

Data Structure - Spring 2022

4. Introduction to DS and Algorithm

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Based on:

Goodrich, Chapter 3-4
Karumanchi, Chapter 1-2
Slides by Prof. Yung Yi, KAIST



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Variables and Data Types

- **Variable**

- Name that hold value (data)

- **Data type**

- Set of data with predefined values
- Ex) integer, floating point, character, string, etc.

- **System-defined data types**

- Also known as primitive data types
- int, float, str, etc.

- **User-defined data types**

```
class Student:
    def __init__(self):
        self.id = 20212345
        self.name = "Mr"
        self.cgpa = 4.20
        self.year = 2
```

Data Structure

- **Data structure**
 - Particular way of storing and organizing data
 - so that it can be used efficiently
- **Linear data structures**
 - Elements are accessed sequentially
 - Linked lists, stacks, queues, etc.
- **Non-linear data structures**
 - Stored/accessed in non-sequential order
 - Trees, graphs, etc.
- **Abstract data type (ADT)**
 - Data structure + its operation

Algorithm and its Performance

- **Algorithm**
 - Step-by-step unambiguous instructions to solve a problem
- **Program**
 - Data structure + algorithm
- **Performance analysis**
 - Running time
 - (also, memory, etc.)

```
Algorithm. Sum of Array  
sum  $\leftarrow$  0  
for i = 1 to N  
    sum  $\leftarrow$  sum + Array[i]
```

Running Time Analysis

- **Goal of analysis**

- How processing time increases as the size of problem (input size) increases
- Input size: size of an array, elements in a matrix, degree of polynomial, etc.

- **Experimental analysis**

- Execution time
- Specific to a computer

- **Number of statements**

- Specific to a language

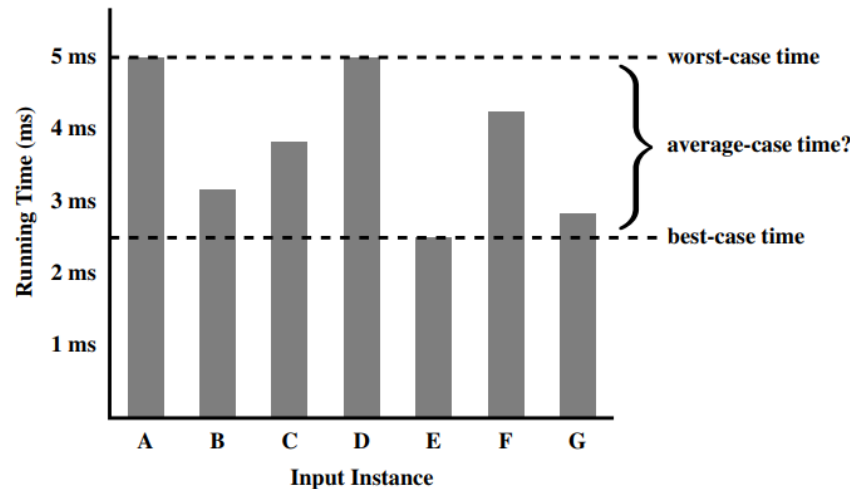
- **Number of primitive operations**

- Generally applicable
- Assignment, arithmetic-ops, function-call, return, etc.

```
from time import time
start_time = time( )
run algorithm
end_time = time( )
elapsed = end_time - start_time
```

