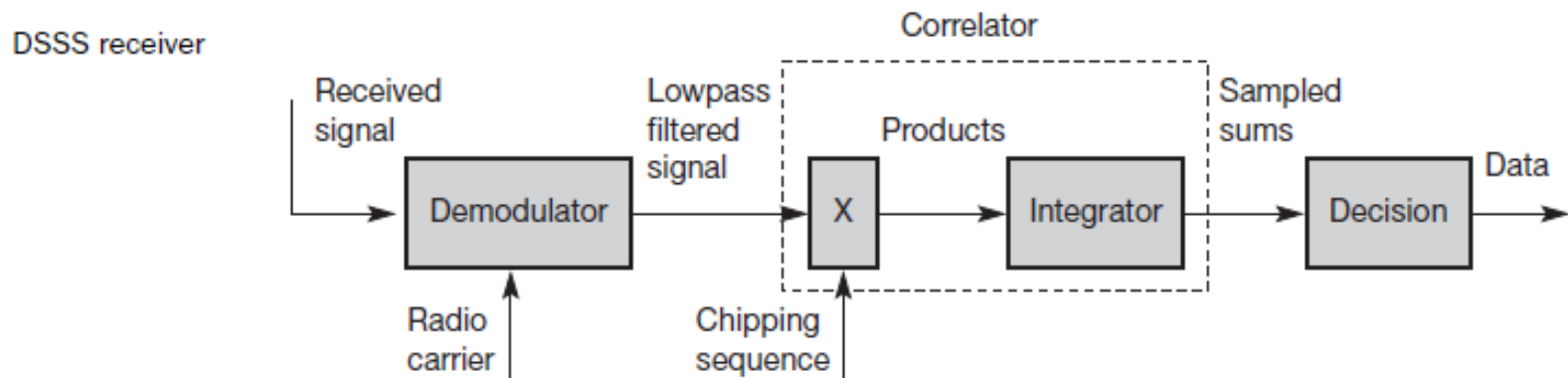
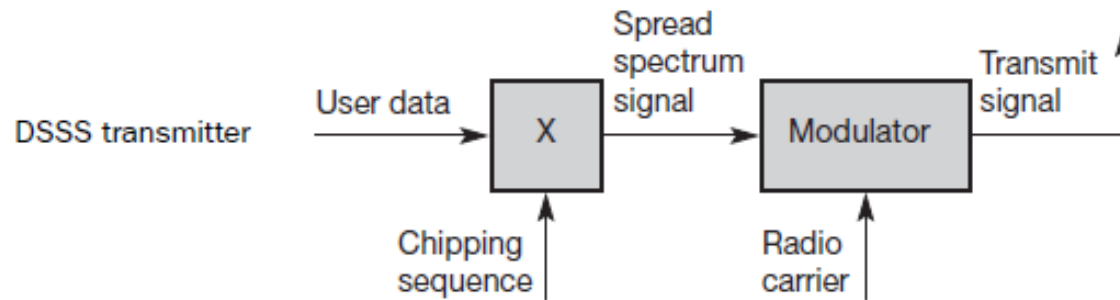
❑ Hence sender and receiver has to be exact synchronization



DSSS Spreading with a Chip Sequence (bipolar representation)



DSSS Transmission and Reception process



Correlation of Code Sequences

Correlation:

- Concept of determining how much **similarity** one set of data has with another
- Correlation is defined with a range of between -1 and 1 with the following meanings

<u>Correlation Value</u>	<u>Interpretation</u>
1	The second sequence matches the first sequence exactly
0	There is not relation at all between the two sequences
-1	The two sequences are mirror images of each other

Note: A spread spectrum sequence should have good correlation properties

Correlation of Code Sequences Contd

Auto Correlation:

Correlation of a code sequence C_i of N elements with all phase shifts of itself

$$\Phi_{ii}[n] = \frac{1}{N} \sum_{m=1}^N c_i[m] c_i[m+n]$$

with $n = 1..N$

Cross Correlation:

Correlation of two code sequences C_i and C_j

$$\Phi_{ij}[n] = \frac{1}{N} \sum_{m=1}^N c_i[m] c_j[m+n]$$

with $n = 1..N$

Auto-correlation

- A code has a good auto-correlation if the inner product with itself is large and its inner product with the same, but shifted code is small
- Good auto-correlation is essential to achieve synchronization between sender and receiver

