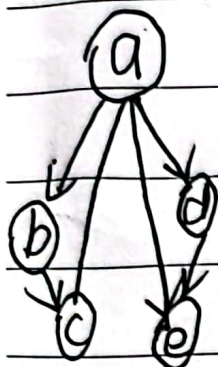


1.



2. 设其形式定义为 $DS = (D, R)$

$D = \{1, 2, 3, 4, 5\}, R = \{(1, 2), (1, 3), (2, 5), (3, 4), (3, 5)\}$

3. $O(n \log n)$

4. $O(n^2)$

5. $2 < \log n < n^{\frac{1}{2}} < 20n < 4n^2 < 3^n < n!$

6. (1) 使 $2^n > n^6$

②. $n/n^2 \geq 6/n$

$n \geq 6 \log_2 n$

当 $n \geq 32, n \geq 6 \log_2 n, 2^n \geq n^6$

\therefore 当 $n \geq 32, C_1 2^n + C_2 n^6 \leq (C_1 + C_2) n^6$

$\therefore C_1 2^n + C_2 n^6 = O(2^n)$, 其中 $n_0 = 32, C = C_1 + C_2$

同理: 当 $n \geq 32, C_1 2^n + C_2 n^6 \geq (C_1 + C_2) n^6$

$\therefore C_1 2^n + C_2 n^6 = \Omega(n^6)$, 其中 $n_0 = 32, C = C_1 + C_2$

(2). 易知, 当 $n \geq 10, n \lg n \geq n$

\therefore 当 $n \geq 10, C_4 n \lg n + C_5 n \leq (C_4 + C_5) n \lg n$

$\therefore C_4 n \lg n + C_5 n = O(n \lg n)$, 其中 $n_0 = 10, C = C_4 + C_5$

当 $n \geq 10, C_4 n \lg n + C_5 n \geq (C_4 + C_5) n$

$\therefore C_4 n \log n + C_5 n = \Omega(n)$, 其中 $n_0 = 10$, $C = C_4 + C_5$

7. (1)

$$T(n) = 1 + 1 + \sum_{i=0}^{n^2-1} 4 + 1$$
$$= 4n^2 + 3$$

$$\therefore T(n) = \Theta(n^2)$$

(2)

~~设~~ 设 Random 时间常数函数

$$T(n) = 1 + \sum_{i=0}^{n-1} \left(3 + \sum_{j=0}^{n-1} (4 + k) + [n \log n] \right) + 1$$
$$= n^2(4+k) + 4n + n^2 \log n + 2$$

$$\therefore T(n) = \Theta(n^2 \log n)$$