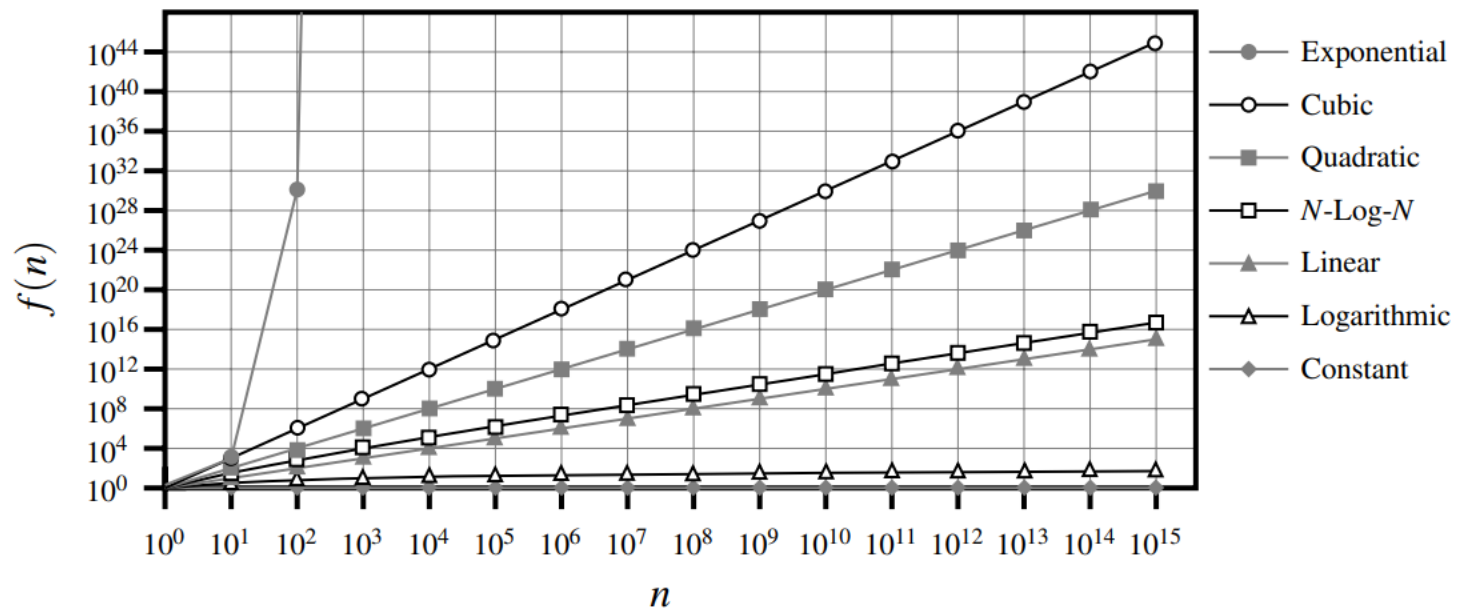
Theoretical Analysis

- **Number of operations as a function of input size**
 - $f(n)$: n is the input size
 - Captures the growth rate of running time
- **Focuses on the worst-case input**
 - Problems with average-case analysis?



Common Functions

constant	logarithm	linear	n -log- n	quadratic	cubic	exponential
1	$\log n$	n	$n \log n$	n^2	n^3	a^n



Running Time Analysis: Example

```
1 def unique1(S):
2     """ Return True if there are no duplicate elements in sequence S."""
3     for j in range(len(S)):
4         for k in range(j+1, len(S)):
5             if S[j] == S[k]:
6                 return False           # found duplicate pair
7     return True                       # if we reach this, elements were unique
```


Some useful summation series

Arithmetic series

$$\sum_{k=1}^n k = 1 + 2 + \dots + n = \frac{n(n+1)}{2}$$

Geometric series

$$\sum_{k=0}^n x^k = 1 + x + x^2 + \dots + x^n = \frac{x^{n+1} - 1}{x - 1} \quad (x \neq 1)$$

Harmonic series

$$\sum_{k=1}^n \frac{1}{k} = 1 + \frac{1}{2} + \dots + \frac{1}{n} \approx \log n$$