Example: **Barker code** +1,-1,+1,+1,-1,+1,+1,+1,-1,-1,-1

n	$c_i[0], c_i[0+n]$	$\Phi_{ii}[n]$
0	(+1, -1,+1,+1, -1,+1,+1,+1, -1, -1, -1) (+1, -1,+1,+1, -1,+1,+1,+1, -1, -1, -1)	1
1	(+1, -1,+1,+1, -1,+1,+1,+1, -1, -1, -1) (-1,+1,+1, -1,+1,+1,+1, -1, -1, -1,+1)	-1/11
2	(+1, -1,+1,+1, -1,+1,+1,+1, -1, -1, -1) (+1,+1, -1,+1,+1,+1, -1, -1, -1,+1, -1)	-1/11

The Barker code has good auto-correlation. This is used for spreading in 802.11

Cross-correlation

- ❑ Two codes have a low cross-correlation if their product is low for all shift combinations
- ❑ Low cross-correlation between a sequence and noise is useful to the receiver in filtering out noise
- ❑ Low cross-correlation between two sequences is useful to the receiver to discriminate among signals generated by different users
- ❑ If the cross-correlation between two sequences is **0**, the sequences are said to be **full orthogonal**

Cross-correlation between c_1 and c_2 is $4/8 = 0.5$

$$\begin{array}{rcccccccc}
 c_1: & +1 & +1 & +1 & +1 & -1 & +1 & +1 & -1 \\
 \oplus & c_2: & -1 & +1 & -1 & +1 & +1 & -1 & -1 & +1 \\
 \hline
 & & -1 & +1 & -1 & +1 & -1 & -1 & -1 & =4
 \end{array}$$

*c_3 and c_4 are **orthogonal**, because their cross-correlation is 0*

$$\begin{array}{rcccccccc}
 c_3: & -1 & -1 & -1 & +1 & +1 & -1 & +1 & +1 \\
 \oplus & c_4: & -1 & -1 & +1 & -1 & +1 & +1 & -1 \\
 \hline
 & & +1 & +1 & -1 & -1 & +1 & -1 & -1 & =0
 \end{array}$$

DSSS CDMA system

CDMA systems use codes with certain characteristics to separate different users in code space and to enable access to a shared medium without interference.

But how to find such good codes ?

These code for a certain user should have a **good auto-correlation** and should be **orthogonal** to other codes (**cross-correlation is 0**)

$$\sum_{k=0}^{M-1} c_i[k] c_i[k] = 1$$

$$\sum_{k=0}^{M-1} c_i[k] c_j[k] = 0 \quad \text{for } i \neq j$$

Note: For CDM, each sequence as a spreading code, thereby providing zero cross correlations among all users.

Example of such codes are **Walsh codes**, **OVSF**(Orthogonal Variable Spreading Factor)

CDMA transmission and reception analogy

Tx 1 : Information Signal **S1**, Spreading Code **C1**. Transmits (**S1.C1**) over the air

Tx 2 : Information Signal **S2**, Spreading Code **C2**. Transmits (**S2.C2**) over the air

From Spreading code's auto-correlation and orthogonal property -

$$\mathbf{C1.C1 = 1}$$

$$\mathbf{C2.C2 = 1}$$

$$\mathbf{C1.C2 = 0}$$

When signal transmitted from two source, over the air it becomes -

$$\mathbf{(S1.C1 + S2.C2)}$$

Rx 1 : Despreading on reception $\mathbf{(S1.C1 + S2.C2).C1 = S1.C1.C1 + S2.C2.C1 = S1}$

Rx 2 : Despreading on reception $\mathbf{(S1.C1 + S2.C2).C2 = S1.C1.C2 + S2.C2.C2 = S2}$

Walsh Codes

- ❑ Set of Walsh codes of length n consists of the rows of a n×n Hadamard matrix
- ❑ Matrix is recursively defined as : $W_1 = (+1) \quad W_{2n} = \begin{pmatrix} W_n & W_n \\ W_n & \overline{W_n} \end{pmatrix}$
 - Where n is the dimension of the matrix and the overscore denotes the logical NOT of the bits in the matrix

Example for n=2,4 :

$$W_2 = \begin{pmatrix} +1 & +1 \\ +1 & -1 \end{pmatrix} \quad W_4 = \begin{pmatrix} +1 & +1 & +1 & +1 \\ +1 & -1 & +1 & -1 \\ +1 & +1 & -1 & -1 \\ +1 & -1 & -1 & +1 \end{pmatrix}$$

All users using a set of orthogonal codes must be synchronized, because the cross-correlation between different shifts of Walsh sequences is not zero (refer last two rows of W4 matrix)