Other important formulae

$$\sum_{k=1}^n \log k \approx n \log n$$

$$\sum_{k=1}^n k^p = 1^p + 2^p + \dots + n^p \approx \frac{1}{p+1} n^{p+1}$$

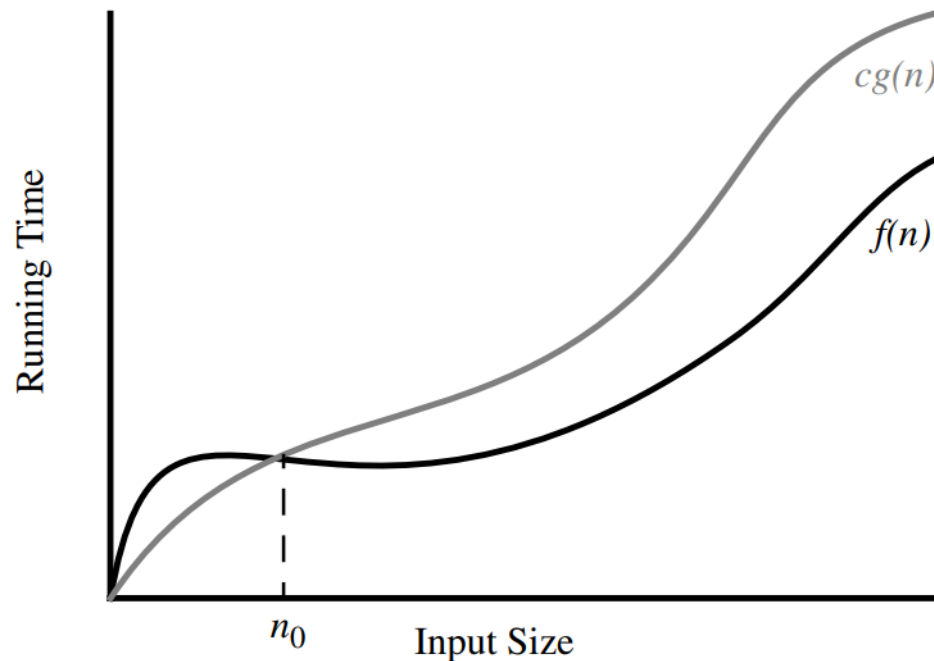
Asymptotic Analysis

- **The “big-picture” approach**
 - Only focuses on the main factor determining the growth
 - $2n^2 \approx n^2, 2n + 8 \approx n$
- **Running time of the following function**
 - Is proportional to n

```
1 def find_max(data):
2     """Return the maximum element from a nonempty Python list."""
3     biggest = data[0]           # The initial value to beat
4     for val in data:           # For each value:
5         if val > biggest:       # if it is greater than the best so far,
6             biggest = val       # we have found a new best (so far)
7     return biggest             # When loop ends, biggest is the max
```

“Big-Oh” Notation

- **$f(n)$ is $O(g(n))$**
 - if $f(n) \leq c \cdot g(n)$, for $n \geq n_0$
 - ($c > 0$)



Some properties of Big-Oh

$5n^4 + 3n^3 + 2n^2 + 4n + 1$ is $O(n^4)$. How can we justify?

If $f(n)$ is a polynomial of degree d , that is,

$$f(n) = a_0 + a_1n + \cdots + a_dn^d,$$

and $a_d > 0$, then $f(n)$ is $O(n^d)$.

Big-Oh representation practice

$$5n^2 + 3n \log n + 2n + 5$$

$$3 \log n + 2$$

$$2^{n+2}$$

$$2n + 100 \log n$$

Example

```
1 def unique1(S):
2     """ Return True if there are no duplicate elements in sequence S."""
3     for j in range(len(S)):
4         for k in range(j+1, len(S)):
5             if S[j] == S[k]:
6                 return False           # found duplicate pair
7     return True                       # if we reach this, elements were unique
```

Example

```
1 def unique2(S):
2     """ Return True if there are no duplicate elements in sequence S."""
3     temp = sorted(S)           # create a sorted copy of S
4     for j in range(1, len(temp)):
5         if S[j-1] == S[j]:
6             return False      # found duplicate pair
7     return True              # if we reach this, elements were unique
```

Big-Oh comparison

- **Algorithm 1. unique1**
 - $O(n^2)$
- **Algorithm2. unique2**
 - $O(n \log n)$
- **Which algorithm is better/faster?**