IS-95 uses 64x64 Walsh code

Generating chip sequences using Walsh Code

$$W_1 = \begin{bmatrix} +1 \end{bmatrix} \qquad W_{2N} = \begin{bmatrix} W_N & W_N \\ W_N & \overline{W_N} \end{bmatrix}$$

$$W_1 = \begin{bmatrix} +1 \end{bmatrix}$$
$$W_2 = \begin{bmatrix} +1 & +1 \\ +1 & -1 \end{bmatrix}$$
$$W_4 = \begin{bmatrix} +1 & +1 & +1 & +1 \\ +1 & -1 & +1 & -1 \\ +1 & +1 & -1 & -1 \\ +1 & -1 & -1 & +1 \end{bmatrix}$$

Orthogonal chip sequences

- Encoding rules:

- If station has to send data bit 0: transmit signal -1
- If station has to send data bit 1: transmit signal +1
- If station does not have data to send: transmit signal 0

+1, +1, +1, +1

A

+1, -1, +1, -1

B

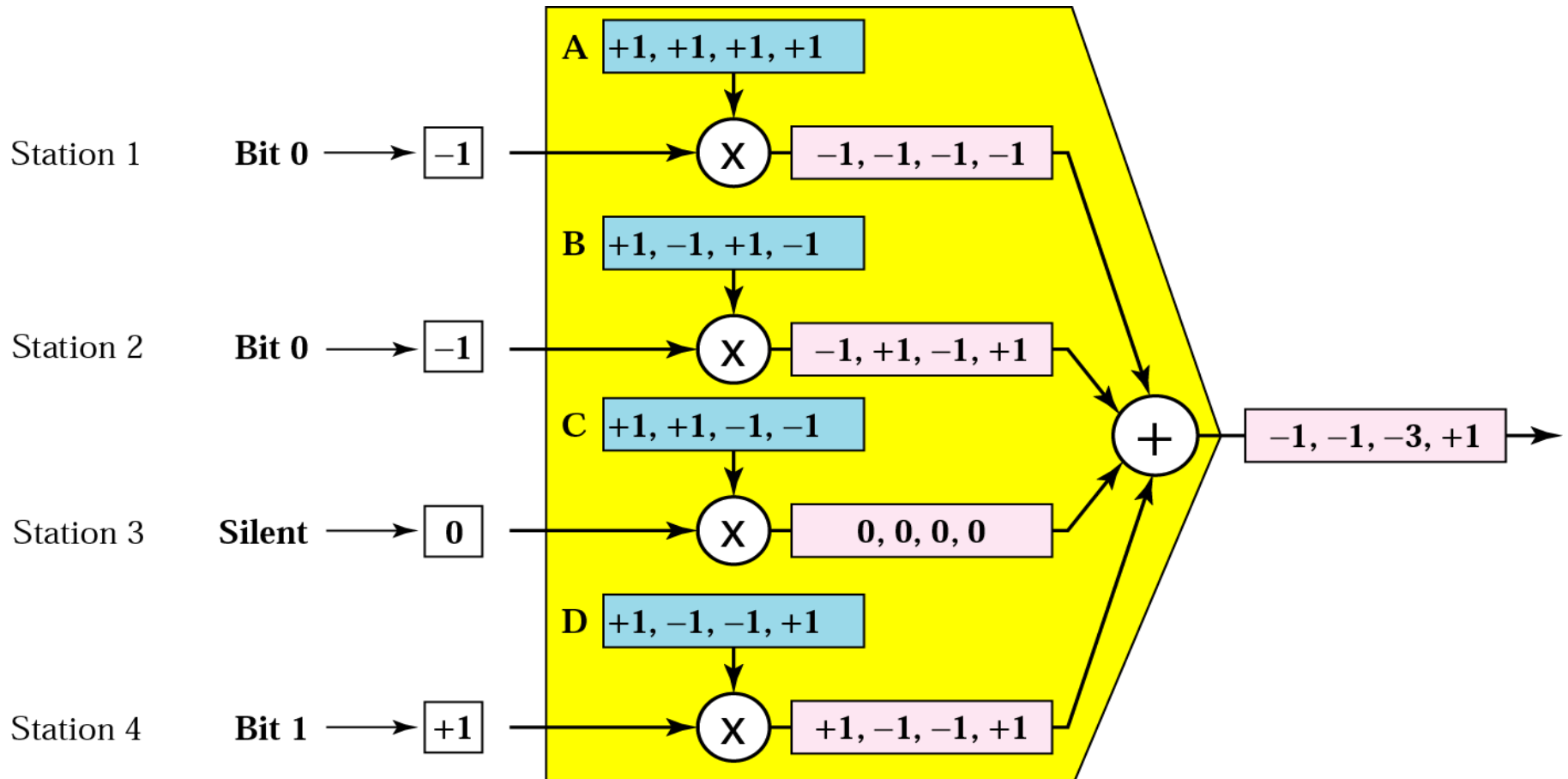
+1, +1, -1, -1

C

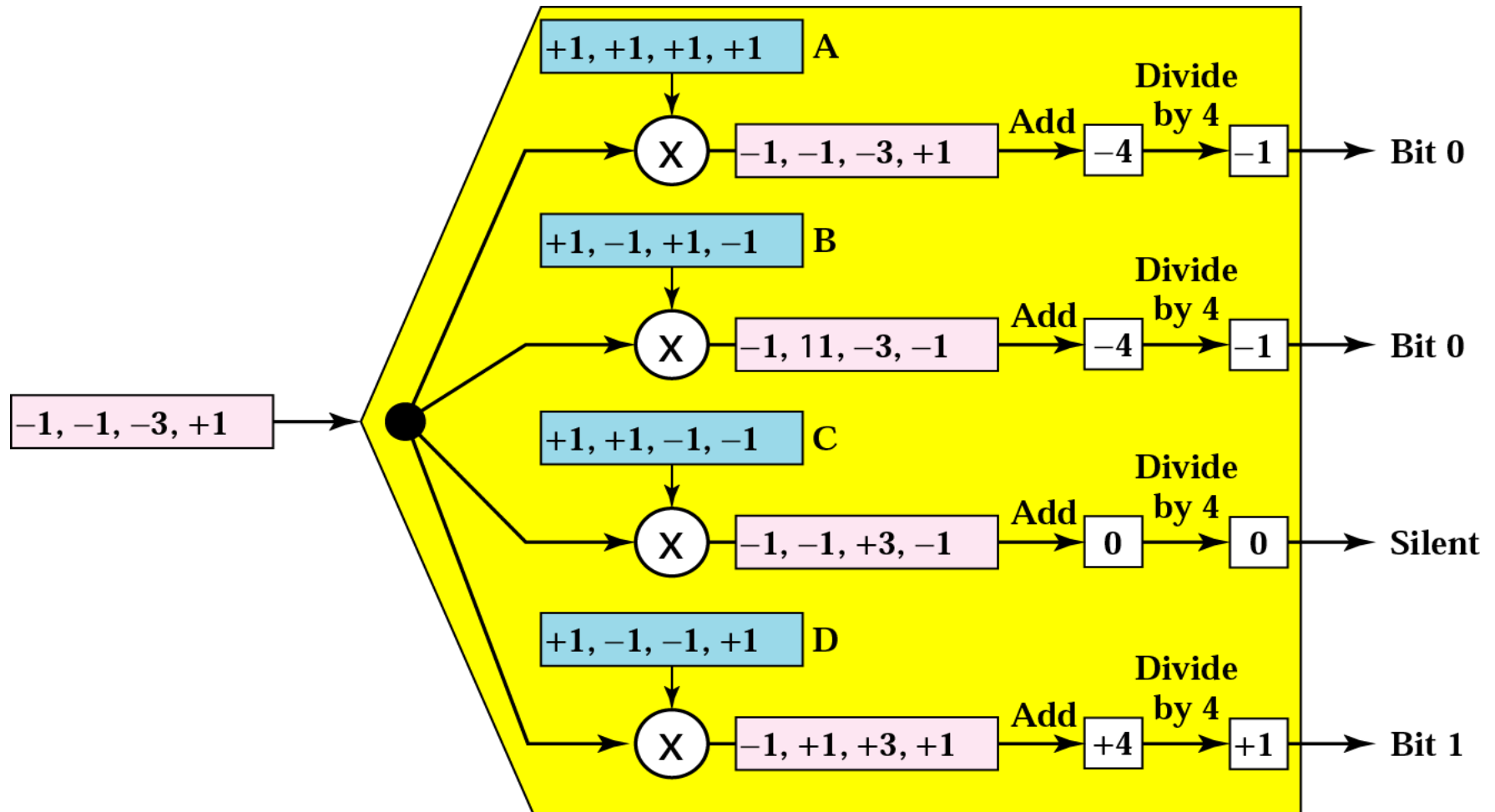
+1, -1, -1, +1

D

CDMA Multiplexer



CDMA De-Multiplexer



For your study

- ✓ Orthogonal Variable Spreading Factor (OVSF)
- ✓ Frequency Hopping Spread Spectrum (FHSS)

Reference

- ❑ *Mobile Communication, Jochen Schiller*
- ❑ *Data Communications & Networking, 5th Edition, Behrouz A. Forouzan*